

**KCG College of Technology, Chennai-96**  
**Computer Science and Engineering**

**CS 2304 SYSTEM SOFTWARE                      III Sem CSE**  
**QUESTION BANK   -   UNIT-II   ASSEMBLERS**

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**1) Why an Assembly Language is needed?**

Programming in machine code, by supplying the computer with the numbers of the operations it must perform, can be quite a burden, because for every operation the corresponding number must be looked up or remembered. Looking up all numbers takes a lot of time, and mis-remembering a number may introduce [computer bugs](#).

So Assembly Languages are evolved which contains mnemonic instructions corresponding to the Machine codes using which the program can be written easily.

Therefore a set of mnemonics was devised. Each number was represented by an alphabetic code. So instead of entering the number corresponding to addition to add two numbers one can enter "add".

Although mnemonics differ between different [CPU designs](#) some are common, for instance: "sub" (subtract), "div" (divide), "add" (add) and "mul" (multiply).

**2) What is an Assembler?**

An assembler is a program that accepts an *assembly language program* as input and produces its *machine language equivalent* along with *information for the loader* (An Assembler translates a program written in an assembly language to its machine language equivalent)

**3) Explain the terms a)Label,b)Opcode,c)Operand,and d)Comment  
(What is the format in which the assembly language program is written?).**

➤ **Label field.**

- The label is a symbolic name that represents the memory address of an executable statement or a variable.

➤ **Opcode/directive fields.**

- The opcode (e.g. operation code) specifies the symbolic name for a machine instruction.
- The directive specifies commands to the assembler about the way to assemble the program.

➤ **Operand field.**

- The operand specifies the data that is needed by a statement.

➤ **Comment field.**

- The comment provides clear explanation for a statement.

#### 4) What are the basic functions of an assembler?

##### Functions of a Basic Assembler

- Convert *mnemonic operation codes* to their *machine language equivalents*  
E.g. STL -> 14 (line 10)
- Convert *symbolic operands* to their equivalent *machine addresses*  
E.g. RETADR -> 1033 (line 10)
- Build the machine instructions in the proper format
- Convert the *data constants* to *internal machine representations*  
E.g. EOF -> 454F46 (line 80)
- Write the *object program* and the *assembly listing*

#### 5) What are assembler Directives?

Assembler directives are **Pseudo-instructions that are not** translated into machine instructions and they provide instructions to the assembler itself.

- **The SIC assembler directives.**
  - START
    - Specification of the name and start address of the program.
  - END
    - Indication of the end of the program and optionally the address of the first executable instruction.
  - BYTE
    - Declaration of character or string constants.
  - WORD
    - Declaration of integer constants.
  - RESB
    - Declaration of character variables or arrays.
  - RESW
    - Declaration of integer variables or arrays.

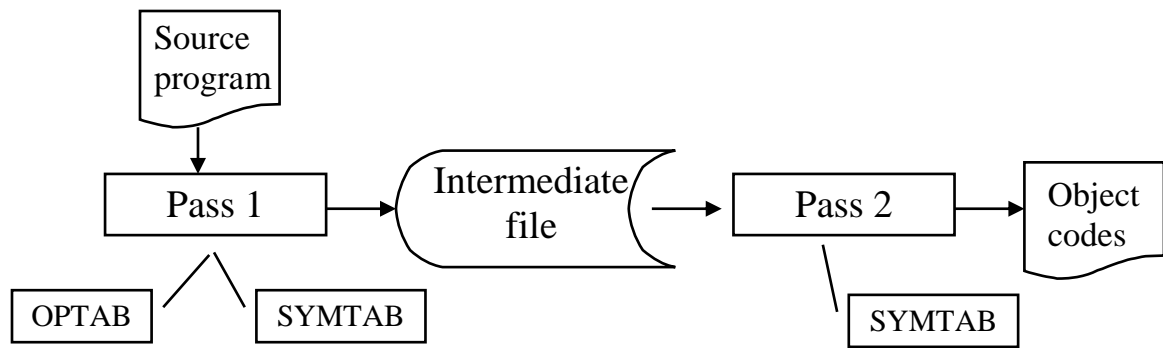
#### 6) What are the functions of two pass assembler?

##### Functions of Two Pass Assembler

- **Pass 1 - define symbols (assign addresses)**
  - Assign addresses to all statements in the program
  - Save the values assigned to all labels for use in Pass 2
  - Process some assembler directives
- **Pass 2 - assemble instructions and generate object program**

Assemble instructions

  - Generate data values defined by BYTE, WORD, etc.
  - Process the assembler directives not done in Pass 1
  - Write the object program and the assembly listing



**7) What is the format of the Object Program generated by the Assembler?**

Contains 3 types of records:

**Header record:**

Col. 1 H

Col. 2-7 Program name

Col. 8-13 Starting address (hex)

Col. 14-19 Length of object program in bytes (hex)

**Text record**

Col.1 T

Col.2-7 Starting address in this record (hex)

Col. 8-9 Length of object code in this record in bytes (hex)

Col. 10-69 Object code (hex) (2 columns per byte)

**End record**

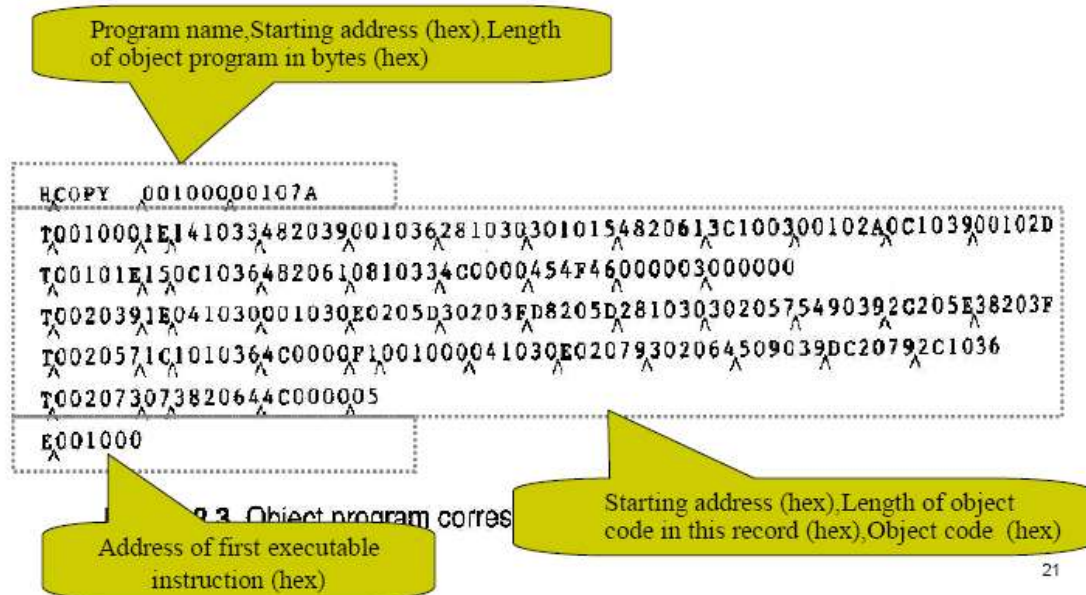
Col.1 E

Col.2~7 Address of first executable instruction (hex)

(END program\_name)

**8) Give an example of object program generated by an Assembler.**

## Object Program for Fig 2.2 (Fig 2.3)



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### 9) What is forward reference?

*Forward reference* is a reference to a label that is defined later in the program.

#### Example

10 STL RETADR

- o **RETA****DR** is not yet defined when we encounter **STL** instruction
- o So it is called *forward reference*

### 10) Give an example of Assembly language along with the objectcode generated.

Line	Loc	Source statement			Object code
5	1000	COPY	START	1000	
10	1000	FIRST	STL	RETA	DR 141033
15	1003	CLOOP	JSUB	RDREC	482039
20	1006		LDA	LENGTH	001036
25	1009		COMP	ZERO	281030
30	100C		JEQ	ENDFIL	301015
35	100F		JSUB	WRREC	482061
40	1012		J	CLOOP	3C1003
45	1015	ENDFIL	LDA	EOF	00102A
50	1018		STA	BUFFER	0C1039
55	101B		LDA	THREE	00102D
60	101E		STA	LENGTH	0C1036
65	1021		JSUB	WRREC	482061
70	1024		LDL	RETA	DR 081033

75	1027		RSUB		4C0000
80	102A	EOF	BYTE	C'EOF'	454F46
85	102D	THREE	WORD	3	000003
90	1030	ZERO	WORD	0	000000
95	1033	RETADR	RESW	1	
100	1036	LENGTH	RESW	1	
105	1039	BUFFER	RESB	4096	
110		.			
115		.	SUBROUTINE TO READ RECORD INTO BUFFER		
120		.			
125	2039	RDREC	LDX	ZERO	041030
130	203C		LDA	ZERO	001030
135	203F	RLOOP	TD	INPUT	E0205D
140	2042		JEQ	RLOOP	30203D
145	2045		RD	INPUT	D8205D
150	2048		COMP	ZERO	281030
155	204B		JEQ	EXIT	302057
160	204E		STCH	BUFFER,X	549039
165	2051		TIX	MAXLEN	2C205E
170	2054		JLT	RLOOP	38203F
175	2057	EXIT	STX	LENGTH	101036
180	205A		RSUB		4C0000
185	205D	INPUT	BYTE X'F1' F1		
190	205E	MAXLEN	WORD	4096	001000
195		.			
200		.	SUBROUTINE TO WRITE RECORD FROM BUFFER		
205		.			
210	2061	WRREC	LDX	ZERO	041030
215	2064	WLOOP	TD	OUTPUT	E02079
220	2067		JEQ	WLOOP	302064
225	206A		LDCH	BUFFER,X	509039
230	206D		WD	OUTPUT	DC2079
235	2070		TIX	LENGTH	2C1036
240	2073		JLT	WLOOP	382064
245	2076		RSUB		4C0000
250	2079	OUTPUT	BYTE X'05' 05		
255			END FIRST		

10) Write an Algorithm for pass 1 of SIC Assembler.

Pass 1:

```
begin
  read first input line
  if OPCODE = 'START' then
    begin
      save #[OPERAND] as starting address
      initialize LOCCTR to starting address
      write line to intermediate file
      read next input line
    end {if START}
  else
    initialize LOCCTR to 0
  while OPCODE ≠ 'END' do
    begin
      if this is not a comment line then
        begin
          if there is a symbol in the LABEL field then
            begin
              search SYMTAB for LABEL
              if found then
                set error flag (duplicate symbol)
              else
                insert (LABEL,LOCCTR) into SYMTAB
              end {if symbol}
            end
          search OPTAB for OPCODE
          if found then
            add 3 (instruction length) to LOCCTR
          else if OPCODE = 'WORD' then
            add 3 to LOCCTR
          else if OPCODE = 'RESW' then
            add 3 * #[OPERAND] to LOCCTR
          else if OPCODE = 'RESEB' then
            add #[OPERAND] to LOCCTR
          else if OPCODE = 'BYTE' then
            begin
              find length of constant in bytes
              add length to LOCCTR
            end {if BYTE}
          else
            set error flag (invalid operation code)
          end {if not a comment}
          write line to intermediate file
          read next input line
        end {while not END}
      write last line to intermediate file
      save (LOCCTR - starting address) as program length
    end {Pass 1}
```

Figure 2.4(a) Algorithm for Pass 1 of assembler.

11) Write an algorithm for pass 2 of SIC assembler.

Pass 2:

```
begin
  read first input line {from intermediate file}
  if OP CODE = 'START' then
    begin
      write listing line
      read next input line
    end {if START}
  write Header record to object program
  initialize first Text record
  while OP CODE ≠ 'END' do
    begin
      if this is not a comment line then
        begin
          search OPTAB for OP CODE
          if found then
            begin
              if there is a symbol in OPERAND field then
                begin
                  search SYMTAB for OPERAND
                  if found then
                    store symbol value as operand address
                  else
                    begin
                      store 0 as operand address
                      set error flag (undefined symbol)
                    end
                  end {if symbol}
                else
                  store 0 as operand address
                  assemble the object code instruction
                end {if opcode found}
              else if OP CODE = 'BYTE' or 'WORD' then
                convert constant to object code
              if object code will not fit into the current Text record then
                begin
                  write Text record to object program
                  initialize new Text record
                end
              add object code to Text record
            end {if not comment}
          write listing line
          read next input line
        end {while not END}
      write last Text record to object program
      write End record to object program
      write last listing line
    end {Pass 2}
```

Figure 2.4(b) Algorithm for Pass 2 of assembler.

12) What are the Data Structures used in an Assembler?

Data Structures:

- Operation Code Table (OPTAB)
- Symbol Table (SYMTAB)
- Location Counter(LOCCTR)

**13) Explain the features of a Symbol Table.**

- **SYMTAB (symbol table)**
- **Content**  
Label name and its value (address)  
May also include flag (type, length) etc.
- **Usage**  
**Pass 1:** labels are entered into SYMTAB with their address (from LOCCTR) as they are encountered in the source program  
**Pass 2:** symbols used as operands are looked up in SYMTAB to obtain the address to be inserted in the assembled instruction
- **Characteristic**  
Dynamic table (insert, delete, search)
- **Implementation**  
Hash table for efficiency of *insertion* and *retrieval*

<b>COPY</b>	<b>1000</b>
<b>FIRST</b>	<b>1000</b>
<b>CLOOP</b>	<b>1003</b>
<b>ENDFIL</b>	<b>1015</b>
<b>EOF</b>	<b>1024</b>
<b>THREE</b>	<b>102D</b>
<b>ZERO</b>	<b>1030</b>
<b>RETADR</b>	<b>1033</b>
<b>LENGTH</b>	<b>1036</b>
<b>BUFFER</b>	<b>1039</b>
<b>RDREC</b>	<b>2039</b>

SYMBOL TABLE(SYMTAB)

**14) What is Location Counter?**

- Location Counter**
- A variable used to help in *assignment of addresses*
- Initialized to the beginning address specified in the START statement
- Counted in bytes



**15) What are the machine dependant fetures of a SIC/XE Assembler?**

**Machine-dependent features of assemblers**

**Features of the SIC/XE machine**

**Programming features.**

- a. # symbol.
  - i. Indication of the immediate addressing mode.
  - ii. Immediate addressing provides a faster access to an operand reference.
- b. @ symbol.
  - i. Indication of the indirect addressing mode.
  - ii. Indirect addressing reduces the number of instructions.
- c. + symbol.
  - i. Explicit selection of the format 4 instruction with a direct addressing mode.
  - ii. Format 4 is selected when the 12-bit displacement of format 3 is too small.
- d. BASE directive.
  - i. Indication that the base register B holds a base address used in a base addressing.
  - ii. NOBASE directive disables the base register.
  - iii. LDB instruction loads the base register with a base address.
- e. Register-to-register addressing.
  - i. Register addressing reduces the size of a machine instruction and speeds up a computation

**Assembling features.**

- f. Multiprogramming.
  - i. Larger memory allows us to load many programs.
  - ii. The object code is relative to zero because the load address is variable.
  - iii. Program must be relocated when it is loaded in memory.
- g. Register set mapping.
  - i. A separate register table can store the numeric values of the registers.
  - ii. The numeric values of the registers can be preloaded with the symbol table.
- h. Relative (PC and base) addressing mode.
  - i. Operand value is subtracted from PC or base register value.
  - ii. PC relative addressing provides a displacement from -2048 to +2047.
  - iii. Base relative addressing provides a displacement from 0 to 4095.

## 16) What is Program Relocation?

### Program relocation

#### ➤ Principles.

- The load address of an object program is unknown at assembly time if the system implements the multiprogramming feature.
- The assembler generates addresses relative to zero in the object program.
- At load time, relocation is performed by adding the load address to the relative addresses.
- Operands of instructions that use direct addressing must be relocated, and the assembler provides the relocation information in the object program.
- Operands of instructions that use relative addressing do not need to be relocated.
- Relocation can be processed by the loader or by the CPU using relocation registers.

## 17) What are the advantages of program relocation?

### **Program Relocation**

- The larger main memory of SIC/XE
  - Several programs can be loaded and run at the same time.
  - This kind of sharing of the machine between programs
  - is called ***multiprogramming***
- To take full advantage
  - Load programs into memory wherever there is room
  - Not specifying a fixed address at assembly time
  - Called ***program relocation***

## 18) What are program blocks?

### **Program Blocks**

- Refer to segments of code that are rearranged within a single object program unit
- USE [blockname]
- At the beginning, statements are assumed to be part of the unnamed (default) block
- If no USE statements are included, the entire program belongs to this single block
- Each program block may actually contain several separate segments of the source program

## 19) How the program blocks are assembled?

### **Program Blocks - Implementation**

#### ■ Pass 1

- Each program block has a separate location counter
- Each label is assigned an address that is relative to the start of the block that contains it
- At the end of Pass 1, the latest value of the location counter for each block indicates the length of that block
- The assembler can then assign to each block a starting address in the object program

#### ■ Pass 2

- The address of each symbol can be computed by adding the assigned block starting address and the relative address of the symbol to that block

Line	Source statement				
5	COPY	START	0		COPY FILE FROM INPUT TO OUTPUT
10	FIRST	STL	RETADR		SAVE RETURN ADDRESS
15	CLOOP	JSUB	RDREC		READ INPUT RECORD
20		LDA	LENGTH		TEST FOR EOF (LENGTH = 0)
25		COMP	#0		
30		JEQ	ENDFIL		EXIT IF EOF FOUND
35		JSUB	WRREC		WRITE OUTPUT RECORD
40		J	CLOOP		LOOP
45	ENDFIL	LDA	=C'EOF'		INSERT END OF FILE MARKER
50		STA	BUFFER		
55		LDA	#3		SET LENGTH = 3
60		STA	LENGTH		
65		JSUB	WRREC		WRITE EOF
70		J	@RETADR		RETURN TO CALLER
92		USE	CDATA		
95	RETADR	RESW	1		
100	LENGTH	RESW	1		LENGTH OF RECORD
103		USE	CBLKS		
105	BUFFER	RESB	4096		4096-BYTE BUFFER AREA
106	BUFEND	EQU	*		FIRST LOCATION AFTER BUFFER
107	MAXLEN	EQU	BUFEND-BUFFER		MAXIMUM RECORD LENGTH
110					
115					
120					
123		USE			
125	RDREC	CLEAR	X		CLEAR LOOP COUNTER
130		CLEAR	A		CLEAR A TO ZERO
132		CLEAR	S		CLEAR S TO ZERO
133		+LDT	#MAXLEN		
135	RLOOP	TD	INPUT		TEST INPUT DEVICE
140		JEQ	RLOOP		LOOP UNTIL READY
145		RD	INPUT		READ CHARACTER INTO REGISTER A
150		COMPR	A,S		TEST FOR END OF RECORD (X'00')
155		JEQ	EXIT		EXIT LOOP IF BOR
160		STCH	BUFFER,X		STORE CHARACTER IN BUFFER
165		TXR	T		LOOP UNLESS MAX LENGTH
170		JLT	RLOOP		HAS BEEN REACHED
175	EXIT	STX	LENGTH		SAVE RECORD LENGTH
180		RSUB			RETURN TO CALLER
183		USE	CDATA		
185	INPUT	BYTE	X'F1'		CODE FOR INPUT DEVICE
195					
200					
205					
208		USE			
210	WRREC	CLEAR	X		CLEAR LOOP COUNTER
212		LDT	LENGTH		
215	WLOOP	TD	=X'05'		TEST OUTPUT DEVICE
220		JEQ	WLOOP		LOOP UNTIL READY
225		LDCH	BUFFER,X		GET CHARACTER FROM BUFFER
230		WD	=X'05'		WRITE CHARACTER
235		TXR	T		LOOP UNTIL ALL CHARACTERS
240		JLT	WLOOP		HAVE BEEN WRITTEN
245		RSUB			RETURN TO CALLER
252		USE	CDATA		
253		LTORG			
255		END	FIRST		

**Figure 2.11** Example of a program with multiple program blocks.

- Each source line is given a relative address assigned and a block number

Block name	Block number	Address	Length
(default)	0	0000	0066
CDATA	1	0066	000B
CBLKS	2	0071	1000

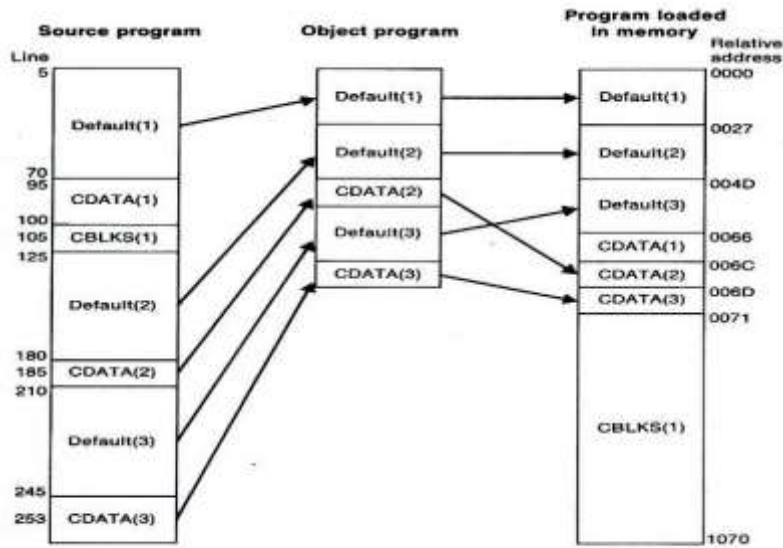


Figure 2.14 Program blocks from Fig. 2.11 traced through the assembly and loading processes.

## 20) What is one pass assembler? Explain the functioning of one-pass assembler.

- One-pass assemblers are used when
  - it is necessary or desirable to avoid a second pass over the source program
  - the external storage for the intermediate file between two passes is slow or is inconvenient to use
- Main problem: forward references to both data and instructions
- One simple way to eliminate this problem: require that all areas be defined before they are referenced.
  - It is possible, although inconvenient, to do so for data items.
  - Forward jump to instruction items cannot be easily eliminated.

Sample Program for a One-Pass Assembler

## Sample Program for a One-Pass Assembler

Line	Loc	Source statement			Object code
0	1000	COPY	START	1000	
1	1000	EOF	BYTE	C'EOF'	454F46
2	1003	THREE	WORD	3	000003
3	1006	ZERO	WORD	0	000000
4	1009	RETADR	RESW	1	
5	100C	LENGTH	RESW	1	
6	100F	BUFFER	RESB	4096	
9		.			
10	200F	FIRST	STL	RETADR	141009
15	2012	CLOOP	<u>JSUB</u>	<u>RDREC</u>	48203D
20	2015		LDA	LENGTH	00100C
25	2018		COMP	ZERO	281006
30	201B		<u>JEQ</u>	<u>ENDFIL</u>	302024
35	201E		<u>JSUB</u>	<u>WRREC</u>	482062
40	2021		J	CLOOP	302012
45	2024	ENDFIL	LDA	EOF	001000
50	2027		STA	BUFFER	0C100F
55	202A		LDA	THREE	001003
60	202D		STA	LENGTH	0C100C
65	2030		<u>JSUB</u>	<u>WRREC</u>	482062
70	2033		LDL	RETADR	081009
75	2036		RSUB		4C0000
110		.			

110	.				
115	.	SUBROUTINE TO READ RECORD INTO BUFFER			
120					
121	2039	INPUT	BYTE	X'F1'	F1
122	203A	MAXLEN	WORD	4096	001000
124		.			
125	203D	RDREC	LDX	ZERO	041006
130	2040		LDA	ZERO	001006
135	2043	RLOOP	TD	INPUT	E02039
140	2046		JEQ	RLOOP	302043
145	2049		RD	INPUT	D82039
150	204C		COMP	ZERO	281006
155	204F		<u>JEQ</u>	<u>EXIT</u>	30205B
160	2052		STCH	BUFFER,X	54900F
165	2055		TIX	MAXLEN	2C203A
170	2058		JLT	RLOOP	382043
175	205B	EXIT	STX	LENGTH	10100C
180	205E		RSUB		4C0000
185		.			

```

195      .
200      .      SUBROUTINE TO WRITE RECORD FROM BUFFER
205
206      2061      OUTPUT      BYTE      X'05'      05
207
210      2062      WRREC      LDX      ZERO      041006
215      2065      WLOOP      TD      OUTPUT      E02061
220      2068      JEQ      WLOOP      302065
225      206B      LDCH      BUFFER,X      50900F
230      206E      WD      OUTPUT      DC2061
235      2071      TIX      LENGTH      2C100C
240      2074      JLT      WLOOP      382065
245      2077      RSUB      4C0000
255      END      FIRST

```

#### Load-and-Go Assembler

- Load-and-go assembler generates their object code in memory for immediate execution.
- No object program is written out, no loader is needed.
- It is useful in a system oriented toward program development and testing such that the efficiency of the assembly process is an important consideration.

#### How to Handle Forward References

- Load-and-go assembler
  - Omits the operand address if the symbol has not yet been defined
  - Enters this undefined symbol into SYMTAB and indicates that it is undefined
  - Adds the address of this operand address to a list of forward references associated with the SYMTAB entry
  - Scans the reference list and inserts the address when the definition for the symbol is encountered.
  - Reports the error if there are still SYMTAB entries indicated undefined symbols at the end of the program
  - Search SYMTAB for the symbol named in the END statement and jumps to this location to begin execution if there is no error

# Object Code in Memory and SYMTAB

After scanning line 40

Memory address	Contents	Symbol Value
1000	454F4600 00030000 00xxxxxx xxxxxxxx	LENGTH 100C
1010	xxxxxxxx xxxxxxxx xxxxxxxx xxxxxxxx	RDREC * → 2013 0
.	.	THREE 1003
.	.	ZERO 1006
2000	xxxxxxxx xxxxxxxx xxxxxxxx xxxxxxxx14	WRREC * → 201F 0
2010	100940 3D0100C 28100630 202408	EOF 1000
2020	780012 0010000C 100F0010 000C100C	ENDFIL * → 201C 0
.	.	RETADR 1009
.	.	BUFFER 100F
.	.	CLOOP 2012
.	.	FIRST 200F

# Object Code in Memory and SYMTAB

After scanning line 160

Memory address	Contents	Symbol Value
1000	454F4600 00030000 00xxxxxx xxxxxxxx	LENGTH 100C
1010	xxxxxxxx xxxxxxxx xxxxxxxx xxxxxxxx	RDREC 203D
.	.	THREE 1003
.	.	ZERO 1006
2000	xxxxxxxx xxxxxxxx xxxxxxxx xxxxxxxx14	WRREC * → 201F → 2031 0
2010	100940 3D0100C 28100630 202408	EOF 1000
2020	780012 0010000C 100F0010 000C100C	ENDFIL 2024
2030	400008 10094C00 00F10010 00041006	RETADR 1009
2040	000005E0 20393020 43D82039 28100630	BUFFER 100F
2050	00000490 00000000 00000000 00000000	CLOOP 2012
.	.	FIRST 200F
.	.	MAXLEN 203A
.	.	INPUT 2039
.	.	EXIT * → 2050 0
.	.	RLOOP 2043



# Object Program from One-Pass Assembler

```
HCOPY  00100000107A
T00100009454F46000003000000
T00200F1514100948000000100C2810063000004800003C2012
T00201C022024
T002024190010000C100F0010030C100C4800000810094C0000F1001000
T00201302203D
T00203D1E041006001006E02039302043D8203928100630000054900F2C203A382043
T00205002205B
T00205B0710100C4C000005
T00201F022062
T002031022062
T00206218041006E0206130206550900FDC20612C100C3820654C0000
E00200F
```

## 21) What is a multi-pass assembler? Explain with an example, the functioning of a multi-pass assembler.

### Multi-Pass Assemblers

- Prohibiting forward references in symbol definition:
  - This restriction is not a serious inconvenience.
  - Forward references tend to create difficulty for a person reading the program.
- Allowing forward references
  - To provide more flexibility
  - Solution:
    - A multi-pass assembler that can make as many passes as are needed to process the definitions of symbols.
    - Only the portions of the program that involve forward references in symbol definition are saved for multi-pass reading.
- For a two pass assembler, forward references in symbol definition are not allowed:

```
ALPHA  EQU  BETA
BETA   EQU  DELTA
DELTA  RESW 1
```

- Reason: symbol definition must be completed in pass 1.
- Motivation for using a multi-pass assembler
  - DELTA can be defined in pass 1
  - BETA can be defined in pass 2
  - ALPHA can be defined in pass 3

### Implementation



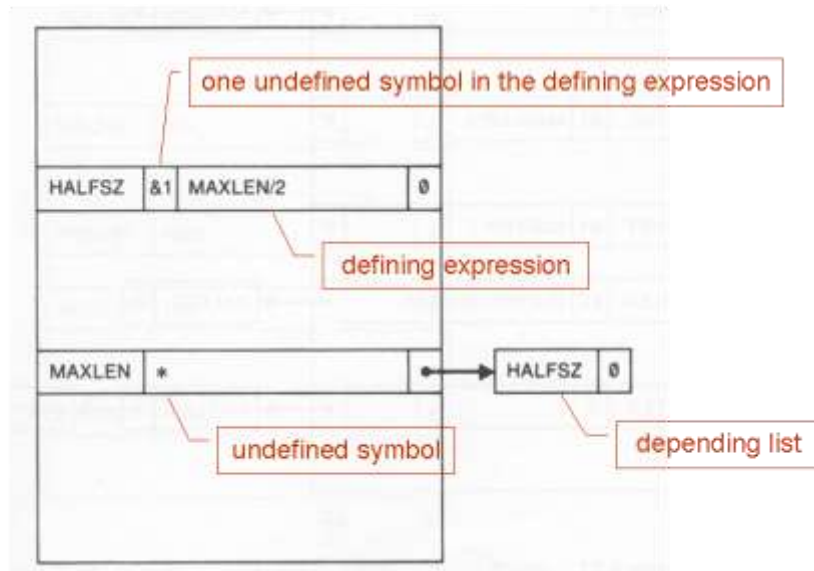
- A symbol table is used
  - to store symbol definitions that involve forward references
  - to indicate which symbols are dependant on the values of others
  - to facilitate symbol evaluation
- For a forward reference in symbol definition, we store in the SYMTAB:
  - the symbol name
  - the defining expression
  - the number of undefined symbols in the defining expression
  - the undefined symbol (marked with a flag \*) associated with a list of symbols depend on this undefined symbol.
- When a symbol is defined, we can recursively evaluate the symbol expressions depending on the newly defined symbol.

## Forward Reference Example

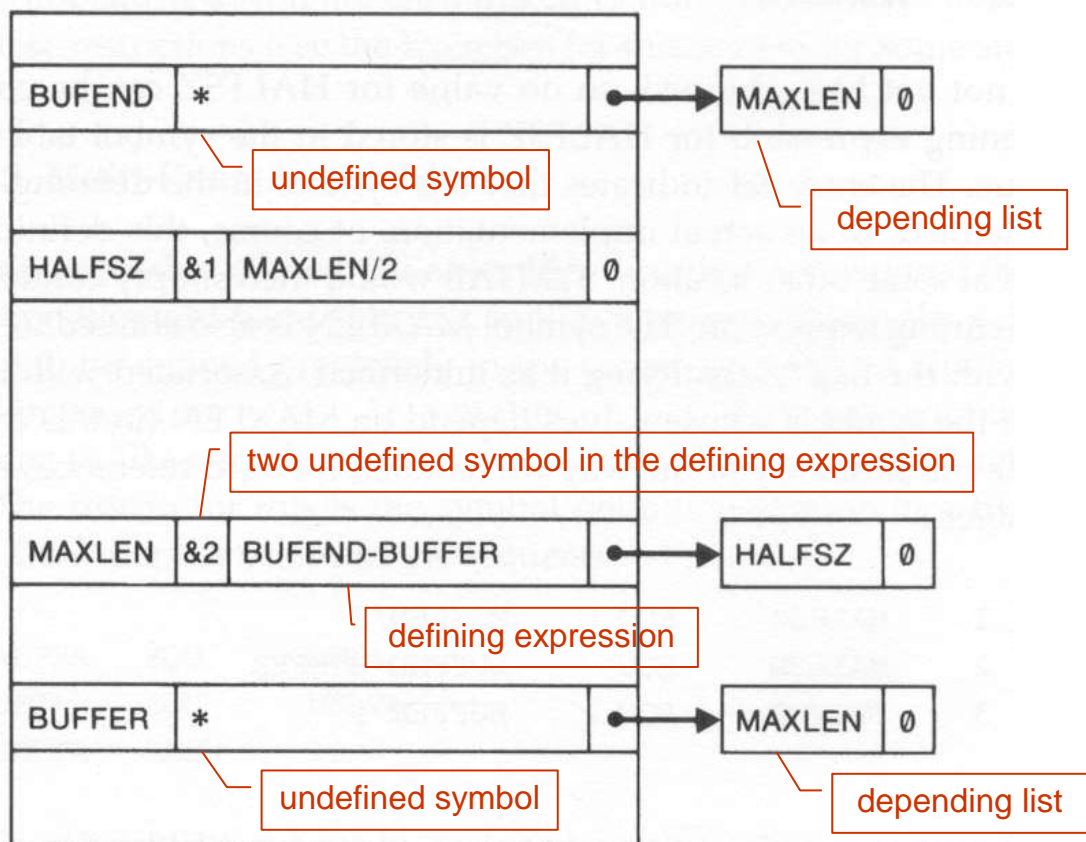
1	HALFSZ	EQU	MAXLEN/2
2	MAXLEN	EQU	BUFEND-BUFFER
3	PREVBT	EQU	BUFFER-1
			.
			.
			.
4	BUFFER	RESB	4096
5	BUFEND	EQU	*

# Forward Reference Example

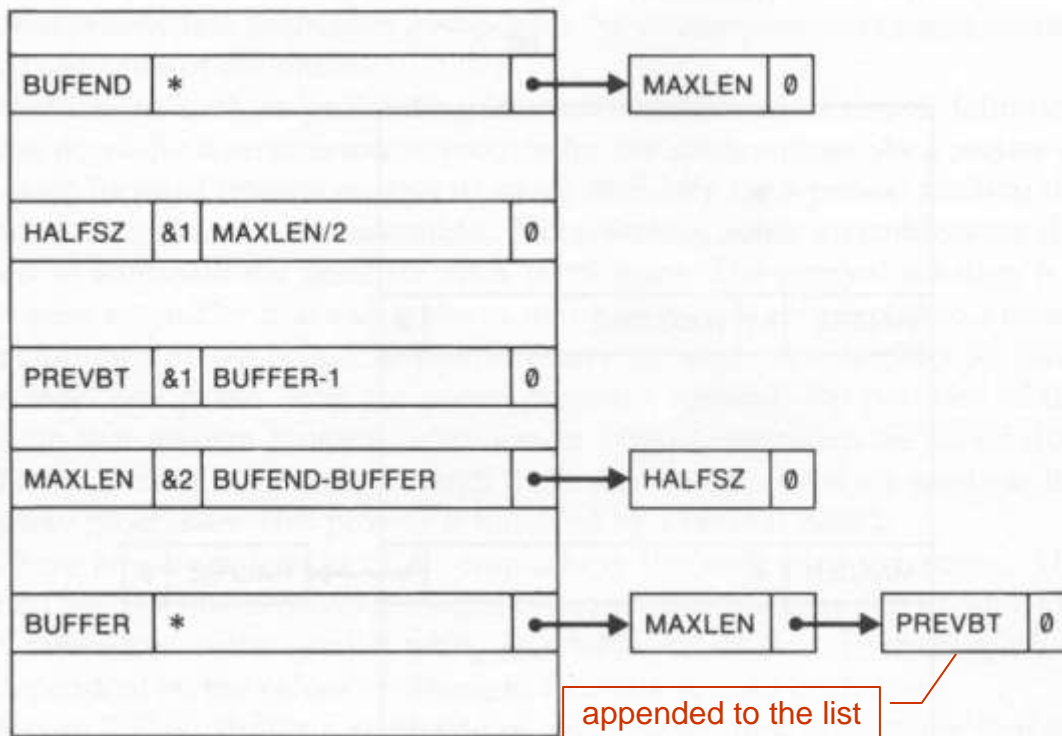
1    HALFSZ    EQU    MAXLEN/2



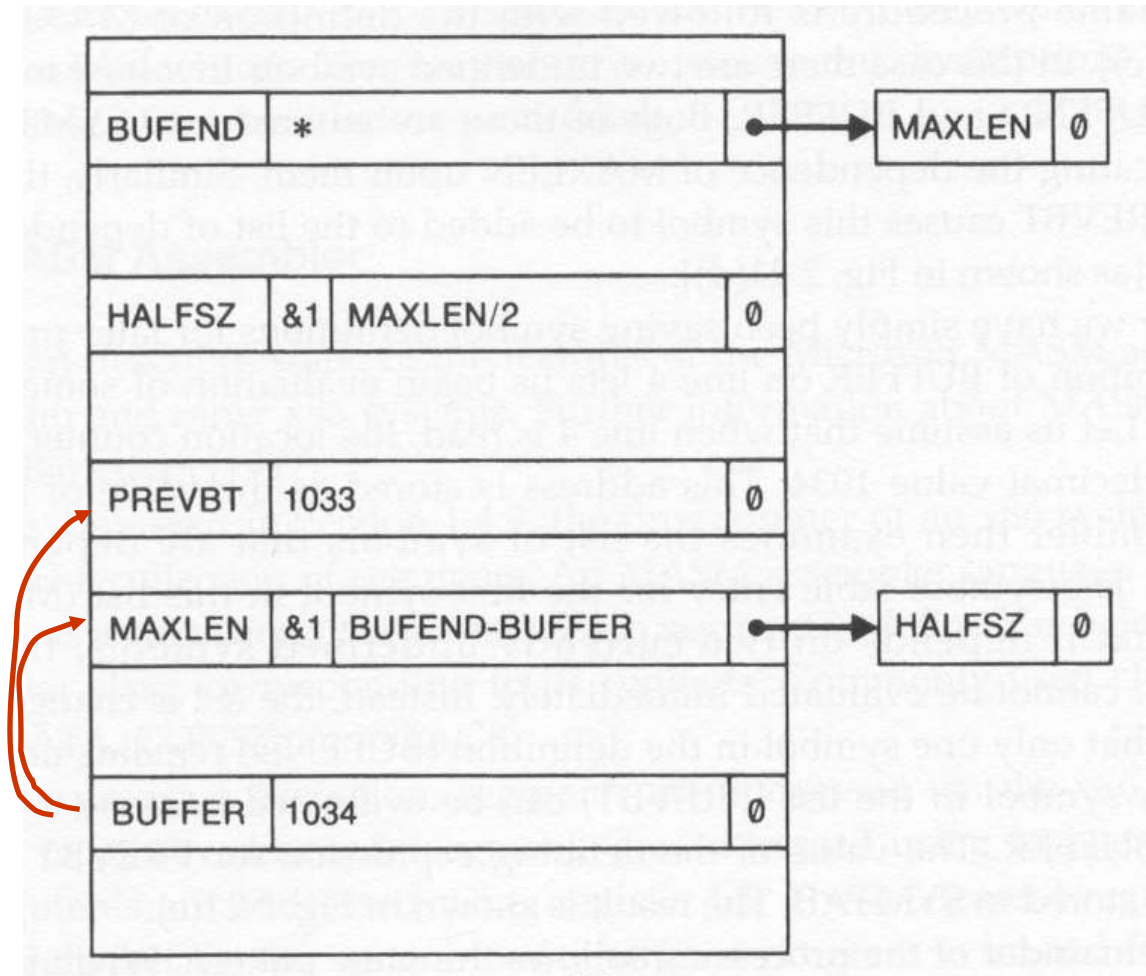
2    MAXLEN    EQU    BUFEND-BUFFER



3    PREVBT    EQU    BUFFER-1



4      BUFFER    RESB    4096



5      BUFEND    EQU    \*

BUFEND	2034	0
HALFSZ	800	0
PREVBT	1033	0
MAXLEN	1000	0
BUFFER	1034	0

```
graph TD; BUFEND --> HALFSZ; HALFSZ --> PREVBT; PREVBT --> MAXLEN; MAXLEN --> BUFFER;
```