

# Any Colony Optimization

Adam DeJans

Oakland University

*addejans@oakland.edu*

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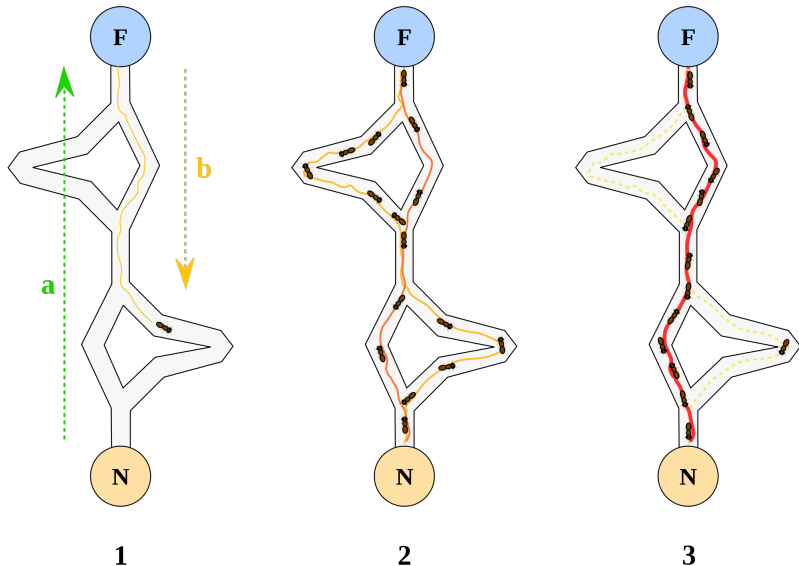
# Overview

- 1 Introduction
  - Pictorial Example
- 2 Ant System
- 3 ACO Metaheuristic
- 4 Applications of ACO Algorithms
  - Advantages and Disadvantages
- 5 Data Network Routing
  - Subsection Example
- 6 Conclusion

# Introduction

Ant algorithms are inspired from the observation of real ant colonies. In particular, the how ants are able to find shortest paths between food sources and their nest.

# Pictorial Example



## Similarities

- Colony of cooperating individuals
- Pheromone trail and stigmergy
- Shortest path searching and local moves
- Stochastic and myopic state transition policy

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## Unique to Artificial Ants:

- live in a discrete world
- have an internal state
- deposit an amount of pheromone which is a function of the quality of the solution found
- in many cases update pheromone trails only after having a generated solution
- have extra capabilities (e.g., looking ahead, local optimization, backtracking, etc.)

*“A new general-purpose heuristic algorithm to solve different optimization problems”*

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- This is a classical NP-hard combinatorial optimization problem.
- TSP is a very important problem in regards to Ant Colony Optimization (ACO) because the original ACO algorithm, Ant System (AS), was first applied to TSP and TSP is often used as a benchmark to test a new algorithmic ideas.

# AS Pseudocode for TSP

```
def AS_TSP():  
    initialize()  
    random_ant_placement()  
    while (stopping_criteria == False):  
        for (ant in Ants):  
            while (more_cities_toVisit == True):  
                transition_ant()  
            return_to_initial_city()  
        pheromone_update()
```

# Tour Construction

- $m$  (artificial) ants are created to concurrently build a tour of the TSP
- Ants are initially put on randomly chosen cities.

How does an artificial ant choose which vertex to visit next?



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This is determined probabilistically

# Path Selection - Transition Probability

The transition probability from city  $i$  to city  $j$  for the  $k$ -th ant is defined as

$$p_{i,j}^k(t) = \begin{cases} \frac{[\tau_{i,j}(t)]^\alpha \cdot [\eta_{i,j}(t)]^\beta}{\sum_{l \in \mathcal{N}_i^k} [\tau_{i,l}(t)]^\alpha \cdot [\eta_{i,l}(t)]^\beta} & \text{if } j \in \mathcal{N}_i^k \\ 0 & \text{otherwise} \end{cases}$$

2

Variables:

- $\tau_{i,j}(t)$  = amount of pheromone trail on edge  $(i,j)$  at time  $t$ .
- $\eta_{i,j}$  = the visibility; expressed as  $\frac{1}{d_{i,j}}$ , where  $d_{i,j}$  is the euclidean distance between two vertices.
- $\alpha$  = parameter controlling the relative weight of pheromone trail
- $\beta$  = parameter controlling the weight of visibility

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<sup>2</sup>Notice: the fraction ensures that  $\sum_{j \in \mathcal{V}} p_{i,j}^k(t) = 1$



# Path Selection - Transition Probability (cont.)

## Data Structures:

- $tabu_k$  – dynamically growing vector containing the tabu list of the  $k$ -th ant (list of vertices visited by the  $k$ -th ant)
- $\mathcal{N}_i^k = V \setminus tabu_k$  – list of vertices not visited by the  $k$ -th ant

# Updating Pheromone

After each ant has completed their tour, pheromone evaporation on all arcs is triggered, and then each ant  $k$  deposits a quantity of pheromone  $\Delta\tau_{ij}^k(t)$  on each arc that it has used.

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<sup>3</sup>was found to have negligible influence

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$$\Delta\tau_{i,j}^k(t) = \begin{cases} \frac{Q}{L^k(t)} & \text{if arc } (i,j) \in \text{tabu}_k(t) \\ 0 & \text{otherwise} \end{cases}$$

$Q$  is a constant<sup>3</sup>

$L^k(t)$  is the length of the tour done by ant  $k$  at iteration  $t$

---

<sup>3</sup>was found to have negligible influence

# Updating Pheromone (cont.)

One can implement the addition of new pheromone and pheromone evaporation to all arcs by

$$\tau_{i,j}(t) \leftarrow (1 - \rho)\tau_{i,j}(t) + \Delta\tau_{i,j}(t)$$

where  $\Delta\tau_{i,j}(t) = \sum_{k=1}^m \Delta\tau_{i,j}^k(t)$  and  $\rho$  is a constant in  $(0, 1]$  representing *evaporation*.

# End of First Run

- 1 Save best ant tour (sequence and length)
- 2 All ants die
- 3 New ants are born

# Stopping Criteria

The algorithm can be set to stop after

- Stagnation; or
- A maximum number of iterations is reached

- Not so great results
- Modified to now have excellent results (Ant Colony System)
- *daemon\_actions()* - the optional function used to implement centralized actions that can't be performed by individual ants.
  - E.g., using global perspective to decide whether or not to deposit additional pheromone to persuade the search process from a non-local aspect.
  - Local search methods to constructed solutions

Three major modifications to enhance results:

- 1 Search experience accumulated from ants is used more strongly
- 2 Pheromone evaporation and deposit take place only on arcs belonging to the best-so-far tour
- 3 Each time an ant uses an arc it removes some pheromone from the arc to increase the exploration of alternative paths.



# ACS Pseudocode for TSP

```
def ACS_TSP():  
    while (stopping_criteria == False):  
        for (city in Cities):  
            for (ant in Ants)  
                transition_ant()  
                local_trail_update()  
    global_trail_update
```

# Tour Construction (ACS)

When ant  $k$  is at vertex  $i$ , city  $j$  is chosen according to the so called *pseudorandom proportional* rule

$$j = \begin{cases} \operatorname{argmax}_{l \in \mathcal{N}_i^k} \{\tau_{i,l} [\eta]^\beta\} & \text{if } q \leq q_o \text{ Exploitation} \\ J & \text{otherwise Exploration} \end{cases}$$

where  $q$  is a random variable uniformly distributed in  $[0, 1]$ ,  $q_o$  ( $0 \leq q_o \leq 1$ ) is a parameter, and  $J$  is a random variable selected according to the probability distribution

$$p_{i,j}^k(t) = \begin{cases} \frac{[\tau_{i,j}(t)]^\alpha \cdot [\eta_{i,j}(t)]^\beta}{\sum_{l \in \mathcal{N}_i^k} [\tau_{i,l}(t)]^\alpha \cdot [\eta_{i,l}(t)]^\beta} & \text{if } j \in \mathcal{N}_i^k \\ 0 & \text{otherwise} \end{cases}$$

Similar to evaporation

$$\tau_{i,j}(t) \leftarrow (1 - \rho)\tau_{i,j}(t) + \Delta\tau_{i,j}(t)$$

but we let  $\Delta\tau_{i,j} = (nL_{nn})^{-1}$ ;

where  $L_{nn}$  is the length of a nearest-neighbour tour.

# Consequence of Local Trail Update

- Desirability of edges changes dynamically
- Edge pheromone level decreases
- As edges are visited by different ants the desirability to visit lessens
- Exploration of unvisited edges becomes more favourable and in turn ants visit vertices at different times in the tour (for example, ant  $m_1$  will not visit the early visited vertices of ant  $m_2$ 's tour.)

After each iteration the best ant is chosen to update it's tour by setting additional pheromone inversely proportional to the length of it's tour

$$\tau_{i,j}(t) \leftarrow (1 - \rho)\tau_{i,j}(t) + \rho \frac{1}{L^{ba}} \quad \forall (i,j) \in T^{ba}$$

where  $L^{ba}$  is the length of the best ant's tour and  $T^{ba}$  is the tour of the best ant.

- AS deposits pheromone only after completing a tour
- In ACS only the ant with the best tour is allowed to deposit pheromone (globally)
- AS deposits pheromone to all edges
- In ACS the ant with the best tour deposits pheromone only on it's associated tour

# Generalized ACO

- Stochastic construction procedure
- Probabilistically build a solution
- Iteratively add solution components to partial solutions via
  - Heuristic information
  - Pheromone trail
- Modify the problem representation at each iteration
- Ants work together and independently
- Leads to good solutions (maybe not optimal but they're good and fast)

```

def ACO_Meta-heuristic():
    while (termination_criterion_not_satisfied):
        schedule_activities:
            ants_generation_and_activity()
            pheromone_evaporation()
            daemon_actions() #optional

def ants_generation_and_activity():
    while (available_resources):
        schedule_the_creation_of_a_new_ant()
        new_active_ant()

```



```

def new_active_ant():
    initialize_ant()
M = update_ant_memory() #ant lifecycle
while (current_state != target_state)
    A = read_local_ant-routing_table()
    P = compute_transition_probabilities \
        (A, M, problem_constraints)
    next_state = apply_ant_decision_policy \
        (P, problem_constraints)
    move_to_next_state(next_state)
    if (online_step-by-step_pheromone_update)
        deposit_pheromone_on_the_visited_arc()
        update_ant-routing_table()
    M = update_internal_state()
if (online_delayed_pheromone_update)
    evaluate_solution()
    deposit_pheromone_on_all_visited_arcs()
    update_ant-routing_table()
die()

```

## Problems Solvable with ACO

- Vehicle Routing Problem  
(and variations)
- Assignment Problems
- Scheduling
- Subset problems

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*\*We see that ACO is especially helpful with dynamic problems*

## Advantages

- Good solutions in practical time
- Works well with dynamic and “ill-structured” problems (e.g., network routing problem)

# Advantages and Disadvantages

## Advantages

- Good solutions in practical time
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## Disadvantage

- Optimal solution may not be found
- Difficult theory
- Based off of experimentation

- The main advantage over other heuristic algorithms such as *simulated annealing* and *genetic algorithm* is that the ant colony algorithm can be run continuously and adapt to changes in real time
- This is especially helpful when the graph may change dynamically such as in *network routing*

# Data Network Routing - AntNet

The main ideas of ACO are applied but in an adaptive way.

## High Level Description of AntNet Algorithm:<sup>4</sup>

- In certain time intervals *forward-ants* are launched towards random destination vertices
- Ants act independently but communicate indirectly by reading and writing locally to vertices visited
- Ants try to find the minimum cost path from its source to destination
- Ants move in accordance to a stochastic policy based on
  - ① local ant generated and maintained information
  - ② local problem dependent heuristic information
  - ③ ant private information
- When arrived at destination, *backward-ants* are launched back to the source along the same path but in opposite direction
- While backtracking, network status and routing table of each visited vertex is modified as a function of the path taken and of its goodness
- Once returned, the ant is destroyed and new ants are born.

<sup>4</sup> AntNet: Distributed Stigmergetic Control for Communications Networks



# Things To Consider When Implementing ACO Algorithms

- How many ants to use?
- How to balance ant exploration vs exploiting the pheromone trails?
- Heuristic choices?
- When to update pheromones?
- Which ants should update pheromones?
- Stopping criteria?

# Final Conclusions

- (summary)

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# The End