1 Contest Setup

1.1 Java template

```java
import java.io.*;
import java.util.*;

public class Main
{
    public static void main(String[] args)
    {
        MyScanner sc = new MyScanner();
        out = new PrintWriter(new BufferedOutputStream(System.out));
        // Start writing your solution here.

        // Stop writing your solution here.
        out.close();
    }

    public static PrintWriter out;

    public static class MyScanner
    {
        BufferedReader br;
        StringTokenizer st;

        public MyScanner()
        {
            br = new BufferedReader(new InputStreamReader(System.in));
        }

        boolean hasNext()
        {
            while (st == null || !st.hasMoreElements()) {
                try {
                    st = new StringTokenizer(br.readLine());
                } catch (Exception e) {
                    return false;
                }
            }
```

```java
            return true;
        }

        String next()
        {
            if (hasNext())
                return st.nextToken();
            return null;
        }

        int nextInt()
        {
            return Integer.parseInt(next());
        }

        long nextLong()
        {
            return Long.parseLong(next());
        }

        double nextDouble()
        {
            return Double.parseDouble(next());
        }

        String nextLine()
        {
            String str = "";
            try {
                str = br.readLine();
            } catch (IOException e) {
                e.printStackTrace();
            }
            return str;
        }
    }
}
```

1.1.1 Java Issues

1. Random Shuffle before sorting:
   `Random rnd = new Random(); rnd.nextInt();`
2. Use `StringBuilder` for large output
3. Java has strict parsing rules. e.g. using `sc.nextInt()` to read a long will result in RE
4. For class sorting, use code `implements Comparable<Class name>`. Or, use code `new Comparator<Interval>() {}` at `Collections.sort()` second argument

2 System Testing

1. Setup vimrc and bashrc
2. Test g++ and Java 8 compiler
3. Look for compilation parameter and code it into bashrc
4. Test if c++ and Java templates work properly on local and judge machine (bits, auto, and other c++11 stuff)
5. Test "divide by 0" → RE/TLE?

6. Make a complete graph and run Floyd warshall, to test time complexity upper bound
7. Make a linear graph and use DFS to test stack size
8. Test output with extra newline and spaces
9. Go to $Eclipse \rightarrow preference \rightarrow Java \rightarrow Editor \rightarrow ContentAssist$, add $.abcdefghijklmnopqrstuvwxyz$ to auto activation triggers for Java in Eclipse

3 Reminder

1. 隊友的建議，要認真聽！要記得心平氣和的小聲討論喔！通常隊友的建議都會突破你盲點。
2. 每一題都要小心讀，尤其是 IO 的格式和限制都要看清楚。
3. 小心估計時間複雜度和 空間複雜度
4. Coding 要兩人一組，要相信你的隊友的實力！
5. 1WA 罰 20 分鐘！放輕鬆，不要急，多產幾組測資後再丟。
6. 範測一定要過！產個幾組極端測資，例如 input 下限、特殊 cases 0, 1, -1、空集合等等
7. 比賽是連續測資，一定要全部讀完再開始 solve 喔！
8. Bus error: 有 `scanf`, `fgets` 但是卻沒東西可以讀取了！可能有 early termination 但是時機不對。
9. 圖論一定要記得檢查連通性。最簡單的做法就是 loop 過所有的點
10. `long long = int * int` 會完蛋
11. `long long int` 的位元運算要記得用 `1LL << 35`
12. 記得清理 Global variable
13. 建圖時要注意有無重邊！

4 Topic list

1. 列舉、窮舉 enumeration
2. 貪心 greedy
3. 排序 sorting, topological sort
4. 二分搜 binary search (數學算式移項合併後查詢)
5. 爬行法（右跑左追）Two Pointer
6. 離散化
7. Dynamic programming, 矩陣快速冪
8. 鴿籠原理 Pigeonhole
9. 最近共同祖先 LCA (倍增法, LCA 轉 RMQ)
10. 折半完全列舉 (能用 vector 就用 vector)
11. 離線查詢 Offline (DFS, LCA)

12. 圖的連通性 Directed graph connectivity -> DFS. Undirected graph -> Union Find
13. 因式分解
14. 從答案推回來
15. 寫出數學式，有時就馬上出現答案了！
16. 奇偶性質

5 Useful code

5.1 Leap year $O(1)$

```
(year % 400 == 0 || (year % 4 == 0 && year % 100 != 0))
```

5.2 Fast Exponentiation $O(log(exp))$

Fermat's little theorem: 若 m 是質數，則 $a^{m-1} \equiv 1 \pmod{m}$

```
1  ll fast_pow(ll a, ll b, ll M) {
2      ll ans = 1;
3      ll base = a % M;
4      while (b) {
5          if (b & 1)
6              ans = ans * base % M;
7          base = base * base % M;
8          b >>= 1;
9      }
10     return ans;
11 }
```

5.3 Mod Inverse $O(logn)$

Case 1: $gcd(a, m) = 1$: ax + my = gcd(a, m) = 1 (use ext_gcd)

Case 2: $m\ is\ prime$: $a^{m-2} \equiv a^{-1} mod\ m$

5.4 GCD $O(log(min(a+b)))$

注意負數的 case! C++ 是看被除數決定正負號的。

```
ll gcd(ll a, ll b)
{
    return b == 0 ? a : gcd(b, a % b);
}
```

5.5 Extended Euclidean Algorithm GCD $O(log(min(a+b)))$

Bezout identity $ax + by = gcd(a, b)$, where $| x |\leq \frac{b}{d}$ and $| y |\leq \frac{a}{d}$.

```
1  ll extgcd(ll a, ll b, ll& x, ll&y) {
2      if (b == 0) {
3          x = 1;
```

```
4          y = 0;
5          return a;
6      }
7      else {
8          ll d = extgcd(b, a % b, y, x);
9          y -= (a / b) * x;
10         return d;
11     }
12 }
```

5.6 Prime Generator $O(nloglogn)$

```
1  const ll MAX_NUM = 1e6; // 要是合數
2  bool is_prime[MAX_NUM];
3  vector<ll> primes;
4
5  void init_primes() {
6      fill(is_prime, is_prime + MAX_NUM, true);
7      is_prime[0] = is_prime[1] = false;
8      for (ll i = 2; i < MAX_NUM; i++) {
9          if (is_prime[i]) {
10             primes.push_back(i);
11             for (ll j = i * i; j < MAX_NUM; j += i)
12                 is_prime[j] = false;
13         }
14     }
15 }
```

5.7 C++ Reference

```
1  algorithm
2      ::find: [it s, it t, val] -> it
3      ::count: [it s, it t, val] -> int
4      ::unique: [it s, it t] -> it (it = new end)
5      ::merge: [it s1, it t1, it s2, it t2, it o] -> void (o allocated)
6
7  string::
8      ::replace(idx, len, string) -> void
9      ::find (str, pos = 0) -> idx
10     ::substr (pos = 0, len = npos) -> string
11 string <-> int
12     ::stringstream; // remember to clear
13     ::sscanf(s.c_str(), "%d", &i);
14     ::sprintf(result, "%d", i); string s = result;
15
16 math/cstdlib
17     ::atan2(y=0, x=-1) -> pi
18
19 io printf/scanf
20     ::int:                "%d"      /    "%d"
21     ::double:             "%lf","f" /    "%lf"
22     ::string:             "%s"      /    "%s"
23     ::long long:          "%lld"    /    "%lld"
24     ::long double:        "%Lf"     /    "%Lf"
25     ::unsigned int:       "%u"      /    "%u"
26     ::unsigned long long: "%ull"    /    "%ull"
27     ::oct:                "0%o"
28     ::hex:                "0x%x"
29     ::scientific:         "%e"
30     ::width:              "%05d"
31     ::precision:          "%.5f"
32     ::adjust left:        "%-5d"
33
34 io cin/cout
35     ::oct:            cout << oct << showbase;
36     ::hex:            cout << hex << showbase;
37     ::scientific:     cout << scientific;
38     ::width:          cout << setw(5);
39     ::precision:      cout << fixed << setprecision(5);
40     ::adjust left:    cout << setw(5) << left;
```

6 Search

6.1 Ternary Search $O(nlogn)$

```
double l = ..., r = ....; // input
for(int i = 0; i < 100; i++) {
    double m1 = l + (r - l) / 3, m2 = r - (r - l) / 3;
    if (f (m1) < f (m2)) // f - convex function
        l = m1;
    else
        r = m2;
}
f(r) - maximum of function
```

7 Basic data structure

7.1 1D BIT

```
// BIT is 1-based
const int MAX_N = 20000; //這個記得改!
ll bit[MAX_N + 1];

ll sum(int i) {
    int s = 0;
    while (i > 0) {
        s += bit[i];
        i -= (i & -i);
    }
    return s;
}

void add(int i, ll x) {
    while (i <= MAX_N) {
        bit[i] += x;
        i += (i & -i);
    }
}
```

7.2 2D BIT

```
// BIT is 1-based
const int MAX_N = 20000, MAX_M = 20000; //這個記得改!
ll bit[MAX_N + 1][MAX_M + 1];
```

```
ll sum(int a, int b) {
    ll s = 0;
    for (int i = a; i > 0; i -= (i & -i))
        for (int j = b; j > 0; j -= (j & -j))
            s += bit[i][j];
    return s;
}

void add(int a, int b, ll x) {
    // MAX_N, MAX_M 須適時調整!
    for (int i = a; i <= MAX_N; i += (i & -i))
        for (int j = b; j <= MAX_M; j += (j & -j))
            bit[i][j] += x;
}
```

7.3   Union Find

```
const int MAX_N = 20000; // 記得改
struct UFDS {
    int par[MAX_N];

    void init(int n) {
        memset(par, -1, sizeof(int) * n);
    }

    int root(int x) {
        return par[x] < 0 ? x : par[x] = root(par[x]);
    }

    void merge(int x, int y) {
        x = root(x);
        y = root(y);

        if (x != y) {
            if (par[x] > par[y])
                swap(x, y);
            par[x] += par[y];
            par[y] = x;
        }
    }
};
```

7.4   Segment Tree

```
 1  typedef long long ll;
 2  const int MAX_N = 100000;
 3  const int MAX_NN = (1 << 20); // bigger than MAX_N
 4
 5  struct SegTree {
 6      int NN;                  // size of tree
 7      ll dflt;                 // default val
 8      ll seg[2 * MAX_NN];  // 0-based index, 2 * MAX_NN - 1 in fact
 9      ll lazy[2 * MAX_NN]; // 0-based index, 2 * MAX_NN - 1 in fact
10      // lazy[u] != 0 <->
11      // substree of u (u not inclued) is not up-to-date (it's dirty)
12
13      void init(int n, ll val)
14      {
15          dflt = val;
16          NN = 1;
17          while (NN < n)
18              NN <<= 1;
19          fill(seg, seg + 2 * NN, dflt);
20          fill(lazy, lazy + 2 * NN, dflt);
21      }
22
23      void gather(int u, int l, int r)
24      {
25          seg[u] = seg[u * 2 + 1] + seg[u * 2 + 2];
26      }
27
28      void push(int u, int l, int r)
29      {
30          if (lazy[u] != 0) {
31              int m = (l + r) / 2;
32
33              seg[u * 2 + 1] += (m - l) * lazy[u];
34              seg[u * 2 + 2] += (r - m) * lazy[u];
35
36              lazy[u * 2 + 1] += lazy[u];
37              lazy[u * 2 + 2] += lazy[u];
38              lazy[u] = 0;
39          }
40      }
41
42      void build(int u, int l, int r)
43      {
44          if (r - l == 1)
45              return;
46          int m = (l + r) / 2;
47          build(u * 2 + 1, l, m);
48          build(u * 2 + 2, m, r);
49          gather(u, l, r);
50      }
51
52      ll query(int a, int b, int u, int l, int r)
53      {
54          if (l >= b || r <= a)
55              return dflt;
56          if (l >= a && r <= b)
57              return seg[u];
58          int m = (l + r) / 2;
59          push(u, l, r);
60          ll res1 = query(a, b, u * 2 + 1, l, m);
61          ll res2 = query(a, b, u * 2 + 2, m, r);
62          gather(u, l, r); // data is dirty since previous push
63          return res1 + res2;
64      }
65
66      void update(int a, int b, int x, int u, int l, int r)
67      {
68          if (l >= b || r <= a)
69              return;
70          if (l >= a && r <= b) {
71              seg[u] += (r - l) * x; // update u and
72              lazy[u] += x;          // set subtree u is not up-to-date
73              return;
```

```
74          }
75          int m = (l + r) / 2;
76          push(u, l, r);
77          update(a, b, x, u * 2 + 1, l, m);
78          update(a, b, x, u * 2 + 2, m, r);
79          gather(u, l, r); // remember this
80      }
81  };
```

7.5 Sparse Table

```
1  struct Sptb {
2      int sp[MAX_LOG_N][MAX_N]; // MAX_LOG_N = ceil(lg(MAX_N))
3
4      void build(int inp[], int n)
5      {
6          for (int j = 0; j < n; j++)
7              sp[0][j] = inp[j];
8
9          for (int i = 1; (1 << i) <= n; i++)
10             for (int j = 0; j + (1 << i) <= n; j++)
11                 sp[i][j] = min(sp[i - 1][j], sp[i - 1][j + (1 << (i -
    ↪  1))]);
12     }
13
14     int query(int l, int r) // [l, r)
15     {
16         int k = floor(log2(r - l));
17         return min(sp[k][l], sp[k][r - (1 << k)]);
18     }
19 };
```

8 Tree

8.1 LCA

```
1  const int MAX_N = 10000;
2  const int MAX_LOG_N = 14; // (1 << MAX_LOG_N) > MAX_N
3
4  int N;
5  int root;
6  int dep[MAX_N];
7  int par[MAX_LOG_N][MAX_N];
8
9  vector<int> child[MAX_N];
10
11 void dfs(int u, int p, int d) {
12     dep[u] = d;
13     for (int i = 0; i < int(child[u].size()); i++) {
14         int v = child[u][i];
15         if (v != p) {
16             dfs(v, u, d + 1);
17         }
18     }
19 }
```

```
20
21 void build() {
22     // par[0][u] and dep[u]
23     dfs(root, -1, 0);
24
25     // par[i][u]
26     for (int i = 0; i + 1 < MAX_LOG_N; i++) {
27         for (int u = 0; u < N; u++) {
28             if (par[i][u] == -1)
29                 par[i + 1][u] = -1;
30             else
31                 par[i + 1][u] = par[i][par[i][u]];
32         }
33     }
34 }
35
36 int lca(int u, int v) {
37     if (dep[u] > dep[v]) swap(u, v); // 讓 v 較深
38     int diff = dep[v] - dep[u]; // 將 v 上移到與 u 同層
39     for (int i = 0; i < MAX_LOG_N; i++) {
40         if (diff & (1 << i)) {
41             v = par[i][v];
42         }
43     }
44
45     if (u == v) return u;
46
47     for (int i = MAX_LOG_N - 1; i >= 0; i--) { // 必需倒序
48         if (par[i][u] != par[i][v]) {
49             u = par[i][u];
50             v = par[i][v];
51         }
52     }
53     return par[0][u];
54 }
```

8.2 Tree Center

```
1  int diameter = 0, radius[N], deg[N]; // deg = in + out degree
2  int findRadius()
3  {
4      queue<int> q; // add all leaves in this group
5      for (auto i : group)
6          if (deg[i] == 1)
7              q.push(i);
8
9      int mx = 0;
10     while (q.empty() == false) {
11         int u = q.front();
12         q.pop();
13
14         for (int v : g[u]) {
15             deg[v]--;
16             if (deg[v] == 1) {
17                 q.push(v);
```

```
18              radius[v] = radius[u] + 1;
19              mx = max(mx, radius[v]);
20          }
21      }
22  }
23
24  int cnt = 0; // crucial for knowing if there are 2 centers or not
25  for (auto j : group)
26      if (radius[j] == mx)
27          cnt++;
28
29  // add 1 if there are 2 centers (radius, diameter)
30  diameter = max(diameter, mx * 2 + (cnt == 2));
31  return mx + (cnt == 2);
32  }
```

9 Graph

9.1 Articulation point / Bridge

```
1  const int MAX_N = 1111;
2  vector<int> g[MAX_N];
3
4  // for bridge
5  typedef pair<int, int> ii;
6  vector<ii> ans;
7
8  // for articulation point
9  int root;                 // set it before dfs() call
10 bool isCutVertex[MAX_N]; // init to false
11
12 int tt = 0, dfn[MAX_N], low[MAX_N]; // init array to -1
13 void dfs(int u, int p)
14 {
15     dfn[u] = low[u] = tt++;
16
17     // for articulation point, root needs to have >= 2 childrens
18     int child = 0;
19     for (auto v : g[u]) {
20         if (v == p)
21             continue;
22         child++;
23
24         if (dfn[v] == -1) {
25             dfs(v, u);
26             low[u] = min(low[u], low[v]);
27
28             if (low[v] > dfn[u]) // bridge
29                 ans.push_back(ii(min(u, v), max(u, v)));
30
31             if (u != root && low[v] >= dfn[u]) { // articulation point
32                 isCutVertex[u] = true;
33             } else if (u == root && child >= 2) { // articulation point
34                 isCutVertex[u] = true;
35             }
```

```
36         } else {
37             // u -> v, u has direct access to v -> back edge
38             low[u] = min(low[u], dfn[v]);
39         }
40     }
41 }
```

9.2 2-SAT

$$p \vee (q \wedge r)$$
$$= ((p \wedge q) \vee (p \wedge r))$$
$$p \oplus q$$
$$= \neg((p \wedge q) \vee (\neg p \wedge \neg q))$$
$$= (\neg p \vee \neg q) \wedge (p \vee q)$$

```
// 建圖
// (x1 or x2) and ... and (xi or xj)
// (xi or xj) 建邊
// ~xi -> xj
// ~xj -> xi

tarjan(); // scc 建立的順序是倒序的拓璞排序
for (int i = 0; i < 2 * N; i += 2) {
    if (belong[i] == belong[i ^ 1]) {
        // 無解
    }
}
for (int i = 0; i < 2 * N; i += 2) { // 迭代所有變數
    if (belong[i] < belong[i ^ 1]) { // i 的拓璞排序比 ~i 的拓璞排序大
        // i = T
    }
    else {
        // i = F
    }
}
```

9.3 CC

9.3.1 BCC

以 Edge 做分界的話，stack 要裝入 (u - v)，並 pop 終止條件為!= (u - v)
以 Articulation point 做為分界 (code below)，注意有無坑人的重邊
注意，用 SCC 的 code 的話，只要多判一個 u 是否為 p，如果是的話就直接 return
(加在第 21 行之後)

9.3.2 SCC

First of all we run DFS on the graph and sort the vertices in decreasing of their finishing time (we can use a stack).
Then, we start from the vertex with the greatest finishing time, and for each vertex v that is not yet in any SCC, do : for each u that v is reachable by u and u is not yet in any SCC, put it in the SCC of vertex v. The code is quite simple.

```
1   const int MAX_V = ...;
2   const int INF = 0x3f3f3f3f;
3   int V;
4   vector<int> g[MAX_V];
5
6   int dfn_idx = 0;
7   int scc_cnt = 0;
8   int dfn[MAX_V];
9   int low[MAX_V];
10  int belong[MAX_V];
11  bool in_st[MAX_V];
12  vector<int> st;
13
14  void scc(int v)
15  {
16      dfn[v] = low[v] = dfn_idx++;
17      st.push_back(v);
18      in_st[v] = true;
19
20      for (int i = 0; i < int(g[v].size()); i++) {
21          const int u = g[v][i];
22          if (dfn[u] == -1) {
23              scc(u);
24              low[v] = min(low[v], low[u]);
25          } else if (in_st[u]) {
26              low[v] = min(low[v], dfn[u]);
27          }
28      }
29
30      if (dfn[v] == low[v]) {
31          int k;
32          do {
33              k = st.back();
34              st.pop_back();
35              in_st[k] = false;
36              belong[k] = scc_cnt;
37          } while (k != v);
38          scc_cnt++;
39      }
40  }
41
42  void tarjan() // scc 建立的順序即為反向的拓璞排序
43  {
44      st.clear();
45      fill(dfn, dfn + V, -1);
46      fill(low, low + V, INF);
47      dfn_idx = 0;
48      scc_cnt = 0;
49      for (int v = 0; v < V; v++) {
50          if (dfn[v] == -1) {
51              scc(v);
52          }
53      }
54  }
```

9.4 Shortest Path

Time complexity notations: V = vertex, E = edge
Minimax: $dp[u][v] = min(dp[u][v], max(dp[u][k], dp[k][v]))$

9.4.1 Dijkatra (next-to-shortest path) $O(V log E)$

密集圖別用 priority queue!

```
1   struct Edge {
2       int to, cost;
3   };
4
5   typedef pair<int, int> P; // <d, v>
6   const int INF = 0x3f3f3f3f;
7
8   int N, R;
9   vector<Edge> g[5000];
10
11  int d[5000];
12  int sd[5000];
13
14  int solve()
15  {
16      fill(d, d + N, INF);
17      fill(sd, sd + N, INF);
18      priority_queue<P, vector<P>, greater<P>> pq;
19
20      d[0] = 0;
21      pq.push(P(0, 0));
22
23      while (!pq.empty()) {
24          P p = pq.top();
25          pq.pop();
26          int v = p.second;
27
28          if (sd[v] < p.first) // 比次短距離還大, 沒用, 跳過
29              continue;
30
31          for (size_t i = 0; i < g[v].size(); i++) {
32              Edge &e = g[v][i];
33              int nd = p.first + e.cost;
34              if (nd < d[e.to]) { // 更新最短距離
35                  swap(d[e.to], nd);
36                  pq.push(P(d[e.to], e.to));
37              }
38              if (d[e.to] < nd && nd < sd[e.to]) { // 更新次短距離
39                  sd[e.to] = nd;
40                  pq.push(P(sd[e.to], e.to));
41              }
42          }
43      }
44
45      return sd[N - 1];
46  }
```

9.4.2 SPFA

```
1   #define sz(x) (int(x.size()))
2   const int MAX_V = ...;
3   const int INF = 0x3f3f3f3f;
4
5   struct Edge {
6       int to, w;
7   };
8   vector<Edge> g[MAX_V];
9   int d[MAX_V];
10
11  bool spfa(int V, int S) {
12      fill(d, d + V, INF);
13      queue<int> q;
14      vector<bool> inq(V, false);
15      vector<int> cnt(V, 0);
16
17      d[S] = 0;
18      cnt[S]++;
19      inq[S] = true;
20      q.push(S);
21
22      while (!q.empty()) {
23          int u = q.front(); q.pop(); inq[u] = false;
24          for (int i = 0; i < sz(g[u]); i++) {
25              const Edge& e = g[u][i];
26              if (d[e.to] > d[u] + e.w) {
27                  d[e.to] = d[u] + e.w;
28                  if (++cnt[e.to] >= V) {
29                      return true;
30                  }
31                  if (!inq[e.to]) {
32                      inq[e.to] = true;
33                      q.push(e.to);
34                  }
35              }
36          }
37      }
38
39      return false;
40  }
```

9.4.3 Bellman-Ford $O(VE)$

```
1   struct Edge {
2       int from, to, cost;
3   };
4
5   const int MAX_V = ...;
6   const int MAX_E = ...;
7   const int INF = 0x3f3f3f3f;
8   int V, E;
9   Edge edges[MAX_E];
10  int d[MAX_V];
11
12  bool bellman_ford()
```

```
13  {
14      fill(d, d + V, INF);
15
16      d[0] = 0;
17      for (int i = 0; i < V; i++) {
18          for (int j = 0; j < E; j++) {
19              Edge &e = edges[j];
20              if (d[e.to] > d[e.from] + e.cost) {
21                  d[e.to] = d[e.from] + e.cost;
22
23                  if (i == V - 1) // negative cycle
24                      return true;
25              }
26          }
27      }
28
29      return false;
30  }
```

9.4.4 Floyd-Warshall $O(V^3)$

The graph is stored using adjacency matrix. The initial state is $diagnal = 0$ and $others = INF$. (If INF is int, use long long for the matrix)
If diagonal numbers are negative ← cycle .

```
for(int k = 0; k < N; k++)
    for(int i = 0; i < N; i++)
        for(int j = 0; j < N; j++)
            dp[i][j] = min(dp[i][j], dp[i][k] + dp[k][j]);
```

9.5 MST

9.5.1 Kruskal

1. Store the graph by $(weight, (from, to))$
2. Sort the graph by $weight$
3. Start from the smallest weight, and keep adding edges that won't form a cycle with the current MST set
4. Early termination condition: $n - 1$ edges has been added, NOT size of the union-find set

9.5.2 Second MST

```
1   const int INF = 0x3f3f3f3f;
2   const int MAX_V = 100;
3   const int MAX_LOG_V = 7;
4   int V, E; // 記得初使化
5
6   struct Edge {
7       int u, v, w;
8   };
9   vector<Edge> edges;
10
```

```
11  // btn[i][u] = u 前往它 2^i parent 的路上經過的最大權重
12  // par[i][u] = u 的 2^i parent 是誰
13  int dep[MAX_V]; // should be init to -1
14  int btn[MAX_LOG_V][MAX_V];
15  int par[MAX_LOG_V][MAX_V];
16
17  // mst
18  struct AdjE {
19      int to, w;
20  };
21  vector<AdjE> g[MAX_V];
22
23  void dfs(int u, int p, int d) {
24      dep[u] = d;
25      par[0][u] = p;
26      for (auto e : g[u]) {
27          if (e.to != p) {
28              btn[0][e.to] = e.w;
29              dfs(e.to, u, d + 1);
30          }
31      }
32  }
33
34  void build() {
35      for (int u = 0; u < V; u++) {
36          if (dep[u] == -1) {
37              dfs(u, -1, 0);
38          }
39      }
40
41      for (int i = 0; i + 1 < MAX_LOG_V; i++) {
42          for (int u = 0; u < V; u++) {
43              if (par[i][u] == -1 || par[i][par[i][u]] == -1) {
44                  par[i + 1][u] = -1;
45                  btn[i + 1][u] = 0;
46              }
47              else {
48                  par[i + 1][u] = par[i][par[i][u]];
49                  btn[i + 1][u] = max(btn[i][u], btn[i][par[i][u]]);
50              }
51          }
52      }
53  }
54
55  int lca(int u, int v) { // 回傳 u, v 之間的最大權重
56      int mx = -INF; // u, v 之間的最大權重
57
58      if (dep[u] > dep[v]) swap(u, v);
59      int diff = dep[v] - dep[u];
60      for (int i = MAX_LOG_V - 1; i >= 0; i--) {
61          if (diff & (1 << i)) {
62              mx = max(mx, btn[i][v]);
63              v = par[i][v];
64          }
65      }
66
67      if (u == v) return mx;
68
69      for (int i = MAX_LOG_V - 1; i >= 0; i--) {
70          if (par[i][u] != par[i][v]) {
71              mx = max(mx, btn[i][u]);
72              mx = max(mx, btn[i][v]);
73              u = par[i][u];
74              v = par[i][v];
75          }
76      }
77      // lca = par[0][u] = par[0][v];
78      mx = max(mx, max(btn[0][u], btn[0][v]));
79
80      return mx;
81  }
82
83  // second mst
84  build();
85  int ans = INF;
86  for (auto e: non_mst_edges) {
87      int mx_w = lca(e.u, e.v);
88      ans = min(ans, (total_w + e.w - mx_w));
89  }
```

10 Flow

10.1 Max Flow (Dinic)

```
1   struct Edge {
2       int to, cap, rev;
3       Edge(int a, int b, int c) {
4           to = a;
5           cap = b;
6           rev = c;
7       }
8   };
9
10  const int INF = 0x3f3f3f3f;
11  const int MAX_V = 20000 + 10;
12  // vector<Edge> g[MAX_V];
13  vector< vector<Edge> > g(MAX_V);
14  int level[MAX_V];
15  int iter[MAX_V];
16
17  inline void add_edge(int u, int v, int cap) {
18      g[u].push_back((Edge){v, cap, (int)g[v].size()});
19      g[v].push_back((Edge){u, 0, (int)g[u].size() - 1});
20  }
21
22  void bfs(int s) {
23      memset(level, -1, sizeof(level)); // 用 fill
24      queue<int> q;
25
26      level[s] = 0;
27      q.push(s);
```

```
28      while (!q.empty()) {
29          int v = q.front(); q.pop();
30          for (int i = 0; i < int(g[v].size()); i++) {
31              const Edge& e = g[v][i];
32              if (e.cap > 0 && level[e.to] < 0) {
33                  level[e.to] = level[v] + 1;
34                  q.push(e.to);
35              }
36          }
37      }
38  }
39
40  int dfs(int v, int t, int f) {
41      if (v == t) return f;
42      for (int& i = iter[v]; i < int(g[v].size()); i++) { // & 很重要
43          Edge& e = g[v][i];
44          if (e.cap > 0 && level[v] < level[e.to]) {
45              int d = dfs(e.to, t, min(f, e.cap));
46              if (d > 0) {
47                  e.cap -= d;
48                  g[e.to][e.rev].cap += d;
49                  return d;
50              }
51          }
52      }
53      return 0;
54  }
55
56  int max_flow(int s, int t) { // dinic
57      int flow = 0;
58      for (;;) {
59          bfs(s);
60          if (level[t] < 0) return flow;
61          memset(iter, 0, sizeof(iter));
62          int f;
63          while ((f = dfs(s, t, INF)) > 0) {
64              flow += f;
65          }
66      }
67  }
```

10.2 Min Cost Flow

```
1  #define st first
2  #define nd second
3
4  typedef pair<double, int> pii; // 改成用 int
5  const double INF = 1e10;
6
7  struct Edge {
8      int to, cap;
9      double cost;
10     int rev;
11 };
12
```

```
13  const int MAX_V = 2 * 100 + 10;
14  int V;
15  vector<Edge> g[MAX_V];
16  double h[MAX_V];
17  double d[MAX_V];
18  int prevv[MAX_V];
19  int preve[MAX_V];
20  // int match[MAX_V];
21
22  void add_edge(int u, int v, int cap, double cost) {
23      g[u].push_back((Edge){v, cap, cost, (int)g[v].size()});
24      g[v].push_back((Edge){u, 0, -cost, (int)g[u].size() - 1});
25  }
26
27  double min_cost_flow(int s, int t, int f) {
28      double res = 0;
29      fill(h, h + V, 0);
30      fill(match, match + V, -1);
31      while (f > 0) {
32          // dijkstra 找最小成本增廣路徑
33          // without h will reduce to SPFA = O(V*E)
34          fill(d, d + V, INF);
35          priority_queue< pii, vector<pii>, greater<pii> > pq;
36
37          d[s] = 0;
38          pq.push(pii(d[s], s));
39
40          while (!pq.empty()) {
41              pii p = pq.top(); pq.pop();
42              int v = p.nd;
43              if (d[v] < p.st) continue;
44              for (size_t i = 0; i < g[v].size(); i++) {
45                  const Edge& e = g[v][i];
46                  if (e.cap > 0 && d[e.to] > d[v] + e.cost + h[v] -
↪   h[e.to]) {
47                      d[e.to] = d[v] + e.cost + h[v] - h[e.to];
48                      prevv[e.to] = v;
49                      preve[e.to] = i;
50                      pq.push(pii(d[e.to], e.to));
51                  }
52              }
53          }
54
55          // 找不到增廣路徑
56          if (d[t] == INF) return -1; // double 時不能這樣判
57
58          // 維護 h[v]
59          for (int v = 0; v < V; v++)
60              h[v] += d[v];
61
62          // 找瓶頸
63          int bn = f;
64          for (int v = t; v != s; v = prevv[v])
65              bn = min(bn, g[prevv[v]][preve[v]].cap);
66
```

```cpp
67      // // find match
68      // for (int v = prevv[t]; v != s; v = prevv[prevv[v]]) {
69      //     int u = prevv[v];
70      //     match[v] = u;
71      //     match[u] = v;
72      // }
73
74      // 更新剩餘圖
75      f -= bn;
76      res += bn * h[t]; // SPFA: res += bn * d[t]
77      for (int v = t; v != s; v = prevv[v]) {
78          Edge& e = g[prevv[v]][preve[v]];
79          e.cap -= bn;
80          g[v][e.rev].cap += bn;
81      }
82    }
83    return res;
84 }
```

10.3   Bipartite Matching, Unweighted

```cpp
1  const int MAX_V = ...;
2  int V;
3  vector<int> g[MAX_V];
4  int match[MAX_V];
5  bool used[MAX_V];
6
7  void add_edge(int u, int v) {
8      g[u].push_back(v);
9      g[v].push_back(u);
10 }
11
12 // 回傳有無找到從 v 出發的增廣路徑
13 // (首尾都為未匹配點的交錯路徑)
14 // [待確認] 每次遞迴都找一個未匹配點 v 及匹配點 u
15 bool dfs(int v) {
16     used[v] = true;
17     for (size_t i = 0; i < g[v].size(); i++) {
18         int u = g[v][i], w = match[u];
19         // 尚未配對或可從 w 找到增廣路徑 (即路徑繼續增長)
20         if (w < 0 || (!used[w] && dfs(w))) {
21             // 交錯配對
22             match[v] = u;
23             match[u] = v;
24             return true;
25         }
26     }
27     return false;
28 }
29
30 int bipartite_matching() { // 匈牙利演算法
31     int res = 0;
32     memset(match, -1, sizeof(match));
33     for (int v = 0; v < V; v++) {
34         if (match[v] == -1) {
35             memset(used, false, sizeof(used));
```

```cpp
36         if (dfs(v)) {
37             res++;
38         }
39     }
40   }
41   return res;
42 }
```

11   String

11.1   Rolling Hash

1. Use two rolling hashes if needed.
2. The prime for pre-calculation can be 137 and 257, for modulo can be $1e9 + 7$ and $0xdefaced$

```cpp
1  #define N 1000100
2  #define B 137
3  #define M 1000000007
4
5  typedef long long ll;
6
7  char inp[N];
8  int len;
9  ll p[N], h[N];
10
11 void init()
12 { // build polynomial table and hash value
13     p[0] = 1; // b to the ith power
14     for (int i = 1; i <= len; i++) {
15         h[i] = (h[i - 1] * B % M + inp[i - 1]) % M; // hash value
16         p[i] = p[i - 1] * B % M;
17     }
18 }
19
20 ll get_hash(int l, int r) // [l, r] of the inp string array
21 {
22     return ((h[r + 1] - (h[l] * p[r - l + 1])) % M + M) % M;
23 }
```

11.2   KMP

```cpp
1  void fail()
2  {
3      int len = strlen(pat);
4
5      f[0] = 0;
6      int j = 0;
7      for (int i = 1; i < len; i++) {
8          while (j != 0 && pat[i] != pat[j])
9              j = f[j - 1];
10
11         if (pat[i] == pat[j])
```

```
12              j++;
13
14          f[i] = j;
15      }
16  }
17
18  int match()
19  {
20      int res = 0;
21      int j = 0, plen = strlen(pat), tlen = strlen(text);
22
23      for (int i = 0; i < tlen; i++) {
24          while (j != 0 && text[i] != pat[j])
25              j = f[j - 1];
26
27          if (text[i] == pat[j]) {
28              if (j == plen - 1) { // find match
29                  res++;
30                  j = f[j];
31              } else {
32                  j++;
33              }
34          }
35      }
36
37      return res;
38  }
```

11.3   Z Algorithm

```
1   int len = strlen(inp), z[len];
2   z[0] = 0; // initial
3
4   int l = 0, r = 0; // z box bound [l, r]
5   for (int i = 1; i < len; i++)
6   {
7       if (i > r) { // i not in z box
8           l = r = i; // z box contains itself only
9           while (r < len && inp[r - l] == inp[r])
10              r++;
11          z[i] = r - l;
12          r--;
13      } else { // i in z box
14          if (z[i - l] + i < r) // over shoot R bound
15              z[i] = z[i - l];
16          else {
17              l = i;
18              while (r < len && inp[r - l] == inp[r])
19                  r++;
20              z[i] = r - l;
21              r--;
22          }
23      }
24  }
```

11.4   Trie

注意 count 的擺放位置，視題意可以擺在迴圈外

```
1   struct Node {
2       int cnt;
3       Node* nxt[2];
4       Node() {
5           cnt = 0;
6           fill(nxt, nxt + 2, nullptr);
7       }
8   };
9
10  const int MAX_Q = 200000;
11  int Q;
12
13  int NN = 0;
14  Node data[MAX_Q * 30];
15  Node* root = &data[NN++];
16
17  void insert(Node* u, int x) {
18      for (int i = 30; i >= 0; i--) {
19          int t = ((x >> i) & 1);
20          if (u->nxt[t] == nullptr) {
21              u->nxt[t] = &data[NN++];
22          }
23
24          u = u->nxt[t];
25          u->cnt++;
26      }
27  }
28
29  void remove(Node* u, int x) {
30      for (int i = 30; i >= 0; i--) {
31          int t = ((x >> i) & 1);
32          u = u->nxt[t];
33          u->cnt--;
34      }
35  }
36
37  int query(Node* u, int x) {
38      int res = 0;
39      for (int i = 30; i >= 0; i--) {
40          int t = ((x >> i) & 1);
41          // if it is possible to go the another branch
42          // then the result of this bit is 1
43          if (u->nxt[t ^ 1] != nullptr && u->nxt[t ^ 1]->cnt > 0) {
44              u = u->nxt[t ^ 1];
45              res |= (1 << i);
46          }
47          else {
48              u = u->nxt[t];
49          }
50      }
51      return res;
52  }
```

12 Matrix

12.1 Gauss Jordan Elimination

```
1   typedef long long ll;
2   typedef vector<ll> vec;
3   typedef vector<vec> mat;
4
5   vec gauss_jordan(mat A) {
6       int n = A.size(), m = A[0].size(); // 增廣矩陣
7       for (int i = 0; i < n; i++) {
8           // float: find j s.t. A[j][i] is max
9           // mod: find min j s.t. A[j][i] is not 0
10          int pivot = i;
11          for (int j = i; j < n; j++) {
12              // if (fabs(A[j][i]) > fabs(A[pivot])) {
13              //     pivot = j;
14              // }
15              if (A[pivot][i] != 0) {
16                  pivot = j;
17                  break;
18              }
19          }
20
21          swap(A[i], A[pivot]);
22          if (A[i][i] == 0) { // if (fabs(A[i][i]) < eps)
23              // 無解或無限多組解
24              // 可改成 continue，全部做完後再判
25              return vec();
26          }
27
28          ll divi = inv(A[i][i]);
29          for (int j = i; j < m; j++) {
30              // A[i][j] /= A[i][i];
31              A[i][j] = (A[i][j] * divi) % MOD;
32          }
33
34          for (int j = 0; j < n; j++) {
35              if (j != i) {
36                  for (int k = i + 1; k < m; k++) {
37                      // A[j][k] -= A[j][i] * A[i][k];
38                      ll p = (A[j][i] * A[i][k]) % MOD;
39                      A[j][k] = (A[j][k] - p + MOD) % MOD;
40                  }
41              }
42          }
43      }
44
45      vec x(n);
46      for (int i = 0; i < n; i++)
47          x[i] = A[i][m - 1];
48      return x;
49  }
```

12.2 Determinant

整數版本

```
1   typedef long long ll;
2   typedef vector<ll> vec;
3   typedef vector<vec> mat;
4
5   ll determinant(mat m) { // square matrix
6       const int n = m.size();
7       ll det = 1;
8       for (int i = 0; i < n; i++) {
9           for (int j = i + 1; j < n; j++) {
10              int a = i, b = j;
11              while (m[b][i]) {
12                  ll q = m[a][i] / m[b][i];
13                  for (int k = 0; k < n; k++)
14                      m[a][k] = m[a][k] - m[b][k] * q;
15                  swap(a, b);
16              }
17
18              if (a != i) {
19                  swap(m[i], m[j]);
20                  det = -det;
21              }
22          }
23
24          if (m[i][i] == 0)
25              return 0;
26          else
27              det *= m[i][i];
28      }
29      return det;
30  }
```

13 Geometry

1. Keep things in integers as much as possible!
2. Try not to divide
3. If you have decimals, if they are fixed precision, you can usually just multiply all the input and use integers instead

13.1 EPS

$= 0$: $fabs \leq eps$
< 0: $< -eps$
> 0: $> +eps$

```
1   // if the points are given in doubles form, change the code accordingly
2
3   typedef long long ll;
4
5   typedef pair<ll, ll> pt;   // points are stored using long long
6   typedef pair<pt, pt> seg; // segments are a pair of points
7
8   #define x first
```

```
9    #define y second
10
11   #define EPS 1e-9
12
13   pt operator+(pt a, pt b)
14   {
15       return pt(a.x + b.x, a.y + b.y);
16   }
17
18   pt operator-(pt a, pt b)
19   {
20       return pt(a.x - b.x, a.y - b.y);
21   }
22
23   pt operator*(pt a, int d)
24   {
25       return pt(a.x * d, a.y * d);
26   }
27
28   ll cross(pt a, pt b)
29   {
30       return a.x * b.y - a.y * b.x;
31   }
32
33   int ccw(pt a, pt b, pt c)
34   {
35       ll res = cross(b - a, c - a);
36       if (res > 0) // left turn
37           return 1;
38       else if (res == 0) // straight
39           return 0;
40       else // right turn
41           return -1;
42   }
43
44   double dist(pt a, pt b)
45   {
46       double dx = a.x - b.x;
47       double dy = a.y - b.y;
48       return sqrt(dx * dx + dy * dy);
49   }
50
51   bool zero(double x)
52   {
53       return fabs(x) <= EPS;
54   }
55
56   bool overlap(seg a, seg b)
57   {
58       return ccw(a.x, a.y, b.x) == 0 && ccw(a.x, a.y, b.y) == 0;
59   }
60
61   bool intersect(seg a, seg b)
62   {
63       if (overlap(a, b) == true) { // non-proper intersection
64           double d = 0;
```

```
65           d = max(d, dist(a.x, a.y));
66           d = max(d, dist(a.x, b.x));
67           d = max(d, dist(a.x, b.y));
68           d = max(d, dist(a.y, b.x));
69           d = max(d, dist(a.y, b.y));
70           d = max(d, dist(b.x, b.y));
71
72           // d > dist(a.x, a.y) + dist(b.x, b.y)
73           if (d - (dist(a.x, a.y) + dist(b.x, b.y)) > EPS)
74               return false;
75           return true;
76       }
77       //                      |
78       // Equal sign for -----| case
79       // non qeual sign => proper intersection
80       if (ccw(a.x, a.y, b.x) * ccw(a.x, a.y, b.y) <= 0 &&
81           ccw(b.x, b.y, a.x) * ccw(b.x, b.y, a.y) <= 0)
82           return true;
83       return false;
84   }
85
86   double area(vector<pt> pts)
87   {
88       double res = 0;
89       int n = pts.size();
90       for (int i = 0; i < n; i++)
91           res += (pts[i].y + pts[(i + 1) % n].y) * (pts[(i + 1) % n].x -
     ↪ pts[i].x);
92       return res / 2.0;
93   }
94
95   vector<pt> halfHull(vector<pt> &points)
96   {
97       vector<pt> res;
98
99       for (int i = 0; i < (int)points.size(); i++) {
100          while ((int)res.size() >= 2 &&
101              ccw(res[res.size() - 2], res[res.size() - 1], points[i]) <
     ↪ 0)
102              res.pop_back(); // res.size() - 2 can't be assign before
     ↪ size() >= 2
103              // check, bitch
104
105          res.push_back(points[i]);
106      }
107
108      return res;
109  }
110
111  vector<pt> convexHull(vector<pt> &points)
112  {
113      vector<pt> upper, lower;
114
115      // make upper hull
116      sort(points.begin(), points.end());
117
```

```
118        upper = halfHull(points);
119        // make lower hull
120        reverse(points.begin(), points.end());
121        lower = halfHull(points);
122
123        // merge hulls
124        if ((int)upper.size() > 0) // yes sir~
125            upper.pop_back();
126        if ((int)lower.size() > 0)
127            lower.pop_back();
128
129        vector<pt> res(upper.begin(), upper.end());
130        res.insert(res.end(), lower.begin(), lower.end());
131
132        return res;
133 }
```

13.2 Rectangle area

```
1  #define sz(x) (int(x.size()))
2
3  const int MAX_NN = (1 << 17);
4
5  struct Rect {
6      double x1, y1, x2, y2;
7  };
8
9  struct Event {
10     double y; int x1, x2, type;
11     bool operator < (const Event& e) const {
12         if (y == e.y)
13             return type < e.type;
14         return y < e.y;
15     }
16 };
17
18 vector<double> xs;
19
20 struct SegTree {
21     int NN;
22     int cnt[MAX_NN];
23     double len[MAX_NN];
24
25     void init(int n) {
26         NN = 1;
27         while (NN < n)
28             NN <<= 1;
29         fill(cnt, cnt + 2 * NN, 0);
30         fill(len, len + 2 * NN, double(0.0));
31     }
32
33     void maintain(int u, int l, int r) {
34         if (cnt[u] > 0) len[u] = xs[r] - xs[l];
35         else {
36             if (u >= NN - 1)
37                 len[u] = 0;
```

```
38             else
39                 len[u] = len[u * 2 + 1] + len[u * 2 + 2];
40         }
41     }
42
43     void update(int a, int b, int x, int u, int l, int r) { // [a, b),
    ↪ [l, r)
44         if (r <= a || l >= b) return;
45         if (a <= l && r <= b) {
46             cnt[u] += x;
47             maintain(u, l, r);
48             return;
49         }
50         int m = (l + r) / 2;
51         update(a, b, x, u * 2 + 1, l, m);
52         update(a, b, x, u * 2 + 2, m, r);
53         maintain(u, l, r);
54     }
55 };
56
57 double get_union_area(const vector<Rect>& rect) {
58     // 離散化 x
59     xs.clear();
60     for (int i = 0; i < sz(rect); i++) {
61         xs.push_back(rect[i].x1);
62         xs.push_back(rect[i].x2);
63     }
64     sort(xs.begin(), xs.end());
65     xs.resize(unique(xs.begin(), xs.end()) - xs.begin());
66
67     // sweep line events
68     vector<Event> es;
69     for (int i = 0; i < sz(rect); i++) {
70         int x1 = lower_bound(xs.begin(), xs.end(), rect[i].x1) -
    ↪ xs.begin();
71         int x2 = lower_bound(xs.begin(), xs.end(), rect[i].x2) -
    ↪ xs.begin();
72         es.push_back((Event) {rect[i].y1, x1, x2, +1}); // bottom
73         es.push_back((Event) {rect[i].y2, x1, x2, -1}); // top
74     }
75     sort(es.begin(), es.end());
76
77     // find total area
78     SegTree seg;
79     seg.init(sz(xs));
80     seg.update(es[0].x1, es[0].x2, es[0].type, 0, 0, seg.NN);
81
82     double res = 0;
83     for (int i = 1; i < sz(es); i++) {
84         res += seg.len[0] * (es[i].y - es[i - 1].y);
85         seg.update(es[i].x1, es[i].x2, es[i].type, 0, 0, seg.NN);
86     }
87
88     return res;
89 }
```

14 Math

14.1 Euclid's formula (Pythagorean Triples)

$a = p^2 - q^2$
$b = 2pq$ **(always even)**
$c = p^2 + q^2$

14.2 Difference between two consecutive numbers' square is odd

$(k+1)^2 - k^2 = 2k + 1$

14.3 Summation

$\sum_{k=1}^{n} 1 = n$
$\sum_{k=1}^{n} k = \frac{n(n+1)}{2}$
$\sum_{k=1}^{n} k^2 = \frac{n(n+1)(2n+1)}{6}$
$\sum_{k=1}^{n} k^3 = \frac{n^2(n+1)^2}{4}$

14.4 Combination

14.4.1 Pascal triangle

```
#define N 210
ll C[N][N];

void Combination() {
    for(ll i=0; i<N; i++) {
        C[i][0] = 1;
        C[i][i] = 1;
    }

    for(ll i=2; i<N; i++) {
        for(ll j=1; j<=i; j++) {
            C[i][j] = (C[i-1][j] + C[i-1][j-1])%M; // if needed, mod it
        }
    }
}
```

14.4.2 Lucus

$$\binom{n}{m} \equiv \prod_{i=0}^{k} \binom{n_i}{m_i} \pmod{p}$$

where

$n = n_k p^k + n_{k-1} p^{k-1} + \cdots + n_1 p + n_0,$

$m = m_k p^k + m_{k-1} p^{k-1} + \cdots + m_1 p + m_0$

p is prime$c$

```
1  typedef long long ll;
2
3  ll fast_pow(ll a, ll b, ll p) {
4      ll ans = 1;
5      ll base = a % p;
6      b = b % (p - 1); // Fermat's little theorem
7      while (b) {
8          if (b & 1) {
9              ans = (ans * base) % p;
10         }
11         base = (base * base) % p;
12         b >>= 1;
13     }
14     return ans;
15 }
16
17 ll inv(ll a, ll p) {
18     return fast_pow(a, p - 2, p);
19 }
20
21 ll C(ll n, ll m, ll p) {
22     if (n < m) return 0;
23     m = min(m, n - m);
24     ll nom = 1, den = 1;
25     for (ll i = 1; i <= m; i++) {
26         nom = (nom * (n - i + 1)) % p;
27         den = (den * i) % p;
28     }
29     return (nom * inv(den, p)) % p;
30 }
31
32 // To make C(n, m) % p computed in O(log(p, n) * p) instead of O(m)
33 // https://en.wikipedia.org/wiki/Lucas's_theorem
34 ll lucas(ll n, ll m, ll p) {
35     if (m == 0) return 1;
36     return C(n % p, m % p, p) * lucas(n / p, m / p, p) % p;
37 }
```

14.4.3 線性

```
ll binomialCoeff(ll n, ll k)
{
    ll res = 1;

    if ( k > n - k ) // Since C(n, k) = C(n, n-k)
        k = n - k;

    for (int i = 0; i < k; ++i) // n...n-k / 1...k
    {
        res *= (n - i);
        res /= (i + 1);
    }

    return res;
}
```

(1)

14.5 Chinese remainder theorem

```
1  typedef long long ll;
2
3  struct Item {
```

```
 4        ll m, r;
 5    };
 6
 7    Item extcrt(const vector<Item> &v)
 8    {
 9        ll m1 = v[0].m, r1 = v[0].r, x, y;
10
11        for (int i = 1; i < int(v.size()); i++) {
12            ll m2 = v[i].m, r2 = v[i].r;
13            ll g = extgcd(m1, m2, x, y); // now x = (m/g)^(-1)
14
15            if ((r2 - r1) % g != 0)
16                return {-1, -1};
17
18            ll k = (r2 - r1) / g * x % (m2 / g);
19            k = (k + m2 / g) % (m2 / g); // for the case k is negative
20
21            ll m = m1 * m2 / g;
22            ll r = (m1 * k + r1) % m;
23
24            m1 = m;
25            r1 = (r + m) % m; // for the case r is negative
26        }
27
28        return (Item) {
29            m1, r1
30        };
31    }
```

14.6   2-Circle relations

$d =$ 圓心距, R, r 為半徑 $(R \geq r)$
內切: $d = R - r$
外切: $d = R + r$
內離: $d < R - r$
外離: $d > R + r$
相交: $d < R + r$ 且 $d > R - r$

14.7   Fun Facts

1. 如果 $\frac{b}{a}$ 是最簡分數, 則 $1 - \frac{b}{a}$ 也是

15   Dynamic Programming - Problems collection

```
# 零一背包 (poj 1276)
fill(dp, dp + W + 1, 0);
for (int i = 0; i < N; i++)
    for (int j = W; j >= items[i].w; j--)
        dp[j] = max(dp[j], dp[j - w[i]] + v[i]);
return dp[W];

# 多重背包二進位拆解 (poj 1276)
for_each(ll v, w, num) {
    for (ll k = 1; k <= num; k *= 2) {
        items.push_back((Item) {k * v, k * w});
        num -= k;
```

```
    }
    if (num > 0)
        items.push_back((Item) {num * v, num * w});
}

# 完全背包
dp[i][j] = 前 i + 1 個物品, 在重量 j 下所能組出的最大價值
第 i 個物品, 不放或至少放一個
dp[i][j] = max(dp[i - 1][j], dp[i][j - w[i]] + v[i])
fill(dp, dp + W + 1, 0);
for (int i = 0; i < N; i++)
    for (int j = w[i]; j <= W; j++)
        dp[j] = max(dp[j], dp[j - w[i]] + v[i]);
return dp[W];

# Coin Change (2015 桂冠賽 E)
dp[i][j] = 前 i + 1 個物品, 組出 j 元的方法數
第 i 個物品, 不用或用至少一個
dp[i][j] = dp[i - 1][j] + dp[i][j - coin[i]]

# Cutting Sticks (2015 桂冠賽 F)
補上二個切點在最左與最右
dp[i][j] = 使 (i, j) 區間中的所有切點都被切的最小成本
dp[i][j] = min(dp[i][c] + dp[c][j] + (p[j] - p[i]) for i < c < j)
dp[i][i + 1] = 0
ans = dp[0][N + 1]

# Throwing a Party (itsa dp 06)
給定一棵有根樹, 代表公司職位層級圖, 每個人有其權重, 現從中選一個點集合出來,
且一個人不能與其上司一都在集合中, 並最大化集合的權重和, 輸出該總和。
dp[u][0/1] = u 在或不在集合中, 以 u 為根的子樹最大權重和
dp[u][0] = max(max(dp[c][0], dp[c][1]) for children c of u) + val[u]
dp[u][1] = max(dp[c][0] for children c of u)
bottom up dp

# LIS (O(N^2))
dp[i] = 以 i 為結尾的 LIS 的長度
dp[i] = max(dp[j] for 0 <= j < i) + 1
ans = max(dp)

# LIS (O(nlgn)), poj 1631
dp[i] = 長度為 i + 1 的 LIS 的最後一項的最小值, 不存在時為 INF
fill(dp, dp + N, INF);
for (int i = 0; i < N; i++)
    *lower_bound(dp, dp + N, A[i]) = A[i];
ans = lower_bound(dp, dp + N, INF) - dp;

# Maximum Subarray

# Not equal on a Segment (cf edu7 C)
給定長度為 n 的陣列 a[] 與 m 個詢問。
針對每個詢問 l, r, x 請輸出 a[l, r] 中不等於 x 的任一位置。
不存在時輸出 -1
dp[i] = max j such that j < i and a[j] != a[i]
dp[0] = -1
dp[i] = dp[i - 1] if a[i] == a[i - 1] else i - 1
針對每筆詢問 l, r, x
1. a[r] != x            -> 輸出 r
2. a[r] = x && dp[r] >= l   -> 輸出 dp[r]
3. a[r] = x && dp[r] < l    -> 輸出 -1

# bitmask dp, poj 2686
給定一個無向帶權圖, 代表 M 個城市之間的路, 與 N 張車票,
每張車票有一個數值 t[i], 若欲使用車票 t[i] 從城市 u 經由路徑 d[u][v] 走到城市 v,
所花的時間為 d[u][v] / t[i]。請問, 從城市 A 走到城市 B 最快要多久?
dp[S][v] = 從城市 A 到城市 v 的最少時間, 其中 S 為用過的車票的集合
考慮前一個城市 u 是誰, 使用哪個車票 t[i] 而來, 可以得到轉移方程式:
```

```
dp[S][v] = min([
    dp[S - {v}][u] + d[u][v] / t[i]
    for all city u has edge to v, for all ticket in S
])
```

Tug of War
N 個人參加拔河比賽，每個人有其重量 w[i]，欲使二隊的人數最多只差一，雙方的重量和越接近越好
請問二隊的重量和分別是多少?
dp[i][j][k] = 只考慮前 i + 1 個人，可不可以使左堆的重量為 j，且左堆的人數為 k
dp[i][j][k] = dp[i - 1][j - w[i]][k - 1] or dp[i - 1][j][k]
dp[i][j] = (dp[i - 1][j - w[i]] << 1) | (dp[i - 1][j])

Modulo Sum (cf 319 B)
給定長度為 N 的序列 A 與一正整數 M，請問該序列中有無一個子序列，子序列的總合是 M 的倍數
若 N > M，則根據鴿籠原理，必有至少兩個前綴和的值 mod M 為相同值，解必定存在
dp[i][j] = 前 i + 1 個數可否組出 mod m = j 的數
dp[i][j] = true if
    dp[i - 1][(j - (a[i] mod m)) mod m] or
    dp[i - 1][j] or
    j = a[i] % m

POJ 2229
給定正整數 N，請問將 N 拆成一堆 2^x 之和的方法數
dp[i] = 拆解 N 的方法數
dp[i] = dp[i / 2] if i is odd
      = dp[i - 1] + dp[i / 2] if i is even

POJ 3616
給定 N 個區間 [s, t)，每個區間有權重 w[i]，從中選出一些不相交的區間，使權重和最大
dp[i] = 考慮前 i + 1 個區間，且必選第 i 個區間的最大權重和
dp[i] = max(dp[j] | 0 <= j < i) + w[i]
ans = max(dp)

POJ 2184
N 隻牛每隻牛有權重 <s, f>，從中選出一些牛的集合，
使得 sum(s) + sum(f) 最大，且 sum(s) > 0, sum(f) > 0。
枚舉 sum(s)，將 sum(s) 視為重量對 f 做零一背包。

POJ 3666
給定長度為 N 的序列，請問最少要加多少值，使得序列單調遞增
dp[i][j] = 使序列前 i+1 項變為單調，且將 A[i] 變為「第 j 小的數」的最小成本
dp[i][j] = min(dp[i - 1][k] | 0 <= k <= j) + abs(S[j] - A[i])
min(dp[i - 1][k] | 0 <= k <= j) 動態維護
for (int j = 0; j < N; j++)
```
dp[0][j] = abs(S[j] - A[0]);
for (int i = 1; i < N; i++) {
    int pre_min_cost = dp[i][0];
    for (int j = 0; j < N; j++) {
        pre_min_cost = min(pre_min_cost, dp[i-1][j]);
        dp[i][j] = pre_min_cost + abs(S[j] - A[i]);
    }
}
ans = min(dp[N - 1])
```

POJ 3734
N 個 blocks 上色，R, G, Y, B，上完色後紅色的數量與綠色的數量都要是偶數。請問方法數。
dp[i][0/1/2/3] = 前 i 個 blocks 上完色，紅色數量為奇數/偶數，綠色數量為數/偶數
用遞推，考慮第 i + 1 個 block 的顏色，找出個狀態的轉移，整理可發現
dp[i + 1][0] = dp[i][2] + dp[i][1] + 2 * dp[i][0]
dp[i + 1][1] = dp[i][3] + dp[i][0] + 2 * dp[i][1]
dp[i + 1][2] = dp[i][0] + dp[i][3] + 2 * dp[i][2]
dp[i + 1][3] = dp[i][1] + dp[i][2] + 2 * dp[i][3]
矩陣快速冪加速求 dp[N - 1][0][0]

POJ 3171
數線上，給定 N 個區間 [s[i], t[i]]，每個區間有其代價，求覆蓋區間 [M, E] 的最小代價。
dp[i][j] = 最多使用前 i + 1 個區間，使 [M, j] 被覆蓋的最小代價
考慮第 i 個區間用或不用，可得:
dp[i][j] =
    1. min(dp[i - 1][k] for k in [s[i] - 1, t[i]]) + cost[i] if j = t[i]
    2. dp[i - 1][j] if j ≠ t[i]
壓空間，使用線段樹加速。
dp[t[i]] = min(dp[t[i]],
    min(dp[i - 1][k] for k in [s[i] - 1, t[i]]) + cost[i]
)
fill(dp, dp + E + 1, INF);
seg.init(E + 1, INF);
int idx = 0;
while (idx < N && A[idx].s == 0) {
    dp[A[idx].t] = min(dp[A[idx].t], A[idx].cost);
    seg.update(A[idx].t, A[idx].cost);
    idx++;
}
for (int i = idx; i < N; i++) {
    ll v = min(dp[A[i].t], seg.query(A[i].s - 1, A[i].t + 1) + A[i].cost);
    dp[A[i].t] = v;
    seg.update(A[i].t, v);
}
```