

ITP 30002 Operating System

Process

(OSTEP:Ch.4)

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# Process

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- a running instance of a program
  - program vs. process
- time sharing of a CPU provides the illusion of many CPUs
  - concurrency vs. parallelism
  - mechanism: context switching
  - policies: scheduling

# Constitution of Program Execution Context

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- memory states
  - address space
- CPU states
  - registers: general-purpose and special-purpose
- I/O information

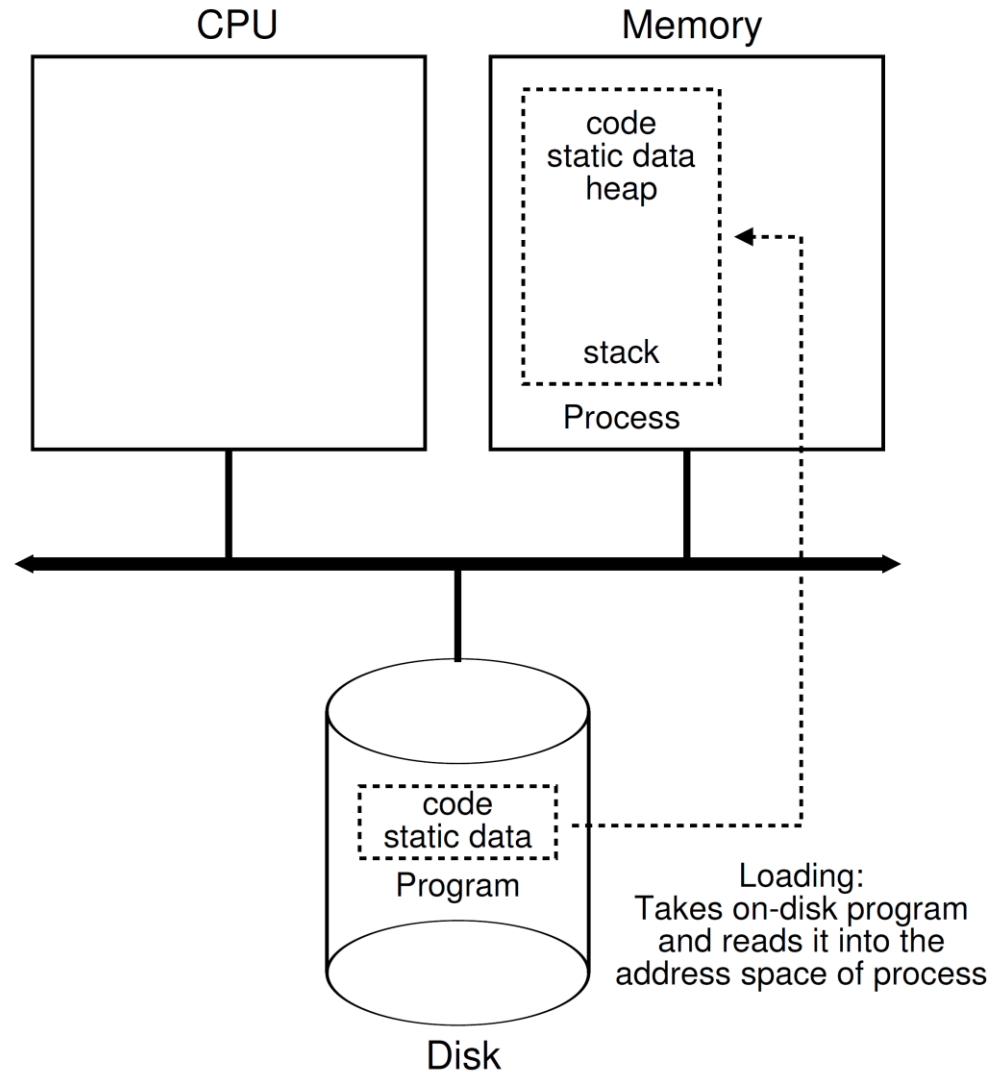
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# Life Cycle of a Process

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- process creation
  - resource allocation
  - loading
    - eager manner
    - lazy manner



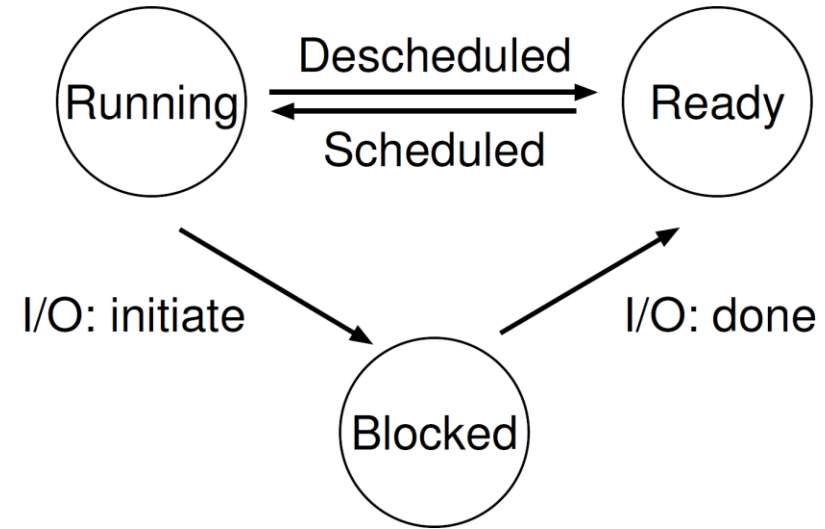
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# Life Cycle of a Process

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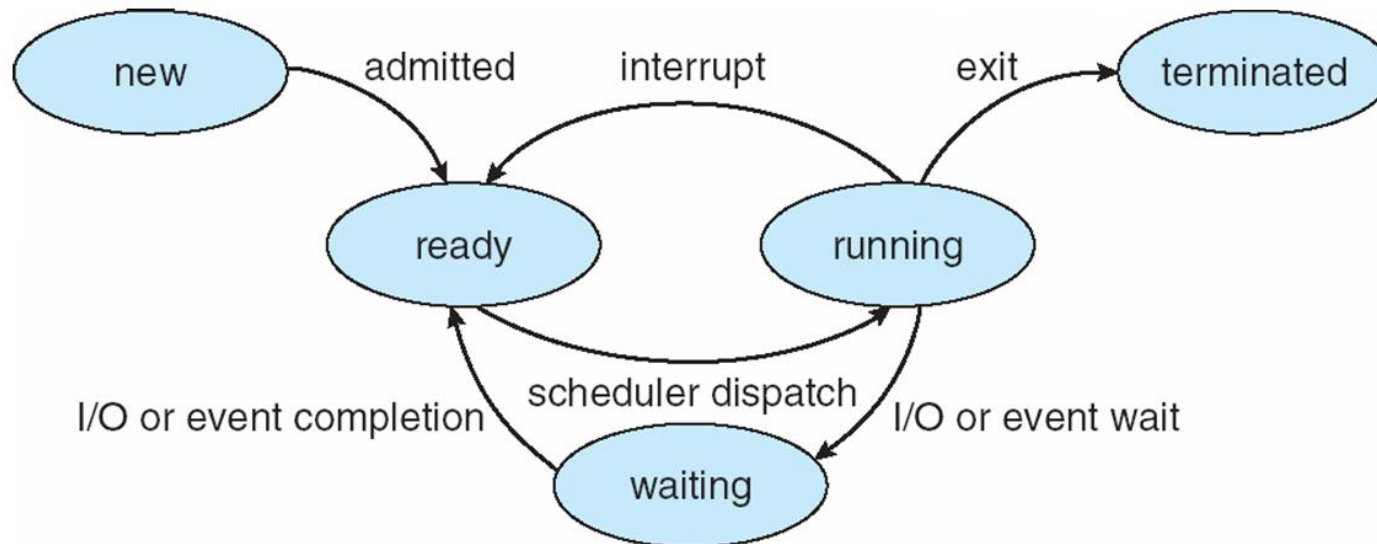
- Running state
  - hold a CPU and execute instructions
- Ready state
  - can make a progress, but cannot hold a CPU
- Blocked state
  - cannot make a progress since it needs to wait for a certain condition (i.e., I/O)
- CPU scheduler makes a decision for process state transitions



# Another Representation of Process Life Cycle

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- As a process executes, the process state changes
  - **new**: The process is being created
  - **ready**: The process is waiting to be assigned to a processor
  - **running**: Instructions are being executed
  - **waiting**: The process is waiting for some event to occur
  - **terminated**: The process has finished execution



# Example

| Time | Process <sub>0</sub> | Process <sub>1</sub> | Notes                         |
|------|----------------------|----------------------|-------------------------------|
| 1    | Running              | Ready                |                               |
| 2    | Running              | Ready                |                               |
| 3    | Running              | Ready                |                               |
| 4    | Running              | Ready                | Process <sub>0</sub> now done |
| 5    | –                    | Running              |                               |
| 6    | –                    | Running              |                               |
| 7    | –                    | Running              |                               |
| 8    | –                    | Running              | Process <sub>1</sub> now done |

| Time | Process <sub>0</sub> | Process <sub>1</sub> | Notes                              |
|------|----------------------|----------------------|------------------------------------|
| 1    | Running              | Ready                |                                    |
| 2    | Running              | Ready                |                                    |
| 3    | Running              | Ready                | Process <sub>0</sub> initiates I/O |
| 4    | Blocked              | Running              | Process <sub>0</sub> is blocked,   |
| 5    | Blocked              | Running              | so Process <sub>1</sub> runs       |
| 6    | Blocked              | Running              |                                    |
| 7    | Ready                | Running              | I/O done                           |
| 8    | Ready                | Running              | Process <sub>1</sub> now done      |
| 9    | Running              | –                    |                                    |
| 10   | Running              | –                    | Process <sub>0</sub> now done      |

# Process Data Structure in Kernel: xv6 Example

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- Called as Process Control Block
- Example of the xv6 kernel

```
// the registers xv6 will save and restore
// to stop and subsequently restart a process
struct context {
    int eip;
    int esp;
    int ebx;
    int ecx;
    int edx;
    int esi;
    int edi;
    int ebp;
};

// the different states a process can be in
enum proc_state { UNUSED, EMBRYO, SLEEPING,
                  RUNNABLE, RUNNING, ZOMBIE };
```

```
// the information xv6 tracks about each process
// including its register context and state
struct proc {
    char *mem;                // Start of process memory
    uint sz;                  // Size of process memory
    char *kstack;             // Bottom of kernel stack
                               // for this process
    enum proc_state state;    // Process state
    int pid;                  // Process ID
    struct proc *parent;      // Parent process
    void *chan;               // If !zero, sleeping on chan
    int killed;               // If !zero, has been killed
    struct file *ofile[NOFILE]; // Open files
    struct inode *cwd;        // Current directory
    struct context context;   // Switch here to run process
    struct trapframe *tf;     // Trap frame for the
                               // current interrupt
};
```



# Remainig Issues

- What's the overhead for virtualizing CPU?
- How to implement context switching?
- How to manage multiple processes?
- How a scheduler determines a next process to execute?