

Problem	Recurrence Formula
<b>Fibonacci Numbers</b>	$dp[i] = dp[i-1] + dp[i-2]$
<b>Climbing Stairs</b>	$dp[i] = dp[i-1] + dp[i-2]$
<b>Min Cost Climbing Stairs</b>	$dp[i] = cost[i] + \min(dp[i-1], dp[i-2])$
<b>House Robber</b>	$dp[i] = \max(dp[i-1], dp[i-2] + nums[i])$
<b>Decode Ways</b>	$dp[i] = dp[i-1] + dp[i-2]$ (if valid)
<b>Tiling Problem</b>	$dp[i] = dp[i-1] + dp[i-2]$
<b>Generalized k-step</b>	$dp[i] = dp[i-1] + dp[i-2] + \dots + dp[i-k]$

Cheat Sheet Table : Fibonacci Family Tree (1D DP Problems)

#### That "1D DP" Means

- The DP state has **one dimension** (like  $dp[i]$ ).
- Each state depends only on previous states in that same dimension.

$dp[0] = 0$

$dp[1] = 1$

for  $i$  in range(2,  $n+1$ ):

$dp[i] = dp[i-1] + dp[i-2]$

Cheat Sheet Table : Knapsack Variant

#### ◆ Base Problem: 0/1 Knapsack

- **Given:**  $n$  items with  $value[i]$ ,  $weight[i]$  and capacity  $W$ .
- **Goal:** Maximize total value without exceeding  $W$ .
- **Choice:** For each item  $i$ , take it (if  $weight \leq capacity$ ) or skip it.

$dp[i][w] = \max($

$dp[i-1][w],$  # skip item  $i$

$value[i-1] + dp[i-1][w - weight[i-1]]$  # take item  $i$

)

Variant	State	Recurrence
<b>0/1 Knapsack</b>	$dp[i][w]$	$\max(dp[i-1][w], val + dp[i-1][w - wt])$
<b>Subset Sum</b>	$dp[i][s]$ (True/False)	$dp[i-1][s]$ or $dp[i-1][s - arr[i]]$

<b>Partition Equal Subset Sum</b>	Same as Subset Sum	Check sum/2
<b>Count of Subsets with Given Sum</b>	Count ways	$dp[i-1][s] + dp[i-1][s-arr[i]]$
<b>Min Subset Sum Diff</b>	Based on achievable sums	Pick closest to sum/2
<b>Unbounded Knapsack</b>	Unlimited copies	$\max(dp[i-1][w], val + dp[i][w-wt])$
<b>Rod Cutting</b>	Like Unbounded Knapsack	Max profit by lengths
<b>Coin Change</b>	Ways/Min coins (Unbounded)	Ways: $dp[i][amt] = dp[i-1][amt] + dp[i][amt-coin[i]]$
<b>Bounded Knapsack</b>	Limited copies	Convert to multiple 0/1
<b>Multi-Dimensional Knapsack</b>	Multi-constraints	$dp[i][w1][w2]$ etc.
<b>Fractional Knapsack</b>	Greedy	Sort by value/weight

### Grid / Matrix DP Problem Family

The **base structure**:

You have a **2D grid/matrix**.

Your state =  $dp[i][j]$  (usually row, col).

You fill the grid based on **neighbors** or **previous rows**.

Problem	State dp	Recurrence
<b>Unique Paths</b>	Ways to reach (i,j)	$dp[i][j] = dp[i-1][j] + dp[i][j-1]$
<b>Min Path Sum</b>	Min cost to (i,j)	$grid[i][j] + \min(dp[i-1][j], dp[i][j-1])$
<b>Falling Path Sum</b>	Min cost from top to bottom	$grid[i][j] + \min(dp[i-1][j-1], dp[i-1][j], dp[i-1][j+1])$
<b>Triangle Min Path Sum</b>	Min cost in triangle	$triangle[i][j] + \min(dp[i+1][j], dp[i+1][j+1])$
<b>Maximal Square</b>	Largest square of 1's	$1 + \min(dp[i-1][j], dp[i][j-1], dp[i-1][j-1])$
<b>LCS</b>	LCS length up to i,j	match: $1 + dp[i-1][j-1]$ else $\max(dp[i-1][j], dp[i][j-1])$
<b>Edit Distance</b>	Min edits up to i,j	if match: $dp[i-1][j-1]$ else $1 + \min(dp[i-1][j], dp[i][j-1], dp[i-1][j-1])$
<b>Cherry Pickup</b>	Two travelers DP	$dp[x1][y1][x2]$

Problem	State dp	Recurrence
<b>Unique Paths</b>	Ways to reach (i,j)	$dp[i][j] = dp[i-1][j] + dp[i][j-1]$
<b>Min Path Sum</b>	Min cost to (i,j)	$grid[i][j] + \min(dp[i-1][j], dp[i][j-1])$

<b>Falling Path Sum</b>	Min cost from top to bottom	$\text{grid}[i][j] + \min(\text{dp}[i-1][j-1], \text{dp}[i-1][j], \text{dp}[i-1][j+1])$
<b>Triangle Min Path Sum</b>	Min cost in triangle	$\text{triangle}[i][j] + \min(\text{dp}[i+1][j], \text{dp}[i+1][j+1])$
<b>Maximal Square</b>	Largest square of 1's	$1 + \min(\text{dp}[i-1][j], \text{dp}[i][j-1], \text{dp}[i-1][j-1])$
<b>LCS</b>	LCS length up to i,j	match: $1 + \text{dp}[i-1][j-1]$ else $\max(\text{dp}[i-1][j], \text{dp}[i][j-1])$
<b>Edit Distance</b>	Min edits up to i,j	if match: $\text{dp}[i-1][j-1]$ else $1 + \min(\text{dp}[i-1][j], \text{dp}[i][j-1], \text{dp}[i-1][j-1])$
<b>Cherry Pickup</b>	Two travelers DP	$\text{dp}[x1][y1][x2]$

## Strings DP

Problem	State & Recurrence
<b>LCS</b>	if match: $1 + \text{dp}[i-1][j-1]$ else $\max(\text{dp}[i-1][j], \text{dp}[i][j-1])$
<b>Longest Common Substring</b>	if match: $1 + \text{dp}[i-1][j-1]$ else 0
<b>Edit Distance</b>	if match: $\text{dp}[i-1][j-1]$ else $1 + \min(\text{dp}[i-1][j], \text{dp}[i][j-1], \text{dp}[i-1][j-1])$
<b>Longest Palindromic Subsequence</b>	if $s[i] == s[j]$ : $2 + \text{dp}[i+1][j-1]$ else $\max(\text{dp}[i+1][j], \text{dp}[i][j-1])$
<b>Palindrome Partitioning</b>	$\text{dp}[i] = \min(\text{dp}[j-1] + 1)$ if $s[j..i]$ palindrome
<b>Decode Ways</b>	$\text{dp}[i] = \text{dp}[i-1] + \text{dp}[i-2]$ if valid
<b>Distinct Subsequences</b>	if match: $\text{dp}[i-1][j-1] + \text{dp}[i-1][j]$ else $\text{dp}[i-1][j]$
<b>Interleaving String</b>	$\text{dp}[i][j] = (\text{dp}[i-1][j] \text{ or } \text{dp}[i][j-1])$ if chars match
<b>Regex Matching</b>	Handles . and *
<b>Wildcard Matching</b>	Handles ? and *
Pattern	Typical DP Table Shape
<b>Two strings (compare)</b>	$\text{dp}[i][j]$ from diagonal/top/left
<b>One string (palindrome)</b>	$\text{dp}[i][j]$ on substrings (i..j)
<b>One string (count ways)</b>	$\text{dp}[i]$ for prefix length
<b>Pattern matching</b>	$\text{dp}[i][j]$ for string vs pattern