Attay Vs Linker List 1) Cost of Accessing an Element: Address 200 Contex A[47; \* Algorithmic Complexity -> Constant Time

Linked List head nocle \* To access an element in the linked list of a Particular Position, we first need to start at the head nocle, there go to the second node and see the address of the third nocle. \* In the worst case, to in the list, we will be + saversing all the element in the list. \* Algorithmic Complexity -> Linear Time 0/n).

Linked List 283° A8899 Memory Resuisement. -> No unused Meanosy -> Fixed Size -> Extra memory regularment For Pointer Variable. \* We need to know the Size of array before Creating it. 12/0/ x4/0/> Null int A[5] [2/4/-1int. 4bytes Pointex. 4 bytes 8 bytes A node takes 2×8=166 85ytes -> 5 x 4 = 20 bytes \* it is Greated as one Contigous clock of memory \* Stored in scattered \* Memory should be Menory Block available as one large

Linked List 3) Cost of Inserting an head 2/0/ /4/0/ /6/0/ 18/ /2 Element: -> 3 Scenatio For Insertion int A[2] [2/4/6/8/-/-]

> If assay is not full.

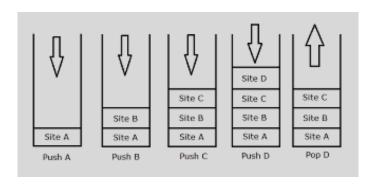
> 3 Scenario For Insection Position Algorithmic Complexity a) at beginning o(1) b) at end o(n) a) at beginning o(n) c) at it Position O(N) b) at end o(1) 72 + 54 - 1 76 - 1 78/ 1> c) at ith Position 0(n) \* New head \* we have to shift all element for insection at beginning. y) Ease of use -> Attay is simple -> Linked List is Complex

## What Is a Stack?

A stack is a linear data structure which allows the adding and removing of elements in a particular order.

## Implementing Stack Functionalities Using a Linked List

Stack can be implemented using both arrays and linked lists. The limitation, in the case of an array, is that we need to define the size at the beginning of the implementation. This makes our stack static. It can also result in "stack overflow" if we try to add elements after the array is full. So, to alleviate this problem, we use a linked list to implement the stack so that it can grow in real time.



Stack: LIFO