# Data Structures Chapter 4

- 1. Singly Linked List
  - Pointer Reviewed & Linked
  - Linked List (1)
  - Linked List (2)
  - Reverse Singly Linked List
    - in-place O(n)
    - using stack O(n)
    - sub-list reverse  $O(n^2)$ ,
    - sub-list reverse O(n)

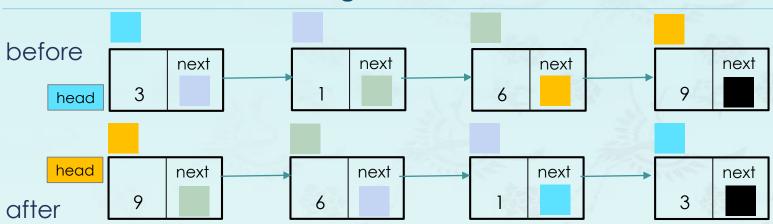


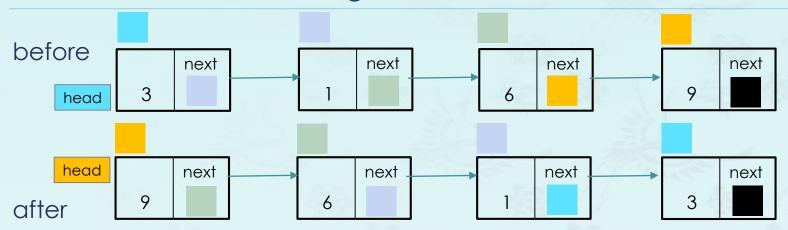
하나님이 우리를 구원하사 거룩하신 소명으로 부르심은 우리의 행위대로 하심이 아니요 오직자기의 뜻과 영원 전부터 그리스도 예수 안에서 우리에게 주신 은혜대로 하심이라 (딤후1:9)



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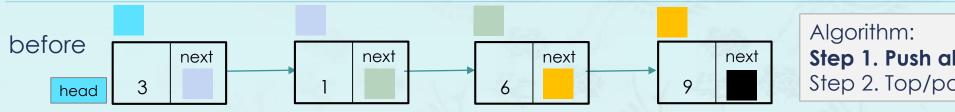
God has saved us and called us to a holy life – not because of anything we have done but because of his own purpose and grace. This grace was given us in Christ Jesus before the beginning of time. 2 Tim1:9



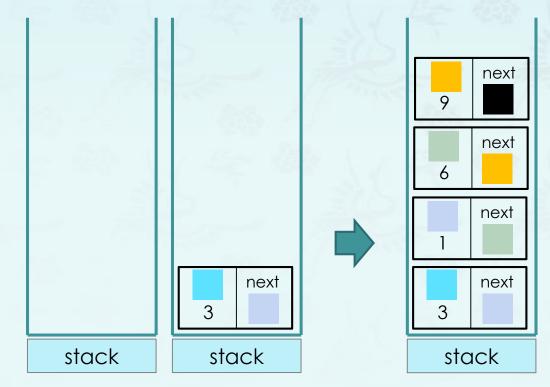


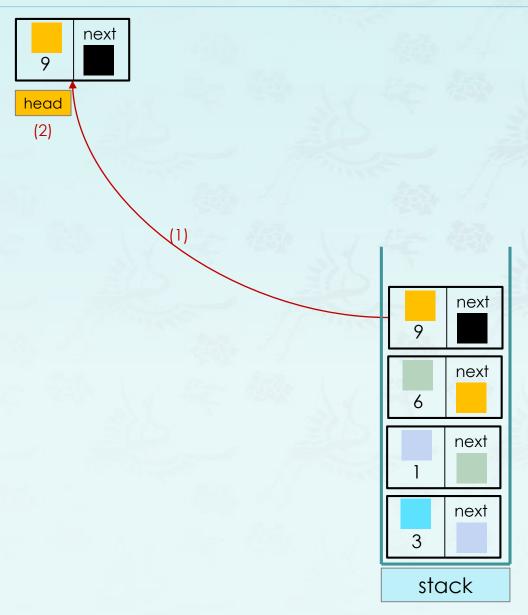
#### Algorithm:

Step 1. Push all nodes onto the stack. Step 2. Top/pop all nodes and relink.



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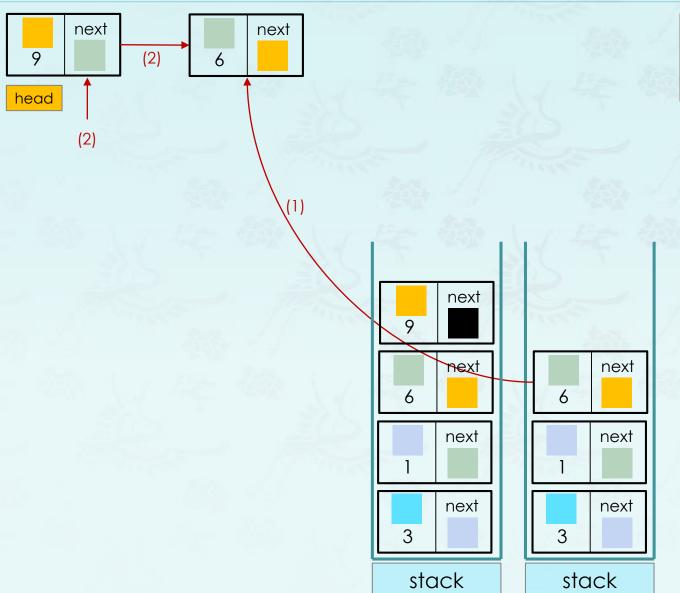




Algorithm:

Step 1. Push all nodes onto the stack.

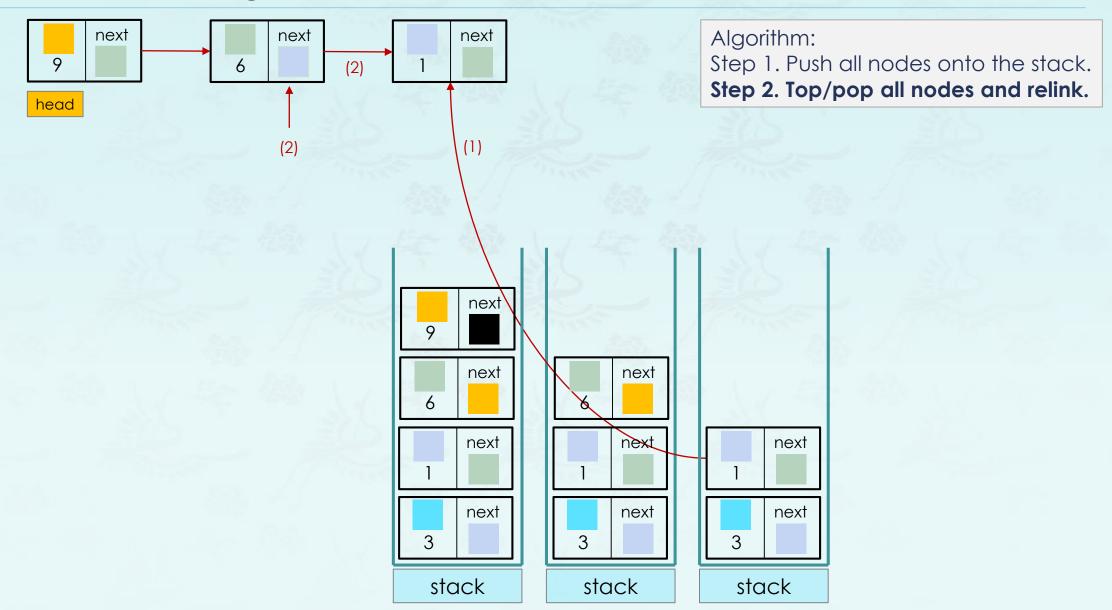
Step 2. Top/pop all nodes and relink.

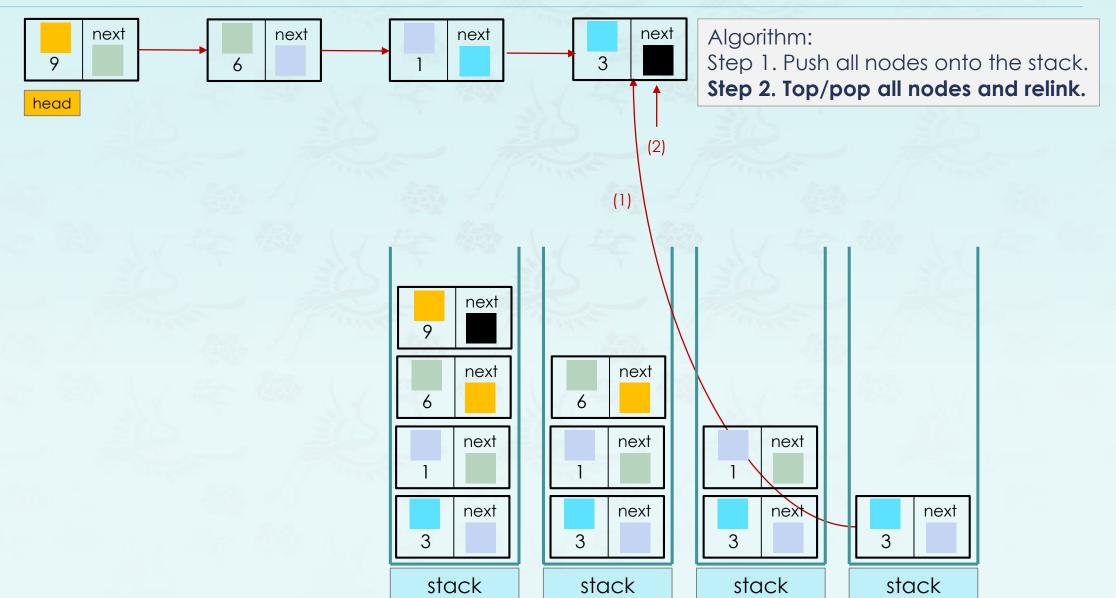


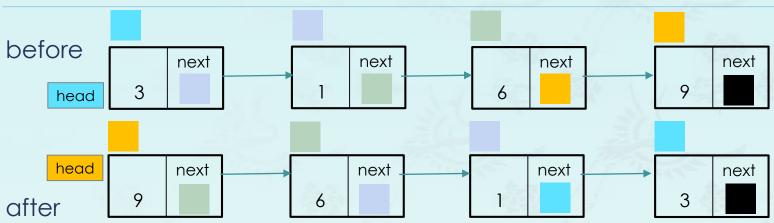
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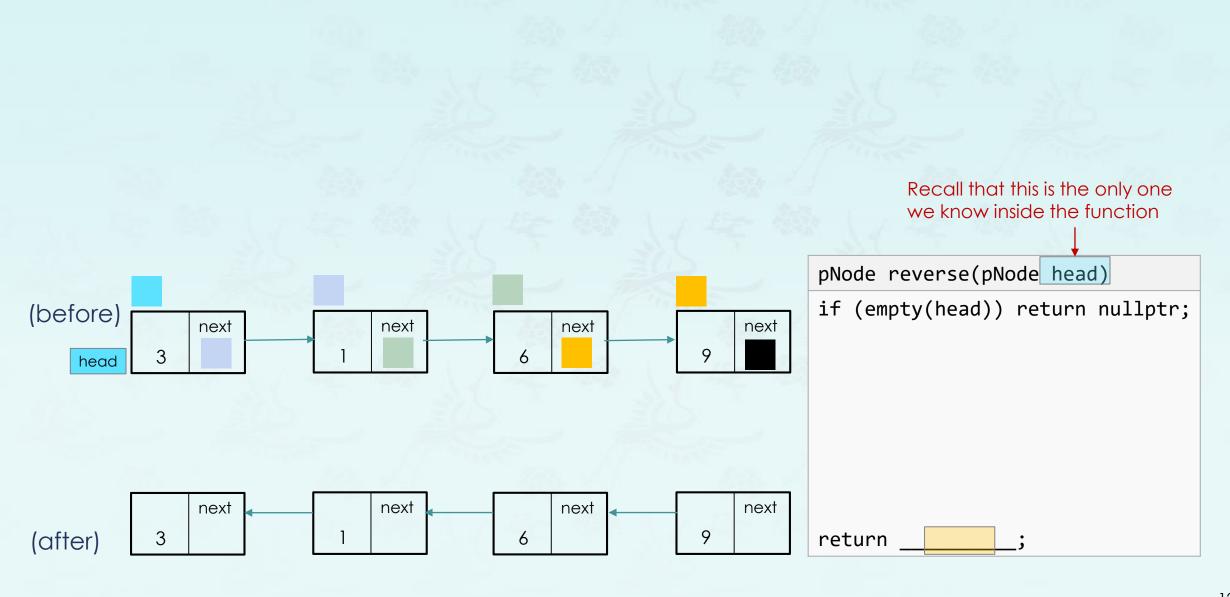


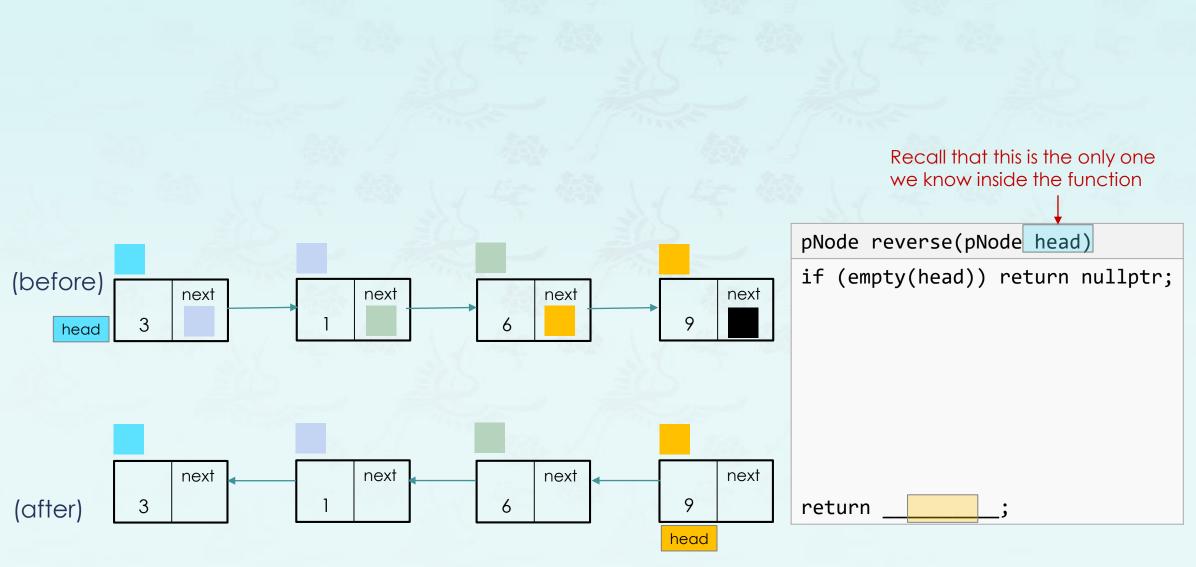


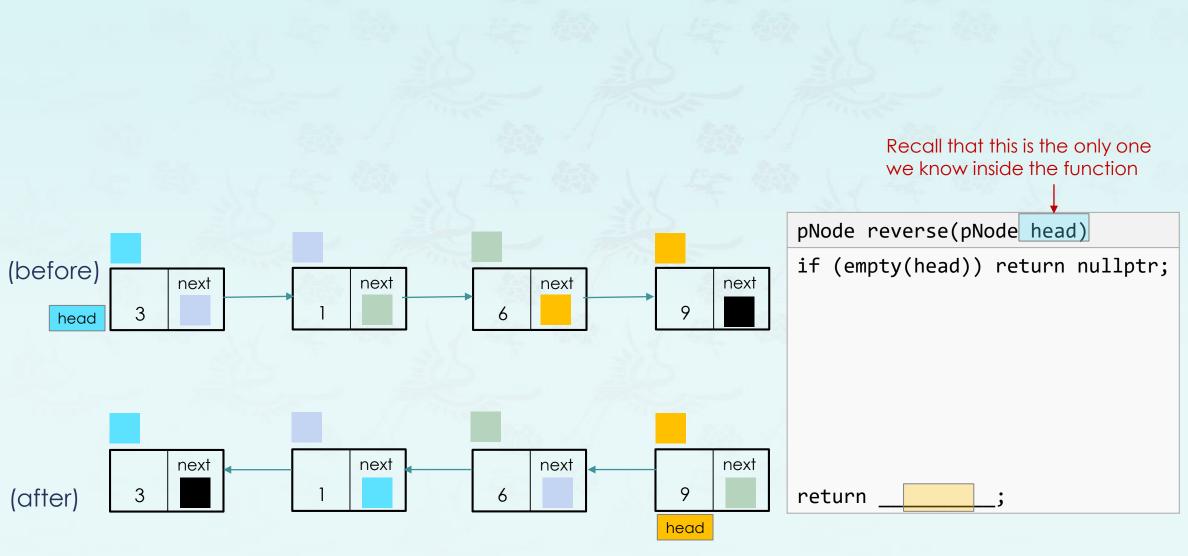


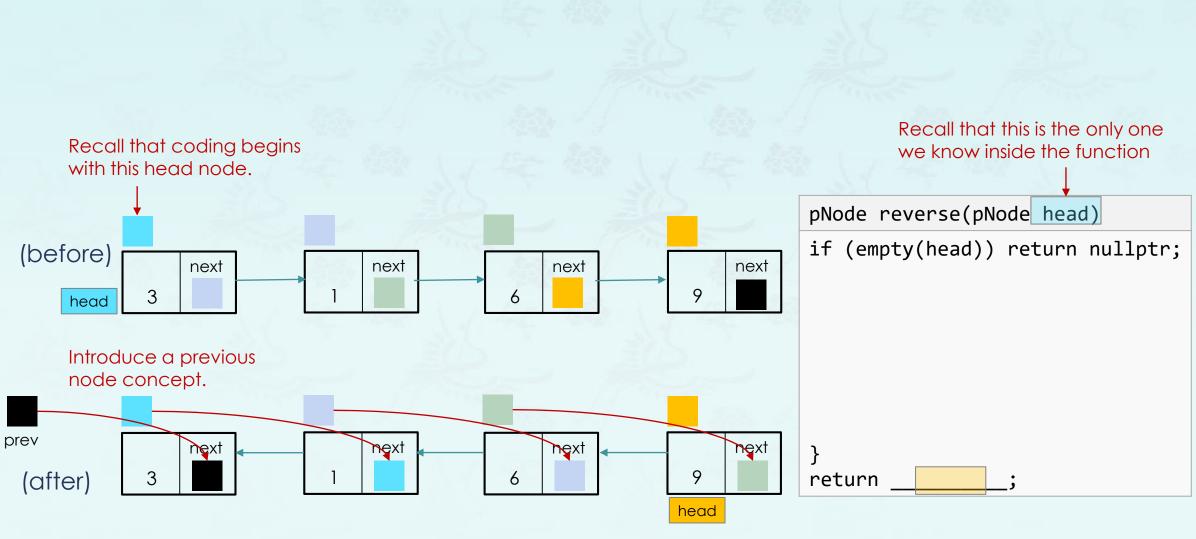
```
pNode reverse(pNode head)

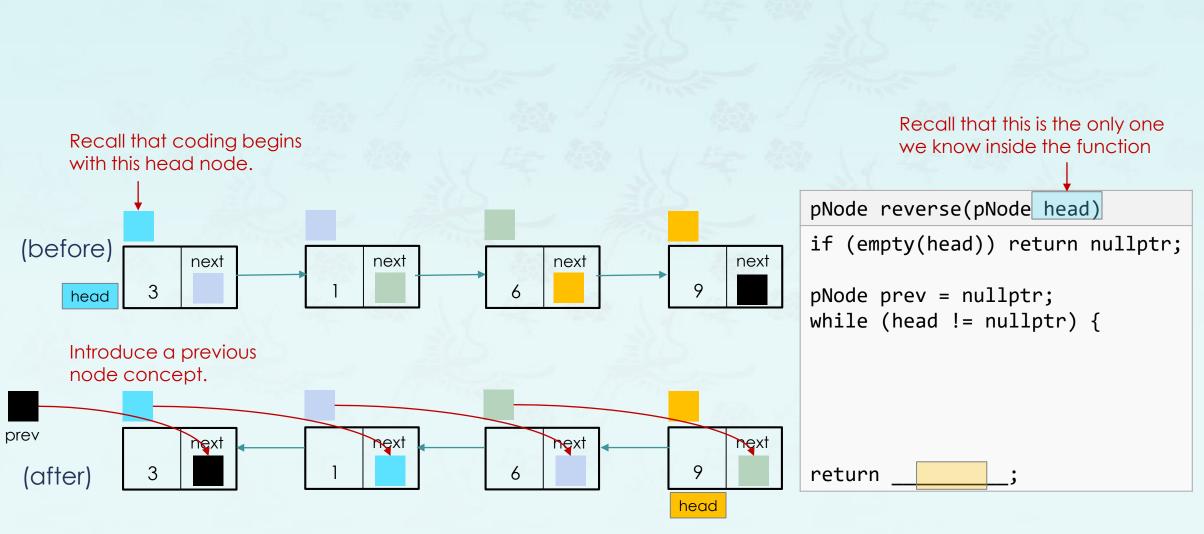
if (empty(head)) return nullptr;
while( list is not empty )
  get a node from list
  push it onto the stack
}
while( stack is not empty )
  get a node from the stack
  relink it back the new list
}
return head; // new head
```











#### Linked List - insert()

**TASK:** Code a function that inserts a node(5) at a node position x specified by a value(7).

- If the first node(or head) is the position, then just invoke push\_front().
- As observed below, we must to know **the pointer x** which is stored in the **previous node** of node x.

```
pNode insert(pNode head, int val, int x)
if (head->data == x)
  return push_front(val, head);
pNode curr = head;
pNode prev = nullptr;
while (curr != nullptr) {
  prev = curr;
  curr = curr->next;
return head;
```

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if (head->data == x)
  return push_front(val, head);
pNode curr = head;
pNode prev = nullptr;
while (curr != nullptr) {
  if (curr->data == x) {
               = new Node{
                                          };
    return head;
  prev = curr;
  curr = curr->next;
return head;
```

#### Linked List - insert()

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```
head prev curr

7 (before)

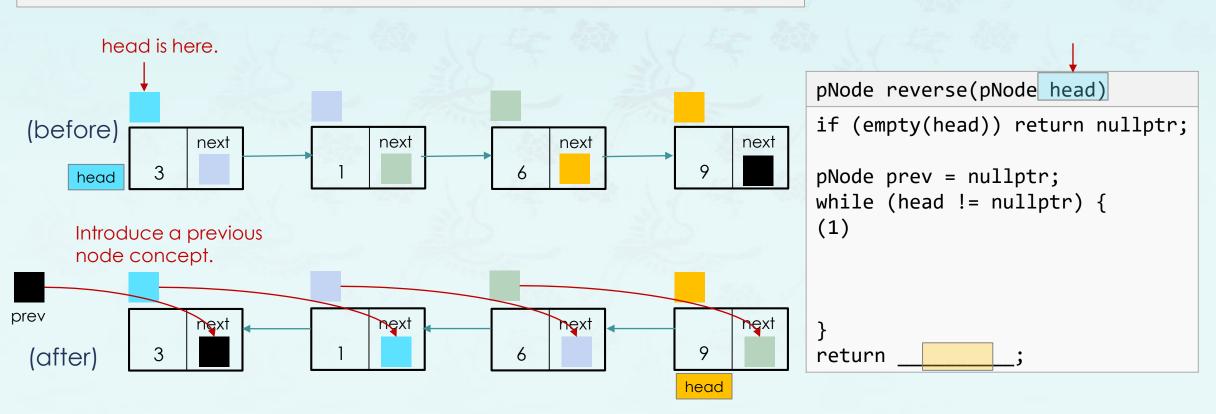
X

head prev new curr (after)

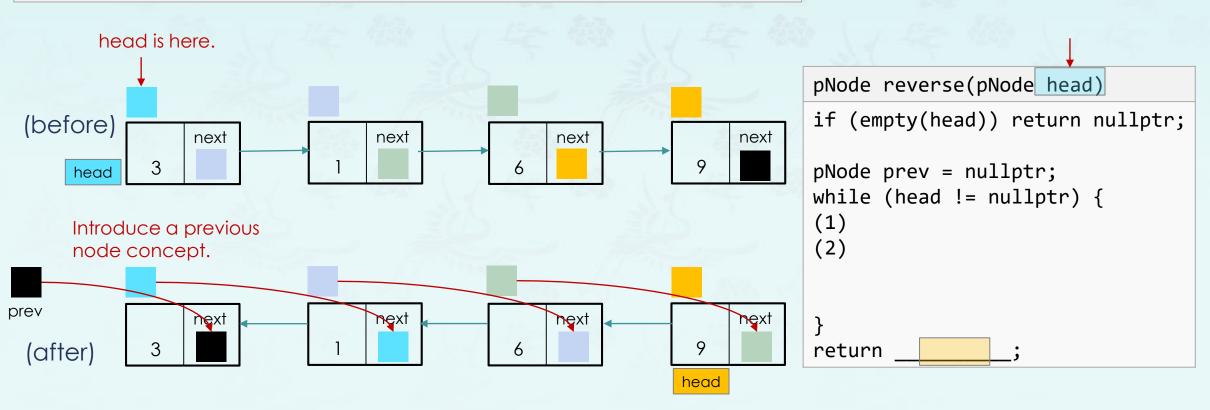
5 7 7
```

```
pNode insert(pNode head, int val, int x)
if (head->data == x)
  return push_front(val, head);
pNode curr = head;
pNode prev = nullptr;
while (curr != nullptr) {
 if (curr->data == x) {
    prev->next = new Node{val, prev->next};
   return head;
  prev = curr;
  curr = curr->next;
return head;
```

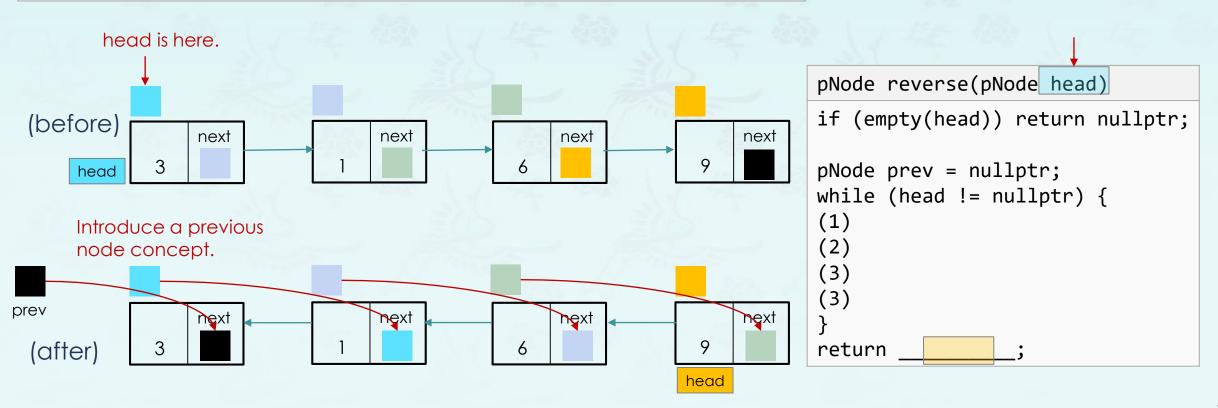
1. We want to overwrite head->next with prev
But before overwriting we must save head->next .
because we it is the next node we need to process.



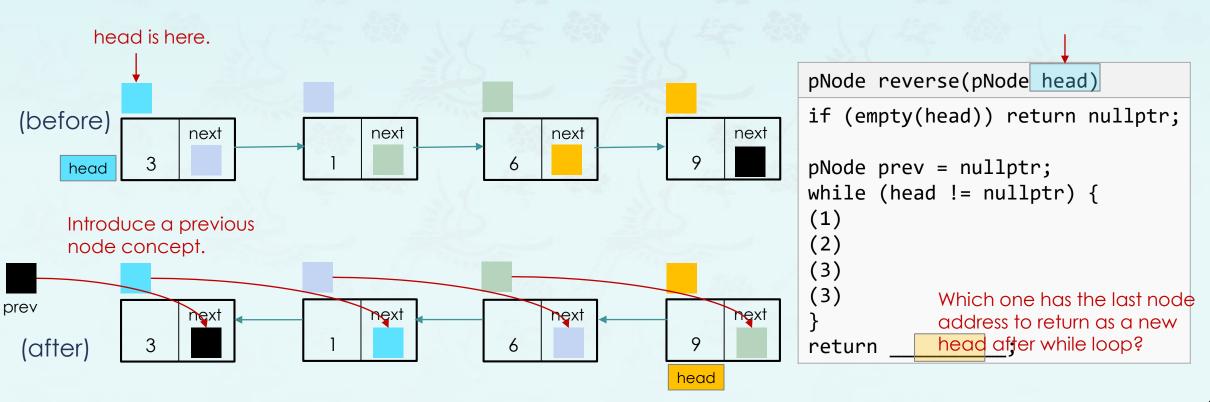
- 1. We want to overwrite **head->next** with **prev**But before overwriting we must save **head->next**because we it is the next node we need to process.
- 2. Once we overwrite **head->next** with **prev**



- 1. We want to overwrite **head->next** with **prev**But before overwriting we must save **head->next**because we it is the next node we need to process.
- 2. Once we overwrite **head->next** with **prev**
- 3. be ready to process the next node by setting prev as and head as .



- 1. We want to overwrite **head->next** with **prev**But before overwriting we must save **head->next**because we it is the next node we need to process.
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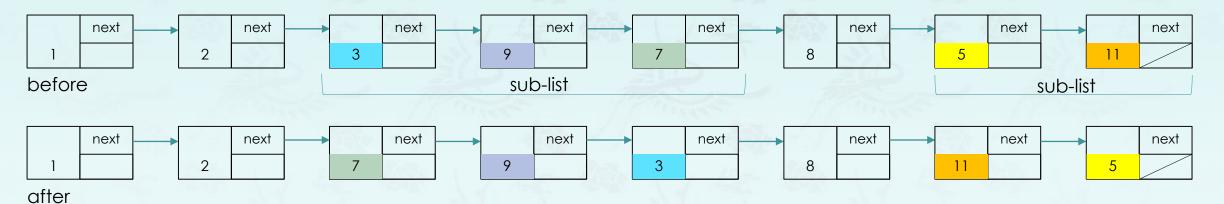


#### Linked List – reverse elements in sub-lists of odd numbers

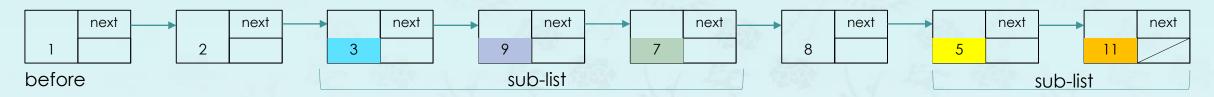
**TASK:** Reverse elements in sub-lists of odd numbers only in a singly-linked list.

Given a linked list that contains N integers, select all the sub-lists contain only odd integers. Reverse elements in those sub lists only.

**For example,** if the list is {1, 2, 3, 9, 7, 8, 5, 11}, then the selected sub-lists are {3, 9, 7} and {5, 11}. Reverse elements in those list such as {7, 9, 3} and {11, 5}. Now, this function returns the original list except odd numbers reversed in the sub-lists. In this example, it returns {1, 2, 7, 9, 3, 8, 11, 5}



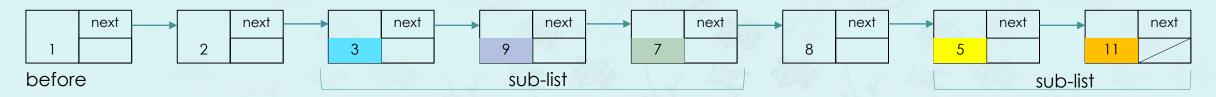
#### Linked List – reverse elements in sub-lists of odd numbers (version 1 – $O(n^2)$ )



```
while (head != nullptr) {
                             version 1 - O(n^2)
  if the node is odd {
    push it to odd_stack
    go for the next node
    continue
  // even node encountered
 while (odd stack is not empty) {
    get top of odd_stack & pop
    push_back to the head2
  add even node to head2
                                 added
 go for the next node
while( odd_stack is not empty) {
 get top of odd_stack & pop
  push back to the head2
clear head
return head2
```

- For the sake of the simplicity of coding, we use push\_back().
- head is the original list head.
- head2 is the new list as a result.
- odd stack stacks up odd(s) until an even shows up.
- You may use either stack<Node\*> or stack<int> but
   recall that push\_back() takes a data item, not a node itself.

#### Linked List – reverse elements in sub-lists of odd numbers (version 2 – O(n))



```
while (head != nullptr) {
                            version 2 – O(n)
  if the node is odd {
    push it to odd_stack
   go for the next node
    continue
  } // even node encountered
 while (odd stack is not empty) {
    get top of odd_stack & pop
    head2 is null, set head2
    add it to the tail2 & set tail2
 head2 is null, set head2
  add even node to tail2 & set tail2
  go for the next node
while( odd_stack is not empty) {
 get top of odd_stack & pop
 head2 is null, set head2
  add it to the tail2 & set tail2
 // no clear head necessary
return head2
```

- For the sake of the speed of the code, do not use push\_back().
  Use almost the same algorithm, but manage to add a node at the head2 and tail2 by yourself instead of calling push\_back().
- head is the original list head.
- head2 is the new list as a result.
- tail2 is the tail node of the head2.
- odd\_stack stacks up odd(s) until an even shows up.
- Do not use stack<int>, but stack<Node\*> to reuse the nodes.
- Do not clear head since all nodes are reused.

## Data Structures Chapter 4

- 1. Singly Linked List
  - Pointer Reviewed & Linked
  - Linked List (1)
  - Linked List (2)
  - Reverse
- 2. Doubly Linked List

