

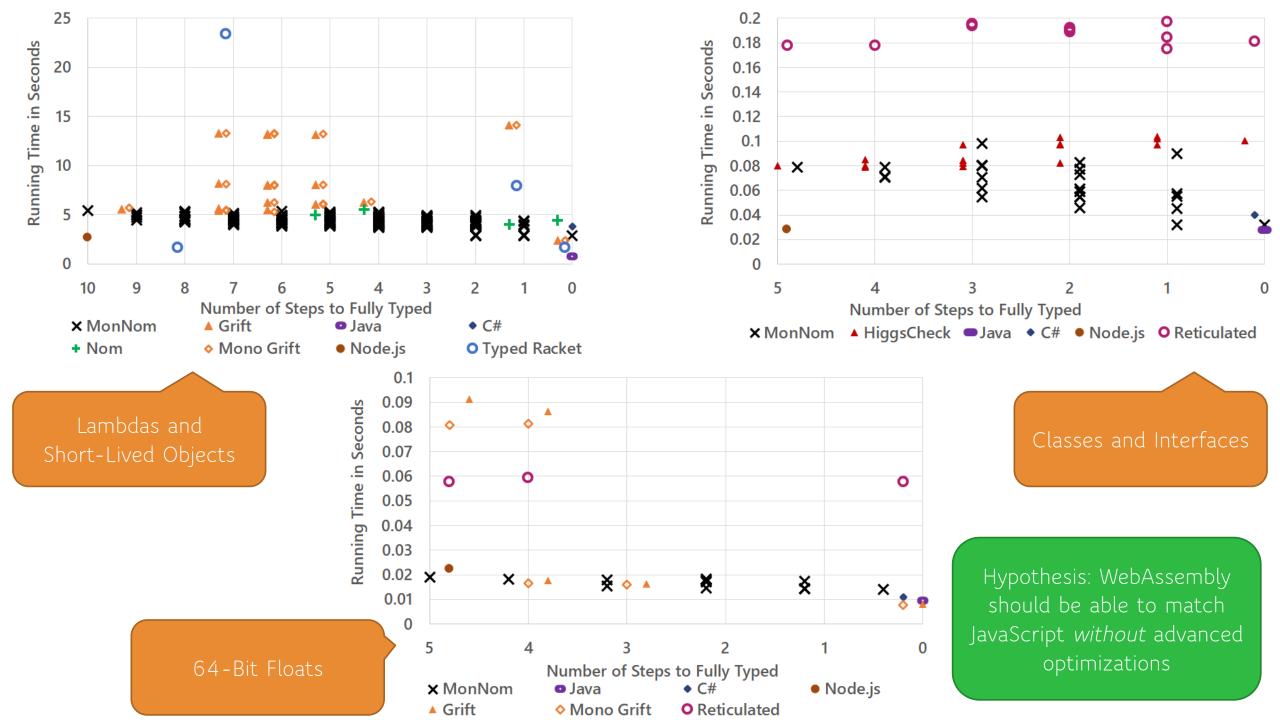
Experiments with AOT Compilation

ROSS TATE

main.mn function main() { twice(fooey()); twice.mn function twice(f : dyn) : dyn { Α return f.foo(f.foo(0)); function twice(f : dyn) : dyn { Foo F = f; return F.foo(F.foo(0)); function twice(f : Foo) : int { return f.foo(f.foo(0)); foo.mn A interface Foo {

foo(i : int) : int;

```
fooey.mn
     function fooey() : dyn {
         dyn f = new {};
A
        f.foo = \lambda(i : dyn) : dyn \{ return i+1; \};
         return f;
     function fooey() : dyn {
         return new {
B
             foo(i : dyn) : dyn { return i+1; };
         };
     function fooey() : dyn {
         return new {
             foo(i : int) : int { return i+1; };
         };
     class Fooey() implements Foo {
         foo(i : int) : int { return i+1; };
     function fooey() : dyn { return new Fooey(); }
     class Fooey() implements Foo {
E
         foo(i : int) : int { return i+1; };
     function fooey() : Foo { return new Fooey(); }
```



Caching

Observation: The L1 cache is amazing

- Typically no performance benefit gained from eliminating a load from the L1 cache
 - Even chained loads

Hypothesis: Moduleinstance data should be quickly accessible if in L1

Call Tags

We implemented interface methods and lambdas using call tags

- Used call-tag switching to handle conflicts between interface methods
- Used defaults to support typed-untyped interop (e.g. casting and representation conversion)

Class methods implemented using standard typed code pointers

Experiment:

- Replace all interfaces with classes in benchmark where each interface is implemented by only one class
- Effect on implementation: replaces uses of call tags with uses of typed code pointers
- Impact on performance: None

Hypothesis: Typed-function references add complexity without adding performance or expressiveness benefits

Hypothesis: Call tags add both performance and expressiveness benefits