COT 5615 Math for Intelligent Systems, Fall 2021

Homework 9

Due Dec. 7

1. Solving regularized least-squares. Consider the least-squares problem

$$\underset{\boldsymbol{x}}{\text{minimize}} \quad \|\boldsymbol{A}\boldsymbol{x} - \boldsymbol{b}\|^2 + \lambda \|\boldsymbol{x}\|^2.$$

The solution is

$$\widehat{m{x}} = egin{bmatrix} m{A} \ \sqrt{\lambda} m{I} \end{bmatrix}^\dagger egin{bmatrix} m{b} \ 0 \end{bmatrix} = (m{A}^{\! op} \! m{A} + \lambda m{I})^{-1} m{A}^{\! op} m{b}.$$

This may suggest that the complexity is at least $\mathcal{O}(mn^2+n^3)$. Suppose \boldsymbol{A} is wide, i.e., m < n. Explain how to solve it using $\mathcal{O}(m^2n+m^3)$ flops. *Hint*. Use the "kernel trick" in §15.5.2 of the textbook.

2. Companion matrix. The matrix C below is called a companion matrix. Consider the companion matrix C (left) and the Vandermonde matrix V (right):

$$C = \begin{bmatrix} 0 & 0 & \cdots & 0 & -c_0 \\ 1 & 0 & \cdots & 0 & -c_1 \\ 0 & 1 & 0 & 0 & -c_2 \\ \vdots & \ddots & & \vdots \\ 0 & 0 & \cdots & 1 & -c_{n-1} \end{bmatrix} \qquad V = \begin{bmatrix} 1 & \lambda_1 & \lambda_1^2 & \cdots & \lambda_1^{n-1} \\ 1 & \lambda_2 & \lambda_2^2 & \cdots & \lambda_2^{n-1} \\ 1 & \lambda_3 & \lambda_3^2 & \cdots & \lambda_3^{n-1} \\ \vdots & \vdots & \vdots & & \vdots \\ 1 & \lambda_n & \lambda_n^2 & \cdots & \lambda_n^{n-1} \end{bmatrix}.$$

- (a) What is the characteristic polynomial of C?
- (b) Assume that C has n distinct eigenvalues $\lambda_1, \lambda_2, \ldots, \lambda_n$. Show that V is the matrix of left eigenvectors of C, i.e., that we have $VC = \Lambda V$ where $\Lambda = \text{Diag}(\lambda_1, \ldots, \lambda_n)$. (Recall that the eigenvalues are the roots of the characteristic polynomial. Some of you may know that polynomials of degree 5 or higher does not have analytical expressions of their roots. This is how their roots are found numerically in practice.)

3. Circulant matrix. The matrix C below is called a circulant matrix. The matrix F below is called the DFT matrix (discrete Fourier transform).

$$C = \begin{bmatrix} c_1 & c_n & \cdots & c_3 & c_2 \\ c_2 & c_1 & \cdots & c_4 & c_3 \\ c_3 & c_2 & \cdots & c_5 & c_4 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ c_n & c_{n-1} & \cdots & c_2 & c_1 \end{bmatrix} \qquad F = \frac{1}{\sqrt{n}} \begin{bmatrix} 1 & 1 & 1 & \cdots & 1 \\ 1 & \omega & \omega^2 & \cdots & \omega^{n-1} \\ 1 & \omega^2 & \omega^4 & \cdots & \omega^{2(n-1)} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & \omega^{n-1} & \omega^{2(n-1)} & \cdots & \omega^{(n-1)(n-1)} \end{bmatrix},$$

where $\omega = e^{-2\pi i/n}$ and $i = \sqrt{-1}$. Denote the first column of C as c.

- (a) Verify that $C = F^* \Lambda F$ where $\Lambda = \text{Diag}(Fc)$. The operation Fc is called the discrete Fourier transform (DFT) of c. Since $F^{-1} = F^*$, this shows the eigenvectors of a circulant matrix are columns of the DFT matrix.
- (b) Suppose you want to solve the linear equation Cx = b where C is a circulant matrix. How would you do it and what is the complexity? *Hint*. For the $n \times n$ DFT matrix F, the matrix-vector multiplication can be done in $n \log(n)$ flops (rather than n^2 flops for general $n \times n$ matrices) using the fast Fourier transform (FFT) algorithm, and similarly for its conjugate transpose using the inverse FFT.
- 4. General least-squares. Consider the least-squares problem

$$\underset{\boldsymbol{x}}{\text{minimize}} \quad \|\boldsymbol{A}\boldsymbol{x} - \boldsymbol{b}\|^2.$$

If the columns of \boldsymbol{A} are linearly independent, we have shown that the solution is unique and is

$$\widehat{\boldsymbol{x}} = \boldsymbol{A}^{\dagger} \boldsymbol{b} = (\boldsymbol{A}^{\mathsf{T}} \boldsymbol{A})^{-1} \boldsymbol{A}^{\mathsf{T}} \boldsymbol{b}.$$

Now suppose that the columns may not be linearly independent. If we are given the 'thin' SVD of $A = U \Sigma V^{\top}$, then we can show that a solution is

$$\widehat{oldsymbol{x}} = oldsymbol{A}^\dagger oldsymbol{b} = oldsymbol{V} oldsymbol{\Sigma}^{-1} oldsymbol{U}^ op oldsymbol{b}.$$

In this question we try to verify this result.

- (a) Setting the gradient equal to zero, we obtain the equation $\mathbf{A}^{\top} \mathbf{A} \hat{\mathbf{x}} = \mathbf{A}^{\top} \mathbf{b}$. Verify that the solution obtained by the SVD satisfy this equation.
- (b) Follow the steps on page 230 of the textbook, verify that it indeed minimizes the objective function.
- 5. QRSVD. For a $m \times n$ matrix \boldsymbol{A} (assume $m \geq n$) the QR SVD algorithm is as follows
 - 1: $A = Q_0 R_0$

 \triangleright QR factorization of \boldsymbol{A}

- 2: **for** $k = 0, 1, 2, \dots$ **do**
- 3: $oldsymbol{R}_k^{ op} = oldsymbol{Q}_{k+1} oldsymbol{R}_{k+1}$ ightharpoonup QR factorization of $oldsymbol{R}_k$
- 4: end for

In the QR factorization it is assumed that diagonal entries of R are positive. If $U\Sigma V^{\top}$ is the 'thin' SVD of A, then it is observed that the product $Q_1Q_3\cdots Q_{2k+1}\cdots$ converges to V, the product $Q_0Q_2\cdots Q_{2k}\cdots$ converges to U, and R_{2k} converges to Σ .

- (a) Try the algorithm on a random 6×3 and compare it with svd(A).
- (b) Show that $Q_1(R_1R_0)$ is the QR decomposition of $R_0^{\top}R_0$ (Q_1 is the orthonormal factor and R_1R_0 is the upper triangular factor).
- (c) Show that $(\mathbf{R}_1\mathbf{R}_0)\mathbf{Q}_1 = \mathbf{R}_1\mathbf{R}_1^{\mathsf{T}}$ (Hint: focus first on the middle matrix \mathbf{R}_0 on the left hand side) and then that $\mathbf{R}_1\mathbf{R}_1^{\mathsf{T}} = \mathbf{R}_2^{\mathsf{T}}\mathbf{R}_2$. Show that $\mathbf{B}_0 = \mathbf{R}_0^{\mathsf{T}}\mathbf{R}_0$ and $\mathbf{B}_1 = \mathbf{R}_2^{\mathsf{T}}\mathbf{R}_2$ have the same set of eigenvalues.
- (d) More generally consider the matrices $\boldsymbol{B}_k = \boldsymbol{R}_{2k}^{\top} \boldsymbol{R}_{2k}$ for $k = 1, 2, 3, \ldots$ Show that these are the QR iterations applied to the matrix $\boldsymbol{B}_0 = \boldsymbol{R}_0^{\top} \boldsymbol{R}_0$.