Secure Your Programming Future!

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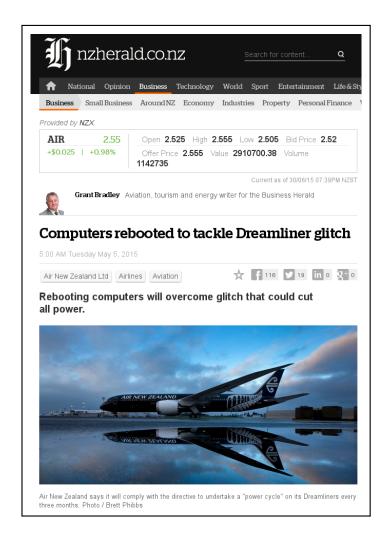
@WhileyDave

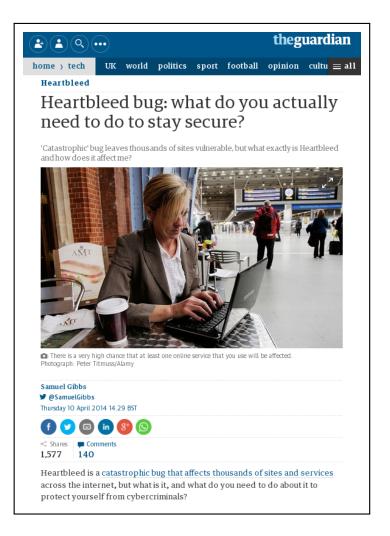
http://whiley.org

http://github.com/Whiley

Background

Verification: Who Cares?





Verification: A Challenge for Computer Science

"A verifying compiler uses automated mathematical and logical reasoning methods to check the correctness of the programs that it compiles"

-Hoare'03

Whiley

Overview: What is Whiley?

```
function max(int x, int y) \rightarrow (int z)

// result must be one of the arguments

ensures x == z \mid \mid y == z

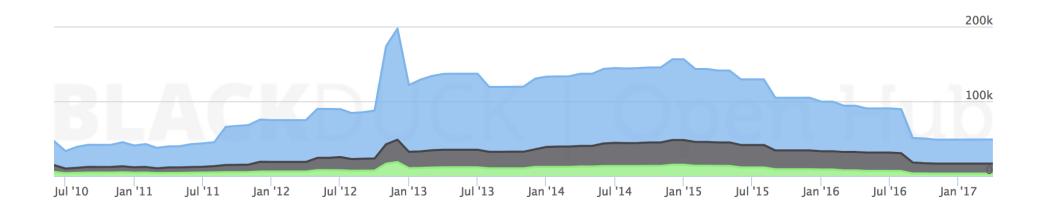
// result must be greater-or-equal than arguments

ensures x <= z \& y <= z:

...
```

- A language designed specifically to simplify verifying software
- Several trade offs e.g. performance for verifiability
 - Unbounded Arithmetic, value semantics, etc
- Goal: to statically verify functions meet their specifications

History of Whiley



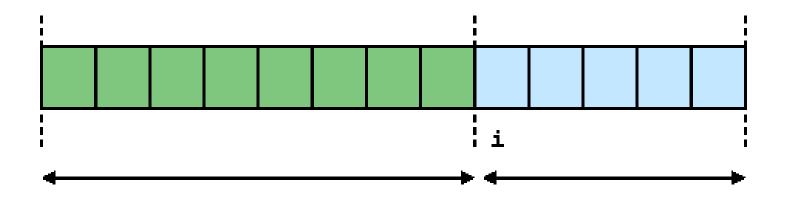
- 2009 Initial version of Whiley released (GPL Licence)
- 2010 GitHub repository and http://whiley.org go live
- 2010 **Version 0.3.0** released (BSD Licence)
- 2016 Version 0.4.0 released
- 2017 **Version 0.4.1** released

Demo Time...

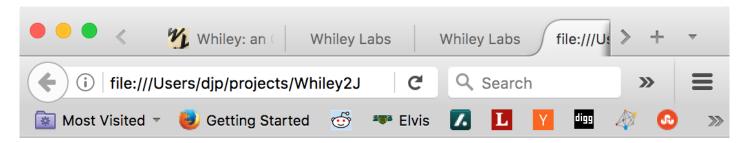
Example: max(int[])

```
// Returns index of largest item in array
function max(int[] items) → (int r)
```

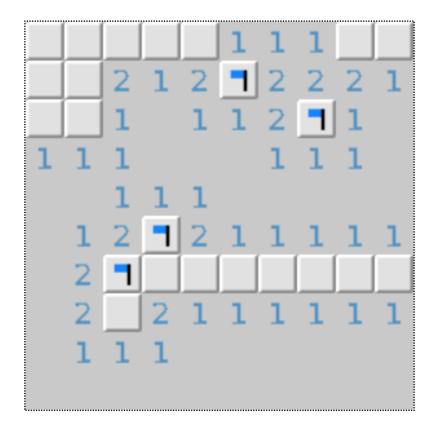
Diagram!



Diagram!



Minesweeper (in Whiley)



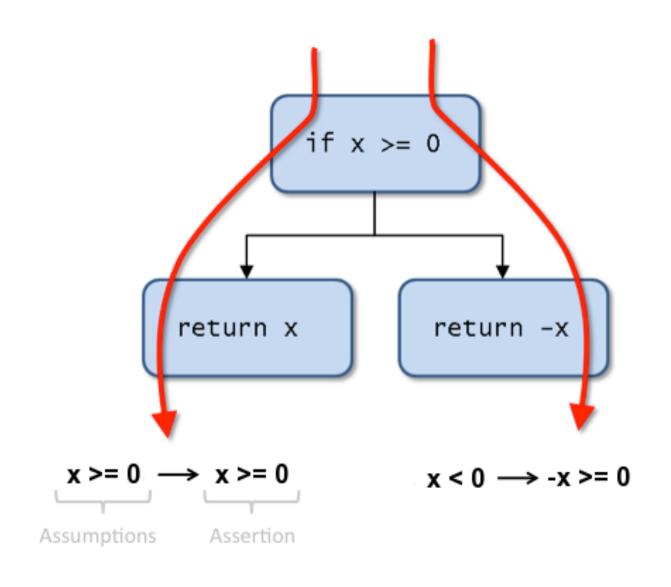
How does it work?

Verification: How does it work?

```
function abs(int x) => (int r)
// return value cannot be negative
ensures r >= 0:
    //
    if x >= 0:
        return x
    else:
        return -x
```

- To verify above function, compiler generates verification conditions
- Verification conditions are (roughly) first-order logic formulas

Verification: Verification Condition Generation



Verification: Assertion Language

Whiley compiler emits verification conditions in assertion language

```
assert:
  forall (int x):
    x >= 0 ==> x >= 0

assert:
  forall (int x):
    x < 0 ==> -x >= 0
```

- Verification conditions from | abs () | example shown above
- In principle, can hook up different automatic theorem provers

People (so far)



Art
(built C backend, 2012)



Melby
(built GPGPU backend,
2013)



Daniel
(helping with WhileyWeb)



Matt
(compiling for a QuadCopter,
2014)



Henry
(improving verification, 2014)



Sam
(started PhD on
Parallelisation, 2014)



Lindsay
(A/Prof, Victoria University)



Mark
(A/Prof, University of
Waikato)

http://whiley.org

@WhileyDave
http://github.com/Whiley

Verification: Constrained Types

```
type N is (T x) where e
```

- Above defines constrained type
- Invariant: for any variable of type N, follows that e always holds
- Constrained types can simplify specifications / invariants
- Example: natural numbers

```
type nat is (int n) where n >= 0
```

Verification: Structural Typing

```
type nat is (int n) where n >= 0

function cut(int x) → (nat y):
   if x >= 0:
      return x
   else:
      return 0
```

Variable types in Whiley are ephemeral ...

... and determined by what is **known** (not what was declared)

Verification: Flow Typing

```
function indexOf(int[] items, int item) → (int|null r)
// If integer value returned, must be index of item
ensures r is int ==> items[r] == item
// No element before integer r matches item
ensures r is int ==> all { k in 0..r | items[k] != item }
// If null returned, no matching item
ensures r is null ==> all { k in 0..|items| | items[k] != item }:
    int i = 0
    while i < |items|</pre>
    where i >= 0 && i <= |items|
    where all { j in 0..i | items[j] != item }:
         if items[i] == item:
             return i
         i = i + 1
    return null
```