ITP 30002 Operating System

Conditional Variable & Semaphore

OSTEP Chapters 30 and 31

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Motivation

 There are many times where a thread wishes to check whether a condition is satisfied before continuing its execution

Example: thread join

```
volatile int done = 0;
   void *child(void *arg) {
       printf("child\n");
       done = 1;
       return NULL;
7
8
   int main(int argc, char *argv[]) {
9
       printf("parent: begin\n");
10
       pthread_t c;
11
       Pthread_create(&c, NULL, child, NULL); // create child
12
       while (done == 0)
13
            ; // spin
14
       printf("parent: end\n");
15
       return 0;
16
17
```

Conditional Variables & Semaphores

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Conditional Variable

- A conditional variable is an explicit queue that threads can put themselves on when some state of execution is not desired
 - -when it said the condition might be changed, a waiting thread is awken up to continue its execution
 - -operation
 - wait () pthread_cond_wait (pthread_cond_t *c, pthread_mutex_t *m);
 signal () pthread_cond_signal (pthread_cond_t *c);
- a conditional variable is used together with a mutex for ensuring mutual exclusion of the condition checking

4

Example

```
int done = 0;
   pthread_mutex_t m = PTHREAD_MUTEX_INITIALIZER;
   pthread_cond_t c = PTHREAD_COND_INITIALIZER;
4
   void thr exit() {
       Pthread_mutex_lock(&m);
       done = 1;
7
       Pthread_cond_signal(&c);
       Pthread_mutex_unlock(&m);
9
10
11
   void *child(void *arg) {
12
       printf("child\n");
13
       thr_exit();
14
       return NULL;
15
16
17
   void thr_join() {
18
       Pthread mutex lock(&m);
19
       while (done == 0)
20
            Pthread_cond_wait(&c, &m);
21
       Pthread_mutex_unlock(&m);
22
23
24
   int main(int argc, char *argv[]) {
25
       printf("parent: begin\n");
26
       pthread_t p;
27
       Pthread_create(&p, NULL, child, NULL);
28
       thr_join();
29
       printf("parent: end\n");
       return 0;
31
32
```

Conditional Variables & Semaphores

Producer Consumer Problem

• There are one or more **producer threads** each of which place a task with data item at a buffer, and one or more **consumer threads** that takes a placed task and process it one by one

- Requirements
 - -a producer must be blocked if the buffer is full
 - a consumer must be blocked if there is no task found in buffer
 - -accesses to a shared buffer must be synchronized

Conditional Variables & Semaphores

Producer Consumer with Conditional Variable

```
cond_t empty, fill;
  int buffer[MAX];
  int fill_ptr = 0;
                                                       mutex_t mutex;
  int use_ptr = 0;
   int count
               = 0;
                                                       void *producer(void *arg) {
                                                           int i;
   void put(int value) {
                                                           for (i = 0; i < loops; i++) {
       buffer[fill_ptr] = value;
                                                                Pthread_mutex_lock(&mutex);
       fill_ptr = (fill_ptr + 1) % MAX;
                                                                while (count == MAX)
       count++;
                                                                    Pthread_cond_wait(&empty, &mutex);
10
11
                                                                put(i);
                                                    10
   int get() {
12
                                                                Pthread_cond_signal(&fill);
      int tmp = buffer[use_ptr];
13
                                                                Pthread_mutex_unlock(&mutex);
       use_ptr = (use_ptr + 1) % MAX;
14
                                                    13
       count--;
15
                                                    14
       return tmp;
16
17
                                                    15
                                                       void *consumer(void *arg) {
   void *producer(void *arg) {
                                                           int i;
       int i;
                                                           for (i = 0; i < loops; i++) {
       int loops = (int) arg;
                                                                Pthread_mutex_lock(&mutex);
       for (i = 0; i < loops; i++) {
                                                    19
                                                                while (count == 0)
           put(i);
                                                    20
                                                                    Pthread_cond_wait(&fill, &mutex);
                                                    21
                                                                int tmp = get();
                                                    22
7
                                                                Pthread_cond_signal(&empty);
                                                    23
   void *consumer(void *arg) {
                                                                Pthread_mutex_unlock(&mutex);
                                                    24
       while (1) {
                                                                printf("%d\n", tmp);
10
                                                    25
            int tmp = get();
11
           printf("%d\n", tmp);
12
13
14
```

Semaphore

- A counter-like synchronization primitive
 - represents the number of available shared resources
 - provides increase and decrease operations
 - decrease the counter when a thread is checking out a resource
 - block the caller thread if the counter became negative
 - increase the counter when a thread is putting back a held resource
 - wake up a blocked thread if the counter became zero
- POSIX semaphore
 - sem_init ()
 - sem_wait (), sem_post ()

Conditional Variables & Semaphores

Semaphore As Mutex

```
sem_t m;
sem_init(&m, 0, 1);

sem_wait(&m);
// critical section here
sem_post(&m);
```

Val	Thread 0	State	Thread 1	State
1		Run		Ready
1	call sem_wait()	Run		Ready
0	sem_wait() returns	Run		Ready
0	(crit sect begin)	Run		Ready
0	Interrupt; Switch $ ightarrow$ T1	Ready		Run
0	•	Ready	call sem_wait()	Run
-1		Ready	decr sem	Run
-1		Ready	$(sem<0) \rightarrow sleep$	Sleep
-1		Run	Switch \rightarrow T0	Sleep
-1	(crit sect end)	Run		Sleep
-1	call sem_post()	Run		Sleep
0	incr sem	Run		Sleep
0	wake(T1)	Run		Ready
0	sem_post() returns	Run		Ready
0	Interrupt; Switch $ ightarrow$ T1	Ready		Run
0	•	Ready	sem_wait() returns	Run
0		Ready	(crit sect)	Run
0		Ready	call sem_post()	Run
1		Ready	sem_post() returns	Run

Semaphore as Ordering Primitive

```
sem_t s;
2
   void *child(void *arg) {
       printf("child\n");
       sem_post(&s); // signal here: child is done
       return NULL;
7
8
   int main(int argc, char *argv[]) {
       sem_init(&s, 0, X); // what should X be?
10
       printf("parent: begin\n");
11
       pthread_t c;
12
       Pthread_create(&c, NULL, child, NULL);
13
       sem_wait(&s); // wait here for child
14
       printf("parent: end\n");
15
       return 0;
16
17
```

Conditional Variables & Semaphores

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Implementing Semaphore with Mutext and Conditional Variables

```
typedef struct __Zem_t {
   int value;
   pthread_cond_t cond;
   pthread_mutex_t lock;
} Zem_t;
```

```
// only one thread can call this
   void Zem init(Zem t *s, int value) {
       s->value = value;
       Cond init (&s->cond);
10
       Mutex init(&s->lock);
11
12
   void Zem_wait(Zem_t *s) {
       Mutex lock(&s->lock);
15
       while (s->value <= 0)
16
            Cond wait (&s->cond, &s->lock);
17
        s->value--;
18
       Mutex_unlock(&s->lock);
19
20
   void Zem_post(Zem_t *s) {
22
        Mutex_lock(&s->lock);
23
        s->value++;
24
       Cond_signal(&s->cond);
25
        Mutex unlock (&s->lock);
26
27
```

Producer and Consumer with Semaphore

Version 1

```
sem_t empty;
  sem_t full;
  void *producer(void *arg) {
     int i;
     for (i = 0; i < loops; i++) {
        // Line P2
        put(i);
        sem_post(&full);  // Line P3
10
11
12
  void *consumer(void *arg) {
     int tmp = 0;
14
     while (tmp != -1) {
        16
        tmp = get();
                        // Line C2
17
        printf("%d\n", tmp);
19
20
21
22
  int main(int argc, char *argv[]) {
     // ...
24
     sem_init(&empty, 0, MAX); // MAX are empty
     // ...
28
```

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Producer and Consumer with Semaphore

Version 2

```
void *producer(void *arg) {
   int i;
   for (i = 0; i < loops; i++) {
     put(i);
                // Line P2
     sem_post(&full);
                // Line P3
7
     9
10
11
 void *consumer(void *arg) {
   int i;
13
   for (i = 0; i < loops; i++) {
     15
     16
     int tmp = get();  // Line C2
17
     18
     19
     printf("%d\n", tmp);
20
21
22
```

Conditional Variables & Semaphores

Producer and Consumer with Semaphore

Version 3

```
void *producer(void *arg) {
      int i;
      for (i = 0; i < loops; i++) {
          sem_wait(&empty);
                          // Line P1
          sem_wait(&mutex);
                               // Line P1.5 (MUTEX HERE)
                               // Line P2
         put(i);
                               // Line P2.5 (AND HERE)
          sem_post(&mutex);
                               // Line P3
          sem_post(&full);
10
11
  void *consumer(void *arg) {
      int i;
13
      for (i = 0; i < loops; i++) {
          15
          sem_wait(&mutex);
                               // Line C1.5 (MUTEX HERE)
16
          int tmp = get();
                               // Line C2
17
          sem_post(&mutex);
                               // Line C2.5 (AND HERE)
18
          sem_post(&empty);
                               // Line C3
19
         printf("%d\n", tmp);
20
21
22
```

Conditional Variables & Semaphores