- 1 entry per page in virtual address space

even; I the address space is intised

- Entry contains frame # + bits that hell

us; I the page is accessable (in

memory or not, dirty etc.).

large enough to select a frame.

Virtual memory 3

Virtual memory 3

Page size of entry (95, les = 325, ls)

+ypical

4-Kb

232

= 2 32 2 12 * 4 byles = 2 20 * 4 = 4 MByle

Problem; space bable is large even; A process uses very 1846 memory

so our 2 Mb appressos a page table of 4 Mbyles

What about 64 bit address agrace?



= 3.6 x 10 16 byles => bit more Hance have (Smithings)

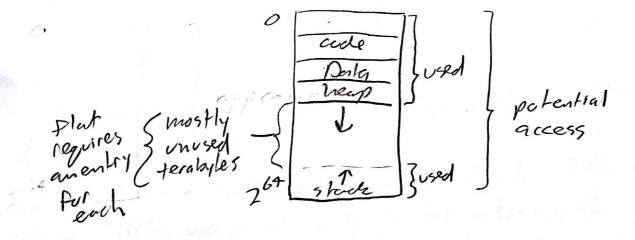
Multi level page table

- reduce the FIZE of Plat page bubles

- Size propertional to how much mem application can address, not how much its using!

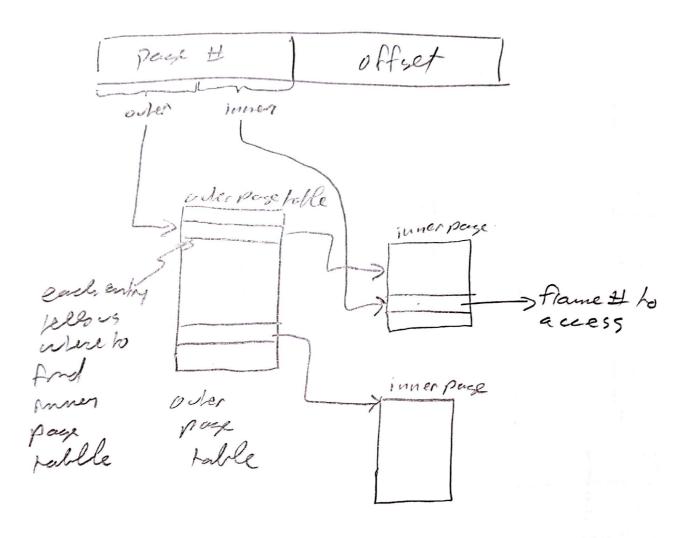
32 bit 2 4-8 Mbyles 64 bit 2 hoobig

how to solve



MolH level

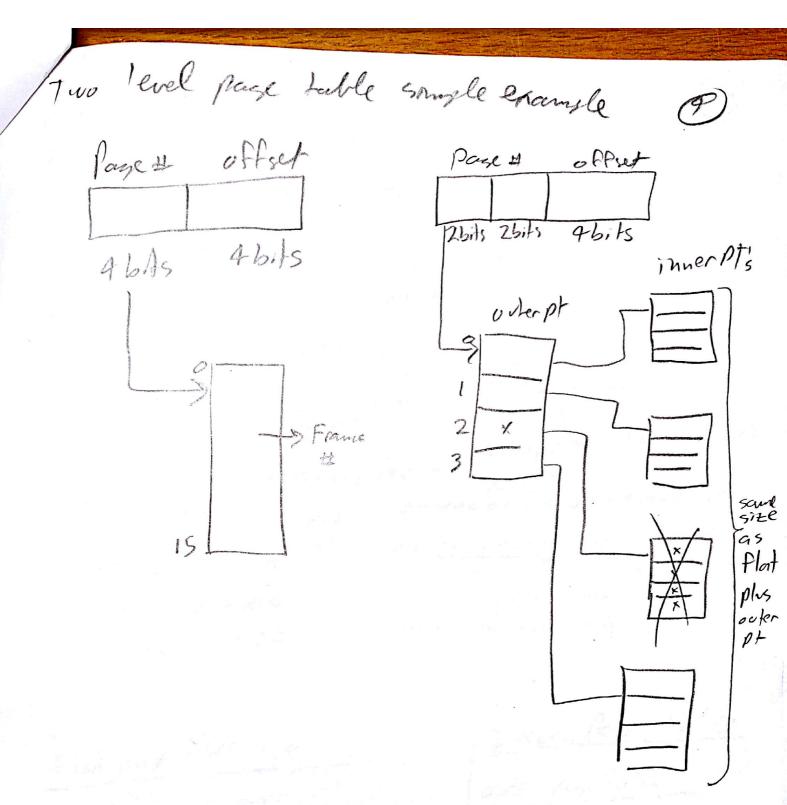
use bits to ordex tables avoid entries for unused entries



outer page # tello us which of the small muer page 4's to use.

page tables we need to look for a frame #

seems like are garned nothing have just as many page bables (plus the order page bable) as flat bable



Savings? Luben outer page table points
to miner page tables that have no entries
don't need to have those miner page trables
trave! outer page table is young miner page tables

Where are savings? I herel page table size 32 bit address space 212 4 Kb per page 1029 entry outer page table Pt enly 4 byles offsel 10 bits 10 bits 12 6/15 program uses memery from (bin bin odoooooo) outer juner offset 0,000 10000 1111 1111 1111 1111 0000 0x FFFF 0000 Fin her morning - " oxftft ffft 2 level Page bable flat page publisize 232 = 220 x4 (+MB) over page habe 2 10 x 9 = 4 Khyles inver page bulles? for lower address space outer page is a lucys con orter= 2 10 x 4] (2Kb)
inner 2 x 210 x 4 } (2Kb)

this is when multi level peop lables are ver!