Binary Search Tree

- Recursion Revisited
- binary search tree Implementation
 - traversal inorder, preorder, postorder, levelorder
 - minimum, maximum,
 - predecessor, successor
 - height
 - clear
 - contains
 - grow
 - trim

bunnyEars(): counting bunny ears in recursion

```
// each bunny has two ears.
int bunnyEars(int bunnies) {
    return 2 + bunnyEars(bunnies-1);
}
```

funnyEars(): counting funny ears in recursion

```
// even numbered funny has two ears, odd numbered funny three.
int funnyEars(int funnies) {
  if (bunnies == 0) return 0;

  if (funnies % 2 == 0)
    return
  else
    return
}
```

bunnyEars(): counting bunny ears in recursion

```
// each bunny has two ears.
int bunnyEars(int bunnies) {
  if (bunnies == 0) return 0;
  return 2 + bunnyEars(bunnies-1);
}
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funnyEars(): counting funny ears in recursion

```
// even numbered funny has two ears, odd numbered funny three.
int funnyEars(int funnies) {
  if (bunnies == 0) return 0;

  if (funnies % 2 == 0)
    return 2 + funnyEars(funnies - 1);
  else
    return 3 + funnyEars(funnies - 1);
}
```

size(): in doubly linked list

```
int size(pList p) {
  int count = 0;
  for (pNode c = begin(p); c != end(p); c = c->next)
     count++;
  return count;
}
```

size(): in singly linked list

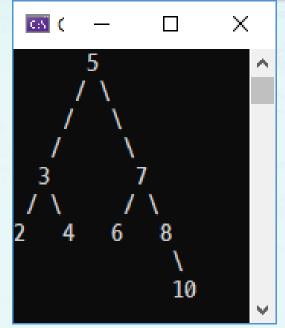
```
int size(pNode node) {
  if (node->next == nullptr) return 0;
  return 1 + size(node->next);
}
```

```
int size(tree node) {
  if (node == nullptr) return 0;
  return
}
```

```
int size(tree node) {
  if (node == nullptr) return 0;
  return 1 + size(node->left)
}
```

```
int size(tree node) {
  if (node == nullptr) return 0;
  return 1 + size(node->left) + size(node->right);
}
```

```
int size(tree node) {
   if (node == nullptr) return 0;
   cout << " size at: " << node->key << endl;
   return 1 + size(node->left) + size(node->right);
}
```



```
int size(tree node) {
  if (node == nullptr) return 0;
  cout << " size at: " << node->key << endl;</pre>
  return 1 + size(node->left) + size(node->right);
CS (
                                                                     X
                                                   size at:
                                                   size at:
            10
```

```
int size(tree node) {
  if (node == nullptr) return 0;
  cout << " size at: " << node->key << endl;</pre>
  return 1 + size(node->left) + size(node->right);
CS (
                                                                     X
                                                   size at: 5
                                                   size at: 3
                                                   size at: 2
                                                   size at: 4
                                                   size at: 7
                                                   size at: 6
                                                   size at: 8
                                                   size at: 10
            10
```

size: Count the number of nodes in the binary tree recursively.

```
int size(tree node) {
  if (node == nullptr) return 0;
  return 1 + size(node->left) + size(node->right);
}
```

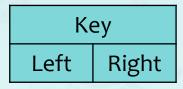
height: compute the max height(or depth) of a tree.

// It is the number of nodes along the longest path from the root node

// down to the farthest leaf node.

```
int height(tree node) {
}
```

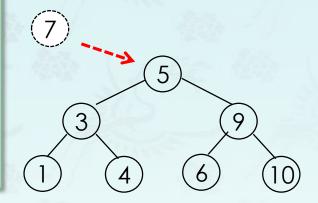
BST Node structure:





Operations: grow

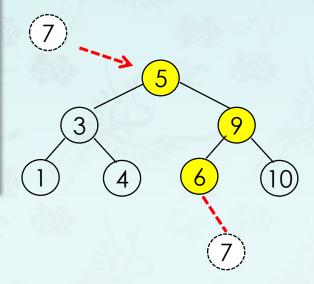
- grow(T, k)
 - Insert a node with Key = k into BST T
 - Time complexity? O(h)
- Step 1:
 if the tree is empty, then Root(T) = k
- Step 2: Pretending we are searching for k in BST, until we meet a nullptr node
- Step 3: Insert k



Q: Where is it inserted at?

Operations: grow

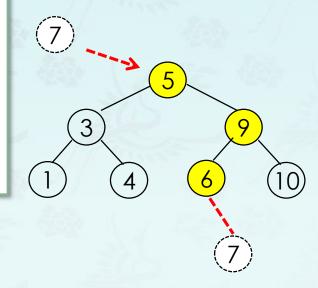
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The light nodes are compared with key.

Operations: grow

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 - Insert a node with Key = k into BST T
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The light nodes are compared with key.

Q: Do you see the difference between the complete binary tree and binary search tree?

```
tree grow(tree node, int key) {
  if (node == nullptr)
  if (key < node->key) // recur down the tree
    grow(node->left, key);
  else
    grow(node->right, key);
  else
    cout << "grow: the same key " << key << " is ignored.\n";</pre>
  return node;
```

```
tree grow(tree node, int key) {
  if (node == nullptr) return new Tree(key);
  if (key < node->key) // recur down the tree
    grow(node->left, key);
  else
    grow(node->right, key);
  else
    cout << "grow: the same key " << key << " is ignored.\n";</pre>
  return node;
```

```
tree grow (tree node, int key) {
  if (node == nullptr) return new Tree(key);
  if (key < node->key) // recur down the tree
    grow(node->left, key);
  else if (key > node->key)
    grow(node->right, key);
                                               Something wrong?
  else
    cout << "grow: the same key " << key << " is ignored.\n";</pre>
  return node;
```

```
tree grow(tree node, int key) {
  if (node == nullptr) return new Tree(key);
  if (key < node->key) // recur down the tree
    grow(node->left, key);
  else if (key > node->key)
    grow(node->right, key);
  else
  return node;
```

```
tree grow (tree node, int key) {
  if (node == nullptr) return new Tree(key);
  if (key < node->key) // recur down the tree
    node->left = grow(node->left, key);
  else if (key > node->key)
    node->right = grow(node->right, key);
  else
  return node;
```

inorder traversal: do inorder traversal of BST.

```
void inorder(tree node) {
   if (node == nullptr) return;

   inorder(node->left);
   cout << node->key;
   inorder(node->right);
}
```

inorder traversal: do inorder traversal of BST.

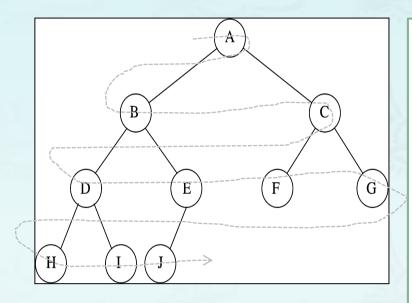
```
void inorder(tree node) {
    if (node == nullptr) return;
    inorder(node->left);
    cout << node->key;
    inorder(node->right);
void inorder(tree node, vector<int>& vec) {
  if (node == nullptr) return;
  inorder(node->left, vec);
  inorder(node->right, vec);
```

inorder traversal: do inorder traversal of BST.

```
void inorder(tree node) {
    if (node == nullptr) return;
    inorder (node->left);
    cout << node->key;
    inorder(node->right);
void inorder(tree node, vector<int>& vec) {
  if (node == nullptr) return;
                                         case '1':
  inorder(node->left, vec);
                                           cout << "\n\tinorder:</pre>
                                           vec.clear();
                                           inorder(root, vec);
  inorder(node->right, vec);
                                           for (int i : vec)
                                             cout << i << " ";
```

Level-order traversal

- 1. **Depth first** search(DFS) preorder, inorder, postorder traversal
- 2. Breadth first search (BFS) level-order traversal



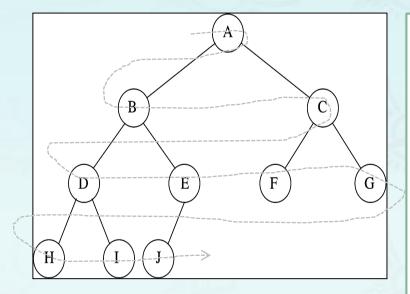
```
#include <queue>
#include <vector>
```

void levelorder(tree root, vector<int>& vec)

- Visit the root. if it is not null, enqueue it.
- while queue is not empty
 - 1. que.front() get the node in the queue
 - 2. save the key in vec.
 - 3. if its left child is not null, enqueue it.
 - 4. if its right child is not null, enqueue it.
 - 5. que.pop() remove the node in the queue.

Level-order traversal

- 1. **Depth first** search(DFS) preorder, inorder, postorder traversal
- 2. Breadth first search(BFS) level-order traversal



```
#include <queue>
#include <vector>
void levelorder(tree root, vector<int>& vec) {
  queue<tree> que;
  if (!root) return;
  que.push(root);
  while ...{
     cout << "your code here\n";</pre>
```

minimum, maximum:

returns the node with min or max key.

Note that the entire tree does not need to be searched.

```
tree minimum(tree node) { // returns left-most node key
}
tree maximum(tree node) { // returns right-most node key
}
```

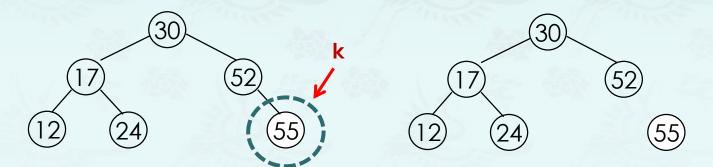
pred(), succ() - predecessor, successor:

```
Input: root node, key
output: predecessor node, successor node
1. If root is nullptr, then return
2. if key is found then
    a. If its left subtree is not nullptr
        Then predecessor will be the right most
        child of left subtree or left child itself.
    b. If its right subtree is not nullptr
        The successor will be the left most child
        of right subtree or right child itself.
    return
```

Operations: trim

- trim(**T**, k)
 - trim a node with Key = k into BST T
 - Time complexity: O(h)

Case 1: k has no child

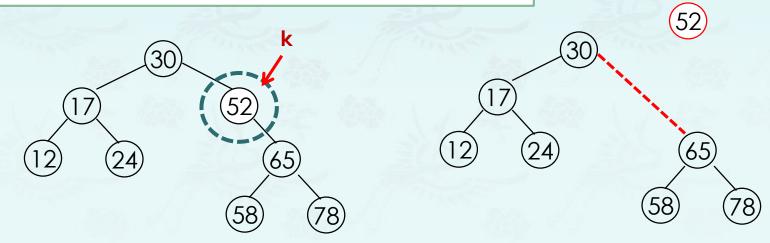


We can simply trim it from the tree

Operations: trim

- trim(**T**, k)
 - trim a node with Key = k into BST T
 - Time complexity: O(h)

Case 2: k has one child

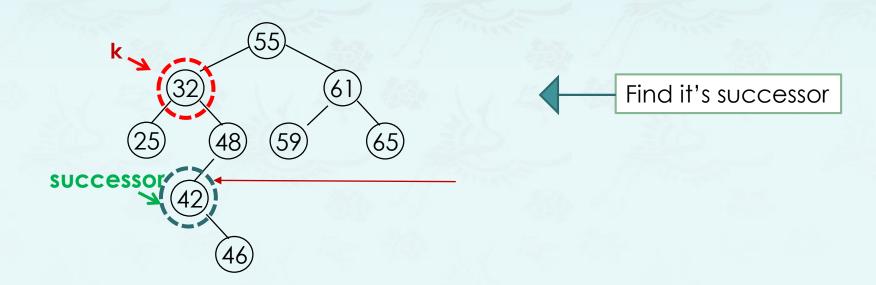


After removing it, connect it's subtree to it's parent node.

Operations: trim

- trim(**T**, k)
 - trim a node with Key = k into BST T
 - Time complexity: O(h)

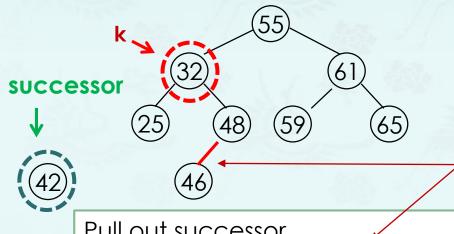
Case 3: k has two children



Operations: trim

- trim(**T**, k)
 - trim a node with Key = k into BST T
 - Time complexity: O(h)

Case 3: k has two children



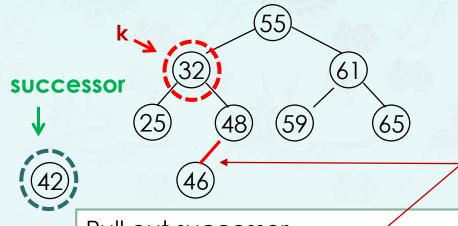
This is done by calling another trim() with succ key, recursively.

Pull out successor, and connect the tree with it's child

Operations: trim

- trim(T, k)
 - trim a node with Key = k into BST T
 - Time complexity: O(h)

Case 3: k has two children



```
int succ = value(minimum(node->right));
node->key = succ;
node->right = trim(node->right, succ);
```

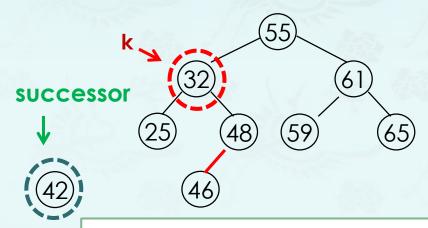
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Operations: trim

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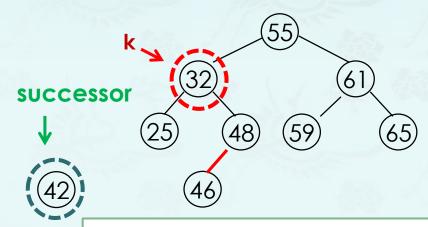
Pull out successor, and connect the tree with it's child

Q: What if successor has two children?

Operations: trim

- trim(**T**, k)
 - trim a node with Key = k into BST T
 - Time complexity: O(h)

Case 3: k has two children



A: Not possible!

Because if it has two nodes, at least one of them is less than it, then in the process of finding successor, we won't pick it!

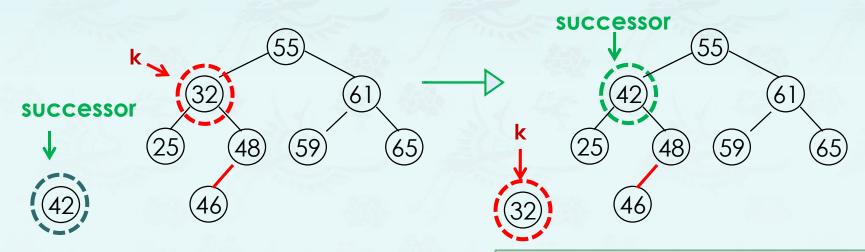
Pull out successor, and connect the tree with it's child

Q: What if successor has two children?

Operations: trim

- trim(**T**, k)
 - trim a node with Key = k into BST T
 - Time complexity: O(h)

Case 3: k has two children



Replace the **key** with it's successor

trim**: trim node with the key and return the new root.

```
tree trim(tree node, int key) {
 if (node == nullptr) return node;// base case
 if (key < node->key)
   node->left = trim(node->left, key);
 else if (key > node->key) {
   node->right = trim(node->right, key);
 else {
   if (node->left == nullptr) {
      // your code here - trim the right one, return node
   else if (node->right == nullptr) {
       // your code here - trim the left one, return node
   else {// two children case
     // get the successor: smallest in right subtree
     // copy the successor's content to this "node" node
     // node->right = trim the successor recursively, which has one or no child case
 return node;
```

http://visualgo.net/bst