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$$25.(a) L = \sum_{i=1}^m P_i l_i = \sum_{i=1}^m P_i \lceil \log \frac{1}{P_i} \rceil \quad \text{since } \log \frac{1}{P_i} \leq \lceil \log \frac{1}{P_i} \rceil < \log \frac{1}{P_i} + 1$$

$$\text{we have } \sum_{i=1}^m P_i \log \frac{1}{P_i} \leq \sum_{i=1}^m P_i \lceil \log \frac{1}{P_i} \rceil < \sum_{i=1}^m P_i \log \frac{1}{P_i} + \sum_{i=1}^m P_i = \sum_{i=1}^m P_i \log \frac{1}{P_i} + 1$$

$$\text{Hence } H(X) \leq L < H(X) + 1$$

$$\text{since } \log \frac{1}{P_i} \leq l_i < \log \frac{1}{P_i} + 1 \quad 2 \leq P_i \leq 2^{l_i+1}$$

Consider F_i and F_j , where $j > i$, F_j is different from F_i at least from the first i bit. $\sum_{i=1}^m 2^{-l_i} \leq \sum_{i=1}^m P_i = 1$ by Kraft inequality

$$(b) X = \{1, 2, 3, 4\}$$

Hence F_i can't be the prefix of F_j , the code is ① prefix codes,

$$F_1 = 0, F_2 = 0.1, F_3 = 0.11, F_4 = 0.111$$

$$= (0)_2, = (0.1)_2, = (0.11)_2, = (0.111)_2$$

Hence, the code constructed by ① this process is prefix-free.

$$l_1 = \log 2 = 1, l_2 = \log 4 = 2, l_3 = l_4 = \log 8 = 3$$

Hence, the code for the probability distribution $(0.5, 0.25, 0.125, 0.125)$

$$\text{is } C(1) = 0, C(2) = 10, C(3) = 110, C(4) = 111$$

verification:

$$H(X) = 0.5 + 0.25 \times 2 + 0.125 \times 3 \times 2 = \frac{7}{4} \quad \text{Here,}$$

$$L = E(L(X)) = 0.5 + 0.25 \times 2 + 0.125 \times 3 \times 2 = \frac{7}{4}$$

$H(X) = L$ Let $f(D) = 2(D^{-1} + D^{-2} + D^{-3})$ which is monotonic decreasing

$$2 \cdot \sum_{i=1}^6 D^{-l_i} = 2(D^{-1} + D^{-2} + D^{-3}) \leq 1 \rightarrow D = 3, 4, 5, \dots$$

Hence, a good lower bound on D is $D=3$.

explain the title of the problem: since the cardinality of this code alphabet is 3. we guess a Martian has 3 fingers?