



WHO IS THIS GUY?

- hex0punk
- Father & husband
- Senior AppSec Consultant at nVisium
- Fan of Go and Rust
- Likes playing with binaries

WHAT ARE WE DOING?

- Learning about what makes go binaries different than C and C++ binaries
- Identifying techniques for recognizing and conducting analysis of go binaries.
- Tips for finding vulnerabilities in go binaries.
- Identifying common patterns found in go binaries
- Learning about protections that can be added to go binaries



ONCE UPON A TIME

- Found go binaries during an IoT assessment
- Wait, this isn't C or C++!
- This is Go!
- Now what?

THE GO ASSEMBLER

- The go compiler is based on the plan9 compiler
- Semi-abstract instruction set
- Pseudo-Assembly
- Not a direct representation of the underlying machine (i.e. a MOV may be a LD)
- It also introduces a set of pseudo registers

GO TOOL OBJDUMP go tool objdump -s <func> <bin>

```
TEXT main.main(SB) /Users/alexuseche/Projects/CTF-Source/go-bin/main.go
  main.go:11
                        0x4deaa0
                                                64488b0c25f8ffffff
                                                                         MOVQ FS:0xfffffff8, CX
 main.go:11
                                                                         LEAQ -0x50(SP), AX
                        0x4deaa9
                                                488d4424b0
 main.go:11
                        0x4deaae
                                                483b4110
                                                                         CMPQ 0x10(CX), AX
 main.go:11
                                                                         JBE 0x4ded08
                        0x4deab2
                                                0f8650020000
 main.go:11
                        0x4deab8
                                                4881ecd0000000
                                                                         SUBQ $0xd0, SP
 main.go:11
                        0x4deabf
                                                                         MOVQ BP, 0xc8(SP)
                                                4889ac24c8000000
 main.go:11
                                                                         LEAQ 0xc8(SP), BP
                        0x4deac7
                                                488dac24c8000000
 main.go:13
                                                                         XORPS X0, X0
                        0x4deacf
                                                0f57c0
 main.go:13
                        0x4dead2
                                                0f118424b8000000
                                                                         MOVUPS X0, 0xb8(SP)
```

RADARE2

```
; var int64_t var_8h @ rsp+0xc8
                                 mov rcx, qword fs:[0xffffffffffffff8]
 0x004deaa0
                  64488b0c25f8.
  0x004deaa9
                 488d4424b0
                                 lea rax, [rsp - 0x50]
  0x004deaae
                  483b4110
                                 cmp rax, qword [rcx + 0x10]
< 0x004deab2</pre>
                  0f8650020000
                                 jbe 0x4ded08
                                 sub rsp, 0xd0
  0x004deab8
                  4881ecd00000.
                                 mov qword [var_8h], rbp
 0x004deabf
                  4889ac24c800.
                                 lea rbp, [var_8h]
 0x004deac7
                  488dac24c800.
 0x004deacf
                                 xorps xmm0, xmm0
                  0f57c0
                                 movups xmmword [var_18h], xmm0
  0x004dead2
                  0f118424b800.
```

go tool objdump vs. radare2

- Use both!
- Differences are easy to understand:
- A big advantage of using go tool
 objdump is that you get line numbers in
 code. This can be help you group
 instructions by operations

```
false
                        true for MachO
canary
class
          ELF64
          false
crypto
          little
endian
havecode true
laddr
          0 \times 0
lang
          go
linenum
          true
lsyms
          true
machine
          AMD x86-
maxopsz
          16
minopsz
nx
          true
          linux
OS
pcalign
oic
                        Got ROP?
          false
relocs
          true
rpath
          NONE
sanitiz
          false
static
          true
stripped false
          linux
suheve
```

THE GO ASSEMBLER DEFAULT PROTECTIONS

- Stack Protection: export CGO_LDFLAGS='fstack-protector'
- Stripped: GOOS=linux go build -ldflags="-s -w"
- PIE: export GOFLAGS='-buildmode=pie'
- Use a tool like UPX (https://upx.github.io/) to strip function names and reduce size

rabin2 -zz go-bin-goo | grep 1.__TEXT.__rodata | grep intruder

-zz go-bin-goo | grep 1.__TEXT.__rodata | grep intruder

33526 0x00135e7f 0x01135e7f 2040 2040 1.__TEXT.__rodata ffdirecto ascii ry not emptydisc quota exceededfile already closedfile already existsfile does not exi stm not found in allmmarking free objectmarkroot: bad indexmspan.sweep: state=no STREA M resourcesnotesleep not on g0nwait > work.nprocsoperation timed outpanic during mallo cpanic during panic\npanic holding lockspanicwrap: no (in panicwrap: no) in previous owner diedreflect.Value.Fieldreflect.Value.Floatreflect.Value.Indexreflect.Value.IsNi lreflect.Value.Sliceruntime: insert t= runtime: pcdata is runtime: preempt g0semaRoot rotateLeftskip this directorystopm holding lockstoo many open filesunknown wait reason unsupported messagezero length segment markroot jobs done\n to unallocated span, Recur sionDesired: /usr/share/zoneinfo/37252902984619140625EMULTIHOP (Reserved)Egyptian_Hier oglyphsEnter the password: IDS_Trinary_OperatorIntruder! intruder!\nMeroitic_Hieroglyp hsSIGALRM: alarm clockSIGTERM: terminationSTREAM ioctl timeoutSeek: invalid offsetSeek : invalid whenceTerminal_Punctuationauthentication errorbad defer size classbad system page sizebad use of bucket.bpbad use of bucket.mpchan send (nil chan)close of nil cha nnelfloating point errorforcegc: phase errorgc_trigger underflowgo of nil func valuego park: bad g statusinvalid DNS responsemSpanList.insertBackmalloc during signalnon-empt y swept listnotetsleep not on g0p mcache not flushedpacer: assist ratio=preempt off re ason: reflect.makeFuncStubruntime: casgstatus runtime: double waitruntime: unknown pc semaRoot rotateRightsysctl kern.hostnametime: invalid numbertrace: out of memoryunexpe cted IP lengthunexpected network: wirep: already in goworkbuf is not emptywrite of Go pointer gp.gcscanvalid=true\n of unexported method previous allocCount=18626451492309 5703125931322574615478515625Anatolian_HieroglyphsInscriptional_PahlaviOther_Grapheme_E xtend[+] Checking password_cgo_unsetenv missingbad type in compare: block device requi redbufio: negative countcheckdead: runnable gconcurrent map writesdefer on system stac kdevice not configuredfindrunnable: wrong pillegal byte s

SEARCHING FOR STRINGS



SEARCHING FOR STRINGS

- Go does not store null terminated strings.
- Strings are clumped together, while keeping a separate table with length information.
- This can make it difficult to look for string cross-references
- We can use a projects like https://github.com/carvesystems/gostringsr2 to parse strings
- When working with MachO binaries, you'd have to list strings from .rodata or entire binary
 - rabin2 -zz <binary>
 - izz using r2

SEARCHING FOR FUNCTIONS

sym.local.runtime.modulesinit

- Most of the times, functions are easy to find, even in stripped binaries.
- Functions follow with the following naming conventions:

```
package.function
package.receiver struct.function
Package.__receiver struct__.function
```

func (m *Modules) InitProcessors(mods []base.ModuleConfig) error

```
sym.local.runtime.moduledataverify
sym.local.runtime.moduledataverify1
sym.local.runtime.findmoduledatap
sym.local.github.com_DharmaOfCode_gorp_modules.__Modules_.InitProcessors
sym.local.github.com_DharmaOfCode_gorp_modules.setModuleOption
sym.local.github.com_DharmaOfCode_gorp_modules.printOptions
sym.local.github.com_DharmaOfCode_gorp_modules.__Modules_.GetProcessor
sym.local.github.com_DharmaOfCode_gorp_modules.__Modules_.InitInspectors
```

THE EASY PART FINDING MAIN

```
nm go-bin-elf | grep main
 5f6ec0 D main..inittask
 4dee40 T main.authenticate
 54fa60 R main.authenticate.stkobj
 4dfae0 T main.check
 54cd30 R main.check.stkobj
 4df7a0 T main.checkName
 54cd50 R main.checkName.stkobj
 4df8a0 T main.checkPassword
 54cd70 R main.checkPassword.stkobj
 4ded20 T main.handleConnection
 54cd90 R main.handleConnection.stkobj
4df620 T main.jackpot
4deaa0 T main.main
```

main.main main.handleConnection main.authenticate main.jackpot main.checkName main.checkPassword main.check

- Go routines have small stacks by default (2 kibibyte = 1024 bytes stack)
- Many goroutines will calls
 morestack to grow the stack (in powers of 2) as needed using stack copying
- This is called because go can't be sure the function will outgrow the stack (i.e. recursive functions) given non-deterministic goroutines
- When this occurs, stack grows, pointers in the stack are updated.
- Additionally, each function compares its stack pointer against g->stackguard to check for overflow.

GO STACKS

GO FUNCTION PROLOGUE

```
mov byte [arg_18h], 0
mov rbp, qword [var_8h]
add rsp, 0x58
ret

Name @ 0x4df7ad
call sym.runtime.morestack_noctxt
jmp sym.main.checkName
```

GO FUNCTION PROLOGUE

OR FINDING OPPORTUNITIES TO OVERFLOW THE STACK

- Developers call tell go to skip this check to test performance issues with pragma //go:nosplit
- Some library functions may call //go:nosplit
- Look for opportunities to corrupt the stack
- Callers pass arguments in the stack

```
//go:nosplit
func checkName(name string) bool{
   fmt.Println( a: "[+] Checking name")

   if len(name) != 12{
      return false
   }
```

```
(int64_t arg_8h, int64_t arg_10h, int64_t
     int64_t var_50h @ rsp+0x8
 var int64_t var_48h @ rsp+0x10
 var int64_t var_40h @ rsp+0x18
 var int64_t var_38h @ rsp+0x20
 var int64_t var_18h @ rsp+0x40
 var int64_t var_10h @ rsp+0x48
 var int64_t var_8h @ rsp+0x50
; arg int64_t arg_8h @ rsp+0x60
 arg int64_t arg_10h @ rsp+0x68
 arg int64_t arg_18h @ rsp+0x70
               4883ec58
                              sub rsp, 0x58
0x010e72a0
0x010e72a4
               48896c2450
                              mov qword [var_8h], rbp
               488d6c2450
0x010e72a9
                              lea rbp, [var_8h]
```

CONVENTIONS ARGUMENTS AND RETURN VALUES

- Return values are placed in the stack
- A opposed to C where return values are placed in registers
- Arguments are also moved to the stack rather than registers.

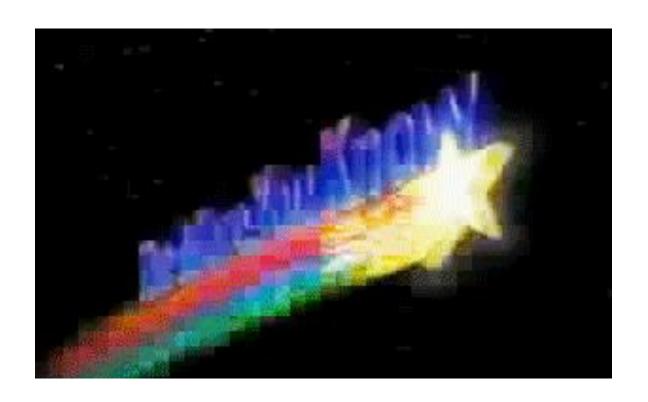
Placing return values in the stack

```
mov byte [arg_18h], 0
mov rbp, qword [var_8h]
add rsp, 0x58
ret_
```

Passing arguments in the stack

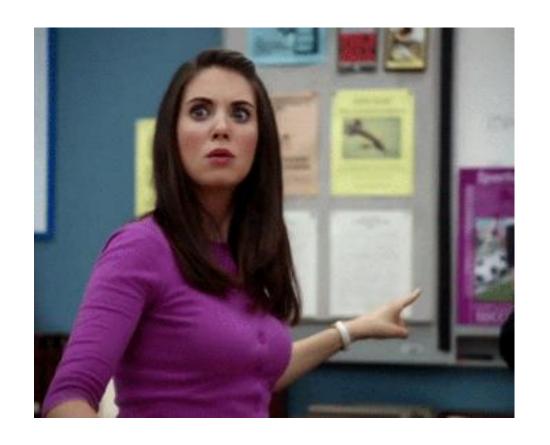
```
e8026dfcff call sym.strings.TrimRight
488b442420 mov rax, qword [var_240h]
488b4c2428 mov rcx, qword [var_238h]
48890424 mov qword [rsp], rax
48894c2408 mov qword [var_258h], rcx
e89a070000 call sym.main.checkName
```

IT'S ALL ABOUT THE GO LIBRARY FUNCTIONS



USE GO LIBRARIES AS A GUIDE TO UNDERSTAND THE BINARY

- Understanding go internal libraries can significantly help us understand what is going on the assembly code
- Read the go docs!



CALLING A GOROUTINE

```
go student.DerefCopyMethod(&students[0])
core.RefCopy(&student, students)
```

```
MOVL $0x10, 0(SP)

LEAQ go.func.*+95(SB), AX

MOVQ AX, 0x8(SP)

MOVQ 0x58(SP), AX

MOVQ AX, 0x10(SP)

MOVQ 0x68(SP), CX

MOVQ CX, 0x18(SP)

CALL runtime.newproc(SB)
```

DEREFENCING RECEIVER

```
func (s *Student) DerefCopyMethod(student *Student){
    *s = *student
}
```

```
//go:nosplit
func typedmemmove(typ *_type, dst, src unsafe.Pointer) {
        if dst == src {
                return
        if typ.ptrdata != 0 {
                bulkBarrierPreWrite(uintptr(dst), uintptr(srd), typ.size)
        // There's a race here: if some other goroutine can write to
        // src, it may change some pointer in src after we've
        // performed the write barrier but before we perform the
        // memory copy. This safe because the write performed by that
        // other goroutine must also be accompanied by a write
        // barrier, so at worst we've unnecessarily greyed the old
       // pointer that was in src.
       memmove(dst, src, typ.size)
        if writeBarrier.cgo {
                cgoCheckMemmove(typ, dst, src, 0, typ.size)
```

```
MOVQ 0x30(SP), CX
TESTB AL, 0(CX)
CMPL $0x0, runtime.writeBarrier(SB)
JNE 0x109d584
MOVQ O(CX), DX
MOVQ DX, 0(AX)
MOVUPS 0x8(CX), X0
MOVUPS X0, 0x8(AX)
MOVUPS 0x18(CX), X0
MOVUPS X0, 0x18(AX)
MOVUPS 0x28(CX), X0
MOVUPS X0, 0x28(AX)
MOVUPS 0x38(CX), X0
MOVUPS X0, 0 \times 38(AX)
MOVQ 0x18(SP), BP
ADDQ $0x20, SP
RET
LEAQ runtime.rodata+159488(SB), DX
MOVQ DX, 0(SP)
MOVQ AX, 0x8(SP)
CALL runtime.typedmemmove(SB)
```

SO WHAT?

- We have identified a race condition
- Note //go:nosplit for typedmemmove



```
//go:nosplit
func typedmemmove(typ *_type, dst, src unsafe.Pointer) {
       if dst == src {
                return
       if typ.ptrdata != 0 {
                bulkBarrierPreWrite(uintptr(dst), uintptr(src), typ.size)
       // There's a race here: if some other goroutine can write to
       // src, it may change some pointer in src after we've
       // performed the write barrier but before we perform the
       // memory copy. This safe because the write performed by that
       // other goroutine must also be accompanied by a write
       // barrier, so at worst we've unnecessarily greyed the old
       // pointer that was in src.
       memmove(dst, src, typ.size)
       if writeBarrier.cgo {
                cgoCheckMemmove(typ, dst, src, 0, typ.size)
```

DEFER

- Function responsible for concurrency are easy to locate
- In the case of a defer call:
 - Functions are registered with deferprocStack
 - Check whether a panic has been caught
 - Prepare return values and call deferred functions before exiting function
 - Assure to defer if a panic is caught

```
lea rax, sym.os.__File_.Close.f;
mov qword [var_80h], rax
mov qword [var_68h], rcx
lea rax, [var_98h]
mov qword [rsp], rax
call sym.runtime.deferprocStack
test eax, eax
jne 0xc0f90
```

```
xorps xmm0, xmm0
movups xmmword [arg_58h], xmm0
nop
call sym.runtime.deferreturn
mov rbp, qword [var_8h]
add rsp, 0x1c8
ret
```

```
call sym.runtime.deferreturn
mov rbp, qword [var_8h]
add rsp, 0x1c8
ret
```

GO INTERFACES

- Interfaces consist of:
 - A pointer to a vtable (which contains function pointers that point to the virtual functions)
 - A value pointer

GO ERROR HANDLING

- error is an an interface
- Error handling in clumsy in go
- Bugs due to unhandled errors are common
- When checking for error != nil we load the error vtable and error value.
- Then we test if the value is nil
- And branch depending on the result

```
s, err := ioutil.ReadFile( filename:
   if err != nil{
      panic(err)
}
```

```
call sym.io_ioutil.ReadFile
mov rax, qword [var_48h]
mov rcx, qword [var_40h]
mov rdx, qword [var_60h]
mov rbx, qword [var_58h]
test rax, rax
jne 0x4df789
```

SO WHAT? WELL...

集CVE-2018-1002105 Detail

Current Description

In all Kubernetes versions prior to v1.10.11, v1.11.5, and v1.12.3, incorrect handling of error responses to proxied upgrade requests in the kube-apiserver allowed specially crafted requests to establish a connection through the Kubernetes API server to backend servers, then send arbitrary requests over the same connection directly to the backend, authenticated with the Kubernetes API server's TLS credentials used to establish the backend connection.

WHAT ELSE SHOULD I LOOK FOR?

LEAQ 0x30(SP), AX
MOVQ AX, 0x18(SP)
MOVQ main._cgo_8c2b83f3b113_Cfunc_memcpy(SB), AX
MOVQ AX, 0(SP)

- cgo functions!
- Hacking like its 1995(98)(2000)



EL FIN

- Go is everywhere
- We are likely to see see more go in IoT with projects like TinyGo
- Fancy, yet exploitable under the right conditions
- Look for programmer errors (i.e. error handling)
- Understand stack handling
- Look for cgo usage
- Leverage go tools
- RTFM

- For developers:
 - Test your code with gosec, govet
 - Fuzz it with go-fuzz

QUESTIONS?