SortsUnderstanding Array Sorting Algorithms, Efficiency, Comparison and Implementation

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Runtime and Memory Complexity: Basics

- Runtime Complexity
 - o Time an algorithm takes relative to input length.
- Memory Complexity
 - o Memory needed by an algorithm relative to input size.
- Both are crucial for:
 - o Comparing algorithm efficiency.
 - Choosing the right algorithm for the job.

Measuring Program Execution Time: Challenges Types of Time Measurements

• Time Measurement Issues

- Real Time: Wall-clock time for program execution.
- User Time: CPU time for executing user program code.
 - · Excludes system operations.
 - · Reflects direct program execution time.
- System Time¹: CPU time for system operations for the program.
 - · File operations, I/O tasks.
 - · Essential for resource-intensive operations.
- Variability in measurements due to:
 - o System load, resources, hardware.
 - o Inconsistencies across environments.

¹See https://stackoverflow.com/questions/556405/ what-do-real-user-and-sys-mean-in-the-output-of-time1

Measuring Program Execution Time: Theoretical Approach Abstracting Time Measurements

• Theoretical Approach

- Approximate with **input size** (n) and **operation count**.
- o $n \in \mathbb{N}^*$: Number of loops or iterations main driver of complexity.
- o $k \in \mathbb{N}^*$: Parameters affecting complexity, aside from input size.
 - · Here intended as range of the non-negative key values
- Focus on growth trends rather than exact times.

• Conclusion:

- o Theoretical focus helps identify scalability issues.
- o Prioritizes relative efficiency over absolute timing.

Understanding Big O Notation

- **Big O Notation**: Describes the upper bound of complexity.
 - o Focuses on worst-case scenario.
 - o Ignores constant factors and lower order terms.

• Basic Rules

- Linear Terms: $O(\alpha n + \beta) = O(n)$.
 - · Constants α , β don't affect growth rate.
- Sum Rule: $O(f(n)) + O(g(n)) = O(\max(f(n), g(n))).$
- Product Rule: $O(f(n)) \cdot O(g(n)) = O(f(n) \cdot g(n))$.

Implications

- o Simplifies comparing algorithms.
- o Emphasizes dominant factors affecting growth.

• Example

- $\circ \ \mathcal{O}(3n^2 + 10n + 100) = \mathcal{O}(n^2).$
 - · n^2 term dominates as n grows.



Classic Sorting Algorithms and Their Complexities

Algo.	Best	Average	Worst	Mem.
Selection Sort	$O(n^2)$	$O(n^2)$	$O(n^2)$	0(1)
Insertion Sort	O(n)	$O(n^2)$	$O(n^2)$	O(1)
Bubble Sort	O(n)	$O(n^2)$	$O(n^2)$	O(1)
Merge Sort	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$	O(n)
Quick Sort	$O(n \log n)$	$O(n \log n)$	$O(n^2)$	$O(\log n)$



More Algorithms and Their Complexities²

Algorithm	Best	Average	Worst	Memory
Selection Sort	$\Omega(n^2)$	$\Theta(n^2)$	$O(n^2)$	O(1)
Insertion Sort	$\Omega(n)$	$\Theta(n^2)$	$O(n^2)$	O(1)
Bubble Sort	$\Omega(n)$	$\Theta(n^2)$	$O(n^2)$	O(1)
Merge Sort	$\Omega(n \log n)$	$\Theta(n \log n)$	$O(n \log n)$	O(n)
Quick Sort	$\Omega(n \log n)$	$\Theta(n \log n)$	$O(n^2)$	$O(\log n)$
Timsort	$\Omega(n)$	$\Theta(n \log n)$	$O(n \log n)$	O(n)
Heapsort	$\Omega(n \log n)$	$\Theta(n \log n)$	$O(n \log n)$	O(1)
Tree Sort	$\Omega(n \log n)$	$\Theta(n \log n)$	$O(n^2)$	O(n)
Shell Sort	$\Omega(n \log n)$	$\Theta(n(\log n)^2)$	$O(n(\log n)^2)$	O(1)
Bucket Sort	$\Omega(n+k)$	$\Theta(n+k)$	$O(n^2)$	O(n)
Radix Sort	$\Omega(nk)$	$\Theta(nk)$	O(nk)	O(n+k)
Counting Sort	$\Omega(n+k)$	$\Theta(n+k)$	O(n+k)	O(k)
Cubesort	$\Omega(n)$	$\Theta(n \log n)$	$O(n \log n)$	O(n)

²See https://www.bigocheatsheet.com/