

Mobile Platform Security:
OS Hardening and Trusted Execution Environment

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# Agenda

Attack Types on a Platform
Platform Security Components
Operating System Security
OS Security: Linux
Questions on OS security
Motivation for Isolation
HW Components for Secure Code Execution
Trusted Execution Environment
TEE Services and Use Cases (Authentication, RKP, Key Management, DRM and SIM Lock)

## Attack Types on a Platform

### Hack Attack

- Attacking via software using malwares, spywares and vulnerabilities
- Does not need hands on the device.
   Remotely done and effective on masses.

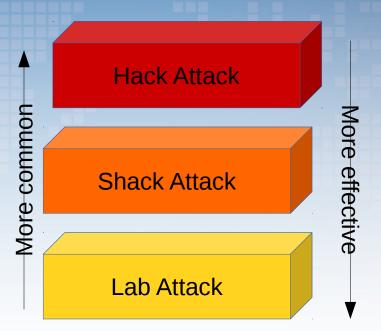
### Shack Attack

- Targets certain platforms using vulnerabilities with complex attack techniques. Sometimes requires proximity to the device. Attack cost does not include hardware tampering.
- Some side channel attacks are used as shack attacks.

#### Lab Attack

 Requires possession of the device and a facility to perform the attack. The costliest and most complex of the attacks. Bus probing etc.





# **Platform Security Components**



Virtualizable

Web and Mobile Applications

SW Secure Element

REE (OS) Kernel

Trusted App 1

Trusted App 2

Trusted Applet

TEE Kernel

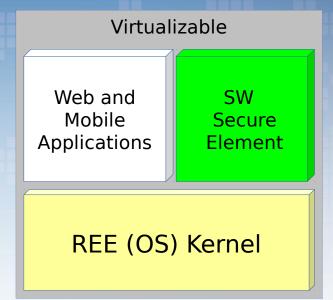
HW Secure Element TPM or Secure Co-Processor

Hardware



# **Operating System Security**

- Process Isolation (Virtual Address Spaces, Sand-boxing, Shared Memory, Limited IPC Traffic, Sockets)
- Discretionary Access Control (Object Access Control) – Identity based access restriction for objects. Access permission is transferable between objects.
- Mandatory Access Control (Policy Based Access Control)
- Memory Configuration (Non Executable Stack, Heap)





# OS Security: Linux 1/3

### Linux Security Extensions & Features

- Address Space Layout Randomization A Loader feature which loads the parts of the executable in non-contiguous order.
- Support for extended file attributes. (e.g. SELinux, SMACK extended file attributes.)
- Mandatory Access Control systems
- Integrity Measurement IMA / EVM



OS Security: Linux 2/3

Package and Kernel Module Verification

- Kernel Module White Lists Limiting allowed modules in kernel.
- Disabling Dynamic Kernel Module Loading Disabling a kernel module from loading if attacker slips its own binary.
- Signed Application Package Managers (RPM, DPKG ...) Integrity and authenticity of the package is protected. Even confidentiality if sensitive information is included in the package.



OS Security: Linux 3/3

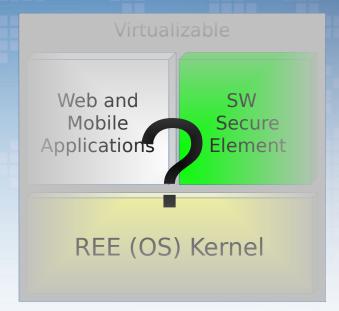
## **Compiler Options**

- Position Independent Execution Libraries, heap stack are randomly located using ASLR but randomization of executable layout is done by PIE parameter.
- Stack Smashing Protection Canary values to pin the return address of the function on the stack.
- Fortify Source Calls for boundary checking functions instead of standard c functions without boundary check such as "strcpy".

## Food for Thoughts

- Is REE always safe with constant installation of new applications?
- Are attack surfaces on an OS easy and cheap to guard?
- Are both runtime and boot time security provided?
- Are sensitive operations (crypto, decoding, encoding, sign, verify) safe?
- Do all the parties in a mobile platform (i.e. OEM, Operator, OS vendor, Application Providers) provide their services in isolation?
- Is user's sensitive data (Bank id, fingerprint, Login information) safe ?







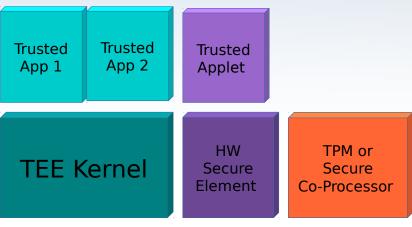
# Motivation for Isolation of Sensitive Code Execution

- REE if full of attack surfaces from kernel to Applications.
- Security frameworks have high complexity and dependencies. The weakest link can risk the entire system
- Maintenance of a complex security system is difficult and multiple bugs can complement to a successful attack while they are harmless when alone.
- Security policies are complex to design. Circular permission chains can cause policy breaches
- Sensitive operations need to be maintained independent from the rest of the applications.



# Hardware Components for Secure Code Execution

- Trusted Platform Module Simple crypto capabilities with limited key storage capability. Used in laptops and desktops. Tamper resistant.
- Secure Element Subscriber Identity Module. Useful for platform independent identification. Crypto operations with tamper resistance. Low power and performance.
- Secure Co-Processors- External Hardware Entity with relatively high performance. Crpyto operations, limited secure storage capabilities. (Crypto Accelerators, Big Number Operation Accelerators such as NXP ZigBee, NFC solutions and AMD PSP). Tamper resistant.
- Trusted Execution Environment
   Microkernel controlled ARM® TrustZone®
   Special instruction based Intel® SGX



### Trusted Execution Environment

- Corresponds to a secure area of main processor of the system. Two virtual systems in one physical.
- SoC architectures highly benefit because of cost and power management ease. i.e. big.LITTLE, MP models.
- Capable of mastering system peripherals so that a system-wide isolation is facilitated
- Relatively high performance and support for 64 systems.
- Has to share one single core with the REE.
- Hack attacks and Shack attacks are protected in TEE. But timing attacks and side channel attacks (still possible to mitigate some) are not avoided.
- TEE Side can access REE side memory if mapping is possible.



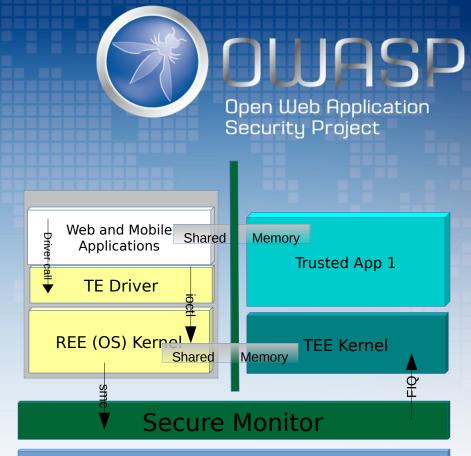
### **Secure Monitor**

### Hardware

ARM® TrustZone® Design

## Trusted Execution Environment: How Does ARM® TrustZone® Work

- User space application establishes a session and world shared memory for data sharing
- Once data is in place, the application notifies the driver via device api.
- Driver translates the notification message and lets related kernel module to issue a Secure Monitor Call.
- Secure Monitor handles the interrupt, configures some critical registers and switches the context to counterparting trusted application.
- The trusted application copies the data into a non-shared memory block, processes and returns the response to the shared memory.
- Trusted application issues a secure interrupt to switch context to normal world.



### Hardware

ARM® TrustZone® Design



## **Exisiting TEE Services**

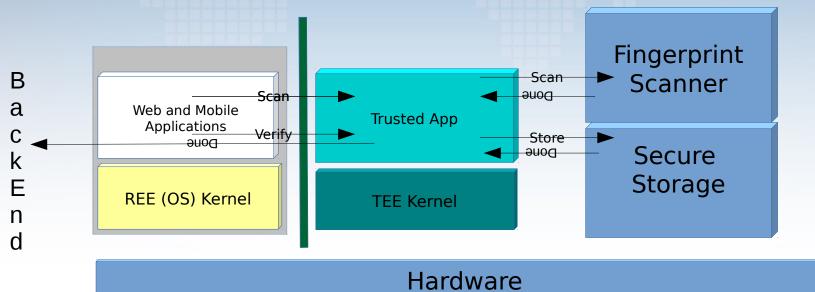
- Key management
- Secure crypto operations
- Verification and signing operations
- Biometrics and Simple authentication features.(FIDO like authentication)
- REE runtime memory protection
- Secure Code Execution
- Secure boot chain
- Secure storage
- Trusted User Interface
- Digital Rights Management

## **Authentication Service**

OUASP
Open Web Application
Security Project

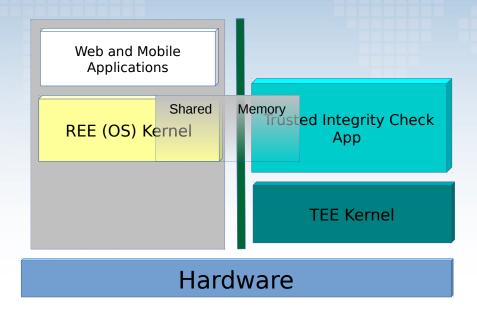
With Hardware Support

- FIDO like password free authentication is possible.
- Device authentication is possible.
- Fingerprint Authentication is possible: Scanning, Storing and Verification





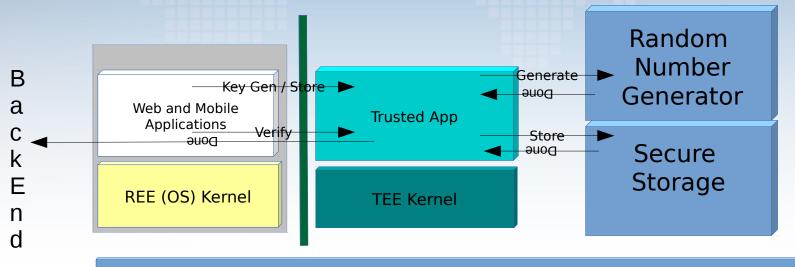
- Runtime integrity check for REE kernel performed by TEE application.
- Code block of the kernel has to remain the same as boot time.
- Checks are performed periodically (watchdog timer) or patagonix (PTE non-exec, when attempted exception triggers to check)





# Key Management

- ODUASP
  Open Web Application
  Security Project
- Keys are stored and used respective to the applications owning them
- All of the keys are derived using RNG
- Keys are never exposed to the REE as well as cypto operations
- Android Key Master compatible technology(Jbean)
- Heartbleed could be avoided

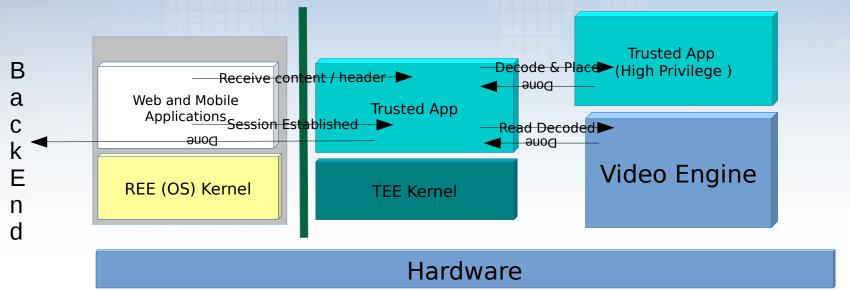


Hardware



# Digital Rights Management

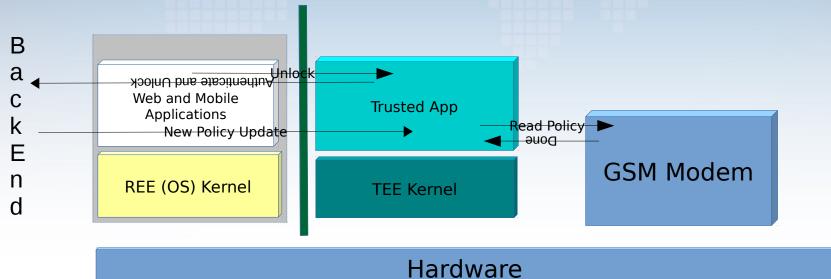
- Content sharing and session key generation
- Encoding and decoding support
- Securing video engine and its bus.
- Storing and reusing session keys.
- Authentication of device





## SIM Lock

- Sim lock operations are enabled via TEE
- Proxying new unlock policy
- Device Authentication is done by TEE
- Security is directly related to GSM modem software





Thanks for your time. Any questions?

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