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Prove for $i \neq 1$, $\sum_{c \in \mathcal{C}} \Pr(c) \Pr[e_i(c) | W=1] \leq 2^{-n[I(X;Y) - 3\epsilon]}$

$$\begin{aligned}
 \text{Proof. } & \sum_{c \in \mathcal{C}} \Pr(c) \Pr[e_i(c) | W=1] \\
 &= \sum_{c \in \mathcal{C}} \left(\prod_{i=1}^{2^{nR}} \Pr(x_c^n(i)) \right) \Pr[e_i(c) | W=1] \\
 &= \sum_{x_i^n} \sum_{c: x_c^n(i) = x_i^n} \prod_{i=1}^{2^{nR}} \Pr(x_c^n(i)) \Pr(x_i^n \text{ and } Y^n \text{ are joint typical} | W=1) \\
 &= \sum_{x_i^n} \cancel{\Pr(x_c^n(i))} \sum_{x_i^n} \Pr(x_i^n) \Pr(x_i^n \text{ and } Y^n \text{ are joint typical} | W=1) \\
 &\quad \times \sum_{c: x_c^n(i) = x_i^n} \prod_{j \neq i}^{2^{nR}} \Pr(x_c^n(j)) \\
 &= \sum_{x_i^n} \Pr(x_i^n) \Pr(x_i^n \text{ and } Y^n \text{ are joint typical} | W=1) \\
 &= \Pr(x_i^n \text{ and } Y^n \text{ are joint typical} | W=1) = \Pr(E_i | W=1)
 \end{aligned}$$

By the code generation process, $X^n(1)$ and $X^n(i)$ are independent for $i \neq 1$, so are Y^n and $X^n(i)$. Hence the probability that $X^n(i)$ and Y^n are jointly typical, i.e. $\Pr(E_i | W=1) =$

$$\Pr[(X^n(i), Y^n) \in A_{\epsilon}^{(n)} | W=1] \leq 2^{-n[I(X;Y) - 3\epsilon]}$$