



Ashima Garg

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## MULTIPLE SUBJECT : DIGITAL LOGIC + DISCRETE MATHEMATICS (GATE - 2019) - REPORTS

[OVERALL ANALYSIS](#)
[COMPARISON REPORT](#)
[SOLUTION REPORT](#)

ALL(33)

CORRECT(0)

INCORRECT(0)

SKIPPED(33)

### Q. 1

Consider the following two statements:

 $S_1$ : If  $\bar{G}$  is connected, then  $G$  will be disconnected.

 $S_2$ : If  $\bar{G}$  is disconnected, then  $G$  will be connected.

Which of the above statements is correct?

[Solution Video](#) | [Have any Doubt ?](#)

 A  
Both  $S_1$  and  $S_2$ 

 B  
Only  $S_1$ 

 C  
Only  $S_2$ 

Correct Option

**Solution :**

(c)

 $S_2$  is correct and  $S_1$  is incorrect.

 This is because of a theorem which says 'at least one of  $G$  and  $\bar{G}$  must be connected'.

 Anyways for the sake of clarity, here's a counter example for  $S_1$ .

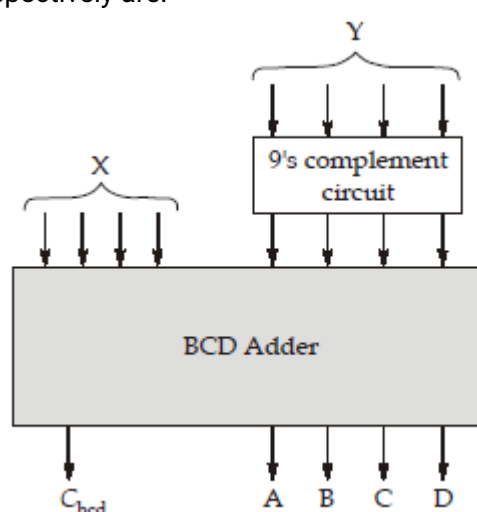
 Clearly  $G_1$  is connected and  $\bar{G}_1$  is also connected.

 Hence  $S_1$  is false.

 D  
None of these

[QUESTION ANALYTICS](#)

### Q. 2

 In the given below circuit input  $X = 1001$  and  $Y = 0011$  the output  $ABCD$  ( $8 - 4 - 2 - 1$ ) and  $C_{bcd}$  of the circuit respectively are:


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B

1011, 0

C

0100, 1

D

0101, 1

Correct Option

**Solution :**

(d)

$$X = 1001$$

$$9's \text{ complement of } Y(0011) = 0110$$

Now

$$X + Y = \begin{array}{r} 1001 \\ + 0110 \\ \hline 1111 \end{array}$$

$$\begin{array}{r} 1111 \\ + 0110 \\ \hline \end{array}$$

$$\begin{array}{r} 1111 \\ + 0110 \\ \hline \end{array}$$

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$$\begin{array}{r} 1111 \\ + 0110 \\ \hline \end{array}$$

$$\begin{array}{r} 1111 \\ + 0110 \\ \hline \end{array}$$

Option (d) is correct.

QUESTION ANALYTICS

**Q. 3**

Consider the following statements:

$S_1$ : If  $n$  pigeons occupy  $m$  pigeonholes such that  $n < m$ , at least  $\left(\left\lfloor \frac{n-1}{m} \right\rfloor + 1\right)$  pigeons will occupy the same pigeonhole.

$S_2$ : If  $m$  pigeons occupy  $n$  pigeonholes such that  $n < m$ , at least  $\left\lfloor \frac{m}{n} \right\rfloor$  pigeons will occupy the same pigeonhole.

Which of the above statements is/are true?

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A

Both  $S_1$  and  $S_2$ 

B

Only  $S_1$ 

C

Only  $S_2$ 

Correct Option

**Solution :**

(c)

In  $S_1$ , no of pigeons  $<$  no of pigeonholes. Hence pigeonhole principle doesn't apply and then  $S_1$  is false.

$S_2$  is true, because


$$\left(\left\lfloor \frac{m}{n} \right\rfloor\right) = \left(\left\lfloor \frac{m-1}{n} \right\rfloor + 1\right)$$


This is actually another version of pigeonhole principle and both are one and the same. Hence  $S_2$  is true.




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QUESTION ANALYTICS

Q. 4

The 6's complement of the number  $(73)_8$ , when it is expressed with the base of 6

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A  
155

B  
312

C  
135

D  
421

Correct Option

**Solution :**  
(d)  
 $(73)_8 = (59)_{10} = (135)_6$   
Now, 5's complement of  $(135)_6$  is

	5 5 5
	<u>- 1 3 5</u>
5's complement:	4 2 0
Add 1 to make for 6's complement:	<u>+ 1</u>
6's complement:	4 2 1

QUESTION ANALYTICS

Q. 5

Let A be a finite set with  $m$  elements. Then the number of elements in the largest equivalence relation of A is

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A  
 $m$

B  
 $m^2$

Correct Option

**Solution :**  
(b)  
The largest equivalence relation is formed by relating all the elements with each other. Hence there will be  $m \times m = m^2$  ordered pairs which means that the size of largest equivalence relation will be  $m^2$ .

C  
 $\frac{m(m-1)}{2}$

D  
1

QUESTION ANALYTICS



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A	B	$Q(t+1)$
0	0	$Q(t)$
0	1	$\bar{Q}(t)$
1	0	0
1	1	1

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A

$\bar{A}\bar{B}Q(t) + B\bar{Q}(t) + AB$

Correct Option

Solution :

(a)

Using K-map:

AB \ Q(t)	00	01	11	10
0		1	1	
1	1		1	

$Q(t+1) = \bar{A}\bar{B}Q(t) + B\bar{Q}(t) + AB$

B

$\bar{A}\bar{B}Q(t) + AB$

C

$\bar{A}\bar{B}\bar{Q}(t) + \bar{A}\bar{B}$

D

$\bar{A}\bar{B}\bar{Q}(t) + \bar{B}Q(t) + AB$

QUESTION ANALYTICS

Q.7

The statement  $(p \Rightarrow p \vee q) \wedge (q \Rightarrow \sim p \vee q)$  is equivalent to which of the following?

- I.  $p \Rightarrow p \vee q$
- II.  $q \Rightarrow (p \Rightarrow q)$
- III. True
- IV. False

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A

I only

B

I and IV only

C

I and III only

D

I, II and III only

Correct Option

Solution :

(d)

The statement is,  $(p \Rightarrow p \vee q) \wedge (q \Rightarrow \sim p \vee q)$

$= (p' + p + q) \wedge (q' + p' + q)$



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$\equiv \text{True}$

Consider  $(p \Rightarrow p \vee q) \wedge (q \Rightarrow \sim p \vee q)$  as

$A \wedge B$

So if  $A \wedge B$  is true, both A and B will be true.

$\Rightarrow$  I is equivalent to the original statement.

II is  $q \Rightarrow (p \Rightarrow q) \equiv p \Rightarrow (\sim p \vee q)$

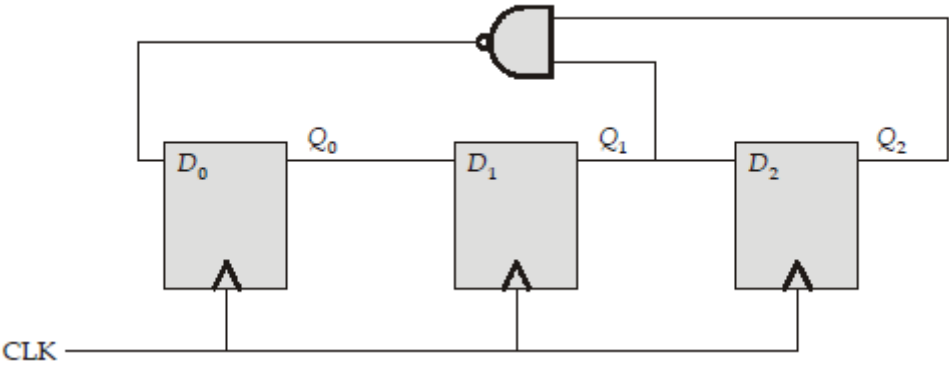
III is also equivalent.

Hence I, II and III all are equivalent to the original statement.

QUESTION ANALYTICS

Q. 8

In the circuit shown below, initially all flip-flops are cleared. The output  $Q_2Q_1Q_0$  after four clock pulses is:



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A

111

B

110

Correct Option

Solution :

(b)

$$D_0 = \overline{Q_1 Q_2} = \overline{Q_1} + \overline{Q_2}$$
$$D_1 = Q_0$$
$$D_2 = Q_1$$

CLK	$D_2$	$D_1$	$D_0$	$Q_2$	$Q_1$	$Q_0$
	0	0	1	0	0	0
1	0	1	1	0	0	1
2	1	1	1	0	1	1
3	1	1	0	1	1	1
4				1	1	0

After 4-clock pulse  $(Q_2Q_1Q_0) = (110)$

C

100

D

000

QUESTION ANALYTICS

Q. 9

Let M and N be two sets respectively. The cardinality of M is  $x$  and the cardinality of N is  $y$ . It is also given that  $y > x$ . Then the number of possible mappings from M to N that are not one one when  $x = 3$  and  $y = 4$  is



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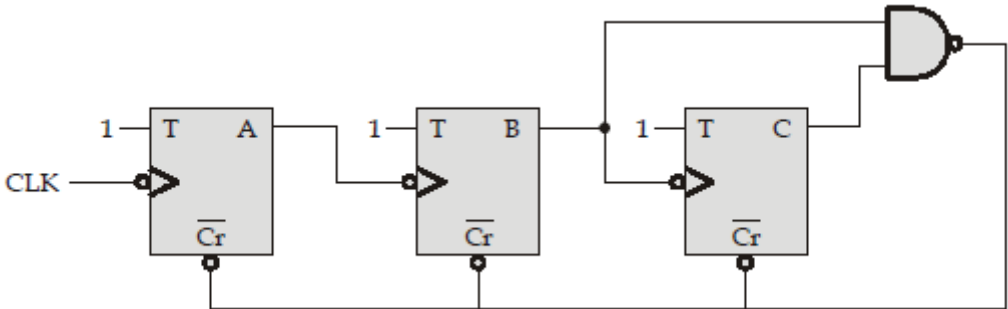
Correct Option

Solution :  
40  
Number of one one mappings =  $({}^yP_x)$   
Total number of functions possible =  $(y^x)$   
 $\Rightarrow$  Number of functions that are NOT one one  
 $= y^x - ({}^yP_x)$   
 $= 4^3 - ({}^4P_3)$   
 $= 64 - (4 \times 3!)$   
 $= 64 - 24$   
 $= 40$

QUESTION ANALYTICS

Q. 10

Consider the asynchronous circuit given below:



The modulus value of circuit is \_\_\_\_\_.

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6

Correct Option

Solution :  
6

A	B	C	$\overline{Cr} = \overline{BC}$
0	0	0	1
1	0	0	1
0	1	0	1
1	1	0	1
0	0	1	1
1	0	1	1
0	1	1	0

QUESTION ANALYTICS

Q. 11

The ratio of chromatic number to the diameter for  $C_{40}$ , where  $C_n$  represents the cycle graph with  $n$  vertices is \_\_\_\_\_. (Upto 1 decimal place)

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0.1


Correct Option


Solution :  
0.1  
$$[\text{Chromatic number for } C_n] = \begin{cases} 2, & n \text{ is even} \\ 3, & n \text{ is odd} \end{cases}$$
$$\Rightarrow \text{Chromatic number } (C_{40}) = 2$$
$$\frac{2}{|n|}$$




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Hence the ratio =  $\frac{2}{20} = \frac{1}{10} = 0.1$

QUESTION ANALYTICS

Q. 12

The 2's compliment representation in 8-bit format is 11010000. The equivalent decimal value of the original number is \_\_\_\_\_.

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48

Correct Option

**Solution :**  
48

$$\begin{aligned} &11010000 \\ \text{2's compliment} &= 00110000 \\ \text{Magnitude} &= 48 \end{aligned}$$

The correct answer is (–48) because MSB bit is 1 in 2's compliment form.

QUESTION ANALYTICS

Q. 13

Number of integers from 1 to 2100 are divisible by 3 or 5 or 7 \_\_\_\_\_.

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1140

Correct Option

**Solution :**  
1140

$$\begin{aligned} n(3 \text{ or } 5 \text{ or } 7) &= n(3) + n(5) + n(7) - [n(3 \cap 5) + n(5 \cap 7) + n(3 \cap 7)] + n(3 \cap 5 \cap 7) \\ &= \left\lfloor \frac{2100}{3} \right\rfloor + \left\lfloor \frac{2100}{5} \right\rfloor + \left\lfloor \frac{2100}{7} \right\rfloor - \left( \left\lfloor \frac{2100}{15} \right\rfloor + \left\lfloor \frac{2100}{35} \right\rfloor + \left\lfloor \frac{2100}{21} \right\rfloor \right) + \left\lfloor \frac{2100}{105} \right\rfloor \end{aligned}$$

Solve to get,  $n = 1140$

QUESTION ANALYTICS

Q. 14

Suppose a synchronous counter is designed that counts the sequence (1, 3, 1, 0, 1, 2 and repeats). The minimum number of T flip-flop are required to construct this sequential circuit are \_\_\_\_\_.

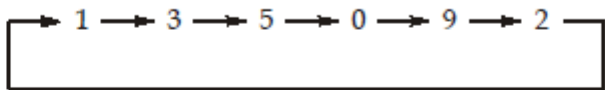
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4

Correct Option

**Solution :**  
4

First design synchronous counter using 4 flip-flops for the sequence.



The sequence will be as follows:

FF3 FF2 FF1 FF0





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0	0	0	0	0	→	0
9	1	0	0	1	→	1
2	0	0	1	0	→	2

Now, take output from FF1 and FF0 to get the desired sequence i.e. 1, 3, 1, 0, 1, 2 as shown dotted box.

QUESTION ANALYTICS

**Q. 15**

 The number of terms in the expansion of  $(x + y + z + w)^3$  is \_\_\_\_\_.

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20

Correct Option

Solution :

20

$$4 - 1 + {}^3C_3 = {}^6C_3 = 20$$

QUESTION ANALYTICS

**Q. 16**

 Simplify the function  $F = \sum m(1, 2, 3, 4, 5, 8, 9, 10)$ . The number of literals in the minimal product of sum form are \_\_\_\_\_.

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11

Correct Option

Solution :

11

		CD	00	01	11	10
AB	00		0	1	3	2
	01		4	5	0	0
	11		0	0	0	0
	10		8	9	0	10

$$F = \overset{(2)}{(\bar{B} + \bar{C})} \overset{(2)}{(\bar{A} + \bar{B})} \overset{(3)}{(\bar{A} + \bar{C} + \bar{D})} \overset{(4)}{(A + B + C + D)}$$

Hence, total number of literals in the minimal POS form = 11.

QUESTION ANALYTICS

**Q. 17**

Consider the sentence 'All the odd and even numbers are integers'. Let the domain be the set of all real numbers. Now consider the following first order formulae.

 I.  $\forall x(\text{Odd}(x) \wedge \text{Even}(x) \Rightarrow \text{Integer}(x))$ 

 II.  $\forall x(\sim \text{Integer}(x) \Rightarrow \sim \text{Odd}(x) \wedge \sim \text{Even}(x))$ 

 III.  $\forall x(\text{Odd}(x) \wedge \text{Even}(x) \wedge \text{Integer}(x))$





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Only I

B

Only II

Correct Option

Solution :

(b)

I is incorrect as a number can't be both odd and even at the same time.

To see how II is correct, take its contrapositive.

$$\forall x[(\text{Odd}(x) \vee \text{Even}(x)) \Rightarrow \text{Integer}(x)]$$

If a number is odd or even, then surely it's an integer.

Hence II is correct.

III is clearly incorrect.

Hence correct choice is (b).

C

Only III

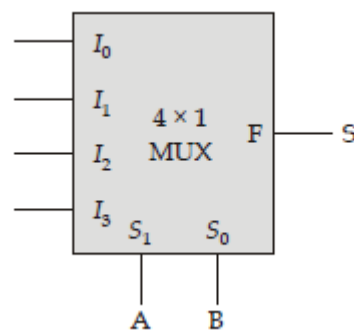
D

None of these

QUESTION ANALYTICS

### Q. 18

Consider a  $4 \times 1$  MUX to be used to implement the sum of a 1-bit full adder with input bits A and B and the carry input  $C_{in}$ . Which of the following combination of inputs to  $I_0, I_1, I_2$  and  $I_3$  of the MUX will realize the sum S?



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A

$$I_0 = I_3 = C_{in}; I_1 = I_2 = \bar{C}_{in}$$

Correct Option

Solution :

(a)

For,  $4 \times 1$  MUX

$$F = \bar{A}\bar{B} \cdot I_0 + \bar{A}BI_1 + A\bar{B}I_2 + ABI_3$$

Now, truth table for full adder.

	A	B	$C_{in}$	Sum	
$I_0$	0	0	0	0	$I_0 = C_{in}$
	0	0	1	1	
$I_1$	0	1	0	1	$I_1 = \bar{C}_{in}$
	0	1	1	0	
$I_2$	1	0	0	1	$I_2 = \bar{C}_{in}$
	1	0	1	0	
$I_3$	1	1	0	0	$I_3 = C_{in}$
	1	1	1	1	

B



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$$I_0 = I_1 = C_{in}; I_2 = I_3 = \bar{C}_{in}$$

D

$$I_0 = I_1 = \bar{C}_{in}; I_2 = I_3 = C_{in}$$

QUESTION ANALYTICS

Q. 19

The value of the following expression:

$$\left[ \frac{1}{81^n} - \frac{10}{81^n} {}^{2n}C_1 + \frac{10^2}{81^n} {}^{2n}C_2 - \frac{10^3}{81^n} {}^{2n}C_3 + \dots + \frac{10^{2n}}{81^n} \right] \text{ is}$$

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A

1024

B

256

C

180

D

1

Correct Option

Solution :  
(d)

Simply the expression first, it's quite easy.

$$\frac{1}{81^n} \left[ 1 - 10 \cdot {}^{2n}C_1 + 10^2 \cdot {}^{2n}C_2 - 10^3 \cdot {}^{2n}C_3 + \dots + (-1)^{2n} 10^{2n} \right]$$

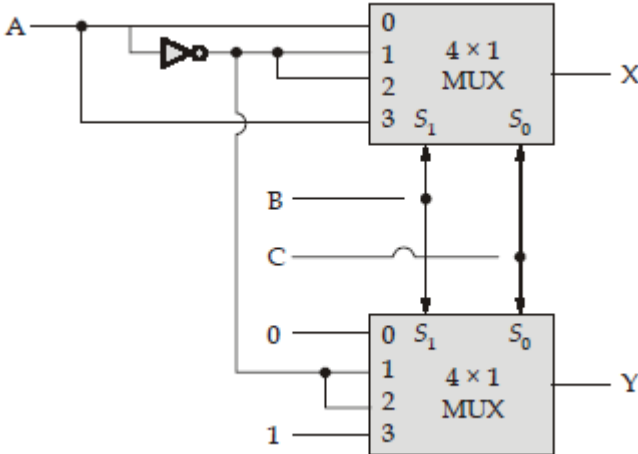
$$\frac{1}{81^n} \left[ {}^{2n}C_0 (10)^0 - {}^{2n}C_1 (10)^1 + {}^{2n}C_2 (10)^2 - {}^{2n}C_3 (10)^3 + \dots + {}^{2n}C_{2n} (10)^{2n} \right]$$

$$\begin{aligned} \frac{1}{81^n} (1 - 10)^{2n} &= \frac{(-9)^{2n}}{81^n} = \frac{[(-9)^2]^n}{81^n} \\ &= \frac{81^n}{81^n} = 1 \end{aligned}$$

QUESTION ANALYTICS

Q. 20

Consider the circuit given below:



The circuit act as

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B  
Full subtractor

Correct Option

**Solution :**  
(b)  
The circuit act as full subtractor.

B	C	X	Y
0	0	A	0
0	1	$\overline{A}$	$\overline{A}$
1	0	$\overline{A}$	$\overline{A}$
1	1	A	1

$$X(A, B, C) = \sum_m(1, 2, 4, 7)$$
$$Y(A, B, C) = \sum_m(1, 2, 3, 7)$$
  
Which is clearly the minterms of full subtractor. X for difference and Y for borrow.

C  
Priority encoder

D  
Half adder sum and half subtractor difference

QUESTION ANALYTICS

Q. 21

How many ways can we distribute at most 10 identical balls to 3 boxes?

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A  
 ${}^{13}C_3$

Correct Option

**Solution :**  
(a)  
Number of ways to distribute  $\leq 10$  identical balls to 3 boxes is equivalent to no of non neg integral solutions to this equation.  
$$x_1 + x_2 + x_3 \leq 10$$
  
Which is same as,  
$$x_1 + x_2 + x_3 + x_4 = 10 \quad \text{[Box method]}$$
  
$$\Rightarrow {}^{4-1+10}C_{10} = {}^{13}C_{10}$$
  
Hence the answer is (a)

B  
 ${}^{13}C_5$

C  
 ${}^{23}C_{13}$

D  
None of these

QUESTION ANALYTICS

Q. 22

A synchronous counter switching sequence is given below:


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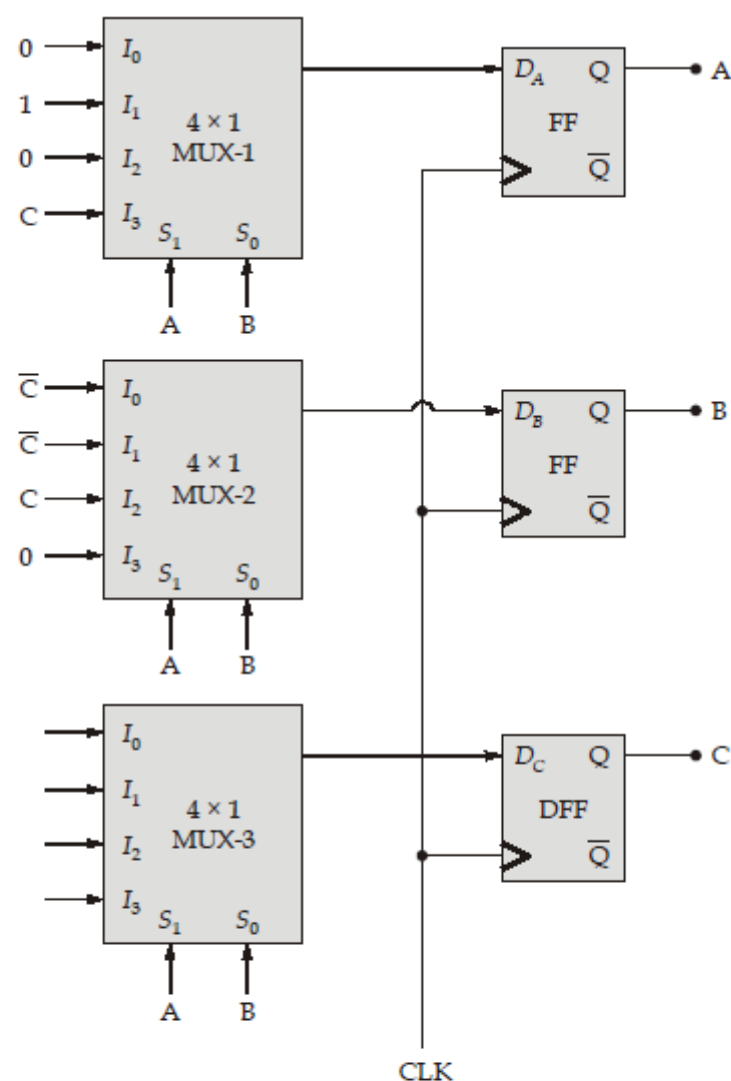
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	$I_0$	$I_1$	$I_2$	$I_3$
(a)	1	C	$\bar{C}$	0
(b)	0	$\bar{C}$	C	C
(c)	$\bar{C}$	C	C	1
(d)	C	0	0	1

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A

a

B

b

**Correct Option**
**Solution :**

(b)

CLK		A	B	C	FF-C
	0	0	0	0	0
	2	0	1	0	①
	7	1	1	1	①
	5	1	0	1	①
	3	0	1	1	0
	4	1	0	0	0
		0	0	0	

Now,

		AB →			
C		00	01	10	11
$\bar{C}$	0	0	②	4	6
C	1	1	3	⑤	⑦
		0	$\bar{C}$	C	C

$$I_0 = 0$$

$$I_1 = \bar{C}$$

$$I_2 = C$$

$$I_3 = C$$

So correct option is (b).


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 D  
d

QUESTION ANALYTICS

**Q. 23**

 Consider the function  $\Gamma : n \rightarrow \mathbb{N}$ , where  $\mathbb{N}$  is the set of natural numbers defined as

$$\Gamma(n) = \begin{cases} n^2, & n \text{ is odd} \\ 2n+1, & n \text{ is even} \end{cases}$$

 For  $n \in \mathbb{N}$ , which of the following is true for  $\Gamma$ ?

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 A  
Surjective but not injective

 B  
Injective but not surjective

 C  
Bijective

 D  
Neither surjective nor injective

**Correct Option**
**Solution :**

(d)

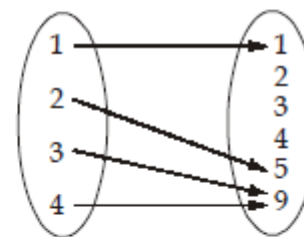
**Check for Injectivity:**

 Put  $n = 3$ 

$$\Gamma(3) = 3^2 = 9$$

 And  $n = 4, \Gamma(4) = 2(4) + 1 = 9$ 

 Since  $\Gamma(3) = \Gamma(4)$  and  $3 \neq 4 \Rightarrow \Gamma$  is not injective.

**Check for Surjectivity:**

 To check for surjectivity, we need to see if Codomain = Range. But here we see that 2, 3, don't have any preimage, thus making  $\text{range} \subset \text{Codomain}$ .

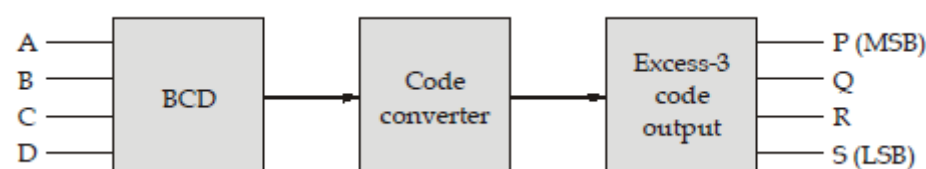
 $\Rightarrow \Gamma$  is INTO, not ONTO.

Hence (d) is the appropriate choice.

QUESTION ANALYTICS

**Q. 24**

Consider a code converter as shown below that converts BCD to excess-3 code for the decimal digits.



The simplified Boolean function output excess-3 code for R will be

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 A  
 $AB + (\overline{C} + \overline{D})$


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C

 $C \odot D$ 

Correct Option

**Solution :**

(c)

**Truth table for BCD to excess-3 code output.**

Decimal value	Input				Output				Decimal value
	A	B	C	D	P	Q	R	S	
0	0	0	0	0	0	0	1	1	3
1	0	0	0	1	0	1	0	0	4
2	0	0	1	0	0	1	0	1	5
3	0	0	1	1	0	1	1	0	6
4	0	1	0	0	0	1	1	1	7
5	0	1	0	1	1	0	0	0	8
6	0	1	1	0	1	0	0	1	9
7	0	1	1	1	1	0	1	0	10
8	1	0	0	0	1	0	1	1	11
9	1	0	0	1	1	1	0	0	12

**K-map for R:**

AB \ CD	CD			
	00	01	11	10
00	1		1	
01	1		1	
11	X	X	X	X
10	1		X	X

$$R = \overline{C}\overline{D} + CD$$

$$= C \odot D$$

D

 $AB + (\overline{C \odot D})$ 
[QUESTION ANALYTICS](#)
**Q. 25**

Out of all boolean matrices of size  $n \times n$  possible, a matrix is picked at random. Let us call this matrix as  $M_A$ . The probability that the matrix chosen obeys the property  $M_A = (M_A)^t$ , where  $(M_A)^t$  refers to the matrix obtained by taking transpose of the original matrix is \_\_\_\_\_. (Take  $n = 3$  and upto 3 decimal places)

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0.125 (0.125 - 0.125)

Correct Option

**Solution :**

0.125 (0.125 - 0.125)

The question is actually asking the probability of picking a symmetric relation out of all possible relations.

$$\text{Required probability, } P(\text{symmetric}) = \frac{2^n \cdot 2^{n(n-1)/2}}{2^{n^2}}$$

 Substitute  $n = 3$ 

$$\text{To get, } \frac{1}{8} = 0.125$$


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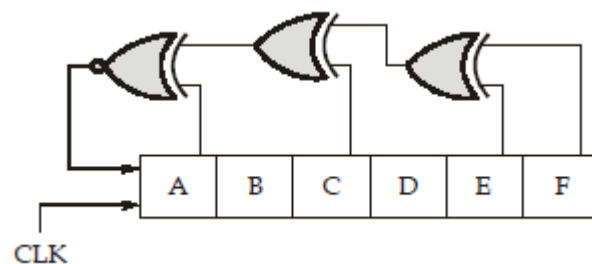
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**Q. 26**

A six bit right shift register is initialized to a value of (A, B, C, D, E, F) = (010000). The minimum number of clock pulses are required to produce 101101 are \_\_\_\_\_.


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4

**Correct Option**
**Solution :**

4

Assume output from XNOR gate is X.

$$X = (C \oplus (E \oplus F)) \odot A$$

CLK	X	A	B	C	D	E	F
0	1	0	1	0	0	0	0
1	1	1	0	1	0	0	0
2	0	1	1	0	1	0	0
3	1	0	1	1	0	1	0
4		1	0	1	1	0	1

Required output

So, total 4 clock cycles are required.

QUESTION ANALYTICS

**Q. 27**

We define a new measure, called GoldIndex(G, C). It takes two arguments as input, namely a graph G, and a set of colours, C respectively. The subroutine outputs an integer denoting the number of ways of assigning colours to vertices in G such that at least two vertices in G have the same colour. Let  $K_n$  denote the complete graph having n vertices respectively, and  $C = \{\text{Red, Green, Blue, Yellow}\}$ . Then the GoldIndex( $K_3$ , C) will be equal to \_\_\_\_\_.

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40

**Correct Option**
**Solution :**

40

Let's do this problem by complementary counting.

Let X = Total number ways of colouring each vertex in G

Let Y = Number of ways of colouring G such that no two vertices have the same colour

$$\text{GoldIndex}(K_3, C) = X - Y$$

Let's first find  $X = 4 \cdot 4 \cdot 4 = 4^3 = 64$

(4 choices i.e. Red, Green, Blue, Yellow for colouring each vertex, and 3 such vertices in  $K_3$ )

$$Y = {}^4C_3 \cdot 3!$$

(Since Chromatic Number of  $K_n$  is n, first choose 3 out of 4 colours and then assign them to vertices in 3! ways)

Therefore GoldIndex( $K_3$ , C) = 64 - 24 = 40

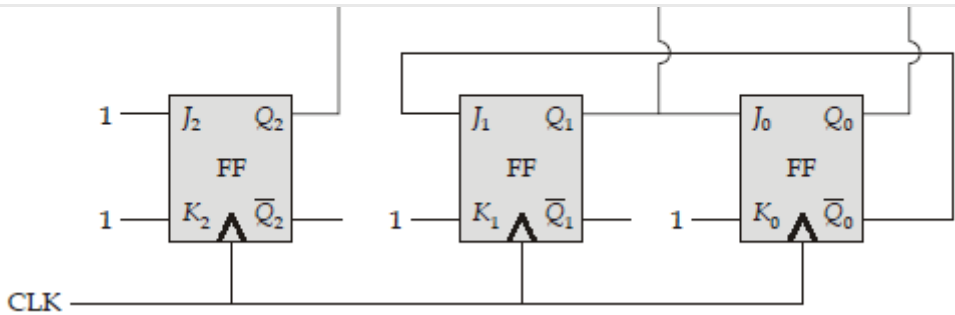
QUESTION ANALYTICS





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The number of clock cycles required to get at  $Q_2Q_1Q_0 = 000$  is \_\_\_\_\_.

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4

Correct Option

Solution :  
4

CLK	Present State			$J_2 = 1 \ K_2 = 1$		$J_1 = \overline{Q_0} \ K_1 = 1$		$J_0 = Q_1 \ K_0 = 1$	
	$Q_2$	$Q_1$	$Q_0$						
	0	0	1	1	1	0	1	0	1
1	1	0	0	1	1	1	1	0	1
2	0	1	0	1	1	1	1	1	1
3	1	0	1	1	1	0	1	0	1
4	0	0	0						

Hence total 4 clocks are required to get the desired state 000.

QUESTION ANALYTICS

Q. 29

The value of the summation,  $\sum_{r=2}^5 {}^5C_r {}^5C_r$  is equal to \_\_\_\_\_.

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226

Correct Option

Solution :  
226

We know,  $\sum_{r=0}^n ({}^nC_r) = 2^n$

Therefore,  $\sum_{r=0}^5 ({}^5C_r)^2 = {}^{10}C_5$

$$= ({}^5C_0)^2 + ({}^5C_1)^2 + \sum_{r=2}^5 ({}^5C_r)^2 = {}^{10}C_5$$
$$\sum_{r=2}^5 ({}^5C_r)^2 = 252 - 26$$
$$= 226$$

QUESTION ANALYTICS

Q. 30

The number of minterms after minimizing the following Boolean expression is \_\_\_\_\_.  
 $[D' + AB' + A'C + AC'D + A'C'D']'$

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1


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$$\begin{aligned}
 &= [D' + AB' + A'C + AC'D + A'C'D]' \\
 &= [D' + AC' + AB' + A' [C + C'D]]' \\
 &= [\underline{D' + AC' + AB' + A'C + A'D}]' \\
 &= [D' + \underline{A' + AC' + AB' + AC}]' \\
 &= [D' + \underline{A' + C' + AB' + A'C}]' \\
 &= [D' + \underline{A' + C' + AB'}] \\
 &= [A' + B' + C' + D']' \\
 &= ABCD
 \end{aligned}$$

Hence, only 1 minterm after minimizing.

QUESTION ANALYTICS

**Q. 31**

 Let X be a set containing n elements. We define a relation R on  $X^4$ , such that

$$R = \{(\alpha, \beta, \Gamma, \delta) \mid (\alpha < \beta \text{ and } \beta < \alpha) \text{ or } (\Gamma < \delta \text{ and } \delta < \Gamma)\}$$

Now consider the following statements:

1. R is reflexive.
2. R is symmetric.
3. R is antisymmetric.
4. R is asymmetric.
5. R is transitive.
6. R is a partial order relation.
7. R is irreflexive.

How many statements is/are correct \_\_\_\_\_.

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5

Correct Option

**Solution :**

5

R if observed carefully, is actually empty set. And we know that the empty relation is irreflexive, symmetric, antisymmetric, asymmetric and transitive.

However it is not reflexive and hence cannot be partial order relation.

Hence 5 is the answer.

QUESTION ANALYTICS

**Q. 32**

Consider the input A = 10110100 and B = 01110000 is feeded as input as shown in the below diagram:



The value of X is \_\_\_\_\_.

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38

Correct Option

**Solution :**

38

$$\begin{array}{rcl}
 A \odot B & [A \text{ EX-NOR } B] & \\
 \Rightarrow & 10110100 & \\
 & 01110000 & \\
 \hline
 & 00111011 & 
 \end{array}$$

Now convert above binary data to gray code.

Gray code of (0011 1011) is 0010 0110.









$$(0010 \ 0110)_2 = (38)_{10}$$



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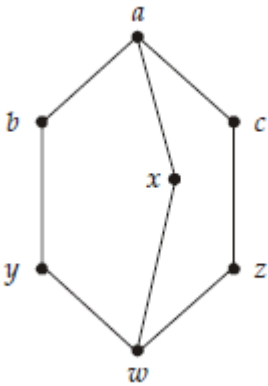
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Q. 33

Consider the Hasse diagram of a lattice given below:



Let X be the number of complements of the element z. Also let Y be the number of complements of y. Then  $(X - Y)^2$  will be equal to \_\_\_\_\_.

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0

Correct Option

**Solution :**  
0  
Complements of element  $y = \{x, c, z\}$   
 $\Rightarrow X = 3$   
Complements of element  $z = \{x, b, y\}$   
 $\Rightarrow Y = 3$   
Hence  $(X - Y)^2 = (3 - 3)^2 = 0$

QUESTION ANALYTICS