

# Embedded System and Microcomputer Principle

LAB3 ARMv7 Assembly Programming

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01

Cortex-M3 introduction



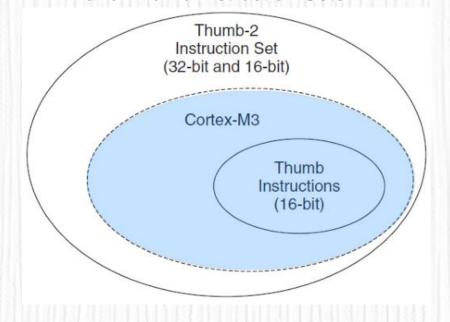
#### Instruction Set Development

- Historically (since ARM7TDMI), two different instruction sets are supported on the ARM processor: the ARM instructions that are 32-bit and Thumb instructions that are 16-bit.
- During program execution, the processor can be dynamically switched between the ARM state or the Thumb state to use either one of the instruction sets.
- The Thumb instruction set provides only a subset of the ARM instructions, but it can provide higher code density.
- The Thumb instruction set is useful for products with tight memory requirements.



- The Thumb-2 Instruction Set Architecture (ISA)
  - In 2003, ARM announced the Thumb-2 instruction set, which is a new superset of Thumb instructions that contains both 16-bit and 32-bit instructions.
  - The Cortex-M3 supports only the Thumb-2 (and traditional Thumb) instruction set.

The Relationship between Thumb-2 and Thumb Instruction Sets



With support for both 16-bit and 32-bit instructions in the Thumb-2 instructions set, there is no need to switch the processor between Thumb state (16-bit instructions) and ARM state (32-bit instructions).



#### Registers (R0 to R15)

- General-Purpose Registers R0–R7 (low registers)
  - They can be accessed by all 16-bit Thumb and all 32-bit Thumb-2 instructions.
  - They are all 32-bit; the reset value is unpredictable.
- General-Purpose Registers R8–R12 (high registers)
  - They are accessible by all Thumb-2 but not by all 16-bit Thumb instructions.
  - These registers are all 32-bit; the reset value is unpredictable.
- R13: Stack Pointers
  - The Cortex-M3 contains two stack pointers, R13. They are banked so that only
    one is visible at a time.
- R14: The Link Register
  - When a subroutine is called, the return address is stored in the link register.
- R15: The Program Counter
  - The program counter is the current program address. This register can be written to control the program flow.



#### Special Registers

- They have special functions and can be accessed only by special instructions.
- They cannot be used for normal data processing.
- The Cortex-M3 processor also has a number of special registers:
  - Program Status Registers (PSRs: APSR, IPSR, EPSR)
  - Interrupt Mask Registers (PRIMASK, FAULTMASK, BASEPRI)

Register	Function
xPSR	Provide ALU flags (Z, C), execution status, and current executing interrupt number
PRIMASK	Disable all interrupts except the nonmaskable interrupt (NMI) and HardFault
FAULTMASK	Disable all interrupts except the NMI
BASEPRI	Disable all interrupts of specifi c priority level or lower priority level
CONTROL	Defi ne privileged status and stack pointer selection



02

CPUlator system simulator

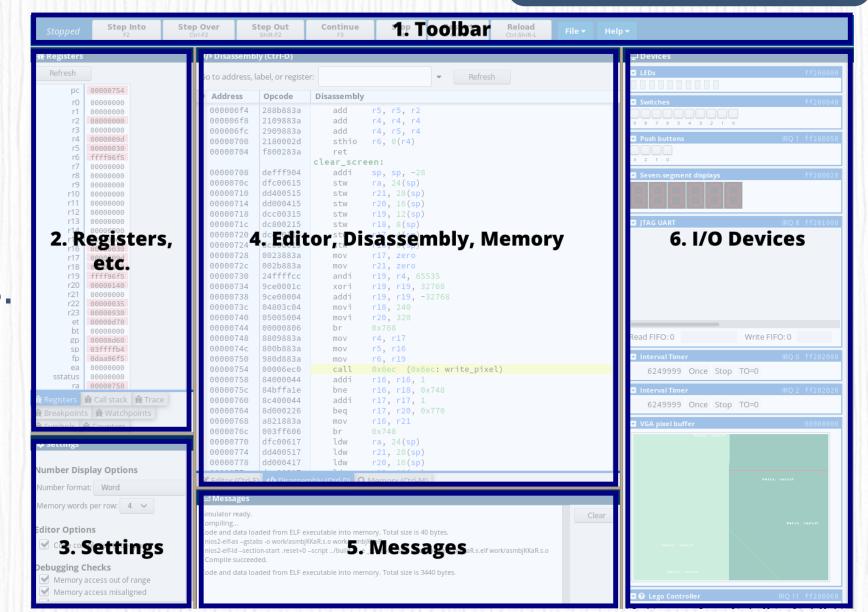




- CPUlator is a simulator and debugger of a computer system (CPU, memory, and I/O devices) that runs inside a web browser.
- https://cpulator.01xz.net/?sys=arm
- Features
  - No download
  - No account sign up
  - Only Internet

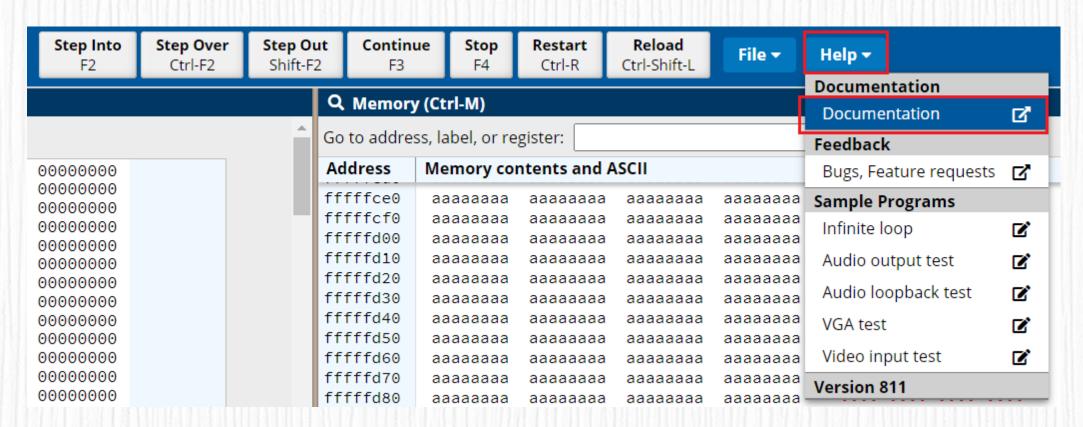


 The simulator interface consists of the toolbar and a collection of movable panels.



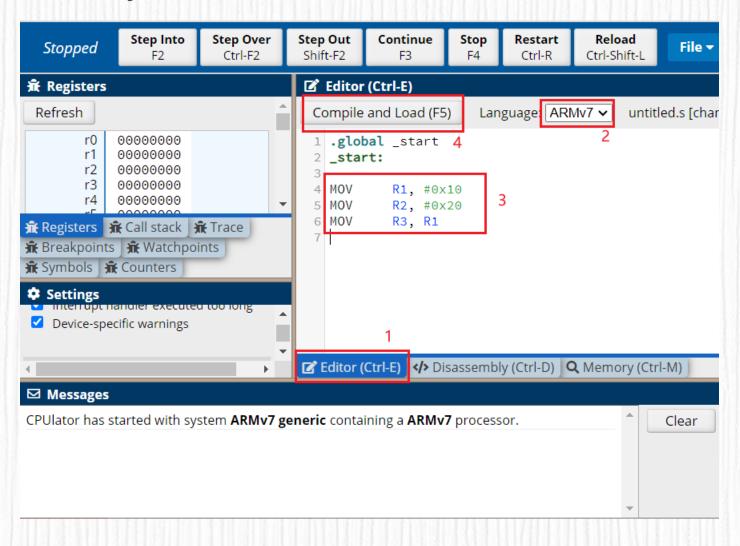


The help documentation can provide more information.

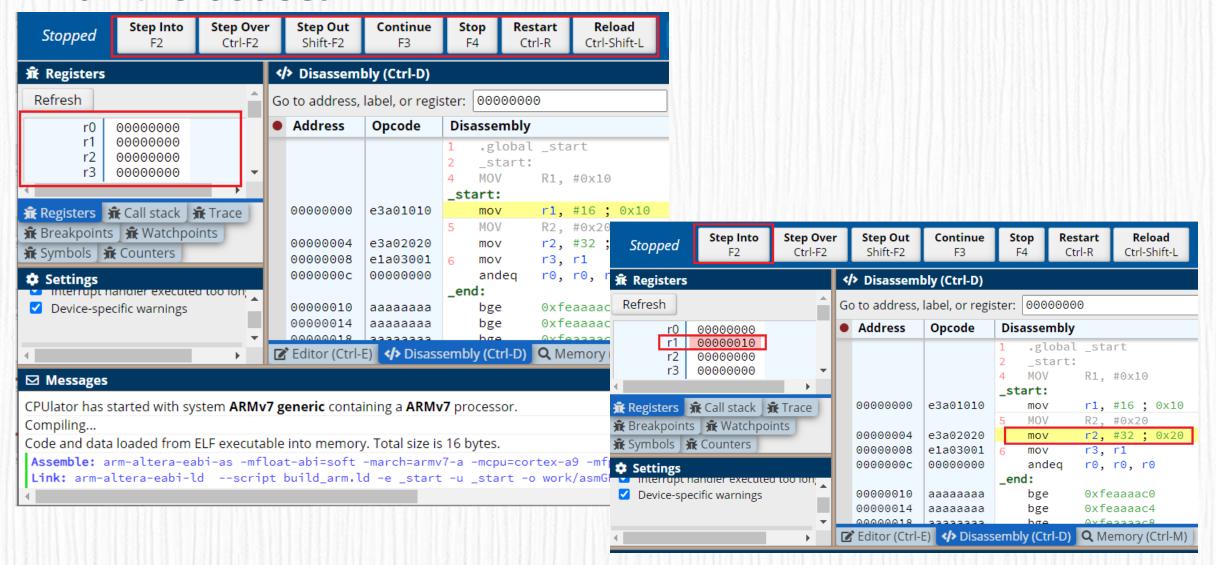


Load assembly codes.





Run the codes.







03

Basic assembly instructions



Bit Fields in Cortex-M3
 Program Status Registers

N: Negative

- Z: Zero

C: Carry / Borrow

V: Overflow

Defreeb			
Refresh			
r0	00000000		
r1	00000010		
r2	00000000		
r3	00000000		
r4	00000000		
r5	00000000		
r6	00000000		
r7	00000000		
r8	00000000		
r9	00000000		
r10	00000000		
r11	00000000		
r12	00000000		
sp	00000000		
lr	00000000		
рс	00000000		
cpsr	000001d3	NZCVI	SVC
spsr	00000000	NZCVI	?

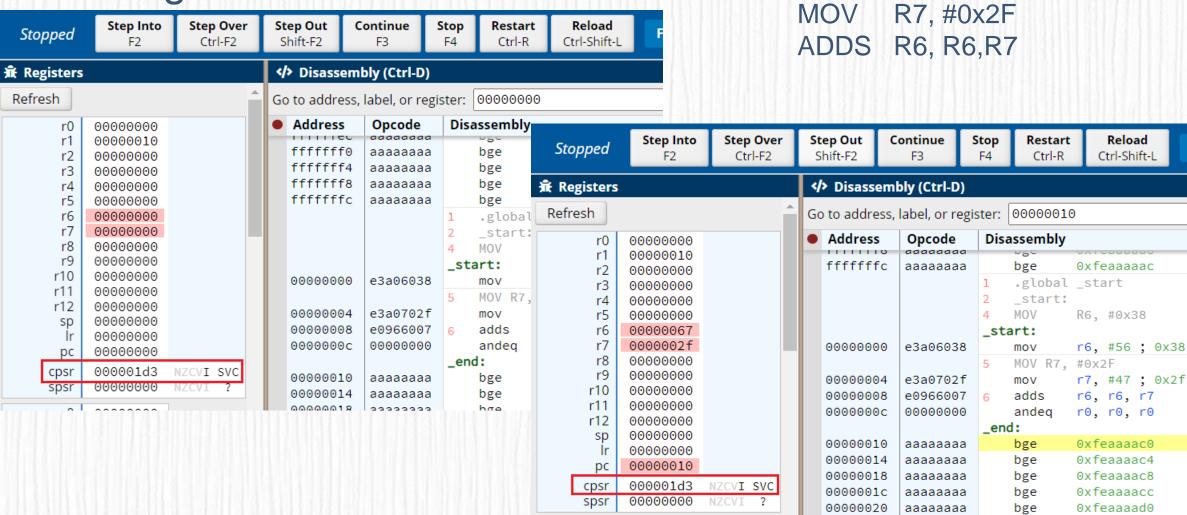
	31	30	29	28	27	26:25	24	23:20	19:16	15:10	9	8	7	6	5	4:0
APSR	N	Z	С	V	Q											
IPSR	Exception Number															
EPSR						ICI/IT	Т			ICI/IT						



R6, #0x38

MOV

Flag related instructions demo1

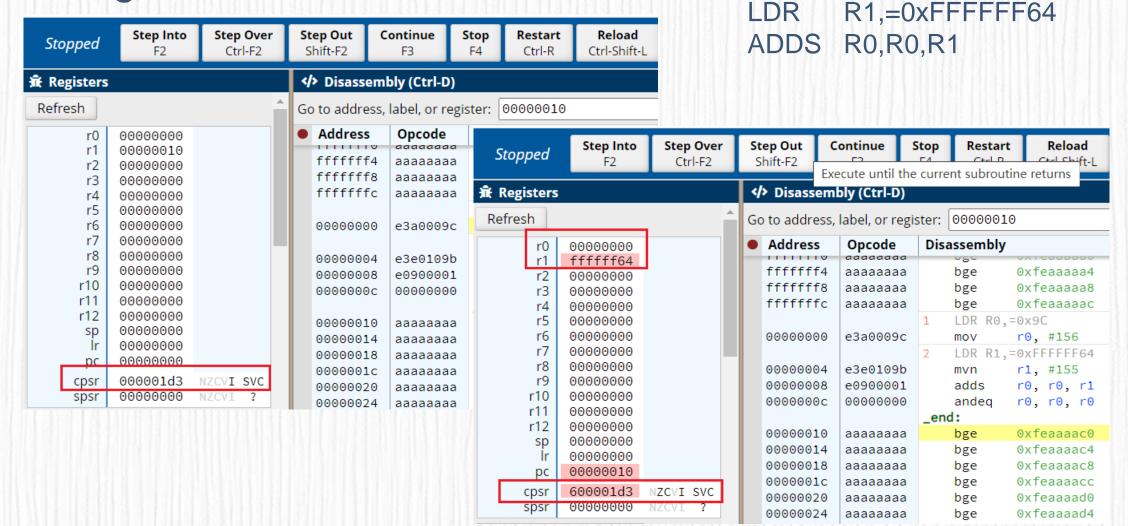




R0,=0x9C

LDR

Flag related instructions demo2





Flag related instructions demo2 LDR R0,=0x9C\_DR R1,=0xFFFFF64 Step Into Step Over Step Out Continue Stop P Juad Restart ADDS R0,R0,R1 Stopped Ctrl-R Ctrl-Shift-L Ctrl-F2 Shift-F2 Registers ♦ Disassembly (Ctrl-D) Go to address, label or register: 00000010 Refresh Address Opcode r0 00000000 Step Over Step Out Continue Stop Step Into Restart Reload 00000010 spped Ctrl-F2 Shift-F2 aaaaaaaa F2 00000000 Execute until the current subroutine returns 0000009c aaaaaaaa 0000000 **₩** Registers Disassembly (Ctrl-D) 00000010 aaaaaa .a 00000000 00000000 Refresh Go to address, label, or register: 00000010 0000009c r0 00000000 00000 ffffff64 00000000 00000000 Address Opcode Disassembly 00000000 00000004 e3e0109b ffffff64 aaaaaaaa r1 00000000 fffffff4 aaaaaaaa 0xfeaaaaa4 00000008 e0900001 bge 00000000 r10 0000000 fffffff8 aaaaaaaa bge 0xfeaaaaa8 0000000c 00000000 00000000 r11 00000000 fffffffc aaaaaaaa bge 0xfeaaaaaac r4 00000000 r12 00000000 LDR Ro,=0x9C 00000000 00000010 aaaaaaaa 00000000 sp 00000000 00000000 e3a0009c mov r0, #156 00000014 aaaaaaaa 0000000 r7 0000000 LDR R1,=0xFFFFF64 00000018 aaaaaaaa 0000000 00000000 00000004 e3e0109b r1, #155 0000001c mvn aaaaaaaa 000001d3 NZCVI SVC 00000000 cpsr 80000000 e0900001 adds r0, r0, r1 00000020 aaaaaaaa 00000000 r10 00000000 000000c 0000000 andeq r0, r0, r0 00000024 aaaaaaaa r11 00000000 end: r12 00000000 00000010 0xfeaaaac0 aaaaaaaa bge 00000000 0xfeaaaac4 00000014 aaaaaaaa bge 0000000 00000018 0xfeaaaac8 aaaaaaaa bge 00000010 0000001c 0xfeaaaaacc bge aaaaaaaa 600001d3

cpsr

00000000

NZCVI SVC

00000020

00000024

aaaaaaaa

aaaaaaaa

bge

bge

0xfeaaaad0

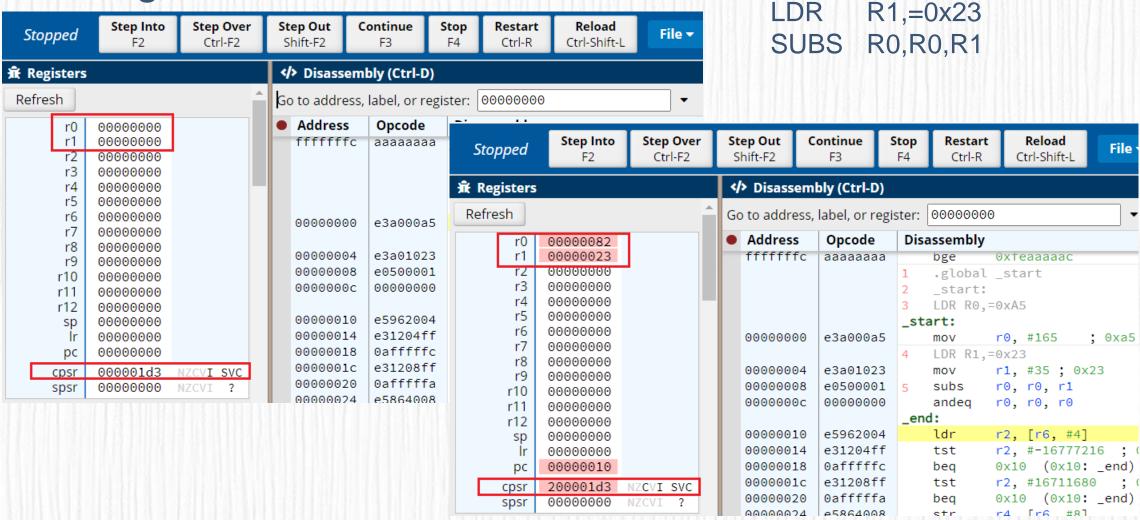
0xfeaaaad4



R0,=0xA5

LDR

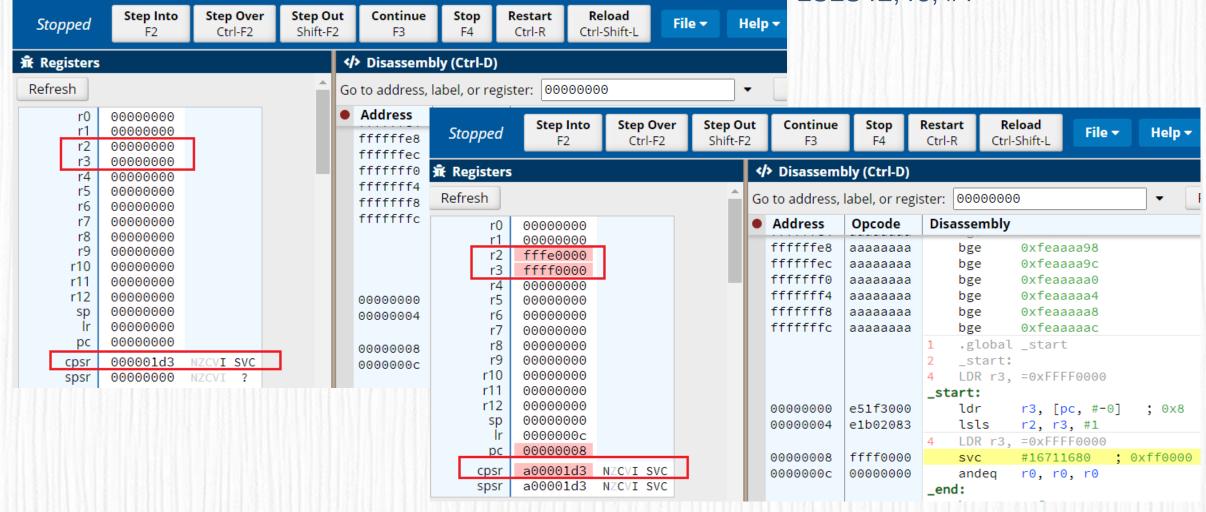
Flag related instructions demo3





Flag related instructions demo4

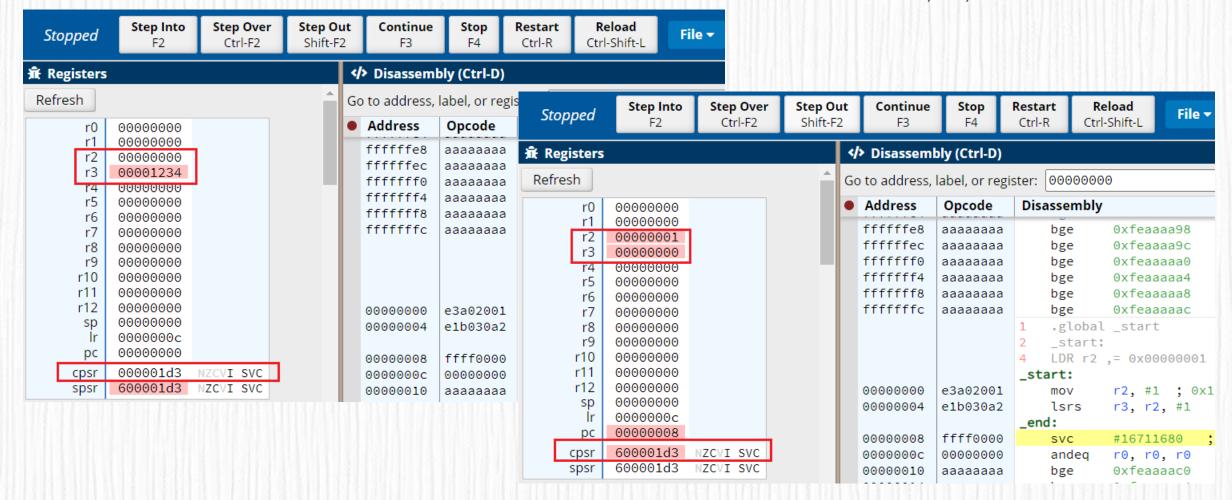
LDR r3, =0xFFFF0000 LSLS r2, r3, #1



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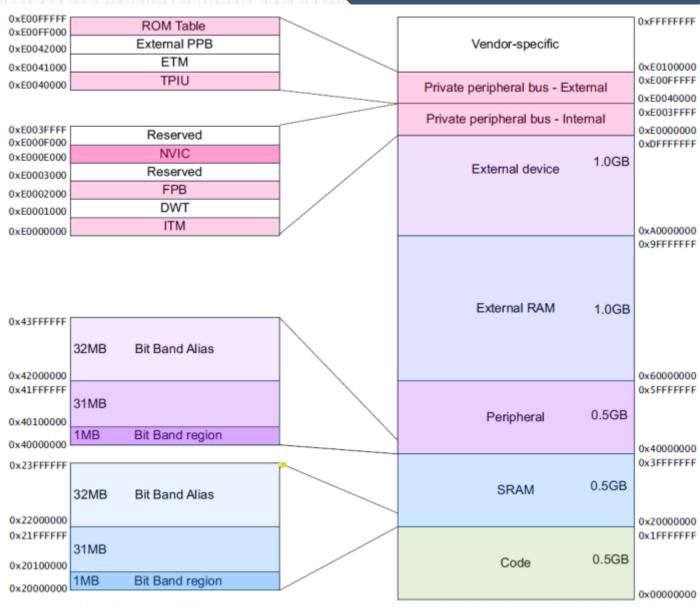
Flag related instructions demo5

LDR r2 ,= 0x00000001 LSRS r3, r2, #1





- The Cortex-M3 processor has a total of 4 GB of address space.
- Program code can be located in the Code region, the SRAM region, or the External RAM region.
- Another 0.5 GB block of address range is allocated to on-chip peripherals.
- Two slots of 1 GB memory space are allocated for external RAM and external devices.
- The last 0.5 GB of memory is for the system-level omponents, internal peripheral buses, external peripheral bus, and vendorspecific system peripherals.







# Common memory access instructions

Example	Description					
LDRB Rd, [Rn, #offset]	Read byte from memory location Rn+offset					
LDRH Rd, [Rn, #offset]	Read half-word from memory location Rn+offset					
LDR Rd, [Rn, #offset]	Read word from memory location Rn+offset					
LDRD Rd1, Rd2, [Rn, #offset]	Read double word (64-bits) from memory location Rn+offset, set Rd1 as lower 32-bits and Rd2 as higher 32-bits					
STRB Rd, [Rn, #offset]	Store byte to memory location Rn+offset					
STRH Rd, [Rn, #offset]	Store half-word to memory location Rn+offset					
STR Rd, [Rn, #offset]	Store word to memory location Rn+offset					
LDRD Rd1, Rd2, [Rn, #offset]	Store double word to memory location Rn+offset					



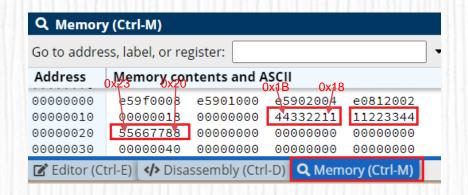
Memory access instruction demo

```
.global _start
_start:

LDR R0,=list
LDR R1,[R0]
LDR R2,[R0, #4]
ADD R2, R1, R2
.data
list:

.byte 0x11, 0x22, 0x33, 0x44
.word 0x11223344, 0x55667788
```

Read this piece of codes, tell the result of R2, describe how to load data from memory?



After loading, we can observe data in memory window.

- Q1: What is the starting address of list?
- Q2: Does byte 0x11, 0x22, 0x33, 0x44 and word 0x11223344 be stored the same in memory?
- Q3: Does ARMv7 adopt large endianness or small endianness storage mode?





#### Common conditional branch instructions

Instruction	Function
В	Branch
BL	Branch and link
BEQ	Branch if two values are equal
BNE	Branch if two values are not equal
ВСС	Branch if Carry = 0
BCS	Branch if Carry = 1
TBB	Table branch byte; forward branch using a table of single byte offset
TBH	Table branch half word; forward branch using a table of half word offset



#### Conditional branch instruction demo

.global \_start \_start: \_MOV R1, #0xa

MOV R2, #0xb CMP R1, #0x1

BNE else //go to else if \_?\_

then:

MOV R2, #3 //R2 = 3

B endif //go to endif

else:

MOV R2, #4 //b = 4

endif:

LDR R1, =dst

STR R2, [R1]

dst:

.asciz "00000000"

Read this piece of codes, tell the results of following questions.

- Q1: What condition(s) should be filled in blank \_?\_ according to the codes?
- Q2: What are the address of lables then,
   else, endif and dst individually?
- Q3: Which address will be accessed after BNE instruction?
- Q4: What value will be stored in memory?
- Q5: If we want to store another value into memory, how should we change the codes?



#### Conditional branch instruction demo -- loop

```
.global _start
start:
        MOV r0, #0 // i
        MOV r1, #0 // sum
loop:
        CMP r0, \#10 // check whether i < 10
        BGE endloop // skip if _?_
        ADD r1, r1, r0 // sum += i
        ADD r0, r0, #1 // i++
        B loop
endloop:
```

Read this piece of codes, tell the results of following questions.

- Q1: What condition(s) should be filled in blank \_?\_ according to the codes?
- Q2: How many times did the loop code block execute?
- Q3: What's the function of the codes?
- Q4: What are the final values of R0, R1?



04

Practice

#### 4. Practice

- 1. Write your own codes and test in emulator, which sets C=1, and V=1;
- 2. Read the codes on right
  - Answer the questions1~4
  - What's the usage of the code?
  - NOTES: this piece of codes can not run in CPUlator.

```
RCC_APB2ENR EQU 0x40021018 ; set RCC_APB2ENR as constant 0x40021018
     GPIOA CRL EQU 0x40010800
     GPIOA_CRH EQU 0x40010804
     GPIOA IDR EQU 0x40010808
     GPIOA_ODR EQU 0x4001080C
     EXPORT main
     AREA MAIN, CODE, READONLY
     __main
         LDR R1,=RCC_APB2ENR
         LDR R0, [R1]
         ORR R0, R0, #0xFC
         STR R0,[R1] ;___1. Purpose of the previous lines?_
12
         LDR R1,=GPIOA_CRH
13
         LDR R0,=0x44444443
14
         STR R0,[R1] ;___2. Purpose of the previous lines?_
15
         LOOP:
16
         LDR R1,=GPIOA ODR
17
         LDR R0, [R1]
         EOR R0, R0, 0x00000100
20
         STR R0, [R1] ;___3. Purpose of the previous lines?
         BL delay ; jump to line 24 to execute function delay
22
         B LOOP
     delay:
         MOV R4, #0
24
                      _{-4}. value of R5 = ?
         MVN R5, #1;
     d loop:
         ADD R4, R4, #1
         CMP R5, R4
         BGE d_loop ; if >=, jump to d_loop
         BX LR ; jump to line 22
     END
```