

# HANDLING ADVANCED THREATS

SANS 2020 ONLINE SUMMIT

(EXTENDED VERSION)



by Alexandre Borges



# Agenda:

- ❖ Introduction
- ❖ Reversing
- ❖ Anti-Reversing
- ❖ De-obfuscation

- Cyber Security Researcher
  - Speaker at DEF CON USA 2019
  - Speaker at DEF CON USA 2018
  - Speaker at DEF CON CHINA 2019
  - Speaker at DC2711 (Johannesburg)
  - Speaker at NO HAT 2019 (Italy)
  - Speaker at HITB 2019 (Amsterdam)
  - Speaker at CONFidence 2019 (Poland)
  - Speaker at DevOpsDays BH 2019
  - Speaker at BSIDES  
2019/2018/2017/2016
  - Speaker at H2HC 2016/2015
  - Speaker at BHACK 2019/2018
  - Researcher on Android/iOS Reversing,  
Rootkits and Digital Forensics.
  - Referee on Digital Investigation: The  
International Journal of Digital  
Forensics & Incident Response
- [HTTP://WWW.BLACKSTORMSECURITY.COM](http://www.blackstormsecurity.com)

# □ Last conferences:

- BHACK 2019 (Belo Horizonte/Brazil)
  - DC2711 (Johannesburg/South Africa)
  - NO HAT Conference 2019 (Bergamo/Italy)
  - DEF CON USA 2019 (Las Vegas / USA)
  - CONFidence Conference 2019 (Krakow / Poland)
  - DEF CON China 2019 (Beijing / China)
  - HITB Security Conference 2019 (Amsterdam)
  - BSIDES Sao Paulo 2019 (Sao Paulo / Brazil)
  - DEF CON USA 2018 (Las Vegas / USA)
- 
- **Malwoverview Tool:** <https://github.com/alexandreborges/malwoverview>

# INTRODUCTION



# INTRODUCTION

- During reverse engineering of a malware sample, we need to understand few aspects on the threat:
  - Is the malware **packed**?
  - Are **DLLs and functions resolved dynamically**?
  - Are **strings encrypted**?
  - Are there any anti-forensic techniques such as **anti-vm, anti-debugging or anti-disassembly**?
  - Is there any **obfuscation technique** being used?
- Unpacking malware is usually easy, but you might find sophisticated packers...



# INTRODUCTION

- **Advanced threats** are different from any daily malware because:
  - They **don't use common packers**.
  - Most of the time, they bring **malicious device drives** (rootkits).
  - Sometimes they try to **compromise the platform (bootkits)**
  - They can use **0-days** to exploit the infrastructure and systems.
  - They **bypass most of defenses** and run under the radar.
  - **C2** transmits beacons and data **once per week with short duration**.
  - They might implement tricks to prevent any memory acquisition.
  - Most certainly, there'll be **anti-forensic techniques**.
  - **It's hard and take so much time to reverse it.**
- If you have luck, so you'll have the opportunity to analyze them. 😊



# INTRODUCTION

- There're many packers that we know about to unpack them or, at least, how to manage them (but it can be hard...):
  - **Native code:** ASPack, Armadillo, Petite, FSG, UPX, MPRESS, NSPack, PECompact, WinUnpack and so on...
  - **.NET packers and obfuscators:** .NET Reactor, Salamander .NET Obfuscator, Dotfuscator, Smart Assembly, CryptoObfuscator for .NET, Agile, ArmDot, babelfor.NET, Eazfuscator.NET, Spice.Net, Skater.NET, VM Protect 3.40+ and so on...
  - **Android Packers/Obfuscators:** DexGuard, DexProtect, DevGuard, Arxan, ApkGyard, and so on...



# INTRODUCTION

- **.NET packers** use similar tricks of native code:
  - **Control flow obfuscation** and **dead/junk code insertion**.
  - **Renaming**: methods signatures, fields, methods implementation, namespaces, metadata and external references.
  - **Re-encoding**: changing printable to unprintable characters
  - Simple **encryption of methods and strings**.
  - **Cross reference obfuscation**.
- **In native code**, there're well know memory APIs: **VirtualAlloc/Ex( )**, **HeapCreate( ) / RtlCreateHeap( )**, **HeapReAlloc( )**, **GlobalAlloc( )**, **RtlAllocateHeap( )**





# INTRODUCTION

- Most packed binaries can be unpacked using debuggers, breakpoints and dumping unpacked content from memory.
- Even when a binary uses customized packing techniques, it is still possible:
  - dumping the unpacked code from memory using Volatility.
  - fixing the ImageAddress field using few lines in Python its respective IAT using impscan plugin to analyze it in IDA Pro:
    - `python vol.py -f memory.vmem procdump -p 2096 -D . --memory`  
(to keep slack space)
    - `python vol.py -f memory.vmem impscan --output=idc -p 2096`

# REVERSING



# INTRODUCTION

```
's' .rdata:0040... 0000000E C (1... (null)
's' .rdata:0040... 00000007 C (null)
's' .rdata:0040... 00000005 C ('8PW
's' .rdata:0040... 00000005 C 700PP
's' .rdata:0040... 00000012 C ``hhh\b\b\axppwpp\b\b
's' .rdata:0040... 0000000F C bad allocation
's' .rdata:0040... 00000011 C =ÂUf>pěÂ?Mš\"x1B_0\t_
's' .rdata:0040... 00000005 C Úzz\\-
's' .rdata:0040... 00000006 C ĚCÁ!dÍ
's' .rdata:0040... 00000006 C =°UÈ;É
's' .rdata:0040... 00000005 C Î/~ØÕ
's' .rdata:0040... 00000007 C ãØòšÄ&\b
's' .rdata:0040... 00000007 C Žf{ØÀPý
's' .rdata:0040... 00000009 C a}A.{D|;Î
's' .rdata:0040... 00000009 C 1ē!Fà\\x1BŪ
's' .rdata:0040... 00000005 C @à_Ãç
's' .rdata:0040... 00000005 C Ě6ÓàY
's' .rdata:0040... 0000000C C Qië.ôóĚËSS%{
's' .rdata:0040... 00000007 C Û2K\v$/Ø
's' .rdata:0040... 0000000A C !À7D~(!9q
's' .rdata:0040... 0000000E C Noe(°W~\n8YYhùœç
's' .rdata:0040... 00000008 C uèEĚ|Ěøâ
's' .rdata:0040... 00000007 C tĚ\æâĚµV
's' .rdata:0040... 00000005 C }pl#Y
's' .rdata:0040... 00000006 C J[Ěf|ă
's' .rdata:0040... 00000005 C Y;sVĐ
```

- ❖ As we've mentioned previously, **strings** are one of first references.
- ❖ However, they are all **encrypted** and writing **YARA rules** using them could not be so interesting 😊



# INTRODUCTION

```
text:00408B90  push    ebp
text:00408B91  mov     ebp, esp
text:00408B93  push    0FFFFFFFh
text:00408B95  push    offset SEH_408B90
text:00408B9A  mov     eax, large fs:0 ; Remember:
text:00408B9A                               ;
text:00408B9A                               ; FS --> TEB --> TIB --> SEH
text:00408B9A                               ;
text:00408B9A                               ; push Exception Handler (0xFFFFFFFF to end of handler list)
text:00408B9A                               ; push next record
text:00408B9A                               ; mov fs:[0], esp
text:00408BA0  push    eax
text:00408BA1  sub     esp, 110h
text:00408BA7  mov     eax, ___security_cookie
text:00408BAC  xor     eax, ebp
text:00408BAE  mov     [ebp+var_14], eax
text:00408BB1  push    ebx
text:00408BB2  push    esi
text:00408BB3  push    edi
text:00408BB4  push    eax
text:00408BB5  lea     eax, [ebp+var_C]
text:00408BB8  mov     large fs:0, eax
text:00408BBE  mov     [ebp+var_10], esp
text:00408BC1  push    256             ; Size
text:00408BC6  lea     eax, [ebp+Dst]
text:00408BCC  push    0               ; Val
text:00408BCE  push    eax             ; Dst
text:00408BCF  call    memset           ; void *memset(void *dest, int c,size_t count);
text:00408BCF                               ; Sets buffers to a specified character.
text:00408BD4  add     esp, 0Ch
text:00408BD7  call    ab_resolve_function_addresses
text:00408BDC  test    eax, eax
```

Setup an exception framework

DLLs seem to be “obfuscated” 😊



# INTRODUCTION

Old anti-debugger tricks...



```
jz      loc_408D35
lea     eax, [ebp+Dst]
call    sub_408770
call    sub_409340
test    eax, eax
jnz     loc_408D35
push    2 ; _DWORD
call    dword_4DA870
```

```
mov     edi, ds:SetUnhandledExceptionFilter ; Enables an application to supersede the top-level
                                              ; exception handler of each thread of a process.
                                              ; After calling this function, if an exception occurs
                                              ; in a process that is not being debugged, and the
                                              ; exception makes it to the unhandled exception filter,
                                              ; that filter will call the exception filter function
                                              ; specified by the lpTopLevelExceptionFilter parameter
                                              ;
                                              ; LPTOP_LEVEL_EXCEPTION_FILTER SetUnhandledExceptionFilter(
                                              ;   LPTOP_LEVEL_EXCEPTION_FILTER lpTopLevelExceptionFilter
                                              ; );
```

```
push    offset sub_408760 ; lpTopLevelExceptionFilter
call    edi ; SetUnhandledExceptionFilter
call    sub_409AB0
test    al, al
jz      short loc_408C32
call    sub_409B50
test    al, al
jnz     short loc_408C32
call    sub_4095C0
mov     eax, 1
jmp     loc_408D37
```



# INTRODUCTION

```
.text:00402D60 ab_resolve_function_addresses proc near ; CODE XREF: ab_possible_decode_fn_1+47!p
.text:00402D60
.text:00402D60 var_4          = dword ptr -4
.text:00402D60
.text:00402D60     push     ebp
.text:00402D61     mov      ebp, esp
.text:00402D63     push     ecx
.text:00402D64     push     ebx
.text:00402D65     push     edi
.text:00402D66     xor      eax, eax
.text:00402D68     mov      [ebp+var_4], 0
.text:00402D6F     call     ab_decode_dll_names
.text:00402D74     mov      ebx, ds:GetModuleHandleA
.text:00402D7A     push     eax                ; lpModuleName_GetModuleHandleA
.text:00402D7B     call     ebx ; GetModuleHandleA
.text:00402D7D     mov      edi, eax
.text:00402D7F     test     edi, edi
.text:00402D81     jz       loc_403400
.text:00402D87     push     esi
.text:00402D88     mov      eax, 1
.text:00402D8D     call     ab_decode_dll_names
.text:00402D92     mov      esi, ds:GetProcAddress
.text:00402D98     push     eax                ; lpProcName_GetProcAddress
.text:00402D99     push     edi                ; hModule_GetProcAddress
.text:00402D9A     call     esi ; GetProcAddress
.text:00402D9C     mov      dword_4DA824, eax
.text:00402DA1     test     eax, eax
.text:00402DA3     jz       loc_4033FF
```

Dynamic DLL name resolution being executed before resolving the function addresses.



# INTRODUCTION

```
text:00405380
text:00405380 ab_decode_dll_names proc near          ; CODE XREF: ab_resolve_function_addresses+F↑p
text:00405380                                     ; ab_resolve function addresses+2D↑p ...
text:00405380                                     jmp     ds:ab_dll_resolver_switch_cases[eax*4] → Jump table
text:00405387 ; -----
text:00405387
text:00405387 loc_405387:                                ; CODE XREF: ab_decode_dll_names↑j
text:00405387                                     ; DATA XREF: .text:ab_dll_resolver_switch_cases↓o
text:00405387     xor     edx, edx
text:00405389     cmp     var_2, edx
text:0040538F     jnz     short loc_405405
text:00405391     mov     var_4, 204C2044h
text:0040539B     mov     var_1, 50EA7350h
text:004053A5     mov     var_3, 6A586B07h → DLL name (obfuscated format)
text:004053AF     mov     var_2, 620C7E1Bh
text:004053B9     mov     var_5, 17E21920h
text:004053C3     mov     eax, 13h
text:004053C8     jmp     short loc_4053D0
text:004053C8 ; -----
text:004053CA     align 10h
text:004053D0
text:004053D0 loc_4053D0:                                ; CODE XREF: ab_decode_dll_names+48↑j
text:004053D0                                     ; ab_decode_dll_names+5D↓j
text:004053D0     mov     cl, [eax+4DAF43h]
text:004053D6     xor     [eax+4DAF44h], cl
text:004053DC     dec     eax
text:004053DD     jnz     short loc_4053D0
text:004053DF     xor     ecx, ecx
text:004053E1
```





# INTRODUCTION

```
text:004053E1 loc_4053E1:                                ; CODE XREF: ab_decode_dll_names+83↓j
text:004053E1      cmp     eax, 4
text:004053E4      jb      short loc_4053F9
text:004053E6      cmp     [eax+4DAF44h], dl
text:004053EC      jnz     short loc_4053F5
text:004053EE      mov     ecx, 1
text:004053F3      jmp     short loc_4053F9
text:004053F5 ; -----
text:004053F5
text:004053F5 loc_4053F5:                                ; CODE XREF: ab_decode_dll_names+6C↑j
text:004053F5      cmp     ecx, edx
text:004053F7      jz      short loc_4053FF
text:004053F9
text:004053F9 loc_4053F9:                                ; CODE XREF: ab_decode_dll_names+64↑j
text:004053F9      ; ab_decode_dll_names+73↑j
text:004053F9      mov     byte ptr var_1[eax], dl
text:004053FF
text:004053FF loc_4053FF:                                ; CODE XREF: ab_decode_dll_names+77↑j
text:004053FF      inc     eax
text:00405400      cmp     eax, 14h
text:00405403      jb      short loc_4053E1
text:00405405
text:00405405 loc_405405:                                ; CODE XREF: ab_decode_dll_names+F↑j
text:00405405      mov     eax, 4DAF48h
text:0040540A      retn
```





# INTRODUCTION

```
text:0040677C ab_dll_resolver_switch_cases dd offset loc_405387
text:0040677C                                     ; DATA XREF: ab_decode_dll_names↑r
text:00406780 dd offset loc_40540B
text:00406784 dd offset loc_40548B
text:00406788 dd offset loc_40550B
text:0040678C dd offset loc_40558B
text:00406790 dd offset loc_40560B
text:00406794 dd offset loc_4056AB
text:00406798 dd offset loc_40573B
text:0040679C dd offset loc_4057CB
text:004067A0 dd offset loc_40585B
text:004067A4 dd offset loc_4058FB
text:004067A8 dd offset loc_40599B
text:004067AC dd offset loc_405A2B
text:004067B0 dd offset loc_405ABB
text:004067B4 dd offset loc_405B4B
text:004067B8 dd offset loc_405BCB
text:004067BC dd offset loc_405C3D
text:004067C0 dd offset loc_405CBB
text:004067C4 dd offset loc_405D4B
text:004067C8 dd offset loc_405DCB
text:004067CC dd offset loc_405E4B
text:004067D0 dd offset loc_405ECB
text:004067D4 dd offset loc_405F4B
text:004067D8 dd offset loc_405FCB
```

❖ Each offset takes us to a **different switch case**, which is a **different DLL name resolution function**. 😊



# INTRODUCTION

Direction	Type	Address	Text
	p	ab_resolve_function_addresses+F	call ab_decode_dll_names
Down	p	ab_resolve_function_addresses+2D	call ab_decode_dll_names
Down	p	ab_resolve_function_addresses+4E	call ab_decode_dll_names
Down	p	ab_resolve_function_addresses+69	call ab_decode_dll_names
Down	p	ab_resolve_function_addresses+84	call ab_decode_dll_names
Down	p	ab_resolve_function_addresses+9F	call ab_decode_dll_names
Down	p	ab_resolve_function_addresses+BA	call ab_decode_dll_names
Down	p	ab_resolve_function_addresses+D5	call ab_decode_dll_names
Down	p	ab_resolve_function_addresses+F0	call ab_decode_dll_names
Down	p	ab_resolve_function_addresses+10B	call ab_decode_dll_names
Down	p	ab_resolve_function_addresses+126	call ab_decode_dll_names
Down	p	ab_resolve_function_addresses+141	call ab_decode_dll_names
Down	p	ab_resolve_function_addresses+15C	call ab_decode_dll_names
Down	p	ab_resolve_function_addresses+177	call ab_decode_dll_names
Down	p	ab_resolve_function_addresses+192	call ab_decode_dll_names
Down	p	ab_resolve_function_addresses+1AD	call ab_decode_dll_names
Down	p	ab_resolve_function_addresses+1C8	call ab_decode_dll_names
Down	p	ab_resolve_function_addresses+1E3	call ab_decode_dll_names
Down	p	ab_resolve_function_addresses+1FE	call ab_decode_dll_names
Down	p	ab_resolve_function_addresses+219	call ab_decode_dll_names
Down	p	ab_resolve_function_addresses+234	call ab_decode_dll_names
Down	p	ab_resolve_function_addresses+24F	call ab_decode_dll_names
Down	p	ab_resolve_function_addresses+26A	call ab_decode_dll_names
Down	p	ab_resolve_function_addresses+285	call ab_decode_dll_names
Down	p	ab_resolve_function_addresses+2A0	call ab_decode_dll_names
Down	p	ab_resolve_function_addresses+2D3	call ab_decode_dll_names

❖ As expected, there're many calls to the **same function** for “decrypting” the DLL names 😊



# INTRODUCTION

```
1  from binascii import *
2
3  var_1 = ['50', '73', 'EA', '50']
4  var_2 = ['1B', '7E', '0C', '62']
5  var_3 = ['07', '6B', '58', '6A']
6  var_4 = ['44', '20', '4C', '20']
7  var_5 = ['20', '19', 'E2', '17']
8
9  mylist = var_1 + var_2 + var_3 + var_4 + var_5
10
11 ▼ def mydecrypt(hexdata):
12     max = len(hexdata) - 1
13     counter = max
14     output = ""
15 ▼     while(True):
16         hexdata[counter] = ord(unhexlify(hexdata[counter])) ^ ord(unhexlify(hexdata[counter - 1]))
17         counter -= 1
18         if counter == 0:
19             break
20     return hexdata
21
22 final = mydecrypt(mylist)
23 for x in range(0,4):
24     final[x] = 0
25 final1 = ''.join([chr(w) for w in final])
26 print("The output is %s" % final1)
```

➤ output: kernel32.dll



# INTRODUCTION

- Of course, we could try to **improve and automatize the de-obfuscation** of all **functions names** by using:
  - **IDA Python**: using IDA Python you can de-obfuscated function names and save them into the idb.
  - **IDC**: it's a bit more complicated, but very powerful.
- If your time is short, so you could try emulation tools such as **Floss** (<https://github.com/fireeye/flare-floss>) to decode possible obfuscated strings and create an IDA script to decorate the reversed code:
  - `floss --no-static-strings -x malware.bin --ida=floss_ida.py`



# INTRODUCTION

```
E:\malware_samples>floss --no-static-strings -x malware.bin --ida=floss_ida.py
```

```
...
```

```
Decoding function at 0x405380 (decoded 38 strings)
```

Offset	Called At	String
-----	-----	-----
0x4DAF48	0x402D6F	Kernel32.dll
0x4DAF5C	0x402D8D	CloseHandle
0x4DAF70	0x402DAE	CreateFileA
0x4DAF84	0x402DC9	CreateMutexA
0x4DAFAC	0x402DE4	CreateToolhelp32Snapshot
0x4DAFCC	0x402DFF	DeviceIoControl
0x4DAFE4	0x402E1A	GetCurrentThread
0x4DAFFC	0x402E35	GetLongPathNameA
0x4DB014	0x402E50	GetModuleFileNameA
0x4DB030	0x402E6B	GetNativeSystemInfo
0x4DB04C	0x402E86	GetProcessHeap
0x4DB064	0x402EA1	GetSystemInfo
0x4DB07C	0x402EBC	GetThreadContext
0x4DB094	0x402ED7	HeapAlloc
0x4DB0A8	0x402EF2	HeapFree
0x4DAF98	0x402F0D	HeapReAlloc
0x4DB0B8	0x402F28	IsBadReadPtr
0x4DB0CC	0x402F43	Module32First
0x4DB0E4	0x402F5E	Module32Next



# INTRODUCTION

```
.text:00405380 ab_decode_dll_names proc near          ; CODE XREF: ab_resolve_function_addresses+F↑p
.text:00405380                                     ; ab_resolve_function_addresses+2D↑p ...
.text:00405380             jmp     ds:ab_dll_resolver_switch_cases[eax*4]
.text:00405387 ; -----
.text:00405387 loc_405387:                          ; CODE XREF: ab_decode_dll_names↑j
.text:00405387                                     ; DATA XREF: .text:ab_dll_resolver_switch_cases↓o
.text:00405387             xor     edx, edx
.text:00405389             cmp     var_2, edx          ; FLOSS: Kernel32.dll
.text:0040538F             jnz     short loc_405405
.text:00405391             mov     var_4, 204C2044h
.text:0040539B             mov     var_1, 50EA7350h
.text:004053A5             mov     var_3, 6A586B07h
.text:004053AF             mov     var_2, 620C7E1Bh      ; FLOSS: Kernel32.dll
.text:004053B9             mov     var_5, 17E21920h
.text:004053C3             mov     eax, 13h
.text:004053C8             jmp     short loc_4053D0
.text:004053C8 ; -----
.text:004053CA             align 10h
.text:004053D0 loc_4053D0:                          ; CODE XREF: ab_decode_dll_names+48↑j
.text:004053D0                                     ; ab_decode_dll_names+5D↓j
.text:004053D0             mov     cl, [eax+4DAF43h]
.text:004053D6             xor     [eax+4DAF44h], cl
.text:004053DC             dec     eax
.text:004053DD             jnz     short loc_4053D0
.text:004053DF             xor     ecx, ecx
```

Same result, of course. ☺



# INTRODUCTION

- IDA Python is ALWAYS very handy and we can use IDA Pro SDK to write plugins for:
  - extending the IDA Pro functionalities
  - analyzing piece of obfuscated code and data flow
  - automatizing unpacking of strange malicious files
  - decoding and loading encrypted / modified MBRs
- It is quick to create a simple IDA Pro plugin.
  - Download the IDA SDK from <https://www.hex-rays.com/products/ida/support/download.shtml> (likely, you will need a professional account).
  - Copy it to a folder ([idasdk695/](#)) within the IDA Pro installation directory.



# ➤ INTRODUCTION

- Create a project in Visual Studio 2017/2019 (File → New → Create Project → Visual C++ → Windows Desktop → Dynamic-Link Library (DLL) ).
- Change **few project properties** as shown in this slide and next ones.

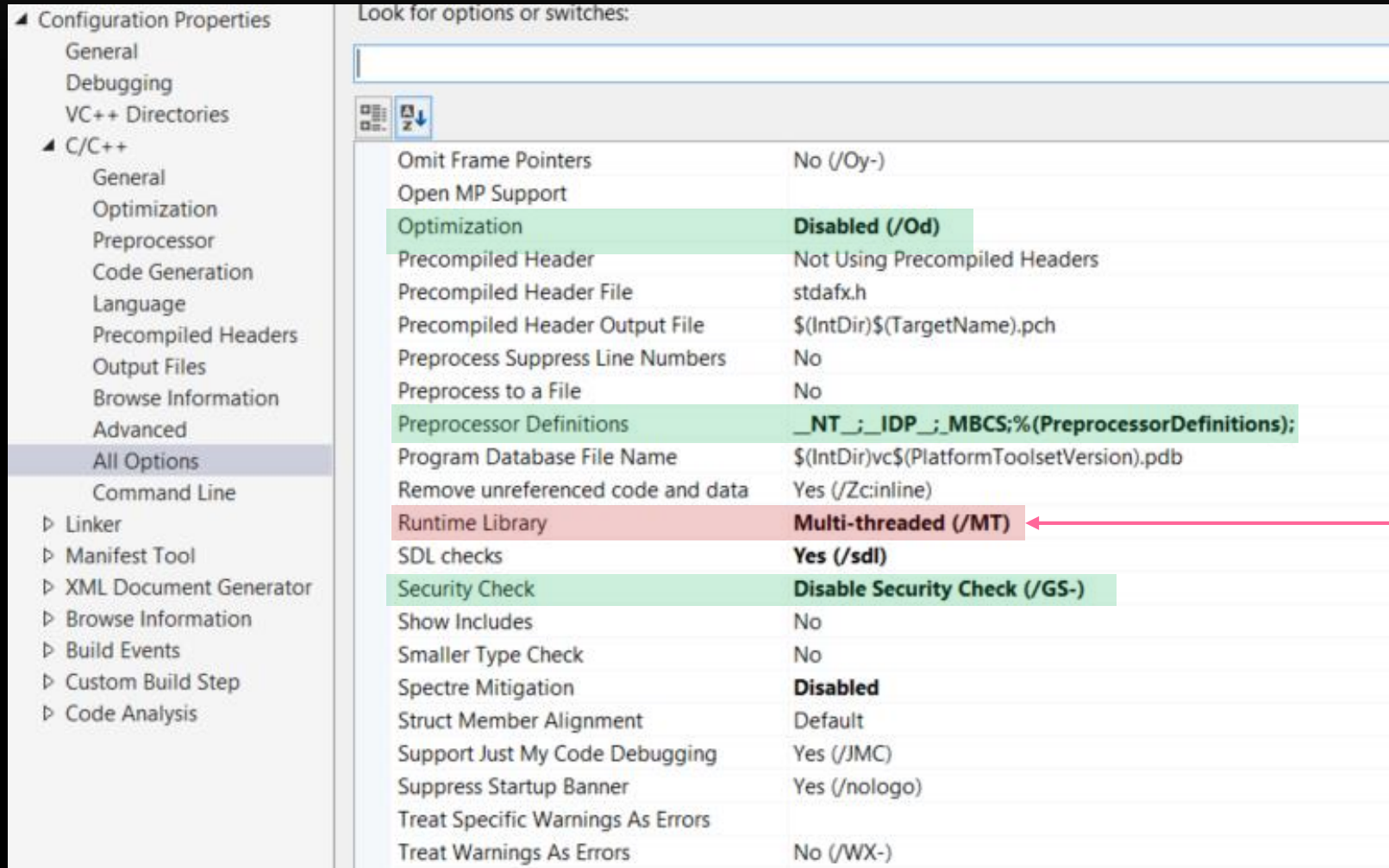
Look for options or switches:

Additional #using Directories	
Additional Include Directories	C:\Program Files (x86)\IDA 6.95\idasdk695\include
Additional Options	
ASM List Location	\$(IntDir)
Assembler Output	No Listing
Basic Runtime Checks	Both (/RTC1, equiv. to /RTCsu) (/RTC1)
Browse Information File	\$(IntDir)
C++ Language Standard	
Calling Convention	_stdcall (/Gz)
Common Language RunTime Support	
Compile As	Default
Conformance mode	Yes (/permissive-)



# ➤ INTRODUCTION

- Include the “`__NT__`; `__IDP__`” in **Processor Definitions** and change **Runtime Library** to “**Multi-threaded**” (MT).



**Warning!**  
It is NOT /MTd

**ida.lib: C:\Program Files (x86)\IDA  
6.95\idasdk695\lib\x86\_win\_vc\_32**

```

1  #include <ida.hpp>
2  #include <idp.hpp>
3  #include <loader.hpp>
4  #include <allins.hpp>
5  #include <strlist.hpp>
6  #include <search.hpp>
7
8
9  int IDAP_init()
10 {
11     return PLUGIN_KEEP;
12 }
13
14 void IDAP_term(void)
15 {
16 }
17
18
19 void IDAP_run(int arg)
20 {
21     msg("Blackstorm Security: basic plugin example :)\n\n");
22
23     char blackstorm[MAXSTR];
24     string_info_t strinfo;
25     char s[] = "[a-z0-9_-]+[\\.]\\{1,\\}([a-zA-Z0-9_-]+[\\.]\\{1,\\}[a-z0-9]{2,}\\>";
26     auto last = BADADDR;
27     auto ea = 0;
28     auto urlcount = 1;

```

Needed headers. 😊

Initialization function.

Make the plugin available to this idb and keep the plugin loaded in memory.

Clean-up tasks.

This is called when users activates the plugin.

Simple (and incomplete) URL regex. 😊

```

29
30 for (int x = 0; x < get_strlist_qty(); x++) {
31
32     get_strlist_item(x, &strinfo);
33     if (strinfo.length < sizeof(blackstorm)) {
34
35         get_many_bytes(strinfo.ea, blackstorm, strinfo.length);
36
37         {
38             ea = 0;
39             ea = find_text(strinfo.ea, 0, 0, s, SEARCH_REGEX);
40
41             if (ea == strinfo.ea) {
42                 msg("Address 0x%x - URL %d: %s\n", strinfo.ea, urlcount,
43                     blackstorm);
44                 urlcount++;
45             }
46         }
47     }
48
49     return;
50
51 }

```

It gets the number of strings from "Strings view".

It gets strings.

The core logic is only it. It checks whether the string matches to the URL regex.

If checks, so `ea == strinfo.ea`. 😊



# INTRODUCTION

```
52
53 char IDAP_comment[] = "The simplest possible plugin";
54 char IDAP_help[] = "Blackstorm plugin";
55 char IDAP_name[] = "Blackstorm plugin";
56 char IDAP_hotkey[] = "ALT-C";
```

Plugin will be activated by  
combination ALT-C. 😊

```
57
58 plugin_t PLUGIN =
59 {
60     IDP_INTERFACE_VERSION,
61     0,
62     IDAP_init,
63     IDAP_term,
64     IDAP_run,
65     IDAP_comment,
66     IDAP_help,
67     IDAP_name,
68     IDAP_hotkey
69 };
70
```

Plugin structure.

```

sub_99F8D380
sub_99F8D3B0
sub_99F8D430
sub_99F8D4A0
sub_99F8D5A0
sub_99F8D650
sub_99F8D680
sub_99F8D6A0
sub_99F8D6D0
sub_99F8D700
sub_99F8D830
sub_99F8D910
sub_99F8D950
sub_99F8D970
nullsub_1
sub_99F8D9D0
sub_99F8DA30
sub_99F8DB50
sub_99F8DC40
sub_99F8DCB0
sub_99F8DCD0
sub_99F8DD30
sub_99F8DEE0
sub_99F8E610

.text:99F8D000 ; int __stdcall sub_99F8D000(PVOID P)
.text:99F8D000 sub_99F8D000 proc near ; CODE XREF: .text:99F96737↓p
.text:99F8D000
.text:99F8D000 P = dword ptr 4
.text:99F8D000
    .text:99F8D000 mov     eax, [esp+P]
    .text:99F8D004 push    0 ; Tag
    .text:99F8D006 push    eax ; P
    .text:99F8D007 call    ds:ExFreePoolWithTag
    .text:99F8D00D retn    4
.text:99F8D00D sub_99F8D000 endp
.text:99F8D00D
.text:99F8D010 ; ===== S U B R O U T I N E =====
.text:99F8D010
.text:99F8D010
.text:99F8D010 ; int __stdcall sub_99F8D010(PVOID P, ULONG AllocationSize, PVOID *BaseAddress, int, int)
.text:99F8D010 sub_99F8D010 proc near ; CODE XREF: .text:99F96763↓p
.text:99F8D010
.text:99F8D010 P = dword ptr 4
.text:99F8D010 AllocationSize = dword ptr 8

Line 1 of 206 00000407| 99F8D007: sub_99F8D000+7| (Synchronized with Hex View-1)

```

Output window

StructTypers init() called!

Python 2.7.14 (v2.7.14:84471935ed, Sep 16 2017, 20:19:30) [MSC v.1500 32 bit (Intel)]  
 IDAPython v1.7.0 final (serial 0) (c) The IDAPython Team <idapython@googlegroups.com>

Blackstorm Security: basic plugin example :)

```

Address 0x99F990d8 - URL 1: ntp2.usno.navy.mil
Address 0x99F990eb - URL 2: ntp.adc.am
Address 0x99F990f6 - URL 3: tock.usask.ca
Address 0x99F99104 - URL 4: ntp.crifo.org
Address 0x99F99112 - URL 5: ntp1.arnes.si
Address 0x99F99120 - URL 6: ntp.ucsd.edu
Address 0x99F9912d - URL 7: ntp.duckcorp.org
Address 0x99F9913e - URL 8: www.nist.gov
Address 0x99F9914b - URL 9: clock.isc.org
Address 0x99F99159 - URL 10: time.windows.com
Address 0x99F9916a - URL 11: time2.one4vision.de
Address 0x99F9917e - URL 12: time.cerias.purdue.edu

```

URLs found within this malicious driver. 😊

ALT + C

# ANTI-REVERSING





# ANTI-REVERSING

- Advanced threats don't use standard tricks and there're many lots of facts about them:
  - They use similar techniques from modern packers such as Themida, Arxan, Agile .NET, Tigress, Obfuscator-LLVM and so on.
  - Most of them are focused on 64-bit code.
  - Obviously, almost all functions are removed from IAT. Remember that Themida packer usually keeps only TlsSetValue( ).
  - String encryption is also a common technique.





# ANTI-REVERSING

- Advanced threats are concerned to integrity:
  - They **protect and check the memory integrity**.
  - Thus, **it is not possible to dump a clean executable from the memory** (using Volatility, for example) because **original instructions are not completely decoded in the memory**.
  - There could be **checksum functions** verifying the integrity of the code itself. Therefore, any attempt of changing the code may break it up.
  - Additionally, advanced threats may use **watermark** to control the “ownership”. This is the same technique used to program copyright. 😊

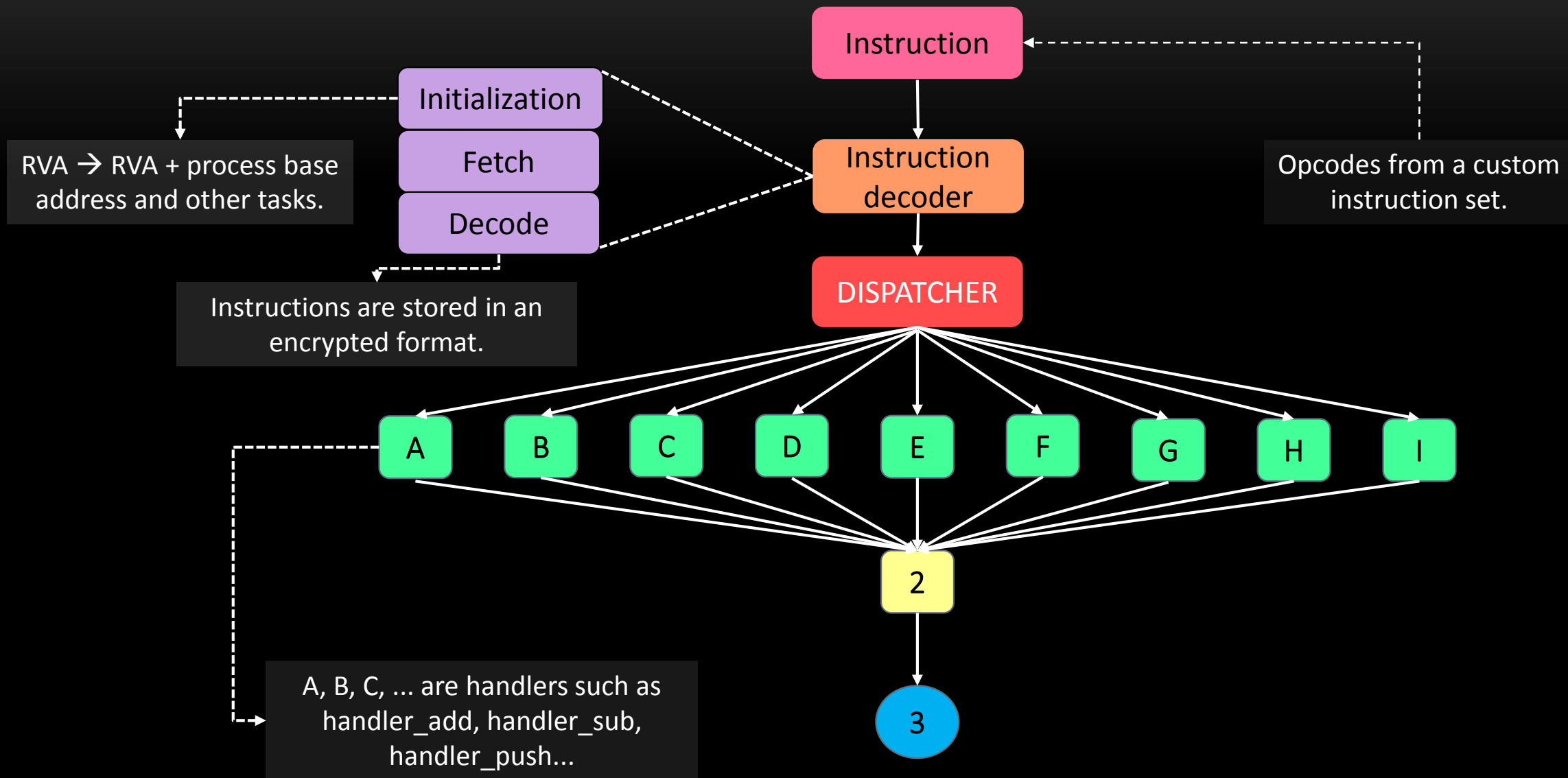


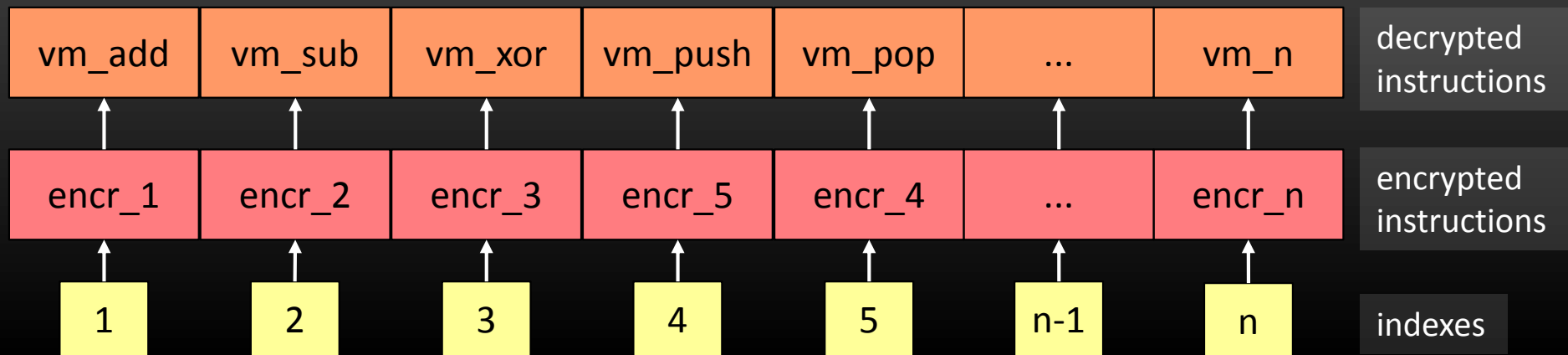
# ANTI-REVERSING

- Many additional techniques are used:
  - There are also **fake push instructions**.
  - There are many **dead and useless pieces of code**.
  - There is some **code reordering** using unconditional jumps.
  - All obfuscators use **code flattening**.
  - Packers have **few anti-debugger and anti-vm tricks**. Weird **anti-vm methods** based on **temperature**, for example.

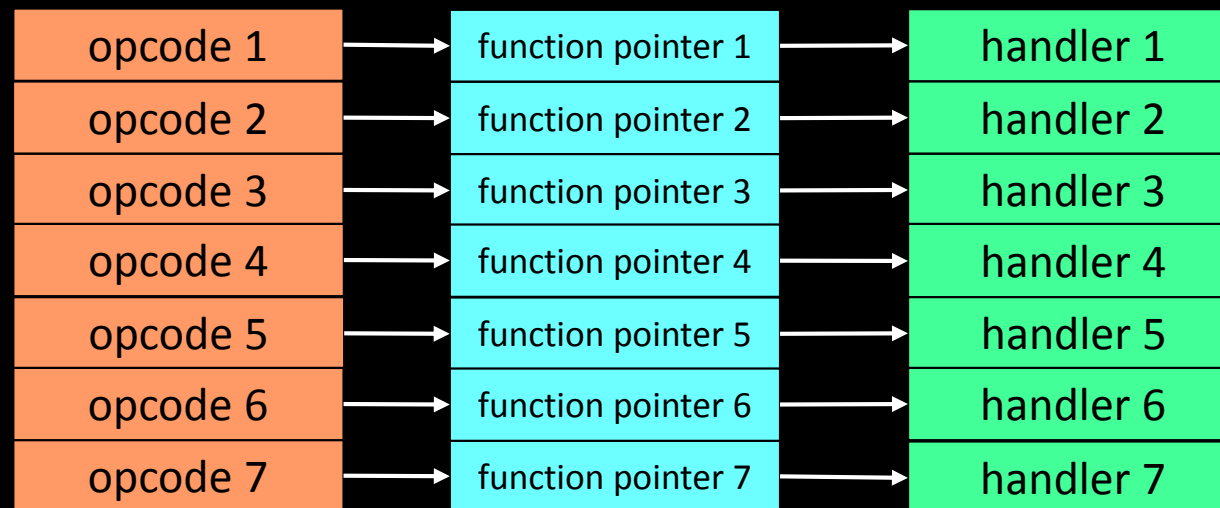


# ANTI-REVERSING





recovering and  
decrypting functions



function pointer table  
(likely encrypted)



# ANTI-REVERSING

**Remember: obfuscating** is transforming a code from A to B by using different techniques (including **virtualization**).

What're the **transition points** from **native code** to **virtualized code** and vice-versa?

**Prologues and epilogues** from each function **could not be virtualized**. Take care. 😊

Have you tried to open an advanced packer in **IDA Pro**? First sight: **only red and grey blocks (non-functions and data)**. 😞



# ANTI-REVERSING

- In few cases the **VM handlers** come **from data blocks**.

- Original **code section** could be “**split**” and “**scatter**” around the program. In this case, **data and instructions are mixed in the binary**, without having just one instruction block.

- Instructions which reference **imported functions** could have been either **zeroed or replaced by NOP**. ☹

- Most certainly, they will be **restored (re-inserted) dynamically** by the packer later.

# ➤ ANTI-REVERSING

- If **references** are not zeroed, so they are usually **translated to short jumps using RVA**, for the **same import address** (“IAT obfuscation”) 😊
- **API names** could be **hashed** (as used in shellcodes). 😊
- Custom packers usually **don't virtualize all native (x86/x64) instructions**.
- There's a **mix** between **virtualized, native instructions** and data after the packing procedure.



# ANTI-REVERSING

- **Native APIs could be redirected to stub code (proxy), which forwards the call to (copied) native DLLs (from the respective APIs).**
- **The “hidden” function code could be copied (memcpy( )) to memory allocated by VirtualAlloc( ) 😊 Of course, there must be a fixup in the code to get these instructions.**





# ANTI-REVERSING

- By the way, **how many virtualized instructions** exist in the binary?

- It is recommended to try to **find handlers to native x86 instructions** (non-virtualized instruction)

- Try to **classify virtualized instructions in groups** according to operands and their purpose such as **memory access, conditional/unconditional jumps, arithmetic, general...**

- Try to understand the **size of virtualized instructions** which we might fit into a **structure that represents encryption key, data, RVA (location), opcode (type)** and so on.



# ANTI-REVERSING

- Pay attention to **instruction's stem** to put **similar classes of instructions** together as, for example, **jump instructions, direct calls, indirect calls** and so on.
- Find the **transition instructions from native mode to virtualized mode** and vice versa.
- Find **similarity between virtualized instructions and x86 instructions**.
- **x86 instructions** are also kept **encrypted and compressed together with the virtualized instructions**.



# ANTI-REVERSING

- **Constant unfolding:** technique used by obfuscators to replace a constant by a bunch of code that produces the same resulting constant's value.
- **Pattern-based obfuscation:** exchange of one instruction by a set of equivalent instructions.
- **Abusing inline functions.**
- **Anti-VM techniques:** prevents the malware sample to run inside a VM.
- **Dead (garbage) code:** this technique is implemented by inserting codes whose results will be overwritten in next lines of code or, worse, they won't be used anymore.



# ANTI-REVERSING

- **Code duplication:** different paths coming into the same destination (used by virtualization obfuscators).
- **Control indirection 1:** call instruction → stack pointer update → return skipping some junk code after the call instruction (RET x).
- **Control indirection 2:** malware trigger an exception → registered exception is called → new branch of instructions.
- **Anti-debugging:** used as irritating techniques to slow the process analysis.



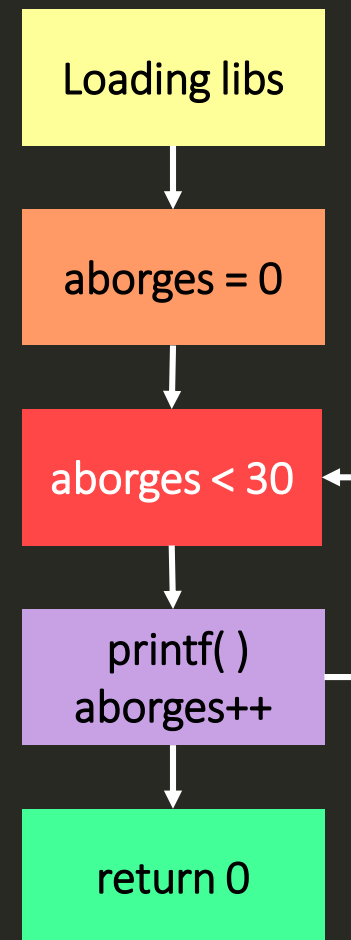
# ANTI-REVERSING

- **Opaque predicate:** Although apparently there is an evaluation (**conditional jump: jz/jnz**), the result is always evaluated to true (or false), which means an **unconditional jump**. Thus, **there is a dead branch**.
- **Polymorphism:** it is produced by **self-modification code** (like shellcodes) and by **encrypting resources** (similar most malware samples).



# ANTI-REVERSING

```
1  #include <stdio.h>
2
3  int main (void)
4  {
5      int aborges = 0;
6      while (aborges < 30)
7      {
8          printf("%d\n", aborges);
9          aborges++;
10     }
11
12     return 0;
13 }
14
15
16
```



```

; Attributes: bp-based frame

; int __cdecl main(int argc, const char **argv, const char **envp)
public main
main proc near

var_4= dword ptr -4

; __unwind {
push    rbp
mov     rbp, rsp
sub     rsp, 10h
mov     [rbp+var_4], 0
jmp     short loc_1160

```

Original Program

```

loc_1160:
cmp     [rbp+var_4], 1Dh
jle     short loc_1146

```

```

loc_1146:
mov     eax, [rbp+var_4]
mov     esi, eax
lea     rdi, format      ; "%d\n"
mov     eax, 0
call    _printf
add     [rbp+var_4], 1

```

```

mov     eax, 0
leave
retn
; } // starts at 1135
main endp

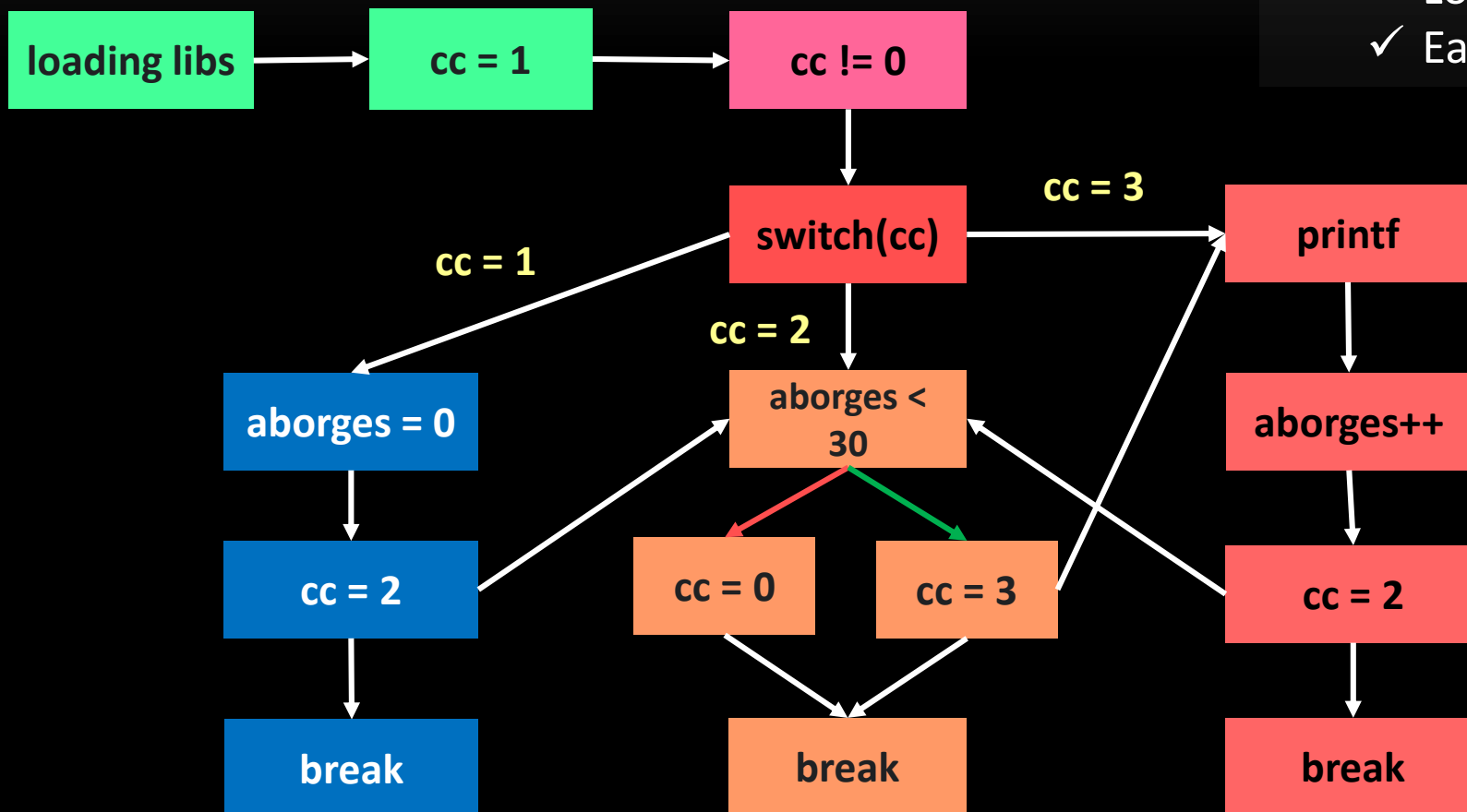
```



# ANTI-REVERSING

## ❖ Disadvantages:

- ✓ Loss of performance
- ✓ Easy to identify the Code flattening







# ANTI-REVERSING

- The **obfuscator-llvm** is an excellent project to be used for **code obfuscation**. To install it, it is recommended to add a swap file first (because the linkage stage):
  - `fallocate -l 8GB /swapfile ; chmod 600 /swapfile`
  - `mkswap /swapfile ; swapon /swapfile ; swapon --show`
  - `apt-get install llvm-4.0`
  - `apt-get install gcc-multilib` (install gcc lib support to 32 bit)
  - `git clone -b llvm-4.0 https://github.com/obfuscator-llvm/obfuscator.git`
  - `mkdir build ; cd build/`
  - `cmake -DCMAKE_BUILD_TYPE=Release -DLLVM_INCLUDE_TESTS=OFF ../obfuscator/`
  - `make -j7`



# ANTI-REVERSING

- Possible usages:
  - `./build/bin/clang alexborges.c -o alexborges -mllvm -fla`
  - `./build/bin/clang alexborges.c -m32 -o alexborges -mllvm -fla`
  - `./build/bin/clang alexborges.c -o alexborges -mllvm -fla -mllvm -sub`
- Where:
  - **fla**: Control Flow Flattening
  - **sub**: Instruction Substitution
  - **bcf**: Opaque Predicate



# ANTI-REVERSING

- A better option would be using Tigress.
- Download Tigress binary from <https://tigress.wtf/download.html>
- Install Tigress is pretty easy:
  - `unzip tigress-3.1-bin.zip`
  - Export the `TIGRESS_HOME` environment variable:
    - `export TIGRESS_HOME=/root/Downloads/tigress/3.1`
  - Add the Tigress installation directory to the `PATH` variable:
    - `export PATH=$PATH:/root/Downloads/tigress/3.1`



# ANTI-REVERSING

```
1  #include <stdio.h>
2  #include "/root/Downloads/tigress/3.1/tigress.h"
3
4  int main (void)
5
6  {
7      int aborges = 0;
8      while (aborges < 30)
9      {
10         printf("%d\n", aborges);
11         aborges++;
12     }
13
14     return 0;
15
16 }
```



# ANTI-REVERSING

- To transform a C source using Tigress:
  - `tigress --Environment=x86_64:Linux:Gcc:4.6 --Transform=Flatten --Functions=main --out=aleborges_obfuscated.c aleborges_trigress.c`
- There're many notes about the command above:
  - We should pick up one or more functions to be transformed. Of course, I've chosen the `main( )` only for educational purposes.
  - The argument for Environment must be according to your environment (`x86_64:Linux:Gcc:4.6`, `x86_64:Darwin:Clang:5.1`, `armv7:Linux:Gcc:4.6`, `armv8:Linux:Gcc:4.6`)



# ANTI-REVERSING

- Additional notes:
  - We could use **multiple transformations** (specifying the --Transform option and --Function option multiple times) such as:
    - Opaque Predicate (InitOpaque)
    - Virtualization Obfuscation (Virtualize)
    - Split and Merge
    - Encode Literals
    - Encode Branches
    - AntiTaintAnalysis
    - ...and much more...



# ANTI-REVERSING

```
76 int main(int _formal_argc , char **_formal_argv , char **_formal_envp )
77 {
78     int aborges ;
79     int _BARRIER_0 ;
80     unsigned long _1_main_next ;
81
82     {
83     megaInit();
84     _global_argc = _formal_argc;
85     _global_argv = _formal_argv;
86     _global_envp = _formal_envp;
87     _BARRIER_0 = 1;
88     {
89     _1_main_next = 2UL;
90     }
91     while (1) {
92         switch (_1_main_next) {
93             case 4: ;
94             return (0);
95             break;
96             case 3:
97                 printf((char const  /* __restrict */) "%d\n", aborges);
```

Program obfuscated with Tigress



# ANTI-REVERSING

```
98     aborges ++;
99     {
100     _1_main_next = 0UL;
101     }
102     break;
103     case 0: ;
104     if (aborges < 30) {
105     {
106     _1_main_next = 3UL;
107     }
108     } else {
109     {
110     _1_main_next = 4UL;
111     }
112     }
113     break;
114     case 2:
115     aborges = 0;
116     {
117     _1_main_next = 0UL;
118     }
119     break;
120     }
121 }
122 }
123 }
124 void megaInit(void)
125 {
```

Program obfuscated with  
Tigress (remaining part)





# ANTI-REVERSING

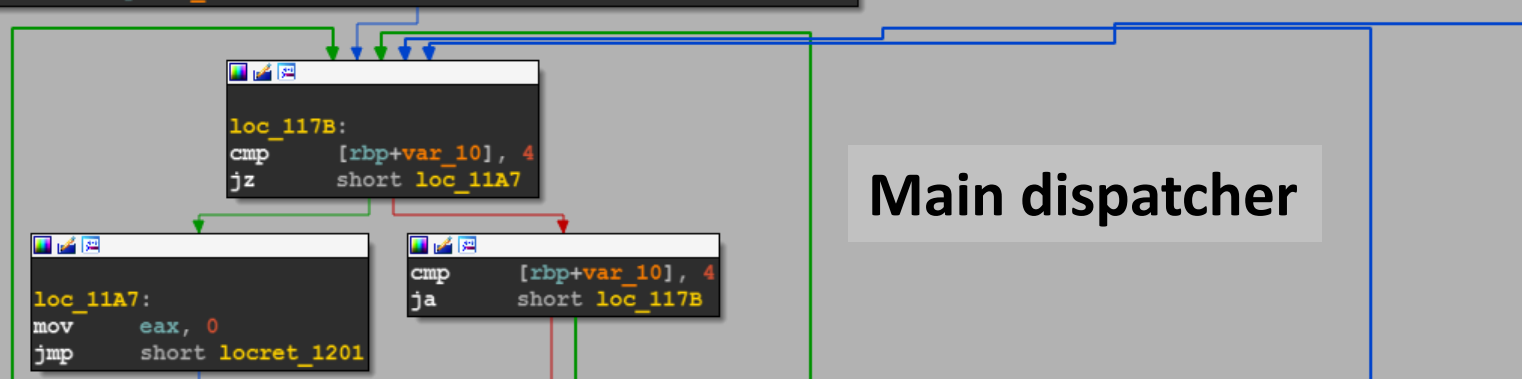
```
; Attributes: bp-based frame

; int __cdecl main(int argc, const char **argv, const char **envp)
public main
main proc near

var_38= qword ptr -38h
var_30= qword ptr -30h
var_24= dword ptr -24h
var_14= dword ptr -14h
var_10= qword ptr -10h
var_4= dword ptr -4

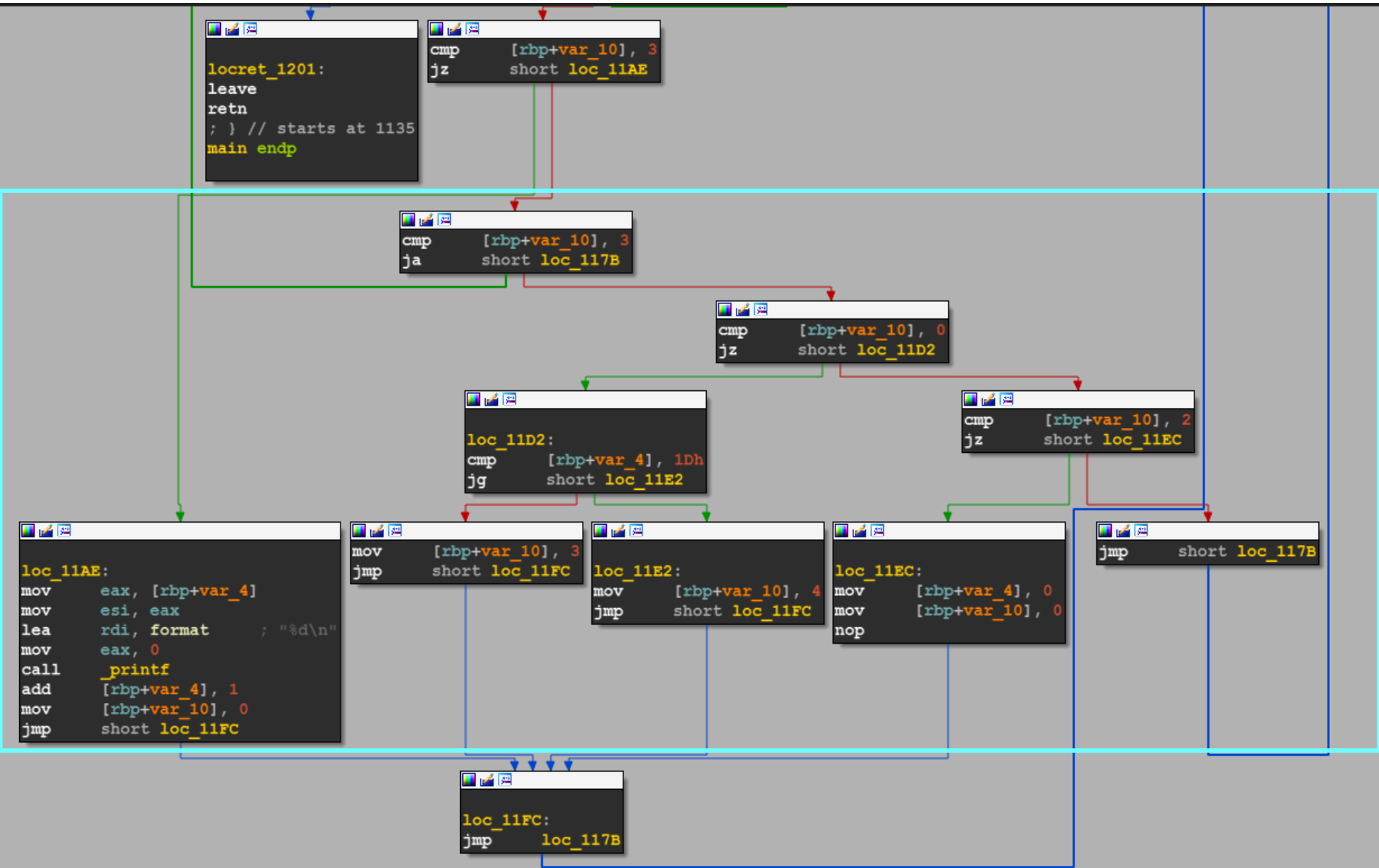
; __unwind {
push    rbp
mov     rbp, rsp
sub     rsp, 40h
mov     [rbp+var_24], edi
mov     [rbp+var_30], rsi
mov     [rbp+var_38], rdx
call    megaInit
mov     eax, [rbp+var_24]
mov     cs:_global_argc, eax
mov     rax, [rbp+var_30]
mov     cs:_global_argv, rax
mov     rax, [rbp+var_38]
mov     cs:_global_envp, rax
mov     [rbp+var_14], 1
mov     [rbp+var_10], 2
mov     [rbp+var_4], 0
}
```

Prologue and  
initial assignment





# ANTI-REVERSING





# ANTI-REVERSING

Simple opaque predicate and anti-disassembly technique

.text:00401000 loc\_401000: ; CODE XREF: \_main+Fp

.text:00401000 push ebp

.text:00401001 mov ebp, esp

.text:00401003 xor eax, eax

.text:00401005 jz short near ptr loc\_40100D+1

.text:00401007 jnz near ptr loc\_40100D+4

.text:0040100D

.text:0040100D loc\_40100D: ; CODE XREF: .text:00401005j

.text:0040100D ; .text:00401007j

.text:0040100D jmp near ptr 0D0A8837h



# ANTI-REVERSING

00401040	call + \$5
00401045	pop ecx
00401046	inc ecx
00401047	inc ecx
00401048	add ecx, 4
00401049	add ecx, 4
0040104A	push ecx
0040104B	ret
0040104C	sub ecx, 6
0040104D	dec ecx
0040104E	dec ecx
0040104F	jmp 0x401320

- Call stack manipulation:
  - Do you know what's happening here? 😊



# ANTI-REVERSING

```
.text:0053CDF1 sub_53CDF1      proc near                ; CODE XREF: sub_53CD66+75↑p
.text:0053CDF1                                     ; sub_54221E+1F↓p ...
.text:0053CDF1 If it's running on
.text:0053CDF3 Windows 8, so die
.text:0053CDF8 immediately. If it's
.text:0053CDFA running on Window 7,
.text:0053CDFC so there's a chance with
.text:0053CDFE Exception Handlers
.text:0053CDFF below. ☺
.text:0053CE01 ; -----
.text:0053CE01
.text:0053CE01 loc_53CE01:                ; CODE XREF: sub_53CDF1+9↑j
.text:0053CE01
.text:0053CE02 push     esi
.text:0053CE04 push     1
.text:0053CE04 mov     esi, 0C0000417h
.text:0053CE09 push     esi
.text:0053CE0A push     2
.text:0053CE0C call    sub_53CC17
.text:0053CE11 add     esp, 0Ch
.text:0053CE14 push     esi                ; uExitCode
.text:0053CE15 call    ds:GetCurrentProcess
.text:0053CE1B push     eax                ; hProcess
.text:0053CE1C call    ds:TerminateProcess
.text:0053CE22 pop     esi
.text:0053CE23 retn
.text:0053CE23 sub_53CDF1      endp
```

Run "lidt 29" in both Windows 7 and 8 to confirm interrupt handlers or not. ☺

The anti-debugging techniques are within this function, which set up exception handlers. The easiest one is IsDebuggerPresent( ), of course.

We usually avoid using ExitProcess( ) or TerminateProcess( ) because possible leakage, but malware's authors don't have this concern. ☺

- Few lines of code could have several anti-debugging tricks. Maybe some non-recommended functions too ☺



# ANTI-REVERSING

- Are there only these **anti-forensic techniques**? No.... 😊
  - Instruction substitution
  - Garbage insertion
  - Self-modifying code (like shellcode)
  - Method renaming
  - Cryptography
  - Variable reordering
  - Instruction reordering
  - Randomized Stack frame
  - Inheritance manipulation
  - And more...much more 😊

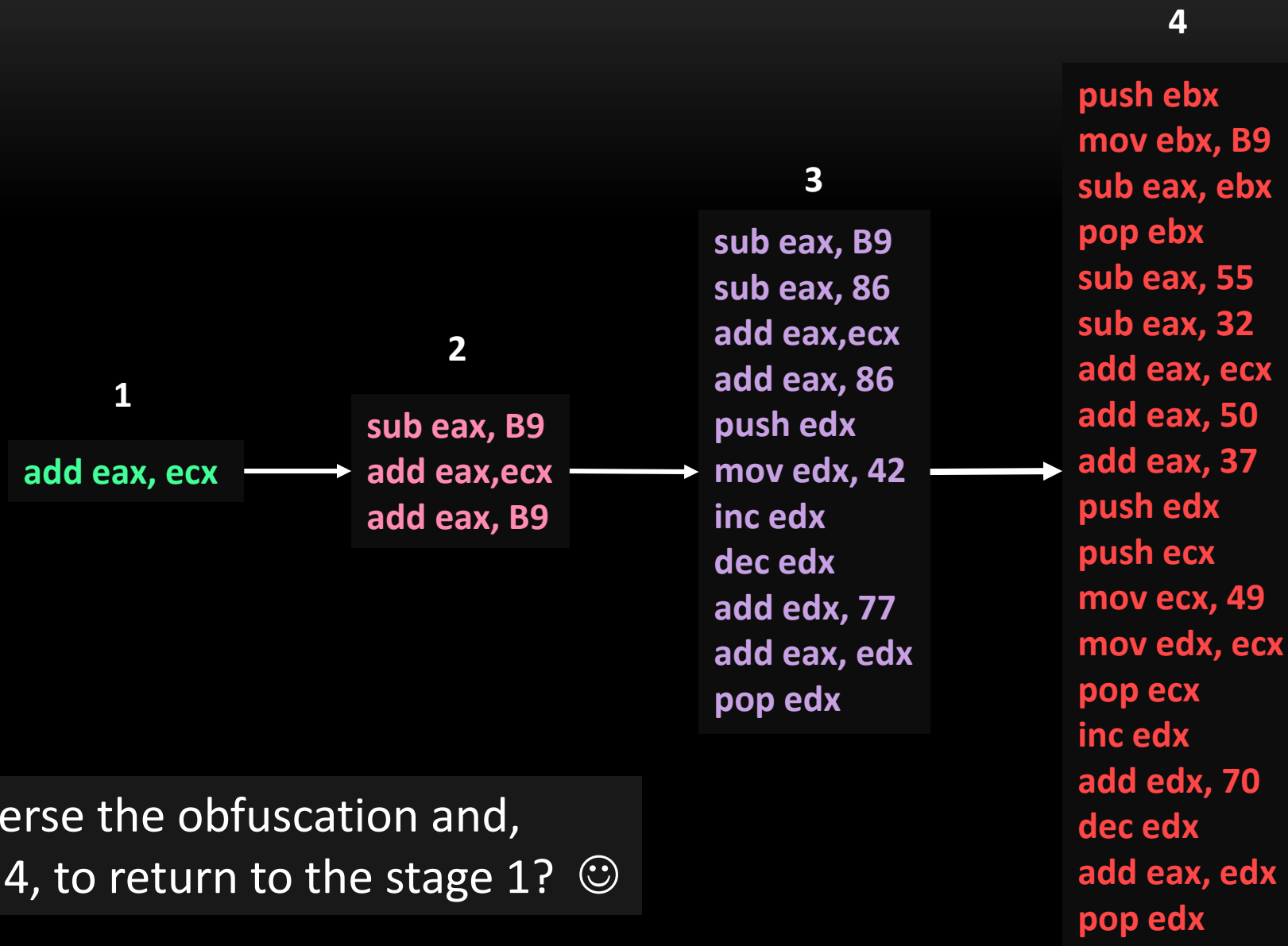
# DEOBFUSCATION

# METASM





# DEOBFUSCATION



- How to reverse the obfuscation and, from stage 4, to return to the stage 1? 😊



# DEOBFUSCATION

- METASM works as disassembler, assembler, debugger, compiler and linker.
- Key features:
  - Written in Ruby
  - C compiler and decompiler
  - Automatic backtracking
  - Live process manipulation
  - Supports the following architecture:
    - Intel IA32 (16/32/64 bits)
    - PPC
    - MIPS



# DEOBFUSCATION

- `root@kali:~/github# git clone https://github.com/jjyg/metasm.git`
- Include the following line into `.bashrc` file to indicate the Metasm directory installation:
  - `export RUBYLIB=$RUBYLIB:~/github/metasm`
- Test metasm:
  - `ruby -r metasm -e 'p Metasm::VERSION'`
- You should see a number in the output. It's done 😊



# DEOBFUSCATION

```
1  #!/usr/bin/env ruby
2  #
3
4  require "metasm"
5  include Metasm
6
7  mycode = Metasm::Shellcode.assemble(Metasm::Ia32.new, <<EOB)
8
9  entry:
10
11     push ebx
12     mov ebx, 0xb9
13     sub eax, ebx
14     pop ebx
15     sub eax, 0x55
16     sub eax, 0x32
17     add eax, ecx
18     add eax, 0x50
19     add eax, 0x37
20     push edx
21     push ecx
22     mov ecx, 0x49
23     mov edx, ecx
24     pop ecx
25     inc edx
26     add edx, 0x70
27     dec edx
28     add eax, edx
29     pop edx
30     jmp eax
31
32  EOB
```

- Based on metasm.rb file and Bruce Dang's code.

- This instruction was inserted to make the eax register evaluation easier. 😊



# DEOBFUSCATION

```
33
34 addrstart = 0
35 asmcode = mycode.init_disassembler
36 asmcode.disassemble(addrstart)
37 conference_di = asmcode.di_at(addrstart)
38 conference = conference_di.block
39 puts "\n<!!!> Blackstorm Security 2020:\n "
40 puts conference.list
41
42 conference.list.each{|aborges|
43     puts "\n<!!!> #{aborges.instruction}"
44     back = aborges.backtrace_binding()
45     v = back.values
46     k = back.keys
47     j = k.zip(v)
48     puts "Our data flow follows below:\n"
49     j.each do |mykeys, myvalues|
50         puts "Processing: #{mykeys} ==> #{myvalues}"
51
52         if aborges.opcode.props[:setip]
53             puts "\nOur control flow follows below:\n"
54             puts ">>> #{asmcode.get_xrefs_x(aborges)}"
55         end
56     end
57 }
```

- initialize and disassemble code since beginning (start).

- list the assemble code

- initialize the backtracking engine

- determines which is the final instruction to walk back from there



# DEOBFUSCATION

- Backtracing from the last instruction

- Logs the sequence of backtraced instructions.

```
58
59 addrstart2 = 0
60 asmcode2 = mycode.init_disassembler
61 asmcode2.disassemble(addrstart2)
62 dd = asmcode2.block_at(addrstart2)
63 final = asmcode2.get_xrefs_x(dd.list.last).first
64 puts "\n[+] final output: #{final}"
65
66 values = asmcode2.backtrace(final, dd.list.last.address, {:log => backtracing_log = [] , :include_start => true})
67 backtracing_log.each{|record|
68   case type = record.first
69   when :start
70     record, expression, addresses = record
71     puts "[start] Here is the sequence of expression evaluations #{expression} from 0x#{addresses.to_s(16)}\n"
72
73   when :di
74     record, new, old, instruction = record
75     puts "[new update] instruction #{instruction},\n --> updating expression once again from #{old} to #{new}\n"
76
77   end
78 }
79
80 effective = backtracing_log.select{|y| y.first==:di}.map{|y| y[3]}.reverse
81 puts "\nThe effective instructions are:\n\n"
82 puts effective
83
84
85
86
```

- Shows only the effective instructions, which really can alter the final result.



# DEOBFUSCATION

```
root@kali:~# ./metasmtest.rb
```

```
<!!!> Blackstorm Security 2020:
```

```
0 push ebx
1 mov ebx, 0b9h
6 sub eax, ebx
8 pop ebx
9 sub eax, 55h
0ch sub eax, 32h
0fh add eax, ecx
11h add eax, 50h
14h add eax, 37h
17h push edx
18h push ecx
19h mov ecx, 49h
1eh mov edx, ecx
20h pop ecx
21h inc edx
22h add edx, 70h
25h dec edx
26h add eax, edx
28h pop edx
29h jmp eax
```

Remember: this is our obfuscated code. 😊



# DEOBFUSCATION

```
<!!!> push ebx
```

Our data flow follows below:

Processing: esp  $\Rightarrow$  esp-4

Processing: dword ptr [esp]  $\Rightarrow$  ebx

```
<!!!> mov ebx, 0b9h
```

Our data flow follows below:

Processing: ebx  $\Rightarrow$  0b9h

```
<!!!> sub eax, ebx
```

Our data flow follows below:

Processing: eax  $\Rightarrow$  eax-ebx

Processing: eflag\_z  $\Rightarrow$  (((eax&0xffffffff)-(ebx&0xffffffff))&0xffffffff)=0

Processing: eflag\_s  $\Rightarrow$  (((eax&0xffffffff)-(ebx&0xffffffff))&0xffffffff)>>1fh) $\neq$ 0

Processing: eflag\_c  $\Rightarrow$  (eax&0xffffffff)<(ebx&0xffffffff)

Processing: eflag\_o  $\Rightarrow$  (((eax&0xffffffff)>>1fh) $\neq$ 0)=(!(((ebx&0xffffffff)>>1fh) $\neq$ 0)))&&(((eax&0xffffffff)>>1fh) $\neq$ 0) $\neq$ (((eax&0xffffffff)-(ebx&0xffffffff))&0xffffffff)>>1fh) $\neq$ 0))

```
<!!!> pop ebx
```

Our data flow follows below:

Processing: esp  $\Rightarrow$  esp+4

Processing: ebx  $\Rightarrow$  dword ptr [esp]

```
<!!!> sub eax, 55h
```

Our data flow follows below:

Processing: eax  $\Rightarrow$  eax-55h





# DEOBFUSCATION

Our control flow follows below:

```
>>> [Expression[:eax]]
```

[+] final output: eax

```
[start] Here is the sequence of expression evaluations eax from 0x29
[new update] instruction 26h add eax, edx,
→ updating expression once again from eax to eax+edx
[new update] instruction 25h dec edx,
→ updating expression once again from eax+edx to eax+edx-1
[new update] instruction 22h add edx, 70h,
→ updating expression once again from eax+edx-1 to eax+edx+6fh
[new update] instruction 21h inc edx,
→ updating expression once again from eax+edx+6fh to eax+edx+70h
[new update] instruction 1eh mov edx, ecx,
→ updating expression once again from eax+edx+70h to eax+ecx+70h
[new update] instruction 19h mov ecx, 49h,
→ updating expression once again from eax+ecx+70h to eax+0b9h
[new update] instruction 14h add eax, 37h,
→ updating expression once again from eax+0b9h to eax+0f0h
[new update] instruction 11h add eax, 50h,
→ updating expression once again from eax+0f0h to eax+140h
[new update] instruction 0fh add eax, ecx,
→ updating expression once again from eax+140h to eax+ecx+140h
[new update] instruction 0ch sub eax, 32h,
→ updating expression once again from eax+ecx+140h to eax+ecx+10eh
```



# DEOBFUSCATION

[new update] instruction 9 sub eax, 55h,  
→ updating expression once again from `eax+ecx+10eh` to `eax+ecx+0b9h`  
[new update] instruction 6 sub eax, ebx,  
→ updating expression once again from `eax+ecx+0b9h` to `eax-ebx+ecx+0b9h`  
[new update] instruction 1 mov ebx, 0b9h,  
→ updating expression once again from `eax-ebx+ecx+0b9h` to `eax+ecx`

Game over 😊

The effective instructions are:

```
1 mov ebx, 0b9h
6 sub eax, ebx
9 sub eax, 55h
0ch sub eax, 32h
0fh add eax, ecx
11h add eax, 50h
14h add eax, 37h
19h mov ecx, 49h
1eh mov edx, ecx
21h inc edx
22h add edx, 70h
25h dec edx
26h add eax, edx
```

■ Output originated from `backtracing_log.select` command (in reverse)

# MIASM



# DEOBFUSCATION

- **MIASM** is one of most impressive framework for reverse engineering, which is able to **analyze, generate and modify** several different types of programs.
- **MIASM** supports **assembling and disassembling** programs from different platforms such as **ARM, x86, MIPS** and so on, and it also is able to **emulate by using JIT**.
- Therefore, **MIASM** is excellent to de-obfuscation.
- Installing MIASM (python 2.7.x):
  - `git clone https://github.com/serpilliere/elfesteem.git elfesteem`
  - `cd elfesteem/`



# DEOBFUSCATION

- `python setup.py build`
- `python setup.py install`
- `apt-get install clang texinfo texi2html`
- `apt-get remove libtcc-dev`
- `apt-get install llvm`
- `cd ..`
- `git clone http://repo.or.cz/tinycc.git`
- `cd tinycc/`
- `git checkout release_0_9_26`
- `./configure --disable-static`
- `make`
- `make install`



# DEOBFUSCATION

- pip install llvmlite
- pip install future
- apt-get install z3
- apt-get install python-pycparser
- pip install pyparsing
- cd ..
- git clone <https://github.com/cea-sec/miasm.git>
- cd miasm
- python setup.py build
- python setup.py install
- cd test/
- python test\_all.py



# DEOBFUSCATION

```
root@kali:~/github/miasm/test# python2.7 test_all.py
[LLVM] Python'llvmlite' module is required for llvm tests
[Z3] Z3 and its python binding are necessary for TranslatorZ3.
TEST/ARCH msp430/arch.py
TEST/ARCH ppc32/arch.py
TEST/ARCH x86/arch.py
TEST/ARCH aarch64/arch.py
DONE msp430/arch.py 0s
TEST/ARCH x86/sem.py python
DONE ppc32/arch.py 0s
TEST/ARCH x86/unit/mn_strings.py python
DONE x86/sem.py python 1s
TEST/ARCH x86/unit/mn_stack.py gcc
DONE x86/unit/mn_strings.py python 1s
TEST/ARCH x86/unit/mn_daa.py gcc
DONE x86/unit/mn_daa.py gcc 3s
TEST/ARCH x86/unit/mn_das.py gcc
DONE x86/unit/mn_stack.py gcc 3s
TEST/ARCH x86/unit/mn_int.py gcc
DONE x86/unit/mn_int.py gcc 1s
TEST/ARCH x86/unit/mn_pshufb.py gcc
DONE x86/unit/mn_pshufb.py gcc 1s
TEST/ARCH x86/unit/mn_psrl_psll.py gcc
```



# DEOBFUSCATION

- Before proceeding with Miasm, we need to create a binary containing our code, so we need an assembler and Keystone is great.
- Keystone Engine acts as an assembler and:
  - Supports x86, Mips, Arm and many other architectures.
  - It is implemented in C/C++ and has bindings to Python, Ruby, Powershell and C# (among other languages).
- Installing Keystone:
  - `root@kali:~/Desktop# wget https://github.com/keystone-engine/keystone/archive/0.9.1.tar.gz`





# DEOBFUSCATION

- root@kali:~/programs# `cp /root/Desktop/keystone-0.9.1.tar.gz .`
- root@kali:~/programs# `tar -zxvf keystone-0.9.1.tar.gz`
- root@kali:~/programs/keystone-0.9.1# `apt-get install cmake`
- root@kali:~/programs/keystone-0.9.1# `mkdir build ; cd build`
- root@kali:~/programs/keystone-0.9.1/build# `apt-get install time`
- root@kali:~/programs/keystone-0.9.1/build# `../make-share.sh`
- root@kali:~/programs/keystone-0.9.1/build# `make install`
- root@kali:~/programs/keystone-0.9.1/build# `ldconfig`
- root@kali:~/programs/keystone-0.9.1/build# `tail -2 /root/.bashrc`

```
export RUBYLIB=$RUBYLIB:~/programs/metasm
```

```
export LD_LIBRARY_PATH=$LD_LIBRARY_PATH:/usr/local/lib
```



# DEOBFUSCATION

```
1  #include <stdio.h>
2  #include <keystone/keystone.h>
3
4  #define SANS "push ebx; mov ebx,0xb9; sub eax,ebx; pop ebx; sub eax,0x55; sub
5  #eax,0x32; add eax,ecx; add eax,0x50; add eax,0x37; push edx; push ecx; mov
6  #ecx,0x49; mov edx,ecx; pop ecx; inc edx; add edx,0x70; dec edx; add eax,edx;
7  #pop edx"
8
9  int main(int argc, char **argv)
10 {
11     ks_engine *keyeng;
12     ks_err keyerr = KS_ERR_ARCH;
13     size_t count;
14     unsigned char *encode;
15     size_t size;
16
17     keyerr = ks_open(KS_ARCH_X86, KS_MODE_32, &keyeng);
18     if (keyerr != KS_ERR_OK) {
19         printf("ERROR: A fail has ocurred while calling ks_open(), quit\n");
20         return -1;
21     }
22
23     if (ks_asm(keyeng, SANS, 0, &encode, &size, &count)) { printf("ERROR: A fail has occurred while calling ks_asm() with
24     count = %lu, error code = %u\n", count, ks_errno(keyeng)); } else { size_t i;
25
26         for (i=0; i < size; i++) {
27             printf("%02x ", encode[i]);
28         }
29     }
30
31     ks_free(encode);
32     ks_close(keyeng);
33
34     return 0;
35 }
```

- instructions from the original obfuscated code

- Creating a keystone engine

- Assembling our instructions using Keystone engine.

- Freeing memory and closing engine.



# DEOBFUSCATION

```
root@kali:~/conference# more Makefile
.PHONY: all clean

KEYSTONE_LDFLAGS = -lkeystone -lstdc++ -lm

all:
    ${CC} -o conference2020 conference2020.c ${KEYSTONE_LDFLAGS}

clean:
    rm -rf *.o conference2020

root@kali:~/conference#
root@kali:~/conference# make
cc -o conference2020 conference2020.c -lkeystone -lstdc++ -lm
root@kali:~/conference#
root@kali:~/conference# ./conference2020
53 bb b9 00 00 00 29 d8 5b 83 e8 55 83 e8 32 01 c8 83 c0 50 83 c0 37 52 51 b9 49 00 00 00
89 ca 59 42 83 c2 70 4a 01 d0 5a root@kali:~/conference#
root@kali:~/conference#
root@kali:~/conference# ./conference2020 | xxd -r -p - > conference2020.bin
root@kali:~/conference#
root@kali:~/conference# hexdump -C conference2020.bin
00000000  53 bb b9 00 00 00 29 d8 5b 83 e8 55 83 e8 32 01 |S.....).[..U..2.|
00000010  c8 83 c0 50 83 c0 37 52 51 b9 49 00 00 00 89 ca |...P..7RQ.I.....|
00000020  59 42 83 c2 70 4a 01 d0 5a                        |YB..pJ..Z|
00000029
```



# DEOBFUSCATION

```
seg000:00000000 ; Format      : Binary file
seg000:00000000 ; Base Address: 0000h Range: 0000h - 0029h Loaded length: 0029h
seg000:00000000
seg000:00000000      .686p
seg000:00000000      .mmx
seg000:00000000      .model flat
seg000:00000000
seg000:00000000 ; =====
seg000:00000000
seg000:00000000 ; Segment type: Pure code
seg000:00000000 seg000      segment byte public 'CODE' use32
seg000:00000000      assume cs:seg000
seg000:00000000      assume es:nothing, ss:nothing, ds:nothing, fs:nothing, gs:nothing
seg000:00000000      push     ebx
seg000:00000001      mov      ebx, 0B9h
seg000:00000006      sub      eax, ebx
seg000:00000008      pop      ebx
seg000:00000009      sub      eax, 55h
seg000:0000000C      sub      eax, 32h
seg000:0000000F      add      eax, ecx
seg000:00000011      add      eax, 50h
seg000:00000014      add      eax, 37h
seg000:00000017      push     edx
seg000:00000018      push     ecx
seg000:00000019      mov      ecx, 49h
seg000:0000001E      mov      edx, ecx
seg000:00000020      pop      ecx
seg000:00000021      inc      edx
seg000:00000022      add      edx, 70h
seg000:00000025      dec      edx
seg000:00000026      add      eax, edx
seg000:00000028      pop      edx
seg000:00000028 seg000      ends
```

IDA Pro confirms: it's our content 😊



# DEOBFUSCATION

```
1 from colorama import init, Fore, Back, Style
2 from miasm.analysis.binary import Container
3 from miasm.analysis.machine import Machine
4 from miasm.jitter.csts import PAGE_READ, PAGE_WRITE
```

Open the binary file generated through Keystone. The Container provides the byte source to the disasm engine.

```
6 with open("/root/conference/conference2020.bin") as fdesc:
7     cont=Container.from_stream(fdesc)
```

Instantiates the assemble engine using the x86 32-bits architecture.

```
9 offset = 0
```

```
10 confrencemach=Machine('x86_32')
```

```
11 confrencedis=confrencemach.dis_engine(cont.bin_stream)
```

```
12 asmcfg = confrencedis.dis_multiblock(offset)
```

```
13 ira = confrencemach.ira(confrencedis.loc_db)
```

Initialize and run the Symbolic Execution Engine, setting all registers to an initial value.

```
15 ircfg = ira.new_ircfg_from_asmcfg(asmcfg)
```

```
16 open('out.dot', 'w').write(ircfg.dot())
```

```
18 from miasm.ir.symbexec import SymbolicExecutionEngine
```

```
19 symb = SymbolicExecutionEngine(ira, confrencemach.mn.regs.regs_init)
```

```
20 symbolic_pc = symb.run_at(ircfg, 0, step=True)
```

```
21 print(Fore.YELLOW + "\nThe final value of EAX register is: %s" %
```

```
22 symb.symbols[confrencemach.mn.regs.EAX])
```

Print the final value of EAX.

```
23 print (Fore.RESET)
```



# DEOBFUSCATION

```
root@kali:~# python conference_symbolic.py | more
```

```
[WARNING ]: not enough bytes in str
```

```
[WARNING ]: cannot disasm at 29
```

```
[WARNING ]: not enough bytes in str
```

```
[WARNING ]: cannot disasm at 29
```

```
Instr PUSH      EBX
```

```
Assignblk:
```

```
ESP = ESP + -0x4
```

```
@32[ESP + -0x4] = EBX
```

---

```
R12      = R12_init
MM2      = MM2_init
FS       = FS_init
XMM8     = XMM8_init
AL       = AL_init
float_c3 = float_c3_init
ST(0)    = ST(0)_init
EAX      = EAX_init
cf       = cf_init
R8W      = R8W_init
float_st6 = float_st6_init
MM1      = MM1_init
```





# DEOBFUSCATION

Instr MOV EBX, 0xB9

Assignblk:

EBX = 0xB9

---

R12	= R12_init
MM2	= MM2_init
FS	= FS_init
XMM8	= XMM8_init
AL	= AL_init
float_c3	= float_c3_init
ST(0)	= ST(0)_init
EAX	= EAX_init
cf	= cf_init
R8W	= R8W_init
float_st6	= float_st6_init
MM1	= MM1_init
pf	= pf_init
R15B	= R15B_init
XMM10	= XMM10_init
SP	= SP_init
zf	= zf_init
XMM15	= XMM15_init
BL	= BL_init



# DEOBFUSCATION

```
DR3      = DR3_init
R14      = R14_init
XMM11    = XMM11_init
CX       = CX_init
R13B     = R13B_init
DR2      = DR2_init
RCX      = RCX_init
i_d      = i_d_init
XMM1     = XMM1_init
RSI      = RSI_init
R8       = R8_init
DI       = DI_init
RAX      = RAX_init
float_st4 = float_st4_init
XMM6     = XMM6_init
R15      = R15_init
MM7      = MM7_init
SPL      = SPL_init
@32[ESP_init + 0xFFFFFFFF8] = ECX_init
@32[ESP_init + 0xFFFFFFFFC] = EDX_init
```

---

- I've skipped a very long output, which shows all instruction being execute. We are interested in the final value of EAX. 😊
- If you want, view the graph:
  - `apt-get install graphviz`
  - `apt-get install xdot`
  - `xdot out.dot`

The final value of EAX register is: `EAX_init + ECX_init` ← Finally.... 😊



# ➤ Closing thoughts and Acknowledgments...

- Honestly, I didn't even scratch the surface of this topic....
- There're tons of obfuscation techniques and de-obfuscation/reversing techniques to explain... it would take one or two entire weeks...and probaly you wouldn't to know about it... 😊
- I'd like to thank SANS for the event and, in special, my friend Stephen Sims for the invite. 😊
- And, of course, I'd like to thank you (the audience) for attending my talk.



# THANK YOU FOR ATTENDING MY TALK 😊

## ➤ Twitter:

@ale\_sp\_brazil  
@blackstormsecbr

➤ Website: <http://www.blackstormsecurity.com>

➤ LinkedIn: <http://www.linkedin.com/in/aleborges>

➤ E-mail: [alexandreborges@blackstormsecurity.com](mailto:alexandreborges@blackstormsecurity.com)

- Security Researcher
- Speaker at DEF CON USA 2019
- Speaker at DEF CON USA 2018
- Speaker at DEF CON CHINA 2019
- Speaker at DC2711 (Johannesburg)
- Speaker at NO HAT 2019 (Italy)
- Speaker at HITB 2019 (Amsterdam)
- Speaker at CONFidence 2019 (Poland)
- Speaker at DevOpsDays BH 2019
- Speaker at BSIDES 2019/2018/2017/2016
- Speaker at H2HC 2016/2015
- Speaker at BHACK 2018
- Researcher on Android/iOS Reversing, Rootkits and Digital Forensics.
- Referee on Digital Investigation: The International Journal of Digital Forensics & Incident Response