

C++ STL

```
#include <bits/stdc++.h>
using namespace std;
```

→ C++ has lots of libraries math.h, string.h, everything covered by it.

- namespace `say`

```
int val = 50;
```

```
int main()
```

```
double val = 10.0;
```

```
cout << val << endl; → o/p: 10.0
```

```
cout << say::val << endl; → o/p: 50
```

- struct node

```
{
```

```
string str;
```

```
int num;
```

```
double doub;
```

```
char x;
```

```
node(str_, num_, doub_, x_)
```

```
str = str_;
```

```
num = num_;
```

```
doub = doub_;
```

```
x = x_;
```

```
}
```

```
int main()
```

```
{
```

```
node say = node("SAM", 79, 10.0, "0");
```

```
}
```

→ for Arrays -

```
int arr[5] = { }; → {1, 0, 0, 0, 0}  
0 initialized
```

```
int arr[5] = {0}; → {0, 0, 0, 0, 0}  
best way to initialize array w/ 0.
```

but if you want to
fill entire array with 1 value -

Container of STL [array<int, 3> arr; → if inside main(), + gq, gq, gq]
behaves as int arr[3];

```
∴ arr.fill(10); → {10, 10, 10} ←
```

now,

```
array<int, 4> arr;
arr.fill(1);
```

for this we can use - arr.at(index)

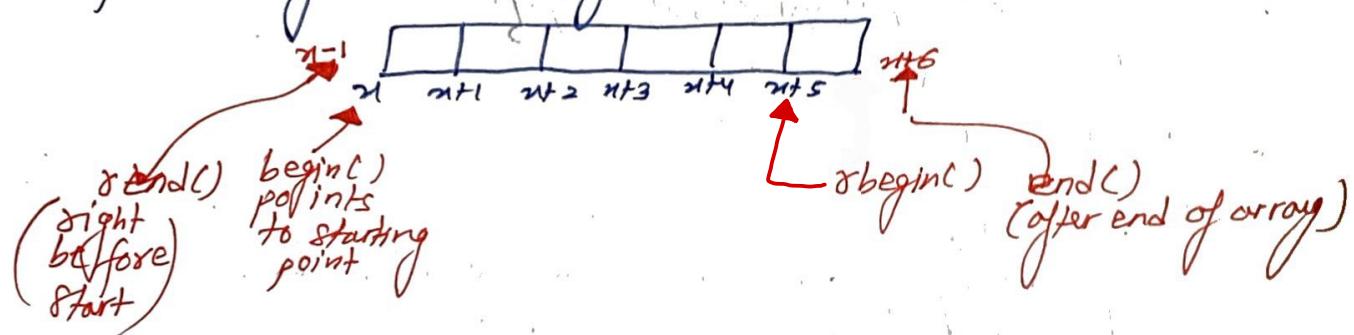
gives element at index

```
for (int i=0; i<4; i++)
    cout << arr.at(i) << endl; → o/p: 1, 1, 1, 1
```

Iterators

→ begin(), end(), ^{reverse begin} rbegin(), rend()] these iterators can be used in all DS.

for array, it is contiguous



to print

```
array<int, 5> arr = {1, 3, 4, 5, 6};
```

```
for (auto it = arr.begin(); it != arr.end(); it++)
    cout << *it << " "; → o/p: 1 3 4 5 6
```

rev print

```
for (auto it = arr.rbegin(); it != arr.rend(); it++)
    cout << *it << " "; → 6 5 4 3 1
```

→ 1, 3, 4, 5, 6
arr.begin()

not it--
as rbegin()
is for reverse
∴ wo automatically
it++ karte
rev jata rhega

(Q)

```
for (auto it = arr.end() - 1; it >= arr.begin(); it--)
    cout << *it << " "; → 6 5 4 3 1
```

`for(auto it : arr) // goes forward
 ↓
 element itself (assume the datatype)`

`cout << it << " "; → 1 3 4 5 6`

for string also -

`String s = "x heg";
 for(auto it : s)
 ↓
 cout << it << " "; → o/p: x h e g`

// size

`cout << arr.size();`

// front [to print front element]

`cout << arr.front(); → arr, arr.at(0);`

// back [to print last element]

`cout << arr.back(); → arr, arr.at(arr.size() - 1)`

Vector :

Note: Max size of array we can define in C++

`int arr[107]; // globally
 int main()
 {
 int arr[105]; // locally
 }`

for int, double, char array

→ 10⁷
 → 10⁸

`vector<int> arr; → a size of 0 or empty vector created`

`arr.push_back(1); → f 1`

`arr.size();
 for this
 → o/p = 0`

`arr.pop_back(); // pop out last element pushed in arr vector`

→ for max size of vector [same as above: for int, double, char → 10⁸
 ↴ 10⁷ locally
 ↴ 10⁸ globally

for bool → 10⁷ locally
 → 10⁸ globally

- `vector<int> v = {1, 2, 3, 4, 5};`

(Q8)

- `vector<int> v {1, 2, 3, 4, 5};`

`vec1.clear(); // erase all elements at once i.e. 4`
in vector

- `vector<int> vec(4, 0); // {0, 0, 0, 0}`

[after this also, we can push back]

`vector<int> vec(4, 8); // {8, 8, 8, 8}`

- `vector<int> vec1(4, 0);`

`vector<int> vec2(4, 8);`

`vector<int> vec3(vec2.begin(), vec2.end()); // for vec2, from begin() to end(), all elements of
vec2 copied to vec3`

- `vector<int> vec1;`

`vec1.push_back(1);`

`vec1.push_back(2);`

`vec1.push_back(3);`

`vec1.push_back(4);`

→ {1, 2, 3, 4}

now, we want vec1 elements from index = 0 to 1 i.e. 1 & 2 elements

∴ `vector<int> vec2 (vec1.begin(), vec1.begin() + 2);`

points to 1, 2, 3, 4

vec1.end() - 2

- `vec.empty();`

↳ tells vector is empty → returns T

- to define 2-D vector — not empty → " F

`vector<vector<int>> vec;`

`vector<int> vec1(2, 1); // {1, 1}`

`vector<int> vec2(3, 8); // {8, 8, 8}`

`vec.push_back(vec1);`

`vec.push_back(vec2);`

	0	1	2
0	1	1	1
1	8	8	8
2	1	1	1

→ here we have
`vec[i][j]` ↗
but not
`vec[i][j][k]` ✗

`for (auto it : vec)`
↳ here it is the vector itself

`for (auto it1 : it)`
↳ element of vector

`cout << it1 << "`

`cout << endl;`

↳ iterator val
Karte hai, at it
Know ends of every
vector of vector

→ o/p: 1 1
8 8 8

Q8, other way -

```
for(int i=0; i<vec.size(); i++)  
{  
    for(int j=0; j<vec[i].size(); j++)  
    {  
        cout << vec[i][j] << " ";  
    }  
    cout << endl;  
}
```

- to define 10×20 array [2D-vector]
 \leftrightarrow dimension

```
vector<vector<int>> vec(10, vector<int>(20, 0));
```

↳ we can do like -

```
vec.push-back (vector<int>(20, 0));
```

0	0 1 ... 19
1	0 0 0 0 ...
2	0 0 0 0 ...
3	0 0 0 0 ...
4	0 0 0 0 ...
5	0 0 0 0 ...
6	0 0 0 0 ...
7	0 0 0 0 ...
8	0 0 0 0 ...
9	0 0 0 0 ...

- vector<int> arr[4];

fairie
child cannot
do for it
we can do -

arr[1].push-back(0);
ie here, size of
array fixed here
[# of vectors = 4 fixed]

↳ array of vectors

arr	Vector 1 3	vector 0	vector 1 2	vector 3
	0	1	2	3

- to define 3D vector -

say, $10 \times 20 \times 30$

```
vector<vector<vector<int>>> vec(10, vector<vector<int>>(20,  
vector<int>(30, 0))
```

- rec.front();

rec.back();

rec.clear(); it clears all elements in vector, making it size=0

↳ rec.erase() -

To erase 1 element (I) rec has {1, 2, 3, 4, 5}
vector<int>::iterator it;
it = rec.begin();
rec.erase(it);

now rec has {2, 3, 4, 5}

To erase range of elements:

(II) rec has {1, 2, 3, 4, 5}
vector<int>::iterator it1, it2;
it2 --; { → here
it2 --; }
vec.erase(it1, it2);
1 2 3 4 5
it1 it2

if we do -
if for(auto it : rec)
{ cout << it << " "; }

Set:

(I)

ordered set -

given n -elements, tell me the no. of unique elements?

$\text{arr}[] = \{2, 5, 2, 1, 5\} \rightarrow$ there are 3 unique elements $\{1, 2, 5\}$

to get unique elements -

`set<int> st;`

`int n;`

`cin >> n;`

`for (int i=0; i<n; i++)`

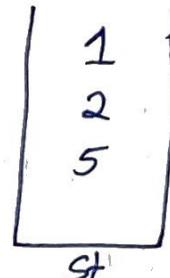
`{ int x;`

`cin >> x;`

`st.insert(x);`

$\hookrightarrow O(\log N)$

\downarrow
size of
set



here, st has $\{1, 2, 5\}$ in sorted order
 first element of set and element of set "3rd"

we cannot access set as:

$st[0] \times$

$st[1] \times$

to access set,

* $st.begin() \rightarrow$ gives 1

* $st.begin() + 1 \rightarrow$ 2

(Q8)

can use iterators.

• $st \{1, 2, 5\}$ say

① $st.erase(st.begin());$ // delete first element

from set st

Syntax:
 $st.erase(iterator);$

$\therefore st$ is $\{2, 5\} \bowtie$

② $st.erase(st.begin(), st.begin() + 2);$ // $st.erase(start_iterator, end_iterator)$

for $st \{1, 2, 5\}$

$\begin{matrix} \nearrow \\ \text{begin} \end{matrix}$ $\begin{matrix} \nearrow \\ \text{begin} + 2 \end{matrix}$

$\therefore \begin{matrix} \nearrow \\ 1 \end{matrix} \quad \begin{matrix} \nearrow \\ 2 \end{matrix} \quad \begin{matrix} \nearrow \\ 5 \end{matrix} \rightarrow \text{erased}$

$\rightarrow st$ is $\{5\}$

$\left[\begin{matrix} st.erase(\dots) \\ O(\log n) \end{matrix} \right]$

③ for $st \{1, 3, 5\}$ if 5 not there, then it'll erase nothing

$st.erase(5);$ // Syntax: $st.erase(key);$

$\hookrightarrow st$ is $\{1, 3\}$

- to declare set - [whatever we learnt in vector is applicable over here]

① `set<int> st = {1, 5, 7, 8}`

`set<int> st1(st);` //copy set to st1

- to find iterator of an element/key

$st \{ 1, 5, 7, 8 \}$
↑
points to 7

`auto it = st.find(7);`
↳ $O(\log n)$

Note: `auto it = st.find(key)`

if key does not exist in set st
then $it = st.end();$, i.e. $1, 5, 7, 8$

- //size

`st.size()` → gives size of set st

- to store 'node' → first page main like has

∴ `set<node> st;`

st will have unique node

- to print set -

```
for(auto it = st.begin(); it != st.end(); it++)  
    cout << *it << "
```

(Q8)

```
for(auto it : st)  
    cout << it << endl;
```

- to delete entire set -

`st.erase(st.begin(), st.end())`

(Q8)

`st.clear()`

II Unordered set - (will have unique elements)

unordered_set<int> st;

st.insert(2);
st.insert(3);
st.insert(1);

→ don't give guarantee set will be sorted, set can be any order.

all the operations is same as ordered set.

for Unordered set → Avg TC is $O(1)$ but WC is $O(n)$
Ordered set → Avg TC is $O(\log n)$ set size

Note: if you don't want elements in ascending order
use unordered set else ordered set
↳ if using this TLE occurs,
then samjhao kch test cases
 $WC = O(n)$ layout has
∴ use ordered set

III Multiset - (Stores Duplicates also)

multiset<int> ms;

ms.insert(1);
ms.insert(1);
ms.insert(2);
ms.insert(2);
ms.insert(3);

→ ms is {1, 1, 2, 2, 3}

↳ multiset helps to store
elements in sorted fashion.

all functions used here is same as in set ordered.

↳ Note: for erase:

ms.erase(2); → {1, 1, 3}

↳ all instances
of 2 is erased

- auto it = ms.find(2);
↳ returns iterator
pointing to first 2 in ms.

- ms.clear(); → delete entire set

ms.erase(ms.begin(), ms.end());

TC for every function in multiset = $O(\log n)$

size of set
↑
multiset

```

• for(auto it : st)
  {
    cout << it << " ";
  }
  } to print
  elements
  in multiset

```

Note : In ① set → ordered, unordered, multiset
 st.erase()

↳ can be ① iterator of an element
 ② range : start, end iterator
 ③ Key

② for vector, v.erase()

↳ can be ①

② ✓

③ X

• for ms → {1, 1, 2, 3, 3}

ms.erase(3); → delete all instances of 3 i.e. {1, 1, 2}

↳ to delete/erase only 1 instance of 3, use

ms.erase(ms.find(3))

↳ gives iterator of first instance of 3

→ to delete 2, 3, 3

∴ ms.erase(ms.find(2), ms.find(2) + 3))

• ms.count(3); ↳ O/P = 2

↳ tells how many times 3 occurs

→ MAP (default ordered)

↳ use key value concept

eg raj → 27
 hima → 31
 Sam → 1
 tank → 89

map<string, int> m;

m["raj"] = 27;

m["Hima"] = 31;

m["Sam"] = 1;

m["tank"] = 89

(Hima, 31)
(raj, 27)
(Sam, 1)
(tank, 89)

sorted acc
to the
keys

here, in lexicographical order
 ↳ string
 key is

(Same as ordered set, maps store in sorted order)

→ if we write again: m["raj"] = 29 (maps stores only unique keys)
 ↳ overwrites 27 (see table)

- we can write -

$$m["raj"] = 29$$

(or)

$m.\text{insert}(\text{"raj"}, 45);$

- $m.\text{erase}(\text{"raj"}); // m.\text{erase}(\text{key})$

↓ just give the key name
only single instance in m
↳ Deletes raj

- $m.\text{erase}(m.\text{begin}()); // Deletes first key$
↳ iterator position

- $m.\text{clear}(); // Deletes entire map$
(or)

$m.\text{erase}(m.\text{begin}(), m.\text{end}())$

- $m.\text{erase}(m.\text{begin}(), m.\text{begin}() + 2);$

here, ~~(Hima, 31)~~
deleted ~~(raj, 37)~~
~~(Sam, 1)~~
(Tank, 89)

- $\text{auto it} = m.\text{find}(\text{"raj"});$ → if raj does not exist, then it points to last key i.e. to "Tank" in m
↳ gives pointer to raj
i.e. first second
it → (raj, 37)
∴ string s = it → first;
int n = it → second;

- To check map is empty or not -

$\text{if}(m.\text{empty}())$

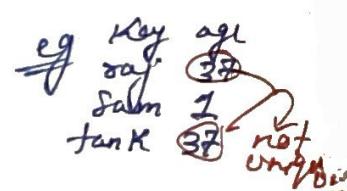
cout << "Yes it is empty";

- $m.\text{count}(\text{"raj"});$

↳ always returns 1 as only 1 key as "raj" is in map

TC of map = $O(\log n)$

Note - to delete a row from map →
use key & not any other values
as it is unique



• printing map

↳ maps store \langle string, int \rangle → a pair is pair \langle int, int \rangle p;
p.first = 1;
p.second = 20;

```
for(auto it : m)
    cout << it.first << " " << it.second << endl;
}
```

(08)
 for (auto it = m.begin(); it != m.end(); it++)
 { cout << it->first << " " << it->second << endl;

Unordered map Ordered

ordered map (ordered) \rightarrow $\{ \text{key} : \text{value} \}$

map ↗
↳ same as map, it uses same function

only thing $O(1)$ in all cases
 $\& O(n)$ in WC (bare)

$O(n^2)$ in worst case
size of map

- • `Unordered_map<int, int> m;`

↓
does not store in
any order (random
order)

→ Pair class-

• pair<int, int> p = {1, 2};
p.first is 1 → p.second is 2

- pair<pair<int,int>,int> p = {{1,2},3}

to access this

p-first-second ↵

- pair<pair<int,int>, pair<int,int>> p = { {1,2}, {3,4} };

- We can consider pair as ^{an} element - $\text{p} \cdot \text{first} \cdot \text{first} \Rightarrow 1$

```
vector<pair<int,int>> v;
```

```
set<pair<int,int>> st;
```

map < pair <int,int>, int> m

Key

Note: for unordered-map,

unordered_map<pair<int,int>, int> m;

X
not possible

** to store pair \rightarrow use map

→ store pair → ordered map
→ Unordered map can only store single keys.
→ same for unordered sets

III multimap → stores everything in sorted order

- $\text{multimap} < \text{string, int} \rangle m;$

↳ $m["raj"] = 2;$
 $m["raj"] = 4;$ ↳ stores duplicates.

[functions used here same as set]

Key	val
raj	2
raj	4

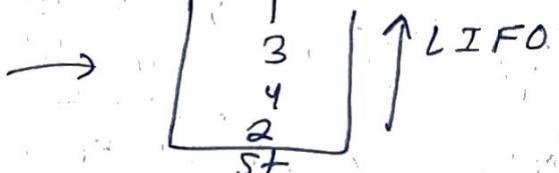
Stack & Queue

I
Stack

- $\text{stack} < \text{int} \rangle st; // LIFO DS$

all operations $TC = O(1)$

↳ $st.push(2);$
 $st.push(4);$
 $st.push(3);$
 $st.push(1);$



- $\text{int } x = st.top();$
↳ $x = 1$

- $st.pop();$



int $x = st.top();$
now
= 3

#elements in stack

- to delete entire stack $\rightarrow TC = O(n)$

Note : $\text{bool flag} = st.empty();$ ↳ returns T or F
empty ↳ not empty

while (!st.empty())
{
 st.pop();
}

- $st.size();$ // gives number of elements as stack

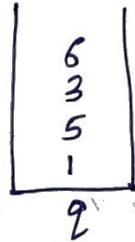
Note : $\text{stack} < \text{int} \rangle st;$
 $\text{cout} \ll st.top();$
↳ throw error as stack is empty

∴ use -

if (!st.empty()) // always have a check b
{
 cout << st.top();
}

II) Queue → All operations $O(1)$

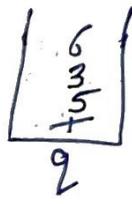
- queue<int> q;
- q.push(1);
- q.push(5);
- q.push(3);
- q.push(6);



- int x = q.front();

$$= 1$$

- q.pop(); \rightarrow



- to clear queue - $TC = O(n)$ elements in queue

```
while (!q.empty())
{
    q.pop();
}
```

Priority-Queue - (Heaps)

- set \rightarrow stored unique elements & sorted order

$\rightarrow TC = O(\log n)$ for every function

- unordered set \rightarrow stored unique elements

$\rightarrow TC = O(1)$ ($\& WC = O(n)$ same)

that's why we use priority queue

\hookrightarrow stored all in sorted order & $TC = O(\log n)$

priority queue has

- push()
- pop()
- top()
- empty()
- size()

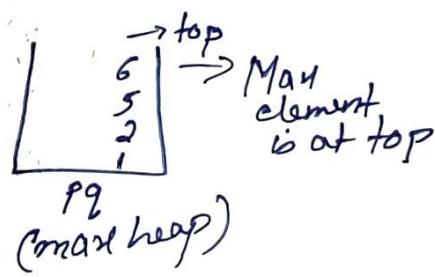
uses concept of heap sort

- priority-queue<int> pq;

- pq.push(1);
- pq.push(5);
- pq.push(2);
- pq.push(6);

\rightarrow pq [6 5 2 1] i.e.,

(in descending order)



- int $x = pq.top();$
= 6

- $pq.pop();$

```
int x = pq.top();
= 5
```



- if storing pairs in priority queue-

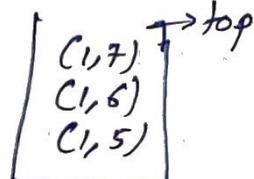
priority-queue<pair<int, int>> pq;

pq.push(1, 5);

pq.push(1, 6);

pq.push(1, 7);

priority queue
store acc to pairs



for min heap

(if you want pair, replace int w/ pair) here, first is same for all
:: sort acc to second

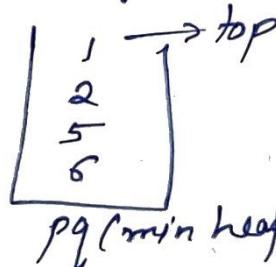
↳ priority-queue<int, vector<int>, greater<int>> pq; // syntax for min heap

pq.push(1);

pq.push(5);

pq.push(2);

pq.push(6);



- int $x = pq.top();$
= 1

→ a Doubly LL

List (behaves same as vectors but here we can push from front & back & pop from front & back)

- list<int> ls;

↳ 1) push-front()

2) push-back()

3) pop-front()

4) pop-back()

5) begin(), end(), rbegin(), rend() // for iterators

6) size()

7) clear()

8) at()

9) remove() → TC = O(1)

↳ to remove any element

eg) ls.push-front(1)

ls.push-front(2)

ls.push-front(3);

↳ (ls.remove(2));
all instances of removes 2 of

(PTO)

- `ls.remove(key)` \rightarrow remove all instances of value = key in list
 eg) $\{1, 2, 2, 3, 2\}$
 $\text{ls.remove}(2);$
 \rightarrow now $\{1, 3\}$

Question \rightarrow N elements take from Keyboard. Print the element that occurs max no. of times.

```

int n;
cin >> n;
map<int, int> m;
int maxi = 0;
for (int i = 0; i < n; i++) {
    int x;
    cin >> x;
    m[x]++;
    if (m[x] > m[maxi]) {
        maxi = x;
    }
}
cout << maxi;
    
```

Note

if map used

$TC = O(n) * O(\log n)$
 $= O(n \log n)$

if $m[0]$ does not exist
 returns NULL

if compared to int, converted to 0

if not exist
 returns NULL ie 0

if unordered_map used

$$TC = O(n) * O(1)$$
 $= O(n)$

\hookrightarrow in WC = $O(n^2)$ // gar TLE aya used map

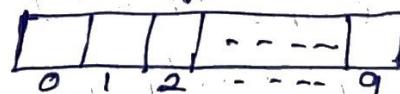
Bit Set (only stores 1 & 0)

generally, $\text{int} \rightarrow 16 \text{ bits}$
 $\text{char} \rightarrow 8 \text{ bits}$

$$\therefore \text{int } a[100]; \rightarrow 16b * 100$$
 $\text{char } a[100]; \rightarrow 8b * 100$

// bitset $\rightarrow 1 \text{ bit}$

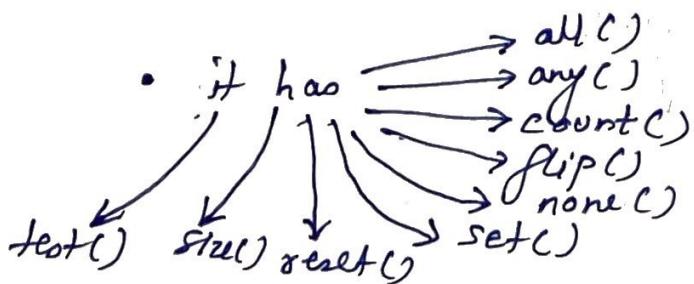
`bitset<10> bt;` \rightarrow creates array like structure
 takes less space



It is generally used in segment trees problems as it uses less space.

- `bitset<10> bt;`
`cin >> bt;`
 \hookrightarrow give: 101010001 (possible)

⑧



① say, $bt \{ 1, 1, 1, 1, 1 \}$

`cout << bt.all();` \rightarrow O/P: 1

\hookrightarrow returns true \rightarrow if all bits are set
 false \rightarrow if any of bits are not set

② $bt \{ 1, 0, 0, 0, 0 \}$

`cout << bt.any();`

\hookrightarrow if any of one bit is set \rightarrow returns true
 else false

③ $bt \{ 1, 0, 1, 1, 0 \}$

$bt.count();$ \rightarrow gives no. of bits that are set
 O/P = 3

④ $bt \{ 1, 0, 1, 0, 0 \}$

⑤ $bt.flip(2);$ \rightarrow O/P:
 bt becomes $\{ 1, 0, \underline{1}, 0, 0 \}$

$(\begin{matrix} 1 & \rightarrow & 0 \\ 0 & \rightarrow & 1 \end{matrix})$ \downarrow flip the bit at
 2nd index

⑥ $bt.flip();$ \rightarrow converts everything
 as $1 \rightarrow 0$
 $0 \rightarrow 1$
 \hookrightarrow i.e. 1's complement

⑤ $bt \{ 1, 0, 0, 0 \}$

$bt.none();$ \rightarrow O/P = 0

\hookrightarrow returns true if no bit is set

\hookrightarrow returns false if any of bits is set

⑥ set()

does not matter whether
 bt has 0 or 1,
 it will replace
 them with 1

\rightarrow to set entire bt with 1 \rightarrow $bt.set();$

\hookrightarrow bt $\{ 1, 0, 1, 0 \}$
 $bt.set(); \rightarrow \{ 1, 1, 1, 1 \}$

bt $\{ 1, 0, 1, 0 \}$

$bt.set(3); \rightarrow \{ 1, 0, 1, 1, 0 \}$

but this

$bt.set(2, 0);$ \rightarrow change whatever present at
 index = 2 with 0.

- reset()
 - `bt.reset();` → turns all indexes to 0
 - `bt.reset(2);` → turns the 2nd index to 0.
 - ~~`bt.reset(2,0)`~~ X (no such format)
- bt.size(); → #elements in bt
 $\text{bt}\{1,0,1,0\} \rightarrow \text{bt.size() is } 5$ [but `bt.count()` tells no. of 1's]
- test()
 - checks if a bit is set or not at index i & returns true → if set false → if not set
 - `bt.test(ind);`

→ Algorithms -

Sorting -

↳ for array, vector

- `int n;`
`cin >> n;`
`int arr[n];` → $TC = O(n \log n)$
`sort(arr, arr+n); // sort in ↑ order`
- `vector<int> vec{1, 5, 6, 7, 2};`
`sort(vec.begin(), vec.end()); // []`
 ↳ start iterator ↳ last iterator
- say `vec → {1, 6, 2, 7, 4}`
 sort it from indexes 1 to 3 → final o/p: 1, 2, 6, 7, 4
`sort(vec.begin() + 1, vec.begin() + 4);`

Reverse → $TC = O(n)$

↳ array, vector

- `reverse(arr, arr+n);`
- `reverse(start iterator, end iterator);`
 ↳ `reverse(vec.begin(), vec.end());`
- to reverse element in index from 1 to 3 i.e. $\{1, 5, 6, 7, 2\} \xrightarrow{\text{o/p}} \{1, 7, 6, 5, 2\}$
`reverse(vec.begin(), vec.begin() + 4);`

→ To find the max elements in any index range say i to j

normal method

```
int maxi = INT_MIN;  
for (int k = i; k <= j; k++)  
    if (arr[k] > maxi)  
        maxi = arr[k];
```

$T.C = O(n)$

→ int maxi = *max_element(arr, arr+n);

for min element,

```
int mini = *min_element(arr, arr+n);
```

$T.C = O(n)$

returns iterator

for vectors: (vec.begin(), vec.end());

for vectors:
(vec.begin(), vec.end())

→ Given a range, find sum in that range. (i to j)

normal method

```
int sum = 0;  
for (int k = i; k <= j; k++)  
    sum += arr[k];
```

→ int sum = accumulate(arr, arr+n, 0);

$T.C = O(n)$

for vector-

(vec.begin(), vec.end(), 0)

(initial sum value
can be 10, 20, ...)

→ Given arr[7] → {1, 6, 7, 1, 2, 1, 3} & $x = 1$.

tell how many times 'x' (or 1) occurs in in array.

normal method

```
int cnt = 0;  
for (int i = 0; i < n; i++)  
    if (arr[i] == x)  
        ++cnt;
```

$\rightarrow T.C = O(n)$

→ count(first iterator, last iterator, x);

\therefore int cnt = count(arr, arr+n, x);

for vector :
(vec.begin(), vec.end(), x);

→ given arr[] → {1, 2, 5, 1, 2, 4, 4}

find the first occurrence of 2:

↳ at index = 2

```

int ind = -1;
for(int i=0; i<n; i++)
{
    if(arr[i] == x)
    {
        ind = i;
        break;
    }
}

```

→ auto it = find(arr, arr+n, x) // return pointer pointing to first instance of x if x is present else returning pointer pointing to end() if not there

we don't get the index
but since array is contiguous

to get index no. -
int ind = (it - arr); → &*ind → get element

for vector:

auto it = find(vec.begin(), vec.end(), x);

int ind = it - vec.begin();

↳ *ind gives element

say, vec → {1, 5, 6, 2, 3, 5, 6}

x = 4

auto it = find(vec.begin(), vec.end(), 4);

vec [1|5|6)2|3|5|6]

∴ if(it == vec.end())

cout << "element not present";

iterator will be at end()

else cout << "element first occurs at:" << it - vec.begin();

Binary Search

$TC = O(\log n)$

↳ works on sorted arrays/vectors [but `find()` works on every array]
 $\text{arr}[] \rightarrow \{1, 5, 7, 9, 10\}$ sorted/not sorted
 $x = 9$ $TC = O(n)$]

Syntax: `binary-search(first iterator, last iterator, x);`
 ↳ returns true → if x present
 false → if not present

`bool res = binary-search(arr, arr+n, 8);`
 ↳ returns false;

↳ implement in BS $TC = O(\log n)$ for vector:

→ lower_bound function (for sorted) $(\text{vec.begin()}, \text{vec.end()}, x)$:

↳ returns an iterator pointing to first element which is not less than x .

$\text{arr}[] \rightarrow \{1, 5, 7, 7, 8, 10, 10, 11, 11, 12\}$

for $x = 10$;

for $x = 6$;

for $x = 13$;

points to `end()`

`auto it = lower_bound(arr, arr+n, x);`

`int ind = it - arr;`

↳ for vector:

$(\text{vec.begin()}, \text{vec.end()}, x)$;

→ upper_bound [implemented in BS & $TC = O(\log n)$]

↳ returns an iterator which points to an element which is just greater than x .

$\text{arr}[] \rightarrow \{1, 5, 7, 7, 8, 10, 10, 10, 11, 11, 12\}$

$x = 7$

$x = 6$

$x = 12$

$x = 15$ same

`end()`

`auto it = upper_bound(arr, arr+n, x);`

`int ind = it - arr;`

↳ for vector:

$(\text{vec.begin()}, \text{vec.end()}, x)$;

vec.begin()

① if (binary-search(a_{xx}, a_{xx}+n, x) == true)
 { cout << lower-bound(a_{xx}, a_{xx}+n, x) - a_{xx};
 } else { cout << "does not exist"; } // O(log n)
 ② int ind = lower-bound(a_{xx}, a_{xx}+n, x) - a_{xx};
 if (ind != n && a_{xx}[ind] == x)
 cout << ind;
 else cout << "not found"; // x not present

Q> find me the last occurrence of 11 in arr & arr is sorted.
 say arr[7] = {1, 5, 7, 7, 8, 10, 10, 10, 11, 11, 12}.

Soln

int ind = upper_bound(arr, arr+ n , x) - arr;

if (ind > 0 && arr[ind-1] == x) \Rightarrow if $x=0$
 cout << ind-1;
 else cout << "not found";
 \Rightarrow x not present

ans
 if $x=0$
 s: 7, 7, 8, ...
 UB
 $\therefore (ind-1)$
 here
 if $x=14$
 1, 5, 7, ..., 13, 12, 11

Q) Tell me the no. of times 2 appears in 288?

eg $arr[7] \rightarrow \{1, 5, 7, 7, 8, 10, 10, 10, 11, 11, 12\}$
 $\text{if } x = 10 \rightarrow \text{ans} = 3$

Soln 0 5 7 7 8 10 10 10 11 11 12
 ↑ ↑
 U_b U_b
 for $x=10$ $\therefore U_b - U_b = 3.44$

$\therefore \text{int } \text{cnt} = \text{upper_bound}(\text{arr}, \text{arr} + n, n);$ (∴ no need to write external checks)

→ Next permutation → $TC = O(n)$

String $s = "abc"$
all permutations are as follows:

abc
acb
bac
bca
cab
cba

all are in lexicographical sorted order.

we can have array abc

i.e. $arr[7] \rightarrow [2, 1, 3]$,

∴ $\begin{matrix} 1 & 2 & 3 \\ 1 & 3 & 2 \\ 2 & 1 & 3 \\ 2 & 3 & 1 \\ 3 & 1 & 2 \\ 3 & 2 & 1 \end{matrix}$ ↴ sorted

say, $s = "bca"$

bool res = next_permutation(s.begin(), s.end());

↳ after this: bca ki baad cab hai ↳ it returns true as it gets to next permutation.
∴ s becomes $s = "cab"$

if $s = "cba"$

∴ bool res = next_permutation(s.begin(), s.end());

here, after cba hai naikai

∴ no next permutation & s remains same but ↴ res = false

Q) given any random string $s = "bca"$
print all the permutations.

soln

String $s = "bca";$

fout(s.begin(), s.end()); // merges the string "abc"

do {

cout << s << endl;

} while(next_permutation(s.begin(), s.end()));

∴ abc
acb
bac
bca
cab
cba

↓ lexicographically sorted

→ #permutations = $O(n!)$

$TC = O(n \log n) + O(n!)$
to sort ↴ to which loop

Similarly we have prev-permutation()
prev-permutation(s.begin(), s.end())

- sort(arr, arr+n) // sorts in ascending order

for descending order

↳ we have to write comparator functn

↳ returns true or false

$\begin{array}{l} \{abc \\ acb \\ bac \\ bca \\ cab \\ cba \} \\ \downarrow \\ \text{if } s = bac \\ \text{if op} = acb \\ \therefore \text{in previous} \\ \text{page,} \\ \text{for prev-pm} \\ \text{don't sort} \end{array}$

bool comp(int ele1, int ele2)

consider 2 elements

ele1 ele2

if ele1 is before ele2 return true
else false

(soche 2 element ko sort karna ha)

int main()

sort(arr, arr+n, greater<int>())

- pair<int, int> arr[] = {{1, 4}, {5, 2}, {5, 9}};

// sort in such a way that the element who have first element in pair is smaller

[if first is equal, sort according to second
& keep larger second]

↳ op: 1, 4

as elements are pairs

5, 4 } **

5, 2 }

bool comp(pair<int, int> ele1, pair<int, int> ele2)

if (ele1.first < ele2.first)
return true;

if (ele1.first == ele2.first) // if equal

if (ele1.second > ele2.second)

return true;

} return false;

int main()

{

sort(arr, arr+3, comp);

}

but if we do-
sort (arr, arr+3);
↳ sorts everything in
ascending order
i.e. (1, 4)
(5, 2)
(5, 9)

We can sort: array, vector, string (Linear O.S)