### INFORMATION THEORY & CODING

Part 4 : Asymptotic Equipartition Property (AEP)

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October 10, 2023



### Stock Market

 Initial investment Y<sub>0</sub>, daily return ratio r<sub>i</sub>, in t-th day, your money is

$$Y_t = Y_0 r_1 \cdot \ldots \cdot r_t$$
.

• Now if returns ratio  $r_i$  are i.i.d., with

$$r_i = \begin{cases} 4, & \text{w.p. } 1/2\\ 0, & \text{w.p. } 1/2 \end{cases}$$

- So you think the expected return ratio is  $E[r_i] = 2$ .
- And then

$$E[Y_t] = E[Y_0r_1 \cdot ... \cdot r_t] = Y_0(E[r_i])^t = Y_02^t$$
???



### Stock Market

- With  $Y_0 = 1$ , actual return  $Y_t$  goes like
  - 1 4 16 0 0 ...

- Why?
  - The 'typical' sequences will end up with 0 return.
  - Occasionally, we got high return.
  - The expected return is increasing.
  - Expectation does not show the typical feature of this random sequence. We can turn to typical set.



## Weak Law of Large Numbers

### Theorem (Weak Law of Large Numbers)

Suppose that  $X_1, X_2, \dots, X_n$  are n independent, identically distributed (i.i.d.) random variables, then

$$\frac{1}{n}\sum_{i=1}^{n}X_{i}\to E[X] \qquad \text{in probability},$$

i.e. for every number  $\epsilon > 0$ ,

$$\lim_{n\to\infty} \Pr\left[\left|\frac{1}{n}\sum_{i=1}^n X_i - E[X]\right| \le \epsilon\right] = 1.$$



## Definition (Convergence of random variables)

Given a sequence of random variables,  $X_1, X_2, ...$ , we say that the sequence  $X_1, X_2, ...$  converges to a random variable X:

- **1** In probability if for every  $\epsilon > 0$ ,  $\Pr[|X_n X| \ge \epsilon] \to 0$
- ② In mean square if  $E[(X_n X)^2] \rightarrow 0$
- **3** With probability 1 (a.k.a. almost surely) if  $\Pr[\lim_{n\to\infty} X_n = X] = 1$



### Theorem 3.1.1 (AEP)

If 
$$X_1, X_2, \ldots$$
 are i.i.d.  $\sim p(x)$ , then 
$$-\frac{1}{n}\log p(X_1, X_2, \ldots, X_n) \to H(X) \qquad \text{in probability}.$$

### Proof.

Since  $X_i$  are i.i.d., so are  $\log p(X_i)$ . Hence, by the weak law of large numbers,

$$-rac{1}{n}\log p\left(X_1,X_2,\ldots,X_n
ight) = -rac{1}{n}\sum_i\log p\left(X_i
ight) \ 
ightarrow -E[\log p(X)] \qquad ext{in probability} \ = H(X)$$

## Typical Set

#### **Definition**

A *typical set*  $A_{\epsilon}^{(n)}$  contains all sequence realizations

$$(x_1, x_2, \dots, x_n) \in \mathcal{X}^n$$
 with

$$2^{-n(H(X)+\epsilon)} \le p(x_1, x_2, \dots, x_n) \le 2^{-n(H(X)-\epsilon)}$$
.



#### Theorem 3.1.2

- If  $(x_1, x_2, \dots, x_n) \in A_{\epsilon}^{(n)}$ , then  $H(X) \epsilon \le -\frac{1}{n} \log p(x_1, x_2, \dots, x_n) \le H(X) + \epsilon.$
- $\Pr[(X_1, X_2, ..., X_n) \in A_{\epsilon}^{(n)}] > 1 \epsilon$  for n sufficiently large.
- $|A_{\epsilon}^{(n)}| \leq 2^{n(H(X)+\epsilon)}$ , where |A| denotes the cardinality of the set A.
- $|A_{\epsilon}^{(n)}| \ge (1 \epsilon)2^{n(H(X) \epsilon)}$  for n sufficiently large.

### Proof.

1. Immediate from the definition of  $A_{\epsilon}^{(n)}$ .



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#### Theorem 3.1.2

- If  $(x_1, x_2, \dots, x_n) \in A_{\epsilon}^{(n)}$ , then  $H(X) \epsilon \le -\frac{1}{n} \log p(x_1, x_2, \dots, x_n) \le H(X) + \epsilon.$
- $\Pr[A_{\epsilon}^{(n)}] > 1 \epsilon$  for n sufficiently large.
- $|A_{\epsilon}^{(n)}| \leq 2^{n(H(X)+\epsilon)}$ , where |A| denotes the cardinality of the set A.
- $|A_{\epsilon}^{(n)}| \ge (1 \epsilon)2^{n(H(X) \epsilon)}$  for n sufficiently large.

#### Proof.

2. By Theorem 3.1.1, the probability of the event  $(X_1, X_2, \ldots, X_n) \in A_{\epsilon}^{(n)}$  tends to 1 as  $n \to \infty$ . Thus, for any  $\delta > 0$ , there exists an  $n_0$  such that for all  $n \ge n_0$ , we have

$$\Pr\left\{\left|-\frac{1}{n}\log p\left(X_1,X_2,\ldots,X_n\right)-H(X)\right|<\epsilon\right\}>1-\delta.$$

Setting  $\delta = \epsilon$ , the conclusion follows.

#### Theorem 3.1.2

- If  $(x_1, x_2, \dots, x_n) \in A_{\epsilon}^{(n)}$ , then  $H(X) \epsilon \le -\frac{1}{n} \log p(x_1, x_2, \dots, x_n) \le H(X) + \epsilon.$
- $\Pr[A_{\epsilon}^{(n)}] > 1 \epsilon$  for n sufficiently large.
- $|A_{\epsilon}^{(n)}| \leq 2^{n(H(X)+\epsilon)}$ , where |A| denotes the cardinality of the set A.
- $|A_{\epsilon}^{(n)}| \ge (1 \epsilon)2^{n(H(X) \epsilon)}$  for n sufficiently large.

### Proof.

3.

$$1 = \sum_{\mathbf{x} \in \mathcal{X}^n} p(\mathbf{x}) \ge \sum_{\mathbf{x} \in A_{\epsilon}^{(n)}} p(\mathbf{x})$$

$$\ge \sum_{\mathbf{x} \in A_{\epsilon}^{(n)}} 2^{-n(H(X) + \epsilon)}$$

$$= 2^{-n(H(X) + \epsilon)} |A_{\epsilon}^{(n)}|.$$

#### Theorem 3.1.2

- If  $(x_1, x_2, \dots, x_n) \in A_{\epsilon}^{(n)}$ , then  $H(X) \epsilon \le -\frac{1}{n} \log p(x_1, x_2, \dots, x_n) \le H(X) + \epsilon.$
- $\Pr[A_{\epsilon}^{(n)}] > 1 \epsilon$  for n sufficiently large.
- $|A_{\epsilon}^{(n)}| \leq 2^{n(H(X)+\epsilon)}$ , where |A| denotes the cardinality of the set A.
- $|A_{\epsilon}^{(n)}| \ge (1 \epsilon)2^{n(H(X) \epsilon)}$  for n sufficiently large.

### Proof.

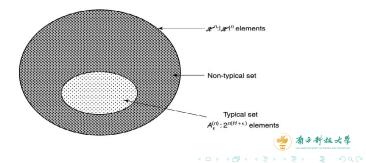
4. For sufficiently large n,  $\Pr[A_{\epsilon}^{(n)}] > 1 - \epsilon$ , so that

$$\begin{split} 1 - \epsilon &< \Pr\left[A_{\epsilon}^{(n)}\right] \\ &\leq \sum_{\mathbf{x} \in A_{\epsilon}^{(n)}} 2^{-n(H(X) - \epsilon)} \\ &= 2^{-n(H(X) - \epsilon)} \left|A_{\epsilon}^{(n)}\right|. \end{split}$$

## Typical set diagram

This enables us to divide all sequences into two sets

- Typical set: high probability to occur, sample entropy is close to true entropy
  - so we will focus on analyzing sequences in typical set
- Non-typical set: small probability, can ignore in general



#### Theorem 3.2.1

Let  $X_1, X_2, ..., X_n$  be i.i.d. random variables with distribution p(x), and  $X^n = X_1 X_2 ... X_n$ . For arbitrarily small  $\epsilon > 0$ , there exists a code that maps every realization  $x^n = x_1 x_2 ... x_n$  of  $X^n$  into one binary string, such that the mapping is one-to-one (and therefore invertible) and

$$E\left[\frac{1}{n}\ell(X^n)\right] \le H(X) + \epsilon$$

for a sufficiently large n.



#### Theorem 3.2.1

$$E\left[\frac{1}{n}\ell(X^n)\right] \leq H(X) + \epsilon.$$

for n sufficiently large.

#### Proof.

Description in typical set requires no more than  $n(H(X) + \epsilon) + 1$  bits (correction of 1 bit because of integrality).

Description in atypical set  $A_{\epsilon}^{(n)^C}$  requires no more than  $n \log |\mathcal{X}| + 1$  bits.

Add another bit to indicate whether in  $A_{\epsilon}^{(n)}$  or not to get whole description.

#### Theorem 3.2.1

$$E[\frac{1}{n}\ell(X^n)] \le H(X) + \epsilon.$$

for n sufficiently large.

#### Proof.

Let  $\ell(x^n)$  be the length of the binary description of  $x^n$ . Then,  $\forall \epsilon > 0$ , there exists  $n_0$  s.t.  $\forall n > n_0$ ,  $P(x^n) = \sum_{x^n \in A_\epsilon^{(n)}} p(x^n) \ell(x^n)$   $= \sum_{x^n \in A_\epsilon^{(n)}} p(x^n) \ell(x^n) + \sum_{x^n \in A_\epsilon^{(n)}} p(x^n) \ell(x^n)$   $\leq \sum_{x^n \leq A_\epsilon^{(n)}} p(x^n) (n(H+\epsilon)+2) + \sum_{x^n \in A_\epsilon^{(n)}} p(x^n) (n \log |\mathcal{X}|+2)$   $= \Pr[A_\epsilon^{(n)}](n(H+\epsilon)+2) + \Pr[A_\epsilon^{(n)}](n \log |\mathcal{X}|+2)$   $\leq n(H+\epsilon) + \epsilon n(\log |\mathcal{X}|) + 2$   $= n(H+\epsilon')$ 

where  $\epsilon' = \epsilon + \epsilon \log |\mathcal{X}| + \frac{2}{r}$  can be made arbitrarily small by choosing *n* properly.

## Reading

Reading: Whole Chapter 3

