

CALIFORNIA STATE UNIVERSITY, NORTHRIDGE

Implementation of (255, 223) Reed Solomon Code
Encoder /Decoder

A graduate project submitted in partial fulfillment of the requirements
For the degree of Master of Science in
Electrical Engineering

By

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Dedication

To my parents and
my brothers

Acknowledgements

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Abstract

Implementation of (255,223) Reed Solomon code

Encoder /Decoder

By

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Master of Science in Electrical Engineering

In this project, the algorithm of the Reed-Solomon code encoder and decoder is discussed, followed by the implementations of (255, 223) Reed-Solomon codes. The encoder takes a 223-byte data block and generates a 255-byte code block to be transmitted on a digital communication channel. This code is defined over a Galois Field $GF(2^8)$ and has the capability of correcting up to sixteen short bursts of errors. The software part of the design is written in C language and tested using Simulink, and the hardware part is implemented using a microcontroller.

The encoding process is simpler than the decoding process and requires less complex hardware. The decoder, however, is implemented using a parallel structure to speed up the Galois Field computations necessary to find out the locations of the errors and their values.

The decoder logic circuit comprises multiple units in order to perform a parallel computation for the syndrome calculation. The Berlekamp's iterative algorithm is used to determine the coefficients of the error location polynomial. At the same time, the Chien's searching algorithm is used to calculate the roots of the error location polynomial, their error values, and perform the correction of the found errors.

Chapter 1

Introduction

1.1. Introduction

In 1948, Claude Shannon, considered to be the founding father of electronic communications, began to prove that codes with special properties can accomplish unfailing data communications over a noisy channel, if the capacity of the channel is larger than the data rate of the message source. This proof has opened the door to further discoveries in that field, where substantial work has been done not only in digital communications, but also in error control coding theory. Since that time, many engineers have been trying to construct such code to prove the Shannon theorem without any mathematical abstracts. The engineers have ended up, however, being able to detect the error but not correct it; one code appending the parity check bit to the message bits results in a single bit error.

The demands of achieving reliable coding have continued over time due to the high volume of data exchanged. The new code needed to be constructed on form where it has enough redundancy to be able to detect, or detect and correct, multiple bit errors up to a definite probability that relates to the characteristics of the digital channel. The code rates (message length/ code length in bits) are close to unity and are easy to implement in hardware.

The cyclic code is one of the codes that has been developed and has become a useful code in data communications. This code is very reliable and relatively easy to implement. In practical applications, the cyclic code is used in systematic forms which have the advantage of the parity data, which can simply be appended to the source block, and receivers do not need to do any work to recover the original source symbols if received correctly.

The Bose-Chaudhuri-Hocquenghem (BCH) codes are a large class of cyclic code. The BCH comes in two forms: binary BCH code and nonbinary BCH code. Among the nonbinary

BCH code, the Reed Solomon code contains code symbols that are elements of a finite field (Galois Field) with an order of 2^m . All information, voice, video, pictures, and data are represented by a long string of zeroes and ones, also known as bits. All digital systems use binary symbols at the hardware level to represent any type of data; therefore, the representation of elements from $GF(2^m)$ is straightforward.

1.2. Objective

The objective of this project is to design and implement an error detection and correction system using Reed Solomon code. The first part discusses the general procedure for encoding and decoding Reed Solomon codes.

The block code length in the code has been selected as 255 and the message length is 223, which is known as (255, 223) Reed Solomon codes. The code rate is $223/255 = 0.875$ and this code is defined over $GF(2^8)$, where each symbol is represented by an 8-bit (one byte) to make $255 \times 8 = 2040$ binary bits transmitted serially. This code is capable of detecting and correcting up to sixteen errors bytes in every block of 255 bytes. In other words, errors in up to 16 bytes anywhere in the codeword can be automatically corrected. The choice over $GF(2^8)$ is justified by the large spread of the popular 8-bit data buses used in the industry today. This code is not hard to implement in hardware which is not a part of this project. However, testing the code in software using both Simulink and C language makes this code more reliable and the most popular code used for any type of data communication.

1.3. Outline

The project consists of eight chapters and three appendixes. The second chapter is an introduction to information theory and coding. There is also a brief introduction about Hamming code as well as Reed Solomon code. Please note that the basic algebraic field structure should be understood by the reader. Chapter three contains the theory and the theoretical description of the Reed Solomon code encoder and decoder. Chapter four presents the general concept of the project from the hardware point of view. Chapter five is the implementation of the encoder and chapter six is the implementation of the decoder. Chapter seven includes the test of the design and an explanation of the use of simulation of (255, 223) Reed Solomon code. Chapter eight is the conclusion. The appendixes contain assisting materials.

The simulation software used in this project is Simulink, representing the project as blocks where the parameters are specified. The Matlab is used to generate nonzero elements in the field and C language source code is used in the implementation of the encoder and the decoder.

Chapter 2

Hamming codes

2.1. Hamming codes

The early 1950s is the beginning of error correcting codes. Richard W. Hamming discovered the first class of linear block code for error correction, which was later named after him and is known as Hamming code. Hamming code is a set of error-correction codes used to detect and correct bit errors occurring in digital communication or data storage systems and can be defined for any positive integer $m \geq 3$ having the following properties:

Code length	$n = 2^m - 1$
Number of parity check bits	$n - k = m$
Number of information symbols	$k = (2^m - 1) - m$
Minimum distance	$d_{\min} = 3$
Error correcting capability	$t = 1$

The Hamming code's single-error-correcting (SEC) and double-error-detecting (DED) extended versions marked the beginning of coding theory. These codes remain important for both theoretical and practical reasons.

2.2. BCH codes

The generalization of the Hamming codes for multiple error correction codes is known as BCH codes, named and discovered independently by A. Hocquenghem in 1959 and R.C Bose and D.K. Ray-Chaudhuri in 1960. BCH codes are cyclic codes and are, by far, the most extensive and powerful codes. BCH codes are first defined in binary symbols and then generalized to codes in p^m for any codes where p is a prime number and m is any positive

number. In this section, binary BCH codes will be described and the following section will address nonbinary BCH code.

For any positive integer m and t where $t < 2^{m-1}$, there are binary BCH codes which have these properties:

Block length $n = 2^m - 1$

Number of parity check digits $n - k \leq mt$

Minimum distance $d_{\min} \geq 2t + 1$

Binary BCH codes are capable of correcting any combination of t or fewer errors in a code block of $n = 2^m - 1$ digits, which is a t -error correcting BCH codes.

Letting α be a primitive element in $GF(2^m)$ and consider the following sequence:

$$\alpha^1, \alpha^2, \alpha^3, \dots, \alpha^{2t}$$

If $m_i(x)$ is the minimum polynomial of α^i , then the generator $g(x)$ polynomial of the t -error correcting BCH code is the least common multiple of $m_1(x), m_2(x), \dots, m_{2t}(x)$

$$\text{i.e. } g(x) = \text{LCM} \{m_1(x), m_2(x), \dots, m_{2t}(x)\}$$

$\alpha^1, \alpha^2, \alpha^3, \dots, \alpha^{2t}$ respectively, they are roots of $g(\alpha^i) = 0$. If i is an even integer, then $i = i' 2^j$ where i' is an odd integer and $j \geq 1$, and they have the same polynomial i.e. $m_i(x) = m_{i'}(x)$ which means that $g(x)$ for t -correcting errors can be reduced to

$$g(x) = \text{LCM} \{m_1(x), m_2(x), \dots, m_{2t}(x)\}$$

where the degree of any minimum polynomial is m or less and the degree of $g(x)$ is mt or less; if t is small, then $n-k$ exactly equals mt .

2.3. Nonbinary BCH codes and Reed Solomon codes

Nonbinary BCH codes can be defined as, if p is a prime number and q is any power of p , then codes with code symbols from the elements of Galois Field $GF(q)$ exist. These codes are called q -ary codes and for any positive integers s and t there exist a q -ary BCH code with block length $n = q^s - 1$ which can correct any combination of t or fewer errors and require no more than $2st$ parity check bits. However, when $s = 1$ is special subclass of q -ary BCH and known as Reed Solomon codes. Reed Solomon codes are t -error correction codes and have the following properties:

Block length	$n = q - 1$
Number of parity bits	$n - k = 2t$
Dimension	$k = q - 1 - 2t$
Minimum distance	$d = 2t + 1$

Letting $q = 2^m$ then the generator polynomial of a primitive t error correcting Reed Solomon codes will be of length $2^m - 1$ and can be expressed as:

$$g(x) = (x + \alpha)(x + \alpha^2) \dots (x + \alpha^{2t})$$

Where α is a primitive element of $GF(2^m)$.

In Reed Solomon codes, the length of the code is one less than the size of the code alphabet and the minimum distance is one greater than the number of parity check symbol $n - k$. So, correcting t symbol errors requires no more than $2t$ parity symbols. Reed Solomon codes, constructed by Reed and Solomon in 1960, are therefore very effective in correcting random symbol errors and random burst errors in communication, as well as data storage systems, ranging from deep space telecommunications to compact discs.

Chapter 3

Reed Solomon encoder/decoder

3.1. Reed-Solomon encoding

The generating polynomial for an R-S code is driven as:

$$g(X) = g_0 + g_1X + g_2X^2 + \dots + g_{2t-1}X^{2t-1} + X^{2t}$$

With degree equal to the number of parity symbols (n-k). Since the generator polynomial of degree n-k, there must be precisely n-k successive powers of α ($\alpha, \alpha^2, \alpha^3, \dots, \alpha^{2t}$) that are roots of the polynomial. It is not necessary to start with the root α ; starting with any power of α is possible. Reed Solomon codes are cyclic codes and encoding in systematic form is analogous to the binary encoding procedure. By shifting a message polynomial $m(X)$, which is to be encoded into the rightmost k stages of a codeword register, and then appending a parity polynomial, $p(X)$ by placing it in the leftmost n-k stages. Then, multiplying $m(X)$ by X^{n-k} is done to manipulate the message polynomial algebraically so that it is right-shifted n - k positions. Next, dividing $X^{n-k} m(X)$ by the generator polynomial $g(X)$, which can be written as:

$$X^{n-k} m(X) = q(X) g(X) + p(X)$$

Where $q(X)$ is the quotient polynomial and $p(X)$ is the remainder polynomial with degree of $2t-1$ or less. This is the same as in the binary case where the remainder is the parity which can be written as:

$$P(X) = X^{n-k} m(X) \text{ mod } g(X)$$

So, the resulting codeword polynomial $U(X)$ can be written as:

$$U(X) = P(X) + X^{n-k} m(X)$$

3.2. Reed-Solomon decoding

Let

$v(X) = v_0 + v_1X + \dots + v_{n-1}X^{n-1}$ is the transmitted code vector and let

$r(X) = r_0 + r_1X + \dots + r_{n-1}X^{n-1}$ is the received vector

This assumes that, during transmission, the codeword becomes corrupted so that some symbols are received in error (this number of errors corresponds to the maximum error-correcting capability of the code). Then, the error pattern $e(X)$ caused by the channel noise during the transmitter can be written as:

$$e(X) = r(X) - v(X)$$

$$e(X) = e_0 + e_1X + \dots + e_{n-1}X^{n-1} = \sum_{i=0}^{n-1} e_iX^i$$

In binary BCH code decoding, the decoder only needs to find the error locations. However, the nonbinary BCH codes and R-S codes symbols require not only the identification of the error locations, but also determining the correct symbol values at those locations. So, the decoder will calculate the syndrome from the received vector $r(X)$, which is the first step of decoding R-S process. The syndrome S with $2t$ components can be defined as the following:

$$S_i = r(\alpha^i)$$

$$S_i = r_0 + r_1\alpha^i + \dots + r_{n-1}(\alpha^i)^{(n-1)} \quad \text{for } i = 1, 2, \dots, 2t$$

$$S_i = v(\alpha^i) + e(\alpha^i) \quad \text{from the above equations for } i = 1, 2, \dots, 2t$$

The roots of the code polynomial $v(X)$ are $\alpha, \alpha^2, \dots, \alpha^{2t}$

Therefore, $v(\alpha^i) = 0$ for $i = 1, 2, \dots, 2t$ which makes the syndrome equation written as:

$$S_i = e(\alpha^i) \quad \text{for } i = 1, 2, \dots, 2t$$

Assuming $e(X)$ is an error pattern of z errors then,

$$e(X) = X^{j1} + X^{j2} + \dots + X^{jz} \quad (2.1)$$

And $2t$ syndrome equations have been formed as following:

$$\begin{aligned} S_1 &= \alpha^{j_1} + \alpha^{j_2} + \dots + \alpha^{j_z} \\ S_2 &= (\alpha^{j_1})^2 + (\alpha^{j_2})^2 + \dots + (\alpha^{j_z})^2 \\ &\vdots \\ S_{2t} &= (\alpha^{j_1})^{2t} + (\alpha^{j_2})^{2t} + \dots + (\alpha^{j_z})^{2t} \end{aligned}$$

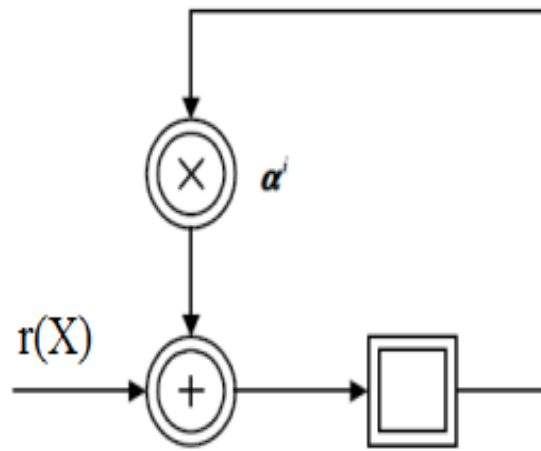
Any error correction process is a method of solving the above equations and finding $\alpha^{j_1}, \alpha^{j_2}, \dots, \alpha^{j_z}$ then the powers j_1, j_2, \dots, j_z which are the error locations in $e(X)$ at eq.(2.1). These above equations have many finite possible solutions and each one will yield a different error pattern. When the number of errors in the actual error pattern $e(X)$ is t or less ($z \leq t$), then the solution with the smallest number of errors is the correct solution and the error pattern corresponding to that solution will be the actual error pattern. On the other hand, for large t , solving the above equations will be more difficult and ineffective. However, determining α^{j_l} from the syndrome components S_i is done as the following:

Let

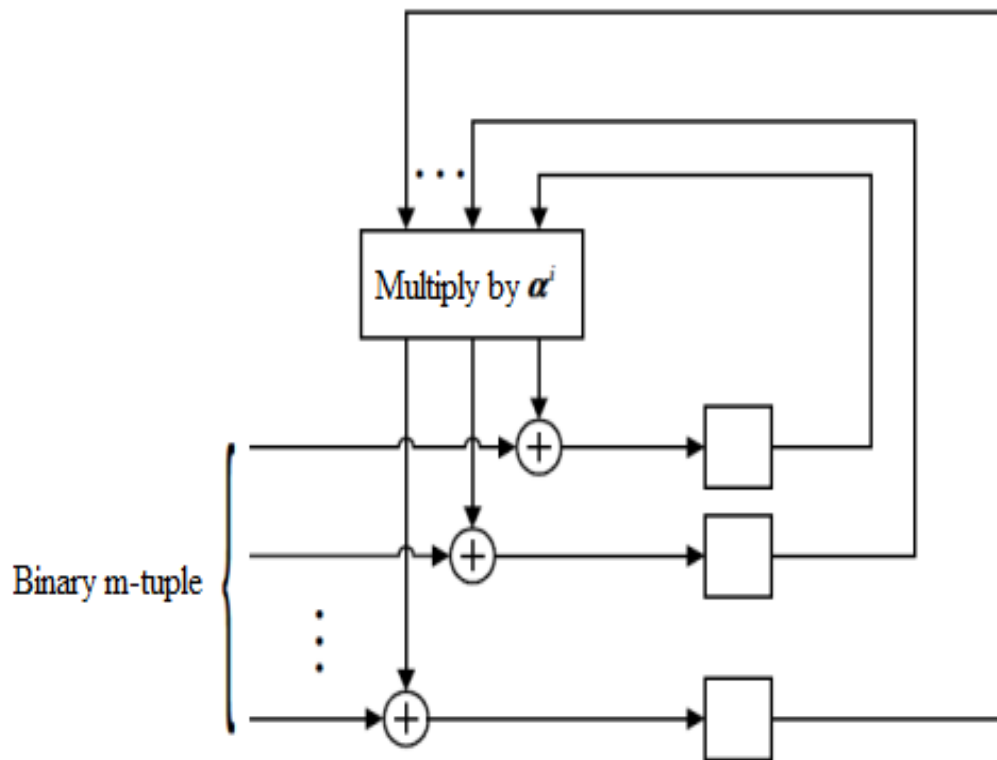
$$\beta_l = \alpha^{j_l} \quad \text{for } 1 \leq l \leq z$$

Be the error location numbers which means the $2t$ syndrome equations above are symmetric functions in $\beta_1, \beta_2, \dots, \beta_z$. This is also known as power sum symmetric functions and can be written as:

$$\begin{aligned} S_1 &= \beta_1 + \beta_2 + \dots + \beta_z \\ S_2 &= \beta_1^2 + \beta_2^2 + \dots + \beta_z^2 \\ &\vdots \\ S_{2t} &= \beta_1^{2t} + \beta_2^{2t} + \dots + \beta_z^{2t} \end{aligned}$$



(a)



(b)

Figure 3.2. Syndrome computation for Reed Solomon code in
(a)- $GF(2^m)$ and (b)- binary representation form

Next, the error location polynomial $\sigma(X)$ is defined as the following:

$$\sigma(X) = (1 + \beta_1 X)(1 + \beta_2 X) \dots (1 + \beta_z X)$$

$$\sigma(X) = \sigma_0 + \sigma_1 X + \dots + \sigma_z X^z$$

Where

$$\sigma_1 = 1$$

$$\sigma_2 = \beta_1 + \beta_2 + \dots + \beta_z$$

.

.

.

$$\sigma_z = \beta_1 \beta_2 \dots \beta_z$$

The roots of $\sigma(X)$ are the inverses of error location numbers: $\beta_1^{-1}, \beta_2^{-1}, \dots, \beta_z^{-1}$, so the

σ_i 's are elementary symmetric functions of β_z 's can be related to the syndrome components S_i 's

by the following Newton identities:

$$S_1 + \sigma_1 = 0$$

$$S_2 + \sigma_1 S_1 + 2\sigma_2 = 0$$

.

.

.

$$S_{z+1} + \sigma_1 S_z + \dots + \sigma_{z-1} S_2 + \sigma_z S_1 = 0$$

At this point, the error correcting procedure for R-S codes decoder has these four steps:

1. Compute the syndrome $S = S_1, S_2, \dots, S_{2t}$ from the received polynomial $r(X)$.
2. Determine the error location polynomial $\sigma(X)$ from the syndrome components.
3. Determine the error location numbers $\beta_1, \beta_2, \dots, \beta_{2t}$ by finding the roots of $\sigma(X)$.
4. Substitute the error location numbers into the error polynomials and solve for the corresponding error values e_{zi} . Knowledge of the values of β_i and e_{zi} is sufficient for error correction.

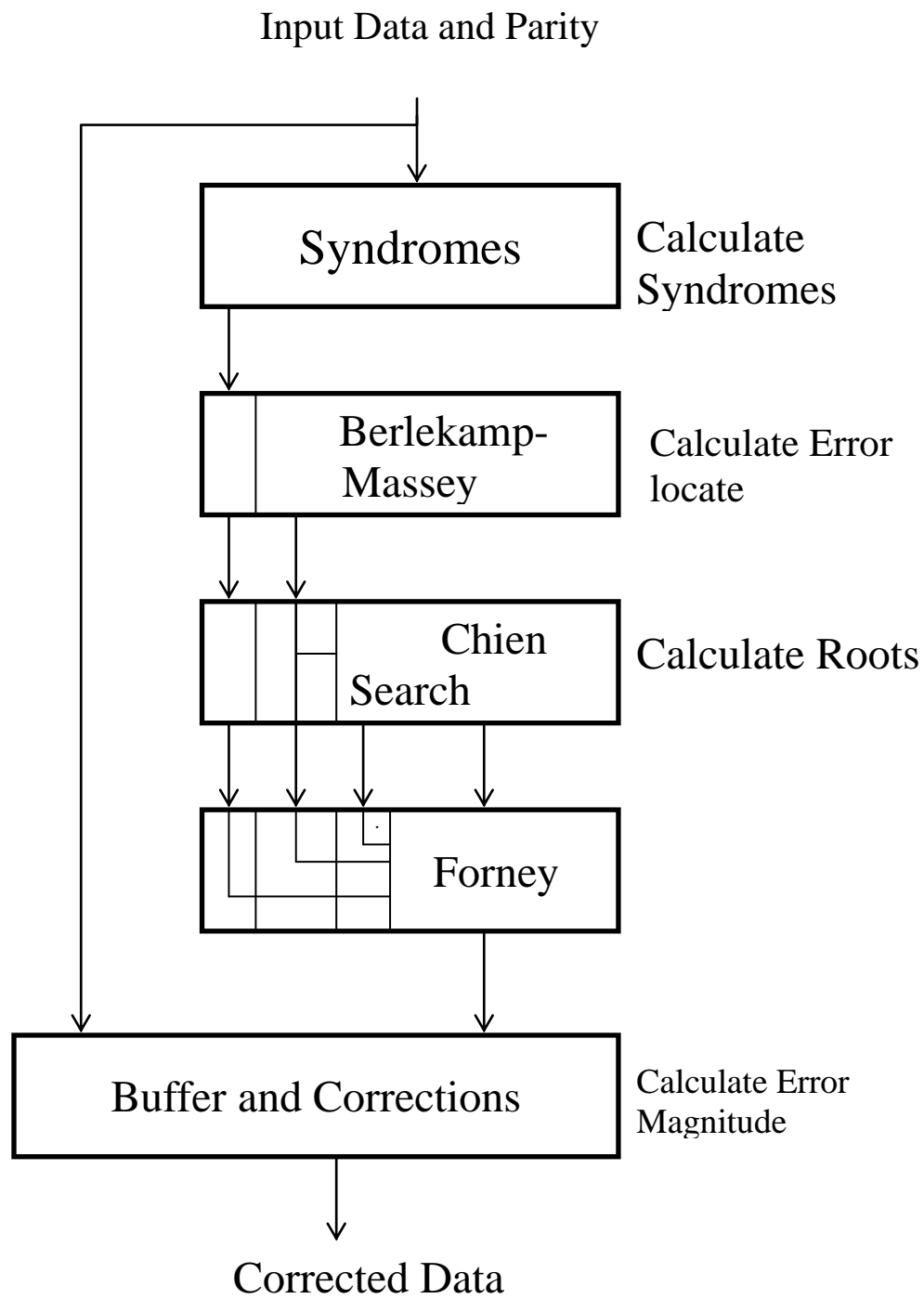


Figure 3.3 Reed Solomon decoder block diagram

Chapter 4

System specification

4.1. System specification

The system will have (255, 223) Reed Solomon code where the symbols of the code are elements of $GF(2^8)$ and have the following characteristics:

Degree of the polynomial $m = 8$

Block length $n = 2^8 - 1 = 255$

Number of parity check digits $n - k = 2t = 32$

Error correction capability $t = 16$

However, each symbol is represented by eight binary digits or one byte. Also, each data block contains 223 information symbols. This code is capable of correcting up to sixteen short burst errors of one byte or any burst error combination of up to a total length of eight bytes, providing that they only affect a maximum of sixteen individual symbols.

The elements of $GF(2^8)$ are generated by primitive polynomial of degree 8:

$$p(X) = 1 + X^2 + X^3 + X^4 + X^8$$

let α be the primitive element in $GF(2^8)$ and the root of $p(X)$

then $p(X) = 1 + X^2 + X^3 + X^4 + X^8 = 0$

or

$$p(X) = 1 + X^2 + X^3 + X^4 = X^8$$

So, the elements can be represented in an 8-tuple with 8 components being 0 or 1 and represent code word. The zero element of $GF(2^8)$ appears as an all zero 8-tuple. Matlab is used to represent nonzero elements as shown in Appendix A.

Also, if α is a primitive element in $GF(2^m)$, then the root of $p(x)$ is only the first thirty-two powers of α and are the roots of the generator polynomial. Meanwhile, the generator polynomial for (255, 223) code is given by:

$$\begin{aligned}
g(X) &= (X + \alpha)(X + \alpha^2)(X + \alpha^3)(X + \alpha^4)(X + \alpha^5)(X + \alpha^6)(X + \alpha^7)(X + \alpha^8)(X + \alpha^9)(X + \alpha^{10}) \\
&\quad (X + \alpha^{11})(X + \alpha^{12})(X + \alpha^{13})(X + \alpha^{14})(X + \alpha^{15})(X + \alpha^{16})(X + \alpha^{17})(X + \alpha^{18})(X + \alpha^{19}) \\
&\quad (X + \alpha^{20})(X + \alpha^{21})(X + \alpha^{22})(X + \alpha^{23})(X + \alpha^{24})(X + \alpha^{25})(X + \alpha^{26})(X + \alpha^{27})(X + \alpha^{28}) \\
&\quad (X + \alpha^{29})(X + \alpha^{30})(X + \alpha^{31})(X + \alpha^{32}) \\
g(X) &= 45 + 216X + 239X^2 + 24X^3 + 253X^4 + 104X^5 + 27X^6 + 40X^7 + 107X^8 + 50X^9 + 163X^{10} + 210X^{11} \\
&\quad + 227X^{12} + 134X^{13} + 224X^{14} + 158X^{15} + 119X^{16} + 13X^{17} + 158X^{18} + X^{19} + 238X^{20} + 164X^{21} \\
&\quad + 82X^{22} + 43X^{23} + 15X^{24} + 232X^{25} + 246X^{26} + 142X^{27} + 50X^{28} + 189X^{29} + 29X^{30} + 232X^{31} + X^{32}
\end{aligned}$$

therefore, the coefficients of $g(x)$ used in the encoder multiplication are :

$$\begin{aligned}
g_0 &= \alpha^{18} = 45, & g_1 &= \alpha^{251} = 216, & g_2 &= \alpha^{215} = 239, & g_3 &= \alpha^{28} = 24, & g_4 &= \alpha^{80} = 253, \\
g_5 &= \alpha^{107} = 104, & g_6 &= \alpha^{248} = 27, & g_7 &= \alpha^{53} = 40, & g_8 &= \alpha^{84} = 107, & g_9 &= \alpha^{194} = 50, \\
g_{10} &= \alpha^{91} = 163, & g_{11} &= \alpha^{59} = 210, & g_{12} &= \alpha^{176} = 227, & g_{13} &= \alpha^{99} = 134, & g_{14} &= \alpha^{203} = 224, \\
g_{15} &= \alpha^{137} = 158, & g_{16} &= \alpha^{43} = 119, & g_{17} &= \alpha^{104} = 13, & g_{18} &= \alpha^{137} = 158, & g_{19} &= \alpha^0 = 1, \\
g_{20} &= \alpha^{44} = 238, & g_{21} &= \alpha^{149} = 164, & g_{22} &= \alpha^{148} = 82, & g_{23} &= \alpha^{218} = 43, & g_{24} &= \alpha^{75} = 15, \\
g_{25} &= \alpha^{11} = 232, & g_{26} &= \alpha^{173} = 246, & g_{27} &= \alpha^{254} = 142, & g_{28} &= \alpha^{194} = 50, & g_{29} &= \alpha^{109} = 189, \\
g_{30} &= \alpha^8 = 29, & g_{31} &= \alpha^{11} = 232, & g_{32} &= 1
\end{aligned}$$

At the same time, a Reed Solomon code with smaller dimensions can be chosen. (255,223) code has a suitable symbol size that matches the 8-bit data byte commonly used today. Additionally, the length of this code of 255 symbols can be shortend to any length so that it can match any system specification without any major change in the encoder/decoder hardware circuitry.

4.2.Encoder architecture

Encoders are designed as feedback shift register, Figure 3.1. The data message blocks of 223 symbols shift sequentially as an input to the encoder and when the last message symbol is loaded, the feedback register contains the thirty-two parity check symbols. These symbols will then be shifted out following the 223 information symbols to generate a code word of 255 symbols as an output of the encoder.

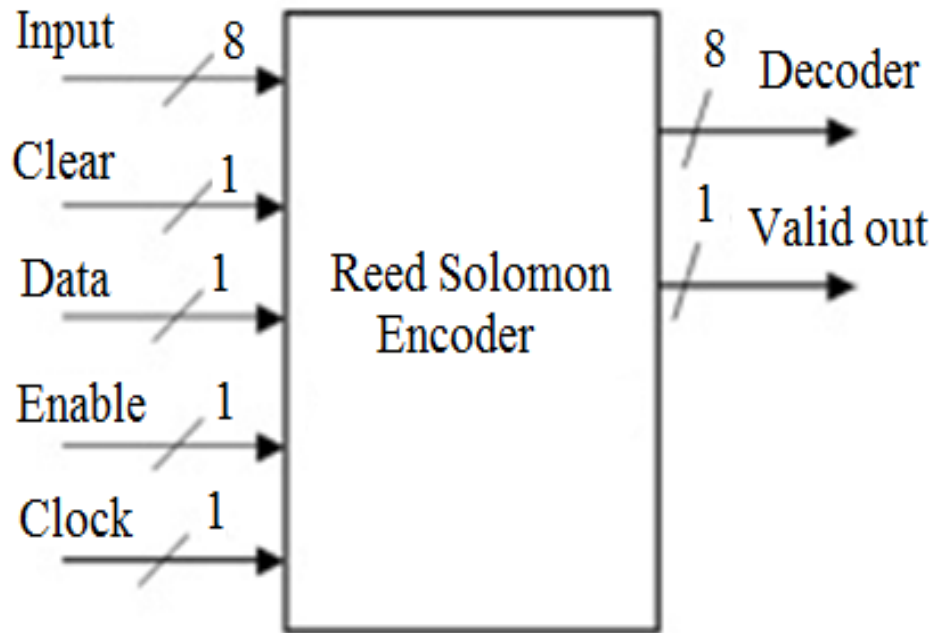


Figure 4.1 Reed Solomon encoder block diagram

4.3.Decoder architecture

The decoder operation is carried out as follows:

1- Compute the syndrome $S = (S_1, S_2, \dots, S_{2t})$.

In this step, (255,223) Reed Solomon code syndrome calculation can be done as follows:

Let $r(\alpha)$ be the received code polynomial. Then,

$$r(\alpha) = S_1 = e_{j1}\alpha^{j1} + e_{j2}\alpha^{j2} + \dots + e_{jz}\alpha^{jz}$$

$$r(\alpha^2) = S_2 = e_{j1}(\alpha^{j1})^2 + e_{j2}(\alpha^{j2})^2 + \dots + e_{jz}(\alpha^{jz})^2$$

·
·
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$$r(\alpha^{32}) = S_{32} = e_{j1}(\alpha^{j1})^{32} + e_{j2}(\alpha^{j2})^{32} + \dots + e_{jz}(\alpha^{jz})^{32}$$

All the received data block will shift into the syndrome processor using the same circuit to compute the thirty-two syndrome components S_1, S_2, \dots, S_{32} which will be done in sequence form. The only difference is the multiplication values in the feedback loop for each syndrome.

2- Determine the error location polynomial $\sigma(X)$.

In this step, Berlekamp's iterative algorithm method (Table 3.1.) has been used, without being given any proof, to determine the coefficients of the error location polynomial $\sigma(X)$. There are $(n-k) = 2t = (255-223) = 32$ iterations at step μ to be performed and the algorithm will yield the minimum degree polynomial that will satisfy the first μ Newton equations.

μ	$\sigma^{(\mu)}(X)$	d_μ	1_μ	$\mu - 1_\mu$
-1	1	1	0	-1
0	1	S_1	0	0
1				
2				
.				
.				
2t				

At any iteration step, the entries in the table start with evaluating the discrepancy value d_μ from the previous row and:

If $d_\mu = 0$, we set $\sigma^{(\mu+1)}(X) = \sigma^{(\mu)}(X)$, which is the entered as in the previous step.

If $d_\mu \neq 0$, a correction factor must be added. So,

$$\sigma^{(\mu+1)}(X) = \sigma^{(\mu)}(X) + d_\mu d_\rho^{-1} X^{(\mu-\rho)} \sigma^{(\rho)}(X)$$

Where ρ is a previous iteration step such that $d_\rho \neq 0$, and $\rho - 1_\rho$ in the last column of the table have the largest value. The next entry in the last two columns are:

$$1_{\mu+1} = \max [1_\mu, 1_\rho + \mu - \rho]$$

The last entry for the current iteration step involves computing the new discrepancy value as follows:

$$d_{\mu+1} = S_{\mu+2} + \sigma_1^{\mu+1} S_{\mu+1} + \dots + \sigma_{1_{\mu+1}}^{(\mu+1)} S_{\mu+2-1_{\mu+1}}$$

After the last step in the iteration process, which is the last row in this case or step number thirty two, the $\sigma^{(\mu)}(X)$ column has the correct error location polynomial if there are less than sixteen errors in the data block.

3- Determine the error value evaluator.

In this step, the roots of error location can be found by Chien's search, which initially starts by loading the register with the coefficients of the error location polynomial $\sigma(X)$. Then, at each clock pulse, every register is multiplied by the feedback loop value. After each clock, all the registers are XOR-ed together with unity element of $GF(2^8)$; if the XOR value is zero, then the root has been found. The search is run for 256 clock cycles and $t+1 = 17$ stages in the search which each consists of multiplexer, Adder (XOR), and a register. Meanwhile, the feedback values have been chosen to be powers of the primitive element α , such that in 254 trials, all nonzero elements of $GF(2^8)$ are checked as the roots of the error location polynomial. The resulting finding of a root is saved as the inverses of the error location numbers.

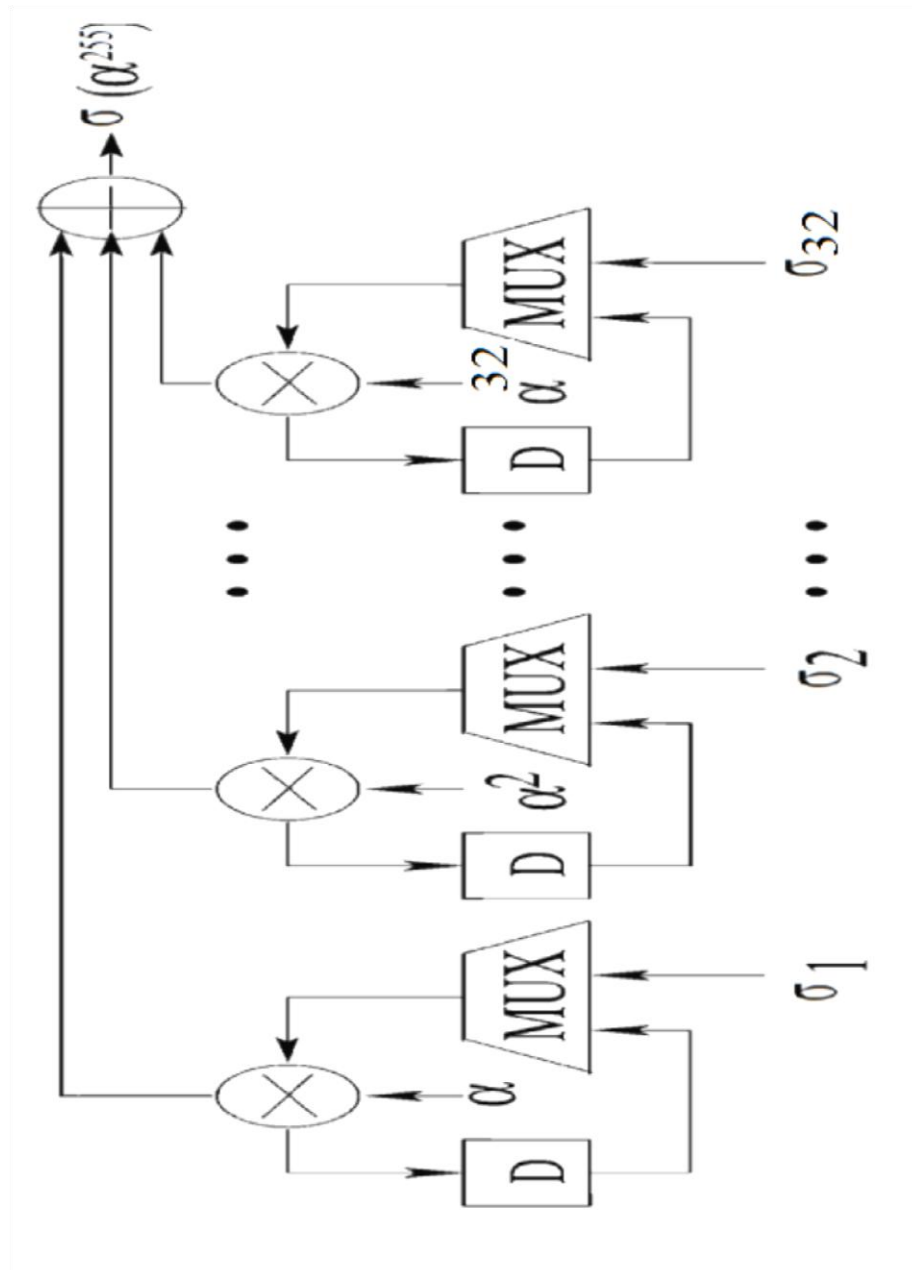


Figure 4.2 Chien's search architectures for (255, 223) Reed Solomon code

4- Evaluate error location numbers and error values and perform error correction.

In this step, the following polynomial is constructed:

$$Z(X) = 1 + (S_1 + \sigma_1)X + (S_2 + \sigma_1 S_1 + \sigma_2)X^2 + \dots + (S_z + \sigma_1 S_{z-1} + \dots + \sigma_{z-1} S_1 + \sigma_z) X^z$$

Then, the error value for the location given by $\beta_l = \alpha^{jl}$ is:

$$e_{j1} = \frac{Z(\beta_1^{n-1})}{\prod_{\substack{i=1 \\ i \neq 1}}^z (1 + \beta_i \beta_1^{n-1})}$$

The Chien search calculates the error value and adds it to the received codeword to make the correction. This step can be implemented in a straightforward form.

Chapter 5

The hardware implementation encoder

5.1.The hardware implementation encoder

Encoding RS- code can be done as follows. Let,

$$m(X) = m_0 + m_1X + m_2X^2 + \dots + m_{k-1}X^{k-1}$$

be the message polynomial. In systematic form, the $2t$ parity check symbols are the coefficients of the remainder $b(X) = b_0 + b_1X + \dots + b_{2t-1}X^{2t-1}$ resulting from dividing the message polynomial $X^{2t}m(X)$ by the generator polynomial:

$$b(X) + X^{2t} m(X) = a(X) g(X)$$

Where:

$g(X)$ is the generator polynomial

$a(X)$ is the quotient from dividing the message polynomial by the generator polynomial

$b(X)$ is the remainder

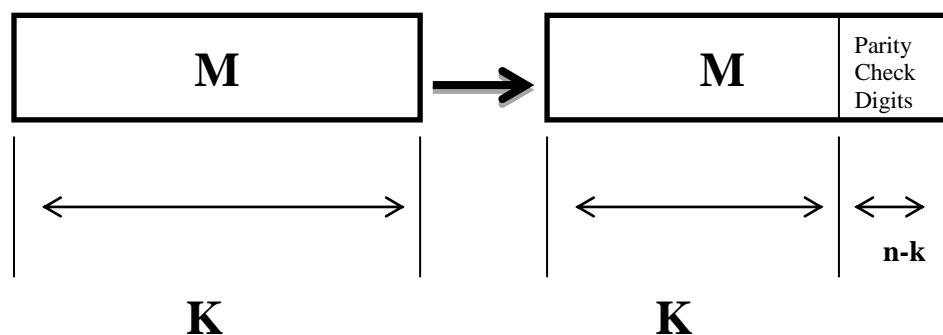


Figure 5.1 Systematic form

In hardware implementation for (255,223) R-S code, the generator polynomial $g(X)$ is:

$$g(X) = 45 + 216X + 239X^2 + 24X^3 + 253X^4 + 104X^5 + 27X^6 + 40X^7 + 107X^8 + 50X^9 + 163X^{10} + 210X^{11} \\ + 227X^{12} + 134X^{13} + 224X^{14} + 158X^{15} + 119X^{16} + 13X^{17} + 158X^{18} + X^{19} + 238X^{20} + 164X^{21} \\ + 82X^{22} + 43X^{23} + 15X^{24} + 232X^{25} + 246X^{26} + 142X^{27} + 50X^{28} + 189X^{29} + 29X^{30} + 232X^{31} + X^{32}$$

and the code polynomial $v(X)$ of $m(X)$ is:

$$v(X) = b(X) + X^{32} m(X) = a(X) g(X)$$

$$v(X) = b_0 + b_1X + \dots + b_{31}X^{31} + m_0X^{32} + m_1X^{33} + \dots + m_{222}X^{254}$$

The encoder used for this Reed Solomon code has the following blocks:

- Thirty-two 8-bit wide adder that adds two field elements from $GF(2^8)$.
- Thirty-two multiplier that multiplies any field elements of $G(2^8)$ by a fixed elements from same field.
- Thirty-two 8-bit registers to store a field element from the $G(2^8)$.
- One 8-bit wild multiplexer.
- One 8-bit wild AND gate.

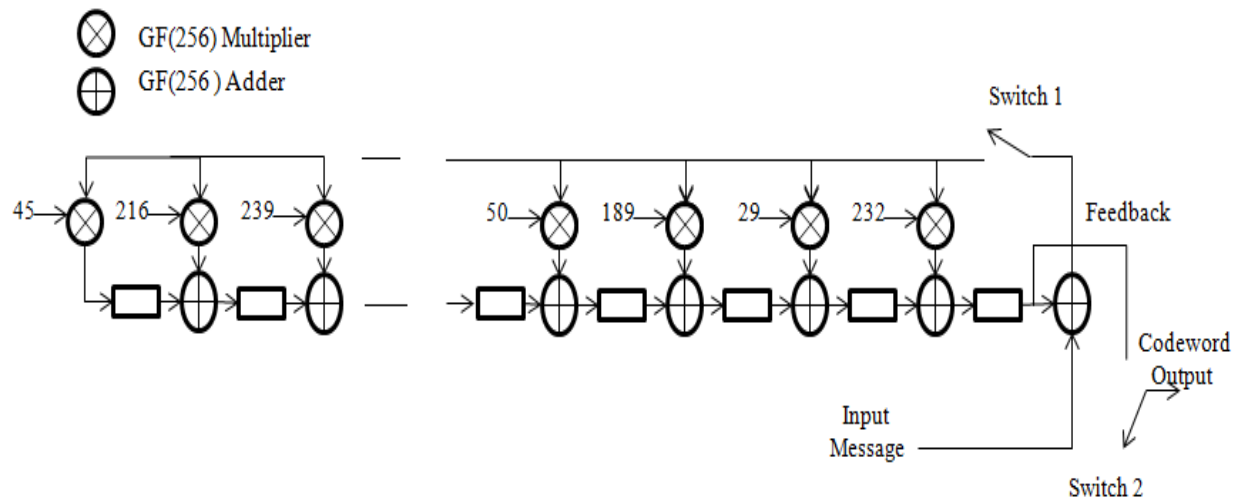


Figure 5.2 feedback shift register configuration Reed Solomon (255,223) encoder

The encoder operation is carried out as follows:

- 1- AND gate ON, so the 223 information digits are shifted into the shift register and simultaneously into the communication channel. As soon as the complete message has entered the shift register, the thirty-two digits in the register form the parity check digits.
- 2- AND gate OFF, so the feedback connection is disable.

The parity check digits are shifted out and sent into the channel so that the thirty two parity check digits together with the 223 information digits form a complete codeword. There will be 223 clock pulses to shift the information message into the shift register and the encoder output. Also, 32 clock-pulses shift the thirty-two parity check digits to the output, which means that there are 255 clock pulses in total.

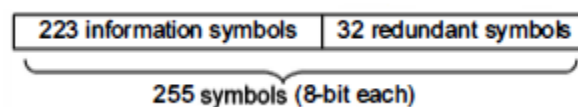


Figure 5.3 Codeword generated by Reed Solomon encoder

The encoder has three blocks, an input register, a processor register and an output register, as shown in Figure 5.4 . The encoder circuit is a single high-performance reduced instruction set computing RISC processor; as an example, 8-bit CMOS EPROM-ROM microcontroller PIC17C44 from Microchip Technology. This chip has a Harvard architecture as well as these features: a maximum frequency of operation of 33 MHz, the instruction cycle consists of four oscillator periods or the normal instruction execution time is 160 ns, and a transfer rate more than 100 Kbyte/sec on the communication channel .

The encoder transfers the input bytes of the message as well as computes the thirty two bytes parity check and, to have minimum bandwidth for the transmission channel, the output data rate has to be constant. As a result, the encoder is designed to be interrupt-driven and an external clock generates each interrupt. To speed up the system, the thirty two parity digits are implemented with one level ROM multiplication. For every multiplication, a table with 256 entries are calculated by multiplying any elements of $GF(2^8)$ with the coefficients of the generator polynomial $g(X)$. The thirty two tables are stored in the program memory and then read through the table instructions of PIC17C44. The entries of the table are calculated in Matlab as well as C and shown in Appendix B. Figure 5.3 shows the flowcharts of the software implementation for the main module in the encoder, which only performs initializing the special purpose register, setting up the timer and performing the interrupt. Next, the processor will wait at a certain address until the external clock generates an interrupt. As the interrupt occurs, the processor will read the input then send the data byte to the output. After that, it will update all the thirty two parity check digits. The next step is to check if all the 223 bytes read and, if yes, shift out the thirty two parity bytes at the same rate as other bytes. If not, then return from the interrupt subroutine and wait for the next input.

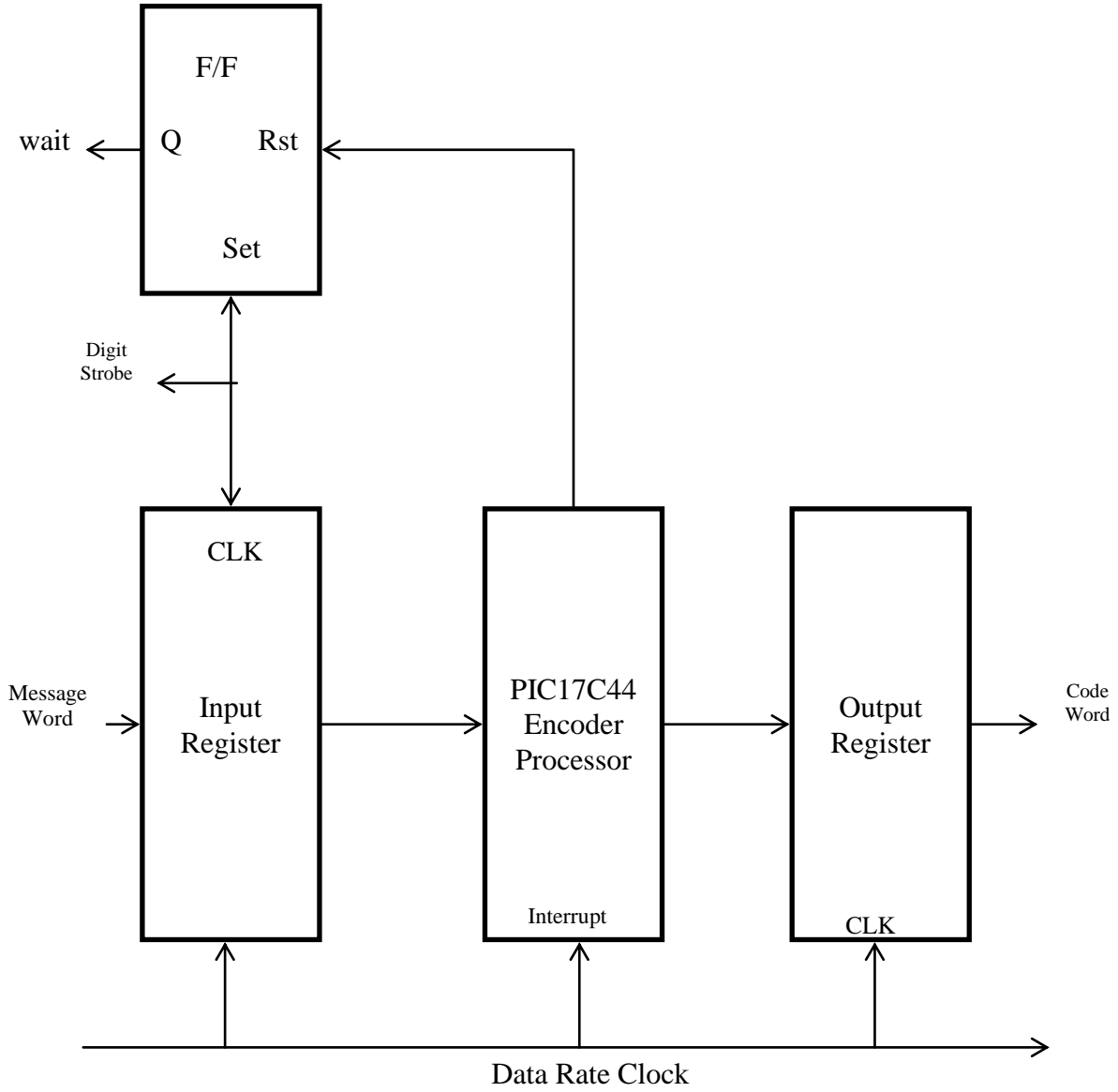


Figure 5.4 Encoder Block Diagram

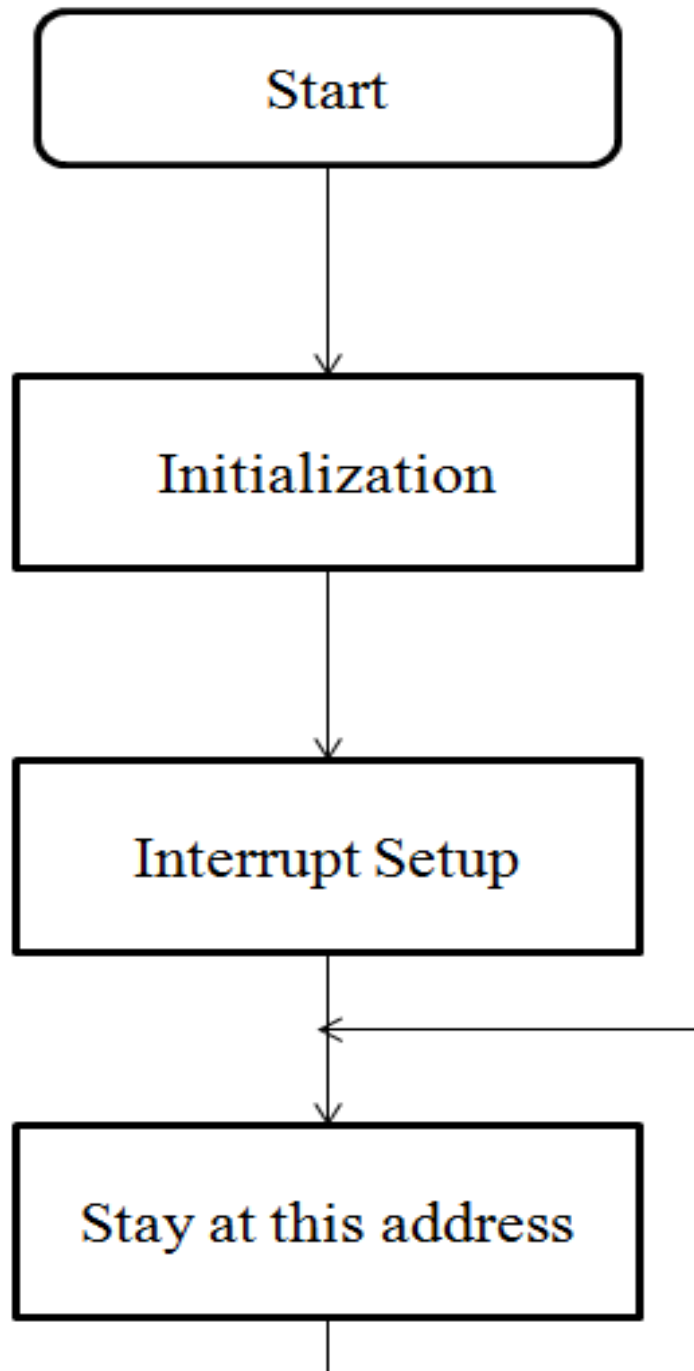


Figure 5.6 Encoder main module flowchart

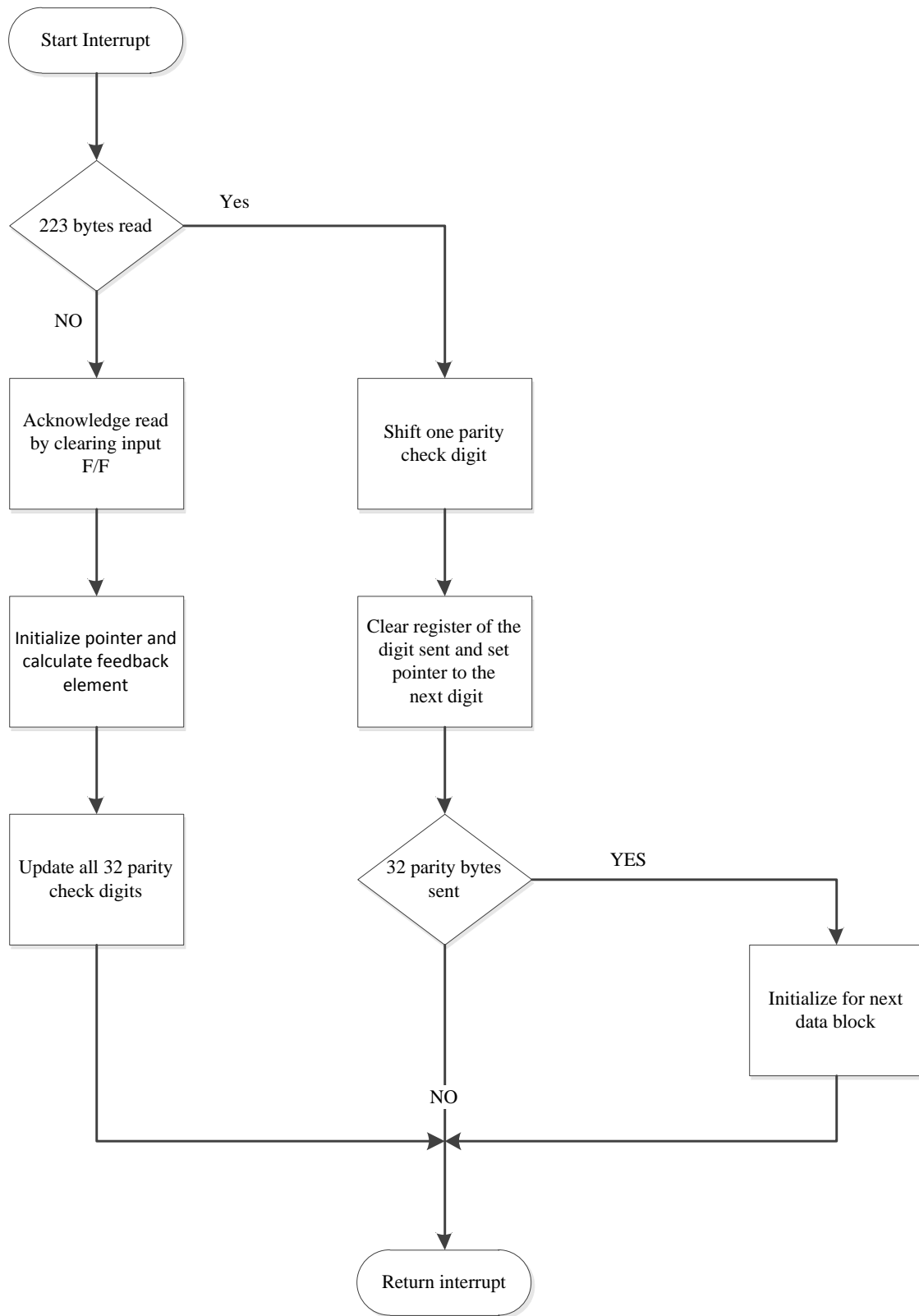


Figure 5.7 Encoder interrupt subroutine flowchart

Chapter 6

The hardware implementation decoder

6.1. The hardware implementation decoder

The decoder is more complex than the encoder due to the multiple stages of the decoding procedure outlined in chapter four.

The hardware implementation is done by using three single high-performance RISC microcontrollers. Three PIC17C44 microcontroller chips work in parallel for the decoder as shown in Figure 6.1. The first data block will be entered into the first microcontroller one byte at a time and saved in the internal RAM register. After the last byte is entered, the second microcontroller will take the next 255 bytes and then the third microcontroller will follow the same process. This process will give the first microcontroller the time to compute the error location, error location number, error value, and output the correct 223 byte message.

The syndrome computation is the stage in the decoder that will use more time as well as finding the roots of the error location polynomial. Therefore, sending out the entire syndrome in parallel will be the best. Since the data rate is uniform in using the channel effectively, it will give minimum bandwidth in the communication channel as well as a lower throughput. The flowcharts below show the decoding process outlined in chapter four. The code is presented in outline form to make it easier for the reader.

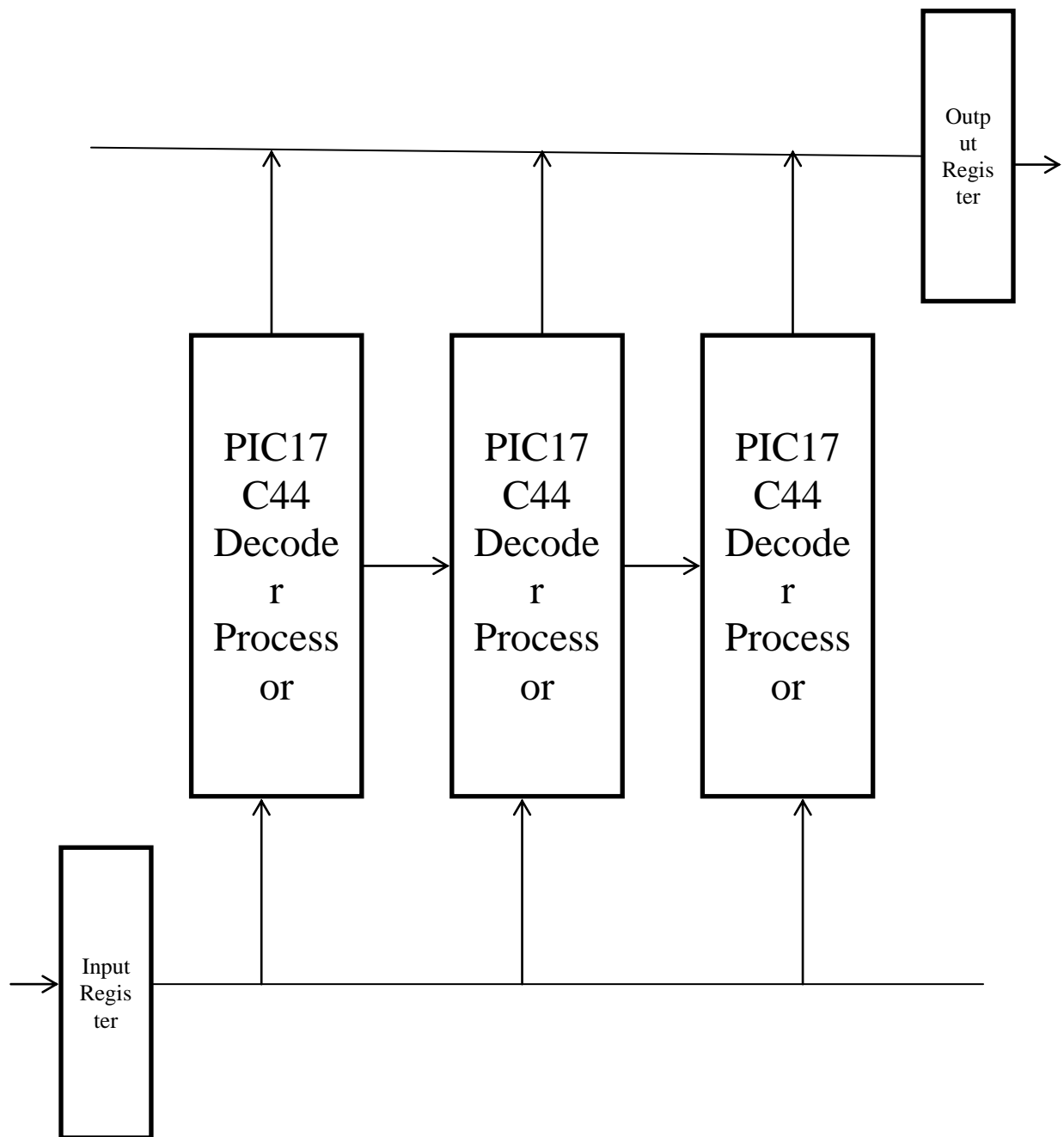


Figure 6.1 Decoder block diagram

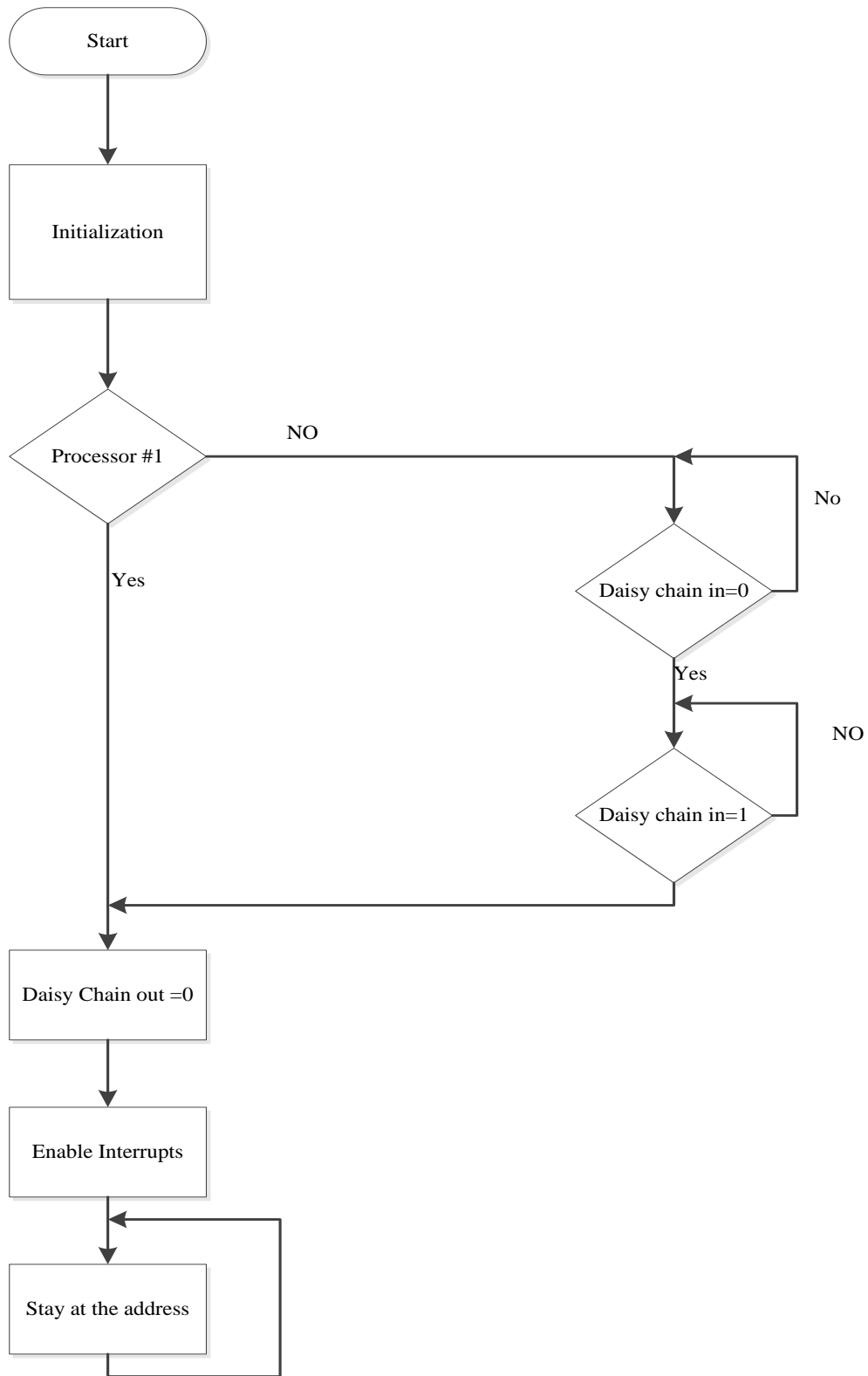


Figure 6.2 Decoder main module flowcharts

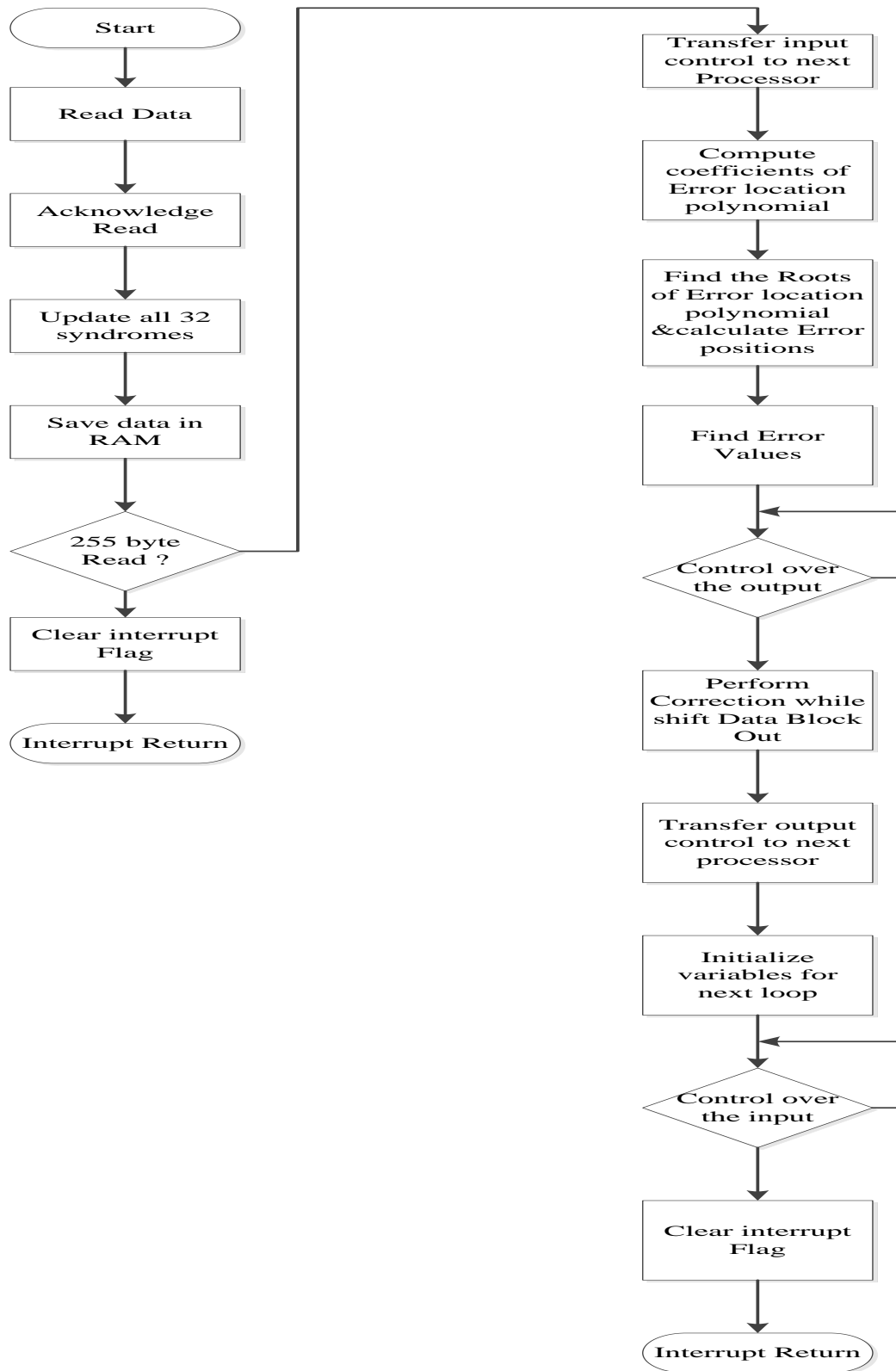


Figure 6.3 Decoder interrupt subroutine

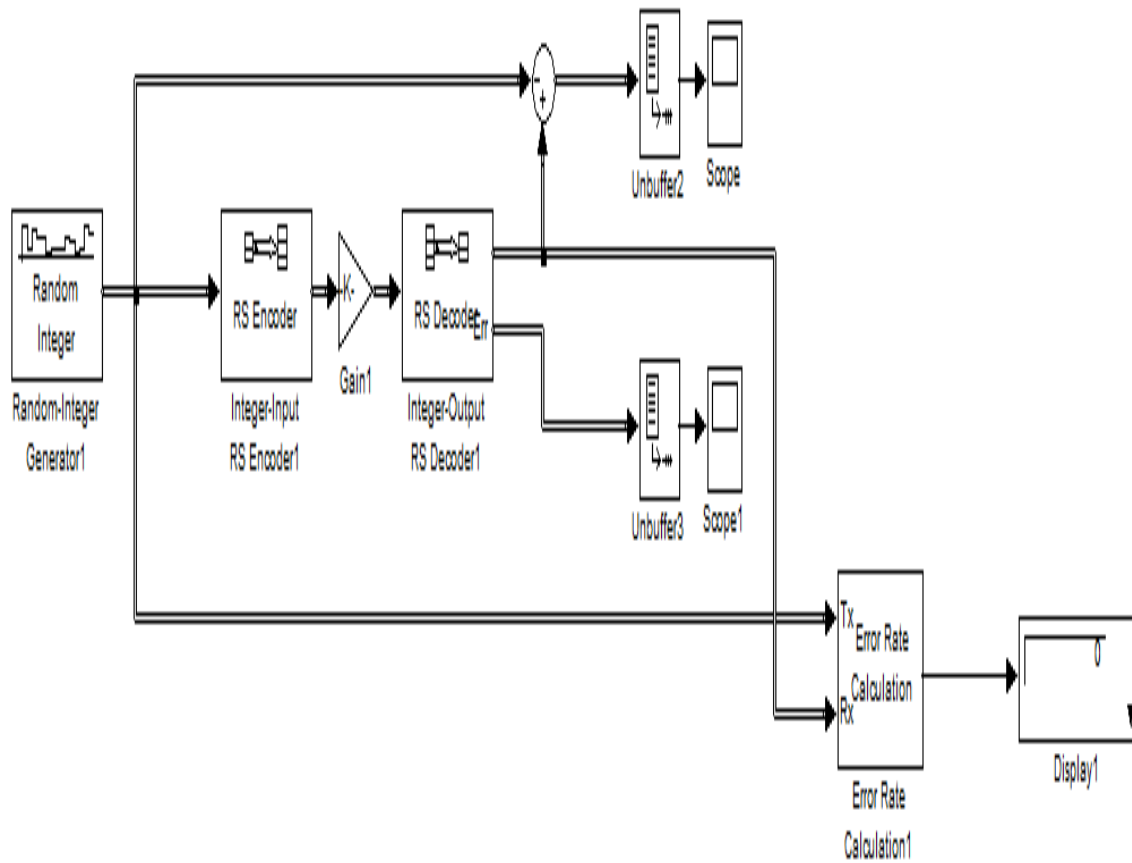


Figure 7.1 (255, 223) Reed Solomon Code in Simulink

After running the design first scope, Figure 7.2 gives the difference between the original message and the recovered message. Since the decoder was able to correct all errors that occurred, each of the sixteen data streams in the plot is zero as shown in Figure 7.2 and the error rate calculation between the input and the output message is zero as shown in Figure 7.3.

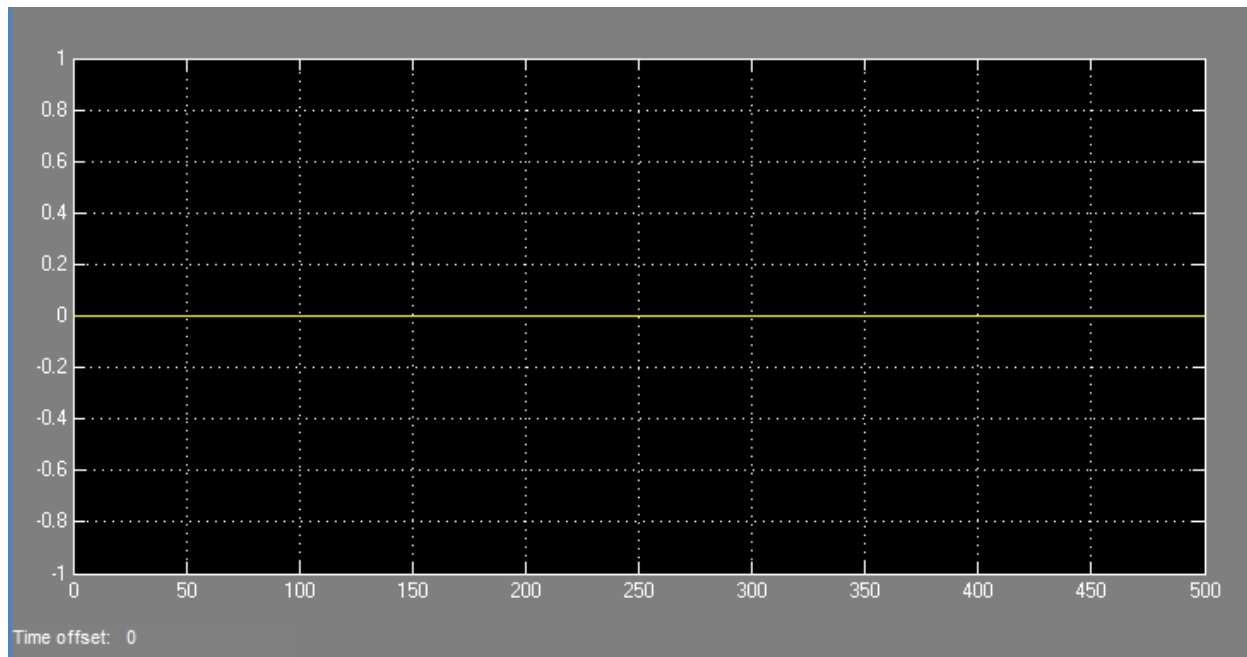


Figure 7.2 Output of scope 1

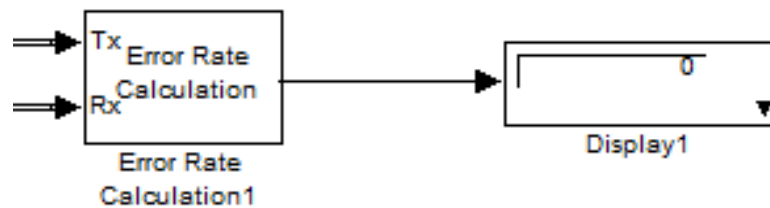


Figure 7.3 Error rate between the input and output message no error

The second plot Figure 7.4 gives the number of errors that the decoder detected while trying to recover the message. Often the number is thirty two because the gain replaces the thirty two first symbols in each codeword with zeros. However, the number of errors is less than thirty two whenever a correct codeword contains one or more zeros in the first thirty two places.

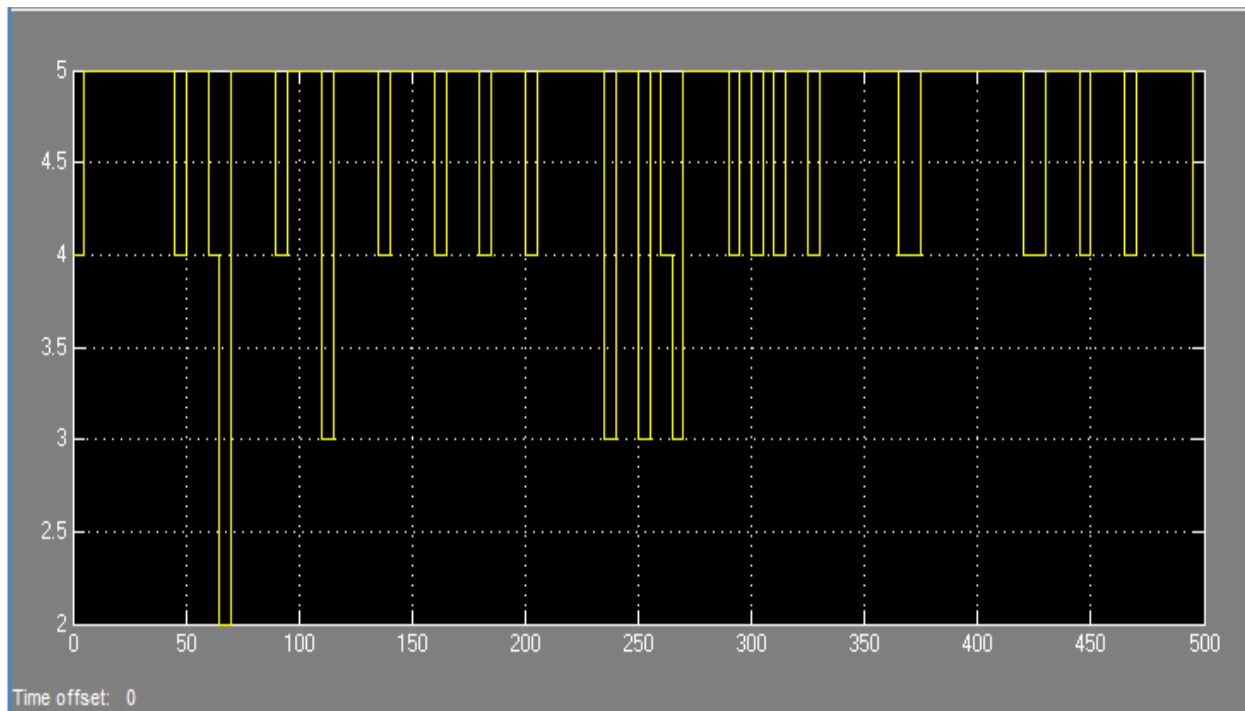


Figure 7.4 Output of scope 2

On the other hand, the decoder cannot correct all the errors in the codeword; this may be because the error received in the message happens to match a codeword. This will cause the decoder to clear as no error occur or the decoder cannot correct all the errors on the message. Figure 7.5 shows the error rate between the input and the output when the decoder is not able to correct all the errors.

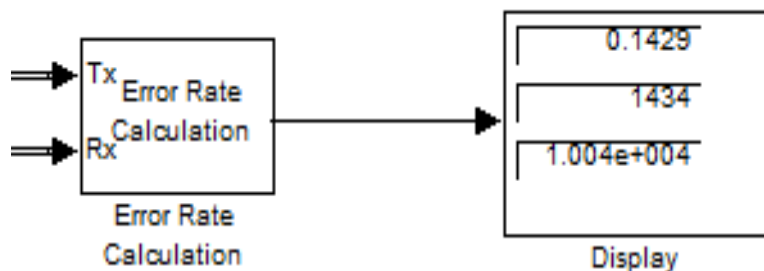


Figure 7.5 Error rate between the input and output message error occur

7.2.1. Using C program source code

This program is an encoder/decoder for (255, 223) Reed-Solomon codes. Encoding is in systematic form, decoding via the Berlekamp iterative algorithm. In the present form, the constants $mm = 8$, $nn = 255$, $tt = 16$, and $kk = nn - 2tt = 223$ specified. Also, the irreducible polynomial used to generate $GF(2^{mm})$ has been entered, $p(X) = 1 + X^2 + X^3 + X^4 + X^8$. The representation of the elements of $GF(2^m)$ is either in index form, where the number is the power of the primitive element α , which is convenient for multiplication (add the powers modulo $2^m - 1$) or in polynomial form, where the bits represent the coefficients of the polynomial representation of the number, which is the most convenient form for addition. The two forms are swapped via lookup tables. This leads to fairly messy-looking expressions, but unfortunately, there is no easy alternative when working with Galois arithmetic. The code is not written in the most elegant way [9]. However, when including it into a simulation program, you may want to do some conversion of global variables to local variables where appropriate, and passing parameters, e.g. array addresses, to the functions may be a sensible strategy to reduce the number of global variables and thus decrease the chance of a bug being introduced.

The program has four main boxes: generator polynomial, encoder, decoder, and the main. However, the decoder by itself has four boxes as shown below where the first is finding the syndrome, the second is finding the error location polynomial, the third is finding the roots of the error location polynomial and finally the four is to evaluate errors. The output of each box is shown in Appendix B and the total time to run the software is 887ms.

Void generate_gf ()

Generate GF(2**mm) from the irreducible polynomial p(X) in pp[0]..pp[mm]

Lookup tables: index->polynomial form alpha_to[] contains j=alpha**i;

polynomial form -> index form index_of [j=alpha**i] = i

alpha=2 is the primitive element of GF (2**mm)

```
#include <math.h>
#include <stdio.h>
#define mm 8                                /* RS code over GF(2**8) - change to suit */
#define nn 255                             /* nn=2**mm -1  length of codeword */
#define tt 16                              /* number of errors that can be corrected */
#define kk 223                             /* kk = nn-2*tt */
int pp [mm+1] = { 1, 0, 1, 1, 1, 0, 0, 0, 1 } ; /* specify irreducible polynomial coeffs */
int alpha_to [nn+1], index_of [nn+1], gg [nn-kk+1] ;
int recd [nn], data [kk], bb [nn-kk];
{
    register int i, mask ;
    printf("void generate_gf()");
    mask = 1 ;
    alpha_to[mm] = 0 ;
    for (i=0; i<mm ; i++)
    {
        alpha_to[i] = mask ;
        index_of[alpha_to[i]] = i ;
        printf("\n i=%3d mask=%3d mm=%3d alpha_to[mm]=%3d index_of[ alpha_to[i] ]=%3d",
               i, mask, mm, alpha_to[mm], index_of[alpha_to[i]] );
        if (pp[i]!=0)
            alpha_to[mm] ^= mask ;
        printf("\n mask=%3d pp[%3d]=%3d mm=%3d alpha_to[mm]=%3d ",
               mask, i, pp[i], mm, alpha_to[mm] );
        mask <<= 1 ;
    }
    index_of[alpha_to[mm]] = mm ;
    mask >>= 1 ;
    for (i=mm+1; i<nn; i++)
    {
        if (alpha_to[i-1] >= mask)
            alpha_to[i] = alpha_to[mm] ^ ((alpha_to[i-1]^mask)<<1) ;
        else alpha_to[i] = alpha_to[i-1]<<1 ;
        index_of[alpha_to[i]] = i ;
        printf("\nmask=%3d i=%3d alpha_to[i]=%3d index_of[ alpha_to[i] ]=%3d ",
               mask, i, alpha_to[i], index_of[alpha_to[i]] );
    }
    index_of[0] = -1 ;
}
```

Void gen_poly()

Obtain the generator polynomial of the tt-error correcting, length

nn=(2**mm -1) Reed Solomon code from the product of (X+alpha**i), i=1..2*tt

```
{
register int i,j ;
printf("\n void gen_poly()");
gg[0] = 2 ;          /* primitive element alpha = 2 for GF(2**mm) */
gg[1] = 1 ;          /* g(x) = (X+alpha) initially */
for (i=0; i<=nn-kk; i++)
    printf("\n alpha_to[%2d]=%3d",i,alpha_to[i]);
for (i=2; i<=nn-kk; i++)
{
    gg[i] = 1 ;
for (j=i-1; j>0; j--)
    if (gg[j] != 0) gg[j] = gg[j-1]^ alpha_to[(index_of[gg[j]]+i)%nn] ;
    else gg[j] = gg[j-1] ;
    gg[0] = alpha_to[(index_of[gg[0]]+i)%nn] ;          /* gg[0] can never be zero */
}
/* convert gg[] to index form for quicker encoding */
printf("\n convert gg[] to index form for quicker encoding");
for (i=0; i<=nn-kk; i++)
    printf("\nindex_of[%2d]=%3d",i,index_of[i]);
for (i=0; i<=nn-kk; i++)
{
    printf("\n gp gg[%2d ]=%3d ", i, gg[i]); /*display generator polynomial array */
    gg[i] = index_of[gg[i]] ;
}
}
```


Void encode_rs()

Take the string of symbols in data[i], i=0.....(k-1) and encode systematically to produce 2*tt parity symbols in bb[0].....bb[2*tt-1] data[] is input and bb[] is output in polynomial form.

Encoding is done by using a feedback shift register with appropriate connections specified by the elements of gg[], which was generated above. Codeword is $c(X) = data(X)*X^{nn-kk} + b(X)$

```
{
    register int i,j ;
    printf("\n void encode_rs()");
    int feedback ;
    for (i=0; i<nn-kk; i++)    bb[i] = 0 ;
    for (i=kk-1; i>=0; i--)
    {
        feedback = index_of[data[i]^bb[nn-kk-1]] ;
        if (feedback != -1)
        {
            for (j=nn-kk-1; j>0; j--)
                if (gg[j] != -1)
                    bb[j] = bb[j-1]^alpha_to[(gg[j]+feedback)%nn] ;
            else
                bb[j] = bb[j-1] ;
            bb[0] = alpha_to[(gg[0]+feedback)%nn] ;
        }
        else
        { for (j=nn-kk-1; j>0; j--)
            bb[j] = bb[j-1] ;
            bb[0] = 0 ;
        } ;
    } ;
};
```

Void decode_rs()

Assume we have received bits grouped into mm-bit symbols in recd[i], $i=0..(nn-1)$, and recd[i] is index form (ie as powers of alpha). We first compute the $2*tt$ syndromes by substituting α^{*i} into rec(X) and evaluating, storing the syndromes in s[i], $i=1.....2tt$ (leave s[0] zero). Then we use the Berlekamp iteration to find the error location polynomial elp[i]. If the degree of the elp is $>tt$, we cannot correct all the errors and just put out the information symbols uncorrected. If the degree of elp is $\leq tt$, we substitute α^{*i} , $i=1.....n$ into the elp to get the roots, the inverse roots, and the error location numbers. If the number of errors located does not equal the degree of the elp, we have more than tt errors and cannot correct them. Otherwise, we then solve for the error value at the error location and correct the error. For the cases where the number of errors is known to be too large to correct, the information symbols as received are output (the advantage of systematic encoding is that hopefully some of the information symbols will be okay and that the errors are in the parity part of the transmitted codeword). Of course, these insoluble cases can be returned as error flags to the calling routine if desired.

```
{
  register int i,j,u,q ;
  printf("\n void decode_rs()");
  int elp[nn-kk+2][nn-kk], d[nn-kk+2], l[nn-kk+2], u_lu[nn-kk+2], s[nn-kk+1] ;
  int count=0, syn_error=0, root[tt], loc[tt], z[tt+1], err[nn], reg[tt+1] ;

  /* first form the syndromes */
  for (i=1; i<=nn-kk; i++)
  { s[i] = 0 ;
    for (j=0; j<nn; j++)
      if (recd[j]!=-1)
        s[i] ^= alpha_to[(recd[j]+i*j)%nn] ;    /* recd[j] in index form */
  /* convert syndrome from polynomial form to index form */
    if (s[i]!=0) syn_error=1 ;                  /* set flag if non-zero syndrome => error */
    s[i] = index_of[s[i]] ;
  } ;
  if (syn_error)                                /* if errors, try and correct */
  {
```

Compute the error location polynomial via the Berlekamp iterative algorithm:

$d[u]$ is the ' μ 'th discrepancy, where $u = \mu + 1$ and ' μ ' (the Greek letter) is the step number ranging from -1 to $2 \cdot t$ (see 4.3), $l[u]$ is the degree of the elp at that step, and $u_l[u]$ is the difference between the step number and the degree of the elp.

```

/* initialise table entries */
d[0] = 0 ;          /* index form */
d[1] = s[1] ;       /* index form */
elp[0][0] = 0 ;     /* index form */
elp[1][0] = 1 ;     /* polynomial form */
for (i=1; i<nn-kk; i++)
{ elp[0][i] = -1 ; /* index form */
  elp[1][i] = 0 ; /* polynomial form */
}
l[0] = 0 ;
l[1] = 0 ;
u_lu[0] = -1 ;
u_lu[1] = 0 ;
u = 0 ;
do
{
  u++ ;
  if (d[u]==-1)
  { l[u+1] = l[u] ;
    for (i=0; i<=l[u]; i++)
    { elp[u+1][i] = elp[u][i] ;
      elp[u][i] = index_of[elp[u][i]] ;
    }
  }
}
else
/* search for words with greatest u_lu[q] for which d[q]!=0 */
{ q = u-1 ;
  while ((d[q]==-1) && (q>0)) q-- ;
/* have found first non-zero d[q] */
if (q>0)
{ j=q ;
  do
  { j-- ;
    if ((d[j]!=-1) && (u_lu[q]<u_lu[j]))
      q = j ;
  } while (j>0) ;
} ;

```

```

/* have now found q such that d[u]!=0 and u_lu[q] is maximum */
/* store degree of new elp polynomial */
    if (l[u]>l[q]+u-q) l[u+1] = l[u] ;
    else l[u+1] = l[q]+u-q ;
/* form new elp(x) */
    for (i=0; i<nn-kk; i++) elp[u+1][i] = 0 ;
    for (i=0; i<=l[q]; i++)
        if (elp[q][i]!=-1)
            elp[u+1][i+u-q] = alpha_to[(d[u]+nn-d[q]+elp[q][i])%nn] ;
    for (i=0; i<=l[u]; i++)
        { elp[u+1][i] ^= elp[u][i] ;
          elp[u][i] = index_of[elp[u][i]] ; /*convert old elp value to index*/
        }
    }
    u_lu[u+1] = u-l[u+1] ;
/* form (u+1)th discrepancy */
    if (u<nn-kk) /* no discrepancy computed on last iteration */
    {
        if (s[u+1]!=-1)
            d[u+1] = alpha_to[s[u+1]] ;
        else
            d[u+1] = 0 ;
        for (i=1; i<=l[u+1]; i++)
            if ((s[u+1-i]!=-1) && (elp[u+1][i]!=0))
                d[u+1] ^= alpha_to[(s[u+1-i]+index_of[elp[u+1][i]])%nn] ;
        d[u+1] = index_of[d[u+1]] ; /* put d[u+1] into index form */
    }
} while ((u<nn-kk) && (l[u+1]<=tt)) ;
u++ ;
if (l[u]<=tt) /* can correct error */
{
/* put elp into index form */
    for (i=0; i<=l[u]; i++) elp[u][i] = index_of[elp[u][i]] ;
}

```

```

/* find roots of the error location polynomial */
for (i=1; i<=l[u]; i++)
    reg[i] = elp[u][i] ;
count = 0 ;
for (i=1; i<=nn; i++)
{
    q = 1 ;
    for (j=1; j<=l[u]; j++)
        if (reg[j]!=-1)
        {
            reg[j] = (reg[j]+j)%nn ;
            q ^= alpha_to[reg[j]] ;
        } ;
    if (!q) /* store root and error location number indices */
    {
        root[count] = i;
        loc[count] = nn-i ;
        count++ ;
    } ;
} ;
if (count==l[u]) /* no. roots = degree of elp hence <= tt errors */
{
/* form polynomial z(x) */
for (i=1; i<=l[u]; i++) /* Z[0] = 1 always - do not need */
{
    if ((s[i]!=-1) && (elp[u][i]!=-1))
        z[i] = alpha_to[s[i]] ^ alpha_to[elp[u][i]] ;
    else if ((s[i]!=-1) && (elp[u][i]==-1))
        z[i] = alpha_to[s[i]] ;
    else if ((s[i]==-1) && (elp[u][i]!=-1))
        z[i] = alpha_to[elp[u][i]] ;
    else
        z[i] = 0 ;
    for (j=1; j<i; j++)
        if ((s[j]!=-1) && (elp[u][i-j]!=-1))
            z[i] ^= alpha_to[(elp[u][i-j] + s[j])%nn] ;
    z[i] = index_of[z[i]] ; /* put into index form */
} ;
} ;

```

```

/* evaluate errors at locations given by error location numbers loc[i] */
for (i=0; i<nn; i++)
{ err[i] = 0 ;
  if (recd[i]!=-1) /* convert recd[] to polynomial form */
    recd[i] = alpha_to[recd[i]] ;
  else recd[i] = 0 ;
}
for (i=0; i<l[u]; i++) /* compute numerator of error term first */
{ err[loc[i]] = 1; /* accounts for z[0] */
  for (j=1; j<=l[u]; j++)
    if (z[j]!=-1)
      err[loc[i]] ^= alpha_to[(z[j]+j*root[i])%nn] ;
  if (err[loc[i]]!=0)
  { err[loc[i]] = index_of[err[loc[i]]] ;
    q = 0 ; /* form denominator of error term */
    for (j=0; j<l[u]; j++)
      if (j!=i)
        q += index_of[1^alpha_to[(loc[j]+root[i])%nn]] ;
    q = q % nn ;
    err[loc[i]] = alpha_to[(err[loc[i]]-q+nn)%nn] ;
    recd[loc[i]] ^= err[loc[i]] ; /*recd[i] must be in polynomial form */
  }
}
}
else /* no. roots != degree of elp => >tt errors and cannot solve */
for (i=0; i<nn; i++) /* could return error flag if desired */
  if (recd[i]!=-1) /* convert recd[] to polynomial form */
    recd[i] = alpha_to[recd[i]] ;
  else recd[i] = 0 ; /* just output received codeword as is */
}
else /* elp has degree has degree >tt hence cannot solve */
for (i=0; i<nn; i++) /* could return error flag if desired */
  if (recd[i]!=-1) /* convert recd[] to polynomial form */
    recd[i] = alpha_to[recd[i]] ;
  else recd[i] = 0 ; /* just output received codeword as is */
}
else /* no non-zero syndromes => no errors: output received codeword */
for (i=0; i<nn; i++)
  if (recd[i]!=-1) /* convert recd[] to polynomial form */
    recd[i] = alpha_to[recd[i]] ;
  else recd[i] = 0 ;
}

```

```

main()
{
    register int i;
/* generate the Galois Field GF(2**mm) */
    generate_gf() ;
    printf("\n Look-up tables for GF(2**%d)\n",mm) ;
    printf(" i  alpha_to[i]  index_of[i]\n") ;
    for (i=0; i<=nn; i++)
        printf("%3X    %3X    %3X\n",i,alpha_to[i],index_of[i]) ;
    printf("\n\n") ;
    for (i=0; i<=nn; i++)
        printf("%3d,",alpha_to[i]) ;
/* compute the generator polynomial for this RS code */
    gen_poly() ;
/* for known data, stick a few numbers into a zero codeword. Data is in polynomial form.*/
    for (i=0; i<kk; i++)  data[i] = kk - i ;
/* for example, say we transmit the following message (nothing special!) */
    data[i] = kk - i
/* encode data[] to produce parity in bb[].  Data input and parity output is in polynomial
form*/
    encode_rs() ;
/* put the transmitted codeword, made up of data plus parity, in recd[] */
    for (i=0; i<nn-kk; i++)  recd[i] = bb[i] ;
    for (i=0; i<kk; i++)  recd[i+nn-kk] = data[i] ;
/* if you want to test the program, corrupt some of the elements of recd[] here. This can also
be done easily in a debugger. */
/* Again, lets say that a middle element is changed */
    data[nn-nn/2] = 16 ;
    for( i=0; i<255; i++)
        printf("\nMessage   %3d   Output   %3d ", i, codeword[ i ] );
    for (i=0; i<nn; i++)
    {
        printf("\ni= %3d  recd[i]=%3d index_of[recd[i]]=%3d ", i, recd[i],
index_of[recd[i]] );
        recd[i] = index_of[recd[i]] ;      /* put recd[i] into index form */
        printf("\ni=%3d  recd[i]=%3d ", i, recd[i] );
    }
/* decode recv[] */
    decode_rs() ;      /* recd[] is returned in polynomial form */
/* print out the relevant stuff - initial and decoded {parity and message} */
    printf("\n Results for Reed-Solomon code (n=%3d, k=%3d, t= %3d)\n\n",nn,kk,tt) ;
    printf(" i  data[i]  recd[i](decoded)  (data, recd in polynomial form)\n");
    for (i=0; i<nn-kk; i++)
        printf("%3d   %3d   %3d\n",i, bb[i], recd[i]) ;
    for (i=nn-kk; i<nn; i++)
        printf("%3d   %3d   %3d\n",i, data[i-nn+kk], recd[i]) ;
}

```

Chapter 8

Conclusion

8.1. Conclusion

The communication channel in modern digital and data storage systems requires error detecting and correcting codes to correct the errors that occur during the transmission of data.

In this project, the algorithms of (255, 223) Reed Solomon code were first discussed and then designed and implemented using RISC microcontroller. The (255, 223) Reed Solomon code is defined over Galois Field $GF(2^8)$ and is capable of correcting up to sixteen errors. A Reed Solomon code with different dimensions could have been chosen, but for the very common 8-bit data byte used today for many applications this code has been chosen.

The encoder is designed and implemented using linear feedback shift register (LFSR) and has been tested by simulation.

The decoder is more complicated than the encoder and, therefore, two level ROMs have been used and implemented in parallel to achieve a faster data rate. The decoder design is based on Berlekamp's iterative algorithm for finding the coefficients of the error location polynomial and Chien's searching algorithm for finding its roots.

The functionality of the encoder/decoder system was fully verified by simulation using Simulink and written in C language.

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Appendix A: Non zero elements of GF (2^8)

Filed element in powers Of alpha	Filed element in DEC form	Filed element in HEX form	Filed element in BINARY Form “nonzero elements”
-	0	00	0 0 0 0 0 0 0 0
0	1	01	1 0 0 0 0 0 0 0
1	2	02	0 1 0 0 0 0 0 0
2	4	04	0 0 1 0 0 0 0 0
3	8	08	0 0 0 1 0 0 0 0
4	16	10	0 0 0 0 1 0 0 0
5	32	20	0 0 0 0 0 1 0 0
6	64	40	0 0 0 0 0 0 1 0
7	128	80	0 0 0 0 0 0 0 1
8	29	1D	1 0 1 1 1 0 0 0
9	58	3A	0 1 0 1 1 1 0 0
10	116	74	0 0 1 0 1 1 1 0
11	232	E8	0 0 0 1 0 1 1 1
12	205	CD	1 0 1 1 0 0 1 1
13	135	87	1 1 1 0 0 0 0 1
14	19	13	1 1 0 0 1 0 0 0
15	38	26	0 1 1 0 0 1 0 0
16	76	4C	0 0 1 1 0 0 1 0
17	152	98	0 0 0 1 1 0 0 1
18	45	2D	1 0 1 1 0 1 0 0
19	90	5A	0 1 0 1 1 0 1 0
20	180	B4	0 0 1 0 1 1 0 1
21	117	75	1 0 1 0 1 1 1 0
22	234	EA	0 1 0 1 0 1 1 1
23	201	C9	1 0 0 1 0 0 1 1
24	143	8F	1 1 1 1 0 0 0 1
25	3	3	1 1 0 0 0 0 0 0
26	6	6	0 1 1 0 0 0 0 0
27	12	C	0 0 1 1 0 0 0 0
28	24	18	0 0 0 1 1 0 0 0
29	48	30	0 0 0 0 1 1 0 0
30	96	60	0 0 0 0 0 1 1 0
31	192	C0	0 0 0 0 0 0 1 1
32	157	9D	1 0 1 1 1 0 0 1
33	39	27	1 1 1 0 0 1 0 0
34	78	4E	0 1 1 1 0 0 1 0
35	156	9C	0 0 1 1 1 0 0 1
36	37	25	1 0 1 0 0 1 0 0

37	74	4A	0	1	0	1	0	0	1	0
38	148	94	0	0	1	0	1	0	0	1
39	53	35	1	0	1	0	1	1	0	0
40	106	6A	0	1	0	1	0	1	1	0
41	212	D4	0	0	1	0	1	0	1	1
42	181	B5	1	0	1	0	1	1	0	1
43	119	77	1	1	1	0	1	1	1	0
44	238	EE	0	1	1	1	0	1	1	1
45	193	C1	1	0	0	0	0	0	1	1
46	159	9F	1	1	1	1	1	0	0	1
47	35	23	1	1	0	0	0	1	0	0
48	70	46	0	1	1	0	0	0	1	0
49	140	8C	0	0	1	1	0	0	0	1
50	5	5	1	0	1	0	0	0	0	0
51	10	A	0	1	0	1	0	0	0	0
52	20	14	0	0	1	0	1	0	0	0
53	40	28	0	0	0	1	0	1	0	0
54	80	50	0	0	0	0	1	0	1	0
55	160	A0	0	0	0	0	0	1	0	1
56	93	5D	1	0	1	1	1	0	1	0
57	186	BA	0	1	0	1	1	1	0	1
58	105	69	1	0	0	1	0	1	1	0
59	210	D2	0	1	0	0	1	0	1	1
60	185	B9	1	0	0	1	1	1	0	1
61	111	6F	1	1	1	1	0	1	1	0
62	222	DE	0	1	1	1	1	0	1	1
63	161	A1	1	0	0	0	0	1	0	1
64	95	5F	1	1	1	1	1	0	1	0
65	190	BE	0	1	1	1	1	1	0	1
66	97	61	1	0	0	0	0	1	1	0
67	194	C2	0	1	0	0	0	0	1	1
68	153	99	1	0	0	1	1	0	0	1
69	47	2F	1	1	1	1	0	1	0	0
70	94	5E	0	1	1	1	1	0	1	0
71	188	BC	0	0	1	1	1	1	0	1
72	101	65	1	0	1	0	0	1	1	0
73	202	CA	0	1	0	1	0	0	1	1
74	137	89	1	0	0	1	0	0	0	1
75	15	F	1	1	1	1	0	0	0	0
76	30	1E	0	1	1	1	1	0	0	0
77	60	3C	0	0	1	1	1	1	0	0
78	120	78	0	0	0	1	1	1	1	0
79	240	F0	0	0	0	0	1	1	1	1
80	253	FD	1	0	1	1	1	1	1	1
81	231	E7	1	1	1	0	0	1	1	1

82	211	D3	1	1	0	0	1	0	1	1
83	187	BB	1	1	0	1	1	1	0	1
84	107	6B	1	1	0	1	0	1	1	0
85	214	D6	0	1	1	0	1	0	1	1
86	177	B1	1	0	0	0	1	1	0	1
87	127	7F	1	1	1	1	1	1	1	0
88	254	FE	0	1	1	1	1	1	1	1
89	225	E1	1	0	0	0	0	1	1	1
90	223	DF	1	1	1	1	1	0	1	1
91	163	A3	1	1	0	0	0	1	0	1
92	91	5B	1	1	0	1	1	0	1	0
93	182	B6	0	1	1	0	1	1	0	1
94	113	71	1	0	0	0	1	1	1	0
95	226	E2	0	1	0	0	0	1	1	1
96	217	D9	1	0	0	1	1	0	1	1
97	175	AF	1	1	1	1	0	1	0	1
98	67	43	1	1	0	0	0	0	1	0
99	134	86	0	1	1	0	0	0	0	1
100	17	11	1	0	0	0	1	0	0	0
101	34	22	0	1	0	0	0	1	0	0
102	68	44	0	0	1	0	0	0	1	0
103	136	88	0	0	0	1	0	0	0	1
104	13	D	1	0	1	1	0	0	0	0
105	26	1A	0	1	0	1	1	0	0	0
106	52	34	0	0	1	0	1	1	0	0
107	104	68	0	0	0	1	0	1	1	0
108	208	D0	0	0	0	0	1	0	1	1
109	189	BD	1	0	1	1	1	1	0	1
110	103	67	1	1	1	0	0	1	1	0
111	206	CE	0	1	1	1	0	0	1	1
112	129	81	1	0	0	0	0	0	0	1
113	31	1F	1	1	1	1	1	0	0	0
114	62	3E	0	1	1	1	1	1	0	0
115	124	7C	0	0	1	1	1	1	1	0
116	248	F8	0	0	0	1	1	1	1	1
117	237	ED	1	0	1	1	0	1	1	1
118	199	C7	1	1	1	0	0	0	1	1
119	147	93	1	1	0	0	1	0	0	1
120	59	3B	1	1	0	1	1	1	0	0
121	118	76	0	1	1	0	1	1	1	0
122	236	EC	0	0	1	1	0	1	1	1
123	197	C5	1	0	1	0	0	0	1	1
124	151	97	1	1	1	0	1	0	0	1
125	51	33	1	1	0	0	1	1	0	0
126	102	66	0	1	1	0	0	1	1	0

127	204	CC	0	0	1	1	0	0	1	1
128	133	85	1	0	1	0	0	0	0	1
129	23	17	1	1	1	0	1	0	0	0
130	46	2E	0	1	1	1	0	1	0	0
131	92	5C	0	0	1	1	1	0	1	0
132	184	B8	0	0	0	1	1	1	0	1
133	109	6D	1	0	1	1	0	1	1	0
134	218	DA	0	1	0	1	1	0	1	1
135	169	A9	1	0	0	1	0	1	0	1
136	79	4F	1	1	1	1	0	0	1	0
137	158	9E	0	1	1	1	1	0	0	1
138	33	21	1	0	0	0	0	1	0	0
139	66	42	0	1	0	0	0	0	1	0
140	132	84	0	0	1	0	0	0	0	1
141	21	15	1	0	1	0	1	0	0	0
142	42	2A	0	1	0	1	0	1	0	0
143	84	54	0	0	1	0	1	0	1	0
144	168	A8	0	0	0	1	0	1	0	1
145	77	4D	1	0	1	1	0	0	1	0
146	154	9A	0	1	0	1	1	0	0	1
147	41	29	1	0	0	1	0	1	0	0
148	82	52	0	1	0	0	1	0	1	0
149	164	A4	0	0	1	0	0	1	0	1
150	85	55	1	0	1	0	1	0	1	0
151	170	AA	0	1	0	1	0	1	0	1
152	73	49	1	0	0	1	0	0	1	0
153	146	92	0	1	0	0	1	0	0	1
154	57	39	1	0	0	1	1	1	0	0
155	114	72	0	1	0	0	1	1	1	0
156	228	E4	0	0	1	0	0	1	1	1
157	213	D5	1	0	1	0	1	0	1	1
158	183	B7	1	1	1	0	1	1	0	1
159	115	73	1	1	0	0	1	1	1	0
160	230	E6	0	1	1	0	0	1	1	1
161	209	D1	1	0	0	0	1	0	1	1
162	191	BF	1	1	1	1	1	1	0	1
163	99	63	1	1	0	0	0	1	1	0
164	198	C6	0	1	1	0	0	0	1	1
165	145	91	1	0	0	0	1	0	0	1
166	63	3F	1	1	1	1	1	1	0	0
167	126	7E	0	1	1	1	1	1	1	0
168	252	FC	0	0	1	1	1	1	1	1
169	229	E5	1	0	1	0	0	1	1	1
170	215	D7	1	1	1	0	1	0	1	1
171	179	B3	1	1	0	0	1	1	0	1

172	123	7B	1	1	0	1	1	1	1	0
173	246	F6	0	1	1	0	1	1	1	1
174	241	F1	1	0	0	0	1	1	1	1
175	255	FF	1	1	1	1	1	1	1	1
176	227	E3	1	1	0	0	0	1	1	1
177	219	DB	1	1	0	1	1	0	1	1
178	171	AB	1	1	0	1	0	1	0	1
179	75	4B	1	1	0	1	0	0	1	0
180	150	96	0	1	1	0	1	0	0	1
181	49	31	1	0	0	0	1	1	0	0
182	98	62	0	1	0	0	0	1	1	0
183	196	C4	0	0	1	0	0	0	1	1
184	149	95	1	0	1	0	1	0	0	1
185	55	37	1	1	1	0	1	1	0	0
186	110	6E	0	1	1	1	0	1	1	0
187	220	DC	0	0	1	1	1	0	1	1
188	165	A5	1	0	1	0	0	1	0	1
189	87	57	1	1	1	0	1	0	1	0
190	174	AE	0	1	1	1	0	1	0	1
191	65	41	1	0	0	0	0	0	1	0
192	130	82	0	1	0	0	0	0	0	1
193	25	19	1	0	0	1	1	0	0	0
194	50	32	0	1	0	0	1	1	0	0
195	100	64	0	0	1	0	0	1	1	0
196	200	C8	0	0	0	1	0	0	1	1
197	141	8D	1	0	1	1	0	0	0	1
198	7	7	1	1	1	0	0	0	0	0
199	14	E	0	1	1	1	0	0	0	0
200	28	1C	0	0	1	1	1	0	0	0
201	56	38	0	0	0	1	1	1	0	0
202	112	70	0	0	0	0	1	1	1	0
203	224	E0	0	0	0	0	0	1	1	1
204	221	DD	1	0	1	1	1	0	1	1
205	167	A7	1	1	1	0	0	1	0	1
206	83	53	1	1	0	0	1	0	1	0
207	166	A6	0	1	1	0	0	1	0	1
208	81	51	1	0	0	0	1	0	1	0
209	162	A2	0	1	0	0	0	1	0	1
210	89	59	1	0	0	1	1	0	1	0
211	178	B2	0	1	0	0	1	1	0	1
212	121	79	1	0	0	1	1	1	1	0
213	242	F2	0	1	0	0	1	1	1	1
214	249	F9	1	0	0	1	1	1	1	1
215	239	EF	1	1	1	1	0	1	1	1
216	195	C3	1	1	0	0	0	0	1	1

217	155	9B	1	1	0	1	1	0	0	1
218	43	2B	1	1	0	1	0	1	0	0
219	86	56	0	1	1	0	1	0	1	0
220	172	AC	0	0	1	1	0	1	0	1
221	69	45	1	0	1	0	0	0	1	0
222	138	8A	0	1	0	1	0	0	0	1
223	9	9	1	0	0	1	0	0	0	0
224	18	12	0	1	0	0	1	0	0	0
225	36	24	0	0	1	0	0	1	0	0
226	72	48	0	0	0	1	0	0	1	0
227	144	90	0	0	0	0	1	0	0	1
228	61	3D	1	0	1	1	1	1	0	0
229	122	7A	0	1	0	1	1	1	1	0
230	244	F4	0	0	1	0	1	1	1	1
231	245	F5	1	0	1	0	1	1	1	1
232	247	F7	1	1	1	0	1	1	1	1
233	243	F3	1	1	0	0	1	1	1	1
234	251	FB	1	1	0	1	1	1	1	1
235	235	EB	1	1	0	1	0	1	1	1
236	203	CB	1	1	0	1	0	0	1	1
237	139	8B	1	1	0	1	0	0	0	1
238	11	B	1	1	0	1	0	0	0	0
239	22	16	0	1	1	0	1	0	0	0
240	44	2C	0	0	1	1	0	1	0	0
241	88	58	0	0	0	1	1	0	1	0
242	176	B0	0	0	0	0	1	1	0	1
243	125	7D	1	0	1	1	1	1	1	0
244	250	FA	0	1	0	1	1	1	1	1
245	233	E9	1	0	0	1	0	1	1	1
246	207	CF	1	1	1	1	0	0	1	1
247	131	83	1	1	0	0	0	0	0	1
248	27	1B	1	1	0	1	1	0	0	0
249	54	36	0	1	1	0	1	1	0	0
250	108	6C	0	0	1	1	0	1	1	0
251	216	D8	0	0	0	1	1	0	1	1
252	173	AD	1	0	1	1	0	1	0	1
253	71	47	1	1	1	0	0	0	1	0
254	142	8E	0	1	1	1	0	0	0	1

Table A1. The field elements for GF (255) with $P(X) = 1+X^2+X^3+X^4+X^8$

Appendix B: The ROMs tables

Elem. GF (255) Of α^i	The ROMs Tables “HEX Form” Multiplying any elements from GF(255) with the relevant coefficients of the generator polynomial $g(X)$										
	g_0 = α^{18}	g_1 = α^{251}	g_2 = α^{215}	g_3 = α^{28}	g_4 = α^{80}	g_5 = α^{107}	g_6 = α^{248}	g_7 = α^{53}	g_8 = α^{84}	g_9 = α^{194}	g_{10} = α^{91}
-	00	00	00	00	00	00	00	00	00	00	00
0	2D	D8	EF	18	FD	68	1B	28	6B	32	A3
1	5A	AD	C3	30	E7	D0	36	50	D6	64	5B
2	B4	47	9B	60	D3	BD	6C	A0	B1	C8	B6
3	75	8E	2B	C0	BB	67	D8	5D	7F	8D	71
4	EA	02	56	9D	6B	CE	AD	BA	FE	7	E2
5	C9	04	AC	27	D6	81	47	69	E1	E	D9
6	8F	08	45	4E	B1	1F	8E	D2	DF	1C	AF
7	3	10	8A	9C	7F	3E	02	B9	A3	38	43
8	6	20	9	25	FE	7C	04	6F	5B	70	86
9	C	40	12	4A	E1	F8	02	DE	B6	E0	11
10	18	80	24	94	DF	ED	04	A1	71	DD	22
11	30	1D	48	35	A3	C7	08	5F	E2	A7	44
12	60	3A	90	6A	5B	93	10	BE	D9	53	88
13	C0	74	3D	D4	B6	3B	20	61	AF	A6	D
14	9D	E8	7A	B5	71	76	40	C2	43	51	1A
15	27	CD	F4	77	E2	EC	80	99	86	A2	34
16	4E	87	F5	EE	D9	C5	1D	2F	11	59	68
17	9C	13	F7	C1	AF	97	3A	5E	22	B2	D0
18	25	26	F3	9F	43	33	74	BC	44	79	BD
19	4A	4C	FB	23	86	66	E8	65	88	F2	67
20	94	98	EB	46	11	CC	CD	CA	D	F9	CE
21	35	2D	CB	8C	22	85	87	89	1A	EF	81
22	6A	5A	8B	5	44	17	13	F	34	C3	1F
23	D4	B4	B	A	88	2E	26	1E	68	9B	3E
24	B5	75	16	14	D	5C	4C	3C	D0	2B	7C
25	77	EA	2C	28	1A	B8	98	78	BD	56	F8
26	EE	C9	58	50	34	6D	2D	F0	67	AC	ED
27	C1	8F	B0	A0	68	DA	5A	FD	CE	45	C7
28	9F	3	7D	5D	D0	A9	B4	E7	81	8A	93
29	23	6	FA	BA	BD	4F	75	D3	1F	9	3B
30	46	C	E9	69	67	9E	EA	BB	3E	12	76
31	8C	18	CF	D2	CE	21	C9	6B	7C	24	EC
32	5	30	83	B9	81	42	8F	D6	F8	48	C5
33	A	60	1B	6F	1F	84	3	B1	ED	90	97
34	14	C0	36	DE	3E	15	6	7F	C7	3D	33
35	28	9D	6C	A1	7C	2A	C	FE	93	7A	66

36	50	27	D8	5F	F8	54	18	E1	3B	F4	CC
37	A0	4E	AD	BE	ED	A8	30	DF	76	F5	85
38	5D	9C	47	61	C7	4D	60	A3	EC	F7	17
39	BA	25	8E	C2	93	9A	C0	5B	C5	F3	2E
40	69	4A	02	99	3B	29	9D	B6	97	FB	5C
41	D2	94	04	2F	76	52	27	71	33	EB	B8
42	B9	35	08	5E	EC	A4	4E	E2	66	CB	6D
43	6F	6A	10	BC	C5	55	9C	D9	CC	8B	DA
44	DE	D4	20	65	97	AA	25	AF	85	B	A9
45	A1	B5	40	CA	33	49	4A	43	17	16	4F
46	5F	77	80	89	66	92	94	86	2E	2C	9E
47	BE	EE	1D	F	CC	39	35	11	5C	58	21
48	61	C1	3A	1E	85	72	6A	22	B8	B0	42
49	C2	9F	74	3C	17	E4	D4	44	6D	7D	84
50	99	23	E8	78	2E	D5	B5	88	DA	FA	15
51	2F	46	CD	F0	5C	B7	77	D	A9	E9	2A
52	5E	8C	87	FD	B8	73	EE	1A	4F	CF	54
53	BC	5	13	E7	6D	E6	C1	34	9E	83	A8
54	65	A	26	D3	DA	D1	9F	68	21	1B	4D
55	CA	14	4C	BB	A9	BF	23	D0	42	36	9A
56	89	28	98	6B	4F	63	46	BD	84	6C	29
57	F	50	2D	D6	9E	C6	8C	67	15	D8	52
58	1E	A0	5A	B1	21	91	5	CE	2A	AD	A4
59	3C	5D	B4	7F	42	3F	A	81	54	47	55
60	78	BA	75	FE	84	7E	14	1F	A8	8E	AA
61	F0	69	EA	E1	15	FC	28	3E	4D	02	49
62	FD	D2	C9	DF	2A	E5	50	7C	9A	04	92
63	E7	B9	8F	A3	54	D7	A0	F8	29	08	39
64	D3	6F	3	5B	A8	B3	5D	ED	52	10	72
65	BB	DE	6	B6	4D	7B	BA	C7	A4	20	E4
66	6B	A1	C	71	9A	F6	69	93	55	40	D5
67	D6	5F	18	E2	29	F1	D2	3B	AA	80	B7
68	B1	BE	30	D9	52	FF	B9	76	49	1D	73
69	7F	61	60	AF	A4	E3	6F	EC	92	3A	E6
70	FE	C2	C0	43	55	DB	DE	C5	39	74	D1
71	E1	99	9D	86	AA	AB	A1	97	72	E8	BF
72	DF	2F	27	11	49	4B	5F	33	E4	CD	63
73	A3	5E	4E	22	92	96	BE	66	D5	87	C6
74	5B	BC	9C	44	39	31	61	CC	B7	13	91
75	B6	65	25	88	72	62	C2	85	73	26	3F
76	71	CA	4A	D	E4	C4	99	17	E6	4C	7E
77	E2	89	94	1A	D5	95	2F	2E	D1	98	FC
78	D9	F	35	34	B7	37	5E	5C	BF	2D	E5
79	AF	1E	6A	68	73	6E	BC	B8	63	5A	D7
80	43	3C	D4	D0	E6	DC	65	6D	C6	B4	B3

81	86	78	B5	BD	D1	A5	CA	DA	91	75	7B
82	11	F0	77	67	BF	57	89	A9	3F	EA	F6
83	22	FD	EE	CE	63	AE	F	4F	7E	C9	F1
84	44	E7	C1	81	C6	41	1E	9E	FC	8F	FF
85	88	D3	9F	1F	91	82	3C	21	E5	3	E3
86	D	BB	23	3E	3F	19	78	42	D7	6	DB
87	1A	6B	46	7C	7E	32	F0	84	B3	C	AB
88	34	D6	8C	F8	FC	64	FD	15	7B	18	4B
89	68	B1	5	ED	E5	C8	E7	2A	F6	30	96
90	D0	7F	A	C7	D7	8D	D3	54	F1	60	31
91	BD	FE	14	93	B3	7	BB	A8	FF	C0	62
92	67	E1	28	3B	7B	E	6B	4D	E3	9D	C4
93	CE	DF	50	76	F6	1C	D6	9A	DB	27	95
94	81	A3	A0	EC	F1	38	B1	29	AB	4E	37
95	1F	5B	5D	C5	FF	70	7F	52	4B	9C	6E
96	3E	B6	BA	97	E3	E0	FE	A4	96	25	DC
97	7C	71	69	33	DB	DD	E1	55	31	4A	A5
98	F8	E2	D2	66	AB	A7	DF	AA	62	94	57
99	ED	D9	B9	CC	4B	53	A3	49	C4	35	AE
100	C7	AF	6F	85	96	A6	5B	92	95	6A	41
101	93	43	DE	17	31	51	B6	39	37	D4	82
102	3B	86	A1	2E	62	A2	71	72	6E	B5	19
103	76	11	5F	5C	C4	59	E2	E4	DC	77	32
104	EC	22	BE	B8	95	B2	D9	D5	A5	EE	64
105	C5	44	61	6D	37	79	AF	B7	57	C1	C8
106	97	88	C2	DA	6E	F2	43	73	AE	9F	8D
107	33	D	99	A9	DC	F9	86	E6	41	23	7
108	66	1A	2F	4F	A5	EF	11	D1	82	46	E
109	CC	34	5E	9E	57	C3	22	BF	19	8C	1C
110	85	68	BC	21	AE	9B	44	63	32	5	38
111	17	D0	65	42	41	2B	88	C6	64	A	70
112	2E	BD	CA	84	82	56	D	91	C8	14	E0
113	5C	67	89	15	19	AC	1A	3F	8D	28	DD
114	B8	CE	F	2A	32	45	34	7E	7	50	A7
115	6D	81	1E	54	64	8A	68	FC	E	A0	53
116	DA	1F	3C	A8	C8	9	D0	E5	1C	5D	A6
117	A9	3E	78	4D	8D	12	BD	D7	38	BA	51
118	4F	7C	F0	9A	7	24	67	B3	70	69	A2
119	9E	F8	FD	29	E	48	CE	7B	E0	D2	59
120	21	ED	E7	52	1C	90	81	F6	DD	B9	B2
121	42	C7	D3	A4	38	3D	1F	F1	A7	6F	79
122	84	93	BB	55	70	7A	3E	FF	53	DE	F2
123	15	3B	6B	AA	E0	F4	7C	E3	A6	A1	F9
124	2A	76	D6	49	DD	F5	F8	DB	51	5F	EF
125	54	EC	B1	92	A7	F7	ED	AB	A2	BE	C3

126	A8	C5	7F	39	53	F3	C7	4B	59	61	9B
127	4D	97	FE	72	A6	FB	93	96	B2	C2	2B
128	9A	33	E1	E4	51	EB	3B	31	79	99	56
129	29	66	DF	D5	A2	CB	76	62	F2	2F	AC
130	52	CC	A3	B7	59	8B	EC	C4	F9	5E	45
131	A4	85	5B	73	B2	B	C5	95	EF	BC	8A
132	55	17	B6	E6	79	16	97	37	C3	65	9
133	AA	2E	71	D1	F2	2C	33	6E	9B	CA	12
134	49	5C	E2	BF	F9	58	66	DC	2B	89	24
135	92	B8	D9	63	EF	B0	CC	A5	56	F	48
136	39	6D	AF	C6	C3	7D	85	57	AC	1E	90
137	72	DA	43	91	9B	FA	17	AE	45	3C	3D
138	E4	A9	86	3F	2B	E9	2E	41	8A	78	7A
139	D5	4F	11	7E	56	CF	5C	82	9	F0	F4
140	B7	9E	22	FC	AC	83	B8	19	12	FD	F5
141	73	21	44	E5	45	1B	6D	32	24	E7	F7
142	E6	42	88	D7	8A	36	DA	64	48	D3	F3
143	D1	84	D	B3	9	6C	A9	C8	90	BB	FB
144	BF	15	1A	7B	12	D8	4F	8D	3D	6B	EB
145	63	2A	34	F6	24	AD	9E	7	7A	D6	CB
146	C6	54	68	F1	48	47	21	E	F4	B1	8B
147	91	A8	D0	FF	90	8E	42	1C	F5	7F	B
148	3F	4D	BD	E3	3D	02	84	38	F7	FE	16
149	7E	9A	67	DB	7A	04	15	70	F3	E1	2C
150	FC	29	CE	AB	F4	08	2A	E0	FB	DF	58
151	E5	52	81	4B	F5	10	54	DD	EB	A3	B0
152	D7	A4	1F	96	F7	20	A8	A7	CB	5B	7D
153	B3	55	3E	31	F3	40	4D	53	8B	B6	FA
154	7B	AA	7C	62	FB	80	9A	A6	B	71	E9
155	F6	49	F8	C4	EB	1D	29	51	16	E2	CF
156	F1	92	ED	95	CB	3A	52	A2	2C	D9	83
157	FF	39	C7	37	8B	74	A4	59	58	AF	1B
158	E3	72	93	6E	B	E8	55	B2	B0	43	36
159	DB	E4	3B	DC	16	CD	AA	79	7D	86	6C
160	AB	D5	76	A5	2C	87	49	F2	FA	11	D8
161	4B	B7	EC	57	58	13	92	F9	E9	22	AD
162	96	73	C5	AE	B0	26	39	EF	CF	44	47
163	31	E6	97	41	7D	4C	72	C3	83	88	8E
164	62	D1	33	82	FA	98	E4	9B	1B	D	02
165	C4	BF	66	19	E9	2D	D5	2B	36	1A	04
166	95	63	CC	32	CF	5A	B7	56	6C	34	08
167	37	C6	85	64	83	B4	73	AC	D8	68	10
168	6E	91	17	C8	1B	75	E6	45	AD	D0	20
169	DC	3F	2E	8D	36	EA	D1	8A	47	BD	40
170	A5	7E	5C	7	6C	C9	BF	9	8E	67	80

171	57	FC	B8	E	D8	8F	63	12	02	CE	1D
172	AE	E5	6D	1C	AD	3	C6	24	04	81	3A
173	41	D7	DA	38	47	6	91	48	08	1F	74
174	82	B3	A9	70	8E	C	3F	90	10	3E	E8
175	19	7B	4F	E0	02	18	7E	3D	20	7C	CD
176	32	F6	9E	DD	04	30	FC	7A	40	F8	87
177	64	F1	21	A7	08	60	E5	F4	80	ED	13
178	C8	FF	42	53	10	C0	D7	F5	1D	C7	26
179	8D	E3	84	A6	20	9D	B3	F7	3A	93	4C
180	7	DB	15	51	40	27	7B	F3	74	3B	98
181	E	AB	2A	A2	80	4E	F6	FB	E8	76	2D
182	1C	4B	54	59	1D	9C	F1	EB	CD	EC	5A
183	38	96	A8	B2	3A	25	FF	CB	87	C5	B4
184	70	31	4D	79	74	4A	E3	8B	13	97	75
185	E0	62	9A	F2	E8	94	DB	B	26	33	EA
186	DD	C4	29	F9	CD	35	AB	16	4C	66	C9
187	A7	95	52	EF	87	6A	4B	2C	98	CC	8F
188	53	37	A4	C3	13	D4	96	58	2D	85	3
189	A6	6E	55	9B	26	B5	31	B0	5A	17	6
190	51	DC	AA	2B	4C	77	62	7D	B4	2E	C
191	A2	A5	49	56	98	EE	C4	FA	75	5C	18
192	59	57	92	AC	2D	C1	95	E9	EA	B8	30
193	B2	AE	39	45	5A	9F	37	CF	C9	6D	60
194	79	41	72	8A	B4	23	6E	83	8F	DA	C0
195	F2	82	E4	9	75	46	DC	1B	3	A9	9D
196	F9	19	D5	12	EA	8C	A5	36	6	4F	27
197	EF	32	B7	24	C9	5	57	6C	C	9E	4E
198	C3	64	73	48	8F	A	AE	D8	18	21	9C
199	9B	C8	E6	90	3	14	41	AD	30	42	25
200	2B	8D	D1	3D	6	28	82	47	60	84	4A
201	56	7	BF	7A	C	50	19	8E	C0	15	94
202	AC	E	63	F4	18	A0	32	02	9D	2A	35
203	45	1C	C6	F5	30	5D	64	04	27	54	6A
204	8A	38	91	F7	60	BA	C8	08	4E	A8	D4
205	9	70	3F	F3	C0	69	8D	10	9C	4D	B5
206	12	E0	7E	FB	9D	D2	7	20	25	9A	77
207	24	DD	FC	EB	27	B9	E	40	4A	29	EE
208	48	A7	E5	CB	4E	6F	1C	80	94	52	C1
209	90	53	D7	8B	9C	DE	38	1D	35	A4	9F
210	3D	A6	B3	B	25	A1	70	3A	6A	55	23
211	7A	51	7B	16	4A	5F	E0	74	D4	AA	46
212	F4	A2	F6	2C	94	BE	DD	E8	B5	49	8C
213	F5	59	F1	58	35	61	A7	CD	77	92	5
214	F7	B2	FF	B0	6A	C2	53	87	EE	39	A
215	F3	79	E3	7D	D4	99	A6	13	C1	72	14

216	FB	F2	DB	FA	B5	2F	51	26	9F	E4	28
217	EB	F9	AB	E9	77	5E	A2	4C	23	D5	50
218	CB	EF	4B	CF	EE	BC	59	98	46	B7	A0
219	8B	C3	96	83	C1	65	B2	2D	8C	73	5D
220	B	9B	31	1B	9F	CA	79	5A	5	E6	BA
221	16	2B	62	36	23	89	F2	B4	A	D1	69
222	2C	56	C4	6C	46	F	F9	75	14	BF	D2
223	58	AC	95	D8	8C	1E	EF	EA	28	63	B9
224	B0	45	37	AD	5	3C	C3	C9	50	C6	6F
225	7D	8A	6E	47	A	78	9B	8F	A0	91	DE
226	FA	9	DC	8E	14	F0	2B	3	5D	3F	A1
227	E9	12	A5	02	28	FD	56	6	BA	7E	5F
228	CF	24	57	04	50	E7	AC	C	69	FC	BE
229	83	48	AE	08	A0	D3	45	18	D2	E5	61
230	1B	90	41	10	5D	BB	8A	30	B9	D7	C2
231	36	3D	82	20	BA	6B	9	60	6F	B3	99
232	6C	7A	19	40	69	D6	12	C0	DE	7B	2F
233	D8	F4	32	80	D2	B1	24	9D	A1	F6	5E
234	AD	F5	64	1D	B9	7F	48	27	5F	F1	BC
235	47	F7	C8	3A	6F	FE	90	4E	BE	FF	65
236	8E	F3	8D	74	DE	E1	3D	9C	61	E3	CA
237	02	FB	7	E8	A1	DF	7A	25	C2	DB	89
238	04	EB	E	CD	5F	A3	F4	4A	99	AB	F
239	08	CB	1C	87	BE	5B	F5	94	2F	4B	1E
240	10	8B	38	13	61	B6	F7	35	5E	96	3C
241	20	B	70	26	C2	71	F3	6A	BC	31	78
242	40	16	E0	4C	99	E2	FB	D4	65	62	F0
243	80	2C	DD	98	2F	D9	EB	B5	CA	C4	FD
244	1D	58	A7	2D	5E	AF	CB	77	89	95	E7
245	3A	B0	53	5A	BC	43	8B	EE	F	37	D3
246	74	7D	A6	B4	65	86	B	C1	1E	6E	BB
247	E8	FA	51	75	CA	11	16	9F	3C	DC	6B
248	CD	E9	A2	EA	89	22	2C	23	78	A5	D6
249	87	CF	59	C9	F	44	58	46	F0	57	B1
250	13	83	B2	8F	1E	88	B0	8C	FD	AE	7F
251	26	1B	79	3	3C	D	7D	5	E7	41	FE
252	4C	36	F2	6	78	1A	FA	A	D3	82	E1
253	98	6C	F9	C	F0	34	E9	14	BB	19	DF
254	2D	D8	EF	18	FD	68	CF	28	6B	32	A3

.....→Continue→.....

Elem. GF (255) Of α^i	The ROMs Tables “HEX Form” Multiplying any elements from GF(255) with the relevant coefficients of the generator polynomial g(X)										
	g_{11} = α^{59}	g_{12} = α^{176}	g_{13} = α^{99}	g_{14} = α^{203}	g_{15} = α^{137}	g_{16} = α^{43}	g_{17} = α^{104}	g_{18} = α^{137}	g_{19} = α^0	g_{20} = α^{44}	g_{21} = α^{149}
-	00	00	00	00	00	00	00	00	00	00	00
0	D2	E3	86	E0	9E	77	D	9E	1	EE	A4
1	B9	DB	11	DD	21	EE	1A	21	02	C1	55
2	6F	AB	22	A7	42	C1	34	42	04	9F	AA
3	DE	4B	44	53	84	9F	68	84	08	23	49
4	A1	96	88	A6	15	23	D0	15	10	46	92
5	5F	31	D	51	2A	46	BD	2A	20	8C	39
6	BE	62	1A	A2	54	8C	67	54	40	5	72
7	61	C4	34	59	A8	5	CE	A8	80	A	E4
8	C2	95	68	B2	4D	A	81	4D	1D	14	D5
9	99	37	D0	79	9A	14	1F	9A	3A	28	B7
10	2F	6E	BD	F2	29	28	3E	29	74	50	73
11	5E	DC	67	F9	52	50	7C	52	E8	A0	E6
12	BC	A5	CE	EF	A4	A0	F8	A4	CD	5D	D1
13	65	57	81	C3	55	5D	ED	55	87	BA	BF
14	CA	AE	1F	9B	AA	BA	C7	AA	13	69	63
15	89	41	3E	2B	49	69	93	49	26	D2	C6
16	F	82	7C	56	92	D2	3B	92	4C	B9	91
17	1E	19	F8	AC	39	B9	76	39	98	6F	3F
18	3C	32	ED	45	72	6F	EC	72	2D	DE	7E
19	78	64	C7	8A	E4	DE	C5	E4	5A	A1	FC
20	F0	C8	93	9	D5	A1	97	D5	B4	5F	E5
21	FD	8D	3B	12	B7	5F	33	B7	75	BE	D7
22	E7	7	76	24	73	BE	66	73	EA	61	B3
23	D3	E	EC	48	E6	61	CC	E6	C9	C2	7B
24	BB	1C	C5	90	D1	C2	85	D1	8F	99	F6
25	6B	38	97	3D	BF	99	17	BF	3	2F	F1
26	D6	70	33	7A	63	2F	2E	63	6	5E	FF
27	B1	E0	66	F4	C6	5E	5C	C6	C	BC	E3
28	7F	DD	CC	F5	91	BC	B8	91	18	65	DB
29	FE	A7	85	F7	3F	65	6D	3F	30	CA	AB
30	E1	53	17	F3	7E	CA	DA	7E	60	89	4B
31	DF	A6	2E	FB	FC	89	A9	FC	C0	F	96
32	A3	51	5C	EB	E5	F	4F	E5	9D	1E	31
33	5B	A2	B8	CB	D7	1E	9E	D7	27	3C	62
34	B6	59	6D	8B	B3	3C	21	B3	4E	78	C4
35	71	B2	DA	B	7B	78	42	7B	9C	F0	95
36	E2	79	A9	16	F6	F0	84	F6	25	FD	37
37	D9	F2	4F	2C	F1	FD	15	F1	4A	E7	6E

38	AF	F9	9E	58	FF	E7	2A	FF	94	D3	DC
39	43	EF	21	B0	E3	D3	54	E3	35	BB	A5
40	86	C3	42	7D	DB	BB	A8	DB	6A	6B	57
41	11	9B	84	FA	AB	6B	4D	AB	D4	D6	AE
42	22	2B	15	E9	4B	D6	9A	4B	B5	B1	41
43	44	56	2A	CF	96	B1	29	96	77	7F	82
44	88	AC	54	83	31	7F	52	31	EE	FE	19
45	D	45	A8	1B	62	FE	A4	62	C1	E1	32
46	1A	8A	4D	36	C4	E1	55	C4	9F	DF	64
47	34	9	9A	6C	95	DF	AA	95	23	A3	C8
48	68	12	29	D8	37	A3	49	37	46	5B	8D
49	D0	24	52	AD	6E	5B	92	6E	8C	B6	7
50	BD	48	A4	47	DC	B6	39	DC	5	71	E
51	67	90	55	8E	A5	71	72	A5	A	E2	1C
52	CE	3D	AA	02	57	E2	E4	57	14	D9	38
53	81	7A	49	04	AE	D9	D5	AE	28	AF	70
54	1F	F4	92	08	41	AF	B7	41	50	43	E0
55	3E	F5	39	10	82	43	73	82	A0	86	DD
56	7C	F7	72	20	19	86	E6	19	5D	11	A7
57	F8	F3	E4	40	32	11	D1	32	BA	22	53
58	ED	FB	D5	80	64	22	BF	64	69	44	A6
59	C7	EB	B7	1D	C8	44	63	C8	D2	88	51
60	93	CB	73	3A	8D	88	C6	8D	B9	D	A2
61	3B	8B	E6	74	7	D	91	7	6F	1A	59
62	76	B	D1	E8	E	1A	3F	E	DE	34	B2
63	EC	16	BF	CD	1C	34	7E	1C	A1	68	79
64	C5	2C	63	87	38	68	FC	38	5F	D0	F2
65	97	58	C6	13	70	D0	E5	70	BE	BD	F9
66	33	B0	91	26	E0	BD	D7	E0	61	67	EF
67	66	7D	3F	4C	DD	67	B3	DD	C2	CE	C3
68	CC	FA	7E	98	A7	CE	7B	A7	99	81	9B
69	85	E9	FC	2D	53	81	F6	53	2F	1F	2B
70	17	CF	E5	5A	A6	1F	F1	A6	5E	3E	56
71	2E	83	D7	B4	51	3E	FF	51	BC	7C	AC
72	5C	1B	B3	75	A2	7C	E3	A2	65	F8	45
73	B8	36	7B	EA	59	F8	DB	59	CA	ED	8A
74	6D	6C	F6	C9	B2	ED	AB	B2	89	C7	9
75	DA	D8	F1	8F	79	C7	4B	79	F	93	12
76	A9	AD	FF	3	F2	93	96	F2	1E	3B	24
77	4F	47	E3	6	F9	3B	31	F9	3C	76	48
78	9E	8E	DB	C	EF	76	62	EF	78	EC	90
79	21	02	AB	18	C3	EC	C4	C3	F0	C5	3D
80	42	04	4B	30	9B	C5	95	9B	FD	97	7A
81	84	08	96	60	2B	97	37	2B	E7	33	F4
82	15	10	31	C0	56	33	6E	56	D3	66	F5

83	2A	20	62	9D	AC	66	DC	AC	BB	CC	F7
84	54	40	C4	27	45	CC	A5	45	6B	85	F3
85	A8	80	95	4E	8A	85	57	8A	D6	17	FB
86	4D	1D	37	9C	9	17	AE	9	B1	2E	EB
87	9A	3A	6E	25	12	2E	41	12	7F	5C	CB
88	29	74	DC	4A	24	5C	82	24	FE	B8	8B
89	52	E8	A5	94	48	B8	19	48	E1	6D	B
90	A4	CD	57	35	90	6D	32	90	DF	DA	16
91	55	87	AE	6A	3D	DA	64	3D	A3	A9	2C
92	AA	13	41	D4	7A	A9	C8	7A	5B	4F	58
93	49	26	82	B5	F4	4F	8D	F4	B6	9E	B0
94	92	4C	19	77	F5	9E	7	F5	71	21	7D
95	39	98	32	EE	F7	21	E	F7	E2	42	FA
96	72	2D	64	C1	F3	42	1C	F3	D9	84	E9
97	E4	5A	C8	9F	FB	84	38	FB	AF	15	CF
98	D5	B4	8D	23	EB	15	70	EB	43	2A	83
99	B7	75	7	46	CB	2A	E0	CB	86	54	1B
100	73	EA	E	8C	8B	54	DD	8B	11	A8	36
101	E6	C9	1C	5	B	A8	A7	B	22	4D	6C
102	D1	8F	38	A	16	4D	53	16	44	9A	D8
103	BF	3	70	14	2C	9A	A6	2C	88	29	AD
104	63	6	E0	28	58	29	51	58	D	52	47
105	C6	C	DD	50	B0	52	A2	B0	1A	A4	8E
106	91	18	A7	A0	7D	A4	59	7D	34	55	02
107	3F	30	53	5D	FA	55	B2	FA	68	AA	04
108	7E	60	A6	BA	E9	AA	79	E9	D0	49	08
109	FC	C0	51	69	CF	49	F2	CF	BD	92	10
110	E5	9D	A2	D2	83	92	F9	83	67	39	20
111	D7	27	59	B9	1B	39	EF	1B	CE	72	40
112	B3	4E	B2	6F	36	72	C3	36	81	E4	80
113	7B	9C	79	DE	6C	E4	9B	6C	1F	D5	1D
114	F6	25	F2	A1	D8	D5	2B	D8	3E	B7	3A
115	F1	4A	F9	5F	AD	B7	56	AD	7C	73	74
116	FF	94	EF	BE	47	73	AC	47	F8	E6	E8
117	E3	35	C3	61	8E	E6	45	8E	ED	D1	CD
118	DB	6A	9B	C2	02	D1	8A	02	C7	BF	87
119	AB	D4	2B	99	04	BF	9	04	93	63	13
120	4B	B5	56	2F	08	63	12	08	3B	C6	26
121	96	77	AC	5E	10	C6	24	10	76	91	4C
122	31	EE	45	BC	20	91	48	20	EC	3F	98
123	62	C1	8A	65	40	3F	90	40	C5	7E	2D
124	C4	9F	9	CA	80	7E	3D	80	97	FC	5A
125	95	23	12	89	1D	FC	7A	1D	33	E5	B4
126	37	46	24	F	3A	E5	F4	3A	66	D7	75
127	6E	8C	48	1E	74	D7	F5	74	CC	B3	EA

128	DC	5	90	3C	E8	B3	F7	E8	85	7B	C9
129	A5	A	3D	78	CD	7B	F3	CD	17	F6	8F
130	57	14	7A	F0	87	F6	FB	87	2E	F1	3
131	AE	28	F4	FD	13	F1	EB	13	5C	FF	6
132	41	50	F5	E7	26	FF	CB	26	B8	E3	C
133	82	A0	F7	D3	4C	E3	8B	4C	6D	DB	18
134	19	5D	F3	BB	98	DB	B	98	DA	AB	30
135	32	BA	FB	6B	2D	AB	16	2D	A9	4B	60
136	64	69	EB	D6	5A	4B	2C	5A	4F	96	C0
137	C8	D2	CB	B1	B4	96	58	B4	9E	31	9D
138	8D	B9	8B	7F	75	31	B0	75	21	62	27
139	7	6F	B	FE	EA	62	7D	EA	42	C4	4E
140	E	DE	16	E1	C9	C4	FA	C9	84	95	9C
141	1C	A1	2C	DF	8F	95	E9	8F	15	37	25
142	38	5F	58	A3	3	37	CF	3	2A	6E	4A
143	70	BE	B0	5B	6	6E	83	6	54	DC	94
144	E0	61	7D	B6	C	DC	1B	C	A8	A5	35
145	DD	C2	FA	71	18	A5	36	18	4D	57	6A
146	A7	99	E9	E2	30	57	6C	30	9A	AE	D4
147	53	2F	CF	D9	60	AE	D8	60	29	41	B5
148	A6	5E	83	AF	C0	41	AD	C0	52	82	77
149	51	BC	1B	43	9D	82	47	9D	A4	19	EE
150	A2	65	36	86	27	19	8E	27	55	32	C1
151	59	CA	6C	11	4E	32	02	4E	AA	64	9F
152	B2	89	D8	22	9C	64	04	9C	49	C8	23
153	79	F	AD	44	25	C8	08	25	92	8D	46
154	F2	1E	47	88	4A	8D	10	4A	39	7	8C
155	F9	3C	8E	D	94	7	20	94	72	E	5
156	EF	78	02	1A	35	E	40	35	E4	1C	A
157	C3	F0	04	34	6A	1C	80	6A	D5	38	14
158	9B	FD	08	68	D4	38	1D	D4	B7	70	28
159	2B	E7	10	D0	B5	70	3A	B5	73	E0	50
160	56	D3	20	BD	77	E0	74	77	E6	DD	A0
161	AC	BB	40	67	EE	DD	E8	EE	D1	A7	5D
162	45	6B	80	CE	C1	A7	CD	C1	BF	53	BA
163	8A	D6	1D	81	9F	53	87	9F	63	A6	69
164	9	B1	3A	1F	23	A6	13	23	C6	51	D2
165	12	7F	74	3E	46	51	26	46	91	A2	B9
166	24	FE	E8	7C	8C	A2	4C	8C	3F	59	6F
167	48	E1	CD	F8	5	59	98	5	7E	B2	DE
168	90	DF	87	ED	A	B2	2D	A	FC	79	A1
169	3D	A3	13	C7	14	79	5A	14	E5	F2	5F
170	7A	5B	26	93	28	F2	B4	28	D7	F9	BE
171	F4	B6	4C	3B	50	F9	75	50	B3	EF	61
172	F5	71	98	76	A0	EF	EA	A0	7B	C3	C2

173	F7	E2	2D	EC	5D	C3	C9	5D	F6	9B	99
174	F3	D9	5A	C5	BA	9B	8F	BA	F1	2B	2F
175	FB	AF	B4	97	69	2B	3	69	FF	56	5E
176	EB	43	75	33	D2	56	6	D2	E3	AC	BC
177	CB	86	EA	66	B9	AC	C	B9	DB	45	65
178	8B	11	C9	CC	6F	45	18	6F	AB	8A	CA
179	B	22	8F	85	DE	8A	30	DE	4B	9	89
180	16	44	3	17	A1	9	60	A1	96	12	F
181	2C	88	6	2E	5F	12	C0	5F	31	24	1E
182	58	D	C	5C	BE	24	9D	BE	62	48	3C
183	B0	1A	18	B8	61	48	27	61	C4	90	78
184	7D	34	30	6D	C2	90	4E	C2	95	3D	F0
185	FA	68	60	DA	99	3D	9C	99	37	7A	FD
186	E9	D0	C0	A9	2F	7A	25	2F	6E	F4	E7
187	CF	BD	9D	4F	5E	F4	4A	5E	DC	F5	D3
188	83	67	27	9E	BC	F5	94	BC	A5	F7	BB
189	1B	CE	4E	21	65	F7	35	65	57	F3	6B
190	36	81	9C	42	CA	F3	6A	CA	AE	FB	D6
191	6C	1F	25	84	89	FB	D4	89	41	EB	B1
192	D8	3E	4A	15	F	EB	B5	F	82	CB	7F
193	AD	7C	94	2A	1E	CB	77	1E	19	8B	FE
194	47	F8	35	54	3C	8B	EE	3C	32	B	E1
195	8E	ED	6A	A8	78	B	C1	78	64	16	DF
196	02	C7	D4	4D	F0	16	9F	F0	C8	2C	A3
197	04	93	B5	9A	FD	2C	23	FD	8D	58	5B
198	08	3B	77	29	E7	58	46	E7	7	B0	B6
199	10	76	EE	52	D3	B0	8C	D3	E	7D	71
200	20	EC	C1	A4	BB	7D	5	BB	1C	FA	E2
201	40	C5	9F	55	6B	FA	A	6B	38	E9	D9
202	80	97	23	AA	D6	E9	14	D6	70	CF	AF
203	1D	33	46	49	B1	CF	28	B1	E0	83	43
204	3A	66	8C	92	7F	83	50	7F	DD	1B	86
205	74	CC	5	39	FE	1B	A0	FE	A7	36	11
206	E8	85	A	72	E1	36	5D	E1	53	6C	22
207	CD	17	14	E4	DF	6C	BA	DF	A6	D8	44
208	87	2E	28	D5	A3	D8	69	A3	51	AD	88
209	13	5C	50	B7	5B	AD	D2	5B	A2	47	D
210	26	B8	A0	73	B6	47	B9	B6	59	8E	1A
211	4C	6D	5D	E6	71	8E	6F	71	B2	02	34
212	98	DA	BA	D1	E2	02	DE	E2	79	04	68
213	2D	A9	69	BF	D9	04	A1	D9	F2	08	D0
214	5A	4F	D2	63	AF	08	5F	AF	F9	10	BD
215	B4	9E	B9	C6	43	10	BE	43	EF	20	67
216	75	21	6F	91	86	20	61	86	C3	40	CE
217	EA	42	DE	3F	11	40	C2	11	9B	80	81

218	C9	84	A1	7E	22	80	99	22	2B	1D	1F
219	8F	15	5F	FC	44	1D	2F	44	56	3A	3E
220	3	2A	BE	E5	88	3A	5E	88	AC	74	7C
221	6	54	61	D7	D	74	BC	D	45	E8	F8
222	C	A8	C2	B3	1A	E8	65	1A	8A	CD	ED
223	18	4D	99	7B	34	CD	CA	34	9	87	C7
224	30	9A	2F	F6	68	87	89	68	12	13	93
225	60	29	5E	F1	D0	13	F	D0	24	26	3B
226	C0	52	BC	FF	BD	26	1E	BD	48	4C	76
227	9D	A4	65	E3	67	4C	3C	67	90	98	EC
228	27	55	CA	DB	CE	98	78	CE	3D	2D	C5
229	4E	AA	89	AB	81	2D	F0	81	7A	5A	97
230	9C	49	F	4B	1F	5A	FD	1F	F4	B4	33
231	25	92	1E	96	3E	B4	E7	3E	F5	75	66
232	4A	39	3C	31	7C	75	D3	7C	F7	EA	CC
233	94	72	78	62	F8	EA	BB	F8	F3	C9	85
234	35	E4	F0	C4	ED	C9	6B	ED	FB	8F	17
235	6A	D5	FD	95	C7	8F	D6	C7	EB	3	2E
236	D4	B7	E7	37	93	3	B1	93	CB	6	5C
237	B5	73	D3	6E	3B	6	7F	3B	8B	C	B8
238	77	E6	BB	DC	76	C	FE	76	B	18	6D
239	EE	D1	6B	A5	EC	18	E1	EC	16	30	DA
240	C1	BF	D6	57	C5	30	DF	C5	2C	60	A9
241	9F	63	B1	AE	97	60	A3	97	58	C0	4F
242	23	C6	7F	41	33	C0	5B	33	B0	9D	9E
243	46	91	FE	82	66	9D	B6	66	7D	27	21
244	8C	3F	E1	19	CC	27	71	CC	FA	4E	42
245	5	7E	DF	32	85	4E	E2	85	E9	9C	84
246	A	FC	A3	64	17	9C	D9	17	CF	25	15
247	14	E5	5B	C8	2E	25	AF	2E	83	4A	2A
248	28	D7	B6	8D	5C	4A	43	5C	1B	94	54
249	50	B3	71	7	B8	94	86	B8	36	35	A8
250	A0	7B	E2	E	6D	35	11	6D	6C	6A	4D
251	5D	F6	D9	1C	DA	6A	22	DA	D8	D4	9A
252	BA	F1	AF	38	A9	D4	44	A9	AD	B5	29
253	69	FF	43	70	4F	B5	88	4F	47	77	52
254	D2	E3	86	E0	9E	77	D	9E	8E	EE	A4

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Elem. GF (255) Of α^i	The ROMs Tables “HEX Form” Multiplying any elements from GF(255) with the relevant coefficients of the generator polynomial g(X)									
	g_{22} = α^{148}	g_{23} = α^{218}	g_{24} = α^{75}	g_{25} = α^{11}	g_{26} = α^{173}	g_{27} = α^{254}	g_{28} = α^{194}	g_{29} = α^{109}	g_{30} = α^8	g_{31} = α^{11}
-	00	00	00	00	00	00	00	00	00	00
0	52	2B	F	E8	F6	8E	32	BD	1D	E8
1	A4	56	1E	CD	F1	02	64	67	3A	CD
2	55	AC	3C	87	FF	04	C8	CE	74	87
3	AA	45	78	13	E3	08	8D	81	E8	13
4	49	8A	F0	26	DB	10	7	1F	CD	26
5	92	9	FD	4C	AB	20	E	3E	87	4C
6	39	12	E7	98	4B	40	1C	7C	13	98
7	72	24	D3	2D	96	80	38	F8	26	2D
8	E4	48	BB	5A	31	1D	70	ED	4C	5A
9	D5	90	6B	B4	62	3A	E0	C7	98	B4
10	B7	3D	D6	75	C4	74	DD	93	2D	75
11	73	7A	B1	EA	95	E8	A7	3B	5A	EA
12	E6	F4	7F	C9	37	CD	53	76	B4	C9
13	D1	F5	FE	8F	6E	87	A6	EC	75	8F
14	BF	F7	E1	3	DC	13	51	C5	EA	3
15	63	F3	DF	6	A5	26	A2	97	C9	6
16	C6	FB	A3	C	57	4C	59	33	8F	C
17	91	EB	5B	18	AE	98	B2	66	3	18
18	3F	CB	B6	30	41	2D	79	CC	6	30
19	7E	8B	71	60	82	5A	F2	85	C	60
20	FC	B	E2	C0	19	B4	F9	17	18	C0
21	E5	16	D9	9D	32	75	EF	2E	30	9D
22	D7	2C	AF	27	64	EA	C3	5C	60	27
23	B3	58	43	4E	C8	C9	9B	B8	C0	4E
24	7B	B0	86	9C	8D	8F	2B	6D	9D	9C
25	F6	7D	11	25	7	3	56	DA	27	25
26	F1	FA	22	4A	E	6	AC	A9	4E	4A
27	FF	E9	44	94	1C	C	45	4F	9C	94
28	E3	CF	88	35	38	18	8A	9E	25	35
29	DB	83	D	6A	70	30	9	21	4A	6A
30	AB	1B	1A	D4	E0	60	12	42	94	D4
31	4B	36	34	B5	DD	C0	24	84	35	B5
32	96	6C	68	77	A7	9D	48	15	6A	77
33	31	D8	D0	EE	53	27	90	2A	D4	EE
34	62	AD	BD	C1	A6	4E	3D	54	B5	C1
35	C4	47	67	9F	51	9C	7A	A8	77	9F
36	95	8E	CE	23	A2	25	F4	4D	EE	23
37	37	02	81	46	59	4A	F5	9A	C1	46

38	6E	04	1F	8C	B2	94	F7	29	9F	8C
39	DC	08	3E	5	79	35	F3	52	23	5
40	A5	10	7C	A	F2	6A	FB	A4	46	A
41	57	20	F8	14	F9	D4	EB	55	8C	14
42	AE	40	ED	28	EF	B5	CB	AA	5	28
43	41	80	C7	50	C3	77	8B	49	A	50
44	82	1D	93	A0	9B	EE	B	92	14	A0
45	19	3A	3B	5D	2B	C1	16	39	28	5D
46	32	74	76	BA	56	9F	2C	72	50	BA
47	64	E8	EC	69	AC	23	58	E4	A0	69
48	C8	CD	C5	D2	45	46	B0	D5	5D	D2
49	8D	87	97	B9	8A	8C	7D	B7	BA	B9
50	7	13	33	6F	9	5	FA	73	69	6F
51	E	26	66	DE	12	A	E9	E6	D2	DE
52	1C	4C	CC	A1	24	14	CF	D1	B9	A1
53	38	98	85	5F	48	28	83	BF	6F	5F
54	70	2D	17	BE	90	50	1B	63	DE	BE
55	E0	5A	2E	61	3D	A0	36	C6	A1	61
56	DD	B4	5C	C2	7A	5D	6C	91	5F	C2
57	A7	75	B8	99	F4	BA	D8	3F	BE	99
58	53	EA	6D	2F	F5	69	AD	7E	61	2F
59	A6	C9	DA	5E	F7	D2	47	FC	C2	5E
60	51	8F	A9	BC	F3	B9	8E	E5	99	BC
61	A2	3	4F	65	FB	6F	02	D7	2F	65
62	59	6	9E	CA	EB	DE	04	B3	5E	CA
63	B2	C	21	89	CB	A1	08	7B	BC	89
64	79	18	42	F	8B	5F	10	F6	65	F
65	F2	30	84	1E	B	BE	20	F1	CA	1E
66	F9	60	15	3C	16	61	40	FF	89	3C
67	EF	C0	2A	78	2C	C2	80	E3	F	78
68	C3	9D	54	F0	58	99	1D	DB	1E	F0
69	9B	27	A8	FD	B0	2F	3A	AB	3C	FD
70	2B	4E	4D	E7	7D	5E	74	4B	78	E7
71	56	9C	9A	D3	FA	BC	E8	96	F0	D3
72	AC	25	29	BB	E9	65	CD	31	FD	BB
73	45	4A	52	6B	CF	CA	87	62	E7	6B
74	8A	94	A4	D6	83	89	13	C4	D3	D6
75	9	35	55	B1	1B	F	26	95	BB	B1
76	12	6A	AA	7F	36	1E	4C	37	6B	7F
77	24	D4	49	FE	6C	3C	98	6E	D6	FE
78	48	B5	92	E1	D8	78	2D	DC	B1	E1
79	90	77	39	DF	AD	F0	5A	A5	7F	DF
80	3D	EE	72	A3	47	FD	B4	57	FE	A3
81	7A	C1	E4	5B	8E	E7	75	AE	E1	5B
82	F4	9F	D5	B6	02	D3	EA	41	DF	B6

83	F5	23	B7	71	04	BB	C9	82	A3	71
84	F7	46	73	E2	08	6B	8F	19	5B	E2
85	F3	8C	E6	D9	10	D6	3	32	B6	D9
86	FB	5	D1	AF	20	B1	6	64	71	AF
87	EB	A	BF	43	40	7F	C	C8	E2	43
88	CB	14	63	86	80	FE	18	8D	D9	86
89	8B	28	C6	11	1D	E1	30	7	AF	11
90	B	50	91	22	3A	DF	60	E	43	22
91	16	A0	3F	44	74	A3	C0	1C	86	44
92	2C	5D	7E	88	E8	5B	9D	38	11	88
93	58	BA	FC	D	CD	B6	27	70	22	D
94	B0	69	E5	1A	87	71	4E	E0	44	1A
95	7D	D2	D7	34	13	E2	9C	DD	88	34
96	FA	B9	B3	68	26	D9	25	A7	D	68
97	E9	6F	7B	D0	4C	AF	4A	53	1A	D0
98	CF	DE	F6	BD	98	43	94	A6	34	BD
99	83	A1	F1	67	2D	86	35	51	68	67
100	1B	5F	FF	CE	5A	11	6A	A2	D0	CE
101	36	BE	E3	81	B4	22	D4	59	BD	81
102	6C	61	DB	1F	75	44	B5	B2	67	1F
103	D8	C2	AB	3E	EA	88	77	79	CE	3E
104	AD	99	4B	7C	C9	D	EE	F2	81	7C
105	47	2F	96	F8	8F	1A	C1	F9	1F	F8
106	8E	5E	31	ED	3	34	9F	EF	3E	ED
107	02	BC	62	C7	6	68	23	C3	7C	C7
108	04	65	C4	93	C	D0	46	9B	F8	93
109	08	CA	95	3B	18	BD	8C	2B	ED	3B
110	10	89	37	76	30	67	5	56	C7	76
111	20	F	6E	EC	60	CE	A	AC	93	EC
112	40	1E	DC	C5	C0	81	14	45	3B	C5
113	80	3C	A5	97	9D	1F	28	8A	76	97
114	1D	78	57	33	27	3E	50	9	EC	33
115	3A	F0	AE	66	4E	7C	A0	12	C5	66
116	74	FD	41	CC	9C	F8	5D	24	97	CC
117	E8	E7	82	85	25	ED	BA	48	33	85
118	CD	D3	19	17	4A	C7	69	90	66	17
119	87	BB	32	2E	94	93	D2	3D	CC	2E
120	13	6B	64	5C	35	3B	B9	7A	85	5C
121	26	D6	C8	B8	6A	76	6F	F4	17	B8
122	4C	B1	8D	6D	D4	EC	DE	F5	2E	6D
123	98	7F	7	DA	B5	C5	A1	F7	5C	DA
124	2D	FE	E	A9	77	97	5F	F3	B8	A9
125	5A	E1	1C	4F	EE	33	BE	FB	6D	4F
126	B4	DF	38	9E	C1	66	61	EB	DA	9E
127	75	A3	70	21	9F	CC	C2	CB	A9	21

128	EA	5B	E0	42	23	85	99	8B	4F	42
129	C9	B6	DD	84	46	17	2F	B	9E	84
130	8F	71	A7	15	8C	2E	5E	16	21	15
131	3	E2	53	2A	5	5C	BC	2C	42	2A
132	6	D9	A6	54	A	B8	65	58	84	54
133	C	AF	51	A8	14	6D	CA	B0	15	A8
134	18	43	A2	4D	28	DA	89	7D	2A	4D
135	30	86	59	9A	50	A9	F	FA	54	9A
136	60	11	B2	29	A0	4F	1E	E9	A8	29
137	C0	22	79	52	5D	9E	3C	CF	4D	52
138	9D	44	F2	A4	BA	21	78	83	9A	A4
139	27	88	F9	55	69	42	F0	1B	29	55
140	4E	D	EF	AA	D2	84	FD	36	52	AA
141	9C	1A	C3	49	B9	15	E7	6C	A4	49
142	25	34	9B	92	6F	2A	D3	D8	55	92
143	4A	68	2B	39	DE	54	BB	AD	AA	39
144	94	D0	56	72	A1	A8	6B	47	49	72
145	35	BD	AC	E4	5F	4D	D6	8E	92	E4
146	6A	67	45	D5	BE	9A	B1	02	39	D5
147	D4	CE	8A	B7	61	29	7F	04	72	B7
148	B5	81	9	73	C2	52	FE	08	E4	73
149	77	1F	12	E6	99	A4	E1	10	D5	E6
150	EE	3E	24	D1	2F	55	DF	20	B7	D1
151	C1	7C	48	BF	5E	AA	A3	40	73	BF
152	9F	F8	90	63	BC	49	5B	80	E6	63
153	23	ED	3D	C6	65	92	B6	1D	D1	C6
154	46	C7	7A	91	CA	39	71	3A	BF	91
155	8C	93	F4	3F	89	72	E2	74	63	3F
156	5	3B	F5	7E	F	E4	D9	E8	C6	7E
157	A	76	F7	FC	1E	D5	AF	CD	91	FC
158	14	EC	F3	E5	3C	B7	43	87	3F	E5
159	28	C5	FB	D7	78	73	86	13	7E	D7
160	50	97	EB	B3	F0	E6	11	26	FC	B3
161	A0	33	CB	7B	FD	D1	22	4C	E5	7B
162	5D	66	8B	F6	E7	BF	44	98	D7	F6
163	BA	CC	B	F1	D3	63	88	2D	B3	F1
164	69	85	16	FF	BB	C6	D	5A	7B	FF
165	D2	17	2C	E3	6B	91	1A	B4	F6	E3
166	B9	2E	58	DB	D6	3F	34	75	F1	DB
167	6F	5C	B0	AB	B1	7E	68	EA	FF	AB
168	DE	B8	7D	4B	7F	FC	D0	C9	E3	4B
169	A1	6D	FA	96	FE	E5	BD	8F	DB	96
170	5F	DA	E9	31	E1	D7	67	3	AB	31
171	BE	A9	CF	62	DF	B3	CE	6	4B	62
172	61	4F	83	C4	A3	7B	81	C	96	C4

173	C2	9E	1B	95	5B	F6	1F	18	31	95
174	99	21	36	37	B6	F1	3E	30	62	37
175	2F	42	6C	6E	71	FF	7C	60	C4	6E
176	5E	84	D8	DC	E2	E3	F8	C0	95	DC
177	BC	15	AD	A5	D9	DB	ED	9D	37	A5
178	65	2A	47	57	AF	AB	C7	27	6E	57
179	CA	54	8E	AE	43	4B	93	4E	DC	AE
180	89	A8	02	41	86	96	3B	9C	A5	41
181	F	4D	04	82	11	31	76	25	57	82
182	1E	9A	08	19	22	62	EC	4A	AE	19
183	3C	29	10	32	44	C4	C5	94	41	32
184	78	52	20	64	88	95	97	35	82	64
185	F0	A4	40	C8	D	37	33	6A	19	C8
186	FD	55	80	8D	1A	6E	66	D4	32	8D
187	E7	AA	1D	7	34	DC	CC	B5	64	7
188	D3	49	3A	E	68	A5	85	77	C8	E
189	BB	92	74	1C	D0	57	17	EE	8D	1C
190	6B	39	E8	38	BD	AE	2E	C1	7	38
191	D6	72	CD	70	67	41	5C	9F	E	70
192	B1	E4	87	E0	CE	82	B8	23	1C	E0
193	7F	D5	13	DD	81	19	6D	46	38	DD
194	FE	B7	26	A7	1F	32	DA	8C	70	A7
195	E1	73	4C	53	3E	64	A9	5	E0	53
196	DF	E6	98	A6	7C	C8	4F	A	DD	A6
197	A3	D1	2D	51	F8	8D	9E	14	A7	51
198	5B	BF	5A	A2	ED	7	21	28	53	A2
199	B6	63	B4	59	C7	E	42	50	A6	59
200	71	C6	75	B2	93	1C	84	A0	51	B2
201	E2	91	EA	79	3B	38	15	5D	A2	79
202	D9	3F	C9	F2	76	70	2A	BA	59	F2
203	AF	7E	8F	F9	EC	E0	54	69	B2	F9
204	43	FC	3	EF	C5	DD	A8	D2	79	EF
205	86	E5	6	C3	97	A7	4D	B9	F2	C3
206	11	D7	C	9B	33	53	9A	6F	F9	9B
207	22	B3	18	2B	66	A6	29	DE	EF	2B
208	44	7B	30	56	CC	51	52	A1	C3	56
209	88	F6	60	AC	85	A2	A4	5F	9B	AC
210	D	F1	C0	45	17	59	55	BE	2B	45
211	1A	FF	9D	8A	2E	B2	AA	61	56	8A
212	34	E3	27	9	5C	79	49	C2	AC	9
213	68	DB	4E	12	B8	F2	92	99	45	12
214	D0	AB	9C	24	6D	F9	39	2F	8A	24
215	BD	4B	25	48	DA	EF	72	5E	9	48
216	67	96	4A	90	A9	C3	E4	BC	12	90
217	CE	31	94	3D	4F	9B	D5	65	24	3D

218	81	62	35	7A	9E	2B	B7	CA	48	7A
219	1F	C4	6A	F4	21	56	73	89	90	F4
220	3E	95	D4	F5	42	AC	E6	F	3D	F5
221	7C	37	B5	F7	84	45	D1	1E	7A	F7
222	F8	6E	77	F3	15	8A	BF	3C	F4	F3
223	ED	DC	EE	FB	2A	9	63	78	F5	FB
224	C7	A5	C1	EB	54	12	C6	F0	F7	EB
225	93	57	9F	CB	A8	24	91	FD	F3	CB
226	3B	AE	23	8B	4D	48	3F	E7	FB	8B
227	76	41	46	B	9A	90	7E	D3	EB	B
228	EC	82	8C	16	29	3D	FC	BB	CB	16
229	C5	19	5	2C	52	7A	E5	6B	8B	2C
230	97	32	A	58	A4	F4	D7	D6	B	58
231	33	64	14	B0	55	F5	B3	B1	16	B0
232	66	C8	28	7D	AA	F7	7B	7F	2C	7D
233	CC	8D	50	FA	49	F3	F6	FE	58	FA
234	85	7	A0	E9	92	FB	F1	E1	B0	E9
235	17	E	5D	CF	39	EB	FF	DF	7D	CF
236	2E	1C	BA	83	72	CB	E3	A3	FA	83
237	5C	38	69	1B	E4	8B	DB	5B	E9	1B
238	B8	70	D2	36	D5	B	AB	B6	CF	36
239	6D	E0	B9	6C	B7	16	4B	71	83	6C
240	DA	DD	6F	D8	73	2C	96	E2	1B	D8
241	A9	A7	DE	AD	E6	58	31	D9	36	AD
242	4F	53	A1	47	D1	B0	62	AF	6C	47
243	9E	A6	5F	8E	BF	7D	C4	43	D8	8E
244	21	51	BE	02	63	FA	95	86	AD	02
245	42	A2	61	04	C6	E9	37	11	47	04
246	84	59	C2	08	91	CF	6E	22	8E	08
247	15	B2	99	10	3F	83	DC	44	02	10
248	2A	79	2F	20	7E	1B	A5	88	04	20
249	54	F2	5E	40	FC	36	57	D	08	40
250	A8	F9	BC	80	E5	6C	AE	1A	10	80
251	4D	EF	65	1D	D7	D8	41	34	20	1D
252	9A	C3	CA	3A	B3	AD	82	68	40	3A
253	29	9B	89	74	7B	47	19	D0	80	74
254	52	2B	F	E8	F6	8E	32	BD	1D	E8

Table B1. The ROMs Tables

Appendix C: The output of the C code

- The generator polynomial

alpha_to[1]= 2	gp gg[0]= 45
alpha_to[2]= 4	gp gg[1]=216
alpha_to[3]= 8	gp gg[2]=239
alpha_to[4]= 16	gp gg[3]= 24
alpha_to[5]= 32	gp gg[4]=253
alpha_to[6]= 64	gp gg[5]=104
alpha_to[7]=128	gp gg[6]= 27
alpha_to[8]= 29	gp gg[7]= 40
alpha_to[9]= 58	gp gg[8]=107
alpha_to[10]=116	gp gg[9]= 50
alpha_to[11]=232	gp gg[10]=163
alpha_to[12]=205	gp gg[11]=210
alpha_to[13]=135	gp gg[12]=227
alpha_to[14]= 19	gp gg[13]=134
alpha_to[15]= 38	gp gg[14]=224
alpha_to[16]= 76	gp gg[15]=158
alpha_to[17]=152	gp gg[16]=119
alpha_to[18]= 45	gp gg[17]= 13
alpha_to[19]= 90	gp gg[18]=158
alpha_to[20]=180	gp gg[19]= 1
alpha_to[21]=117	gp gg[20]=238
alpha_to[22]=234	gp gg[21]=164
alpha_to[23]=201	gp gg[22]= 82
alpha_to[24]=143	gp gg[23]= 43
alpha_to[25]= 3	gp gg[24]= 15
alpha_to[26]= 6	gp gg[25]=232
alpha_to[27]= 12	gp gg[26]=246
alpha_to[28]= 24	gp gg[27]=142
alpha_to[29]= 48	gp gg[28]= 50
alpha_to[30]= 96	gp gg[29]=189
alpha_to[31]=192	gp gg[30]= 29
alpha_to[32]=157	gp gg[31]=232

- Look-up tables for GF ($2^{**}8$)

i in DEC Form	i in HEX Form	alpha_to[i]	index_of[i] in HEX Form	index_of[i] in DEC Form
0	0	1	FFFFFFFF	-1
1	1	2	0	0
2	2	4	1	1
3	3	8	19	25
4	4	10	2	2
5	5	20	32	50
6	6	40	1A	26
7	7	80	C6	198
8	8	1D	3	3
9	9	3A	DF	223
10	A	74	33	51
11	B	E8	EE	238
12	C	CD	1B	27
13	D	87	68	104
14	E	13	C7	199
15	F	26	4B	75
16	10	4C	4	4
17	11	98	64	100
18	12	2D	E0	224
19	13	5A	E	14
20	14	B4	34	52
21	15	75	8D	141
22	16	EA	EF	239
23	17	C9	81	129
24	18	8F	1C	28
25	19	3	C1	193
26	1A	6	69	105
27	1B	C	F8	248
28	1C	18	C8	200
29	1D	30	8	8
30	1E	60	4C	76
31	1F	C0	71	113
32	20	9D	5	5
33	21	27	8A	138
34	22	4E	65	101
35	23	9C	2F	47
36	24	25	E1	225
37	25	4A	24	36
38	26	94	F	15

39	27	35	21	33
40	28	6A	35	53
41	29	D4	93	147
42	2A	B5	8E	142
43	2B	77	DA	218
44	2C	EE	F0	240
45	2D	C1	12	18
46	2E	9F	82	130
47	2F	23	45	69
48	30	46	1D	29
49	31	8C	B5	181
50	32	5	C2	194
51	33	A	7D	125
52	34	14	6A	106
53	35	28	27	39
54	36	50	F9	249
55	37	A0	B9	185
56	38	5D	C9	201
57	39	BA	9A	154
58	3A	69	9	9
59	3B	D2	78	120
60	3C	B9	4D	77
61	3D	6F	E4	228
62	3E	DE	72	114
63	3F	A1	A6	166
64	40	5F	6	6
65	41	BE	BF	191
66	42	61	8B	139
67	43	C2	62	98
68	44	99	66	102
69	45	2F	DD	221
70	46	5E	30	48
71	47	BC	FD	253
72	48	65	E2	226
73	49	CA	98	152
74	4A	89	25	37
75	4B	F	B3	179
76	4C	1E	10	16
77	4D	3C	91	145
78	4E	78	22	34
79	4F	F0	88	136
80	50	FD	36	54
81	51	E7	D0	208
82	52	D3	94	148
83	53	BB	CE	206

84	54	6B	8F	143
85	55	D6	96	150
86	56	B1	DB	219
87	57	7F	BD	189
88	58	FE	F1	241
89	59	E1	D2	210
90	5A	DF	13	19
91	5B	A3	5C	92
92	5C	5B	83	131
93	5D	B6	38	56
94	5E	71	46	70
95	5F	E2	40	64
96	60	D9	1E	30
97	61	AF	42	66
98	62	43	B6	182
99	63	86	A3	163
100	64	11	C3	195
101	65	22	48	72
102	66	44	7E	126
103	67	88	6E	110
104	68	D	6B	107
105	69	1A	3A	58
106	6A	34	28	40
107	6B	68	54	84
108	6C	D0	FA	250
109	6D	BD	85	133
110	6E	67	BA	186
111	6F	CE	3D	61
112	70	81	CA	202
113	71	1F	5E	94
114	72	3E	9B	155
115	73	7C	9F	159
116	74	F8	A	10
117	75	ED	15	21
118	76	C7	79	121
119	77	93	2B	43
120	78	3B	4E	78
121	79	76	D4	212
122	7A	EC	E5	229
123	7B	C5	AC	172
124	7C	97	73	115
125	7D	33	F3	243
126	7E	66	A7	167
127	7F	CC	57	87
128	80	85	7	7

129	81	17	70	112
130	82	2E	C0	192
131	83	5C	F7	247
132	84	B8	8C	140
133	85	6D	80	128
134	86	DA	63	99
135	87	A9	D	13
136	88	4F	67	103
137	89	9E	4A	74
138	8A	21	DE	222
139	8B	42	ED	237
140	8C	84	31	49
141	8D	15	C5	197
142	8E	2A	FE	254
143	8F	54	18	24
144	90	A8	E3	227
145	91	4D	A5	165
146	92	9A	99	153
147	93	29	77	119
148	94	52	26	38
149	95	A4	B8	184
150	96	55	B4	180
151	97	AA	7C	124
152	98	49	11	17
153	99	92	44	68
154	9A	39	92	146
155	9B	72	D9	217
156	9C	E4	23	35
157	9D	D5	20	32
158	9E	B7	89	137
159	9F	73	2E	46
160	A0	E6	37	55
161	A1	D1	3F	63
162	A2	BF	D1	209
163	A3	63	5B	91
164	A4	C6	95	149
165	A5	91	BC	188
166	A6	3F	CF	207
167	A7	7E	CD	205
168	A8	FC	90	144
169	A9	E5	87	135
170	AA	D7	97	151
171	AB	B3	B2	178
172	AC	7B	DC	220
173	AD	F6	FC	252

174	AE	F1	BE	190
175	AF	FF	61	97
176	B0	E3	F2	242
177	B1	DB	56	86
178	B2	AB	D3	211
179	B3	4B	AB	171
180	B4	96	14	20
181	B5	31	2A	42
182	B6	62	5D	93
183	B7	C4	9E	158
184	B8	95	84	132
185	B9	37	3C	60
186	BA	6E	39	57
187	BB	DC	53	83
188	BC	A5	47	71
189	BD	57	6D	109
190	BE	AE	41	65
191	BF	41	A2	162
192	C0	82	1F	31
193	C1	19	2D	45
194	C2	32	43	67
195	C3	64	D8	216
196	C4	C8	B7	183
197	C5	8D	7B	123
198	C6	7	A4	164
199	C7	E	76	118
200	C8	1C	C4	196
201	C9	38	17	23
202	CA	70	49	73
203	CB	E0	EC	236
204	CC	DD	7F	127
205	CD	A7	C	12
206	CE	53	6F	111
207	CF	A6	F6	246
208	D0	51	6C	108
209	D1	A2	A1	161
210	D2	59	3B	59
211	D3	B2	52	82
212	D4	79	29	41
213	D5	F2	9D	157
214	D6	F9	55	85
215	D7	EF	AA	170
216	D8	C3	FB	251
217	D9	9B	60	96
218	DA	2B	86	134

219	DB	56	B1	177
220	DC	AC	BB	187
221	DD	45	CC	204
222	DE	8A	3E	62
223	DF	9	5A	90
224	E0	12	CB	203
225	E1	24	59	89
226	E2	48	5F	95
227	E3	90	B0	176
228	E4	3D	9C	156
229	E5	7A	A9	169
230	E6	F4	A0	160
231	E7	F5	51	81
232	E8	F7	B	11
233	E9	F3	F5	245
234	EA	FB	16	22
235	EB	EB	EB	235
236	EC	CB	7A	122
237	ED	8B	75	117
238	EE	B	2C	44
239	EF	16	D7	215
240	F0	2C	4F	79
241	F1	58	AE	174
242	F2	B0	D5	213
243	F3	7D	E9	233
244	F4	FA	E6	230
245	F5	E9	E7	231
246	F6	CF	AD	173
247	F7	83	E8	232
248	F8	1B	74	116
249	F9	36	D6	214
250	FA	6C	F4	244
251	FB	D8	EA	234
252	FC	AD	A8	168
253	FD	47	50	80
254	FE	8E	58	88

Results for Reed-Solomon code ($n = 255$, $k = 223$, $t = 16$) (data, recd in polynomial form)

The first 32 is the parity check digits and the rest are the message 223 to form 255 codeword

i	data[i]	recd[i](decoded)
0	120	120
1	41	41
2	190	190
3	87	87
4	41	41
5	197	197
6	214	214
7	196	196
8	192	192
9	17	17
10	239	239
11	31	31
12	208	208
13	221	221
14	2	2
15	196	196
16	251	251
17	31	31
18	17	17
19	171	171
20	240	240
21	225	225
22	164	164
23	61	61
24	184	184
25	155	155
26	22	22
27	239	239
28	17	17
29	65	65
30	237	237
31	104	104
32	223	223
33	222	222
34	221	221
35	220	220
36	219	219
37	218	218

38	217	217
39	216	216
40	215	215
41	214	214
42	213	213
43	212	212
44	211	211
45	210	210
46	209	209
47	208	208
48	207	207
49	206	206
50	205	205
51	204	204
52	203	203
53	202	202
54	201	201
55	200	200
56	199	199
57	198	198
58	197	197
59	196	196
60	195	195
61	194	194
62	193	193
63	192	192
64	191	191
65	190	190
66	189	189
67	188	188
68	187	187
69	186	186
70	185	185
71	184	184
72	183	183
73	182	182
74	181	181
75	180	180
76	179	179
77	178	178
78	177	177
79	176	176
80	175	175
81	174	174
82	173	173

83	172	172
84	171	171
85	170	170
86	169	169
87	168	168
88	167	167
89	166	166
90	165	165
91	164	164
92	163	163
93	162	162
94	161	161
95	160	160
96	159	159
97	158	158
98	157	157
99	156	156
100	155	155
101	154	154
102	153	153
103	152	152
104	151	151
105	150	150
106	149	149
107	148	148
108	147	147
109	146	146
110	145	145
111	144	144
112	143	143
113	142	142
114	141	141
115	140	140
116	139	139
117	138	138
118	137	137
119	136	136
120	135	135
121	134	134
122	133	133
123	132	132
124	131	131
125	130	130
126	129	129
127	128	128

128	127	127
129	126	126
130	125	125
131	124	124
132	123	123
133	122	122
134	121	121
135	120	120
136	119	119
137	118	118
138	117	117
139	116	116
140	115	115
141	114	114
142	113	113
143	112	112
144	111	111
145	110	110
146	109	109
147	108	108
148	107	107
149	106	106
150	105	105
151	104	104
152	103	103
153	102	102
154	101	101
155	100	100
156	99	99
157	98	98
158	97	97
159	96	96
160	16	95
161	94	94
162	93	93
163	92	92
164	91	91
165	90	90
166	89	89
167	88	88
168	87	87
169	86	86
170	85	85
171	84	84
172	83	83

173	82	82
174	81	81
175	80	80
176	79	79
177	78	78
178	77	77
179	76	76
180	75	75
181	74	74
182	73	73
183	72	72
184	71	71
185	70	70
186	69	69
187	68	68
188	67	67
189	66	66
190	65	65
191	64	64
192	63	63
193	62	62
194	61	61
195	60	60
196	59	59
197	58	58
198	57	57
199	56	56
200	55	55
201	54	54
202	53	53
203	52	52
204	51	51
205	50	50
206	49	49
207	48	48
208	47	47
209	46	46
210	45	45
211	44	44
212	43	43
213	42	42
214	41	41
215	40	40
216	39	39
217	38	38

218	37	37
219	36	36
220	35	35
221	34	34
222	33	33
223	32	32
224	31	31
225	30	30
226	29	29
227	28	28
228	27	27
229	26	26
230	25	25
231	24	24
232	23	23
233	22	22
234	21	21
235	20	20
236	19	19
237	18	18
238	17	17
239	16	16
240	15	15
241	14	14
242	13	13
243	12	12
244	11	11
245	10	10
246	9	9
247	8	8
248	7	7
249	6	6
250	5	5
251	4	4
252	3	3
253	2	2
254	1	1