# Machine Independent Assembler Features

Literal, Symbol, Expression

### Literals

#### Motivation

- It is convenient if a programmer can write the value of a constant operand as a part of the instruction that uses it.
- This avoids having to define the constant elsewhere in the program and make up a label for it.
- Such an operand is called a literal because the value is stated literally in the instruction.

# Literal Example (1)

5	COPY	START	0	COPY FILE FROM INPUT TO OUTPUT
10	FIRST	STL	RETADR	SAVE RETURN ADDRESS
13		LDB	#LENGTH	ESTABLISH BASE REGISTER
14		BASE	LENGTH	
15	CLOOP	+JSUB	RDREC	READ INPUT RECORD
20		LDA	LENGTH	TEST FOR EOF (LENGTH = 0)
25		COMP	#0	
30		JEQ	ENDFIL	EXIT IF EOF FOUND
35		+JSUB	WRREC	WRITE OUTPUT RECORD
40		J	CLOOP	LOOP
45	ENDFIL	LDA	=C'EOF'	INSERT END OF FILE MARKER
50		STA	BUFFFR	Use = to represent a literal
55		LDA	#3	SET LENGTH = 3
60		STA	LENGTH	
65		+JSUB	WRREC	WRITE EOF
70		J	@RETADR	RETURN TO CALLER
93		LTORG		
95	RETADR	RESW	1	
100	LENGTH	RESW	1	LENGTH OF RECORD
105	BUFFER	RESB	4096	4096-BYTE BUFFER AREA
106	BUFEND	EQU	*	
107	MAXLEN	EQU	BUFEND-BUFFER	MAXIMUM RECORD LENGTH

# Original Program (1)

_	0000	CODY	CITA DIT	0	
5	0000	COPY	START	0	arin replace Length of the and
10	0000	FIRST	STL	RETADR	17202D
12	0003		LDB	#LENGTH	69202D
13			BASE	LENGTH	
15	0006	CLOOP	+JSUB	RDREC	4B101036
20	000A		LDA	LENGTH	032026
25	000D		COMP	#0	290000
30	0010		JEQ	ENDFIL	332007
35	0013		+JSUB	WRREC	4B10105D
40	0017		J	CLOOP	3F2FEC
45	001A	ENDFIL	LDA	EOF	032010
50	001D		STA	BUFFER	0F2016
55	0020		LDA	#3	010003
60	0023		STA	LENGTH	0F200D
65	0026		+JSUB	WRREC	4B10105D
70	002A	oge deen ort	J	@RETADR	3E2003
80	002D	EOF	BYTE	C'EOF'	454F46
95	0030	RETADR	RESW	1	
100	0033	LENGTH	RESW	1	
105	0036	BUFFER	RESB	4096	
110					

## Literal Program (1)'s Object Code

5	0000	COPY	START	0	
10	0000	FIRST	STL	RETADR	17202D
13	0003		LDB	#LENGTH	69202D
14			BASE	LENGTH	
15	0006	CLOOP	+JSUB	RDREC	4B101036
20	000A		LDA	LENGTH	032026
25	000D		COMP	#0	290000
30	0010		JEQ	ENDFIL	332007
35	0013		+JSUB	WRREC	4B10105D
40	0017		J	CLOOP	3F2FEC
45	001A	ENDFIL	LDA	=C'EOF'	032010
50	001D		STA	BUFFER	0F2016
55	0020		LDA	#3	010003
60	0023		STA	LENGTH	0F200D
65	0026		+JSUB	WRREC	4B10105D
70	002A	and the second second	J	GRETADR	3E2003
93			LTORG		
	002D	*	=C'EOF'		454F46
95	0030	RETADR	RESW	1	The Gasterian tenerge to

Notice that the object code is the same as the previous one.

# Literal Example (2)

195		arm of the	TARREST TO LIDITE	DECORD FROM DITEFFD
200	a editesa l	SUBROU'I	INE TO WRIT	E RECORD FROM BUFFER
205				COLD WILD
210	WRREC	CLEAR	X	CLEAR LOOP COUNTER
212		LDT	LENGTH	
215	WLOOP	TD	=X'05'	TEST OUTPUT DEVICE
220		JEQ	WLOOP	LOOP UNTIL READY
225		LDCH	BUFFER, X	GET CHARACTER FROM BUFFER
230		WD (	=X'05'	WRITE CHARACTER
235		TIXR	т	LOOP UNTIL ALL CHARACTERS
240		JLT	WLOOP	HAVE BEEN WRITTEN
245		RSUB		RETURN TO CALLER
255		END	FIRST	

# Original Program (2)

205						
210	105D	WRREC	CLEAR	X	B410	
212	105F		LDT	LENGTH	774000	
215	1062	WLOOP	TD	OUTPUT	E32011	
220	1065		JEQ	WLOOP	332FFA	
225	1068		LDCH	BUFFER, X	53C003	
230	106B		WD	OUTPUT	DF2008	
235	106E		TIXR	Т	B850	
240	1070		JLT	WLOOP	3B2FEF	
245	1073	HILL CALL THE	RSUB	A 44 0300009 93	4F0000	
250	1076	OUTPUT	BYTE	X'05'	05	
255		n mendery	END	FTRST		

# Literal Program (2)'s Object Code

205		0.0 • C 0.0 0.0 0.0 0.0			
210	105D	WRREC	CLEAR	X	B410
212	105F		LDT	LENGTH	774000
215	1062	WLOOP	TD	=X'05'	E32011
220	1065		JEQ	WLOOP	332FFA
225	1068	DANS TORRES	LDCH	BUFFER, X	53C003
230	106B		WD	=X'05'	DF2008
235	106E		TIXR	T	В850
240	1070		JLT	WLOOP	3B2FEF
245	1073		RSUB		4F0000
255			END	FIRST	
	1076	*	=X'05'		05

Notice that the object code is the same as the previous one.

# Difference between Literal and Immediate Operand

#### Immediate addressing

The operand value is assembled as part of the machine instruction.

#### • Literal

- The assembler generates the specified value as a constant at some other memory location.
- The address of this generated constant is used as the target address for the machine instruction.
- The effect of using a literal is exactly the same as if the programming had defined the constant explicitly and used the label assigned to the constant as the instruction operand.

### Literal Pool

- All of the literal operands used in the program are gathered together into one or more literal pools.
- Normally literals are placed into a pool at the end of the program.
- Sometimes, it is desirable to place literals into a pool at some other location in the object program.
  - LTORG directive is introduced for this purpose.
  - When the assembler encounters a LTORG, it creates a pool that contains all of the literals used since the previous LTORG.

## LTORG Example

- If we do not use a LTORG on line 93, the literal =C'EOF' will be placed at the end of the program.
- This operand will then begin at address 1073 and be too far away from the instruction that uses it. PC-relative addressing mode cannot be used.
- LTORG thus is used when we want to keep the literal operands close to the instruction that uses it.

## **Duplicate Literals**

- Duplicate literals the same literal used in more than one place in the program
- For duplicate literals, we should store only one copy of the specified data value to save space.
- Most assembler can recognize duplicate literals.
  - E.g., There are two uses of =X'05' on lines 215 and 230 respectively.
  - However, just one copy is generated on line 1076.

## Recognize Duplicate Literals

- Comparing the character string defining them
  - For example, =X'05' and =X'05'
- Comparing the generated data value
  - For example, =C'EOF' and =X'454F46'
  - More intelligent
  - However, usually the benefit is not great enough to justify the added complexity.

# The Problem with String-Defining Literals

- We should be careful about the literal whose value depends on their locations in the program.
- E.g., (\* usually denotes the current value of the location counter)

```
BASE *
LDB =*
```

• If \* appears on line 13, it would specify 0003. If it appears on line 55, it would specify an operand with value 0020.

## Literal Processing (1)

- Need a literal table LITTAB. For each literal used, the table contains:
  - Literal name
  - The operand value and length
  - The address assigned to the operand when it is placed in a literal pool.

## Literal Processing (2)

#### • Pass 1:

- When encountering an literal, try to find it in LITTAB.
   If it can be found, do nothing.
- Otherwise, the literal is added to the LITTAB
  - Name, operand value and length can be entered now
  - Only the address field is left blank.
- When encountering a LTORG or the end of the program, the assembler makes a scan of LITTAB. At this time, each literal currently in this table is assigned an address. The location counter then should be updated to reflect the number of bytes occupied by each literal.

## Literal Processing (3)

#### • Pass 2:

- For each literal operand encountered, we search it in LITTAB and find its assigned address.
- The data values specified by the literals in each literal pool are inserted at the appropriate places in the object program exactly as if these values had been generated by BYTE or WORD statements.
- If a literal represents an address in the program (e.g., a location counter value), the assembler must also generate the Modification record.

## Symbols

- We can use the EQU directive to define a symbol's value.
  - So far, the only symbols defined in a program are labels, whose values are the addresses assigned to them.
  - E.g., MAXLEN EQU 4096.
- The value assigned to a symbol may be a constant, or any expression involving constants and previously defined symbols.

## Usages of Symbols (1)

Establish symbolic names to improve readability in place of numeric values.

```
E.g., +LDT #4096 can be changed to:

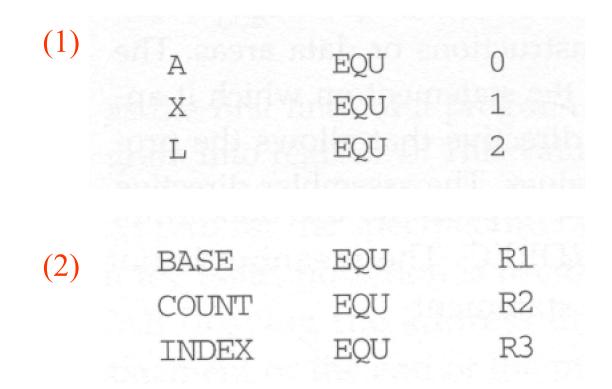
MAXLEN EQU 4096

+LDT #MAXLEN
```

- When the assembler encounters the EQU statement, it enters MAXLEN into SYMTAB (with value 4096).
- During assembly of the LDT instruction, the assembler searches SYMTAB for the symbol MAXLEN, using its value as the operand in the instruction.

## Usages of Symbols (2)

• Define mnemonic names for registers:



## No Forward Reference Allowed

- For EQU, all symbols used on the right hand side of the statement must have been defined previously in the program.
- This is because in the two-pass assembler, we require that all symbols must be defined in pass 1.

ALPHA	RESW	1	Allowed
BETA	EQU	ALPHA	
BETA	EQU	ALPHA	Not allowed
ALPHA	RESW	1	

### **ORG**

- is a assembler directive
- allows the assembler to reset PC to values
  - Syntax: ORG value
- when ORG is encountered, the assembler resets its LOCCTR to the specified value.
- ORG will affect the values of all labels defined until next ORG.
- if previous value of LOCCTR can be automatically remembered, we can return to normal use of LOCCTR by simply writing ORG.
- while using ORG, forward reference is not allowed.

## **ORG**

#### Example: using ORG

If ORG statements are used

STAB	RESB	1100
	ORG	STAB Set LOCCTR to STAB
SYMBOL	RESB	6
VALUE	RESW	1
FLAGS	RESB	2 Size of each field
	ORG	STAB+1100 ← Restore LOCCTR

We can fetch the VALUE field by

LDA VALUE,X X = 0, 11, 22, ... for each entry

## Expression

• So far, when we define the value of a symbol or label, only one term is used.

```
– E.g., 106 BUFEND EQU *
```

• Actually, we can also use expressions which contains many terms to define a symbol's value.

<u> </u>	E.g., 107 M	IAXLEN	EQU	BUFEND - BU	FFER
	002D	*	=C'EOF'		454F46
95	0030	RETADR	RESW	1	
100	0033	LENGTH	RESW	1	
105	0036	BUFFER	RESB	4096	
106	1036	BUFEND	EQU	*	
107	1000	MAXLEN	EQU	BUFEND-BUFFER	
110		20,227,200,200			

### Relative v.s. Absolute Value

- Generally, +, -, \*, /, operations are allowed to be used in an expression.
- Division is defined to produce an integer.
- Regarding program relocation, a symbol's value can be classified as
  - Relative
    - Its value is relative to the beginning of the object program, and thus its value is dependent of program location.
    - E.g., labels or reference to location counter (\*)
  - Absolute
    - Its value is independent of program location
    - E.g., a constant.

## Relative v.s. Absolute Expression

- Depending on the type of value they produce, expressions are classified as:
  - Absolute
    - An expression that contains only absolute terms, or
    - An expression that contains relative terms but the relative terms occur in pairs and the terms in each pair have opposite signs. (/ and \* operations are not allowed)

#### Relative

 An expression in which all relative terms except one can be paired and the remaining unpaired term must have a positive sign. (/ and \* operations are not allowed)

## Absolute Expression Example

- 107 MAXLEN EQU BUFEND BUFFER
- Although BUFEND and BUFFER are relative terms (because their values will change when the program is loaded into a different place in memory), the expression (BUFEND BUFFER) is an absolute expression.
- Why? the value of this expression is 0x1000, which is the same no matter where this program is loaded.
- Why? Because BUFEND and BUFFER can be represented as (startingaddr + x) and (startingaddr + y), BUFEND BUFFER becomes (x y), which is a constant.

## Illegal Expressions

- BUFEND + BUFFER
- 100 BUFFER
- 3 \* BUFFER
- These expressions represent neither absolute nor location within the program
- Therefore, these expression are considered illegal.

## Enhanced Symbol Table

- To determine the type of an expression, we must keep track of the types of all symbols defined in the program,
- Therefore, we need a flag in the symbol table to indicate type of value (absolute or relative) in addition to the value itself.

Symbol	Туре	Value
RETADR	R	0030
BUFFER	R	0036
BUFEND	R	1036
MAXLEN	A	1000