

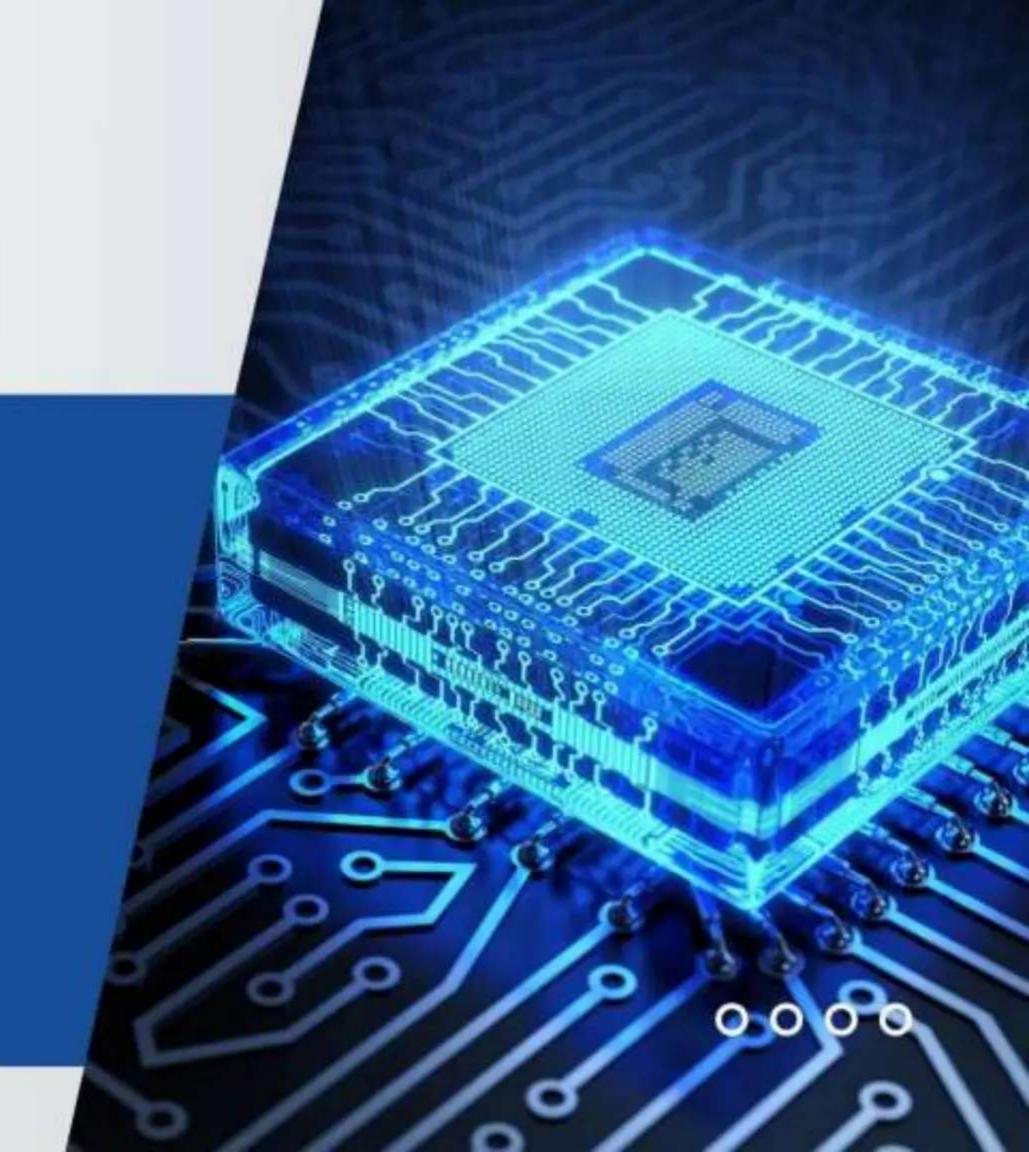


MERL DSU

RISC-V Pipeline Core

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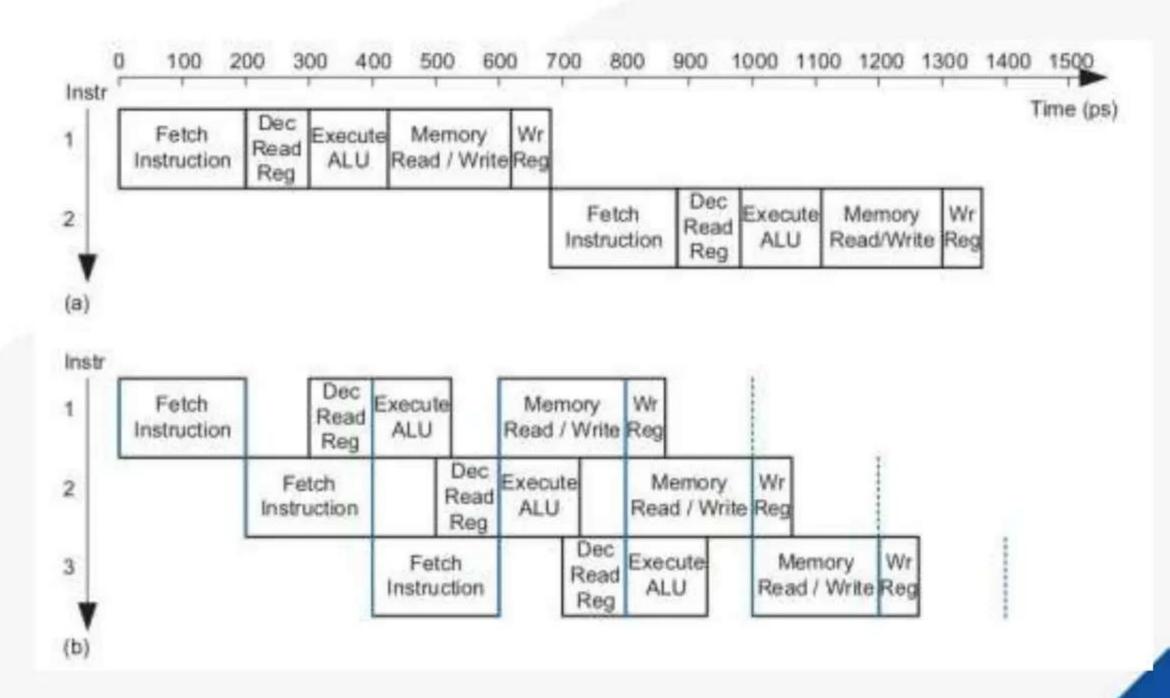
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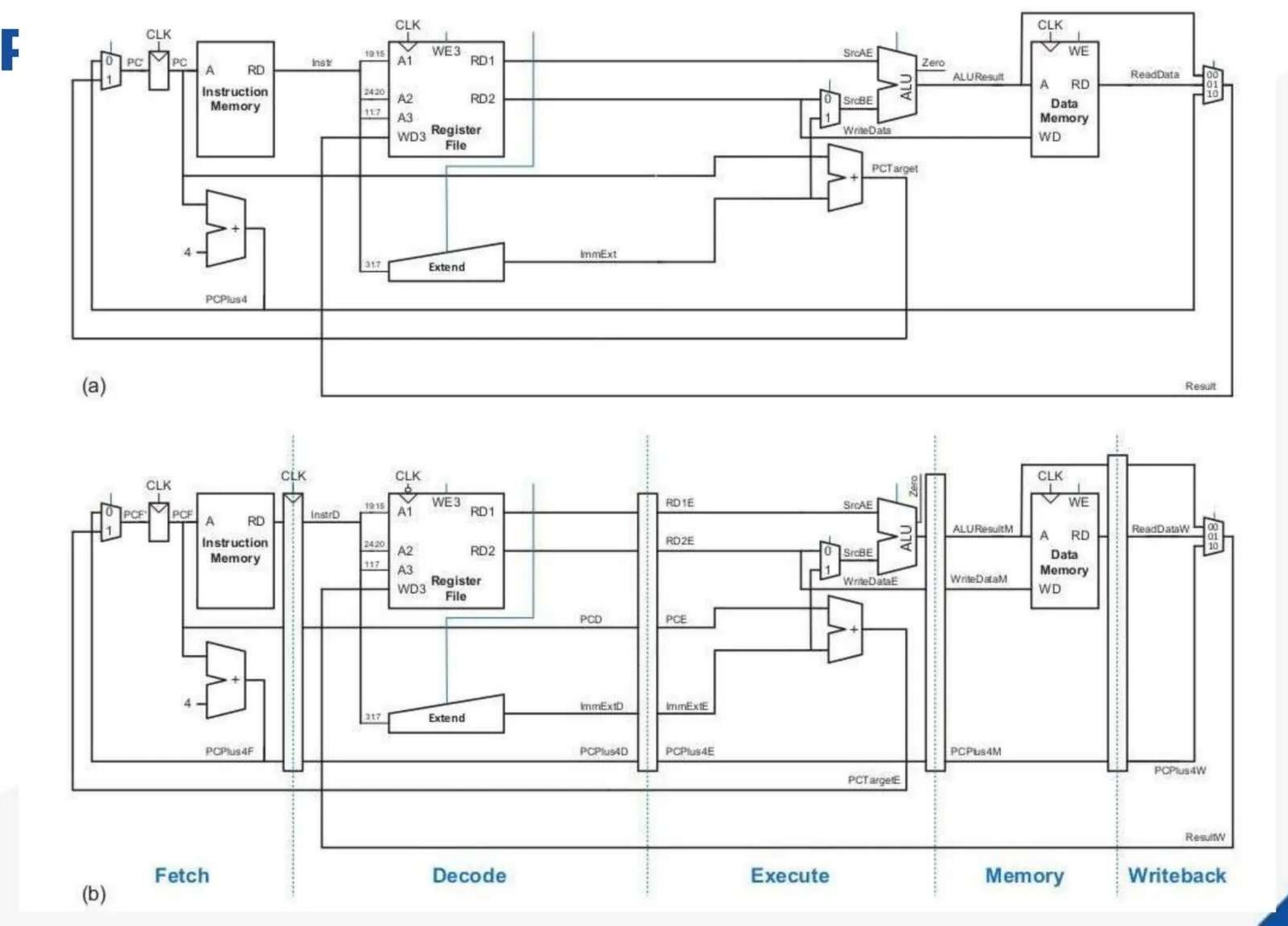


Overview of RISC-V Pipeline Architecture

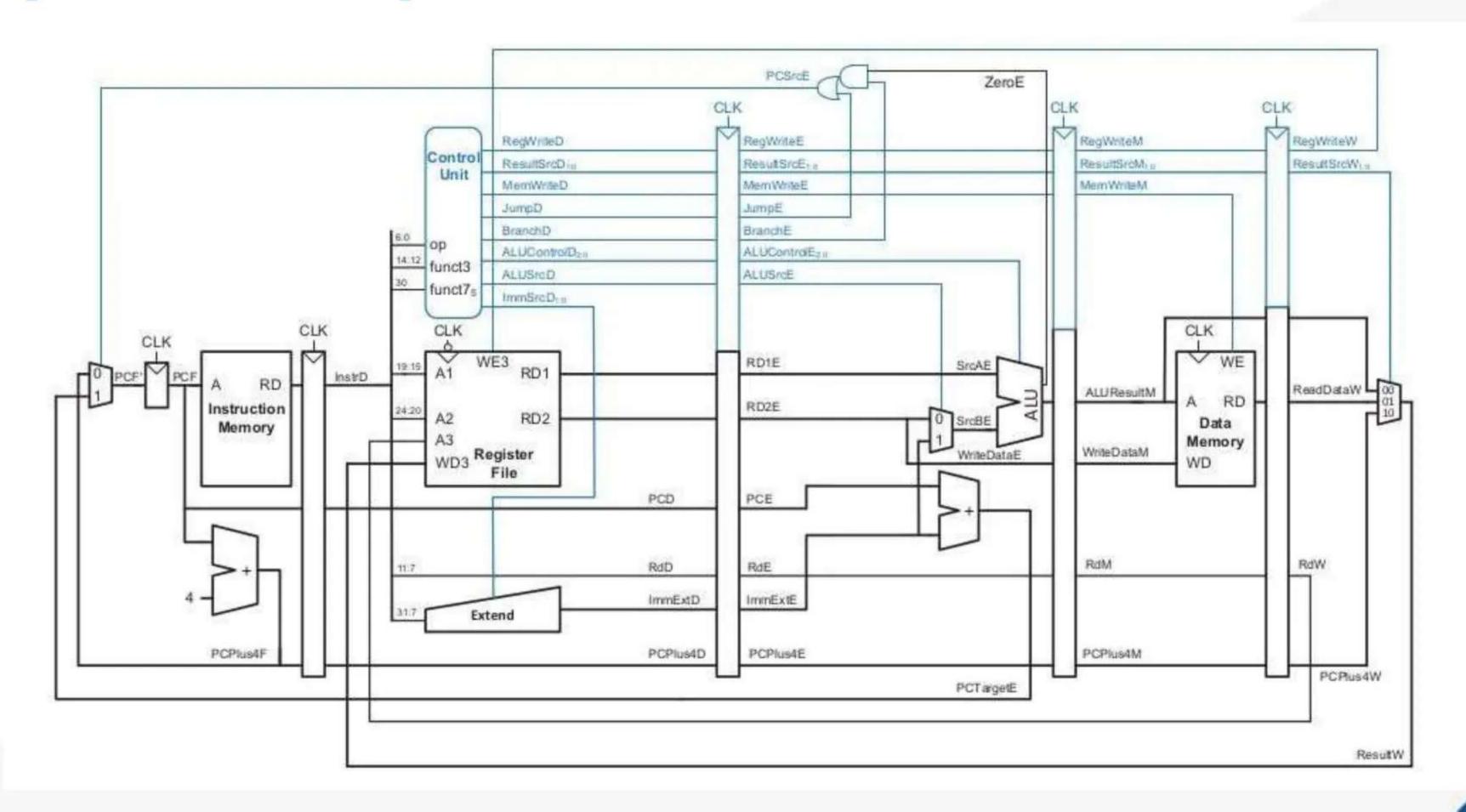
Pipelining

We design a pipelined processor by subdividing the single-cycle processor into five pipeline stages. Thus, five instructions can execute simultaneously, one in each stage. Because each stage has only one-fifth of the entire logic, the clock frequency is approximately five times faster.





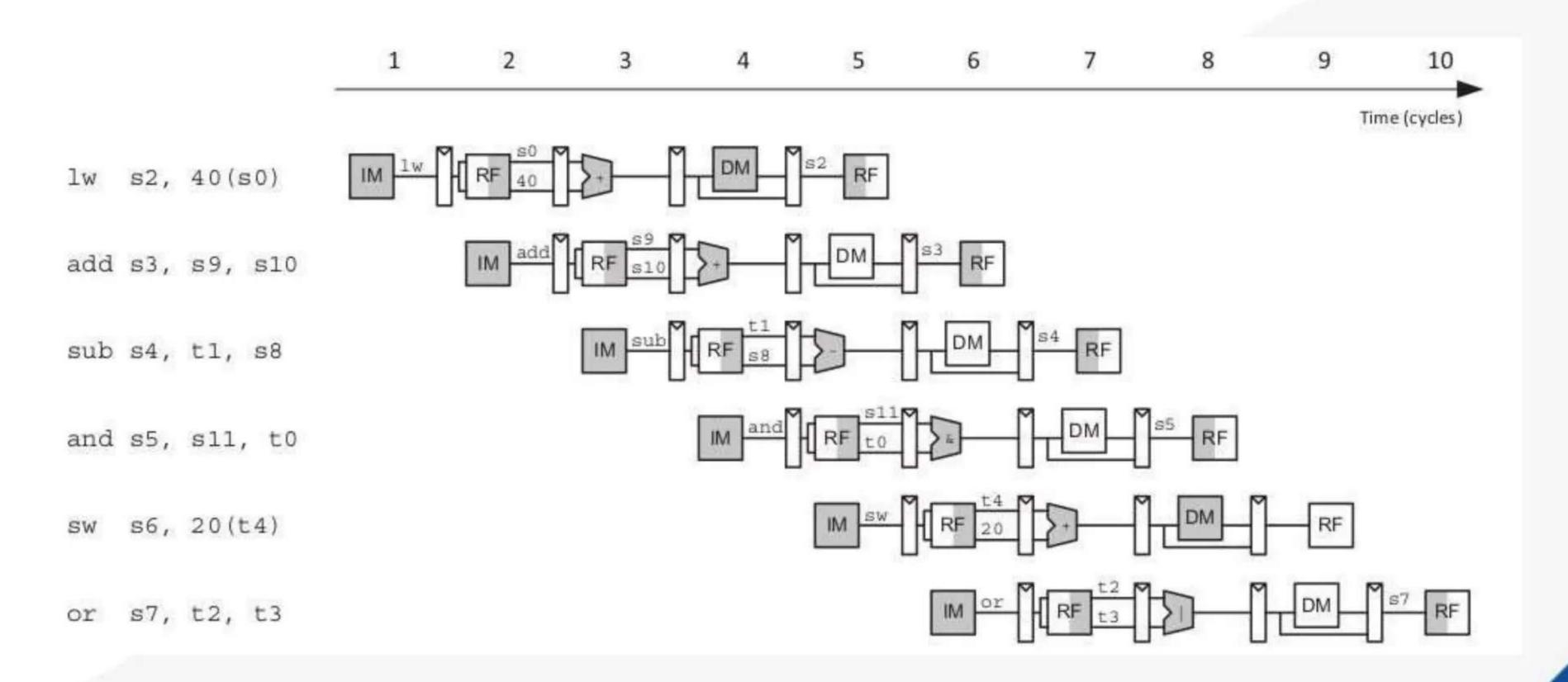
Pipeline Datapath



Implementation of Fetch Cycle



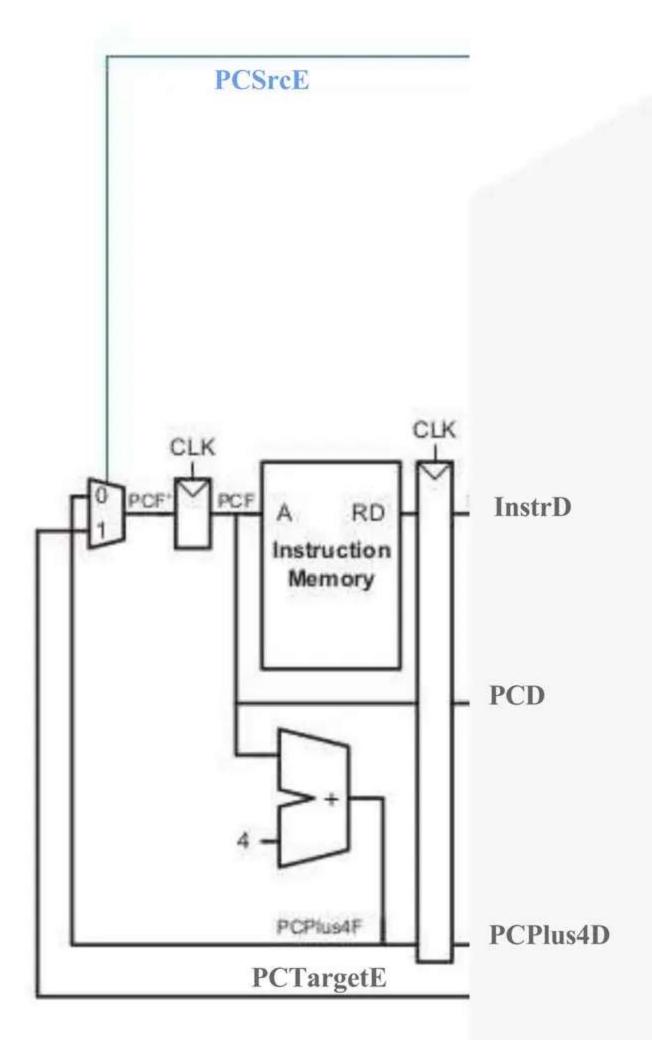
Abstract View of Pipelining



Fetch Cycle Datapath

Modules to be Integrated:

- 1) PC Mux
- 2) Program Counter
- 3) Adder
- 4) Instruction Memory
- 5) Fetch Stage Registers

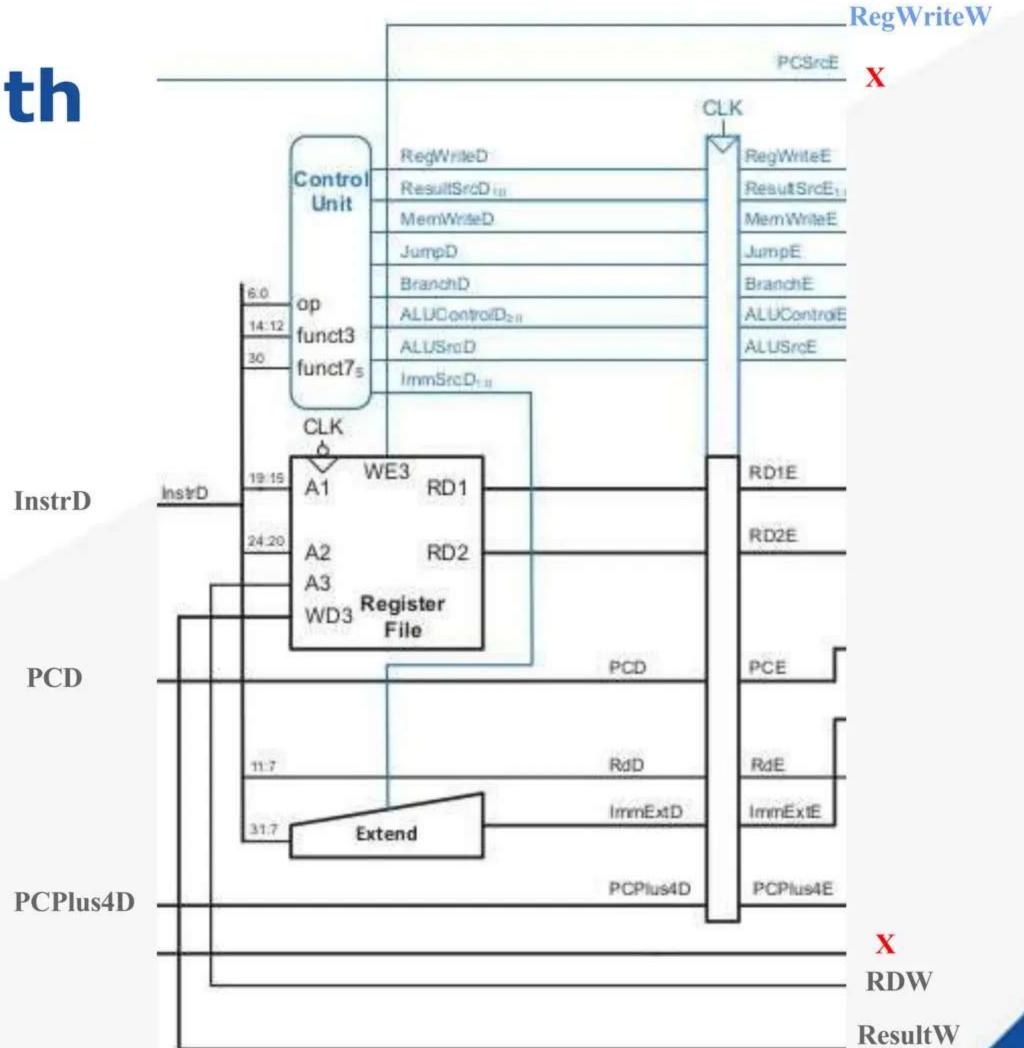


Implementation of Decode Cycle

Decode Cycle Datapath

Modules to be Integrated:

- 1) Control Unit
- 2) Register File
- 3) Extender
- 4) Decode Stage Registers

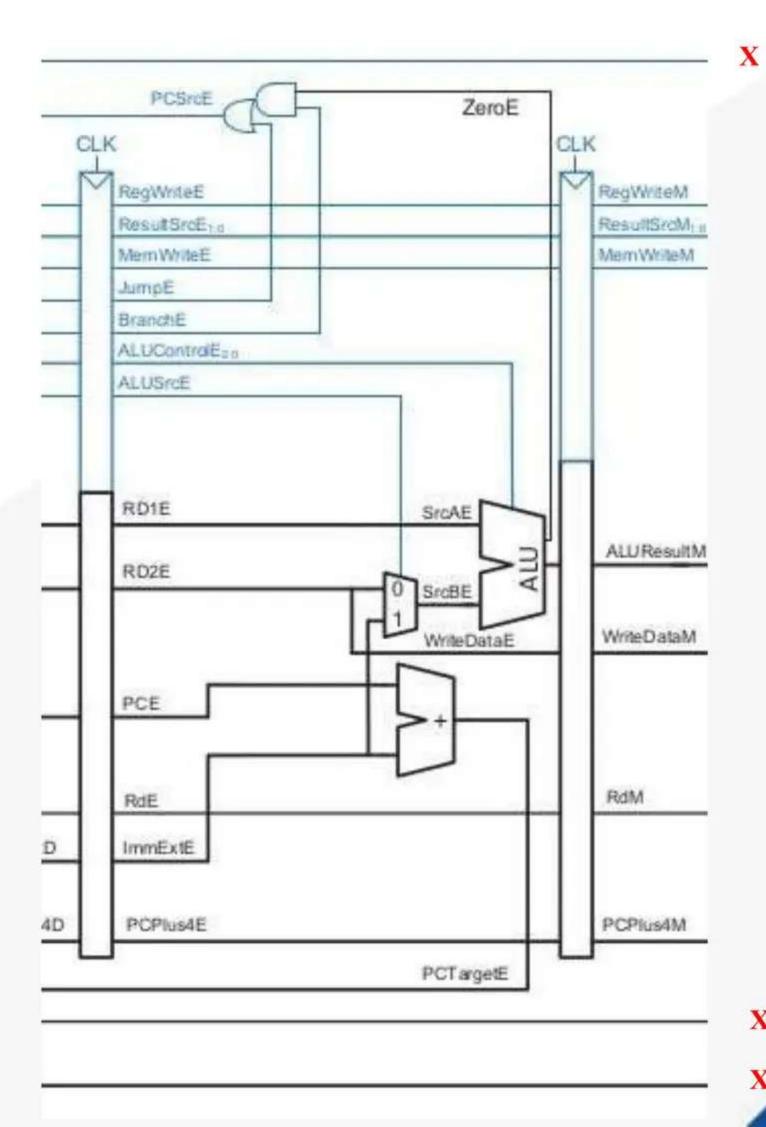


Implementation of Execute Cycle

Execute Cycle Datapath

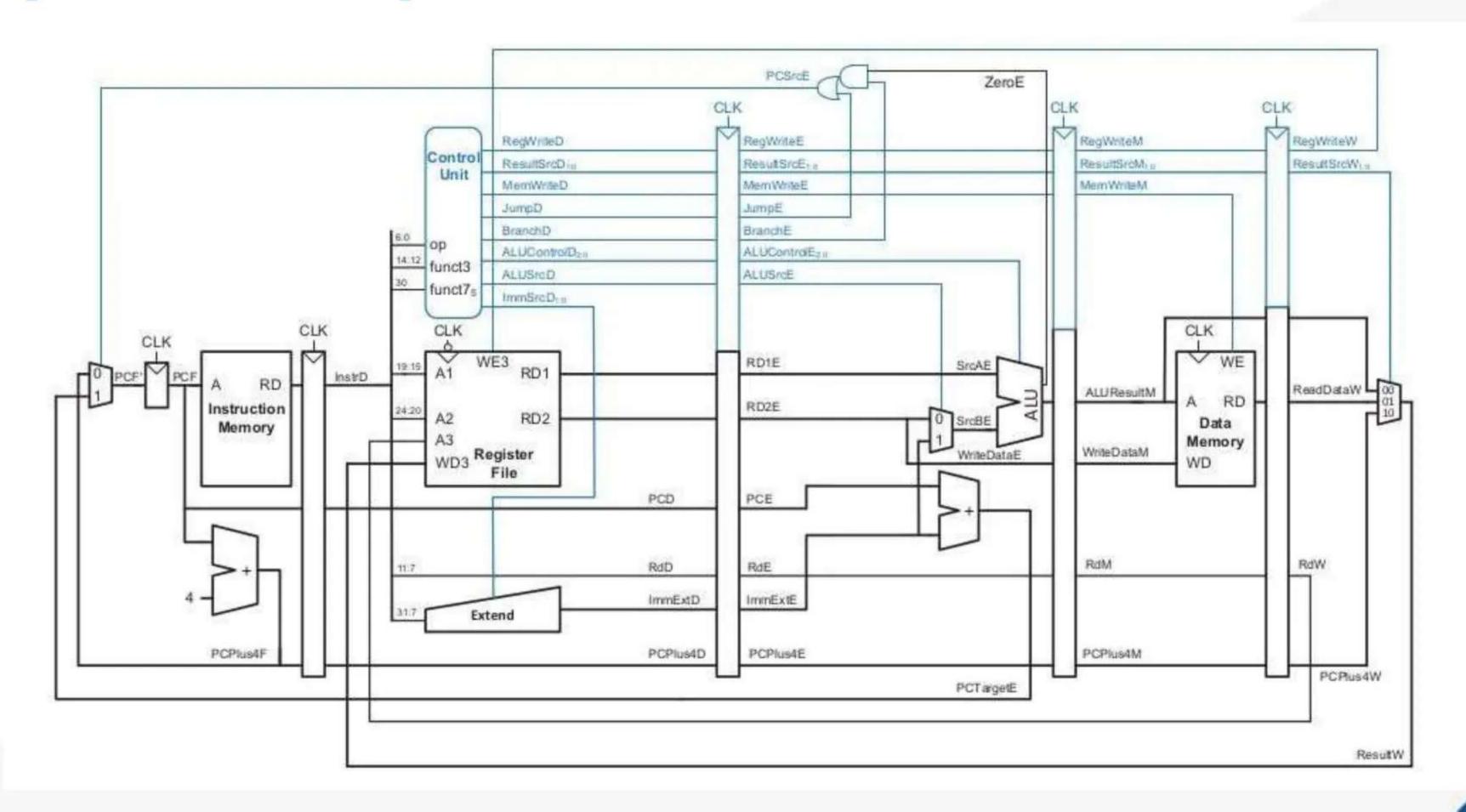
Modules to be Integrated:

- 1) AND Gate
- 2) Mux
- 3) Adder
- 4) ALU
- 5) Execute Stage Registers



Implementation of Memory Cycle

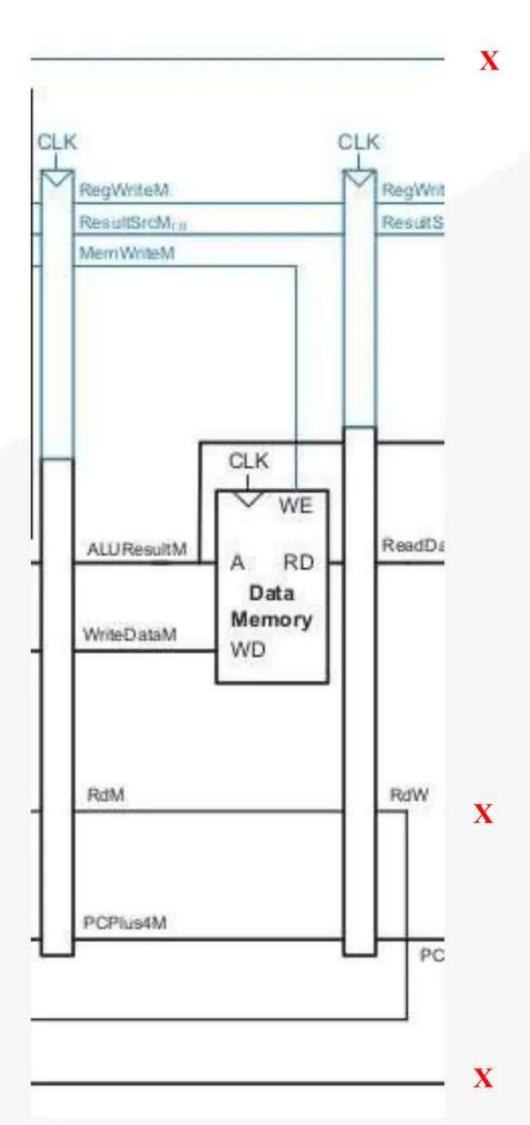
Pipeline Datapath



Memory Cycle Datapath

Modules to be Integrated:

- 1) Data Memory
- 2) Memory Stage Registers



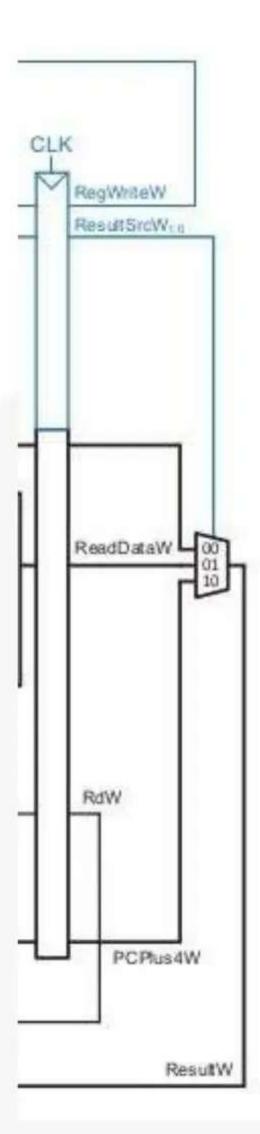
Implementation of Write Bacfi Cycle



Write Bacfi Cycle Datapath

Modules to be Integrated:

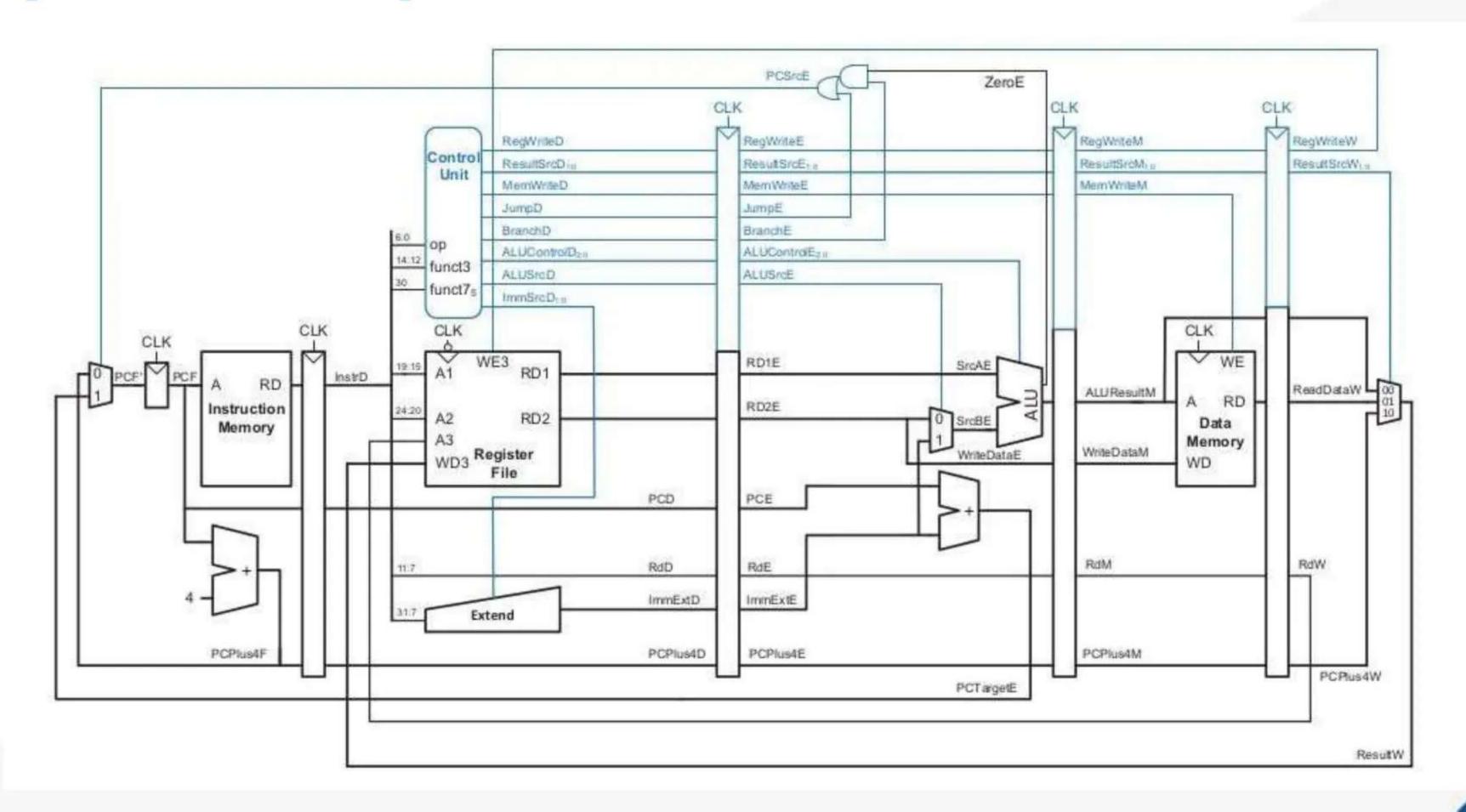
1) Mux



Implementation of Pipeline Top



Pipeline Datapath



Pipeline Hazards

Pipeline Hazard

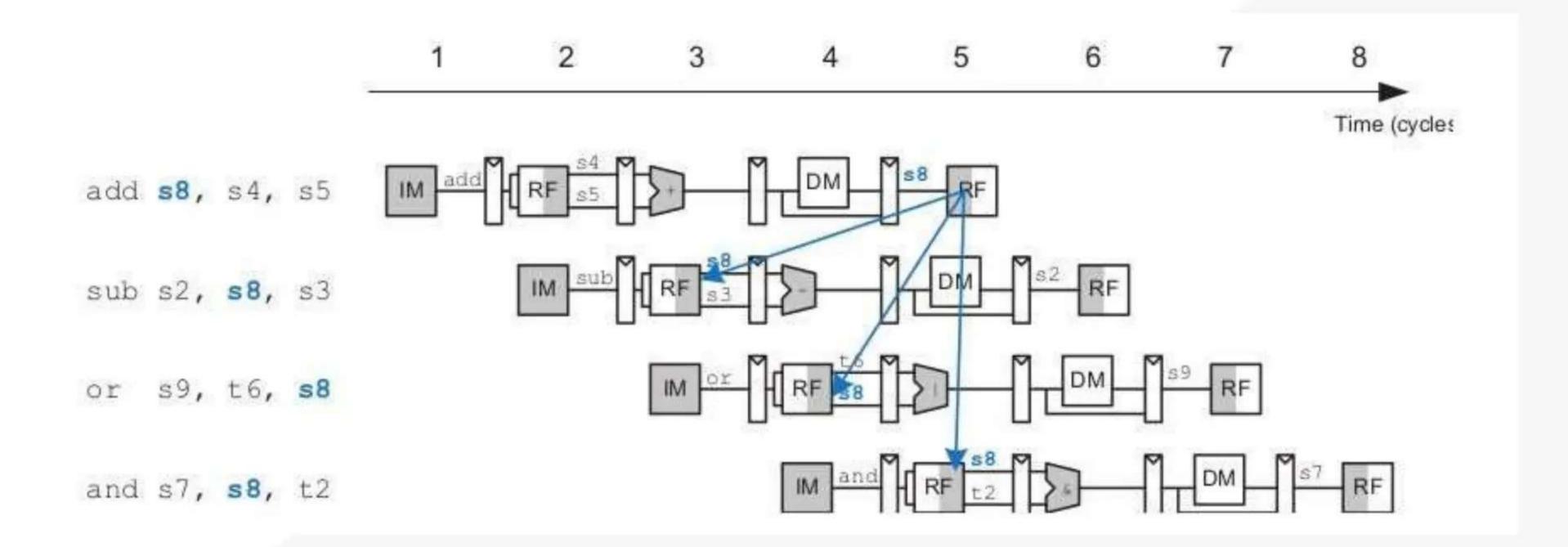
Structural Hazard

- 1. Hardware does not support the execution of instruction in same clock cycle.
- 2. Without having Two memories RISC-V pipelining architecture will have structural hazard.

Data Hazard

- Data to be executed is not available.
- May occur when pipeline is stalled.
- 3. Solve by using **forwarding** or **bypassing** technique.

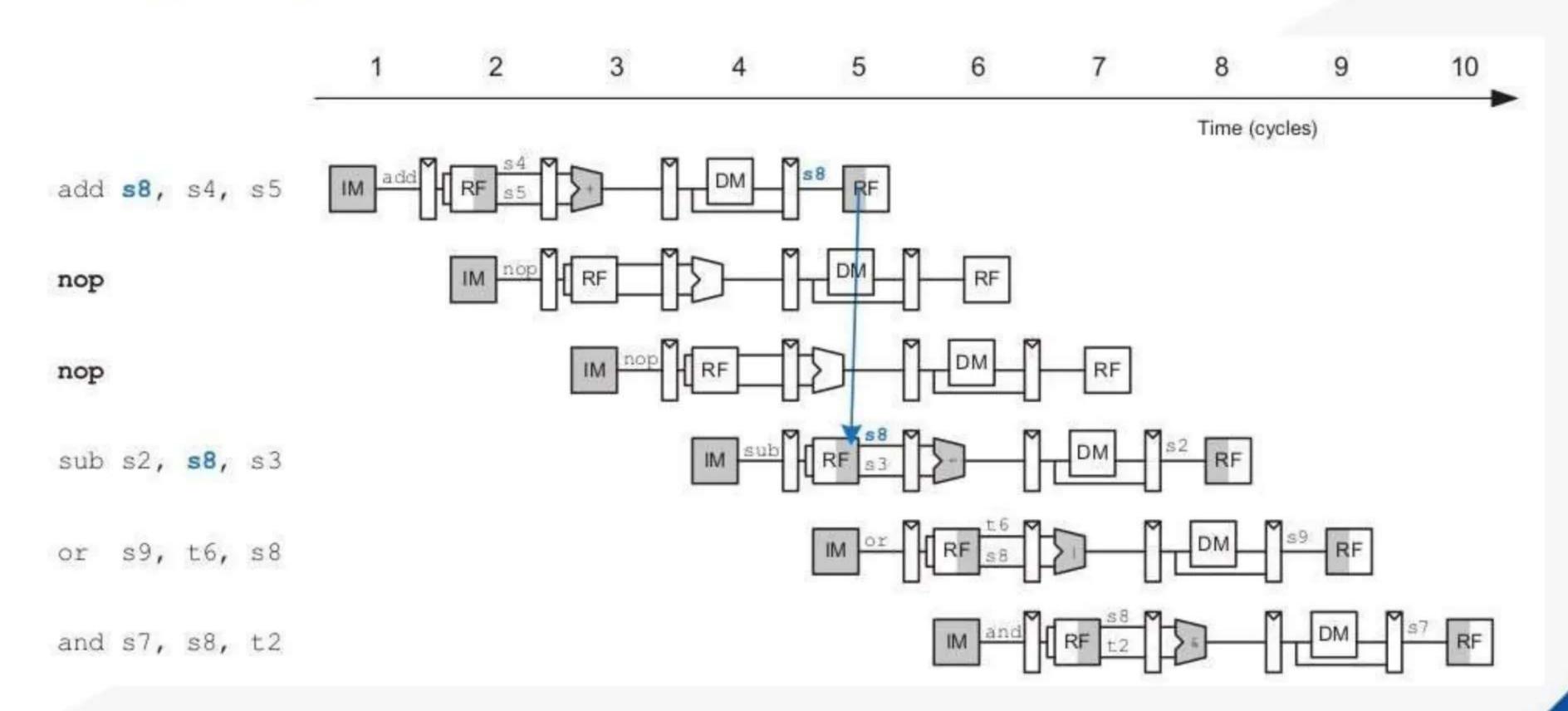
Data Hazard In Pipelining



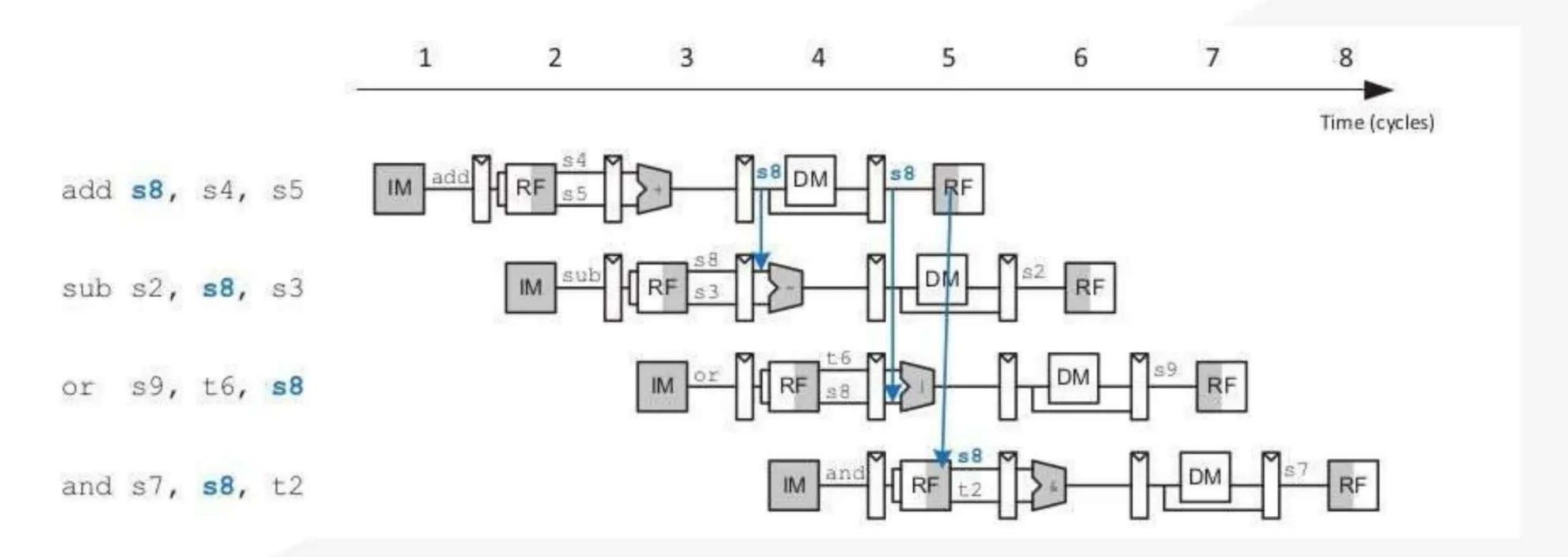
Solution of Data Hazards

- Solving Data Hazards with nops
- Solving Data Hazard with Forwarding / Bypassing

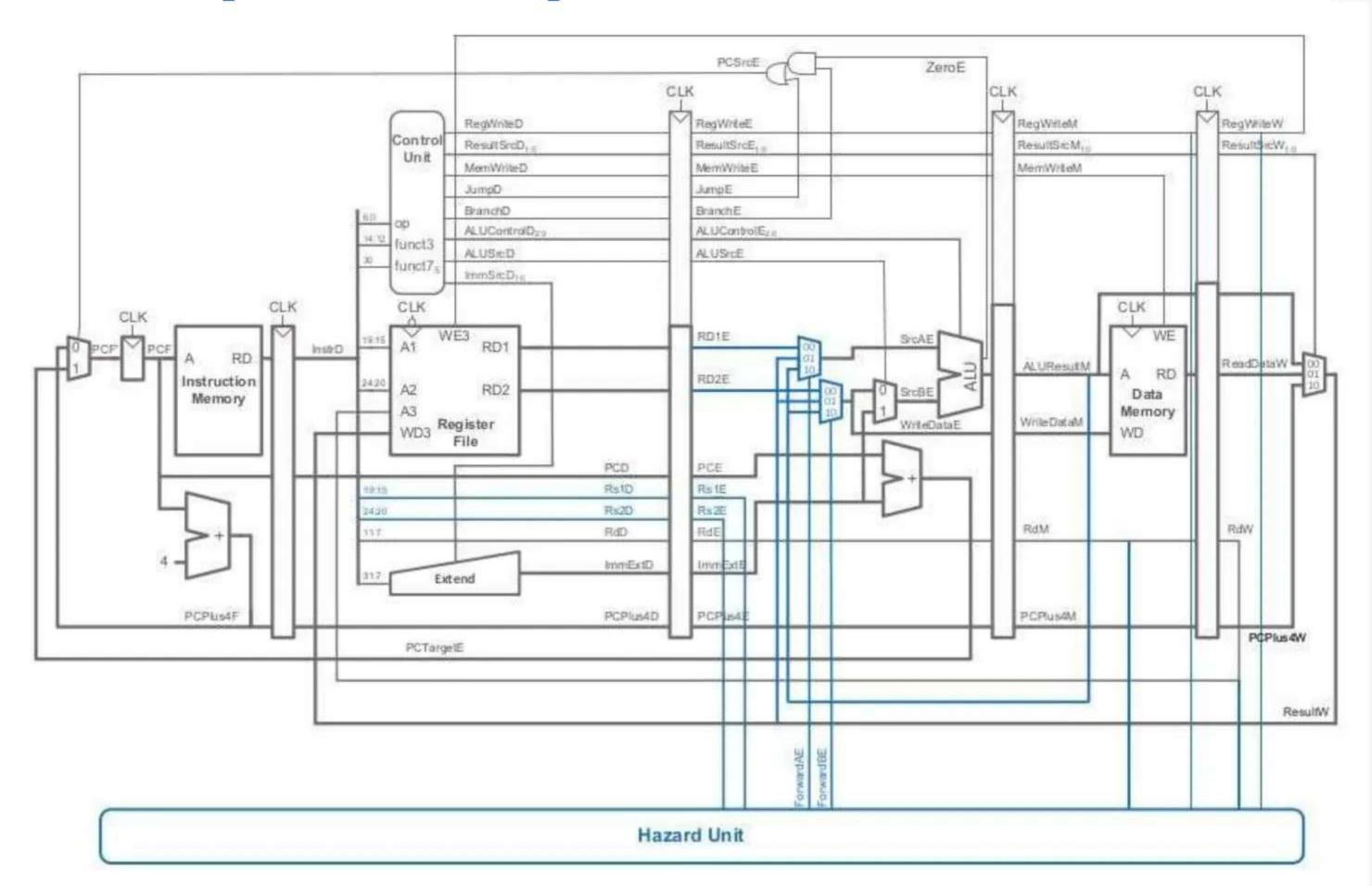
Using Nops



Using Forwarding / Bypassing



Updated Pipeline Top Architecture



Implementation of Hazard Unit



Condition Table

Mux control	Source	Explanation
ForwardA = 00	ID/EX	The first ALU operand comes from the register file.
ForwardA = 10	EX/MEM	The first ALU operand is forwarded from the prior ALU result.
ForwardA = 01	MEM/WB	The first ALU operand is forwarded from data memory or an earlier ALU result.
ForwardB = 00	ID/EX	The second ALU operand comes from the register file.
ForwardB = 10	EX/MEM	The second ALU operand is forwarded from the prior ALU result.
ForwardB = 01	MEM/WB	The second ALU operand is forwarded from data memory or an earlier ALU result.

Condition for Data Hazard

Memory Stage

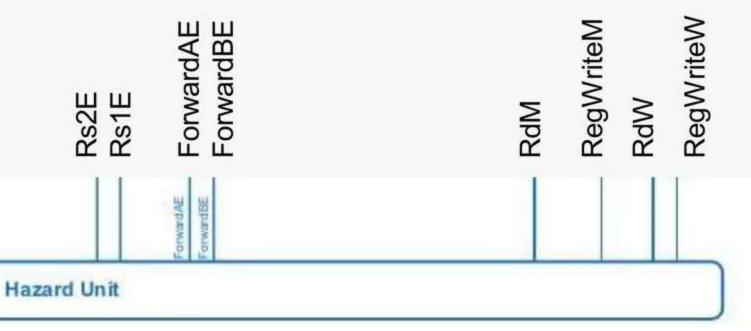
if (RegWriteM and (RdM != 0) and (RdM == Rs1E))
ForwardAE = 10

if (RegWriteM and (RdM != 0) and (RdM == Rs2E))
ForwardBE = 10

WriteBack Stage

if (RegWriteW and (RdW != 0) and (RdW == Rs1E))
ForwardAE = 01

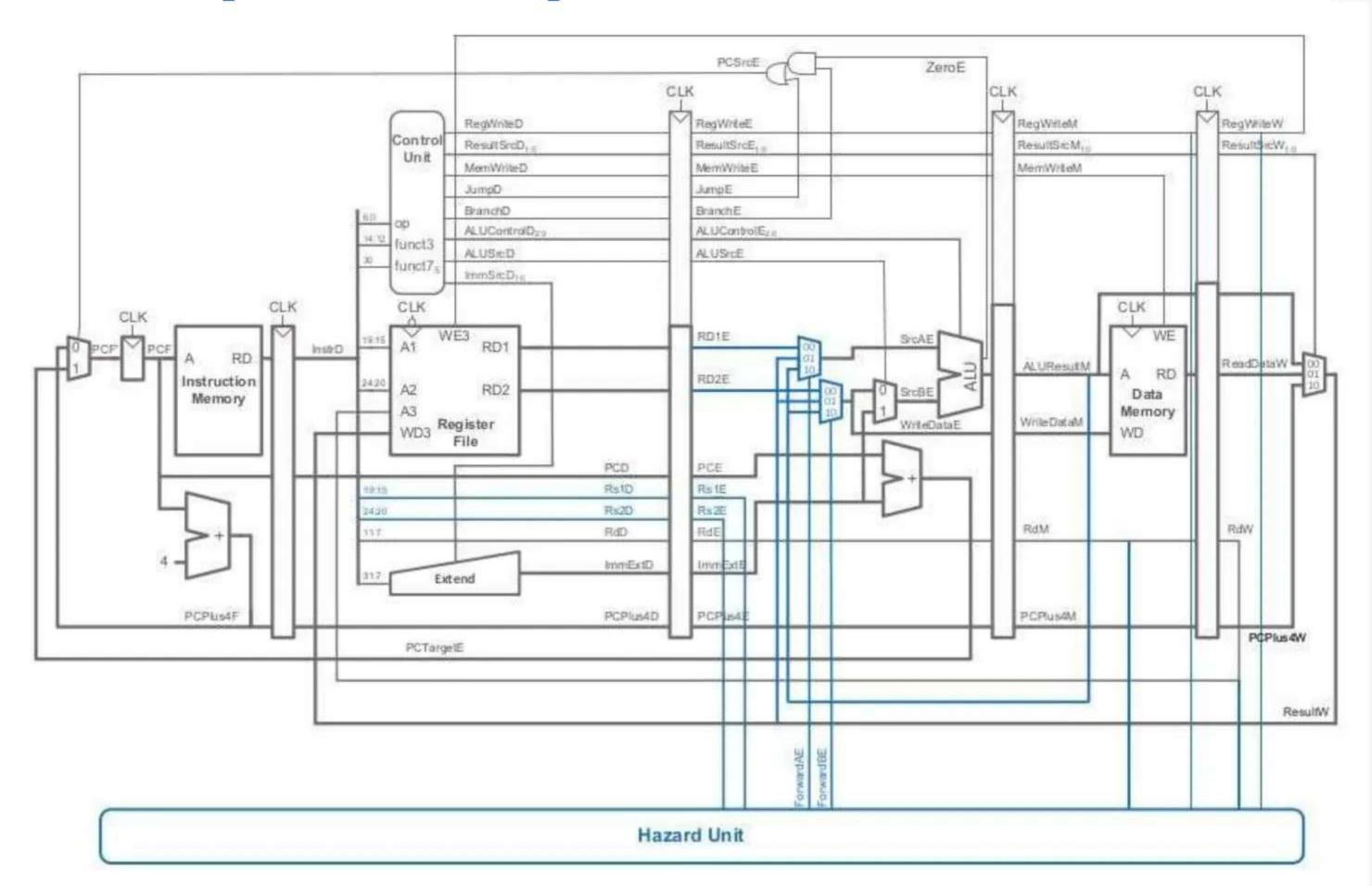
if (RegWriteW and (RdW != 0) and (RdW == Rs2E))
ForwardBE = 01



Implementation of Pipeline Top II



Updated Pipeline Top Architecture



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