

Week 2

Higher-Order Functions

Functions are first-class citizens in functional languages

It means, a function can be passed as a parameter and returned as a result

Functions that can do that are called higher order functions

```
def sum(f: Int => Int, a: Int, b: Int): Int =  
  if (a > b) 0  
  else f(a) + sum(f, a+1, b)
```

```
def sumInts(a: Int, b: Int) = sum(id, a, b)  
def sumCubes(a: Int, b: Int) = sum(cube, a, b)  
def sumFactorials(a: Int, b: Int) = sum(fact, a, b)
```

and

```
def id(x: Int): Int = x  
def cube(x: Int): Int = x * x * x  
def fact(x: Int): Int = if (x == 0) 1 else fact(x-1)
```

$A \Rightarrow B$ - is the type of a function that takes an argument of type A and returns a result of type B

$\text{Int} \Rightarrow \text{Int}$ - from int to int

Anonymous functions

functions without names

$(x: \text{Int}) \Rightarrow x * x * x$
parameter body

type can be
omitted if can be
deduced

$(x: \text{Int}, y: \text{Int}) \Rightarrow x + y$
several parameters

$\text{def sumInts}(a \dots b) = \text{sum}(x \Rightarrow x, a, b)$
shorter!

$\text{def sumCubes}(\dots) = \text{sum}(x \Rightarrow x * x * x, a, b)$

Currying

Can a word ~~be~~ get passed unchanged ~~to~~
Can we make them even shorter?

$\text{def sumInts}(a \dots b) = \text{sum}(x \Rightarrow x, a, b)$

```

def sum(f: Int => Int): (Int, Int) => Int = {
  def sumF(a: Int, b: Int): Int =
    if (a > b) 0
    else f(a) + sum(a + 1, b)
  sumF
}

```

sum now a function that returns another function

the returning function knows f and does the summing

```

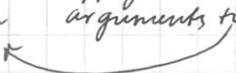
def sumInts = sum(x => x)
def sumCubes = sum(x => x * x * x)
def sumFactorials = sum(fact)

```

sumCubes(1, 10) !

↓
the same as

$\underbrace{\text{sum}(\text{cube})}_{\text{return a function}}(\underbrace{1, 10}_{\text{apply the arguments to}})$
 $= (\text{sum}(\text{cube}))(1, 10)$



and

$$\text{def sum} \left(\overbrace{f: \text{Int} \Rightarrow \text{Int}}^{\text{I parameter list}} \left(\overbrace{a: \text{Int}, b: \text{Int}}^{\text{II parameter list}} \right), \text{Int} \right) =$$

$$\text{if } (a > b) 0 \text{ else } f(a) + \text{sum}(f)(a + 1, b)$$

There's a special syntax in Scala for functions that return functions

general syntax:

- $\text{def } f(\text{args}_1) \dots (\text{args}_n) = E,$

which is equivalent to

- $\text{def } f(\text{args}_1) \dots (\text{args}_{n-1}) = \{$

$$\begin{array}{l} \text{def } g(\text{args}_n) = E \\ g \end{array}$$

$$\}$$

or $f(\text{args}_1) \dots (\text{args}_{n-1}) = (\text{args}_n \Rightarrow E)$

- $\text{def } f(\text{args}_1) \dots (\text{args}_n) = E$

\Rightarrow

$$\text{def } f = (\text{args}_1 \Rightarrow (\text{args}_2 \Rightarrow \dots (\text{args}_n \Rightarrow E) \dots))$$

This style is called currying

What's the type of

$\text{def sum}(f: \text{Int} \Rightarrow \text{Int})(a: \text{Int}, b: \text{Int}): \text{Int} = \dots$
?

Answer:

$$(\text{Int} \Rightarrow \text{Int}) \Rightarrow (\text{Int}, \text{Int}) \Rightarrow \text{Int}$$

Finding a fixed point: Example

x is a fixed point of f if

$$f(x) = x$$

For some functions we can find the fixed point by starting with an initial estimate and then by applying f again and again

$$x \quad f(x) \quad f(f(x)) \quad \dots \quad f(f(f(x))) \quad \dots$$

until it doesn't vary anymore (or the change is very small)

val tolerance = 0.0001

def isCloseEnough(x: Double, y: Double) =
abs((x-y)/x) / x < tolerance

```
def fixedPoint(f: Double => Double) (firstGuess: Double) = {  
  def iterate(guess: Double): Double = {  
    val next = f(guess)  
    if (isCloseEnough(guess, next)) next  
    else iterate(next)  
  }  
  iterate(next firstGuess)  
}
```

Sqrt(x) = the number y such that $y * y = x$
or, (if divide by y)
 $y = \frac{x}{y} \Rightarrow$

Sqrt(x) is a fixed point of the function

$$y = \frac{x}{y}$$

```
def sqrt(x: Double) =  
  fixedPoint(y => x/y)(1.0)
```

↑

but that won't converge

it oscillates 1.0 - 2.0 - 1.0 - 2.0 ...

We can prevent it from varying too much by averaging successive values

```
def sqrt(x: Double) =  
  fixedPoint( y => (x + 1/y) / 2 ) (1.0)
```

This technique of stabilizing by averaging is general enough to be a function

```
def averageDamp(f: Double => Double)(x: Double) =  
  (x + f(x)) / 2.
```

```
def sqrt(x: Double) =  
  fixedPoint(averageDamp( x => 1/y )) (1.0)
```

Summary

- Functions are essential abstractions - they allow us to introduce general methods to perform computations as
 - explicit
 - named components
- } our programming language
- These abstractions can be combined with higher-order functions to create new abstractions

Scala Syntax Summary

Extended Backus-Naur form (EBNF)

- $|$ denotes an alternative
- $[...]$ an option (0 or 1)
- $\{...\}$ a repetition (0 or more)

Type = SimpleType | FunctionType

FunctionType = SimpleType \Rightarrow ' Type |
'(' [Types] ') \Rightarrow ' Type

Simple type = Ident

Types = Type {',' Type}

a type can be

- a numeric type:

Int, Double, Byte, Short, Char, Long, Float

- Boolean

- String

- a function type

Int \Rightarrow Int,

(Int, Int) \Rightarrow Int

Expressions

$\text{Expr} = \text{Infix Expr} \mid \text{Function Expr}$
 $\mid \text{if '(' Expr ')' Expr else }^+ \text{ Expr}$

$\text{Infix Expr} = \text{Prefix Expr} \mid \text{Infix Expr Operator Infix Expr}$

$\text{Operator} = \text{ident}$

$\text{Prefix Expr} = ['+' \mid '-' \mid '!' \mid '^'] \text{ Simple Expr}$

$\text{Simple Expr} = \text{ident} \mid \text{literal} \mid \text{Simple Expr} '.' \text{ident}$
 $\mid \text{Block}$

$\text{Function Expr} = \text{Bindings} '=>' \text{Expr}$

$\text{Bindings} = \text{ident} [':' \text{ Simple Type }] \mid$
 $'(' [\text{Binding} \{ ',' \text{ Binding } \}] ')'$

$\text{Binding} = \text{ident} [':' \text{ Type }]$

$\text{Block} = \{ \{ \text{Def.} ';' \} \text{ Expr} \}$

Expression can be

- An identifier such as x , is Good Enough
- A literal 0 1.0, "abc"
- A function operator, like $\text{sqrt}(x)$
- An operation implication like $-x$, $y+x$
- A selection like math.abs
- A conditional expression $\text{if } (x < 0) -x \text{ else } x$
- A block $\{ \text{val } x = \text{math.abs}(y); x * 2 \}$
- An anonymous function, $x \Rightarrow x+1$

Definitions

Def = Fun Def | Val Def

Fun Def = def ident { '(' [Parameters] ')' }

[':' Type] '=' Expr

Val Def = val ident [':' Type] '=' Expr

Parameter = ident ':' ['='>'] Type

Parameters = Parameter { ',', Parameter }

A definition can be

- a function definition `def square (x: Int) = x * x`
- a value definition, `val y = square(2)`

A parameter can be

- a call-by-value parameter, `(x: Int)`
- a call-by-name parameter, `(x => Double)`

Functions and Data

how functions create and encapsulate data structures

Rational numbers

We want to design a package for doing rational arithmetic

A rational number is represented by 2 integers

$\frac{x}{y}$ - numerator
 - denominator

```
class Rational (x: Int, y: Int) {  
  def numer = x  
  def denom = y  
}
```

This definition introduces

- a new type - Rational
- a constructor to create elements of this type

Elements of a class type - objects

new Rational(1, 2)
Constructor

x.number
x.numer
x.denom

infix operator
"."
for
selecting
members

Rational Arithmetics:

$$\bullet \quad \frac{n_1}{d_1} + \frac{n_2}{d_2} = \frac{n_1 d_2 + n_2 d_1}{d_1 d_2}$$

$$\bullet \quad \frac{n_1}{d_1} - \frac{n_2}{d_2} = \frac{n_1 d_2 - n_2 d_1}{d_1 d_2}$$

$$\bullet \quad \frac{n_1}{d_1} \cdot \frac{n_2}{d_2} = \frac{n_1 n_2}{d_1 d_2}$$

$$\bullet \quad \frac{n_1}{d_1} \mid \frac{n_2}{d_2} \iff \frac{n_1 d_2}{d_1 n_2}$$

$$\bullet \quad \frac{n_1}{d_1} = \frac{n_2}{d_2} \iff n_1 d_2 = d_1 n_2$$

```
def addRational (r: Rational, s: Rational):  
    Rational =
```

```
    new Rational (r.numer * s.denom +  
                  s.numer * r.denom,  
                  r.denom * s.denom)
```

```
def makeString (r: Rational) =  
    r.numer + "/" + r.denom
```

But we can add these functions to the data abstraction itself

Such functions are called methods

```
class Rational (x: Int, y: Int) {
```

```
  def numer = x
```

```
  def denom = y
```

```
  def add (r: Rational) =
```

```
    new Rational (numer * r.denom +  
                  r.numer * denom,  
                  denom * r.denom)
```

```
  def mul (r: Rational) = ...
```

```
  override def toString = numer + "/" + denom
```

```
}
```

↑

declares that this method
redefines a method that
already exists

```
val x = new Rational (1, 3)
```

```
val y = new Rational (5, 7)
```

```
x.add(y).mul(7)
```

But Rational numbers should be
simplified

```
class Rational (...) {
```

```
  private def gcd (a: Int, b: Int): Int =  
    if (b == 0) a else gcd (b, a % b)
```

```
  private val g = gcd (x, y)
```

```
  def numer = x / g
```

```
  ...
```

ged and g are private members - we can access them only from inside the Rational class

The ability to choose different implementation of the data without affecting clients is called data abstraction

```
class Rational (...) {
```

...

```
    def max (that: Rational) =
```

```
        if (this.less(that)) that  
        else this
```

```
}
```

↑
reference to the current object

Preconditions, Assertions

```
class Rational (...) {
```

← to enforce a precondition

```
    require (y > 0, "denominator must be positive")
```

↑

```
        throws IllegalArgumentException  
}
```

```
val x = sqrt(y)
assert(x >= 0)
```

used to check the
code of the function

↑
Assertion Error

Constructors

a class implicitly introduces a
constructor -
this is the primary constructor.

The primary constructor:

- takes the parameters of the class
- executes all statements in the class
body

We also can declare auxiliary constructors -
they are methods called "this"

```
class Rational(...) {  
    def this(x: Int) = this(x, 1)  
    ...  
}
```

new Rational(2) → 2/1

Classes and Substitution

new $C(e_1, \dots, e_m)$

↓

e_i evaluated

new $C(v_1, \dots, v_m)$.

new $C(v_1, \dots, v_m). f(w_1, \dots, w_n)$

↓

$[\{ w_i : y_i \}] [\{ v_i : x_i \}] [\text{new } C \dots / \text{this}] \&$

↑

name-value

↑

name-value

new Rational(1, 2). numer

→ $[1/x, 2/y][][\text{new Rational}(1, 2) / \text{this}] \times$

= 1

...

Operators

it's possible to write

`r.add(s)` instead of `r.add(s)`

`r.less(s)`

And operators can be used as identifiers

An identifier can be:

- alphanumeric - starts with a letter, followed by letters and numbers
- symbolic - started with an operator symbol followed by other operator symbols
- `'_'` counts as a letter
- alphanumeric identifiers can also end in an underscore, followed by some operator symbol

`x1 * e?%& vector_ ++ counter_ =`

```
class Rational (...) {
```

```
    ...
```

```
    def + (r: Rational) = ...
```

```
    def - (r: Rational) = ...
```

```
    def * ...
```

```
} ...
```

val x = new Rational..

val y = new Rational..

$x * x + y * y$

\Downarrow

$(x * x) + (y * y)$

~~Precedence~~ Precedence of an operator is determined by its first character

The priorities:

- (all letters)
- |
- ^
- &
- < >
- = !
- :
- + -
- * / %
- (all other special characters)

$a + b \wedge ? c \wedge ? d \text{ less } a ==> b | c$

\Downarrow

$((a + b) \wedge ? (c \wedge ? d)) \text{ less } ((a ==> b) | c)$

Assignment - Sets.

- type alias

type Set Int => Boolean

↑
takes int

$(x : \text{Int}) \Rightarrow x < 0$

a set with negative integers

just a function that
returns true if
element x is in the set.

def contains (s: Set, elem: Int): Boolean =
s(elem)

- forall (s: Set, $\underbrace{p: \text{Int} \Rightarrow \text{Boolean}}_{\text{predicate}})$: Boolean -

universal quantifier - return true
if all elements of the set s
the predicate p is true

- exists (s: Set, $p: \text{Int} \Rightarrow \text{Boolean}$): Boolean

existential quantifier - returns true
if at least one element matches
the predicate.

Implementation in terms of forall:

$\neg \text{forall}(s, p)$ and $\neg \text{forall}(s, \bar{p})$

\bar{p} = inverse predicate

• map ($s: \text{Set}, f: \text{Int} \Rightarrow \text{Int}$): Set

transforms a set S by applying function f to each element

$$S \xrightarrow{f} Y$$

element $y \in Y$ iff there exists $x \in S$

such that $f(x) = y$ ↑
original set

(so if we are able to ~~find~~ apply a predicate "exists $f(x) = y$ " to S and get true-
 y is in Y).