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Agenda

- User-Mode Injections
 - ➤ Classic Code Injection Methods
 - ➤ Advanced Injections PowerLoader
 - **▶** PowerLoaderEx
 - ➤ PowerLoaderEx64
- Kernel-Mode Injections
 - ➤ Common Injection Methods
 - ► Introducing Trap Frame Injection
 - ➤ Codeless Code Injection
 - **➢** Demo
- Summary



User-Mode Code Injections



Classic Code Injection Methods

- Vanila Injection OpenProcess/ VirtualAllocEx/CreateRemoteThread/LoadLibrary
- GetThreadContext / SetThreadContext
- SetWindowsHookEx
- Applnit_DLLs and APPCERTDLLs
- ShimEng (InjectDll Tag)
- SetWindowLong
- QueueUserApc
- Replace Dependency DLL

•





PowerLoader – POP/POP/INJECT

Back to 2013:

- Loader used in many different dropper families (Gapz / Redyms / Carberp / Vabushky ...)
- Technique was leaked with Carberp source code leak.
- First injection technique via Return Oriented Programming technique (ROP).
- "explorer.exe" is injected using Shell_TrayWnd / NtQueueApcThread (32bit / 64bit)
- Successfully bypassed most HIPSs solutions
- 1. http://www.welivesecurity.com/2013/03/19/gapz-and-redyms-droppers-based-on-power-loader-code/
- 2. http://www.slideshare.net/matrosov/advanced-evasion-techniques-by-win32gapz
- 3. https://www.virusbtn.com/virusbulletin/archive/2012/10/vb201210-code-injection
- 4. http://www.malwaretech.com/2013/08/powerloader-injection-something-truly.html
- > Samples and valuable info http://www.kernelmode.info/forum/



PowerLoader – Abusing Shared Sections

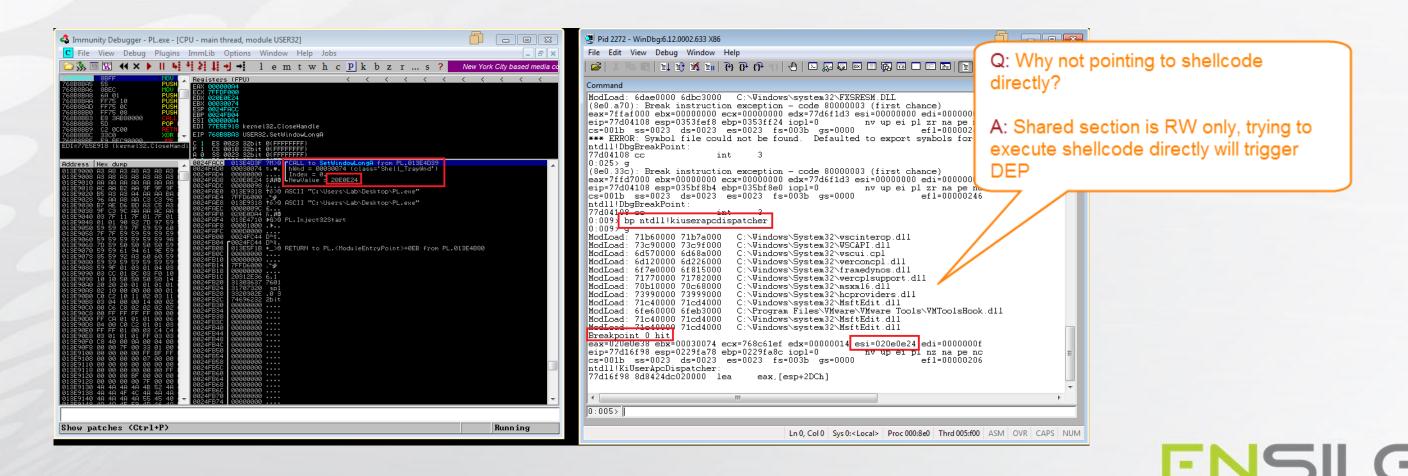
• Explorer.exe has several shared memory sections it uses:

- By calling NtOpenSection and ZwMapViewOfSection PowerLoader maps one of the shared sections to its own address space
- PowerLoader finds the address of the shared section in explorer by calling VirtualQueryEx / ReadProcessMemory / RtlCompareMemory combination
- The shellcode is than written to the shared section. There is no need for VirtualAllocEx / WriteProcessMemory which is monitored by many HIPS



PowerLoader - Triggering Code Execution

- One of the Windows in Explorer.exe is Shell TrayWnd
- The Shell_TrayWnd window has a pointer to CTray class object which handles messages to the window
- Once the Shell_TrayWnd window is found PowerLoader uses SetWindowLongPtr to replace the CTray object with a malicious CTray
 Object in the shared section
- The shared section is not executable thus ROP is needed
- The virtual table of the malicious CTray object point to KiUserApcDispatcher routine which eventually trigger the ROP execution
- To trigger code execution SendNotifyMessage is called with WM_PAINT message



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PowerLoader - ROP n' Roll

- PowerLoader uses a complex ROP chain
- Basic steps by the ROP chain (High-Level):
 - Direction Flag Set
 - Copy ROP chain to stack
 - Direction Flag Clear
 - Use WriteProcessMemory to overwrite ntdll!atan function (this function is probably not used in the context of explorer.exe)
 - Pass control to the overwritten ntdll!atan function



PowerLoader - Quick Summary

- Advanced Injection technique
- Completely bypassed most HIPS at the time
- Using Exploit-Like techniques:
 - The first Injection technique to use ROP (to our knowledge)
 - Overwrites atan function using WriteProcessMemory
 - No similar 64-bit version (to our knowledge)
 - Try to run vs EMET and see what happens...
- Very application specific:
 - Targets Shell_TrayWnd of explorer.exe
 - Uses specific shared sections
- Doesn't use direct code injections



Introducing PowerLoaderEx

Goals:

- Remove dependency in Explorer.exe shared sections (more generic)
- Make a 64-bit version
- Bonus: Do it without reading memory from the target process





UI Shared Memory - Reminder

- UI Objects are stored in one of three places:
 - Session Pool Allocated per user logon
 - Desktop Heap Stores objects specific to a given Desktop
 - Shared Heap Handle Table and object relevant to all Desktops
- The heaps are visible to user-mode via shared-memory
- Commonly used for exploiting privilege escalation vulnerabilities
- Extensively documented in:
 "Windows Hooks Of Death: Kernel Attacks through User-Mode Callbacks" By Tareji Mandt



UI Shared Memory

Kernel Space

User Space

Desktop Heap 1 **Shared Section**

Handle Shared
Table Heap

Desktop
Heap 1

Heap 1

Shared Section
Handle Shared
Heap
Heap
Application 1

Desktop
Heap 1

Shared Section
Handle
Table
Shared
Heap
Application 2



PowerLoaderEx – Desktop Heap as Shared Section

- Desktop Heaps are better as shared section because they are mapped to any process on a given desktop
- Writing arbitrary data to the Desktop Heap is easy:
 - Create a window with enough extra bytes
 - Use SetWindowLongPtr to write the payload
- But how do know where it resides in a target process?





Finding Target Process's Desktop Heap

- Find the Desktop Heap in our process and find its region size
- Enumerate explorer.exe memory regions by calling VirtualQuery.
- The target's Desktop Heap is found if the region characteristics are:
 - State: MEM COMMIT
 - Type: MEM MAPPED
 - Protect: PAGE_READ_ONLY
 - RegionSize is equal to the previously obtained heapSize.

Double Check – Avoiding False Positives

We had NO false positives in finding the shared desktop heap (local process and target process)



Finding The Gadgets

- We do the gadget lookup in our own process
- Load and scan only modules that we know that are always loaded in the target process (i.e. explorer.exe)
- Use only modules that are very likely to have the same base in both processes:
 - ntdll.dll
 - Kernel32.dll
 - Kernelbase.dll
 - User32.dll
 - Shell32.dll





PowerLoaderEx - Look Ma, No Read!

- 1. Create a window and find it in the shared desktop heap
- 2. Find the target process (Explorer.exe) and find where the desktop heap is mapped, requires only PROCESS_QUERY_INFORMATION privileges
- 3. Scan for the required gadgets
- 4. Write the payload to the shared desktop heap using SetWindowLongPtr
- 5. Continue like normal PowerLoader



^{*} Note: If a different target process has more than one Desktop Heap mapped read will be needed



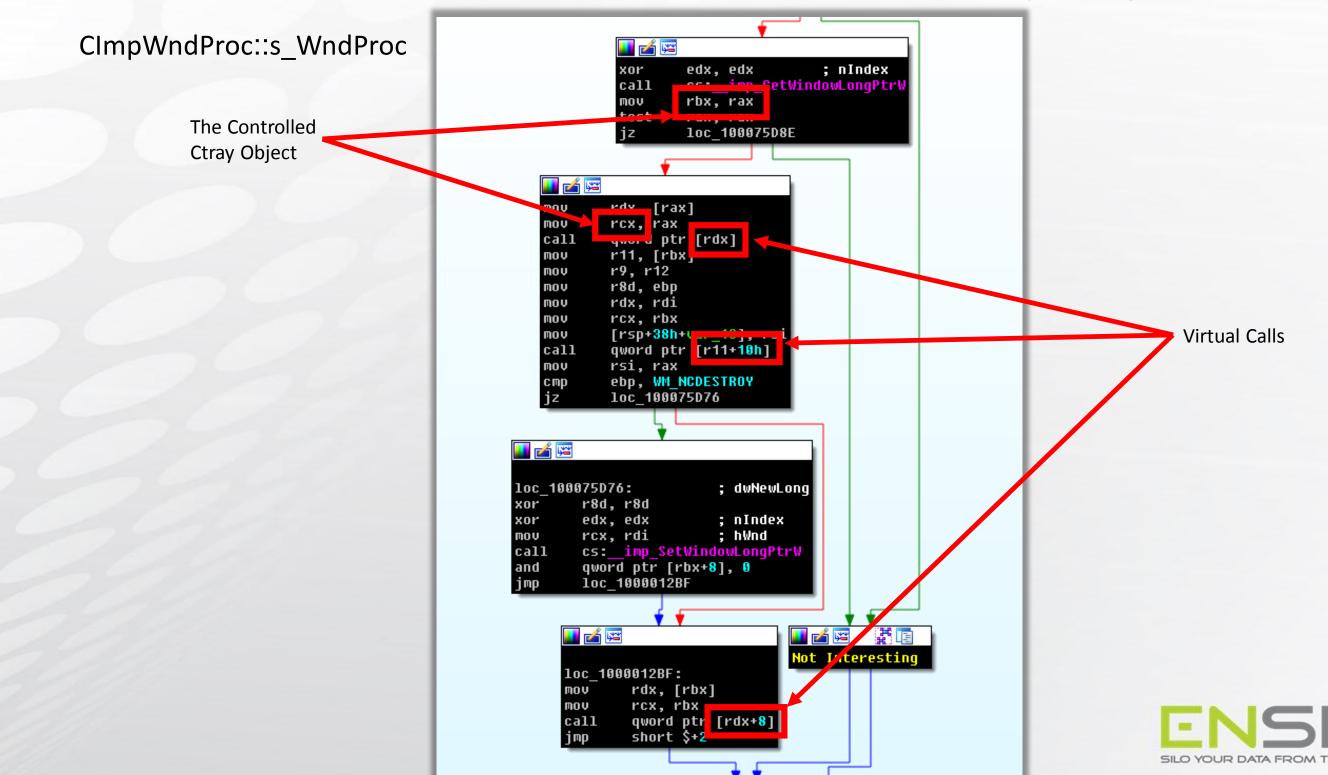
PowerLoaderEx64 - What About 64-bit?

- Challenges:
 - Much harder to find useful gadgets
 - No code writing allowed
 - No reads allowed





PowerLoaderEx64 – Remember the CTray Object?



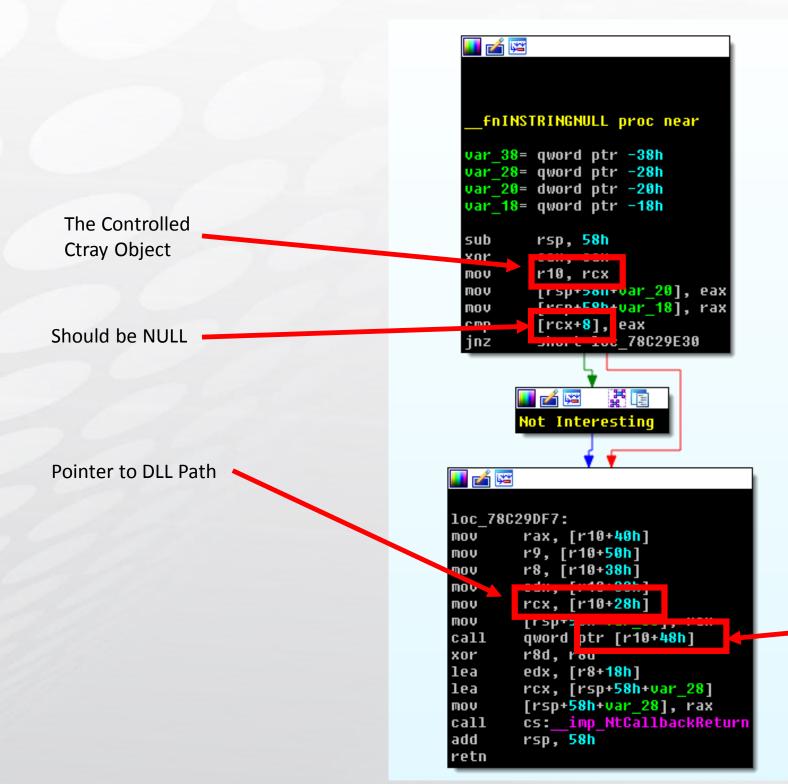
PowerLoaderEx64 - Strategy

- Find some function that can be manipulated to call LoadLibrary with controlled first argument
- Make sure no crashes occur after/during library loading
- To make things simple point 2 of the 3 virtual calls to "Do Nothing" functions (Ret opcode)





PowerLoaderEx64 – User32 Callbacks To Rescue



Pointer to LoadLibraryA



PowerLoaderEx64 - Recap

- 1. Create a window and find it in the shared desktop heap
- 2. Find the target process (Explorer.exe) and find where the desktop heap is mapped, requires only PROCESS_QUERY_INFORMATION privileges
- 3. Write the malicious CTray object to the shared desktop heap using SetWindowLongPtr
- 4. Replace the Shell_TrayWnd window's CTray object using SetWindowLongPtr
- 5. Use SendNotifyMessage to trigger LoadLibrary

PowerLoaderEx64 source code will be on BreakingMalware

http://BreakingMalware.com





Kernel-To-User Code Injections



Introduction - Kernel-To-User Code Injections

- Mainly used for:
 - Injecting DLLs
 - Sandbox escapes After exploiting privilege escalation vulnerability
 - Injecting to protected processes
- Fewer techniques exist than user-mode
- Less documented than user-mode techniques
- Used by both Malware and Software/Security vendors





Common Injection Methods – User APC

- The most common Kernel-To-User injection method
- Used by lots of malwares:
 - TDL
 - ZERO ACCESS
 - Sandbox escape shellcodes
 - •
- Also used by lots of security products:
 - AVG
 - Kaspersky Home Edition
 - Avecto
 - •
- Documented:
 - Blackout: What Really Happened
 - Much more in forums and leaked source codes





Common Injection Methods – User APC

Basic Steps (There are several variations):

- 1. Register load image callback using PsSetLoadImageNotifyRoutine and wait for main module to load
- 2. Write payload that injects a dll using LdrLoadDll (Other variations use LoadLibrary)
- 3. Insert User APC using KelnsertQueueApc

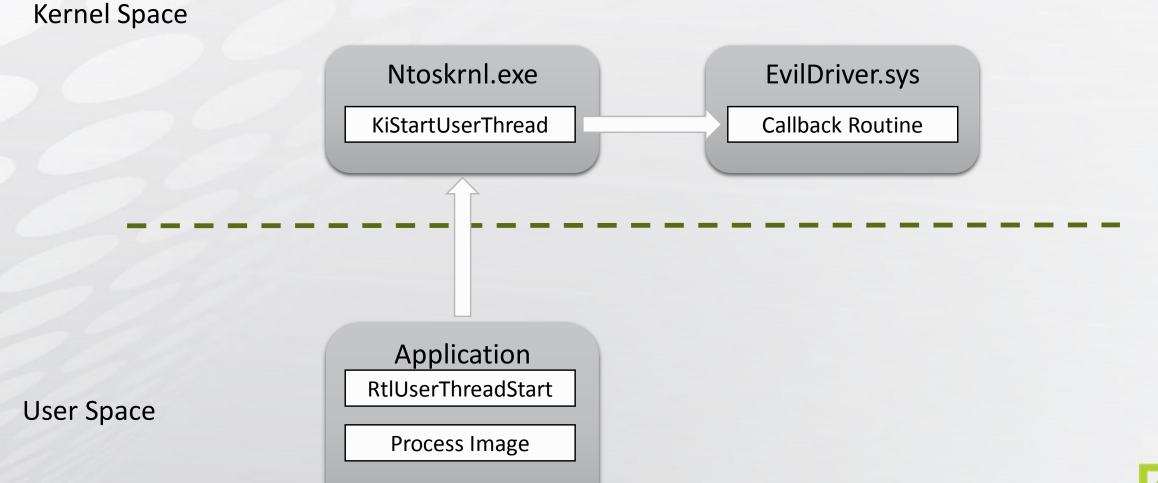


- Not really common but worth mentioning
- Used by Duqu
- Fully documented in: http://binsec.gforge.inria.fr/pdf/Malware2013-Analysis-Diversion-Duqu-paper.pdf



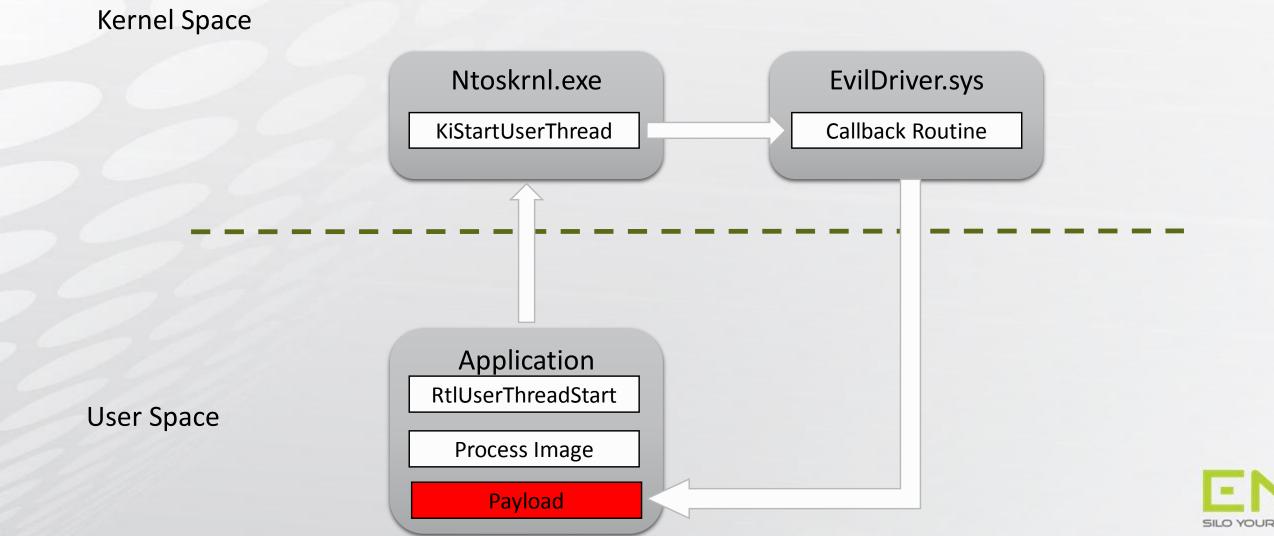


 Register load image callback using PsSetLoadImageNotifyRoutine and wait for main module to load

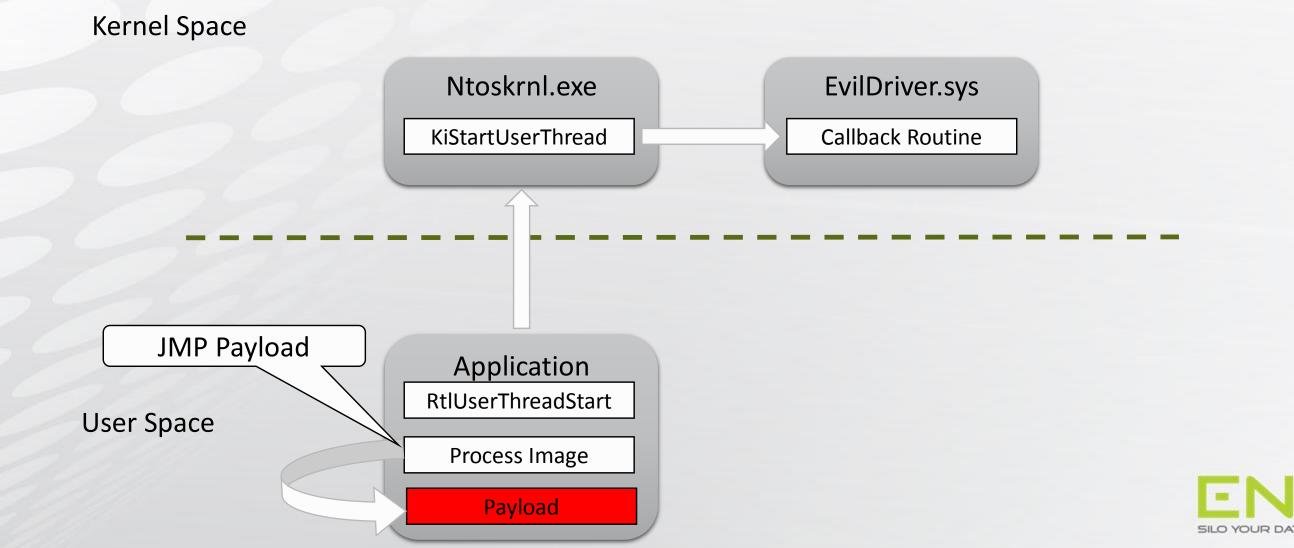




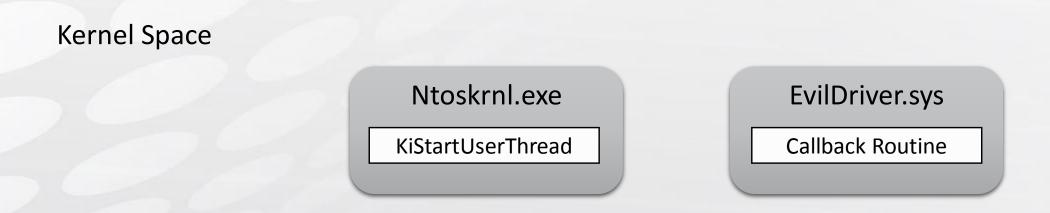
Write the payload to the process address space

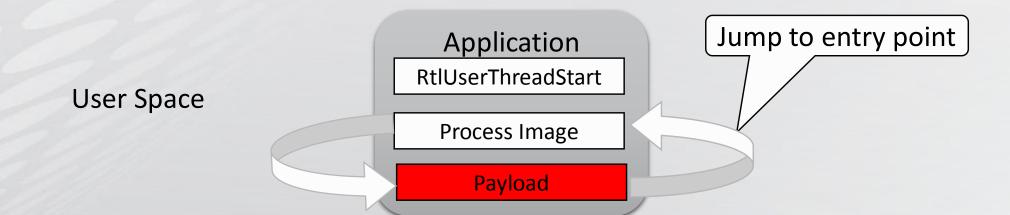


Replace the image entry point with JMP to the new code



• The payload executes, fixes the entry point and jumps to it







Common Injection Methods – Import Table Patching

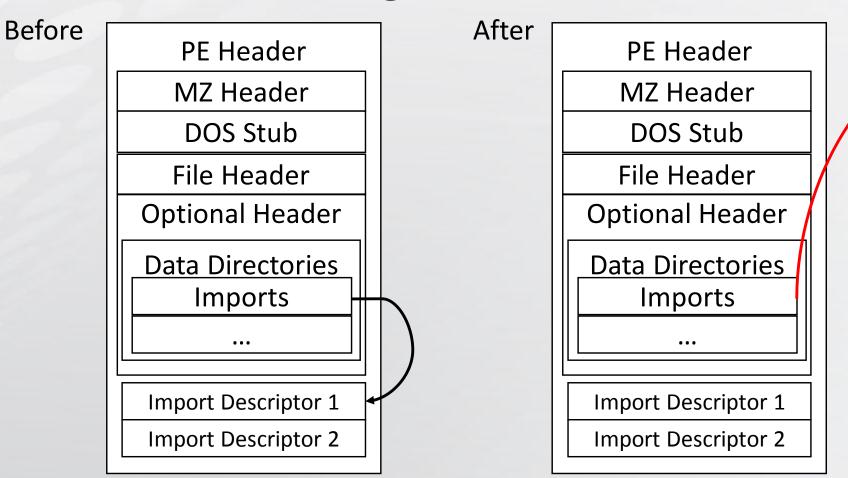
- Undocumented (to our knowledge)
- Never been used by malware (to our knowledge)
- Used by software and security vendors:
 - Symantec
 - Trusteer
 - Microsoft App-V
- Similar method could probably use TLS data directory





Common Injection Methods – Import Table Patching

- 1. Register load image callback using PsSetLoadImageNotifyRoutine and wait for main module to load
- 2. Allocate memory for the new import table and copy old table with a new record for the injected DLL
- 3. Point the import data directory to the new table
- 4. When the DLL is loaded the original PE header is restored



Injected DLL Descriptor
Import Descriptor 1
Import Descriptor 2



Common Injection Methods – Quick Summary

- Kernel-To-User Injections are extensively used by both malware and security/software vendors
- Kernel injections are mainly used to inject a DLL to target(or all) processes
- All these methods require injection of some payload to user mode (except for Import Table Patching)
- All use PsSetLoadImageNotifyRoutine



Introducing Trap Frame Injection

- A new kernel-to-user injection technique
- Trap frames save the CPU user-mode state during exceptions or interrupt handling
- A Trap frame is created each time a system call is made and stored on the kernel stack
- The kernel structure used for trap frames is _KTRAP_FRAME
- The user-mode state is restored when the system call returns
- Current frame is stored in the TrapFrame field of _KTHREAD





KTRAP_FRAME

32-Bit

```
kd> dt _ktrap_frame
nt!_KTRAP_FRAME
   +0x000 DbgEbp
                             : Uint4B
   +0x004 DbgEip
                             : Uint4B
   +0x008 DbgArgMark
                             : Uint4B
   +0x00c DbgArgPointer
+0x010 TempSegCs
                             : Uint4B
                             : Uint2B
   +0x012 Logging
                             : UChar
   +0x013 Reserved
                             : UChar
   +0x014 TempEsp
                             : Uint4B
   +0x018 Dr0
                             : Uint4B
   +0x01c Dr1
                             : Uint4B
   +0x020 Dr2
                             : Uint4B
   +0x024 Dr3
                             : Uint4B
   +0x028 Dr6
                             : Uint4B
   +0x02c Dr7
                             : Uint4B
   +0x030 SegGs
                             : Uint4B
   +0x034 SegEs
                             : Uint4B
   +0x038 SegDs
                             : Uint4B
   +0x03c Edx
                             : Uint4B
   +0x040 Ecx
                             : Uint4B
   +0x044 Eax
                             : Uint4B
   +0x048 PreviousPreviousMode: Uint4B
   +0x04c ExceptionList
                            : Ptr32 _EXCEPTION_REGISTRATION_RECORD
   +0x050 SegFs
                             : Uint4B
   +0x054 Edi
                             : Uint4B
   +0x058 Esi
                             : Uint4B
   +0x05c Ebx
                             : Uint4B
   +0x060 Ebp
                             : Uint4B
   +0x064 ErrCode
                             : Uint4B
   +0x068 Eip
                             : Uint4B
   +0x06c SegCs
                             : Uint4B
   +0x070 EF1ags
                             : Uint4B
   +0x074 HardwareEsp
                             : Uint4B
   +0x078 HardwareSegSs
                             : Uint4B
   +0x07c V86Es
                             : Uint4B
   +0x080 V86Ds
                             : Uint4B
   +0x084 V86Fs
                             : Uint4B
   +0x088 V86Gs
                             : Uint4B
```

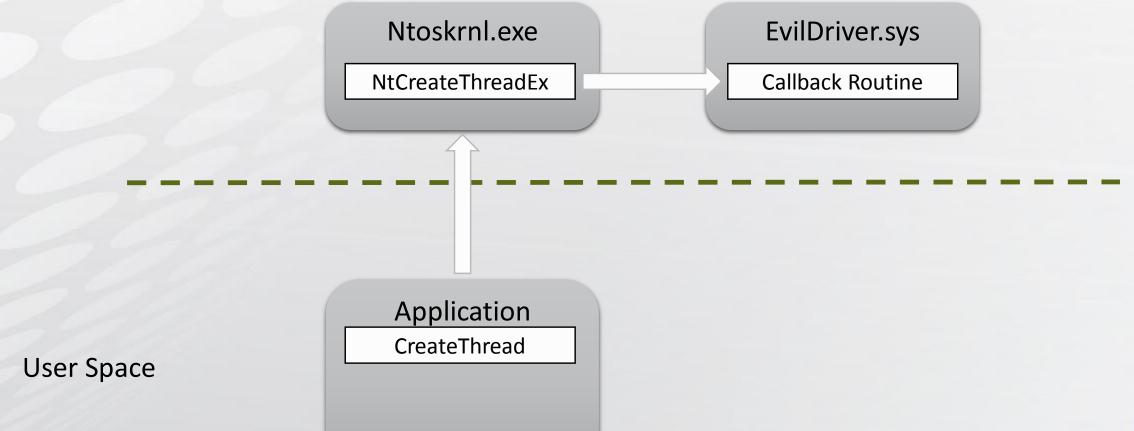
64-Bit

```
kd> dt _ktrap_frame
nt!_KTRAP_FRAME
   +0x030 Rax
                             : Uint8B
   +0x038 Rcx
                             : Uint8B
   +0x040 Rdx
                             : Uint8B
   +0x048 R8
                             : Uint8B
  +0x050 R9
                             : Uint8B
   +0x058 R10
                             : Uint8B
   +0x060 R11
                             : Uint8B
   +0x068 GsBase
                             : Uint8B
   +0x068 GsSwap
                             : Uint8B
  +0x0d0 FaultAddress
                            : Uint8B
   +0x0d0 ContextRecord
                             : Uint8B
  +0x0d0 TimeStampCKCL
                            : Uint8B
   +0x0d8 Dr0
                             : Uint8B
   +0x130 SegDs
                             : Uint2B
   +0x132 SegEs
                             : Uint2B
   +0x134 SegFs
                             : Uint2B
  +0x136 SegGs
                             : Uint2B
  +0x138 TrapFrame
                             : Uint8B
   +0x140 Rbx
                             : Uint8B
  +0x148 Rdi
                             : Uint8B
   +0x150 Rsi
                             : Uint8B
   +0x158 Rbp
                             : Uint8B
   +0x168 Rip
                             : Uint8B
   +0x170 SegCs
                             : Uint2B
                             : Uint8B
   +0x180 Rsp
```



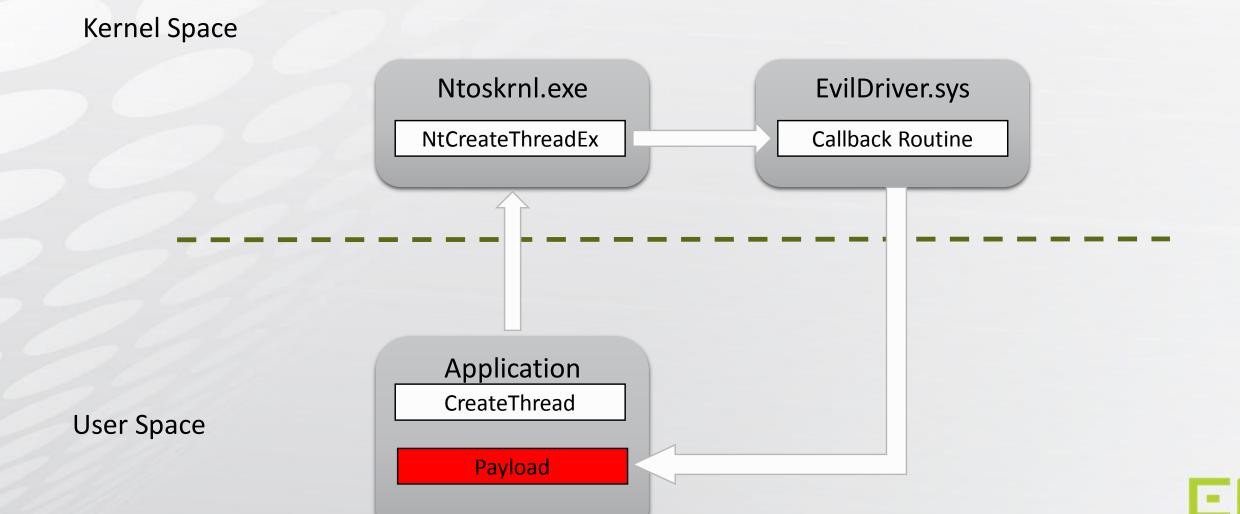
• Wait for some callback that runs in context of a system call (Thread Creation, Registry Access,...)

Kernel Space

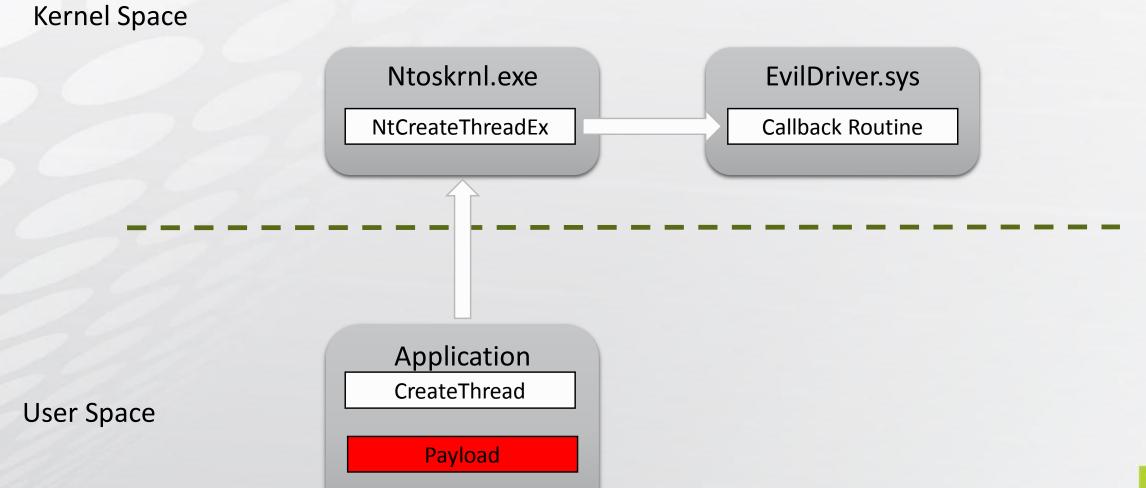




Allocate payload in the target application

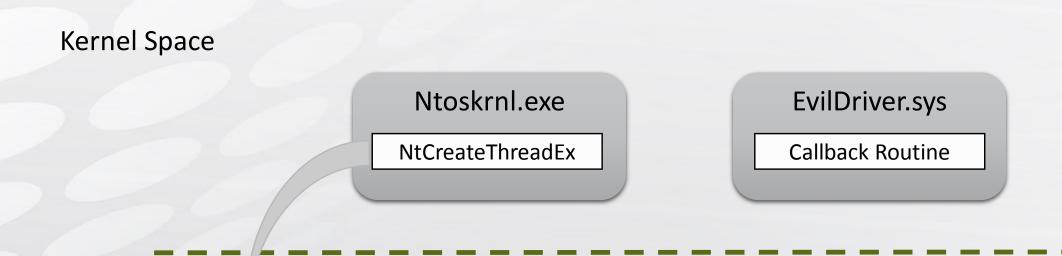


• Fix the trap frame's saved RIP/EIP to the payload





• The payload runs and restores normal execution



User Space

Application
CreateThread

Payload



But We Promised a Codeless Code Injection...

Goals:

- Run code in the context of an arbitrary process:
 - ➤ Force | Explore.exe to send POST
 - ➤ Open remote shell
 - >...
- Easy to write

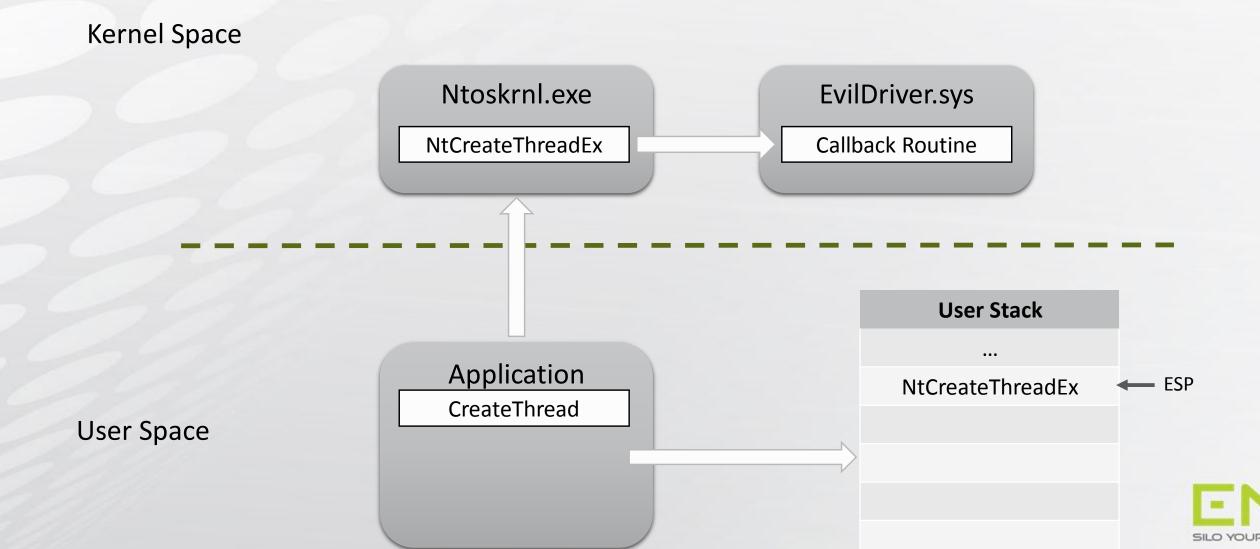
Limitation:

No code injections to the process



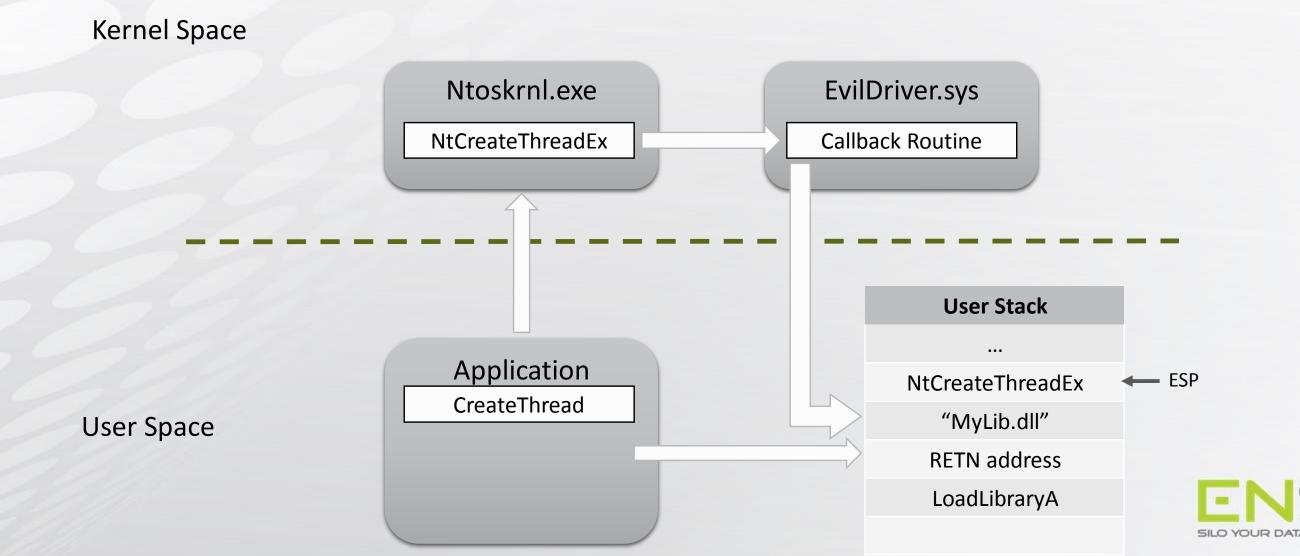
Codeless Code Injection

• Wait for some callback that runs in context of a system call (Thread Creation, Registry Access,...)



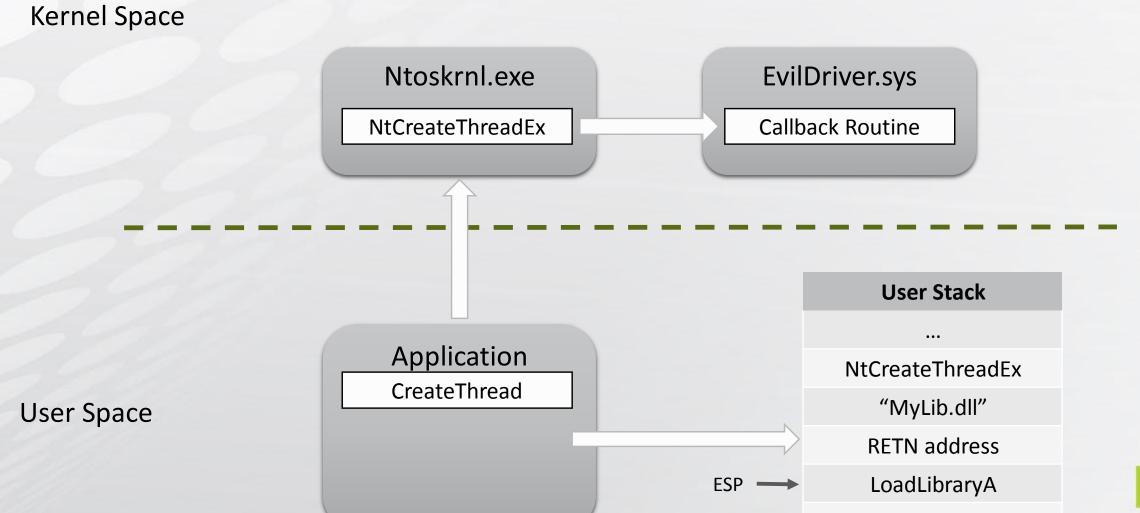
Codeless Code Injection

• Build ROP Chain on the user stack



Codeless Code Injection

Set the stack pointer to the start of the ROP chain





Codeless Code Injection - Challenges

- Getting Return Values how do we get return values back to kernel mode
 - Solution 1: Use device handle to get notifications (deep dive coming up)
 - Solution 2: Save return value somewhere and trigger some hook-able event (Such as registry access)
- Deadlocks if the user mode caller holds a lock that our injected function needs
 - Solution: Use a dedicated thread
- Callback What if we need to a user-mode callback function
 - Solution 1: Use a function that always triggers some hook-able event and fix the context
 - Solution 2: Just don't use such functions ©





Codeless Code Injection - NtClose as a callback

- Create a device using IoCreateDevice
- Whenever we require a callback, create a handle for the device in the target process
- Build the following ROP chain:

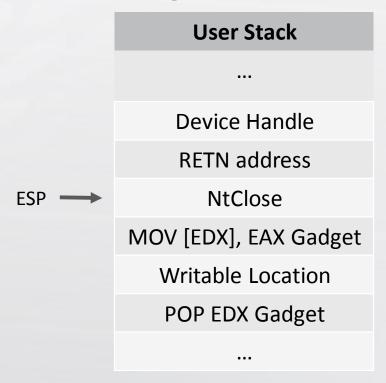
Device Handle RETN address NtClose MOV [EDX], EAX Gadget Writable Location POP EDX Gadget ...

MOV EDX, ADDRESS
MOV [EDX], EAX
NtClose(DeviceHandle)



Codeless Code Injection - NtClose as a callback

- When NtClose(Device Handle) is executed then the IRP_MJ_CLEANUP handler will be called
- The required data will reside in the target address





^{*} moving EAX to some register is not safe because of WOW64

Codeless Code Injection – Creating dedicated thread

```
HANDLE WINAPI CreateThread(
_In_opt_ LPSECURITY_ATTRIBUTES lpThreadAttributes,
In SIZE T
                 dwStackSize,
_In_ LPTHREAD_START_ROUTINE lpStartAddress, 		
                                             NtClose
Device Handle
In DWORD dwCreationFlags,
_Out_opt_ LPDWORD
                      IpThreadId
DWORD WINAPI ThreadProc(_In_ LPVOID lpParameter);
```



Codeless Code Injection – Creating a Dedicated Thread

- Wait for some callback that runs in context of a system call (Thread Creation, Registry Access,...)
- Build the following ROP Chain

ESP

	User Stack
	0
	0
	Device Handle
	NtClose
	0
	NULL
	RETN address
-	CreateThread

```
CreateThread(NULL,

0,

NtClose,

Device Handle,

0,

0);
```

When Handle cleanup is triggered start injecting ROP chains to the new thread

Codeless Code Injection – API

```
ROProgram *CreateProgram(CHAR* TargetName) -
Creates a ROP program that's to be run in a target process
AddStep(ROProgram* Program, ROProgramStepProc StepProc) – Adds a step to the program. A Single ROP chain that ends in callback(NtClose)
RunProgram (ROProgram* Program) -
Wait's for the target process to create a thread and starts a new dedicated thread
UserLibrary OpenUserLibrary (CHAR* LibName) -
Finds a DLL in the target process memory
Call (ROPChain* RopChain, UserLibrary Module, CHAR* Function, u32 NumberOfArgs, ...) -
Add a call to a ROP chain
LoadLibrary (ROPChain* RopChain, char* Name) -
Load a Library into the target process
WriteToUser(ROPChain* RopChain, u8* Buf, u32 BufSize) -
Writes a buffer of data to user-mode
```

Codeless Code Injection – API Example

Load a new library to the target process:

```
ROProgram* Program = CreateProgram("iexplore.exe");
AddStep(Program, LoadSomeLibrary);
RunProgram(Program);
...
bool LoadSomeLibrary(ROProgram* Program, ROPChain* Chain){
    return LoadLibrary(Chain, "MyLibrary.dll");
}
```



Demo

Reverse Shell from WINLOGON



Summary

- Code injection remains an important capability for both malware and security/software vendors
- Like exploits, techniques are becoming less generic (PowerLoader)
- New User and Kernel injection techniques are constantly being invented
- We can probably expect new advanced injection methods in the near future
- Slides and Source-code will be available on BreakingMalware in a few days



Thank You!

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