RS/Conference2020

San Francisco | February 24 – 28 | Moscone Center



SESSION ID: CRYP-F01

Policy-Based Sanitizable Signatures



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Outline

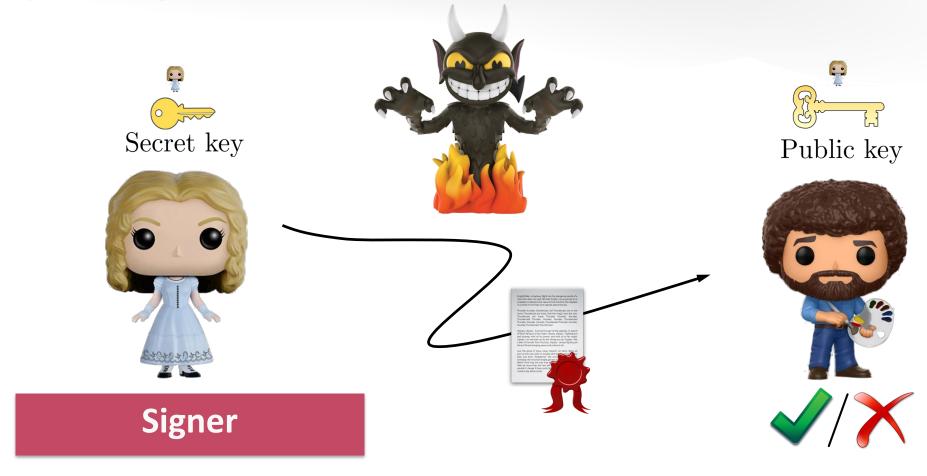
- Digital Signatures
- Sanitizable Signatures
- Policy-Based Sanitizable Signatures
- Conclusions & Take Home



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Digital Signatures

Digital Signatures



Modifications invalidate signature



Controlled Modifications of Signed Messages

- Modify signed messages without invalidating the signature?
 - But that's what we try to prevent?!
 - Can be useful if controlled!

- Controlled modifications
 - Signer determines how signed message can be altered
 - Think of implicitly signing all possible messages
- Control who is allowed to modify
 - Signer specifies entity allowed to perform modifications



Example: Medical Documents

Tokenization for research/accounting

Removing exact diagnosis for sick leave

Etc.

irst Choice

Re-signing after the fact might not be possible (availability, etc.)



Different Types of Schemes

- Redactable Signatures
 - Blacking out/Removal of designated parts by everyone

Sanitizable Signatures
 Replacement of designated parts by designated entity
 Designated entity (=sanitizer) has its own key pair

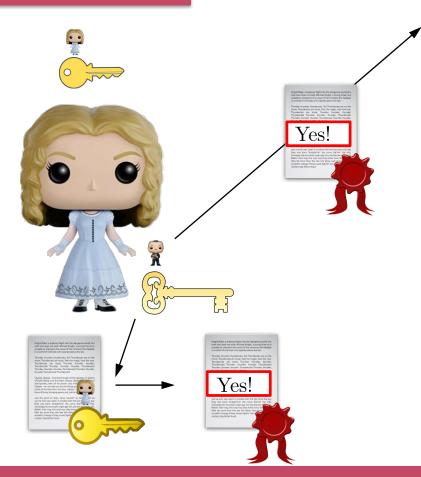


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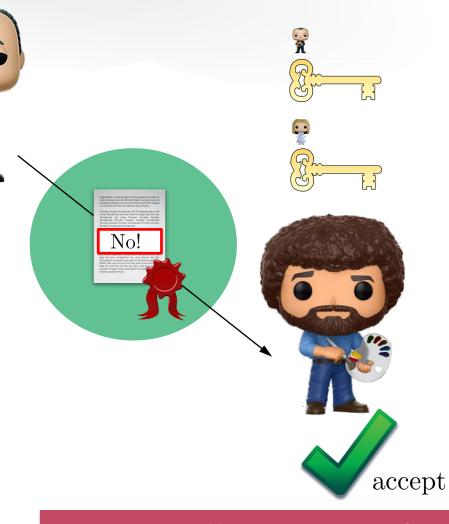
Sanitizable Signatures

Sanitizable Signatures

Signer



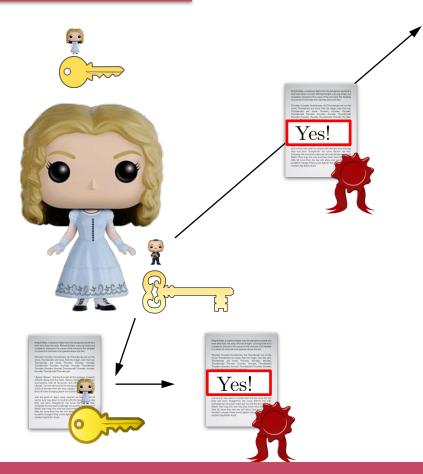
Sanitizer



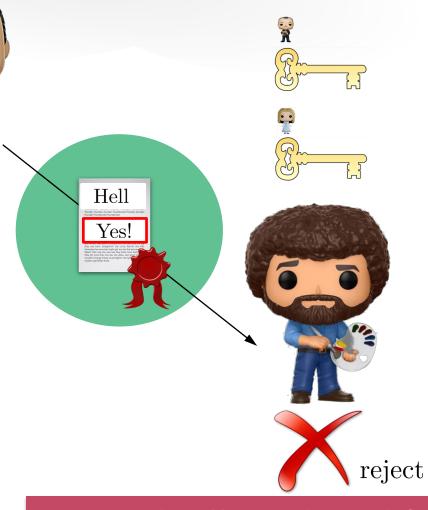
Accept allowed modifications

Sanitizable Signatures

Signer



Sanitizer



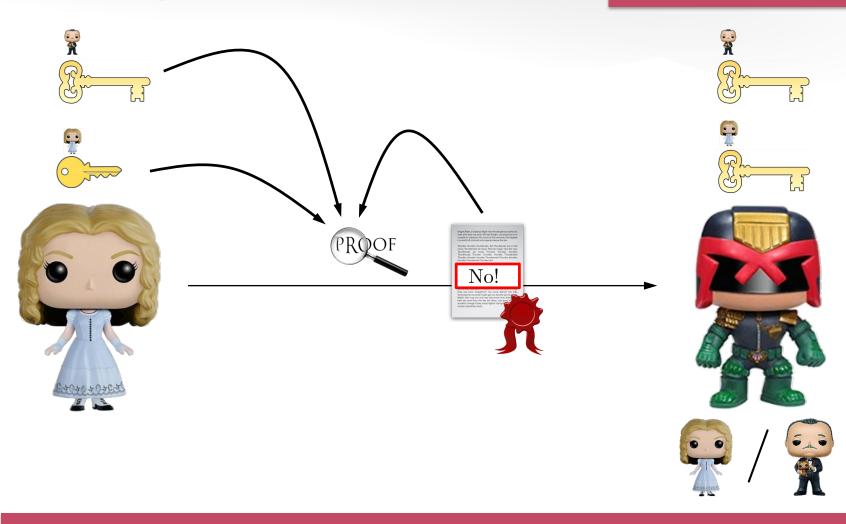
Reject disallowed modifications

Specify modifications and sanitizer

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Sanitizable Signatures

Judge



Generate proof that signature from signer or sanitizer



Security Properties

- Unforgeability
 - Non-sanitizers cannot come up with valid signature for non-signed message
- Immutability
 - Sanitizer cannot come up with valid signature for a message not "derivable" from signed ones
- Privacy
 - No information about sanitized parts can be learned
- Signer/Sanitizer accountability
 - Signer/sanitizer cannot blame the other party for having produced a signature
- Transparancy
 - Freshly signed and sanitized signatures are indistinguishable



Construction Idea I/II

- Originally proposed in [ACMT, ESORICS'05] and rigorous security model in [BFFLP+, PKC'09]
- Generic construction from secure signatures and chameleon-hash functions in [BFFLP+, PKC'09]

- Chameleon Hash [KT, NDSS'00]: Collision-resistant hash keyed with (sk, pk)
 - Hashing: $h \leftarrow CHash(pk,m; r)$
 - Collision: sk allows for any h, m' to compute r' s.t.

CHash(pk,m; r) = CHash(pk, m';r')



Construction Idea II/II

- Simplified construction idea:
 - To sign $m = (m_1, ..., m_n)$
 - Use signature scheme to sign $h = (h_1, ..., h_n)$, where

$$h_i = \begin{cases} CHash(pk, m_i; r_i) & if \ sanitizable \\ m_i & else \end{cases}$$

– As sanitizable signature provide signature on h; additionally include the randomness \mathbf{r}_i of sanitizable parts



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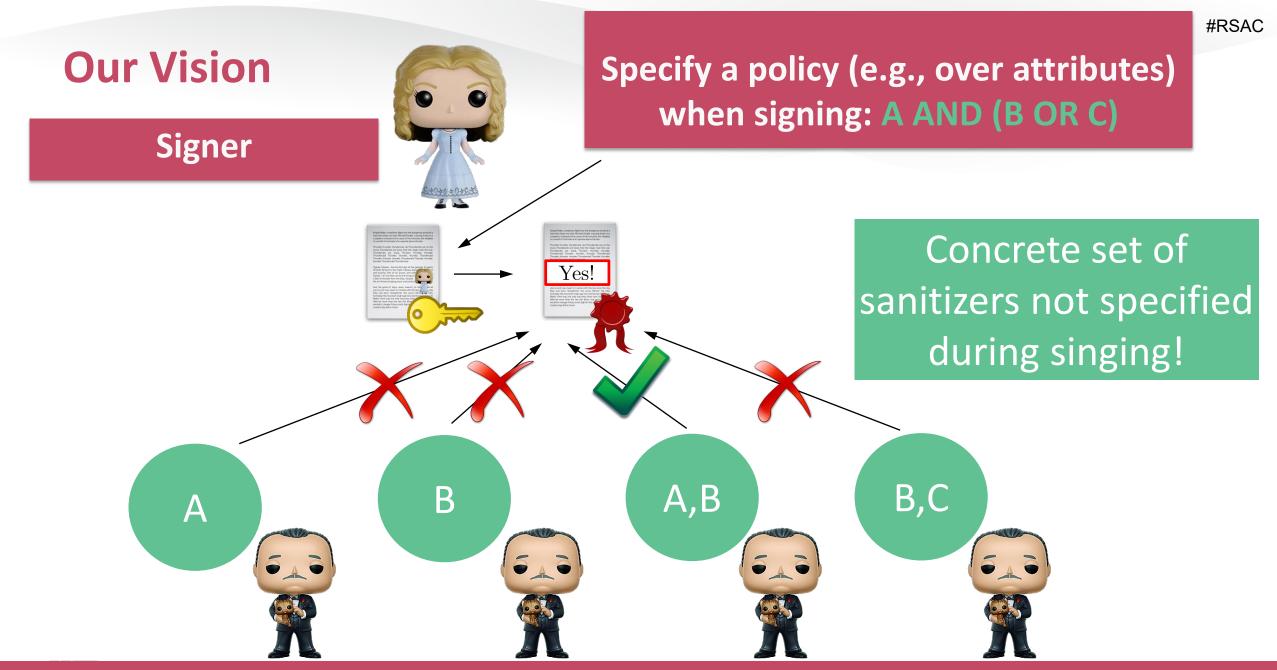
Policy-Based Sanitizable Signatures

Drawbacks of Existing Sanitizable Signatures

- Restricted in flexibility of specifying potential sanitizers
 - Conventional model just considers a single sanitizer specified at signing time
- Some existing works allow multiple sanitizers, but then lose accountability
 - Accountability is a central feature: sanitizers must be traceable

Achieve fine-grained sanitization control with full accountability?





Sanitizers have attributes: if attributes satisfy policy, then sanitizing works

Policy-Based Sanitizable Signatures (PB-SS)

- At signing time signer specifies an access-policy for sanitization and a sanitization group
- Sanitizers with their sanitizer key pair can join sanitization groups dynamically
 - Get then issued sanitization keys for attributes in the group
 - Sanitization works if sanitizer is in the group and attributes satisfy the policy
- Accountability: exact sanitizer can be identified

Achieves fine-grained sanitization control with full accountability!



Security of PB-SS

- Require same properties (but extended) as for sanitizable signatures plus
- Pseudonymity
 - Signature does not leak which party is accountable
- Proof-Soundness
 - Impossible to generate a proof for an adverserially chosen signature/message pair that points to different entities
- Traceability
 - Impossible to generate a verifying signature such that an honest signer cannot identify the accountable party

We also further strengthen (existing) notions whenever possible



Construction Idea: PB-SS

- Follows basically the idea outlined for sanitizable signatures
- Instead of using a chameleon hash we use a strengthened version of a policy-based chameleon hash [DSSS, NDSS'19]
 - Hash is computed with respect to a policy
 - Collisions can be found when policy is satisfied
 - Can be constructed from chameleon hashes with ephemeral trapdoors [CDKPSS, PKC'17] and CCA secure ciphertext-policy attribute-based encryption (CP-ABE) [BSW, Oakland'07]
- Achieving accountability requires some additional technicalities
 - Use of a non-interactive zero-knowledge proof (NIZK) for an OR-language



Potential Application of PB-SS

- Applicable to all existing applications
- A new application can be redactable blockchains
 - Introduced by [AMVA, EuroS&P'17]
 - [DSSS, NDSS'19] introduce the use of PBCH to hash transactions in blockchains
 - Update/rewrite transactions by computing collisions in the PBCH
 - Fine-grained approach by using policies
 - Use of PB-SS instead of PBCH to get additional properties such as accountability



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Conclusions and Take Home

Conclusions & Take Home

- Sanitizable signatures are a tool to realize controlled modifications of signed messages
- Existing schemes are very limited in their expressiveness
- We introduce the notion of policy-based sanitizable signatures
 - Fine-grained sanitization control via policies
- We present a strong security model and a provable secure practical construction

