Roll No.: 027

Experiment no - 02

Aim: Write a program to minimize given DFA.

Theory Explanation:

- DFA refers to deterministic finite automata. Deterministic refers to the uniqueness of the computation. The finite automata are called deterministic finite automata if the machine is read an input string one symbol at a time.
- o In DFA, there is only one path for specific input from the current state to the next state.
- o DFA does not accept the null move, i.e., the DFA cannot change state without any input character.
- o DFA can contain multiple final states. It is used in Lexical Analysis in Compiler.

In the following diagram, we can see that from state q0 for input a, there is only one path which is going to q1. Similarly, from q0, there is only one path for input b going to q2.

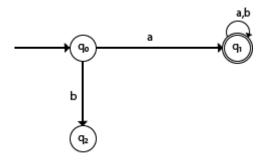


Fig:- DFA

Formal Definition of DFA

A DFA is a collection of 5-tuples same as we described in the definition of FA.

- 1. Q: finite set of states
- 2. Σ : finite set of the input symbol
- 3. q0: initial state
- 4. F: **final** state
- 5. δ : Transition function

Transition function can be defined as:

1.
$$\delta: Q \times \Sigma \rightarrow Q$$

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Graphical Representation of DFA

A DFA can be represented by digraphs called state diagram. In which:

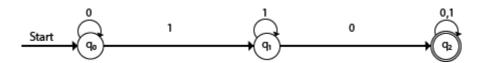
- 1. The state is represented by vertices.
- 2. The arc labeled with an input character show the transitions.
- 3. The initial state is marked with an arrow.
- 4. The final state is denoted by a double circle.

Example 1:

- 1. $Q = \{q0, q1, q2\}$
- 2. $\sum = \{0, 1\}$
- 3. $q0 = \{q0\}$
- 4. $F = \{q2\}$

Solution:

Transition Diagram:



Transition Table:

PRESENT STATE	NEXT STATE FOR INPUT 0	NEXT STATE OF INPUT 1
→Q0	q0	q1
Q1	q2	q1
*Q2	q2	q2

Minimization of DFA

Minimization of DFA means reducing the number of states from given FA. Thus, we get the FSM(finite state machine) with redundant states after minimizing the FSM.

We have to follow the various steps to minimize the DFA. These are as follows:

Step 1: Remove all the states that are unreachable from the initial state via any set of the transition of DFA.

Step 2: Draw the transition table for all pair of states.

Step 3: Now split the transition table into two tables T1 and T2. T1 contains all final states, and T2 contains non-final states.

Step 4: Find similar rows from T1 such that:

- 1. 1. $\delta(q, a) = p$
- 2. $\delta(r, a) = p$

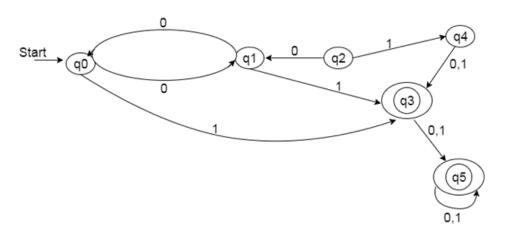
That means, find the two states which have the same value of a and b and remove one of them.

Step 5: Repeat step 3 until we find no similar rows available in the transition table T1.

Step 6: Repeat step 3 and step 4 for table T2 also.

Step 7: Now combine the reduced T1 and T2 tables. The combined transition table is the transition table of minimized DFA.

Example:



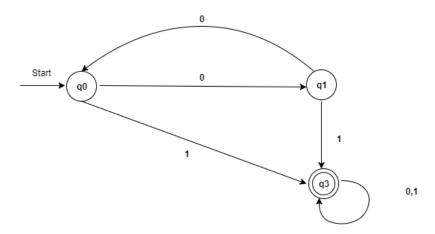
Solution:

Step 1: In the given DFA, q2 and q4 are the unreachable states so remove them.

Step 2: Draw the transition table for the rest of the states.

STATE	0	1
→Q0	q1	q3
Q1	q0	q3

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*Q3		q5	q5
*Q5		q5	q5
Step 3: Now divide	rows of transition table into	two sets as:	
1. One set contains	those rows, which start from	non-final state	es:
STATE	ı	0	1
Q0 Q1		q1 q0	q3 q3
Ų1	1,	qυ	4 5
2 Another set t	sing the gar warre wile talk at the	from Carl 1	to a
	ins those rows, which starts		
STATE		0	1
Q3		q5	q5
Q5		q5	q5
		1	
Step 4: Set 1 has no	o similar rows so set 1 will be	e the same.	
Step 5: In set 2, row		ce q3 and q5 to	ransit to the same state on 0 and
Step 5: In set 2, row	v 1 and row 2 are similar sind	ce q3 and q5 to	ransit to the same state on 0 and
Step 5: In set 2, row 1. So skip q5 and th	v 1 and row 2 are similar sind	ce q3 and q5 to st.	
Step 5: In set 2, row 1. So skip q5 and th STATE Q3	v 1 and row 2 are similar sind en replace q5 by q3 in the re	ce q3 and q5 to st.	1
Step 5: In set 2, row 1. So skip q5 and the STATE Q3 Step 6: Now combine the step in the	v 1 and row 2 are similar sind	ce q3 and q5 to st. 0 q3	1 q3
Step 5: In set 2, row 1. So skip q5 and th STATE Q3 Step 6: Now combi	v 1 and row 2 are similar sind en replace q5 by q3 in the re	ce q3 and q5 to st. 0 q3	1 q3
Step 5: In set 2, row 1. So skip q5 and th STATE Q3 Step 6: Now combi STATE →Q0	v 1 and row 2 are similar sind en replace q5 by q3 in the re	ce q3 and q5 tr st. 0 q3 0 q1	1 q3 1 q3
Step 5: In set 2, row 1. So skip q5 and th STATE Q3 Step 6: Now combi STATE →Q0 Q1	v 1 and row 2 are similar sind en replace q5 by q3 in the re	ce q3 and q5 tr st. 0 q3 0 q1 q0	1 q3 1 q3 q3 q3
Step 5: In set 2, row 1. So skip q5 and th STATE Q3 Step 6: Now combi STATE →Q0	v 1 and row 2 are similar sind en replace q5 by q3 in the re	ce q3 and q5 tr st. 0 q3 0 q1	1 q3 1 q3



Program:

1. DFA.Py

```
from collections import defaultdict
from disjoint_set import DisjointSet
class DFA(object):
  def __init__(self,states_or_filename,terminals=None,start_state=None,
transitions=None,final_states=None):
    if terminals is None:
      self._get_graph_from_file(states_or_filename)
      assert isinstance(states_or_filename,list) or isinstance(states_or_filename,tuple)
      self.states = states_or_filename
     assert isinstance(terminals,list) or isinstance(terminals,tuple)
      self.terminals = terminals
      assert isinstance(start_state,str)
      self.start_state = start_state
      assert is instance(transitions, dict)
      self.transitions = transitions
      assert isinstance(final_states,list) or isinstance(final_states,tuple)
      self.final_states = final_states
  def _remove_unreachable_states(self):
    Removes states that are unreachable from the start state
    g = defaultdict(list)
```

```
for k,v in self.transitions.items():
   g[k[0]].append(v)
 # do DFS
 stack = [self.start state]
 reachable_states = set()
 while stack:
   state = stack.pop()
   if state not in reachable_states:
     stack += g[state]
   reachable_states.add(state)
 self.states = [state for state in self.states \
              if state in reachable_states]
 self.final_states = [state for state in self.final_states \]
              if state in reachable_states]
 self.transitions = { k:v for k,v in self.transitions.items() \
             if k[0] in reachable_states}
def minimize(self):
 self._remove_unreachable_states()
 def order_tuple(a,b):
   return (a,b) if a < b else (b,a)
 table = \{ \}
 sorted_states = sorted(self.states)
 # initialize the table
 for i,item in enumerate(sorted_states):
   for item_2 in sorted_states[i+1:]:
     table[(item,item_2)] = (item in self.final_states) != (item_2\
                          in self.final_states)
 flag = True
 # table filling method
 while flag:
   flag = False
```

```
for i,item in enumerate(sorted_states):
     for item_2 in sorted_states[i+1:]:
       if table[(item,item_2)]:
         continue
       # check if the states are distinguishable
       for w in self.terminals:
         t1 = self.transitions.get((item,w),None)
         t2 = self.transitions.get((item_2, w), None)
         if t1 is not None and t2 is not None and t1 != t2:
           marked = table[order tuple(t1,t2)]
           flag = flag or marked
           table[(item,item_2)] = marked
           if marked:
             break
 d = DisjointSet(self.states)
 # form new states
 for k,v in table.items():
   if not v:
     d.union(k[0],k[1])
 self.states = [str(x) for x in range(1,1+len(d.get()))]
 new_final_states = []
 self.start_state = str(d.find_set(self.start_state))
 for s in d.get():
   for item in s:
     if item in self.final_states:
       new_final_states.append(str(d.find_set(item)))
       break
 self.transitions = \{(str(d.find\_set(k[0])), k[1]): str(d.find\_set(v))\}
             for k,v in self.transitions.items()}
 self.final_states = new_final_states
def __str__(self):
 String representation
 num_of_state = len(self.states)
 start state = self.start state
 num_of_final = len(self.final_states)
```

```
return '{} states. {} final states. start state - {}'.format( \
                   num_of_state,num_of_final,start_state)
 def _get_graph_from_file(self,filename):
    Load the graph from file
    with open(filename, 'r') as f:
     try:
       lines = f.readlines()
       states, terminals, start state, final states = lines[:4]
       if states:
         self.states = states[:-1].split()
       else:
         raise Exception('Invalid file format: cannot read states')
       if terminals:
          self.terminals = terminals[:-1].split()
       else:
         raise Exception('Invalid file format: cannot read terminals')
       if start_state:
          self.start_state = start_state[:-1]
       else:
         raise Exception('Invalid file format: cannot read start state')
       if final_states:
          self.final_states = final_states[:-1].split()
       else:
         raise Exception('Invalid file format: cannot read final states')
       lines = lines[4:]
       self.transitions = {}
       for line in lines:
         current_state,terminal,next_state = line[:-1].split()
         self.transitions[(current_state,terminal)] = next_state
     except Exception as e:
       print("ERROR: ",e)
if name =='' main '':
     filename = 'graph'
     dfa = DFA(filename)
```

```
print(dfa)
dfa.minimize()
print(dfa)
```

2. Disjoint.py

```
class DisjointSet(object):
 def __init__(self,items):
   self._disjoint_set = list()
   if items:
     for item in set(items):
       self._disjoint_set.append([item])
 def _get_index(self,item):
   for s in self._disjoint_set:
     for _item in s:
       if item == item:
         return self._disjoint_set.index(s)
   return None
 def find(self,item):
   for s in self._disjoint_set:
     if item in s:
       return s
   return None
 def find_set(self,item):
   s = self.\_get\_index(item)
   return s+1 if s is not None else None
 def union(self,item1,item2):
   i = self.\_get\_index(item1)
   j = self.\_get\_index(item2)
   if i != j:
     self._disjoint_set[i] += self._disjoint_set[j]
     del self._disjoint_set[j]
 def get(self):
   return self._disjoint_set
```

3. Graph Input

```
1 2 3 4 5

a b

1

1 5

1 a 3

1 b 2

2 b 1

2 a 4

3 b 4

3 a 5

4 a 4

4 b 4

5 a 3

5 b 2
```

OUTPUT:

```
('1',
     '3') : ['a']
('1',
     '2') : ['b']
('2',
     '1') : ['b']
('2',
     '4') : ['a']
('3',
      '4') : ['b']
('3',
      '5') :
             ['a']
('4', '4') : ['a', 'b']
('5',
     '3') : ['a']
('5', '2') : ['b']
5 states. 2 final states. start state - 1
('2', '4') : ['a']
('2', '1') : ['b']
('1', '2') : ['b']
('1', '3') : ['a']
('4', '3') : ['b']
('4',
     '2') : ['a']
('3', '3') : ['a', 'b']
4 states. 1 final states. start state - 2
```

Conclusion : A minimized DFA was successfully observed from 5 total states to 4 states.

References:-

https://github.com/navin-mohan/dfa-minimization

https://www.javatpoint.com/minimization-of-dfa

https://www.javatpoint.com/deterministic-finite-automata