

# Serial form of Data Vs Parallel Form of Data

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- ✓ Data may be available in Parallel form or Serial form.
- ✓ Multi bit data is said to be in parallel form when all the bits are available (accessible) simultaneously.
- ✓ The data is said to be in serial form when data bits appear sequentially (one after another in time) at a single terminal.
- ✓ Data may also be transferred in parallel form or in serial form.

# Data Transmission Serial/Parallel

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- ✓ Parallel data transfer is the simultaneous transmission of all bits of data from one device to another.
- ✓ Serial data transfer is the transmission of one bit of data at time from one device to another.
- ✓ Serial data must be transmitted under the synchronization of a clock, since clock provides the means to specify the time at which each new bit is sampled

# Register

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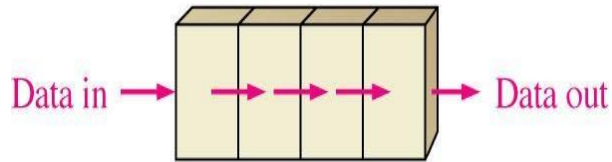
- ✓ As a flip flop can store only one bit of data, a 0 or a 1, it is referred as a single bit register.
- ✓ When more bits of data are to be stored, a number of FFs used.
- ✓ A register is a set of FFs used to store binary data.
- ✓ The storage capacity of a register is the number of bits (1s and 0s) of digital data it can retain.
- ✓ A register may have output data either in serial form or in parallel form.

# Shift Register

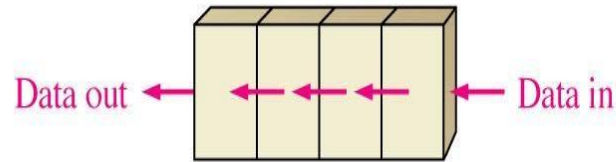
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- ✓ A shift register is a very important digital building blocks. It has innumerable applications.
- ✓ Shift registers are a type of logic circuits closely related to counters.
- ✓ They are used basically for storage and transfer of digital data.
- ✓ The basic difference between a shift register and a counter is that, a shift register has no specified sequence of states whereas a counter has a specified sequence of states.

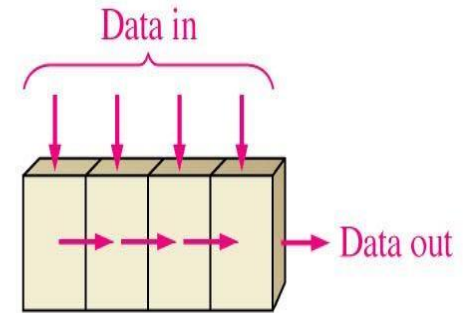
# Basic Data Movements in Shift Registers



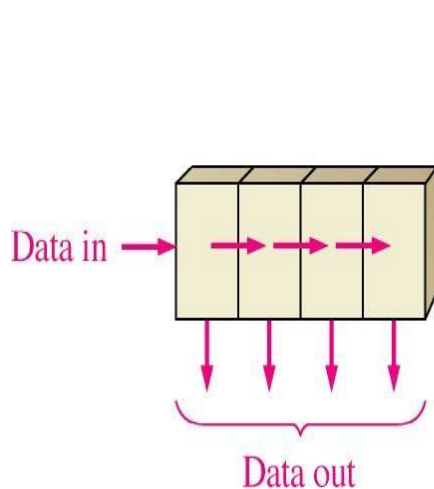
(a) Serial in/shift right/serial out



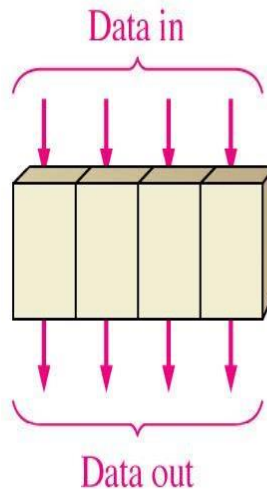
(b) Serial in/shift left/serial out



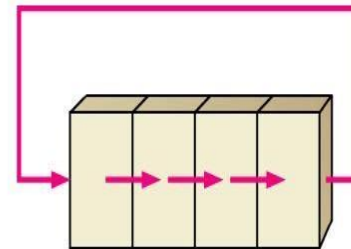
(c) Parallel in/serial out



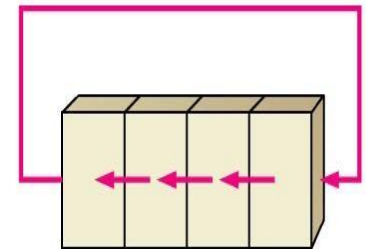
(d) Serial in/parallel out



(e) Parallel in/parallel out



(f) Rotate right



(g) Rotate left

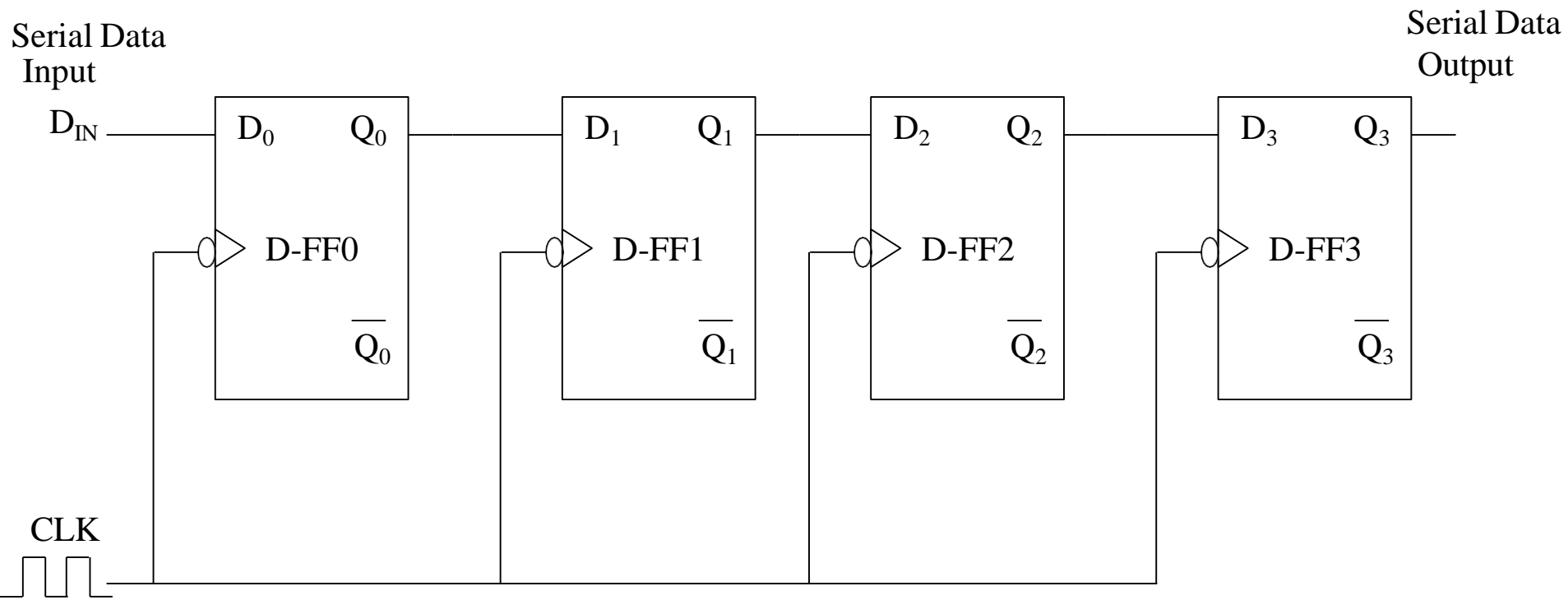
# Types of Shift Registers

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- ✓ SISO – Serial In Serial Out Shift Register
- ✓ SIPO – Serial In Parallel Out Shift Register
- ✓ PISO – Parallel In Serial Out Shift Register
- ✓ PIPO – Parallel In Parallel Out Shift Register
- ✓ Bi-directional Shift Register
- ✓ Universal Shift Register

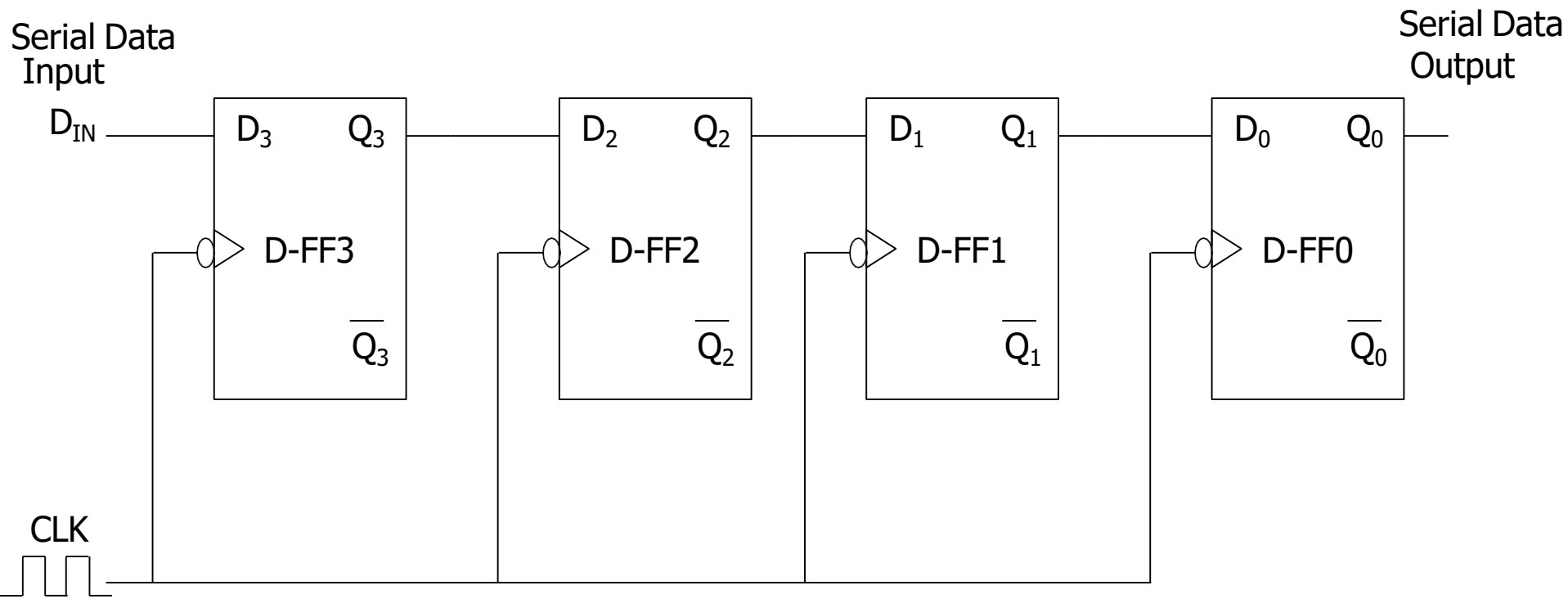
# SISO – Serial In Serial Out Shift Register (Shift Left)

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# SISO – Serial In Serial Out Shift Register (Shift Right)

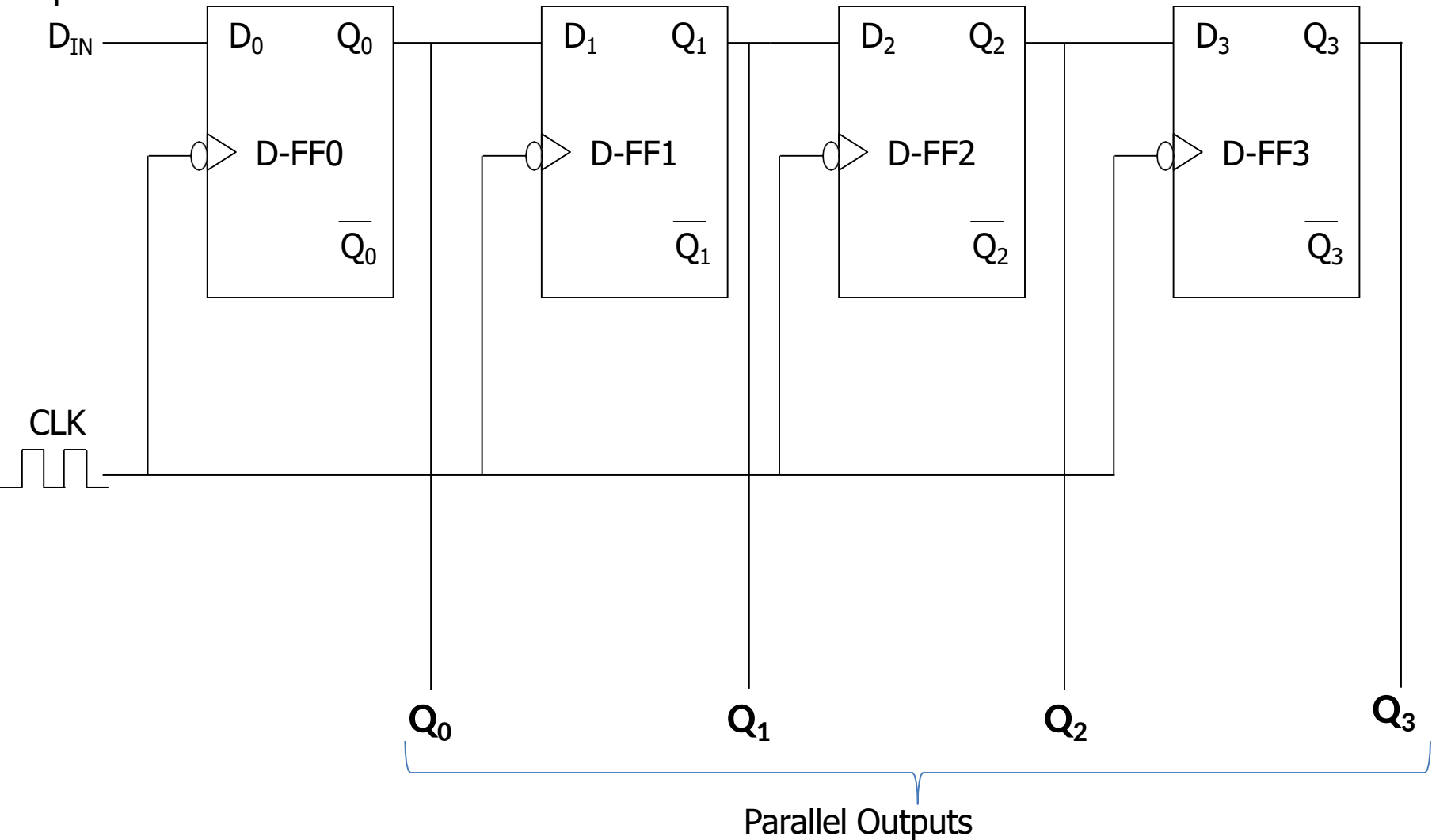
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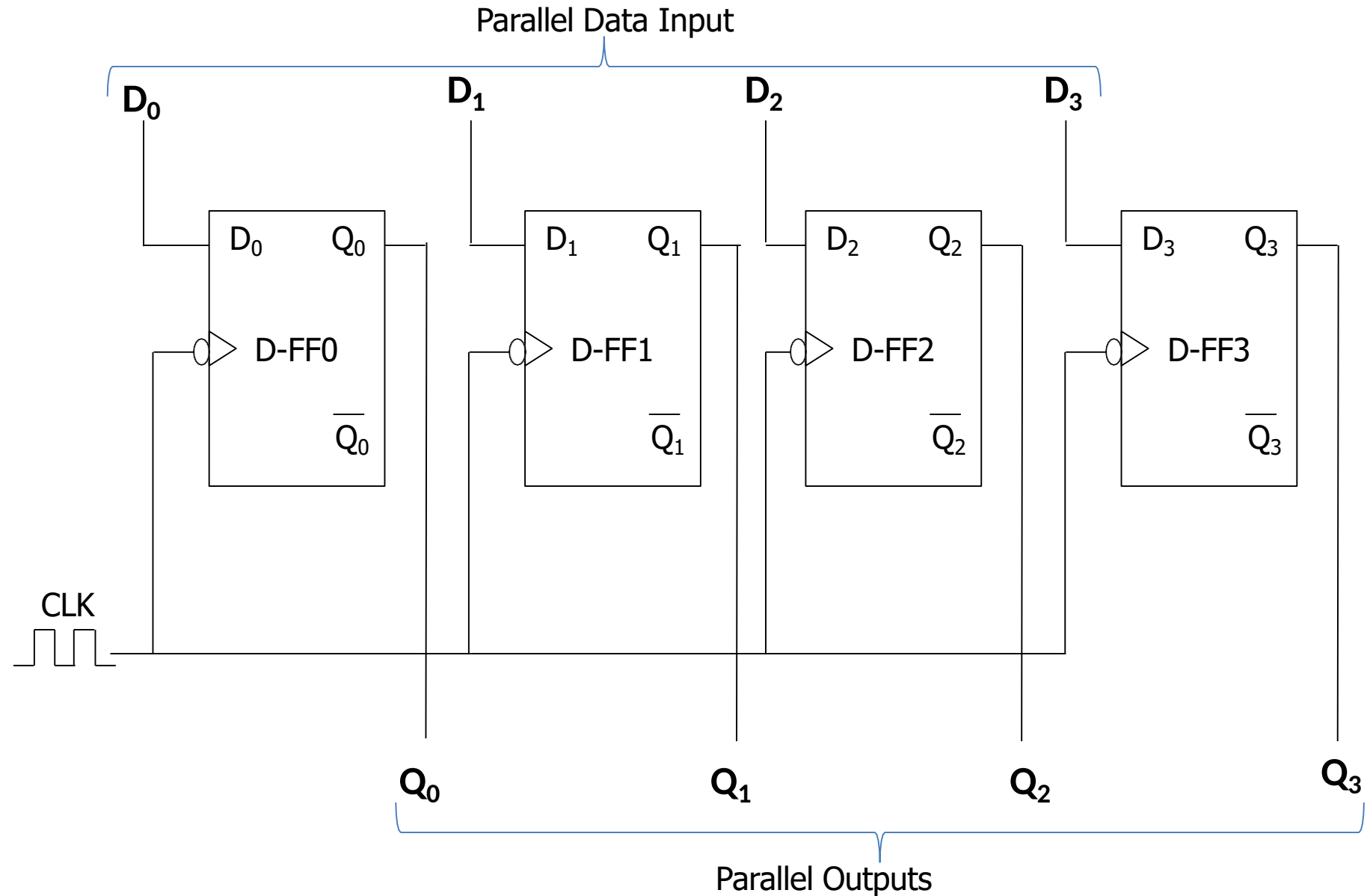


# SIPO – Serial In Parallel Out Shift Register

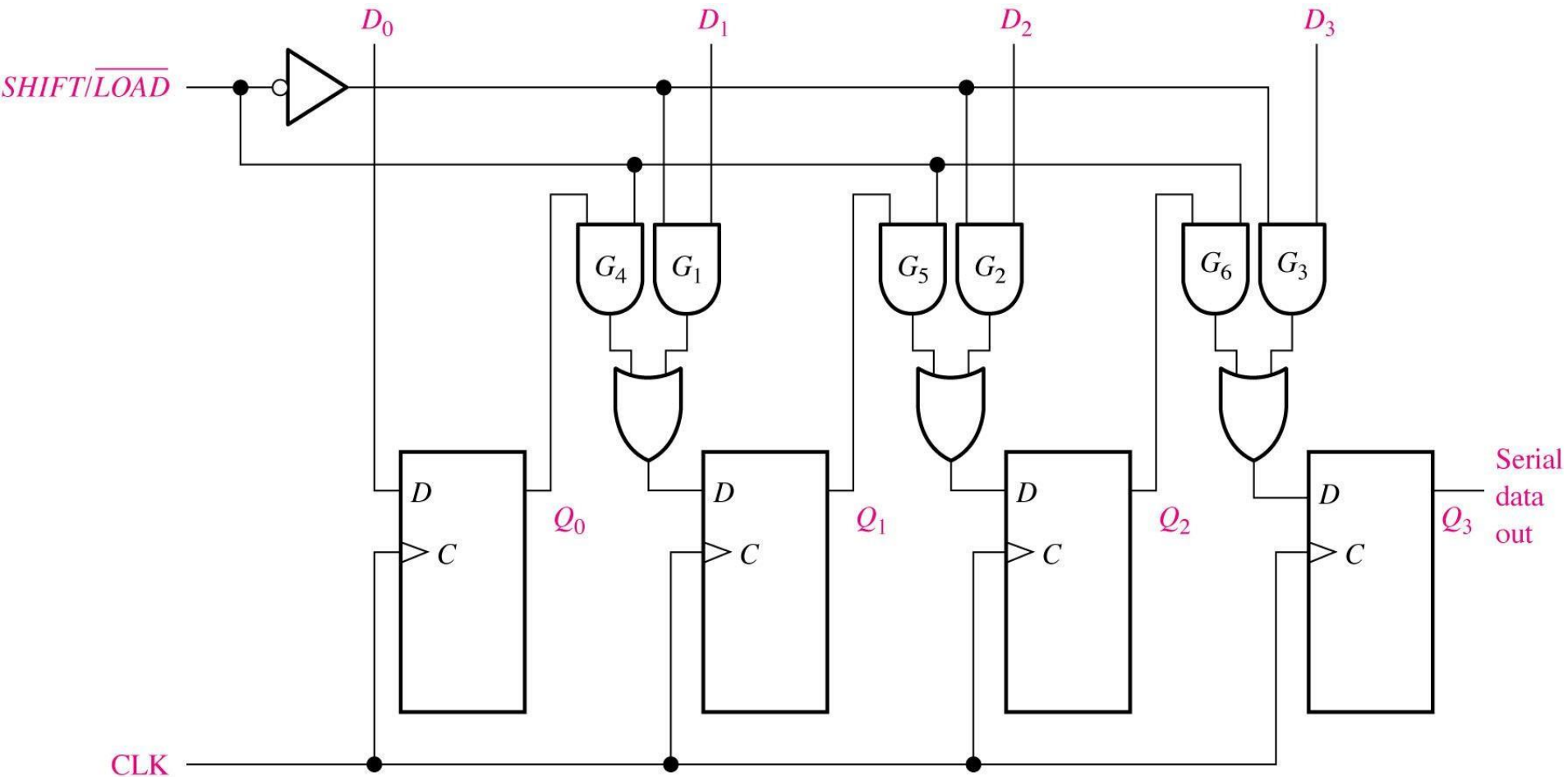
Serial Data  
Input



# PIPO – Parallel In Parallel Out Shift Register



# PISO – Parallel In Serial Out Shift Register



# PISO – Parallel In Serial Out Shift Register

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## Load Mode:

- ✓ When the *Shift / Load* line is Low, the AND gates G1, G2 and G3 become active. They will pass D1, D2, and D3 bits to the corresponding Flip Flops.
- ✓ On the low going edge of clock, the binary inputs D0, D1, D2 and D3 will get loaded into corresponding flip flops. Thus parallel loading takes place.

# PISO – Parallel In Serial Out Shift Register

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## Shift Mode:

- ✓ When the *Shift / Load* line is High, the AND gates G1, G2 and G3 become inactive. Hence parallel loading of data becomes impossible.
- ✓ But AND gates G4, G5 and G6 become active. Therefore the shifting of data from left to right bit by bit on application of clock pulses
- ✓ Thus the parallel in serial out operation takes place

## Shift register:

- A register capable of shifting its binary information either to the right or to the left.
- The logical configuration of shift register consists of a chain of flipflop connected in cascade.
- A flipflop receives a common clock pulse which causes the shift from one stage to another.

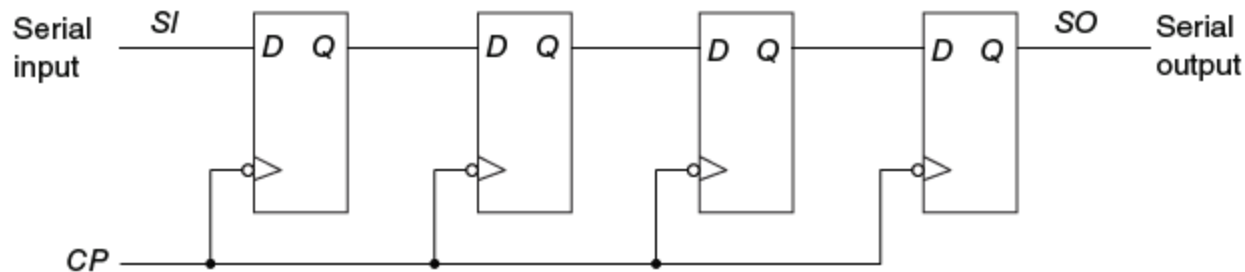
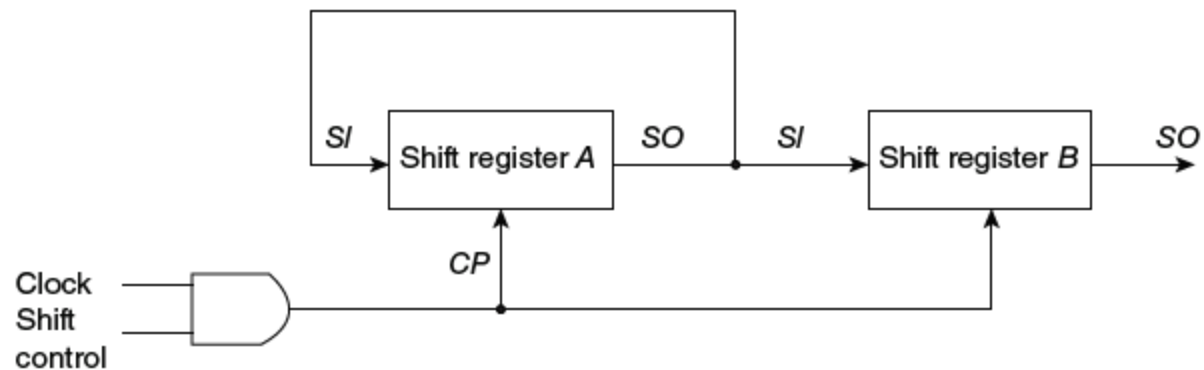


Figure 7-7 Shift register

## Serial transfer:

- A digital system is said to be operate in a serial mode when information is transferred one bit at a time.
- The content of one register is transferred to another by shifting the bit from one register to another.
- The serial transfer of information from register A to register B is is done with shift register as shown in block diagram.
- The serial output of register A ( $SO$ ) goes to serial input of register A ( $Si$ ) of register B
- To prevent the loss of information stored in the source register, the A register is made to circulate the information by connecting the serial output to the input terminal.



(a) Block diagram

- Assume the binary content of A before the shift is 1011 and that of B is 0010.
- The serial transfer from A to B will occur in 4 steps as shown in table.
- After the first pulse  $T_1$  the rightmost bit of A is shifted into the leftmost bit of B and at the same time, this bit is circulated into the leftmost position of A.
- Once the bit of A and B are shifted to right, the previous serial output from B is lost.

Timing pulse	Shift register A				Shift register B				Serial output of B
Initial value	1	0	1	1	0	0	1	0	0
After $T_1$	1	1	0	1	1	0	0	1	1
After $T_2$	1	1	1	0	1	1	0	0	0
After $T_3$	0	1	1	1	0	1	1	0	0
After $T_4$	1	0	1	1	1	0	1	1	1



- Ring counter:
  - A type of counter in which the output of the last flipflop is connected as an input to the first flipflop is known as Ring Counter.
  - The input is shifted between the flipflops in a ring shape so it is called as ring counter.
  - A ring counter is a synchronous counter.

- Johnson counter:
- Johnson counter as a reverse ring counter. In other words, feedback from the last flipflop is fed inversely to the data input of the first flipflop.
- For example for a D- flipflop shift register, the Q' output of the last flipflop is fed to the D input of the first flipflop.
- The Mod of the johnson counter is ' $2n$ ',  $n$  is the bit size of the counter. Mod is the maximum number of states a counter can obtain.

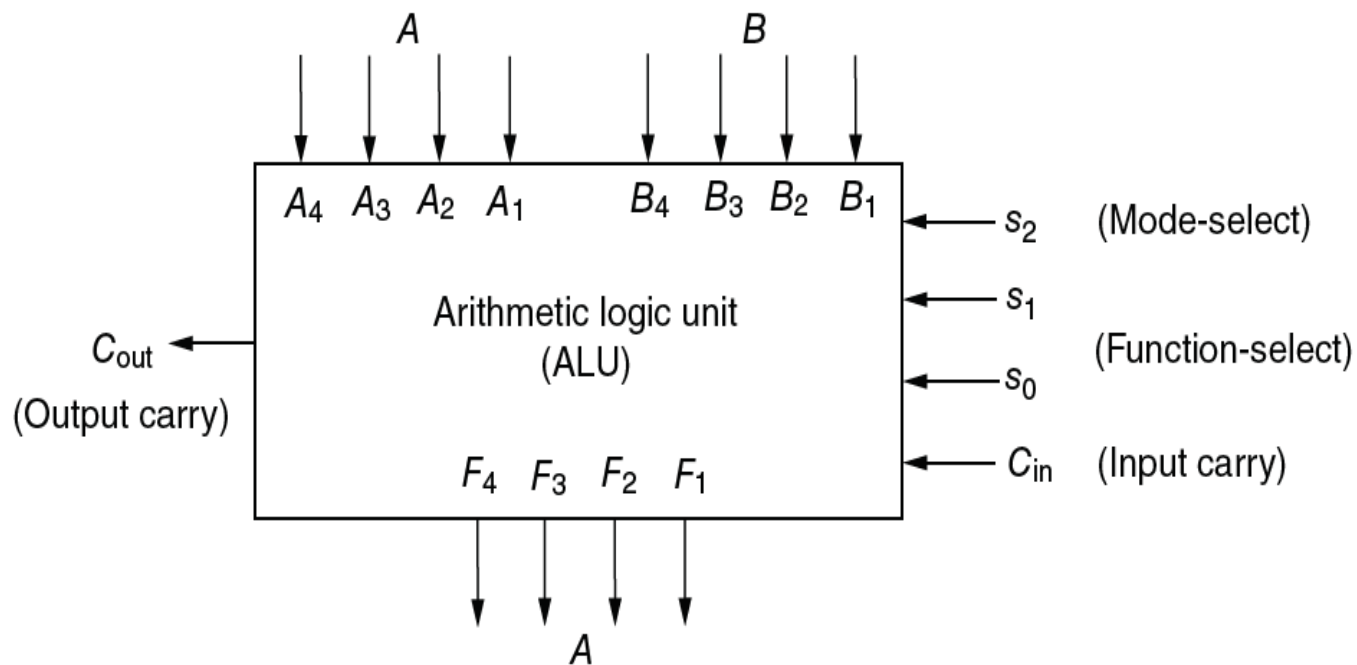
- Advantages:
- More output as compared to ring counter.
- It has the same number of flip flop but it can count twice the number of states the ring counter can count.
- It counts the data in the continuous loop.
- It only needs half the number of flip flops compared to the standard ring counter for the same MOD.

- Disadvantages:
- Only 8 of the 15 states are being used.
- It doesn't count in binary sequence.

# Arithmetic and Logic unit:

Q) Define ALU with block diagram.

- An arithmetic Logic unit (ALU) multi-operation, combinational logic digital function.
- It can perform a set of basic operation and a set of logic operation.
- Figure below shows the block diagram of 4 bit ALU. The four bit data input from 'A' are combined with 4 input from 'B' to generate 8 operation at the F output.



**Figure 9-5** Block diagram of a 4-bit ALU

- Selection line  $s_2$  is used for mode select to distinguish between arithmetic and logic operation.
- Selection line  $s_1$  and  $s_0$  specify the particular arithmetic or logic operation to be generated.
- With three selection variable, it is possible to specify four arithmetic operation and four logic operation.
- The input and output carries have meaning only during an arithmetic operation.
- The fourth selection variable  $c_{in}$  can

# Design of Arithmetic circuit:

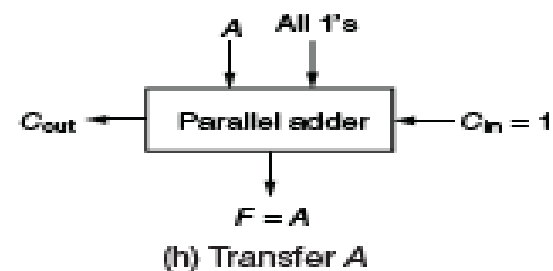
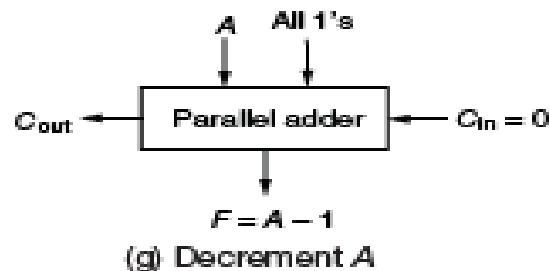
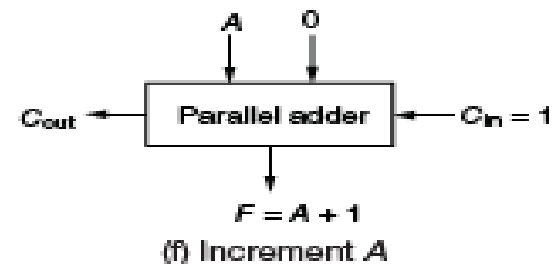
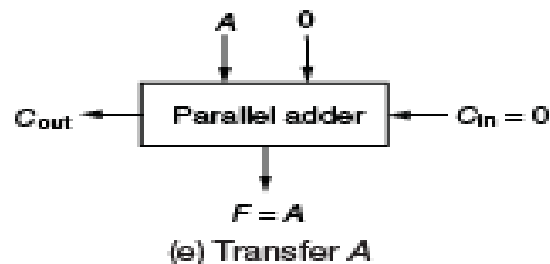
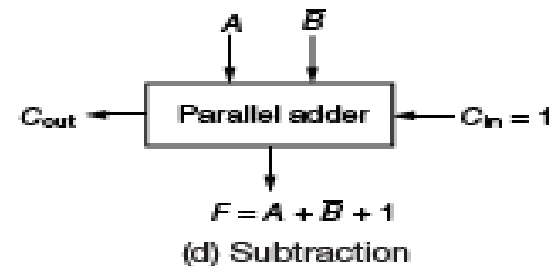
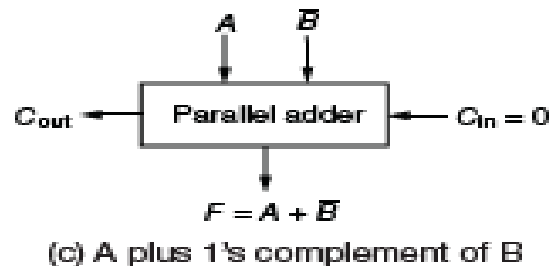
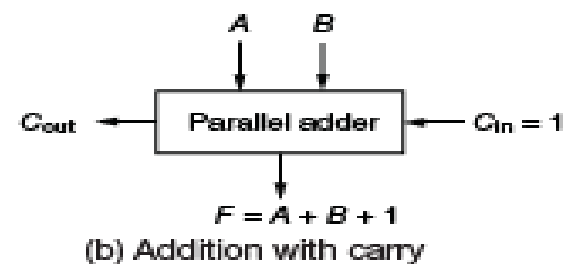
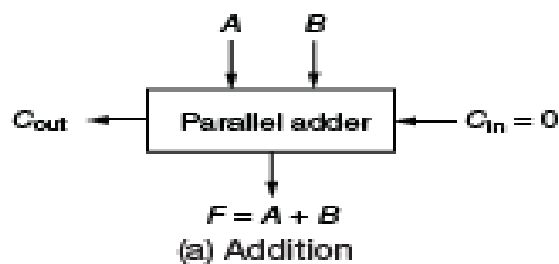
Q) Draw an arithmetic circuit logic diagram with function table that perform 8 different operation.

- The basic component of arithmetic section of ALU is a parallel adder.
- A parallel adder is constructed with number of flipflop. By controlling the data input to the parallel adder, it is possible to obtain 8 different types of arithmetic operation.

**1. Addition:** Arithmetic addition is



2. **Addition with carry:** with  $c_{in} = 1$  in addition.
3. **A plus 1's complement of B:** By complementing all the bit of input B with  $c_{in} = 0$ .
4. **Subtraction:** By making  $c_{in} = 1$ , it is possible to add 1 to the A plus 1's complement of B which produces the sum of A plus 2's complement of B. This operation is similar to subtraction if input carry is discarded.
5. **Transfer A:** if we force all the bit of B terminal to 0 then it transfers input A into output F.
6. **Increment A:** By making  $c_{in} = 1$  and all bit of B terminal to 0, we obtain  $F = A + 1$  which is increment of A.
7. **Decrement A:** If we insert all the bit of B terminal to with  $c_{in} = 0$  then this produce a



**Figure 9-6** Operations obtained by controlling one set of inputs to a parallel adder

# Function table for the arithmetic circuit:

Function select			Y equals	Output equals	Function
$s_1$	$s_0$	$C_{in}$			
0	0	0	0	$F = A$	Transfer $A$
0	0	1	0	$F = A + 1$	Increment $A$
0	1	0	$B$	$F = A + B$	Add $B$ to $A$
0	1	1	$B$	$F = A + B + 1$	Add $B$ to $A$ plus 1
1	0	0	$\bar{B}$	$F = A + \bar{B}$	Add 1's complement of $B$ to $A$
1	0	1	$\bar{B}$	$F = A + \bar{B} + 1$	Add 2's complement of $B$ to $A$
1	1	0	All 1's	$F = A - 1$	Decrement $A$
1	1	1	All 1's	$F = A$	Transfer $A$

When  $s_1s_0 = 00$

$y_i = 0$

When  $s_1s_0 = 01$

$y_i = B_i$

When  $s_1s_0 = 10$

$y_i = B_i'$

When  $s_1s_0 = 11$

$y_i = \text{all } 1\text{'s}$

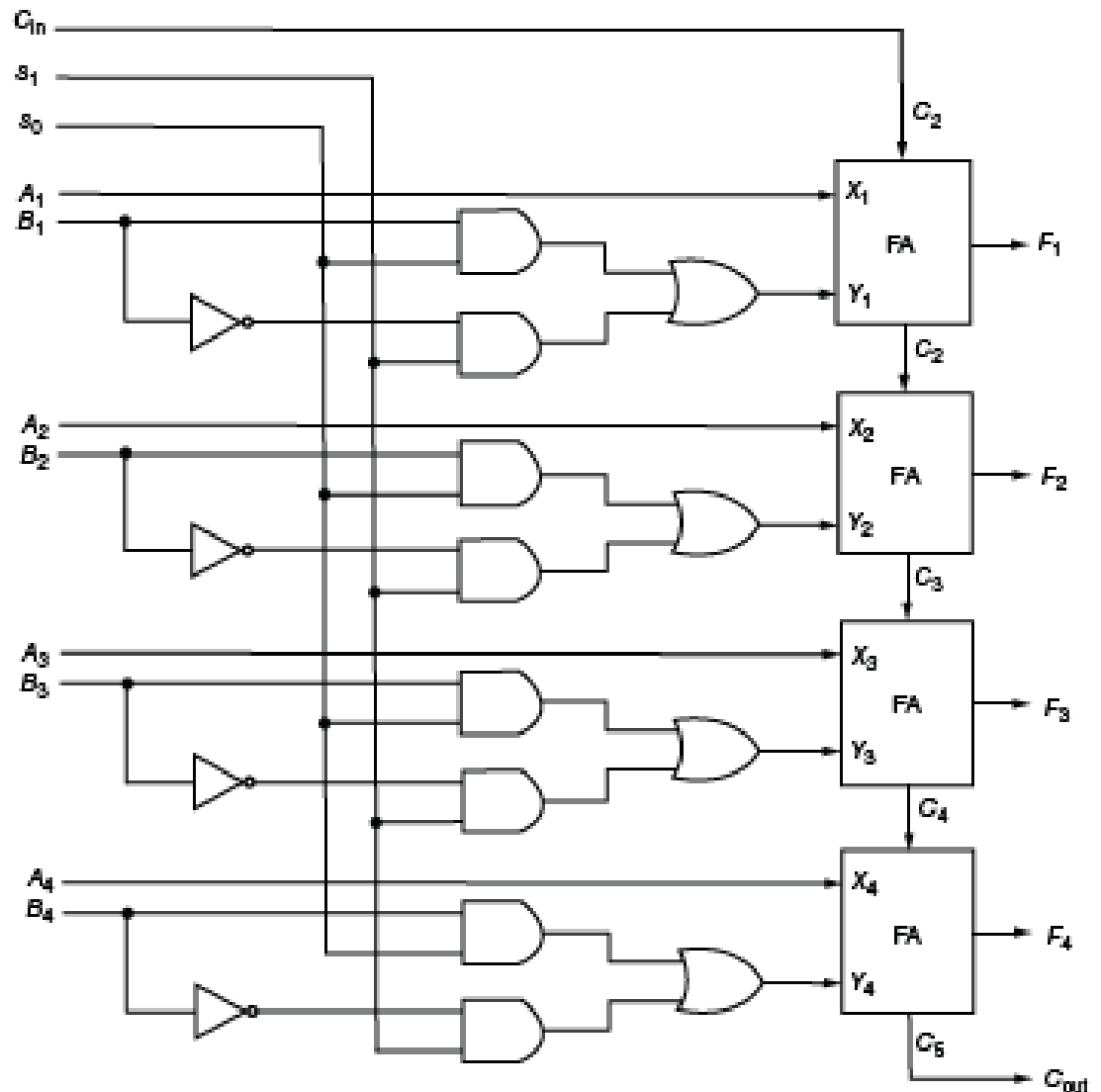
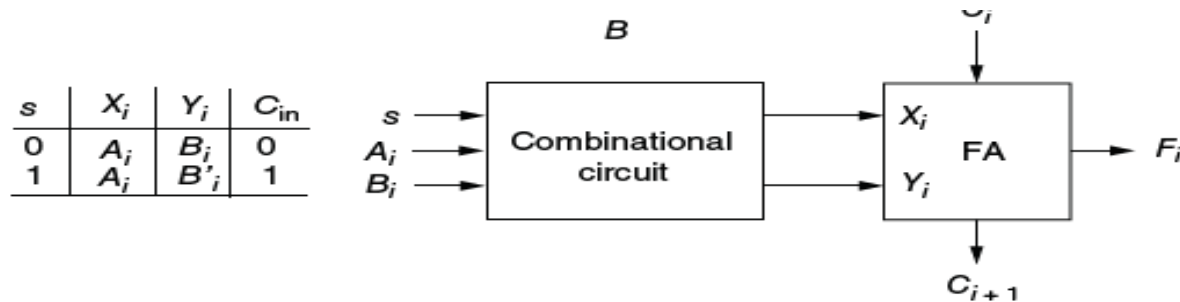


Figure 9-8 Logic diagram of arithmetic circuit

Q) Design an adder/ subtractor circuit with one selection variable and two input A and B, when  $s=0$  the circuit performs  $A+B$ . when  $S=1$  the circuit performs  $A-B$  by taking the 2's complement of B

The arithmetic circuit and truth table for given condition will be as:



(b) Specifying combinational circuit

$s$	$A_i$	$B_i$	$X_i$	$Y_i$	
0	0	0	0	0	
0	0	1	0	1	
0	1	0	1	0	
0	1	1	1	1	
1	0	0	0	1	
1	0	1	0	0	
1	1	0	1	1	
1	1	1	1	0	

$$X_i = A_i$$

$$Y_i = B_i \oplus s$$

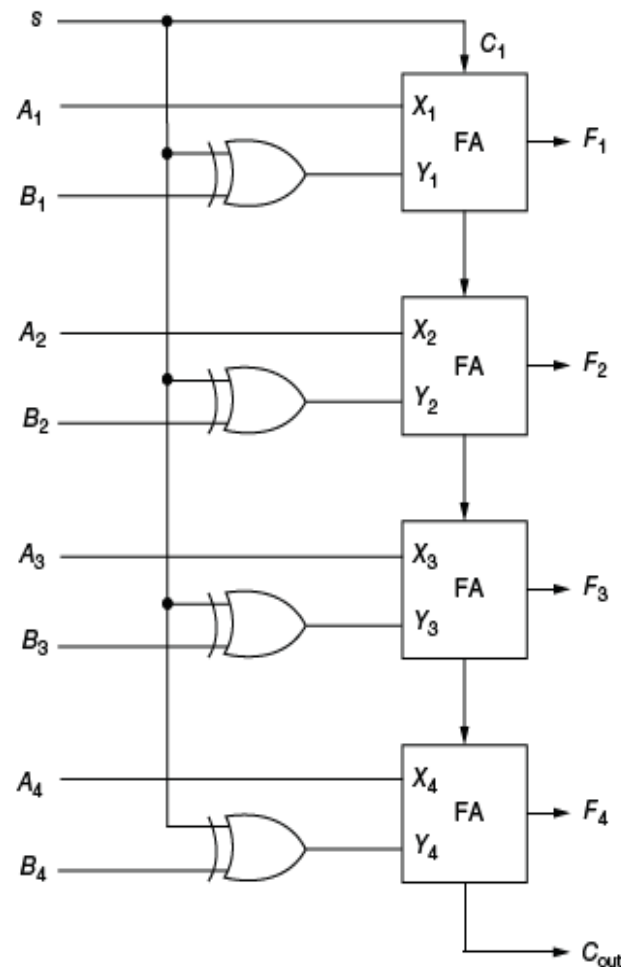
$$C_{in} = s$$

(c) Truth table and simplified equations

**Figure 9.9** Derivation of an adder/subtractor circuit

When  $s=0$   $x_i$  and  $y_i$  for each full adder must be equal to the external input  $A_i$  and  $B_i$  Respectively. When  $s = 1$ ,  $x_i = A_i$  and  $y_i = B_i'$ . The input carry must be equal to value of  $s$ .

The logic diagram will be as:



**Figure 9-10** 4-bit adder/subtractor circuit

## Shift register:

- A register capable of shifting its binary information either to the right or to the left.
- The logical configuration of shift register consists of a chain of flipflop connected in cascade.
- A flipflop receives a common clock pulse which causes the shift from one stage to another.

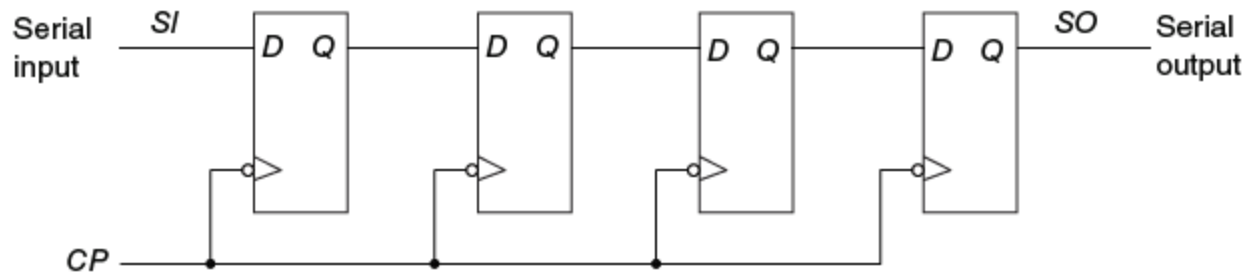
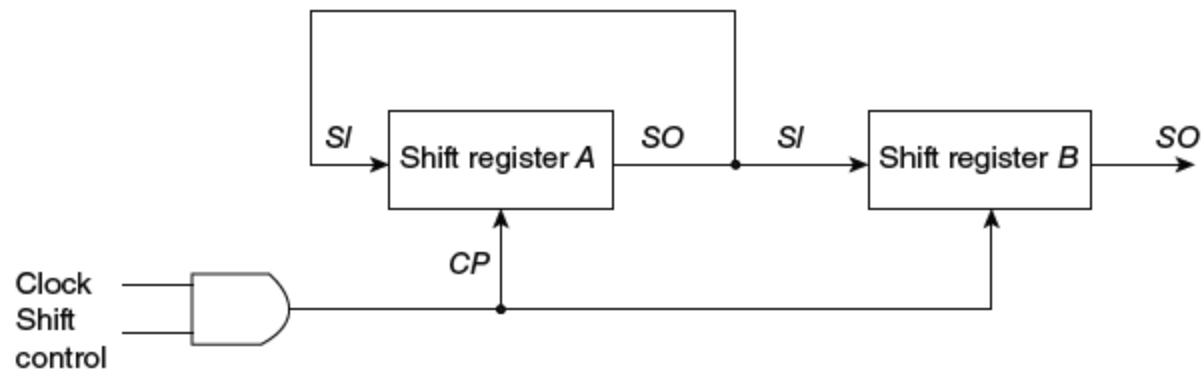


Figure 7-7 Shift register

## Serial transfer:

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(a) Block diagram

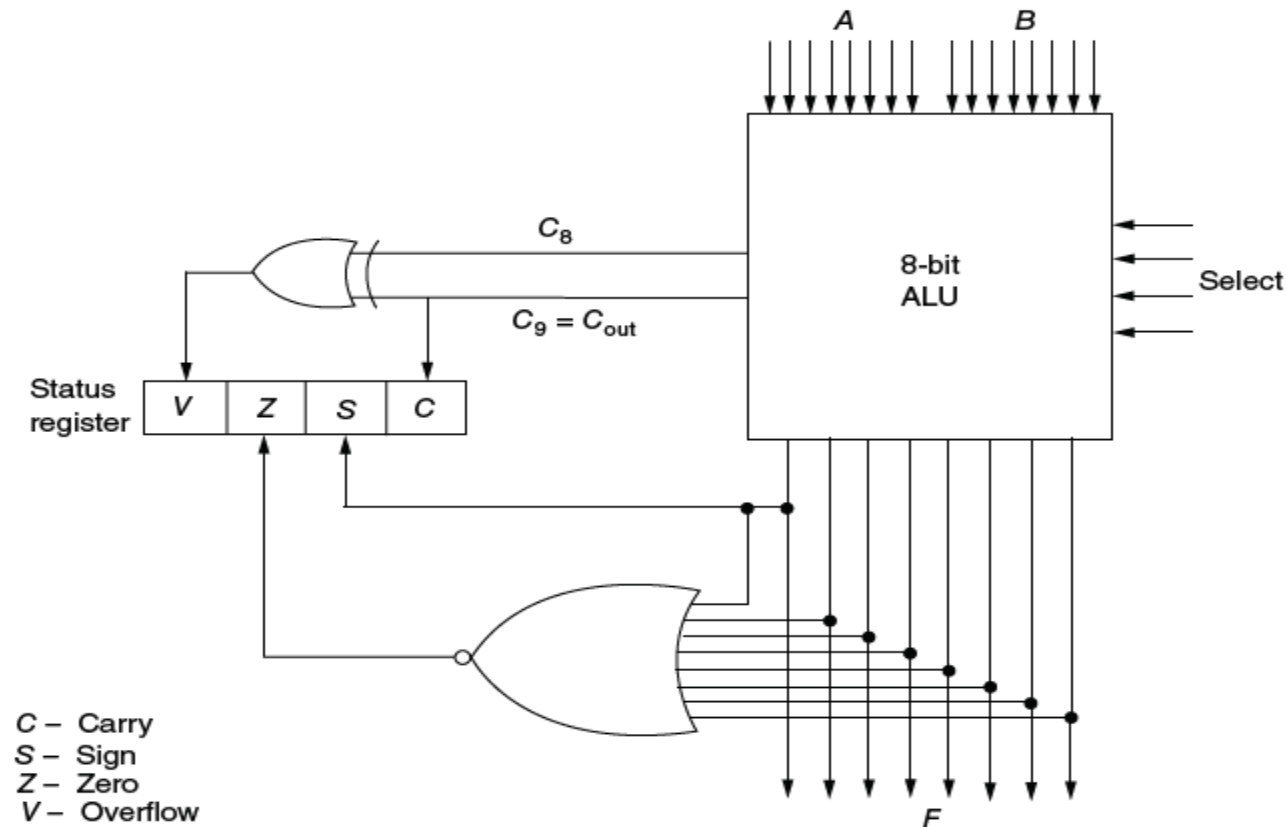


- Assume the binary content of A before the shift is 1011 and that of B is 0010.
- The serial transfer from A to B will occur in 4 steps as shown in table.
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Timing pulse	Shift register A				Shift register B				Serial output of B
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After $T_1$	1	1	0	1	1	0	0	1	1
After $T_2$	1	1	1	0	1	1	0	0	0
After $T_3$	0	1	1	1	0	1	1	0	0
After $T_4$	1	0	1	1	1	0	1	1	1

## Status register:

Figure below shows the block diagram of an 8 bit Alu with a 4 bit status register.



- The four status bits or flag bits are symbolized by  $C$  ,  $S$  ,  $Z$  and  $V$ .
- The bit are set or cleared as a result of an operation performed in the ALU.

1. Bit  $C$  is set if the output carry of the ALU is 1. It is cleared if the output carry is 0.
2. Bit  $S$  is set if the highest-order bit of the result in the output of the ALU (the sign bit) is 1. It is cleared if the highest-order bit is 0.
3. Bit  $Z$  is set if the output of the ALU contains all 0's, and cleared otherwise.  $Z = 1$  if the result is zero, and  $Z = 0$  if the result is nonzero.
4. Bit  $V$  is set if the exclusive-OR of carries  $C_8$  and  $C_9$  is 1, and cleared otherwise. This is the condition for overflow when the numbers are in sign-2's-complement representation (see Section 8-6). For the 8-bit ALU,  $V$  is set if the result is greater than 127 or less than  $-128$ .