

COT 5615 Math for Intelligent Systems Fall 2021 Homework #3

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Problem 5.2**A surprising discovery****Solution**

According to me, the supervisor is wrong and the intern is probably correct. Here because of the independence-dimension inequality rule, any set collection of $n+1$ or more n -vectors is linearly dependent; and she is analysing 400 250-vector stocks, and thus this set is linearly dependent. Thus, return of any stock i.e. Google can be expressed as a linear combination of the return of other stocks. Here, this fact is only valid for present return stock value; and in the near future, this fact might change i.e. Google's return stock might not be expressed as a linear combination of the return of other stocks and thus, is not very useful from monetary perspective.

Problem 5.5**Orthogonalizing vectors****Solution**

Two vectors are orthogonal, if their inner product is zero i.e. we have to find γ such that $(a - \gamma b)^T b = a^T b - \gamma b^T b = 0$. Now, if $b = 0$, then any value of γ will yield $(a - \gamma b) \perp b$ true as all vectors are orthogonal to 0. If $b \neq 0$, $b^T b = \|b\|^2 \neq 0$; then we can take $\gamma = a^T b / b^T b$, which proves that $(a - \gamma b) \perp b$ is true.

Problem 5.9**Solution**

The Gram-Schmidt algorithm requires $n \cdot k^2$ flops, and thus for $n = 10^4$ and $k = 2$, $2 \cdot 10^{10}$ flops are calculated in 2 seconds. Therefore, for $\tilde{n} = 10^3$ and $\tilde{k} = 500$, we can get the run-time of the Gram-Schmidt Algorithm as follows: $(2(2 \cdot 1000 \cdot (500)^2)) / (2 \cdot 10^{10}) = 0.05$ seconds.

Problem 6.17**Stacked matrix****Solution**

- a Let's assume $Sx = 0$, thus $Sx = (Ax, x) = 0$, which implies $x = 0$. In conclusion, S always has linearly independent columns.
- b S has $m+n$ rows and each row is n -dimension wide. Thus, according to the independence-dimension inequality rule, S can never have linearly independent rows i.e. rows are dependent.

Problem A6.8

Solution

- a The columns of matrix A ($m \times n$) may be linearly independent if for $Ax = 0$ implies $x = 0$ when the number of columns is less than or equal to number of rows [$n \leq m$].
- b The rows of matrix A ($m \times n$) may be linearly independent if for $Ax = 0$ implies $x = 0$ when the number of rows is less than or equal to number of columns [$m \leq n$].

Problem 6.18

Vandermonde matrices

Solution

Here Vc vector represents the values of the polynomial at t_1, t_2, \dots, t_m as follows:

$$\begin{aligned} Vc &= (c_1 + c_2 t_1 + c_3 t_1^2 + \dots + c_n t_1^{n-1}, c_1 + c_2 t_2 + c_3 t_2^2 + \dots + c_n t_2^{n-1}, \dots, c_1 + c_2 t_m + c_3 t_m^2 + \dots + c_n t_m^{n-1}) \\ &= (p(t_1), p(t_2), \dots, p(t_m)) \end{aligned}$$

Now, if $Vc = 0$, then $p(t_i)$ is also 0 for $i = 1, 2, \dots, m$; thus $p(t)$ has at least m distinct roots t_1, t_2, \dots, t_m . This is only possible if all the coefficients of p are 0 i.e. $c=0$. Therefore, $Vc = 0$ implies $c = 0$, which proves that the columns of V [Vandermonde Matrix] are linearly independent.

Problem A6.2

Vandermonde matrices in Julia

Solution

```
1 using LinearAlgebra
2 function Vandermonde_generator(n,m)
3     v = ones(length(m),n)
4     for i in 1:n
5         mnew = m.^(i-1)
6         for j in 1:length(m)
7             v[j,i] = mnew[j]
8         end
9     end
10    display(v)
11 end
12 Vandermonde_generator(5,[5,6,7])
13 Vandermonde_generator(5,[7,8,6])
```
