

Compilation Principle Homework 2

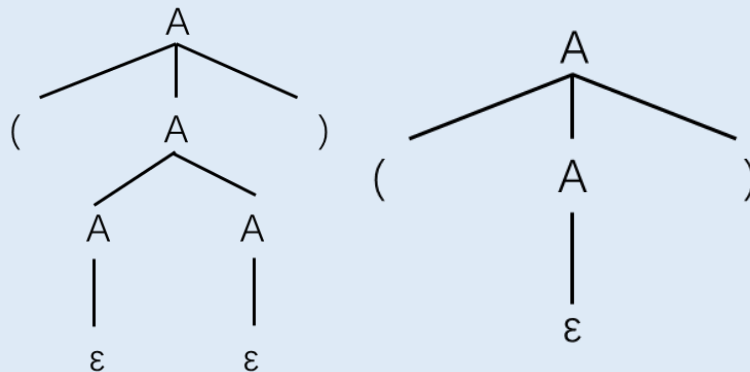
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3.2 Given the grammar $A \rightarrow AA \mid (A) \mid \varepsilon$,

- Describe the language it generates.
- Show it is ambiguous.

Answer:

- It generates a sequence of symmetric parentheses, you can find a pair of parentheses for each one, and this language contains the empty string.
- Ambiguous can be like this:



The string "()" can generate two different parse tree, so this language is ambiguous.

3.3 Given the grammar:

$\text{exp} \rightarrow \text{exp addop term} \mid \text{term}$
 $\text{addop} \rightarrow + \mid -$
 $\text{term} \rightarrow \text{term mulop factor} \mid \text{factor}$
 $\text{mulop} \rightarrow *$
 $\text{factor} \rightarrow (\text{exp}) \mid \text{factor}$

Write down leftmost derivation, parse trees, and abstract syntax trees for the following expressions:

(a) $3+4*5-6$

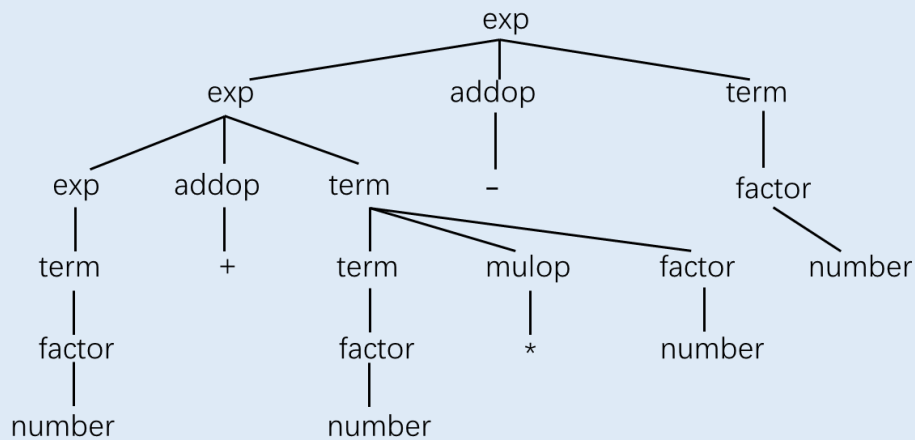
Answer:

Leftmost derivation:

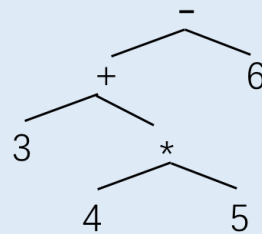
```
exp => exp addop term
    => exp addop term addop term
    => term addop term addop term
    => factor addop term addop term
    => number addop term addop term
    => number + term addop term
    => number + term mulop factor addop term
    => number + factor mulop factor addop term
    => number + number mulop factor addop term
    => number + number * factor addop term
    => number + number * number addop term
```

$\Rightarrow \text{number} + \text{number} * \text{number} - \text{term}$
 $\Rightarrow \text{number} + \text{number} * \text{number} - \text{factor}$
 $\Rightarrow \text{number} + \text{number} * \text{number} - \text{number}$

Parse Tree:



Abstract Syntax Tree:



3.4 The following grammar generates all regular expressions over the alphabet of letters (we have to use quotes to surround operators, since the vertical bar is an operator as well as a metasympol):

$\text{rexp} \rightarrow \text{rexp} \text{ " | " rexp}$
 $\quad \quad \quad | \text{ rexp rexp}$
 $\quad \quad \quad | \text{ rexp " * "}$
 $\quad \quad \quad | \text{ "(" rexp ")"}$
 $\quad \quad \quad | \text{ letter}$

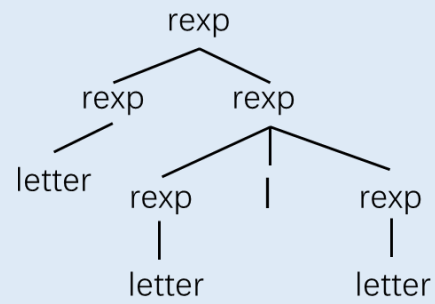
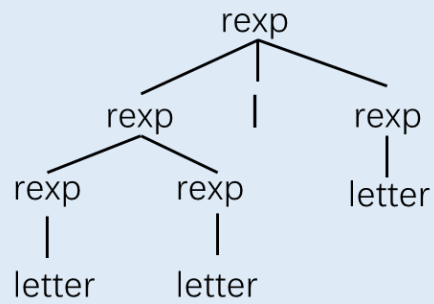
- Give a derivation for the regular expression $(ab | a)^*$ using this grammar.
- Show that this grammar is ambiguous.
- Rewrite this grammar to establish the correct precedences for the operations (see chapter 2).
- What associativity does your answer in part (c) give to the binary operations? Why?

Answer:

a.

$\text{rexp} \Rightarrow \text{rexp} \text{ " * "}$
 $\Rightarrow \text{ "(" rexp ") " " * " }$
 $\Rightarrow \text{ "(" rexp " | " rexp ") " " * " }$
 $\Rightarrow \text{ "(" rexp rexp " | " rexp ") " " * " }$
 $\Rightarrow \text{ "(" letter letter " | " letter ") " " * " }$

b. Ambiguous can be like this:



The string "ab | c" can generate two different parse tree, so this language is ambiguous.

- c.
- rexp \rightarrow rexp "|" rexp_conc | rexp_conc
 - rexp_conc \rightarrow rexp_conc rexp_clos | rexp_clos
 - rexp_clos \rightarrow rexp_clos "*" | rexp_pare
 - rexp_pare \rightarrow "(" rexp ")" | letter

d. Because I use a left recursive in the production, so its operators associate on the left. If not, string like "abc" can have 2 kinds parse trees. But in fact, it's a inessential ambiguity, I mean that $(ab)c = a(bc)$, or $(a|b)|c = a|(b|c)$, their syntax trees are still distinct, but semantic value are the same.