## 浙江大学 2003 — 2004 学年第一学期期终考试

## 《操作系统实验》课程试卷

2	考试时间	]: <u>30</u> 分钟	开课学院:	_计算	机学院	专业:_		<del> </del>
ţ	姓名		学号		任课教	牧师		
题序			1_		=	四		总分
评分								
ì	平阅人							
PART II Operating System Lab Exam  1. Select the CORRECT and BEST answer for each of following questions and fill your answer in following blanks (30marks)  1. ( ) 2. ( ) 3. ( ) 4. ( ) 5. ( )  6. ( ) 7. ( ) 8. ( ) 9. ( ) 10. ( )  11. ( ) 12. ( ) 13. ( ) 14. ( ) 15. ( )								
1.	I/O devices are treated as in Linux. A. common files B. directory files C. index files D. special files							
2.	The A. B. C. D. is not a	command has the creates redirects completely valid of	a the	result new	test" as file output	"	cat called of	test" test test cat
3.	The A. B. ex C. D. none of	"!w" execute ecute the both The above	command the latest con	mmand a	will last that	w started and	with	command a w b
4.	present we all_in_one being used A. cat * > a B. cp * all_	all_in_one in_one all_in_one	(or pwd) into a	single n	ew file ca	ılled "all_iı	n_one"? (a	ssume that

5.	A. print the number of characters in all of the fi B. print the total number of lines in all the files ending in '.c' in the pr	les ending in '.c'. resent working directory.					
6.	If we want to install automatically a file system when power is up, which file can we modify? A. /etc/mtab B. /etc/fastboot C./etc/fstab D./etc/inetd.conf						
7.	A program is reading a standard input from a keyboard. If you wish to terminate input and tell the system the end of input which combined keys can you type?  A. Ctrl+Z B. Ctrl+W C. Ctrl+D D. Ctrl+V						
8.	3. Following messages are displayed when <i>ps</i> command is executed. If you very bash, which command can you use?  PID TTY TIME CMD  336 pts/1 00:00:00 login  337 pts/1 00:00:00 bash  356 pts/1 00:00:00 ps	vant to terminate process					
	A. kill bash B. kill pts/1 C. kill 337 D. kill !337						
9.	P. There are many kind of shell in Linux. Which is the least common used so A. xsh B. bash C. ksh D. csh	shell in following listed?					
10.	10. Which command is often used to build a Linux file system of A. mknod B. fdisk C. format D. mkfs	n a disk partition?					
11.	In order to make the file owner having <i>read</i> and <i>write</i> permissions while go can only <i>read</i> permission for the file, which of following octal number of chmod to change the file permission A.  Solution 566 B. C. 644 C. 655 D. 744						
12.	12. What is the default partition type A. vfat B. ext2 C. swap D.	in Linux?					
13.	13. To prevent to remove files by accident, which of options shall we up A.  B.  -f -R						

```
C.
                                                  -r
D. –i
```

14. Which of following commands can display the amount of disk space available on the file system A.

B.

df

C.

mount

D.

ln

15. What is the process number for following command?

```
$chmod 644 dir.txt&
```

[3] 164

A. 1 3 B. C. 164 D. 644

\_\_\_\_\_ Consider the following LINUX program, please show the possible output on display and in the file test.out. (10 marks)

```
#include <unistd.h>
#include <sys/types.h>
#include <sys/stat.h>
#include <fcntl.h>
int main()
  int fd;
  if( (fd=open("test.out", O_CREAT | O_RDWR, 0644)) == -1)
     printf("Can not open the output file test.out\n");
     exit(-1);
  printf("Begin to test\n");
  close(1);
  dup(fd);
  close(fd);
  printf("End of test\n");
  exit(0);
```

三、The following is the source code of Linux system call **sys\_mount()**, Please briefly describe its functionality and draw the flow chart of do\_add\_mount() or do\_mount() ( 5 marks )

```
struct file_system_type {
     const char *name;
     int fs_flags;
```

```
struct super_block *(*read_super) (struct super_block *, void *, int);
     struct module *owner;
     struct file_system_type * next;
     struct list_head fs_supers;
};
static int do_add_mount(struct nameidata *nd, char *type, int flags,
              int mnt_flags, char *name, void *data)
{
     struct vfsmount *mnt = do_kern_mount(type, flags, name, data);
     int err = PTR\_ERR(mnt);
     if (IS_ERR(mnt))
         goto out;
     down(&mount_sem);
     /* Something was mounted here while we slept */
     while(d_mountpoint(nd->dentry) && follow_down(&nd->mnt, &nd->dentry))
         ;
     err = -EINVAL;
     if (!check_mnt(nd->mnt))
          goto unlock;
     /* Refuse the same filesystem on the same mount point */
     err = -EBUSY;
     if (nd->mnt->mnt_sb == mnt->mnt_sb && nd->mnt->mnt_root == nd->dentry)
          goto unlock;
     mnt->mnt_flags = mnt_flags;
     err = graft_tree(mnt, nd);
unlock:
     up(&mount_sem);
     mntput(mnt);
out:
     return err;
* Flags is a 32-bit value that allows up to 31 non-fs dependent flags to
* be given to the mount() call (ie: read-only, no-dev, no-suid etc).
* data is a (void *) that can point to any structure up to
★ PAGE_SIZE-1 bytes, which can contain arbitrary fs-dependent
* information (or be NULL).
```

```
long do_mount(char * dev_name, char * dir_name, char *type_page,
           unsigned long flags, void *data_page)
{
     struct nameidata nd;
     int retval = 0;
     int mnt_flags = 0;
     /* Discard magic */
     if ((flags & MS_MGC_MSK) == MS_MGC_VAL)
         flags \&= \sim MS\_MGC\_MSK;
     /★ Basic sanity checks ★/
     if (!dir_name | | !*dir_name | | !memchr(dir_name, 0, PAGE_SIZE))
         return -EINVAL;
     if (dev_name && !memchr(dev_name, 0, PAGE_SIZE))
         return -EINVAL;
     /* Separate the per-mountpoint flags */
     if (flags & MS_NOSUID)
         mnt_flags |= MNT_NOSUID;
     if (flags & MS_NODEV)
         mnt_flags |= MNT_NODEV;
     if (flags & MS_NOEXEC)
         mnt_flags |= MNT_NOEXEC;
     flags \&= \sim (MS_NOSUID | MS_NOEXEC | MS_NODEV);
     /* ... and get the mountpoint */
     if (path_init(dir_name, LOOKUP_FOLLOW | LOOKUP_POSITIVE, &nd))
         retval = path_walk(dir_name, &nd);
     if (retval)
         return retval;
     if (flags & MS_REMOUNT)
         retval = do_remount(&nd, flags & ~MS_REMOUNT, mnt_flags,
                     data_page);
     else if (flags & MS_BIND)
         retval = do_loopback(&nd, dev_name, flags & MS_REC);
     else if (flags & MS_MOVE)
         retval = do_move_mount(&nd, dev_name);
     else
         retval = do_add_mount(&nd, type_page, flags, mnt_flags,
```

```
dev_name, data_page);
     path_release(&nd);
     return retval;
}
asmlinkage long sys_mount(char * dev_name, char * dir_name, char * type,
                unsigned long flags, void * data)
{
    int retval;
     unsigned long data_page;
     unsigned long type_page;
     unsigned long dev_page;
     char *dir_page;
     retval = copy_mount_options (type, &type_page);
    if (retval < 0)
          return retval;
     dir_page = getname(dir_name);
     retval = PTR_ERR(dir_page);
    if (IS_ERR(dir_page))
          goto out1;
     retval = copy_mount_options (dev_name, &dev_page);
     if (retval < 0)
          goto out2;
     retval = copy_mount_options (data, &data_page);
    if (retval < 0)
          goto out3;
    lock_kernel();
     retval = do_mount((char*)dev_page, dir_page, (char*)type_page,
                flags, (void*)data_page);
     unlock_kernel();
     free_page(data_page);
out3:
     free_page(dev_page);
out2:
     putname(dir_page);
out1:
     free_page(type_page);
     return retval;
```

```
四、The following is part of the source code for Linux virtual memory allocation function vmalloc().
Please briefly describe its functionality and draw the flow chart depending on given source
code (5marks)
/*
* vmalloc() function
struct vm struct {
      unsigned long flags; /* virtual memory block state flag */
                        /* pointer to the virtual memory block */
      void * addr:
      unsigned long size;
                            /* block size */
      struct vm struct * next; /* pointer to the next block */
static struct vm struct * vmlst = NULL;
rwlock_t vmlist_lock = RW_LOCK_UNLOCKED;
struct vm_struct * vmlist;
inline int vmalloc_area_pages (unsigned long address, unsigned long size,
                    int gfp_mask, pgprot_t prot)
{
    pgd_t * dir;
    unsigned long end = address + size;
    int ret;
    dir = pgd_offset_k(address); /* to find an entry in a kernel page-table-directory */
    spin_lock(&init_mm.page_table_lock);
    do {
         pmd_t *pmd;
         pmd = pmd_alloc(&init_mm, dir, address);
         ret = -ENOMEM;
         if (!pmd)
              break;
         ret = -ENOMEM;
         if (alloc_area_pmd(pmd, address, end - address, gfp_mask, prot))
              break;
         address = (address + PGDIR_SIZE) & PGDIR_MASK;
         dir++;
         ret = 0;
    } while (address && (address < end));
```

}

```
spin_unlock(&init_mm.page_table_lock);
     flush_cache_all();
     return ret;
}
struct vm_struct * get_vm_area(unsigned long size, unsigned long flags)
{
     unsigned long addr;
     struct vm_struct **p, *tmp, *area;
     area = (struct vm_struct *) kmalloc(sizeof(*area), GFP_KERNEL);
     if (!area)
          return NULL;
     size += PAGE_SIZE;
     addr = VMALLOC_START;
     write_lock(&vmlist_lock);/* write spinlock */
     for (p = \&vmlist; (tmp = *p); p = \&tmp->next) {
          if ((size + addr) < addr)
               goto out;
          if (size + addr <= (unsigned long) tmp->addr)
          addr = tmp->size + (unsigned long) tmp->addr;
          if (addr > VMALLOC_END-size)
               goto out;
     }
     area->flags = flags;
     area -> addr = (void *) addr;
     area->size = size;
     area -> next = \star p;
     \star p = area;
     write_unlock(&vmlist_lock); return area;
out:
     write_unlock(&vmlist_lock);
     kfree(area);
     return NULL;
}
void * __vmalloc (unsigned long size, int gfp_mask, pgprot_t prot)
{
     void * addr;
    struct vm_struct *area;
     size = PAGE_ALIGN(size);
```

```
if (!size | | (size >> PAGE_SHIFT) > num_physpages) {
         BUG();
         return NULL;
    area = get_vm_area(size, VM_ALLOC);
    if (!area)
         return NULL;
    addr = area->addr;
    if (vmalloc_area_pages(VMALLOC_VMADDR(addr), size, gfp_mask, prot)) {
         vfree(addr);
         return NULL;
    }
    return addr;
}
    Allocate any pages
static inline void * vmalloc (unsigned long size)
    return __vmalloc(size, GFP_KERNEL | __GFP_HIGHMEM, PAGE_KERNEL);
}
```