Program Reasoning

11. Search Space Pruning

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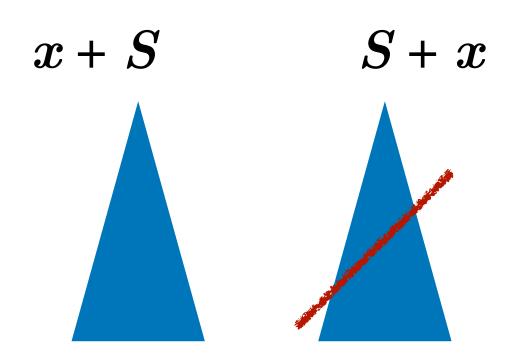


Search Space Pruning

- Problem: too huge search space
- How to prune the search space for program synthesis?
- Common strategies:
 - Equivalence reduction = discard redundant candidates
 - Top-down propagation = discard unpromising candidates

Equivalence Reduction

- Discard redundant subprograms as early as possible
- How to implement this idea for top-down and bottom-up searches?



Eq. Reduction for Top-down

- If two non-ground programs will be expand the same set of programs, discard one of them
- What candidates can be discarded by the following programs in the queue?

if
$$(true)$$
 y S

$$1 * (S + S)$$

sort(S)

Top-down + Eq. Reduction

```
top-down(G = \langle \Sigma, N, R, S \rangle, \phi):
  Q := \{S\}
  while Q != {}:
    p := dequeue(Q)
    if ground(p) \wedge \phi(p): return p
    P' := unroll(R, p)
     forall p' \in P':
       if not equiv(p, p'):
         enqueue(Q, p')
unroll(R, p):
  0' := \{\}
  A := left-most non-terminal in p
  forall (A \rightarrow B) in R:
    p' := p[B/A]
    Q' := Q' \cup \{p'\}
  return Q'
```

Eq. Reduction for Bottom-up

- How to apply equivalence reduction for bottom-up search?
- Problem: candidates are all ground but might not be whole
- Solution for PBE: observation equivalence
 - We only care of equivalence on the given inputs
- E.g., the programs below are observationally equivalent modulo the inputs (1,0) and (2,1)

$$x$$
 if $(x \le y) y x$

Example

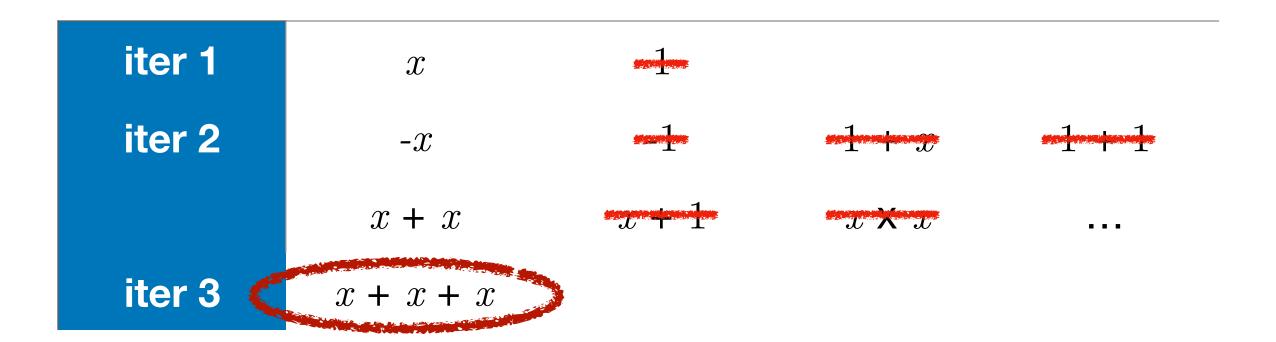
Specification

Find a function f(x) where f(1) = 3

Grammar

$$S \rightarrow x \mid 1 \mid -S \mid S + S \mid S \times S$$

Enumeration



Bottom-up + Eq. Reduction

```
bottom-up(G = <Σ, N, R, S>, φ):
    Q := set of all terminals in G
    while true:
        forall p in Q:
            if φ(p): return p
        Q += grow(R, Q)

grow(R, Q):
    Q' := {}
    forall (A → B) in R:
        Q' += { B[p/C] | p ∈ Q, C →* p }

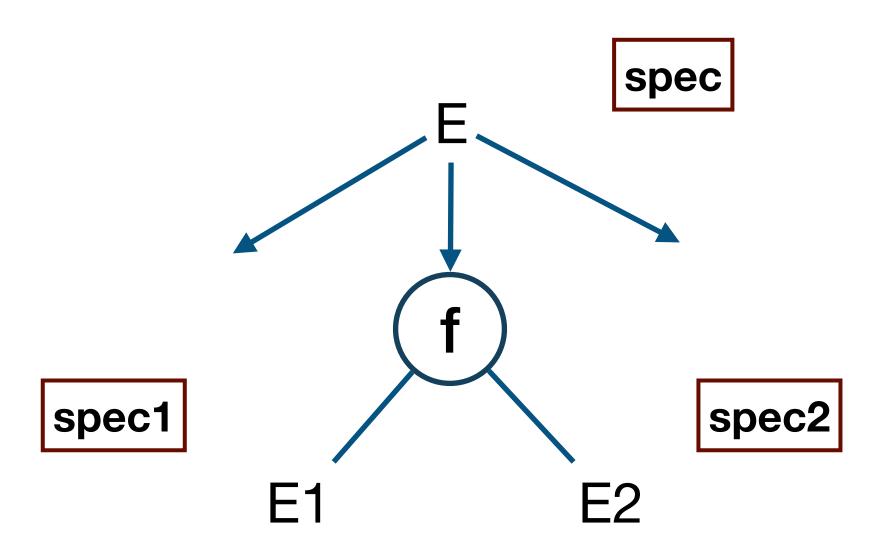
+ Q' := { p'∈ Q' | forall p ∈ Q. ¬equiv(p, p') }
    return Q'
```

Search Space Pruning

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Top-down Propagation

- Discard unpromising subprograms as early as possible
- Given a spec and a production, infer specs for subprograms (divide-and-conquer)
 - When $f < E_1$, E_2 , ..., $E_n > (In) = Out where <math>E_i$ is a subprogram
 - What is the spec for each E_i?
 - If any E_i is undesirable, discard f



Example: String Manipulation

Specification

Find a function f(x) where $f("SA") = "USA" \land f("AE") = "UAE"$

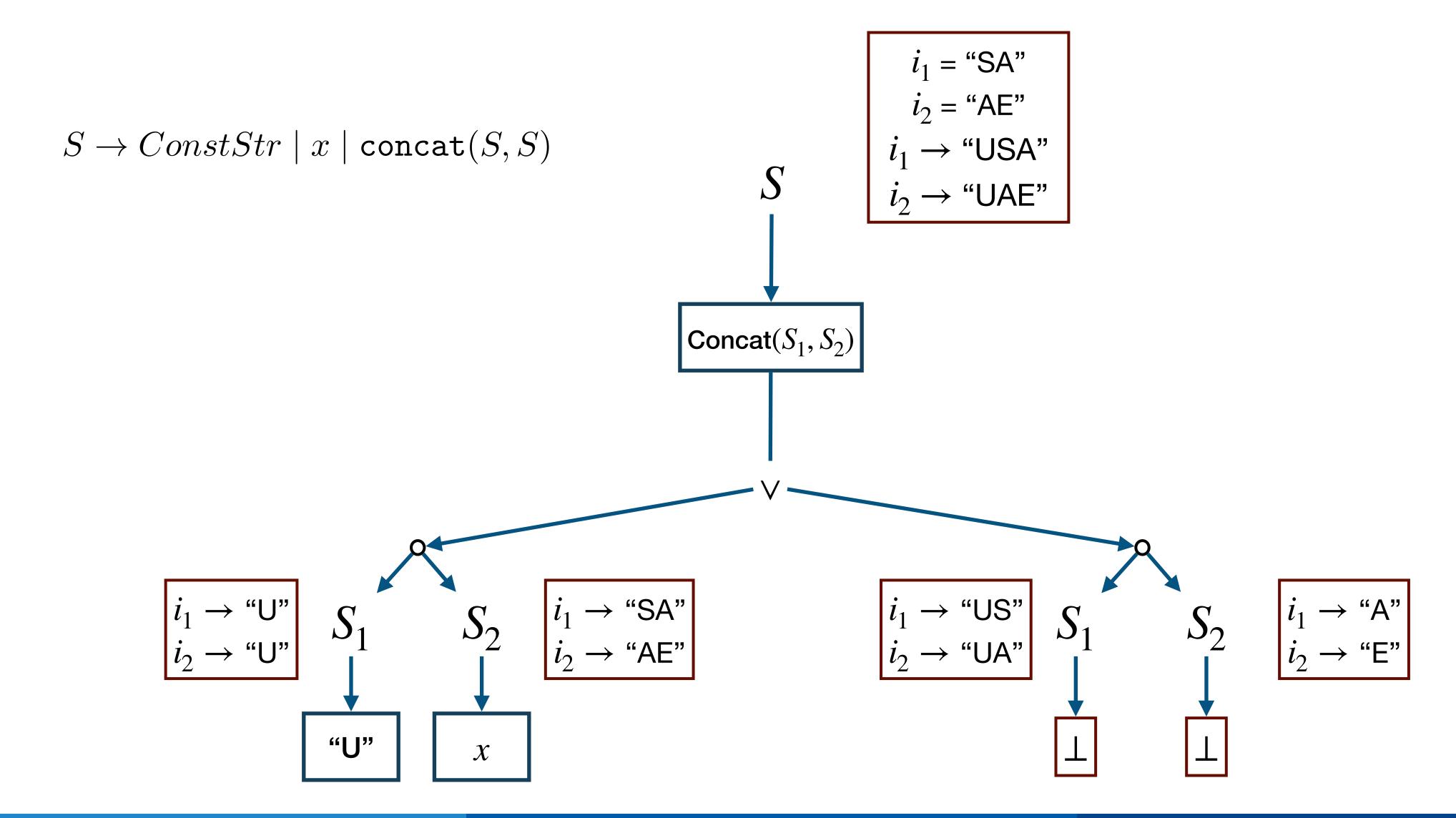
Grammar

$$S \to ConstStr \mid x \mid concat(S, S)$$

Examples

$$concat ("U", "SA") = "USA"$$

TDP for String Manipulation



Example: List Manipulation

Specification

Find a function f(x) where f([1; -3; 1; 7]) = [2; -2; 2; 8]

Grammar

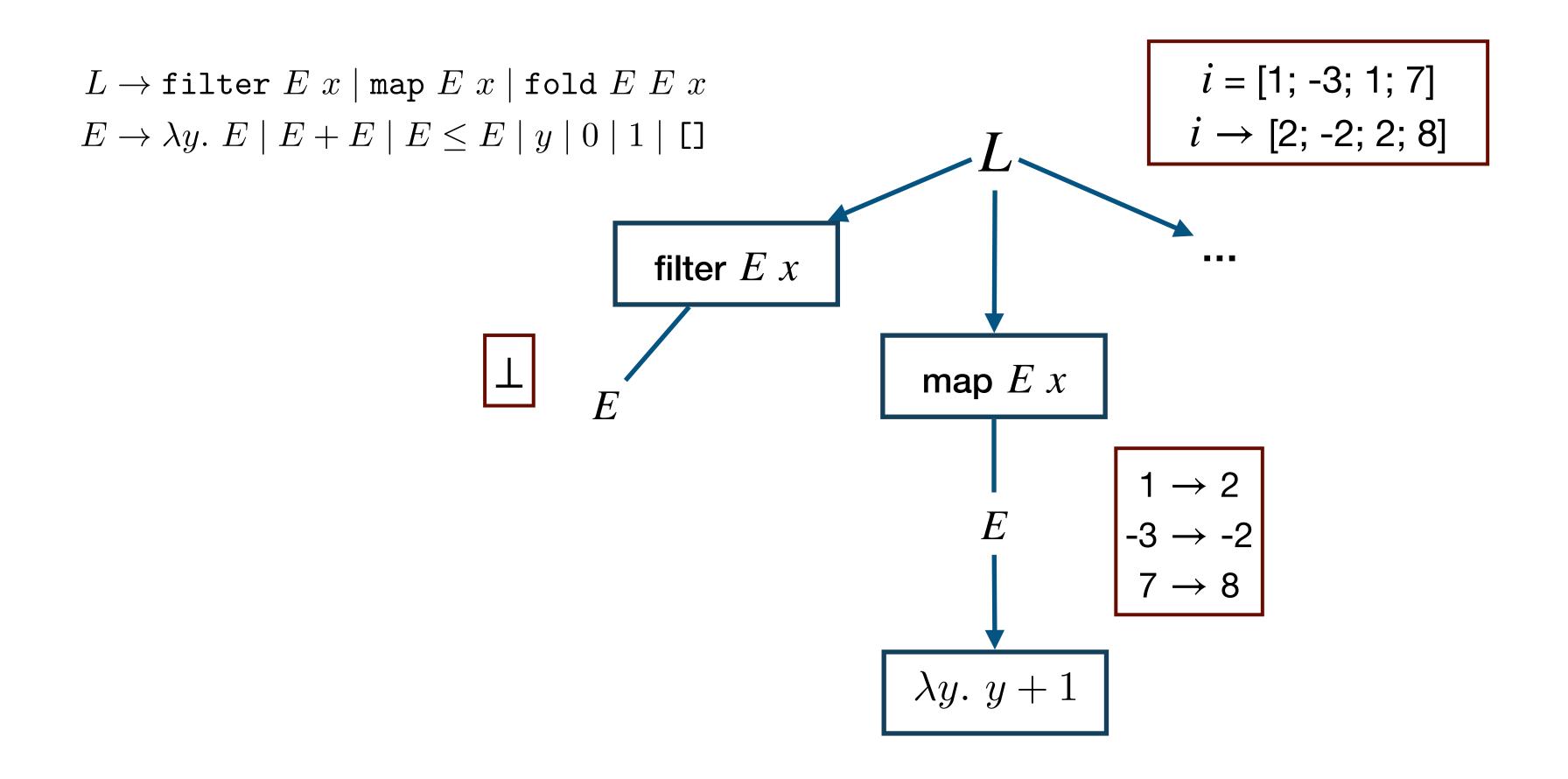
$$L \to \text{filter } E \ x \mid \text{map } E \ x \mid \text{fold } E \ E \ x$$

$$E \to \lambda y. \ E \mid E + E \mid E \le E \mid y \mid 0 \mid 1 \mid []$$

Examples

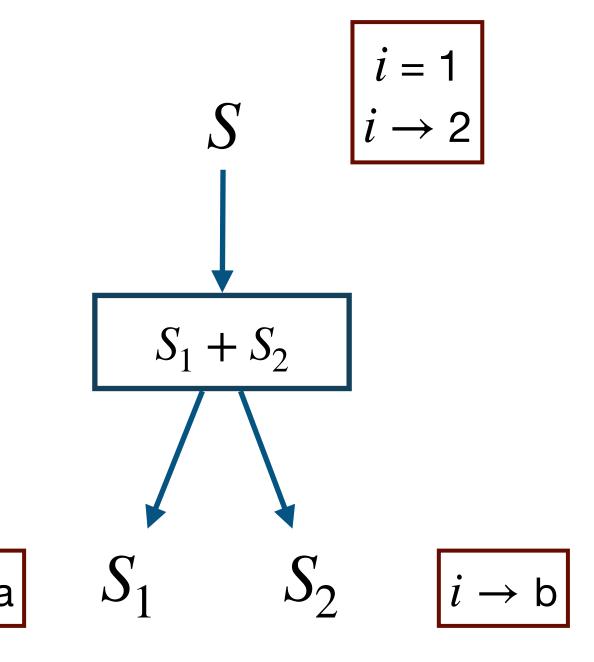
filter
$$(\lambda y.\ y \le 0)$$
 [1; -3; 1; 7] = [-3] map $(\lambda y.\ y+1)$ [1; -3; 1; 7] = [2; -2; 2; 8] fold $(\lambda y.\lambda z.\ y+z)$ 0 [1; -3; 1; 7] = 6

TDP for List Manipulation



Limitation of TDP

- Applicable only when the function is injective so that the inverse exists
- For example, (+): int -> int -> int



What are possible values of a and b?

Summary

- Challenge: huge search space
- Equivalence reduction = discard redundant candidates
 - Applicable to top-down and bottom-up searches
- Top-down propagation = discard unpromising candidates
 - Applicable to top-down search when the inverse functions are computable