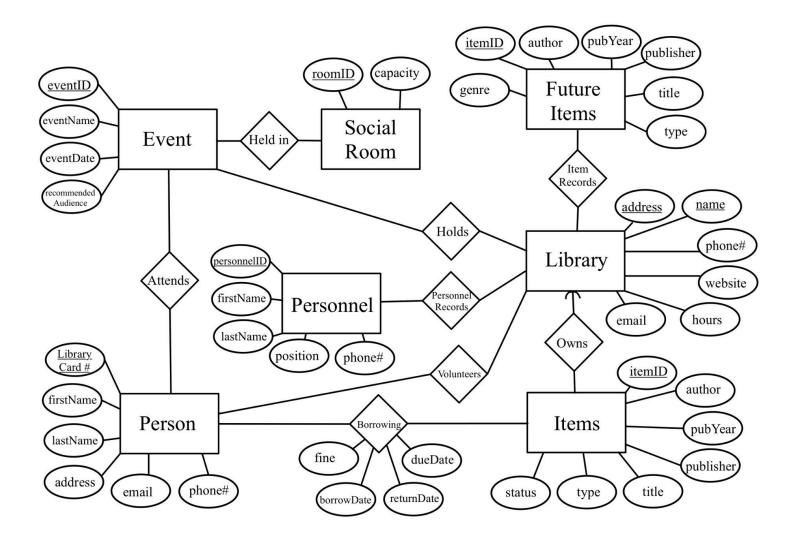
Step (2): Project Specifications

- 1. Person = {libraryCardNumber, firstName, lastName, email, address, phoneNumber}
- 2. Library = {libraryName, address, phoneNumber, email, website, hours}
- 3. Item = {itemID, title, type, publicationYear, authorName, publisherName, status}
- 4. futureItem= {futureItemID, title, type, publicationYear, authorName, publisherName}
- 5. Event = {eventID, eventName, eventDate, recommendedAudience}
- 6. Personnel = {personnelID , firstName, lastName, position, phoneNumber}
- 7. SocialRoom = {roomID , capacity }
- 8. Borrowing = {<u>libraryCardNumber</u>^{FK-Person}, <u>itemID</u>^{FK-Item}, dueDate, returnDate, fine}
- 9. heldIn = {eventID^{FK-Event}, roomID^{FK-SocialRoom}}
- 10. Holds = {<u>libraryName</u>^{FK-Library}, <u>libraryAddress</u>^{FK-Library}, <u>eventID</u>^{FK-Event}}
- 11. Attends = {<u>libraryCardNumber</u> FK-Person, <u>eventID</u>FK-Event}
- 12. Owns = {<u>itemID</u>^{FK-Item}, <u>libraryName</u>^{FK-Library}, <u>libraryAddress</u> (<u>libraryAddress</u>)
- 13. personnelRecords = { $\underline{libraryName}^{FK-Library}$, $\underline{address}^{FK-Library}$, $\underline{personnelID}^{EK-Personnel}$ }
- 14. itemRecords = {libraryName^{FK-Library}, address^{FK-Library}, itemID^{FK-futureItems}}
- 15. volunteers = {\frac{\libraryName}{FK-Library}, \frac{address}{FK-Library}, \frac{\libraryCardNumber}{\libraryCardNumber} \frac{FK-Person}{\libraryCardNumber}}
- 16. $heldIn = \{eventID^{FK-Event}, roomID^{FK-SocialRoom}\}$
- 17. Holds = {libraryName^{FK-Library}, libraryAddress^{FK-Library}, eventID^{FK-Event}}
- 18. Attends = { <u>libraryCardNumber</u> FK-Person, <u>eventID</u> FK-Event }
- 19. Owns = {<u>itemID</u>^{EK-ltem}, <u>libraryName</u>^{EK-Library}, <u>libraryAddress</u>^{EK-Library}}
- 20. personnelRecords = {libraryName^{FK-Library}, address^{FK-Library}, personnelID^{FK-Personnel}}
- 21. itemRecords = {libraryName^{FK-Library}, address^{FK-Library}, itemID^{FK-futureItems}}

Step (3): E/R Diagrams



Step (4): Does your design allow anomalies?

To prove our relation is in BCNF, we must show that for every non-trivial functional dependency $X \to Y$ in the relation, X is a key:

1. libraryCardNumber → firstName, lastName, email, address, phoneNumber

Since every person is assigned a unique library card number, libraryCardNumber is a key for the Person relation, as no two tuples can have the same libraryCardNumber.

2. libraryName, address → phoneNumber, email, website, hours

Since no two libraries share the same name and address simultaneously, {libraryName, address} is a key for the Library relation, as no two tuples can share the same libraryName and address attributes.

3. itemID → title, type, publicationYear, authorName, publisherName, status

Since each library item is assigned a unique ID, itemID is a key for the Item relation, as no two tuples can have the same itemID value.

4. futureItemID → title, type, publicationYear, authorName, publisherName

Since each future library item is assigned a unique ID, futureItemID is a key for the futureItem relation, as no two tuples can have the same futureItemID value.

5. eventID → eventName, eventDate, recommendedAudience

Since each event is assigned a unique eventID, eventID is a key for the Event relation, as no two tuples can have the same eventID value.

6. personnelID → firstName, lastName, position, phoneNumber

Since each personnel member has a unique personnelID, personnelID is a key for the Personnel relation, as no two tuples can have the same personnelID value.

7. $roomID \rightarrow capacity$

Since each room has a unique roomID, roomID is a key for the Personnel relation, as no two tuples can have the same roomID value.

8. libraryCardNumber^{FK-Person}, itemID^{FK-Item} \rightarrow dueDate, returnDate, fine

Since no two people can have the same libraryCardNumber, and no two items have the same itemID, {libraryCardNumber^{FK-Person}, itemID^{FK-Item}} is a key for the Borrowing relation, as no two tuples can have the values for {libraryCardNumber^{FK-Person}, itemID^{FK-Item}}.

Since all our FDs are in BCNF, and there are no bad FDs in the relation as X is a key for every functional dependency $X \rightarrow Y$ in our relation, our design does not allow for anomalies.