

Lab02-Algorithm Analysis

Exercises for Algorithms by Nengjun Zhu, 2022-2023 Fall Semester.

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1. Consider the sorting algorithm shown in Alg.1, which is called BUBBLESORT.

- (a) What is the minimum number of element comparisons? When is this minimum achieved?
- (b) What is the maximum number of element comparisons? When is this maximum achieved?
- (c) Express the running time of Alg.1 in terms of the O and Ω notations.
- (d) Can the running time of the algorithm be expressed in terms of the Θ notation? Explain.

Algorithm 1: BUBBLESORT

```
input : An array  $A[1 \cdots n]$  of  $n$  elements.
output:  $A[1 \cdots n]$  in nondecreasing order.

1  $i \leftarrow 1$ ;  $sorted \leftarrow false$ ;
2 while  $i \leq n - 1$  and not sorted do
3    $sorted \leftarrow true$ ;
4   for  $j \leftarrow n$  downto  $i + 1$  do
5     if  $A[j] < A[j - 1]$  then
6       interchange  $A[j]$  and  $A[j - 1]$ ;
7        $sorted \leftarrow false$ ;
8   end
9 end
10  $i \leftarrow i + 1$ 
11 end
```

解:

- (a) 当数组为升序数组时, 比较次数最少, 最少比较次数为 $n - 1$
- (b) 当数组为降序数组时, 比较次数最多, 最多比较次数为: $(n - 1) + (n - 2) + \dots + 1 = \frac{n(n-1)}{2}$
- (c) 上界为 $O(n^2)$, 下界为 $\Omega(n)$
- (d) 数组逆序度大小就是排序的交换次数, 最大逆序度为 $\frac{n(n-1)}{2}$, 最小逆序度为 0, 因此平均情况的逆序度为 $\frac{n(n-1)}{4}$, 交换次数是 n^2 级的, 而比较次数一定大于等于交换次数, 算法的上界又是 $O(n^2)$, 所以算法的平均复杂度为 $\Theta(n^2)$

2. For Alg.2 and Alg.3 shown below, answer the following questions respectively.

- (a) Give the maximum number of times Line 6 is executed in Alg.2 when n is a power of 3.
- (b) Give the maximum number of times Line 5 is executed in Alg.3 when n is a power of 2.
- (c) What is the time complexity of both algorithms expressed in the O and Θ notations?

Algorithm 2: COUNT1

```

1  count ← 0;
2  for i ← 1 to n do
3      j ← ⌊n/3⌋;
4      while j ≥ 1 do
5          for k ← 1 to i do
6              count ← count + 1;
7              if j is even then
8                  j ← 0;
9              else
10                 j ← ⌊j/3⌋;
11             end
12         end
13     end
14 end

```

Algorithm 3: COUNT2

```

1  count ← 0;
2  for i ← 1 to n do
3      j ← ⌊n/2⌋;
4      while j ≥ 1 do
5          count ← count + 1;
6          if j is odd then
7              j ← 0;
8          else
9              j ← j/2;
10         end
11     end
12 end

```

解:

(a) $(\log_3 n)^2 + \frac{(n + \log_3 n + 1)(n - \log_3 n)}{2}$

(b) $n \log_2 n$

(c) 算法 2: $O(n^2)$ 、

算法 3: $O(n \log_2 n)$ 、 $\Theta(n)$

3. Fill in the blanks with either true or false:

$f(n)$	$g(n)$	$f = O(g)$	$f = \Omega(g)$	$f = \Theta(g)$
$2n^3 + 3n$	$100n^2 + 2n + 100$	false	true	false
$50n + \log n$	$10n + \log \log n$	false	false	true
$50n \log n$	$10n \log \log n$	false	true	false
$\log n$	$\log^2 n$	true	false	false
$n!$	5^n	false	true	false

4. Use the \prec relation to order the following functions by growth rate:

$$n^{1/100}, \sqrt{n}, \log n^{100}, n \log n, 5, \log \log n, \log^2 n, (\sqrt{n})^n, (1/2)^n, 2^{n^2}, n!$$

解:

$$(1/2)^n \prec 5 \prec \log \log n \prec \log n^{100} \prec \log^2 n \prec n^{1/100} \prec \sqrt{n} \prec n \log n \prec 2^{n^2} \prec (\sqrt{n})^n \prec n!$$