

Polymorphism

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Beyond type equality

Defining our own subtypes

Live code demo: person. employee, and student/animal, dog, cat/...

Interfaces and polymorphism

Live code demo: list. empty, and

node.

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Hogeschool Rotterdam Rotterdam, Netherlands



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Lecture topics

- Beyond type equality
- Defining our own subtypes
- Lists and state machines



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Beyond type equality



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Introduction

- In the previous lecture we have seen that Java/C# have a static type system
- Programs that make no sense are outright refused by the compiler
- "Make no sense" means calling methods or making assignments with the wrong types



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So far, we defined a type to be wrong if it is not exactly the same as what we expected.

If we expected an int, but got a Person, then clearly something was off and we expect^a a compiler error.

awelcome, actually



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Wrong types

The two most important typing rules for these violations are those of variable assignment and method call.



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When assigning a variable, we expect the type of the expression and that of the variable to be exactly the same: T = D[x]

$$\frac{\langle \mathtt{e}, D \rangle \to \langle T, D \rangle \wedge \ T = D[x]}{\langle \mathtt{x} = \mathtt{e}, D \rangle \to \langle D[\mathtt{void}], D \rangle}$$



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When calling a method, we expect the type of the parameters and those of the arguments to be exactly the same: $P_i = P'_i$

$$\frac{\langle \mathtt{c}, D \rangle \to \langle \mathtt{C}, D \rangle \wedge \ \langle \mathtt{m}, C \rangle \to \langle (\mathtt{P}_1 \times \mathtt{P}_2 \times \dots \times \mathtt{P}_n \to \mathtt{R}), C \rangle \wedge \ \langle \mathtt{p}_i, D \rangle \to \langle \mathtt{P}_i', D \rangle \wedge \ \mathtt{P}_i = \mathtt{P}_i'}{\langle (\mathtt{c.m}(...\mathtt{p}_i..)), D \rangle \to \langle \mathtt{R}, D \rangle}$$



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For example, the program below violates typing rules:

```
int x = 5;
x = "uh!?";
```



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demo: 6
person, 7
employee, 8
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t/animal, 10
dog, cat/... 11

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node. Coliticiality

For example, the program below violates typing rules:

```
class Counter {
  private int cnt;
  public Counter() {
    this.cnt = 0;
  }
  public void Incr(int diff) {
    this.cnt = (this.cnt + diff);
  }
}
Counter c = new Counter();
c.Incr(c);
```



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int x = 5.5; makes no sense

This will also give a compiler error.



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Always wrong?

- What about a similar expression, such as float x = 10;?
- x is a float
- 10 is an int
- float \neq int, so we should get a type error.



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float x = 10; makes sense, even though it violates the strictest typing rules described so far.

Floating point numbers "contain" integers, so converting the integer 10 to 10.0 loses no information



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Subtyping

This means that the typing rules as seen so far are too restrictive: we should be able to accept an assignment from a more specific data type (such as int) to a less specific data type (such as float).



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Consider any two data types, T and S

We say that S < : T, read "S is a subtype of T", to mean that any value of type S can be safely used where a value of type T is expected.

We also say that, when S <: T, T generalizes S.



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More and less specific data types

For example, int <: float, because any value of type int can be safely used where a value of type float is expected (as the conversion loses no data, and thus can be inserted by the compiler itself).



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More and less specific data types

We can now amend many of our typing rules: instead of type equality, we can use subtyping to preserve safety, but achieve more flexibility.



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When assigning a variable, we expect the type of the expression to be a subtype of that of the variable.

$$\frac{\langle \mathtt{e}, D \rangle \to \langle T, D \rangle \wedge \ T <: D[x]}{\langle \mathtt{x} = \mathtt{e}, D \rangle \to \langle D[\mathtt{void}], D \rangle}$$



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When calling a method, we expect the type of the parameters to be subtypes of the types of the arguments.

$$\frac{\langle \mathtt{c}, D \rangle \to \langle \mathtt{C}, D \rangle \wedge \ \langle \mathtt{m}, C \rangle \to \langle (\mathtt{P}_1 \times \mathtt{P}_2 \times \cdots \times \mathtt{P}_n \to \mathtt{R}), C \rangle \wedge \ \langle \mathtt{p}_i, D \rangle \to \langle \mathtt{P}_i', D \rangle \wedge \ \mathtt{P}_i :> \mathtt{P}_i'}{\langle (\mathtt{c}.\mathtt{m}(...\mathtt{p}_i..)), D \rangle \to \langle \mathtt{R}, D \rangle}$$



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Defining our own subtypes



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Remember that we can define our own classes to extend the data types of the language when they are insufficient for our domain.

We can make use of subtyping for those custom classes.

This makes it possible to capture generalization relationships between our own data types.



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Can you think of a few examples of classes that generalize each other?



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Live code demo: list, empty, and Can you think of a few examples of classes that generalize each other?

Person generalizes Student.

LightEmitter generalizes Lamp.

Animal generalizes Dog.



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Which generalizes which?

Mercedes vs CarBrand?



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Which generalizes which?

Mercedes vs CarBrand? CarBrand :> Mercedes



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LivingSpace vs Apartment?



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LivingSpace vs Apartment? LivingSpace :> Apartment



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LivingSpace vs Apartment? LivingSpace :> Apartment

Cat vs Bird?



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LivingSpace vs Apartment? LivingSpace :> Apartment

Cat vs Bird? Neither.



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LivingSpace vs Apartment? LivingSpace :> Apartment

Cat vs Bird? Neither.

Person vs Employee?



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Mercedes vs CarBrand? CarBrand :> Mercedes

LivingSpace vs Apartment? LivingSpace :> Apartment

Cat vs Bird? Neither.

Person vs Employee? Person :> Employee



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LivingSpace vs Apartment? LivingSpace :> Apartment

Cat vs Bird? Neither.

Person vs Employee? Person :> Employee

Student vs Person?



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LivingSpace vs Apartment? LivingSpace :> Apartment

Cat vs Bird? Neither.

Person vs Employee? Person :> Employee

Student vs Person? Person :> Student.



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Cat vs Bird? Neither.

Person vs Employee? Person :> Employee

Student vs Person? Person :> Student.

Student vs Employee?



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Cat vs Bird? Neither.

Person vs Employee? Person :> Employee

Student vs Person? Person :> Student.

Student vs Employee? Neither.



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LivingSpace vs Apartment? LivingSpace :> Apartment

Cat vs Bird? Neither.

Person vs Employee? Person :> Employee

Student vs Person? Person :> Student.

Student vs Employee? Neither.



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Consider, in particular, the following relationships:

Person :> Employee

Person :> Student



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Consider, in particular, the following relationships:

Person :> Employee

Person :> Student

We could imagine that Person, Employee, and Student are all classes

Moreover, Person and Employee are somehow related

Similarly, Person and Student are related in the same way



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Class inheritance

- This relationships is modeled, in Java/C#, with class inheritance
- A class such as Employee will therefore inherit from class Person
- This means that Employee will automatically have all fields and methods of Person



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Live code demo: list, empty, and When we use inheritance, for example of Person from Employee, then the language automatically infers that Person :> Employee

This means therefore that anywhere in the language (a variable, a parameter, etc.) where we expected a Person, we can give an Employee.

This provides polymorphism, as the same data type (in this case Person) can have multiple shapes: a Person, an Employee, a Student. etc.



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Inheritance requires very little code: in C#, a colon (:) suffices with the name of the inherited class next to that of the defined class.

In Java, we use the keyword extends instead of the colon.

It is possible to inherit at most one class.



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The program below works: why?

```
class A {
  public int M(int x) {
    return (x + x);
  }
}
class B : A {
  public int N(int y) {
    return (y * 10);
  }
}
A b = new B();
Console WriteLine(b.M(5));
```



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node

The program below works: why?

```
class A {
 public int M(int x) {
    return (x + x);
class B : A {
 public int N(int v) {
    return (y * 10);
A b = new B();
Console.WriteLine(b.M(5));
```

Because the declaration of b specifies A as the type, but whenever we expect an A we can use a B thanks to inheritance.



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Live code demo: list, empty, and node. Which in Java then becomes:

```
class A {
  public int M(int x) {
    return (x + x);
  }
}
class B extends A {
  public int N(int y) {
    return (y * 10);
  }
}
A b = new B();
System.out.println(b.M(5));
```



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```
class A {
   public A() {
   }
   public int M(int x) {
      return (x + x);
   }
}
class B : A {
   public B() {
   }
   public int N(int y) {
      return (y * 10);
   }
}
A b = new B();
Console.WriteLine(b.M(5));
```

```
Declarations:
```



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```
class A {
  public A() {
  public int M(int x) {
    return (x + x);
class B : A {
  public B() {
  public int N(int y) {
    return (y * 10);
A b = new B();
Console.WriteLine(b.M(5)):
```

```
PC
Declarations:
                       A=A \rightarrow A
Classes:
                 M=(A \times int) \rightarrow int
```



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```
class A {
  public A() {
  public int M(int x) {
    return (x + x);
class B : A {
  public B() {
  public int N(int y) {
    return (y * 10);
A b = new B();
Console.WriteLine(b.M(5));
```

```
Declarations:
```

Α	В
$\begin{array}{c} A {=} A \to A \\ M {=} (A {\times} int) \to int \end{array}$	$\begin{array}{c} B {=} B \to B \\ M {=} (A {\times} int) \to int \\ N {=} (B {\times} int) \to int \end{array}$



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```
class A {
  public A() {
  public int M(int x) {
    return (x + x);
class B : A {
  public B() {
  public int N(int v) {
    return (y * 10);
A b = new B();
Console.WriteLine(b.M(5));
```

```
PC
Declarations:
                16
                       Α
```

```
B
                                                    B=B \rightarrow B
      A=A \rightarrow A
                                              M=(A \times int) \rightarrow int
M=(A \times int) \rightarrow int
                                              N = (B \times int) \rightarrow int
```



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```
class A {
  public A() {
  public int M(int x) {
    return (x + x);
class B : A {
  public B() {
  public int N(int y) {
    return (y * 10);
A b = new B():
Console.WriteLine(b.M(5));
```

 $N = (B \times int) \rightarrow int$

Declarati	one:	L D		٠. د	16		aigl	LIIIS	
Declarati	Oli 5.	Α		16	nu		int	Α	
	A				В				
Classes:	:s: A=A → A				$B=B \rightarrow B$				
Classes.	١,			→ A) → int			$M=(A \times$		
	"	/I—(A	\wedge IIII	<i>)</i> \rightarrow 1111			M (D)		



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```
class A {
  public A() {
  public int M(int x) {
    return (x + x);
class B : A {
  public B() {
  public int N(int v) {
    return (y * 10);
A b = new B();
Console.WriteLine(b.M(5));
```

```
PC
                                                                      this
                                                ret
                                                           arg_1
Declarations:
                                      16
                                                int
                                                            int
                                                                        Α
                                                                B=B \rightarrow B
Classes:
                        A=A \rightarrow A
                                                          M=(A \times int) \rightarrow int
                  M=(A \times int) \rightarrow int
                                                          N=(B\times int) \rightarrow int
```



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```
class A {
  public A() {
  public int M(int x) {
    return (x + x);
class B : A {
  public B() {
  public int N(int y) {
    return (y * 10);
A b = new B();
Console.WriteLine(b.M(5));
```

```
Declarations:
                17
```

Α	В
$\begin{array}{c} A = A \rightarrow A \\ M = (A \times int) \rightarrow int \end{array}$	$\begin{array}{c} B {=} B \to B \\ M {=} (A {\times} int) \to int \\ N {=} (B {\times} int) \to int \end{array}$



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```
class A {
  public A() {
  }
  public int M(int x) {
    return (x + x);
  }
}
class B : A {
  public B() {
    }
  public int N(int y) {
     return (y * 10);
  }
}
A b = new B();
Console.WriteLine(b.M(5));
```

```
Stack: PC
```



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Interfaces and polymorphism

```
class A {
  public A() {
  }
  public int M(int x) {
    return (x + x);
  }
}
class B : A {
  public B() {
  }
  public int N(int y) {
    return (y * 10);
  }
}
A b = new B();
Console.WriteLine(b.M(5));
```

```
Stack: PC 15
Heap: 1
```



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Interfaces and polymorphism

```
class A {
   public A() {
   }
   public int M(int x) {
      return (x + x);
   }
}
class B : A {
   public B() {
   }
   public int N(int y) {
      return (y * 10);
   }
}
A b = new B();
Console.WriteLine(b.M(5));
```

Stack:	PC		PC	ret	this
Stack:	15		10	null	ref 1
u	1				



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```
class A {
  public A() {
  public int M(int x) {
    return (x + x);
class B : A {
  public B() {
  public int N(int y) {
    return (y * 10);
A b = new B();
Console.WriteLine(b.M(5));
```

Stack:	PC		PC	ret
Stack.	15		10	ref 1
Heap:	1			



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Interfaces and polymorphism

```
class A {
   public A() {
   }
   public int M(int x) {
      return (x + x);
   }
} class B : A {
   public B() {
   }
   public int N(int y) {
      return (y * 10);
   }
} A b = new B();
Console. WriteLine(b.M(5));
```

```
        PC
        b

        16
        ref 1

        Heap:
        1
```



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Interfaces and polymorphism

```
class A {
  public A() {
  }
  public int M(int x) {
    return (x + x);
  }
}
class B : A {
  public B() {
  }
  public int N(int y) {
    return (y * 10);
  }
}
A b = new B();
Console.WriteLine(b.M(5));
```

Stack:	PC		PC	ret	this	×
Stack.	16		5	null	ref 1	5
Heap:	1					



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Interfaces and polymorphism

```
class A {
   public A() {
   }
   public int M(int x) {
      return (x + x);
   }
}
class B : A {
   public B() {
   }
   public int N(int y) {
      return (y * 10);
   }
}
A b = new B();
Console.WriteLine(b.M(5));
```

Stack:	PC		PC	ret
Stack:	16		5	10
[1			



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Interfaces and polymorphism

```
class A {
  public A() {
  }
  public int M(int x) {
    return (x + x);
  }
}
class B : A {
  public B() {
    }
  public int N(int y) {
    return (y * 10);
  }
}
A b = new B();
Console.WriteLine(b.M(5));
```

```
        Stack:
        PC
        b

        17
        ref 1

        Heap:
        1

        Output:
        10
```



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Class inheritance

Similarly, given a method that expects a parameter of type A could accept a parameter of type C.



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Interfaces and polymorphism

Live code demo: list, empty, and Inheritance does not make all combinations possible. For example, the program below does not work: why?

```
class A {
   public int M(int x) {
      return (x + x);
   }
}
class B : A {
   public int N(int y) {
      return (y * 10);
   }
}
class C : A {
   public int O(int z) {
      return (z - 1);
   }
}
B b = new C();
Console.WriteLine(b.M(5));
```



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demo: list. empty, and node

Live code

Inheritance does not make all combinations possible. For example, the program below does not work: why?

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (y * 10);
class C : A {
  public int O(int z) {
    return (z - 1);
B b = new C():
Console.WriteLine(b.M(5));
```

There is no inheritance relationship between B and C!



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Live code demo: list, empty, and node. Which in Java then becomes:

```
class A {
  public int M(int x) {
    return (x + x);
  }
}
class B extends A {
  public int N(int y) {
    return (y * 10);
  }
}
class C extends A {
  public int D(int z) {
    return (z - 1);
  }
}
B b = new C();
System.out.println(b.M(5));
```



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Class inheritance

The subtyping relationship is transitive. This means that given X <: Y and Y <: Z, implies X <: Z.



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Interfaces and polymor-16 phism

Live code demo: list. empty, and

node.

Inheritance can also perform multiple conversion steps:

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (v * 10);
class C : B {
  public int O(int z) {
    return (z - 1):
A a = new C():
Console.WriteLine(a.M(5));
```



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```
class A {
  public int M(int x) {
    return (x + x);
  }
}
class B extends A {
  public int N(int y) {
    return (y * 10);
  }
}
class C extends B {
  public int O(int z) {
    return (z - 1);
  }
}
A a = new C();
System.out.println(a.M(5));
```



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Live code demo: list. empty, and

node.

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (v * 10);
class C : B {
  public int O(int z) {
    return (z - 1);
A a = new C();
Console.WriteLine(a.M(5)):
```

```
Declarations:
```



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```
class A {
  public int M(int x) {
    return (x + x);
  }
}
class B : A {
  public int N(int y) {
    return (y * 10);
  }
}
class C : B {
  public int O(int z) {
    return (z - 1);
  }
}
A a = new C();
Console WriteLine(a.M(5));
```

```
\begin{array}{c|c} \text{Declarations:} & \hline {PC} \\ \hline 5 \\ \\ \text{Classes:} & \hline {A} \\ \hline {M=(A\times int) \rightarrow int} \\ \end{array}
```



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Interfaces and polymorphism

Live code demo: list, empty, and node.

```
class A {
  public int M(int x) {
    return (x + x);
  }
}
class B : A {
  public int N(int y) {
    return (y * 10);
  }
}
class C : B {
  public int O(int z) {
    return (z - 1);
  }
}
A a = new C();
Console.WriteLine(a.M(5));
```

```
Declarations: PC 10
```

Α	В
$M=(A \times int) \rightarrow int$	$M=(A \times int) \rightarrow int$ $N=(B \times int) \rightarrow int$



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Live code demo: list. empty, and

node.

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (y * 10);
class C : B {
  public int O(int z) {
    return (z - 1):
A a = new C();
Console.WriteLine(a.M(5)):
```

PC Declarations: 15

A	В	C
$M=(A \times int) \rightarrow int$	$\begin{array}{c} M{=}(A{\times}int) \to int \\ N{=}(B{\times}int) \to int \end{array}$	$\begin{array}{c} M{=}(A{\times}int) \to int \\ N{=}(B{\times}int) \to int \\ O{=}(C{\times}int) \to int \end{array}$



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Interfaces and polymorphism

Live code demo: list. empty, and

node.

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (y * 10);
class C : B {
  public int O(int z) {
    return (z - 1):
A a = new C();
Console.WriteLine(a.M(5)):
```

Declarations: 17

Α	В	C
$M=(A \times int) \rightarrow int$	$\begin{array}{c} M{=}(A{\times}int) \to int \\ N{=}(B{\times}int) \to int \end{array}$	$\begin{array}{c} M{=}(A{\times}int) \to int \\ N{=}(B{\times}int) \to int \\ O{=}(C{\times}int) \to int \end{array}$



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Live code demo: list, empty, and node. Classes

```
class A {
  public int M(int x) {
    return (x + x);
  }
}
class B : A {
  public int N(int y) {
    return (y * 10);
  }
}
class C : B {
  public int O(int z) {
    return (z - 1);
  }
}
A a = new C();
Console.WriteLine(a.M(5));
```

Declarations:	а	PC	ret	arg 1	this
	Α	17	null	int	Α

Α	В	С
$M=(A \times int) \rightarrow int$	$\begin{array}{c} M{=}(A{\times}int) \to int \\ N{=}(B{\times}int) \to int \end{array}$	$ \begin{array}{c} M{=}(A{\times}int) \to int \\ N{=}(B{\times}int) \to int \\ O{=}(C{\times}int) \to int \end{array} $



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Interfaces and polymorphism

Live code demo: list. empty, and

node.

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (y * 10);
class C : B {
  public int O(int z) {
    return (z - 1):
A a = new C();
Console.WriteLine(a.M(5)):
```

Declarations:	а	PC	ret	${\sf arg}_1$	this
	Α	17	int	int	Α

Α	В	C
M=(A imes int) o int	$\begin{array}{c} M{=}(A{\times}int) \to int \\ N{=}(B{\times}int) \to int \end{array}$	$\begin{array}{c} M{=}(A{\times}int) \to int \\ N{=}(B{\times}int) \to int \\ O{=}(C{\times}int) \to int \end{array}$



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node.

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (y * 10);
class C : B {
  public int O(int z) {
    return (z - 1):
A a = new C();
Console.WriteLine(a.M(5)):
```

```
PC
Declarations:
               18
```

Classes:

A	В	C
$M=(A \times int) \rightarrow int$	$\begin{array}{c} M{=}(A{\times}int) \to int \\ N{=}(B{\times}int) \to int \end{array}$	$\begin{array}{c} M{=}(A{\times}int) \to int \\ N{=}(B{\times}int) \to int \\ O{=}(C{\times}int) \to int \end{array}$



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Interfaces and polymorphism

Live code demo: list, empty, and node.

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (v * 10);
}
class C : B {
  public C() {
  public int O(int z) {
    return (z - 1):
A a = new C();
Console.WriteLine(a.M(5)):
```

Stack: PC 1



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Live code demo: list. empty, and

```
node.
```

Co61-1/41/09

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (y * 10);
}
class C : B {
  public C() {
  public int O(int z) {
    return (z - 1):
A a = new C():
Console.WriteLine(a.M(5));
```

```
Stack:
Heap:
```



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Live code demo: list. empty, and

node.

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (y * 10);
class C : B {
  public C() {
  public int O(int z) {
    return (z - 1):
A a = new C():
Console.WriteLine(a.M(5));
```

Stack:	PC		PC	ret	this
Stack.	18		13		ref 1
Heap:	1				



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Live code demo: list. empty, and

node.

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (y * 10);
}
class C : B {
  public C() {
  public int O(int z) {
    return (z - 1):
A a = new C():
Console.WriteLine(a.M(5));
```

```
PC
                                    ret
Stack:
          18
                            13
                                   ref 1
Heap:
```



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Interfaces and polymorphism

Live code demo: list. empty, and

```
node.
```

Co64/41/09

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (y * 10);
}
class C : B {
  public C() {
  public int O(int z) {
    return (z - 1):
A a = new C():
Console.WriteLine(a.M(5));
```



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node.

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (y * 10);
class C : B {
  public C() {
  public int O(int z) {
    return (z - 1):
A a = new C():
Console.WriteLine(a.M(5));
```

Stack: F	PC		PC	ret	this	х
	19		3	null	ref 1	5
Неар:	1					



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node.

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (y * 10);
}
class C : B {
  public C() {
  public int O(int z) {
    return (z - 1):
A a = new C():
Console.WriteLine(a.M(5));
```



```
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```

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node. Contiluiñon

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (y * 10);
class C : B {
  public C() {
  public int O(int z) {
    return (z - 1);
A a = new C();
Console.WriteLine(a.M(5));
```

```
        PC
        a

        20
        ref 1

Heap: 1
Output: 10
```



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Live code demo: person, employee, and student/animal, dog, cat/...

Interfaces 15 and polymor-16 phism 17

Live code demo: list, empty, and Of course, when inheriting, we can still use all methods available given a variable type.

The code below does indeed work. Why?

```
class A {
   public int M(int x) {
      return (x + x);
   }
}
class B : A {
   public int N(int y) {
      return (y * 10);
   }
}
class C : B {
   public int O(int z) {
      return (z - 1);
   }
}
B b = new C();
Console.WriteLine(b.M(5));
Console.WriteLine(b.N(5));
```



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Interfaces 15 and polymor-16 phism 17

Live code demo: list, empty, and node.

C0681/13/09

Of course, when inheriting, we can still use all methods available given a variable type.

The code below does indeed work. Why?

```
class A {
  public int M(int x) {
    return (x + x);
  }
}
class B : A {
  public int N(int y) {
    return (y * 10);
  }
}
class C : B {
  public int O(int z) {
    return (z - 1);
  }
}
B b = new C();
Console.WriteLine(b.M(5));
Console.WriteLine(b.N(5));
```

It is possible to call both methods ${\tt M}$ and ${\tt N}$ on an instance of B.



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Which in Java then becomes:

```
class A {
  public int M(int x) {
    return (x + x);
class B extends A {
  public int N(int v) {
    return (y * 10);
class C extends B {
  public int O(int z) {
    return (z - 1);
B b = new C():
System.out.println(b.M(5));
System.out.println(b.N(5));
```



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node.

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (y * 10);
class C : B {
  public int O(int z) {
    return (z - 1):
B b = new C():
Console.WriteLine(b.M(5));
Console.WriteLine(b.N(5));
```

Declarations:



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```
class A {
  public int M(int x) {
    return (x + x);
  }
}
class B : A {
  public int N(int y) {
    return (y * 10);
  }
}
class C : B {
  public int O(int z) {
    return (z - 1);
  }
}
B b = new C();
Console.WriteLine(b.M(5));
Console.WriteLine(b.M(5));
```

```
\begin{array}{c} \text{Declarations: } \boxed{\frac{\text{PC}}{5}} \\ \text{Classes: } \boxed{\frac{\text{A}}{\text{M=}(\text{A}\times\text{int}) \rightarrow \text{int}}} \end{array}
```



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```
node.
```

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (y * 10);
class C : B {
  public int O(int z) {
    return (z - 1):
B b = new C();
Console.WriteLine(b.M(5)):
Console.WriteLine(b.N(5));
```

Declarations:

Classes:

	А	В
:	$M=(A \times int) \rightarrow int$	$\begin{array}{c} M{=}(A{\times}int) \to int \\ N{=}(B{\times}int) \to int \end{array}$



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```
node.
```

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (y * 10);
class C : B {
  public int O(int z) {
    return (z - 1):
B b = new C():
Console.WriteLine(b.M(5));
Console.WriteLine(b.N(5));
```

Declarations:

	Α	В	С
Classes:	$M=(A \times int) \rightarrow int$	$\begin{array}{c} M{=}(A{\times}int) \to int \\ N{=}(B{\times}int) \to int \end{array}$	$\begin{array}{c} M{=}(A{\times}int) \to int \\ N{=}(B{\times}int) \to int \\ O{=}(C{\times}int) \to int \end{array}$



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```
node.
```

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (v * 10);
class C : B {
  public int O(int z) {
    return (z - 1):
B b = new C():
Console.WriteLine(b.M(5));
Console.WriteLine(b.N(5));
```

Declarations:	PC	Ь	
	17	В	

В Α $M=(A \times int) \rightarrow int$ Classes: $M=(A \times int) \rightarrow int$ $M=(A \times int) \rightarrow int$ $N=(B\times int) \rightarrow int$ $N=(B\times int) \rightarrow int$ $O=(C\times int) \rightarrow int$



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node.

Classes:

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (y * 10);
class C : B {
  public int O(int z) {
    return (z - 1):
B b = new C():
Console.WriteLine(b.M(5));
Console.WriteLine(b.N(5));
```

	ь	11	IIIII	IIIC	_ D
Declarations:		17	null	int	0
	Ь	PC	ret	arg ₁	this

Α В $M=(A \times int) \rightarrow int$ $M=(A \times int) \rightarrow int$ $M=(A \times int) \rightarrow int$ $N=(B\times int) \rightarrow int$ $N=(B\times int) \rightarrow int$ $O=(C\times int) \rightarrow int$



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```

Classe

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (v * 10);
class C : B {
  public int O(int z) {
    return (z - 1):
B b = new C():
Console.WriteLine(b.M(5));
Console.WriteLine(b.N(5));
```

Declarations:	b	PC	ret	arg 1	this
	В	17	int	int	В

	A	В	С
es:	$M=(A \times int) \rightarrow int$	$\begin{array}{c} M{=}(A{\times}int) \to int \\ N{=}(B{\times}int) \to int \end{array}$	$\begin{array}{c} M{=}(A{\times}int) \to int \\ N{=}(B{\times}int) \to int \\ O{=}(C{\times}int) \to int \end{array}$



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```
class A {
  public int M(int x) {
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  public int N(int y) {
    return (y * 10);
class C : B {
  public int O(int z) {
    return (z - 1):
B b = new C():
Console.WriteLine(b.M(5));
Console.WriteLine(b.N(5));
```

В

	10	. Б	1
Declarations:	PC	b	1

Classes: $M=(A \times int) \rightarrow int$ $M=(A \times int) \rightarrow int$ $N=(B\times int) \rightarrow int$

 $M=(A \times int) \rightarrow int$



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```

```
class A {
  public int M(int x) {
    return (x + x);
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    return (y * 10);
class C : B {
  public int O(int z) {
    return (z - 1):
B b = new C():
Console.WriteLine(b.M(5));
Console.WriteLine(b.N(5));
```

	" В	18	null	int	В
Declarations:	. Ь	PC	ret	arg ₁	this

Α В $M=(A \times int) \rightarrow int$ Classes: $M=(A \times int) \rightarrow int$ $M=(A \times int) \rightarrow int$ $N=(B\times int) \rightarrow int$ $N=(B\times int) \rightarrow int$ $O=(C\times int) \rightarrow int$



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```
class A {
  public int M(int x) {
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  public int N(int y) {
    return (v * 10);
class C : B {
  public int O(int z) {
    return (z - 1):
B b = new C():
Console.WriteLine(b.M(5));
Console.WriteLine(b.N(5));
```

Declarations:	Ь	PC	ret	arg_1	this
	В	18	int	int	В

Γ	Α	В	С
	$M{=}(A{\times}int)\toint$	$\begin{array}{l} M{=}(A{\times}int) \to int \\ N{=}(B{\times}int) \to int \end{array}$	$\begin{array}{c} M{=}(A{\times}int) \to int \\ N{=}(B{\times}int) \to int \\ O{=}(C{\times}int) \to int \end{array}$

Classes:



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```

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class A {
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  public int N(int y) {
    return (y * 10);
class C : B {
  public int O(int z) {
    return (z - 1):
B b = new C():
Console.WriteLine(b.M(5));
Console.WriteLine(b.N(5));
```

		Δ
Declarations:	10	B
_	PC	Ь

В $M=(A \times int) \rightarrow int$ Classes: $M=(A \times int) \rightarrow int$ $M=(A \times int) \rightarrow int$ $N=(B\times int) \rightarrow int$ $N=(B\times int) \rightarrow int$ $O=(C\times int) \rightarrow int$



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```
class A {
  public int M(int x) {
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  public int N(int y) {
    return (y * 10);
class C : B {
  public C() {
  public int O(int z) {
    return (z - 1):
B b = new C():
Console.WriteLine(b.M(5));
Console.WriteLine(b.N(5));
```

```
Stack: PC 1
```



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```
node.
```

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (y * 10);
class C : B {
  public C() {
  public int O(int z) {
    return (z - 1);
B b = new C();
Console.WriteLine(b.M(5));
Console.WriteLine(b.N(5)):
```



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```

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```
node.
```

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (y * 10);
class C : B {
  public C() {
  public int O(int z) {
    return (z - 1);
B b = new C();
Console.WriteLine(b.M(5));
Console.WriteLine(b.N(5)):
```

Stack:	PC		PC	ret	this
Stack.	18		13	null	ref 1
Heap:	1				



```
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```

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```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (y * 10);
class C : B {
  public C() {
  public int O(int z) {
    return (z - 1);
B b = new C();
Console.WriteLine(b.M(5));
Console.WriteLine(b.N(5)):
```

Stack:	PC		PC	ret
	18		13	ref 1
Heap:	1			



```
Polymorphism
```

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```
node.
```

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (y * 10);
class C : B {
  public C() {
  public int O(int z) {
    return (z - 1);
B b = new C();
Console.WriteLine(b.M(5));
Console.WriteLine(b.N(5)):
```

```
Stack:
          19
                 ref 1
Heap:
```



```
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```

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```
node.
```

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (y * 10);
class C : B {
  public C() {
  public int O(int z) {
    return (z - 1);
B b = new C();
Console.WriteLine(b.M(5));
Console.WriteLine(b.N(5)):
```

Stack:	PC		PC	ret	this	х
Stack.	19		3	null	ref 1	5
Heap:	1					



```
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```

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```
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```

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (y * 10);
class C : B {
  public C() {
  public int O(int z) {
    return (z - 1);
B b = new C();
Console.WriteLine(b.M(5));
Console.WriteLine(b.N(5)):
```

Stack:	PC 19		PC 3	ret 10
Неар:	1			



```
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```

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```
node.
```

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (v * 10);
class C : B {
  public C() {
  public int O(int z) {
    return (z - 1):
B b = new C():
Console.WriteLine(b.M(5)):
Console.WriteLine(b.N(5));
```

```
Stack:
         20
               ref 1
Heap:
Output:
```



```
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```

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```
node.
```

Output: 10

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (v * 10);
class C : B {
  public C() {
  public int O(int z) {
    return (z - 1):
B b = new C();
Console.WriteLine(b.M(5)):
Console.WriteLine(b.N(5));
```

Stack:	PC		PC	ret	this	У
Stack.	20		8	null	ref 1	5
Heap:	1					



```
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```

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```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (v * 10);
class C : B {
  public C() {
  public int O(int z) {
    return (z - 1):
B b = new C():
Console.WriteLine(b.M(5)):
Console.WriteLine(b.N(5));
```

```
Stack: PC ... PC ret 20 ... 8 50

Heap: Output: 10
```



```
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```

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```
Stack:
                  Heap:
node.
```

```
class A {
  public int M(int x) {
    return (x + x);
class B : A {
  public int N(int y) {
    return (v * 10);
class C : B {
  public C() {
  public int O(int z) {
    return (z - 1):
B b = new C():
Console.WriteLine(b.M(5)):
Console.WriteLine(b.N(5));
```

```
21
               ref 1
Output:
                 50
```



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and polymor- 14
phism 15

Live code demo: list, empty, and A major difference with Python is that, even if the instance may allow calling some methods, subtyping might disallow it.

The code below does not work. Why?

```
class A {
  public int M(int x) {
    return (x + x);
  }
}
class B : A {
  public int N(int y) {
    return (y * 10);
  }
}
class C : B {
  public int O(int z) {
    return (z - 1);
  }
}
B b = new C();
Console.WriteLine(b.0(5));
```



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A major difference with Python is that, even if the instance may allow calling some methods, subtyping might disallow it.

The code below does not work. Why?

```
class A {
   public int M(int x) {
      return (x + x);
   }
} class B : A {
   public int N(int y) {
      return (y * 10);
   }
} class C : B {
   public int O(int z) {
      return (z - 1);
   }
} b = new C();
Console.WriteLine(b.0(5));
```

b is declared with type B, which has no method O.



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Which in Java then becomes:

```
class A {
  public int M(int x) {
    return (x + x);
  }
}
class B extends A {
  public int N(int y) {
    return (y * 10);
  }
}
class C extends B {
  public int O(int z) {
    return (z - 1);
  }
}
B b = new C();
System.out.println(b.0(5));
```



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Interfaces and polymorphism



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Zero-level data types

- The subtyping relationship can start from a data type so general that it has no concrete implementation.
- This data type is called an interface.
- Interfaces are classes defined with the keyword interface. They have no fields, and no implementation of their methods.
- We say that a class implements, not inherits from, one or more interfaces.



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Zero-level data types

- Interfaces are especially useful to specify what requirements a data type must satisfy to be used in a context.
- With interfaces we are not bound to also giving a "default" implementation.



that of the defined class.

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Implementing interfaces requires very little code: in C#, a colon (:) suffices with the name of the implemented interface next to

In Java, we use the keyword implements instead of the colon.

It is possible to implement multiple interfaces from the same class.



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C 0851/13108

The program below works: why?

```
interface A {
   int M(int x);
}
class B : A {
   public int M(int x) {
      return (x + x);
   }
}
A b = new B();
Console.WriteLine(b.M(5));
```



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The program below works: why?

```
interface A {
  int M(int x);
class B : A {
  public int M(int x) {
    return (x + x);
A b = new B():
Console.WriteLine(b.M(5)):
```

Because the declaration of b specifies A as the type, but whenever we expect an A we can use a B thanks to the subtyping of implemented interfaces.



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College College

Which in Java then becomes:

```
interface A {
   int M(int x);
}
class B implements A {
   public int M(int x) {
      return (x + x);
   }
}
A b = new B();
System.out.println(b.M(5));
```



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Coh@luśib@9

```
The program below does not work: why?
```

```
interface A {
  int M(int x);
}
A a = new A();
Console.WriteLine(a.M(5));
```



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The program below does not work: why?

```
interface A {
  int M(int x);
A a = new A():
Console.WriteLine(a.M(5));
```

Because A has no implementation and so cannot be instantiated: what code could we possibly execute for method M?



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Coll@2ukib09

```
interfaces and polymorphism
```

Which in Java then becomes:

```
interface A {
  int M(int x);
}
A a = new A();
System.out.println(a.M(5));
```



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Zero-level data types

Polymorphism can be used in a lot of contexts, as long as the conversion we expect of the language is provably safe.



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Interfaces and polymorphism

Live code demo: list, empty, and node. The program below works: why?

```
interface A {
   A M(A x);
}
class B : A {
   public A M(A x) {
      return this;
   }
}
A b = new B();
Console.WriteLine(b.M(new B()));
```



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The program below works: why?

```
interface A {
  A M(A x);
class B : A {
  public A M(A x) {
    return this:
A b = new B():
Console.WriteLine(b.M(new B()));
```

Because the argument to M, which should be an A, can safely accept a B as well.



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Coll@Susib@9

Which in Java then becomes:

```
interface A {
   A M(A x);
}
class B implements A {
   public A M(A x) {
      return this;
   }
}
A b = new B();
System.out.println(b.M(new B()));
```



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Conclusion

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Interfaces and polymorphism

Live code demo: list, empty, and Looking back

- Polymorphism makes it possible to pass different data types to other contexts, as long as the conversion is safe
- Inheritance is the basic mechanism of polymorphism
- Interfaces make this even more powerful by allowing the use of polymorphism without a concrete data type



This is it!

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The best of luck, and thanks for the attention!