

Algorithmic C (AC) Datatypes

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Chapter 1: Overview of Algorithmic C Datatypes

This package provides numerical and interface C++ classes aimed at modeling behavior that is targeted to the design of hardware. Modeling of bit-accurate arithmetic is crucial for hardware design. The numerical package provides classes for bit-accurate integer, fixed-point, floating-point and complex numbers. Well-defined semantics and simulation speed are the focus of this package since they are essential for the verification of hardware designs. The package currently includes the following files:

File	Package	Description	Class	AC Dependencies
ac_int.h	Numerical	Integer	ac_int<W,S>	
ac_fixed.h		Fixed-point	ac_fixed<W,I,S,Q,O>	ac_int.h
ac_float.h		Floating-point	ac_float<W,I,E,Q>	ac_fixed.h
ac_complex.h		Complex	ac_complex<T>	ac_float.h
ac_sc.h		AC-SystemC Conversions, Tracing		ac_complex.h
ac_pp.py		Pretty-print for gdb		gdb with python
ac_channel.h	Interface	Channel Fifo	ac_channel<T>	

The parameters for the classes are:

- W: integer representing width of the type. For ac_float it is width of the mantissa.
- S: bool parameter representing signedness of ac_int or ac_fixed type.
- I: integer representing integer width.
- Q,O: enumeration parameter for quantization (rounding) and overflow modes.
- T: type. For ac_complex, T is restricted to AC Numerical types and C++ integer and floating types.

The following sections overview the [Numerical](#) and [Interface](#) Datatypes.

1.1. Overview of Numerical Algorithmic C Datatypes

The advantages of the Algorithmic C numerical datatypes are the following:

- *Arbitrary-Length*: This allows a clean definition of the semantics for all operators that is not tied to an implementation limit. It is also important for writing general IP algorithms that don't have artificial (and often hard to quantify and document) limits for precision. Whenever possible, the return type of operators do not loose any precision. In cases where that is not the case, for example in fixed-point division, what is kept is fully defined and implementation independent.
- *Precise, Uniform and Consistent Definition of Semantics*: Special attention has been paid to define and verify the simulation semantics and to make sure that the semantics are appropriate for synthesis.

No simulation behavior has been left to compiler dependent behavior. Mixed type operators, including mixed operators with C integer types are fully defined to prevent any ambiguity. Also, asserts have been introduced to catch invalid code during simulation. See also “[User Defined Asserts](#)”.

- *Simulation Speed*: The implementation of `ac_int` uses sophisticated template specialization techniques so that a regular C++ compiler can generate optimized assembly language that will run much faster than the equivalent SystemC datatypes. For example, `ac_int` of bit widths in the range 1 to 32 can run 100x faster than the corresponding `sc_bigint/sc_biguint` datatype and 3x faster than the corresponding `sc_int/sc_uint` datatype. In addition the compilation is faster and produces smaller binary executables.
- *Correctness*: The simulation and synthesis semantics have been verified for many size combinations using a combination of simulation and equivalence checking.

1.1.1. Usage of Numerical AC Datatypes

In order to use the Algorithmic C data types, the appropriate header needs to be included in the source. All the definitions are in the AC header files and there are no object files that need to be linked in. Enabling compiler optimizations (for example “-O3” in GCC) is critical to the fastest runtime.

1.1.2. Usage of Numerical AC Datatypes within SystemC

Numerical AC Datatypes can be used in SystemC descriptions. If they are used as types for SystemC `sc_signal`, then the header file `ac_sc.h` needs be included after the `systemc.h` is included. The `ac_sc.h` file includes all the other numerical types by including `ac_complex.h`.

The `ac_sc.h` header file also provides explicit conversion functions to SystemC integer and fixed-point datatypes. The “[Datatype Migration Guide](#)” chapter presents information about how to convert algorithms written with SystemC datatypes to Algorithmic C datatypes.

1.1.3. Definition and Implementation Overview

The numerical datatypes were defined and implemented adhering to the following guiding principles:

- *Static Bit Widths*: all operations and methods return an object with a bit width that is statically determinable from the bit widths of the inputs and “signedness” (signed vs. unsigned) of the inputs. Keeping bit-widths static is essential for fast simulation, as it means that memory allocation is completely avoided. It is also essential for synthesis. For example the left shift operation of an `ac_int` returns an `ac_int` of the same type (width and signedness) as the type of first operand. In contrast, the left shift for `sc_bigint` or `sc_biguint` returns an object with precision that depends on the shift value and has no practical bound on its bit width.
- *Operations Defined Arithmetically*: whenever possible, operations are defined arithmetically, that is, the inputs are treated as arithmetic values and the result value is returned with a type (bitwidth and signedness) that is capable of representing the value without loss of precision. Exceptions to this rule are the shift operators (to maintain static bit widths) and division.

- *Compiler Independent Semantics*: the semantics avoid “implementation dependent” behaviors that are present for some native C integer operations. For example, shift values for a C int needs to be in the range [0,31] and are otherwise implementation dependent.
- *Mixed numerical ac types and C integer type Binary Functions*: all binary operators are defined for mixed *ac_int*, *ac_fixed*, *ac_float* and *ac_complex* and native C integers for consistency. For example the expression “1 + a” where *a* is an *ac_int<36,true>* will compute “*ac_int<32,true>* 1 + a” rather than “1 + (*int*) a”. This is done to ensure that expressions are carried out without unintentional loss of precision and to make sure that compiler errors due to ambiguities are avoided.
- *Uninitialized by Default*: there is no default initialization. This is preferable for both runtime and synthesis where uninitialized variables are treated as don't care. There is a utility function to uninitialized static arrays to undo the initialization due to the “static” qualifier.

The types have a public interface. The base class implementation of these types are private. They are part of the implementation and may be changed. Also any function or class under namespace *ac_private* should not be used as it may be subject to change.

1.1.4. Implementation Guidelines

The following guidelines are followed in the implementation:

1. Minimal inclusion of other header files. For example *systemc.h* should only be included in *ac_sc.h*. No inclusion of std libraries, boost etc.
2. No C++11 features or dependencies.
3. No memory allocation is used. Most code is written to be synthesizable. The exceptions are related to non-synthesizable types such as float and double.
4. No virtual functions used. Given that the types are templated and fully defined in the header files, there is no reason for their use.
5. Warnings that need to be turned off are done so only for the header context using the push/pop mechanisms available in Visual C++ and GCC version 4.6 or higher.

1.1.5. Implementation Assumptions

While the C++ standard does not formally define the bit width and representation of the various integer types, compilers for general software have converged to both widths and representation that are assumed in this package. The most important assumption is that the int type is represented in a 32-bit, 2's complement representation. The width of other integer types are relevant when converting to and from those types and in the definition of mixed AC and C++ integer types. The assumptions are as follows:

1. Two's complement is used for all signed integer versions.
2. The char, signed char and unsigned char have a bitwidth of 8. The type char is treated as signed.
3. The short and unsigned short have a bitwidth of 16 and is represented in 2's complement.

4. The long and unsigned long have a bitwidth as defined in `std::numeric_limits`. In general 32-bit platforms represent them as 32-bit integers and 64-bit platforms represent them as 64-bit integers.
5. The long long and unsigned long long have a bitwidth of 64-bits.

1.2. Overview of Interface Algorithmic C Datatypes

In addition to the numerical datatypes, an interface class `ac_channel` is also provided. This class simplifies the modeling and synthesis of hierarchical designs using C++ function calls. The `ac_channel` class is a C++ template class that enforces a FIFO discipline (reads occur in the same order as writes.) From a modeling perspective, an `ac_channel` is implemented as a simple interface to the C++ standard queue (`std::deque`). That is, for modeling purposes, an `ac_channel` is infinite in length (writes always succeed) and attempting to read from an empty channel generates an assertion failure (reads are nonblocking).

Chapter 2: Arbitrary-Length Bit-Accurate Integer and Fixed-Point Datatypes

The arbitrary-length bit-accurate integer and fixed-point datatypes provide an easy way to model static bit-precision with minimal runtime overhead. Operators and methods on both the integer and fixed-point types are clearly and consistently defined so that they have well defined simulation and synthesis semantics.

The types are named *ac_int* and *ac_fixed* and their numerical ranges are given by their template parameters as shown in Table 1. For both types, the boolean template parameter determines whether the type is constrained to be unsigned or signed. The template parameter *W* specifies the number of bits for the integer or fixed point number and must be a positive integer. For *ac_int*, the value of the integer number $b_{W-1} \dots b_1 b_0$ is interpreted as an unsigned integer or a signed (two's complement) number. The advantage of having the signedness specified by a template parameter rather than having two distinct types is that it makes it possible to write generic functions where signedness is just a parameter.

Table 1: Numerical Ranges of *ac_int* and *ac_fixed*

Type	Description	Numerical Range	Quantum
<i>ac_int</i> <W, false>	unsigned integer	0 to $2^W - 1$	1
<i>ac_int</i> <W, true>	signed integer	-2^{W-1} to $2^{W-1} - 1$	1
<i>ac_fixed</i> <W, I, false>	unsigned fixed-point	0 to $(1 - 2^{-W}) 2^I$	2^{I-W}
<i>ac_fixed</i> <W, I, true>	signed fixed-point	$(-0.5) 2^I$ to $(0.5 - 2^{-W}) 2^I$	2^{I-W}

For *ac_fixed*, the second parameter *I* of an *ac_fixed* is an integer that determines the location of the fixed-point relative to the MSB. The value of the fixed-point number $b_{W-1} \dots b_1 b_0$ is given by $(b_{W-1} \dots b_1 b_0) 2^I$ or equivalently $(b_{W-1} \dots b_1 b_0) 2^{I-W}$ where $b_{W-1} \dots b_1 b_0$ is interpreted as an unsigned integer or a signed (two's complement) number.

Table 2 shows examples for various integer and fixed-point types with their respective numerical ranges and quantum values. The quantum is the smallest difference between two numbers that are represented. Note that an *ac_fixed*<W,W,S> (that is $I=W$) has the same numerical range as an *ac_int*<W,S> where *S* is a boolean value that determines whether the type is signed or unsigned. The numerical range of an *ac_fixed*<W,I,S> is equal to the numerical range of and *ac_int*<W,S> (or an *ac_fixed*<W,W,S>) multiplied by the quantum.

Table 2: Examples for *ac_int* and *ac_fixed*

Type	Numerical Range	Quantum
<i>ac_int</i> <1, false>	0 to 1	1
<i>ac_int</i> <1, true>	-1 to 0	1
<i>ac_int</i> <4, false>	0 to 15	1
<i>ac_int</i> <4, true>	-8 to 7	1

<code>ac_fixed<4, 4, false></code>	0 to 15	1
<code>ac_fixed<4, 4, true></code>	-8 to 7	1
<code>ac_fixed<4, 6, false></code>	0 to 60	4
<code>ac_fixed<4, 6, true></code>	-32 to 28	4
<code>ac_fixed<4, 0, false></code>	0 to 15/16	1/16
<code>ac_fixed<4, 0, true></code>	-0.5 to 7/16	1/16
<code>ac_fixed<4,-1, false></code>	0 to 15/32	1/32
<code>ac_fixed<4,-1, true></code>	-0.25 to 7/32	1/32

It is important to remember when dealing with either an `ac_fixed` or an `ac_int` that in order for both +1 and -1 to be in the numerical range, I and W have to be at least 2. For example, `ac_fixed<6,1,true>` has a range from -1 to +0.96875 (it does not include +1) while `ac_fixed<6,2,true>` has a range -2 to +1.9375 (includes +1).

The fixed-point datatype `ac_fixed` has two additional template parameters that are optional that define the overflow mode (e.g. saturation) and the quantization mode (e.g. rounding):

```
ac_fixed<int W, int I, bool S, ac_q_mode Q, ac_o_mode O>
```

Quantization and overflow occur when assigning (=, +=, etc.) or constructing (including casting) where the target does not represent the source value without loss of precision (this will be covered more precisely in “[Quantization and Overflow](#)”). In all the examples of [Table 2](#) the default quantization and overflow modes AC_TRN and AC_WRAP are implied. The default modes simply throw away bits to the right of LSB and to the left of the MSB which is also the behavior of `ac_int`:

```
ac_fixed<1,1,true> x = 1;    // range is [-1,0], +1 wraps around to -1
ac_int<1,true>      x = 1;    // same as above
ac_fixed<4,4,true> x = 9;    // range is [-8,7], +9 wraps around to -7
ac_int<4,true>      x = 9;    // same as above
ac_fixed<4,4,true> x = 3.7;  // truncated to 3.0
ac_int<4,true>      x = 3.7;  // same as above
ac_fixed<4,4,true> x = -3.2; // truncated to -4.0
ac_int<4,true>      x = -3.2; // same as above
```

[Table 3](#) shows the operators defined for both `ac_int` and `ac_fixed`

Table 3: Operators defined for `ac_int` and `ac_fixed`.

Operators	<code>ac_int</code>	<code>ac_fixed</code>
Two operand +, -, *, /, %, , &, ^	Arithmetic result. First or second arg may be C INT or <code>ac_fixed</code> / truncates towards 0	Arithmetic result. First or second arg may be <code>ac_int</code> or C INT / truncates towards 0. % NOT DEFINED
>>, <<	bidirectional return type is type of first operand Second arg is <code>ac_int</code> or C INT	bidirectional return type is type of first operand Second arg is <code>ac_int</code> or C INT

=	assignment	quantization, then overflow handling specified by target
+=, -=, *=, /=, %=, =, &=, ^=, >>=, <<=	Equiv to op then assign. First arg is <i>ac_int</i>	Equiv to op then assign
==, !=, >, <, >=, <=	First or second arg may be <i>C INT</i> or <i>ac_fixed</i>	First or second arg may be <i>ac_int</i> or <i>C INT</i> or <i>C double</i>
Unary +, -, ~	Arithmetic	Arithmetic
++x, x++, --x, x--	Pre/post incr/dec by 1	Pre/post incr/dec by 2I-W
! x	Equiv to x == 0	Equiv to x == 0
(long long)	defined for <i>ac_int</i> <W,true>, $W \leq 64$	NOT DEFINED
(unsigned long long)	defined for <i>ac_int</i> <W,false>, $W \leq 64$	NOT DEFINED
x[i]	returns <i>ac_int::ac_bitref</i> index: <i>ac_int</i> , <i>unsigned</i> , <i>int</i> asserts for index out of bound	returns <i>ac_fixed::ac_bitref</i> index: <i>ac_int</i> , <i>unsigned</i> , <i>int</i> asserts for index out of bounds

Note that for convenience the conversion operators to (*long long*) and signed *ac_int* and (*unsigned long long*) for unsigned *ac_int* are defined for $W \leq 64$. Among other things, this allows for the use of an *ac_int* as an index to a C array without any explicit conversion call.

Table 4 shows the methods defined for *ac_int* and *ac_fixed* types. The *slc* and *set_slc* methods are templated to get or set a slice respectively. For *slc*, the width needs to be explicitly provided as a template argument. When using the *slc* method in a templated function place the keyword *template* before it as some compilers may error out during parsing.

For example:

```
template<int N> // not important whether or not N is used
int f(int x) {
    ac_int<32,true> t = x;
    ac_int<6,true> r = t.template slc<6>(4); // t.slc<6>(4) could error out
    return r.to_int();
}
```

The *set_slc* method does not need to have a width specified as a template argument since the width is inferred from the width of the argument x. Many of the other methods are conversion functions. The *length* method returns the width of the type. The *set_val* method sets the *ac_int* or *ac_fixed* to a value that depends on the template parameter.

Table 4: Methods defined for *ac_int* and *ac_fixed*.

Methods	<i>ac_int</i> <W,S>	<i>ac_fixed</i> <W,I,S,Q,O>
<code>slc<W2>(int_type i)</code>	Returns slice of width W2 starting at bit index i, in other words slice (W2-1+i downto i). Slice is returned as an <i>ac_int</i> <W2,S>. Parameter i needs to be non-negative and could of any of the following types: <i>ac_int</i> , <i>unsigned</i> , <i>int</i>	
<code>set_slc(int_type i, ac_int<W2,S2> x)</code>	Bits of x are copied at slice with LSB index i. That is, bits (W2-1+i downto i) are set with bits of x. Parameter i needs to be non-negative and could of any of the following types: <i>ac_int</i> , <i>unsigned</i> , <i>int</i>	
<code>to_ac_int()</code>	NOT DEFINED	return an <i>ac_int</i> <W _I ,S> where W _I is max(I, 1). Equiv to AC_TRN quantization. Return type guarantees no overflows.
<code>to_int(), to_uint(), to_long(), to_ulong(), to_int64(), to_uint64()</code>	Conversions to various C INTs	Conversions to various C INTs Equiv to <code>to_ac_int()</code> followed by conversion
<code>to_double()</code>	Conversion to <i>double</i>	Conversion to <i>double</i>
<code>to_string(ac_base_mode base_rep, bool sign_mag = false)</code>	convert to std::string depending on parameters <i>base_rep</i> {AC_HEX, AC_DEC, AC_OCT, AC_BIN} and <i>sign_mag</i>	
<code>length()</code>	Returns bitwidth (value of template parameter W)	
<code>set_val<ac_special_val>()</code>	Set to special value specified by template parameter AC_VAL_DC, AC_VAL_0, AC_VAL_MIN, AC_VAL_MAX, AC_VAL_QUANTUM. See Table 18 for details.	
<code>leading_sign()</code>	If unsigned (S==false) returns number of leading zeros. If signed (S==true) returns number of leading 0/1s minus one. Parameter <i>all_sign</i> is true if return value is W-S (all 0's for S==false, all same for S==true). See Leading Sign and Normalization for details.	
<code>leading_sign(bool &all_sign)</code>		
<code>normalize(ac_int<WE,SE> &exp)</code>	Treats the <i>ac_int/ac_fixed</i> as a mantissa and attempts to normalize. The mantissa shifted partially (not normalized) if exp saturates to its minimum value. See " Leading Sign and Normalization " for details.	

normalize_RME(ac_int<WE,SE> &exp)	Same as above but minimum exponent is reserved (<i>Reserve Minimum Exponent</i>) so <i>exp</i> saturates at minimum exponent plus one. See " Leading Sign and Normalization " for details.
bit_complement()	Returns the raw complemented bits in an unsigned result that has the same width (and integer width for ac_fixed). It is not equivalent to the ~ operator which return an arithmetic result.
bit_fill(const int (&ivec)[Na], bool bigendian=true)	Use the raw bits in the integer vector to fill ac_int/ac_fixed. Missing most significant bits are set to 0. Extra bits are ignored. Element ivec[0] is treated as the most/least significant bits when bigendian is true/false.
bit_fill_hex(const char *str)	Use the hex characters 0-9a-fA-F in string to fill ac_int/ac_fixed. Missing most significant bits are set to 0. Extra bits are ignored. Asserts on non hex characters.

Table 5 shows the reduce methods for *ac_int*. They take the *W* bus of the *ac_int* and apply the respective logical operation on them and return *bool*.

Table 5: Reduce Methods defined for *ac_int*

Reduce Methods	<i>ac_int</i> <W,S>
and_reduce()	AND of bits
or_reduce()	OR of bits
xor_reduce()	XOR of bits

Table 6 shows the constructors that are defined for *ac_int* and *ac_fixed*. When constructing an *ac_fixed*, its quantization/overflow mode is taken into account. Initializing an *ac_int* or *ac_fixed* from floating-point (*float* or *double*) is not as runtime efficient as initializing from integers.

Table 6: Constructors defined for *ac_int* and *ac_fixed*.

Constructor argument	<i>ac_int</i>	<i>ac_fixed</i>
None: Default	does not initialize <i>ac_int</i> declared static are init to 0 by C++ before constructor is called	does not initialize <i>ac_fixed</i> declared static are init to 0 by C++ before constructor is called
bool (1-bit unsigned)		quantization/overflow
char (8-bit signed)		quantization/overflow
signed/unsigned char (8-bit signed/unsigned)		quantization/overflow
signed/unsigned short (16-bit signed/unsigned)		quantization/overflow
signed/unsigned int or long (32-bit signed/unsigned)		quantization/overflow

signed/unsigned long long (64-bit signed/unsigned)		quantization/overflow
<i>double</i>	Not as efficient	quantization/overflow Not as efficient
<i>ac_int</i>		quantization/overflow
<i>ac_fixed</i>	NOT DEFINED Use to <i>ac_int</i> () instead	quantization/overflow

2.1. Quantization and Overflow

The fixed-point type *ac_fixed* provides quantization and overflow modes that determine how to adjust the value when either of the two conditions occur:

- *Quantization*: bits to the right of the LSB of the target type are being lost. The value may be adjusted by the following two strategies:
 - *Rounding*: choose the closest quantization level. When the value is exactly half way two quantization levels, which one is chosen depends on the specific rounding mode as shown in [Table 7](#).
 - *Truncation*: choose the closest quantization level such that result (quantized value) is less than or equal the source value (truncation toward minus infinity) or such that the absolute value of the result is less than or equal the source value (truncation towards zero).

Note that quantization may trigger an overflow so it is always applied before overflow handling.

- *Overflow*: value after quantization is outside the range of the target as defined in [Table 1](#), except when the overflow mode is AC_SAT_SYM where the range is symmetric: $[-2^{W-1}+1, 2^{W-1}-1]$ in which case the most negative number -2^{W-1} triggers an overflow.

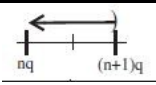
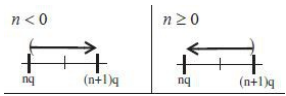
The modes are specified by the 4th and 5th template argument to *ac_fixed*:

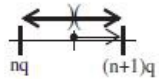
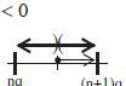
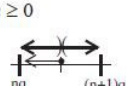
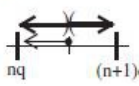
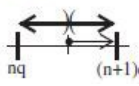
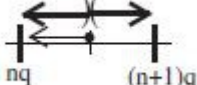
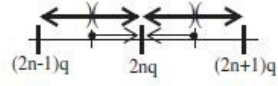
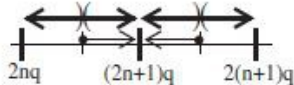
```
ac_fixed<int W, int I, bool S, ac_q_mode Q, ac_o_mode O>
```

that are of enumeration type *ac_q_mode* and *ac_o_mode* respectively. The enumeration values for *ac_q_mode* are shown in [Table 7](#). The enumeration values for *ac_o_modes* are shown in [Table 8](#). The quantization and overflow mode default to AC_TRN and AC_WRAP:

```
ac_fixed<8,4,true> x; // equiv to ac_fixed<8,4,true,AC_TRN,AC_WRAP>
```

Table 7: Quantization modes for *ac_fixed*

Modes	Behavior n is integer, q is 2^{I-W}	Simulation/Synthesis cost
AC_TRN (default) (trunc towards $-\infty$)		Delete bits, no cost except for \neq (div assign) signed
AC_TRN_ZERO (trunc towards 0)		No cost for \neq , or unsigned src For signed: incrementer, OR for deleted bits, AND with sign bit

AC_RND (round towards $+\infty$)		<p>Various forms differ only on the direction of the rounding for values that are exactly at half point.</p> <p>All require an incrementer, some require to OR deleted bits, some only require to look at the MSB of the deleted bits.</p>
AC_RND_ZERO (round towards 0)	<div style="display: flex; justify-content: space-around;"> <div style="text-align: center;"> $n < 0$  </div> <div style="text-align: center;"> $n \geq 0$  </div> </div>	
AC_RND_INF (rounds towards $\pm\infty$)	<div style="display: flex; justify-content: space-around;"> <div style="text-align: center;"> $n < 0$  </div> <div style="text-align: center;"> $n \geq 0$  </div> </div>	
AC_RND_MIN_INF (round towards $-\infty$)		
AC_RND_CONV (round towards even q multiples)		
AC_RND_CONV_ODD (round towards odd q multiples)		

For unsigned *ac_fixed* types, AC_TRN and AC_TRN_ZERO are equivalent, AC_RND and AC_RND_INF are equivalent, and AC_RND_ZERO and AC_RND_MIN_INF are equivalent. The AC_RND_CONV quantization rounds towards even multiples of the quantization while the AC_RND_CONV_ODD quantization rounds towards odd multiples of the quantization. The quantization modes that have two columns (different directions for negative and positive numbers) are symmetric around 0 and are more costly as *ac_fixed* is represented in two's complement arithmetic. On the other hand signed-magnitude representations (for example floating point numbers) are more costly for asymmetric cases.

Quantization and overflow occur when assigning or constructing.

```
ac_fixed<8,1,true,Q,0> x = -0.1; // quantization, no overflow
ac_fixed<8,1,false,Q,0> y = x; // overflow (underflow) as y is unsigned
ac_fixed<4,1,true,Q,0> z = x; // quantization (dropping bits on the right)
(ac_fixed<4,1,true,Q,0>) x; // casting: same as above
ac_fixed<8,4,true,Q,0> a = ...;
ac_fixed<8,4,true,Q,0> b =
```

The behavior of the overflow modes are shown in Table 8. The default is AC_WRAP and requires no special handling (same behavior as with an *ac_int*). The AC_SAT mode saturates to the MIN or MAX limits of the range (as specified in [Table 1](#)) of the target type (different for signed or unsigned targets). The AC_SAT_ZERO sets the value to zero when an overflow is detected. The AC_SAT_SYM makes only sense for signed targets. It saturates to +MAX or -MAX (note that -MAX = MIN+q). It not only saturates on overflow, but also when the value is MIN (it excludes the most negative number that would make the range asymmetric).

The following operators can destroy the symmetric saturation property of an *ac_fixed* with AC_SAT_SYM and should be avoided:

- Changing a bit with operator `[]`.
- Changing bit(s) with the `set_slc` method.

- Performing a shift assign (<<=, >>=) does not trigger quantization or overflow handling.
- Performing a shift (<<, >>) returns a result of type of the first operand. If the first operand is an `ac_fixed` with `AC_SAT_SYM`, the result type will be an `ac_fixed` with `AC_SAT_SYM`, but it is not guaranteed to be symmetrically saturated.

Once the symmetrical saturation property has been destroyed, assignment to the same type will not trigger symmetrical saturation as the following example illustrates:

```
typedef ac_fixed<8,8,true,AC_TRN,AC_SAT_SYM> fx_ss;
typedef ac_fixed<8,8,true> fx;
fx_ss a = 0;
a[7] = 1; // No longer symmetrically saturated

fx_ss b = a; // b remains non symmetrically saturated as a is assumed to be
             // symmetrically saturated and has identical type

a = (fx) a; // forcing saturation by first casting to non-symmetrically
            // saturated type
```

A similar example with shifting:

```
typedef ac_fixed<8,8,true,AC_TRN,AC_SAT_SYM> fx_ss;
fx_ss a = 1;
a <<= 7; // Value of a is not symmetrically saturated
fx_ss b = 1;
b = b << 7; // Value of b is not symmetrically saturated as return type of
            // b << 7 is ac_fixed<8,8,true,AC_TRN,AC_SAT_SYM>
a = b; // Value of a remains non symmetrically saturated
```

As the above example illustrates `a <<= v`, is equivalent to `a = a << v`.

Table 8: Overflow modes for *ac_fixed*

Mode	Behavior all references are to <i>target</i> type MIN, MAX are limits as in Table 1	Simulation/Synthesis cost
AC_WRAP (default)	Drop bits to the left of MSB	No cost
AC_SAT	Saturate to closest of MIN or MAX	Overflow checking and Saturation logic
AC_SAT_ZERO	Set to 0 on overflow	Overflow checking and Saturation logic
AC_SAT_SYM	For unsigned: treat as AC_SAT, For signed: on overflow or number is MIN set to closest of . Note: if source is of type AC_SAT_SYM, it will be treated as already symmetrically saturated.	Overflow checking and Saturation logic

2.2. Using the `ac_int` and `ac_fixed` Datatypes

In order to use the `ac_int` datatype the following file include should be used:

```
#include <ac_int.h>
```

The `ac_int` type is implemented with two template parameters to define its bitwidth and to indicate whether it is signed or unsigned:

```
ac_int<7, true> x;    // x is 7 bits signed
ac_int<19, false> y;  // y is 19 bits unsigned
```

In order to use the `ac_fixed` datatype the following file include should be used:

```
#include <ac_fixed.h>
```

The `ac_fixed.h` includes `ac_int.h` so it is not necessary to include both `ac_int.h` and `ac_fixed.h`.

The `ac_fixed` type is implemented with 5 template parameters that control the behavior of the fixed point type:

```
ac_fixed<int W, int I, bool S, ac_q_mode Q, ac_q_mode O>
```

where `W` is the width of the fixed point type, `I` is the number of integer bits, `S` is a boolean flag that determines whether the fixed-point is signed or unsigned, and `Q` and `O` are the quantization and overflow modes respectively (as shown in [Table 4](#) and [Table 6](#)). The value of the fixed point is given by:

For example:

```
ac_fixed<4,4,true> x; // bbbb signed, AC_TRN, AC_WRAP
ac_fixed<4,0,false> x; // .bbbb unsigned AC_TRN, AC_WRAP
ac_fixed<4,7,false> x; // bbbb000, unsigned, AC_TRN, AC_WRAP
ac_fixed<4,-3,false> x; // .bbbb * pow(2, -3), unsigned, AC_TRN, AC_WRAP
```

2.3. Operators and Methods

This section provides a more detailed specification of the behavior of operators and methods including precisely defining return types. The operators and methods that are defined for `ac_int` and `ac_fixed` can be classified in some broad categories:

- Binary (two operand) operators:
 - Binary Arithmetic and Logical Operators and arithmetic and logical assign operators: `+`, `-`, `*`, `/`, `%`, `&`, `|`, `^`, `+=`, `-=`, `*=`, `/=`, `%=`, `&=`, `|=`, `^=`. The modulo operators `%` and `%=` are not defined for `ac_fixed`. Mixing of `ac_int`, `ac_fixed` and native C integers is allowed.
 - Relational Operators: the result is a boolean value (true/false): `>`, `<`, `>=`, `<=`, `==`, `!=`. Mixing of `ac_int`, `ac_fixed`, native C integers and `double` is allowed.
 - Shift Operators and shift assign operators: `<<`, `>>`, `=<<`, `=>>`. The second argument is an `ac_int` or a native C integer.
- Unary Operators: `+`, `-`, `~` and `!` (one operand). The `!` operator returns `bool`.
- Bit Complement
- Pre/Post Increment and Decrement Operators: `++x`, `--x`, `x++`, `x--`.

- Bit Select Operator: `[]`, returns an `ac_int::ac_bitref` or `ac_fixed::ac_bitref`. Allows reading and modifying bits of an `ac_int` or `ac_fixed`.
- Slice Read Method: `slc` and Slice Write Method: `set_slc` to read and modify a slice in an `ac_int` or `ac_fixed`. A slice of an `ac_fixed` is an `ac_int`.
- Conversion Operators to C Integer Types and Explicit Conversion Methods to C native types.
- The `set_val` Method
- Constructors from `ac_int` and C native types.
- Methods to Fill Bits
- IO Methods
- Mixing `ac_int` and `ac_fixed` with Other Datatypes.

The concatenation operator is not defined for `ac_int`. Bit reversal may be defined in future releases.

2.3.1. Binary Arithmetic and Logical Operators

The two operand arithmetic and logical operators return an `ac_fixed` if either operand is an `ac_fixed`, otherwise the return type is `ac_int`. Binary arithmetic operators “+”, “*”, “/” and “%” and logical operators “&”, “|” and “^” return a signed `ac_int/ac_fixed` if either of the two operands is of type signed. The “-” operator always returns an `ac_int/ac_fixed` of type signed. The result for all operands with the exception of division is computed arithmetically and the bit width (and integer bit width for `ac_fixed`) of the result is such that the result is represented without loss of precision. The “/” operator is defined for both `ac_int` and `ac_fixed` and it returns a type that guarantees that the result does not overflow (see [Table 9](#) and [Table 10](#)). The operator “%” is only defined for `ac_int`. Division by zero is not defined and will generate an exception.

The binary operators “&”, “|” and “^” return the bitwise “and”, “or” and “xor” of the two operands. The return type is signed if either of the two operands is signed. The two operands are treated arithmetically. For instance, if the operands are `ac_fixed`, the fixed point is aligned just as it is done for addition. Then operands are extended, if necessary, so that both operands are represented in the same type which is also the return type.

The arithmetic definition of the “bitwise” operators has the advantage that when mixing `ac_int` (or `ac_fixed`) operands of different lengths and signedness, the operations are associative:

```
(a | b) | c
```

returns the same value (and in this case the same type) as

```
a | (b | c)
```

Also operators are consistent

```
~(a | b) == ~a & ~b
```

[Table 9](#) shows the list of binary (two operand) arithmetic and logical operators for `ac_int` and the return type based on the signedness and bit width of the two input operands. All operators shown in the table are defined arithmetically. The operator `&` could have been defined to return a more constrained type, $S_R = S_1 \& S_2$ and $W_R = \text{abs}(\min(S_1 ? -W_1 : W_1, S_2 ? -W_2 : W_2))$. For instance, the bitwise AND of a `uint1` and an `int5` would return a `uint1`. However, for simplicity it has been defined to be consistent with the other two logical operators.

Regardless of how the operators are defined, synthesis will reduce it to the smallest size that preserves the arithmetic value of the result.

Table 9: Return Types for *ac_int* Binary Arithmetic and Bitwise Logical Operations

Operator	Return Type: <i>ac_int</i> <W _R ,S _R >	
	S _R	Bit Width: W _R
+	S ₁ S ₂	max(W ₁ +!S ₁ &S ₂ , W ₂ +!S ₂ &S ₁)+1
-	true	max(W ₁ +!S ₁ &S ₂ , W ₂ +!S ₂ &S ₁)+1
*	S ₁ S ₂	W ₁ +W ₂
/	S ₁ S ₂	W ₁ +W ₂
%	S ₁	min(W ₁ , W ₂ +!S ₂ &S ₁)
&	S ₁ S ₂	max(W ₁ +!S ₁ &S ₂ , W ₂ +!S ₂ &S ₁)
	S ₁ S ₂	max(W ₁ +!S ₁ &S ₂ , W ₂ +!S ₂ &S ₁)
^	S ₁ S ₂	max(W ₁ +!S ₁ &S ₂ , W ₂ +!S ₂ &S ₁)

Table 10 shows the binary (two operand) arithmetic and logical operators for *ac_fixed* and the return type based on the signedness, bit width and integer bit width of the operands. All operands are defined consistently with *ac_int*: if both *ac_fixed* operands are pure integers (W and I are the same) then the result is an *ac_fixed* that is also a pure integer with the same bitwidth and value as the result of the equivalent *ac_int* operation. For example: a/b where a is an *ac_fixed*<8,8> and b is an *ac_fixed*<5,5> returns an *ac_fixed*<8,8>. In SystemC, on the other hand, the result of a/b returns 64 bits of precision (or SC_FXDIV_WL if defined).

Table 10: Return Types for *ac_fixed* Binary Arithmetic and Bitwise Logical Operations

Operator	Return Type: <i>ac_fixed</i> <W _R ,I _R ,S _R ,AC_TRN,AC_WRAP>		
	S _R	Bit Width: W _R	Integer Bit Width: I _R
+	S ₁ S ₂	I _R +max(W ₁ -I ₁ , W ₂ -I ₂)	max(I ₁ +!S ₁ &S ₂ , I ₂ +!S ₂ &S ₁)+1
-	true	I _R +max(W ₁ -I ₁ , W ₂ -I ₂)	max(I ₁ +!S ₁ &S ₂ , I ₂ +!S ₂ &S ₁)+1
*	S ₁ S ₂	W ₁ +W ₂	I ₁ +I ₂
/	S ₁ S ₂	W ₁ +max(W ₂ -I ₂ , 0)+S ₂	I ₁ +(W ₂ -I ₂)+S ₂
&	S ₁ S ₂	I _R +max(W ₁ -I ₁ , W ₂ -I ₂)	max(I ₁ +!S ₁ &S ₂ , I ₂ +!S ₂ &S ₁)
	S ₁ S ₂	I _R +max(W ₁ -I ₁ , W ₂ -I ₂)	max(I ₁ +!S ₁ &S ₂ , I ₂ +!S ₂ &S ₁)
^	S ₁ S ₂	I _R +max(W ₁ -I ₁ , W ₂ -I ₂)	max(I ₁ +!S ₁ &S ₂ , I ₂ +!S ₂ &S ₁)

The assignment operators +=, -=, *=, /=, %=, &=, |= and ^= have the usual semantics:

A1 @= A2

where @ is any of the operators +, -, *, /, %, &, | and ^ is equivalent in behavior to:

A1 = A1 @ A2

From a simulation speed point of view, the assignment version (for instance `*=`) is more efficient since the target precision can be taken into account to reduce the computation required.

Mixed *ac_int*, *ac_fixed* and C Integer Operators

Binary (two operand) operations that mix *ac_int*, *ac_fixed* and native C integer operands are defined to avoid ambiguity in the semantics or compilation problems due to multiple operators matching an operation. For example, assuming `x` is an *ac_int*, `1+x` gives the same result as `x+1`. The return type is determined by the following rules where `c_int` is a native C type, `width(c_int)` is the bitwidth of the C type, and `signedness(c_int)` is the signedness of the C type:

- If one of the operands is an *ac_fixed* in a binary operation or the first operand is an *ac_fixed* in an assign operation, the other operand is represented as an *ac_fixed*:
 - *ac_int*<W,S> gets represented as *ac_fixed*<W,W,S>
 - `c_int` gets represented as *ac_fixed*<width(`c_int`), width(`c_int`), signedness(`c_int`)>
- Otherwise, if one of the operands is an *ac_int* in a binary operation or the first operand is an *ac_int* in an assign operation, the other operand (native c integer) gets represented as *ac_int*<width(`c_int`), signedness(`c_int`)>

The rules above guarantee that precision is not lost. Note that floating point types are not supported for the operators in this section as the output precision can not be determined by the C compiler. Table 11 shows a few examples of mixed operations. The variables in this table are defined as:

```
ac_int<7,true>      i_s7;
ac_fixed<20,4,false> fx_u20_4;
signed char         c_s8;
```

Table 11: Mixed Expressions Example

Expression	Equivalent Expression
<code>1 + i_s7</code>	<code>(ac_int<32,true>) 1 + i_s7</code>
<code>(bool) 1 + i_s7</code>	<code>(ac_int<1,false>) 1 + i_s7</code>
<code>i_7s + fx_u20_4</code>	<code>(ac_fixed<7,7,true>) i_s7 + fx_u20_4</code>
<code>fx_u20_4 += c_s8</code>	<code>fx_u20_4 += (ac_fixed<8,8,true>) c_s8</code>
<code>c_s8 += fx_u20_4</code>	<code>c_s8 += (signed char) fx_u20_4</code>

Mixed *ac_int* and C pointer for + and - Operators

The operator `+` is defined for *ac_int* and C pointer (and vice versa) so that an *ac_int* can be added to a C pointer. The operator `-` is defined so that an *ac_int* can be subtracted from a C pointer. The result is, in all cases, of the same type as the C pointer.

2.3.2. Relational Operators

Relational operators `!=`, `==`, `>`, `>=`, `<` and `<=` are also binary operations and have some of the same

characteristics described for arithmetic and logical operations: the operations are done arithmetically and mixed *ac_int*, *ac_fixed* and native C integer operators are defined. The return type is bool.

The relational operator for *ac_int* and *ac_fixed* with the C floating type *double* is also defined for convenience, though for simulation performance reasons it is best to store the *double* constant in an appropriate *ac_int* or *ac_fixed* variable outside computation loops so that the overhead of converting the *double* to *ac_fixed* or *ac_int* is minimized.

2.3.3. Shift Operators

Left shift “<<” and right shift “>>” operators return a value of type of the first operand. The left shift operator shifts in zeros. The right shift operator shifts in the MSB bit for *ac_int/ac_fixed* of type signed, 0 for *ac_int/ac_fixed* integers of type unsigned.

If the shift value is negative the first operand is shifted in the opposite direction by the absolute value of the shift value (this is also the semantic of *sc_fixed/sc_ufixed* shifts). Shift values that are greater than W (bitwidth of first operand) are equivalent to shifting by W.

The second operand is an *ac_int* integer of bit width less or equal to 32 bits or a signed or unsigned int.

The shift assign operators “<<=” and “>>=” have the usual semantics:

```
A1 <<= A2; // equiv to A1 = A1 << A2, except for A1 is ac_fixed with
           // AC_SAT_SYM
A1 >>= A2; // equiv to A1 = A1 >> A2, except for A1 is ac_fixed with
           // AC_SAT_SYM and value of A2 < 0
```

Because the return type is the type of the first operand, the shift assign operators do not carry out any quantization or overflow.

Mixed *ac_int*, *ac_fixed* and C Integer

All shift operators are defined for mixed *ac_int*, *ac_fixed* (first operand) and native C integer operands. For example:

```
(short int) x << (ac_int<8,true>) y
```

matches the overloaded operator “<<” that is implemented as follows:

```
(ac_int<16,false>) x << (ac_int<8,true>) y
```

The shift assign operators <<= and >>= are also defined for mixed *ac_int* (first or second operand), *ac_fixed* (first operand) and native C integer (second operand).

Differences with SystemC *sc_bigint/sc_bignint* Types

- The return type of the left shift for *sc_bigint/sc_bignint* or *sc_fixed/sc_ufixed* does not lose bits making the return type of the left shift data dependent (dependent on the shift value). Shift assigns for *sc_fixed/sc_ufixed* may result in quantization or overflow (depending on the mode of the first operand).
- Negative shifts are equivalent to a zero shift value for *sc_bigint/sc_bignint*

Differences with Native C Integer Types

- Shifting occurs on either 32-bit (*int*, *unsigned int*) or 64-bit (*long long*, *unsigned long long*) integrals. If the first operand is an integral type that has less than 32 bits (*bool*, *(un)signed char*, *short*) it is first promoted to *int*. The return type is the type of the first argument after integer promotion (if applicable).
- Shift values are constrained according to the length of the type of the promoted first operand.
 - $0 \leq s < 32$ for 32-bit numbers
 - $0 \leq s < 64$ for 64-bit numbers
- The behavior for shift values outside the allowed ranges is not specified by the C++ ISO standard.

2.3.4. Unary Operators: +, -, ~ and !

Unary “+” and “-” have the usual semantics: “+x” returns x, “-x” returns “0-x”.

The unary operator “~x” returns the arithmetic one’s complement of “x”. The one’s complement is mathematically defined for integers as $-x-1$ (that is $-x+x == -1$). This is equivalent to the bitwise complement of x of a signed representation of x (if x is unsigned, add one bit to represent it as a signed number). The return type is signed and has the bitwidth of x if x is signed and $\text{bitwidth}(x)+1$ if x is unsigned.

The ! operator return true if the *ac_int/ac_fixed* is zero, false otherwise.

Table 12 lists the unary operators and their return types.

Table 12: Unary Operators for *ac_int*<W,S>

Operator	Return Type
+	<i>ac_int</i> <W, S>
-	<i>ac_int</i> <W+1, true>
~	<i>ac_int</i> <W+!S, true>
!	bool

Table 13 lists the unary operators for *ac_fixed* and their return types.

Table 13: Unary Operators for *ac_fixed*<W,I,S,Q,O>

Operator	Return Type
+	<i>ac_fixed</i> <W, I, S>
-	<i>ac_int</i> <W+1, I+1, true>
~	<i>ac_int</i> <W+!S, I+!S, true>
!	bool

2.3.5. Bit Complement

The *bit_complement* member function for *ac_int* and *ac_fixed* are provided as an alternative to the operator \sim which grows by one bit when applied to an unsigned type. The methods are:

```
ac_int<W, false> ac_int<W,S>::bit_complement() const;
ac_fixed<W, I, false> ac_fixed<W,I,S,Q,O>::bit_complement() const;
```

They return an unsigned version of the same *W* (and same *I* for *ac_fixed*). This a bit complement of the raw bits as compared to the complement operator \sim that returns an arithmetic value of $-x-1$ for *ac_int* and $-x-2^{I-W}$ for *ac_fixed*. The following example illustrates the difference:

```
ac_int<3,false> x = 7;    // 111
ac_int<5,true> y;
y = ~x;    // returns -7 - 1 = -8 (1000) as ac_int<4,true>, y = 11000
y = x.bit_complement(); // returns 000 as ac_int<3,false>, y = 00000
ac_int<4,false> x2 = 7;   // 0111
y = ~x2;    // returns -7 - 1 = -8 (11000) as ac_int<5,true>, y = 11000
y = x2.bit_complement(); // returns 1000 as ac_int<4,false>, y = 01000
```

2.3.6. Increment and Decrement Operators

Pre/Post increment/decrement for *ac_int* have the usual semantics as shown in Table 14 (*T_x* is the type of variable *x*).

Table 14: Pre- and Post-Increment/Decrement Operators for *ac_int*

Operator	Equivalent Behavior
<i>x</i> ++	<i>T_x</i> t = <i>x</i> ; <i>x</i> += 1; return t;
++ <i>x</i>	<i>x</i> += 1; return reference to <i>x</i> ;
<i>x</i> --	<i>T_x</i> t = <i>x</i> ; <i>x</i> -= 1; return t;
-- <i>x</i>	<i>x</i> -= 1; return reference to <i>x</i> ;

Pre/Post increment/decrement for *ac_fixed* have the semantics as shown in Table 15 (*T_x* is the type of variable *x*) where *q* is the quantum value of the representation (the smallest difference between two values for *T_x*). This definition is consistent with the definition of *ac_int* where *q* is 1.

Table 15: Pre- and Post-Increment/Decrement Operators for *ac_fixed*<*W,I,S,Q,O*> where $q=2^{I-W}$

Operator	Equivalent Behavior
<i>x</i> ++	<i>T_x</i> t = <i>x</i> ; <i>x</i> += <i>q</i> ; return t;
++ <i>x</i>	<i>x</i> += <i>q</i> ; return reference to <i>x</i> ;
<i>x</i> --	<i>T_x</i> t = <i>x</i> ; <i>x</i> -= <i>q</i> ; return t;
-- <i>x</i>	<i>x</i> -= <i>q</i> ; return reference to <i>x</i> ;

2.3.7. Conversion Operators to C Integer Types

A limited number of conversion operators to C integer types (including bool) are provided by the *ac_int* datatype, as described in the following list. The *ac_fixed* datatype provides no conversion operator to C integer types.

- *ac_int*<W,S> for $W > 64$ has no conversion operators to any C integer type
- *ac_int*<W,true> for $W \leq 64$ has only the "long long" conversion operator
- *ac_int*<W,false> for $W \leq 64$ has only the "unsigned long long" conversion operator

Some coding styles may encounter compilation problems due to the lack of conversion operators. The most common problem is the absence of the conversion to bool for bit widths beyond 64 for *ac_int* and for all bit widths for *ac_fixed*. Table 16 shows some typical scenarios:

Table 16: Conversion to C Integer Types

<i>ac_int</i> <33,true> k = ...;	
if (k)	OK, conversion first to long long then to bool
if ((bool) k)	OK, same as above
switch (k)	OK, operator to long long
switch (2*k)	ERROR: Result of expression is <i>ac_int</i> <65,true> (constant 2 treated as <i>ac_int</i> <32,true>) No conversion to any C integer from <i>ac_int</i> <65,true>
switch (2*(int)k)	OK, conversion first to long long, then to int
a[k]	OK, operator to long long
a[2*k]	ERROR: Result of expression is <i>ac_int</i> <65,true> (constant 2 treated as <i>ac_int</i> <32,true>) No conversion to any C integer from <i>ac_int</i> <65,true>
a[2*(int)k]	OK, conversion first to long long, then to int
<i>ac_int</i> <80, true> k = ...;	
if (k);	ERROR, no conversion operator defined
if ((bool) k)	ERROR, same as above
if (!! k)	OK, operator ! defined
if (k != 0)	OK, operator != defined
switch (k)	ERROR: No conversion operator defined
a[k]	ERROR: No conversion to any C integer from <i>ac_int</i> <80,true>
<i>ac_fixed</i> <3, 3,true> x = ...;	

<code>if (x)</code>	ERROR: No conversion operator defined for any W
<code>if (!! x)</code>	OK, operator ! defined
<code>if (x != 0)</code>	OK, operator != defined

When writing parameterized IP where the bit-widths of some *ac_int* is parameterized, code that may compile for some parameters, may not compile for a different set of parameters. In such cases, it is important to not rely on the conversion operator.

2.3.8. Explicit Conversion Methods

Methods to covert to C signed and unsigned integer types *int*, *long* and *Slong* are provided for both *ac_int* and *ac_fixed* as shown in Table 17. The methods **to_int()**, **to_long()**, **to_int64()**, **to_uint()**, **to_uint64()** and **to_ulong()** are defined for both *ac_int* and *ac_fixed* (same functions are also defined for *sc_bigint/sc_biguint*). The method **to_double()** is also defined for both *ac_int* and *ac_fixed*. The method **to_ac_int()** is defined for *ac_fixed*.

Table 17: Explicit Conversion Methods for *ac_int/ac_fixed*

Method	Types	Return Type
to_int()	<i>ac_int/ac_fixed</i>	<i>int</i>
to_uint()	<i>ac_int/ac_fixed</i>	<i>unsigned int</i>
to_long()	<i>ac_int/ac_fixed</i>	<i>long</i>
to_ulong()	<i>ac_int/ac_fixed</i>	<i>unsigned long</i>
to_int64()	<i>ac_int/ac_fixed</i>	<i>Slong</i>
to_uint64()	<i>ac_int/ac_fixed</i>	<i>Ulong</i>
to_double()	<i>ac_int/ac_fixed</i>	<i>double</i>
to_ac_int()	<i>ac_fixed</i> only	<i>ac_int</i> <max(I,1), S>

2.3.9. Bit Select Operator: []

Bit select is accomplished with the [] operator:

```
y[k] = x[i];
```

The [] operator does not return an *ac_int*, but rather it returns an object of class *ac_int::ac_bitref* that stores the index and a reference to the *ac_int* object that is being indexed.

The conversion function to “bool” (operator *bool*) is defined so that a bit reference may be used where a bool type is required:

```
while( y[k] && z[m] ) {}  
z = y[k] ? a : b;
```

A bit reference may be assigned an integer. The behavior is that the least significant bit of the integer is assigned to the bit reference. For example if *n* is type *int* and *x* is type *ac_int* then the following three assignments have the same behavior:

```
x[k] = n;
x[k] = (ac_int<1,false>) n;
x[k] = 1 & n;
```

The conversion to any *ac_int* is provided and it equivalent to first converting to a bool or to a *ac_int<1,false>*:

```
ac_int<5,false> x = y[0]; // equivalent to x = (bool) y[0]
```

The *ac_bitref::operator=(int val)* returns the bit reference so that assignment chains work as expected:

```
x[k] = z[m] = true; // assigns 1 to z[m] and to x[k]
```

Out of Bounds Behavior

It is invalid to access (read or write) a bit outside the range $[0, W-1]$ where W is the width of the *ac_int* being accessed. Simulation will assert on such cases. See also “[User Defined Asserts](#)” in the *Frequently Asked Questions* section.

2.3.10. Slice Read Method: *slc*

Slice read is accomplished with the template method *slc<W>(int lsb)*:

```
x = y.slc<2>(5);
```

which is equivalent to the VHDL behavior:

```
x := y(6 downto 5);
```

The two arguments to the *slc* method are defined as:

- The bit length of slice W : this is template argument (the length of the slice is constrained to be static so that the length of the slice is known at compile time. The length of the slice must be greater or equal to 1.
- The bit position of the LSB of the slice *slc_lsb*.

The *slc* method returns an *ac_int* of length W and signedness of the *ac_int* being sliced.

Out of Bounds Slice Reads

Accessing a bit to the left of the MSB of the *ac_int<W,S>* (index $\leq W$) is allowed and is defined as if the *ac_int* had been first extended (sign extension for signed, 0 padding for unsigned) so that the index is within range. This is consistent with treating *ac_int* as an arithmetic value.

Attempting to access (read) a bit with a negative index has undefined behavior and is considered to be the product of an erroneous program. If such a negative index read is encountered during execution (simulation) an assert will be triggered. See also “[User Defined Asserts](#)” in the *Frequently Asked Questions* section.

Differences with SystemC *sc_bigint/sc_bignint* Types

The range method and the part select operator in SystemC are fundamentally different than the *ac_int* *slc* and *set_slc* methods in that it allows dynamic length ranges to be specified.

2.3.11. Slice Write Method: `set_slc`

Slices are written with the method:

```
set_slc(int lsb, const ac_int<W,S> &slc)
```

where `lsb` is the index of the LSB of the slice been written and `slc` is `ac_int` slice that is assigned to the `ac_int`:

```
x.set_slc(7, y);
```

Out of Bounds Slice Writes

Attempting to assign to a bit that is outside of the range $[0, W-1]$ of the `ac_int<W,S>` object constitutes an out of bound write. Such a write is regarded as undefined behavior and is the product of an erroneous program. If such a write is encountered during execution (simulation) an assert will be triggered. See also [“User Defined Asserts”](#) in the *Frequently Asked Questions* section.

Differences with Built-in C Integral Types

Accessing a bit or a slice of a C integral type is done by a combination of shift and bit masking. Writing a bit or a slice of a C integral type is done with a combination of shift and bitwise operations.

2.3.12. The `set_val` Method

The `set_val<ac_special_val>()` method sets the `ac_int` or `ac_fixed` to one of a set of “special values” specified by the template parameter as shown in [Table 18](#).

Table 18: Special values

ac_special_val enum	Value for <code>ac_int/ac_fixed</code>
AC_VAL_DC	Used mainly to un-initialized variables that are already initialized (by constructor or by being static). Used for validating that algorithm does not depend on initial value. Synthesis can treat it as a <code>dont_care</code> value.
AC_VAL_0	0
AC_VAL_MIN	Minimum value as specified in Table 1 .
AC_VAL_MAX	Maximum value as specified in Table 1 .
AC_VAL_QUANTUM	Quantum value as specified in Table 1 .

Direct assignment of the enumeration values should not be used since it will assign the integer value of the enumeration to the `ac_int` or `ac_fixed`.

2.3.13. Constructors

Constructors from all C-types are provided. Constructors from *ac_int* are also provided. The default constructor does not initialize *ac_int* or *ac_fixed*. Thus non-static variables of type *ac_int* or *ac_fixed* will not be initialized by default. If the `AC_DEFAULT_IN_RANGE` macro is defined before the first inclusion of the AC datatype header, then the default constructor will adjust the un-initialized value to force it to be in range:

```
#define AC_DEFAULT_IN_RANGE
#include <ac_int.h>
...
ac_int<3,false> n; // n will be in range of 0 to 7
```

Constructors from `char *`, have not been defined/implemented in the current release. See Methods to Fill Bits for alternatives.

2.3.14. Methods to Fill Bits

Utility functions *bit_fill_hex* and *bit_fill* can be used to initialize large bitwidth *ac_int* and *ac_fixed* with raw bits. What is meant by “raw bits” is that its argument is treated as an unsigned bit pattern, without a fixed point and no rounding or overflow handling is performed. The *bit_fill_hex* accepts a hex string. The *bit_fill* accepts and array of integers and it should be the preferred alternative to initialize large *ac_int* and *ac_fixed* with raw bits.

NOTE: The *bit_fill_hex* method should *only* be used to initialize static constants since it is significantly slower than alternative methods.

They are available both as member functions of *ac_int* and *ac_fixed* and as global functions in the `ac` namespace that return the type specified as a template parameter (the type `T` needs to be either an *ac_int* or an *ac_fixed*):

```
void ac_int<W,S>::bit_fill_hex(const char *str);
void ac_fixed<W,I,S,Q,O>::bit_fill_hex(const char *str);
template<typename T> T ac::bit_fill_hex(const char *str);
template<int Na> void ac_int<W,S>::bit_fill(const int (&ivec)[Na], bool
bigendian=true);
template<int Na> void ac_fixed<W,I,S,Q,O>::bit_fill(const int (&ivec)[Na], bool
bigendian=true);
template<typename T, int N> T ac::bit_fill(const int (&ivec)[N], bool bigendian=true);
```

The *bit_fill_hex* function accepts a hex string as an argument which could be shorter or longer than what is required to fill all bits of the *ac_int* or *ac_fixed*. If it is shorter, it is zero padded to fill the remaining most significant bits. If it longer, the extra most significant bits are truncated. The hex string should be a literal constant string and should only contain hex digit characters (0-9, a-f, A-F). Other characters trigger and assert. Because the initialization is done at runtime and this initialization technique is inherently slow, its use to initialize non-static variables is discouraged.

The *bit_fill* function accepts two arguments:

- The first one is an integer array that contains the bit pattern. It could be longer or shorter than what is required to fill all bits. If it is shorter then the remaining most significant bits are zero padded. If it is longer then what would be the extra most significant bits are truncated. The array is not required to be an array of con-

stants.

- The second is a *bool* argument *bigendian* that defaults to *true*.

which means that the bits in the array element with index 0 become the most significant 32 bits of the bit pattern. If the argument is *false*, then the bits in the array element with index 0 become the least significant 32 bits of the bit pattern.

The following example illustrates the use of *bit_fill_hex* and *bit_fill* that do the equivalent functionality:

```
typedef ac_int<80,false> i80_t;
i80_t x;
x.bit_fill_hex("a9876543210fedcba987"); // member function
x = ac::bit_fill_hex<i80_t>("a9876543210fedcba987"); // global function
int vec[] = { 0xa987, 0x6543210f, 0xedcba987 };
x.bit_fill(vec); // member function
x = bit_fill<i80_t>(vec); // global function
// inlining the constant array
x.bit_fill( (int [3]) { 0xa987,0x6543210f,0xedcba987 } ); // member function
x = bit_fill<i80_t>( (int [3]) { 0xa987,0x6543210f,0xedcba987 } ); // global function
```

2.3.15. IO Methods

Methods to performing IO have not been defined/implemented for the current release.

2.3.16. Mixing ac_int and ac_fixed with Other Datatypes

Refer to “[Mixing Datatypes](#)” for information on how to interface *ac_int* and *ac_fixed* to other data types.

2.4. Advanced Utility Functions, Typedefs, etc.

The AC datatypes provide additional utilities such as functions and typedefs. Some of them are available in the ac namespace (ac::), and some of them are available in the scope of the ac datatype itself. The following utility functions/structs/typedefs/static members are described in this section:

- Static members to capture basic parameter information.
- Function for initializing arrays of supported types to a special value.
- Template structs that provide a mechanism to statically compute log2 related functions.
- Typedefs for finding the return type of unary and binary operators.

2.4.1. Accessing Parameter Information

It is often useful to be able to access the value of various template parameters for the datatypes. In some cases it is useful to access the width of type T, where T could be either an *ac_int* or an *ac_fixed*. In this case T::width would provide that information. The various parameters that can be accessed are shown in [Table 19](#).

Table 19: Basic Parameters

Static member	Description for <i>ac_int</i>	Description for <i>ac_fixed</i>
width	Value of W template parameter	Value of W template parameter
i_width	Value of W template parameter	Value of I template parameter
sign	Value of S template parameter	Value of S template parameter
q_mode	AC_TRN	Value of Q template parameter
o_mode	AC_WRAP	Value of O template parameter
e_width	0	0

Note that for generality all the static members are defined for *ac_int* even in the cases where there is no corresponding template parameter involved as they do capture the numerical behavior of *ac_int*.

2.4.2. Using `ac::init_array` for Initializing Arrays

The `ac::init_array` utility function is provided to facilitate the initialization of arrays to zero, or un-initialization (initialization to `dont_care`). The most common usage is to un-initialize an array that is declared static as shown in the following example:

```
void foo( ... ) {
    static int b[200];
    static bool b_dummy = ac::init_array<AC_VAL_DC>(b, 200);
    ...
}
```

The variable `b_dummy` is declared static so that the initialization of array `b` to `dont_care` occurs only once rather than every time the function `foo` is invoked. The return value of `ac::init_array` is always “true”, but in reality only the side effect to array `b` is of interest. A similar example to initialize an array to zero that is not declared static would look like:

```
void foo( ... ) {
    int b[200];
    ac::init_array<AC_VAL_0>(b, 200);
    ...
}
```

The function `ac::init_array` does not check for array bound violations. The template argument to `ac::init_array` is an enumeration that can be any of the following values: `AC_VAL_0`, `AC_VAL_DC`, `AC_VAL_MIN`, `AC_VAL_MAX` or `AC_VAL_QUANTUM` (see [Table 18](#) for details). The function is defined for the integer and fixed point datatypes shown in [Table 20](#):

Table 20: Required Include Files for `ac::init_array` Function

Type	Required include file
C integer types	<i>ac_int.h</i>
<i>ac_int</i> , <i>ac_fixed</i>	No additional include
Supported SystemC types	<i>ac_int.h</i> , <i>ac_sc.h</i>

The first argument is of type pointer to one of the types in [Table 20](#). Arrays of any dimension may be initialized using `ac::init_array` by casting it or taking the address of the first element:

```
static int b[200][200];
static bool b_dummy = ac::init_array<AC_VAL_DC>((int*) b, 200*200);
```

or by taking the address of the first element:

```
static int b[200][200];
static bool b_dummy = ac::init_array<AC_VAL_DC>(&b[0][0], 200*200);
```

The second argument is the number of elements to be initialized. For example:

```
int b[200]; ac::init_array<AC_VAL_0>(b+50, 100);
```

initializes elements `b[50]` to `b[149]` to 0.

Other `ac::init_array` Examples:

```
// Using ac::init_array inside a constructor
class X {
    sc_int<5> a[10][32][5][7];
public:
    X() { ac::init_array<AC_VAL_DC>(&a[0][0][0][0], 10*32*5*7); }
    ...
};

// Will be inlined with initialization loop: b+i, 100+k are not constants
int b[200]; ac::init_array<AC_VAL_0>(b+i, 100+k);

// Will be inlined with initialization loop: mult(n1,n2) not recognized as
// a constant at inlining time
const int n1 = 40;
const int n2 = 5;
int a[n1][n2];
ac::init_array(&a[0][0], mult(n1,n2));

// Uninitialize two ranges of an array
static int b[2][100];
static bool b_dummy = ac::init_array<AC_VAL_DC>(&b[0][0], 50);
static bool b_dummy2 = ac::init_array<AC_VAL_DC>(&b[1][0], 50);

// Alternative to Uninitialize two ranges of an array
static int b[2][100];
static bool b_dummy = ac::init_array<AC_VAL_DC>(&b[0][0], 50) &
    ac::init_array<AC_VAL_DC>(&b[1][0], 50);
```

2.4.3. Static Computation of `log2` Functions

It is statically compute the functions `ceil(log2(x))`, `floor(log2(x))` and `nbits(x)` where `x` is an unsigned integer. The `nbits(x)` function is the minimum number of bits for an unsigned `ac_int` to represent the value `x`.

Static computation of these functions is often useful where `x` is an integer template parameter and the result is meant to be used as a template value (thus it needs to be statically be determined).

For example, lets assume that we have a template class:

```
template<int Size, typename T>
class circular_buffer {
```

```
T _buf[Size];
ac_int< ac::log2_ceil<Size>::val, false> _buf_index;
};
```

for a circular buffer. The minimum bitwidth of the index variable into the buffer is $\text{ceil}(\log_2(\text{Size}))$ where Size is the size of the buffer.

The computation of the \log_2 functions is accomplished using a recursive template class. For the user it suffices to know the syntax on how to retrieve the desired value as shown in [Table 21](#).

Table 21: Syntax for \log_2 functions

Function	Syntax
$\text{ceil}(\log_2(x))$	<code>ac::log2_ceil<x>::val</code>
$\text{floor}(\log_2(x))$	<code>ac::log2_floor<x>::val</code>
$\text{nbits}(x)$	<code>ac::nbits<x>::val</code>

It is important to note that x needs to be a statically determinable constant (constant or template argument).

2.4.4. Return Type for Unary and Binary Operators

It is often useful to find out the return type of an operator. For example, let's assume the following scenario: assume that we have:

```
Ta a = ...;
Tb b = ...;
Tc c = ...;
T res = a * b + c;
```

what should the type T be such that there is no loss of precision?

This section provides the mechanism to find T in terms of Ta , Tb and Tc provided they are AC Datatypes. In addition to return types for binary operations, the return type for unary operators (though the actual operators are not all provided) such as the magnitude (or absolute value), the square, negation is also provided. It is also possible to find out the type required to hold the summation of a set of N values of an algorithmic datatype.

The unary operators are listing in [Table 22](#) (summation is not really an unary operator, but it depends on a single type).

Table 22: Return types for operator on T .

operator on type T	Return type
neg	<code>T::rt_unary::neg</code>
mag	<code>T::rt_unary::mag</code>
mag_sqr	<code>T::rt_unary::mag_sqr</code>
summation of N elements	<code>T::rt_unary::set<N>::sum</code>

The binary operators are shown in [Table 23](#).

Table 23: Return type for (T1(op1) op T2(op2))

operator on types T1, T2	Return Type
op1 * op2	T1::rt_T<T2>::mult
op1 + op2	T1::rt_T<T2>::plus
op1 - op2	T1::rt_T<T2>::minus
op1 / op2	T1::rt_T<T2>::div
op1 (&, , ^) op2	T1::rt_T<T2>::logic
op2 - op1	T1::rt_T<T2>::minus2
op2 / op1	T1::rt_T<T2>::div2

The last two rows in Table 23 are mostly there as helper functions to build up the infrastructure and are not in general needed. For example if both *T1* and *T2* are AC datatypes then instead of using *T1::rt_T<T2>::minus2*, *T2::rt_T<T1>::minus* could be used. These versions are only there for non-commutative operators.

Returning to the `mult_add` example, the type *T* would be expressed as:

```
typedef typename Ta::template rt_T<Tb>::mult a_mult_b;
typename a_mult_b::template rt_T<Tc>::plus res = a * b + c;
```

where the keywords *typename* and *template* are used when *Ta*, *Tb* and *Tc* are template arguments (*dependent-name lookup*). In this case getting to the type was done in two steps by first defining the type `a_mult_b`. In the following version, the it is done in one step so that it can be used directly as the return type of the templated function `mult_add`:

```
template<typename Ta, typename Tb, typename Tc>
typename Ta::template rt_T<Tb>::mult::template rt_T<Tc>::plus mult_add(Ta a, Tb b, Tc
c) {
    typename Ta::template rt_T<Tb>::mult::template rt_T<Tc>::plus res = a * b + c;
    return res;
}
```

Note that additional *template* keywords are used because the lookup of *rt_T* is a dependent-name lookup, that is the parser does not know that *Ta::rt_T* is a templated class until it knows the type *Ta* (this happens only once the function `mult_add` is instantiated).

An example of the use of the type for summation is given below:

```
template <int N, typename T>
typename T::rt_unary::template set<N>::sum accumulate(T arr[N]) {
    typename T::rt_unary::template set<N>::sum acc = 0;
    for(int i=0; i < N; i++)
        acc += arr[i];
    return acc;
}
```

2.5. Methods and Utility Functions for Floating Point

Both *ac_int* and *ac_fixed* provide methods to help implementation of floating point or blocking floating point functionality. The methods implement *leading_sign* and *normalization* and are described in detail in

this section. In addition, utility functions are defined to get the sign, mantissa and exponent of C++ constant *float* and *double* literals. They are named after the commonly used *frexp* function defined in *math.h*.

2.5.1. Leading Sign and Normalization

The *leading_sign* method does not change the *ac_int/ac_fixed* object.

There are two forms:

```
typename rt_unary::leading_sign leading_sign() const;
typename rt_unary::leading_sign leading_sign(bool &all_sign) const;
```

The return type is found so that it has the minimum static bitwidth that can represent the largest value that maybe returned. For unsigned *ac_int/ac_fixed*, *leading_sign* returns the number of leading zeros. For signed *ac_int/ac_fixed*, *leading_sign* returns the number of leading 0/1s minus one. Leading sign thus returns the largest integer such that:

```
y = x.leading_sign() then x == (x << y) * 2-y
```

The second form has the argument *all_sign* that is set to true if all bits are zeros for unsigned or all bits are the same for unsigned *ac_int/ac_fixed* types. The only purpose is to save the extra comparison required to compute that information from the return value. One use for *leading_sign* is for computing the shift left for normalization of a mantissa, though this use is covered by the normalization methods provided for both *ac_int/ac_fixed* described next.

The normalization methods treat the *ac_int/ac_fixed* as a mantissa and has two forms:

```
bool normalize(ac_int<WE,SE> &exp);
bool normalize_RME(ac_int<WE,SE> &exp);
```

The first form shifts the mantissa left as much as possible to normalize it such that:

```
exp = exp_i;
man = man_i;
bool b = man.normalize(exp);
```

then

```
man * 2exp == man_i * 2exp_i
```

which requires that the mantissa does not lose information (see *leading_sign* description) and the exponent does not saturate. Depending on the initial exponent and *man.leading_sign()*, *normalize* may or may not fully normalize the mantissa. The return value is true if the mantissa is normalized after the call.

Note that there is *no implied* leading 1 in the mantissa. Also for sign/magnitude floating representations (such as IEEE floating point representations), an unsigned mantissa should be used. Note that the minimum exponent depends on the value of the template parameters WE and SE.

The second form *normalize_RME* is similar but reserves the minimum exponent. This is useful as the minimum exponent is often used to encode special numbers (NaN etc.) in commonly used floating point representations.

2.5.2. Utility Function to Extract Exponent/Sign/Mantissa from Literal Constants

The following functions in the "ac" namespace provide ways to extract the mantissa, exponent and optionally sign from C *float* and *double* literal constant types:

```
ac_fixed<54,2,true> frexp_d(double d, ac_int<11,true> &exp);  
ac_fixed<25,2,true> frexp_f(float f, ac_int<8,true> &exp);  
ac_fixed<53,1,false> frexp_sm_d(double d, ac_int<11,true> &exp, bool &sign);  
ac_fixed<24,1,false> frexp_sm_f(float f, ac_int<8,true> &exp, bool &sign);
```

The first two return the mantissa into a two's complement signed *ac_fixed*. The mantissa is normalized (`mant[msb] != mant[msb-1]`) unless the exponent is saturated to the IEEE standard exponent minimum of -126 for *float* or -1022 for *double*. The mantissa is provided in full (no implied bit).

The last two ("*_sm_*" in their name) return the result in a sign-magnitude representation: the return value is an unsigned *ac_fixed* and the sign is returned as an argument passed by reference.

Example:

```
ac_int<8,true> exp;  
ac_fixed<54,2,true> mant = ac::frexp_d(0.1, exp);
```

Chapter 3: Arbitrary-Length Bit-Accurate Floating-Point Datatypes

Floating-point types are used to represent a higher range of values than a fixed-point representation having the same number of bits. Floating-point types have both a mantissa and an exponent. The *mantissa* is a value in a fixed-point representation and the *exponent* that shifts the point making it a “floating” point. The type provides a templated class where values are encoded as a signed mantissa m and an exponent e such that:

$$\text{value} = m * 2^e$$

where m is a signed fixed-point quantity and e is a signed integer. The class is templated as:

```
ac_float<W,I,E,Q>
```

where the first two parameters W and I define the mantissa as an *ac_fixed*< $W,I,true$ >, the E defines the exponent as an *ac_int*< $E,true$ > and Q defines the rounding mode. For example:

```
#include <ac_float.h>
...
ac_float<10,1,4,AC_RND> x = ...; // W=10, I=1, E=4, Q=AC_RND
ac_float<10,1,4> y = ...; // Q = AC_TRN (default setting)
...
```

The definition of *ac_float* is targeted for algorithms that are intended to be mapped to hardware where the flexibility of tailoring widths of mantissa and exponents is critical to reduce hardware resources. The I parameter provides a static exponent offset so that the exponent bit width can be minimized. The use of the I parameter and the encoding of the mantissa as an *ac_fixed* provides a more natural transition to and from *ac_fixed*.

Currently the typical optimization of the MSB of the mantissa as an implied ‘1’ bit is not implemented in *ac_float*. This was done to keep the implementation as simple as possible. Supporting the implied ‘1’ bit requires using up one exponent value to encode special conditions such as the fact that the number is denormalized and thus the MSB is not ‘1’. Currently *ac_float* does not reserve any exponent to encode special conditions and therefore does not support *NaN*, *+inf*, *-inf*.

The encoding of the mantissa is in two’s complement. This makes it consistent with *ac_fixed*. Typical implementations for floating-point use a sign-magnitude representation. Such a representation introduces some complexities, but may reduce toggling of bits and be beneficial for lowering the power used in a hardware implementation of the algorithm. Supporting such a sign-magnitude option will be considered as a future enhancement.

As compared to *ac_fixed*, saturation (equivalent to *AC_SAT*) is performed by default. This was done for the following reasons:

- The hardware overhead for saturation logic is relatively small compared to the inherent complexity of floating-point arithmetic
- Finite precision floating-point addition is not associative. This means that optimizations such as tree balancing can not be performed without affecting bit-accuracy. Using saturation rather than wrap around does not worsen this situation.

- The number of template parameters is kept to a minimum on purpose. Additional parameters may be required for future enhancements.

The *ac_float* implementation is built in top of the *ac_fixed* and *ac_int* types. In principle all parameters could be set to arbitrary valid values. However the current implementation requires that the sum of exponent values and offsets (*I* parameter) fit in a 32-bit integer. Note that a C floating *double* has 11 bits of exponents, so it is unlikely that the current limitation is restrictive in most cases. An exponent width of 0 is also not currently supported.

The numeric range of an *ac_float*<*W*,*I*,*E*> (the *Q* parameter has a default of *AC_TRN* and has no impact on the range) has a numeric range of:

$$(-0.5) 2^{I + \max_exp} \text{ to } (0.5 - 2^{-W}) 2^{I + \max_exp} \quad \text{where } \max_exp = 2^{E-1} - 1$$

The smallest increment (quantum) is:

$$2^{I-W+\min_exp} \quad \text{where } \min_exp = -2^{E-1}$$

Based on the ranges a standard *float* is covered by *ac_float*<25,1,8> and *double* is covered by *ac_float*<54,1,11>. Note that the overall number of bits (*W*+*E*) is 33 and 65 as there is no implied ‘1’ bit in *ac_float*. There are differences of rounding between the native C *float* types (*float* and *double*) and the smallest *ac_float* that covers them in range. For example, the division operator for *ac_float* truncates towards zero (the rounding behavior might be enhanced in the future when both addition and subtraction operators will be implemented).

Table 24 summarizes the characteristic of *ac_float* and contrasts it to typical floating-point.

Table 24: Operators and methods defined for *ac_float*

Characteristic	Description and comparison to typical floating-point
Zero	Represented as all zeros for both mantissa and exponent.
Normalized Numbers	No implied ‘1’ bit.
Denormalized Numbers	No exception required as there is no implied ‘1’ for normalized numbers.
Infinity (+inf, -inf)	Not represented. Used by typical floating point types to represent result of <i>a</i> / <i>a</i> !=0.
NaN (not-a-number)	Not represented. Used by typical floating point types to represent inf/inf, inf-inf etc.
Mantissa Representation	Two’s complement (signed). Stored as an <i>ac_fixed</i> < <i>W</i> , <i>I</i> , <i>true</i> >. Parameter <i>I</i> can be regarded as an exponent bias.
Sign bit	Not used as mantissa is signed (two’s complement) already.
Exponent	Stored as two’s complement as an <i>ac_int</i> < <i>E</i> , <i>true</i> >. Different from typical floating point type stored as an unsigned and requiring an offset to get actual value.
Base	Base is 2. Typical floating-point types also have base 2. IEEE standards do allow for other bases (base 10).

Target Use	Algorithmic descriptions that need to be optimized for minimal hardware. Parameterization of precision and range and mixing of different types is essential. This is different than typical floating-point types that are “closed” (result is same type as operands).
Rounding	Done on assignment according to <i>Q</i> parameter.
Overflow	Saturation overflow always performed unless target range covers source range (based on templates).
Operators with non-floats	Operators with <i>ac_fixed</i> , <i>ac_int</i> and C integers, first represent the operand as an <i>ac_float</i> with E=1 (E=0, currently not supported).

Table 25 shows the operators and methods defined for *ac_float*

Table 25: Operators and methods defined for *ac_float*

Operators	<i>ac_float</i>
*	Arithmetic result (no loss of precision). Mantissa type is determined by <i>ac_fixed</i> rules for multiplication. Exponent type is determined by <i>ac_int</i> rules for addition.
/	Mantissa type is determined by <i>ac_fixed</i> rules for division (truncation towards zero). Exponent type is determined by <i>ac_int</i> rules for subtraction.
add(op1, op2) sub(op1, op2)	Methods to add and subtract op2 to/from op1. The result is stored in the object (“this”). This allows to consider the type of the target (including the <i>Q</i> parameter of the object) to perform rounding.
+, -	To be defined. Current definition is placeholder and can be enabled by defining <code>__AC_FLOAT_ENABLE_ALPHA__</code> . It is only used to test compatibility with <i>ac_complex</i> . Challenge is that these operate loose precision and thus need to truncate/round. But it would ideal to avoid two steps of rounding in an assignment “a = b + c”.
>>, <<	Bidirectional. Mantissa unchanged. Shift value added to (<<) or subtracted from (>>) exponent. No loss of precision as exponent grows to accommodate resulting exponent. Second arg is <i>ac_int</i> or C <i>INT</i>
=	quantization (specified by target), then saturation if range of target does not cover range of source
+=, -=, *=, /=, >>=, <<=	Equiv to op then assign where defined. First arg is <i>ac_float</i>
==, !=, >, <, >=, <=	First or second arg may be <i>ac_fixed</i> , <i>ac_int</i> or C <i>INT</i> or C <i>float</i> or <i>double</i>

Unary +, -	Arithmetic
! x	Equiv to x == 0
mantissa(), exp()	Access methods to mantissa and exponent values
Explicit conversion to other types	<i>to_double()</i> , <i>to_float()</i> , <i>to_ac_fixed()</i> , <i>to_ac_int()</i> , <i>to_int()</i> , <i>to_uint()</i> , <i>to_long()</i> , <i>to_ulong()</i> , <i>to_int64()</i> and <i>to_uint64()</i>
Constructors from other types	from <i>ac_fixed</i> , <i>ac_int</i> , <i>float</i> , <i>double</i> , C integers. Perform normalization by default (argument to constructor).
%, %=, &, , ^, &=, =, ^=, ~, ++, --, []	Not defined.

3.1.1. Mixed ac_float and other types

Mixed operators are defined for *ac_int*, *ac_fixed*, C native types to avoid ambiguity in the semantics or compilation problems due to multiple operators matching an operation. The other type is first represented as an *ac_float* of mantissa that can represent the number assuming a zero exponent. The exponent width would in theory be 0, but currently an exponent width of 1 is minimum to compile, so that is used instead.

3.1.2. Shift Operators

Shift operators have a different behavior than for *ac_int* and *ac_fixed*. For *ac_float*, the shift operators are arithmetic (no loss of precision).

For minimal hardware, it is best to constrain the type of the second operand to reduce its range.

3.1.3. The set_val Method

The *set_val<ac_special_val>()* method sets the *ac_int* or *ac_fixed* to one of a set of “special values” specified by the template parameter as shown in Table 26.

Table 26: Special values

ac_special_val enum	Value for ac_float
AC_VAL_DC	Used mainly to un-initialized variables that are already initialized (by constructor or by being static). Used for validating that algorithm does not depend on initial value. Synthesis can treat it as a dont_care value.
AC_VAL_0	0
AC_VAL_MIN	Minimum value
AC_VAL_MAX	Maximum value
AC_VAL_QUANTUM	Quantum value

Direct assignment of the enumeration values should not be used since it will assign the integer value of the

enumeration to the *ac_float*.

3.1.4. Constructors

Constructors from all C-types are provided. Constructors from *ac_int* and *ac_fixed* are also provided. The default constructor does not initialize the *ac_float* variable. Thus non-static variables of type *ac_fixed* will not be initialized by default. If the *AC_DEFAULT_IN_RANGE* macro is defined before the first inclusion of the AC datatype header, then the default constructor will adjust the un-initialized value to force it to be in range:

```
#define AC_DEFAULT_IN_RANGE
#include <ac_float.h>
...
ac_float<4,4,3,false> n; // n will be in range
```

Constructors from *char* *, have not been defined/implemented in the current release.

3.1.5. Accessing Parameter Information

It is often useful to be able to access the value of various template parameters for the datatypes. In some cases it is useful to access the width of type *T*, where *T* is an *ac_float*. In this case *T::width* would provide that information. The various parameters that can be accessed are shown in Table 27.

Table 27: Basic Parameters

Static member	Description for <i>ac_float</i>
<i>width</i>	Value of <i>W</i> template parameter
<i>i_width</i>	Value of <i>I</i> template parameter
<i>sign</i>	true
<i>q_mode</i>	Value of <i>Q</i> template parameter
<i>o_mode</i>	AC_SAT
<i>e_width</i>	Value of <i>E</i> template parameter

3.1.6. Using *ac::init_array* for Initializing Arrays

The utility function *ac::init_array* is provided to facilitate the initialization of arrays of *ac_float* to zero, or un-initialization (initialization to don't_care). The most common usage is to un-initialize an array that is declared static as shown in the following example:

```
void foo( ... ) {
    static ac_float<10,1,3> b[200];
    static bool b_dummy = ac::init_array<AC_VAL_DC>(b,200);
```

```
...  
}
```

Chapter 4: Complex Datatype

The algorithmic datatype *ac_complex* is a templated class for representing complex numbers. The template argument defines the type of the real and imaginary numbers and can be any of the following:

- Algorithmic C integer type: *ac_int*<W,S>
- Algorithmic C fixed-point type: *ac_fixed*<W,I,S,Q,O>
- Native C integer types: *bool*, (un)signed *char*, *short*, *int*, *long* and *long long*
- Native C floating-point types: *float* and *double*

For example, the code:

```
ac_complex<ac_fixed<16,8,true> > x (2.0, -3.0);
```

declares the variable *x* of type *ac_complex* based on *ac_fixed*<16,8,true> and initializes it to have a real part of 2.0 and imaginary part of -3.0 (note: the space between the two ‘>’ is required by C++).

An important feature of the *ac_complex* type is that operators return the types according to the rules of the underlying type. For example, operators on *ac_complex* types based on *ac_int* and *ac_fixed* will return results for the operators ‘+’, ‘-’ and ‘*’ with no loss of precision (‘/’ will follow the rules for *ac_int* and *ac_fixed*). Likewise, operators on *ac_complex* types based on native C integer and floating-point types will return results according to the C rules for arithmetic promotion and conversion.

A second important feature of the *ac_complex* type is that binary operators are defined for *ac_complex* types that are based on different types, provided the underlying types have the necessary operators defined. For instance to implement complex multiplication, it is necessary to have addition and multiplication defined for the underlying types. As the examples below illustrate, the only issue is combining native floating-point types (*float* and *double*) with algorithmic types:

```
ac_complex<ac_int<5,true> > i(2, 1);
ac_complex<ac_fixed<8,3,false> f(1, 5);
ac_complex<unsigned short> s(1, 0);
ac_complex<double> d(3.5, 3.14);

i * f; // OK: ac_int and ac_fixed can be mixed
s * f; // OK: native int type can be mixed with ac_fixed
i * s; // OK: ac_int and native int type can be mixed
s * d; // OK: native int type can be mixed w/native floating-point type
i * d; // ERROR: ac_int and native floating-point types don't mix
i == d; // ERROR: ac_int/double comparison operators is not defined
f * d; // ERROR: ac_fixed/double + and * operators are not defined
f == d; // OK: ac_fixed/double comparison operators is defined
```

Operators for multiplying a variable of type *ac_complex* by a real number also are defined with the same restrictions as outlined above. For example:

```
ac_complex<ac_int<5,true> > i(2, 1);
ac_complex<ac_fixed<8,3,false> f(1, 5);
ac_fixed<8,3,false> f_r = 3;
unsigned short s_r = 5;
```

```
double d_r = 3.5

i * f_r; // OK: ac_int and ac_fixed can be mixed
s_r * i; // OK: native int type can be mixed with ac_int
i * d_r; // ERROR: ac_int/double + and * operators are not defined
i * 0.1; // ERROR: ac_int/float + and * operators are not defined
i == d_r; // ERROR: ac_int/double comparison operator is not defined
f == d_r; // OK: ac_fixed/double comparison operator is defined
i == 0.1; // ERROR: ac_int/float comparison operators is not defined
f == 0.1; // OK: ac_fixed/double comparison operator is defined
```

Table 28 shows the operators defined for both *ac_complex*.

Table 28: Operators defined for *ac_complex*

Operators	<i>ac_complex</i>
Two operand +, -, *, /,	Arithmetic result. First or second arg may be C <i>INT</i> or <i>ac_fixed</i>
=	assignment
+=, -=, *=, /=	Equiv to op then assign. First arg is <i>ac_complex</i>
==, !=	First or second arg may be C <i>INT</i> , <i>ac_int</i> , <i>ac_fixed</i> or C <i>double</i> (comparison of <i>ac_int</i> with <i>double/float</i> not defined)
Unary +, -	Arithmetic
! x	Equiv to x == 0

Table 29 shows the methods defined for the *ac_complex* type.

Table 29: Methods defined for *ac_complex*<T>

Methods	ac_complex
r(), real()	return real part (const T&, or T&)
i(), imag()	return imaginary part (const T& or T&)
set_r(const T2 &r)	assign <i>r</i> to real part
set_i(const T2 &i)	assign <i>i</i> to imaginary part
conj()	complex conjugate
sign_conj()	returns (sign(real), sign(imag))) as an <i>ac_complex</i> < <i>ac_int</i> <2,true> >
mag_sqr()	returns <i>sqr(real) + sqr(imag)</i>
to_string	convert to <i>std::string</i> depending on parameter AC_HEX, AC_DEC, AC_OCT, AC_BIN
type_name()	returns “name” of the type as a <i>std::string</i>

4.1. Usage of ac_complex

In order to use the *ac_complex* datatype the following header file must be included in the C++ source:

```
#include <ac_complex.h>
```

The following list includes recommendations for using the *ac_complex* datatype:

- Do not use native C type *unsigned* (*unsigned int*) as the return type (and the arithmetic) is defined according to the promotion/arithmetic rules of the C language. That is the resulting complex type will be based on the type *unsigned*. For example:

```
ac_complex<unsigned> x(0,1);
cout << x*x; // result is (2^32 - 1, 0)
```

- Pay special attention on the return type when performing division. For example, if two *ac_complex* based on native C type *int* are divided, the result will be an *ac_complex* based on *int* and truncation will take place.

4.2. Advanced utility functions, typedefs, etc for ac_complex

The AC datatypes provide additional utilities such as functions and typedefs. Some of them are available in the *ac* namespace (*ac::*), and some of them are available in the scope of the *ac* datatype itself. The following utility functions/structs/typedefs are described in this section:

- Typedef to capture the underlying type.
- Function for initializing arrays of *ac_complex* to a special value.
- Typedefs for finding the return type of unary and binary operators.

4.2.1. Accessing the Underlying (Element) Type

The type of the real and imaginary elements can be accessed as

```
T::element_type
```

where T is the `ac_complex` type.

4.2.2. Using `ac::init_array` for Initializing Arrays

The utility function “`ac::init_array`” is provided to facilitate the initialization of arrays to zero, or un-initialization (initialization to `dont_care`). For more details about the basic AC Datatypes, refer to the examples in “[Arbitrary-Length, Bit-Accurate Integer and Fixed-Point Datatypes](#)” section. The initialization value is applied to both the real and imaginary components.

4.2.3. Return Type for Unary and Binary Operators

Refer to corresponding sections in “[Arbitrary-Length, Bit-Accurate Integer and Fixed-Point Datatypes](#)” for the basic AC Datatypes.

Chapter 5: Reference Guide for Numerical Algorithmic C Datatypes

This chapter provides a high-level view of the numerical AC Datatype package. It provides a high-level reference for the following items that are exposed to the user:

- [All public functions, member functions](#)
- [Enumerations, Static Constants and Type Definitions](#)
- [Macros](#)

Anything that is not covered in this chapter is implementation specific and is subject to change.

5.1. Functions and Operators

The tables in this section capture all the public functions, member functions and operators that are available for the different types.

5.1.1. Constructors

The following table lists all the constructors available for the AC numerical types.

Table 30: Constructors Available

Constructor	ac_int	ac_fixed	ac_float	ac_complex
Default ¹	✓	✓	✓	✓
Copy	NA	✓	✓	NA
Constructor from same type name but any parameter	✓	✓	✓ ²	✓
Constructor from each bool and integer type	✓	✓	✓	NA ⁴
Constructor from float	NA	NA	✓ ²	NA ⁴
Constructor from double	✓	✓	✓ ²	NA ⁴
Constructor from template T type	NA	NA	NA	✓
Constructor from template T1, T2 type pair	NA	NA	NA	✓
Constructor from other	NA	✓	✓	NA ⁴

AC Numerical Type		ac_int	ac_int, ac_fixed ³ 5.1.2(ac_fixed, ac_int) ³	
-------------------	--	--------	---	--

Notes:

1. The default constructors leave the datatypes un-initialized. If the *AC_DEFAULT_IN_RANGE* macro is defined, it will be un-initialized, but guaranteed to be within range.
2. The *ac_float* constructors have additional arguments to assert on rounding or overflow.
3. The *ac_float* constructors have an additional argument to normalize.
4. The general template constructor from template *T* type covers these cases.

5.1.2. Conversions

Conversions can be classified in conversion operators and explicit conversion methods. The only operators that is available is for *ac_int* with parameter restrictions that guarantee that there is no loss of precision. The explicit conversion methods *ac_float::to_ac_fixed* does not lose precision. Otherwise all other explicit conversions may lose precision depending on the parameter values of the types.

Table 31: Conversion Operators and Methods

Operators and Methods	ac_int	ac_fixed	ac_float	ac_complex
operator long long	✓ W ≤ 64 S = true	NA	NA	NA
operator unsigned long long	✓ W ≤ 64 S = false	NA	NA	NA
to_int	✓	✓	✓	NA
to_uint	✓	✓	✓	NA
to_long	✓	✓	✓	NA
to_ulong	✓	✓	✓	NA
to_int64	✓	✓	✓	NA
to_uint64	✓	✓	✓	NA
to_float	NA	NA	✓	NA
to_double	✓	✓	✓	NA
to_ac_int	NA	✓	✓	NA
to_ac_fixed	NA	NA	✓	NA

to_string	✓	✓	✓	NA
-----------	---	---	---	----

5.1.3. Arithmetic, Relational and Shift Operators and Methods

The operators and methods covers both binary (two operand) operations and unary (single operand) operations. The binary operations have also mixed AC type and C++ integer types global operators defined for them.

Table 32: Arithmetic, Relational and Shift Operators and Methods

Operators and Methods	ac_int	ac_fixed	ac_float	ac_complex
+, -, /, +=, -=, /=	✓ ¹	✓	✓ ^{2,4}	✓ ⁵
add, sub ²	NA	NA	✓	NA
%, %=	✓	NA	NA	NA
==, !=	✓	✓ ³	✓ ⁴	✓ ⁵
>, <, >=, <=	✓	✓ ³	✓ ⁴	NA
>>, <<, >>=, <<=	✓ ⁶	✓ ⁶	✓	NA
Unary +, -, ! ⁷	✓	✓	✓	✓
++, --	✓	✓	NA	NA
abs	NA	NA	✓	NA
conj	NA	NA	NA	✓
mag_sqr	NA	NA	NA	✓
sign_conj	NA	NA	NA	✓

Notes:

1. The + operator is also defined for *ac_int* and C pointer (and vice versa) so that an *ac_int* can be added to a C pointer. The operator - is defined so that an *ac_int* can be subtracted from a C pointer. In all cases, the result is of the same type as the C pointer.
2. The *add* and *sub* methods for *ac_float* provide control over the desired return type and are alternatives for the + and – operators that are still subject to change.
3. The relational operators are defined for mixed *ac_fixed* with C++ double.
4. Operators are defined for mixed *ac_float* with with C++ *float* and *double*.
5. Operators are defined for mixed *ac_complex* with another type that is a template parameter.
6. Shift operators >> and << for *ac_int* and *ac_fixed* return the type of the first argument.

7. The unary \sim operator is covered in Logical Operators and Methods.

5.1.4. Bit and Slice Operators and Methods

The bit and slice operators and methods are defined for `ac_int` and `ac_fixed`. The `slc` method can read a slice out of range.

Table 33: Bit and Slice Operators and Methods

Operators and Methods	<code>ac_int</code>	<code>ac_fixed</code>	<code>ac_float</code>	<code>ac_complex</code>
<code>slc</code>	✓	✓	NA	NA
<code>set_slc</code>	✓	✓	NA	NA
operator <code>[]</code>	✓	✓	NA	NA

There are const and non const versions of the `[]` operator and the index argument can be *int*, *unsigned*, and *ac_int*. There are asserts to check that index is in range.

The `ac_int::operator []` returns an `ac_int::ac_bitref` and the `ac_fixed::operator []` returns an `ac_fixed::ac_bitref`. Both `ac_int::ac_bitref` and `ac_fixed::ac_bitref` have identical definitions as shown in Table 34. The `ac_bitref` classes are not meant to be used explicitly by the user.

Table 34: Operators defined for `ac_int::ac_bitref` and `ac_fixed::ac_bitref`

Operators	Description
operator bool	bit returned as bool
operator <code>ac_int<W2, S2></code>	First call operator bool and then constructs <code>ac_int</code>
<code>ac_bitref operator = (int val)</code>	LSB of int is written to bit
<code>ac_bitref operator = (const ac_bitref &val)</code>	Assignment for <code>ac_bitref</code> for identical type

5.1.5. Logical Operators and Methods

The \sim operator has an arithmetic definition as compared to the `bit_complement` method listed in Table 35 which returns the unsigned bitwise complement of the *ac_int* or *ac_fixed*.

Table 35: Logical Operators and Methods

Operators and Methods	<code>ac_int</code>	<code>ac_fixed</code>	<code>ac_float</code>	<code>ac_complex</code>
<code>&</code> , <code> </code> , <code>^</code> , <code>&=</code> , <code> =</code> , <code>^=</code>	✓	✓	NA	NA

Unary operator ~	✓	✓	NA	NA
bit_complement	✓	✓	NA	NA
and_reduce, or_reduce, xor_reduce	✓	NA	NA	NA

5.1.6. Other Functions and Methods

The table below lists functions string related functions, filling them with bit patterns, setting the types to special values, provide functionality useful for floating-point and for accessing data members for *ac_float* and *ac_complex*. The length function returns the *W* parameter value for *ac_int* and *ac_fixed*.

Table 36: Other functions/methods

Functions and Methods	ac_int	ac_fixed	ac_float	ac_complex
std::ostream & operator <<	✓	✓	✓	✓
type_name	✓	✓	✓	✓
bit_fill	✓	✓	NA	NA
ac::bit_fill	✓	✓	NA	NA
bit_fill_hex	✓	✓	NA	NA
ac::bit_fill_hex	✓	✓	NA	NA
set_val	✓	✓	✓	NA
ac::value	✓	✓	✓	✓
ac::init_array	✓	✓	✓	✓
leading_sign	✓	✓	NA	NA
normalize	✓	✓	✓	NA
normalize_RME	✓	✓	NA	NA
mantissa, exp	NA	NA	✓	NA
set_mantissa, set_exp	NA	NA	✓	NA
i, r, real, imag	NA	NA	NA	✓
set_i, set_r	NA	NA	NA	✓
length	✓	✓	NA	NA

5.1.7. Mantissa/Exponent Extraction of float/double

The following table lists functions for extracting mantissa, exponent and sign information out of C++ *float* and *double* floating-point types.

Table 37: Functions to extract mantissa/exponent from float/double

Function	Header	Description
ac::frexp_f	ac_fixed.h	Exponent and 2's complement mantissa for float
ac::frexp_d	ac_fixed.h	Exponent and 2's complement mantissa for double
ac::frexp_sm_f	ac_fixed.h	Exponent, unsigned mantissa and sign for float
ac::frexp_sm_d	ac_fixed.h	Exponent, unsigned mantissa and sign for double

5.1.8. SystemC Tracing Functions

All numerical types have the *sc_trace* function defined. They are defined in the *ac_sc.h* header file.

5.1.9. Explicit conversions to/from SystemC Types

The explicit conversions functions to and from the SystemC datatypes *sc_bigint*, *sc_biguint*, *sc_fixed* and *sc_ufixed* are shown in Table 38. They are defined in the *ac_sc.h* header file.

Table 38: Explicit conversions to/from AC Datatype from/to SystemC Datatype

From	To	Function
sc_bigint	ac_int, S=true	ac_int<W, true> to_ac(const sc_bigint<W> &)
sc_biguint	ac_int, S=false	ac_int<W, false> to_ac(const sc_biguint<W> &)
sc_fixed	ac_fixed, S=true	ac_fixed<W,I, true> to_ac(const sc_fixed<W,I,Q,O,nbits> &)
sc_ufixed	ac_fixed, S=false	ac_fixed<W,I, false> to_ac(const sc_ufixed<W,I,Q,O,nbits> &)
ac_int, S=true	sc_bigint	sc_bigint<W> to_sc(const ac_int<W,true> &)
ac_int, S=false	sc_biguint	sc_biguint<W> to_sc(const ac_int<W,false> &)
ac_fixed, S=true	sc_fixed	sc_fixed<W,I> to_sc(const ac_fixed<W,I,true,Q,O> &)
ac_fixed, S=false	sc_ufixed	sc_ufixed<W,I> to_sc(const ac_fixed<W,I,false,Q,O> &)

5.2. Enumerations, Static Constants and Type Definitions

This section covers the enumerations, static constants and type definitions (typedefs) that are defined in the numerical AC header types.

5.2.1. General Enumerations

The following table defines general enumeration values and which functions uses them.

Table 39: General Enumerations

Enumeration Value	Enumeration Type	Description	Used by
AC_BIN	ac_base_mode	Binary format	to_string
AC_OCT		Octal format	
AC_DEC		Decimal format	
AC_HEX		Hexadecimal format	
AC_VAL_DC	ac_special_val	Don't Care (un-initialized)	set_val ac::value ac::init_array
AC_VAL_0		Zero	
AC_VAL_MIN		Minimum val for type	
AC_VAL_MAX		Maximum val for type	
AC_VAL_QUANTUM		Smallest increment for type	

5.2.2. Enumerations for Fixed-point Quantization and Overflow Modes

The following enumerations are used to define the quantization and overflow modes for *ac_fixed*, but are forward declared in *ac_int.h*. They are template parameters for *ac_fixed* with default settings as indicated.

Table 40: Enumerations for fixed-point quantization and overflow modes.

Enumeration Value	Enumeration Type	Description
AC_TRN (default)	ac_q_mode (Quantization)	truncate toward minus infinity
AC_RND		round towards plus infinity
AC_TRN_ZERO		truncate towards zero
AC_RND_ZERO		round towards zero
AC_RND_INF		round towards infinity
AC_RND_MIN_INF		round towards minus infinity
AC_RND_CONV		round towards even
AC_RND_CONV_ODD		round towards odd
AC_WRAP (default)	ac_o_mode (Overflow)	wrap
AC_SAT		saturate to max and min

AC_SAT_ZERO		saturate to zero
AC_SAT_SYM		saturate to +/- max

5.2.3. Static Constant Members and Type Definitions to Capture Properties of Types

The following table lists the static data members to query the type for properties that are dependent on template parameters. If a type lacks that parameter it may either depend on another parameter or be implicitly behave as having a constant value for that parameter. For example an *ac_int* behaves as an *ac_fixed* that has identical *width* and *i_width* and has *AC_TRN* and *AC_WRAP*.

Table 41: Static constant members to capture properties of types.

Static Data Member	ac_int	ac_fixed	ac_float	ac_complex
width	✓	✓	✓	NA
i_width	✓ width	✓	✓	NA
e_width	✓ 0	✓ 0	✓	NA
sign	✓	✓	✓ true	NA
q_mode	✓ AC_TRN	✓	✓	NA
o_mode	✓ AC_WRAP	✓	✓ AC_SAT	NA

The following table lists the type definitions that are dependent on the template parameters.

Table 42: Type definitions to capture properties of types.

Static Data Member	ac_int	ac_fixed	ac_float	ac_complex
mant_t	NA	NA	✓	NA
exp_t	NA	NA	✓	NA
element_type	NA	NA	NA	✓

5.2.4. Type Definitions for Signed and Unsigned ac_ints

Under the *ac_intN* namespace, there are typedefs for *ac_int<W,S>* for $1 \leq W \leq 63$ for $S=\{\text{false}, \text{true}\}$. The type names are *int* or *uint* followed by the width. For example:

```
uint6 a = 7;
ac_intN::int17 b = 45676;
```

The *ac_intN* namespace is made available in the global namespace unless the macro *AC_NOT_USING_INTN* is defined.

5.2.5. Utility Enumerations and Type Definitions Based on Template Arguments

Static Computation of Bit-width and Type Based on Range

The following table lists utility enum for the compile-time computation of log2-like functions. They are useful to determine the minimum bit-widths of types that can represent a number or a range of numbers.

Table 43: Utility enums and typedefs based on template parameters

struct enum/typedef	Header	Type or enum	Description
<code>ac::nbits<unsigned>::val</code>	<code>ac_int.h</code>	enum	Number of bits required for unsigned <code>ac_int</code> to store value of template parameter
<code>ac::log2_floor<int>::val</code>	<code>ac_int.h</code>	enum	Floor of log2 of template parameter
<code>ac::log2_ceil<int>::val</code>	<code>ac_int.h</code>	enum	Ceiling of log2 of template parameter
<code>ac::int_range<int, int>::type</code>	<code>ac_int.h</code>	<code>ac_int</code>	<code>ac_int</code> type that can represent close integer range given by template parameters

Minimal Size Destination Type to Represent a Source Type

The type defines `ac_int_represent<T>::type`, `ac_fixed_represent<T>::type` and `ac_float_represent<T>::type` provide the minimal destination `ac_int`, `ac_fixed` and `ac_float` type respectively that is required to store the source type `T` as shown in Table 44. Note that while it would be possible to provide an `ac_fixed` to represent a `float` or `double`, its size would be impractically large so it is not provided.

Table 44: Type definitions for Minimal Size Destination Types.

Source Type	<code>ac_int_represent<T>::type</code>	<code>ac_fixed_represent<T>::type</code>	<code>ac_float_represent<T>::type</code>
C++ integer types	✓	✓	✓
C++ float/double	NA	NA	✓
<code>ac_int</code>	✓	✓	✓
<code>ac_fixed</code>	NA	✓	✓
<code>ac_float</code>	NA	NA	✓

Return Type Infrastructure for Unary Operators and Methods

The *rt_unary_struct* defined in each of the types provides parameter and type information to obtain the return type when it is solely dependent on the template parameters of the type. The enumerations only make sense for scalar values (not for *ac_complex*), but the typedefs are available for all types for which the function that requires that return type is defined.

Table 45: Enumerations defined in `type::rt_unary`.

Enumeration	ac_int	ac_fixed	ac_float	ac_complex
neg_w, neg_s	✓	✓	✓	NA
neg_i	NA	✓	✓	NA
neg_e	NA	NA	✓	NA
mag_sqr_w, mag_sqr_s	✓	✓	✓	NA
mag_sqr_i	NA	✓	✓	NA
mag_sqr_e	NA	NA	✓	NA
mag_w, mag_s	✓	✓	✓	NA
mag_i	NA	✓	✓	NA
mag_e	NA	NA	✓	NA
set<N>::sum_w, set<N>::sum_s	✓	✓	✓	NA
set<N>::sum_i	NA	✓	✓	NA
set<N>::sum_e	NA	NA	✓	NA
leading_sign_w, leading_sign_s	✓	✓	NA	NA
to_fx_w, to_fx_i, to_fx_s	NA	NA	✓	NA
to_i_w, to_i_s	NA	NA	✓	NA

Table 46: Type definitions in `type::rt_unary`.

typedef	ac_int	ac_fixed	ac_float	ac_complex
neg	✓	✓	✓	✓
mag_sqr	✓	✓	✓	✓
mag	✓	✓	✓	✓
set<N>::sum	✓	✓	✓	✓
leading_sign	✓	✓	NA	NA
to_ac_fixed_t	NA	NA	✓	NA
to_ac_int_t	NA	NA	✓	NA

Return Type Infrastructure for Binary Operators and Methods

The *rt* and *rt_T* structs defined in each of the types provides parameter and type information to obtain the return type when it is dependent on the template parameters of the type and of a second operand. The *rt* structure is used for *ac_int*, *ac_fixed* and *ac_float* to determine the return type when the second operand is of the same class, but potentially different parameter values. The *rt_T struct* is used to specify the type directly as a template parameter (T stands for template type). The *ac_complex* type uses the cross type infrastructure as described to figure out the return types where the operands are *ac_complex<T1>* and *ac_complex<T2>*.

The *rt struct* has template parameters as follows

- *ac_int<W,S>::rt<int W2, bool S2>*
- *ac_fixed<W,I,S,Q,O>::rt<int W2, int I2, bool S2>*
- *ac_float<W,I,E,Q>::rt<int W2, int I2, int E2>*

Table 47: Enumerations defined in type::rt.

Enumeration	ac_int	ac_fixed	ac_float
mult_w, mult_s	✓	✓	✓
mult_i	NA	✓	✓
mult_e	NA	NA	✓
plus_w, plus_s	✓	✓	✓
plus_i	NA	✓	✓
plus_e	NA	NA	✓
minus_w, minus_s	✓	✓	✓
minus_i	NA	✓	✓
minus_e	NA	NA	✓
div_w, div_s	✓	✓	✓
div_i	NA	✓	✓
div_e	NA	NA	✓
mod_w, mod_s	✓	NA	NA
logic_w, logic_s	✓	✓	✓
logic_i	NA	✓	✓
logic_e	NA	NA	✓

Table 48: Type definitions in type::rt.

typedefs	ac_int	ac_fixed	ac_float
----------	--------	----------	----------

mult	✓	✓	✓
plus	✓	✓	✓
minus	✓	✓	✓
div	✓	✓	✓
mod	✓	NA	NA
logic	✓	✓	✓
arg1	✓	✓	✓

The *rt_T* struct provides the same typedefs as in Table 48 but exclude the *mod* typedef since *ac_complex* does not provide the % operator. For operations that are not commutative it also provides *minus2* and *div2* which are for the operator types swapped. For example, when *ac_complex* queries the type of an *ac_int* minus an *ac_fixed*, it queries the *ac_int::rt_T<ac_fixed>::div* which gets the type from *ac_fixed::rt_T<ac_int>::div2*.

The *rt_2T* struct has two template parameters *T* and *T2*. For Numerical AC types, it is equivalent to *T::rt_T<T2>*. Use of the *rt_2T* struct is more flexible than the *rt_T struct* in that *T* could be a native C++ integer type.

5.3. Macros

This section covers the macros that are available in the numerical AC header types.

5.3.1. User Definable Macros

Table 49: User Definable Macros

Macro	Description
AC_DEFAULT_IN_RANGE	If defined, default constructors will guarantee a value that is in range even though it is un-initialized.
AC_USER_DEFINED_ASSERT	Can be defined to be a call to user defined assert function with arguments: bool condition, const char *file, int line, const char *msg
AC_NOT_USING_INTN	If defined, disables the “using” of namespace <i>ac_intN</i> . See Type Definitions for Signed and Unsigned ac_ints .

5.3.2. Utility Macros

The following macros are defined in *ac_int.h*. All Numerical AC Datatype headers include *ac_int.h* directly or indirectly so these macro definitions should be available by inclusion of any of the numerical header files.

Table 50: Utility Macros

Macro	Description
AC_MAX(a,b)	Max of arguments
AC_MIN(a,b)	Min of arguments
AC_ABS(a)	Absolute value of argument
AC_ASSERT(cond, msg)	Assert. Affected by AC_USER_DEFINED_ASSERT

Chapter 6: Numerical Datatype Migration Guide

This chapter provides detailed explanations on differences between the Algorithmic numerical datatypes and the built-in C integer types and the SystemC integer and fixed-point types.

6.1. General Compilation Issues

When porting algorithms written with either C integer or SystemC datatypes a compilation error may be encountered when the choices for the question mark operator are *ac_int* or *ac_fixed* types. For instance the expression:

```
b ? x : -x;
```

works when *x* is a C integer or a SystemC data type but will error out when *x* is an *ac_int* or *ac_fixed* because *x* and *-x* don't have the same type (their bitwidths are different). Explicit casting may be needed for the question mark operator so that both choices have the exact same type. For example, in the examples below the expressions in the left (using *sc_int*) are re-coded with *ac_int* as follows:

- (c ? a_5s : b_7u) becomes (c ? (ac_int<8,true>) a_5s : (ac_int<8,true>) b_7u)
- (c ? a_5s : - a_5s) becomes (c ? (ac_int<6,true>) a_5s : - a_5s)
- (c ? a_5s : 1) becomes (c ? a_5s : (ac_int<5,true>) 1), or (c ? (int) a_5s : 1)

where variable *a_5s* is a 5-bit wide signed *sc_int* or *ac_int* and so on.

The SystemC datatypes don't require casting because they share the same base class that contains the actual value of the variable. Note that an integer constant such as 1 is of type *int* and will be represented as an *ac_int<32, true>*, so an expression such as *a_4s + 1* will have type *ac_int<33,true>* instead of *ac_int<5,true>*.

6.2. SystemC Syntax

Table 51 shows the SystemC bit-accurate datatypes that *ac_int* and *ac_fixed* can replace. Using *ac_int* and *ac_fixed* it is possible to write generic algorithms that work for any bitwidth and that simulate faster than the “fast” (but limited) SystemC types *sc_int*, *sc_uint*, *sc_fixed_fast*, and *sc_ufixed_fast*.

Table 51: Relation Between SystemC Datatypes and AC Datatypes

SystemC Datatype	New Datatype	Comments
<i>sc_int</i> <W>	<i>ac_int</i> <W,true>	<i>sc_int</i> limited to 64 bits
<i>sc_uint</i> <W>	<i>ac_int</i> <W,false>	<i>sc_uint</i> limited to 64 bits
<i>sc_bigint</i> <W>	<i>ac_int</i> <W,true>	
<i>sc_bignint</i> <W>	<i>ac_int</i> <W,false>	

<code>sc_fixed_fast<W,I,Q,O></code>	<code>ac_fixed<W,I,true,Q,O></code>	<code>sc_fixed_fast</code> limited to mantissa of <i>double</i>
<code>sc_ufixed_fast<W,I,Q,O></code>	<code>ac_fixed<W,I,false,Q,O></code>	<code>sc_ufixed_fast</code> limited to mantissa of <i>double</i>
<code>sc_fixed<W,I,Q,O></code>	<code>ac_fixed<W,I,true,Q,O></code>	
<code>sc_ufixed<W,I,Q,O></code>	<code>ac_fixed<W,I,false,Q,O></code>	

The `ac_int` and `ac_fixed` types have the same parameters with the same interpretation as the corresponding SystemC type. The difference is an extra boolean parameter *S* that defines whether the type is signed (*S*==true) or unsigned (*S*==false). Using a template parameter instead of different type names makes it easier to write generic algorithms (templated) that can handle both signed and unsigned types. The other difference is that `ac_fixed` does not use the “nbits” parameter that is used for the SystemC fixed-point datatypes.

The template parameters *Q* and *O* are enumerations of type `ac_q_mode` and `ac_o_mode` respectively. All SystemC quantization modes are supported as shown in [Table 52](#). Most commonly used overflow modes are supported as shown in [Table 53](#).

Table 52: Quantization Modes for `ac_fixed` and Their Relation to `sc_fixed/sc_ufixed`

<code>ac_fixed</code>	<code>sc_fixed/sc_ufixed</code>
AC_TRN (default)	SC_TRN (default)
AC_RND	SC_RND
AC_TRN_ZERO	SC_TRN_ZERO
AC_RND_ZERO	SC_RND_ZERO
AC_RND_INF	SC_RND_INF
AC_RND_MIN_INF	SC_RND_MIN_INF
AC_RND_CONV	SC_RND_CONV
AC_RND_CONV_ODD	Not available

Table 53: Overflow Modes for `ac_fixed` and Their Relation to `sc_fixed/sc_ufixed`

<code>ac_fixed</code>	<code>sc_fixed/sc_ufixed</code>
AC_WRAP (default)	SC_WRAP, nbits = 0 (default)
AC_SAT	SC_SAT
AC_SAT_ZERO	SC_SAT_ZERO
AC_SAT_SYM	SC_SAT_SYM

All operands are defined consistently with `ac_int`: if both `ac_fixed` operands are pure integers (*W* and *I* are the same) then the result is an `ac_fixed` that is also a pure integer with the same bitwidth and value as the result of the equivalent `ac_int` operation. For example: *a/b* where *a* is an `ac_fixed<8,8>` and *b* is an `ac_fixed<5,5>` returns an `ac_fixed<8,8>`. In SystemC, on the other hand, the result of *a/b* returns 64 bits of

precision (or SC_FXDIV_WL if defined).

6.2.1. SystemC to AC Differences in Methods/Operators

There are methods that have a different name, syntax and semantic. The main one is the range method `range(int i, int j)` or operator `(int i, int j)`. There are two different methods in `ac_int` for accessing or modifying (assigning to) a range. Note that `ac_int` does not support a dynamic length range.

Methods: range in SystemC to `slc` and `set_slc` in `ac_int` or `ac_fixed`

For accessing a range:

```
x.range(i+W-1, i) (or x(i+W-1, i)) becomes x.slc<W>(i)
```

where `x.slc<W>(i)` returns an `ac_int<W, SX>` where `SX` is the signedness of variable `x`. The slice method returns an `ac_int` for both `ac_int` and `ac_fixed`. Also note that `W` must be a constant. For instance `x.range(i, j)` would translate into `x.slc<i-j+1>(j)` provided both `i` and `j` are constants.

For assigning a range:

```
x.range(i+W-1, i) = y (or x(i+W-1, i) = y) becomes x.set_slc(i, y)
```

this assumes that `y` is of type either `ac_int<W, false>` or `ac_int<W, true>`, otherwise it needs to be cast to either type.

Concatenation

The concatenation operator (the “,” operator in `sc_int/sc_uint` and `sc_bigint/sc_bignint`) is not defined in `ac_int` or `ac_fixed`. The solution is to rewrite it using `set_slc`:

```
y = (x, z); becomes y.set_slc(WZ, x); y.set_slc(0, z);  
where WZ is the width of z.
```

Other Methods

Table 54 shows other less frequently used methods in SystemC datatypes that would require rewriting in `ac_int`.

Table 54: Migration of SystemC Methods to `ac_int`

SystemC	ac_int
iszero	operator !
sign	<code>x < 0</code>
bit	<code>x[i]</code>
reverse	no equivalent
test	<code>x[i]</code>
set	<code>x[i] = 1</code>
clear	<code>x[i] = 0</code>

invert	$x[i] = !x[i]$
--------	----------------

Constructors from *char **, are not defined/implemented for *ac_int* and *ac_fixed*.

6.2.2. Support for SystemC *sc_trace* Methods

The Algorithmic C Datatypes package was updated in 2010 to provide support for using SystemC *sc_trace* methods on the AC datatypes. In order to use the *sc_trace* method in your SystemC design, you must include the following headers in the following order:

```
#include <systemc.h>
#include <ac_fixed.h>    (or ac_int.h or ac_complex.h or ac_float.h)
#include <ac_sc.h>
```

Failing to include them in the above order will result in compile errors. In addition to proper include file ordering, you can only trace using VCD format files (i.e. using the *sc_create_vcd_trace_file()* function in SystemC). Using any other trace file format may result in a crash during simulation.

6.3. Simulation Differences with SystemC types and with C integers

In this section the simulation semantics of the bit-accurate datatypes *sc_int/sc_uint*, *sc_bigint/sc_biguint*, and *ac_int* will be compared and contrasted. For simplicity of discussion the shorthand notation *int<bw>* and *uint<bw>* will be used to denote a signed and unsigned integer of bitwidth *bw* respectively. Also *Slong* and *Ulong* will be used to denote the C 64-bit integer types *long long* and *unsigned long long* respectively.

The differences between limited and arbitrary length integer datatypes can be group in several categories as follows:

- Limited precision (64 bit) vs. arbitrary precision
- Differences due to implementation deficiencies of *sc_int/sc_uint*
- Differences due to definition

6.3.1. Limited Precision vs. Arbitrary Precision

Both *sc_int/sc_uint* use the 64-bit C types *long long* and *unsigned long long* as the underlying type to efficiently their operators. In more mathematical terms, that means that the arithmetic is accurate modulo 2^{64} . As long as every value is representable in 2's complement 64-bit signed, the limited precision should not affect the computation and should agree with the equivalent expression using arbitrary precision integers.

6.3.2. Implementation Deficiencies of *sc_int/sc_uint*

At first glance, it would appear that the only difference between limited and arbitrary length datatypes is that arithmetic is limited to 64-bit. However, there are a number of additional issues that have to do with how the *sc_int/sc_uint* are implemented.

The implementations of the limited precision bit-accurate integer types suffer from a number of deficiencies:

- Mixing signed and unsigned can lead to unexpected results. Many operators are not defined so they fall back to the underlying C types *Slong* and *Ulong*. Conversion rules in C change the signed operand to unsigned when a binary operation has mixed *Slong* and *Ulong* operands. This leads to the following non intuitive results:

```
(uint<8>) 1 / (int<8>) -1 = (Ulong) 1 / (Slong) -1 = (Ulong) 1 / (Ulong) -1 = 0
(uint<6>) 1 > (int<6>) -1 = (Ulong) 1 > (Slong) -1 = (Ulong) 1 / (Ulong) -1 = false
(int<6>) -1 >> (uint<6>) 1 = (Slong) -1 >> (Ulong) 1 = -1
```

Note however, that operations such as `+`, `-`, and `*`, `|`, `&`, `^` provided the result is assigned to an integer type of length 64 or less, or is used in expressions that are not sensitive to the signedness of the result:

```
// OK
w_u20 = a_u8 * b_s9 + x_u13 & y_s4;

// Bad, assigning Ulong to 67 signed
sc_bigint<67> i = a_u8 * b_s9 + x_u13 & y_s4;

// Bad, s/s div should be ok, but numerator is Ulong
w_u20 = (a_u8 * b_s9) / c_s6;

// Bad if both f(Ulong) and f(Slong) are defined
f(a_u8 * b_s9);
```

- Shifting has the same limitations as in C. The C language only defines the behavior of integer shifts on a *Slong* or *Ulong* when the shift value is in the range `[0, 63]`. The behavior outside that range is compiler dependent. Also some compilers (Visual C 6.0 for example) incorrectly convert the shift value from *Slong* to *Ulong* if the first operand is *Ulong*.

6.3.3. Differences Due to Definition

The previous two sections covered the high-level and most often encountered differences among the bit-accurate integer datatypes. This section will cover more detailed differences.

Initialization

The SystemC integer datatypes are initialized by default to 0 by the default constructor, whereas *ac_int* is not initialized by the default constructor. If the algorithm relies on this behavior, the initialization needs to be done explicitly when migrating from SystemC integer datatypes to *ac_int*. This issue is not there for fixed-point datatypes as neither the *sc_fixed/sc_ufixed* nor *ac_fixed* initializes by default.

Note that non local variables (that is global, namespace, and *static* variables) don't have this issue as they are initialized by virtue of how C/C++ is defined.

Shift Operators

The *ac_int* and *ac_fixed* shift operators are described in the [Shift Operators](#) section. Shift operations present the most important differences between the Algorithmic C types and the SystemC types.

SystemC Types

- The return type of the left shift for *sc_bigint/sc_biguint* or *sc_fixed/sc_ufixed* does not lose bits making the return type of the left shift data dependent (dependent on the shift value). Shift assigns for *sc_fixed/sc_ufixed* may result in quantization or overflow (depending on the mode of the first operand).
- Negative shifts are equivalent to a zero shift value for *sc_bigint/sc_biguint*
- The shift operators for the limited precision versions is only defined for shift values in the range [0, 63] (see [Implementation Deficiencies of sc_int/sc_uint](#) on page 59).

Differences with Native C Integer Types

- Shifting occurs on either 32-bit (*int*, *unsigned int*) or 64-bit (*long long*, *unsigned long long*) integrals. If the first operand is an integral type that has less than 32 bits (*bool*, (un)signed *char*, *short*) it is first promoted to *int*. The return type is the type of the first argument after integer promotion (if applicable).
- Shift values are constrained according to the length of the type of the promoted first operand.
 - $0 \leq s < 32$ for 32-bit numbers
 - $0 \leq s < 64$ for 64-bit numbers
- The behavior for shift values outside the allowed ranges is not specified by the C++ ISO standard.

The shift left operator of *ac_int* returns an *ac_int* of the same type (width and signedness) of the first argument and it is not equivalent to the left shift of *sc_int/sc_uint* or *sc_bigint/sc_biguint*. To get the equivalent behavior using *ac_int*, the first argument must be of wide enough so that it does not overflow. For example

```
(ac_int<1,false>) 1 << 1 = 0  
(ac_int<2,false>) 1 << 1 = 2
```

Both the right and left shift operators of *ac_int* return an *ac_fixed* of the same type (width, integer width and signedness) of the first argument and is not equivalent to the corresponding operator in *sc_fixed/sc_ufixed*. Despite the fact that there might be loss of precision when shifting an *ac_fixed*, no quantization or overflow is performed. The first argument must be large enough width and integer width to guarantee that there is no loss of precision.

Differences for the range/slice Methods

It is legal to access bits to the left of the MSB of an *ac_int* or an *ac_fixed* using the *slc* method. The operation is treated arithmetically (as if the value had been represented in the appropriate number of bits).

The following operations are invalid and will generate a runtime error (assert) during C simulation:

- Attempting to access negative indices with the *slc* method
- Attempting to access or modify indices outside the 0 to *W*-1 range for *set_slc* or the *[]* operator

The behavior for indices outside the 0 to $W-1$ range for SystemC datatypes is not consistent. For example *sc_int/sc_uint* and *sc_fixed/sc_ufixed* don't allow it (runtime error) while *sc_bigint/sc_bignint* allow even negative indices (changed to a 0 index).

Conversion Methods

The conversion methods **to_int()**, **to_long()**, **to_int64()**, **to_uint()**, **to_uint64()** and **to_ulong()** for *sc_fixed/sc_ufixed* are implemented by first converting to *double*. For instance:

```
sc_fixed<5,3> x = ...;
int t = x.to_int();           // equiv to (int)(double)x
                               // not equiv to (int)(sc_int<32>)x

ac_fixed<5,3,true> y = ...;
int t = y.to_int();           // equiv to x.to_ac_int().to_int();
```

The difference in most cases will be subtle (*double* has a signed-magnitude representation so it truncates towards zero instead of truncating towards minus infinity) but could be very different if the number would overflow the *int* or *long long* C types.

Neither *ac_int* nor *ac_fixed* provide a conversion operator to *double* (an explicit *to_double* method is provided). SystemC datatypes do provide that conversion. There are a number of cases where that can lead to non intuitive semantics:

```
sc_fixed<7,4> x = ...;
int t = (int) x;               // equiv to (int)(double) x
bool b = !x;                   // equiv to ! (double) x
```

Unary Operators ~, - and Binary Operators &, |, ^

The unary operators *~* and *-* for *ac_int* and *ac_fixed* will return a signed typed. This behavior is consistent with the SystemC integer types but inconsistent with the SystemC fixed-point types.

A common issue when migrating from C/C++ that uses shifting and masking is the following:

```
unsigned int x = 0;
unsigned mask = ~x >> 24; // mask is 0xFF

ac_int<32,false> x = 0;
ac_int<32,false> mask = ~x >> 24; // mask is 0xFFFFFFFF
```

The reason for this discrepancy is that for C integers the return type for the unary operators *~* and *-* is the type of the promoted type for the operand. If the argument is *signed/unsigned int*, *long* or *long long*, integer promotion does not change the type. For example, when the operand is *unsigned int*, then the return type of either *~* or *-* will be *unsigned int*. Note however that *unsigned char* and *unsigned short* get promoted to *int* which makes the behavior consistent with *ac_int*:

```
unsigned short x = 0;
unsigned short mask = ~x >> 8; // mask is 0xFFFF, not 0xFF

ac_int<16,false> x = 0;
ac_int<16,false> mask = ~x >> 8; // mask is 0xFFFF
```

The arithmetic definition of the operator *~* makes the value result independent of the bit-width of the operand:

```
if x == y, then ~x == ~y    // x and y may be different bitwidths
```

Also the arithmetic definition is consistent with the arithmetic definition of the binary (two operand) logical operators `&`, `|`, and `^`. For instance:

```
~(a | b) == ~a & ~b
```

The arithmetic definition of the logical operators `&`, `|`, `^` is necessary since signed and unsigned operands of various bit-widths may be combined.

6.3.4. Mixing Datatypes

This section describes the conversion functions that are used to interface between the bit-accurate integer datatypes.

Conversion Between `sc_int/sc_uint` and `ac_int`

Use the C integer conversions to go from `sc_int/sc_uint` to `ac_int` and vice versa:

```
ac_int<54,true> x = (Slong) y; // y is sc_int<54>
sc_int<43> y = (Slong) x; // x is ac_int<43, true>
```

Conversion Between `sc_bigint/sc_biguint` and `ac_int`

The C integer datatypes can be used to convert between integer datatypes without loss of precision provided the bitwidth does not exceed 64 bits:

```
ac_int<54,true> x = y.to_int64(); // y is sc_bigint<54>
sc_bigint<43> y = (Slong) x; // x is ac_int<43,true>
ac_int<20,true> x = y.to_int(); // y is sc_biguint<20>
```

Explicit conversion functions are provided between the datatypes `sc_bigint/sc_biguint` and `ac_int` and between `sc_fixed/sc_ufixed` and `ac_fixed`. They are provided in a different include file:

```
<ac_datatype_install>/include/ac_sc.h
```

which define the following functions:

```
template<int W> ac_int<W, true> to_ac(const sc_bigint<W> &val);
template<int W> ac_int<W, false> to_ac(const sc_biguint<W> &val);
template<int W> sc_bigint<W> to_sc(const ac_int<W, true> &val);
template<int W> sc_biguint<W> to_sc(const ac_int<W, false> &val);

template<int W, int I, sc_q_mode Q, sc_o_mode O, int nbits>
ac_fixed<W, I, true> to_ac(const sc_fixed<W, I, Q, O, nbits> &val);

template<int W, int I, sc_q_mode Q, sc_o_mode O, int nbits>
ac_fixed<W, I, false> to_ac(const sc_ufixed<W, I, Q, O, nbits> &val);

template<int W, int I, ac_q_mode Q, ac_o_mode O>
sc_fixed<W,I> to_sc(const ac_fixed<W, I, false, Q, O> &val);

template<int W, int I, ac_q_mode Q, ac_o_mode O>
sc_ufixed<W,I> to_sc(const ac_fixed<W, I, true, Q, O> &val);
```

For example:

```
sc_bigint<123> x = to_sc(y); // y is ac_int<123, true>
```

Chapter 7: Frequently Asked Questions on Numerical Datatypes

This section contains answers to some frequently asked questions. These questions are divided into the following sections:

- [Operators \$\sim\$, \$\&\$, \$|\$, \$\wedge\$, \$-\$, \$!\$](#)
- [Conversions to double and Operators with double](#)
- [Constructors from strings](#)
- [Shifting Operators](#)
- [Division Operators](#)
- [Compilation Problems](#)
- [Platform Dependencies](#)
- [Purify Reports](#)
- [User Defined Asserts](#)

7.1.1. Operators \sim , $\&$, $|$, \wedge , $-$, $!$

The following section describes common issues with these operators.

Why does \sim and $-$ for an unsigned return signed?

For a table of return values, refer to the [Unary Operator](#) section.

Why are operators $\&$, $|$, \wedge “arithmetically” defined?

The two operands may have different signedness, have different bit-widths or have non aligned fixed-points (for *ac_fixed*). An arithmetic definition makes the most sense in this case.

Why does operator $!$ return different results for *ac_fixed* and *sc_fixed*?

The $!$ operator is not defined for *sc_fixed* or *sc_ufixed*. The behavior for *sc_fixed* is then equivalent to first casting it to *double* and then applying the $!$ operator which is not correct.

7.1.2. Conversions to double and Operators with double

The following section describes common issues with using *double* datatypes.

Why is the implicit conversion from *ac_fixed* to *double* not defined?

The reason that there is no implicit conversion function to *double* is that it is impractical to define mixed *ac_fixed* and *double* operators. For example, if there was an implicit conversion to *double* the expression “*x* + 0.1” would be computed as (double) *x* + 0.1 even when *x* has more bits of precision than the *double*, thus resulting in an unintended loss of precision.

Why are most binary operations not defined for mixed *ac_fixed* and *double* arguments?

Consider the expression “*x* + 0.1” where *x* is of type *ac_fixed*. Arithmetic operators such as + are defined in such a way that they return a result that does not lose precision. In order to accomplish that with a mixed fixed-point and *double* operator +, the *double* would have to be converted to a fixed-point that is able to represent all values that a *double* can assume. This would require an impractically large *ac_fixed*. Note that the actual value of the constant is not used by a C++ compiler to determine the template parameters for the minimum size *ac_fixed* that can hold it.

Comparison operators are defined for mixed *ac_fixed* and *double* arguments as the result is a *bool* and there is no issue about losing precision. However, the comparison operator is much less efficient in terms of runtime than using a comparison to the equivalent *ac_fixed*:

```
while( ... ) {  
    if( x > 0.5 ) // less efficient  
        ...  
}
```

could be made more efficient by storing the constant in an *ac_fixed* so that the overhead of converting from *double* to *ac_fixed* is incurred once (outside the loop):

```
ac_fixed<1,0,false> c0_5 = 0.5;  
while( ... ) {  
    if( x > c0_5 ) // more efficient  
        ...  
}
```

7.1.3. Constructors from strings

The following section describes common issues with strings.

Why are constructors from strings not defined?

They would be very runtime inefficient.

7.1.4. Shifting Operators

The following section describes common issues with shifting operators.

Why does shifting gives me unexpected results?

The shift operation for *ac_int/ac_fixed* differs from the shift operations in SystemC and native (built-in) C integers. See [Shift Operators](#) section. The main difference is that the shift operation for *ac_int/ac_fixed* returns the type of the first operand.

```
ac_int<2,false> x = 1;
x << 1; // returns ac_int<2,false> (2), value is 2
x << 2; // returns ac_int<2,false> (4), value is 0
(ac_int<3,false>) x << 2; // returns ac_int<3,false> (4), value is 4
```

The main reason for this semantic is that for an arbitrary-length type, a definition that returns a fully arithmetic value requires a floating return type which violates the condition that the return type should be an *ac_int* or an *ac_fixed* type. Supporting a floating return type creates a problem both for simulation speed and synthesis. For example the type of the expression

```
a * ( (x << k) + y)
```

can not be statically determined.

7.1.5. Division Operators

The following section describes common issues with division operators.

Why does division return different results for *ac_fixed* and *sc_fixed*?

Division for *sc_fixed/sc_ufixed* returns 64 bits of precision (or whatever SC_FXDIV_WL is defined as). The return type for *ac_fixed* is defined depending on the parameters of both the dividend and divisor (refer to the [return type table](#) for *ac_fixed*).

7.1.6. Compilation Problems

The following section describes common compilation issues.

Why aren't older compilers supported?

The support of templates is not adequate in older compilers.

Why doesn't the *slc* method compile in some cases?

When using the *slc* method in a templated function place the keyword *template* before it as some compilers may error out during parsing. For example:

```
template<int N>
int f(int x) {
    ac_int<N,true> t = x;
```

```
ac_int<6,true> r = t.template slc<N>(4); // t.slc<N>(4) could error out
return r.to_int();
}
```

Without the keyword `template` the “`t.slc<N>(4)`” is parsed as “`t.slc < N`” since it does not know whether `slc` is a data member or a method (this is known once template function `f` and therefore `ac_int<N,true>` is instantiated).

Why do I get compiler errors related to template parameters?

If this happen while using the GCC compiler, the error might be related to the template bug on GCC that was fixed in version 4.0.2 (http://gcc.gnu.org/bugzilla/show_bug.cgi?id=23789). This compiler bug rarely showed up when using previous versions of `ac_int/ac_fixed` and is even less likely on the current version of `ac_int/ac_fixed`.

7.1.7. Platform Dependencies

The following section describes the supported platforms.

What platforms are supported?

The current implementation assumes that an *int* is 32 bits and that a *long long* is 64-bits, both in 2’s complement representation. These assumptions need to be met for correct simulation of the data types. In addition a *long* is assumed to be 32 bits wide, a *short* is assume to be 16 bits and a *char* is assumed to be 8 bits wide. A plain *char* (neither *signed* or *unsigned*) is assumed to be *signed*. These assumptions are only relevant if the types are used to initialize/construct an `ac_int` or `ac_fixed` or they are used in expressions with `ac_int` or `ac_fixed`.

7.1.8. Purify Reports

The following section describes common issues with using Purify with the AC datatypes.

Why do I get UMRs for `ac_int/ac_fixed` in purify?

The following code will report a UMR in purify:

```
ac_int<2,false> x;
x[0] = 0;          // UMR
x[1] = 1;
```

The UMR occurs because `x` is not initialized, but setting a bit (or a slice) requires accessing the original (un-initialized) value of `x`.

A second source of UMRs is explicit calls to un-initialize an `ac_int/ac_fixed` that were declared static (see section using the [ac::init_array function](#)). This is used mostly for algorithms written for hardware design.

7.1.9. User Defined Asserts

The following section describes common issues with using asserts.

Can I control what happens when an assert is triggered?

Control over what happens when an assert is triggered is accomplished by defining the compiler directive `AC_USER_DEFINED_ASSERT` to a user defined assert function before the inclusion of the AC Datatype header(s). The following example illustrates this:

```
void my_assert(bool condition, const char *file=0, int line=0, const char *msg=0);  
#define AC_USER_DEFINED_ASSERT(cond, file, line, msg) my_assert(cond, file, line, msg)  
#include <ac_int.h>
```

When `AC_USER_DEFINED_ASSERT` is defined, the system header `<ostream>` is included instead of `<iostream>`, as `std::cerr` is no longer required by `ac_assert` in that case. This feature was introduced to reduce the application startup time penalty that can occur when including `iostream` for some compilers that don't support `"#pragma once"`. That startup time penalty is proportional on the number of translation units (static constructor for each *.o file that includes it).

Chapter 8: **ac_channel Datatype**

Several synthesis and modeling problems become apparent when attempting to use C function calls to describe a hierarchical system. The purpose of the AC (Algorithmic C) channel class is to simplify both synthesis and modeling with a minimal impact on coding style and C simulation performance.

The first problem is primarily a synthesis problem. To build a streaming interface, you need to be able to guarantee that the data communicated between two blocks is read and written in the same order. Since an `ac_channel` enforces a FIFO discipline, it guarantees that the hardware can build for streaming interfaces.

The second problem is that systems are often made up of blocks that run using different amounts of data. The first block in a system might produce data 100 values at a time, while the reader may have been written to consume data 128 elements at a time. AC Channels automate the connection of these blocks.

Finally, many systems have feedback and the length of the feedback path must be defined in the C source. Changing the feedback path length during synthesis would result in a design that simulates differently, which would violate the most fundamental rule of Algorithmic Synthesis. AC Channels provide a way to define a feedback path so it can be easily simulated and synthesized.

8.1. The `ac_channel` Class Definition

The fundamental operation that High Level Synthesis performs is to convert a single-threaded C/C++ model into a high-performance, multiple-process RTL model. Each hierarchical function in the C/C++ model becomes a process in the RTL model. To allow for more efficient synchronization of the data among processes, the `ac_channel` class may be used. The restrictions are necessary to insure that the C/C++ model and the RTL model have the same functionality while allowing the RTL to be efficiently implemented.

The `ac_channel` class is a C++ template class that enforces a FIFO discipline (reads occur in the same order as writes.) From a modeling perspective, an `ac_channel` is implemented as a simple interface to the C++ standard queue (`std::deque`). That is, for modeling purposes, an `ac_channel` is infinite in length (writes always succeed) and attempting to read from an empty channel generates an assertion failure (reads are non-blocking).

Within the C/C++ model, an instance of an `ac_channel` may be written by a single hierarchical function and read by a single, different, hierarchical function. In the resulting hardware, the FIFO corresponding to an `ac_channel` will be written from a single process and read by a single process - thus implementing point-to-point communication. For purposes of this discussion, the external world counts as a process. If a channel appears in the top-level interface of a design and is only written by the design, then the external environment is assumed to be the unique reader. Likewise, the external environment is assumed to be the writer of a channel which is only read.

8.2. ac_channel Member Functions

The `ac_channel` constructor includes the ability to add a set of values that are read before the inputs to the channel are read. This is most commonly used to set the size of the feedback path, but also has other applications in the feed-forward path to control the ramp-up of the system.

8.2.1. Member Function: `ac_channel()`

Example usage:

```
ac_channel<ac_fixed<12,6> > my_channel;
```

Description:

Note: You must put a space between any two ">" characters or you will get a compiler error because the parser treats ">>" as a right shift operator.

This will construct a channel with a 6.6 fixed-point datatype. No other information is required to create a channel.

8.2.2. Member Function: `ac_channel(prefill_num)`

Example usage:

```
ac_channel<int> my_channel(16);
```

Description:

This constructs a channel with a 32-bit value and pre-fills the channel with 16 values. For the built-in C++ datatypes *char*, *short*, *int*, *unsigned*, *long*, *float* and *double* the value in the channel is a different random number for each element in the fifo. For all other datatypes, the default constructor is called to set the values.

8.2.3. Member Function: `ac_channel(prefill_num, value)`

Example usage:

```
typedef struct { int re; int im; } complex;
const complex init_complex = {4,5};
ac_channel <complex> my_channel(16, init_complex);
```

Description:

This will construct a channel and pre-fill the channel with 16 complex variables where "*re*" is 4 and "*im*" is 5. This channel is most commonly used for designs with feedback to set the depth of the feedback path while still providing valid data.

8.3. Synthesizable Member Functions

These member functions can be synthesized into hardware. All constructor member functions are also synthesizable. The specific hardware that is created for each of these member functions will be discussed

when the different design styles are covered.

8.3.1. Member Function: val read() or read(&val)

Example usage:

```
ac_channel<int> my_channel(16);
...
int first_value = my_channel.read();
int second_value;
my_channel.read(second_value);
```

Description:

The read member functions are equivalent and two member functions are only provided to allow different coding styles. The code above would read two values, the first would be put into *first_value* and the second would be put into *second_value*.

NOTE: Do not call the read() function multiple times for the same channel in a single statement. This is because the order in which expressions and subexpressions are evaluated (left-to-right, right-to-left or other) is undefined by C/C++ language standards, and consequently varies from one compiler to another. Unexpected results can occur if a code statement contains order dependent (sub)expressions in which a variable is both read and modified. For example, the following statement is potentially problematic because it is left to the compiler to decide the order in which the read() functions are evaluated:

```
output->write( input1->read()*5 + input1->read()*7 )
```

8.3.2. Member Function: bool nb_read(&val)

Example usage:

```
ac_channel<uint8> &pri0, &pri1;
uint8 read_out1, read_out2;
...
if (pri0.nb_read(read_out1))
    out.write(read_out1);
else if (pri1.nb_read(read_out2))
    out.write(read_out2);
```

Description:

The “bool nb_read(T &val)” function will test for the presence of data in the channel and read that data if it is available. This function returns boolean false if no data is available. Otherwise it returns true, pops the first value in the channel and assigns it to the “val” argument. It is capable of completing in one cycle in hardware.

NOTE: When you are verifying the behavior of the C++ design against the resulting RTL netlist, the non-blocking read execution order may be different between the C++ and the RTL. The RTL will implement a viable schedule of the C++ but the C++ testbench is not guaranteed to provide data in a way that causes the C++ execution to exactly match the RTL execution.

8.3.3. Member Function: write(val)

Example usage:

```
ac_channel<int> my_channel(16);
...
int temp = f(x);
my_channel.write(temp);
```

Description:

Write a new value into the channel. Since channels are typed, the value written must be convertible to the type stored in the channel.

8.3.4. Member Function: bool available(num)

Example usage:

```
if (my_channel.available(N)) {
    acc = 0;
    for (int i = 0; i < N; i++)
        acc += my_channel.read();
}
```

Description:

The available member function is used to guard reads from a channel. This check must be done at the start of any function where you cannot predict the amount of data in the input channel.

You may pass an integer variable to the available function. During synthesis, the “available()” function is always considered to return true and the read from the channel is changed to be a blocking read. This transformation allows fast C simulation while still providing good hardware performance.

8.3.5. Member Function: int size()

Example usage:

```
ac_channel<uint8> &pri0;
...
if (pri0.size() > 0)
    out.write(pri0.read());
```

Description:

This member function returns the number of data elements currently in the ac_channel. It synthesizes as non-blocking.

Notes:

- Currently size() is supported only for reading channels, not for writing.
- You may not read the size of the channel and also read data in the channel in a single cycle. For example, the following statement requires two cycles:

```
if (input.size() > 0)
```

```
data = input.read();
```

- Do not use “==” operator to test the return value of size() because it could create a deadlock situation. For example, if the code is testing for “size == 1” and size grows to 2 before the test, then the statement will never be true and the design will deadlock. Use other relational operators (>, >=, !=, and so on) instead.

8.4. Non-synthesizable Member Functions

These member functions may only be used in testbenches or other areas of the C++ code that will not be synthesized.

8.4.1. Member Function: bool empty()

Example usage:

```
ac_channel<int> my_channel;
my_channel.write(10);
my_channel.write(45);
consume_data(my_channel); // This function reads from the channel
if (my_channel.empty())
    cout << "All values were read from my_channel and it is now empty" << endl;
```

Description:

This function returns true if there are no values stored in the channel. This is equivalent to checking if size() returned zero.

8.4.2. Member Function: bool operator ==

Example usage:

```
#include <ac_channel.h>
#include <ac_int.h>

int main (void) {

    ac_channel<int> chan1;
    ac_channel<int> chan2;

    chan1.write(10);
    chan2.write(10);

    if (chan1 == chan2)
        std::cout << "EQUAL" << std::endl;
    else
        std::cout << "NOT EQUAL" << std::endl;

    return 0;
}
```

Description:

This function compares every element in a channel to verify that the two channels have exactly the same contents. This function requires that the base type (in this case, int) be the same for both operands. This function will return false if any element in the two channels is different or if the number of elements in the channels is different.

8.4.3. Member Function: bool operator !=

Example usage:

```
#include <ac_channel.h>
#include <ac_int.h>

int main (void) {

    ac_channel<int> chan1;
    ac_channel<int> chan2;

    chan1.write(10);
    han2.write(15);

    if (chan1 != chan2)
        std::cout << "EQUAL" << std::endl;
    else
        std::cout << "NOT EQUAL" << std::endl;

    return 0;
}
```

Description:

This function compares two channels and returns true if the number of the channels have a different number of elements or if the contents of the channels are different. This is equivalent to the negation of the == operator.

8.4.4. Member Function: val operator[int]

Example usage:

```
#include <ac_channel.h>
#include <ac_int.h>

int main (void) {

    ac_channel<int> chan1;

    ...
    // This loop prints the values from the channel without popping/destroying
    for (int j = 0; j < chan1.size(); j++)
        std::cout << chan1[j];
    return 0;
}
```

Description:

The [] operator returns a value that is stored in the channel without popping the value from the channel (i.e.

it is a peek). If you think of the channel as an array with element zero being the next one to be read, then this operator returns an element from that array. For example, `my_chan[0]` is the value the would be read next from the FIFO and `my_chan[my_chan.size()-1]` is the last value written into the channel.

8.4.5. Member Function: `reset()`

Example usage:

```
ac_channel<int> my_channel(16);
...
if (clear_chan_condition)
    my_channel.reset();
```

Description:

This function will set the channel to the same state that it was in immediately after construction. So, if the channel has a set of initial values, then the same values will be in the channel after `reset()`. All data currently in the channel is lost.

NOTE: This function may cause simulation mismatches when used on any channel implemented in the RTL or directly connected to the RTL. Calling `reset` on these channels will cause the state of the C simulation to mismatch with the hardware simulation and may cause the testbench to report a failure on a correct design.

8.5. Example Design Using Hierarchical Blocks With `ac_channel`

The most common use of `ac_channels` is as a communication pipe between two hierarchical blocks in a design. In this case, the blocks are assumed to run continuously and only stall when there is insufficient input. In the example design below, the block “`block_add_avg()`” reads two input values from each input channel, computes their average and returns the sum of the averages. The second block “`block_mult()`” returns the product of the two input values.

```
void block_add_avg(
    ac_channel<int> &op1,
    ac_channel<int> &op2,
    ac_channel<int> &ret)
{
#ifdef __SYNTHESIS__
    while (op1.available(2) && op2.available(2))
#endif
    {
        int avg1 = (op1.read()+op1.read())/2; // average 2 samples from op1
        int avg2 = (op2.read()+op2.read())/2; // average 2 samples from op2
        ret.write(avg1+avg2); // return sum of averages
    }
}

void block_mult(
    ac_channel<int> &op1,
    ac_channel<int> &op2,
    ac_channel<int> &ret)
{
```

```
#ifndef __SYNTHESIS__
while (op1.available(1) && op2.available(1))
#endif
{
    ret.write(op1.read() * op2.read());
}

void top(
    ac_channel<int> &add1,
    ac_channel<int> &add2,
    ac_channel<int> &coeff,
    ac_channel<int> &result)
{
    static ac_channel<int> sum;
    block_add_avg(add1, add2, sum);
    block_mult(sum, coeff, result);
}
```

In this example, the execution of the body of `block_add_avg()` and `block_mult()` is guarded by calls to the `channel available()` functions. This allows the C++ simulation to continue running a block as long as there is data in the channel. When the channels are empty, the C++ design basically does nothing until the testbench pushes more data into the channels. Likewise, in the resulting hardware the `available()` calls are filtered out and each blocks runs continuously as long as the input channel/pipe has data to feed.

8.6. Example Design Using Non-Blocking `size()` Method

This example arbiter design contains four processes that are all trying to write to a shared resource, the output bus. The arbiter function implements a rotation-based (or round robin) arbitration algorithm to read from each process, one after the other. Using non-blocking reads, if no data is available when the read is attempted, the arbiter moves on to the next process. Illustration 1 illustrates the major features of the design.

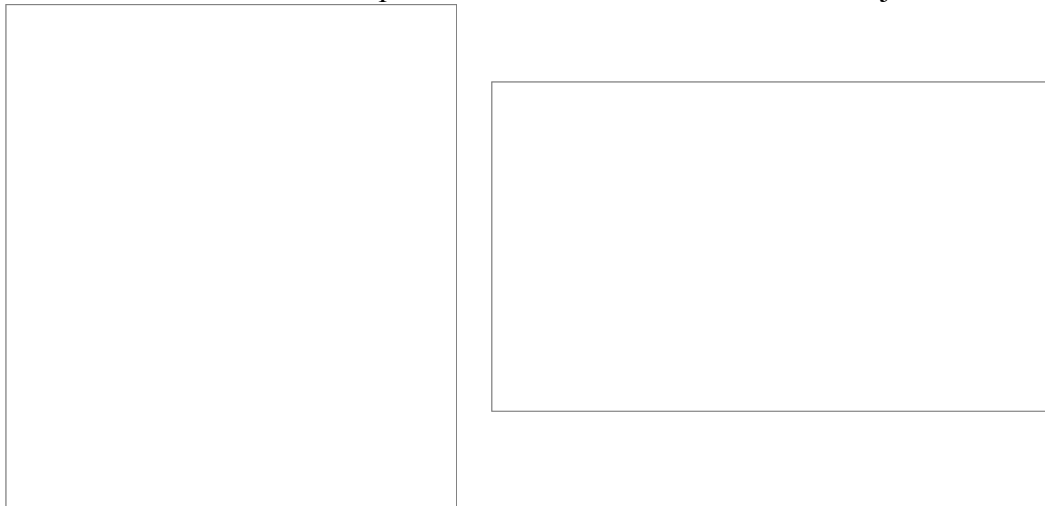


Illustration 1: Graphics View of Design Example

The C++ code implementation is shown in the following example. In, `top.cpp`, blocks A,B,C,D implement the client processes trying to access the output bus. Each of these blocks reads its input `ac_channel` using the

blocking `read()` method, then writes the modified data to its output FIFO on each successful read.

The arbiter block first uses the non-blocking `size()` method to check for data availability in the output FIFO of each process block. Second, for those processes that have data in their FIFOs, the arbiter selects one process to be served based on the rotation arbitration algorithm. Third, it reads the selected process FIFO (blocking method) and writes the data to the bus. Finally, the arbiter updates the serving priority based on rotation algorithm.

Example 1 Arbiter Example Using Non-Blocking `size()` Method

```
#include <ac_channel.h>
#include <ac_int.h>

#pragma hls_design
void block_A(ac_channel<uint8> &in, uint8 &offset, ac_channel<uint8> &out)
{
    #ifndef __SYNTHESIS__
        while (in.available(1))
    #endif
    {
        out.write(in.read()+offset);
    }
}

#pragma hls_design
void block_B(ac_channel<uint8> &in, uint8 &offset, ac_channel<uint8> &out)
{
    #ifndef __SYNTHESIS__
        while (in.available(1))
    #endif
    {
        out.write(in.read()+offset);
    }
}

#pragma hls_design
void block_C(ac_channel<uint8> &in, uint8 &offset, ac_channel<uint8> &out)
{
    #ifndef __SYNTHESIS__
        while (in.available(1))
    #endif
    {
        out.write(in.read()+offset);
    }
}

#pragma hls_design
void block_D(ac_channel<uint8> &in, uint8 &offset, ac_channel<uint8> &out)
{
    #ifndef __SYNTHESIS__
        while (in.available(1))
    #endif
    {
        out.write(in.read()+offset);
    }
}
```

```
#pragma hls_design
void arbiter(ac_channel<uint8> &pri0,
            ac_channel<uint8> &pri1,
            ac_channel<uint8> &pri2,
            ac_channel<uint8> &pri3,
            ac_channel<uint8> &out)
{
    static uint4 priority = 1;
    ac_int<4,false> p_rdy;

    // Non-blocking "size()" function is used here to
    // check for output data in all process block FIFOs */
    p_rdy[0] = pri0.size() > 0;
    p_rdy[1] = pri1.size() > 0;
    p_rdy[2] = pri2.size() > 0;
    p_rdy[3] = pri3.size() > 0;

    uint4 temp_priority=priority;
    for (int i=0;i<4;i++){
        if (p_rdy & temp_priority){
            p_rdy = p_rdy & temp_priority;
            break;
        }
        uint1 temp = temp_priority[3];
        temp_priority = temp_priority << 1;
        temp_priority[0] = temp;
    }

    // Arbiter read data from the selected process block
    // and write it to the bus (arbiter output FIFO)
    if (p_rdy[0])
        { out.write(pri0.read()); }
    else if (p_rdy[1])
        { out.write(pri1.read()); }
    else if (p_rdy[2])
        { out.write(pri2.read()); }
    else if (p_rdy[3])
        { out.write(pri3.read()); }

    // Update priority setting for next cycle.
    if (p_rdy) {
        priority = temp_priority;
        uint1 temp = priority[3];
        priority <<= 1;
        priority[0] = temp;
    }
}

#pragma hls_design top
void top(uint8 & offset1,
        uint8 & offset2,
        uint8 & offset3,
        uint8 & offset4,
        ac_channel <uint8> &in_A,
        ac_channel <uint8> &in_B,
        ac_channel <uint8> &in_C,
        ac_channel <uint8> &in_D,
        ac_channel <uint8> &out)
```

```
{
    static ac_channel<uint8> pri0, pri1, pri2, pri3;
    block_A(in_A, offset1, pri0);
    block_B(in_B, offset2, pri1);
    block_C(in_C, offset3, pri2);
    block_D(in_D, offset4, pri3);
    arbiter(pri0,pri1,pri2,pri3,out);
}
```

Using the following testbench, you can verify the arbiter design.

NOTE: Because this design uses non-blocking reads, the simulation that compares the RTL against the C++ may produce mismatches due to the execution order of the non-blocking reads in the C++ design. The RTL will produce a viable implementation but the simulation may not match the C++ design.

Example 2 Testbench for Arbiter Example Design

```
#include <ac_channel.h>
#include <ac_int.h>

int tb_arbiter(int argc, char *argv[]) {
    static ac_channel<uint8> in_A, in_B, in_C, in_D;
    static ac_channel<uint8> out;

    uint8 tmp;
    uint8 offset = 0;
    for (int i=0; i<10;i++) {
        in_A.write(rand() % 255);
        in_B.write(rand() % 255);
        in_C.write(rand() % 255);
        in_D.write(rand() % 255);

        top(offset, offset, offset, offset,
            in_A, in_B, in_C, in_D, out);

        while(out.available(1)) { // dump outputs
            tmp = out.read();
            printf("out = %d\n",tmp.to_int());
        }
    }
    return 0;
}
```