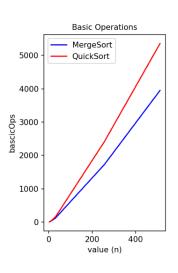
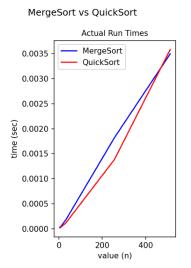
#### Project 1 Divide and Conquer Writeup

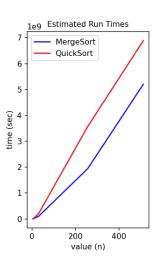
Skyler Ballard and Michael Shipley

#### The Experiment:

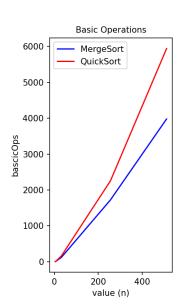
Run 1: this is an example where mergesort was faster than quicksort when n=512, this happens less than 5% of the time

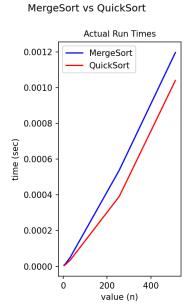


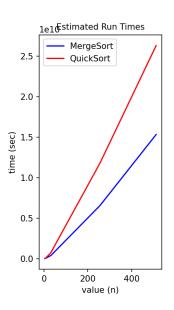




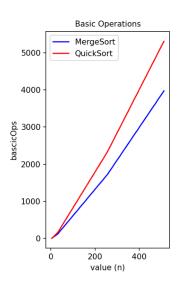
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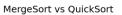


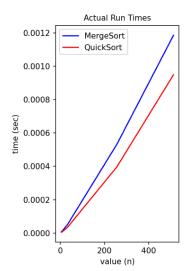


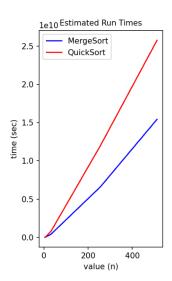


Run 3:









#### The Analysis:

This result surprised me because: even though quicksort did more comparisons (basicOps) and the time estimations for quicksort predicted that it would take longer, it was in fact faster than mergesort. Mergesort was actual slower with less comparisons, and having the better estimated runtime(s). I did run the test over many times and did find every once in a while mergesort would

win a single round during the actual runtimes, but on average 90% + quicksort would win every time.

Regarding the research conducted on QuickSort vs MergeSort, Skyler and Michael have found the following observations regarding the two algorithms aforementioned.

QuickSort is generally faster than MergeSort. In practice, Quicksort performs less allocation than MergeSort, a result of Quicksort being in place. Quicksort also only performs faster when the data is truly random, more on this later.

Mergesort generally has fewer operations than QuickSort. This is due to Mergesort consistently sorting. Quicksort may move a piece of data more than twice. Mergesort only moves a piece of data's order once to Merge. Less basic operations did not translate to a faster algorithm.

Our estimated times were often much lower and inverted to actual times when presented with datasets not of 2^n size. When presented with 2^n size datasets the actual time was magnitudes lower than estimated times. This is likely due to the algorithms being drastically more efficient at datasets in 2^n restriction. Merge sort splits in half and Quicksort compares two at a time, making both naturally more inclined to even data sets. This is shown in our collected data that it does take more time on odd datasets than even datasets. However it seems strange that the difference is in a matter of x1000 and not a few nanoseconds.

Our initial estimates were based on datasets that were multiples of 5. I attempted to reduce the inaccuracy by trying data sets that were of powers of two. The results were better than I hoped to reduce the time, but made our estimates largely inaccurate! The error with datasets as multiples of 5 were off by a few seconds at most, but in sets of powers of 2 the estimate was off by a magnitude of 1000. This is because sets of non powers of 2 are average case, and sets of powers of 2 are best case, for both algorithms. In practice it would be advantageous to fill sets to powers

of two for sorting purposes. The deltas however are fairly similar to the actual result. This is

shown by our side by side data of various tests. The powerset of 2<sup>n</sup> is remarkably more accurate

to the individual algorithms, but not wholly so as it predicts incorrectly which will be more

efficient. With a notable expectation discussed after the proof.

Worst Case QuickSort

Proof by induction: Using the Worst-Case Time Complexity of Quicksort

• Using a list/array of size n

• n is a power of 2  $(n=2^k)$ 

We have a recurrence relation:

• W(1) = 0

• W(n) = W(n-1) + n - 1

Hypothesis: 
$$T(n) = (n * n - n) / 2$$

Base: 
$$T(1) = (1*1 - 1) / 2 = 0/2 = 0$$
 which is true

If it is true for n, then it is also true for (n + 1), substitute (n + 1) for n

Step: 
$$T(n+1) = (((n+1) * (n+1)) - (n+1)) / 2$$

Left Side:

$$=T(n+1)$$

= 
$$T(n) + n$$
 And  $b/c$   $T(n) = (n * n - n) / 2$  We sub it into the equation

$$= ((n*n - n) / 2) + n$$

$$= (n*n - n + 2n) / 2$$

$$= (n^2 + 2n - n) / 2$$

$$= (n^2 + n) / 2$$

$$= (n(n+1))/2$$

#### Right Side:

$$=(((n+1)*(n+1))-(n+1))/2$$

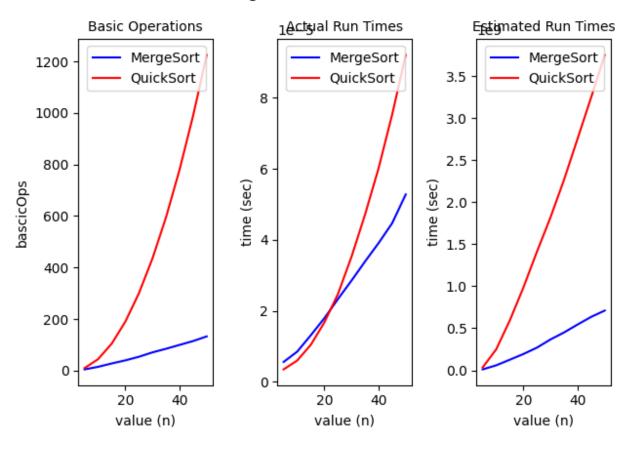
$$= ((n^2 + 2n + 1) - (n + 1)) / 2$$

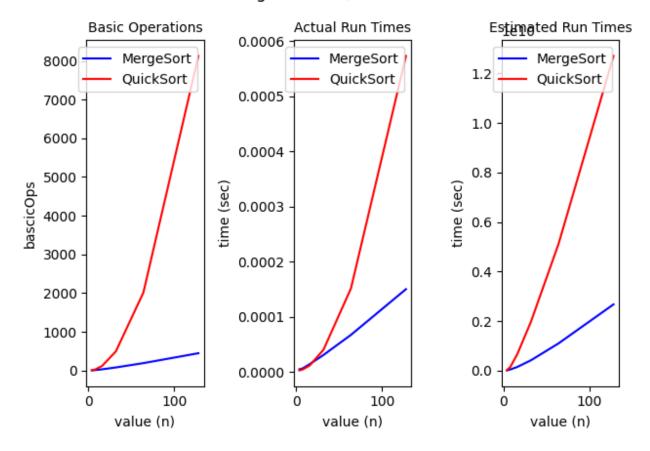
$$= (n^2 + 2n + 1 - n - 1) / 2$$

$$= (n^2 + n) / 2$$

$$= (n(n+1))/2$$

We see that they are equal, therefore the recurrence equals the candidate solution





A note regarding the worst case of QuickSort: Quicksort has a fatal flaw. When the data is already sorted, or well ordered, it performs far worse, as estimated, than MergeSort. Also, in a later project we will make multithreaded MergeSort, which performs significantly better on its time quantum than QuickSort in large sets. Look forward to that data in those projects.

To test the well sorted theory, I implemented a random set that chose its element based on its index, meaning it was always sorted. This showed Mergesort to always perform better than Quicksort except in small datasets around size 4 to 32. This where mergesort still has an efficiency of nlogn and quicksort has an efficiency of n^2. This isn't an uncommon dataset either. Partially ordered sets fail to be fast on Quicksort. This should be considered as it can affect usage uptime of MergeSort, which does no worse on a partially ordered set. In the long run

MergeSort should be used on kept data and Quicksort on new data. This is a lesson in knowing when to use an algorithm. It however does require knowledge regarding the set.

Another note, I did extensive testing and will include the folder of the results of each of the tests. The names on the tests are useful, which is why it had to be in a separate document.

