Overview of SOIL Initiative's GC Proposal

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High-Level Patterns

Our solution:
approximate
subtyping with
abstractions based
on implementation's
architecture

Memory Architecture

"Scheme"

Subtyping based on explicitly stated relationships

Memory Layout

Complex
Description of
Data/Functions

Subtyping based on compatibility of aligned elements

Problem: Needs to be type-checked; known expressiveness and efficiency issues

Key Hypothesis: Archictectures can be described as a coarse abstraction of precise memory layouts

Low-Level Representations

- Problem: wasm module has no control over low-level representation of pointers and casting mechanisms
 - Varying hardware architectures, e.g. 32 vs. 64 bit
 - Varying GC implementations: e.g. summaries vs. bit flags
- Our solution: provide enough information for engine to specialize low-level representation for individual modules
 - E.g. immutable i32 w/o identity can be packed on 64-bit machine
 - E.g. cast with bit flags vs. summary equality vs. array lookup

Modularity and Abstraction

- Key Hypothesis: module coordination is intentional/explicit
 - e.g. references from unrelated modules need not be castable
 - Java modules do not know how to cast Ocaml refs and vice-versa
- Goal: surface-level equivalence implies wasm equivalence
 - e.g. private Java fields should be invisible
 - e.g. public Java fields should be reorderable
 - e.g. names of Java fields should be irrelevant (ignoring reflection)
 - How to achieve this?

Ideas being explored

- Support for engines with simple memory management
 - Incorporate acyclicity for reference counting
 - Invalidatable references for resource release w/o even ref counting
- Async coop between module-manage lin. mem. and host gc
- Eliminating superfluous casts (existential/universal types)
- Interior pointers and nested structures
- Stack-allocated memory (needs to align with heap-alloc mem)
- Complex/cross-module initialization of structures