Key Reinstallation Attacks: Breaking the WPA2 Protocol

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Introduction



PhD Defense, July 2016:

"You recommend WPA2 with AES, but are you sure that's secure?"

Seems so! No attacks in 14 years & proven secure.

ALOT OF BORING MATH LATER...

Introduction

```
/* install the PTK */
if ((*ic->ic_set_key)(ic, ni, k) != 0) {
        reason = IEEE80211_REASON_AUTH_LEAVE;
        goto deauth;
}
ni->ni_flags &= ~IEEE80211_NODE_TXRXPROT;
ni->ni_flags |= IEEE80211_NODE_RXPROT;
```



Key reinstallation when ic_set_key is called again?

Overview



Key reinstalls in 4-way handshake



Practical impact



Misconceptions



Lessons learned

Overview



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Misconceptions



Lessons learned

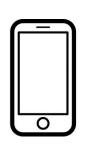
The 4-way handshake

Used to connect to any protected Wi-Fi network

- > Provides mutual authentication
- Negotiates fresh PTK: pairwise temporal key

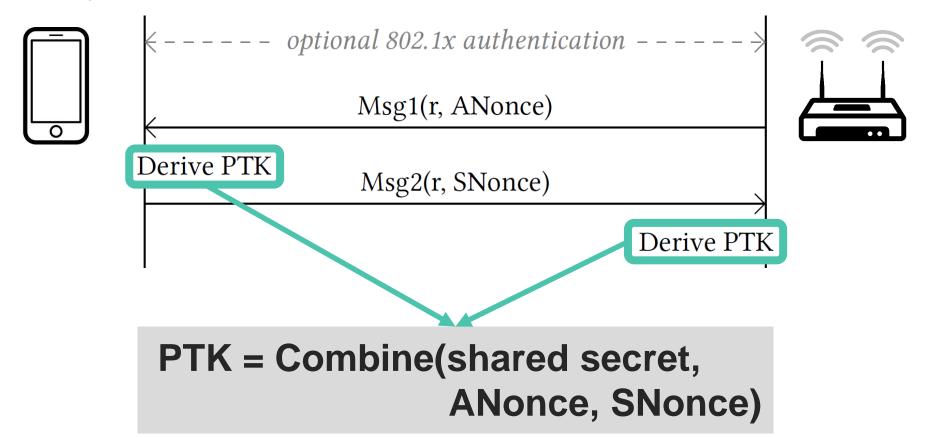
Appeared to be secure:

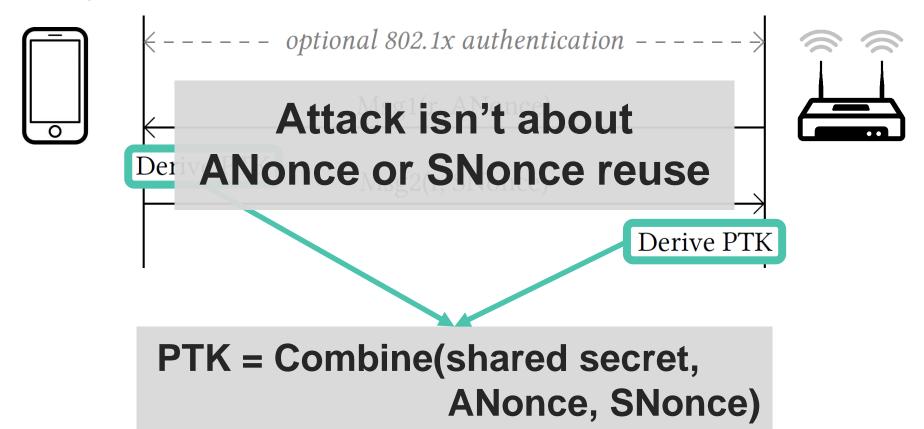
- No attacks in over a decade (apart from password guessing)
- Proven that negotiated key (PTK) is secret¹
- And encryption protocol proven secure⁷

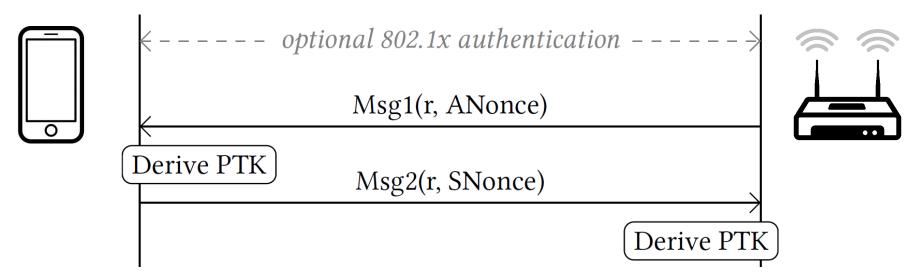


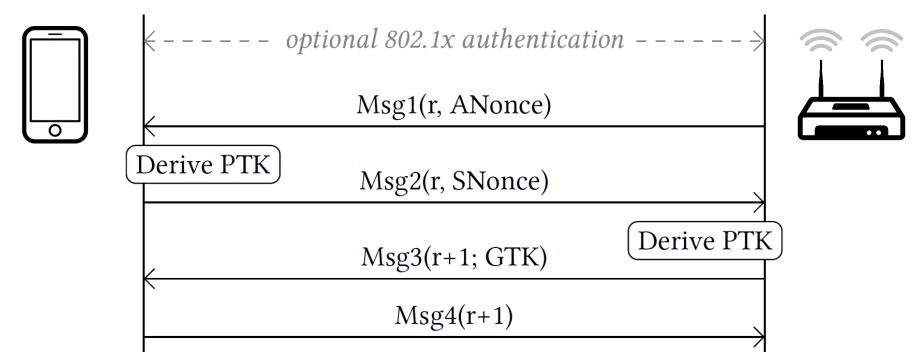
 $\langle -----$ optional 802.1x authentication ----- >

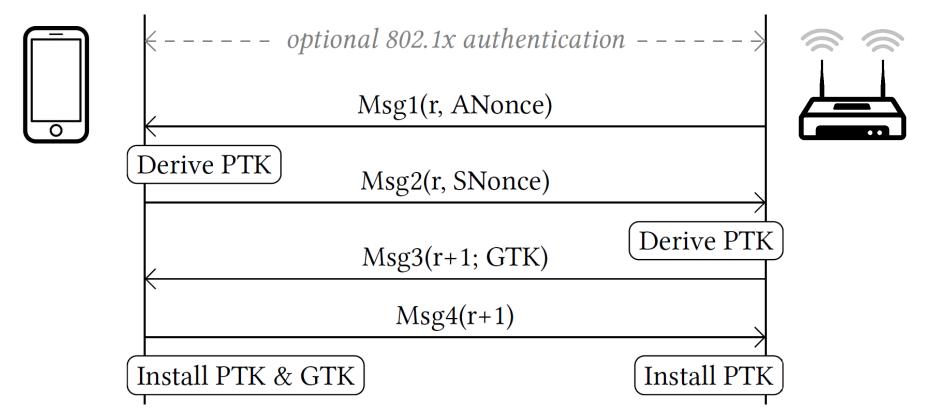


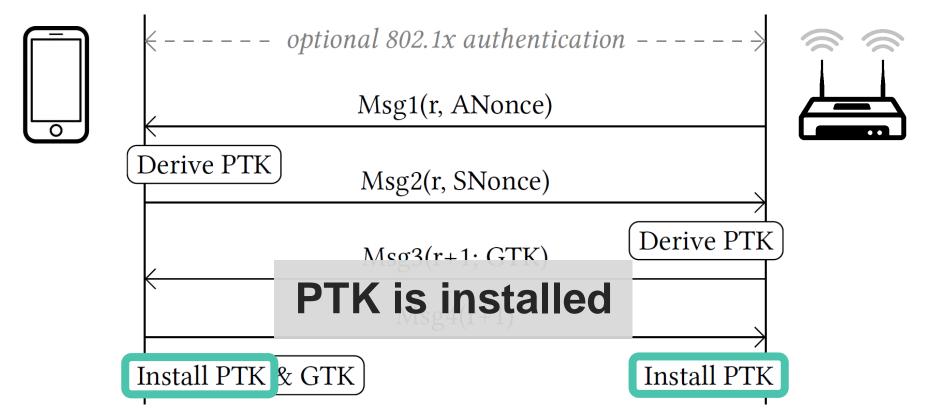


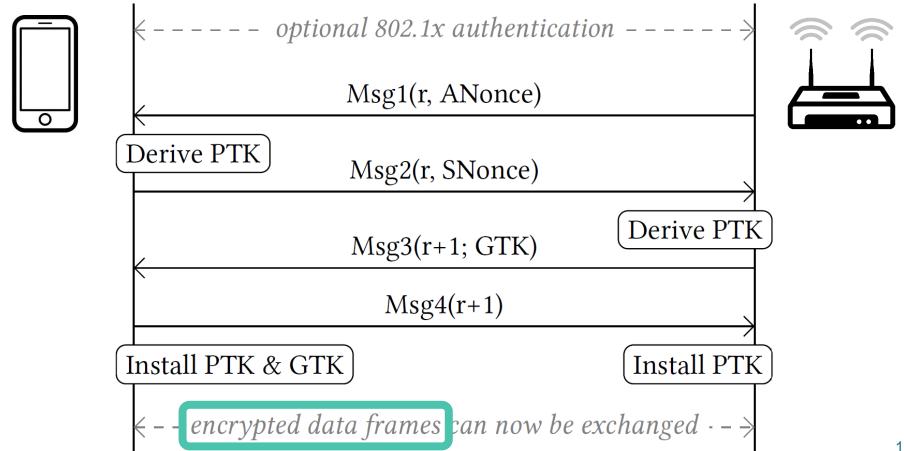




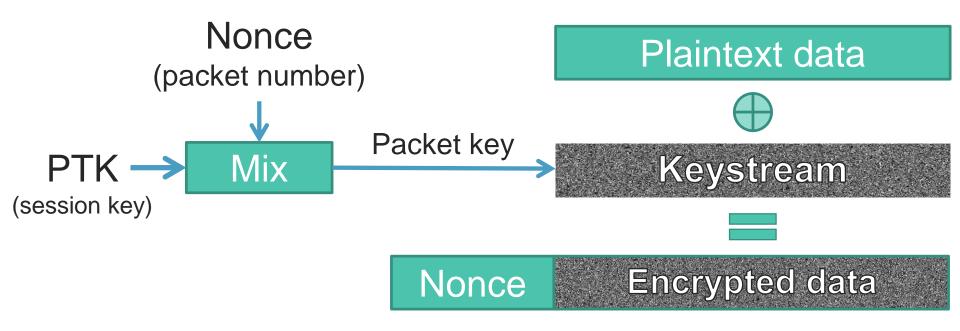




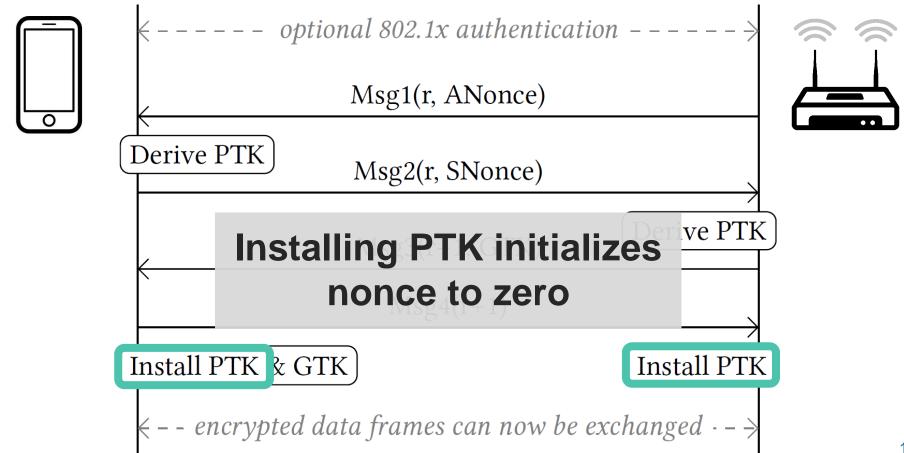




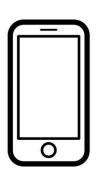
Frame encryption (simplified)



→ Nonce reuse implies keystream reuse (in all WPA2 ciphers)





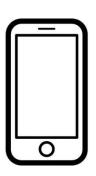


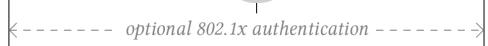
Channel 1

Channel 6



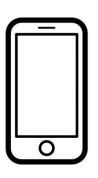








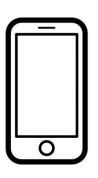


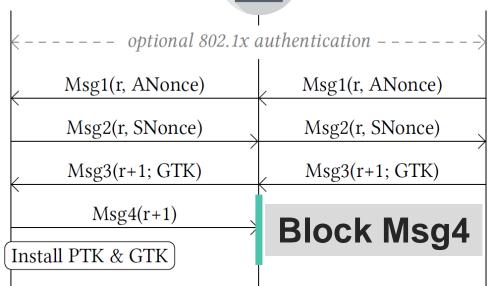


← optional 802.1x authentication >	
Msg1(r, ANonce)	Msg1(r, ANonce)
Msg2(r, SNonce)	Msg2(r, SNonce)
Msg3(r+1; GTK)	Msg3(r+1; GTK)



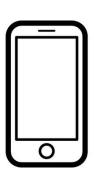


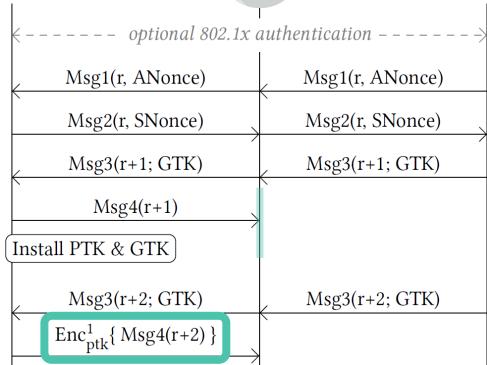






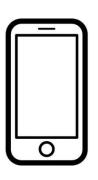


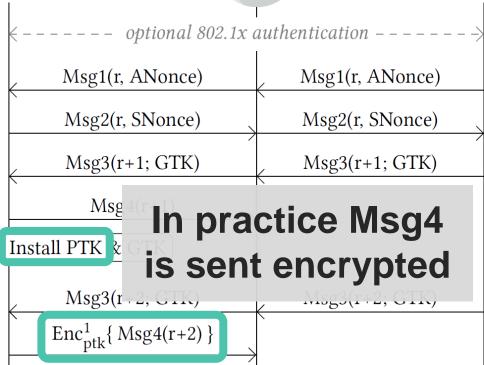






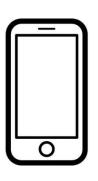


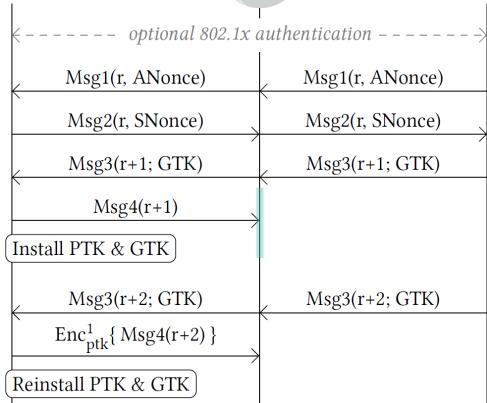






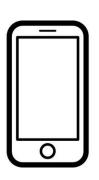


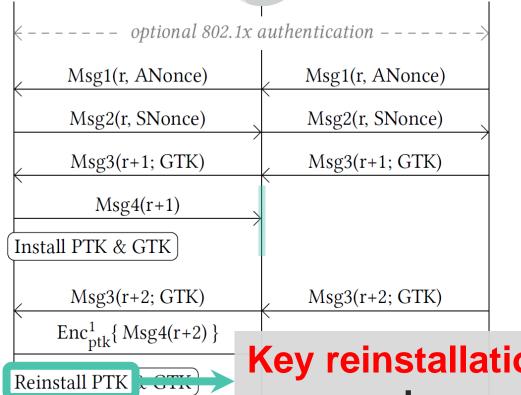








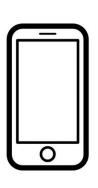


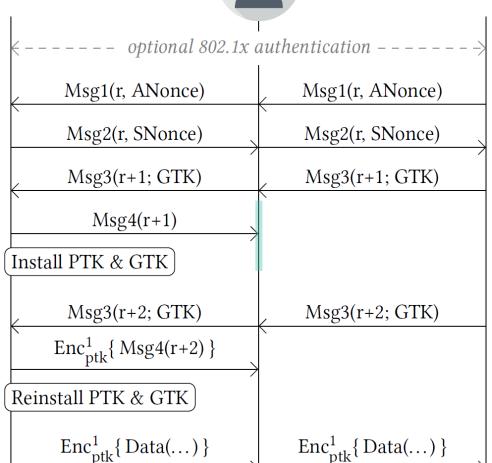




Key reinstallation! nonce is reset

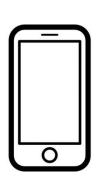


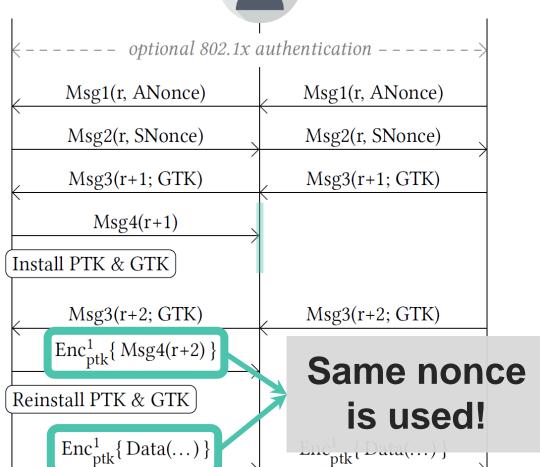






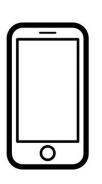


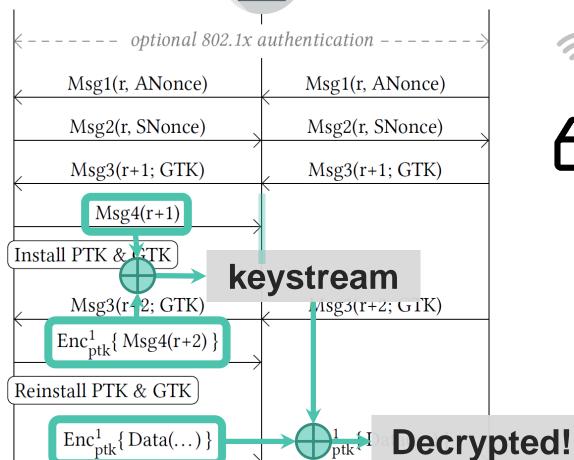






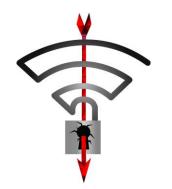








Overview



Key reinstalls in 4-way handshake



Practical impact

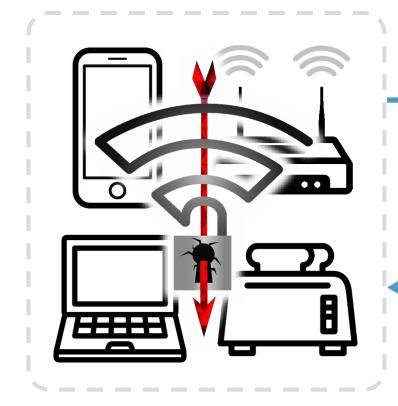


Misconceptions



Lessons learned

General impact



Transmit nonce reset

Decrypt frames sent by victim

Receive replay counter reset

Replay frames towards victim

Cipher suite specific

AES-CCMP: No practical frame forging attacks

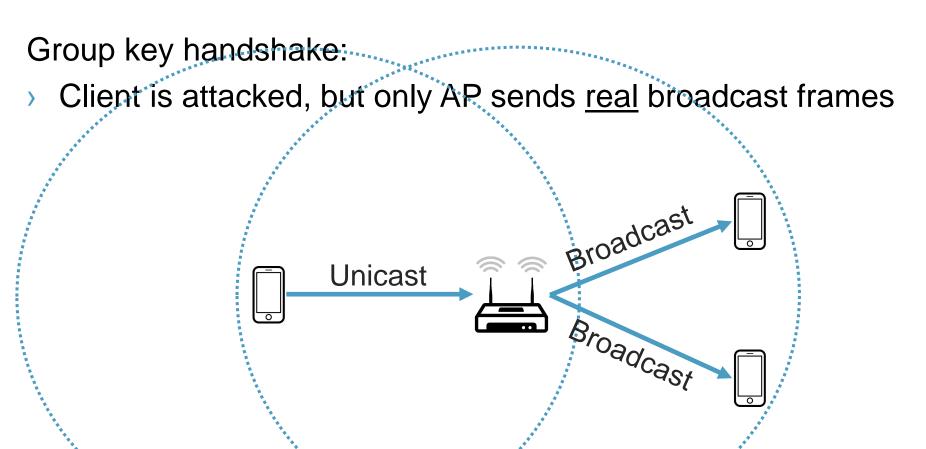
WPA-TKIP:

- Recover Message Integrity Check key from plaintext^{4,5}
- Forge/inject frames sent by the device under attack

GCMP (WiGig):

- Recover GHASH authentication key from nonce reuse⁶
- Forge/inject frames in both directions

Handshake specific



Handshake specific

Group key handshake:

- > Client is attacked, but only AP sends <u>real</u> broadcast frames
- Can only replay broadcast frames to client

4-way handshake:

Client is attacked → replay/decrypt/forge

FT handshake (fast roaming = 802.11r):

- Access Point is attacked → replay/decrypt/forge
- No MitM required, can keep causing nonce resets

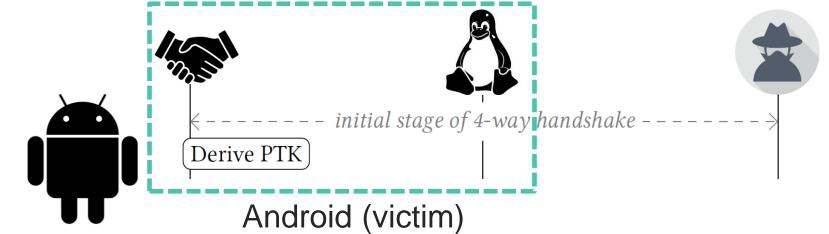
Implementation specific

iOS 10 and Windows: 4-way handshake not affected

- Cannot decrypt unicast traffic (nor replay/decrypt)
- > But group key handshake is affected (replay broadcast)
- Note: iOS 11 does have vulnerable 4-way handshake⁸

wpa_supplicant 2.4+

- Client used on Linux and Android 6.0+
- On retransmitted msg3 will install all-zero key





All-zero key installation







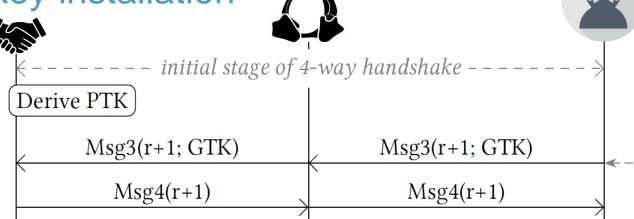










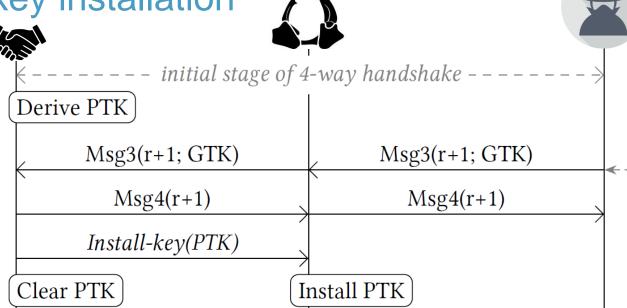










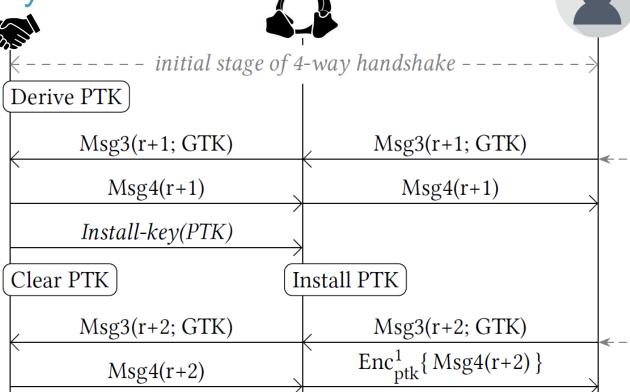










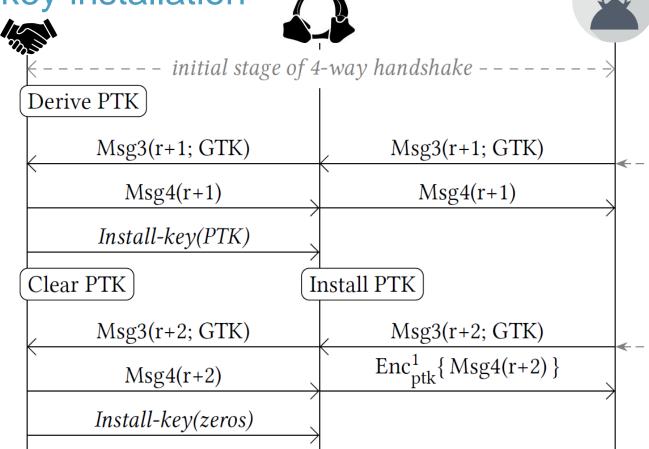












Install all-zero PTK



Reuse of ANonce and SNonce

- Hostapd reused ANonce during rekey
- macOS reused SNonce during rekey

On their own not exploitable, but when combined:

- Current key is again negotiated and (re)used
- And rekeying is recommended with WPA-TKIP

Is dedicated rekeying functionality worth it?

> Practical risks seem to outweigh advantages

Countermeasures

Problem: many clients won't get updates

Solution: AP can prevent (most) attacks on clients!

- Don't retransmit message 3/4
- Don't retransmit group message 1/2

However:

- Impact on reliability unclear
- Clients still vulnerable when connected to unmodified APs

Is your devices affected?

github.com/vanhoefm/krackattacks-scripts



- Test clients and APs
- Works on Kali Linux

Advice:

- Disable hardware encryption
- Use a supported Wi-Fi dongle!

Overview



Key reinstalls in 4-way handshake



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Lessons learned

Misconceptions I

Updating only the client or AP is sufficient

> Both <u>vulnerable</u> clients & <u>vulnerable</u> APs must apply patches

Need to be close to network and victim

Can use special antenna from afar



Must be connected to network as attacker (i.e. have password)

Only need to be nearby victim and network

Misconceptions II

No useful data is transmitted after handshake

Trigger new handshakes during TCP connection

Obtaining channel-based MitM is hard

Nope, can use channel switch announcements

Attack complexity is hard

- Script only needs to be written once ...
- ... and some are already doing this!

Misconceptions III

Using (AES-)CCMP mitigates the attack

Still allows decryption & replay of frames

Enterprise networks (802.1x) aren't affected

Also use 4-way handshake & are affected

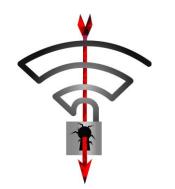
It's the end of the world!

> Let's not get carried away @



Image from "KRACK: Your Wi-Fi is no longer secure" by Kaspersky

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Key reinstalls in 4-way handshake



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Lessons learned

Limitations of formal proofs

- 4-way handshake proven secure
- Encryption protocol proven secure





The combination was not proven secure!

Model vs. implementation

Abstract model ≠ real code

Must assure code matches specification

The wpa_supplicant 2.6 case

- Complex state machine & turned out to still be vulnerable
- Need formal verification of implementations

Need rigorous specifications

Original WPA2 standard (802.11i amendment)

- State machine described in pseudo code
- Doesn't define when messages are accepted

```
StaProcessEAPOL-Key (S, M, A, I, K, RSC, ANonce, RSC, MIC, RSNE, GTK[N], IGTK[M], IPN)

if M = 1 then

if Check MIC(PTK, EAPOL-Key frame) fails then

State \leftarrow FAILED

else

State \leftarrow MICOK

endif

endif

if K = P then

if State \neq FAILED then
```

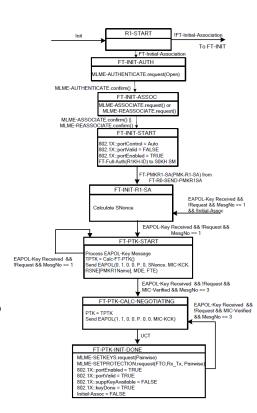
Need rigorous specifications

Original WPA2 standard (802.11i amendment)

- State machine described in pseudo code
- Doesn't define when messages are accepted

802.11r amendment (FT handshake)

- Better defines how/when to handle messages
- > But some terms and cases still unclear



S1KH state machine

On a related note...

Workshop on:

Security Protocol Implementations: Development and Analysis (SPIDA)

Co-located with EuroS&P 2018

"focuses on improving development & analysis of security protocols implementations"

Thank you!

Questions?

krackattacks.com

References

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