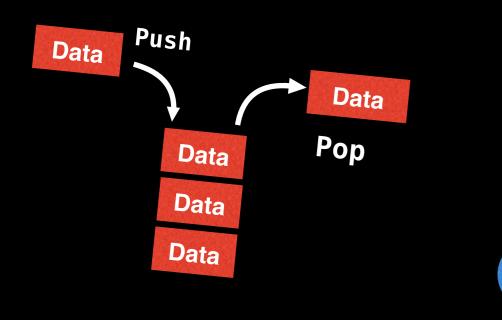
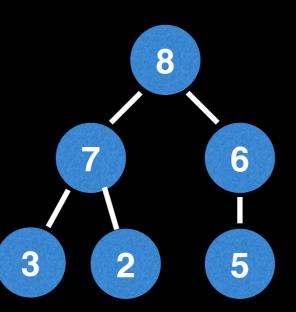
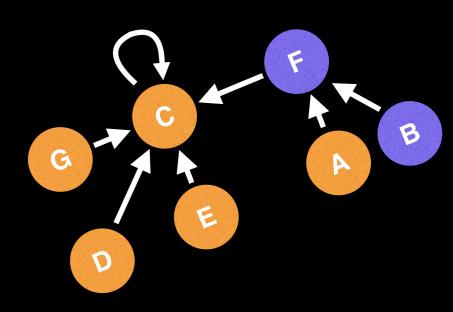


Data Structure Video Series







Indexed Priority Queues

William Fiset

Priority Queue Videos

Display of previous PQ vids

What is an Indexed Priority Queue?

An Indexed Priority Queue is a traditional priority queue variant which on top of the regular PQ operations supports quick updates and deletions of key-value pairs.

Suppose a hospital has a waiting room with N people which need attention with different levels of priority:



Priority: 9



Mary is in labour Akarsh has a paper cut Priority: 1



James has an arrow in his leg Priority: 7



Naida's stomach hurts

Priority: 3

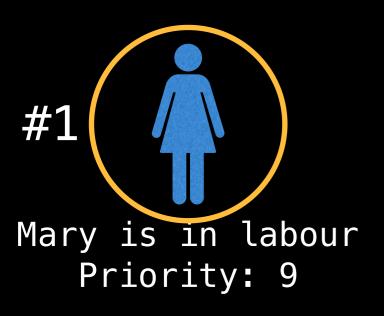


Richard has a fractured wrist Priority: 5

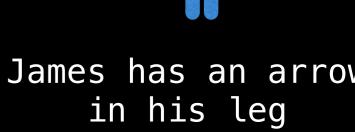


Leah's stomach hurts Priority: 3

The hospital serves Mary first.



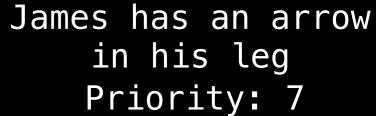




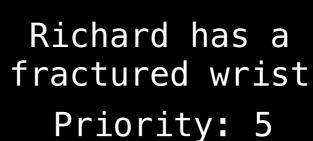




Priority: 1



Naida's stomach hurts Priority: 3





Leah's stomach hurts Priority: 3

Followed by James





Akarsh has a paper cut Priority: 1



James has an arrow in his leg
Priority: 7



Naida's stomach hurts Priority: 3



Richard has a fractured wrist Priority: 5



Leah's stomach hurts
Priority: 3

Then something happens, Naida's condition worsens as she starts vomiting. Her priority gets updated to 6





Mary is in labour Akarsh has a paper cut Priority: 9 Priority: 1



James has an arrow in his leg
Priority: 7



Naida's stomach hurts + vomiting Priority: 6



Richard has a fractured wrist Priority: 5



Leah's stomach hurts
Priority: 3

Naida gets served when the hospital is finished with James.





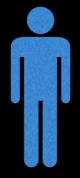
Akarsh has a paper cut Priority: 1



James has an arrow in his leg
Priority: 7



Naida's stomach hurts + vomiting Priority: 6



Richard has a fractured wrist Priority: 5



Leah's stomach hurts
Priority: 3

Richard gets impatient while Naida gets served and leaves to go to the clinic down the street.

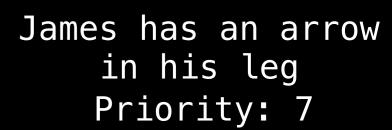






Akarsh has a paper cut Priority: 1







Richard has a fractured wrist Priority: 5



Leah's stomach hurts Priority: 3

While Akarsh goes to take a drink of water he slips and cracks his skull open. His priority is increased to 10



Mary is in labour Priority: 9



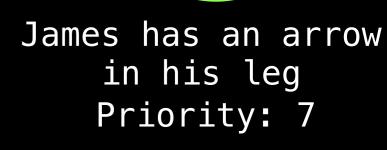
Naida's stomach hurts + vomiting Priority: 6



Akarsh has a paper cut + open skull

Priority: (10)





#2



Richard has a fractured wrist Priority: 5

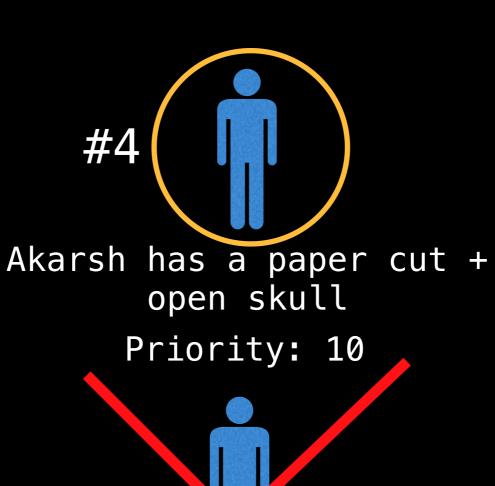
Leah's stomach hurts

Priority: 3

Akarsh gets served next.





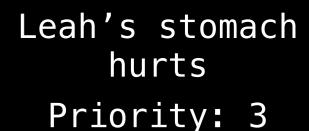






Priority: 7

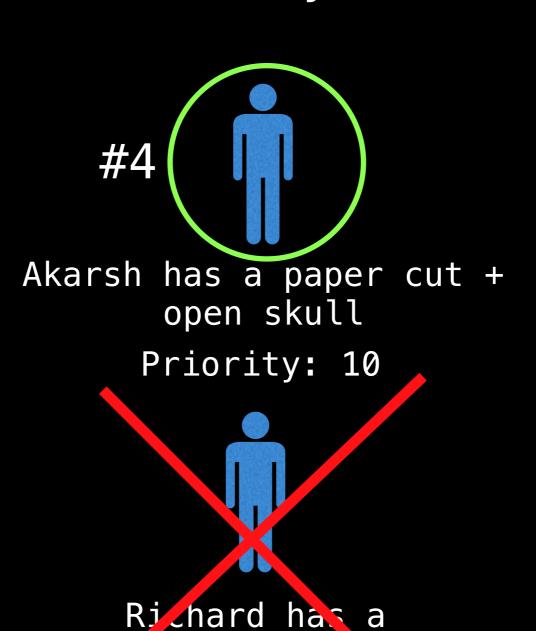




Followed by Leah.







fractured wrist

Priority: 5

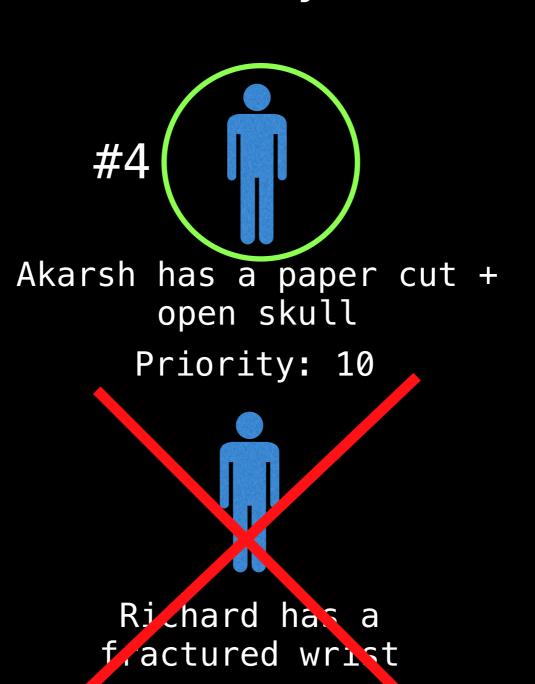


#2

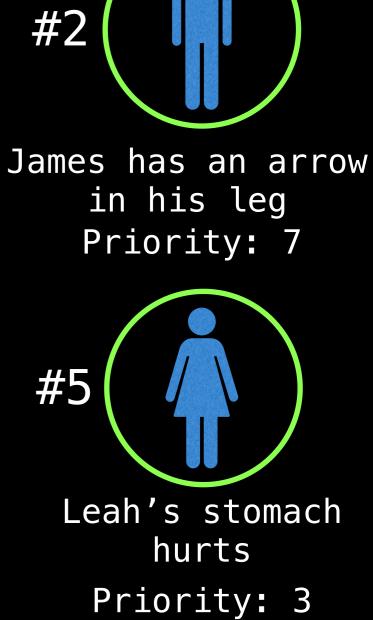
Followed by Leah.







Priority: 5

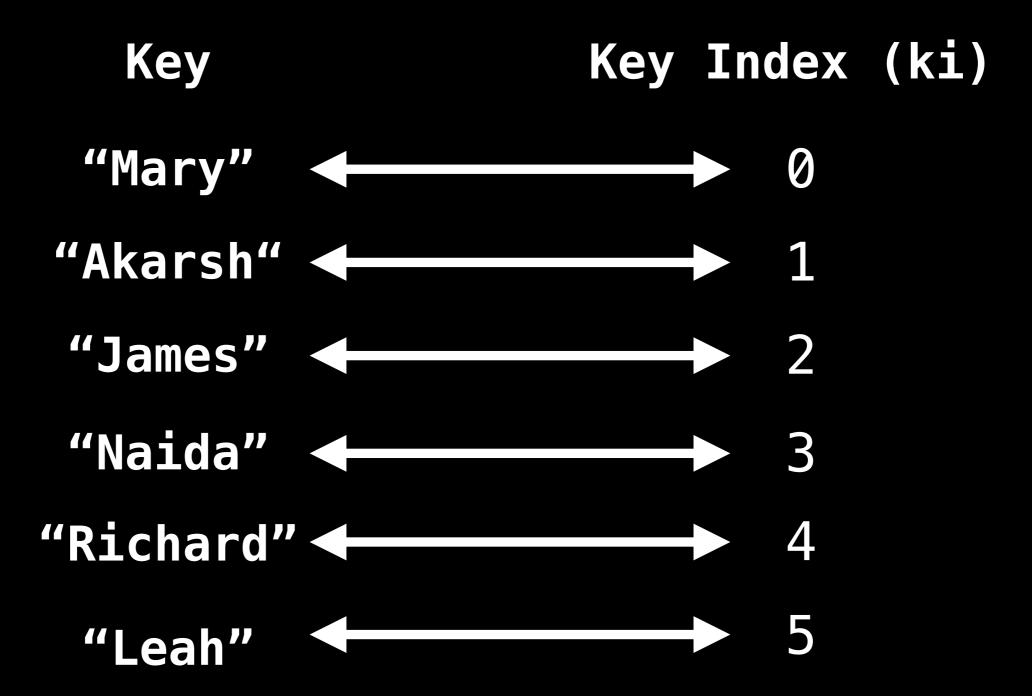


Usefulness of IPQ

In the hospital example, we saw that it was very important to be able to dynamically update the priority (value) of certain people (keys).

The Indexed Priority Queue (IPQ) data structure lets us do this efficiently. The first step to using an IPQ is to assign index values to all the keys forming a bidirectional mapping.

Create a bidirectional mapping between your **N** keys and the domain [0, **N**) using a bidirectional hashtable.



NOTE: This assumes you know how many keys you will have in your IPQ, but this mapping can be constructed dynamically as well

Reason for mapping

Q: Why are we mapping keys to indexes in the domain [0, N)?

A: Typically priority queues are implemented as heaps under the hood which internally use arrays which we want to facilitate indexing into.

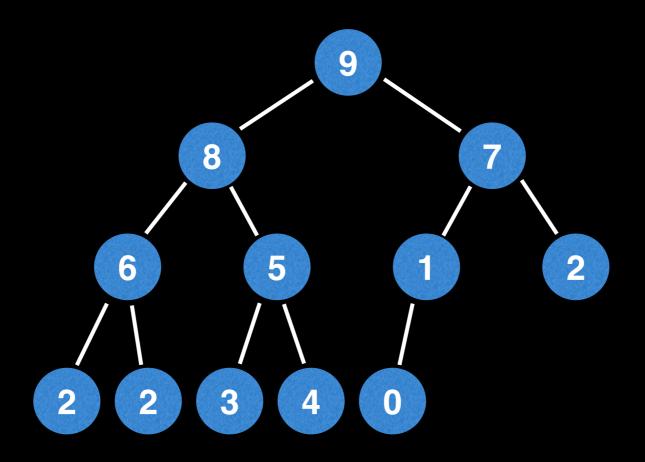
NOTE: Often the keys themselves are integers in the range [0, N) so there's no need for the mapping, but it's handy to be able to support any type of key (like names).

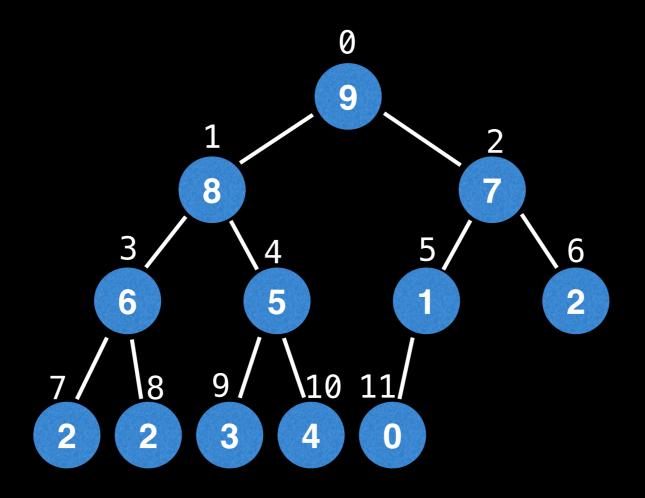
IPQ ADT Interface

```
delete(ki)
valueOf(ki)
contains(ki)
peekMinKeyIndex()
pollMinKeyIndex()
peekMinValue()
pollMinValue()
insert(ki, value)
update(ki, value)
decreaseKey(ki, value)
increaseKey(ki, value)
```

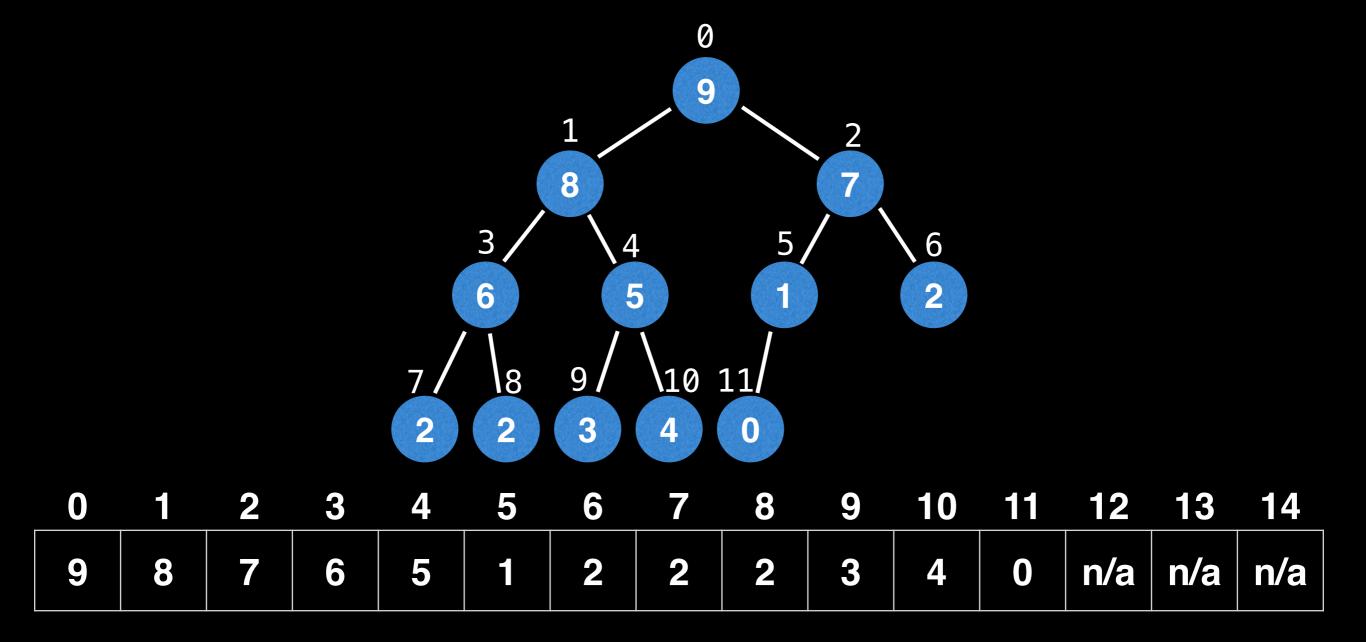
IPQ as a binary heap

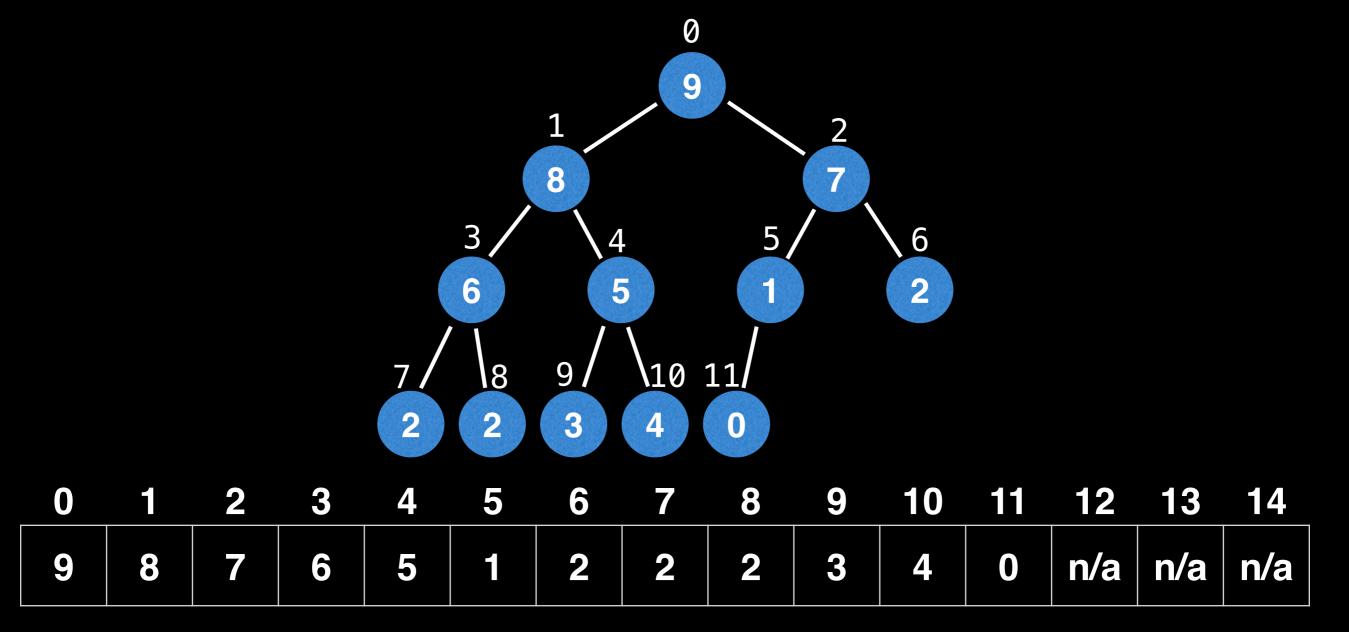
Operation	Indexed Binary Heap PQ
delete(ki)	O(log(n))
valueOf(ki)	O(1)
contains(ki)	O(1)
<pre>peekMinKeyIndex()</pre>	O(1)
<pre>pollMinKeyIndex()</pre>	O(log(n))
peekMinValue()	O(1)
pollMinValue()	O(log(n))
<pre>insert(ki, value)</pre>	O(log(n))
<pre>update(ki, value)</pre>	O(log(n))
<pre>decreaseKey(ki, value)</pre>	O(log(n))
<pre>increaseKey(ki, value)</pre>	O(log(n))





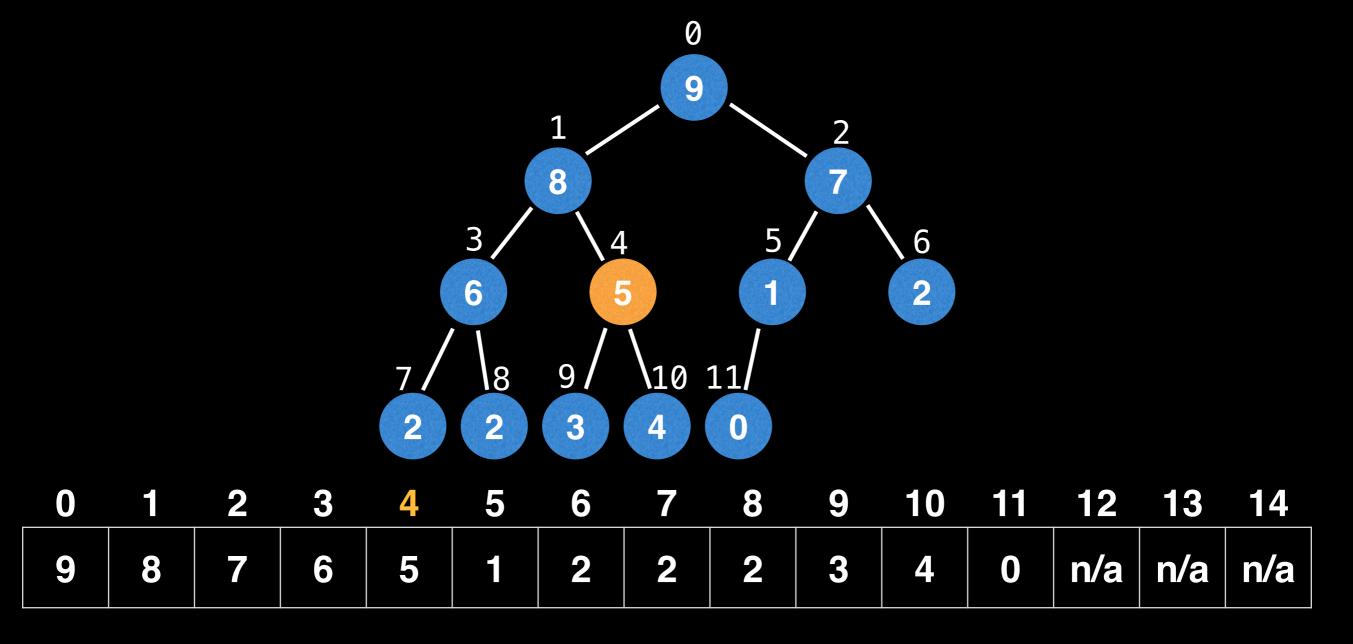
Recall that a very common way to represent a binary heap is with an array since every node is indexed sequentially.



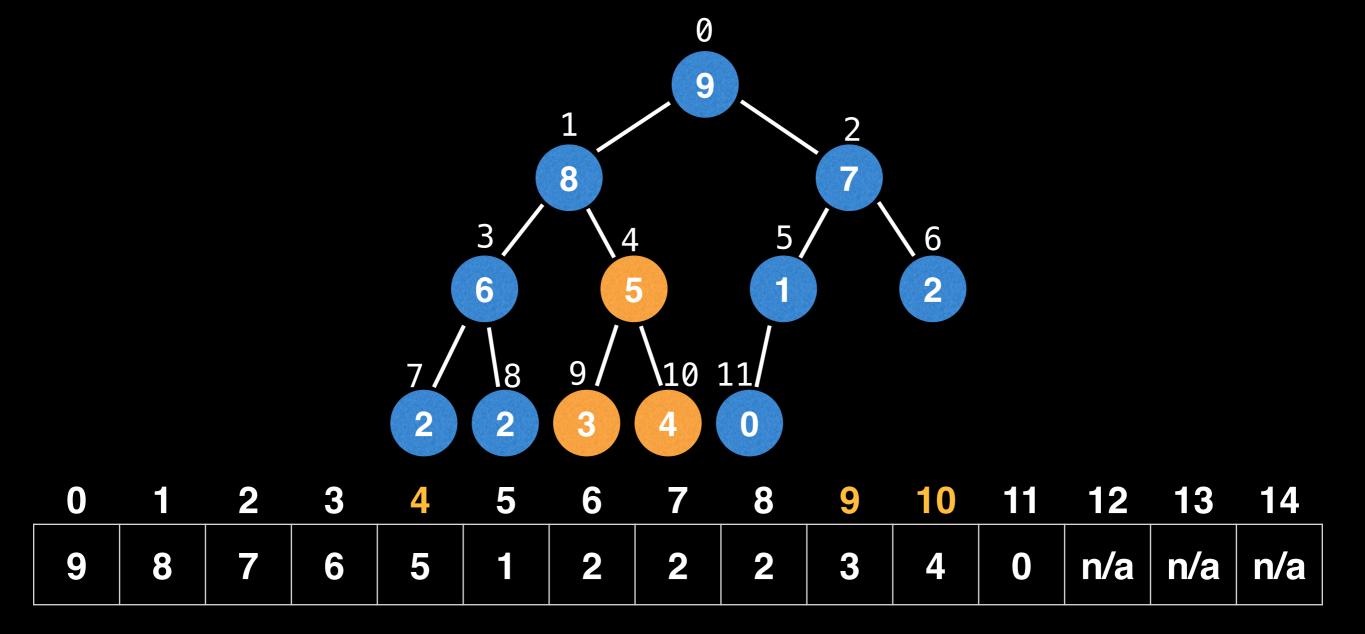


Let i be the current node.

Left child index: 2i + 1
Right child index: 2i + 2
(zero based)



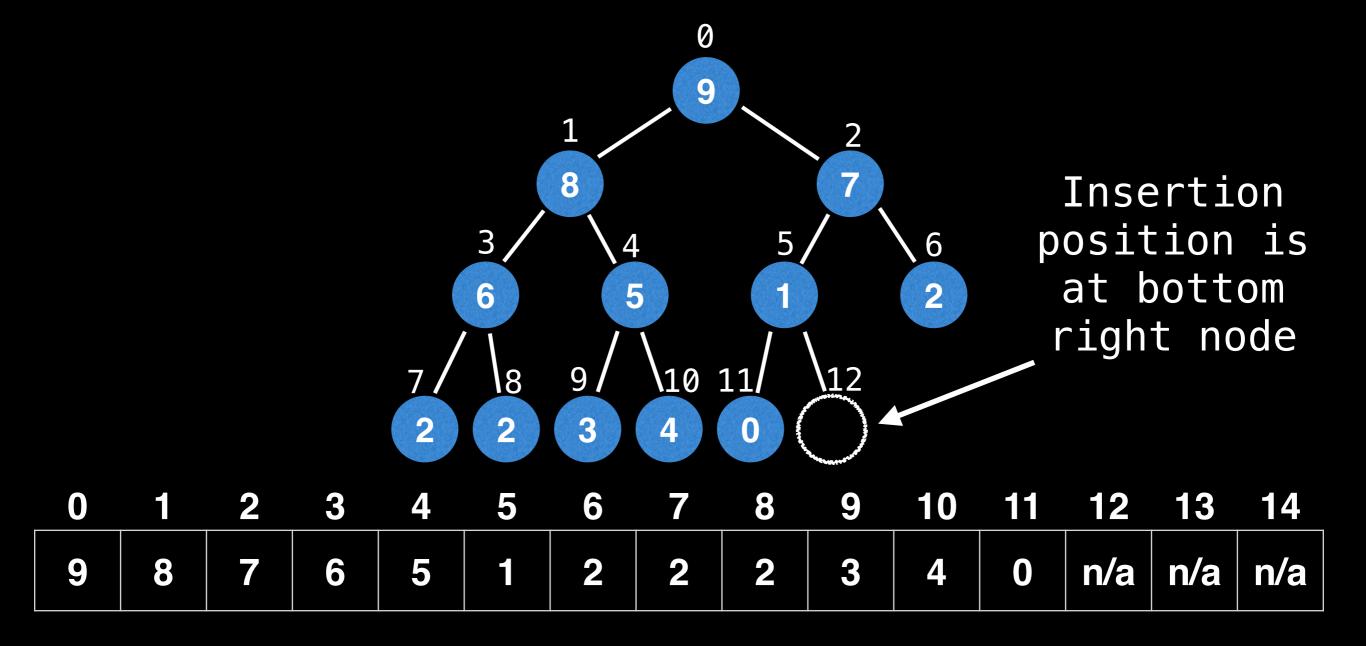
Q: What are the children of the node at index 4?



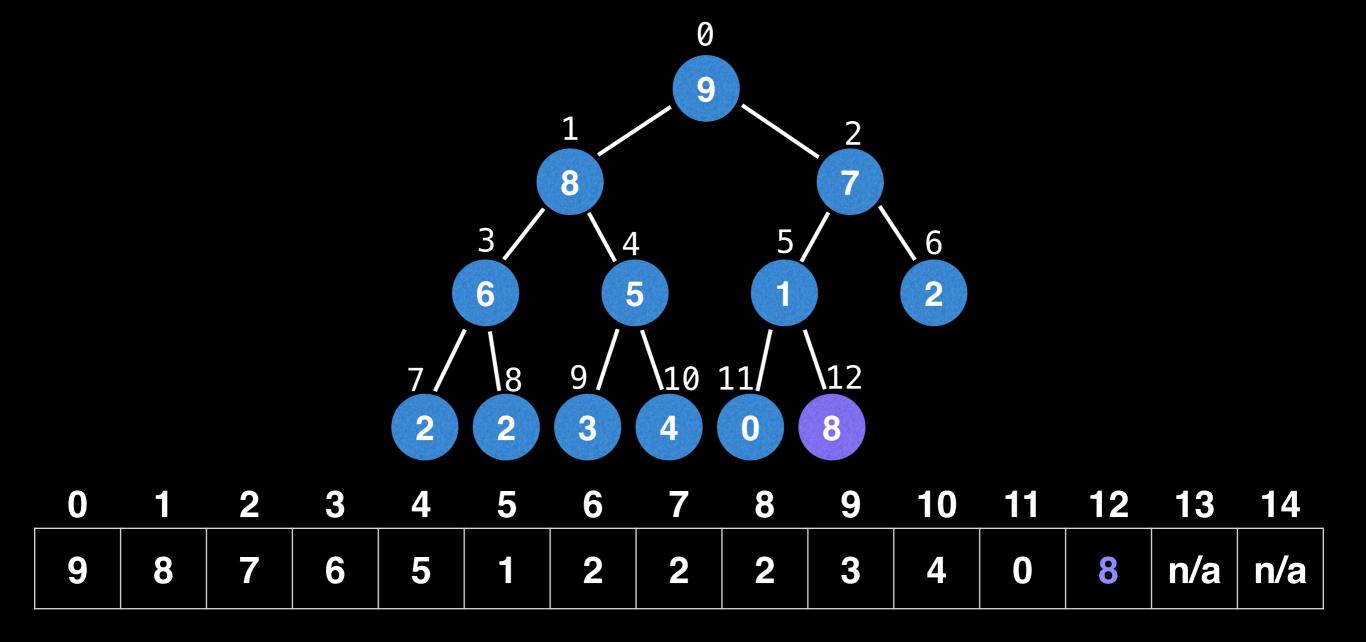
Q: What are the children of the node at index 4?

Left child =
$$2*4 + 1 = 9$$

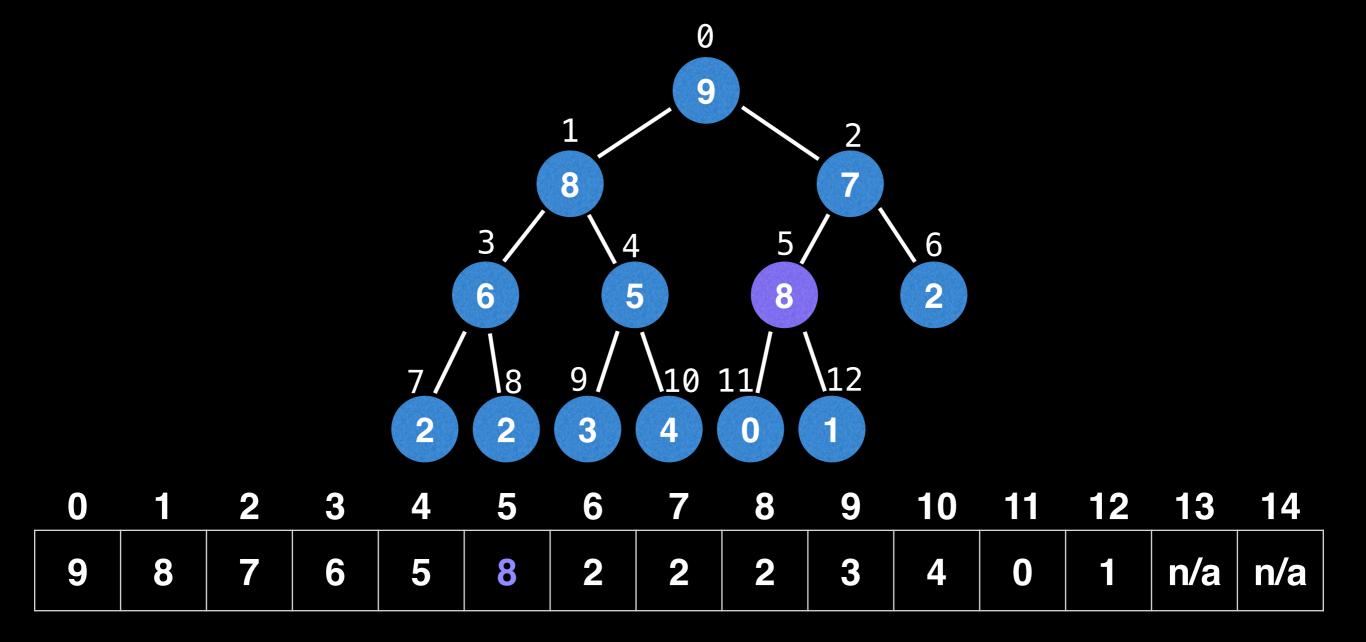
Right child = $2*4 + 2 = 10$



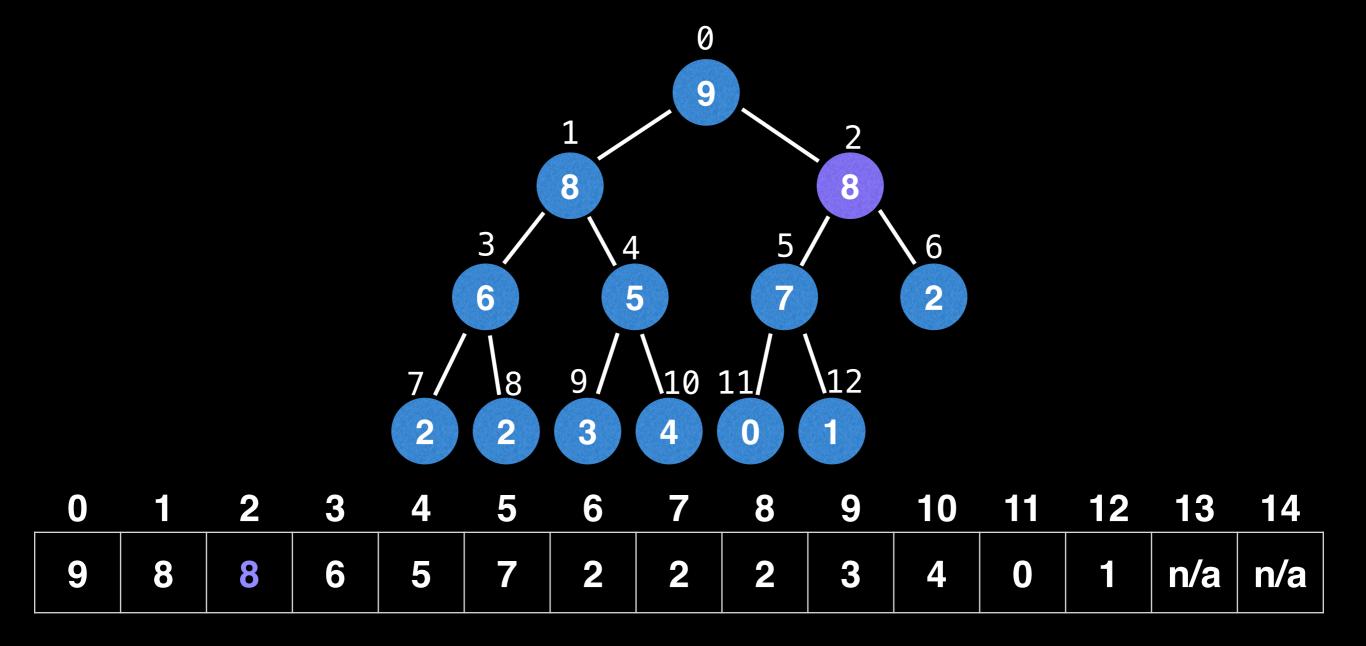
Place values you want to insert into the PQ at the insertion position at the bottom right of the binary tree. Doing this ensures a complete tree structure is always maintained.



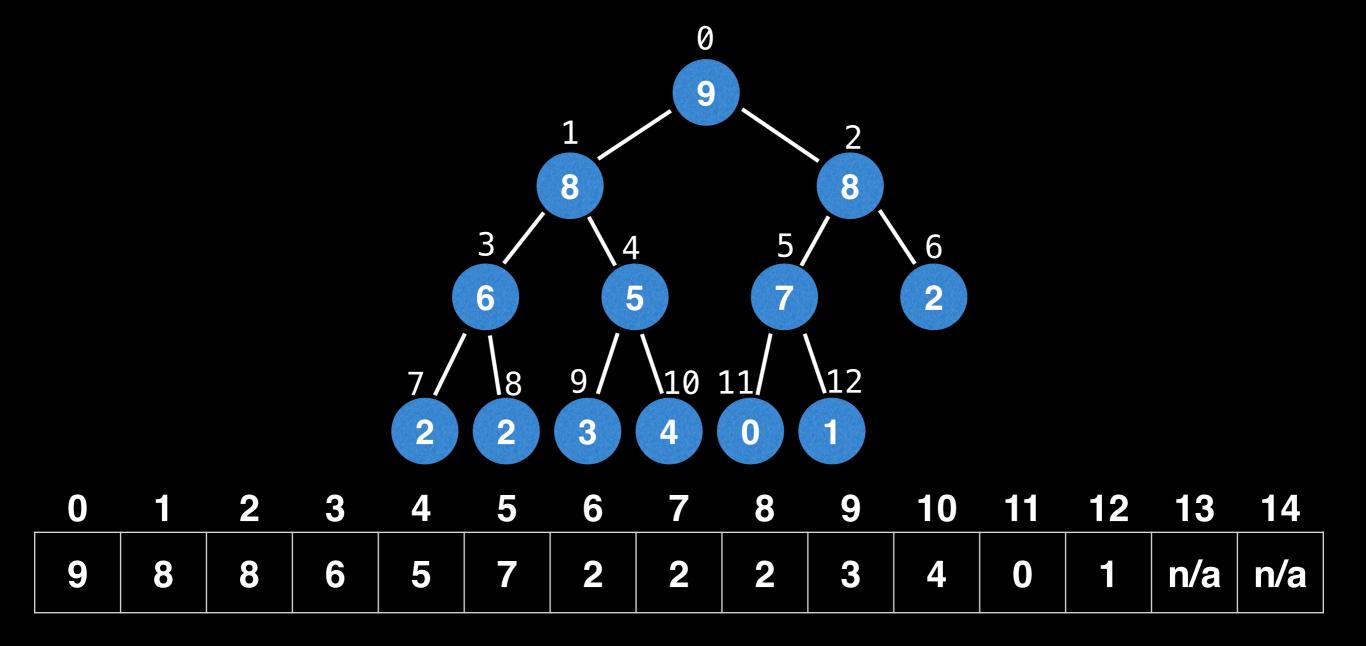
Suppose we insert the value 8. This would violate the heap invariant, so we bubble/sift up the value until the invariant is met.



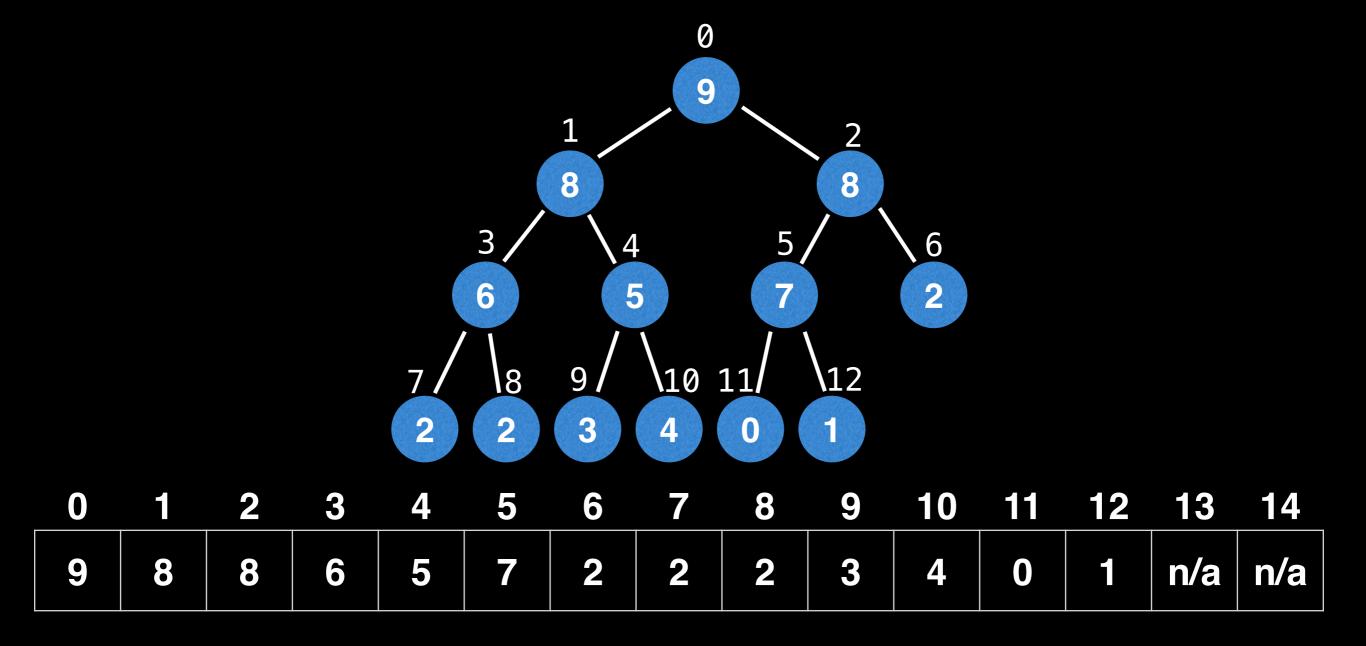
Suppose we insert the value 8. This would violate the heap invariant, so we bubble/sift up the value until the invariant is met.



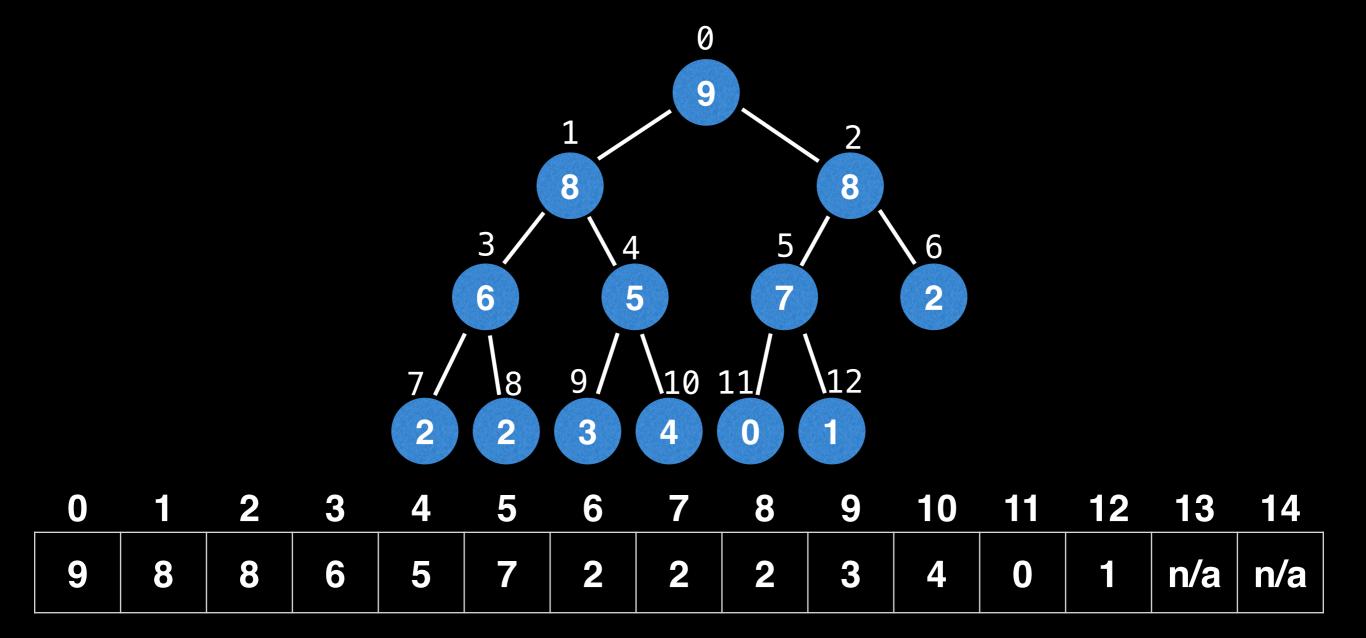
Suppose we insert the value 8. This would violate the heap invariant, so we bubble/sift up the value until the invariant is met.



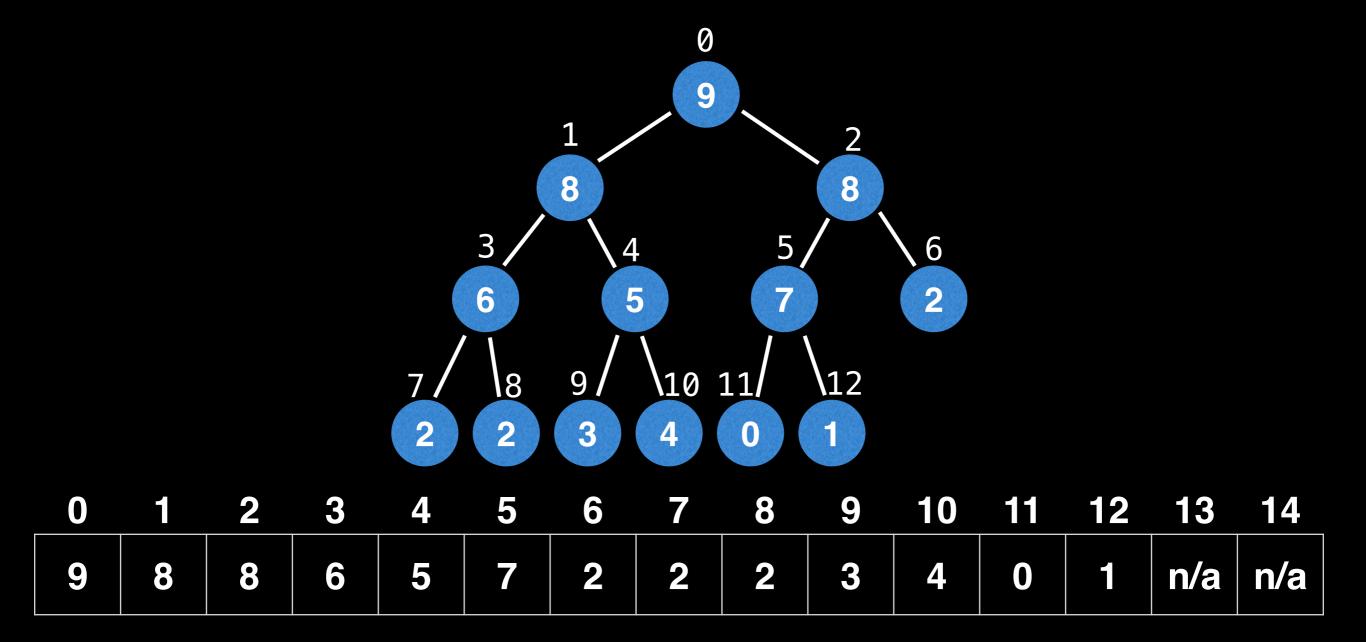
Suppose we insert the value 8. This would violate the heap invariant, so we bubble/sift up the value until the invariant is met.

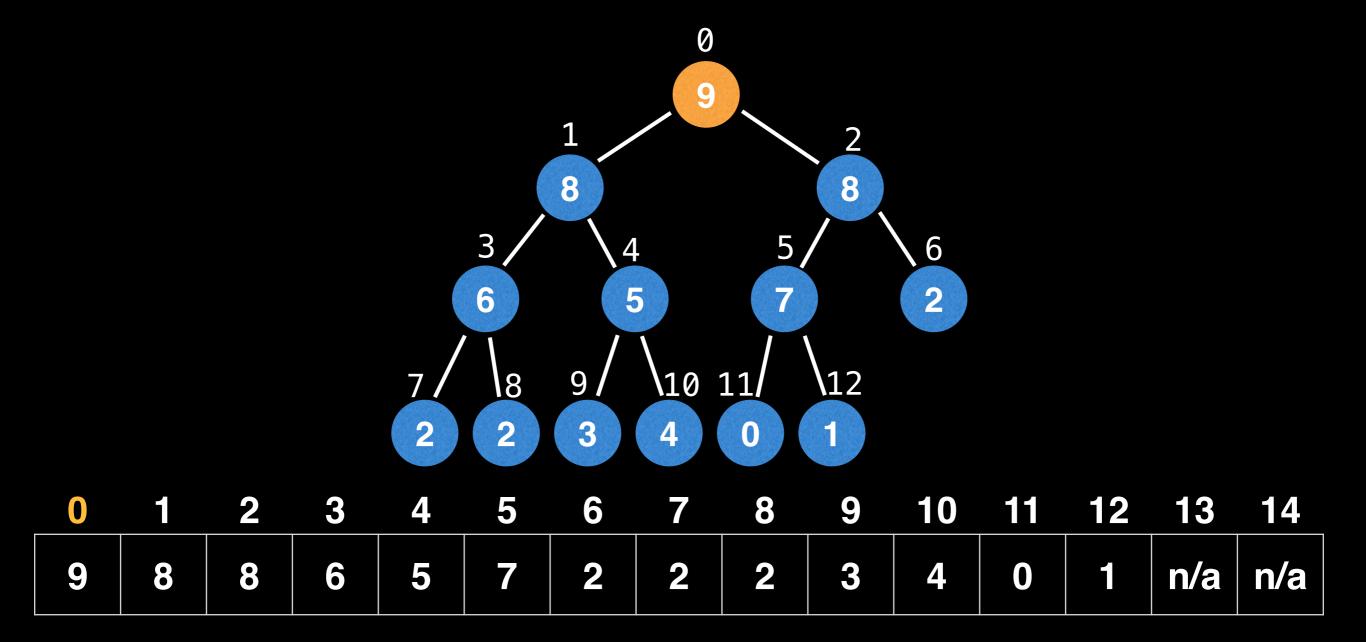


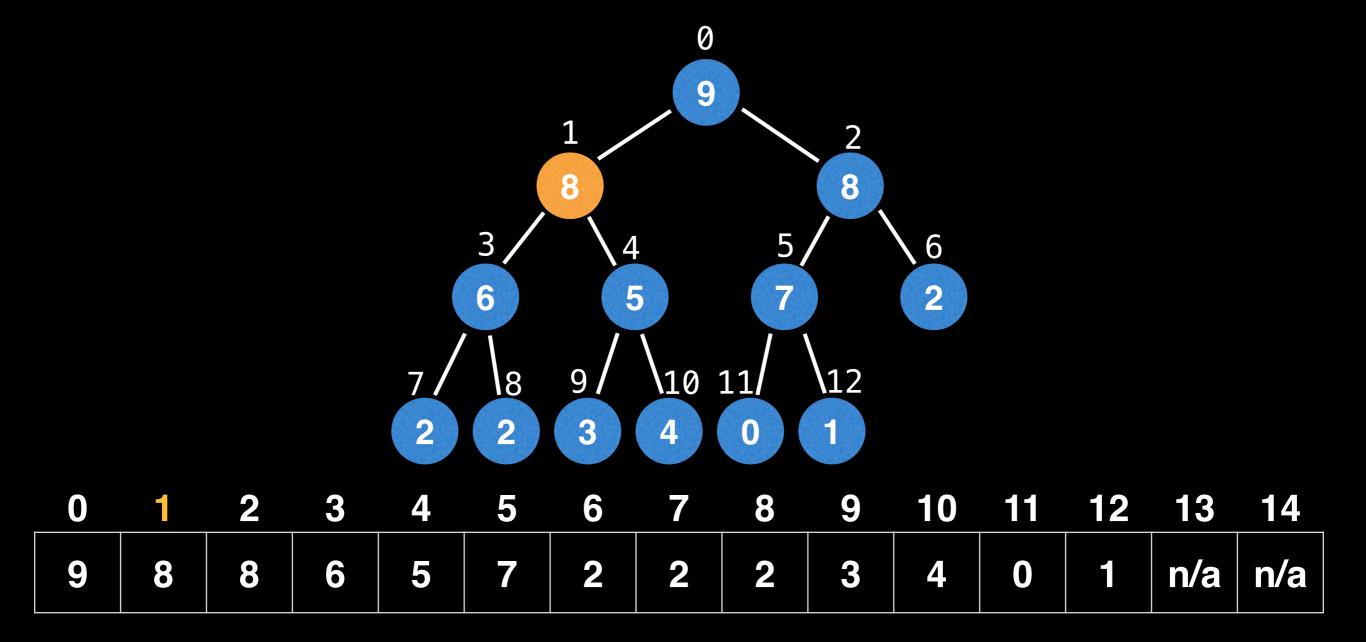
In a traditional PQ, to remove items, search for the element you want to remove and then swap with last node, perform removal and finally bubble up or down the swapped value.

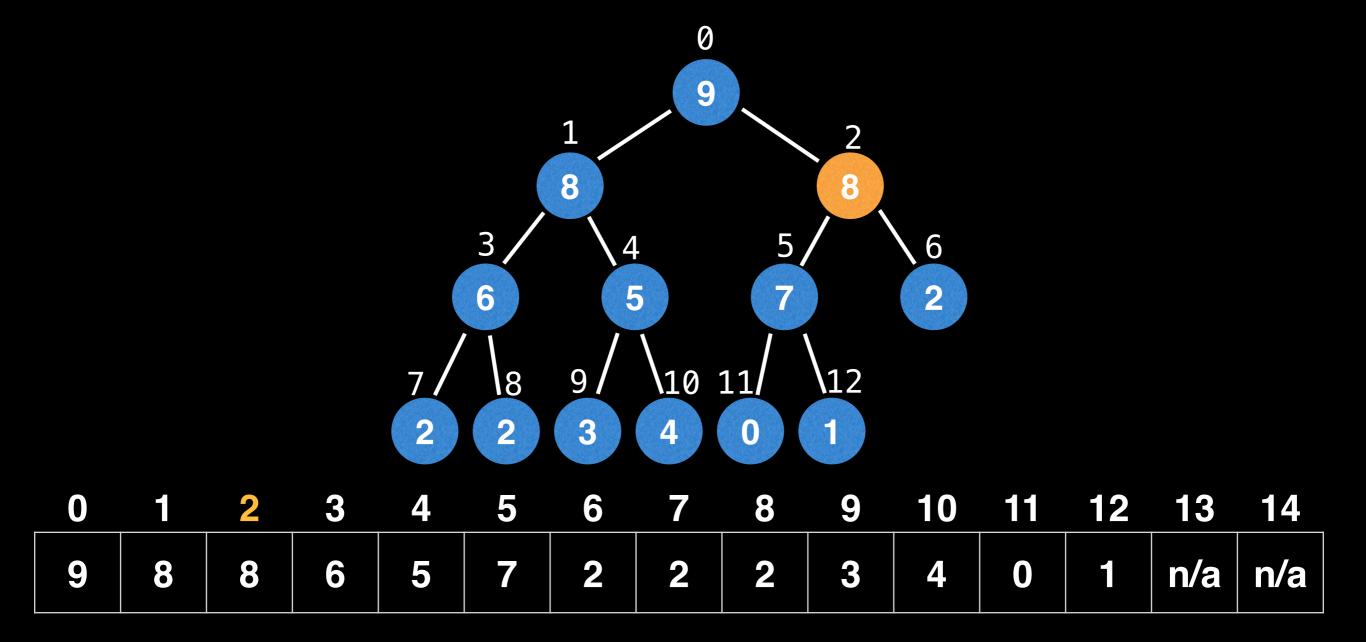


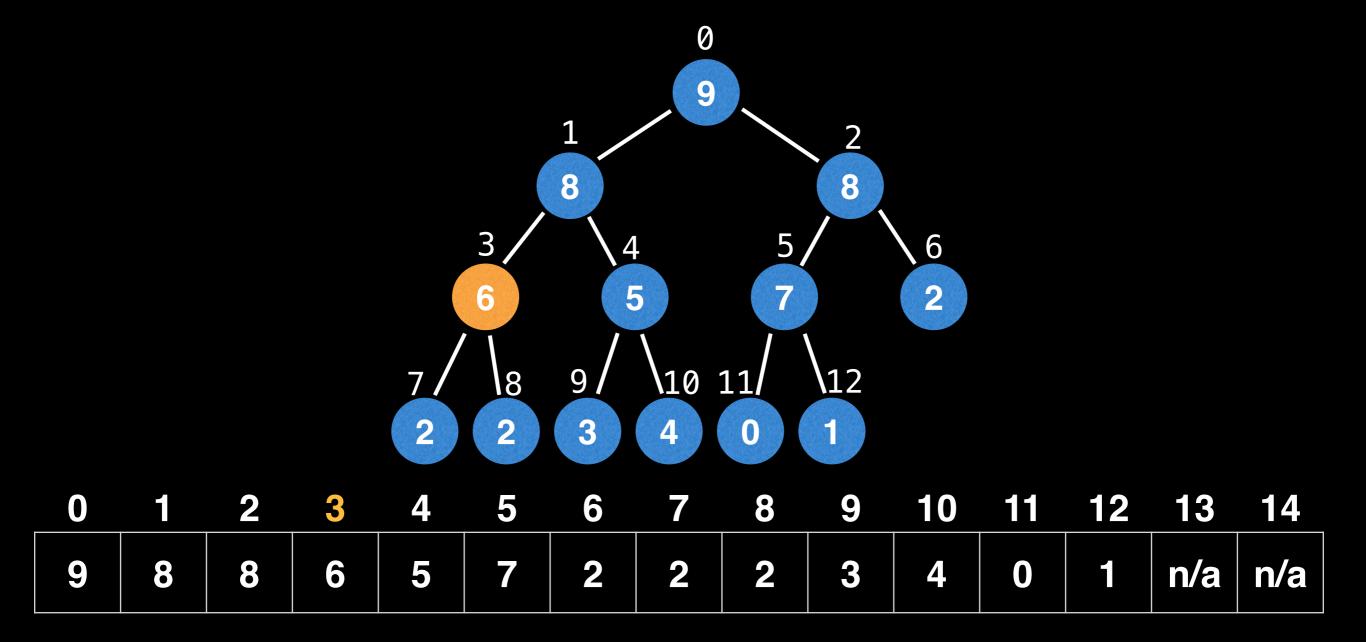
Remove node with value 5.



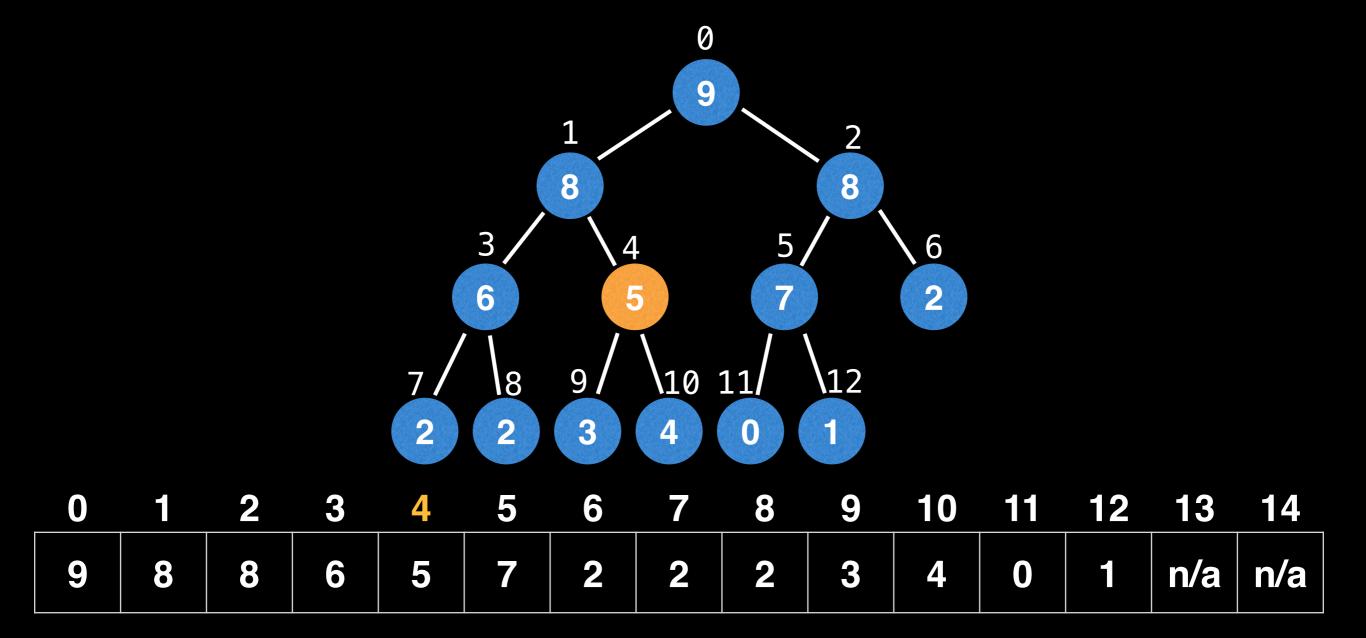




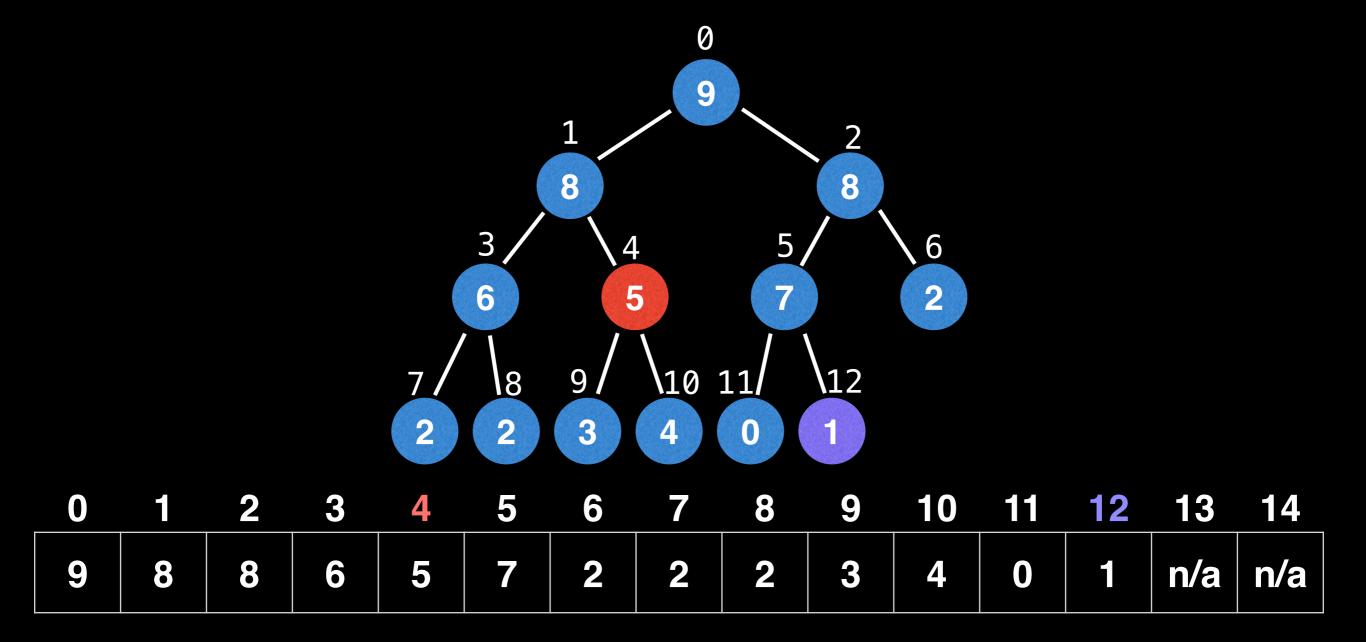




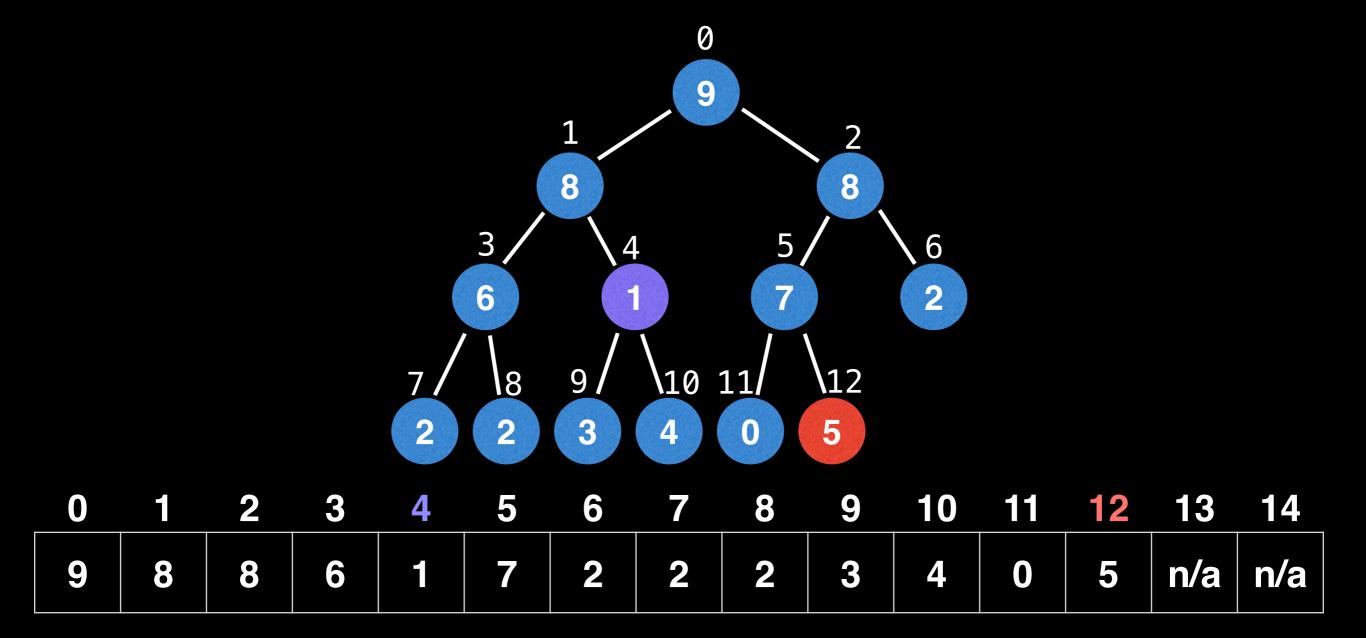
Search for node with value 5.



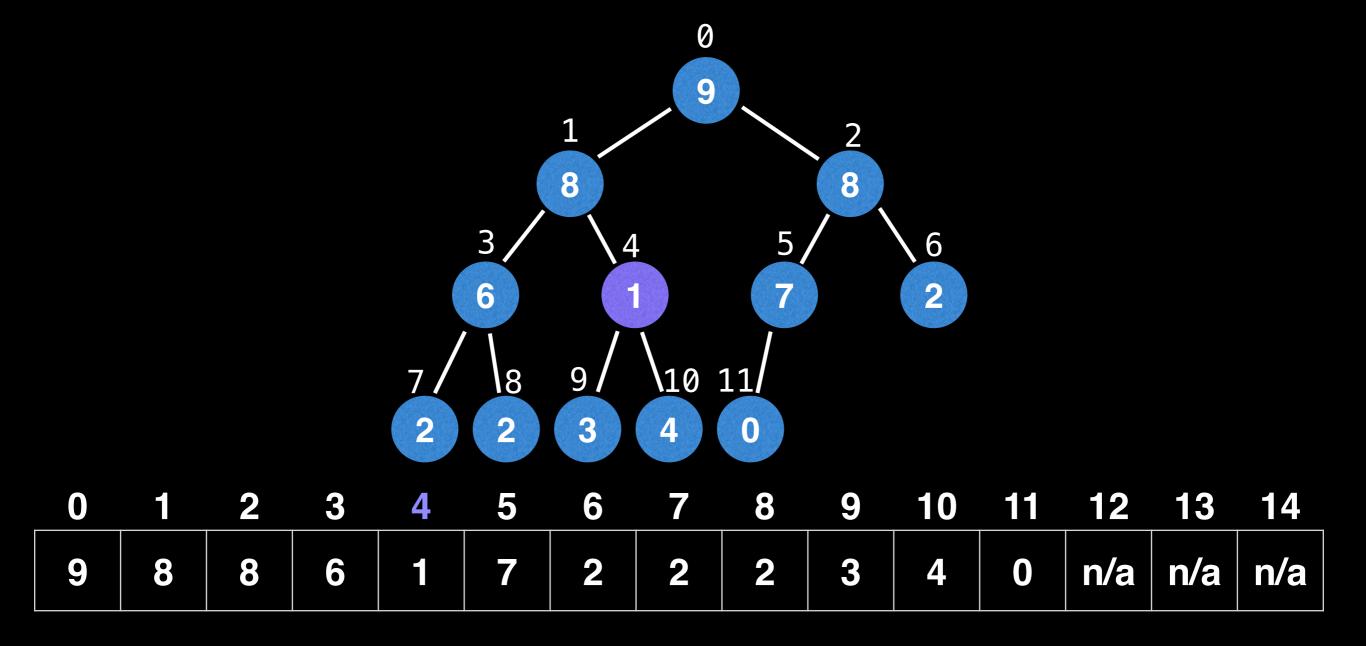
Search for node with value 5.



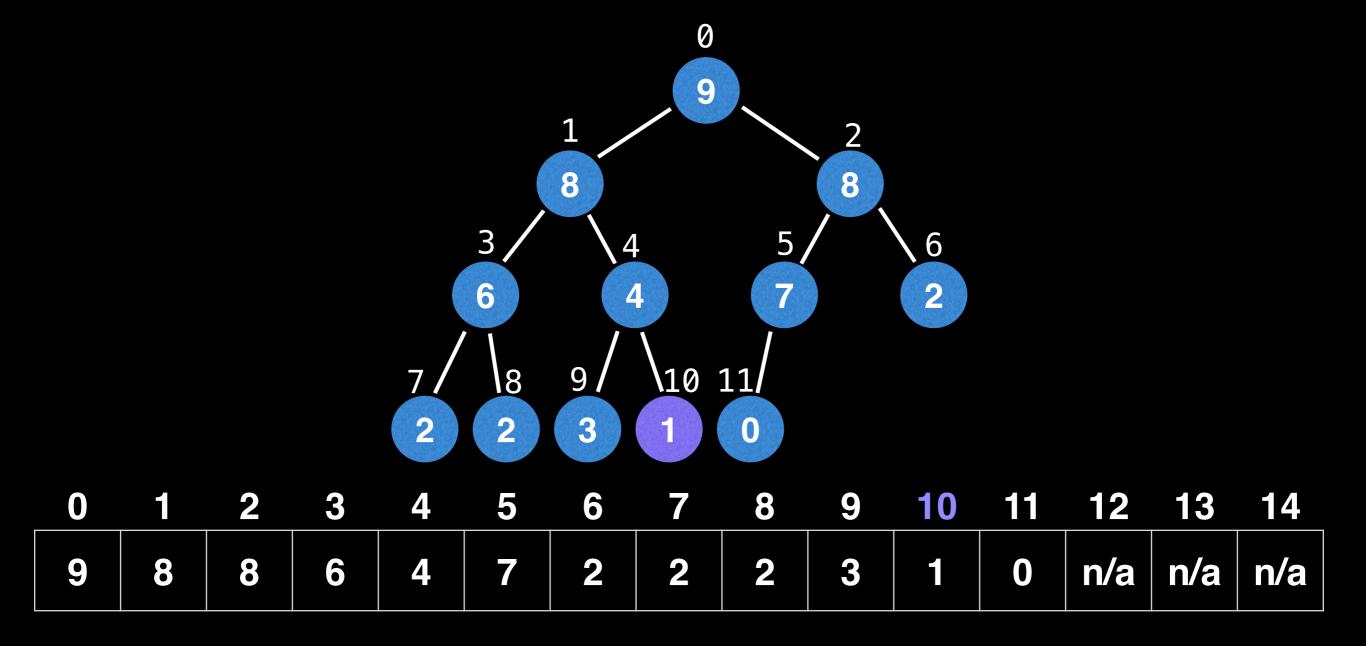
Swap node 5 with rightmost bottom node.



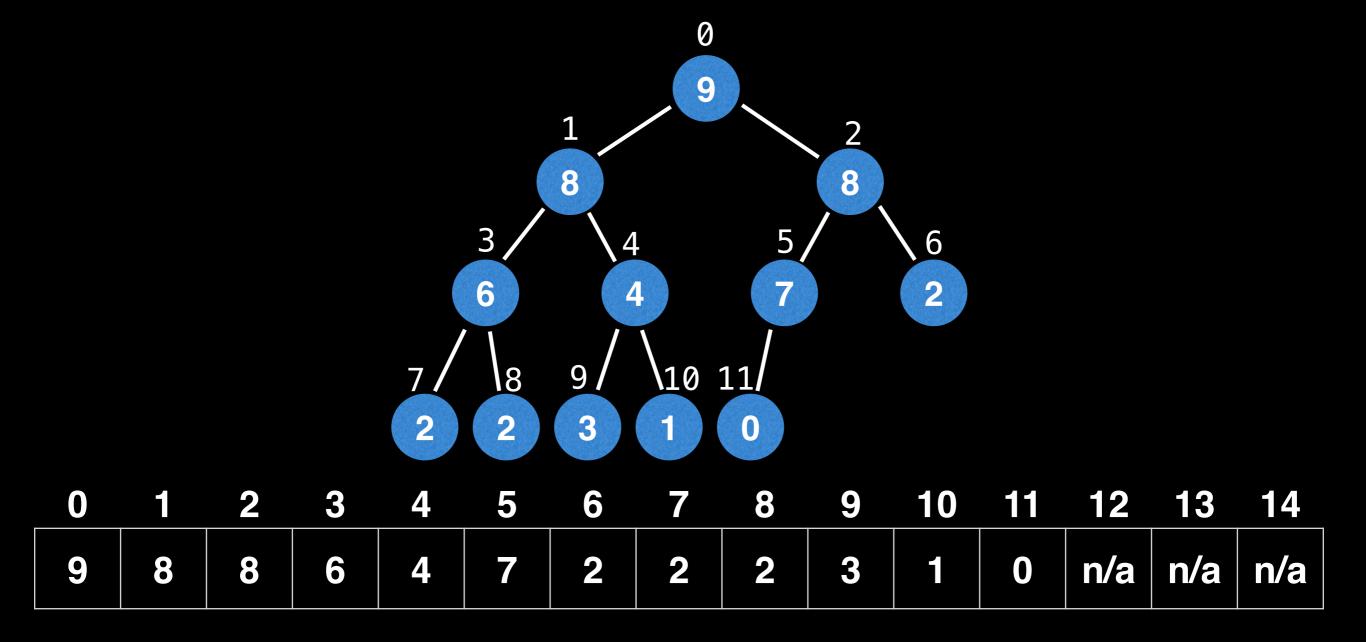
Remove node 5 from tree.



Now the purple node we swapped may not satisfy the heap invariant so we need to either move it up or down the tree.



Now the purple node we swapped may not satisfy the heap invariant so we need to either move it up or down the tree.



Now the purple node we swapped may not satisfy the heap invariant so we need to either move it up or down the tree.

Suppose we have N people with different priorities we need to serve. Assume priorities can dynamically change and we always want to serve the person with the lowest priority.

Name	ki	Value
Anna		
Bella		
Carly		
Dylan		
Emily		
Fred		
George		
Henry		
Isaac		
James		
Kelly		
Laura		

To figure out who to serve next use a Min IPQ to sort by lowest value first.

Name	ki	Value
Anna		
Bella		
Carly		
Dylan		
Emily		
Fred		
George		
Henry		
Isaac		
James		
Kelly		
Laura		

To figure out who to serve next use a Min IPQ to sort by lowest value first.

Arbitrarily assign each person a unique index value between [0, N)

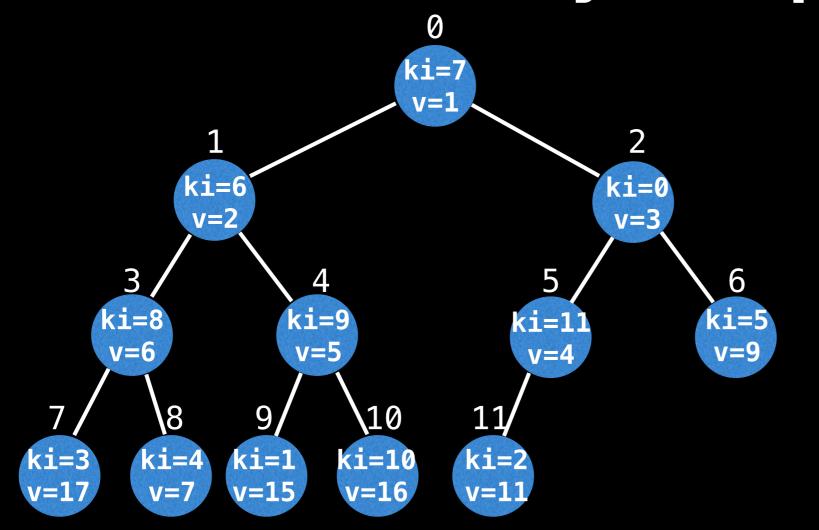
Name	ki	Value
Anna	0	
Bella	1	
Carly	2	
Dylan	3	
Emily	4	
Fred	5	
George	6	
Henry	7	
Isaac	8	
James	9	
Kelly	10	
Laura	11	

Key Index

To figure out who to serve next use a Min IPQ to sort by lowest value first.

Name	ki	Value
Anna	0	3
Bella	1	15
Carly	2	11
Dylan	3	17
Emily	4	7
Fred	5	9
George	6	2
Henry	7	1
Isaac	8	6
James	9	5
Kelly	10	16
Laura	11	4

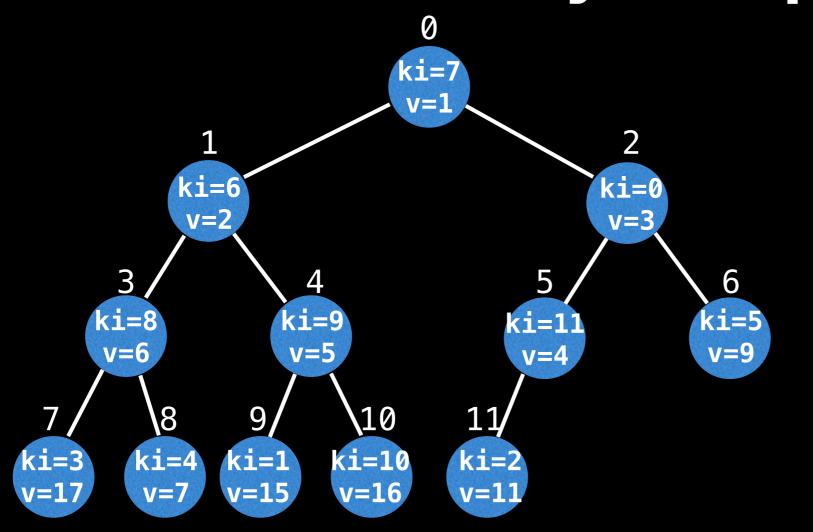
Initial values to place inside IPQ. These will be maintained by the IPQ once inserted. Note that values can be any comparable value not only integers.



Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11

When we insert (ki, v) pairs into an IPQ we sort by the value associated with each key.

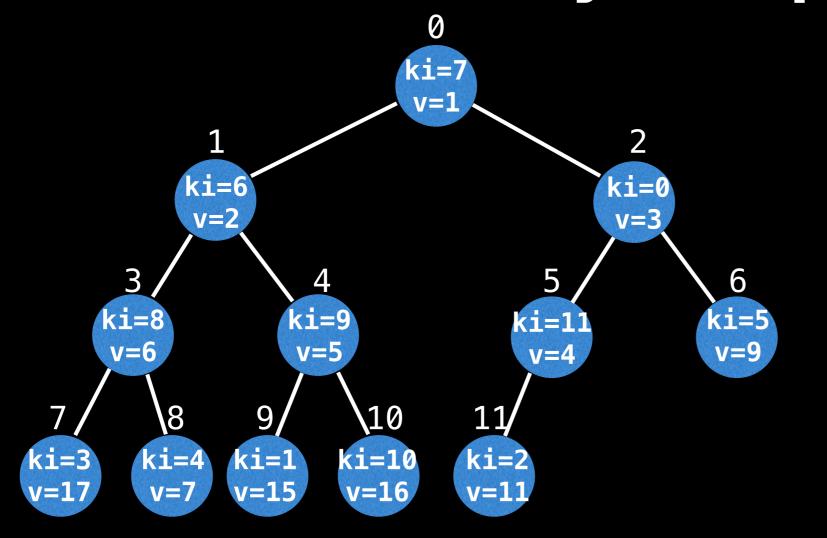
In the heap above we are sorting by smallest value since we're working with a min heap.



Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
vals	3	15	11	17	7	9	2	1	6	5	16	4	-1	-1	-1

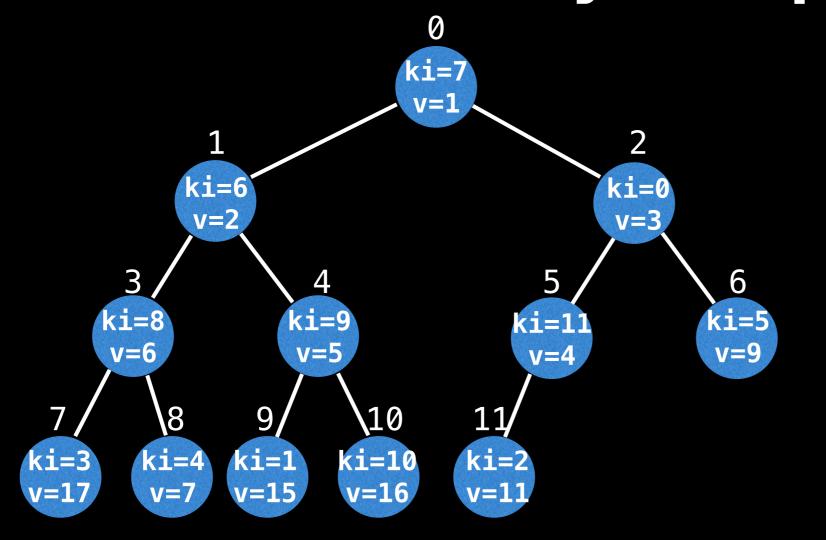
To access the value for any given key k, find its key index (ki) and do a lookup in the vals array maintained by the IPQ.



Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
vals	3	15	11	17	7	9	2	1	6	5	16	4	-1	-1	-1

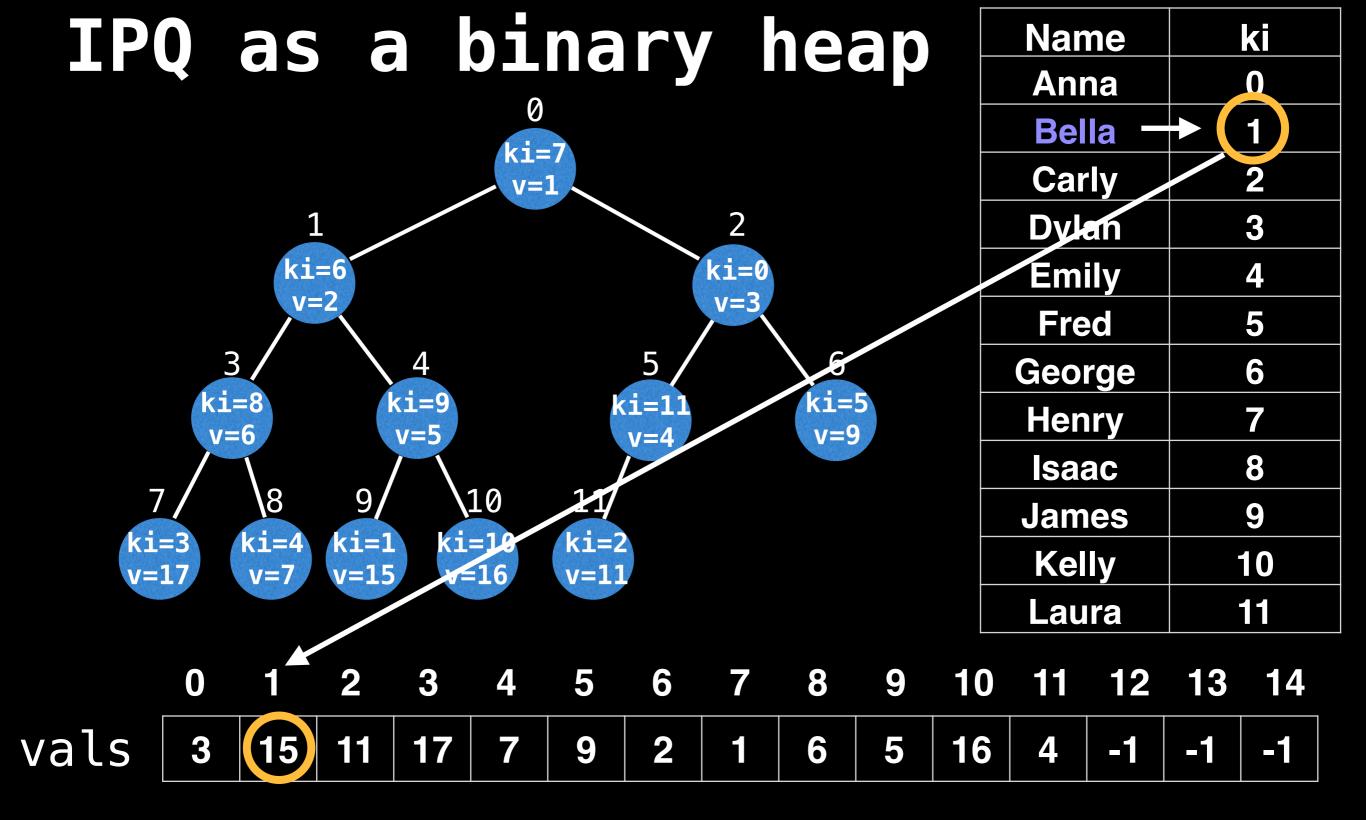
Q: What value does "Bella" have in the IPQ?



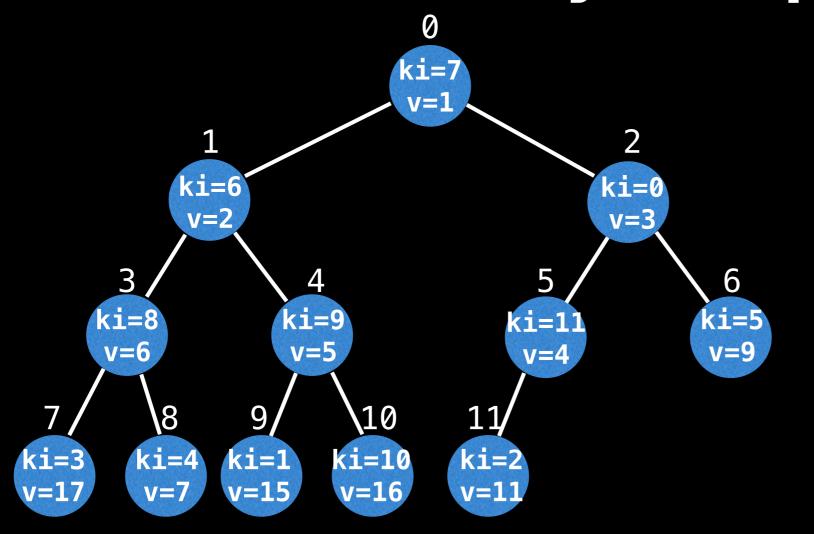
Name	ki
Anna	0
Bella -	(1)
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
vals	3	15	11	17	7	9	2	1	6	5	16	4	-1	-1	-1

Q: What value does "Bella" have in the IPQ?



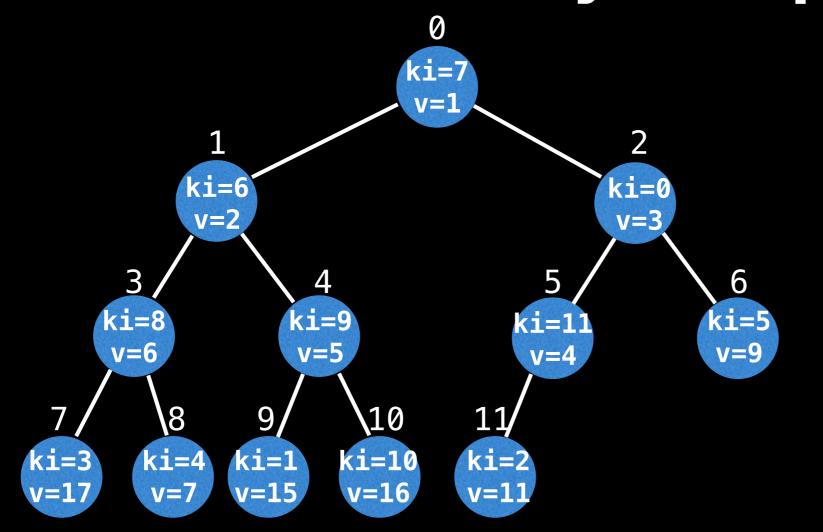
ki = 1, so "Bella" has a value of vals[ki] = 15



Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11

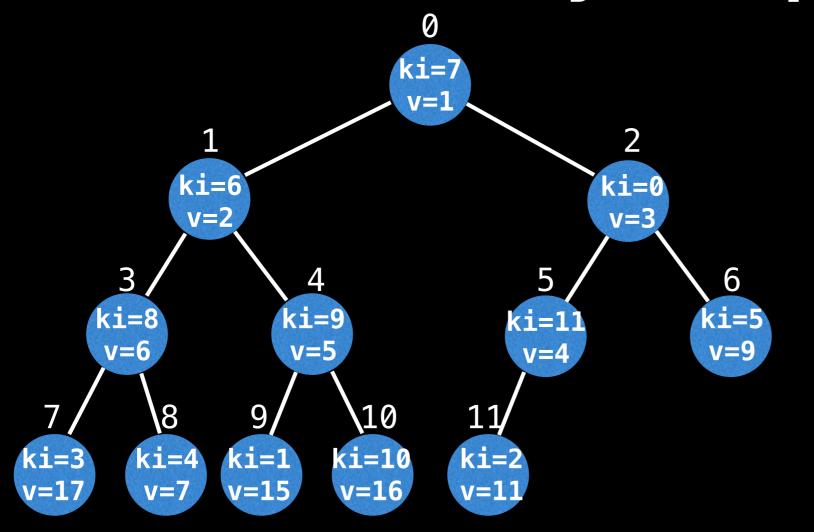
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
vals	3	15	11	17	7	9	2	1	6	5	16	4	-1	-1	-1

ki = 1, so "Bella" has a value of vals[ki] = 15



Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11

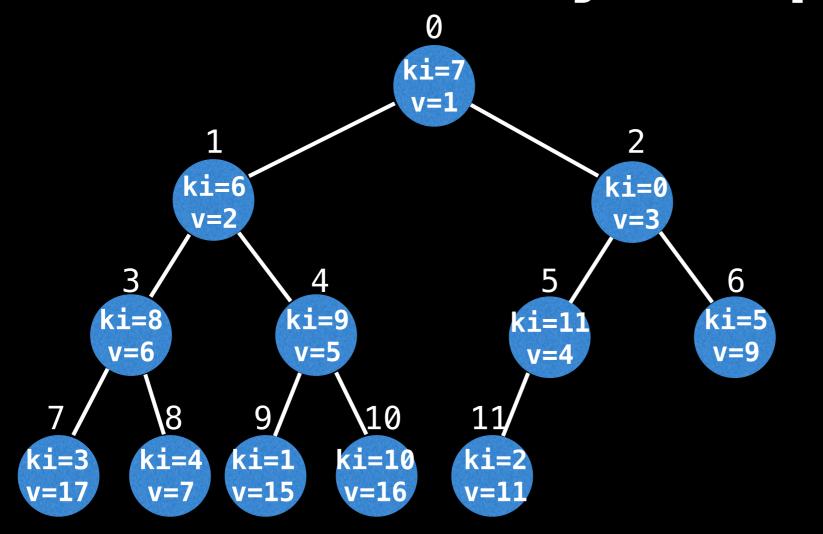
Q: How do I find the index of the node for a particular key?



Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
vals	3	15	11	17	7	9	2	1	6	5	16	4	-1	-1	-1
pm	2	9	11	7	8	6	1	0	3	4	10	5	-1	-1	-1

The array pm is a Position Map we maintain to tell us the index of the node in the heap for a given key index (ki).

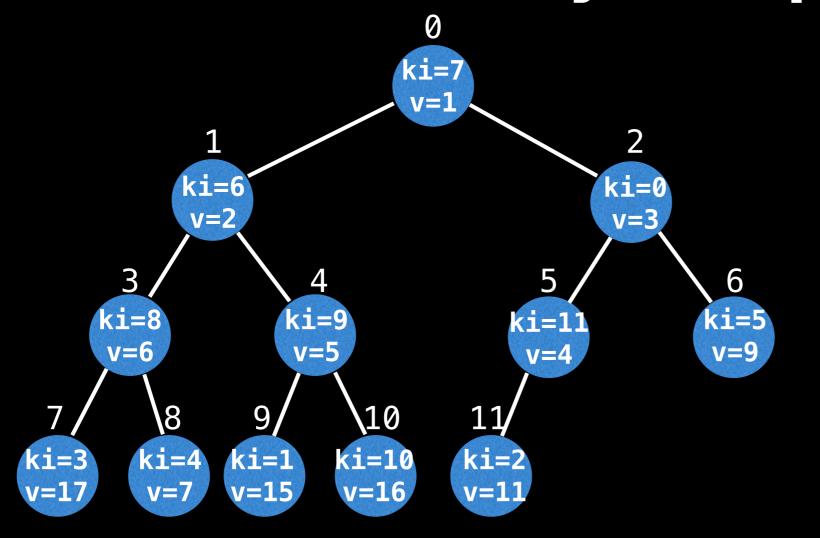


Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11

V	a	J	S	
	D	m		

als	3	15	11	17	7	9	2	1	6	5	16	4	-1	-1	-1
m	2	9	11	7	8	6	1	0	3	4	10	5	-1	-1	-1

Q: Which node represents the key "Dylan"?

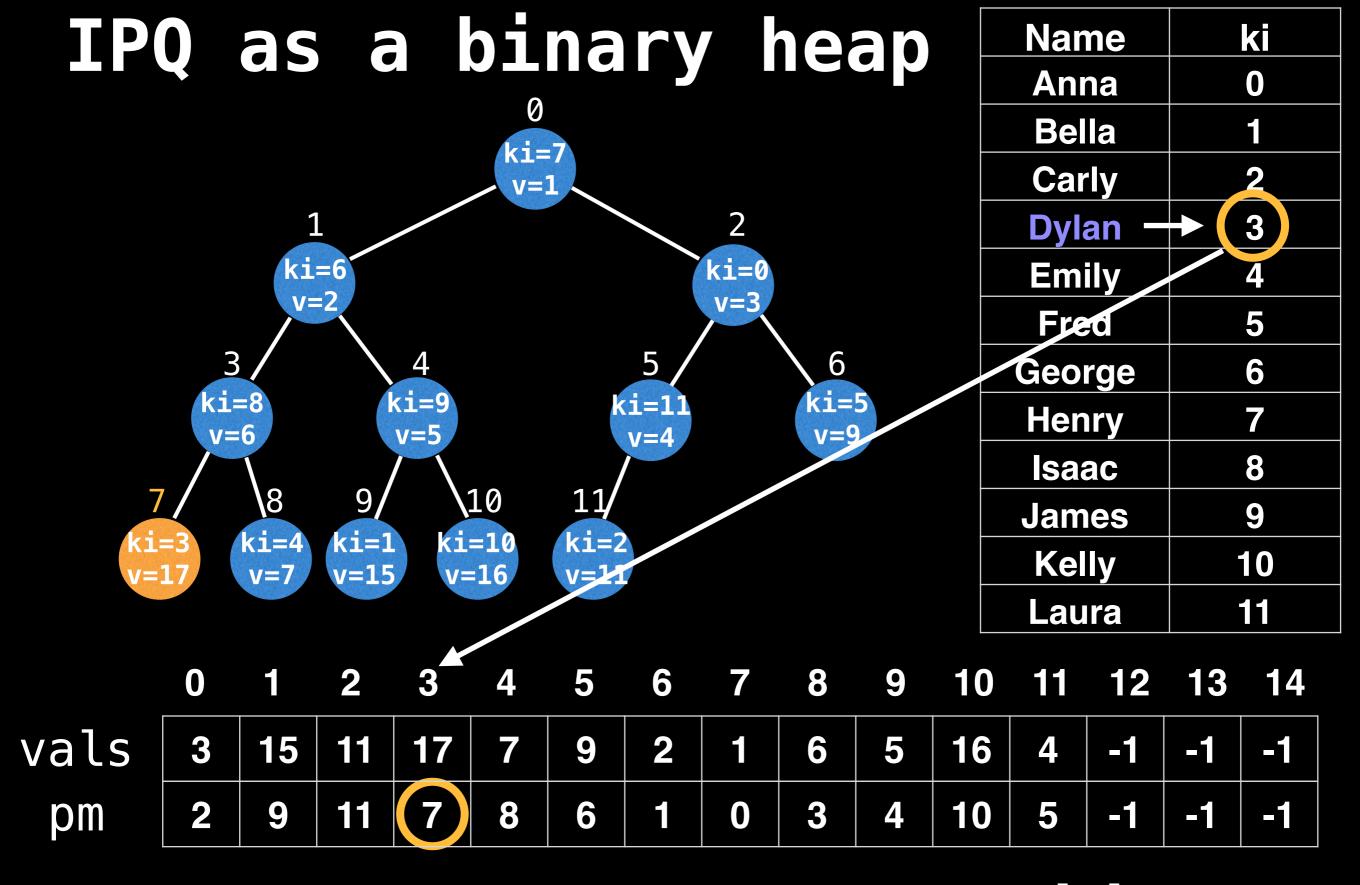


Name	ki
Anna	0
Bella	1
Carly	2
Dylan -	→ (3)
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11

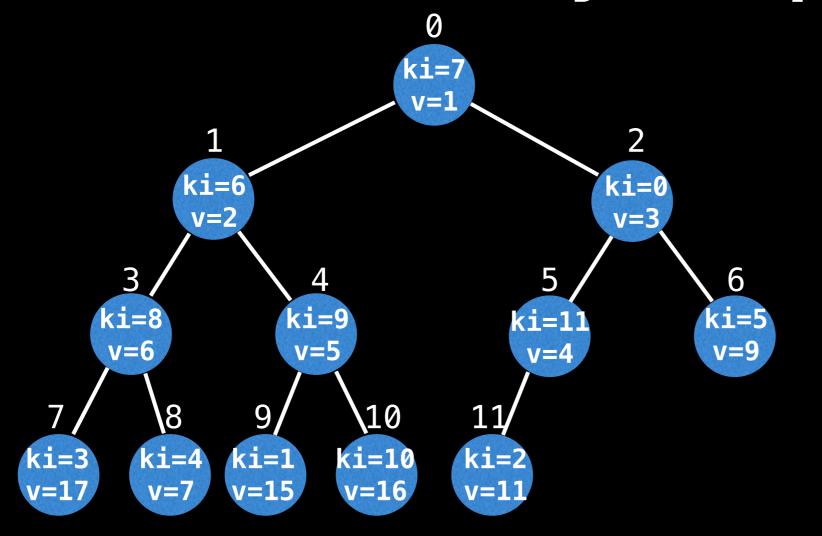
va	ls
р	m

3	15	11	17	7	9	2	1	6	5	16	4	-1	-1	-1
2	9	11	7	8	6	1	0	3	4	10	5	-1	-1	-1

First find the key-index for "Dylan"



"Dylan" has a key index of 3, so pm[3] = 7 is where Dylan is in the heap.

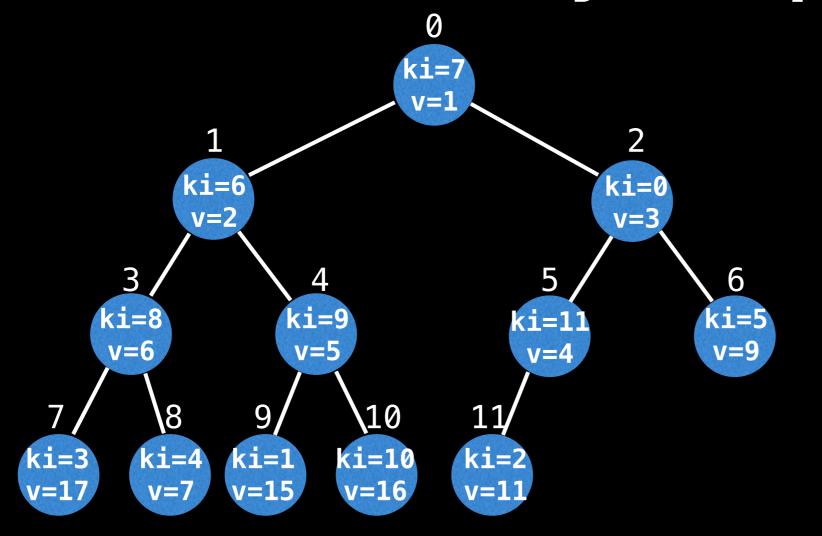


Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11

va	ls
р	m

als	3	15	11	17	7	9	2	1	6	5	16	4	-1	-1	-1
pm	2	9	11	7	8	6	1	0	3	4	10	5	-1	-1	-1

Q: Where is "George" in the heap?

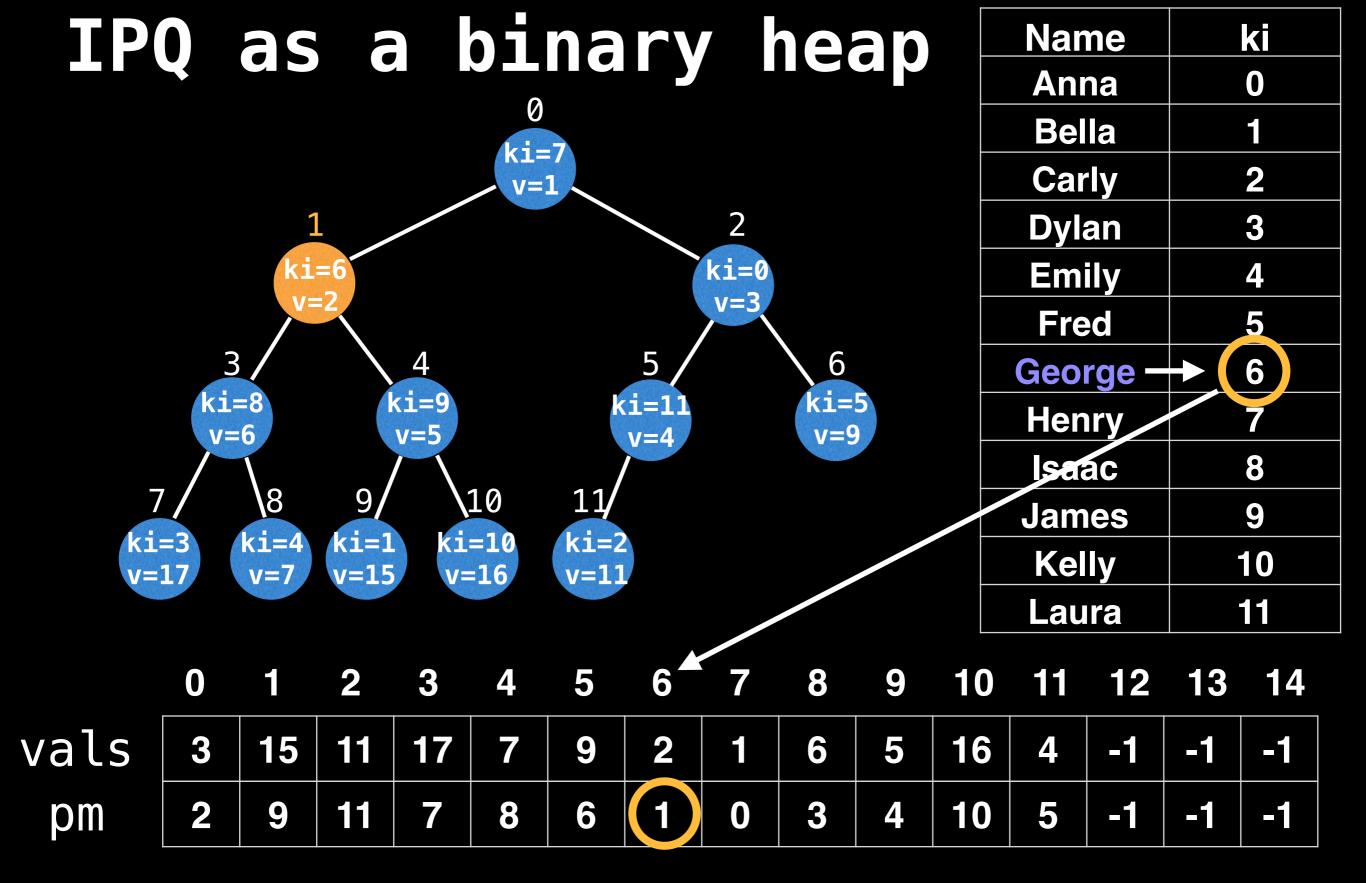


Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George -	→ (6)
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11

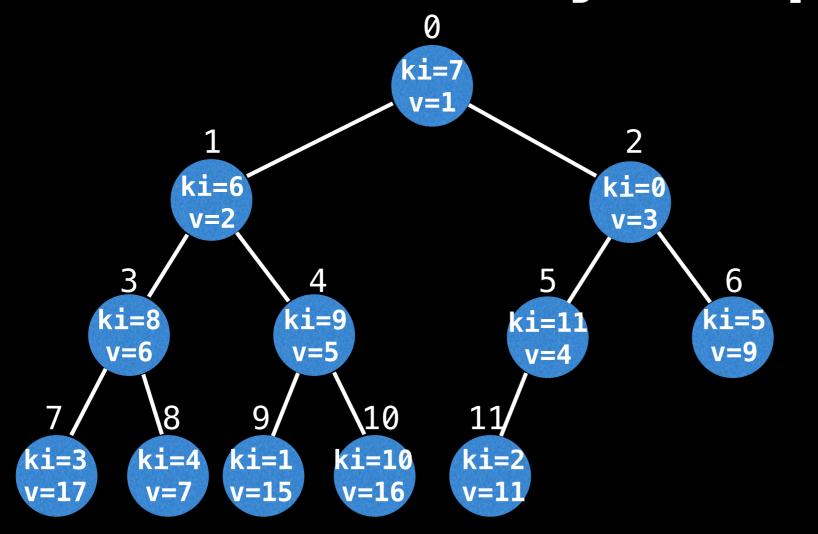
va	ls
p	m

LS	3	15	11	17	7	9	2	1	6	5	16	4	-1	-1	-1
۱ [2	9	11	7	8	6	1	0	3	4	10	5	-1	-1	-1

Q: Where is "George" in the heap?



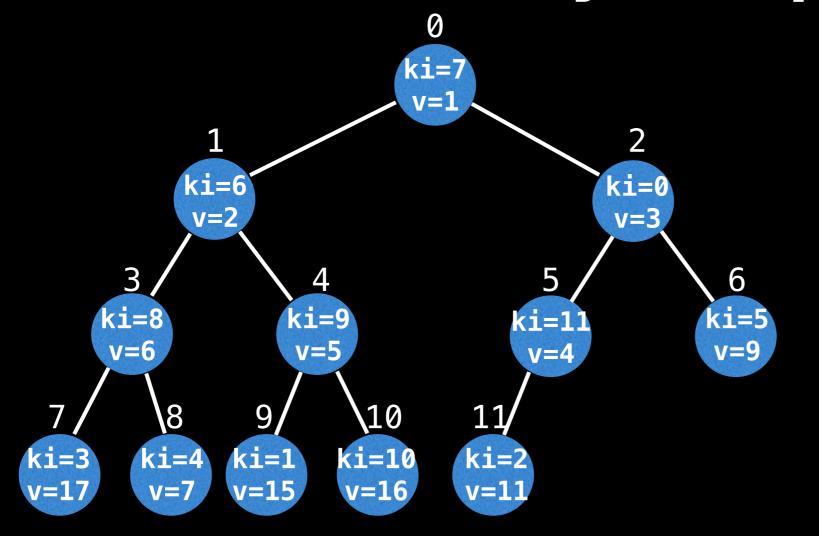
Q: Where is "George" in the heap?



Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11

	U		2	3	4	5	6		8	9	10		12	13	14
vals	3	15	11	17	7	9	2	1	6	5	16	4	-1	-1	-1
pm	2	9	11	7	8	6	1	0	3	4	10	5	-1	-1	-1

Q: How do we go from knowing the position of a node to its key and ki value?



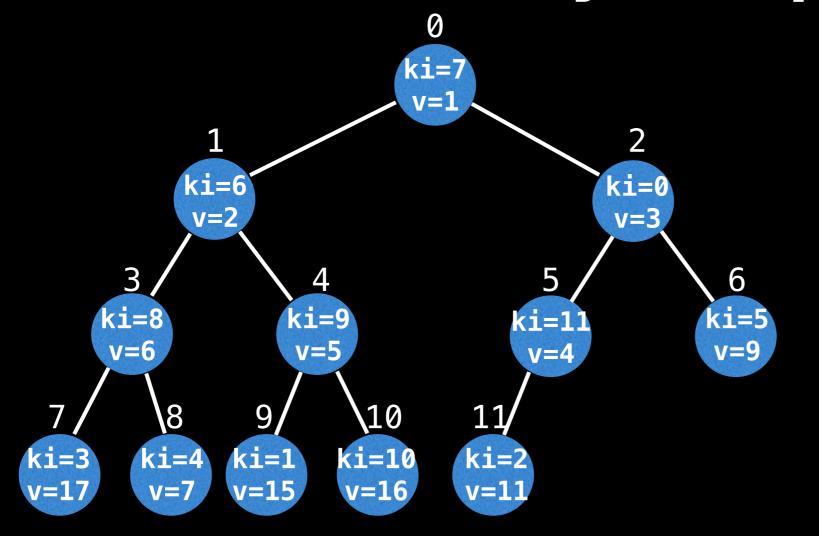
Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11

	O		2	3	4	5	6		8	9	10		12	13	14
valc	2	15	44	17	7	0	9	4	6		16	Л	4	4	1

vals pm im

3	15	11	17	7	9	2	1	6	5	16	4	-1	-1	-1
2	9	11	7	8	6	1	0	3	4	10	5	-1	-1	-1
7	6	0	8	9	11	5	3	4	1	10	2	-1	-1	-1

To do that we also need to maintain an inverse lookup table denoted: im (Inverse Map)



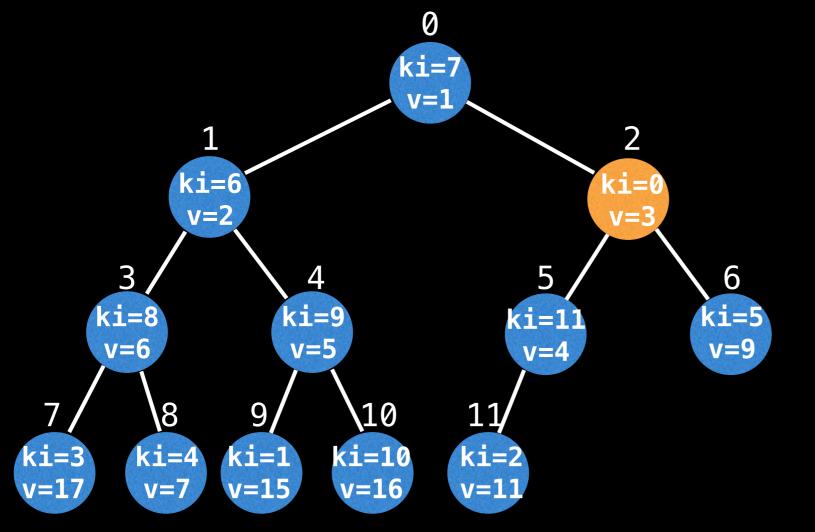
Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11

0	2	3	4	5	6	7	8	9	10	11	12	13	14

vals pm im

3	15	11	17	7	9	2	1	6	5	16	4	-1	-1	1
2	9	11	7	8	6	1	0	3	4	10	5	-1	-1	-1
7	6	0	8	9	11	5	3	4	1	10	2	-1	-1	-1

Q: Which person (key) is represented in the node at index 2?



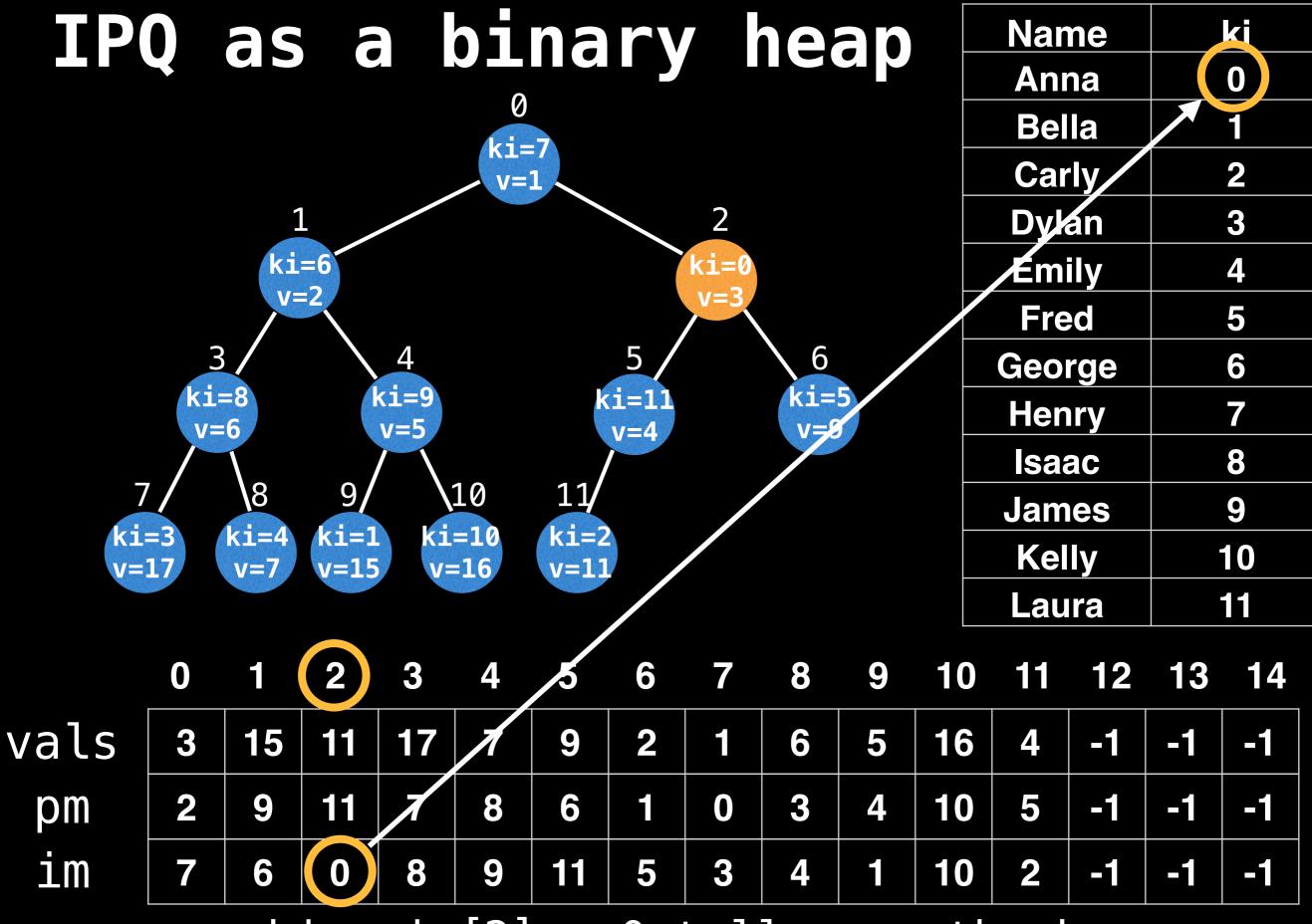
Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11

	1	2)	3	4	5	6	7	8	9	10	11	12	13	14

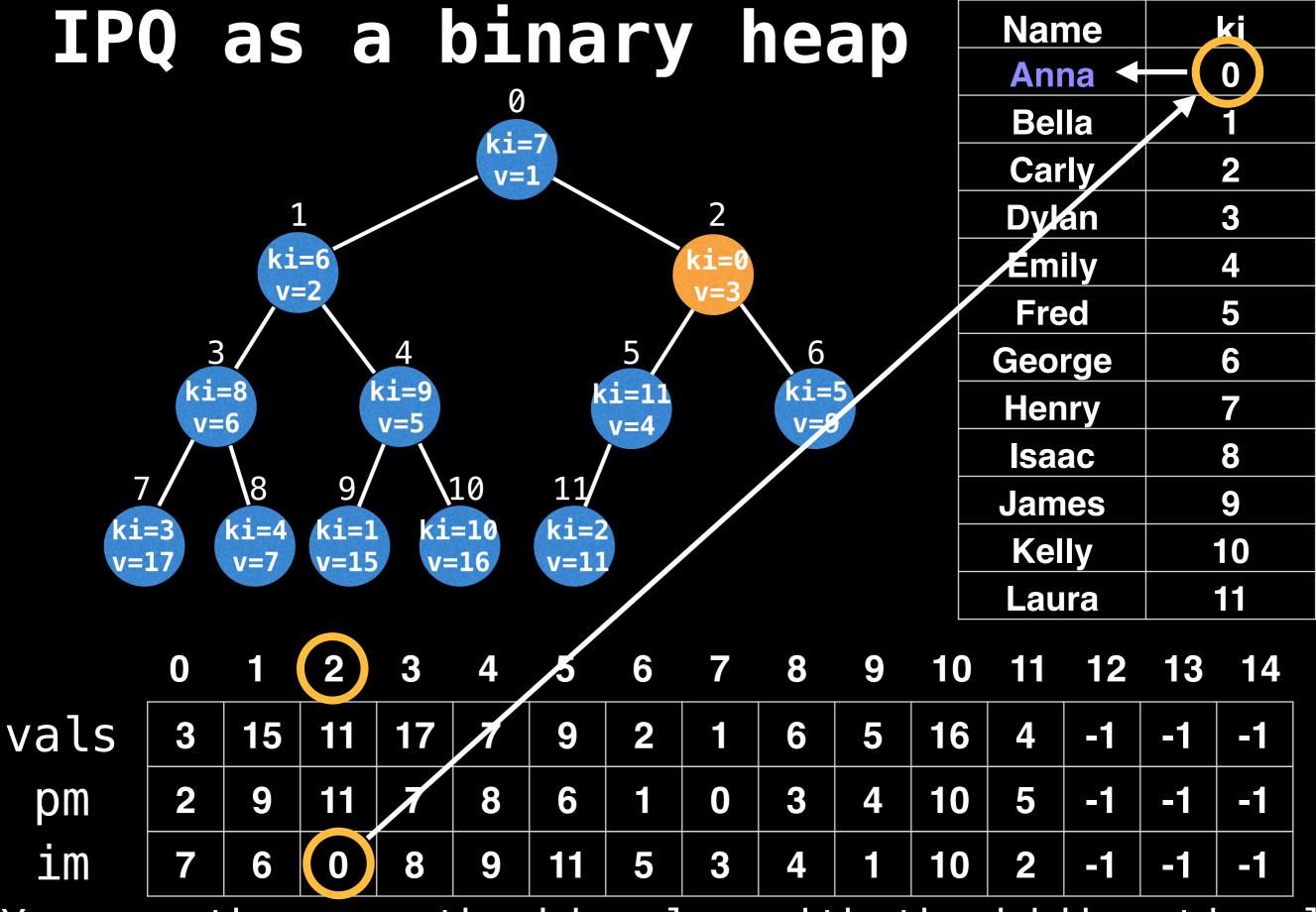
vals pm im

3	15	11	17	7	9	2	1	6	5	16	4	-1	-1	-1
2	9	11	7	8	6	1	0	3	4	10	5	-1	-1	-1
7	6	0	8	9	11	5	3	4	1	10	2	-1	-1	-1

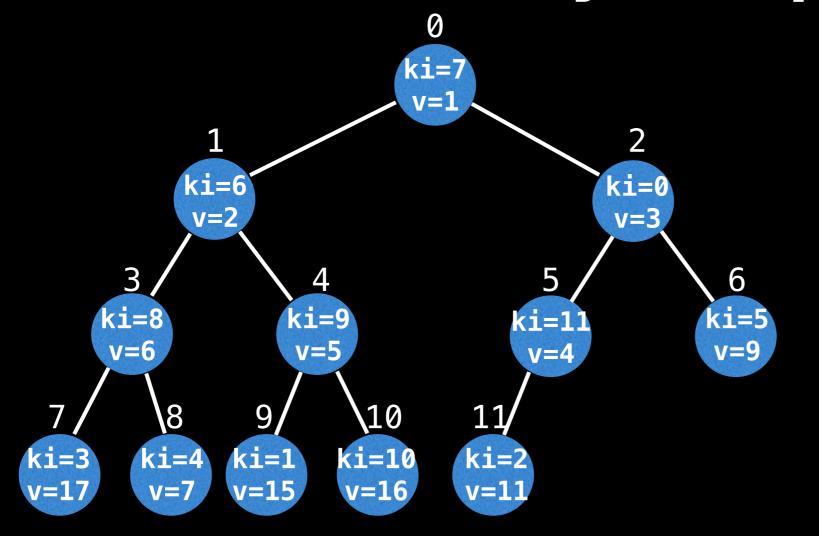
As an example, we can figure out the person represented in the node at index 2



ki = im[2] = 0 tells you the key
index (ki) for that node.



You can then use the ki value with the bidirectional hashtable to retrieve the actual key "Anna"



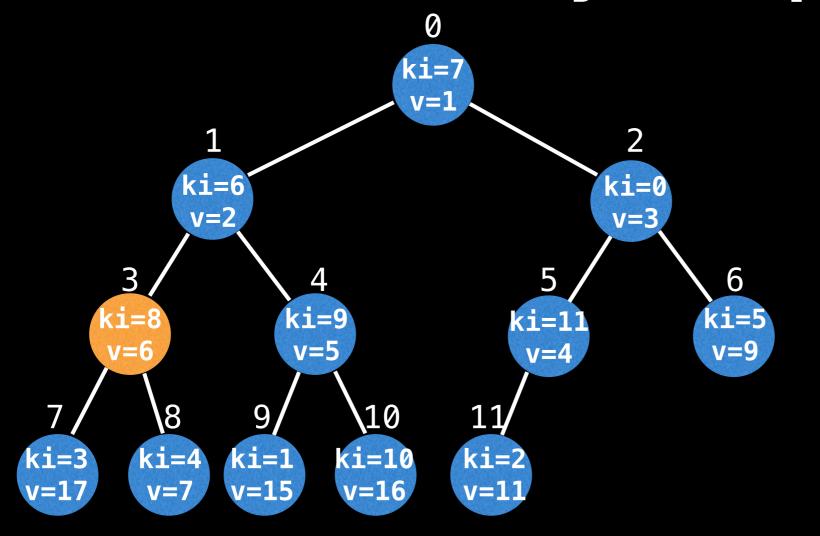
Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11

0		2	3	4	5	6	7	8	9	10	fii	12	13	14
---	--	---	---	---	---	---	---	---	---	----	-----	----	----	----

vals pm im

3	15	11	17	7	9	2	1	6	5	16	4	-1	-1	-1
2	9	11	7	8	6	1	0	3	4	10	5	-1	-1	-1
7	6	0	8	9	11	5	3	4	1	10	2	-1	-1	-1

Q: Which person (key) is being represented in the node at index position 3?



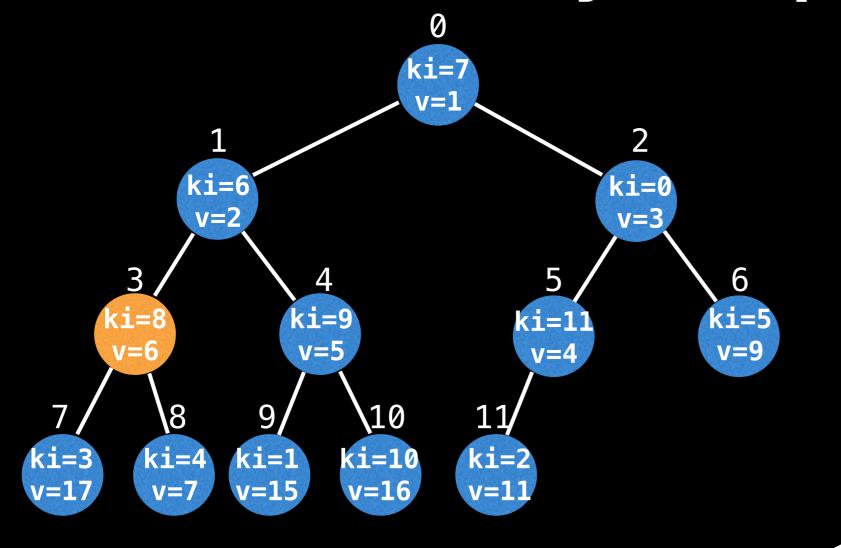
Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11

0 1 2 (3) 4 5 6 7 8 9 10 11 12 13	14
-----------------------------------	----

va	ls
p	m
i	m

3	15	11	17	7	9	2	1	6	5	16	4	-1	-1	-1
2	9	11	7	8	6	1	0	3	4	10	5	-1	-1	-1
7	6	0	8	9	11	5	3	4	1	10	2	-1	-1	-1

Q: Which person (key) is being represented in the node at index position 3?



Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	(8)
James	9
Ketly	10
Laura	11

									3						
vals															
pm	2	9	11	7	8	6	1	0	3	4	10	5	-1	-1	-1
im	7	6	0	8	9	11	5	3	4	1	10	2	-1	-1	-1

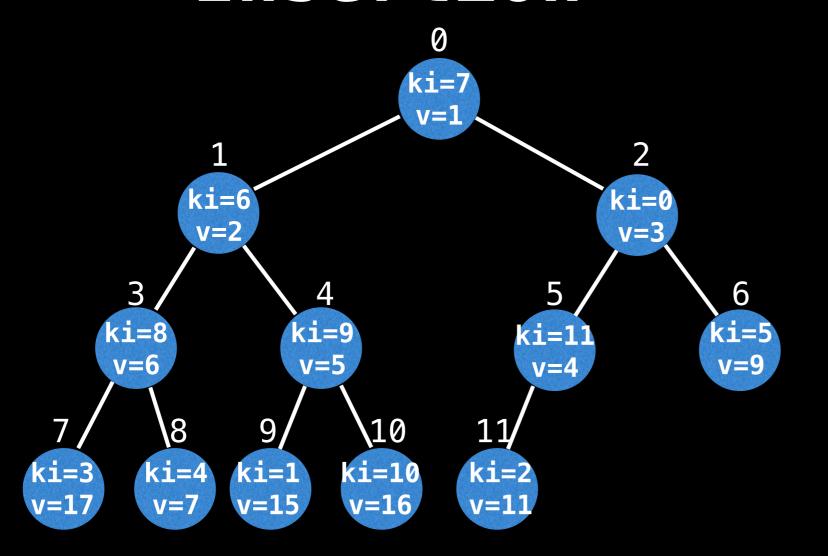
Q: Which person (key) is being represented in the node at index position 3?

IPQ as a binary heap Name ki Anna 0 Bella ki=7 Carly v=1 Dylan 3 ki=6 **ki**=0 **Emily** 4 v=2 v=3 Fred 5 3 5 6 George 6 ki=9 ki=8 ki=11 **ki=5** Henry v=6 v=5 v=9 v=4 8 Isaac 7 8 9 \10 11/ James ki=3 ki=1 **ki=4** ki=10 ki=2 Ketly 10 v=17v=15 v=16 v = 11v=711 Laura

											IU				
vals	3	15	11	17	7	9	2	1	6	5	16	4	-1	-1	-1
pm	2	9	11	7	8	8	1	0	3	4	10	5	-1	-1	-1
im	7	6	0	8	9	11	5	3	4	1	10	2	-1	-1	-1

A: "Isaac"

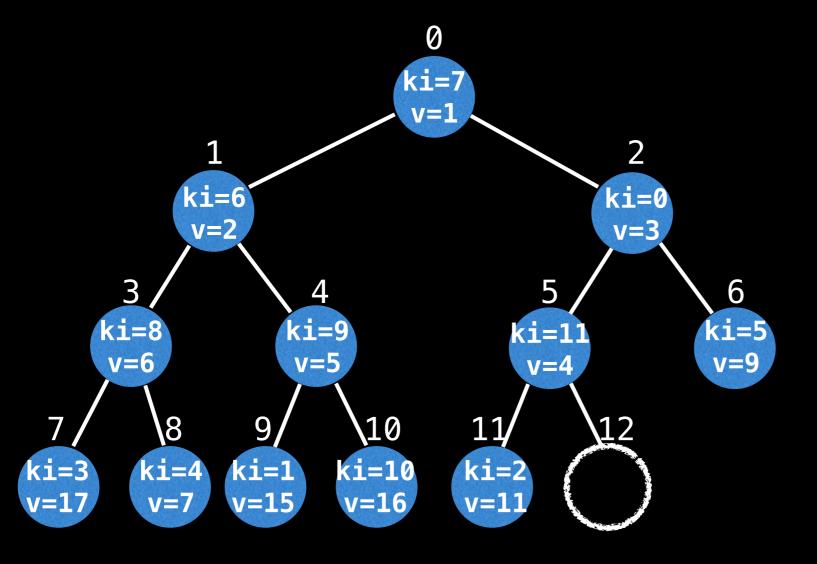
Insertion



Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11

va	ls
p	m
i	m

3	15	11	17	7	9	2	1	6	5	16	4	-1	-1	-1
2	9	11	7	8	6	1	0	3	4	10	5	-1	-1	-1
7	6	0	8	9	11	5	3	4	1	10	2	-1	-1	-1



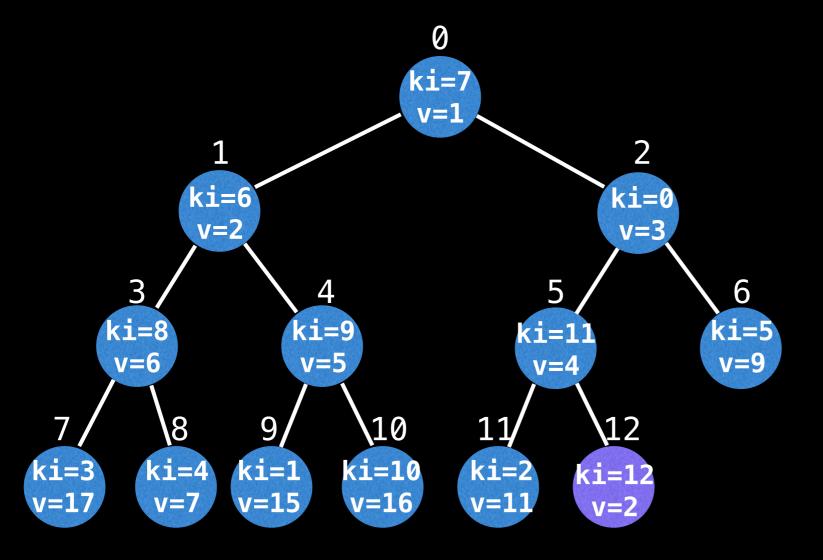
Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12

0 1 2 3 4 5 6 7 8 9 10 11 12 13 14	0		2	3	4	5	6	7	8	9	10	11	12	13	14
------------------------------------	---	--	---	---	---	---	---	---	---	---	----	----	-----------	----	----

va	ls
p	m
i	m

3	15	11	17	7	9	2	1	6	5	16	4	-1	-1	-1
2	9	11	7	8	6	1	0	3	4	10	5	-1	-1	-1
7	6	0	8	9	11	5	3	4	1	10	2	-1	-1	-1

Insert "Mary" with a value of 2



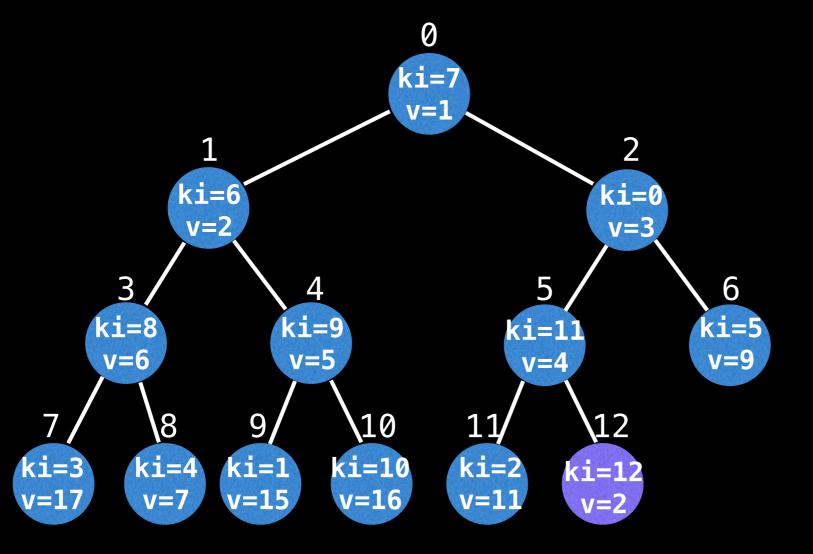
Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12
	10 11

0 1 2 3 4 5 6 7 8 9 10 11 12 13 14	0		2	3	4	5	6	7	8	9	10	11	12	13	14
------------------------------------	---	--	---	---	---	---	---	---	---	---	----	----	-----------	----	----

va	ls
p	m
\mathbf{i}	m

3	15	11	17	7	9	2	1	6	5	16	4	2	-1	-1
2	9	11	7	8	6	1	0	3	4	10	5	12	-1	-1
7	6	0	8	9	11	5	3	4	1	10	2	12	-1	-1

Insert "Mary" with a value of 2



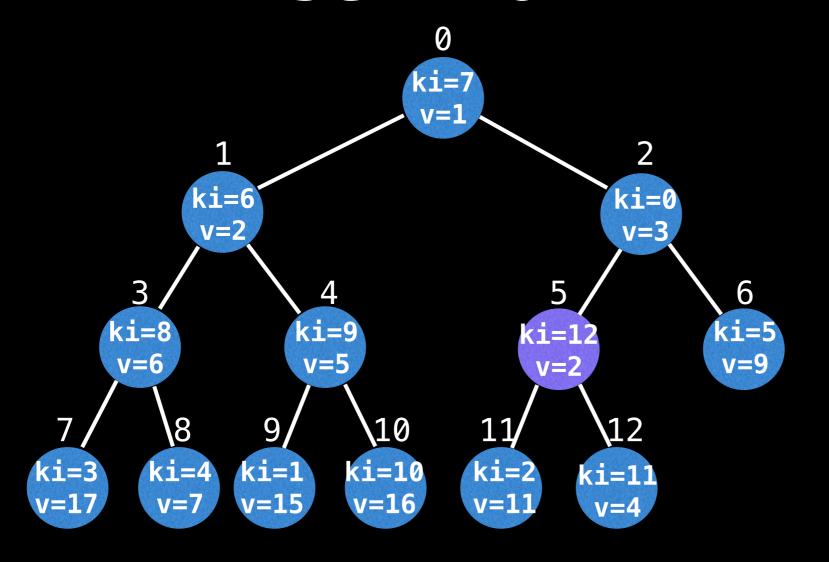
Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12

0	2	3	4	5	6	7	8	9	10	12	13	14

va	J	S
p	m	
i	m	

3	15	11	17	7	9	2	1	6	5	16	4	2	-1	-1
2	9	11	7	8	6	1	0	3	4	10	5	12	-1	-1
7	6	0	8	9	11	5	3	4	1	10	2	12	-1	-1

The heap invariant is not satisfied ATM since node at index 12 has a value less then node 5.

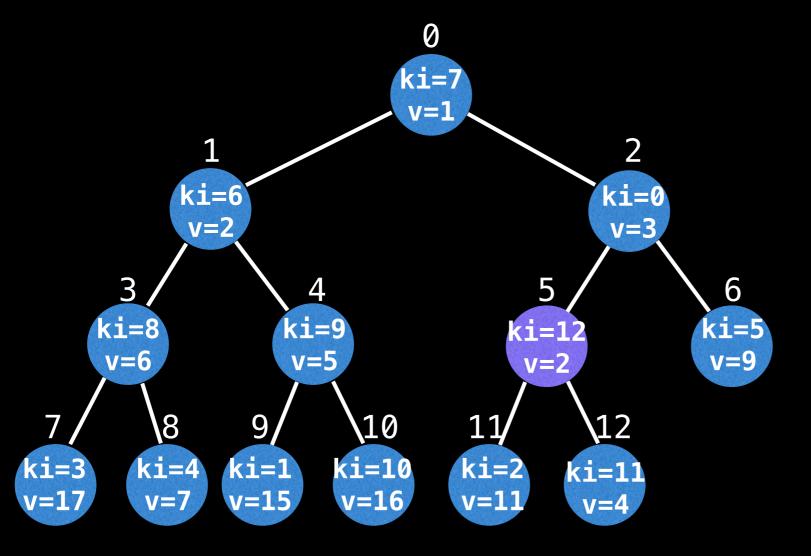


Name	ki						
Anna	0						
Bella	1						
Carly	2						
Dylan	3						
Emily	4						
Fred	5						
George	6						
Henry	7						
Isaac	8						
James	9						
Kelly	10						
Laura	11						
Mary	12						

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
--	---	---	---	---	---	---	---	---	---	---	----	----	----	----	----

V	a	l	S	
	p	m		
	i	m		

3	15	11	17	7	9	2	1	6	5	16	4	2	-1	-1
2	9	11	7	8	6	1	0	3	4	10	12	5	-1	-1
7	6	0	8	9	12	5	3	4	1	10	2	11	-1	-1



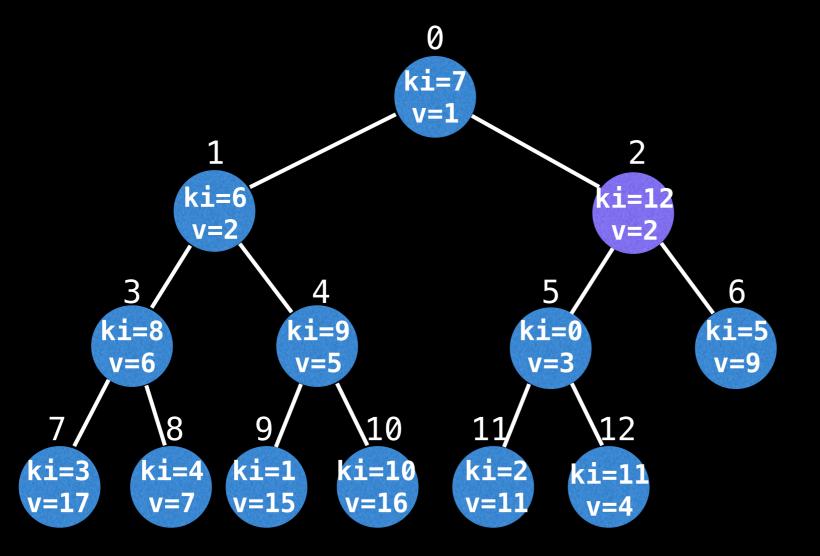
Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12

0 1 2 3 4 5 6 7 8 9 10 11 12 13 14	0		2	3	4	5	6	7	8	9	10	11	12	13	14
------------------------------------	---	--	---	---	---	---	---	---	---	---	----	----	-----------	----	----

va	ls
p	m
i	m

3	15	11	17	7	9	2	1	6	5	16	4	2	-1	1
2	9	11	7	8	6	1	0	3	4	10	12	5	-1	-1
7	6	0	8	9	12	5	3	4	1	10	2	11	-1	-1

The heap invariant is still not satisfied so keep swapping upwards.

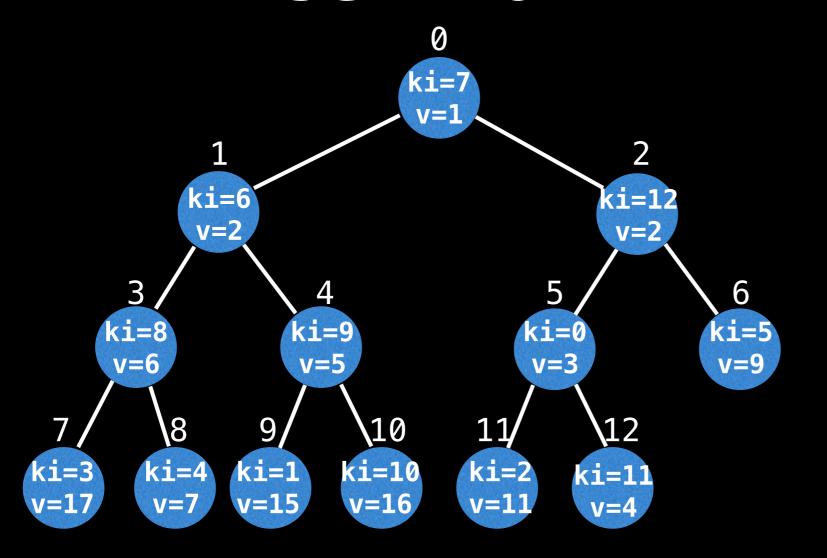


Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12

	0		2	3	4	5	6		8	9	10		12	13	14
vals	3	15	11	17	7	9	2	1	6	5	16	4	2	-1	

vals pm im

3	15	11	17	7	9	2	1	6	5	16	4	2	-1	-1
5	9	11	7	8	6	1	0	3	4	10	12	2	-1	-1
7	8	12	8	9	0	5	3	4	1	10	2	11	-1	-1



Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
--	---	---	---	---	---	---	---	---	---	---	----	----	----	----	----

va	ls
p	m
i	m

3	15	11	17	7	9	2	1	6	5	16	4	2	-1	-1
5	9	11	7	8	6	1	0	3	4	10	12	2	-1	-1
7	6	12	8	9	0	5	3	4	1	10	2	11	-1	-1

```
# Inserts a value into the min indexed binary
# heap. The key index must not already be in
# the heap and the value must not be null.
function insert(ki, value):
  values[ki] = value
  # 'sz' is the current size of the heap
  pm[ki] = sz
  im[sz] = ki
  swim(sz)
  sz = sz + 1
```

```
# Inserts a value into the min indexed binary
# heap. The key index must not already be in
# the heap and the value must not be null.
function insert(ki, value):
 values[ki] = value
  # 'sz' is the current size of the heap
  pm[ki] = sz
  im[sz] = ki
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  sz = sz + 1
```

```
# Inserts a value into the min indexed binary
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function insert(ki, value):
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  swim(sz)
  sz = sz + 1
```

```
# Inserts a value into the min indexed binary
# heap. The key index must not already be in
# the heap and the value must not be null.
function insert(ki, value):
  values[ki] = value
  # 'sz' is the current size of the heap
  pm[ki] = sz
  im[sz] = ki
  swim(sz)
  sz = sz + 1
```

```
# Swims up node i (zero based) until heap
# invariant is satisfied.
function swim(i):
  for (p = (i-1)/2; i > 0 \text{ and less}(i, p)):
    swap(i, p)
    i = p
    p = (i-1)/2
function swap(i, j):
 pm[im[j]] = i
 pm[im[i]] = j
 tmp = im[i]
 im[i] = im[j]
 im[j] = tmp
function less(i, j):
  return values[im[i]] < values[im[j]]</pre>
```

```
# Swims up node i (zero based) until heap
# invariant is satisfied.
function swim(i):
 for (p = (i-1)/2; i > 0 and less(i, p)):
    swap(i, p)
    i = p
    p = (i-1)/2
function swap(i, j):
 pm[im[j]] = i
 pm[im[i]] = j
 tmp = im[i]
 im[i] = im[j]
 im[j] = tmp
function less(i, j):
  return values[im[i]] < values[im[j]]</pre>
```

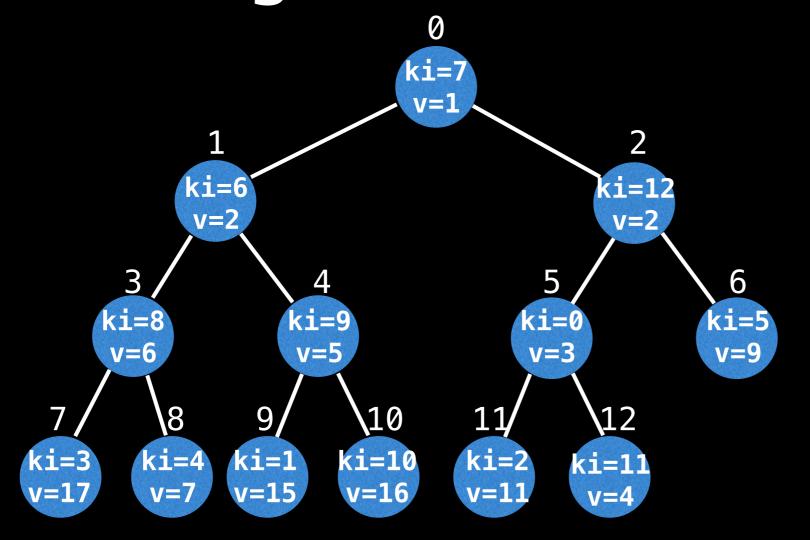
```
# Swims up node i (zero based) until heap
# invariant is satisfied.
function swim(i):
  for (p = (i-1)/2; i > 0 \text{ and less}(i, p)):
    swap(i, p)
    i = p
    p = (i-1)/2
function swap(i, j):
 pm[im[j]] = i
 pm[im[i]] = j
 tmp = im[i]
 im[i] = im[j]
 im[j] = tmp
function less(i, j):
  return values[im[i]] < values[im[j]]</pre>
```

```
# Swims up node i (zero based) until heap
# invariant is satisfied.
function swim(i):
  for (p = (i-1)/2; i > 0 \text{ and less}(i, p)):
   swap(i, p)
   i = p
   p = (i-1)/2
function swap(i, j):
 pm[im[j]] = i
 pm[im[i]] = j
 tmp = im[i]
 im[i] = im[j]
 im[j] = tmp
function less(i, j):
  return values[im[i]] < values[im[j]]</pre>
```

```
# Swims up node i (zero based) until heap
# invariant is satisfied.
function swim(i):
  for (p = (i-1)/2; i > 0 \text{ and less}(i, p)):
    swap(i, p)
    i = p
    p = (i-1)/2
function swap(i, j):
 pm[im[j]] = i
 pm[im[i]] = j
 tmp = im[i]
 im[i] = im[j]
 im[j] = tmp
function less(i, j):
  return values[im[i]] < values[im[j]]</pre>
```

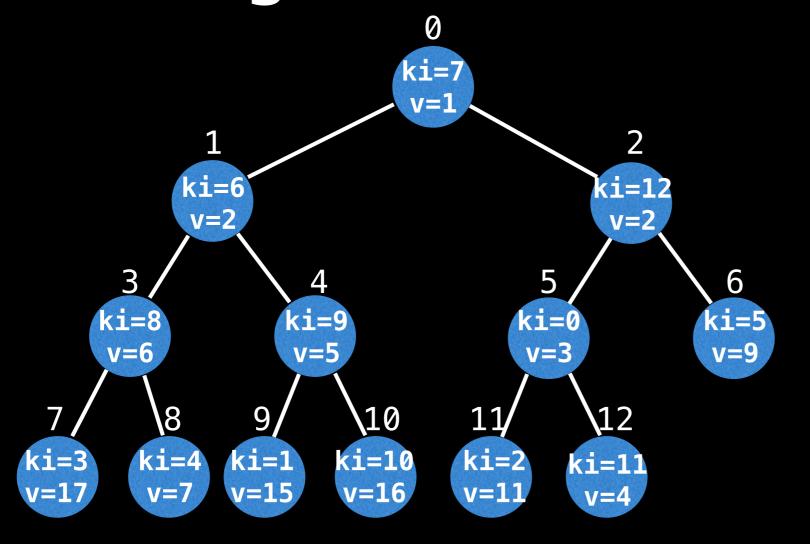
```
# Swims up node i (zero based) until heap
# invariant is satisfied.
function swim(i):
  for (p = (i-1)/2; i > 0 \text{ and less}(i, p)):
    swap(i, p)
    i = p
    p = (i-1)/2
function swap(i, j):
 pm[im[j]] = i
 pm[im[i]] = j
 tmp = im[i]
 im[i] = im[j]
 im[j] = tmp
function less(i, j):
  return values[im[i]] < values[im[j]]</pre>
```

```
# Swims up node i (zero based) until heap
# invariant is satisfied.
function swim(i):
  for (p = (i-1)/2; i > 0 \text{ and less}(i, p)):
    swap(i, p)
    i = p
    p = (i-1)/2
function swap(i, j):
 pm[im[j]] = i
 pm[im[i]] = j
 tmp = im[i]
 im[i] = im[j]
 im[j] = tmp
function less(i, j):
  return values[im[i]] < values[im[j]]</pre>
```



Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12

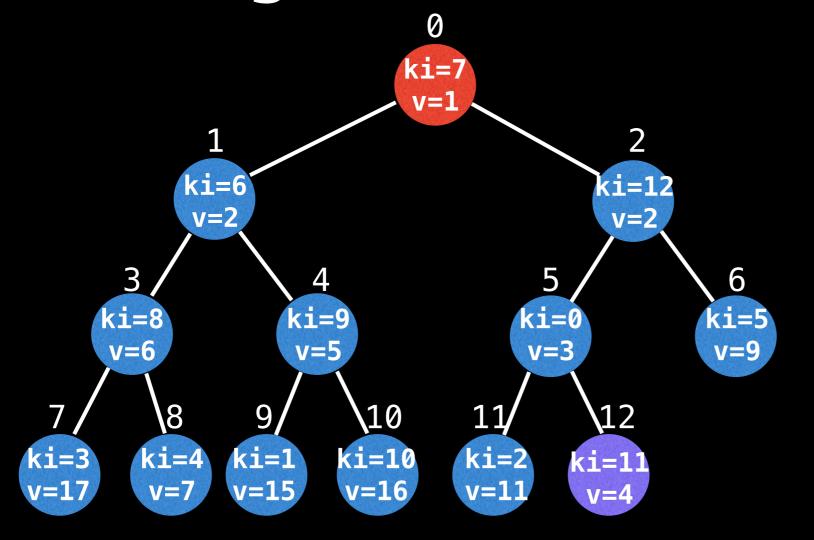
Polling is still O(log(n)) in an IPQ, but removing is improved from O(n) in a traditional PQ to O(log(n)) since node position lookups are O(1) but repositioning is still O(log(n))



Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12

	U			<u> </u>	4	5	O		0	3	10		12	13	14
vals	3	15	11	17	7	9	2	1	6	5	16	4	2	-1	-1
pm	5	9	11	7	8	6	1	0	3	4	10	12	2	-1	-1
im	7	6	12	8	9	0	5	3	4	1	10	2	11	-1	-1

Let's poll the root node. The required steps are almost exactly like a regular binary heap.



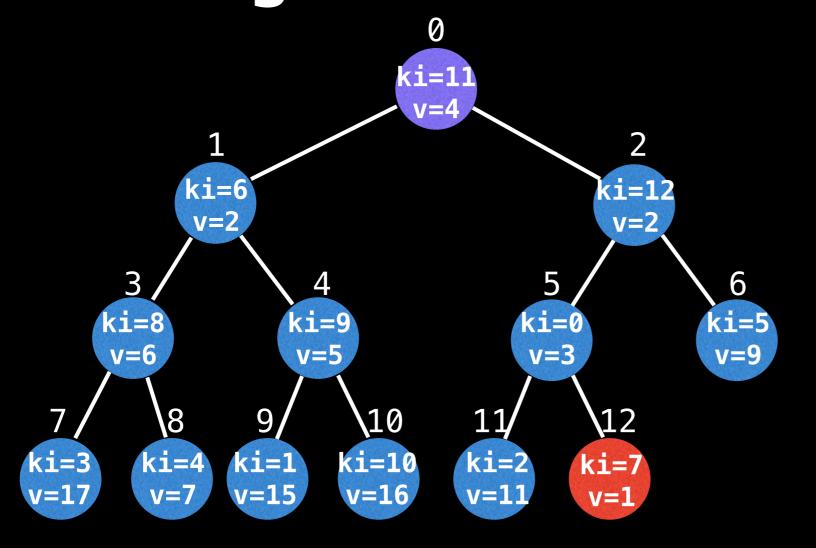
Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12
	40 44

0 1 2 3 4 5 6 7 8 9 10 11 12 13 14	0			2 ;	3 4	4	5	6	7	8	9	10	11	12	13	14
------------------------------------	---	--	--	-----	-----	---	---	---	---	---	---	----	-----------	-----------	-----------	----

va	ls
p	m
i	m

3	15	11	17	7	9	2	1	6	5	16	4	2	-1	-1
5	9	11	7	8	6	1	0	3	4	10	12	2	-1	-1
7	6	12	8	9	0	5	3	4	1	10	2	11	-1	-1

Exchange the root node and the bottom right node.

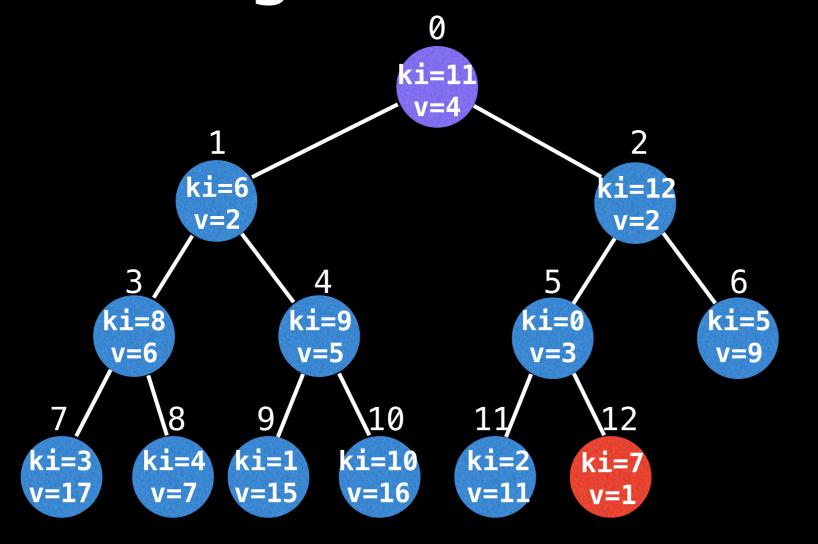


Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12

0		2	3	4	5	6	7	8	9	10		12	13	14
---	--	---	---	---	---	---	---	---	---	----	--	----	-----------	----

va	ls
p	m
i	m

3	15	11	17	7	9	2	1	6	5	16	4	2	-1	-1
5	9	11	7	8	6	1	12	3	4	10	0	2	-1	-1
11	6	12	8	9	0	5	3	4	1	10	2	₹ 7	-1	-1



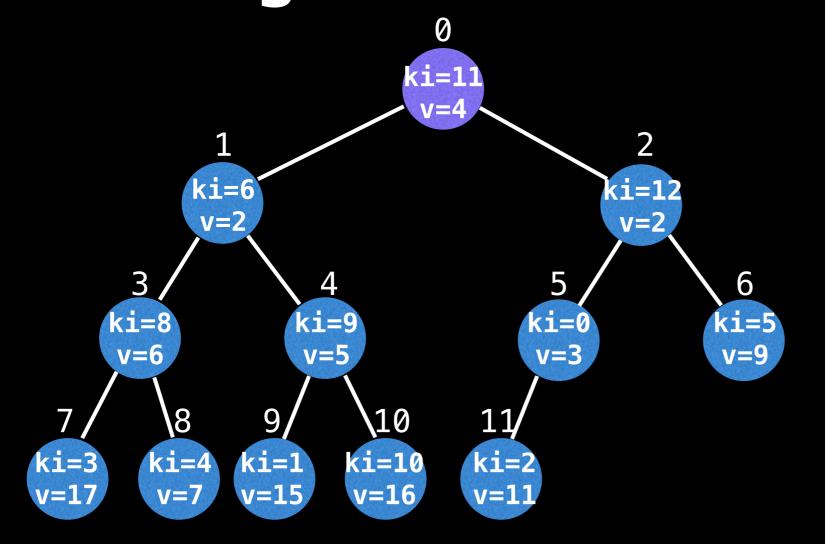
Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12

0	2	3	4	5	6	7	8	9	10	11	12	13	14	

vals pm im

3	15	11	17	7	9	2	1	6	5	16	4	2	-1	-1
5					Ī			Ī				i e		
11	6	12	8	9	0	5	3	4	1	10	2	7	-1	-1

Remove and return the key-value pair (7, 1)

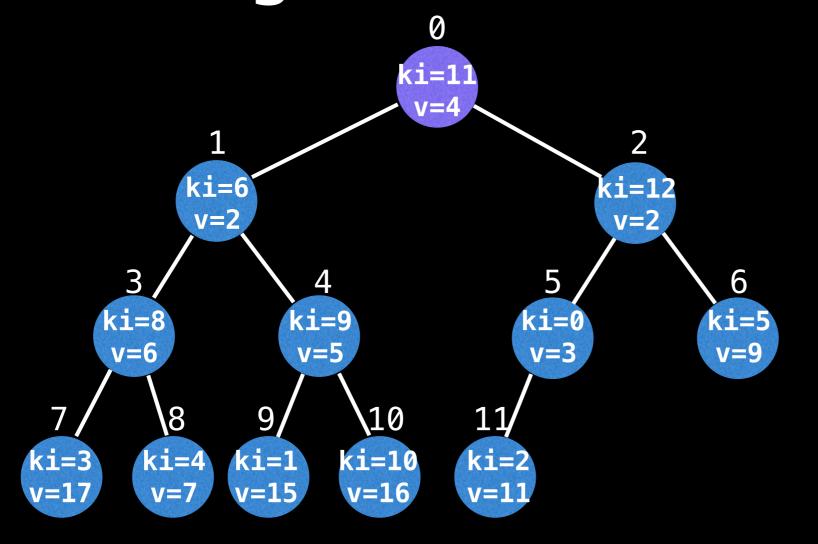


Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	

V	a	J	S
	p	M	
	i	m	

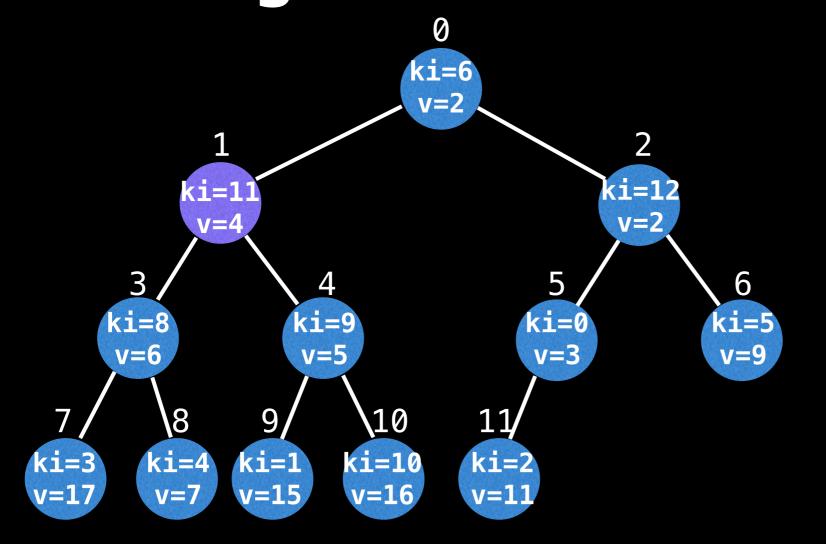
3	15	11	17	7	9	2	Ø	6	5	16	4	2	-1	-1
5	9	11	7	8	6	1	-1	3	4	10	0	2	-1	-1
11	6	12	8	9	0	5	3	4	1	10	2	-1	-1	-1



Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12

	U		_	3	4	5	O		0	9	10		12	13	1 -
vals	3	15	11	17	7	9	2	Ø	6	5	16	4	2	-1	-1
pm	5	9	11	7	8	6	1	-1	3	4	10	0	2	-1	-1
im	11	6	12	8	9	0	5	3	4	1	10	2	-1	-1	-1

Finally restore heap invariant by moving swapped purple node up or down.

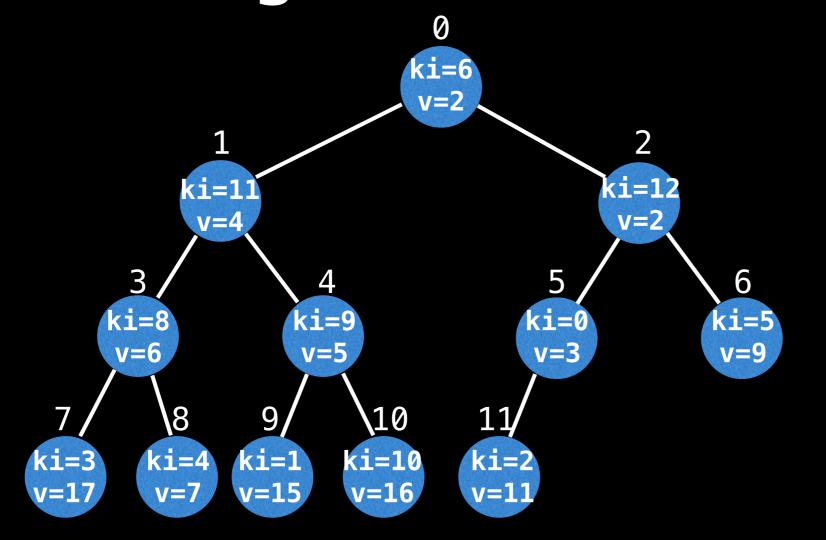


Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14

va	ls
p	m
i	m

3	15	11	17	7	9	2	Ø	6	5	16	4	2	-1	-1
5	9	11	7	8	6	0_	-1	3	4	10	~1	2	-1	-1
6	11	12	8	9	0	5	3	4	1	10	2	-1	-1	-1

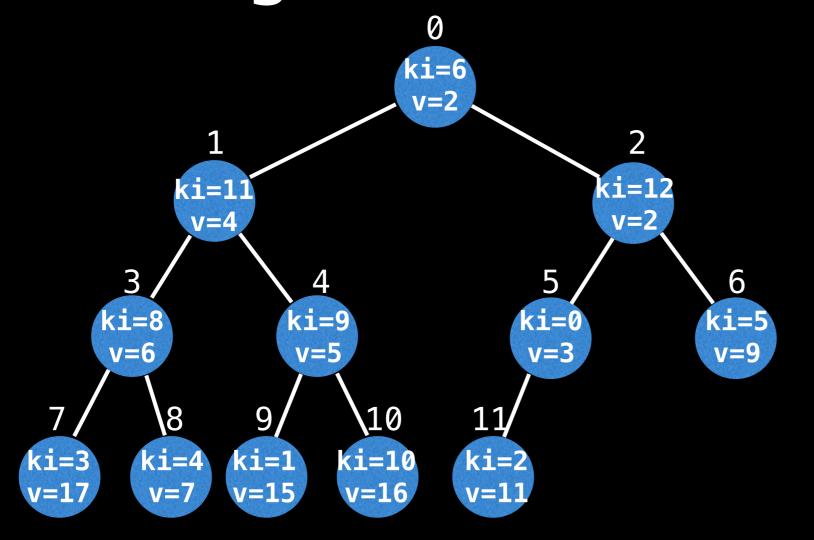


Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12

	U		2	3	4	5	6		8	9	10	12	13	14
valc	2	15	44	17	7	0	2	$\overline{\alpha}$	6		16	 2		1

vals pm im

3	15	11	17	7	9	2	Ø	6	5	16	4	2	-1	-1
5	9	11	7	8	6	0	-1	3	4	10	1	2	-1	-1
6	11	12	8	9	0	5	3	4	1	10	2	-1	-1	-1



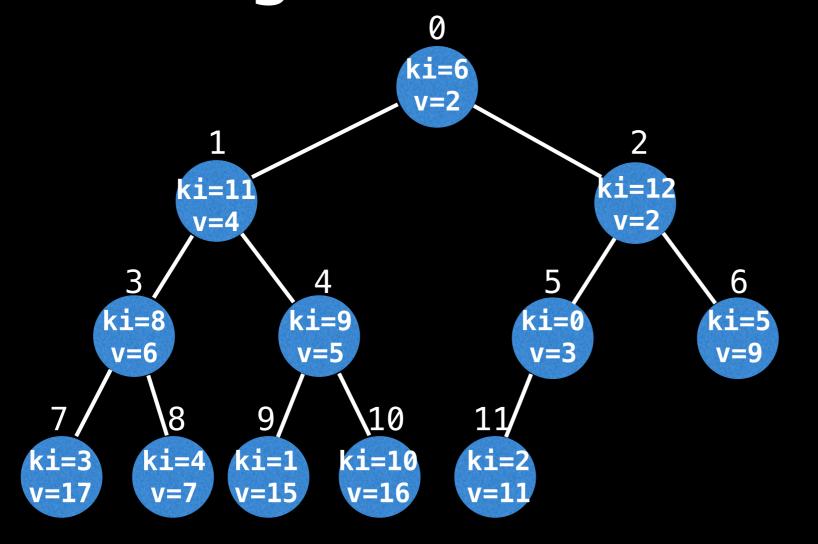
Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	

va	ls
p	m
i	m

3	15	11	17	7	9	2	Ø	6	5	16	4	2	-1	-1
5	9	11	7	8	6	0	-1	3	4	10	1	2	-1	-1
6	11	12	8	9	0	5	3	4	1	10	2	-1	-1	-1

Now let's remove "Laura" from IPQ



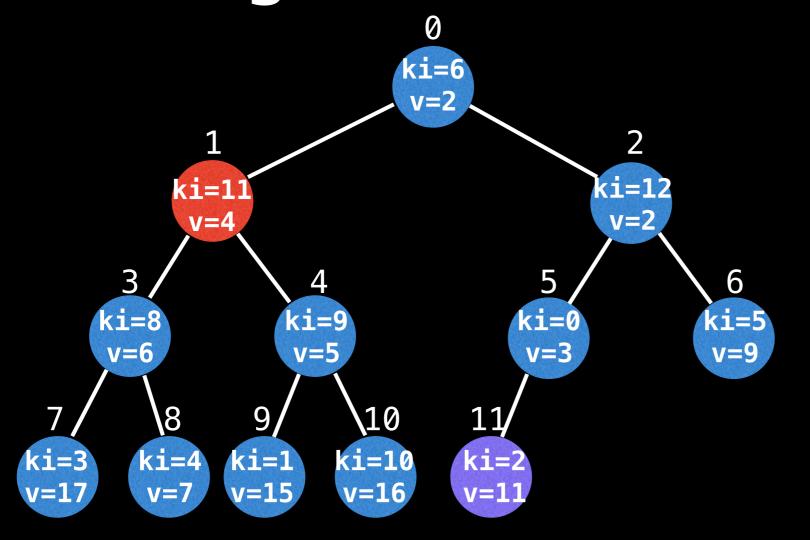
Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura -	→ (11)
Mary	12

0 1 2 3 4 5 6 7 8 9 10 11 12 13 14

vals pm im

3	15	11	17	7	9	2	Ø	6	5	16	4	2	-1	-1
5	9	11	7	8	6	0	-1	3	4	10	1	2	-1	-1
6	11	12	8	9	0	5	3	4	1	10	2	-1	-1	-1

The ki for "Laura" is 11



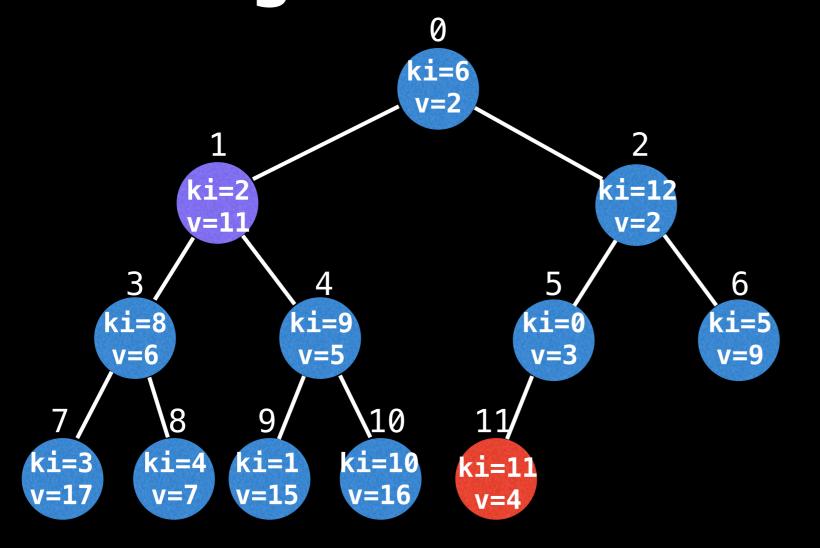
Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura -	→ (11)
Mary	/12

0	1	2	3	4	5	6	7	8	9	10	11	12/13	14
---	---	---	---	---	---	---	---	---	---	----	----	-------	----

va	l	S
p	m	
i	m	

3	15	11	17	7	9	2	Ø	6	5	16	4	/2	-1	-1
5	9	11	7	8	6	0	-1	3	4	10		2	-1	-1
	11													

Look inside the position map to find what node is associated with ki=11 and perform swap.

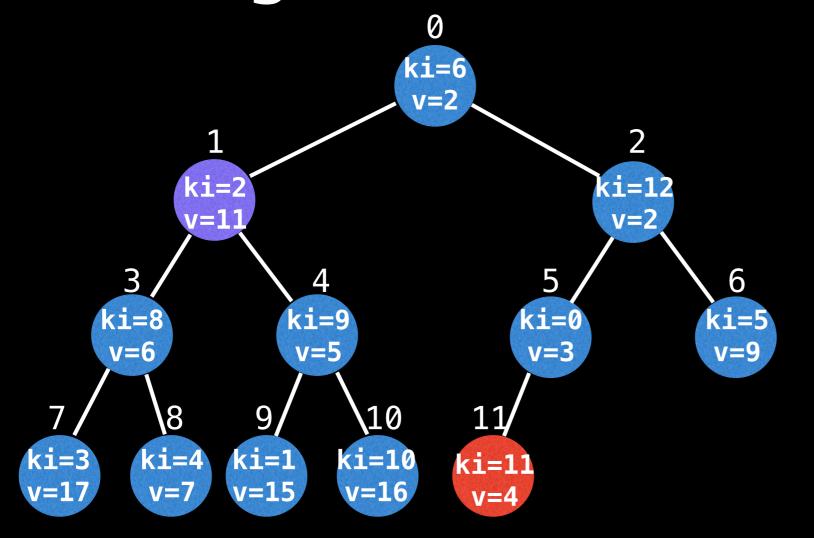


Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12

	O		2	3	4	5	6		8	9	10	12	13	14
valc	2	15	44	17	7	0	2	\bigcirc	6	5	16	 2		

vals pm im

3	15	11	17	7	9	2	Ø	6	5	16	4	2	-1	-1
5	9	1	7	8	6	0	-1	3	4	10	11	2	-1	-1
6	2-	12	8	9	0	5	3	4	1	10	11	-1	-1	-1



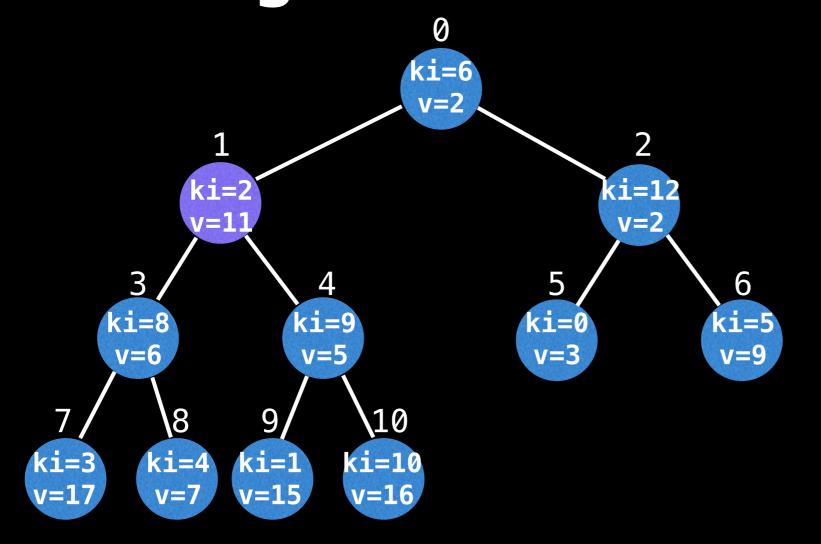
Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12

0	2	3	4	5	6	7	8	9	10	12	13	14
						4						

vals pm im

3	15	11	17	7	9	2	Ø	6	5	16	4	2	-1	-1
5	9	1	7	8	6	0	-1	3	4	10	(11)	2	-1	-1
6	2	12	8	9	0	5	3	4	1	10	11	-1	-1	-1

Remove and return key-value pair (11, 4)

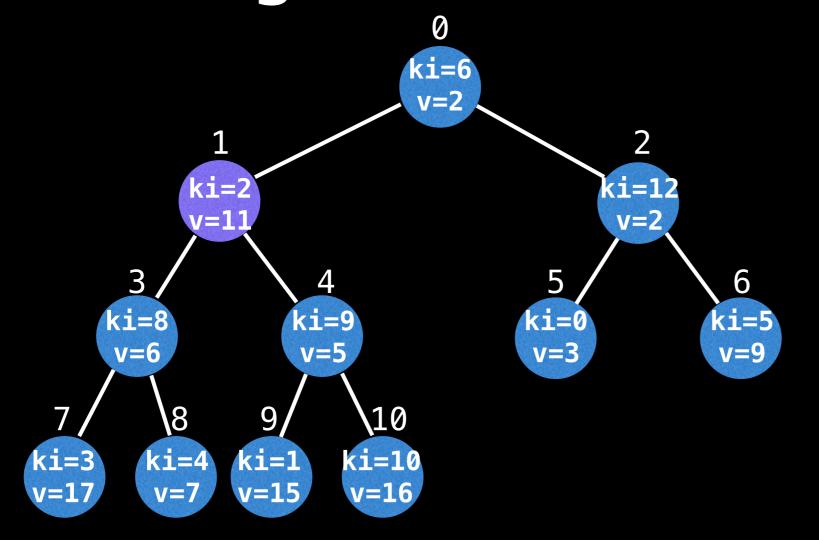


Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12

0	2	3	4	5	6	7	8	9	10	12	13	14

va	ls
p	m
i	m

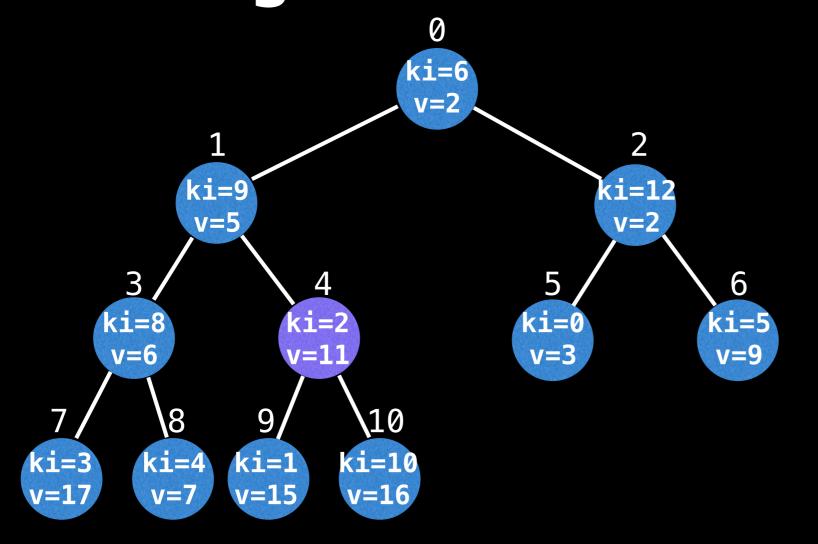
3	15	11	17	7	9	2	Ø	6	5	16	Ø	2	-1	-1
5	9	1	7	8	6	0	-1	3	4	10	-1	2	-1	-1
6	2	12	8	9	0	5	3	4	1	10	-1	-1	-1	-1



Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12

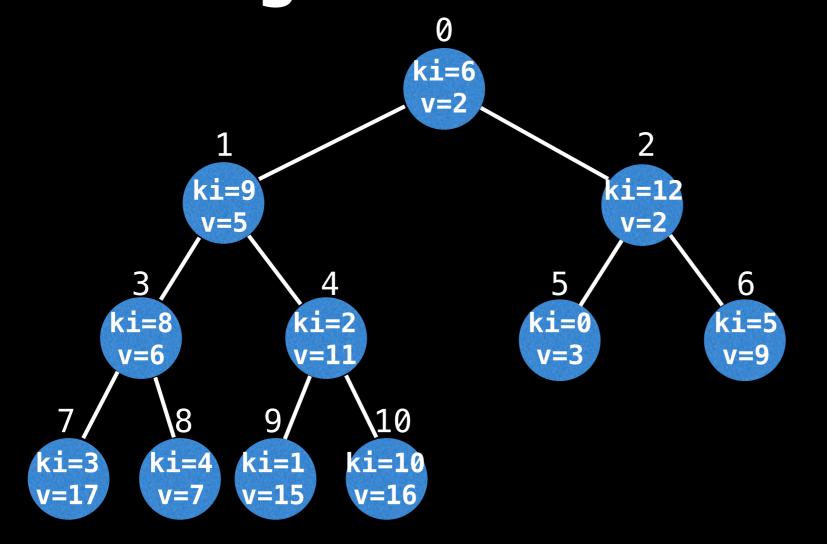
	U		_	3	4	5	O		Ö	9	IU		12	13	14
vals	3	15	11	17	7	9	2	Ø	6	5	16	Ø	2	-1	-1
pm	5	9	1	7	8	6	0	-1	3	4	10	-1	2	-1	-1
im	6	2	12	8	9	0	5	3	4	1	10	-1	-1	-1	-1

Finally restore heap invariant by moving swapped purple node up or down.



Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12

											10				
vals	3	15	11	17	7	9	2	Ø	6	5	16	Ø	2	-1	-1
pm	5	9	4	7	8	6	0	-1	3	1	10	-1	2	-1	-1
im	6	9_	12	8	~2	0	5	3	4/	1	10	-1	-1	-1	-1



Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	

Vä	a	l	S
	Ol	m	
		m	

3	15	11	17	7	9	2	Ø	6	5	16	Ø	2	-1	-1
5	9	4	7	8	6	0	-1	3	1	10	-1	2	-1	-1
6	9	12	8	2	0	5	3	4	1	10	-1	-1	-1	-1

```
# Deletes the node with the key index ki
# in the heap. The key index ki must exist
# and be present in the heap.
function remove(ki):
  i = pm[ki]
  swap(i, sz)
  sz = sz - 1
  sink(i)
  swim(i)
  values[ki] = null
  pm[ki] = -1
  im[sz] = -1
```

```
# Deletes the node with the key index ki
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  swim(i)
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  im[sz] = -1
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  swap(i, sz)
  sz = sz - 1
  sink(i)
 swim(i)
  values[ki] = null
  pm[ki] = -1
  im[sz] = -1
```

```
# Sinks the node at index i by swapping
# itself with the smallest of the left
# or the right child node.
function sink(i):
  while true:
    left = 2*i + 1
    right = 2*i + 2
    smallest = left
    if right < sz and less(right, left):</pre>
      smallest = right
    if left >= sz or less(i, smallest):
      break
    swap(smallest, i)
    i = smallest
```

```
# Sinks the node at index i by swapping
# itself with the smallest of the left
# or the right child node.
function sink(i):
  while true:
    left = 2*i + 1
    right = 2*i + 2
    smallest = left
    if right < sz and less(right, left):</pre>
      smallest = right
    if left >= sz or less(i, smallest):
      break
    swap(smallest, i)
```

i = smallest

```
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    left = 2*i + 1
    right = 2*i + 2
    smallest = left
    if right < sz and less(right, left):</pre>
```

```
smallest = right
```

```
if left >= sz or less(i, smallest):
  break
```

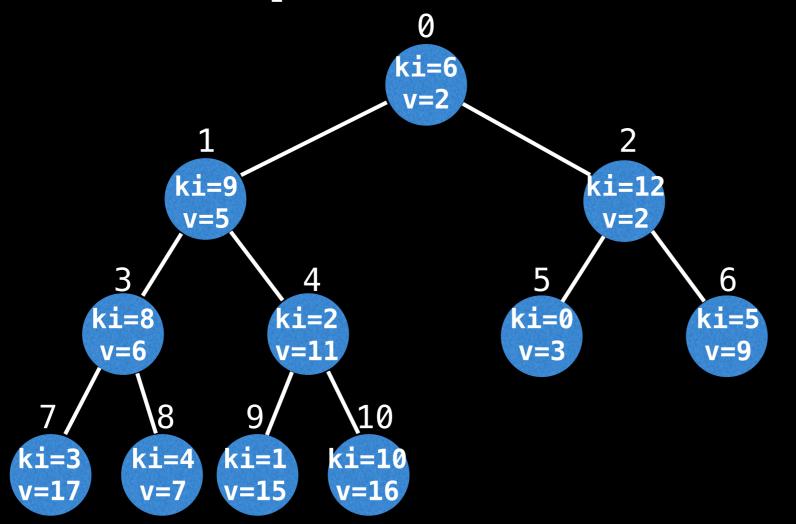
```
swap(smallest, i)
i = smallest
```

```
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# itself with the smallest of the left
# or the right child node.
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    left = 2*i + 1
    right = 2*i + 2
    smallest = left
    if right < sz and less(right, left):</pre>
      smallest = right
    if left >= sz or less(i, smallest):
      break
    swap(smallest, i)
```

i = smallest

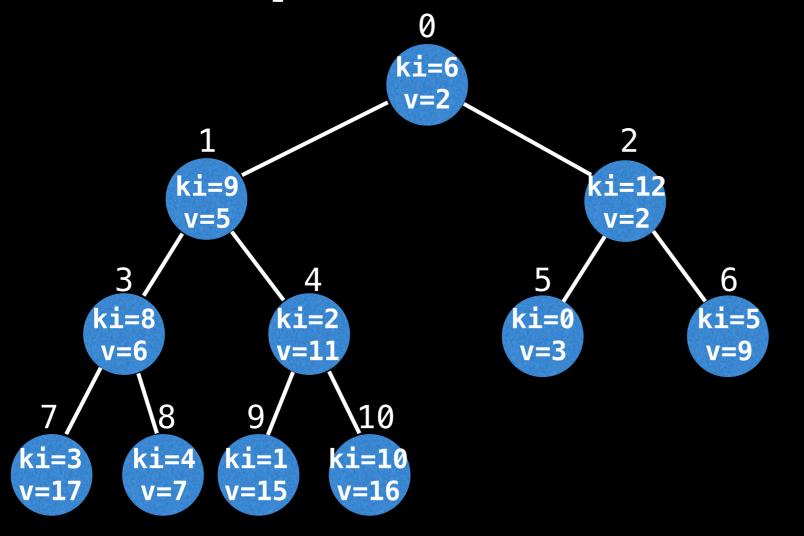
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    left = 2*i + 1
    right = 2*i + 2
    smallest = left
    if right < sz and less(right, left):</pre>
      smallest = right
    if left >= sz or less(i, smallest):
      break
```

swap(smallest, i)
i = smallest



Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12

Similar to removals, updates in a min indexed binary heap also take O(log(n)) due to O(1) lookup time to find the node and O(log(n)) time to adjust where the key-value pair should appear in the heap.



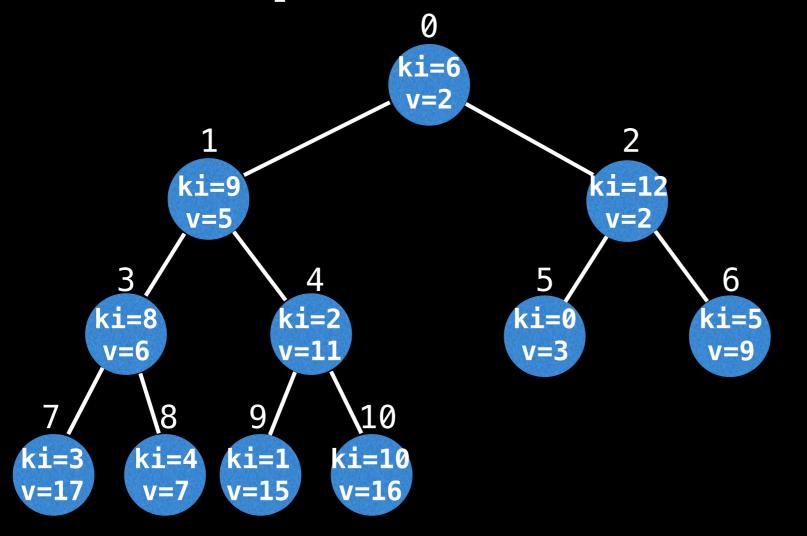
Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14

vals pm im

3	15	11	17	7	9	2	Ø	6	5	16	Ø	2	-1	-1
5	9	4	7	8	6	0	-1	3	1	10	-1	2	-1	-1
6	9	12	8	2	0	5	3	4	1	10	-1	-1	-1	-1

Update "Carly" to have a new value of 1



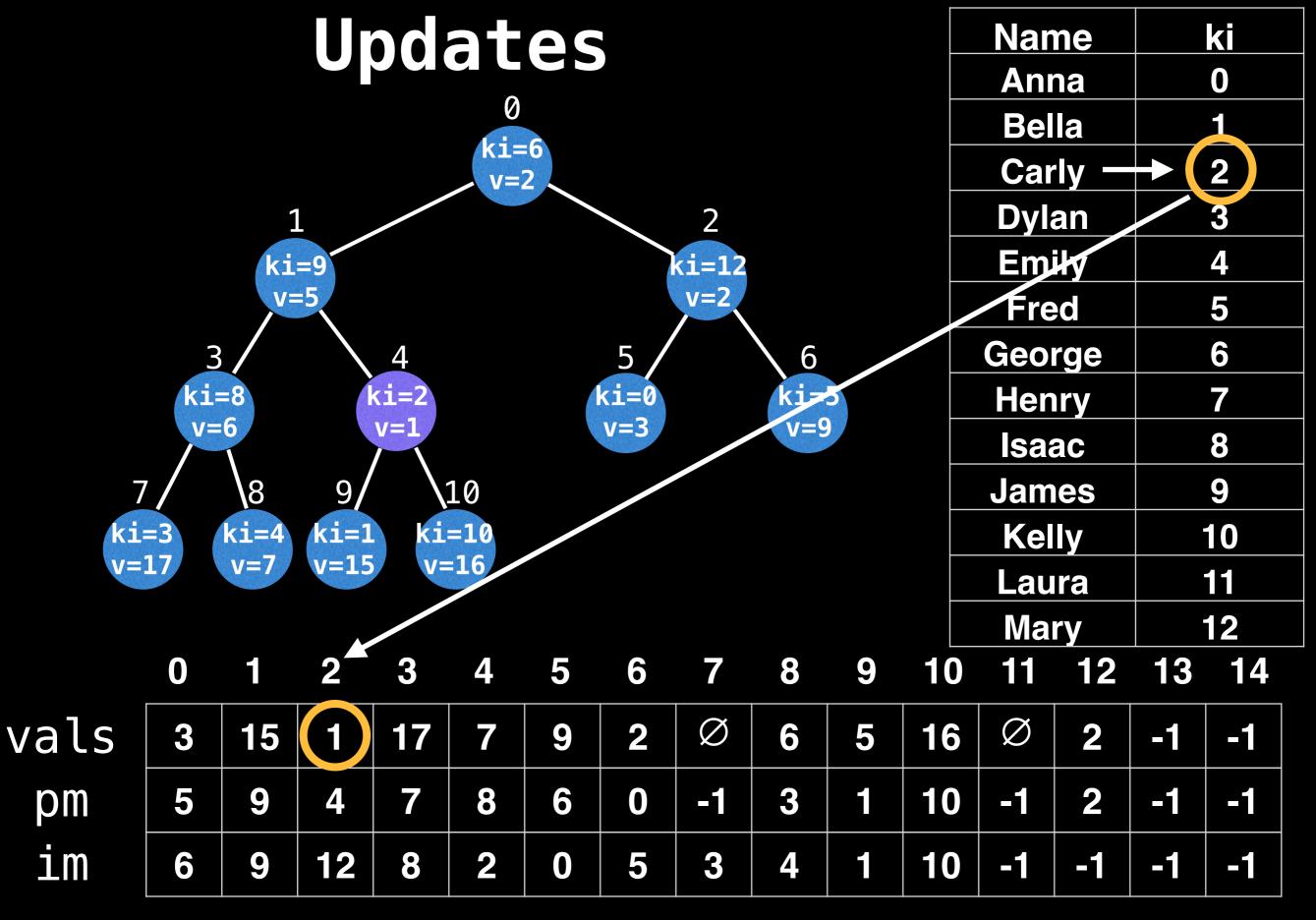
Name	ki
Anna	0
Bella	1
Carly -	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14

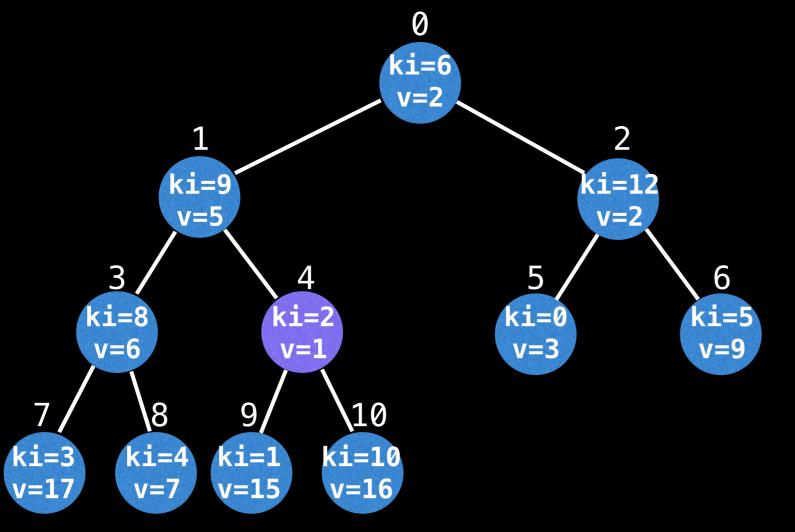
vals pm im

3	15	11	17	7	9	2	Ø	6	5	16	Ø	2	-1	-1
5	9	4	7	8	6	0	-1	3	1	10	-1	2	-1	-1
6	9	12	8	2	0	5	3	4	1	10	-1	-1	-1	-1

Update "Carly" to have a new value of 1



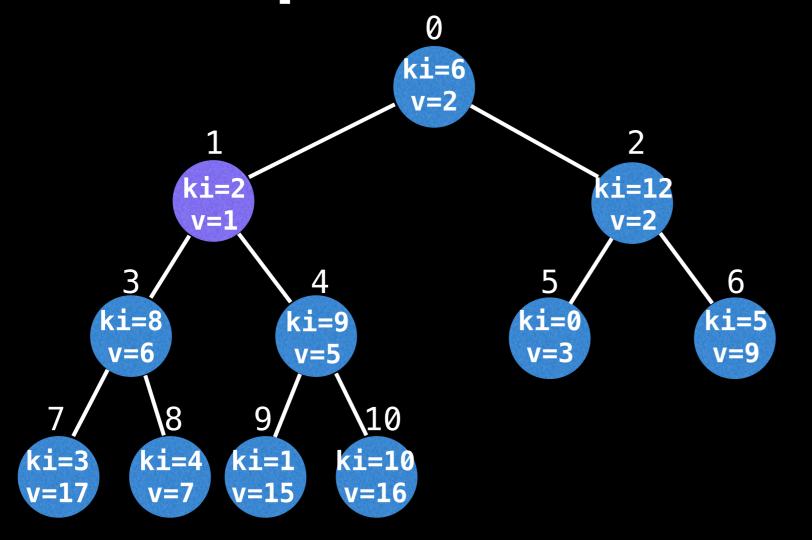
Update values array at ki to have value 1



ki
0
1
2
3
4
5
6
7
8
9
10
11
12

											. •				
vals	3	15	1	17	7	9	2	Ø	6	5	16	Ø	2	-1	-1
pm	5	9	4	7	8	6	0	-1	3	1	10	-1	2	-1	-1
im	6	9	12	8	2	0	5	3	4	1	10	-1	-1	-1	-1

Finally satisfy heap invariant by moving node either up or down the heap.

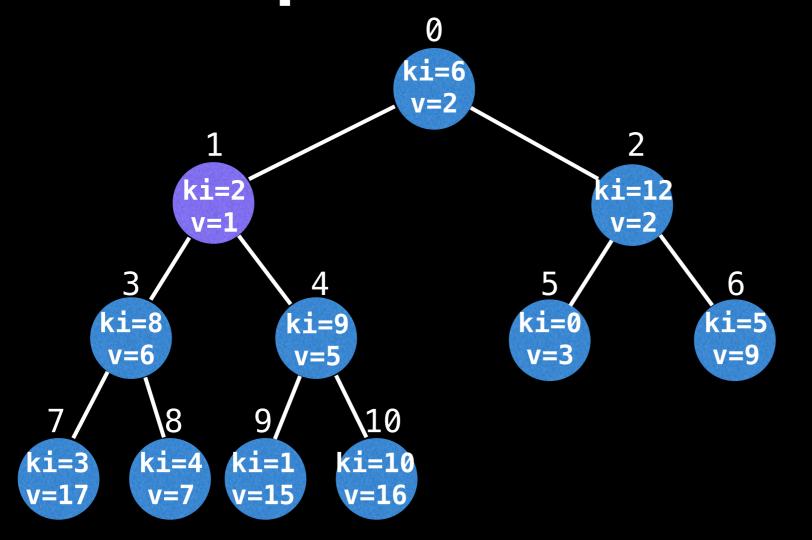


Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12

0	1	2	3	4	5	6	7	8	9	10	111	12	13	14
												1		

vals pm im

3	15	1	17	7	9	2	Ø	6	5	16	Ø	2	-1	-1
5	9	1 🗙	7	8	6	0	-1	3	4	10	-1	2	-1	-1
														-1

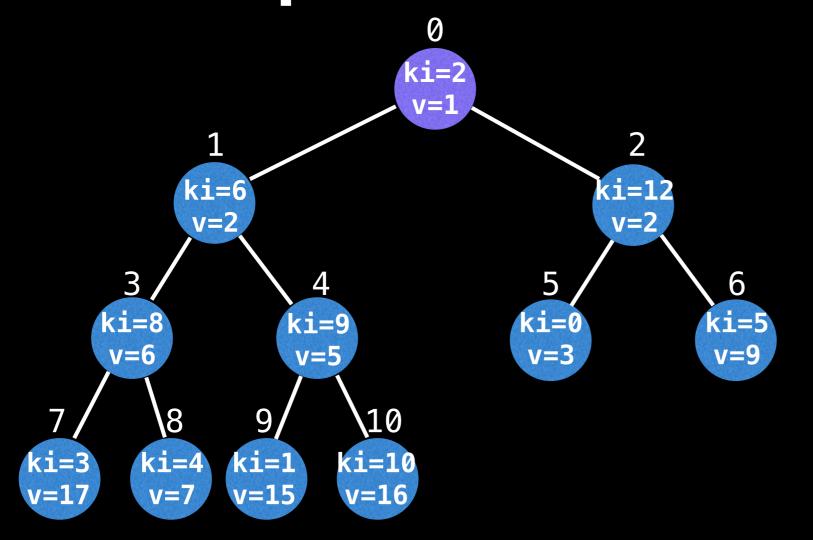


Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12

	0		2	3	4	5	6	7	8	9	10	11	12	13	14
--	---	--	---	---	---	---	---	---	---	---	----	----	----	----	----

Vä	a	l	S
	Ol	m	
		m	

3	15	1	17	7	9	2	Ø	6	5	16	Ø	2	-1	-1
5	9	1	7	8	6	0	-1	3	4	10	-1	2	-1	-1
6	2	12	8	9	0	5	3	4	1	10	-1	-1	-1	-1

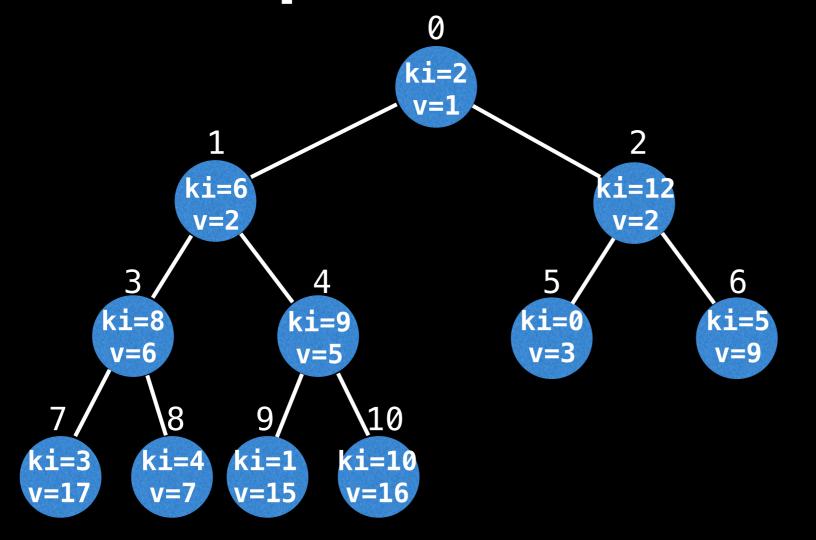


Name	ki
Anna	0
Bella	1
Carly	2
Dylan	3
Emily	4
Fred	5
George	6
Henry	7
Isaac	8
James	9
Kelly	10
Laura	11
Mary	12

0 1 2 3 4 5 6 7 8 9 10 11 12 13 14

va	ls
p	m
i	m

3	15	1	17	7	9	2	Ø	6	5	16	Ø	2	-1	-1
5	9	0_	7	8	6	~1	-1	3	4	10	-1	2	-1	-1
2	6	12	8	9	0	5	3	4	1	10	-1	-1	-1	-1



ki					
0					
1					
2					
3					
4					
5					
6					
7					
8					
9					
10					
11					
12					

0 1 2 3 4 5 6 7 8 9 10 11 12 13 14

Vä	a	l	S
	Ol	m	
		m	

3	15	1	17	7	9	2	Ø	6	5	16	Ø	2	-1	-1
5	9	0	7	8	6	1	-1	3	4	10	-1	2	-1	-1
2	6	12	8	9	0	5	3	4	1	10	-1	-1	-1	-1

Update Pseudo Code

```
# Updates the value of a key in the binary
# heap. The key index must exist and the
# value must not be null.
function update(ki, value):
    i = pm[ki]
    values[ki] = value
    sink(i)
    swim(i)
```

Decrease and Increase key

In many applications (e.g Dijkstra's and Prims
algorithm) it is often useful to only update a
 given key to make its value either always
smaller (or larger). In the event that a worse
value is given the value in the IPQ should not
 be updated.

In such situations it is useful to define a more restrictive form of update operation we call increaseKey(ki, v) and decreaseKey(ki, v)

Increase/Decrease key pseudo code

```
# For both these functions assume ki and value
# are valid inputs and we are dealing with a
# min indexed binary heap.
function decreaseKey(ki, value):
  if less(value, values[ki]):
    values[ki] = value
    swim(pm[ki])
function increaseKey(ki, value):
  if less(values[ki], value):
```

values[ki] = value

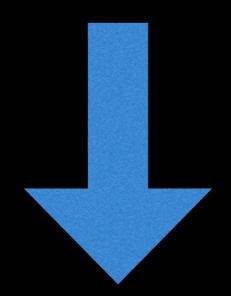
sink(pm[ki])

Source Code Link

Implementation source code and slides can be found at the following link:

github.com/williamfiset/data-structures

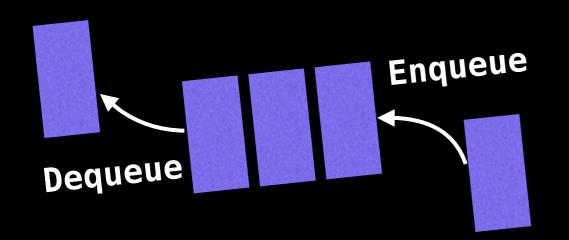
Link in the description:

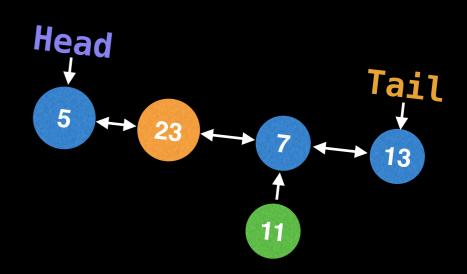


Next Video: IPQ source code

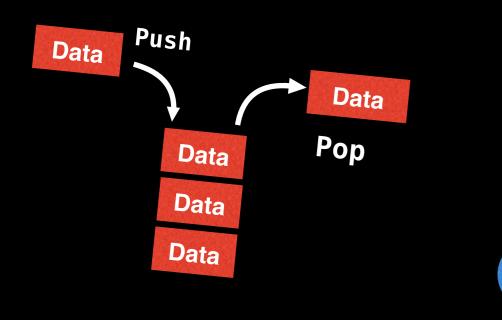
Indexed Priority Queues (UPDATED)

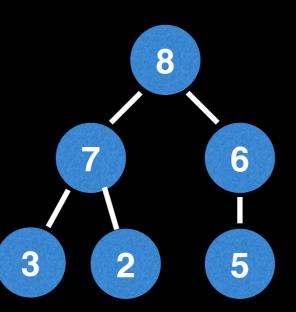
Indexed PQ Source Code

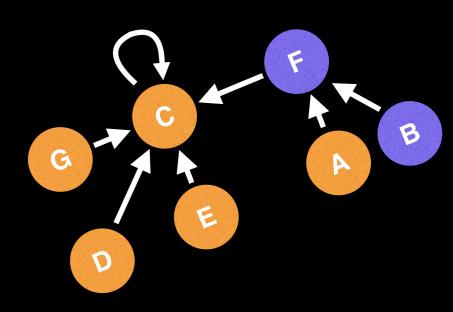




Data Structure Video Series







Indexed Priority Queue Source Code

William Fiset

Source Code Link

Implementation source code and slides can be found at the following link:

github.com/williamfiset/data-structures

Link in the description:

