



DATABASE SYSTEMS

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Lecture 17

Normalization

DB Lifecycle

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- System Requirements
- DB Design
- Normalization "Schema Refinement"
- Relational Model

Problems without Normalization

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normalization is

a technique of organizing the data into multiple related tables, to minimize **DATA REDUNDANCY**.

What is
Data Redundancy?

and why should we **reduce** it?

Data Redundancy

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TABLE		
ROW 1		X
ROW 2		X
ROW 3		X
ROW 4		X

Repetition increases the size of the database

Other issues like:

- Insertion problems
- Deletion problems
- Update problems

Insertion Anomaly

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- To insert redundant data for every new row (of Student data in our case) is a data insertion problem.

STUDENTS TABLE				
rollno	name	branch	hod	office_tel
1	Akon	CSE	Mr. X	53337
2	Bkon	CSE	Mr. X	53337
3	Ckon	CSE	Mr. X	53337
4	Dkon	CSE	Mr. X	53337

Deletion Anomaly

STUDENTS TABLE

rollno	name	branch	hod	office_tel
1	Akon	CSE	Mr. X	53337
2	Bkon	CSE	Mr. X	53337
3	Ckon	CSE	Mr. X	53337
4	Dkon	CSE	Mr. X	53337

STUDENTS TABLE

rollno	name	branch	hod	office_tel
1	Akon	CSE	Mr. X	53337

- ❑ Delete some students
- ❑ We automatically delete the branch information
- ❑ We have not stored the branch information anywhere else

Update Anomaly

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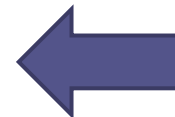
STUDENTS TABLE

rollno	name	branch	hod	office_tel
1	Akon	CSE	Mr. X	53337
2	Bkon	CSE	Mr. X	53337
3	Ckon	CSE	Mr. X	53337
4	Dkon	CSE	Mr. X	53337

Mr. X leaves, and Mr. Y joins
as the new HOD for CSE

STUDENTS TABLE

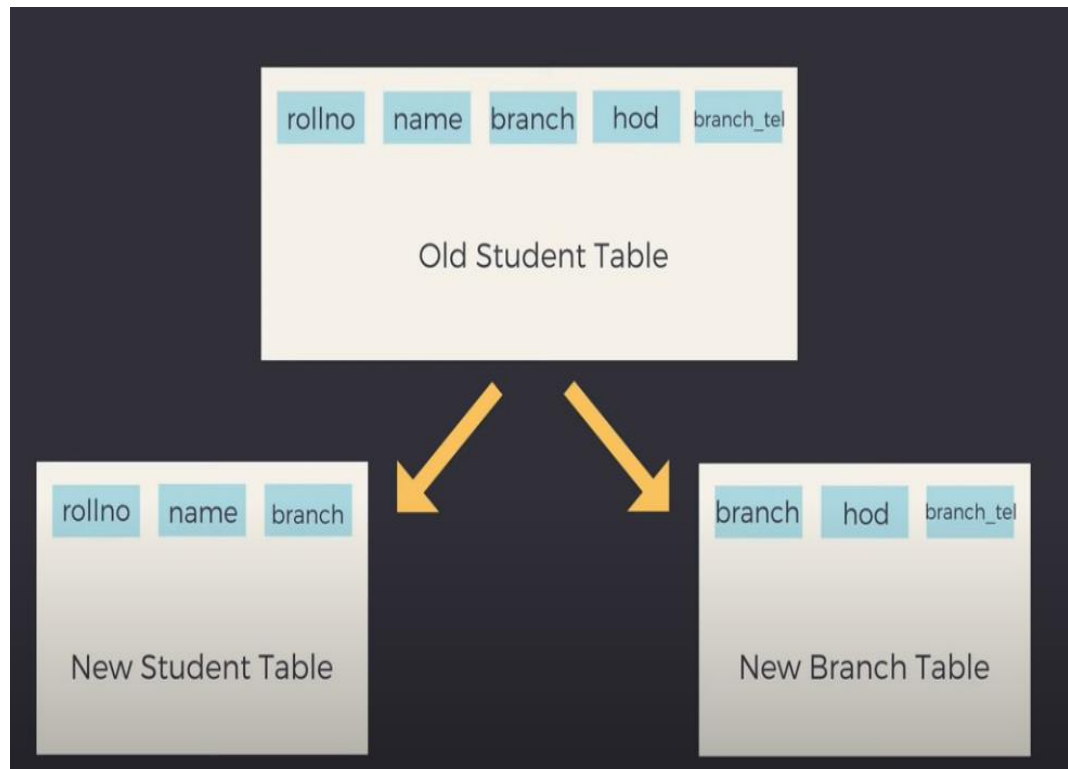
rollno	name	branch	hod	office_tel
1	Akon	CSE	Mr. X Mr. Y	53337
2	Bkon	CSE	Mr. X Mr. Y	53337
3	Ckon	CSE	Mr. X	53337
4	Dkon	CSE	Mr. X Mr. Y	53337



How to solve These Problems?

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□ Use Normalization



STUDENTS TABLE

rollno	name	branch
1	Akon	CSE
2	Bkon	CSE

BRANCH TABLE

branch	hod	office_tel
CSE	Mr. Y	53337

- Normalization not eliminating data redundancy but minimizing it

Database Tables and Normalization

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- **Normalization** is a process for assigning attributes to entities. It reduces data redundancies and helps eliminate the data anomalies.
- Normalization works through a series of stages called normal forms:
 - ▣ **First normal form (1NF)**
 - ▣ **Second normal form (2NF)**
 - ▣ **Third normal form (3NF)**
 - ▣ **Boyce Codd Normal Form (BCNF)**
 - ▣ **Fourth normal form (4NF)**
- The highest level of normalization is not always desirable.

First Normal Form

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1. Each table cell contain atomic value(indivisible)
 2. Each table has a primary key that uniquely identifies records
- Primary key
 - ❑ cannot be NULL
 - ❑ must be unique
 - ❑ The primary key values should rarely be changed



Primary Key

Functional dependency

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Functional Dependencies

- A *Functional Dependency* describes a relationship between *attributes* in a single relation.
- An attribute is *functionally dependant* on another if we can use the value of one attribute to determine the value of another.

Functional dependency

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- We use the symbol ' \rightarrow ' to indicate a functional dependency:

$$X \rightarrow Y$$

- is read "X functionally determines Y" or "Y functionally depends on X"
- more simply "X determines Y"
- or "Y depends on X"

Functional dependency

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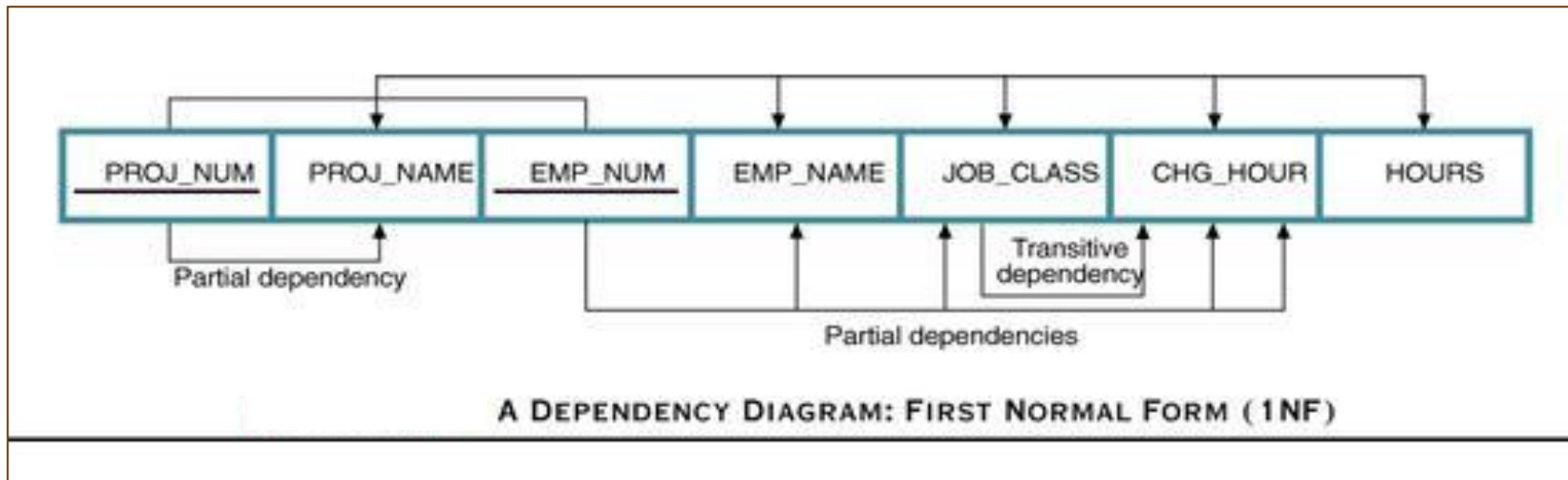
Some Examples

- Student_ID \rightarrow Saddress, SDoB
- Student_ID, Course# \rightarrow Grade
- Model, Year \rightarrow CarPrice
- Course_No, Section \rightarrow Professor, Classroom, Number of Student

Dependency Diagram

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- Dependency Diagram
 - ▣ The primary key components are bold, underlined, and shaded in a different color.
 - ▣ The arrows above entities indicate all desirable dependencies, i.e., dependencies that are based on PK.
 - ▣ The arrows below the dependency diagram indicate less desirable dependencies – **partial dependencies** and **transitive dependencies**.



2NF

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□ 2NF Definition

▣ A table is in 2NF if:

- It is in 1NF and

- It includes no partial dependencies; that is, no attribute is dependent on only a portion of the primary key.

(It is still possible for a table in 2NF to exhibit **transitive dependency**; that is, one or more attributes may be functionally dependent on nonkey attributes.)

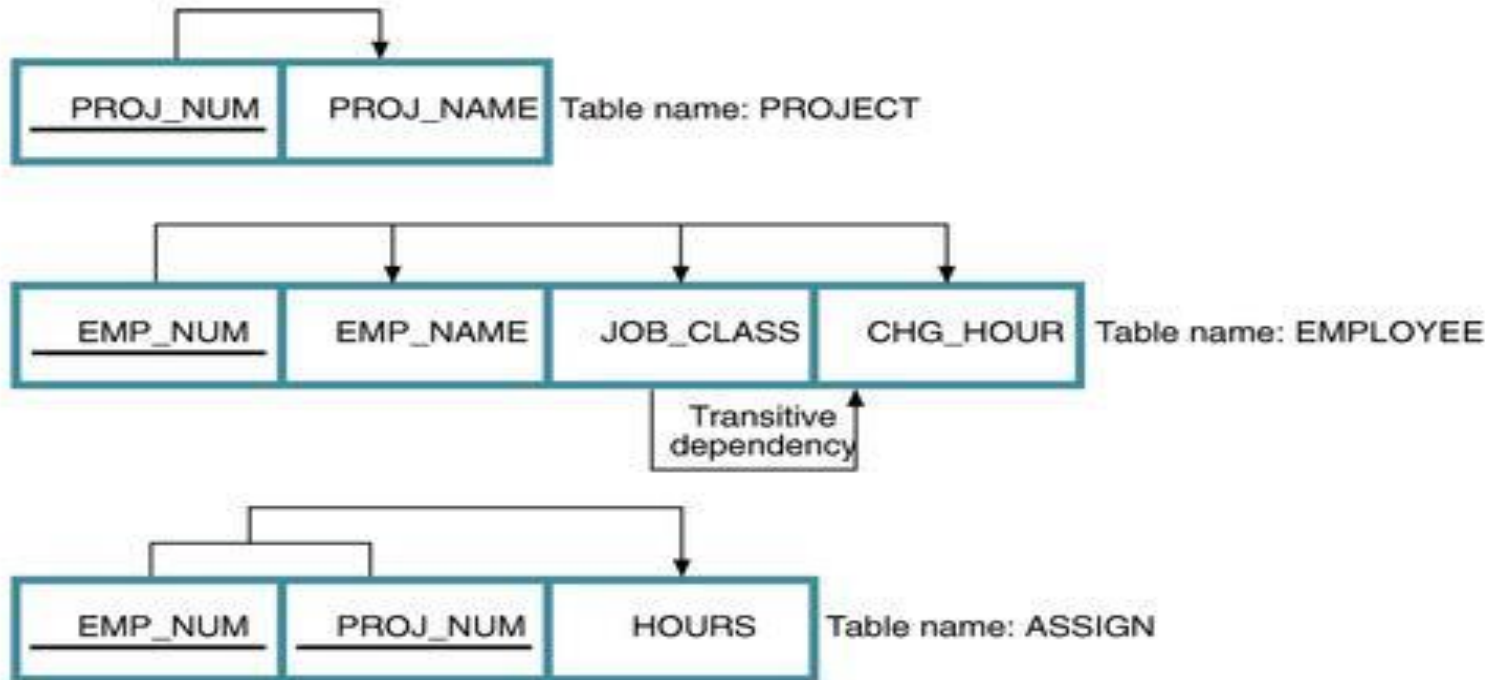
2NF

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□ Conversion to Second Normal Form

- ▣ Starting with the 1NF format, the database can be converted into the 2NF format by
 - Writing each key component on a separate line, and then writing the original key on the last line and
 - Writing the dependent attributes after each new key.

Second Normal Form (2NF) Conversion Results



SECOND NORMAL FORM (2NF) CONVERSION RESULTS

3NF

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□ 3NF Definition

▣ **A table is in 3NF if:**

■ **It is in 2NF and**

■ **It contains no transitive dependencies.**

3NF

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- Conversion to Third Normal Form
 - ▣ **Create a separate table with attributes in a transitive functional dependence relationship.**

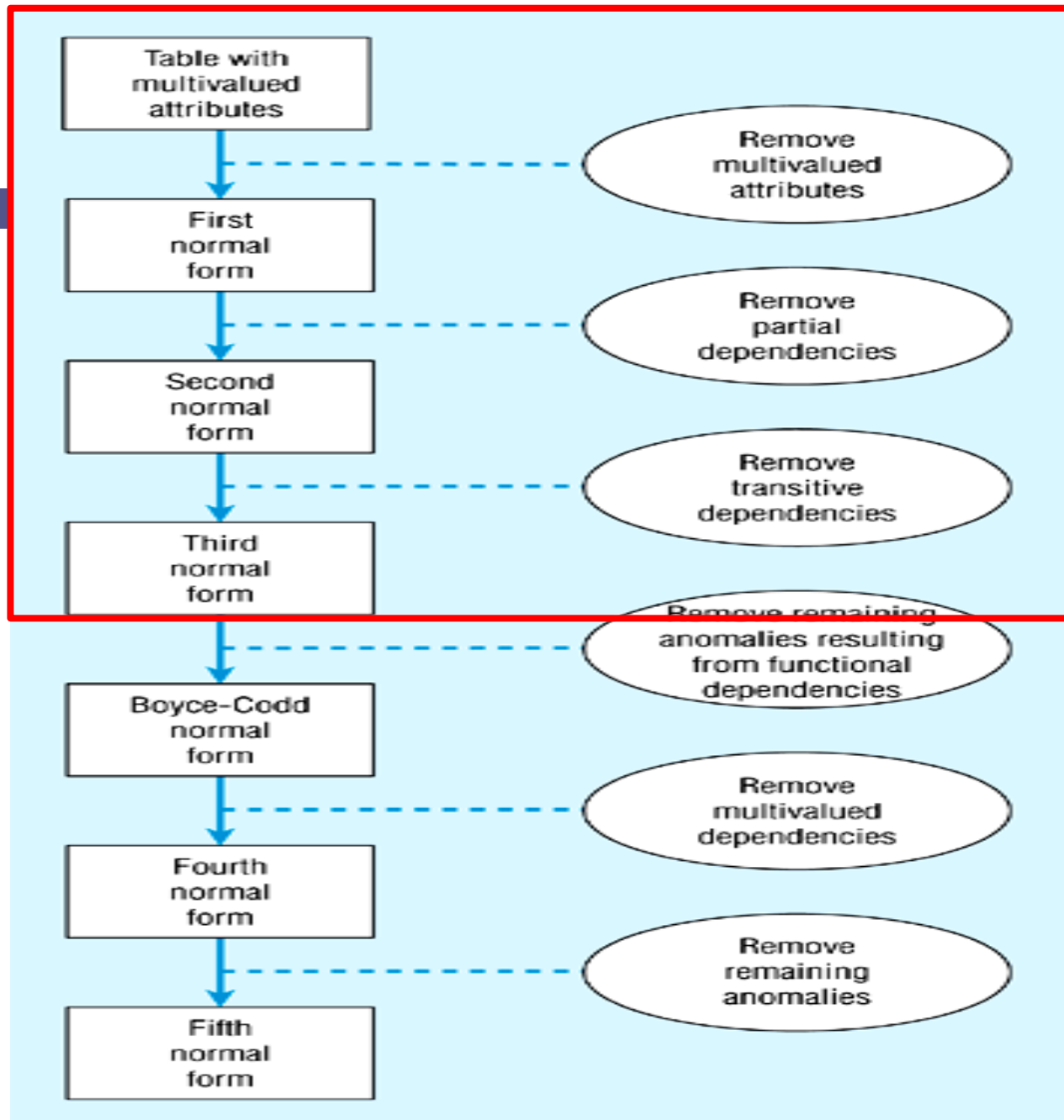
PROJECT (PROJ_NUM, PROJ_NAME)

ASSIGN (PROJ_NUM, EMP_NUM, HOURS)

EMPLOYEE (EMP_NUM, EMP_NAME, JOB_CLASS)

JOB (JOB_CLASS, CHG_HOUR)

Normalization Normal Forms



Normalization Examples

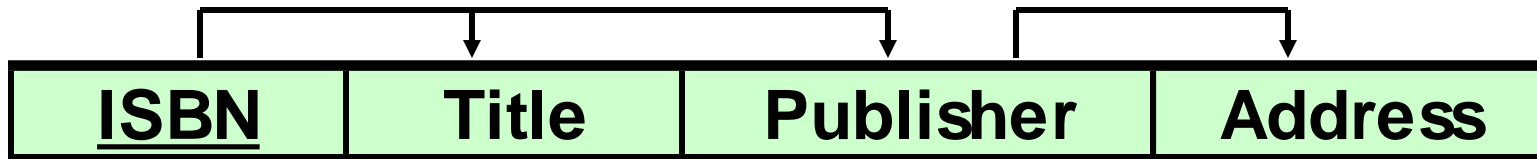
Example 1: Determine NF

23

- ISBN \rightarrow Title
- ISBN \rightarrow Publisher
- Publisher \rightarrow Address

All attributes are directly or indirectly determined by the primary key; therefore, the relation is at least in 1 NF

BOOK



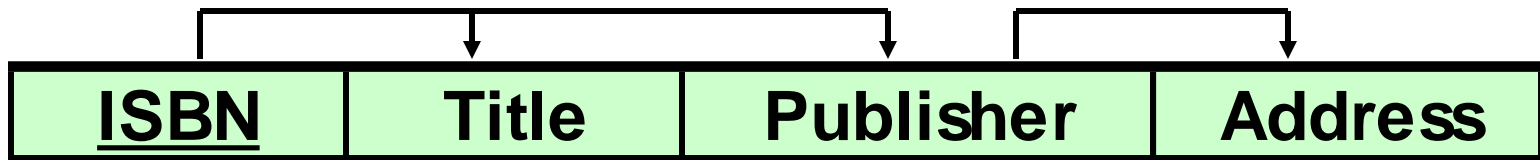
Example 1: Determine NF

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- ISBN \rightarrow Title
- ISBN \rightarrow Publisher
- Publisher \rightarrow Address

The relation is at least in 1NF.
There is no **COMPOSITE**
primary key, therefore there
can't be partial dependencies.
Therefore, the relation is at
least in 2NF

BOOK



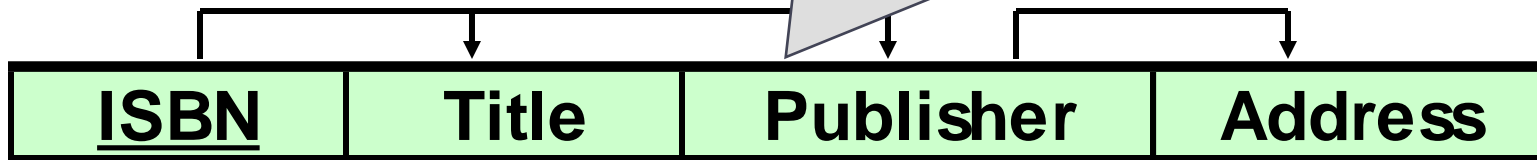
Example 1: Determine NF

25

- ISBN \rightarrow Title
- ISBN \rightarrow Publisher
- Publisher \rightarrow Address

Publisher is a non-key attribute, and it determines Address, another non-key attribute. Therefore, there is a transitive dependency, which means that the relation is **NOT in 3 NF**.

BOOK



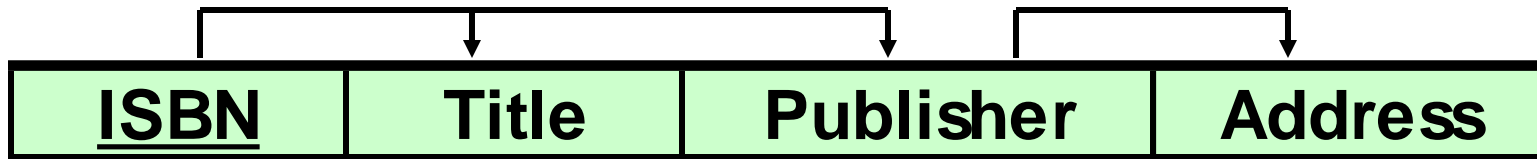
Example 1: Determine NF

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- ISBN \rightarrow Title
- ISBN \rightarrow Publisher
- Publisher \rightarrow Address

We know that the relation is at least in 2NF, and it is not in 3NF. Therefore, we conclude that the relation **is in 2NF**.

BOOK



Example 2: Determine NF

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- $\text{Product_ID} \rightarrow \text{Description}$

ORDER

All attributes are directly or indirectly determined by the primary key; therefore, the relation is at least in 1 NF

<u>Order No</u>	<u>Product ID</u>	Description
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Example 2: Determine NF

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□ Product_ID → Description

The relation is at least in 1NF.
There is a COMPOSITE Primary Key (PK) (Order_No,
Product_ID), therefore there can be partial
dependencies. Product_ID, which is a part of PK,
determines Description; hence, there is a partial
dependency. Therefore, **the relation is not 2NF**. No
sense to check for transitive dependencies!

ORDER



<u>Order No</u>	<u>Product ID</u>	Description
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Example 2: Determine NF

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- Product_ID → Description

We know that the relation is at least in 1NF, and it is not in 2 NF.
Therefore, we conclude that the **relation is in 1 NF.**

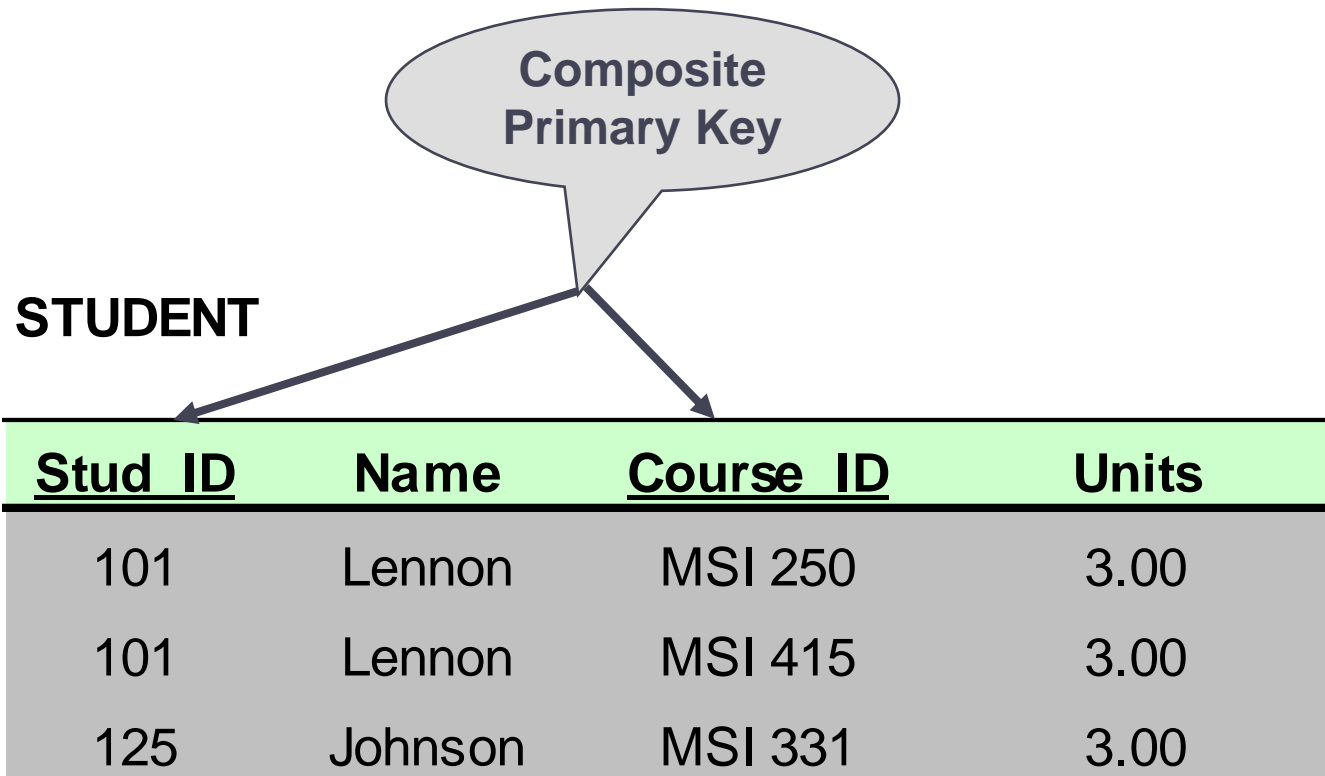
ORDER



<u>Order No</u>	<u>Product ID</u>	Description
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Bringing a Relation to 2NF

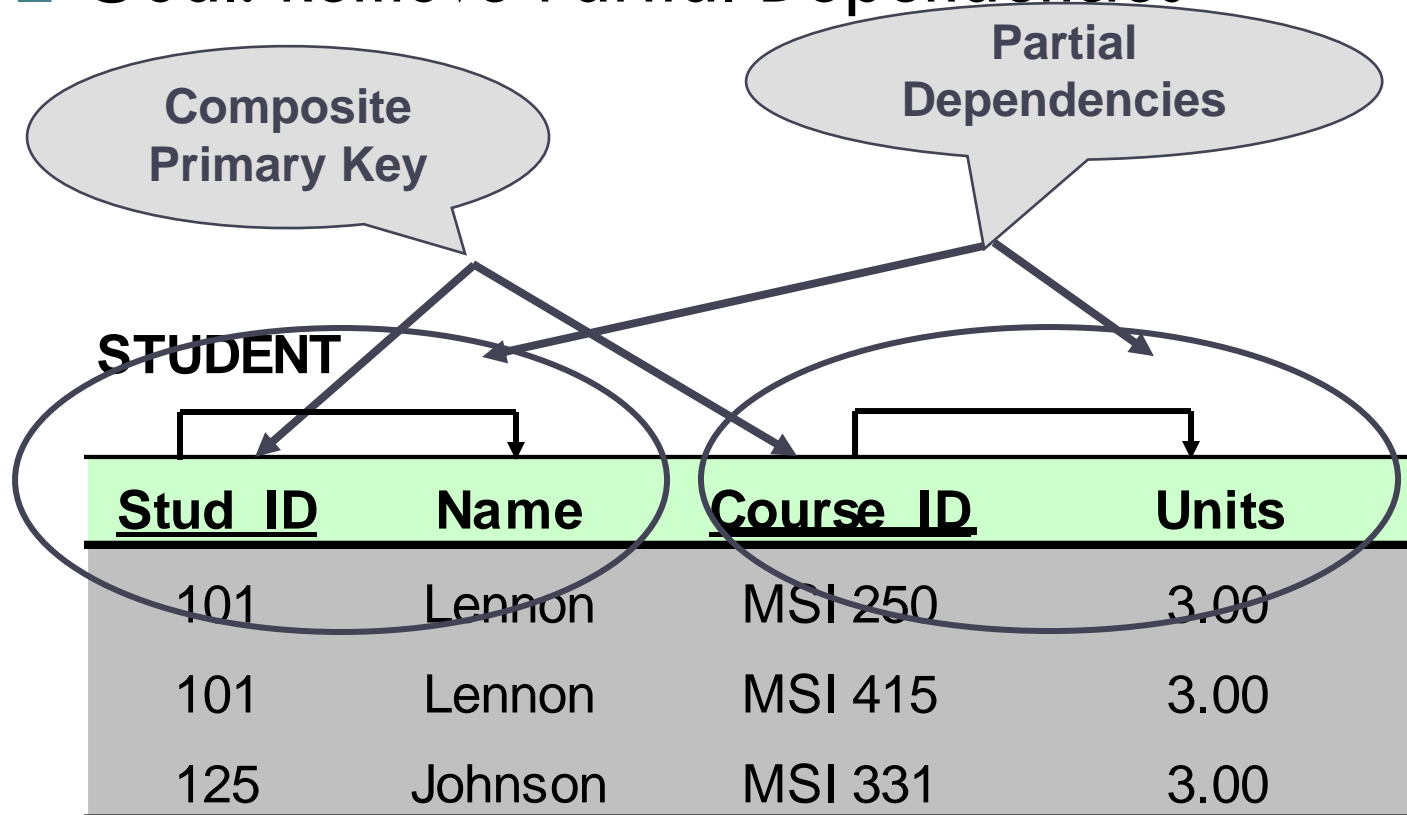
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Bringing a Relation to 2NF

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- Goal: Remove Partial Dependencies



Bringing a Relation to 2NF

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- Remove attributes that are dependent from the part but not the whole of the primary key from the original relation. For each partial dependency, create a new relation, with the corresponding part of the primary key from the original as the primary key.

STUDENT

<u>Stud ID</u>	Name	<u>Course ID</u>	Units
101	Lennon	MSI 250	3.00
101	Lennon	MSI 415	3.00
125	Johnson	MSI 331	3.00

Bringing a Relation to 2NF

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CUSTOMER

<u>Stud ID</u>	Name	<u>Course ID</u>	Units
101	Lennon	MSI 250	3.00
101	Lennon	MSI 415	3.00
125	Johnson	MSI 331	3.00

STUDENT_COURSE

<u>Stud ID</u>	<u>Course ID</u>
101	MSI 250
101	MSI 415
125	MSI 331

STUDENT

<u>Stud ID</u>	Name
101	Lennon
101	Lennon
125	Johnson

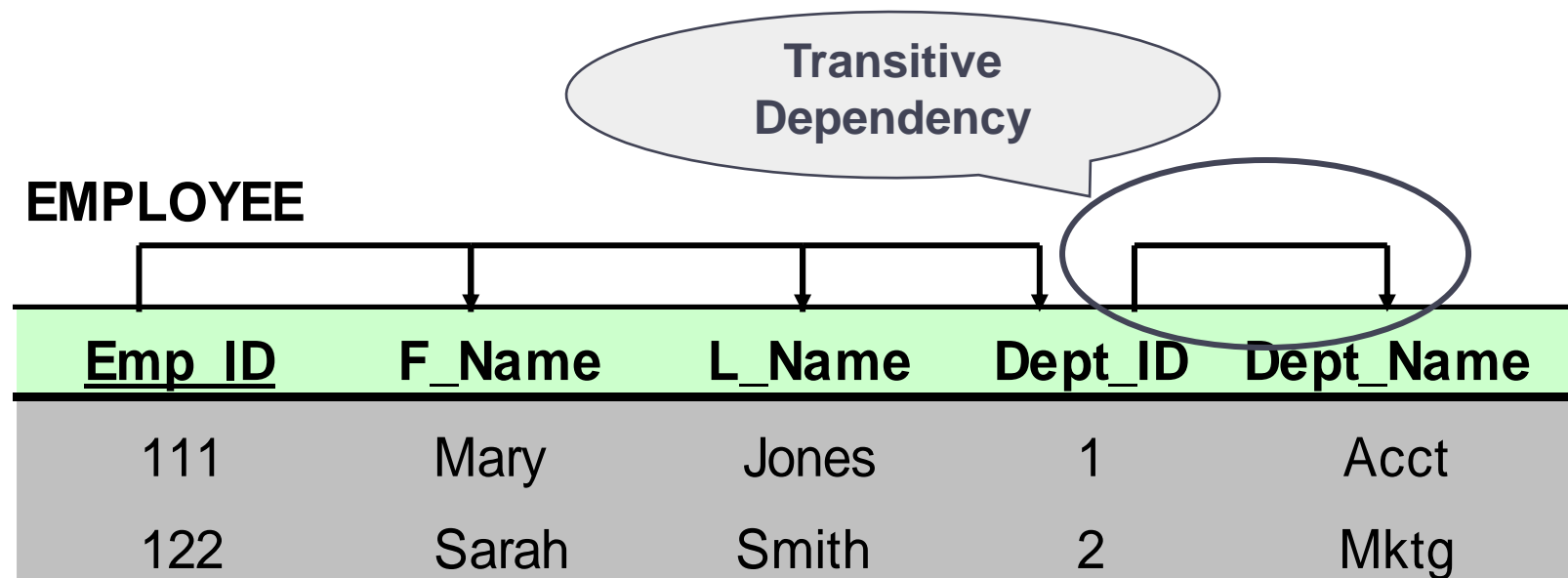
COURSE

<u>Course ID</u>	Units
MSI 250	3.00
MSI 415	3.00
MSI 331	3.00

Bringing a Relation to 3NF

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- Goal: Get rid of transitive dependencies.



Bringing a Relation to 3NF

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- Remove the attributes, which are dependent on a non-key attribute, from the original relation. For each transitive dependency, create a new relation with the non-key attribute which is a determinant in the transitive dependency as a primary key, and the dependent non-key attribute as a dependent.

EMPLOYEE

<u>Emp_ID</u>	F_Name	L_Name	Dept_ID	Dept_Name
111	Mary	Jones	1	Acct
122	Sarah	Smith	2	Mktg

Bringing a Relation to 3NF

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EMPLOYEE

<u>Emp_ID</u>	F_Name	L_Name	Dept_ID	Dept_Name
111	Mary	Jones	1	Acct
122	Sarah	Smith	2	Mktg

EMPLOYEE

<u>Emp_ID</u>	F_Name	L_Name	Dept_ID
111	Mary	Jones	1
122	Sarah	Smith	2

DEPARTMENT

<u>Dept_ID</u>	Dept_Name
1	Acct
2	Mktg