



Ashima Garg

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Computer Science Engineering(CS)

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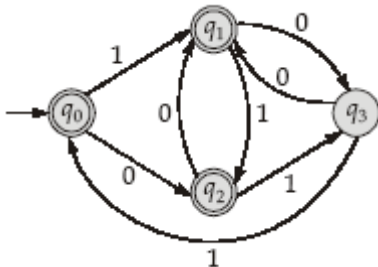
TOPICWISE : THEORY OF COMPUTATION-1 (GATE - 2019) - REPORTS

OVERALL ANALYSIS COMPARISON REPORT SOLUTION REPORT

ALL(17) CORRECT(11) INCORRECT(5) SKIPPED(1)

Q. 1

Consider the following DFA:



Which of the following string is not accepted by above DFA?

FAQ Solution Video Have any Doubt ?

A
011001110

B
11011011001

C
1010101011

Your answer is **Correct**

Solution :
(c)
(a) Accepted
(b) Accepted
(c) Not accepted
(d) Accepted

D
1001110

QUESTION ANALYTICS

Q. 2

Which of the following language is regular?

FAQ Solution Video Have any Doubt ?

A
 $\{wxw^R \mid w \in (a + b)^*, x \in \{a, b\}\}$

Your answer is **Wrong**

B
 $\{wxw^R \mid w, x \in (a + b)^+\}$

Correct Option

Solution :
(b)
• Option (a) represents DCFL since string matching is done i.e. before 'x' and after 'x'.
• Option (b) represents regular for which regular expression $a(a + b)^+ a + b(a + b)^+ b$.
• Option (c) represents CFL since ww^R is done first, hence needs a comparison, which cannot be done by finite automata.
• Option (d) represents CFL since ww^R contains string matching, which can not be done via finite automata.


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D

$$\{ww^R \mid w \in (a+b)^*\}$$

QUESTION ANALYTICS

Q. 3

Let w be any string of length n in $\{a, b\}^*$. Consider L be the set of all strings ending with at least n b's. What is the minimum number of states in a non-deterministic finite automata that accept L ?

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A

$$n$$

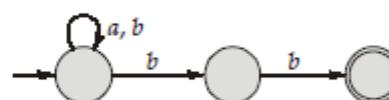
B

$$n+1$$

 Your answer is **Correct**
Solution :

(b)

The number of states for minimum NFA for end with at least n b's is $n+1$ i.e. regular expression $= (a+b)^* b^n$.



So, number of states needed will be $n+1$.

C

$$n+2$$

D

$$2^n$$

QUESTION ANALYTICS

Q. 4

Consider the following statements:

S_1 : DFA for language which contain ' ϵ ' must have initial state as final state.

S_2 : For any language either a language L or its complement L' must be finite.

S_3 : If L is set of all string ending with at least n b's then minimum number of states in non deterministic finite automata that accept L is $n+2$.

S_4 : Non deterministic finite automata is more powerful than deterministic finite automata.

Which of the above statement is true?

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A

 S_1 only

 Your answer is **Correct**
Solution :

(a)

- If DFA accept a null, then initial state must be final state.
- Consider a language $L = \{a^n\}$ on alphabet $\{a, b\}$ and its complement $\Sigma^* - \{a^n\}$ both are infinite. Hence false.
- If L is set of all string ending with at least n b's then minimum number of states in deterministic finite automata that accept L is $n+1$. Hence false.
- Since every language accepted by a NFA is also accepted by some DFA, hence non deterministic finite automata has same power as deterministic finite automata. Hence false.


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 S_2, S_3, S_4 only

D

 S_1, S_3 only

QUESTION ANALYTICS

Q. 5

Which one of the following regular expression describes the language over {a, b} consist of no pair of consecutive a's?

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A

 $(b^*abb^*)(a + \epsilon)$

B

 $(b + ab)^*(a + \epsilon)$

 Your answer is **Correct**
Solution :

(b)

 Option 'a' does not contain 'b' or ' ϵ '.

Option 'c' does not contain 'ba' or 'bba'.

 Option 'b' contain $\epsilon, b, a, ab, ba, \dots$ i.e. no string of pair of consecutive a's.

Option 'd' contain 'aa' i.e. not allowed.

C

 $(b^*abb^*)^*(a + \epsilon) + b^*$

D

 $(b^*ab^*)^*(a + \epsilon) + b^*(a + \epsilon)$

QUESTION ANALYTICS

Q. 6

 The minimum number of states required for DFA that accept the language $L = \{a^n \mid n \text{ is multiple of 3 but not multiple of 5}\}$ are _____.

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15

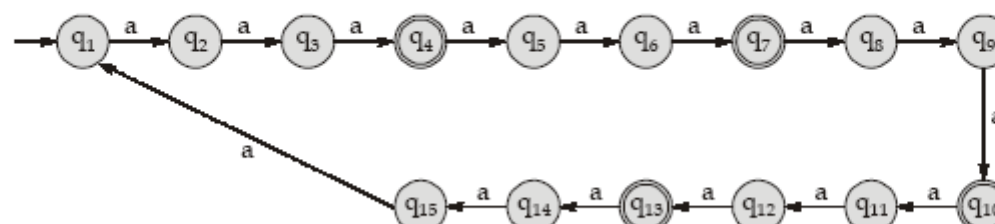
 Your answer is **Correct**15

Solution :

15

 $L = \{a^n \mid n \text{ is multiple of 3 but not multiple of 5}\}$ will contain all strings of a which are divisible by 3 but not a multiple of 15.

So, DFA will be:



So, 15 states are required.

QUESTION ANALYTICS


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The length of the shortest string not in the language (over $\Sigma = \{0, 1\}$) for regular expression $1^*(0 + 10)^*1^*$ is_____.

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4

Correct Option
Solution :

4

Check the string one by one starting from ϵ , 0, 1, 00, 01,... until we reach the first string that is not generated by the given regular expression. In this case, smallest string not generated by the given regular expression is '0110' whose length is 4.

Your Answer is 3

QUESTION ANALYTICS

Q. 8

The number of states in minimal deterministic finite automata for strings starting with ab^2 and ending with b over the alphabet $\{a, b\}$ is _____.

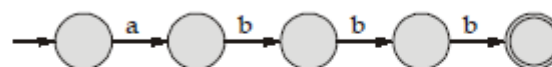
[Solution Video](#) | [Have any Doubt ?](#)

6

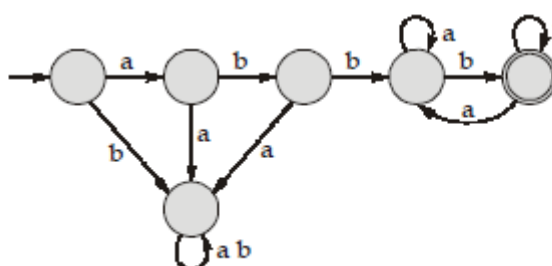
Your answer is Correct6
Solution :

6

Minimum length string accepted by DFA will be ab^2b



Now fill up the left cases, i.e.



QUESTION ANALYTICS

Q. 9

The number of DFA's with four states which can be constructed over the alphabet $\Sigma = \{a, b\}$ with designated initial state are 2^n , then the value of n is _____.

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20

Correct Option
Solution :

20


| | a | b |
|-------|----------|----------|
| q_0 | (1 of 4) | (1 of 4) |
| q_1 | (1 of 4) | (1 of 4) |
| q_2 | (1 of 4) | (1 of 4) |
| q_3 | (1 of 4) | (1 of 4) |


The number of DFA's without any final state are $4^8 \Rightarrow 2^{16}$




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$$= 2^{20}$$

So, n will be 20.

QUESTION ANALYTICS

Q. 10

Consider L_{DF} set of all languages accepted by DPDA by final state and L_{DE} be set of all languages accepted by DPDA by empty stack. Which of the following is true?

FAQ | [Solution Video](#) | [Have any Doubt ?](#) | 

A

$$L_{DF} \subset L_{DE}$$

B

$$L_{DF} = L_{DE}$$

Your answer is **Wrong**

C

$$L_{DF} \supset L_{DE}$$

Correct Option

Solution :
(c)

The set of languages accepted by final state DPDA is proper super set of languages accepted by empty stack DPDA i.e. regular language a^*b^* is accepted by final state but not empty stack.
So, $L_{DF} \supset L_{DE}$.

D

None of these

QUESTION ANALYTICS

Q. 11

Consider P and Q be language over $\Sigma = \{0, 1\}$ represented by the regular expression $0^*(10^*)^*$ and $(0^* + 1^*)^*$ respectively. Which of the following is true?

FAQ | [Solution Video](#) | [Have any Doubt ?](#) | 

A

$$P \subset Q$$

B

$$Q \subset P$$

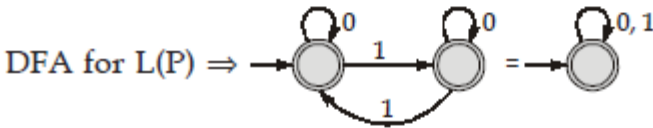
C

$$P = Q$$

Your answer is **Correct**

Solution :
(c)

$$\begin{aligned} L(P) &= 0^*(10^*)^* \\ &= \{\epsilon, 0, 1, 10, 01, 00, 11, \dots\} \\ L(Q) &= (0^* + 1^*)^* \\ &= \{\epsilon, 0, 1, 10, 01, 00, 11, \dots\} \end{aligned}$$




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$$P \cap Q = 0^*1^*$$

QUESTION ANALYTICS

Q. 12

Consider the following statements:

$$S_1 : \{(a^n)^m \mid n \leq m \leq 0\}$$

$$S_2 : \{a^n b^n \mid n \geq 1\} \cup \{a^n b^m \mid n \geq 1, m \geq 1\}$$

Which of the following is regular?

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A

 Only S_1

B

 Only S_2

 Your answer is **Wrong**

C

 Both S_1 and S_2

Correct Option

Solution :

(c)

$$\text{Put } n = 1 \text{ in } S_1 \text{ we get } \{(a^1)^m \mid 1 \leq m\} \cup \{\epsilon\} \\ = \{a^m \mid m \geq 0\} = a^*$$

- Therefore S_1 is regular.
- S_2 represents $a^n b^n \cup a^* b^+ = a^* b^+$ which is regular. Hence regular.

D

 Neither S_1 nor S_2

QUESTION ANALYTICS

Q. 13

 Consider a push down automata (PDA) below which runs over the input alphabet (a, b) . It has the stack alphabet $\{z_0, X\}$ where z_0 is the bottom of stack marker. The set of states of PDA is $\{q_0, q_1\}$ where q_0 is the start state.

$$\delta\{q_0, b, z_0\} = \{(q_0, Xz_0)\}$$

$$\delta\{q_0, b, X\} = \{(q_0, XX)\}$$

$$\delta\{q_0, a, X\} = \{(q_1, X)\}$$

$$\delta\{q_0, \epsilon, z_0\} = \{(q_0, \epsilon)\}$$

$$\delta\{q_1, b, X\} = \{(q_1, \epsilon)\}$$

$$\delta\{q_1, a, z_0\} = \{(q_0, z_0)\}$$

The language accepted by PDA is

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A

$$L = \{(b^n a b^n a)^m \mid m, n \geq 0\}$$

Correct Option

Solution :

(a)

The PDA for given transition function is:

$$\epsilon, z_0, \epsilon$$

$$b, z_0, Xz_0$$










$$b, X, XX$$

$$b, X, \epsilon$$



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a, z_0, z_0

$$L = \{(b^n ab^n a)^m \mid m, n \geq 0\}$$

For clearer understanding, kindly refer the solution video of this question.

B

$$L = \{(b^n ab^n a)^m \mid n, m \geq 0\} \cup \{b^n \mid n \geq 0\}$$

Your answer is **Wrong**

C

$$L = \{(b^n ab^n)^m a \mid n, m \geq 0\}$$

D

None of the above

QUESTION ANALYTICS

Q. 14

Consider the following three languages:

- $L_1 = \{w \mid w \in \{a, b\}^* \text{ and } w = w^R\}$
- $L_2 = \{ww^R \mid w \in \{a, b\}^*\}$
- $L_3 = \{w(a+b)w^R \mid w \in \{a, b\}^*\}$

What is the relation between L_1, L_2 and L_3 ?

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A

$$L_2 \subset L_1 \text{ and } L_3 \subset L_1 \text{ and } L_1 = L_2 \cup L_3$$

Your answer is **Correct**

Solution :
(a)
 L_2 is even palindrome on $\{a, b\}^*$
 L_3 is odd palindrome on $\{a, b\}^*$
 L_1 is any palindrome on $\{a, b\}^*$
Clearly, $L_2 \subset L_1$ and $L_3 \subset L_1$ and $L_1 = L_2 \cup L_3$

B

$$(L_2 = L_3) \subset L_1$$

C

$$L_2 \cap L_1 = L_3$$

D

$$L_2 \subset L_1 \text{ and } L_3 \subset L_1 \text{ but } L_1 \neq L_2 \cup L_3$$

QUESTION ANALYTICS

Q. 15

The number of states in the minimal deterministic finite automata corresponding to the regular expression $(0 + 1)^* (000 + 001)$ is _____.

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5

Your answer is **Correct5**

Solution :


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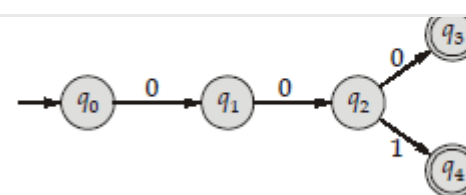
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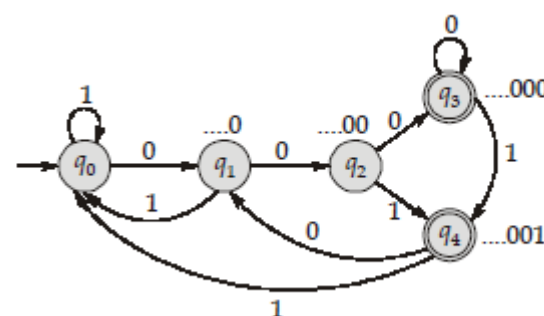
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Now fill the rest of transition:



QUESTION ANALYTICS

Q. 16

 The number of states in minimal NFA, which accepts all strings in which the 3rd last bit is b is _____.
 [Assume $\Sigma = \{a, b\}$]

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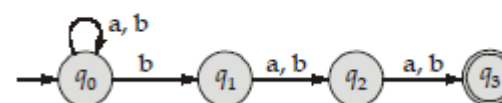
4

 Your answer is **Correct**4

Solution :

4

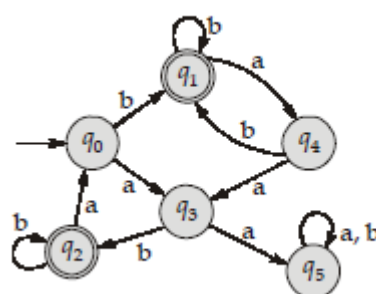
Minimal NFA:



QUESTION ANALYTICS

Q. 17

Consider the following DFA:



The number of states in the minimal DFA obtained by applying minimization algorithm on the above DFA is equal to _____.

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4

 Your answer is **Correct**4

Solution :

4

 Partition-1: $\{q_1, q_2\}, \{q_0, q_3, q_4, q_5\}$

 Partition-2: $\{q_1, q_2\}, \{q_0, q_3, q_4\}, \{q_5\}$

 Partition-3: $\{q_1, q_2\}, \{q_0, q_4\}, \{q_3\}, \{q_5\}$


 Partition-4: $\{q_1, q_2\}, \{q_0, q_4\}, \{q_3\}, \{q_5\}$

Therefore 4 states will be required.




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