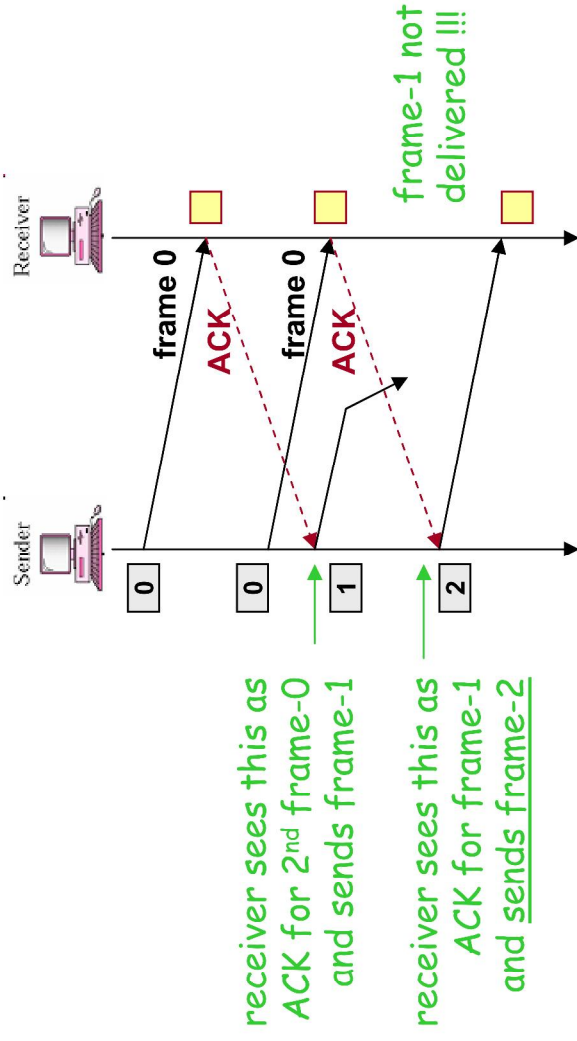
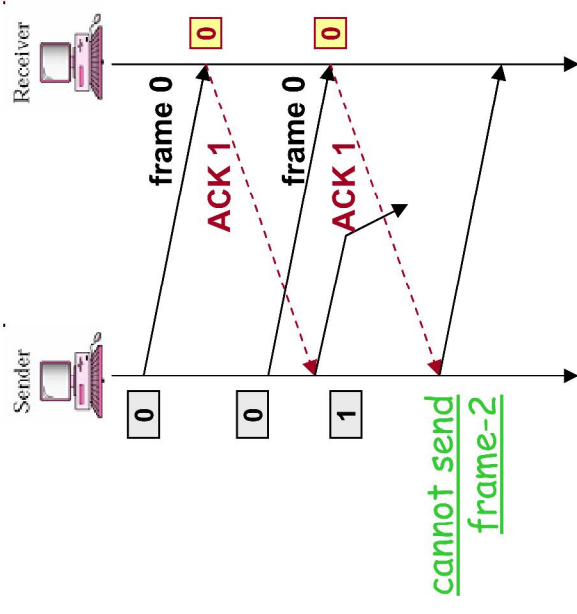


Delayed Acknowledgment (Premature Timeout)

- ACKs can be delayed due to problems with links or network congestion
 - time-out expires early, sender resends frame
 - when delayed ACK arrives, sender assumes that given ACK is for the last frame sent
- **ACKs must be numbered** to prevent gaps in delivered packet sequence



without ACK numbering

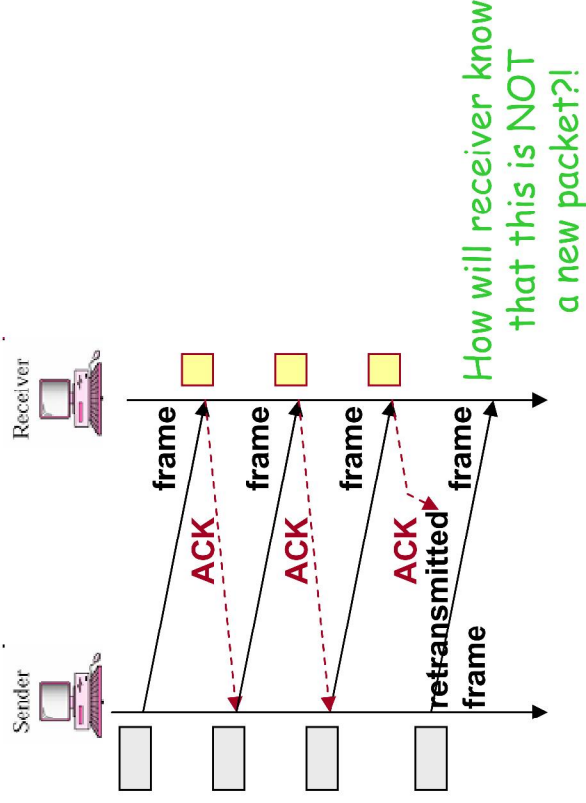


with ACK numbering

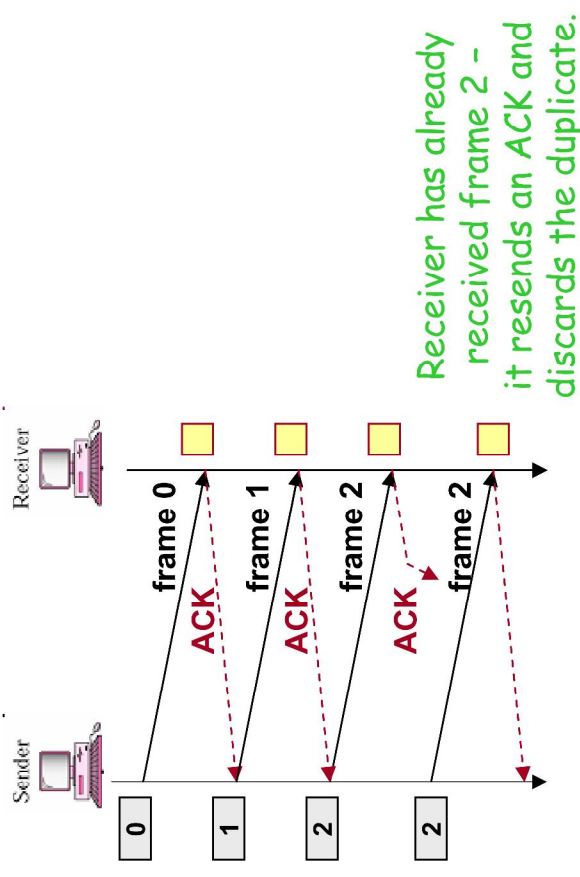
How large should the packet / ACK sequence be? Only 1-bit long !!!

Lost Acknowledgment

- frame received correctly, but ACK undergoes errors / loss
 - after time-out period, sender resends frame
 - receiver receives the same frame twice
- frames must be numbered so that receiver can recognize and discard duplicate frames
 - **sequence # are included in packet header**



without packet numbering

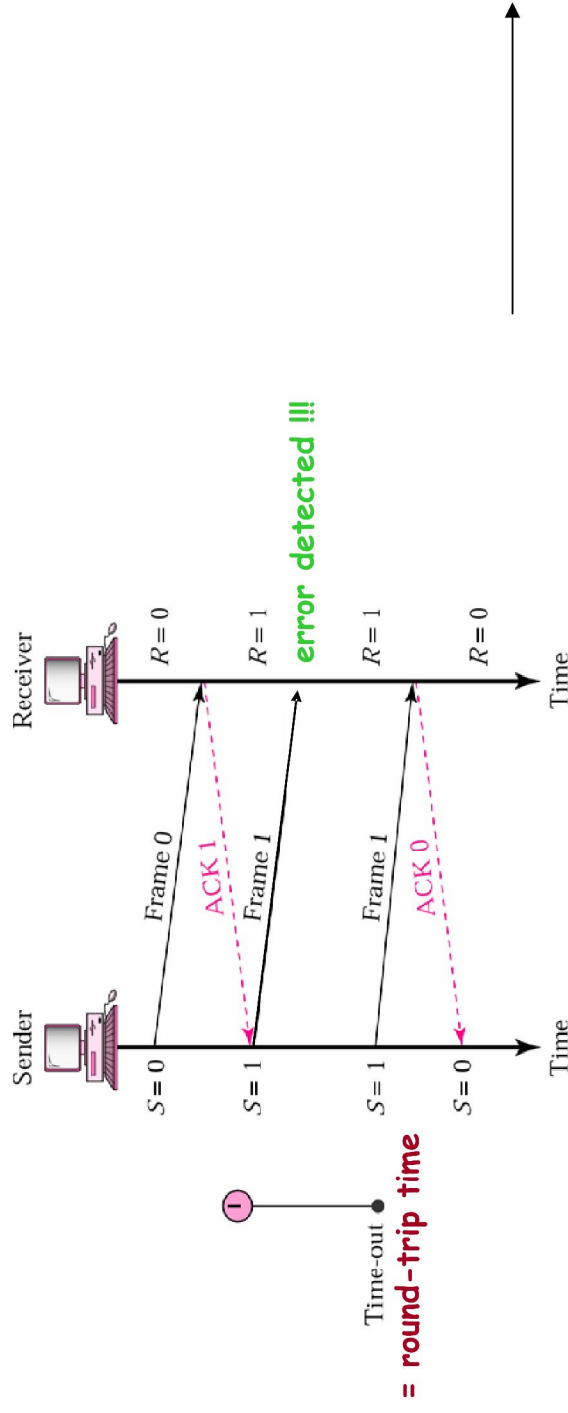


with packet numbering

Stop-and-Wait ARQ

Stop-and-Wait ARQ – simplest flow and error control mechanism

- sender sends an information frame to receiver
- sender, then, stops and waits for an ACK
- if no ACK arrives within **time-out**, sender will resend the frame, and again stop and wait
 - **time-out period > roundtrip time**
- abnormalities (and how to fix them)
 - lost acknowledgment
 - delayed acknowledgment



(1) Stop-and-Wait ARQ

Error and Flow Control

Flow Control – set of procedures used to restrict the amount of data that sender can send while waiting for acknowledgment

- two main strategies

- (1) **Stop-and-Wait**: sender waits until it receives ACK before sending next frame
- (2) **Sliding Window**: sender can send W frames before waiting for ACKs

Error + Flow Control Techniques

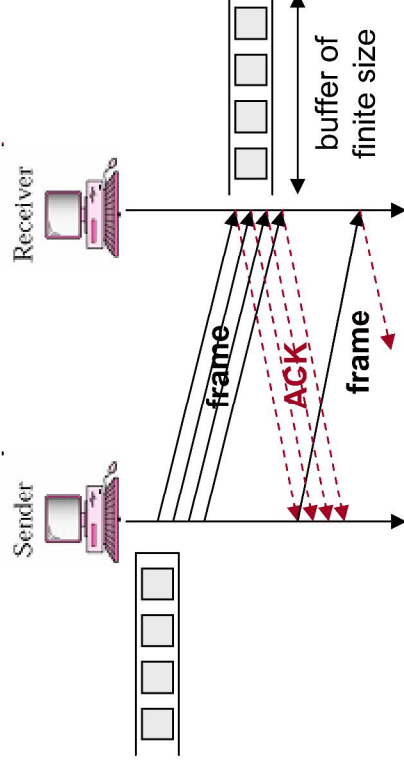
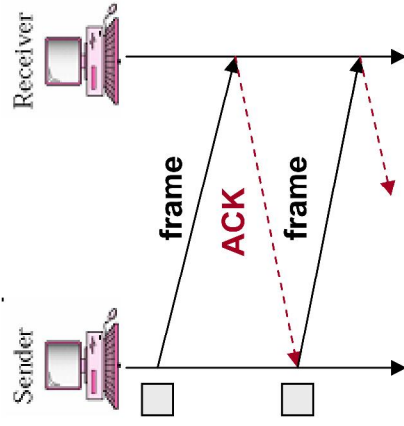
- (1) Stop-and-Wait ARQ
- (2) Go-Back-N ARQ
- (3) Selective Repeat ARQ

Error Detection + ARQ (error detection with retransmissions) must be combined with methods that intelligently limit the number of 'outstanding' (unACKed) frames.

Fewer unACKed frames ⇒ fewer packets buffered at sender and receiver.

Challenges of ARQ-based Error Control

- **send one frame at the time, wait for ACK**
 - easy to implement, but inefficient in terms of channel usage
- **send multiple frames at once**
 - better channel usage, but more complex to implement - sender must keep (all) sent but unACKed frame(s) in a buffer, as such frame(s) may have to be retransmitted



How many frames should be sent
at any point in time?

How should frames be released from
the sending buffer?

Error Control

Error Control (1) Forward Error Correction (FEC)

Approaches

(2) Error Detection + Automatic Retrans. Req. (ARQ)

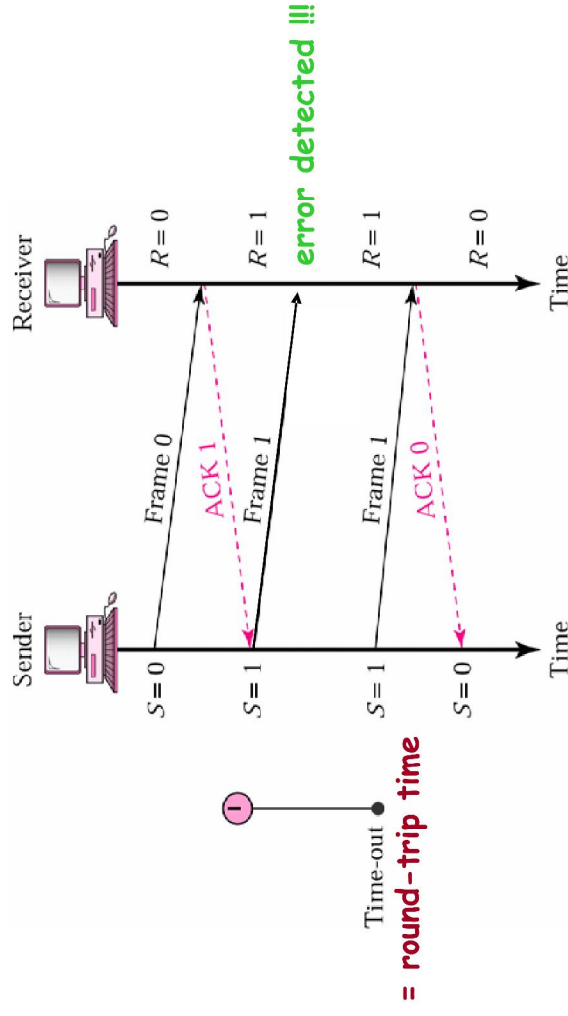
- not enough redundant info to enable error correction

case (a) receiver detects no errors

- an ACK packet is sent back to sender

case (b) receiver detects errors

- no ACK sent back to sender
- sender retransmits frame after a 'time-out'



Error and Flow Control

Required reading:
Garcia 5.2

CSE 3213, Fall 2010
Instructor: N. Vlatjic