The following materials have been collected from the numerous sources such as Stanford CS106 and Harvard CS50 including my own and my students over the years of teaching and experiences of programming. Please help me to keep this tutorial up-to-date by reporting any issues or questions. Please send any comments or criticisms to idebtor@gmail.com. Your assistances and comments will be appreciated..

PSet 02a - Recursion

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Getting Started - Recursion

Recursion in computer science is a method where the solution to a problem depends on solutions to **smaller instances** of the same problem (as opposed to iteration). The approach can be applied to many types of problems, and recursion is one of the central ideas of computer science. An infinite number of computations can be described by a finite recursive program, even if this program contains no explicit repetitions. Most computer programming languages support recursion by allowing a <u>function</u> to call itself within the program text.

(Resources: https://en.wikipedia.org/wiki/Recursion & www.codingBat.com)

Recursive algorithm is expressed in terms of

- 1. base case(s) for which the solution can be stated non-recursively,
- 2. **recursive case(s)** for which the solution can be expressed in terms of a smaller version of itself.

The following figure shows how recursion works by calling itself over and over again.

Implementing Recursion Functions

- 1. Take a look at the sample solution of Example 1 ~ 4 in recursion.cpp and recusrionDriver.cpp.
- 2. Implement the rest of examples while following the patterns of the program including comment sections in **recursion.cpp**. Well written comments are required such that that we can clearly see both how the algorithms work and what the testing results are.

Example 1: Factorial

```
fact(n) = \begin{cases} 1 & \text{if } n = 0\\ n \cdot fact(n-1) & \text{if } n > 0 \end{cases}
```

```
factorial(n) – iterative solution

long long factorial(n) {
  long long total = 1;
  for (int i = n; i > 1; i--) total *= i;
  return total;
}
```

```
factorial(n) – recursive solution

input: integer n such that n >= 0

output: [n \times (n-1) \times (n-2) \times ... \times 1]

1. if n is 0, return 1

2. otherwise, return [n \times factorial(n-1)]

end factorial
```

```
factorial(1) \rightarrow 1
factorial(2) \rightarrow 2
factorial(3) \rightarrow 6
factorial(8) \rightarrow 40320
factorial(12) \rightarrow 479001600
factorial(20) \rightarrow 2432902008176640000
```

Hint:

First, detect the "base case", a case so simple that the answer can be returned immediately (here

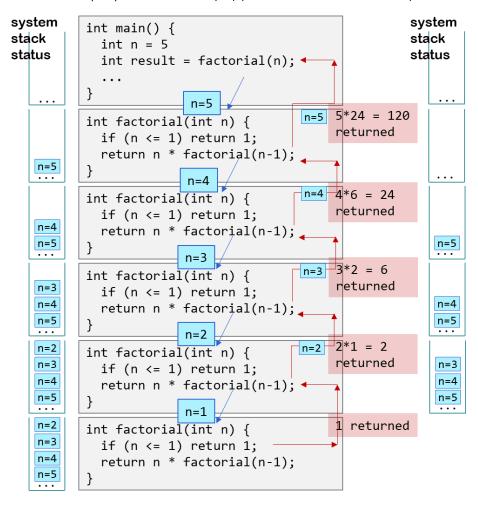
when n==1 or n==0). Otherwise make a recursive call of factorial(n-1) (towards the base case). Assume the recursive call returns a correct value and fix that value up to make our result.

When we think about solving this problem recursively, we need to figure out what our subproblems will be. Let's break it down:

- 1. We know factorial(5) = 5 * factorial(4) aka 5! = 5 * 4!.
- 2. To continue, factorial(5) = 5 * (4 * factorial(3)) which equals 5 * (4 * (3 * factorial(2)) and so on...
- 3. ...Until you get 5 * 4 * 3 * 2 * 1 and the only remaining subproblem is 1!.
- 4. factorial(1) and factorial(0) always equals 1 so this will be our base case.

As we can see in the figure below, the factorial() function is calling itself. However, during each call, we have decreased the value of n by 1. When n is less than 1, the factorial() function ultimately returns the output.

Remember that when a function keeps on calling another without finishing, the current status of function or n value is pushed (or saved) in the system stack. Then as soon as it gets a returned value from factorial(n-1), the n value is popped retrieved for the computation.



Using this line of thinking, we can write a recursive solution to our factorial problem as shown below:

Code Example 1:

```
long long unsigned factorial(int n) {
  if (n == 1 || n == 0) return n;
  return n * factorial(n-1);
}
```

Code Example 2:

```
long long unsigned factorial(int n) {
  return (n == 1 || n == 0) ? 1 : n * factorial(n - 1); // using ternary operator
}
```

Code Example 3:

```
#include <iostream>
long long unsigned factorial(int n) {
   std::cout << "n= " << n << std::endl;
   if (n == 1 || n== 0) return n;
   auto result = n * factorial(n - 1);
   std::cout << "n= " << n << "\tn!= " << result << std::endl;
   return result;
}
int main() {
   factorial(5);
}</pre>
```

Sample Run:

The first function call begins with n = 5, but it never ends until all the factorials from 4, 3, 2, 1 are computed first. In the main (), it calls the function only one, but the function calls itself until n = 1 or the base case is reached.

I recommend that you follow the code by hand or step through the debugger and understand the concept of the recursion.

Practice all three Code Examples and turn in the source code file "Code Example 3" in factorial.cpp.

Example 2: GCD (Great Common Divisor)

```
\gcd(x,y) = \begin{cases} x & \text{if } y = 0\\ \gcd(y, \operatorname{remainder}(x,y)) & \text{if } y > 0 \end{cases}
```

Recurrence relation for greatest common divisor, where x%y expresses the remainder of x/y:

$$gcd(x,y) = gcd(y,x\%y) if y! = 0$$
$$= gcd(x,0) = x if y = 0$$

```
gcd(x, y)
input: integer x, y such that x >= y, y > 0
output: gcd of x and y
1. if y is 0, return x
```

```
2. otherwise, return [ gcd (y, x%y) ] end gcd
```

Ex: Computing the recurrence relation for x = 27 and y = 9

```
gcd(27, 9) = gcd(9, 27 % 9)
= gcd(9, 0)
```

Ex: Computing the recurrence relation for x = 111 and y = 259

```
gcd(111, 259) = gcd(259, 111% 259)

= gcd(259, 111)

= gcd(111, 259% 111)

= gcd(111, 37)

= gcd(37, 111% 37)

= gcd(37, 0)

= 37
```

Code:

```
int gcd(int x, int y) {
  if (y == 0) return x;
  return gcd(y, x % y);
}
```

or

```
int gcd(int x, int y) {
  return y == 0 ? x : gcd(y, x % y);
}
```

Example 3: Fibonacci

The fibonacci sequence is a famous bit of mathematics, and it happens to have a recursive definition. The first two values in the sequence are 0 and 1 (essentially 2 base cases). Each subsequent value is the sum of the previous two values, so the whole sequence is: 0, 1, 1, 2, 3, 5, 8, 13, 21 and so on.

Define a recursive **fibonacci(n)** method that returns the nth fibonacci number, with n=0 representing the start of the sequence.

```
fibonacci(0) \rightarrow 0 fibonacci(1) \rightarrow 1 fibonacci(2) \rightarrow 1 fibonacci(3) \rightarrow 2 fibonacci(4) \rightarrow 3 fibonacci(5) \rightarrow 5 fibonacci(11) \rightarrow 89 fibonacci(33) \rightarrow 3524578 fibonacci(44) \rightarrow 701408733 (It takes some time to compute.)
```

Code:

```
long long fibonacci(int n) {
  if (n == 0 || n == 1) return n;
  return fibonacci(n-1) + fibonacci(n-2);
}
```

or

```
long long fibonacci(int n) {
  return n <= 1 ? n : fibonacci(n - 1) + fibonacci(n - 2);
}</pre>
```

Example 4: Bunny Ears

We have a number of bunnies and each bunny has two big floppy ears. We want to compute the total number of ears across all the bunnies recursively (without loops or multiplication).

```
bunnyEars(0) \rightarrow 0
bunnyEars(1) \rightarrow 2
bunnyEars(2) \rightarrow 4
bunnyEars(3) \rightarrow 6
bunnyEars(234) \rightarrow 468
```

Hint:

First detect the base case (bunnies == 0), and in that case just return 0. Otherwise, make a recursive call to bunnyEars(bunnies-1). Trust that the recursive call returns the correct value, and fix it up by adding 2.

Code:

```
int bunnyEars(int bunnies) {
if (bunnies == 0) return 0;

// Recursive case: otherwise, make a recursive call with bunnies-1
// (towards the base case), and fix up what it returns.
return 2 + bunnyEars(bunnies-1);
}
```

Example 5: Funny Ears

We have bunnies and funnies standing in a line, numbered 1, 2, ... The odd bunnies (1, 3, ...) have the normal 2 ears. The even funnies (2, 4, ...) we'll say have 3 ears, because they each have a raised foot. Recursively return the number of "ears" in the bunny and funny line 1, 2, ... n (without loops or multiplication).

```
funnyEars(0) \rightarrow 0
funnyEars(1) \rightarrow 2
funnyEars(2) \rightarrow 5
funnyEars(3) \rightarrow 7
funnyEars(4) \rightarrow 10
funnyEars(9) \rightarrow 22
funnyEars(10 \rightarrow 25
```

Code:

```
int funnyEars(int funnies) {
  // your code here
}
```

Example 6: Triangle

We have triangle made of blocks. The topmost row has 1 block, the next row down has 2 blocks, the next row has 3 blocks, and so on. Compute recursively (no loops or multiplication) the total number of blocks in such a triangle with the given number of rows.

```
triangle(0) \rightarrow 0
triangle(1) \rightarrow 1
triangle(2) \rightarrow 3
triangle(3) \rightarrow 6
```

```
triangle (4) \rightarrow 10 triangle (7) \rightarrow 28
```

Code:

```
int triangle(int rows) {
// your code here
}
```

Example 7: Sum of digits

Given a non-negative int n, return the sum of its digits recursively (no loops). Note that mod (%) by 10 yields the rightmost digit (126 % 10 is 6), while divide (/) by 10 removes the rightmost digit (126 / 10 is 12).

```
sumDigits(126) \rightarrow 9
sumDigits(12) \rightarrow 3
sumDigits(1) \rightarrow 1
sumDigits(10110) \rightarrow 3
sumDigits(235) \rightarrow 10
```

Code:

```
int sumDigits(int n) {
  // your code here
}
```

Example 8: Count 8

Given a non-negative int n, return the count of the occurrences of 8 as a digit, so for example 818 yields 2. (no loops). Note that mod (%) by 10 yields the rightmost digit (126 % 10 is 6), while divide (/) by 10 removes the rightmost digit (126 / 10 is 12).

```
count8(818) \rightarrow 2
count8(8) \rightarrow 1
count8(123) \rightarrow 0
count8(881238) \rightarrow 3
count8(48581) \rightarrow 2
count8(888586198) \rightarrow 5
count8(99899) \rightarrow 1
```

Code:

```
int count8(int n) {
// your code here
}
```

Example 9: Power N

Given base and n that are both 1 or more, compute recursively (no loops) the value of base to the n power, so powerN(3, 2) is 9 (3 squared).

```
powerN(2, 5) \rightarrow 32
powerN(3, 1) \rightarrow 3
powerN(3, 2) \rightarrow 9
powerN(3, 3) \rightarrow 27
powerN(10, 2) \rightarrow 100
powerN(10, 3) \rightarrow 1000
```

Code:

```
long long powerN(int base, int n) {
```

```
// your code here
}
```

A command line to build

To build this program:

```
g++ recursion.cpp recursionDriver.cpp -I../../include -L../../lib -lnowic -o recursion
```

This command line works at pset2a folder and the following folder structures and files exist.

nowic/lib/libnowic.a
nowic/include/nowic.h
nowic/psets/pset2a/recursion.cpp
nowic/psets/pset2a/recursionDriver.cpp

Submitting your solution

•	Include the following	ine at the top of y	our every source file	e with your name signed.	
	On my honour, I pledo	ge that I have neit	her received nor pro	ovided improper assistance in the	
	completion of this assignment.				
	Signed:	Section:	Student Number:		

- Make sure your code compiles and runs right before you submit it. Every semester, we get dozens of submissions that don't even compile. Don't make "a tiny last-minute change" and assume your code still compiles. You will not get sympathy for code that "almost" works.
- If you only manage to work out the problem sets partially before the deadline, you still need to turn it in. However, don't turn it in if it does not compile and run.
- Place your source files in the folder you and I are sharing.
- After submitting, if you realize one of your programs is flawed, you may fix it and submit again if it is **before the deadline**. You will have to resubmit any related files together, even if you only change one. You may submit as often as you like. **Only the last version** you submit before the deadline will be graded.

Files to submit and Grade

Submit the following file.

- factorial.cpp
- recursion.cpp
- 2.0 points

Upload your file **pset2a** folder in Piazza.

Due

Due: 11:55 pm