ALGORITHM ASSIGNMENT 2

IMPLEMENTATION OF POLYGON TRIANGULATION & EXPLORING DATABASES

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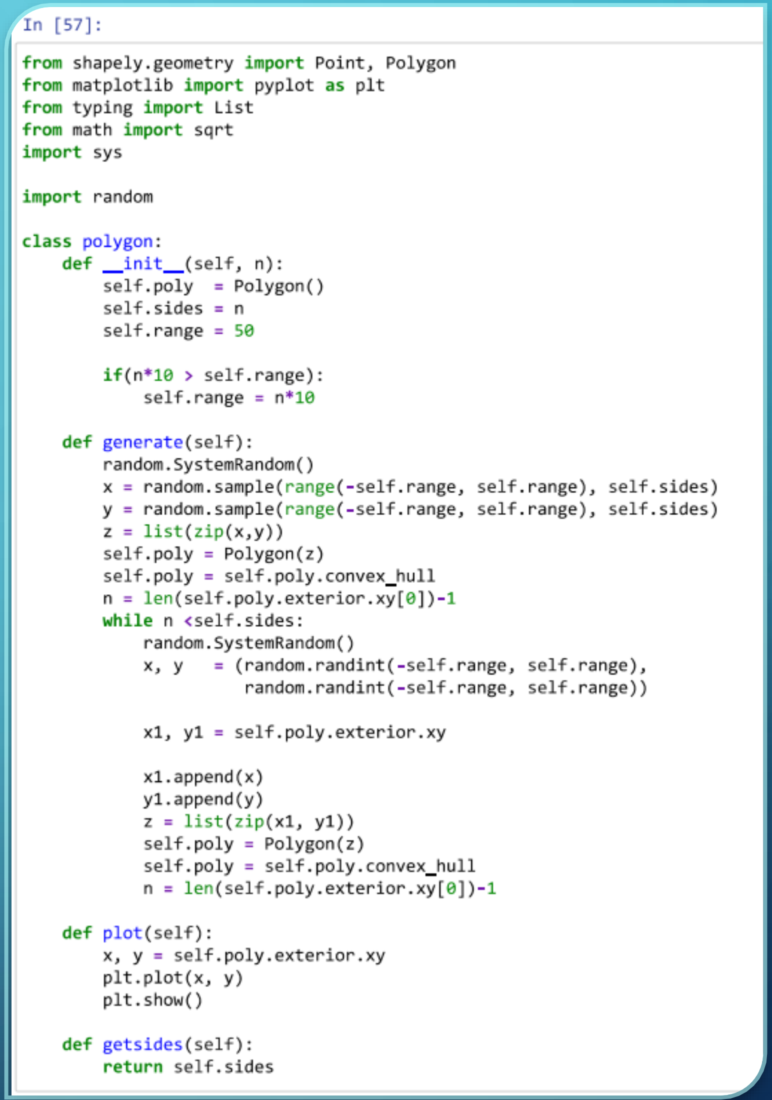
1.

POLYGON TRIANGULATION

Implementing the Polygon Triangulation Problem using different approaches of algorithmic programming

CLASS DESIGN FOR POLYGON TRIANGULATION

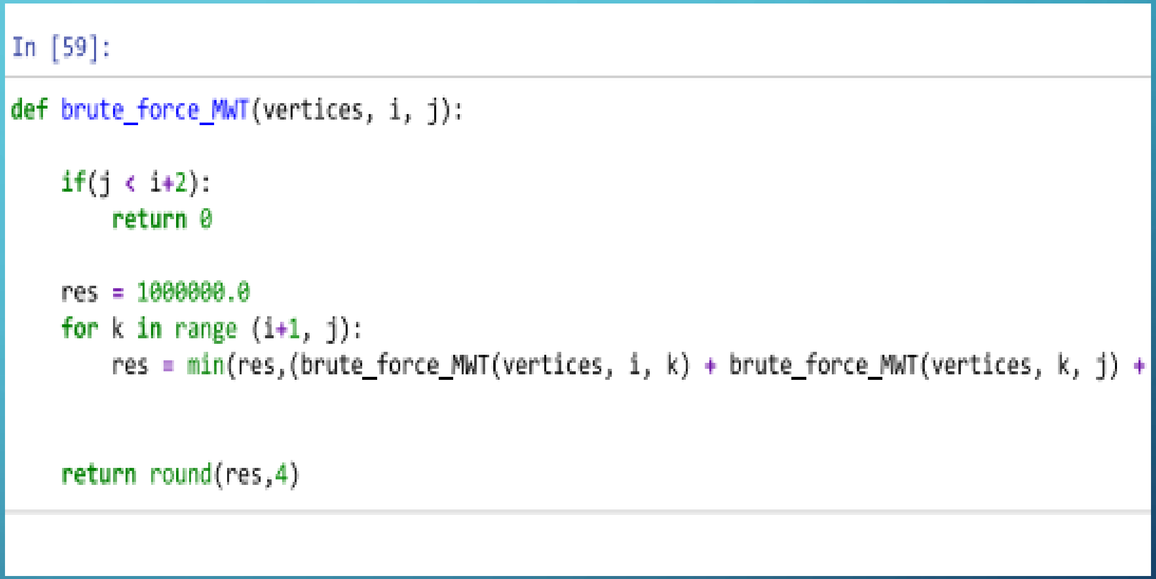
- __init__: Constructor for the class polygon
- generate: generates the random vertices for the polygon
- plot : plots the data in matplotlib
- getsides: returns the no. of sides in the generated convex polygon



```
In [57]:
from shapely.geometry import Point, Polygon
from matplotlib import pyplot as plt
from typing import List
from math import sqrt
import sys

import random

class polygon:
    def __init__(self, n):
        self.poly  = Polygon()
        self.sides = n
        self.range = 50

        if(n*10 > self.range):
            self.range = n*10

    def generate(self):
        random.SystemRandom()
        x = random.sample(range(-self.range, self.range), self.sides)
        y = random.sample(range(-self.range, self.range), self.sides)
        z = list(zip(x,y))
        self.poly = Polygon(z)
        self.poly = self.poly.convex_hull
        n = len(self.poly.exterior.xy[0])-1
        while n <self.sides:
            random.SystemRandom()
            x, y   = (random.randint(-self.range, self.range),
                      random.randint(-self.range, self.range))

            x1, y1 = self.poly.exterior.xy

            x1.append(x)
            y1.append(y)
            z = list(zip(x1, y1))
            self.poly = Polygon(z)
            self.poly = self.poly.convex_hull
            n = len(self.poly.exterior.xy[0])-1

    def plot(self):
        x, y = self.poly.exterior.xy
        plt.plot(x, y)
        plt.show()

    def getsides(self):
        return self.sides
```

BRUTE FORCE APPROACH

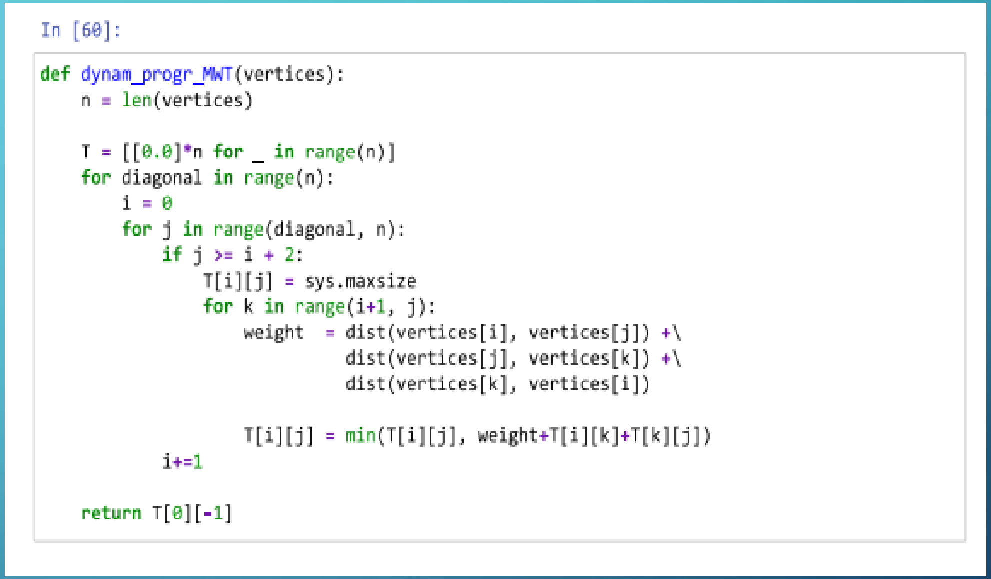
- brute_force_MWT: uses the brute-force approach to calculate the minimum cost of triangulation for the given polygon.



```
In [59]:

def brute_force_MWT(vertices, i, j):

    if(j < i+2):
        return 0

    res = 1000000.0
    for k in range (i+1, j):
        res = min(res,(brute_force_MWT(vertices, i, k) + brute_force_MWT(vertices, k, j) +


    return round(res,4)
```

DYNAMIC PROGRAMMING APPROACH

- dynam_progr_MWT: uses the dynamic programming approach to find the minimum cost of triangulation for the given polygon.

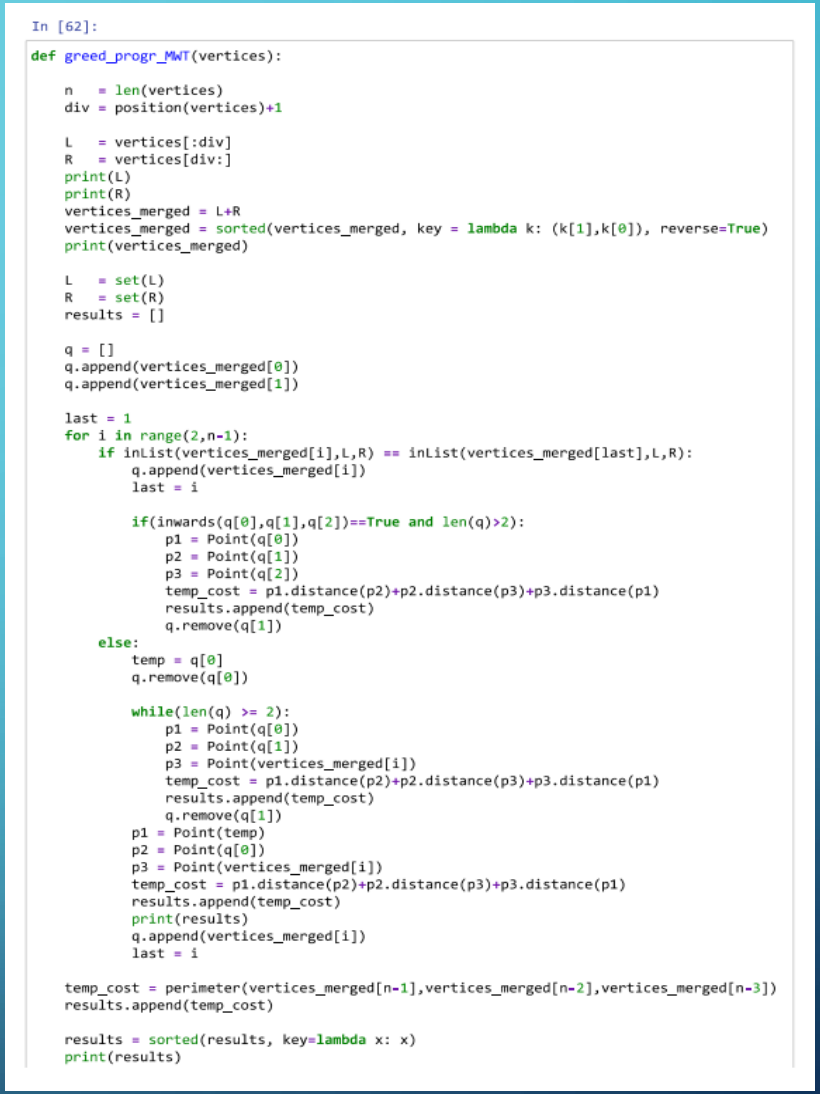


In [60]:

```
def dynam_progr_MWT(vertices):
    n = len(vertices)

    T = [[0.0]*n for _ in range(n)]
    for diagonal in range(n):
        i = 0
        for j in range(diagonal, n):
            if j >= i + 2:
                T[i][j] = sys.maxsize
                for k in range(i+1, j):
                    weight  = dist(vertices[i], vertices[j]) +\
                              dist(vertices[j], vertices[k]) +\
                              dist(vertices[k], vertices[i])

                    T[i][j] = min(T[i][j], weight+T[i][k]+T[k][j])
            i+=1

    return T[0][-1]
```

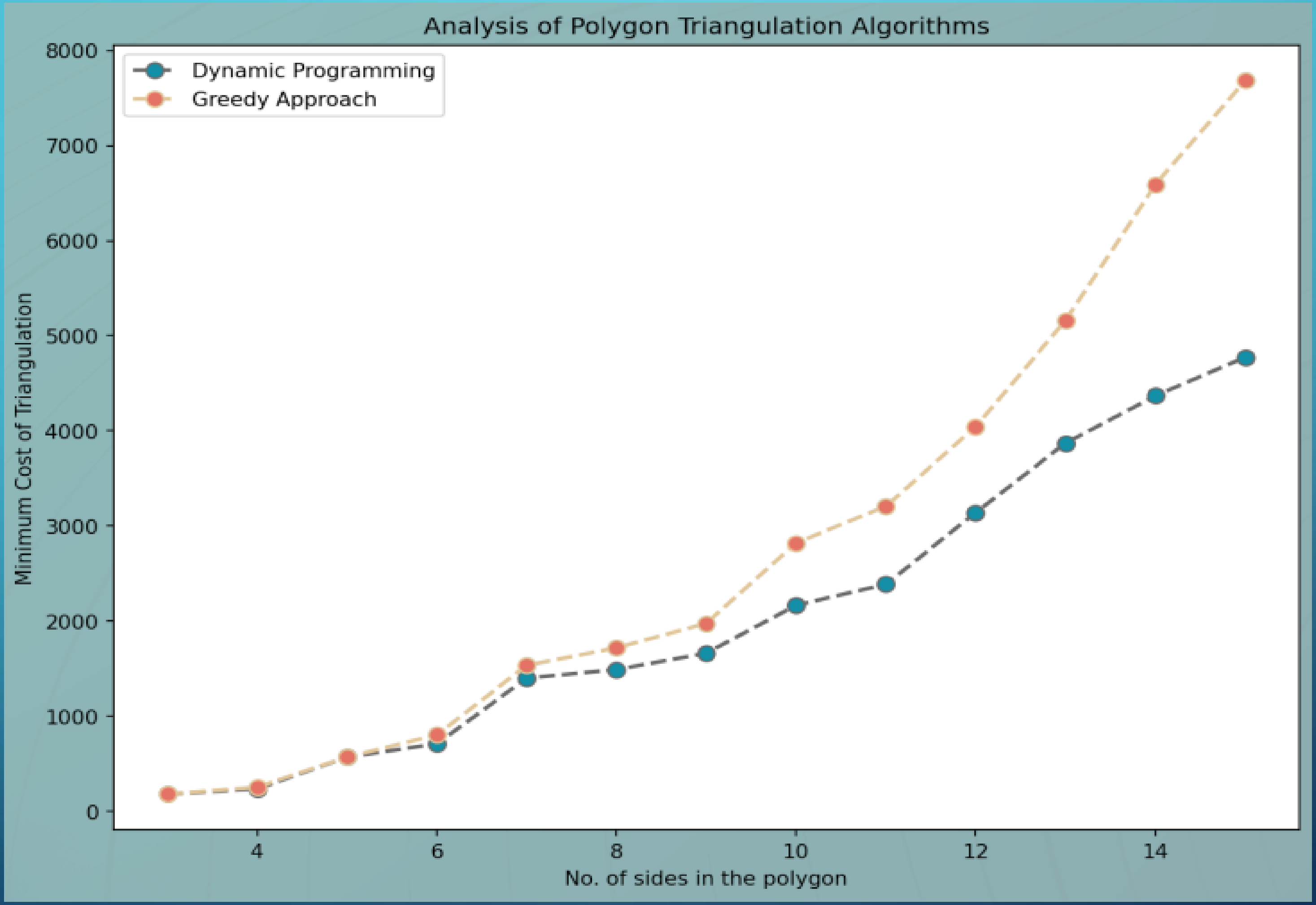
GREEDY PROGRAMMING APPROACH

- greed_progr_MWT: uses the greedy programming to quickly triangulate a polygon. Later on we will check whether whether the triangulation is actually the minimum triangulation.



```
In [62]:
def greed_progr_MWT(vertices):

    n   = len(vertices)
    div = position(vertices)+1

    L   = vertices[:div]
    R   = vertices[div:]
    print(L)
    print(R)
    vertices_merged = L+R
    vertices_merged = sorted(vertices_merged, key = lambda k: (k[1],k[0]), reverse=True)
    print(vertices_merged)

    L   = set(L)
    R   = set(R)
    results = []

    q = []
    q.append(vertices_merged[0])
    q.append(vertices_merged[1])

    last = 1
    for i in range(2,n-1):
        if inList(vertices_merged[i],L,R) == inList(vertices_merged[last],L,R):
            q.append(vertices_merged[i])
            last = i

            if(inwards(q[0],q[1],q[2])==True and len(q)>2):
                p1 = Point(q[0])
                p2 = Point(q[1])
                p3 = Point(q[2])
                temp_cost = p1.distance(p2)+p2.distance(p3)+p3.distance(p1)
                results.append(temp_cost)
                q.remove(q[1])
        else:
            temp = q[0]
            q.remove(q[0])

            while(len(q) >= 2):
                p1 = Point(q[0])
                p2 = Point(q[1])
                p3 = Point(vertices_merged[i])
                temp_cost = p1.distance(p2)+p2.distance(p3)+p3.distance(p1)
                results.append(temp_cost)
                q.remove(q[1])
            p1 = Point(temp)
            p2 = Point(q[0])
            p3 = Point(vertices_merged[i])
            temp_cost = p1.distance(p2)+p2.distance(p3)+p3.distance(p1)
            results.append(temp_cost)
            print(results)
            q.append(vertices_merged[i])
            last = i

    temp_cost = perimeter(vertices_merged[n-1],vertices_merged[n-2],vertices_merged[n-3])
    results.append(temp_cost)

    results = sorted(results, key=lambda x: x)
    print(results)
```

PLOT COMPARISON



ANALYSIS

- The dynamic programming approach takes O(n^3) while in case of Greedy Approach it takes the complexity of O(nlogn) which in this case was O(n^2logn).
- However we can clearly see that the Greedy approach does not always give the minimum triangulation of the polygon, especially for polygons with large number of sides. But it is considerably faster so it can be used as an alternative for the dynamic programming approach if accuracy is not a very important factor.

2.

EXPLORING DATABASES

Exploring the SNAP and KONECT databases and try running different algorithms like MST, Disjoint Sets etc.

STANFORD NETWORK ANALYSIS PROJECT

- STANFORD NETWORK ANALYSIS PLATFORM (SNAP) IS A GENERAL PURPOSE NETWORK ANALYSIS AND GRAPH MINING LIBRARY. IT IS WRITTEN IN C++ AND EASILY SCALES TO MASSIVE NETWORKS WITH HUNDREDS OF MILLIONS OF NODES, AND BILLIONS OF EDGES. IT EFFICIENTLY MANIPULATES LARGE GRAPHS, CALCULATES STRUCTURAL PROPERTIES, GENERATES REGULAR AND RANDOM GRAPHS, AND SUPPORTS ATTRIBUTES ON NODES AND EDGES.
- IT IS ALSO SUPPORTED ON PYTHON UNDER SNAP.PY

KOBLENZ NETWORK COLLECTION

- KOBLENZ NETWORK COLLECTION, AS A PROJECT HAS ROOTS AT THE UNIVERSITY OF KOBLENZ-LANDAU IN GERMANY. ALL SOURCE CODE IS MADE AVAILABLE AS FREE SOFTWARE, AND INCLUDES A NETWORK ANALYSIS TOOLBOX FOR GNU OCTAVE, A NETWORK EXTRACTION LIBRARY, AS WELL AS CODE TO GENERATE WEBPAGES, INCLUDING STATISTICS AND PLOTS. KONECT IS RUN BY RESEARCH GROUP AROUND JÉRÔME KUNEGIS AT THE UNIVERSITY OF NAMUR, IN THE NAMUR CENTER FOR COMPLEX NETWORKS.
- THE KONECT PROJECT HAS 1,326 NETWORK DATASETS IN 24 CATEGORIES.

DISJOINT SET

- GIVEN N VERTICES AND M PREDEFINED EDGES IN A GRAPH , ONE HAS TO FIND A SUBSET OF M EDGES SO THAT THE N VERTICES ARE DIVIDED INTO DISJOINT SETS.
- ALL VERTICES WILL BE FIRST INITIALIZED AS DISJOINT THE EDGES WILL BE USED TO CONNECT THE SETS AS LONG AS IT DOES NOT MAKE A CYCLE. AFTER ALL THE EDGES ARE CONNECTED KEEPING THE PREVIOUS CONDITION IN MIND, WE GET THE FINAL ANSWER.
- THIS ALGORITHM CAN BE IMPLEMENTED USING TWO TYPES OF DATA STRUCTURES
  - LINKED LIST
  - TREE BASED

LINKED LIST IMPLEMENTATION

- TWO LINKED LISTS ARE NEEDED FOR THIS PURPOSE, ONE FOR DYNAMICALLY MANAGING NUMBER OF CONNECTED COMPONENTS OF THE GRAPH (REPRESENTATIVE ELEMENT), AND THE OTHER TO MAINTAIN THE LIST OF VERTICES INSIDE THAT CONNECTED COMPONENT.
- THE CONNECTED COMPONENT LIST IS DOUBLY LINKED WHILE THE VERTICES LIST IS SINGLY LINKED.

TREE IMPLEMENTATION

- WE ONLY USE VERTEX AS NODES, WHICH CONTAIN THE RANK OF THE NODE AND A REFERENCE TO THE PARENT NODE(MOSTLY THE MAIN ROOT NODE) .
- THE RANK OF THE NODE BASICALLY SIGNIFIES HOW BIG THE SUBTREE WITH THE NODE AS ROOT IS.

OBSERVATION

- SNAP FACEBOOK (VERTICES : 4039, EDGES: 88234)
  - TREE APPROACH -> 197.403 MS
  - LINKED LIST APPROACH -> 213.607 MS
- SNAP EPINION (VERTICES: 75888, EDGES: 508837)
  - TREE APPROACH -> 1123.56 MS (1.12 S)
  - LINKED LIST APPROACH -> 600632 MS (10 MIN)
- SNAP JOURNAL (VERTICES: 6262114, EDGES: 15119313)
  - TREE APPROACH -> 204793 MS (3 MIN)
  - LINKED LIST APPROACH -> FINISHED 3% IN ONE HOUR (33 HOURS)

THOUGHTS

- THESE DATABASES PROVIDE NOT ONLY THE FACILITY FOR GENERATING RANDOM DIRECTED AND NON-DIRECTED GRAPHS, THEY ALSO PROVIDE THE MEANS FOR ORDER STATISTICS ON VARIOUS OPERATIONS THAT CAN BE PERFORMED ON SAID GRAPHS.
- ACCORDING TO THE RECORDED OBSERVATIONS, TREE BASED APPROACH IS MUCH FASTER THAN LINKED LIST BASED APPROACH. ALSO THE TIME TAKEN MAINLY DEPENDS ON THE NUMBER OF EDGES, WHILE THE SPACE TAKEN DEPENDS ON THE NUMBER OF VERTICES.
- HOWEVER THE FULL ANALYSIS OF THE DATASETS CANNOT BE POSSIBLY DONE IN A LAPTOP ENVIRONMENT.

Thank You