

CS301

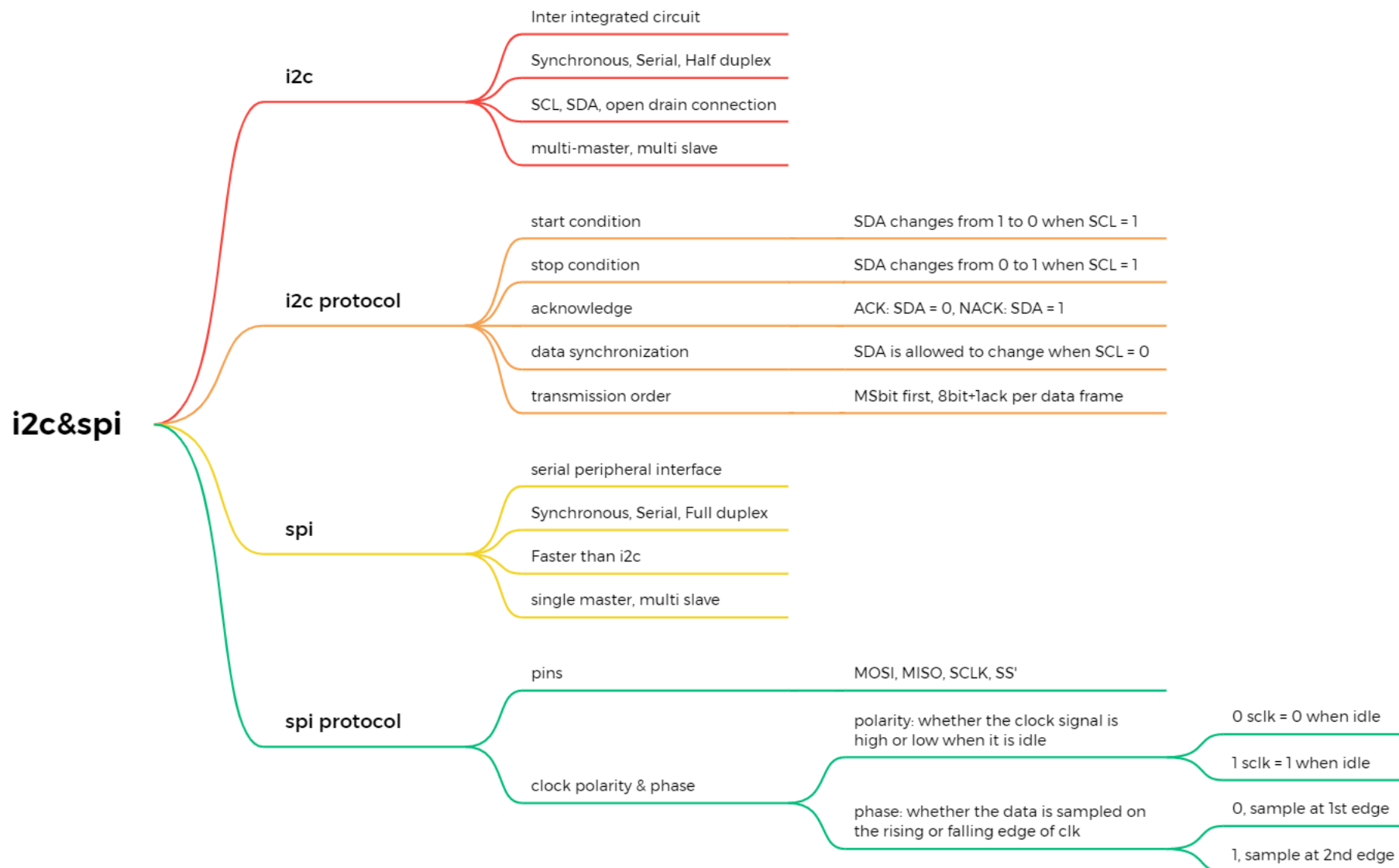
Embedded System and Microcomputer Principle

Lecture 10: Bus

2023 Fall



Recap

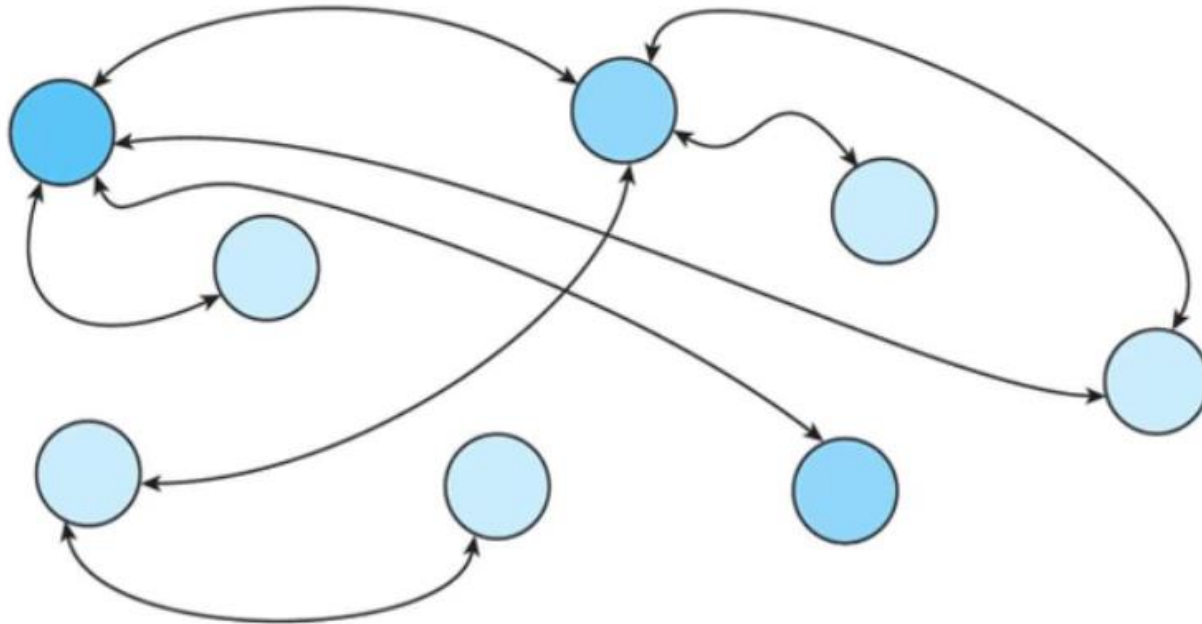


Outline

- **Bus Introduction**
- Bus Performance and Architecture
- Bus Timing
- SoC Bus - AMBA

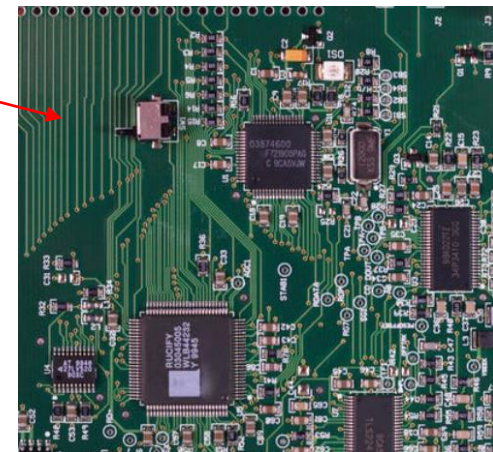
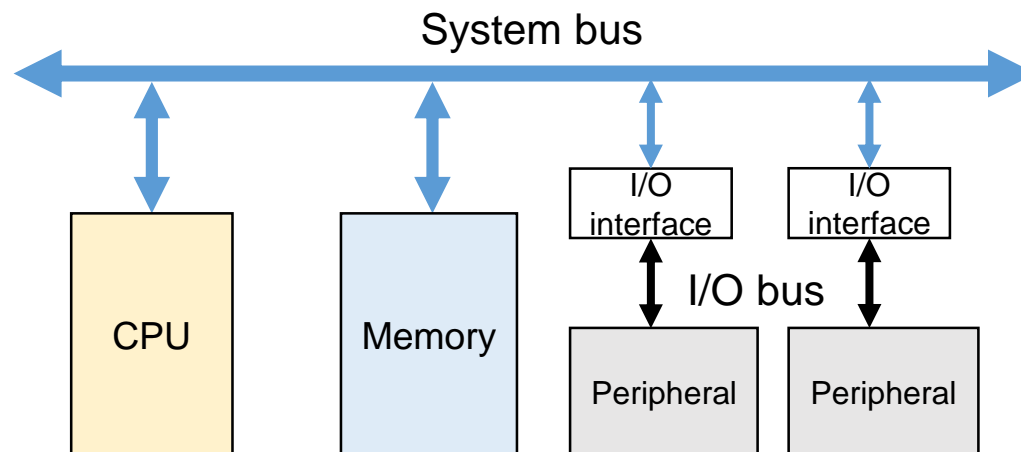
Bus(总线)

- A system without Bus



Bus

- Bus: a shared communication link between subsystems
 - Each line or wire of a bus can at any one time contain a single binary digit.
 - Time Division:
 - At any given moment, only one component is allowed to send information to the bus. Multiple components send information to the bus in different time slots.
 - Shared:
 - Multiple components can be connected to the bus, and each component can exchange information through this set of lines.



What defines a bus?

Transaction Protocol

Timing and Signaling Specification

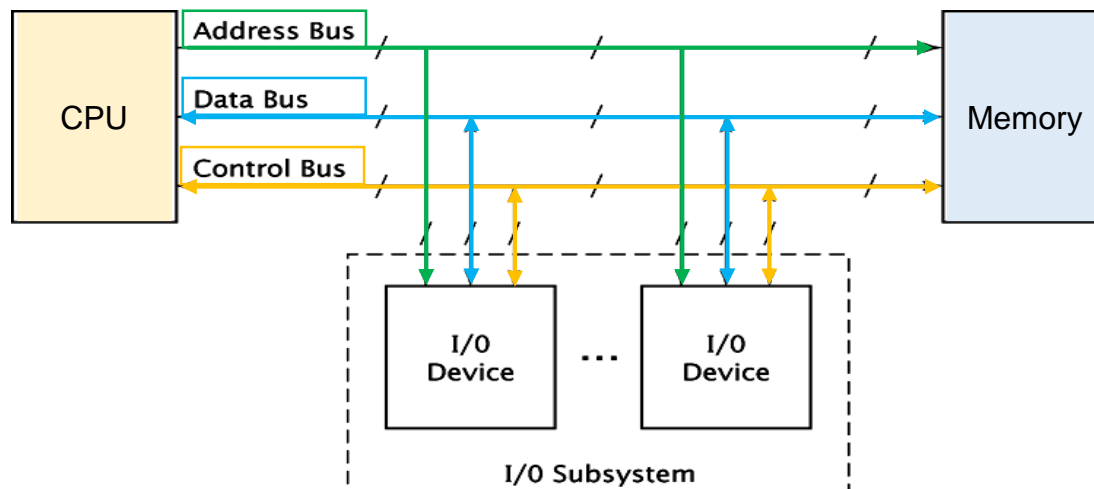
Bunch of Wires

Electrical Specification

Physical/Mechanical characteristics -
the connections

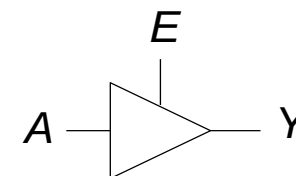
Bus Classification

- Classification based on functionality
 - Address bus (uni-directional)
 - CPU reads/writes data from memory by addressing a unique location
 - Each I/O device has a unique address as well
 - Data bus (bi-directional)
 - Carry information between source and destination
 - Control bus (bi-directional)
 - Signal requests and acknowledgments
 - Indicate what type of information is on the data lines

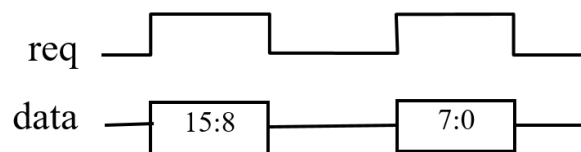
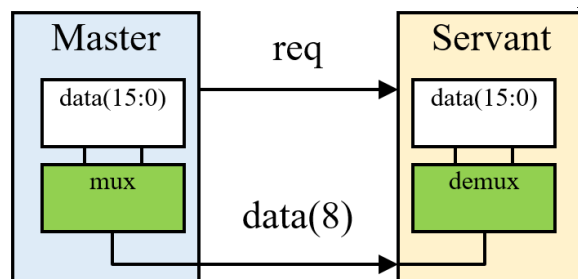


Bus

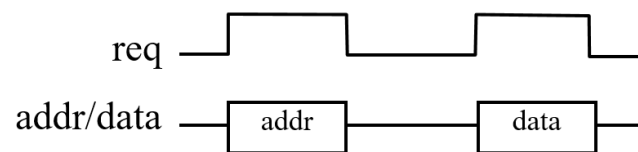
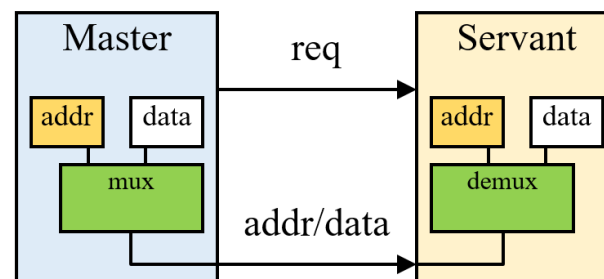
- Bus signals are usually tri-stated
- Multiple pieces of data can be shared on a single set of wires using time multiplexing
- Address and data lines may be multiplexed



E	A	Y
0	0	Z
0	1	Z
1	0	0
1	1	1



data serializing



address/data muxing



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Performance

- Bus Width
 - The number of data bits transmitted simultaneously.
 - Examples: 32-bit, 64-bit, etc.
 - Wider bus width → Higher data transfer rate → Increased bandwidth.
- Bus Bandwidth
 - The amount of data transmitted on a bus within a specific time.
 - Measurement: Max **data transfer rate** (usually in MB/s).
- Bus Frequency
 - Speed at which the bus operates, might be different than clock frequency.
 - Higher clock frequency → Faster bus operation → Increased bandwidth.

$$\text{Bus Bandwidth} = \text{Bus Width} \times \text{Bus Frequency}$$

Example 1

$$\text{Bandwidth} = \text{Width} \times \text{Frequency}$$

- Example: In the 32-bit bus system, the data and address are transferred on a shared bus. The clock frequency is 66MHz, during each clock cycle, two data transfers occur. (one on the rising edge and one on the falling edge).
 1. What is the maximum data transfer rate (bus bandwidth)?
 2. Assume it takes 1 clock cycle to transmit the address. How much time needed for CPU writing 128 bit data to the memory if it's in burst transfer mode.
- Solution Q1:
 - 1 clock cycle: 2 bus transfers \rightarrow bus frequency = $2 * 66 \text{ MHz} = 132 \text{ MHz}$
 - Bus width = 32 bit = 4 byte
 - Bus bandwidth = $4 * 132 = 528 \text{ MB/s}$

Example 1

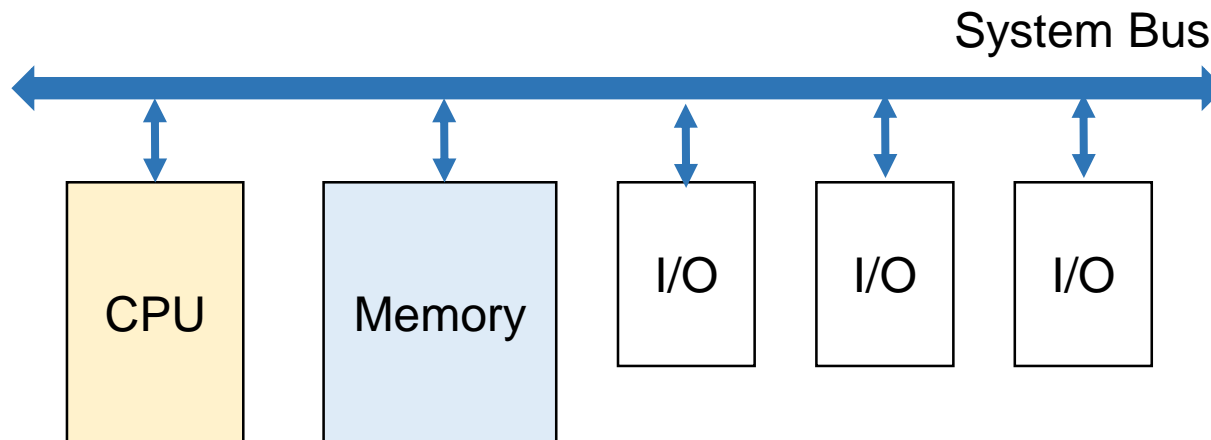
- Example: In the 32-bit bus system, the data and address are transferred on a shared bus. The bus clock frequency is 66MHz, during each clock cycle, two data transfers occur. (one on the rising edge and one on the falling edge).
 1. What is the maximum data transfer rate (bus bandwidth)?
 2. Assume it takes 1 clock cycle to transmit the address. How much time needed for CPU writing 128 bit data to the memory if it's in burst transfer mode.
- Solution Q2:
 - To transfer 128bits \rightarrow 4 words data, it takes 1 clock cycle for address, and 2 clock cycles for data
 - Latency = 3 * clock period = $3/66\text{MHz} = 45\text{ns}$

Example 2

- In a 64-bits bus system, the bus clock frequency of 1GHz, and it takes 1 clock cycle to transmit the address/data. The burst transfer is not supported. The memory needs 6ns to prepare 64-bits data, how much time it takes to read 32 bytes data from memory.
- Solution:
 - Clock period = $1/1\text{GHz} = 1\text{ns}$
 - Bus width = $64/8 = 8\text{bytes}$
 - For each 8 bytes data transfer \rightarrow 1 cycle for transferring address, 1 cycle for transferring data, 6 cycles for preparing data \rightarrow 8 cycles
 - To read 32 bytes data \rightarrow 4 bus transfers \rightarrow 32 clock cycles
 - Latency = $32 * \text{clock period} = 32\text{ns}$

Bus Architecture

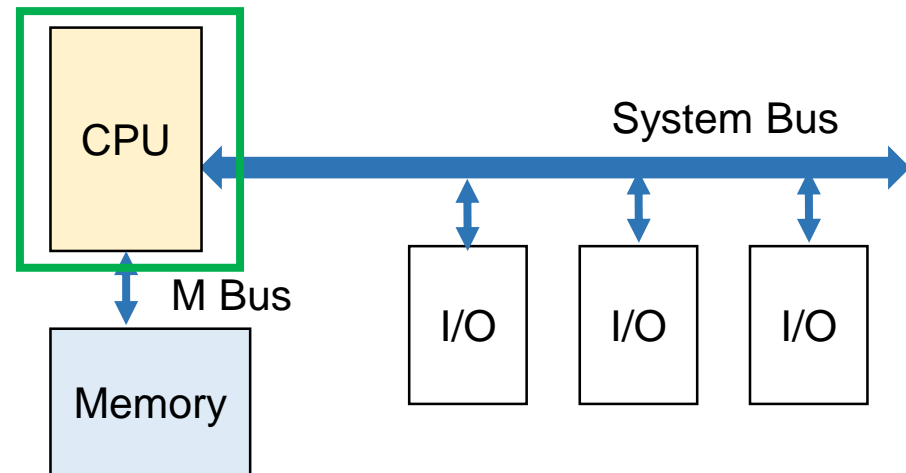
- Single level bus architecture
 - In a single bus structure, one common bus is used to communicate between peripherals and microprocessors.
 - Facilitates adding and removing IO devices.
 - Conflicts occur when multiple components need to use the bus, and priority needs to be determined.
 - Do not support burst transfer(突发传输)



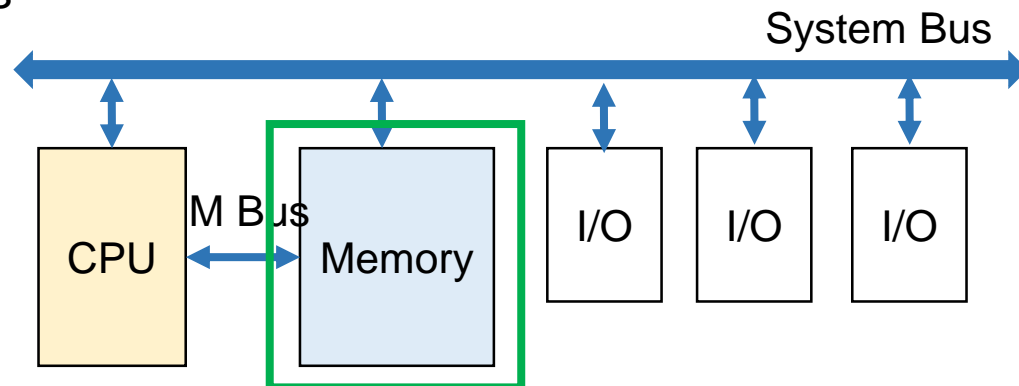
burst transfer refers to the consecutive transfer of multiple data elements, initiated by a single address.

Bus Architecture

- Two level bus architecture
 - CPU centralized
 - Ease of adding or removing IO devices

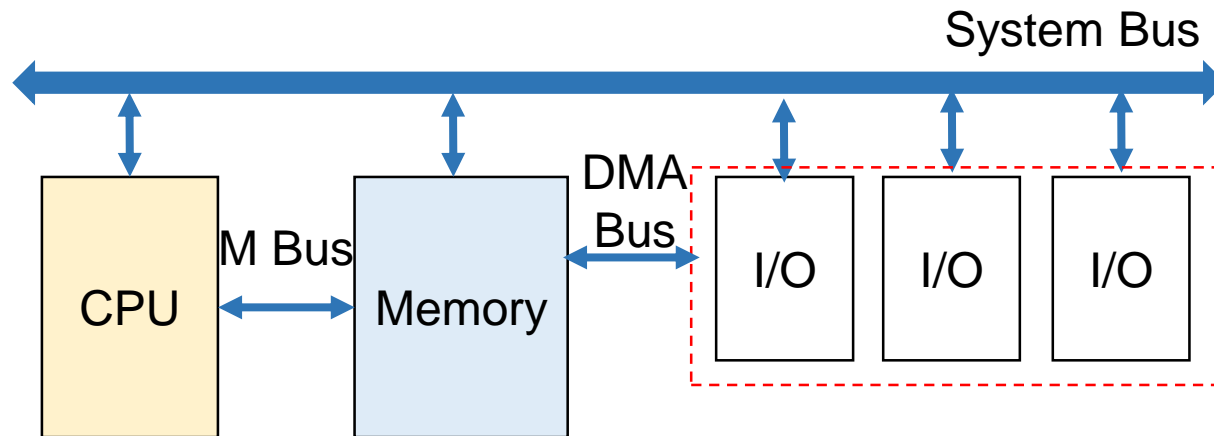


- Memory centralized
 - Memory bus (M-bus) operates at a high speed, reducing the burden on the system bus



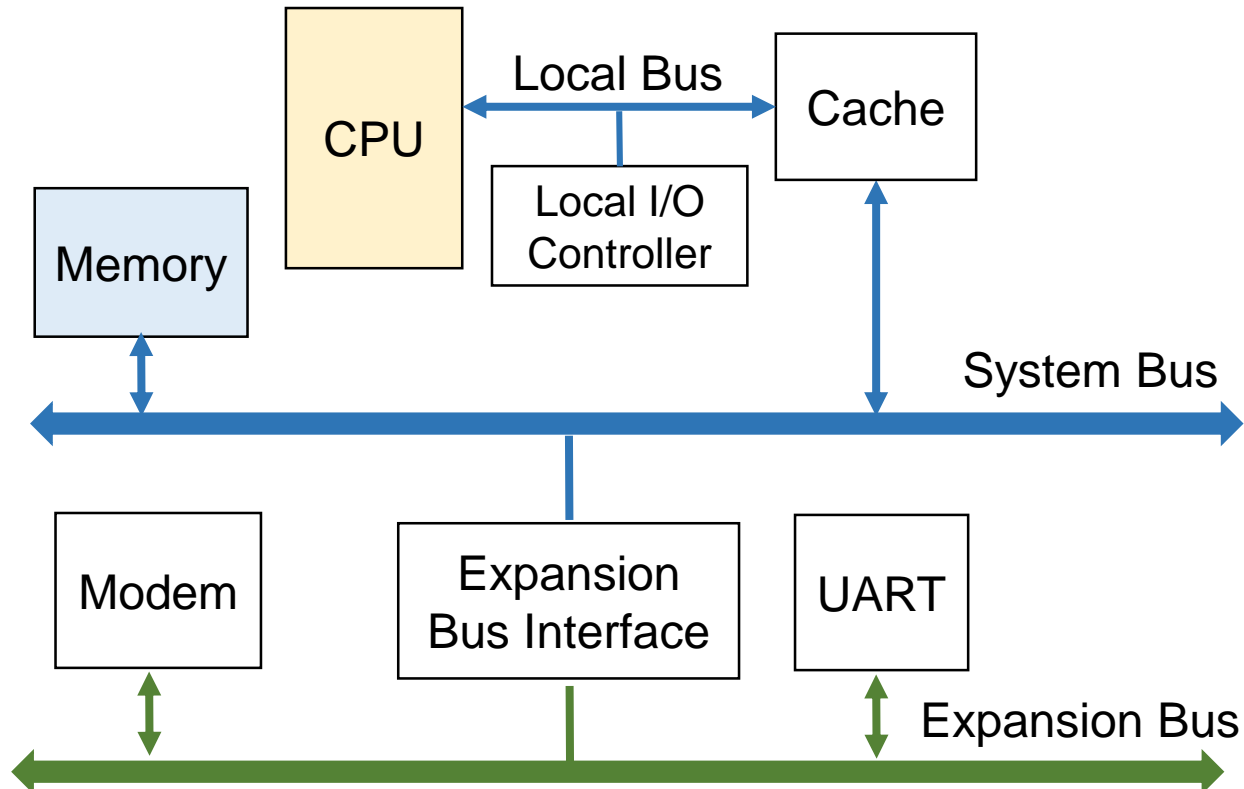
Bus Architecture

- Multilevel bus architecture
 - Memory centralized with DMA
 - Information exchange between IO devices and main memory doesn't require CPU intervention.



Bus Architecture

- Multilevel bus architecture
 - More advanced
 - Allowing various components to connect efficiently while maintaining high performance and data integrity.



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- Bus Introduction
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- **Bus Timing**
- SoC Bus - AMBA

Bus Transaction Steps

1. Request:

- a device that wishes to initiate a transaction on the bus sends a request signal to the bus controller or arbiter
- Arbitration: If there are multiple devices contending for bus access, an arbitration process occurs. The arbiter determines which device should be granted access based on arbitration scheme

2. Addressing:

- Master sends address and control signal to the bus

3. Transfer:

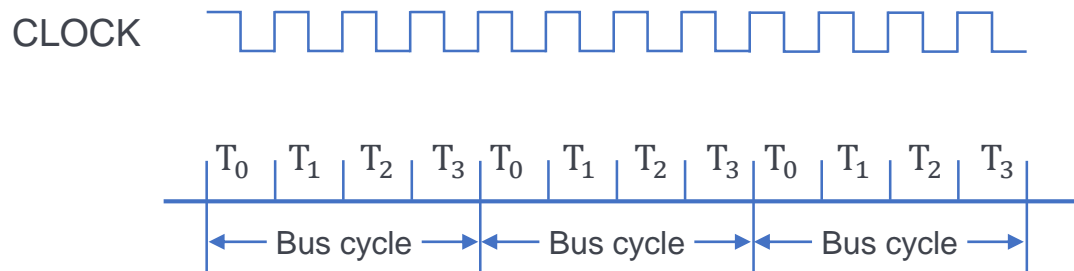
- Master and slave perform data transmission.

4. Complete:

- After the data transfer is complete, master removes relevant information from the bus and release the control from the bus.

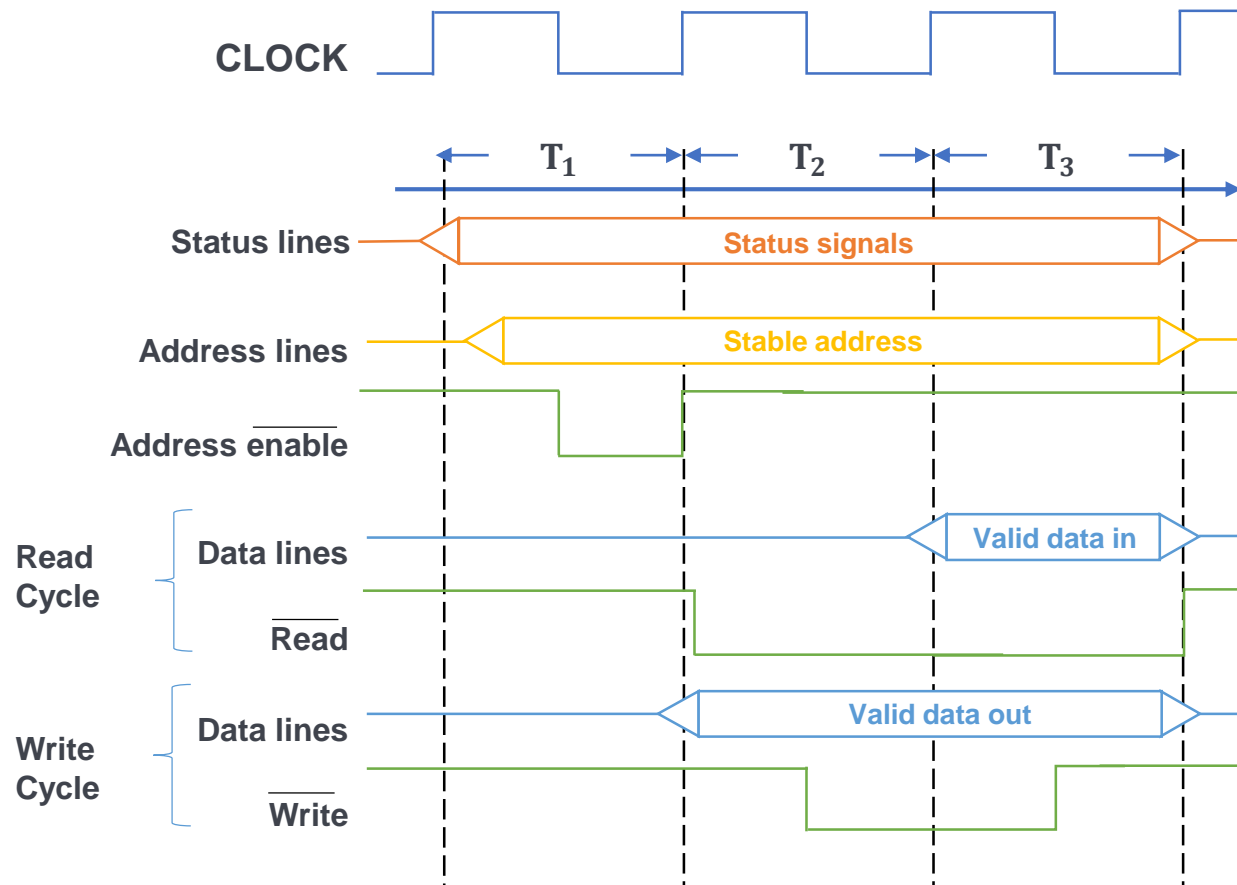
Bus Timing

- Co-ordination of events on bus
 - Synchronous – controlled by a clock
 - Asynchronous – timing is handled by well-defined specifications, i.e., a response is delivered within a specified time after a request
 - An event like CPU Reading from Main Memory:
 - Sending the address to the main memory
 - Sending the read signal to the main memory
 - Delivering the data on the data bus to the CPU



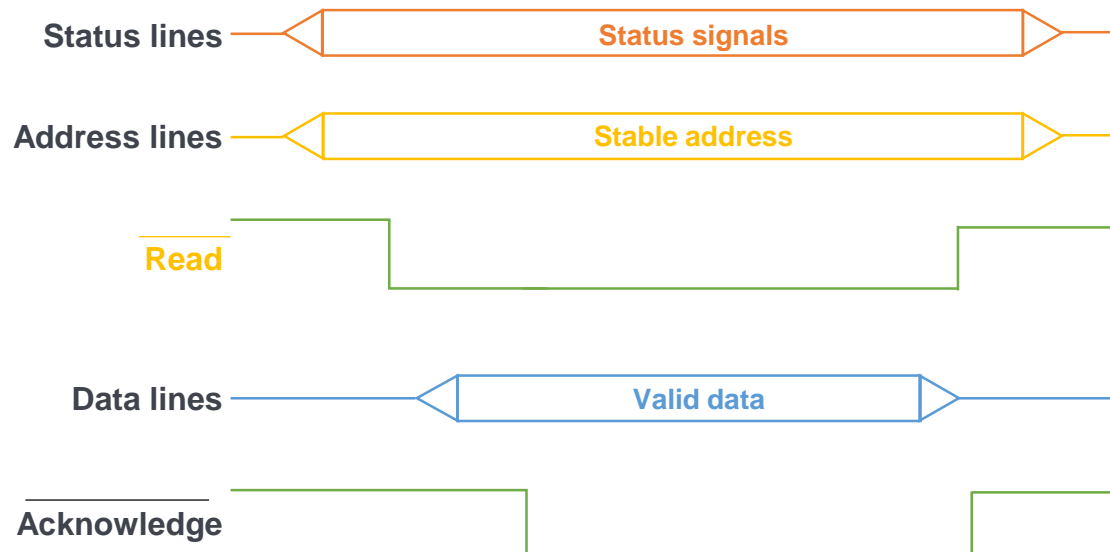
Synchronous Timing Diagram

- CPU(Master) reads/writes data from/to memory(Slave)



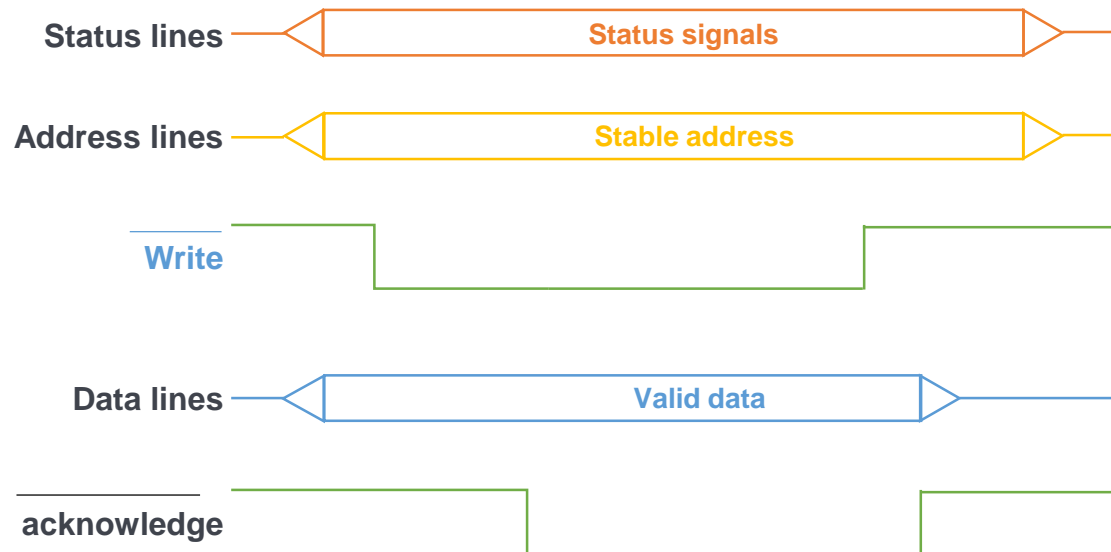
Asynchronous Timing- Read Diagram

- CPU(Master) reads data from memory(Slave)



Asynchronous Timing- Write Diagram

- CPU(Master) writes data to memory(Slave)



Bus Arbitration

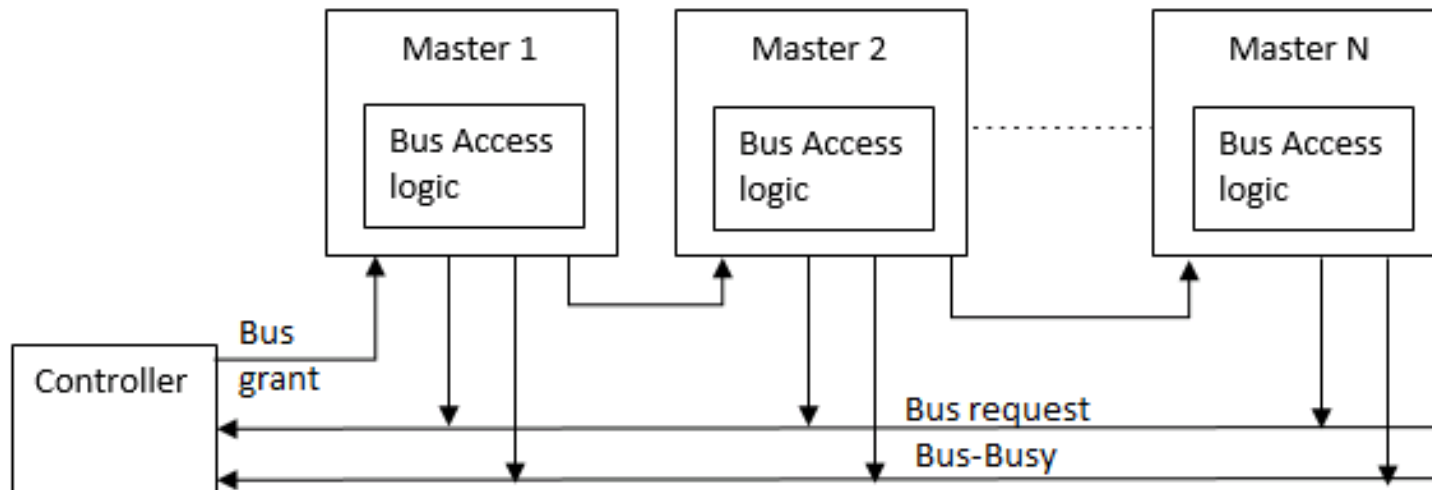
- Conflicts when more than one module controlling the bus
 - e.g. CPU and DMA controller
 - Only one module may control bus at one time
- Arbitration: choosing a master among multiple requests
 - Centralised
 - Single hardware device controlling bus access
 - Bus Controller
 - Arbiter
 - May be part of CPU or separate
 - Distributed
 - Each module may claim the bus
 - Control logic on all modules

Bus Arbitration Schemes

- Centralized Arbitration
 - Serial Arbitration (Daisy Chain)
 - Devices connected in a chain.
 - Sequentially select the bus controller.
 - Parallel Arbitration (Independent Request)
 - Devices independently request bus access.
- Distributed Arbitration
 - Arbitration handled within individual bus devices.
 - Device failures don't affect others.

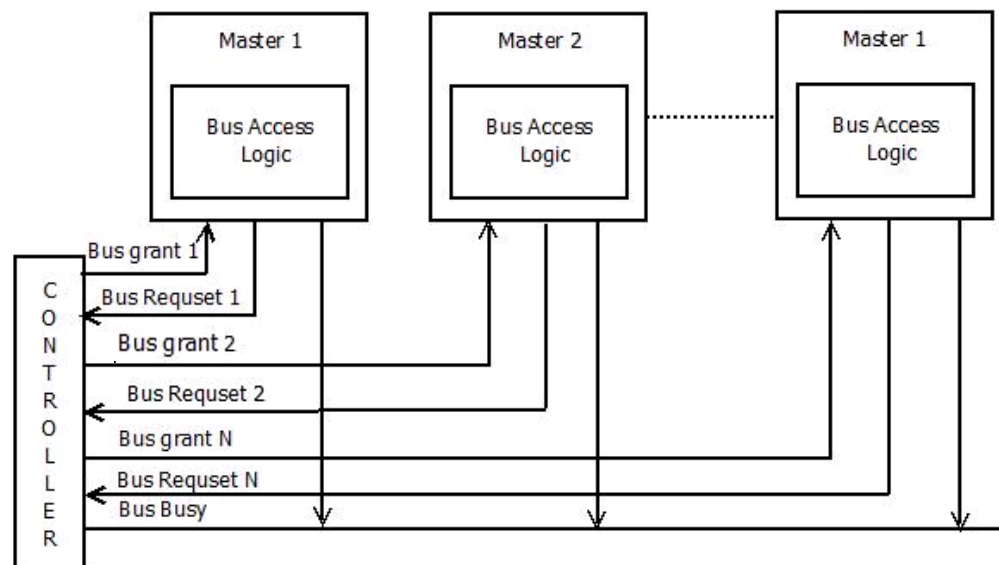
Daisy Chain Arbitration

- All bus master use the same line for bus request.
- If the bus busy line is inactive, the bus controller gives the bus grant signal.
- Bus grant signal is propagated serially through all masters starting from nearest one.



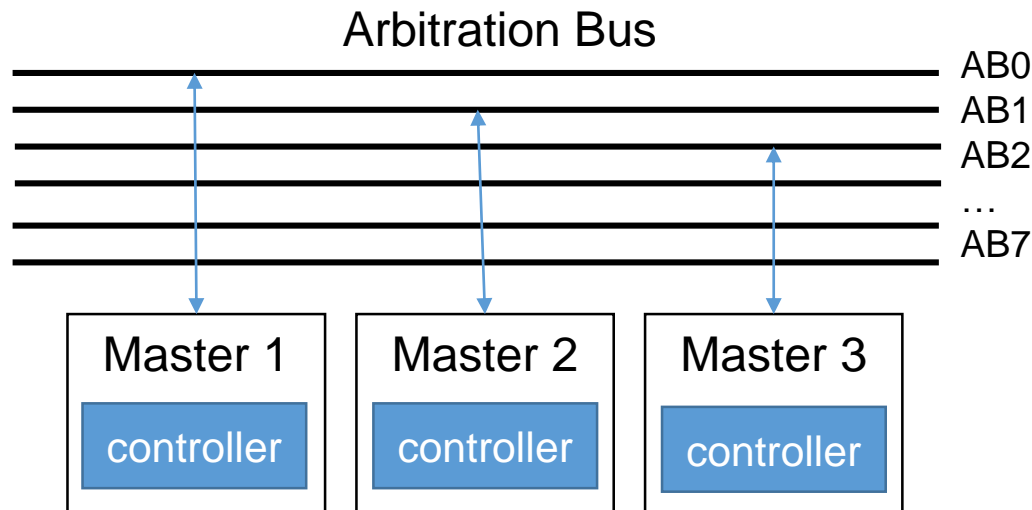
Independent Request Arbitration

- All bus masters have their individual bus request and bus grant lines.
- The controller knows which master has requested, so bus is granted to that master.
- Priorities of the masters are predefined so on simultaneous bus requests, the bus is granted based on the priority.



Distributed Arbitration

- The arbiters are located within each device
 - The bus device arbiter sends its own arbitration number onto the shared bus.
 - The arbiter compares the arbitration number it receives from the bus with its own to determine priority.
 - If its own priority is lower, it withdraws its arbitration number.
 - The highest-priority arbiter remaining on the bus gains control of the bus.



Outline

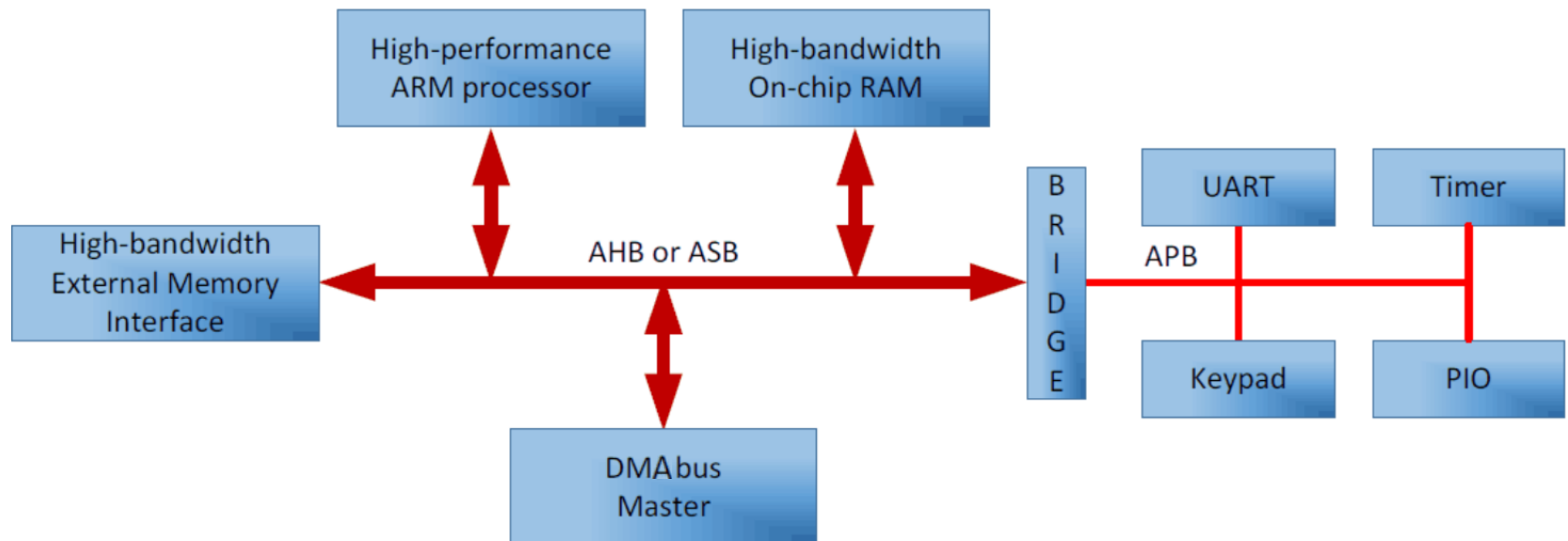
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- **SoC Bus - AMBA**

SOC Bus Architecture

- Advanced Microcontroller Bus Architecture
 - On-chip bus protocol from ARM
 - On-chip interconnect specification for the connection and management of functional blocks including processor and peripheral devices
 - **AMBA2.0**
 - AHB
 - Widely used on ARM7, ARM9 and **ARM Cortex-M** based designs
 - ASB
 - APB
 - **AMBA3.0**
 - AXI3 (or AXI v1.0)
 - widely used on ARM Cortex-A processors including Cortex-A9
 - AHB-Lite v1.0
 - APB3 v1.0
 - ATB v1.0
- AHB: Advanced High-performance Bus
ASB: Advanced System Bus
APB: Advanced Peripheral Bus

AMBA2.0

- ARM Cortex-M processor uses AMBA2.0 to interface to the System Bus
 - Enables IP re-use
 - Flexibility
 - Compatibility
 - Well supported



AHB and APB

- AHB (Advanced High-performance Bus)
 - High Performance
 - Connects CPUs, High-Speed Memory, DMA Controllers
 - High Bandwidth Data Transfer
 - Multiple Masters Supported
 - Ideal for High-Performance Components
- APB (Advanced Peripheral Bus)
 - Low-Power Peripheral Bus
 - Connects Low-Speed Peripherals
 - Suitable for Timers, GPIO Controllers, UART, etc.
 - Designed for Low-Speed, Low-Power Applications
 - Connects Peripheral Devices to System Bus

STM32 System Diagram

