# A Method to Evaluate CFG Comparison Algorithms

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#### EVALUATION METHOLOGY

- Graph Edit Distance 개념을 기초로 함
  - Add a zero-degree node
  - Delete a zero-degree node
  - Add an edge between two existing nodes
  - Delete an existing edge
- scores 대신에 rankings 을 사용

#### EVALUATION METHOLOGY

- Evaluation Algorithm
  - Algorithm 1은 evaluation process 정보를 줌
  - generate test case
  - compare seed CFG with test case
  - rank the results
  - compute a goodness score

```
procedure EVALUATE(\langle \mathcal{A}_0, \mathcal{A}_1, \dots, \mathcal{A}_k \rangle, G_0, d, n)

for a \leftarrow 1, k do

for t \leftarrow 1, n do

for g \leftarrow 1, d do

scores_g \leftarrow \Delta_{\mathcal{A}_a}(G_0, \operatorname{TestCases}_{t,g})

end for

rank \leftarrow \operatorname{rank}(\operatorname{scores})

\operatorname{corr}_{\mathcal{A}_a, t} \leftarrow \operatorname{pearson}((G_{t,1}, G_{t,2}, \dots, G_{t,d}), \operatorname{rank})

end for

avg_{\operatorname{corr}_{\mathcal{A}_a}} = \operatorname{average}(\operatorname{corr}_{\mathcal{A}_a, 1}, \operatorname{corr}_{\mathcal{A}_a, 2}, \dots, \operatorname{corr}_{\mathcal{A}_a, t})

end for

return rank(avg_{\operatorname{corr}})

end procedure
```

Algorithm 1

#### EVALUATION METHOLOGY

- Benchmark Graph Generation
  - 하나의 EDG (Edit Distance Graph)는 seed CFG로부터 생성될 수 있는 가능한 모든 CFG를 나타냄
    - EDG는 무한하기 때문에 depth d를 설정
  - EDG nodes의 level이 생성됨
    - ex) i edit operations로 만들어진 EDG node는 level i에 위치한다
    - $ED(G_0, G_i) = i$

#### CFG SIMILARITY ALGORITHMS

- A. Based on k-subgraphs Mining
- B. Based on Edit Distance

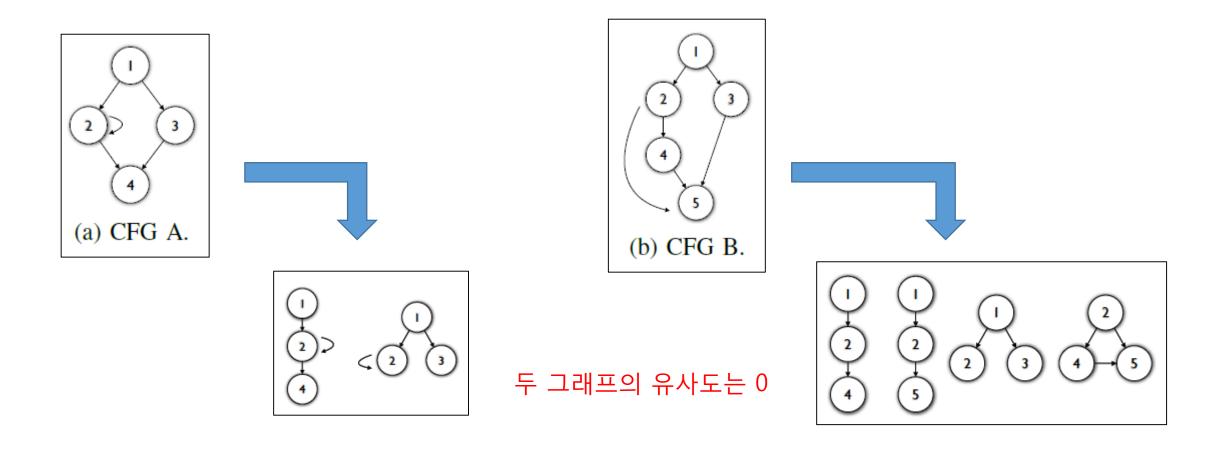
- C. Based on Neighbor Matching
- D. Based on Simulation

### Based on k-subgraphs Mining

- Step1) Generating k-subgraphs
  - maximum outdegree는 2, maximum indegree가 1인 graphs의 set 생성
- Step2) Graph Fingerprinting
  - subgraph들은 숫자로 맵핑
- Step3) Computing Similarity
  - Jaccard index를 사용하여 비교

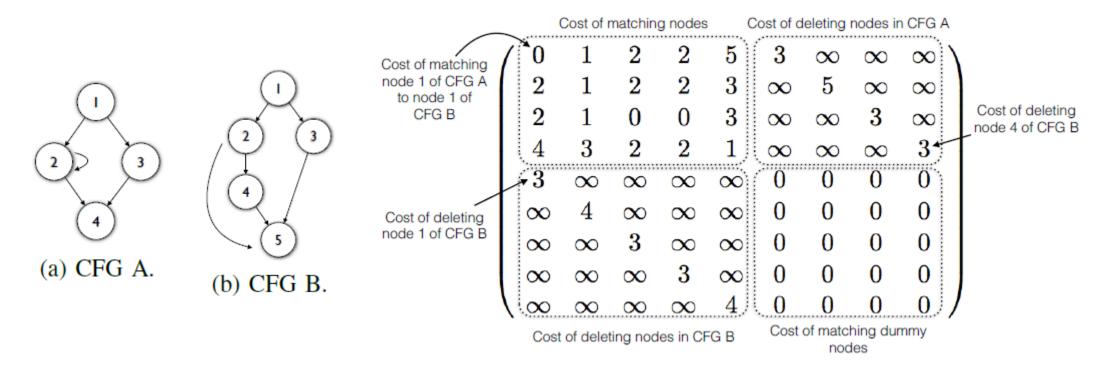
```
\frac{|\operatorname{fingerprint}(A) \cap \operatorname{fingerprint}(B)|}{|\operatorname{fingerprint}(A) \cup \operatorname{fingerprint}(B)|}
\operatorname{Jaccard index}
```

# Based on k-subgraphs Mining



#### Based on Edit Distance

- Step1) Building The Cost Matrix
  - G<sub>1</sub>, G<sub>2</sub>의 Vertex 수 V<sub>1</sub>, V<sub>2</sub>를 가지고 (|V1|+|V2|)\*(|V1|+|V2|) matrix 생성



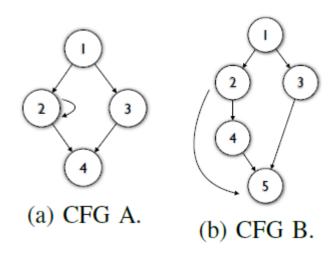
#### Based on Edit Distance

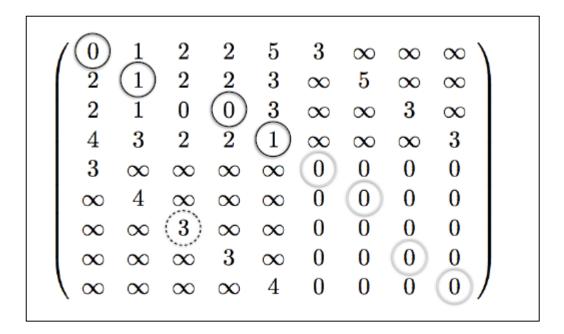
- Step2) Finding Best Match of the Nodes
  - 가장 낮은 코스트를 갖는 매칭 비용을 찾는 것이 목적
  - Hungarian algorithm이 assignment problem의 optimal solution을 찾음
- Step3) Computing Similarity

$$sim(G_1, G_2) = 1 - \frac{edit \ distance}{|V_1| + |E_1| + |V_2| + |E_2|}$$

#### Based on Edit Distance

- Example
  - Hungarian algorithm을 이용하여 best matching 비용을 구함





## Based on Neighbor Matching

• neighbor의 similarity에 근거하여 similarity matrix를 빌드함

- Step1) Topological Similarity
  - k=0,  $a_{ij}^0=1$
- Step2) Node Similarity

$$a_{ij}^{k+1} = \sqrt{y_{ij} \cdot \frac{s_{\text{in}}^{k+1}(i,j) + s_{\text{out}}^{k+1}(i,j)}{2}}$$

$$a_{ij}^{k+1} = \frac{s_{\text{in}}^{k+1}(i,j) + s_{\text{out}}^{k+1}(i,j)}{2}$$
where
$$s_{\text{in}}^{k+1}(i,j) = \frac{1}{m_{\text{in}}} \sum_{l=1}^{n_{\text{in}}} a_{f_{ij}}^{k}(l) g_{ij}^{\text{in}}(l)$$

$$s_{\text{out}}^{k+1}(i,j) = \frac{1}{m_{\text{out}}} \sum_{l=1}^{n_{\text{out}}} a_{f_{ij}^{\text{out}}(l) g_{ij}^{\text{out}}(l)}^{k}$$

$$m_{\text{in}} = \max(|\text{IE}_i|, |\text{IE}_j|)$$

$$n_{\text{in}} = \min(|\text{IE}_i|, |\text{IE}_j|)$$

$$m_{\text{out}} = \max(|\text{OE}_i|, |\text{OE}_j|)$$

$$n_{\text{out}} = \min(|\text{OE}_i|, |\text{OE}_j|)$$

#### Based on Simulation

- Labeled Transition Systems (LTS)를 사용하여 CFG를 만듦
- Recursively match the most similar outgoing nodes starting from the entry nodes
- Sum up the similarity of the matched nodes and edges

#### Based on Simulation

• entry n1, n2를 갖는 두 CFG의 유사도

$$S(n_{1}, n_{2}) = \begin{cases} N(n_{1}, n_{2}) & \text{, if } OD_{n_{1}} = 0 \\ (1 - p) \cdot N(n_{1}, n_{2}) + \\ p \cdot max(W_{1}, W_{2}) & \text{, otherwise,} \end{cases}$$
where
$$p \in (0, 1)$$

$$W_{1} = \max_{n_{2}} \sum_{h n'_{2}} L(b, \epsilon) \cdot S(n_{1}, n'_{2})$$

$$W_{2} = \frac{1}{OD_{n_{1}}} \cdot \sum_{n_{1}} \max(\max_{n_{2}} \sum_{h n'_{2}} (L(a, \epsilon) \cdot S(n'_{1}, n'_{2}))$$

$$b) \cdot S(n'_{1}, n'_{2}), L(a, \epsilon) \cdot S(n'_{1}, n_{2}))$$