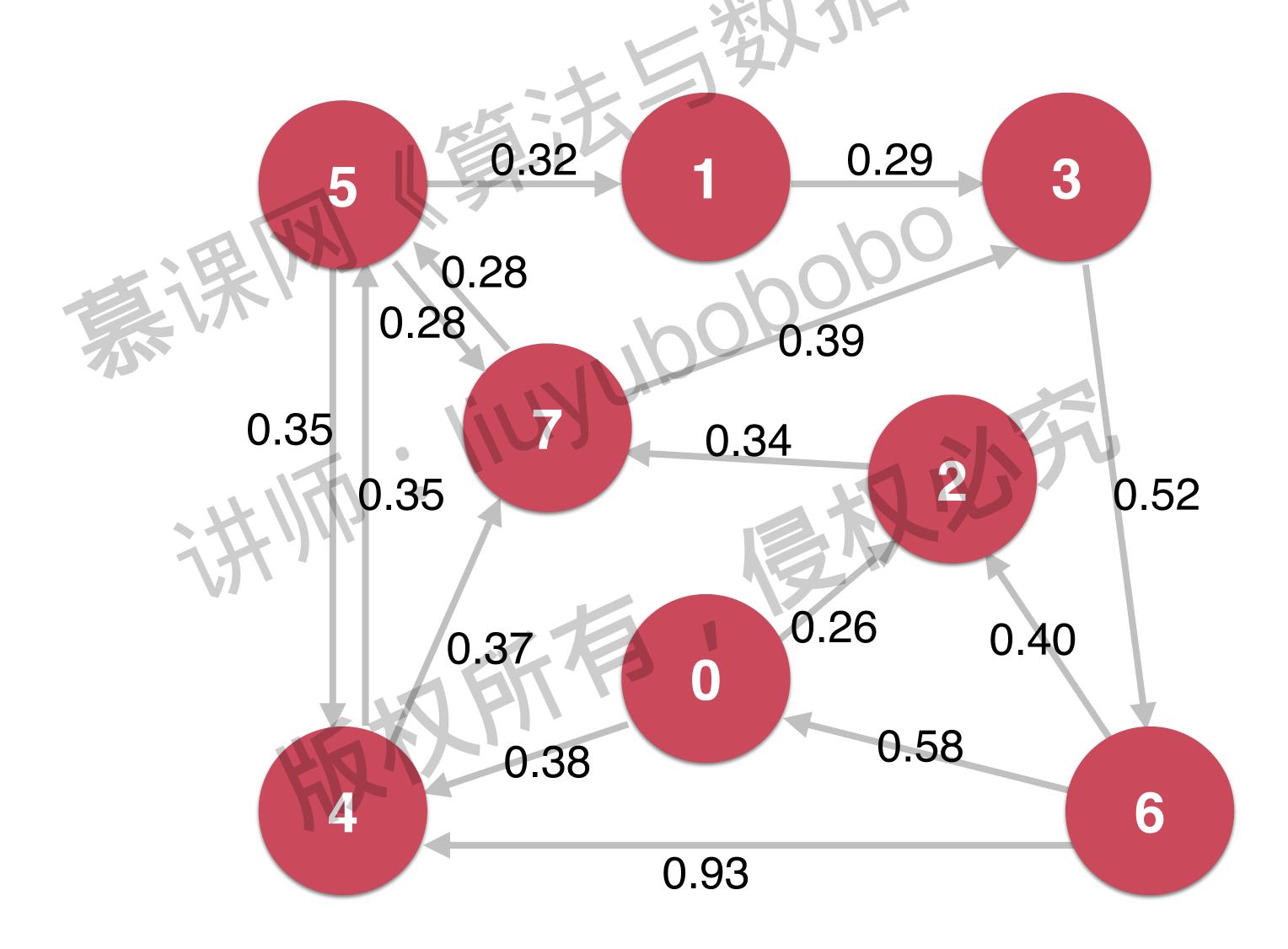
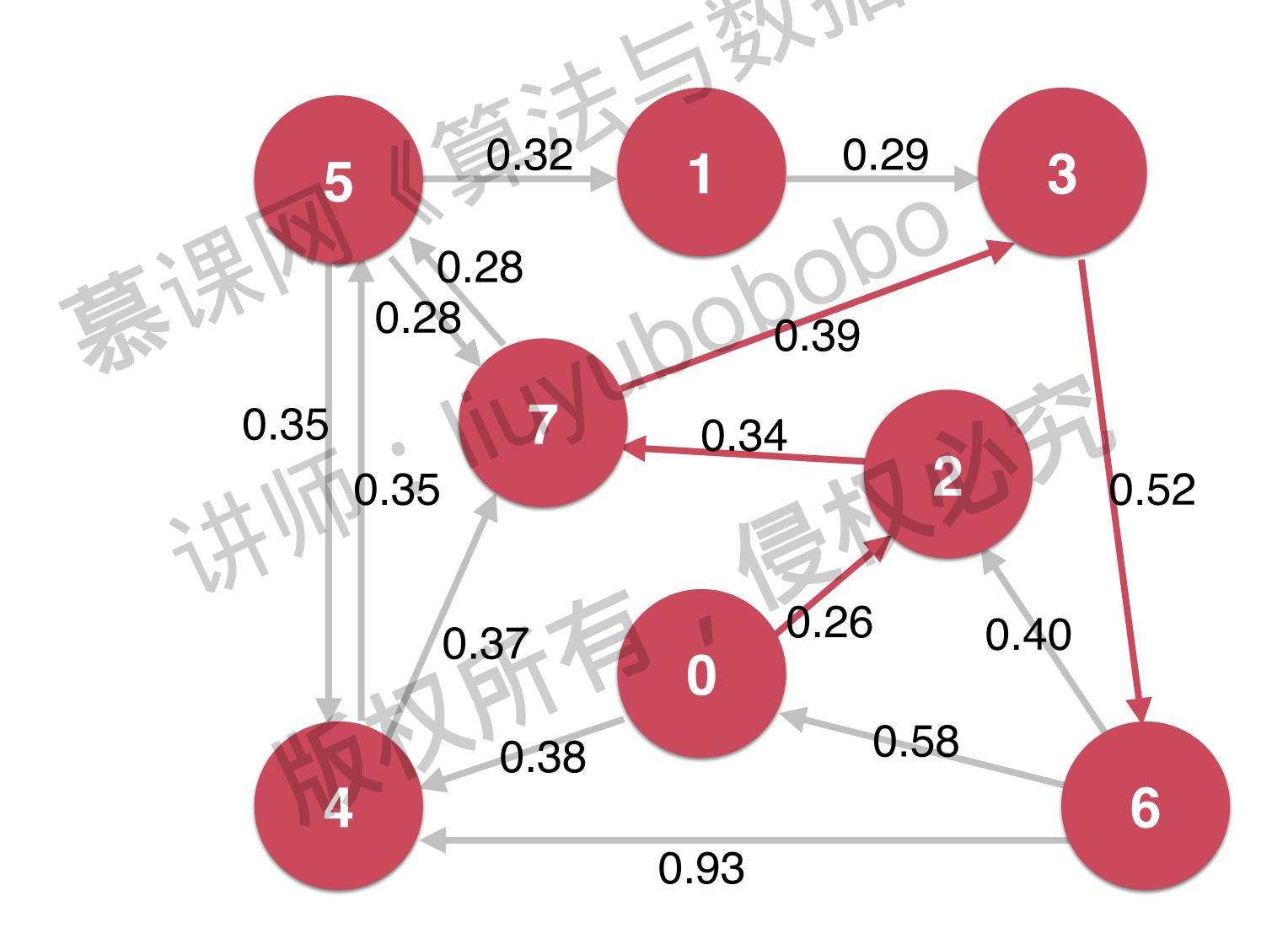
源·果网 (算·法与数据生法构) 神版权所 liuyubobobo

最路径问题。Shortest Path

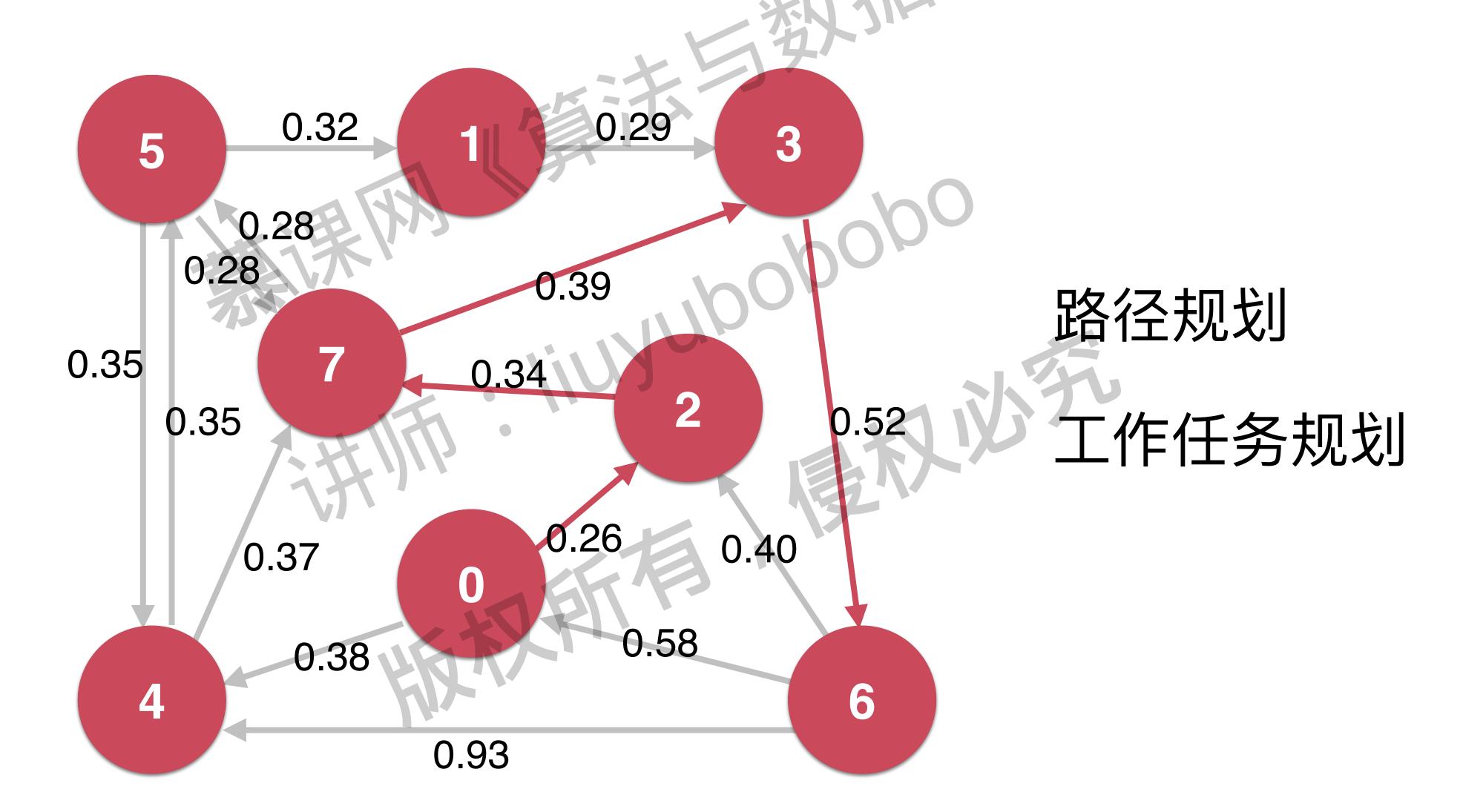
最短路径问题 Shortest Path



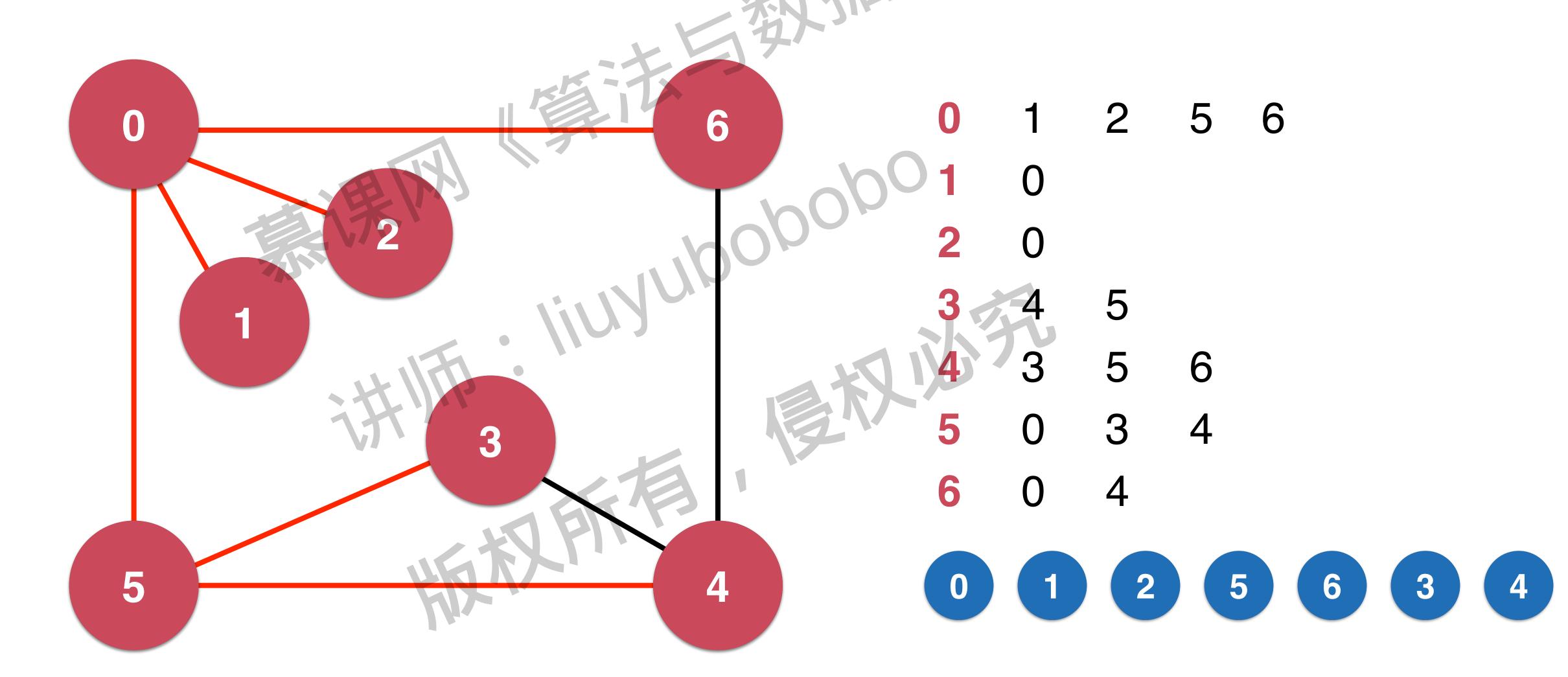
最短路径问题 Shortest Path



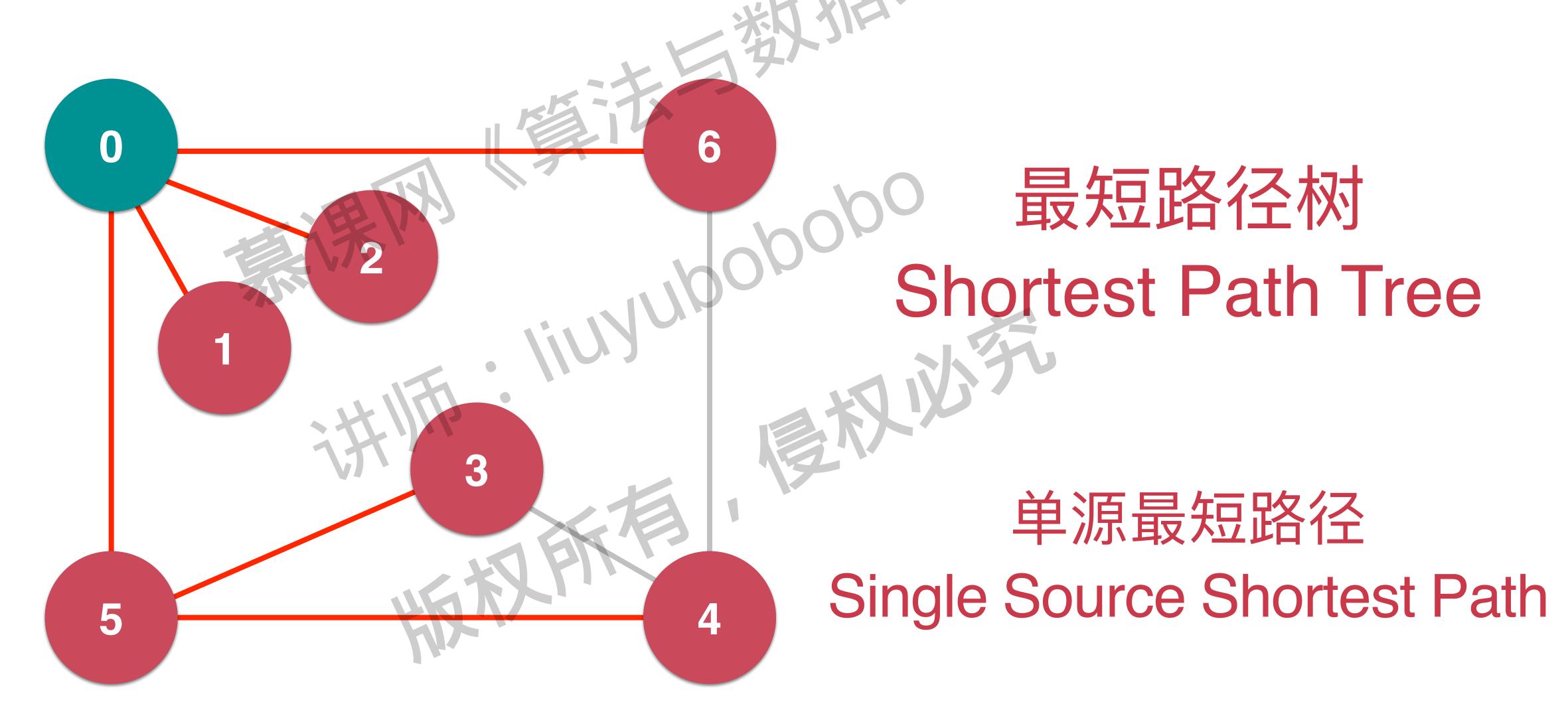
最短路径问题 Shortest Path



广度优先遍历



广度优先遍历



无权图的最短路径

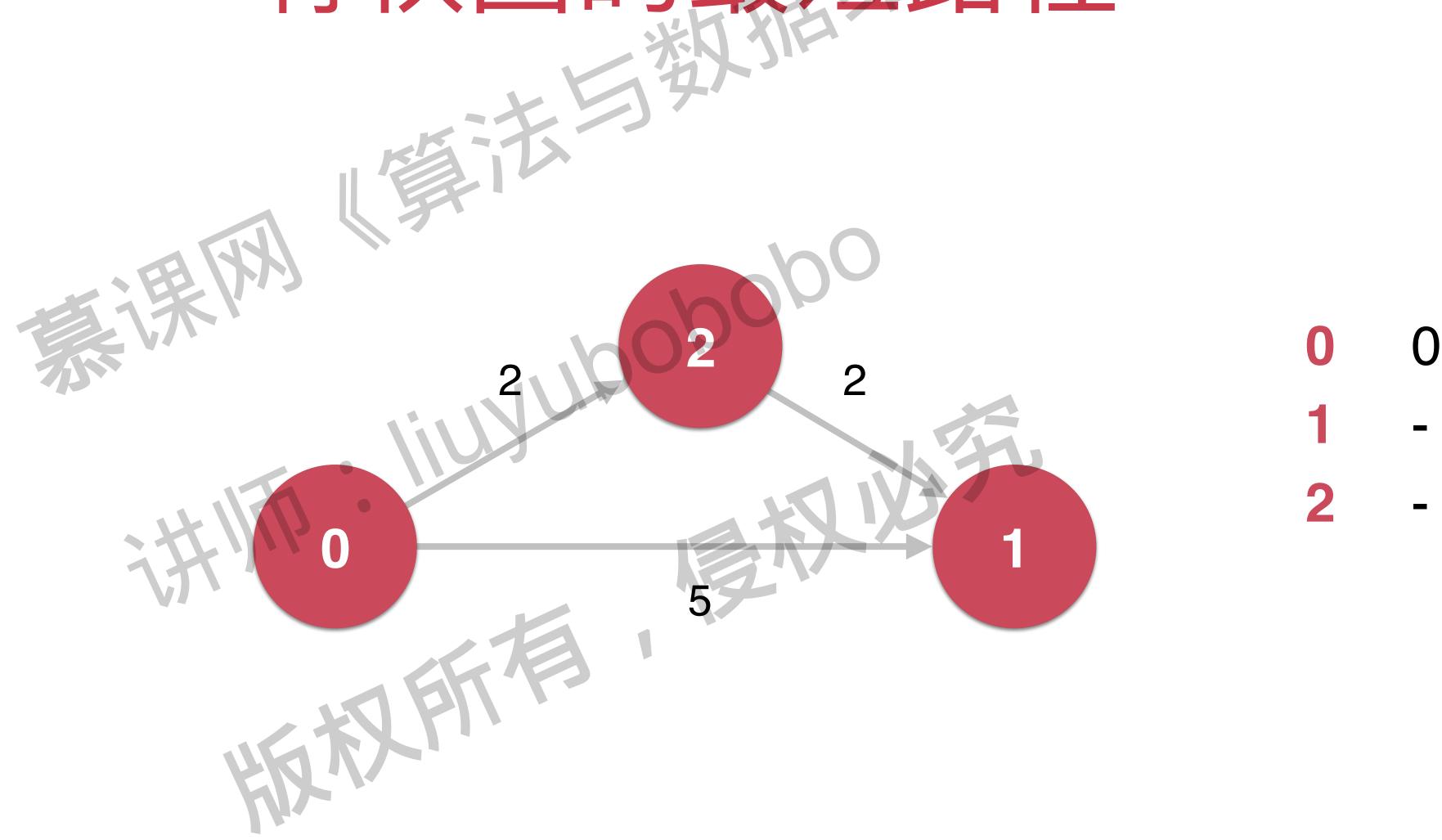


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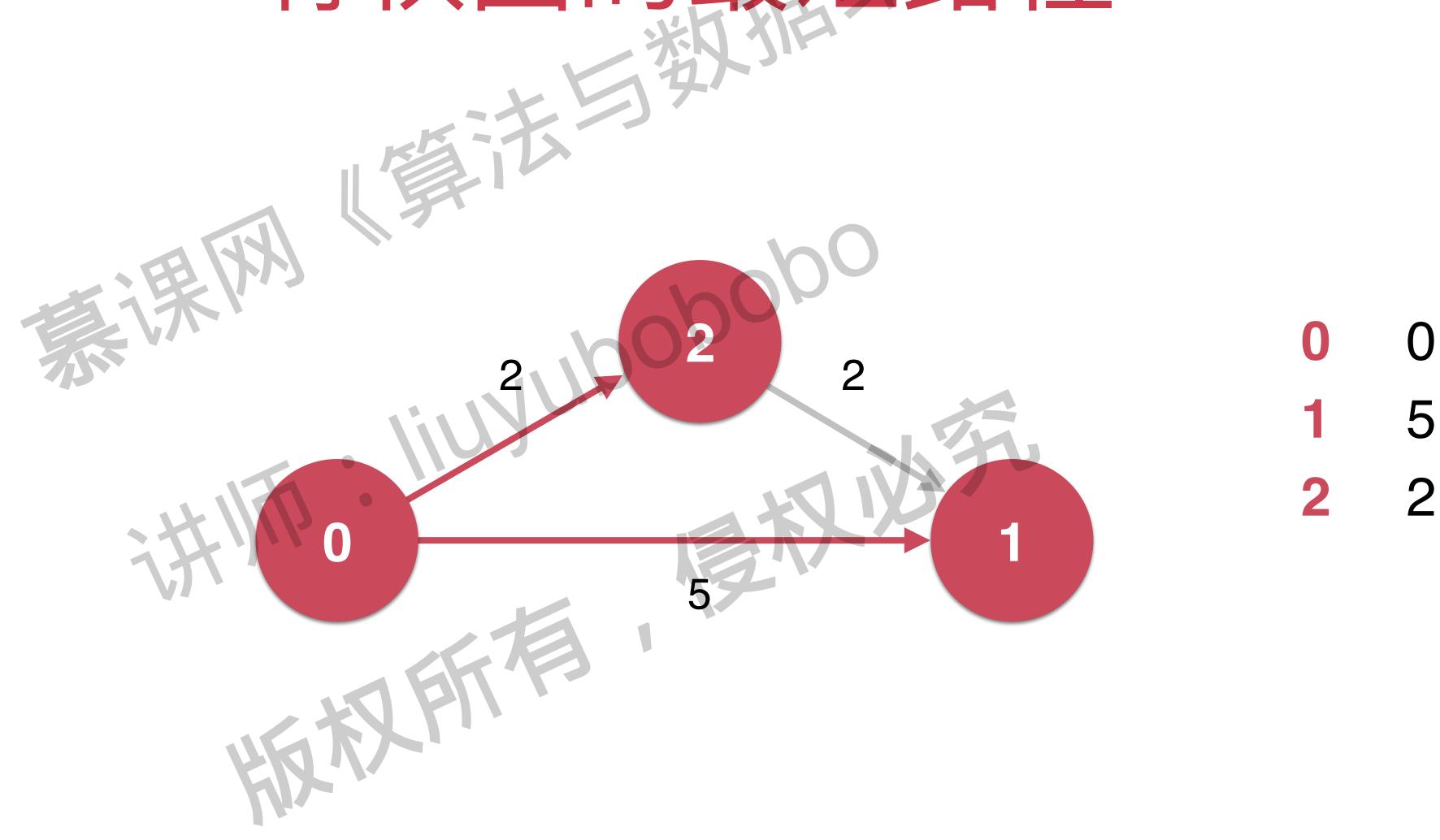
无权图的最短路径



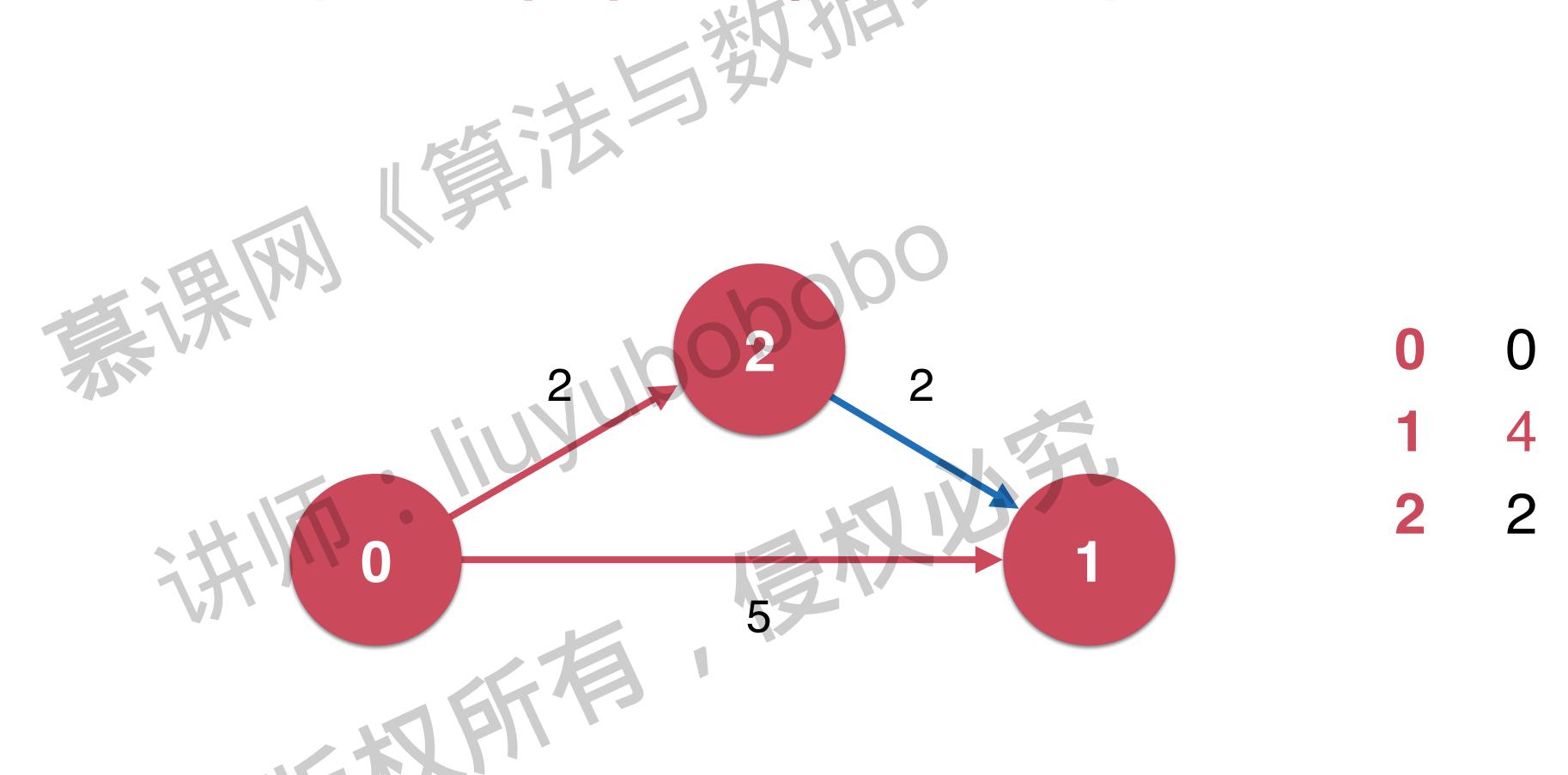
有权图的最短路径



有权图的最短路径



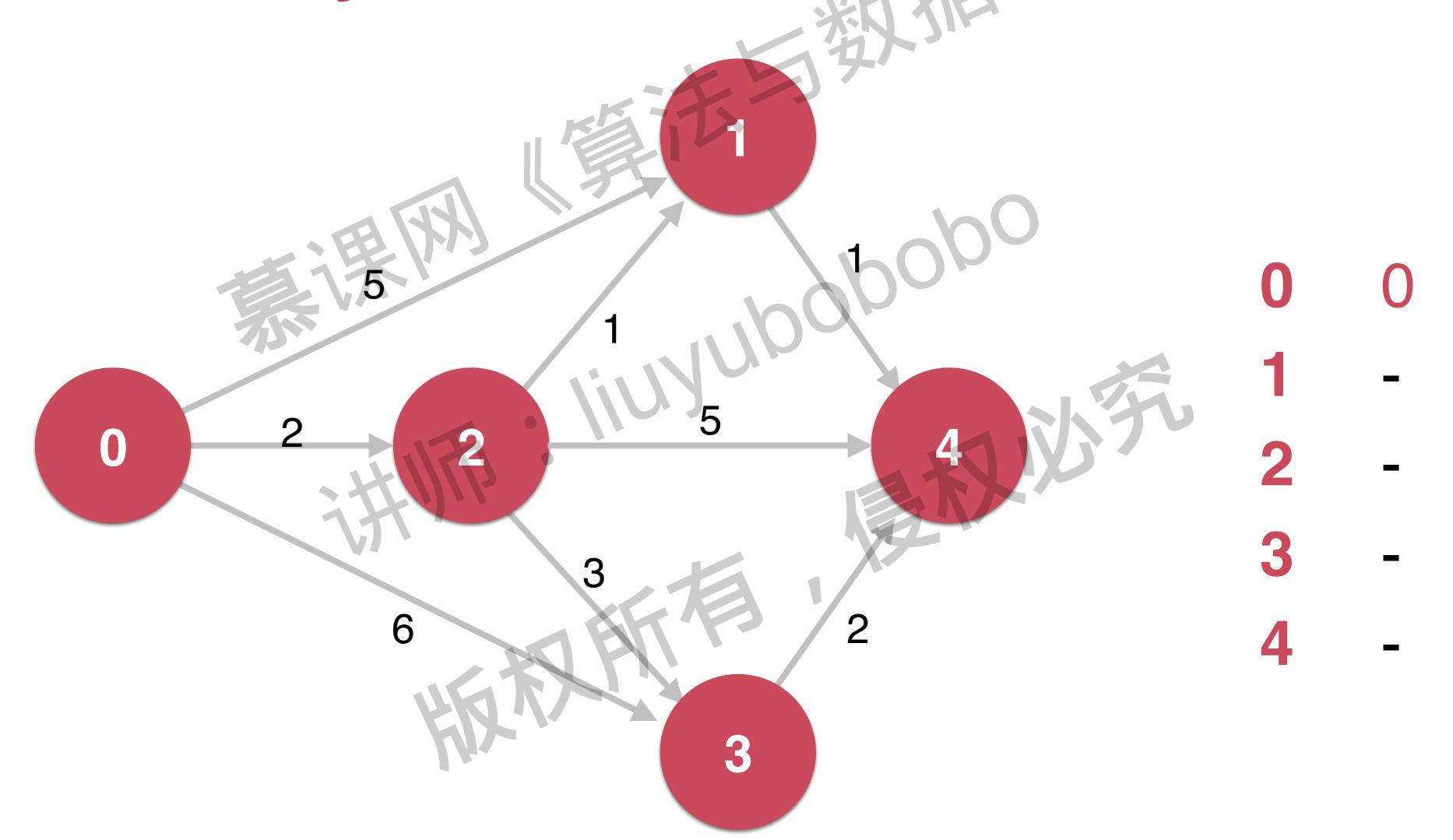
有权图的最短路径

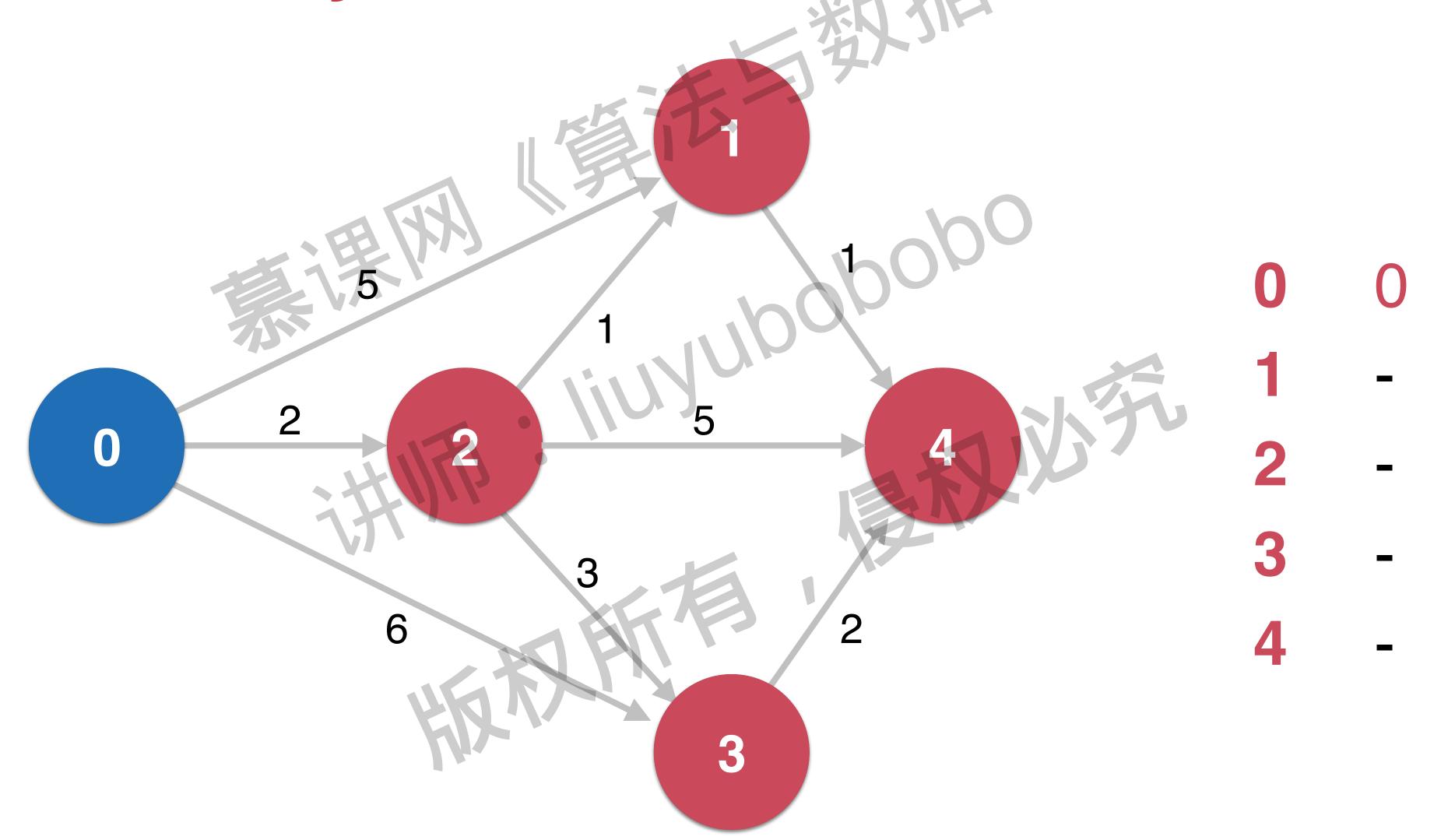


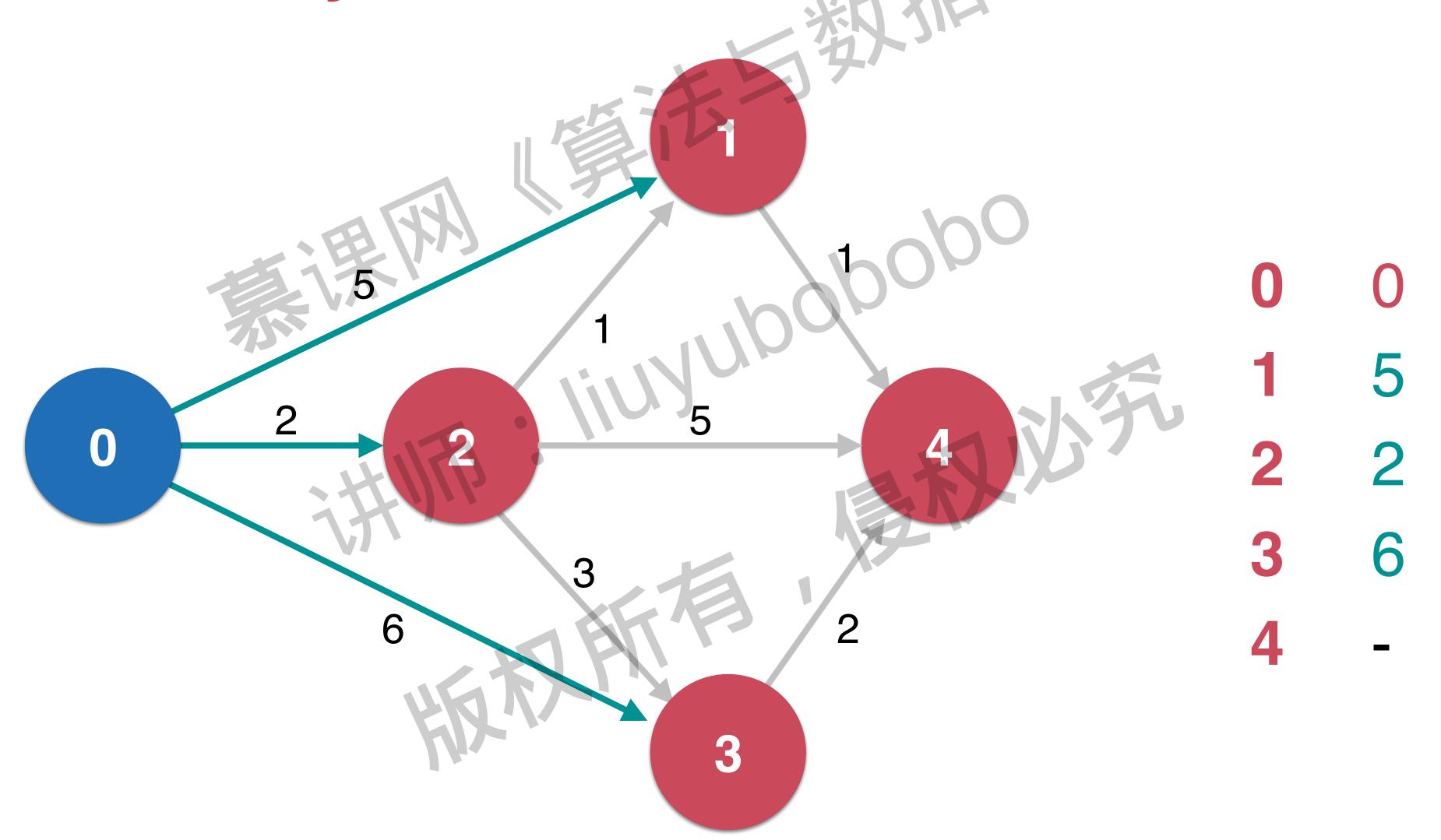
松弛操作 Relaxation

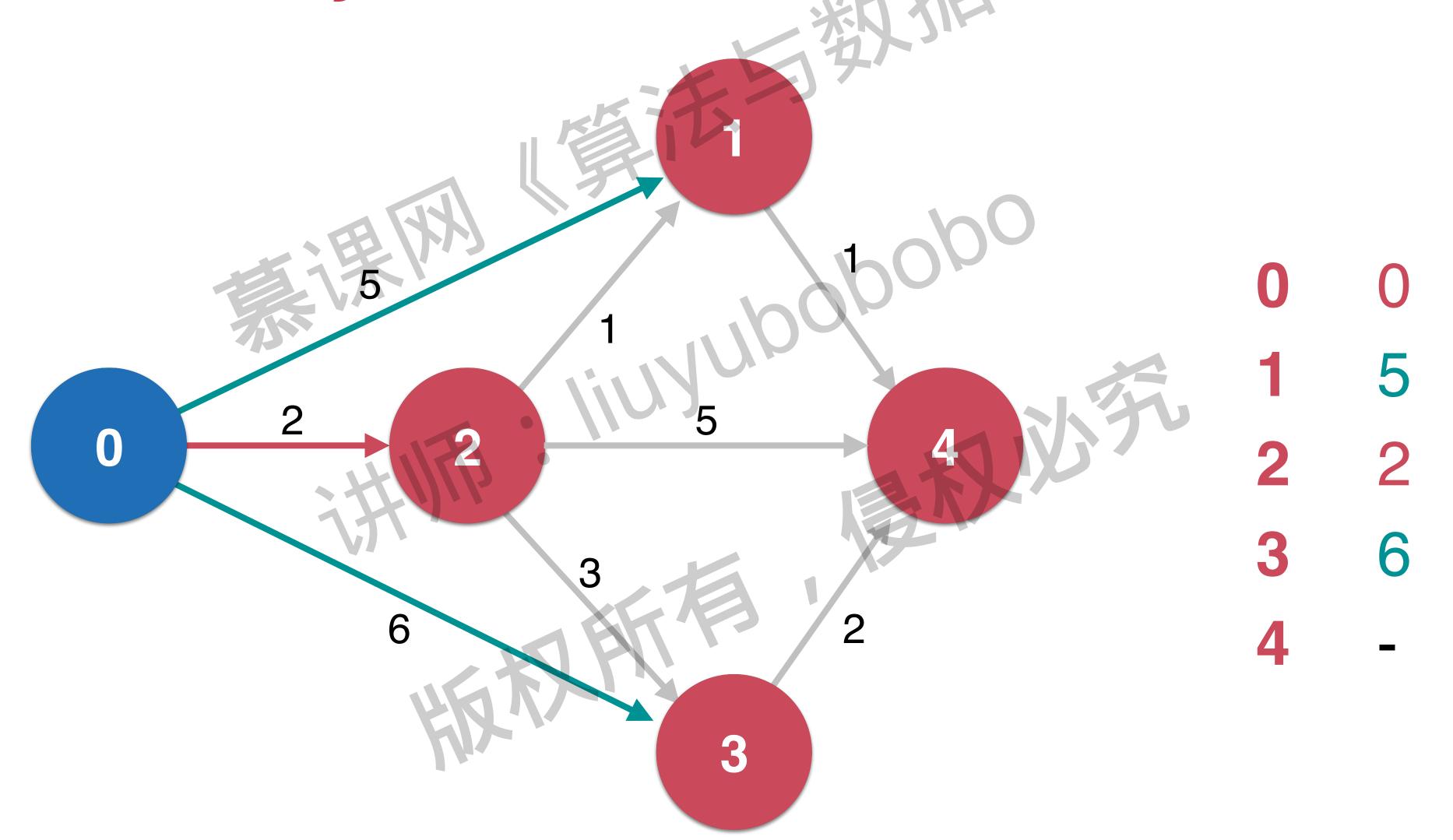
"人" 《第二法》 松弛操作是最短路径求解的核心

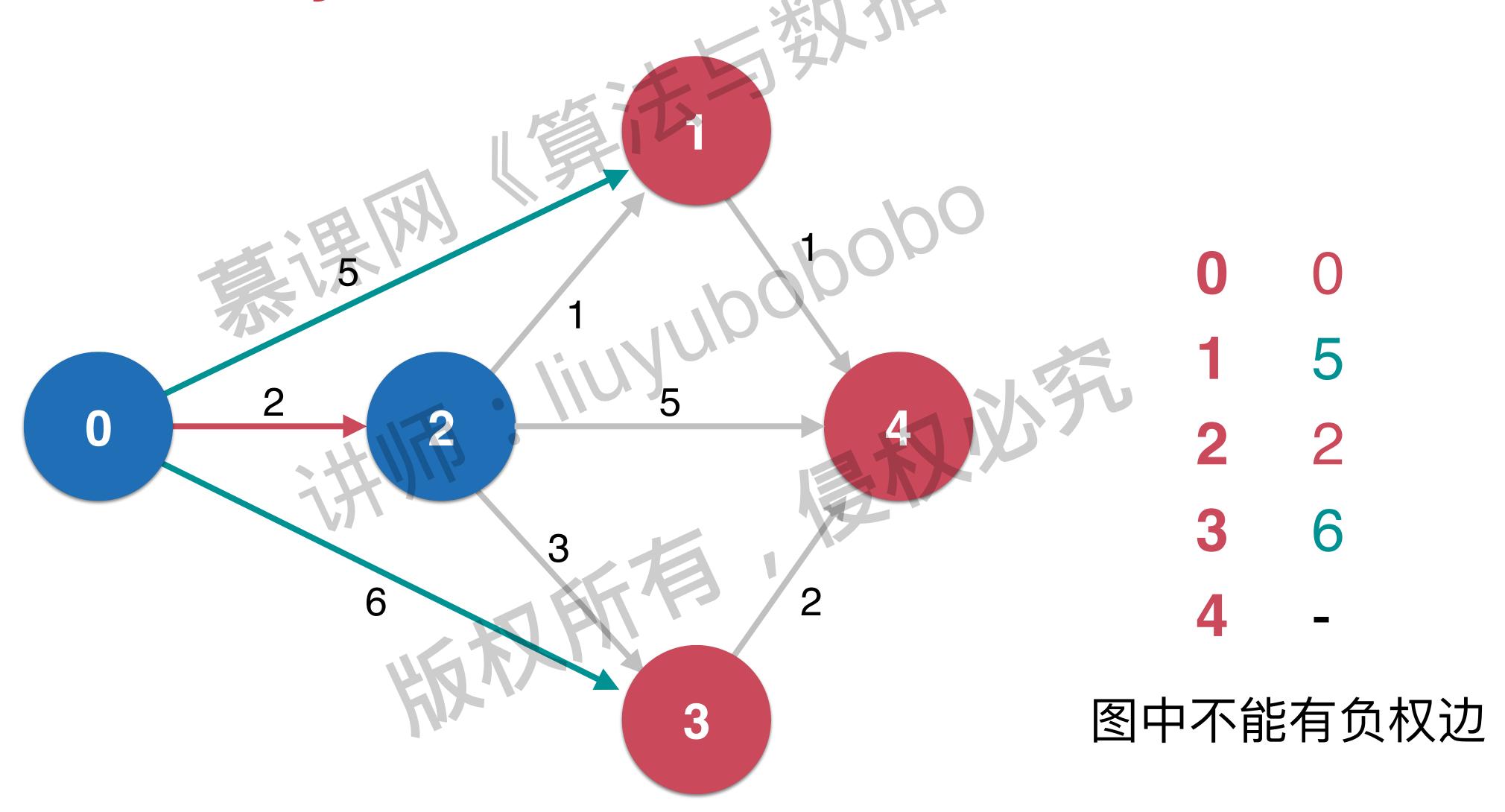
前提:图中不能有负权边 复杂度 O(E log(V))

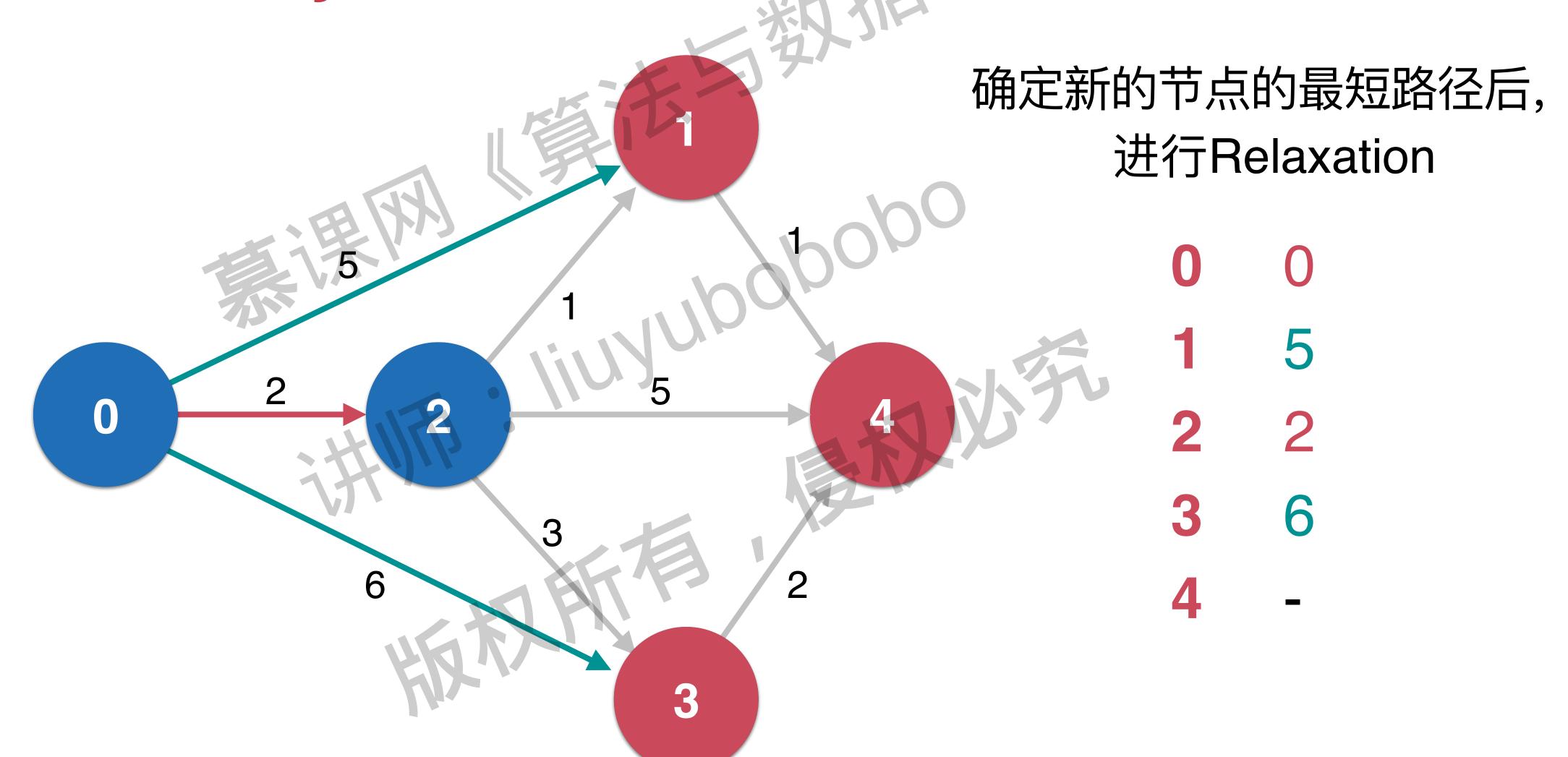


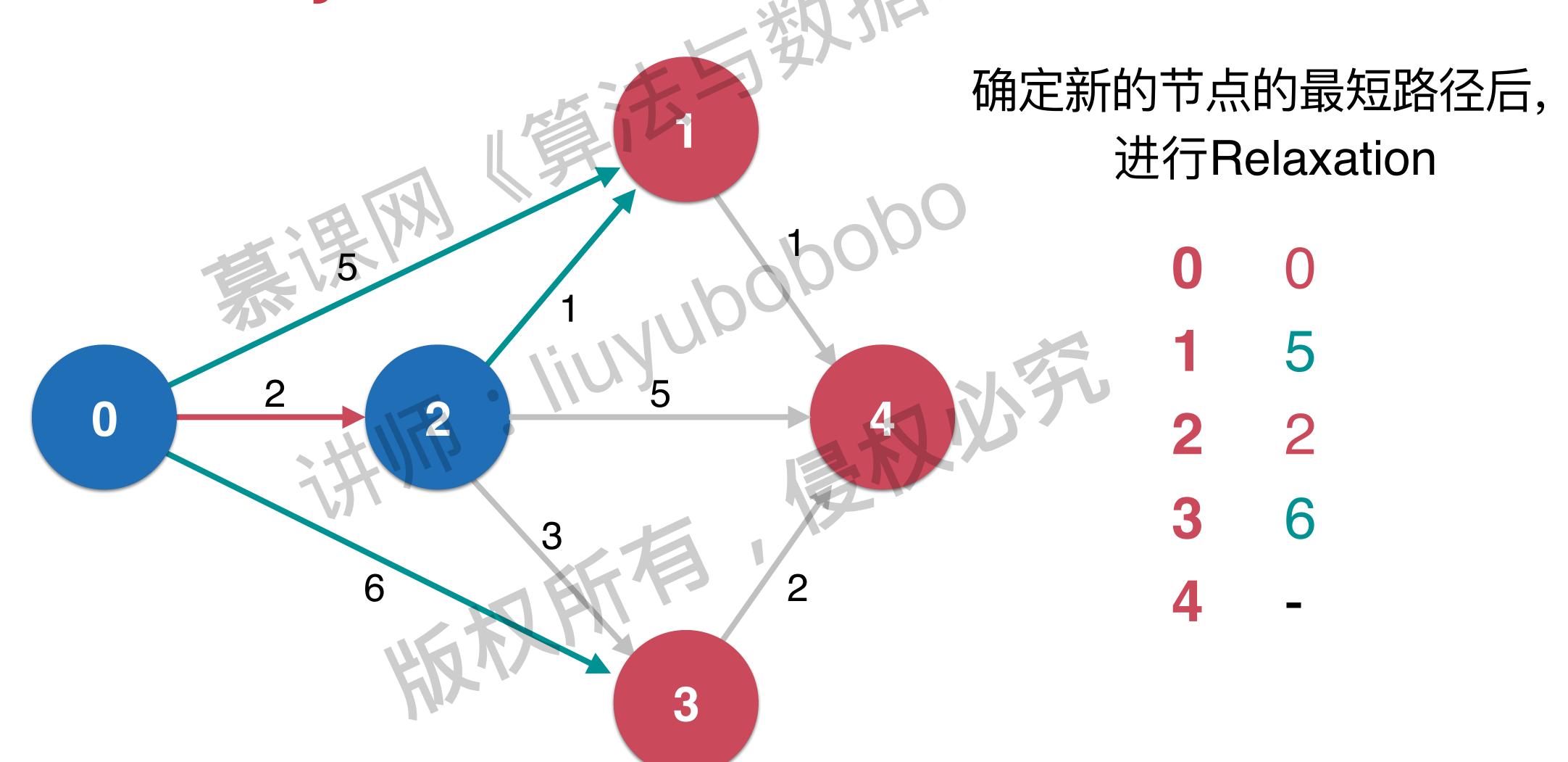


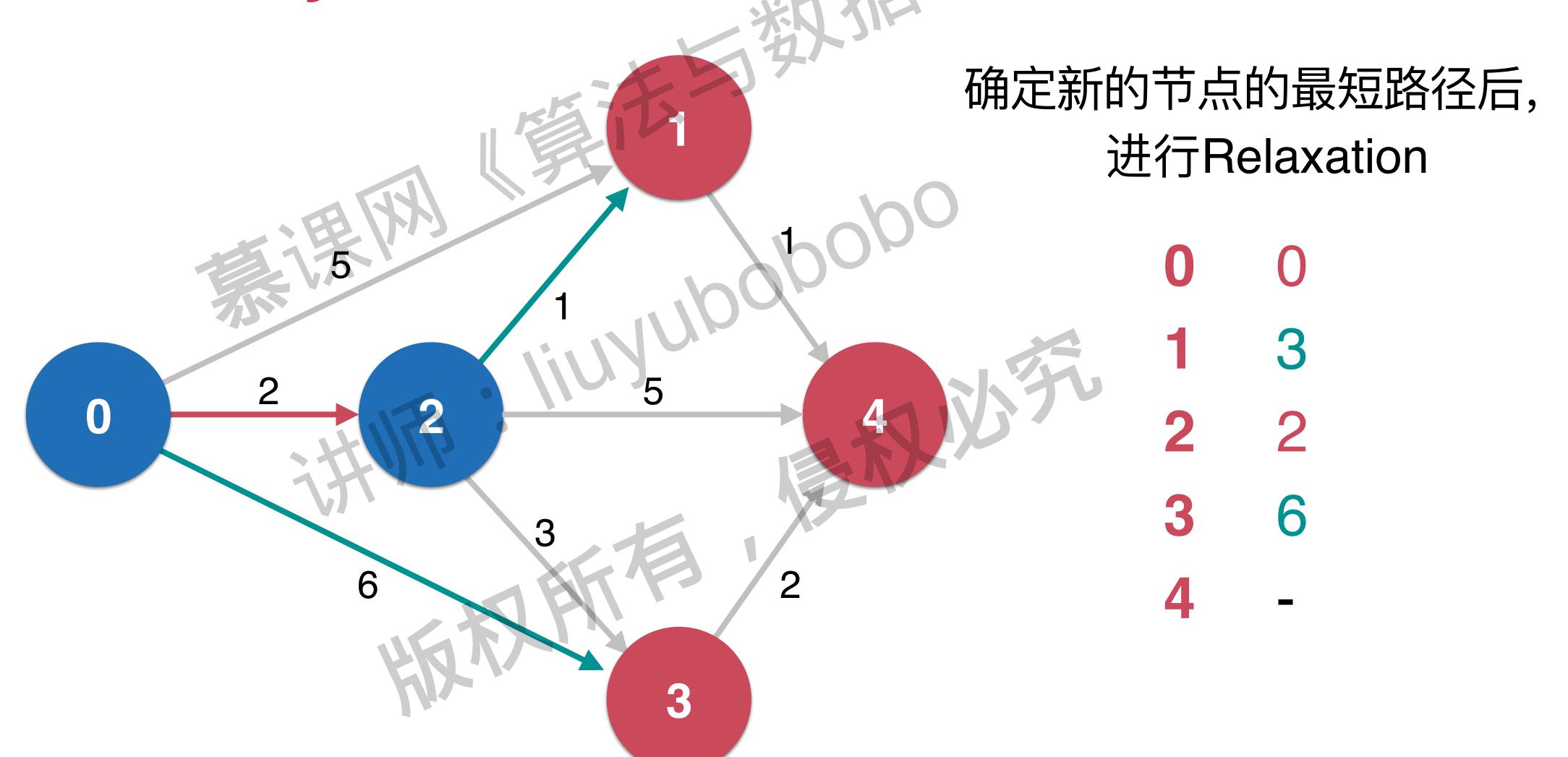


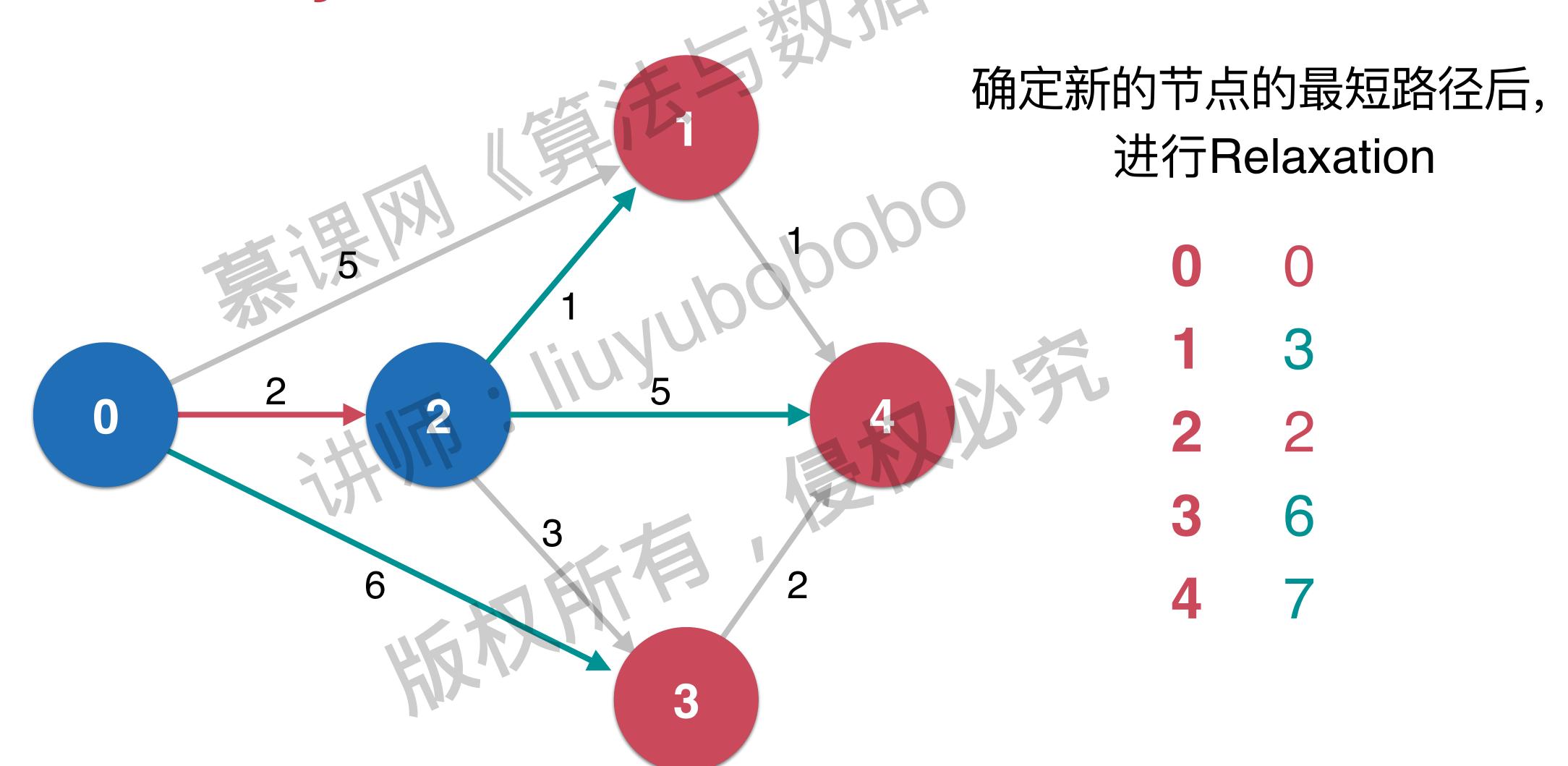


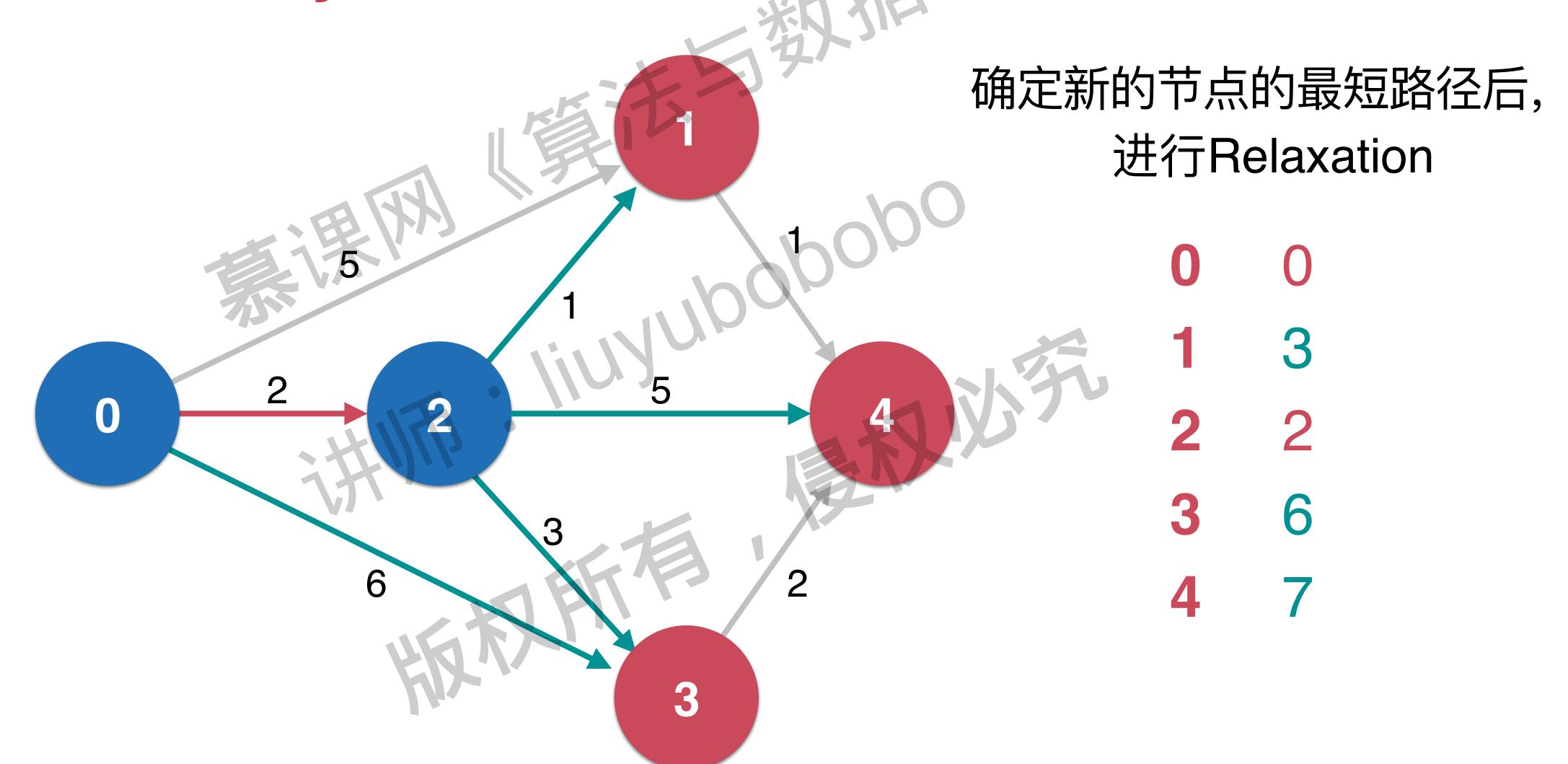


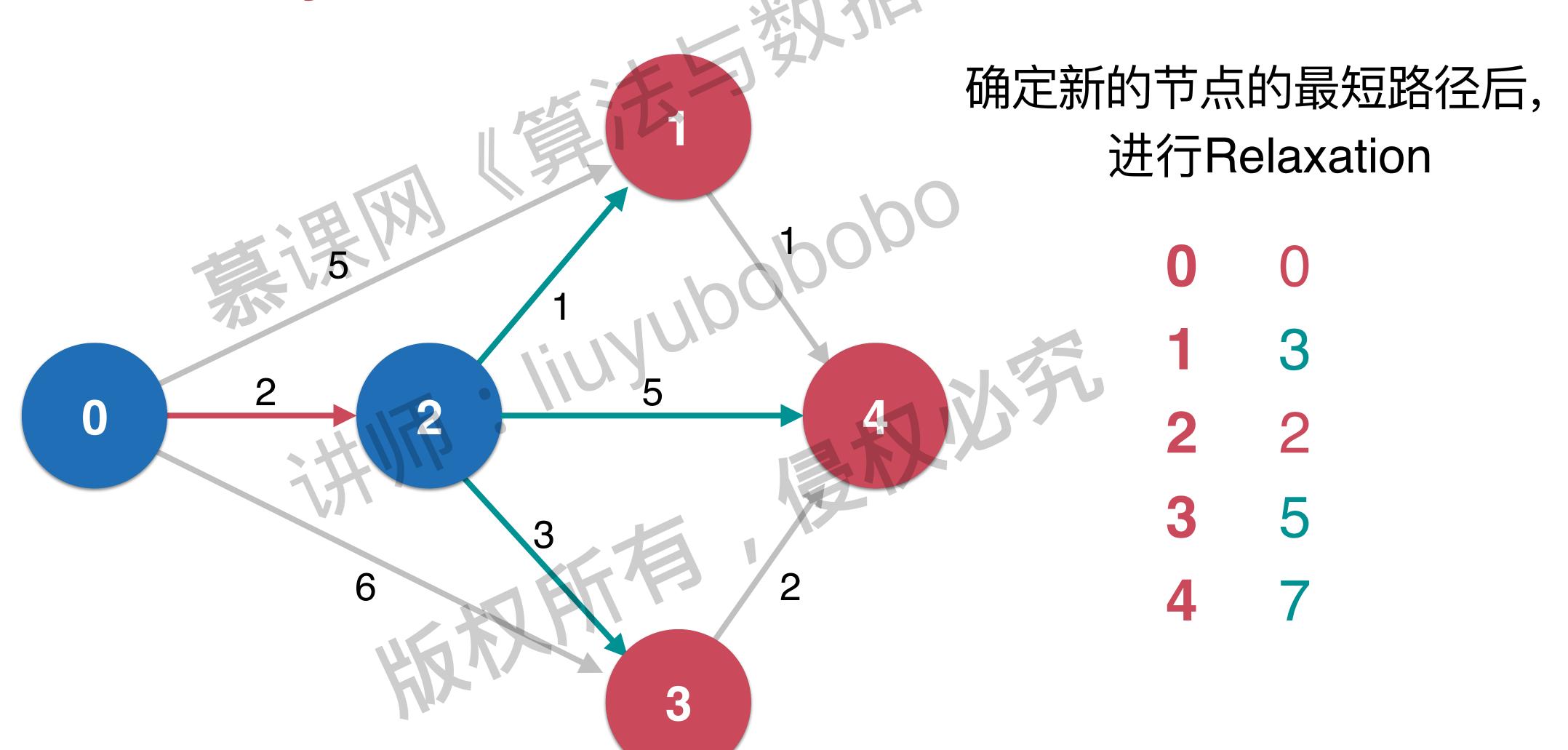


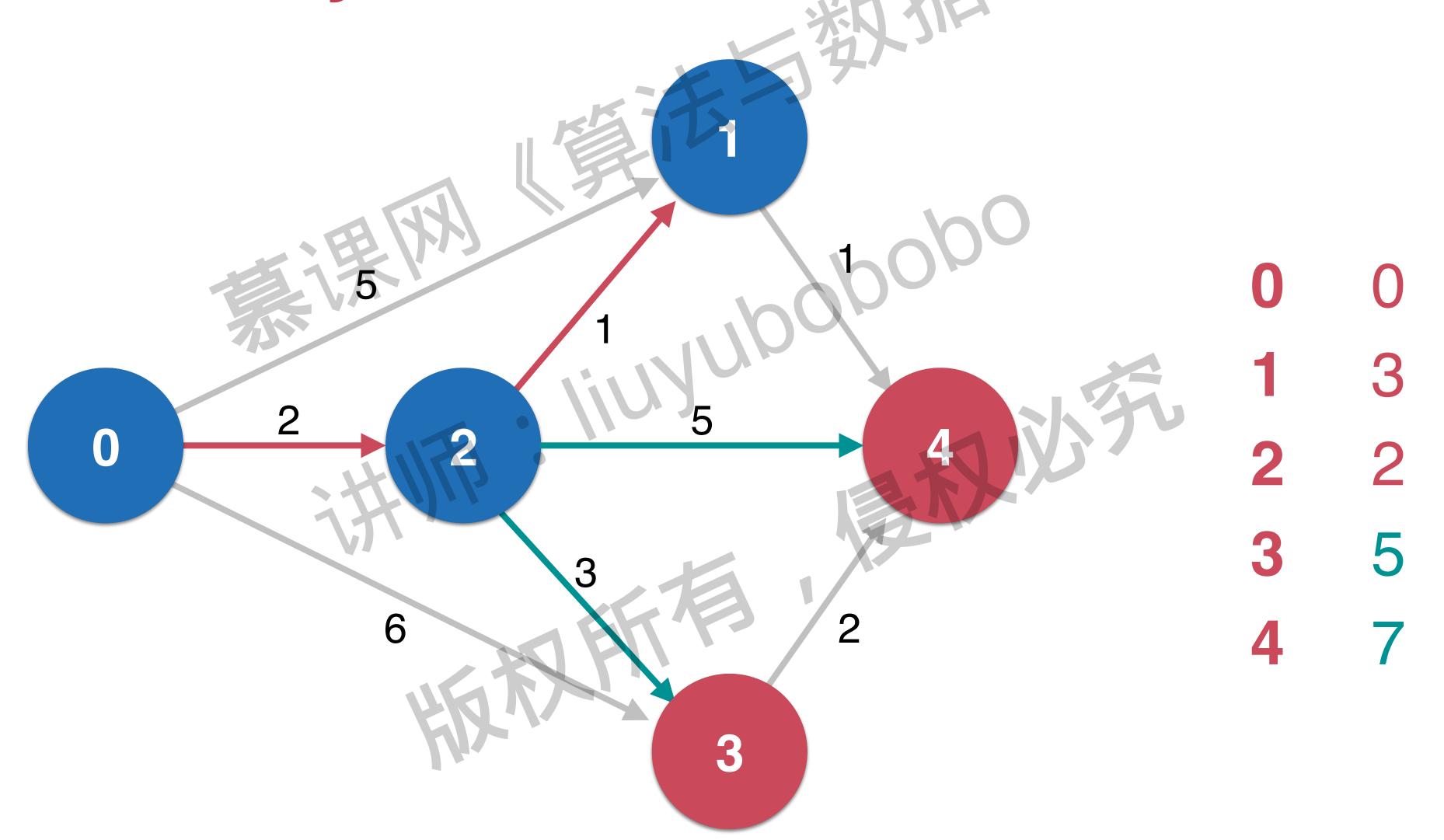


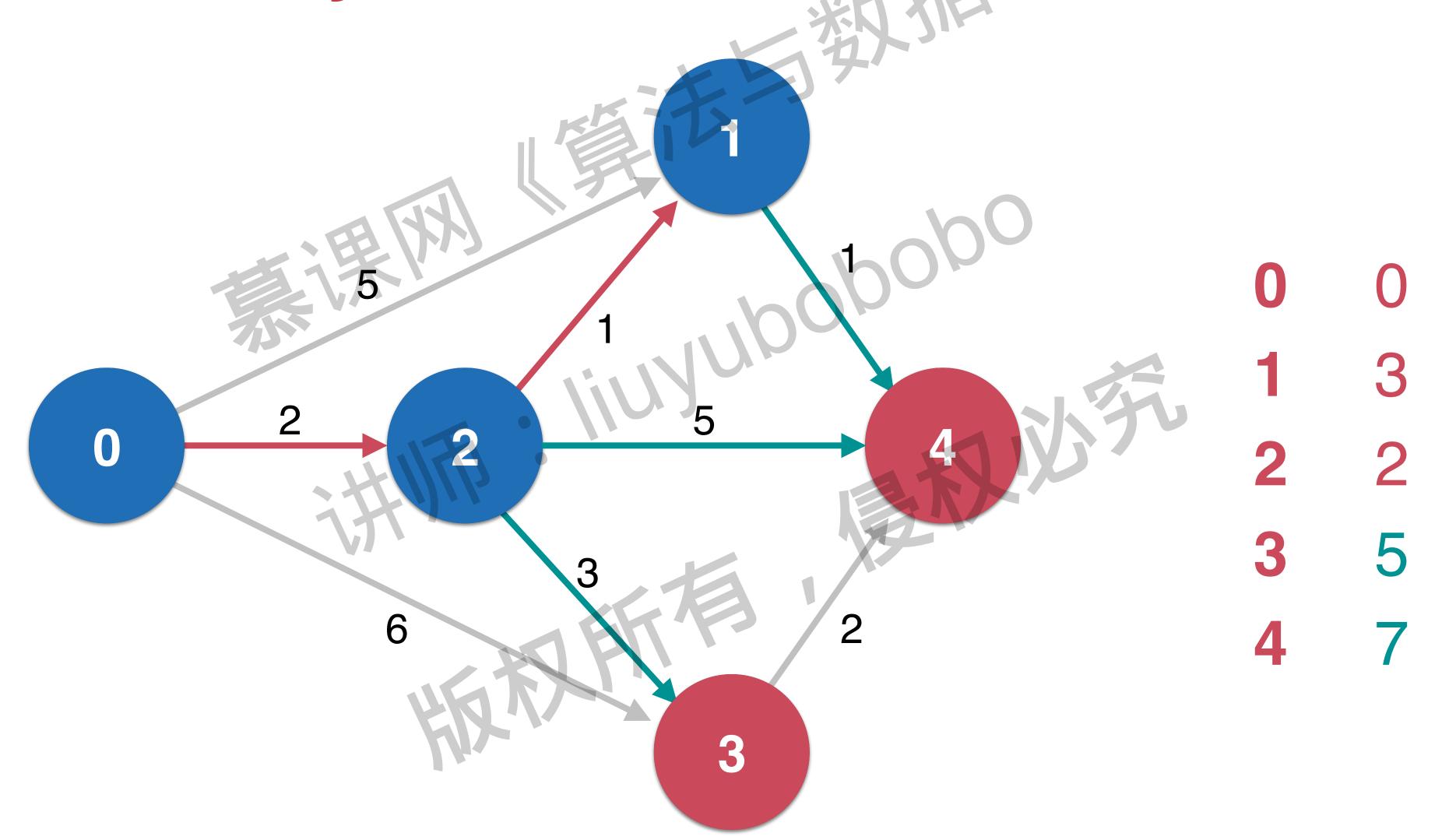


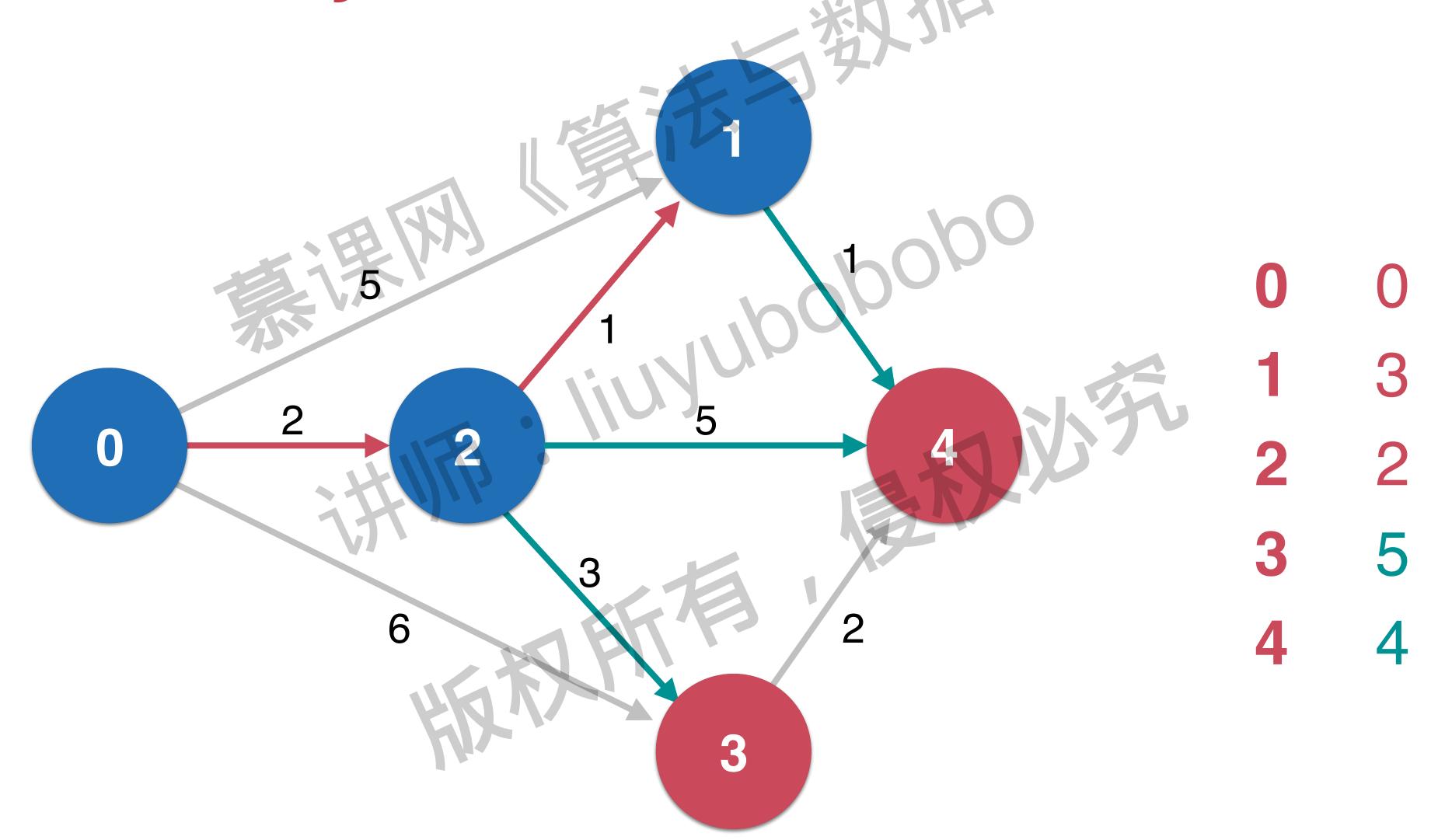


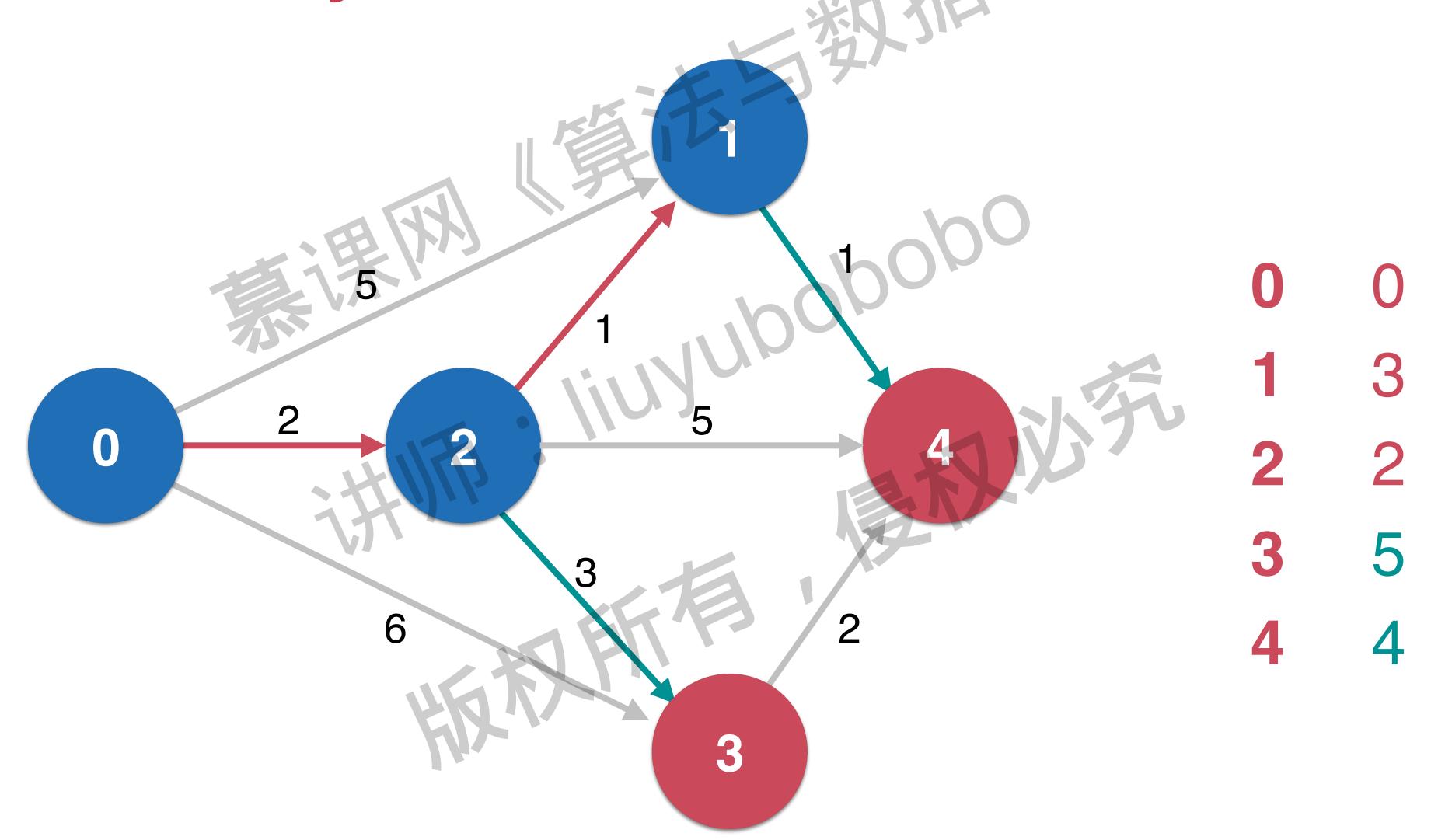


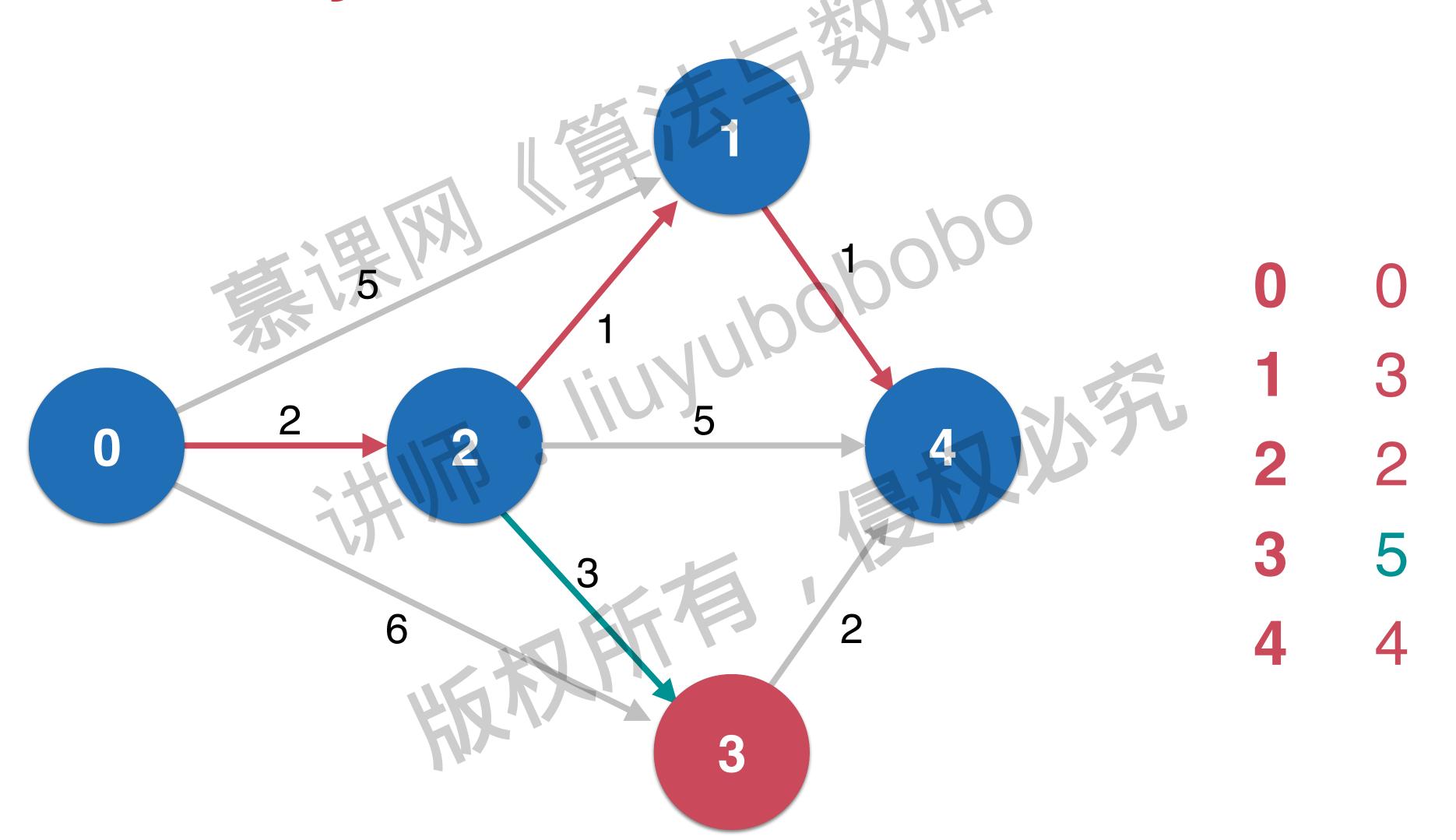


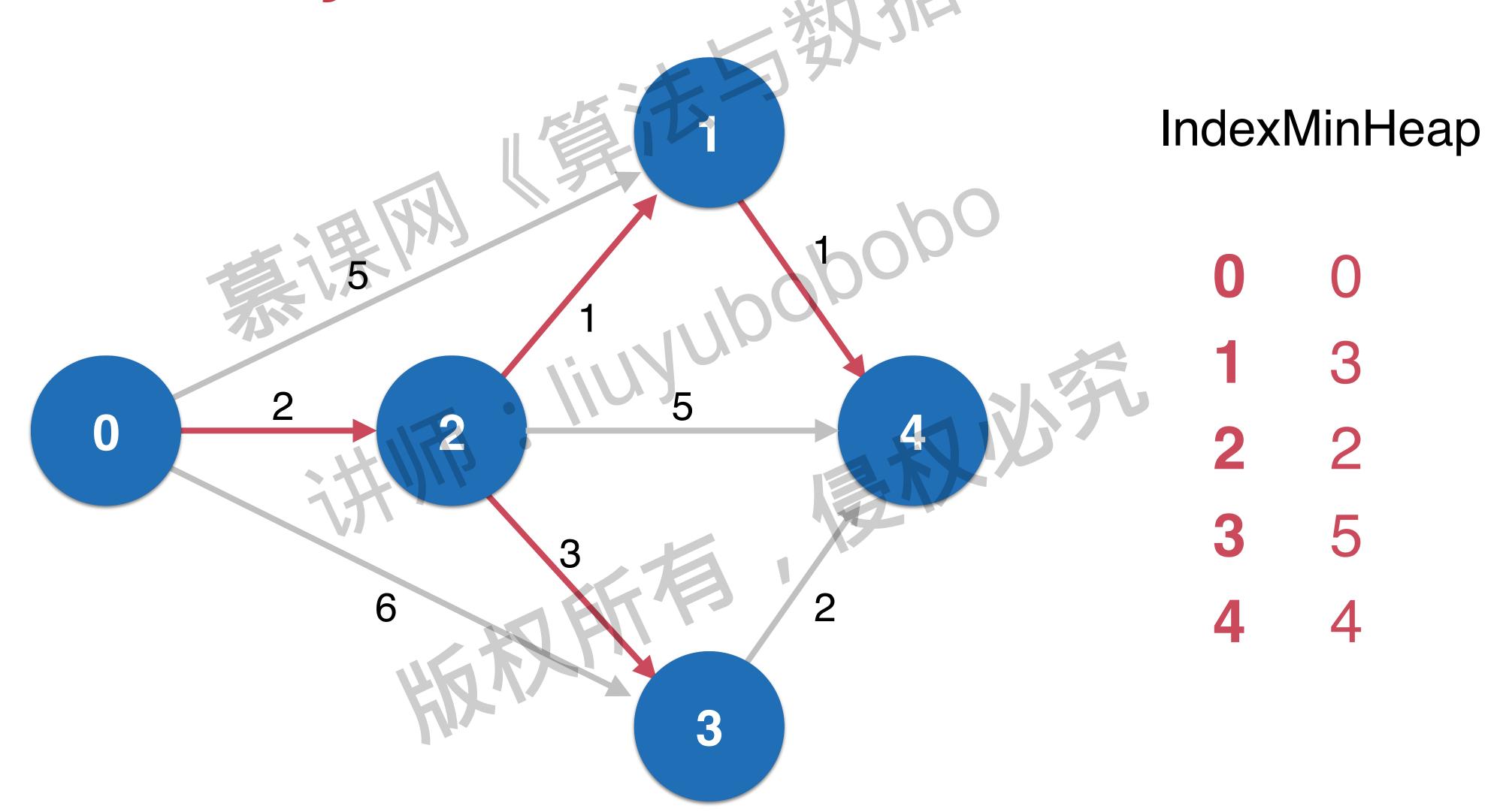






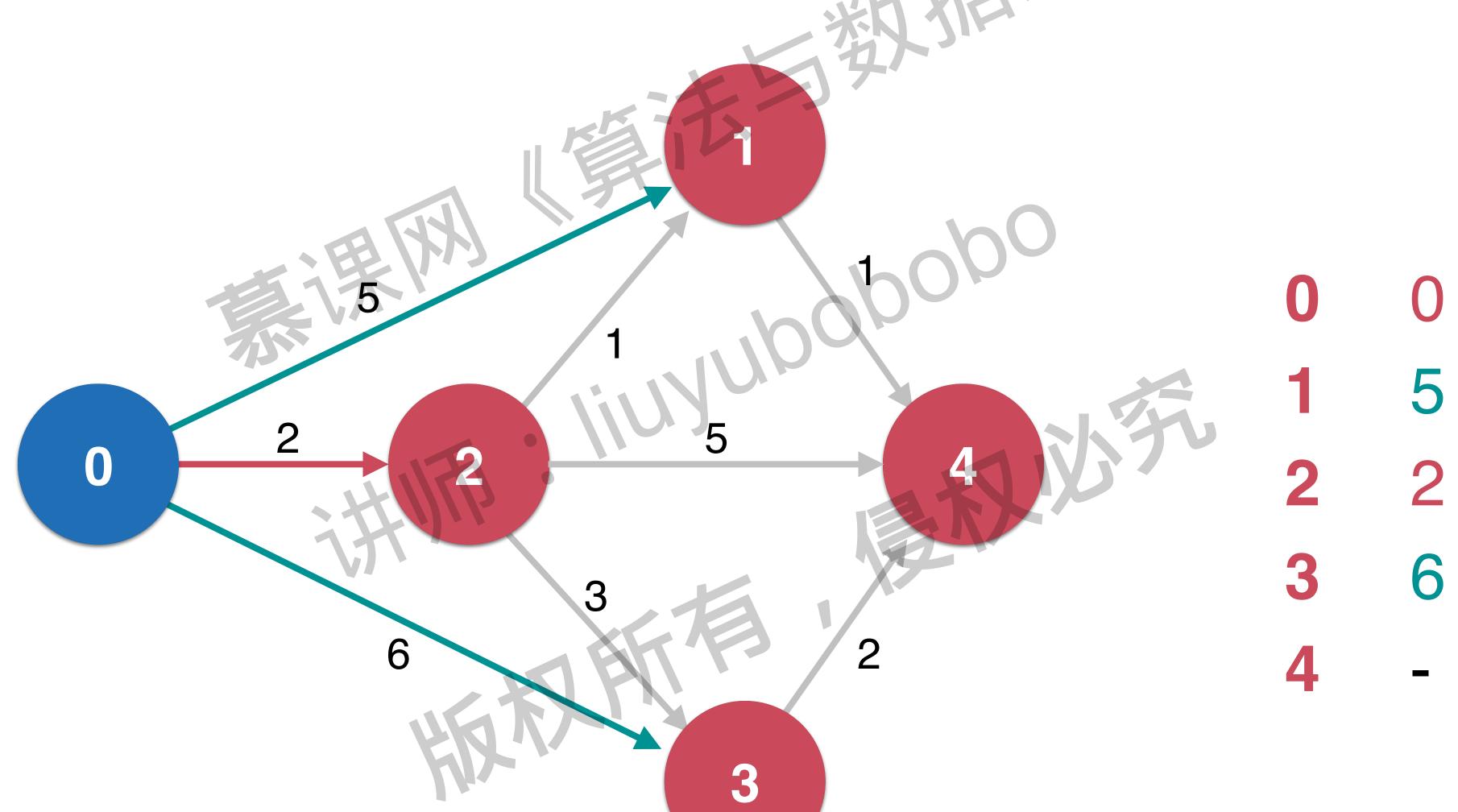




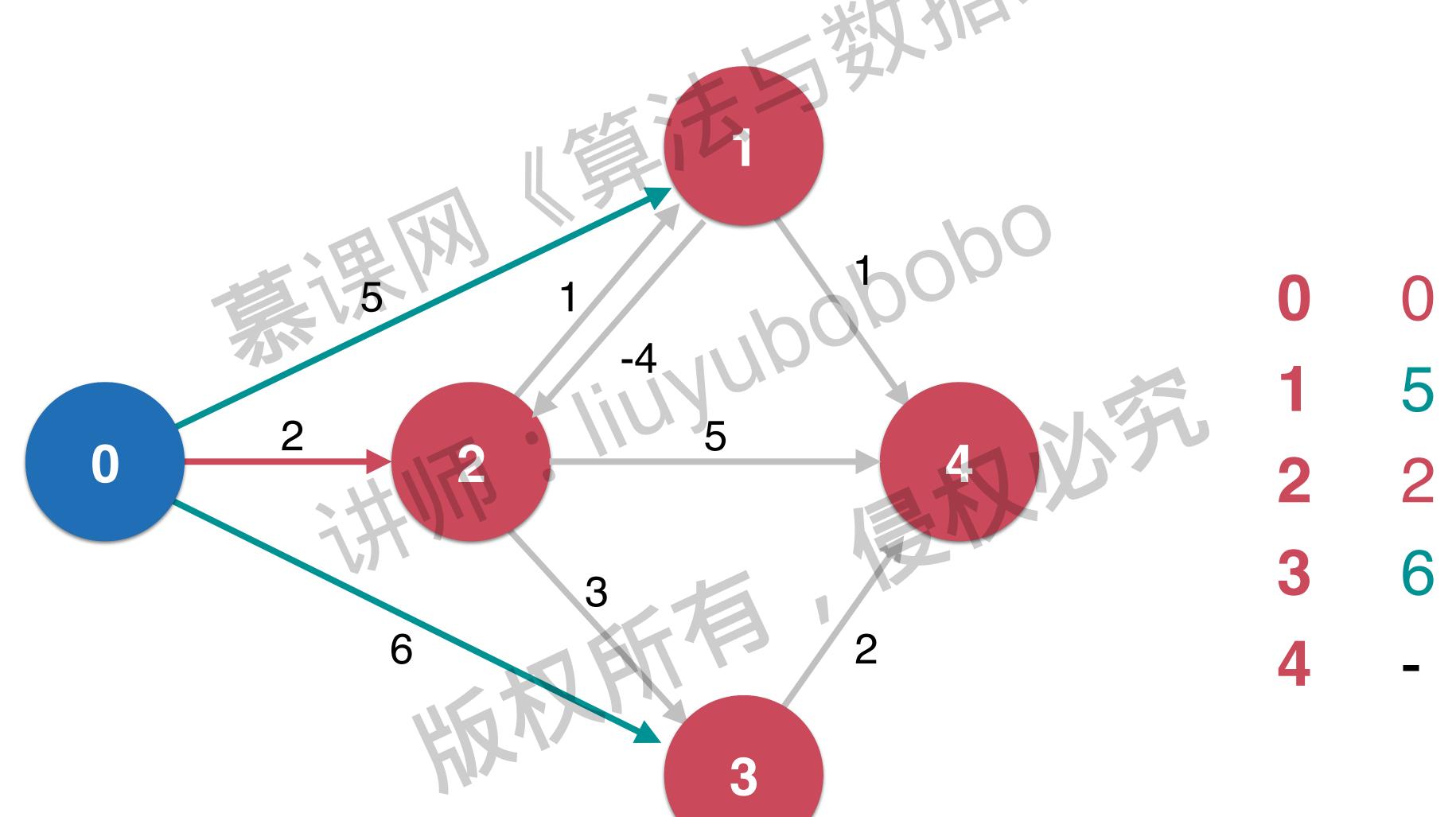


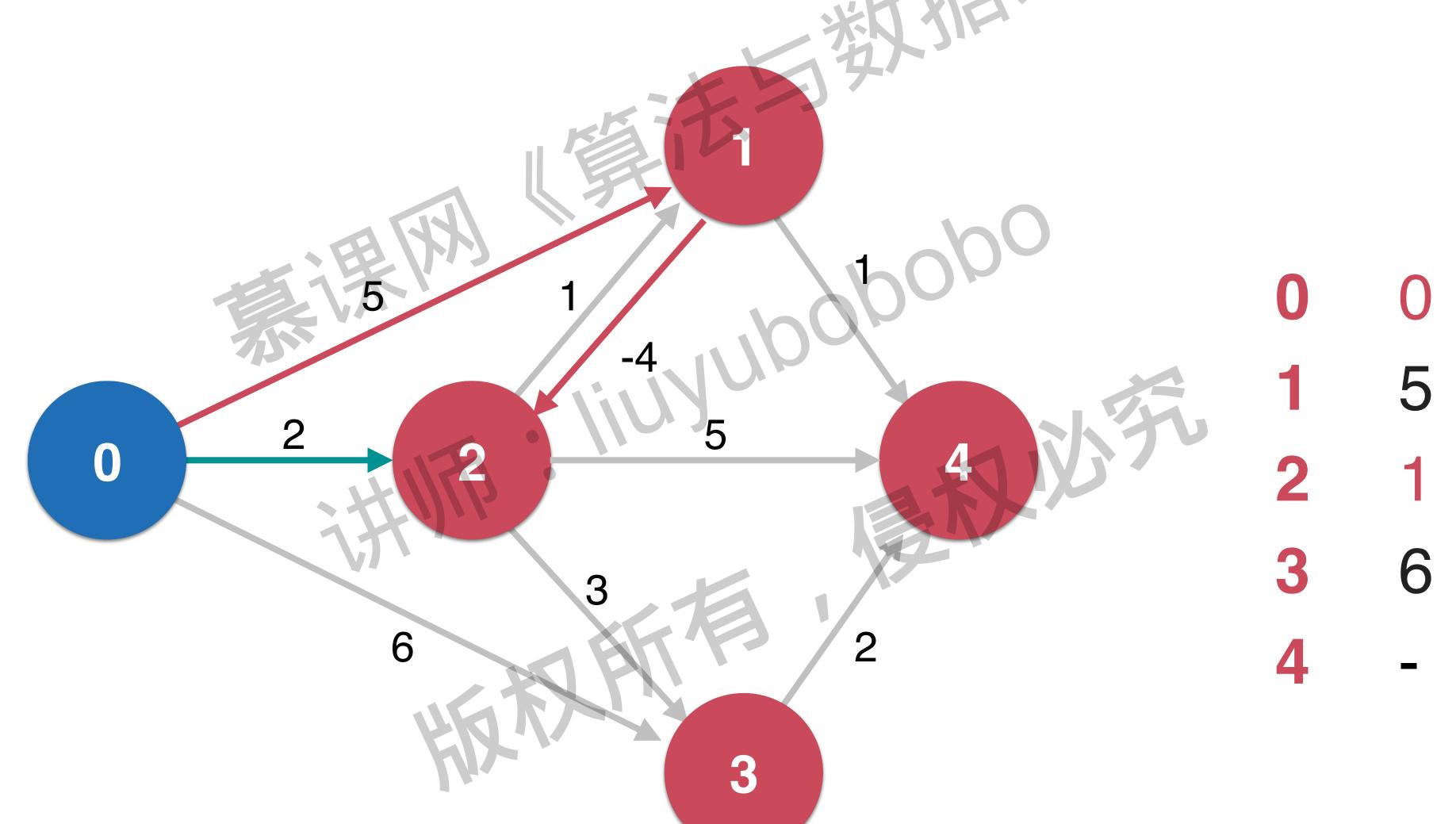
操作:实现dijkstra

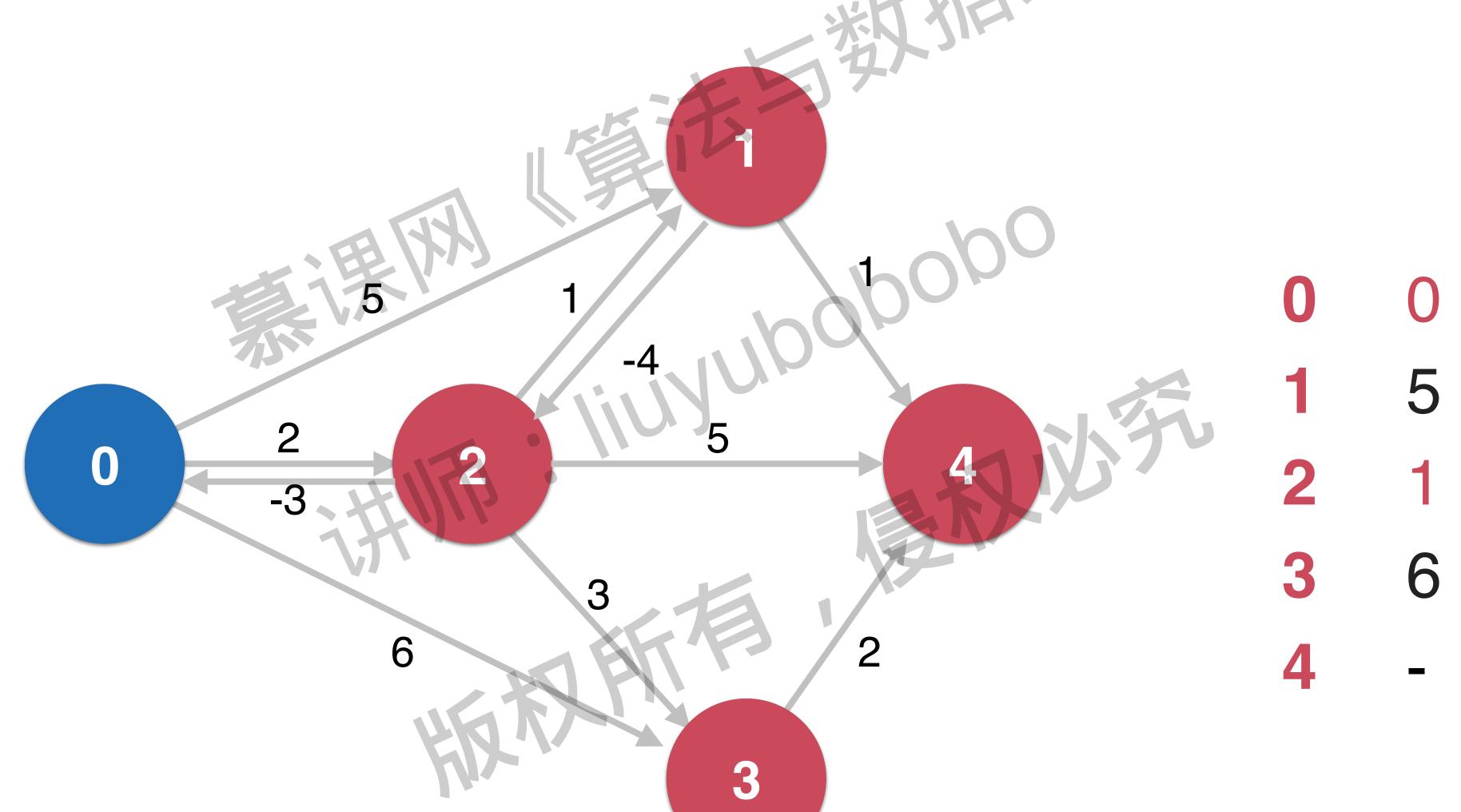
处理负极数

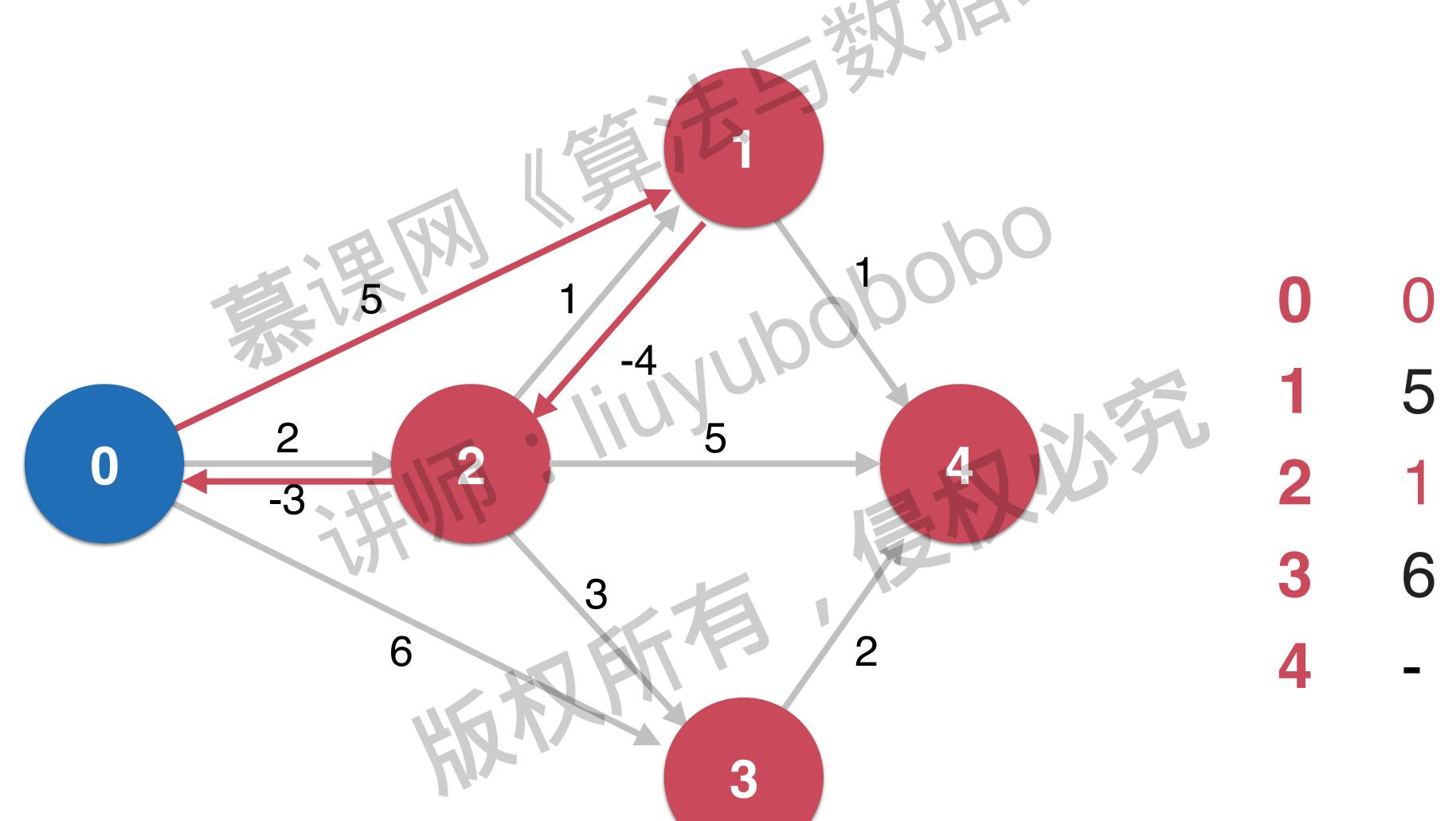


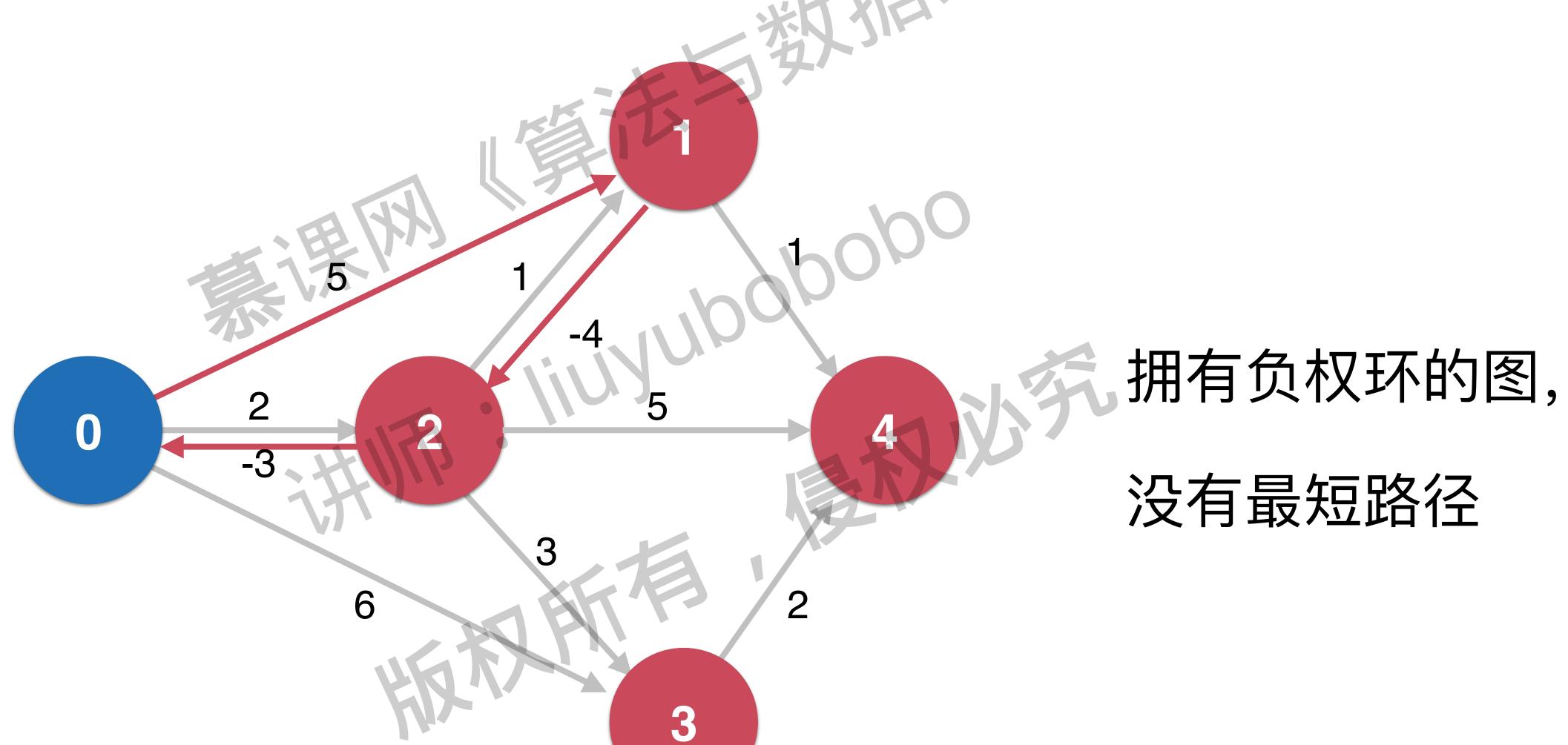
处理负极边

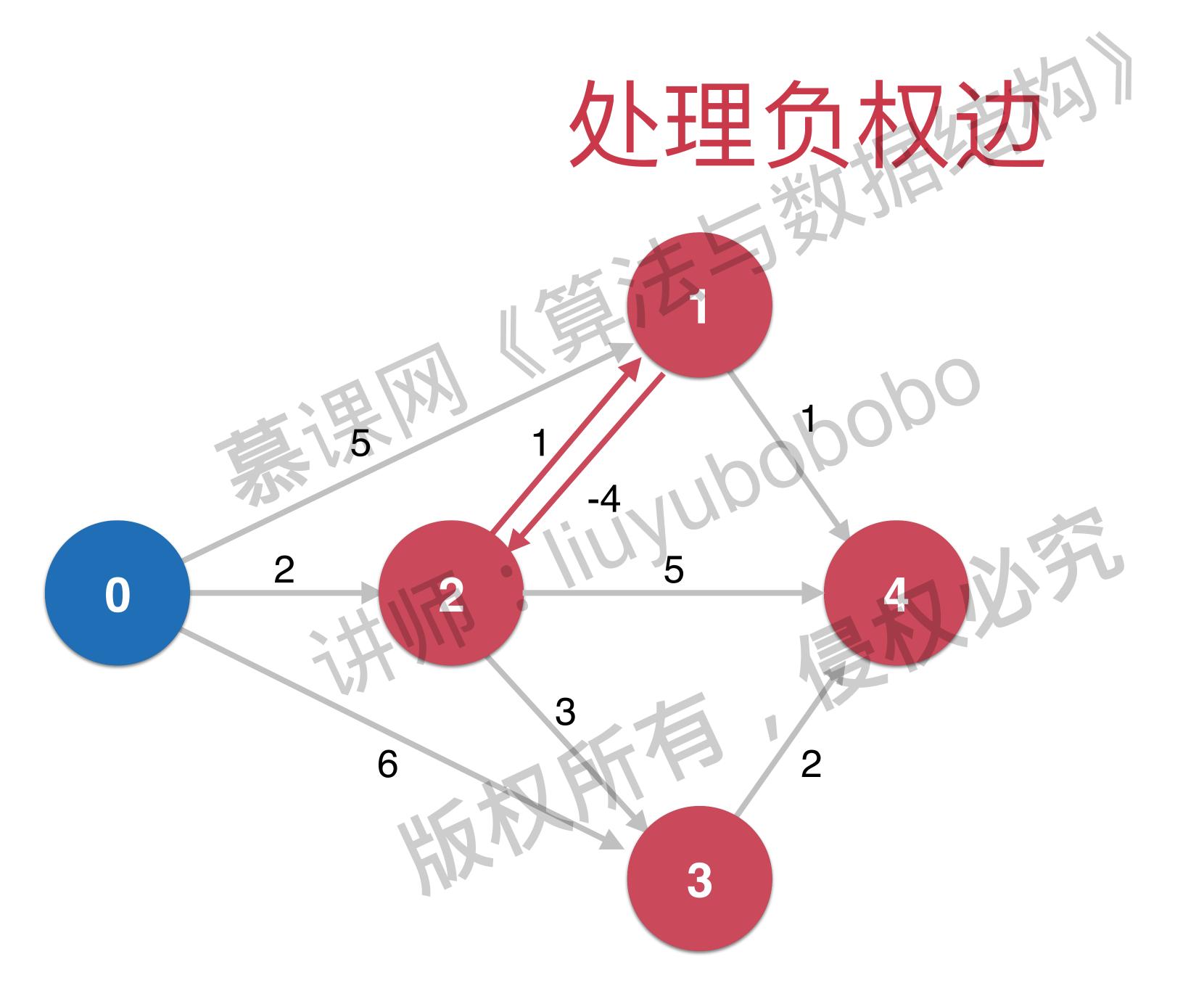














前提: 图中不能有负权环

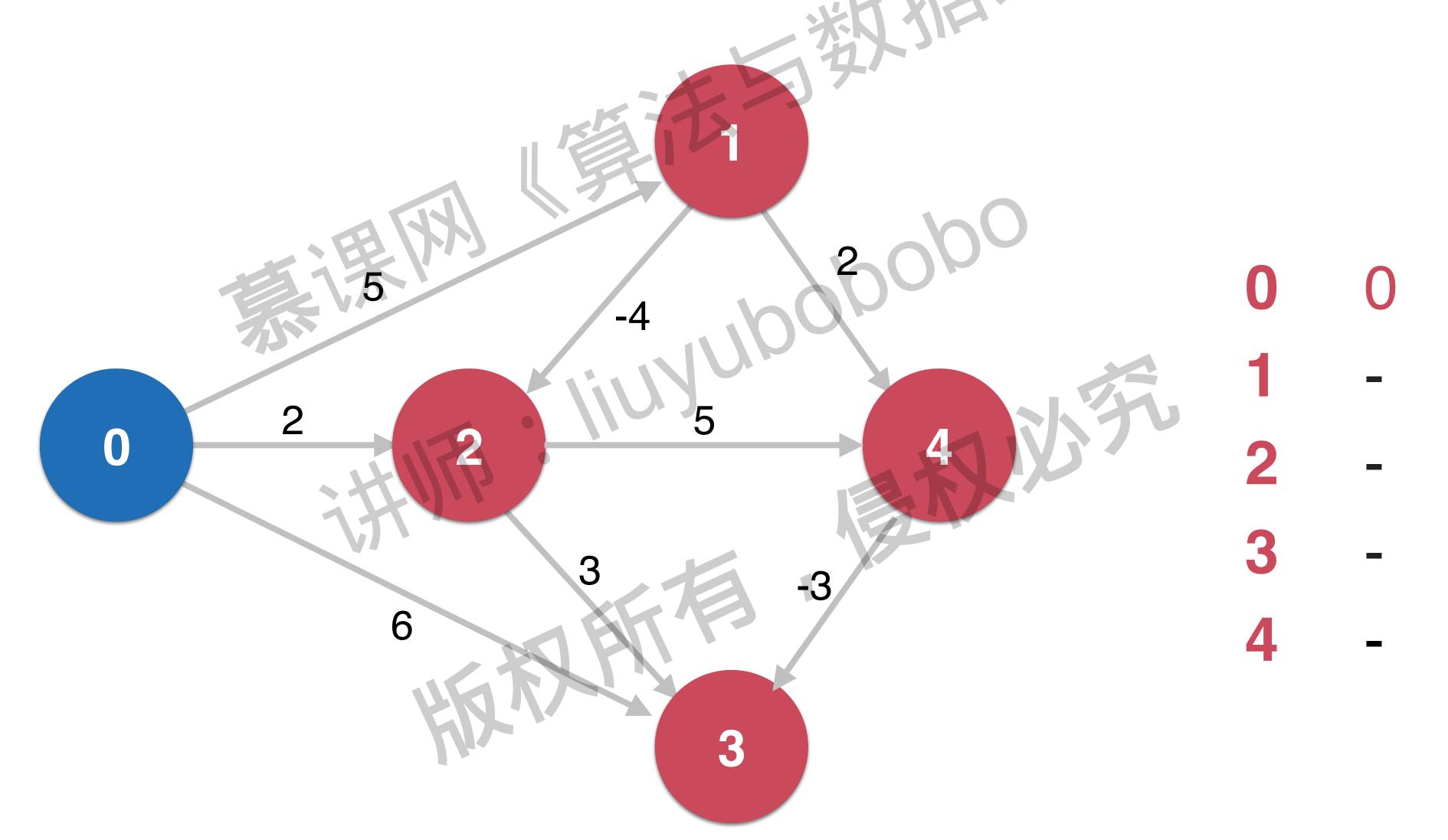
Bellman-Ford可以判断图中是否有负权环

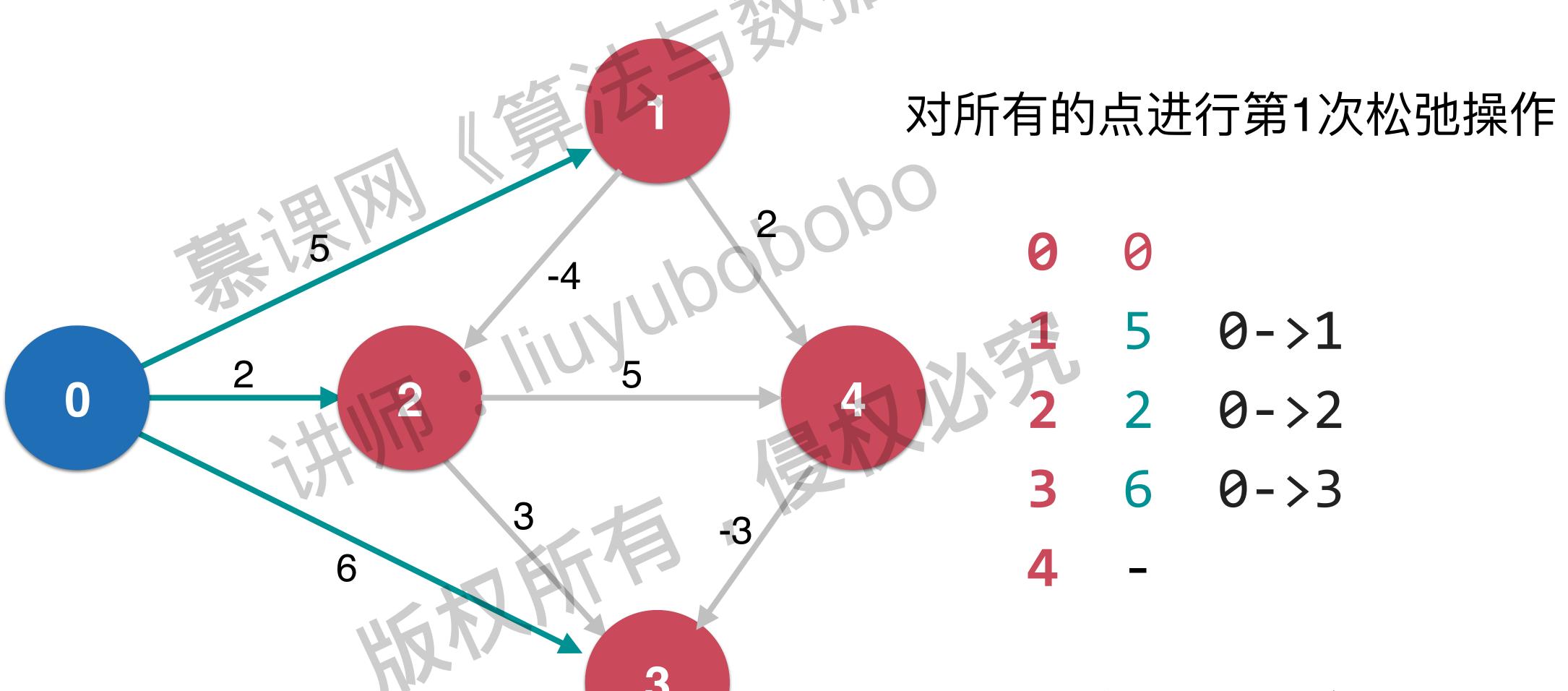
复杂度 O(EV)

如果一个图没有负权环,

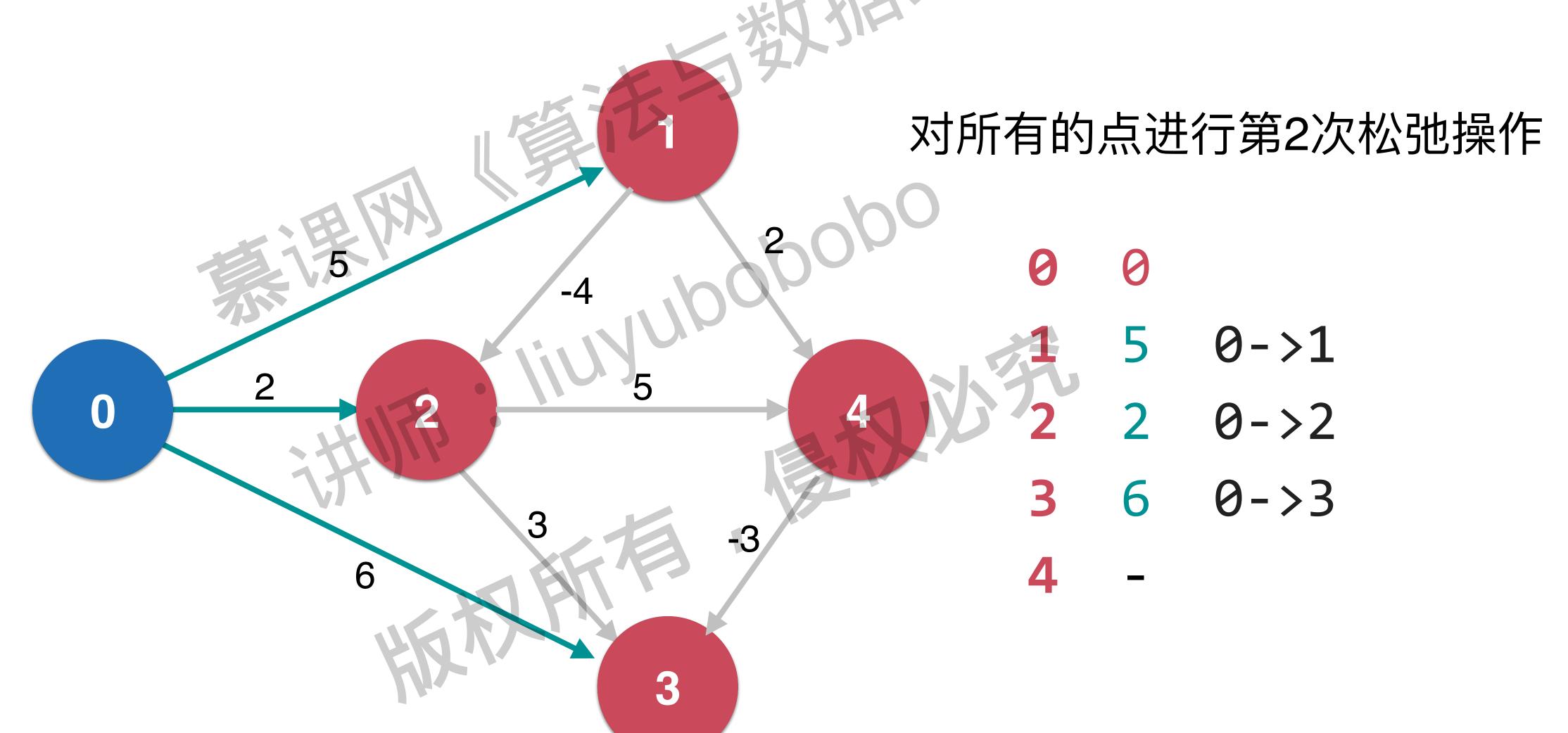
从一点到另外一点的最短路径,最多经过所有的V个顶线,有V-1条边

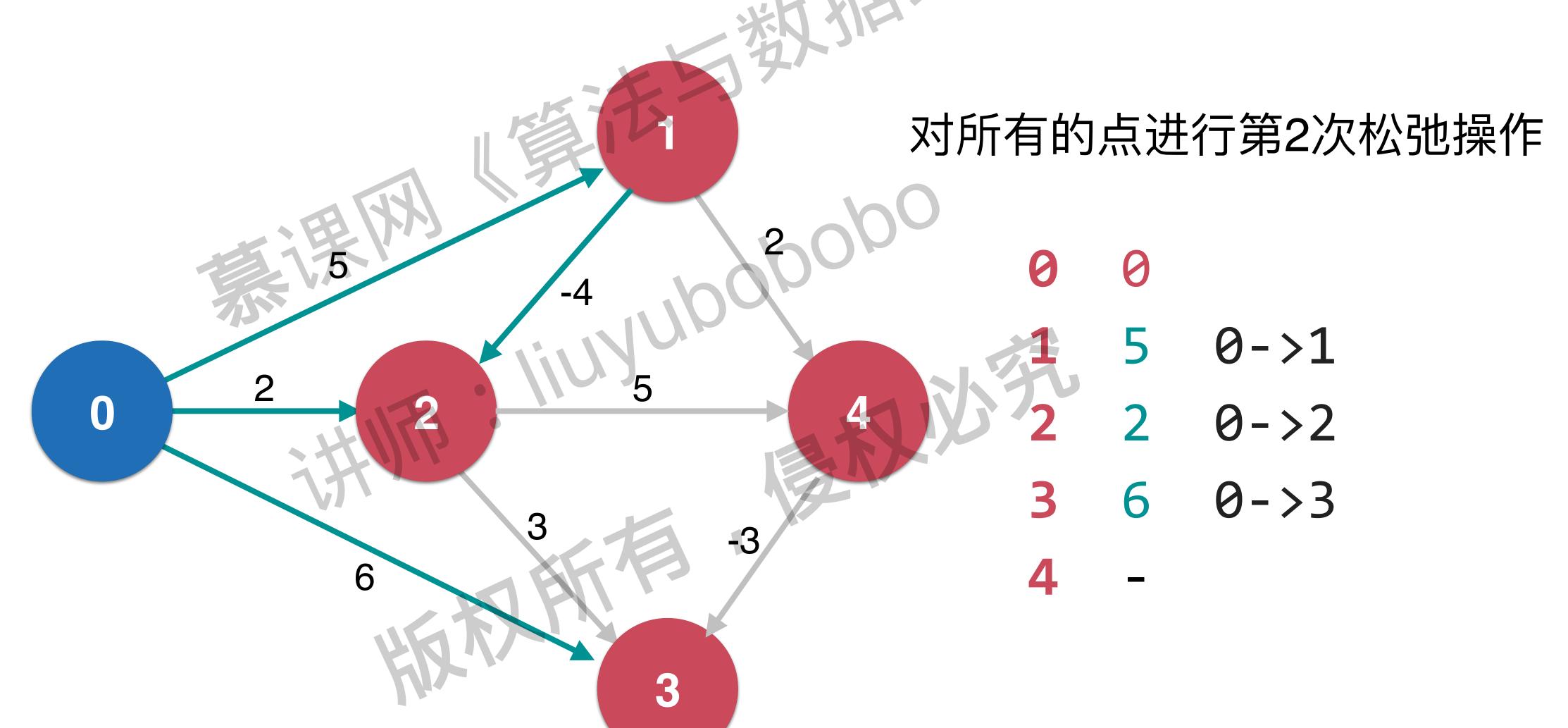
否则,存在顶点经过了两次,既存在负权环

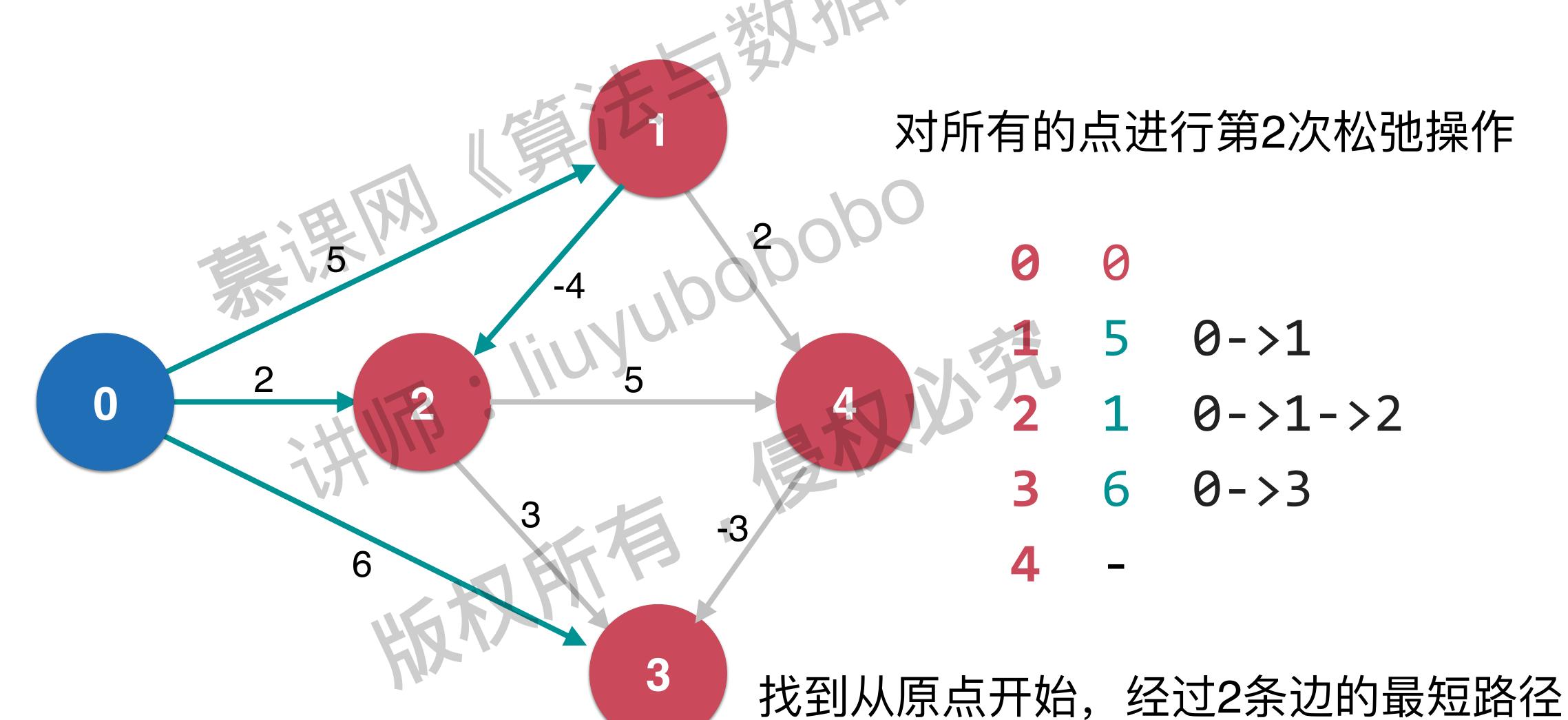




找到从原点开始,经过1条边的最短路径







对一个点的一次松弛操作,就是找到经过这个点的另外一条路径,多一条边,权值更小。

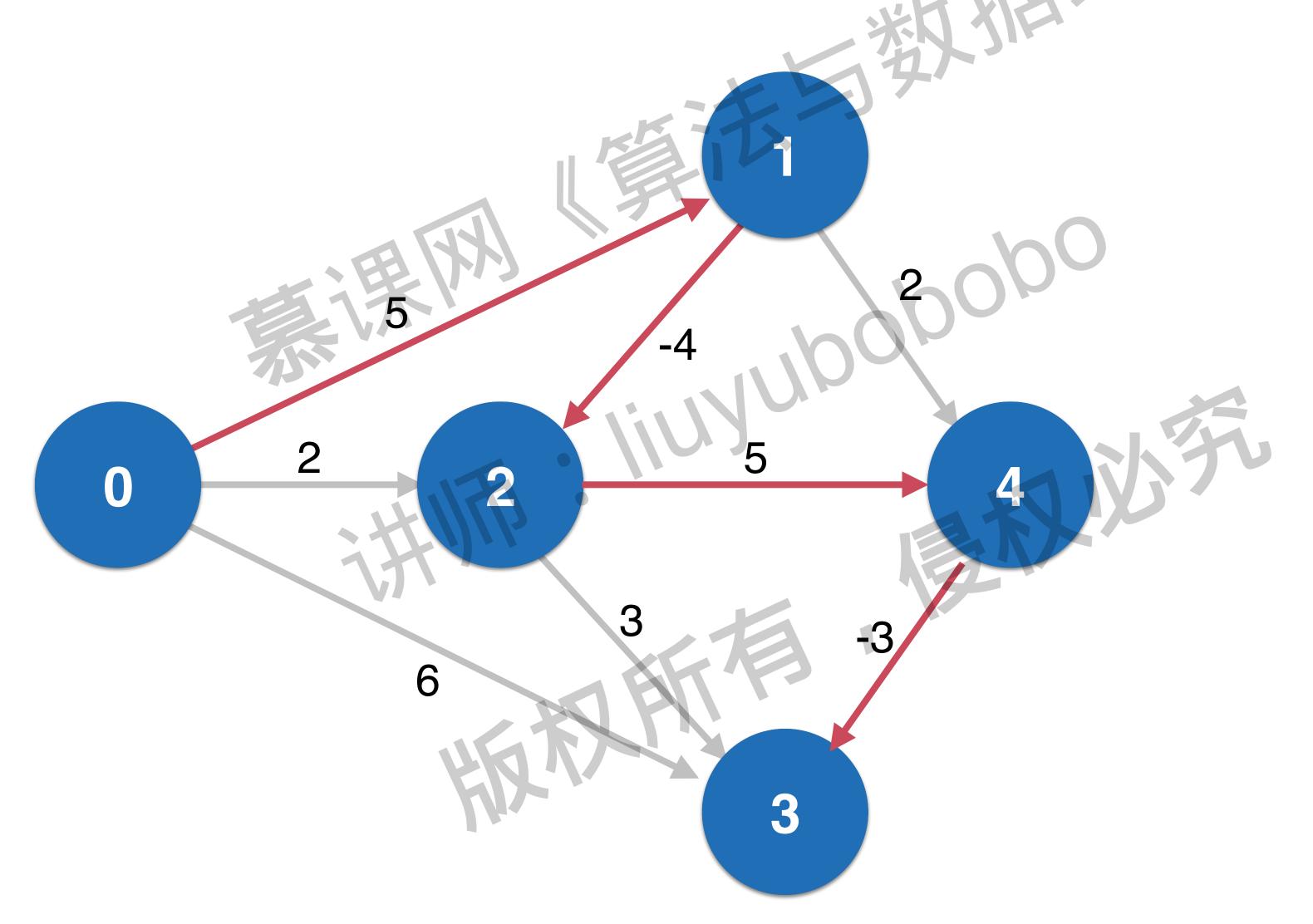
如果一个图没有负权环,从一点到另外一点的最短路径,最多经过所有的V个顶线,有V-1条边

对所有的点进行V-1次松弛操作

对所有的点进行V-1次松弛操作,理论上就找到了从源点到其他所有点的最短路径。

如果还可以继续松弛, 所说原图中有负权环。





更多和最短路径相关的问题

单源最短路径算法

具体实现,distTo[i] 初始化为"正无穷"

Bellman-Ford算法的优化

利用队列数据结构

queue-based bellman-ford算法

单源最短路径算法

dijkstra

无负权边

有向无向图均可

O(ElogV)

Bellman-Ford

无负权环

有向图

O(VE)

利用拓扑排序

有向无环图 DAG

有向图

O(V + E)

所有对最短路徑算法

Floyed算法,处理无负权环的图 O(V^3)

最长路径鎮洗

最长路径问题不能有正权环。

无权图的最长路径问题是指数级难度的。

对于有权图,不能使用Dijkstra求最长路径问题。

可以使用 Bellman-Ford算法。