

Q.1.Pg ①Version (2) (Monte Carlo Roulette I)

Total possibilities = 37

No. of reds = 18

$$P(\text{win}) = \frac{18}{37}$$

Note:- In any roll, red \rightarrow wins
gets you
black - lose - money back

$$\begin{aligned} P(\text{lose}) &= P(\text{lose} | \text{black}) P(\text{black}) + P(\text{lose} | P_1) P(P_1) \\ &\quad \xleftarrow{\text{Law of total probability}} \\ &= 1 \times \frac{18}{37} + \frac{19}{37} \times \frac{1}{37} = 0.50036523 \end{aligned}$$

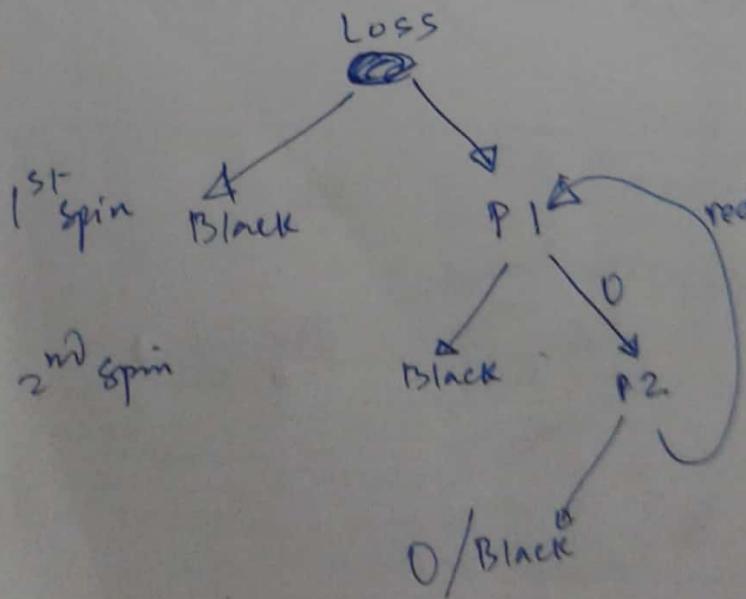
Version (3)

$$P(\text{lose}) = P(\text{lose} | \text{black}) P(\text{black}) + P(\text{lose} | P_1) P(P_1)$$

Same as before

$$\begin{aligned} \text{2nd spin Black} &= B_2 \\ \text{2nd spin } O &= O_2 \quad (\text{sent to } P_2) \end{aligned}$$

$$\begin{aligned} &\left\{ P(\text{lose} \& \text{2nd spin in black} | P_1) \right. \\ &\quad \left. + P(\text{lose} \& \text{2nd spin in } O | P_1) \right\} P(P_1) \\ &= \left\{ P(\text{lose} | B_2, P_1) \times P(B_2 | P_1) \right. \\ &\quad \left. + P(\text{lose} | O_2, P_1) \times P(O_2 | P_1) \right\} \end{aligned}$$



$$\begin{aligned} P(\text{lose}) &= \left\{ 1 \times \frac{18}{37} + \left[\frac{19}{37} + P(\text{lose} | P_1) \times \frac{1}{37} \right] \times \right. \\ &\quad \left. \frac{P(\text{lose} | P_2)}{P(P_2)} \times P(P_1) \right\} \times P(P_1) \\ &= \frac{18}{37} + \left[\frac{19}{37} + P(\text{lose} | P_1) \times \frac{18}{37} \right] \times \frac{1}{37} \end{aligned}$$

$$\therefore P(\text{lose} | P_1) \times \frac{1}{37} = \frac{1}{37} \left\{ \underbrace{\frac{18}{37}}_{1} + \left[\frac{19}{37} + P(\text{lose} | P_1) \times \frac{18}{37} \right] \right\}$$

$$\Rightarrow P(\text{lose} | P_1) = \frac{18}{37} + \frac{19}{37} \times \frac{1}{37} + P(\text{lose} | P_1) \frac{18}{37 \times 37}$$

$$\Rightarrow P(\text{lose} | P_1) \left\{ 1 - \frac{18}{37 \times 37} \right\} = \frac{18 \times 37 + 19}{37 \times 37}$$

$$\Rightarrow P(\text{lose} | P_1) \left\{ \frac{37 \times 37 - 18}{37 \times 37} \right\} = \frac{666 + 19}{37 \times 37}$$

$$\Rightarrow 1351 P(\text{lose} | P_1) = 685$$

$$\Rightarrow P(\text{lose} | P_1) = \frac{685}{1351}$$

$$\therefore P(\text{lose}) = P(\text{lose} | \text{Black}) P(\text{Black}) + P(\text{lose} | P_1) P(P_1)$$

$$= \frac{18}{37} + \frac{685}{1351} \times \frac{1}{37}$$

$$= 0.50019004$$

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W/ Q2)

Pg ②

Bayes' theorem (Witness reliability)

$$P(\text{taxi} = \text{Black}) = 0.8 = P(\text{taxi} = B)$$

$$P(\text{taxi} = \text{Yellow}) = 0.2 = P(\text{taxi} = Y)$$

$$P(\text{Witness} = \text{true} \mid \text{taxi color}) = \frac{8}{10} = 0.8$$

$$\text{i.e. } P(\text{Witness taxi} = Y \mid \text{taxi} = Y) = 0.8$$

$$P(\text{Witness taxi} = B \mid \text{taxi} = B) = 0.8$$

for reliability ^{of info}, we want $P(\text{taxi} = Y \mid \text{Witness taxi} = Y)$.

i.e. We need to find $P(\text{taxi} = Y \mid \text{Witness taxi} = Y)$

Bayes' thm

$$P(\text{taxi} = Y \mid \text{Witness taxi} = Y) = \frac{P(\text{taxi} = Y, \text{Witness taxi} = Y)}{P(\text{Witness taxi} = Y)}$$

$$= \frac{P(\text{Witness taxi} = Y \mid \text{taxi} = Y)P(\text{taxi} = Y)}{P(\text{Witness taxi} = Y)}$$

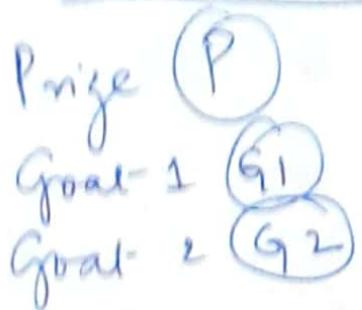
law of
total prob.
compute
Denominator

$$= \frac{0.8 \times 0.2}{P(\text{Witness taxi} = Y \mid \text{taxi} = Y) + P(\text{Witness taxi} = Y \mid \text{taxi} = B)}$$

Conclusion :-

This means it is a 50-50 situation (chance)
& hence witness report = $\frac{0.8 \times 0.2}{0.8 \times 0.2 + 0.2 \times 0.8} = \frac{1}{2}$

Q.3. (Monty Hall problem)



Some assumptions, that you as a Contestant² must make:-

Before the game,

- (i) The latch is placed randomly behind any door i.e. the contestant upon choosing a door has $\text{prob} = \frac{1}{3}$ of winning.
- (ii) The host show knows behind which door(s) the prize is & always opens an empty door. If he has 2 empty doors he can open; he choose one of them at random.

options :-

Door 1	Door 2	Door 3	
P	G1	G2	Arrangement 1
P	G2	G1	Arrangement 2
G1	P	G2	Arrangement 3
G2	P	G1	Arrangement 4
G1	G2	P	Arrangement 5
G2	G1	P	Arrangement 6

Soln :-

Theorem(1) :- Elementary Probability.

B/c of assumption (i); all 6 arrangements have equal probability = $\frac{1}{6}$

In 2nd arrangement - he will lose if he switches.

In 3rd arrangement - host will open door 3;
he will win if he switches

Same w/ 4th, 5th & 6th arrangement } → he will win if he
switches

∴ 4 out of 6 cases; he will win if he
(or $\frac{2}{3}$ prob) switches

i.e. He should switch !!

Method (2) :- Baye's Th

Let us start w/ arrangement ① :-

$$P(\text{keep} \& \text{win}) = P(\text{prize door 1} \mid \text{host door 3})$$

$$\text{Baye's} = \frac{P(\text{host door 3} \mid \text{prize door 1})P(\text{prize door 1})}{P(\text{host door 3})}$$

Assumption
(i&ii)

$$= \frac{\frac{1}{2} \times \frac{1}{3}}{P(\text{host door 3} \mid \text{prize door 1})P(\text{prize door 1}) + P(\text{host door 3} \mid \text{prize door 2})P(\text{prize door 2}) + P(\text{host door 3} \mid \text{prize door 3})P(\text{prize door 3})}$$

$$\text{i.e. } P(\text{switch} \& \text{win}) = \frac{\frac{1}{2} \times \frac{1}{3}}{\frac{1}{2} \times \frac{1}{3} + 1 \times \frac{1}{3} + 0 \times \frac{1}{3}} = \frac{\frac{1}{2} \times \frac{1}{3}}{\frac{1}{3} \left(\frac{1}{2} + 1 \right)} = \frac{\frac{1}{2} \times \frac{1}{3}}{\frac{1}{3} \times \frac{3}{2}} = \frac{1}{3}$$

$$= 1 - \frac{1}{3} = \frac{2}{3}$$

Q.4. Refer Class Notes

Q.5 $P(\text{Success}) = p = 0.2$

$$X \sim \text{geom}_0(p)$$

$$f_X(x) = \begin{cases} (1-p)^x p & ; x = 0, 1, 2, \dots \\ 0 & ; \text{o.w.} \end{cases}$$

$$E(X) = \frac{1-p}{p} = \frac{0.8}{0.2} = 4.$$

Q.6. Let $X \sim \text{geom}_0(p)$ (All calculations will be similar for $X \sim \text{geom}(p)$)

$$\begin{aligned} P(X \geq x) &= 1 - P(X < x) \\ &= 1 - P(X \leq x-1) \\ &= 1 - \left\{ \sum_{m=0}^{x-1} P(X = m) \right\} \\ &= 1 - \left\{ p \sum_{m=0}^{x-1} (1-p)^m \right\} \\ &\stackrel{\text{Geom. Series}}{=} 1 - p \left\{ \frac{1 - (1-p)^x}{1 - (1-p)} \right\} \\ &= 1 - 1 + (1-p)^x = (1-p)^x \\ P(X \geq i+j | X \geq i) &= \frac{P(X \geq i+j, X \geq i)}{P(X \geq i)} \end{aligned}$$

$$= \frac{P(X \geq i+j)}{(1-p)^i}$$

$$= \frac{(1-p)^{i+j}}{(1-p)^i} = (1-p)^j = P(X \geq j)$$

this means $X \geq i$ is forgotten & the event $X \geq i+j$ is just as likely as $X \geq j$ (as if the info $X \geq i$ was not known to us).

We will again see this kind of behavior when we study Markov chains.

(Q.7) $f_X(x) = {}^n C_x p^x (1-p)^{n-x}$; given $X \sim \text{Bin}(n, p)$

$$= \frac{n!}{(n-x)! x!} p^x (1-p)^{n-x}$$

$$\lambda = np = \frac{n!}{(n-x)! x!} \frac{\lambda^x}{n^x} (1 - \frac{\lambda}{n})^{n-x}$$

$$\lim_{n \rightarrow \infty} \frac{n!}{(n-x)! n^x} = \lim_{n \rightarrow \infty} \frac{n(n-1) \dots (n-x+1)}{n^x}$$

Numerator has $n-(n-x)=x$ terms
 \Rightarrow highest order of n is x

$$\lim_{n \rightarrow \infty} \left(1 - \frac{\lambda}{n}\right)^n = e^{-\lambda}$$

$$\lim_{n \rightarrow \infty} \left(1 - \frac{\lambda}{n}\right)^{-x} = \lim_{n \rightarrow \infty} \underbrace{\frac{1}{\left(1 - \frac{\lambda}{n}\right) \dots \left(1 - \frac{\lambda}{n}\right)}}_{x \text{ times}} = 1$$

$$\therefore f_X(x) \xrightarrow{n \rightarrow \infty} \frac{e^{-\lambda} \lambda^x}{x!} \quad ; \quad e^{-\lambda} \cdot \lambda^x \xrightarrow{n \rightarrow \infty} \text{Poisson}(n)$$

Q. 8. by defⁿ: - $E(C) = \sum_{i=1}^{\infty} i \cdot P(C=i)$
 b/c C runs from 1 to ∞

Method (i)

$$\begin{aligned}
 \text{R.H.S.} &= \sum_{i=0}^{\infty} P(C > i) \\
 &= P(C > 0) + P(C > 1) + P(C > 2) + \dots \\
 &= \sum_{i=1}^{\infty} P(C = i) + \sum_{i=2}^{\infty} P(C = i) + \sum_{i=3}^{\infty} P(C = i) + \dots \\
 &= \left[\sum_{i=1}^2 P(C = i) + \sum_{i=2}^{\infty} P(C = i) \right] \\
 &\quad + \left(\sum_{i=2}^3 P(C = i) + \sum_{i=3}^{\infty} P(C = i) \right) \\
 &\quad + \left(\sum_{i=3}^4 P(C = i) + \sum_{i=4}^{\infty} P(C = i) \right) \\
 &= \left[1 \cdot P(C=1) + \sum_{i=2}^2 P(C = i) \right] + \dots \\
 &\quad + \left[\sum_{i=2}^3 P(C = i) \right] \\
 &\quad + \left[\sum_{i=2}^4 P(C = i) + \sum_{i=3}^4 P(C = i) \right] \\
 &= 1 \cdot P(C=1) + 2 \cdot P(C=2) + 3 \cdot P(C=3) + \dots
 \end{aligned}$$

Fix writing

$$E(C) = \sum_{i=0}^{\infty} P(C \geq i) = \sum_{i=1}^{\infty} i \cdot P(C=i)$$

$$\begin{aligned} &= \sum_{i=0}^{\infty} [1 - P(C \leq i)] \\ &= \sum_{i=0}^{\infty} \left[1 - \sum_{j=1}^i P(C=j) \right] \\ &\approx \sum_{i=0}^{\infty} \left[1 - \sum_{j=1}^i (1-P)^{j-1} P \right] \\ &\text{finite Geom series} \sum_{i=0}^{\infty} \left[1 - P^i \frac{1 - (1-P)^i}{1 - (1-P)} \right] \end{aligned}$$

$$\text{infinite Geom series} \sum_{i=0}^{\infty} (1-P)^i = \frac{1}{1 - (1-P)} = \frac{1}{P}$$

$$\begin{aligned} &= \frac{1}{0.4} \\ &= \frac{10}{4} = \frac{5}{2} \end{aligned}$$

Method(2)

Using Law of total expectation.Let A be the event that the system failsin the 1st hr
Recall squirrel problem!!

$$\begin{aligned} E(C) &= E(C|A)P(A) + E(C|A')P(A') \\ &= 1 \times P(A) + \{1 + E(c)\}P(A') \end{aligned}$$

$$\begin{aligned}
 E(c) &= p \cdot E(c) (1-p) \\
 &= p + 1 + E(c) \cancel{- p} = pE(c) \\
 &= 1 + E(c)(1-p)
 \end{aligned}$$

$$\begin{aligned}
 \Rightarrow E(c)[1 - (1-p)] &\approx 1 \\
 \Rightarrow E(c) &= \frac{1}{p} = \frac{1}{0.4} \\
 &\approx \frac{10}{4} \\
 &= 2.5 \text{ ms.}
 \end{aligned}$$

Q. 9. (Power of Indicator RVs; what
ever if we don't know the probability Dⁿ of a
random phenomenon? when we get
let Y be the no. of men who get
their own hats back.

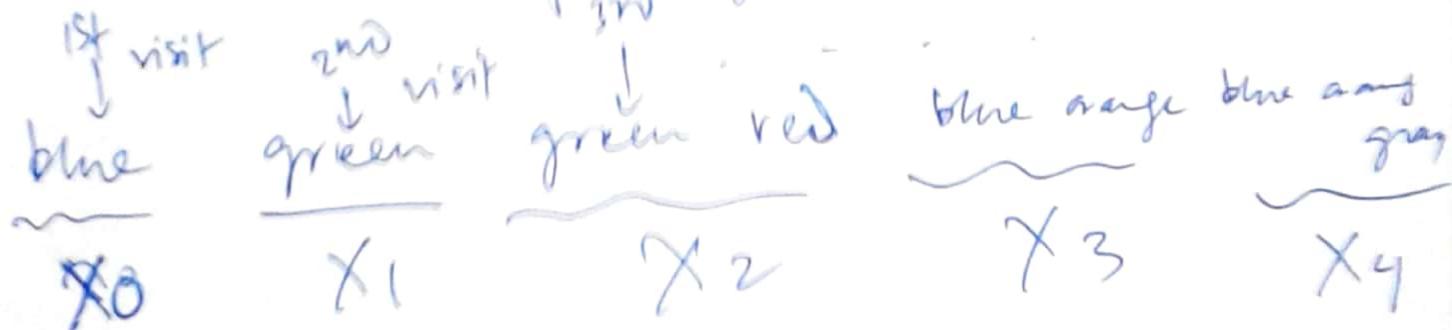
$$Y_i = \begin{cases} 1; & \text{if man gets his hat back} \\ 0; & \text{if man does NOT get his hat back.} \end{cases}$$

$$\text{then } Y = \sum_{i=1}^{100} Y_i$$

$$\begin{aligned}
 E(Y) &= \sum_{i=1}^{100} E(Y_i) \\
 \text{class Notes} &= \sum_{i=1}^{100} P(Y_i = 1) = \sum_{i=1}^{100} \frac{1}{100} \\
 &= 100 \times \frac{1}{100} \\
 &= 1.
 \end{aligned}$$

(10)

Let us consider the coupons received in the corresponding visit as follows.



Let us partition the above outcomes into zones; s.t. a new zone terminates upon receiving a "new" type of coupon.

This way we have a seq. $X_0, X_1, X_2, \dots, X_{10}$ each of varying length

$$\text{here } X_0 = 1$$

$$\text{eg. } X_1 = 1$$

$$X_2 = 2$$

$$X_3 = 2$$

$$X_4 = 3$$

 \vdots

Let $T = X_0 + X_1 + X_2 + \dots + X_{10}$
we need to find $E(T)$.

Note X_k is the length of k^{th} gone

At the beginning of k^{th} gone; we already have $k \leq 10$ different types of coupons

When we have k types; each visit contains a type w/ probability $\frac{k}{10}$

\Rightarrow Each visit contains a new type

w/ probability $1 - \frac{k}{10} = \frac{10-k}{10} = p_{\text{new}}$

- Expected no. of visits until we get the $(k+1)^{\text{th}}$ new type is the same as expected/mean time to failure

w/ probability $\frac{10-k}{10}$ (Q.E.D.)

$$\Rightarrow E(X_k) = \frac{1}{\frac{10-k}{10}} = \frac{10}{10-k}$$

$$E(T) = \sum_{i=1}^{10} E(X_i) = \frac{10}{10-0} + \frac{10}{10-1} + \frac{10}{10-2} + \dots + \frac{10}{10-10} \\ = 10 \left\{ \frac{1}{10} + \frac{1}{9} + \frac{1}{8} + \dots + \frac{1}{2} \right\} + 10$$

= 10 H_{10}

Harmonic

more
true for large
no. of types

$\rightarrow \sim 10 \log_{10} 10$

Series
terminated at
 10^{th} term