Lecture 6 Word vectors Recurrent Neural Networks

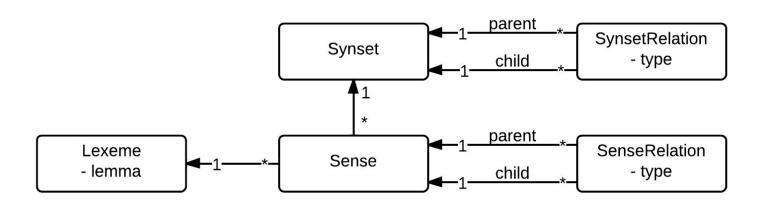
Vladislav Goncharenko

Outline

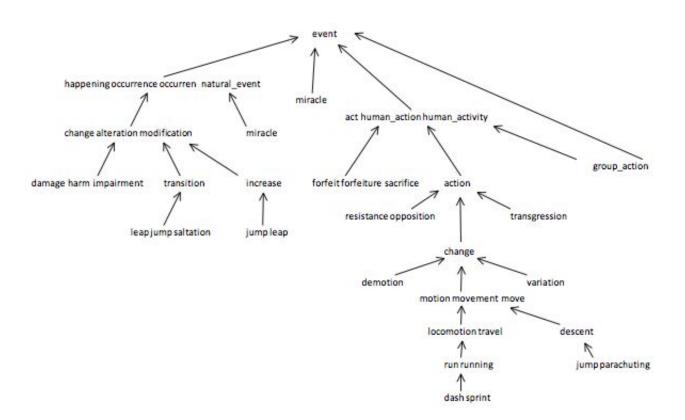
- 1. Discrete representations.
- 2. Matrix of co-occurrence.
- 3. Embeddings (GloVe, word2vec).
- 4. Examples.
- Recurrent Neural Networks

How to represent text in a computer?

Use a taxonomy like WordNet that has hypernyms (is-a) relationships and synonym sets



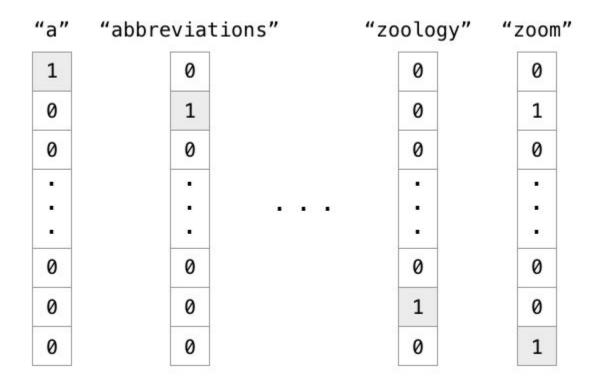
How to represent text in a computer: WordNet



Discrete representations: problems

- Missing new words
- Subjective
- Requires human labor to create and adapt
- Hard to compute accurate word similarity

Discrete representations: one-hot encoding



TF-IDF

$$\mathbf{tf}(t,d) = \frac{f_d(t)}{\max_{w \in d} f_d(w)}$$

$$\mathbf{idf}(t, D) = \ln\left(\frac{|D|}{|\{d \in D : t \in d\}|}\right)$$

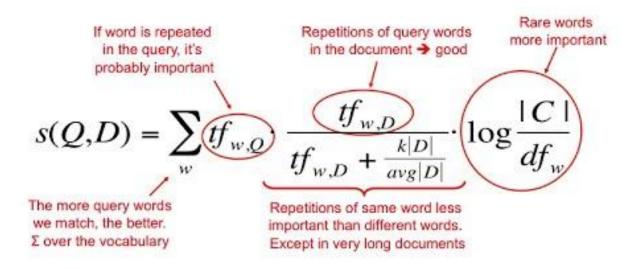
$$\mathbf{tfidf}(t, d, D) = \mathbf{tf}(t, d) \cdot \mathbf{idf}(t, D)$$

$$\mathbf{tfidf}'(t,d,D) = \frac{\mathbf{idf}(t,D)}{|D|} + \mathbf{tfidf}(t,d,D)$$

 $f_d(t) :=$ frequency of term t in document d

D := corpus of documents

TF-IDF



TF - term frequency

IDF - Inverse Document Frequency

TF-IDF: make it simple

$$ext{tf("this",}\ d_1)=rac{1}{5}=0.2$$
 $ext{tf("this",}\ d_2)=rac{1}{7}pprox 0.14$ $ext{idf("this",}\ D)=\log\Bigl(rac{2}{2}\Bigr)=0$

 $\operatorname{tfidf}("\mathsf{this}", d_1, D) = 0.2 \times 0 = 0$

 $\operatorname{tfidf}("\mathsf{this}", d_2, D) = 0.14 \times 0 = 0$

-	

Document 1

Term	Term Count			
this	1			
is	1			
a	2			
sample	1			

Document 2

Term	Term Count				
this	1				
is	1				
another	2				
example	3				

Word 'this' is not very informative

Words cooccurrences

One of the most successful ideas of statistical NLP:

"You shall know a word by the company it keeps"

(J. R. Firth 1957: 11)

Words cooccurrences





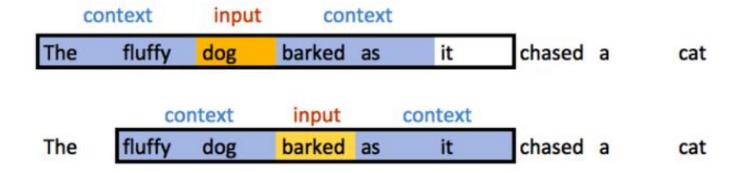
Word-document cooccurrence matrix

Window around each word

Word-document cooccurrence matrix

		I	like	enjoy	deep	learning	NLP	flying	•
X =	I	0	2	1	0	0	0	0	0]
	like	2	0	0	1	0	1	0	0
	enjoy	1	0	0	0	0	0	1	0
	deep	0	1	0	0	1	0	0	0
	learning	0	0	0	1	0	0	0	1
	NLP	0	1	0	0	0	0	0	1
	flying	0	0	1	0	0	0	0	1
	1.	0	0	0	0	1	1	1	0

Words cooccurrences: sliding window



Words cooccurrences: n-grams

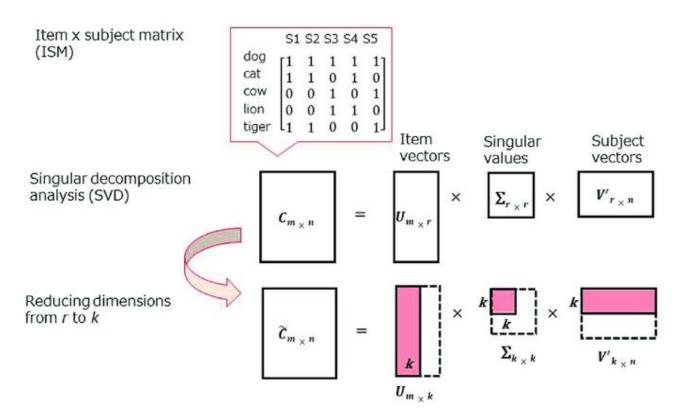
Co-occurrence vectors: problems

- Increase in size with vocabulary
- Very high dimensional: require a lot of storage
- Subsequent classification models have sparsity issues

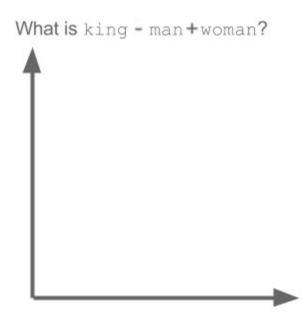


Models are less robust

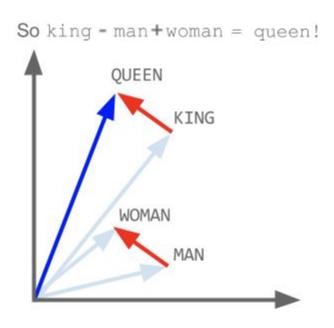
Reducing dimensionality: SVD of cooccurrence matrix

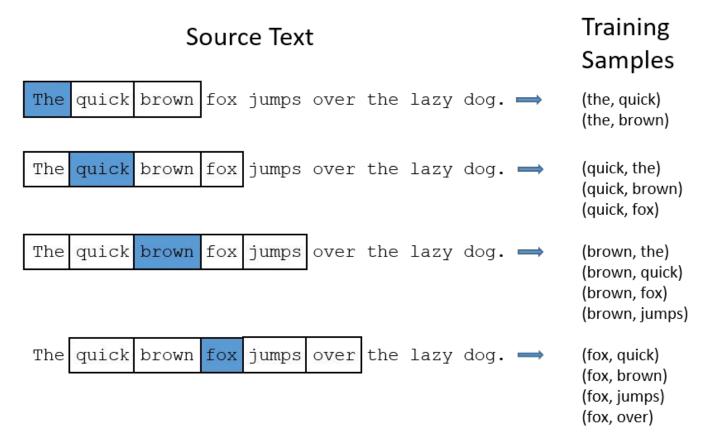


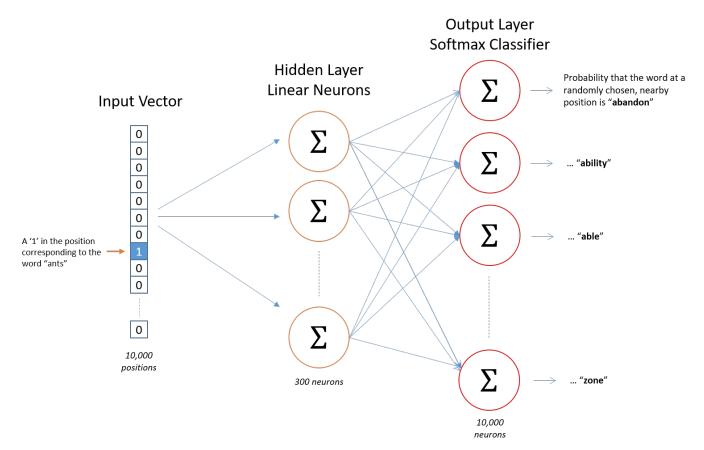
Embeddings: intuition

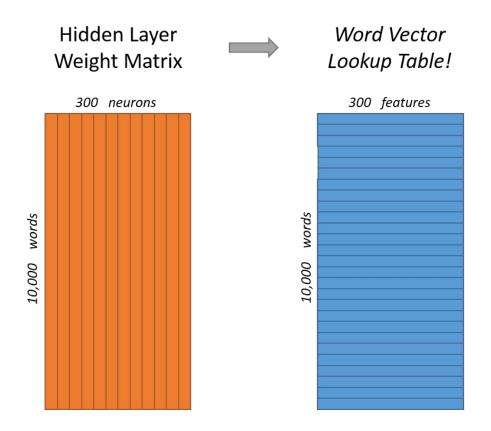


Embeddings: intuition









- Word vectors with 300 components
- Vocabulary of 10,000 words.
- Weight matrix with 300 x 10,000 = 3 million weights each!

Training is too long and computationally expensive

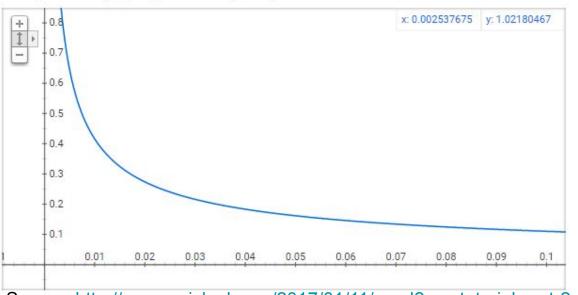
How to fix this?

Basic approaches:

- 1. Treating common word pairs or phrases as single "words" in their model.
- 2. Subsampling frequent words to decrease the number of training examples.
- Modifying the optimization objective with a technique they called "Negative Sampling", which causes each training sample to update only a small percentage of the model's weights.

Subsampling frequent words.

 w_i is the word, $z(w_i)$ is the fraction of this word in the whole Graph for $(\sqrt{(x/0.001)+1})*0.001/x$



 $P(w_i)$ is the probability of *keeping* the word:

$$P(w_i) = (\sqrt{\frac{z(w_i)}{0.001}} + 1) \cdot \frac{0.001}{z(w_i)}$$

Source: http://mccormickml.com/2017/01/11/word2vec-tutorial-part-2-negative-sampling/

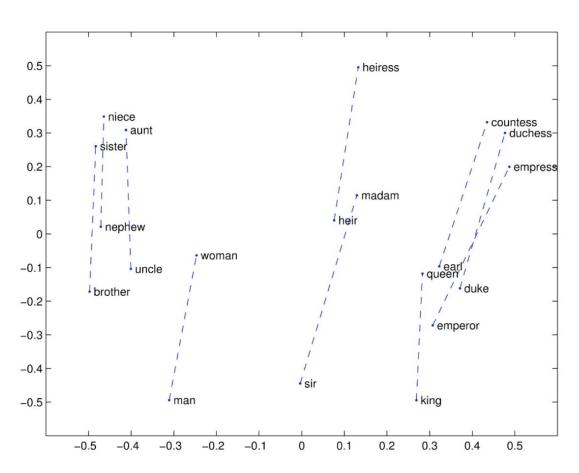
Embeddings: negative sampling

Negative Sampling idea: only few words error is computed. All other words has zero error, so no updates by the backprop mechanism.

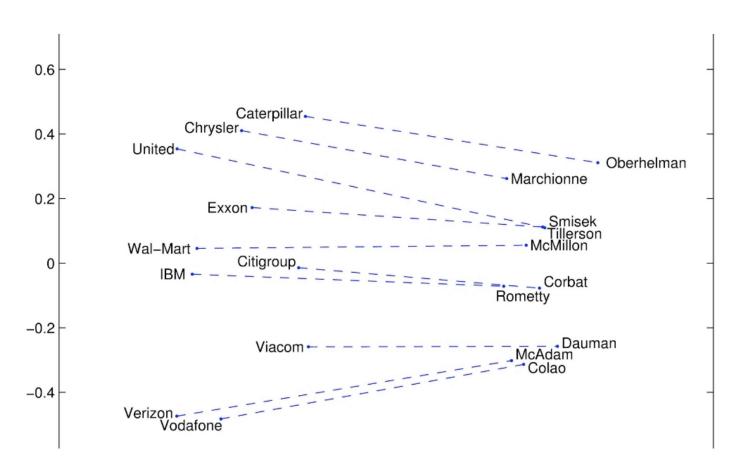
More frequent words are selected to be negative samples more often. The probability for a selecting a word is just it's weight divided by the sum of weights for all words.

$$P(w_i) = \frac{f(w_i)^{3/4}}{\sum_{i=0}^{n} (f(w_i)^{3/4})}$$

GloVe Visualizations



GloVe Visualizations: Company - CEO



Conclusion

Word vectors are simply vectors of numbers that represent the meaning of a word

Approaches:

- One-hot encoding
- Bag-of-words models
- Counts of word / context co-occurrences
- TF-IDF
- Predictions of context given word (skip-gram neural network models, e.g. word2vec)

RNNs generating...

Shakespeare

PANDARUS:

Alas, I think he shall be come approached and the day When little srain would be attain'd into being never fed, And who is but a chain and subjects of his death, I should not sleep.

Second Senator:

They are away this miseries, produced upon my soul, Breaking and strongly should be buried, when I perish The earth and thoughts of many states.

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Clown:

Come, sir, I will make did behold your worship.

VIOLA:

I'll drink it.

Algebraic Geometry (Latex)

```
Proof. Omitted.
Lemma 0.1. Let C be a set of the construction.
   Let C be a gerber covering. Let F be a quasi-coherent sheaves of O-modules. We
have to show that
                                     \mathcal{O}_{\mathcal{O}_{+}} = \mathcal{O}_{X}(\mathcal{L})
Proof. This is an algebraic space with the composition of sheaves \mathcal{F} on X_{Oute} we
                           \mathcal{O}_X(\mathcal{F}) = \{morph_1 \times_{\mathcal{O}_X} (\mathcal{G}, \mathcal{F})\}
where G defines an isomorphism F \to F of O-modules.
Lemma 0.2. This is an integer Z is injective.
Proof. See Spaces, Lemma ??.
Lemma 0.3. Let S be a scheme. Let X be a scheme and X is an affine open
covering. Let U \subset X be a canonical and locally of finite type. Let X be a scheme.
Let X be a scheme which is equal to the formal complex.
The following to the construction of the lemma follows.
Let X be a scheme. Let X be a scheme covering. Let
                      b: X \rightarrow Y' \rightarrow Y \rightarrow Y \rightarrow Y' \times_X Y \rightarrow X.
be a morphism of algebraic spaces over S and Y.
Proof. Let X be a nonzero scheme of X. Let X be an algebraic space. Let F be a
quasi-coherent sheaf of O_X-modules. The following are equivalent

 F is an algebraic space over S.

    (2) If X is an affine open covering.
Consider a common structure on X and X the functor O_X(U) which is locally of
finite type.
```

Linux kernel (source code)

```
* If this error is set, we will need anything right after that BSD.
static void action new function(struct s stat info *wb)
 unsigned long flags;
 int lel idx bit = e->edd, *sys & -((unsigned long) *FIRST COMPAT);
 buf[0] = 0xFFFFFFFF & (bit << 4);
 min(inc, slist->bytes);
 printk(KERN WARNING "Memory allocated %02x/%02x, "
   "original MLL instead\n"),
   min(min(multi run - s->len, max) * num data in),
   frame pos, sz + first seg);
 div u64 w(val, inb p);
 spin unlock(&disk->queue lock);
 mutex unlock(&s->sock->mutex);
 mutex unlock(&func->mutex);
 return disassemble(info->pending bh);
```

Proof. Omitted.

Lemma 0.1. Let C be a set of the construction.

Let C be a gerber covering. Let F be a quasi-coherent sheaves of O-modules. We have to show that

$$\mathcal{O}_{\mathcal{O}_X} = \mathcal{O}_X(\mathcal{L})$$

.

Proof. This is an algebraic space with the composition of sheaves F on $X_{\acute{e}tale}$ we have

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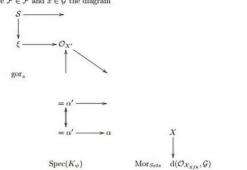
be a morphism of algebraic spaces over S and Y.

Proof. Let X be a nonzero scheme of X. Let X be an algebraic space. Let \mathcal{F} be a quasi-coherent sheaf of \mathcal{O}_X -modules. The following are equivalent

- F is an algebraic space over S.
- (2) If X is an affine open covering.

Consider a common structure on X and X the functor $\mathcal{O}_X(U)$ which is locally of finite type.

This since $F \in F$ and $x \in G$ the diagram



is a limit. Then G is a finite type and assume S is a flat and F and G is a finite type f_* . This is of finite type diagrams, and

- the composition of G is a regular sequence,
- O_{X'} is a sheaf of rings.

Proof. We have see that $X = \operatorname{Spec}(R)$ and \mathcal{F} is a finite type representable by algebraic space. The property \mathcal{F} is a finite morphism of algebraic stacks. Then the cohomology of X is an open neighbourhood of U.

Proof. This is clear that G is a finite presentation, see Lemmas ??. A reduced above we conclude that U is an open covering of C. The functor F is a "field

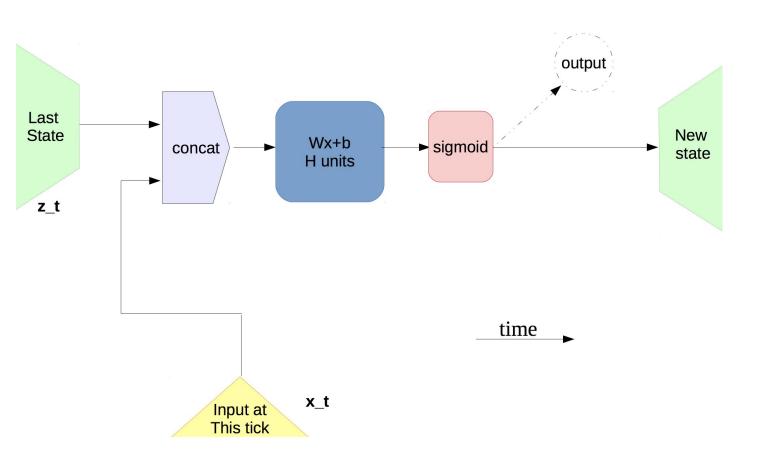
$$\mathcal{O}_{X,x} \longrightarrow \mathcal{F}_{\overline{x}} -1(\mathcal{O}_{X_{\ell tale}}) \longrightarrow \mathcal{O}_{X_{\ell}}^{-1}\mathcal{O}_{X_{\lambda}}(\mathcal{O}_{X_{\eta}}^{\overline{v}})$$

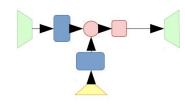
is an isomorphism of covering of \mathcal{O}_{X_i} . If \mathcal{F} is the unique element of \mathcal{F} such that X is an isomorphism.

The property \mathcal{F} is a disjoint union of Proposition ?? and we can filtered set of presentations of a scheme \mathcal{O}_{X} -algebra with \mathcal{F} are opens of finite type over S. If \mathcal{F} is a scheme theoretic image points.

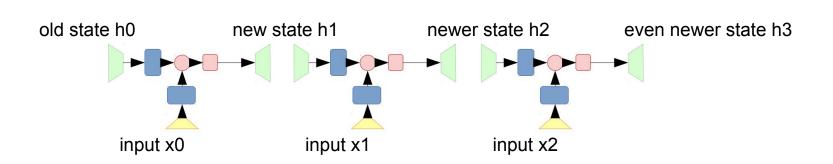
If \mathcal{F} is a finite direct sum $\mathcal{O}_{X_{\lambda}}$ is a closed immersion, see Lemma ??. This is a sequence of \mathcal{F} is a similar morphism.

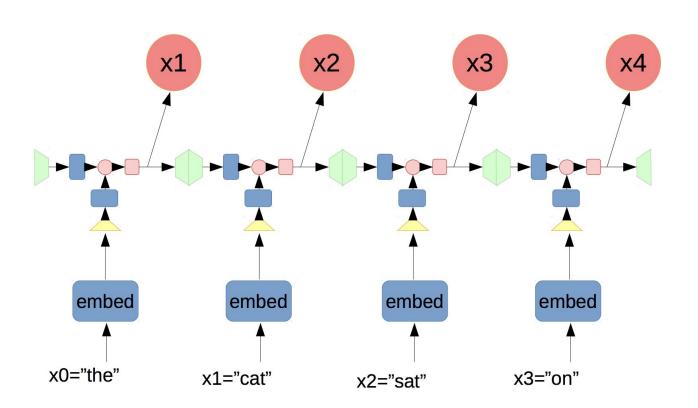
```
#include <asm/io.h>
#include <asm/prom.h>
#include <asm/e820.h>
#include <asm/system info.h>
#include <asm/setew.h>
#include <asm/pgproto.h>
#define REG PG vesa slot addr pack
#define PFM NOCOMP AFSR(0, load)
#define STACK DDR(type) (func)
#define SWAP ALLOCATE(nr)
                            (e)
#define emulate sigs() arch get unaligned child()
#define access rw(TST) asm volatile("movd %%esp, %0, %3" :: "r" (0)); \
 if ( type & DO READ)
static void stat PC SEC read mostly offsetof(struct seq argsqueue, \
         pC>[1]);
static void
os prefix(unsigned long sys)
#ifdef CONFIG PREEMPT
 PUT PARAM RAID(2, sel) = get state state();
  set_pid_sum((unsigned long)state, current_state_str(),
          (unsigned long)-1->lr full; low;
```





We use same weight matrices for all steps





Now with formulas

$$h_{0} = \bar{0}$$

$$h_{1} = \sigma(\langle W_{\text{hid}}[h_{0}, x_{0}] \rangle + b)$$

$$h_{2} = \sigma(\langle W_{\text{hid}}[h_{1}, x_{1}] \rangle + b) = \sigma(\langle W_{\text{hid}}[\sigma(\langle W_{\text{hid}}[h_{0}, x_{0}] \rangle + b, x_{1}] \rangle + b)$$

$$h_{i+1} = \sigma(\langle W_{\text{hid}}[h_{i}, x_{i}] \rangle + b)$$

$$P(x_{i+1}) = \operatorname{softmax}(\langle W_{\text{out}}, h_{i} \rangle + b_{\text{out}})$$

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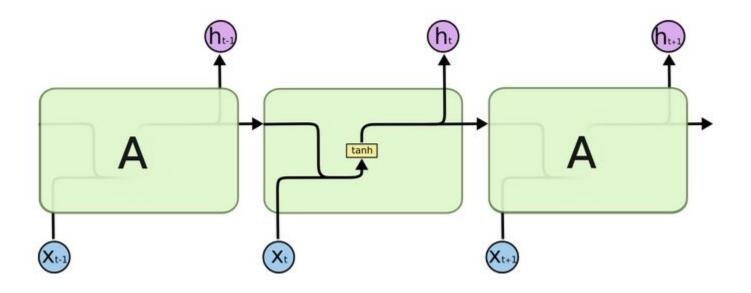
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 min(inc, slist->bytes);
 printk(KERN WARNING "Memory allocated %02x/%02x, "
   "original MLL instead\n"),
   min(min(multi run - s->len, max) * num data in),
   frame pos, sz + first seg);
 div u64 w(val, inb p);
 spin unlock(&disk->queue lock);
 mutex unlock(&s->sock->mutex);
 mutex unlock(&func->mutex);
 return disassemble(info->pending bh);
```

Vanilla RNN



That's all. Feel free to ask any questions.

Remember: a well-defined problem is halfway to being solved