SortsUnderstanding Array Sorting Algorithms, Efficiency, Comparison and Implementation

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Runtime and Memory Complexity: Basics

- Runtime Complexity
 - Time an algorithm takes relative to input length.
- Memory Complexity
 - Memory needed by an algorithm relative to input size.
- Both are crucial for:
 - o Comparing algorithm efficiency.
 - o Choosing the right algorithm for the job.

Measuring Program Execution Time: Challenges Types of Time Measurements

• Time Measurement Issues

- Real Time: Wall-clock time for program execution.
- o User Time: CPU time for executing user program code.
 - · Excludes system operations.
 - · Reflects direct program execution time.
- System Time¹: CPU time for system operations for the program.
 - · File operations, I/O tasks.
 - · Essential for resource-intensive operations.
- Variability in measurements due to:
 - o System load, resources, hardware.
 - o Inconsistencies across environments.

¹See https://stackoverflow.com/questions/556405/ what-do-real-user-and-sys-mean-in-the-output-of-time1

Measuring Program Execution Time: Theoretical Approach Abstracting Time Measurements

• Theoretical Approach

- Approximate with input size (n) and operation count.
- o $n \in \mathbb{N}^*$: Number of loops or iterations main driver of complexity.
- $k \in \mathbb{N}^*$: Parameters affecting complexity, aside from input size.
 - · Here intended as range of the non-negative key values
- Focus on growth trends rather than exact times.

• Conclusion:

- o Theoretical focus helps identify scalability issues.
- o Prioritizes relative efficiency over absolute timing.

Understanding Big O Notation

- **Big O Notation**: Describes the upper bound of complexity.
 - Focuses on worst-case scenario.
 - Ignores constant factors and lower order terms.

Basic Rules

- Linear Terms: $O(\alpha n + \beta) = O(n)$.
 - · Constants α , β don't affect growth rate.
- Sum Rule: $O(f(n)) + O(g(n)) = O(\max(f(n), g(n)))$.
- $\circ \ \textit{Product Rule:} \ \mathfrak{O}(f(n)) \cdot \mathfrak{O}(g(n)) = \mathfrak{O}(f(n) \cdot g(n)).$

Implications

- o Simplifies comparing algorithms.
- o Emphasizes dominant factors affecting growth.

• Example

- $\circ \ \mathcal{O}(3n^2 + 10n + 100) = \mathcal{O}(n^2).$
 - · n^2 term dominates as n grows.

Classic Sorting Algorithms and Their Complexities

Algo.	Best	Average	Worst	Mem.
Selection Sort	$O(n^2)$	$O(n^2)$	$O(n^2)$	O(1)
Insertion Sort	O(n)	$O(n^2)$	$O(n^2)$	O(1)
Bubble Sort	O(n)	$O(n^2)$	$O(n^2)$	O(1)
Merge Sort	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$	O(n)
Quick Sort	$O(n \log n)$	$O(n \log n)$	$O(n^2)$	$O(\log n)$

More Algorithms and Their Complexities²

Algorithm	Best	Average	Worst	Memory
Selection Sort	$\Omega(n^2)$	$\Theta(n^2)$	$O(n^2)$	O(1)
Insertion Sort	$\Omega(n)$	$\Theta(n^2)$	$O(n^2)$	O(1)
Bubble Sort	$\Omega(n)$	$\Theta(n^2)$	$O(n^2)$	O(1)
Merge Sort	$\Omega(n \log n)$	$\Theta(n \log n)$	$O(n \log n)$	O(n)
Quick Sort	$\Omega(n \log n)$	$\Theta(n \log n)$	$O(n^2)$	$O(\log n)$
Timsort	$\Omega(n)$	$\Theta(n \log n)$	$O(n \log n)$	O(n)
Heapsort	$\Omega(n \log n)$	$\Theta(n \log n)$	$O(n \log n)$	O(1)
Tree Sort	$\Omega(n \log n)$	$\Theta(n \log n)$	$O(n^2)$	O(n)
Shell Sort	$\Omega(n \log n)$	$\Theta(n(\log n)^2)$	$O(n(\log n)^2)$	O(1)
Bucket Sort	$\Omega(n+k)$	$\Theta(n+k)$	$O(n^2)$	O(n)
Radix Sort	$\Omega(nk)$	$\Theta(nk)$	O(nk)	O(n+k)
Counting Sort	$\Omega(n+k)$	$\Theta(n+k)$	O(n+k)	O(k)
Cubesort	$\Omega(n)$	$\Theta(n \log n)$	$O(n \log n)$	O(n)

²See https://www.bigocheatsheet.com/

Selection Sort Algorithm

Simplicity in Action

- Iteratively selects the smallest element from the unsorted portion and swaps it with the element at the current position.
- Continues until the entire array is sorted.

Insertion Sort Algorithm Building the Sorted Array

- Builds the sorted array one element at a time.
- Iterates through the array, shifting elements to the right to make space for the current element.

```
def insertion_sort(arr):
    for i in range(1, len(arr)):
        key = arr[i]
        j = i-1
        while j >= 0 and key < arr[j]:
        arr[j+1] = arr[j]
        j -= 1
        arr[j+1] = key
    return arr</pre>
```

Bubble Sort Algorithm

Bubbling Up the Largest

Element

- Repeatedly steps through the list, compares adjacent elements and swaps them if they are in the wrong order.
- o Continues until no more swaps are needed.

```
def bubble_sort(arr):
    n = len(arr)
    for i in range(n):
        for j in range(0, n-i-1):
            if arr[j] > arr[j+1]:
            arr[j], arr[j+1] = arr[j+1], arr[j]
    return arr
```

Quicksort Algorithm

Divide and Conquer

- Employs a divide-and-conquer strategy to sort the array.
- Selects a 'pivot' element and partitions the array into sub-arrays of elements less than and greater than the pivot.
- o Recursively applies the same strategy to the sub-arrays.
- Highly efficient for large datasets, with average time complexity of O(n log n).

```
def quicksort(arr):
    if len(arr) <= 1:
        return arr
    pivot = arr[len(arr) // 2]
    left = [x for x in arr if x < pivot]
    middle = [x for x in arr if x == pivot]
    right = [x for x in arr if x > pivot]
    return quicksort(left) + middle + quicksort(right)
```

Merge Sort Algorithm

How It Works

- Another divide-and-conquer algorithm that divides the array into halves, sorts each half, and then merges the sorted halves together.
- Begins with the division of the array into smallest manageable units, then merges units in a sorted manner to form larger sorted sections until the whole array is merged back together.
- Consistently performs with a time complexity of O(n log n), making it highly efficient for large data sets.

• Key Features

- o Predictable performance.
- Stable sorting algorithm.
- Popular. Default sorting algorithm for many programming languages, but tends to be replaced by Timsort in practice.

Merge Sort Algorithm

```
def merge_sort(arr):
    if len(arr) > 1:
        mid = len(arr) // 2
        L = arr[:mid]
        R = arr[mid:]
        merge_sort(L)
        merge_sort(R)
        i = j = k = 0
        while i < len(L) and j < len(R):
            if L[i] < R[j]:
                arr[k] = L[i]
                i += 1
            else:
                arr[k] = R[i]
                i += 1
            k += 1
        while i < len(L):
            arr[k] = L[i]
            i, k = i + 1, k + 1
        while j < len(R):
            arr[k] = R[j]
            j, k = j + 1, k + 1
    return arr
```

Thank You for Your Attention! Further Resources

Useful Links

- o Big O Cheat Sheet: https://www.bigocheatsheet.com/
 - · A handy reference for complexity of common data structures and algorithms.
- Sorting Algorithms with Animations: https: //www.toptal.com/developers/sorting-algorithms/bubble-sort
 - Explore how different sorting algorithms work with interactive animations.

Contact & Feedback

- My GitHub: https://github.com/onyr
- This presentation: https://github.com/onyr/sorting_algorithms

Once again, thank you and have a great day!