

## 练习6 (第8、9章)

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#1 Points possible: 3

\_\_\_\_\_ memory allocation scheme may produce external fragmentation.

- ☐ Demand
  - ☐ system halts
  - ☒ Multiple-partition
  - ☐ None of above
- 

#2 Points possible: 3

A demand paging system adopts the LRU page replacement algorithm. Consider a reference string 1 8 1 7 8 2 7 2 1 8 3 8 2 1 3 1 7 1 3 7. The total number of page faults given 4 initially empty page frames is \_\_\_\_\_.

- ☐ 4
  - ☐ 5
  - ☒ 6
  - ☐ 7
- 

#3 Points possible: 3

After a page fault handled, \_\_\_\_\_ should be executed.

- ☐ the instruction just before interruption
  - ☒ the instruction caused interruption
  - ☐ the instruction just after interruption
  - ☐ The first instruction of this process
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#4 Points possible: 3

Assume that the probability of page fault is 0.1%, memory access time is 100ns, and the average page fault service time is 25 ms, then the effective access time is \_\_\_\_.

- ☐ 125 $\mu$ s
- ☐ 115ms
- ☒ 25 $\mu$ s
- ☐ 25ms

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#5 Points possible: 3

Considering a system, which uses virtual memory. At what point can address binding be done?

- ☐ compile time
- ☐ load time
- ☒ execution time
- ☐ can be any of the above

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#6 Points possible: 3

Consider a paging system that maps logical address space of 8 pages with 1024 bytes each page to a physical memory of 32 frames, the logical address is of \_\_\_ and the physical address is of \_\_\_.

- ☐ 10 bits, 5 bits
- ☐ 3 bits, 15 bits
- ☐ 13 bits ,5 bits
- ☒ 13 bits, 15 bits

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#7 Points possible: 3

Dynamic relocation relies on \_\_\_\_.

- ☒ a relocation register

- ☐ object code
- ☐ dynamic link libraries
- ☐ relocation program

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#8 Points possible: 3

In a paging memory management system, there is a page table as

Page No.	Frame No.
0	2
1	1
2	6
3	3
4	7

following:

If the page size is 4KB, then paging address hardware will convert logical address 10 into physical address \_\_\_\_.

- ☒ 8202
- ☐ 4106
- ☐ 2058
- ☐ 1034

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#9 Points possible: 3

assume that a task is divided into 4 equal-sized segments, and that the system builds an 8-entry page table for each segment. Therefore, the system has a combination of segmentation and paging. Assume also that the page size is 2Kbytes.

What is the maximum size of each segment?

- ☐ 2Kbytes
- ☐ 4Kbytes

- ☐ 8Kbytes
- ☒ 16Kbytes

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#10 Points possible: 3

assume that a task is divided into 4 equal-sized segments, and that the system builds an 8-entry page table for each segment. Therefore, the system has a combination of segmentation and paging. Assume also that the page size is 2Kbytes.

What is the maximum logical address space for the task?

- ☐ 8Kbytes
- ☐ 16Kbytes
- ☐ 32Kbytes
- ☒ 64Kbytes

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#11 Points possible: 3

Implementing LRU precisely in an OS is expensive, so practical implementations often use an approximation called .

- ☐ MRU
- ☐ MFU
- ☐ LFU
- ☒ NRU

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#12 Points possible: 3

The second-chance (clock) algorithm is an efficient approximation technique for \_\_\_\_\_.

- ☒ LRU page replacement
  - ☐ LFU page replacement
  - ☐ benchmarking file system performance
  - ☐ benchmarking raw disk I/O performance
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#13 Points possible: 3

Which of the following memory management is not suitable for a multi-programming environment?

- ☒ single contiguous memory allocation
  - ☐ fix-sized partitions allocation
  - ☐ variable-sized partitions allocation
  - ☐ segmentation with paging
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#14 Points possible: 3

Suppose that the TLB has a 90% hit ratio, if the times for TLB searching is 20 nanoseconds, access memory is 100 nanoseconds, what is the effective memory-access time?

- ☐ 120 nanoseconds
  - ☒ 130 nanoseconds
  - ☐ 140 nanoseconds
  - ☐ 220 nanoseconds
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#15 Points possible: 3

\_\_\_\_\_ may have not internal fragmentation.

- ☐ Paging memory management
  - ☒ Segmentation memory management
  - ☐ Fix-sized partition memory management
  - ☐ Segmentation with paging memory management
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#16 Points possible: 3

Which of the following page replacement algorithms may produce Belady's anomaly?

- ☒ FIFO
- ☐ LRU

- ☐ OPT
  - ☐ None of the above
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#17 Points possible: 3

The fundamental basis for virtual memory management is \_\_\_\_ .

- ☐ virtuality
  - ☒ locality
  - ☐ globality
  - ☐ dynamics
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#18 Points possible: 3

To fetch a data from main memory in a demand paging system requires \_\_\_\_ accesses to the physical memory.

- ☐ 1
  - ☒ 2
  - ☐ 3
  - ☐ 4
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#19 Points possible: 3

With demand paging, \_\_\_\_\_ have worst system performance.

- ☐ stacks
  - ☐ lists
  - ☒ hash tables
  - ☐ arrays
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#20 Points possible: 3

In a demanding paging system, the size of a page is 4KB. A process access the logical address 12345 (0x3039) will \_\_\_\_ if the page table is as the following:

Page #	Frame #	Validity
0	3	V
1	4	V
2	2	V
3	—	I

- ☐ access physical address  $4 \times 4096 + 57$
- ☐ access physical address  $3 \times 4096 + 57$
- ☐ access physical address  $2 \times 4096 + 57$
- ☒ cause a page-fault interrupt

#21 Points possible: 2

为使虚存系统有效地发挥其预期的作用，所运行的程序应具有的特性是——。

- ☐ 该程序不应含有过多的I/O操作
- ☐ 该程序的大小不应超过实际的内存容量
- ☒ 该程序应具有较好的局部性 ( Locality )
- ☐ 该程序的指令相关不应过多

#22 Points possible: 2

总体上说，请求分页(demand-paging)是个很好的虚拟内存管理策略。但是，有些程序设计技术并不适合于这种环境。例如，\_\_\_\_\_。

- ☐ 堆栈
- ☐ 线性搜索
- ☐ 矢量运算
- ☒ 二分法搜索

#23 Points possible: 2

考虑页面置换算法，系统有m个页帧供调度，初始时全空；引用串长度为p，包含了n个不同的页面，无论用什么缺页算法，缺页次数不会少于\_\_\_\_\_。

- ☐ m
- ☐ p
- ☒ n
- ☐  $\min(m,n)$

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#24 Points possible: 2

首次适应算法的空闲区是\_\_\_\_\_。

- ☒ 按地址递增顺序连在一起
- ☐ 始端指针表指向最大空闲区
- ☐ 按大小递增顺序连在一起
- ☐ 寻找从最大空闲区开始

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#25 Points possible: 2

下述\_\_\_\_\_页淘汰算法会产生Belady现象。

- ☒ 先进先出
- ☐ 最近最少使用
- ☐ 最不经常使用
- ☐ 最佳页面置换

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#26 Points possible: 2

在虚拟分页存储管理系统中，若进程访问的页面不在主存，且主存中没有可用的空闲帧时，系统正确的处理顺序为\_\_\_\_\_。

- ☐ 决定淘汰页→页面调出→缺页中断→页面调入
  - ☐ 决定淘汰页→页面调入→缺页中断→页面调出
  - ☒ 缺页中断→决定淘汰页→页面调出→页面调入
  - ☐ 缺页中断→决定淘汰页→页面调入→页面调出
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