

Data Compression

Lecture 6

Dr. Mona M.Soliman

Faculty of Computers and Artificial Intelligence

Cairo University

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Arithmetic Coding

Drawback of Huffman Coding

The main drawback of Huffman scheme is that it has problems when there is a symbol with very high probability

Minimum Codeword Length for any Symbol is “ONE” bit, there is no “Half Bit” or “Quarter Bit”

Arithmetic coding completely bypasses the idea of replacing an input symbol with a specific code.

Instead, it takes a stream of input symbols and replaces it with a single floating point number

The longer and more complex the message, the more bits are needed to represent the output number

Arithmetic Coding Algorithm

- Find the probability for each Symbol,
- Calculate **Low_Range**, **High_Range** of Symbol
- Divide the initial subinterval based on the distribution
(**Low_Range**, **High_Range** of Symbols)
- The first Symbol is coded by extracting the subinterval corresponding to it, and subdivide again based on the same relative distribution
- Repeat last step for each Symbol until a final interval is determined
- Any value within this subinterval can be used to represent the sequence of Symbols

Example(1) Arithmetic Coding

Given a long Sequence of characters A, B, and C

the probabilities of Characters are

- $P(A)=0.8$
- $P(B)=0.02$
- $P(C)= 0.18$
- Compress the following part of the sequence “**ACBA**” using Arithmetic Coding

Arithmetic Coding

Rules To solve

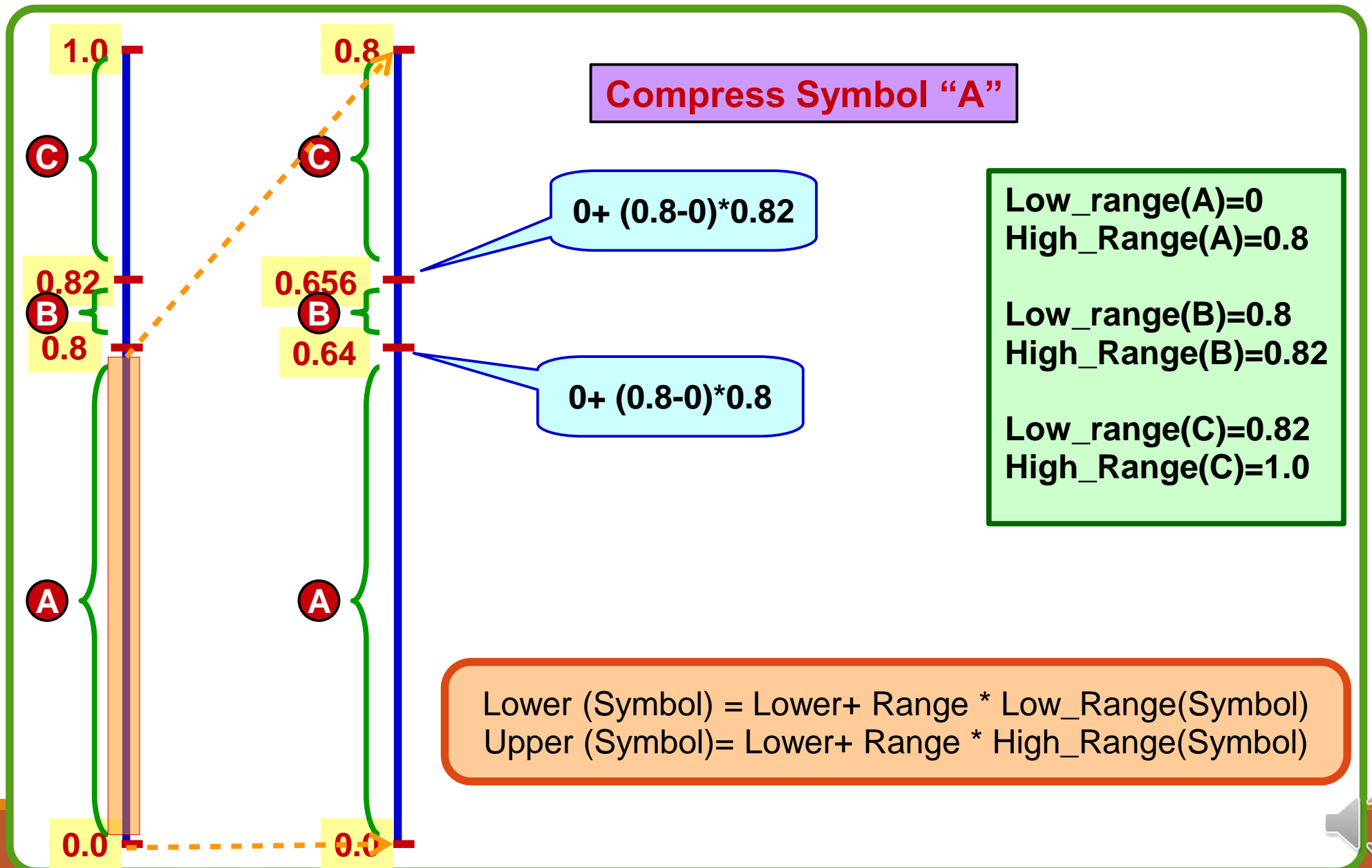
**Low_range(A)=0
High_Range(A)=0.8**

**Low_range(B)=0.8
High_Range(B)=0.82**

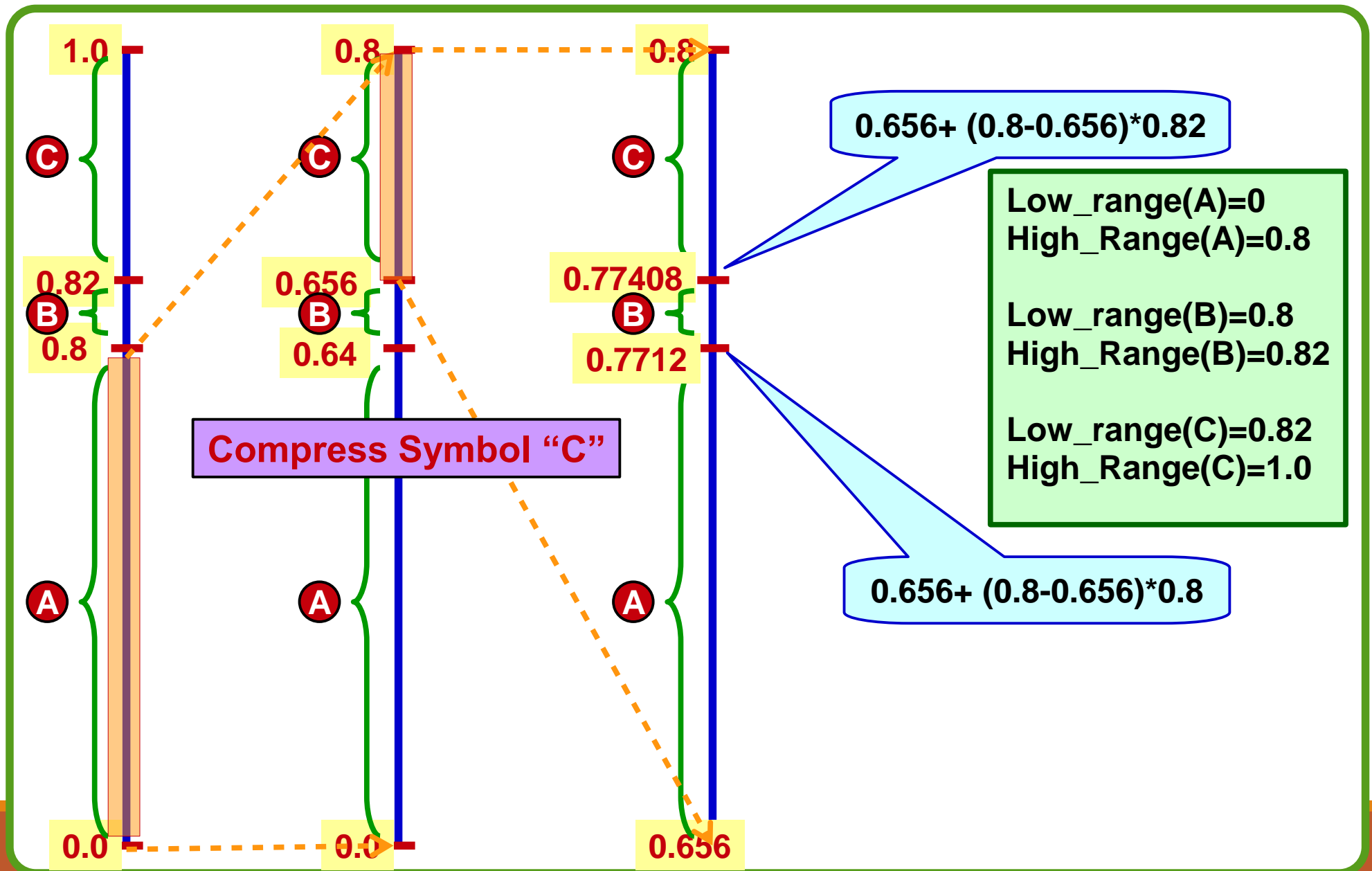
**Low_range(C)=0.82
High_Range(C)=1.0**

**Lower (Symbol) = Lower+ Range * Low_Range(Symbol)
Upper (Symbol)= Lower+ Range * High_Range(Symbol)**

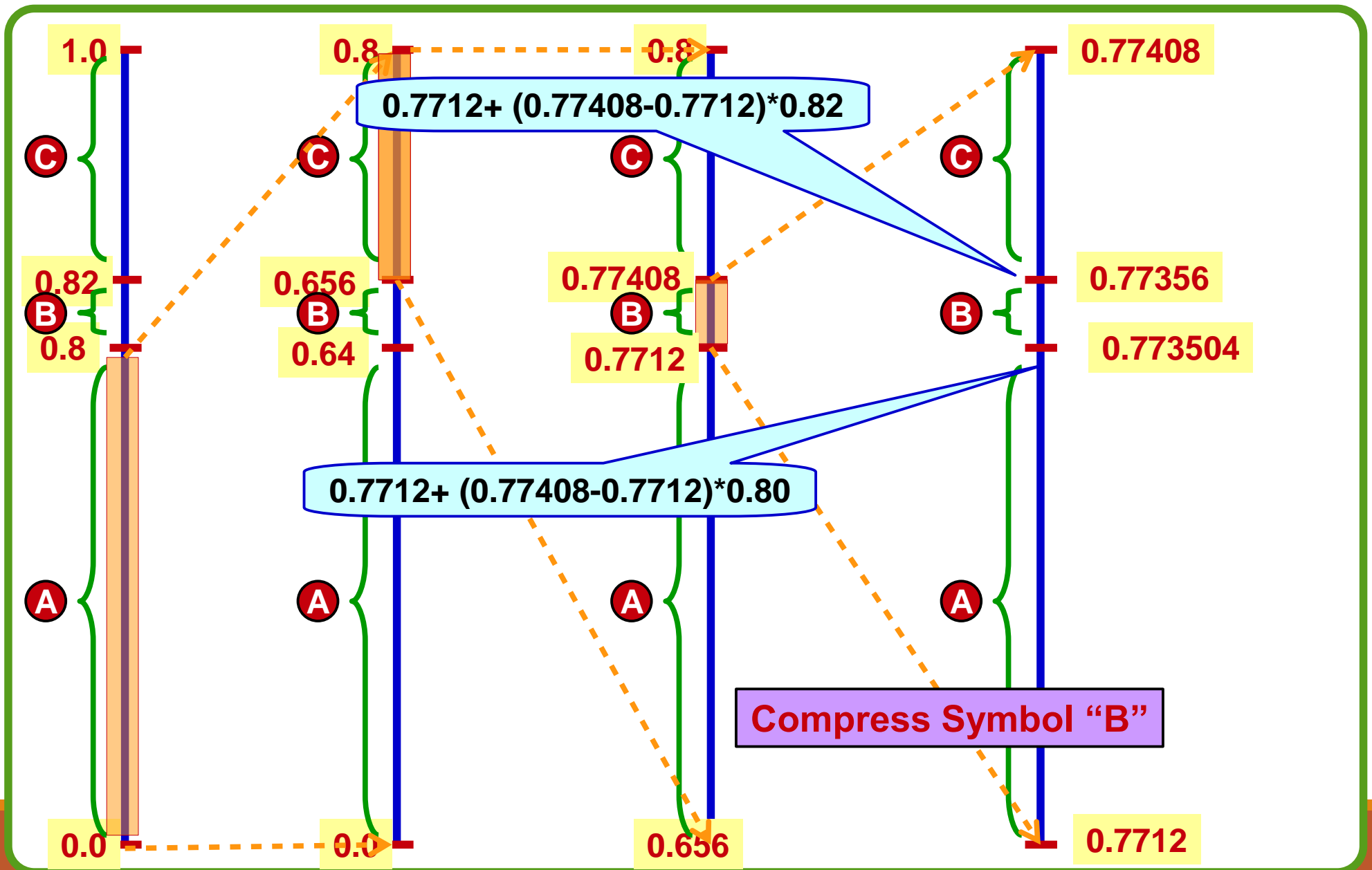
Example(1) Arithmetic Coding



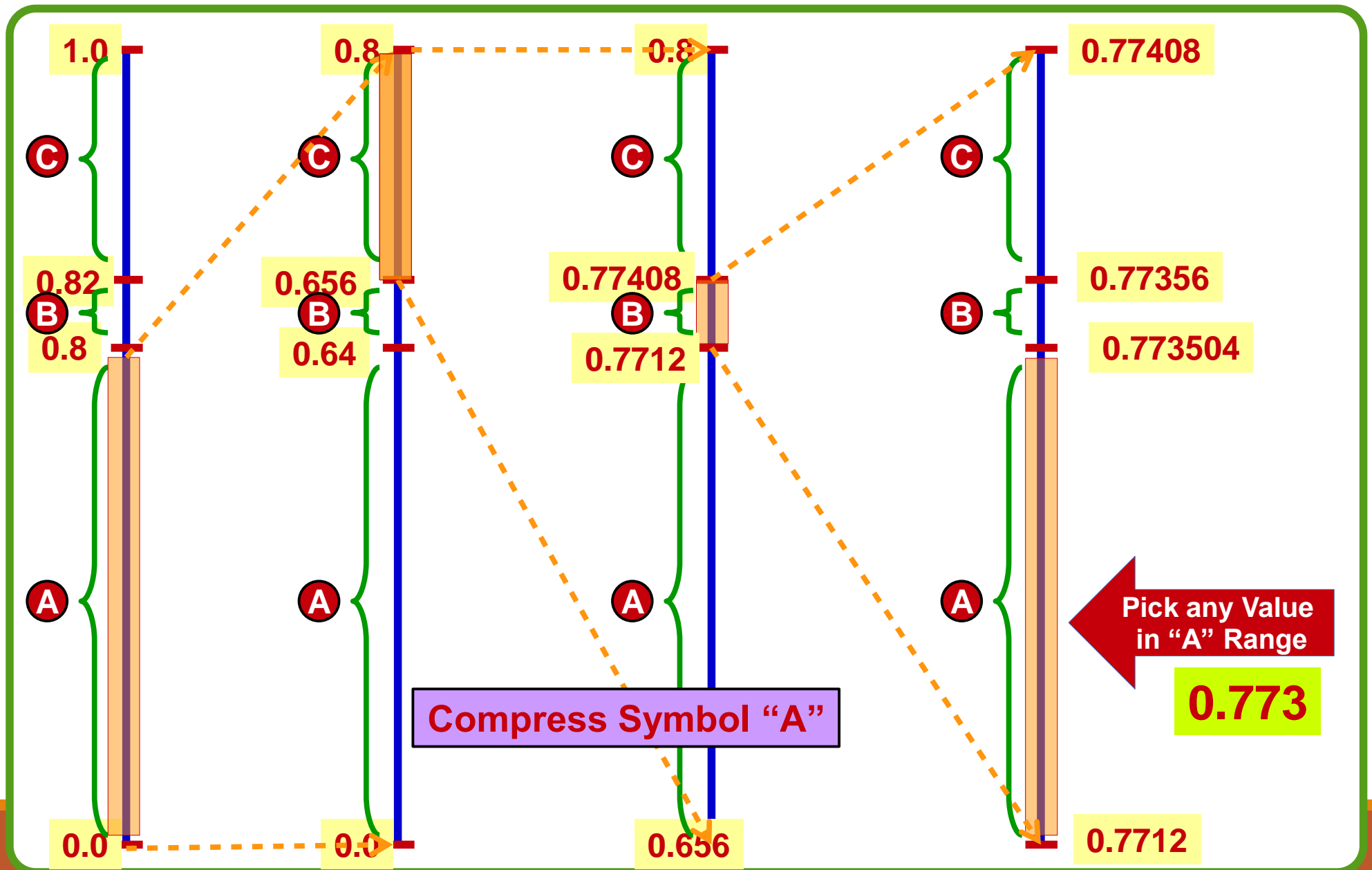
Example(1) Arithmetic Coding



Example(1) Arithmetic Coding



Example(1) Arithmetic Coding



Example(1) Arithmetic Coding

Lower (Symbol) = Lower+ Range * Low_Range(Symbol)
Upper (Symbol)= Lower+ Range * High_Range(Symbol)

Method 1: Using Equations

First Symbol is "A"

Lower(A)=0

Upper(A)=0.8

Second Symbol is "C"

Lower(C)=0+ (0.8-0)*0.82=0.656

Upper(C)=0+ (0.8- 0)*1= 0.8

Third Symbol is "B"

Lower (B)= 0.656+ (0.8-0.656)*0.8=0.7712

Upper(B)=0.656+(0.8-0.656)*0.82=0.77408

Fourth Symbol is "A"

Lower (A)= 0.7712+ (0.77408-0.7712)*0=0.7712

Upper(A)=0.7712+(0.77408-0.7712)*0.8=0.773504

Low_range(A)=0
High_Range(A)=0.8

Low_range(B)=0.8
High_Range(B)=0.82

Low_range(C)=0.82
High_Range(C)=1.0

Pick any Value
in "A" Range

0.773

Decoding Algorithm

Algorithm Arithmetic Decoding

/*

Input: code: binary code

Low_Range(Symbol):Lower Range of Symbol (in Accumulative Probabilities Scale)

High_Range(Symbol):Upper Range of Symbol (in Accumulative Probabilities Scale)

Output: The decoded message

*/

value = convert2decimal(code);

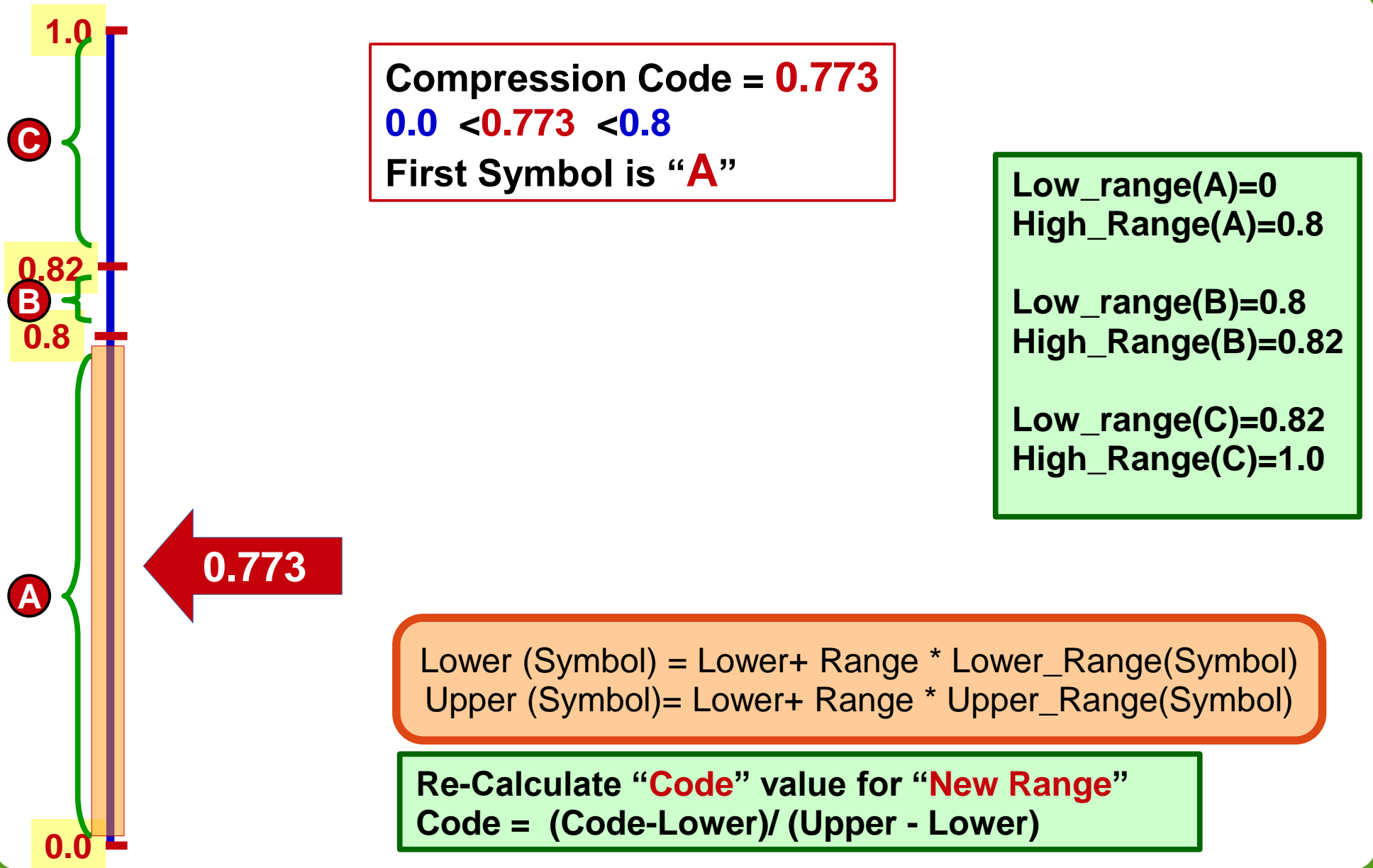
Do {

find a symbol S so that

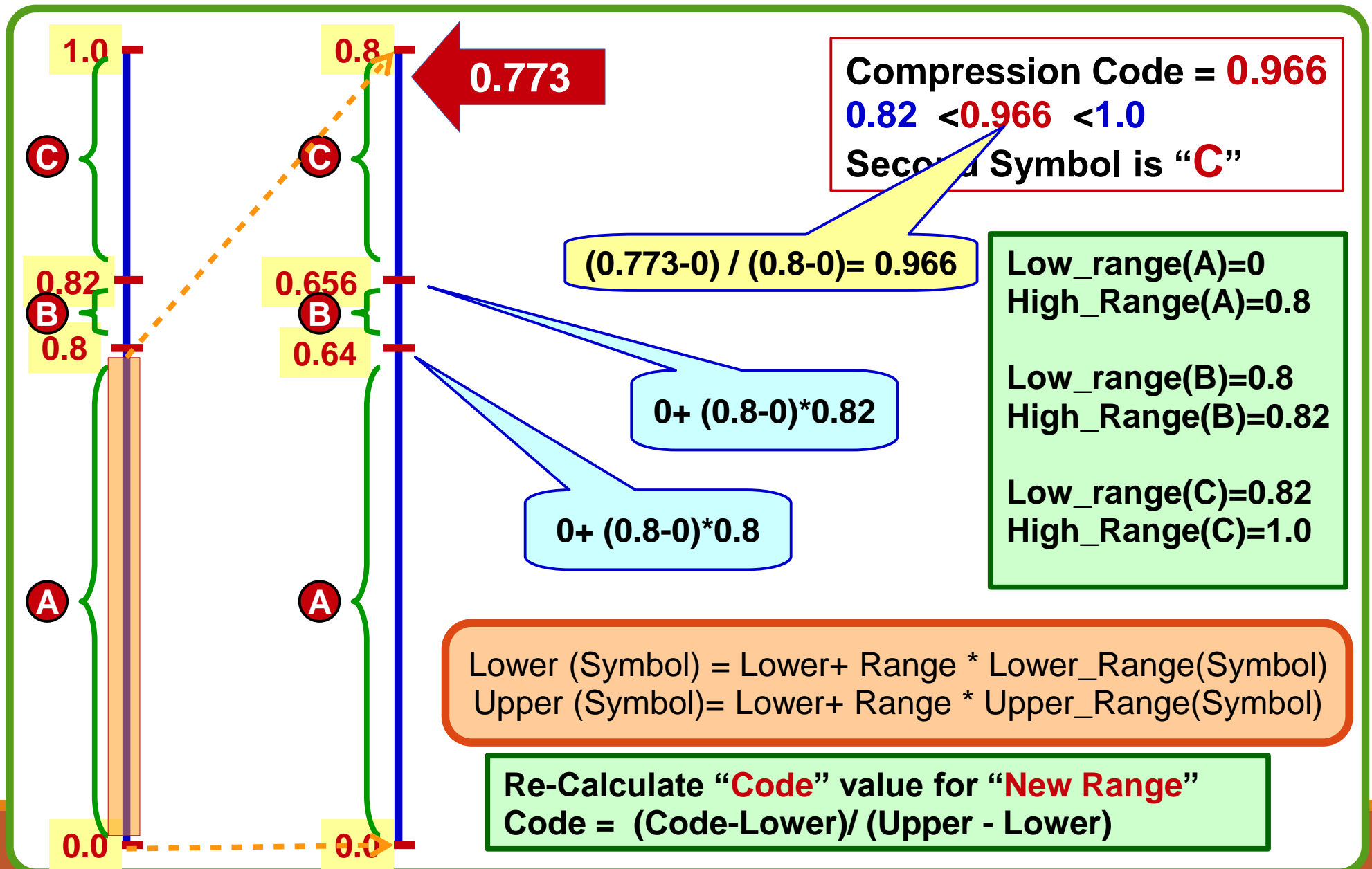
Example(1) Arithmetic Coding

Decoding

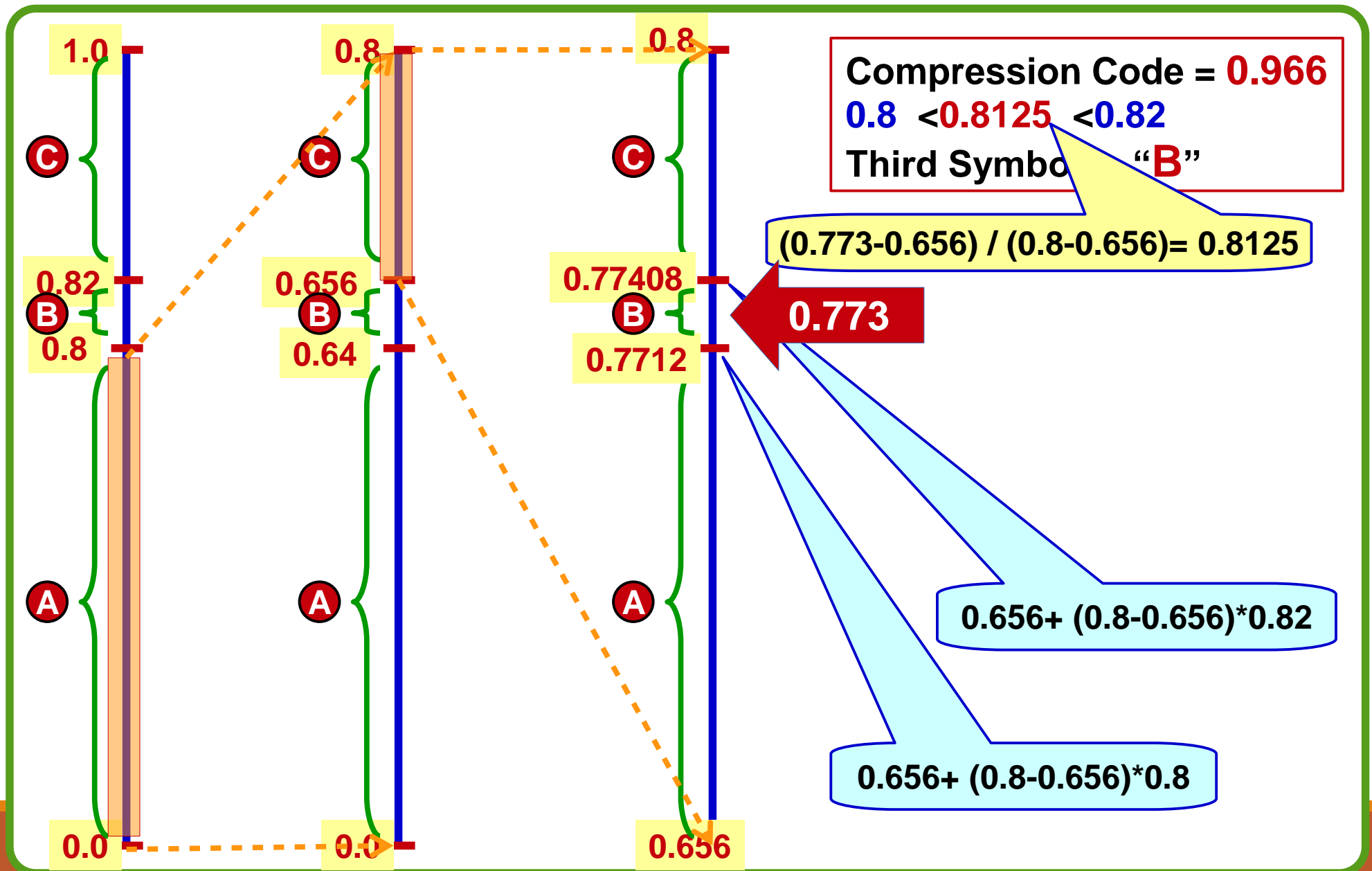
Method 1: Using Graph



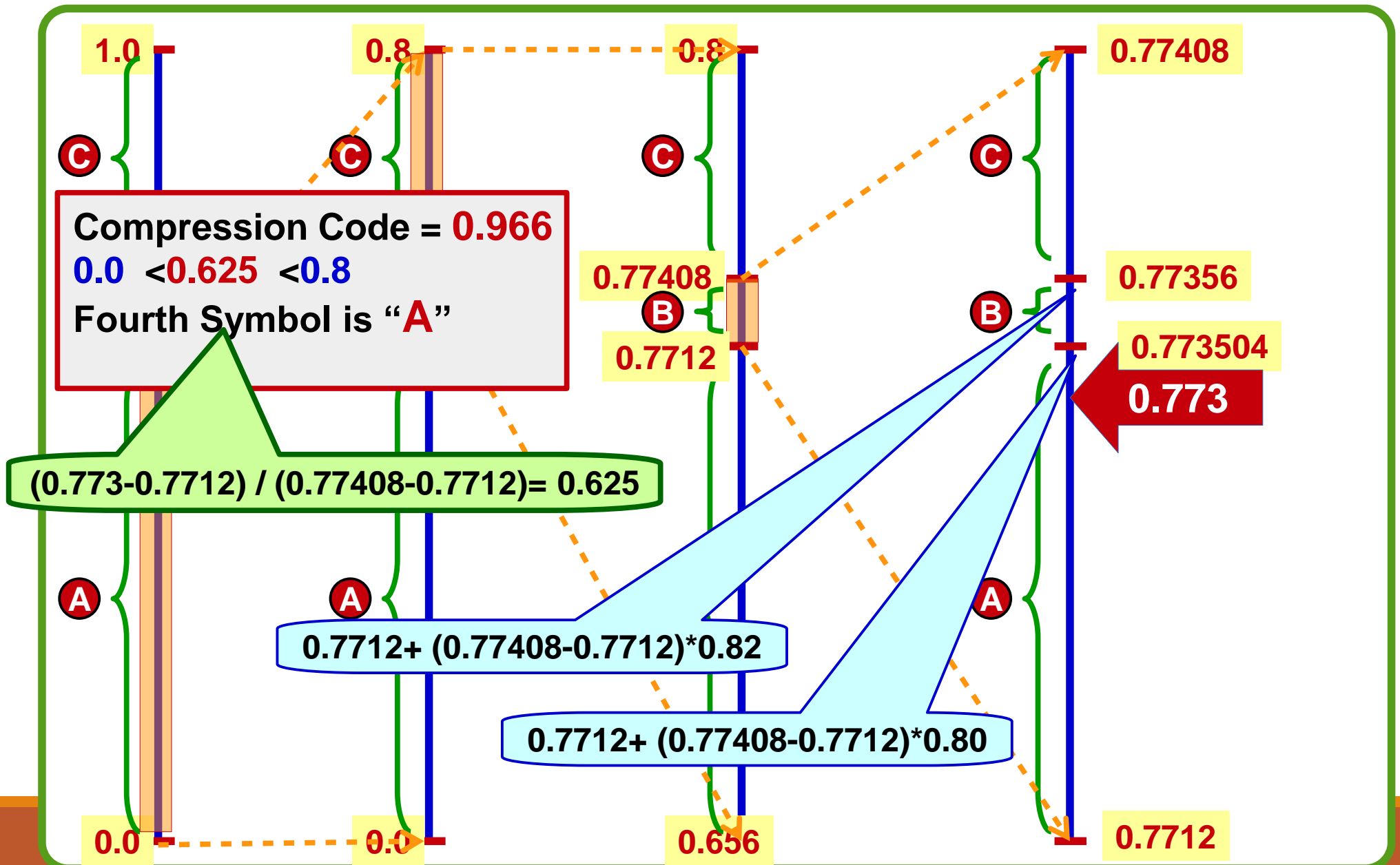
Example(1) Arithmetic Coding



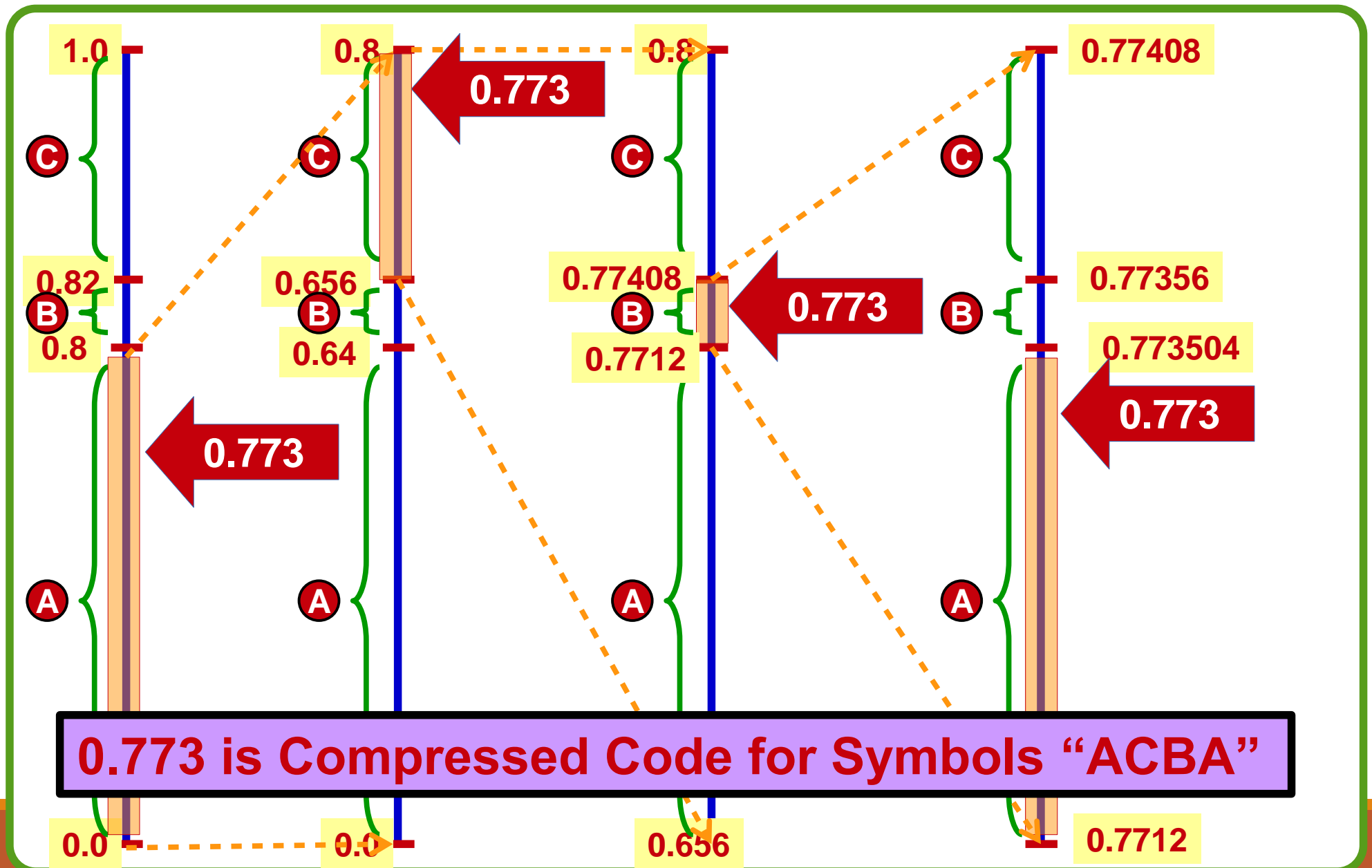
Example(1) Arithmetic Coding



Example(1) Arithmetic Coding



Example(1) Arithmetic Coding



Example(1) Arithmetic Coding (Decompression)

Method 1: Using
Equations

Compression Code = **0.773**

0.0 < **0.773** < **0.8**

First Symbol is “**A**”

Lower(A)=0

Upper(A)=0.8

Code=(0.773-0) / (0.8-0)= 0.966

0.82 < **0.966** < **1.0**

Second Symbol is “**C**”

Lower(C)=0+ (0.8-0)*0.82=0.656

Upper(C)=0+ (0.8- 0)*1= 0.8

$$\text{Code} = (0.733 - 0.656) / (0.8 - 0.656) = 0.8125$$

$$0.8 < 0.8125 < 0.82$$

Third Symbol is "B"

$$\text{Lower (B)} = 0.656 + (0.8 - 0.656) * 0.8 = 0.7712$$

$$\text{Upper(B)} = 0.656 + (0.8 - 0.656) * 0.82 = 0.77408$$

$$\text{Code} = (0.773 - 0.7712) / (0.77408 - 0.7712) = 0.625$$

$$0.0 < 0.625 < 0.8$$

Fourth Symbol is "A"