Linux Binary Exploitation

Return-oriented Programing angelboy@chroot.org

Outline

- ROP
- Using ROP bypass ASLR
- Stack migration

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- Stack migration

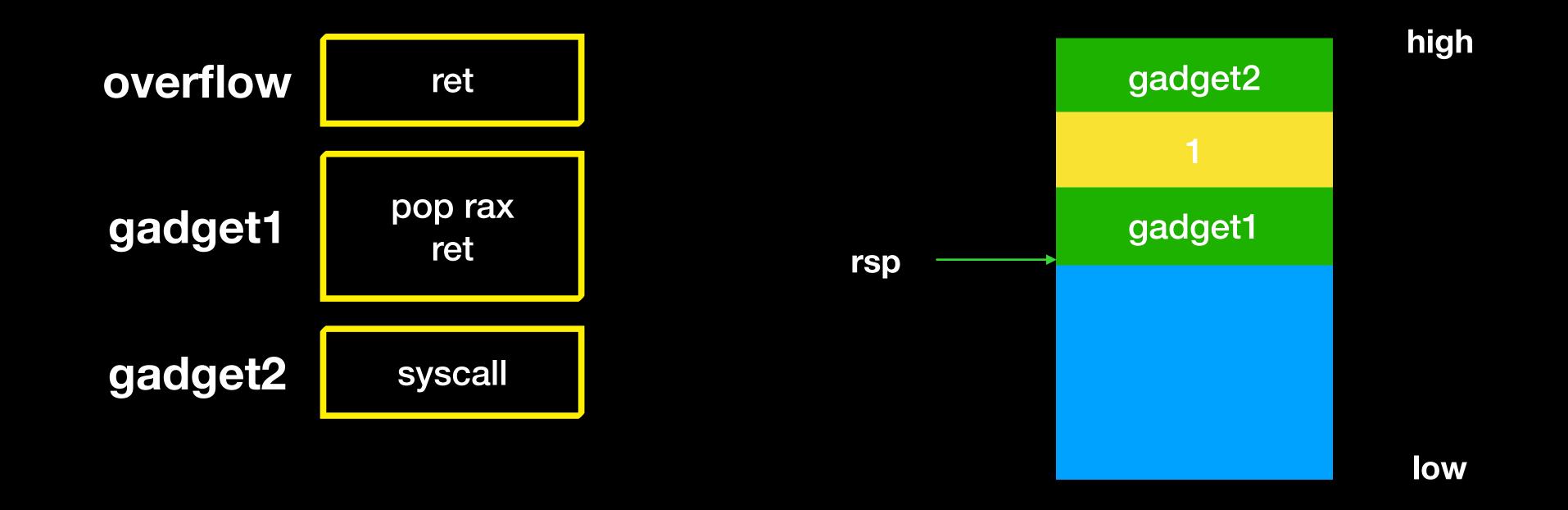
• 透過 ret 去執行其他包含 ret 的程式碼片段

• 這些片段又稱為 gadget

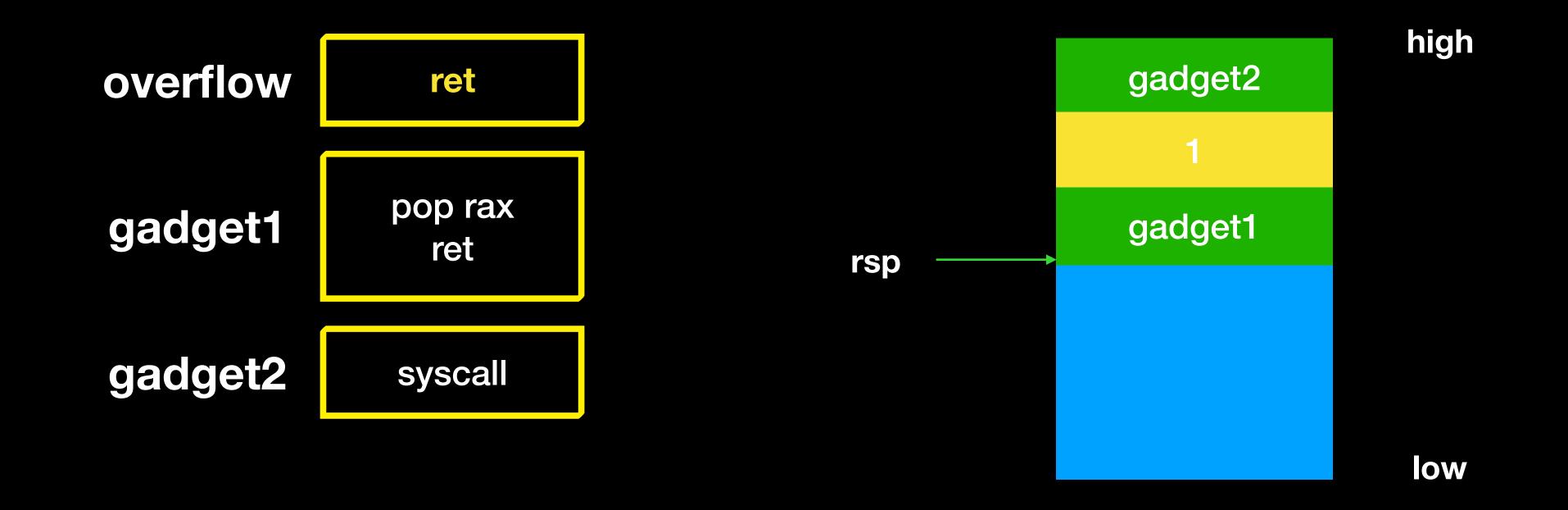
```
4027ba:
              5b
                                              rbx
                                       pop
              5d
4027bb:
                                              rbp
                                       pop
4027bc:
              41 5c
                                              r12
                                       pop
4027be:
              41 5d
                                              r13
                                       pop
4027c0:
              41 5e
                                              r14
                                       pop
4027c2:
              41 5f
                                              r15
                                       pop
4027c4:
              c3
                                       ret
                                               rdi,0x600
40274c:
              48 81 c7 00 06 00 00
                                        add
402753:
              48 89 f8
                                               rax,rdi
                                        mov
402756:
              c3
                                        ret
4026c5:
             5a
                                          rdx
                                   pop
4026c6:
             58
                                   pop
                                          rax
4026c7:
             48 83 c4 08
                                   add
                                          rsp,0x8
4026cb:
             5d
                                   pop
                                          rbp
4026cc:
             с3
                                   ret
```

- Why do we need ROP?
 - Bypass DEP
 - Static linking can do more thing

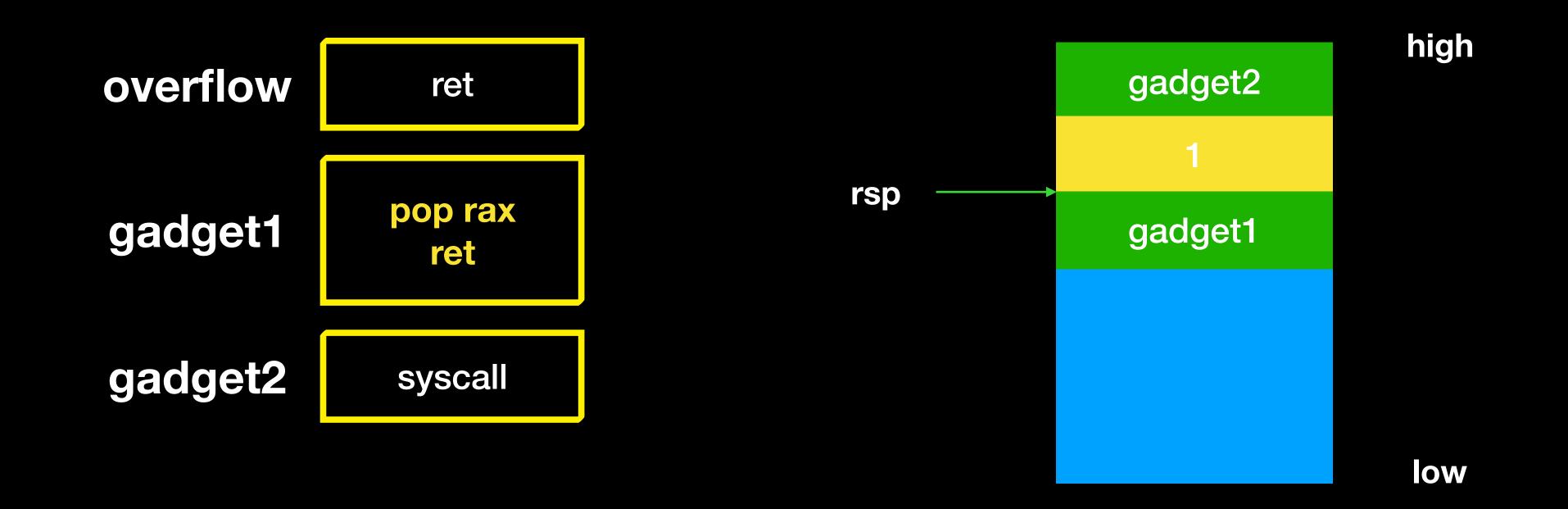
- ROP chain
 - 在夠長的 buffer overflow 後 stack 內容幾乎都由我們控制,我們可以藉由 ret = pop rip 指令,持續的控制 rip



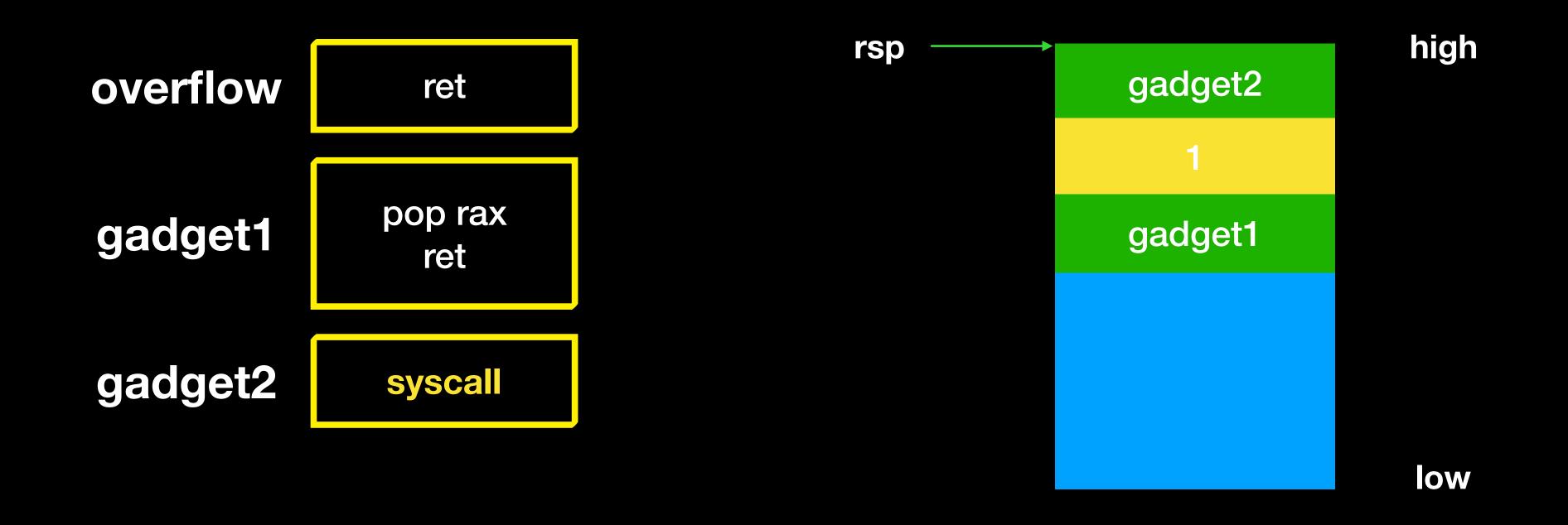
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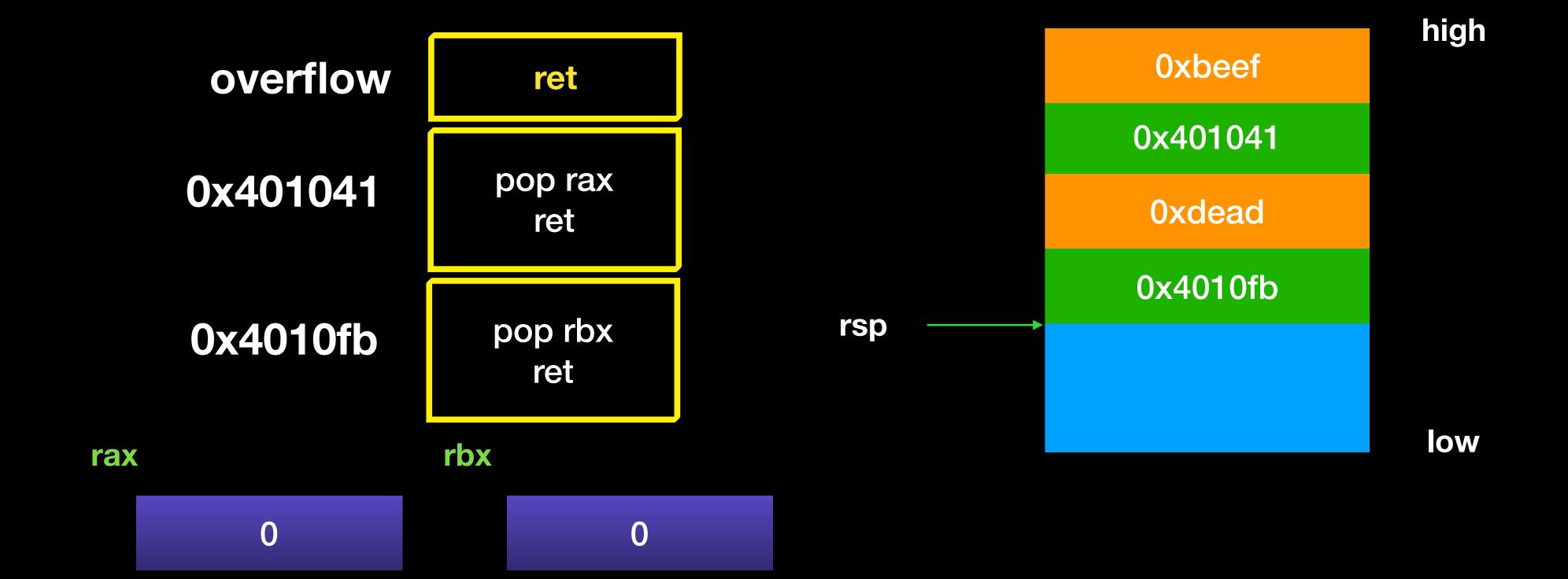
exit

- ROP chain
 - 由眾多的 ROP gadget 組成
 - 藉由不同的 register 及記憶體操作,呼叫 system call 達成任意代碼執行
 - 基本上就是利用 ROP gadget 來串出我們之前寫的 shellcode 的效果

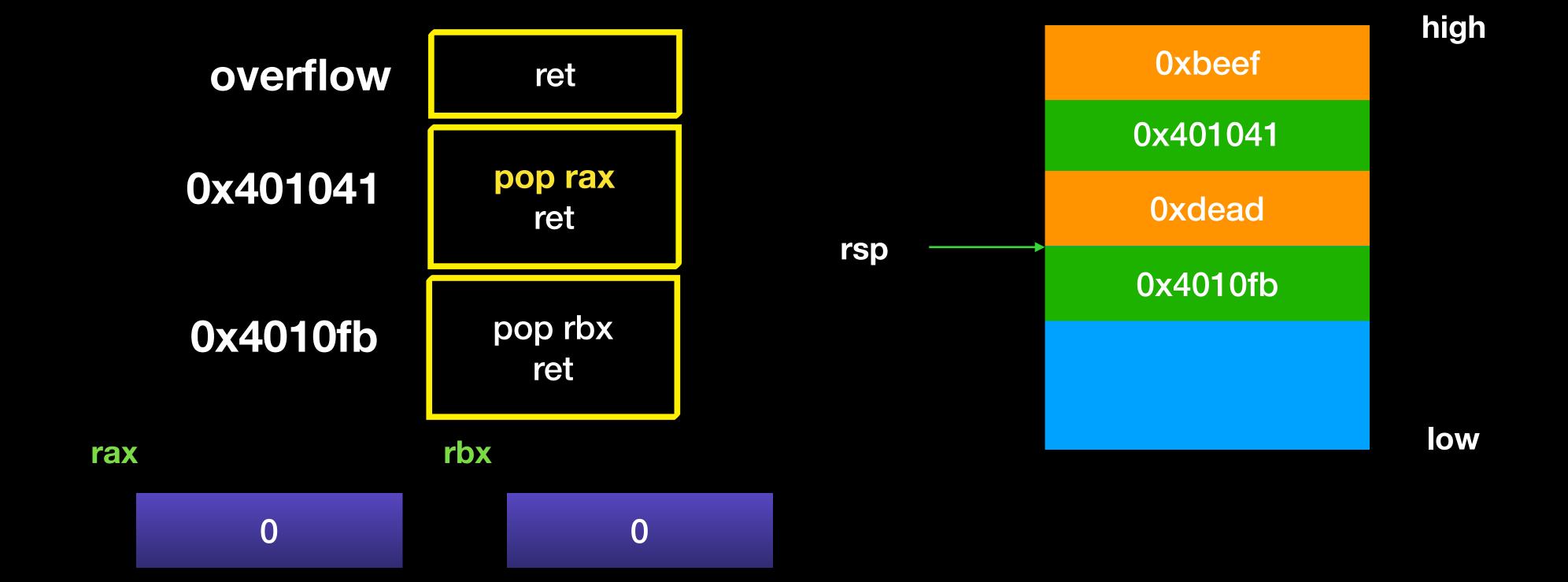
- Gadget
 - read/write register/memory
 - pop rax;pop rcx; ret
 - mov [rax],rcx; ret
 - system call
 - syscall
 - change rsp
 - pop rsp; ret
 - leave ; ret

- Write to Register
 - pop reg; ret
 - mov reg, reg; ret
 - •

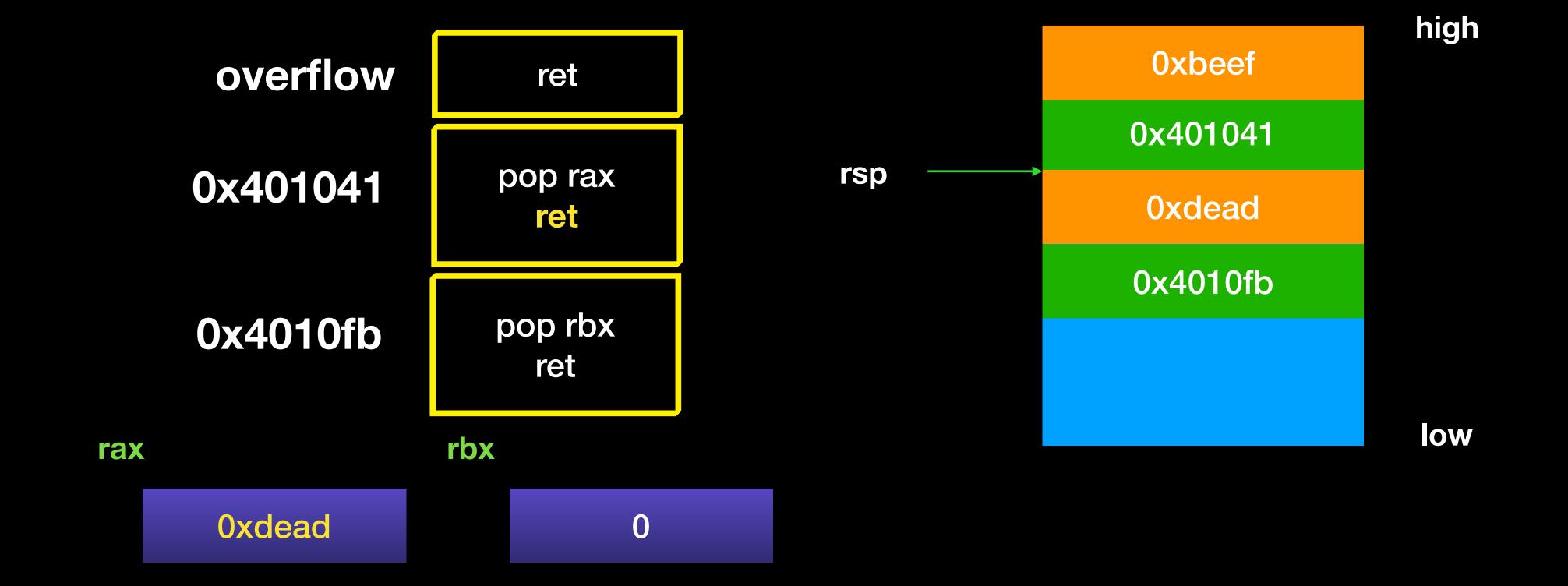
- Write to Register
 - let rax = 0xdead rbx = 0xbeef



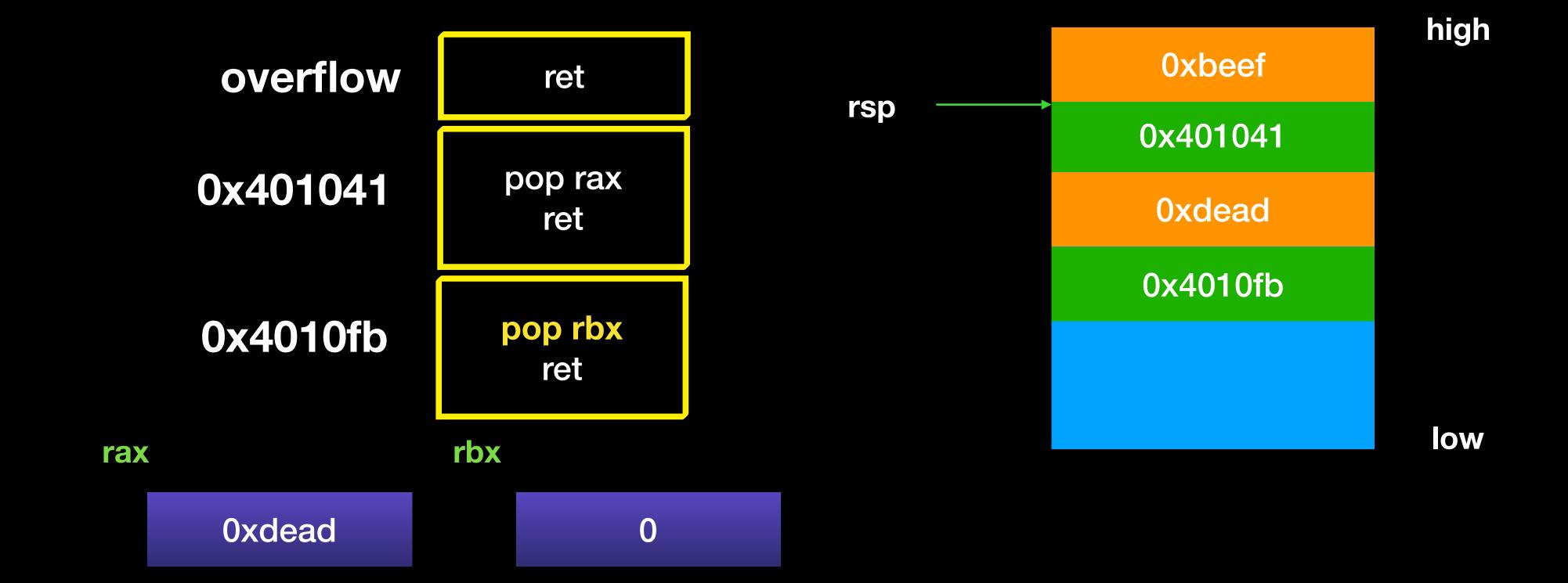
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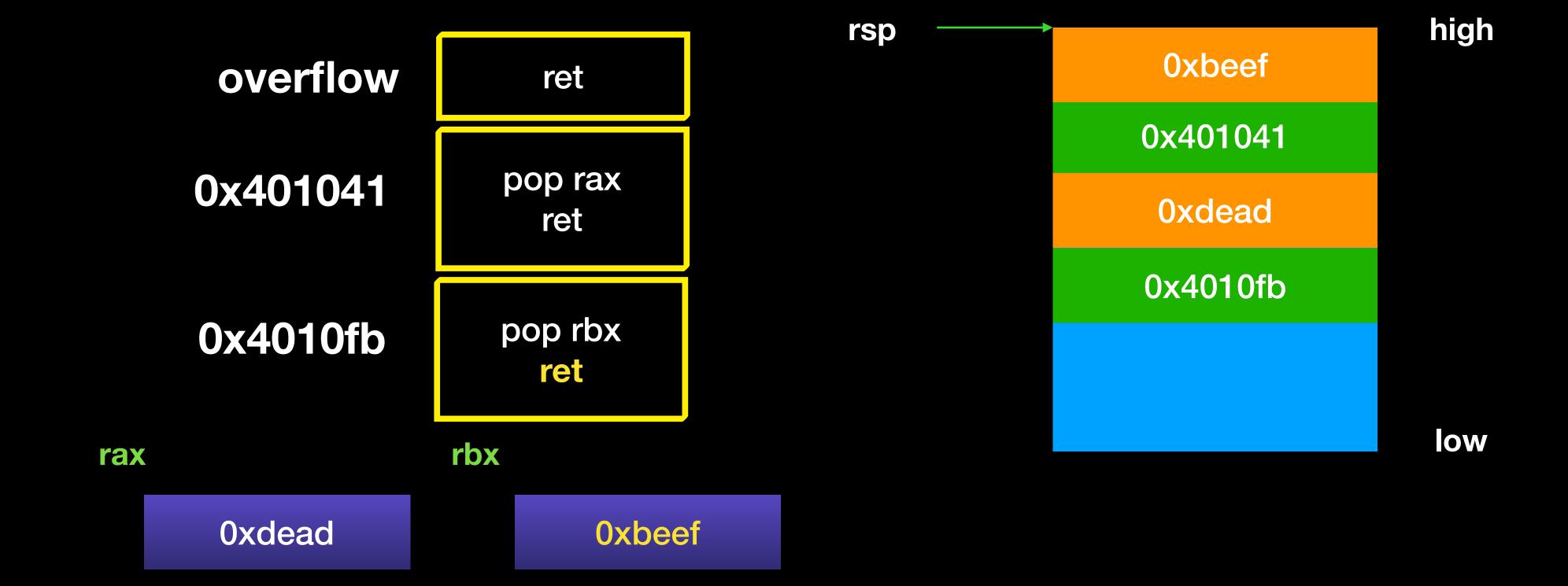
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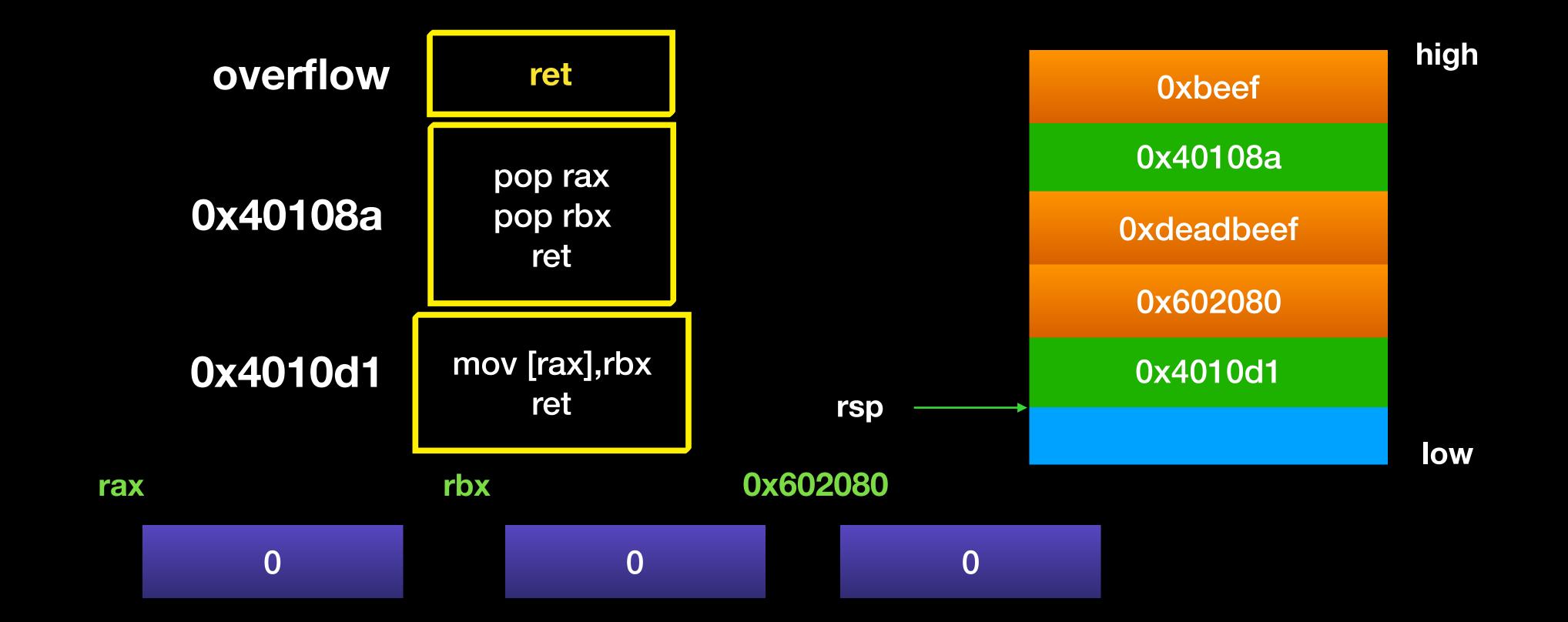


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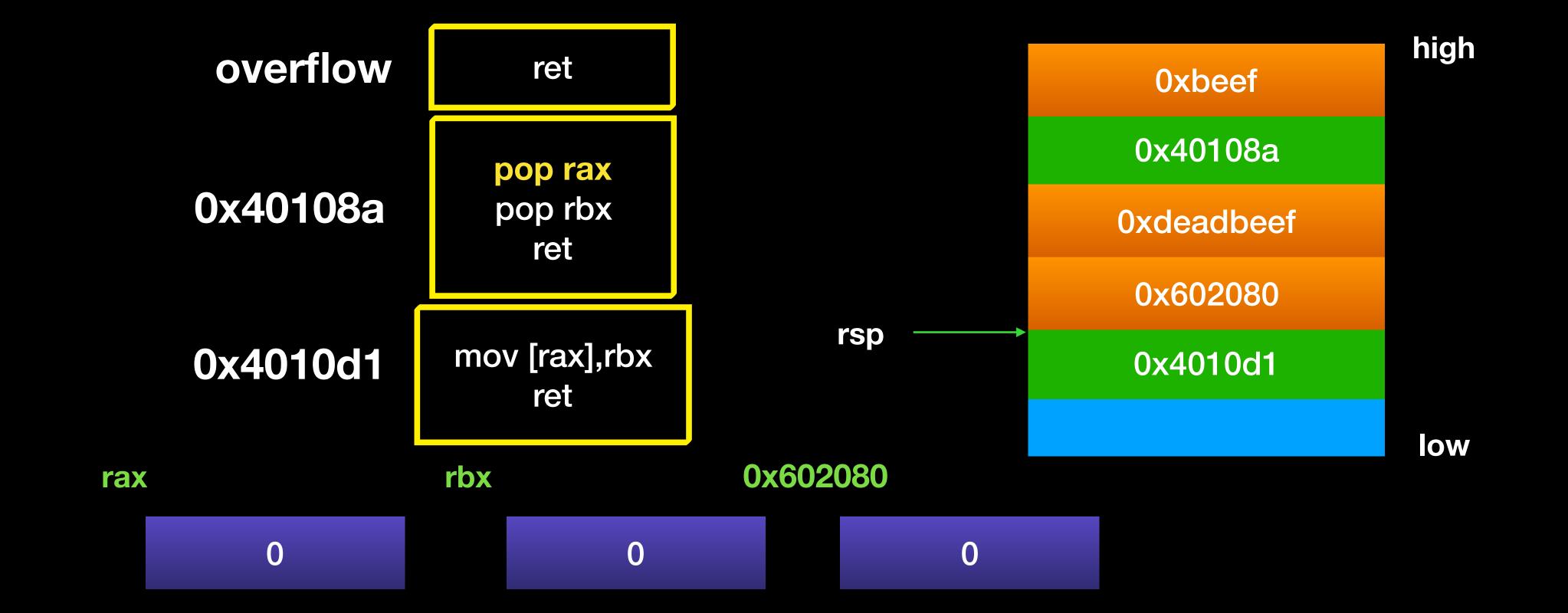


- Write to Memory
 - mov [reg],reg
 - mov [reg+xx], reg
 - •

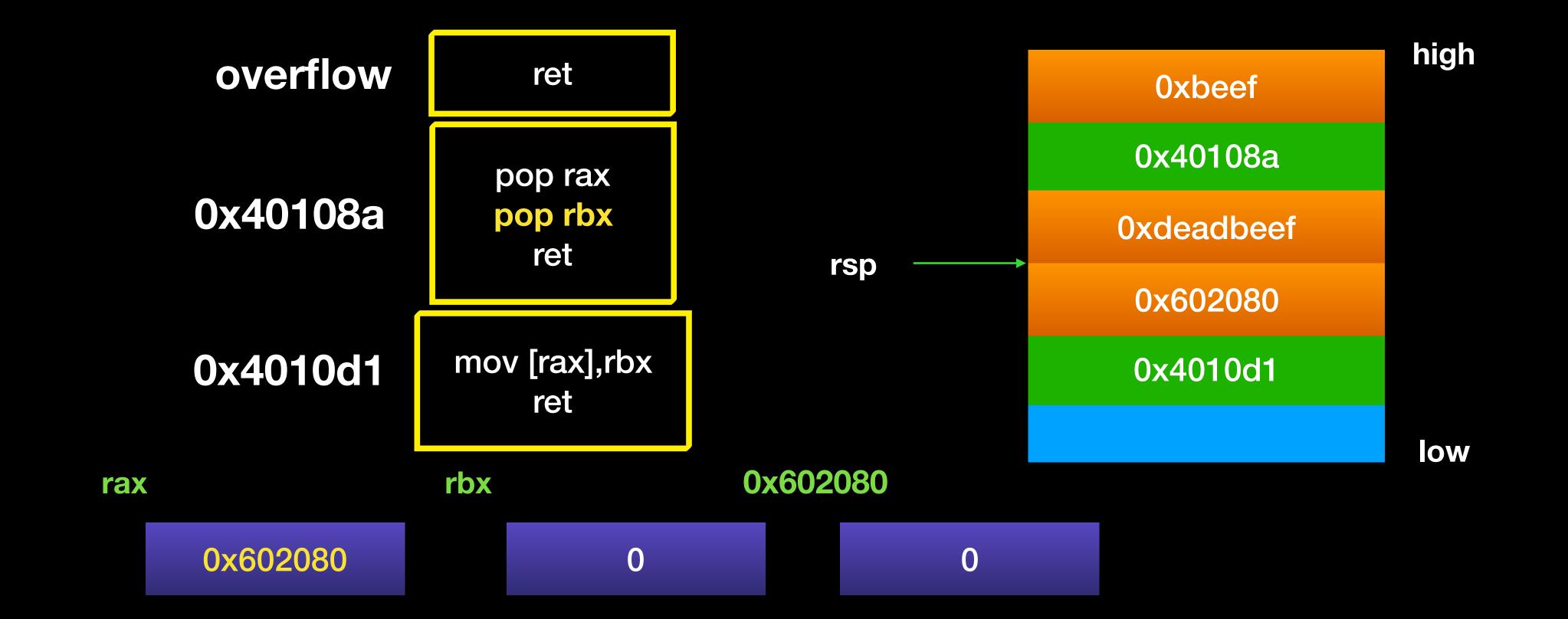
- Write to Memory
 - let *0x602080 = 0xdeadbeef



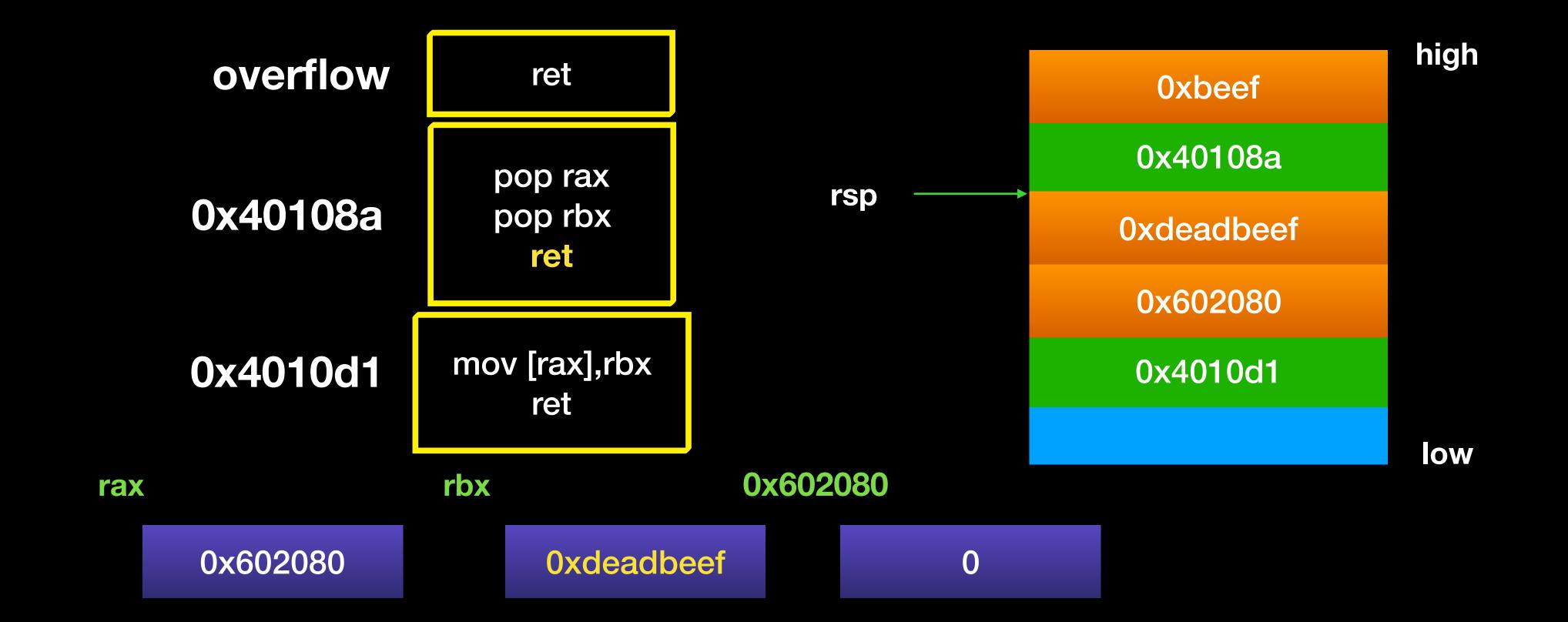
- Write to Memory
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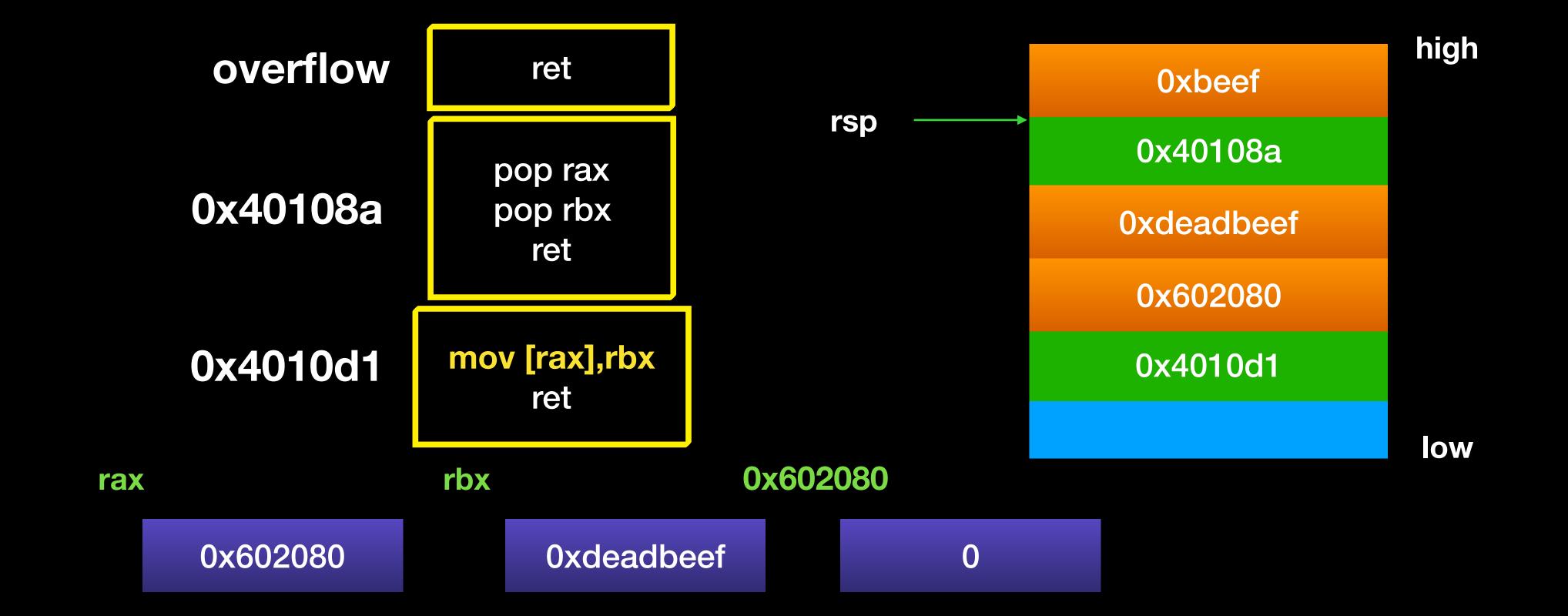
- Write to Memory
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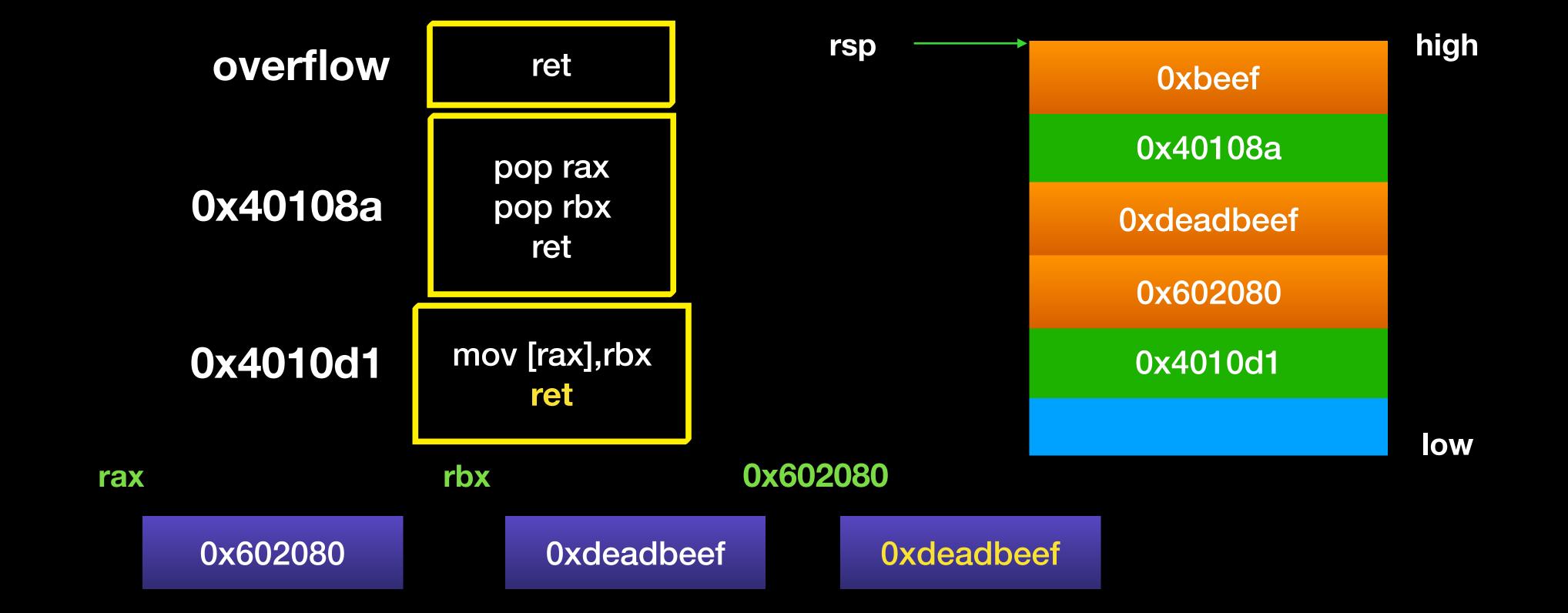
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- execve("/bin/sh",NULL,NULL)
 - write to memory
 - 將 "/bin/sh" 寫入已知位置記憶體中
 - 可分多次將所需字串寫入記憶體中

0x602080 0x602088

/bin/das h\x00\x00\x00...

- execve("/bin/sh",NULL,NULL)
 - write to register
 - rax = 0x3b , rdi = address of "/bin/sh"
 - rsi = 0, rdx = 0
 - syscall

- find gadget
 - https://github.com/JonathanSalwan/ROPgadget

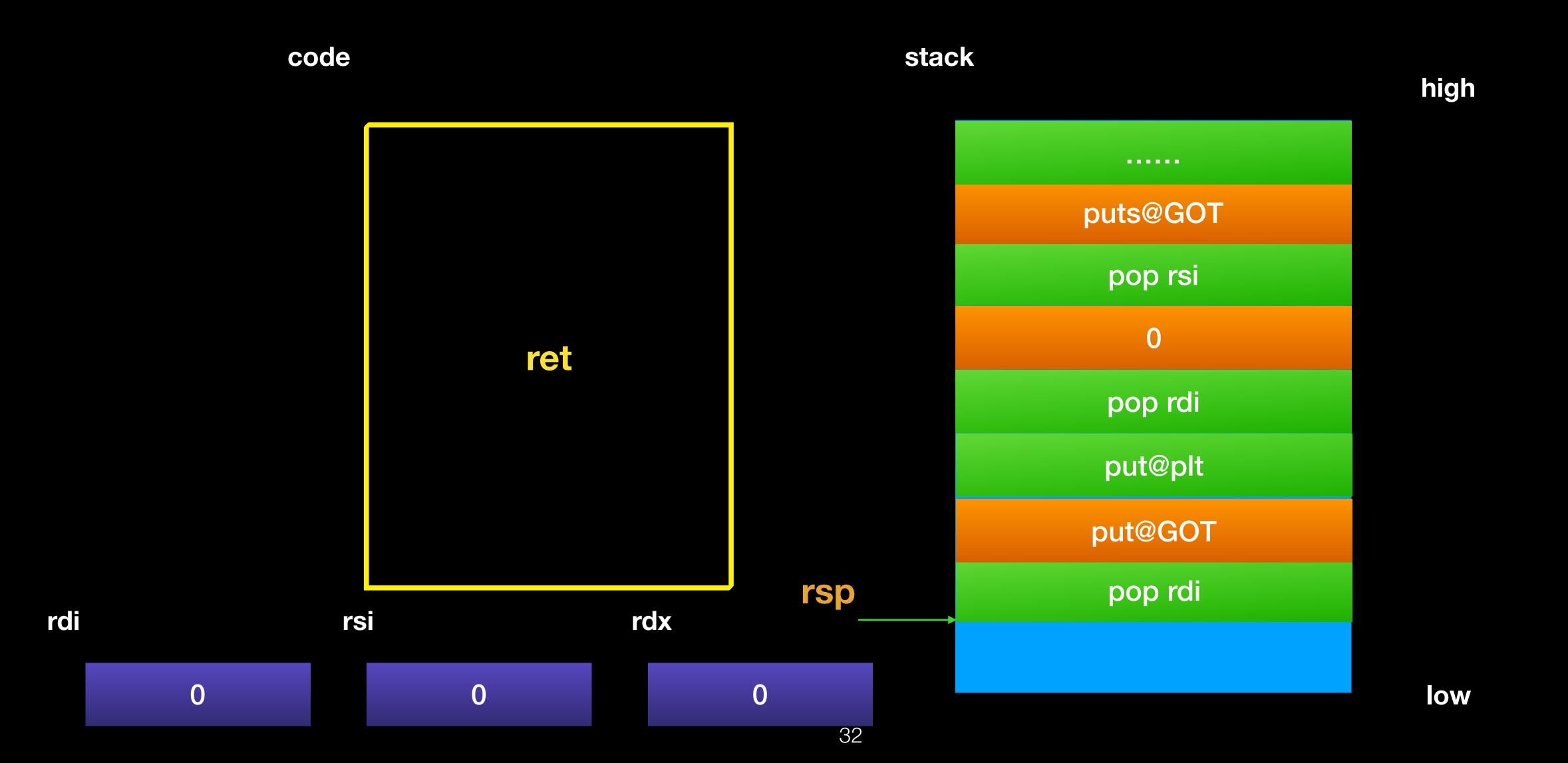
```
0x000000000001947 : pop r12 ; pop r13 ; pop r14 ; pop r15 ; ret
0x000000000004016da : pop r12 ; pop r13 ; pop r14 ; ret
0x00000000000401ee0 : pop r12 ; pop r13 ; ret
0x0000000000001949 : pop r13 ; pop r14 ; pop r15 ; ret
0x000000000004016dc : pop r13 ; pop r14 ; ret
0x000000000000401ee2 : pop r13 ; ret
0x000000000000194b : pop r14 ; pop r15 ; ret
0x00000000000016de : pop r14 ; ret
0x00000000000040194d : pop r15 ; ret
0x000000000004026c6 : pop rax ; add rsp, 8 ; pop rbp ; ret
0x00000000000040260d : pop rax ; ret
0x00000000000001ff : pop rbp ; mov edi, 0x604018 ; jmp rax
0x000000000001946 : pop rbp ; pop r12 ; pop r13 ; pop r14 ; pop r15 ; ret
0x00000000000016d9 : pop rbp ; pop r12 ; pop r13 ; pop r14 ; ret
0x0000000000001edf : pop rbp ; pop r12 ; pop r13 ; ret
0x000000000000194a : pop rbp ; pop r14 ; pop r15 ; ret
0x0000000000004016dd : pop rbp ; pop r14 ; ret
0x000000000000400f30 : pop rbp ; ret
0x0000000000001ede : pop rbx ; pop rbp ; pop r12 ; pop r13 ; ret
0x0000000000001540 : pop rbx ; pop rbp ; ret
0x00000000000040228e : pop rbx ; ret
0x00000000000040194e : pop rdi ; ret
```

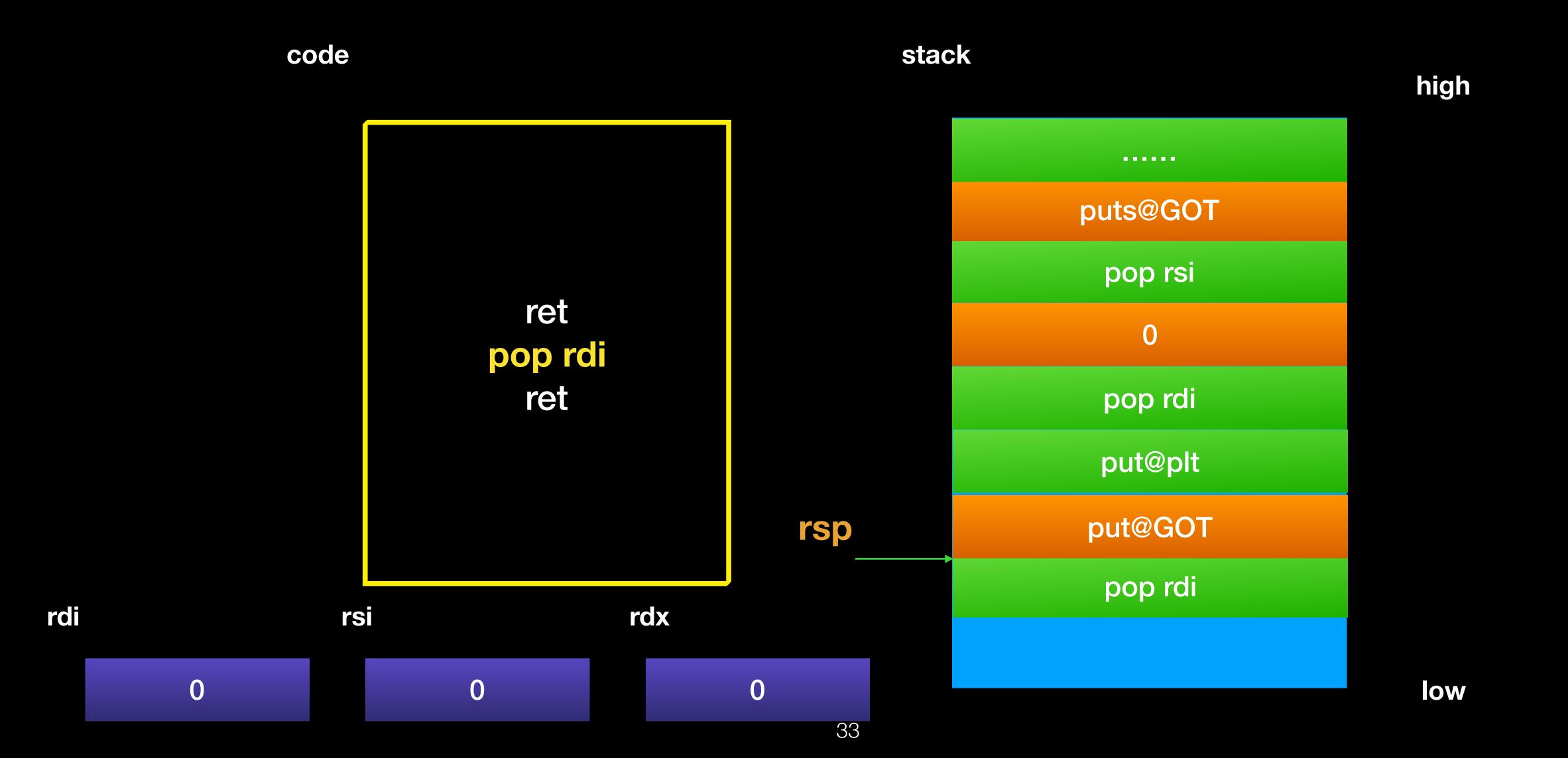
- find gadget
 - ROPgadget - binary binary
 - ROPgadget - ropchain - binary binary
 - 在 Static linking 通常可以組成功 execve 的 rop chain 但通常都很長,需要自己找更短的 gadget 來改短一點

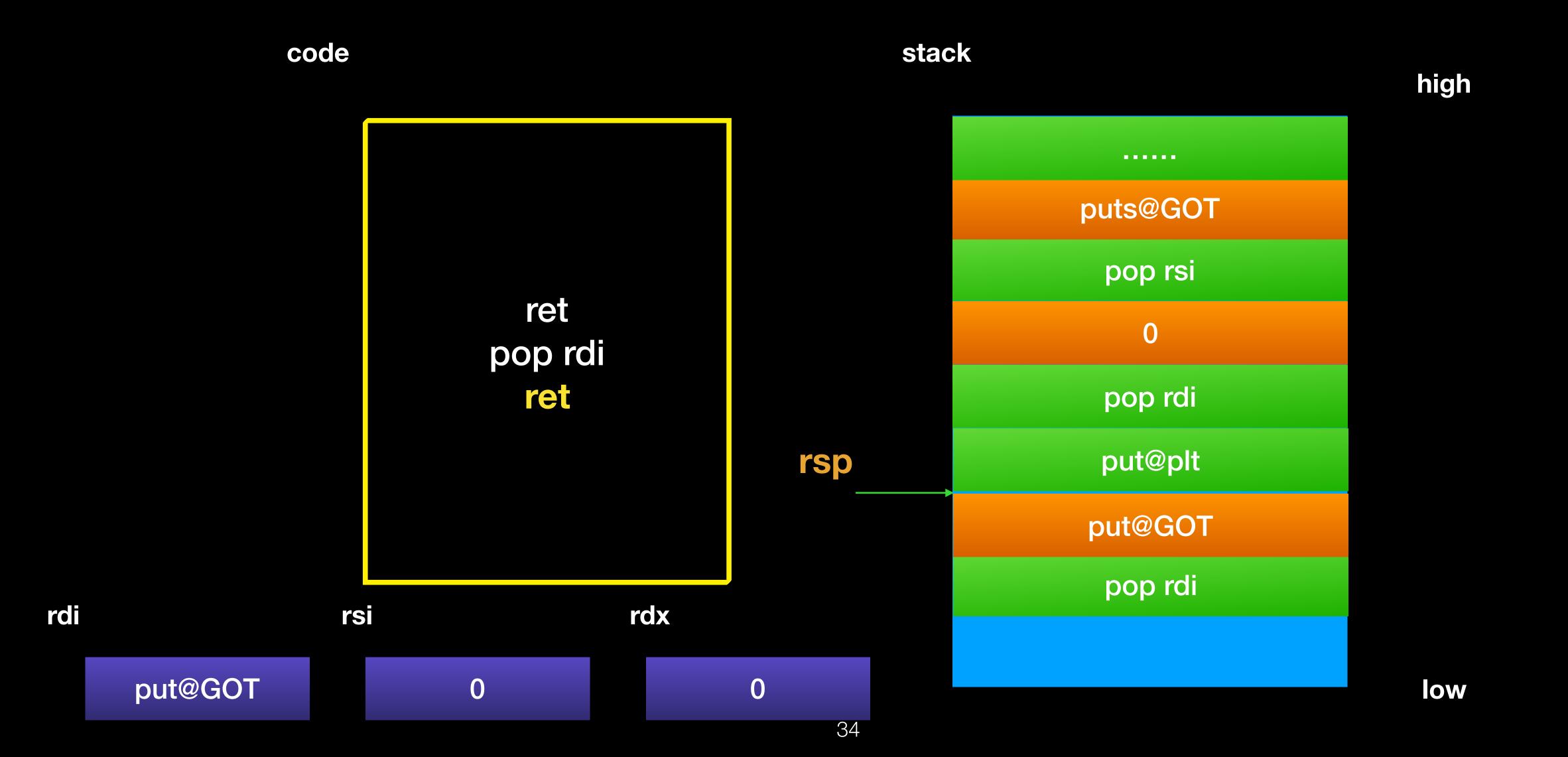
Outline

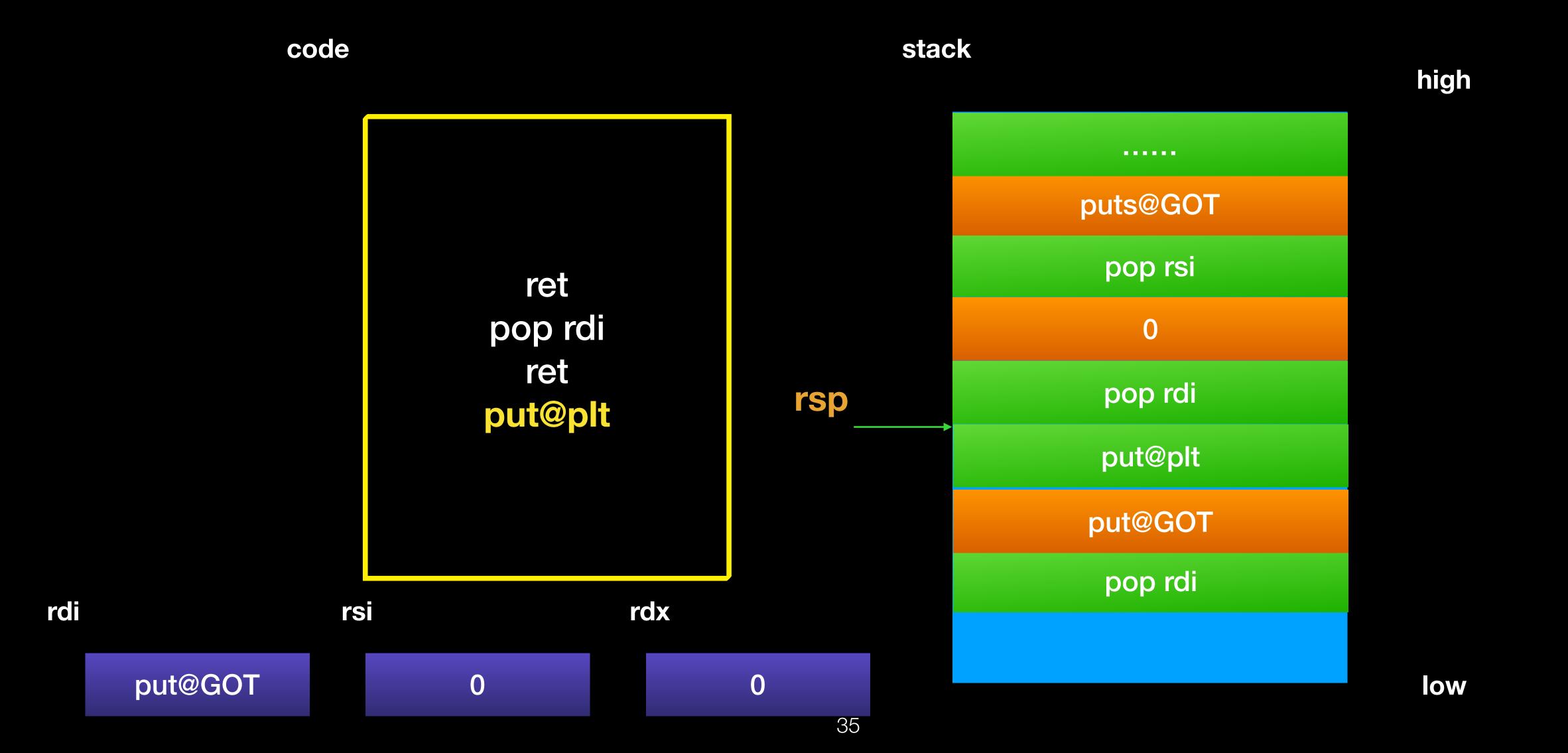
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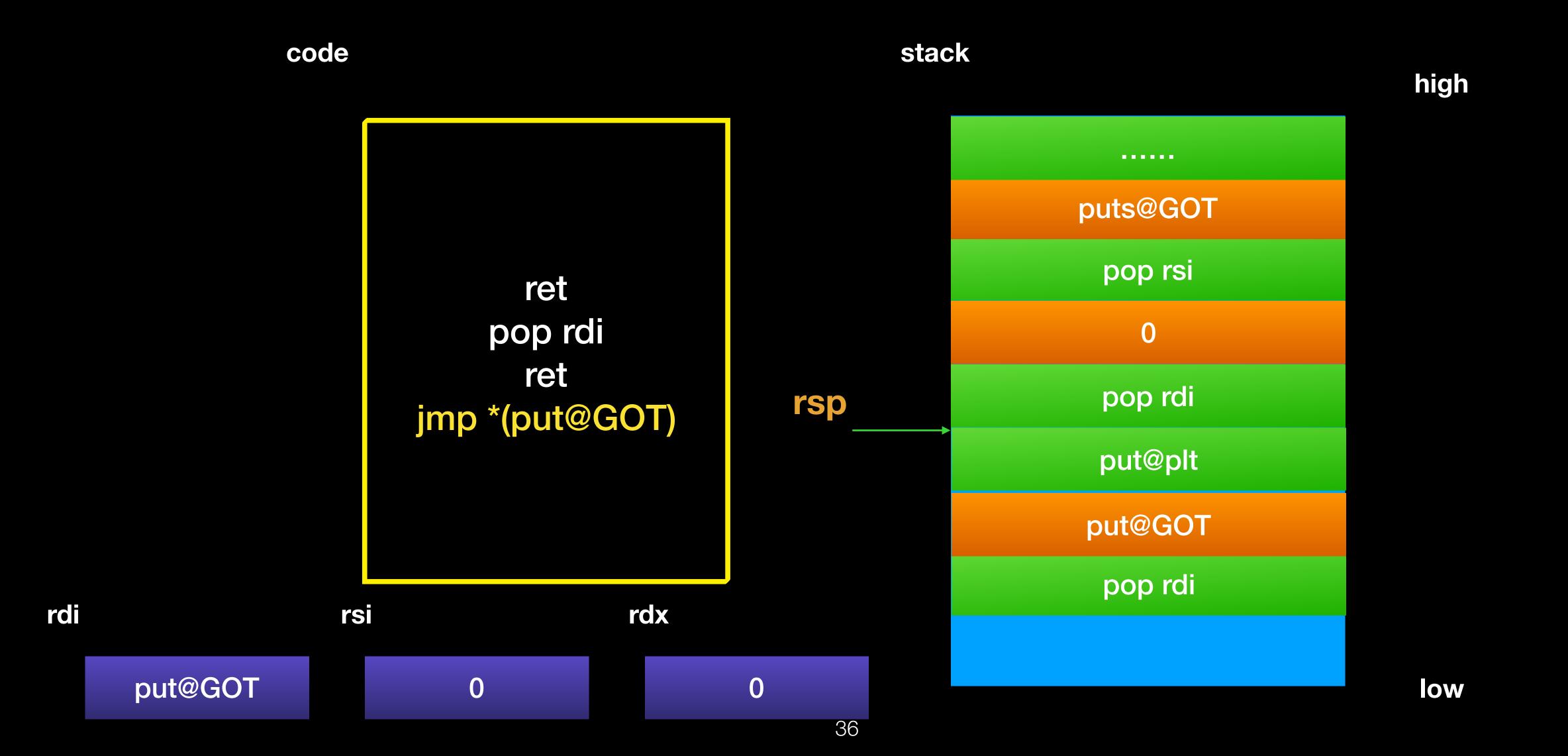
- 假設 dynamic 編譯的程式中有 Buffer Overflow 的漏洞且在沒 PIE 情況下(先不考慮 StackGuard 的情況)
- How to bypass ASLR and DEP?
 - Use .plt section to leak some information
 - ret2plt
 - 通常一般的程式中都會有 put 、 send 、 write 等 output function

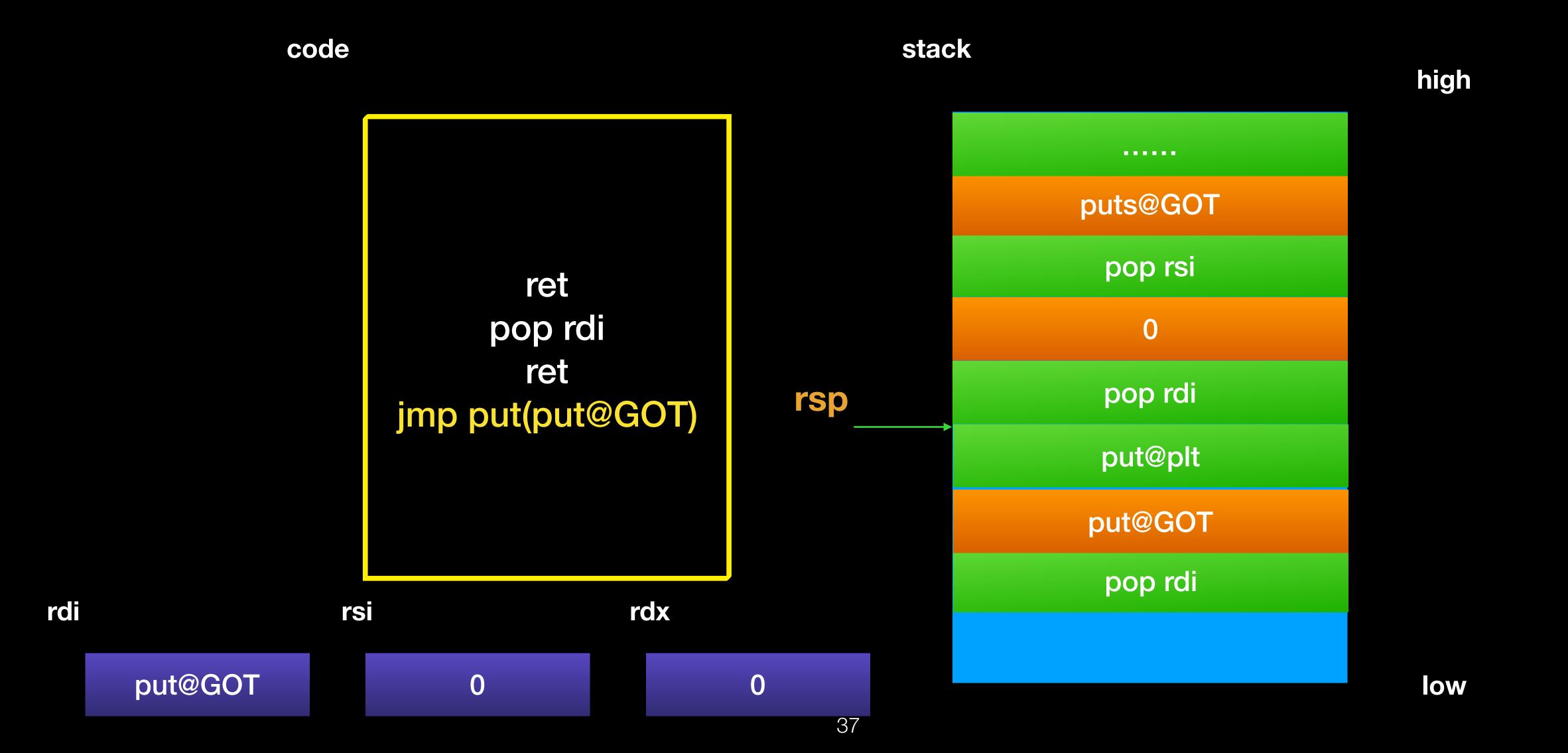


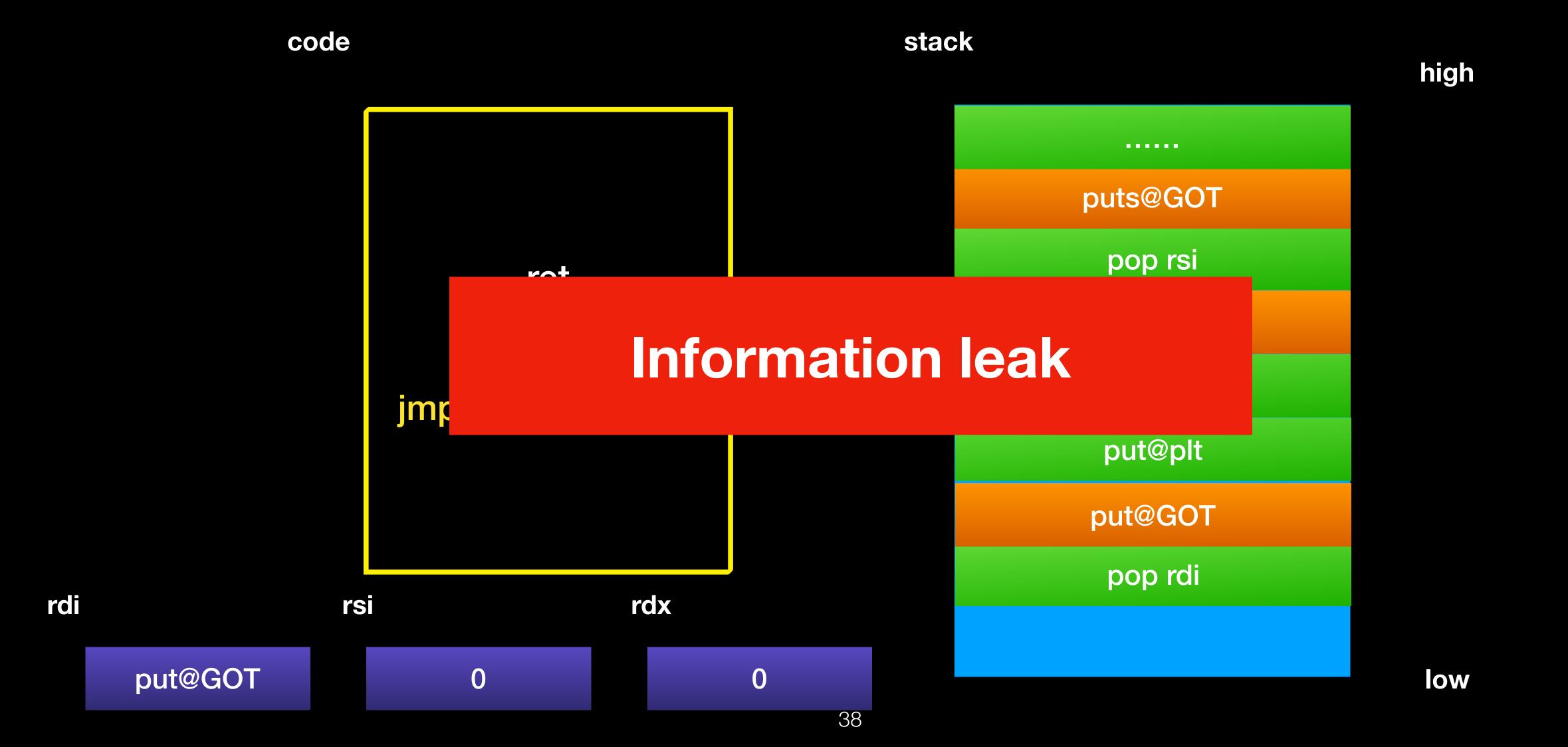


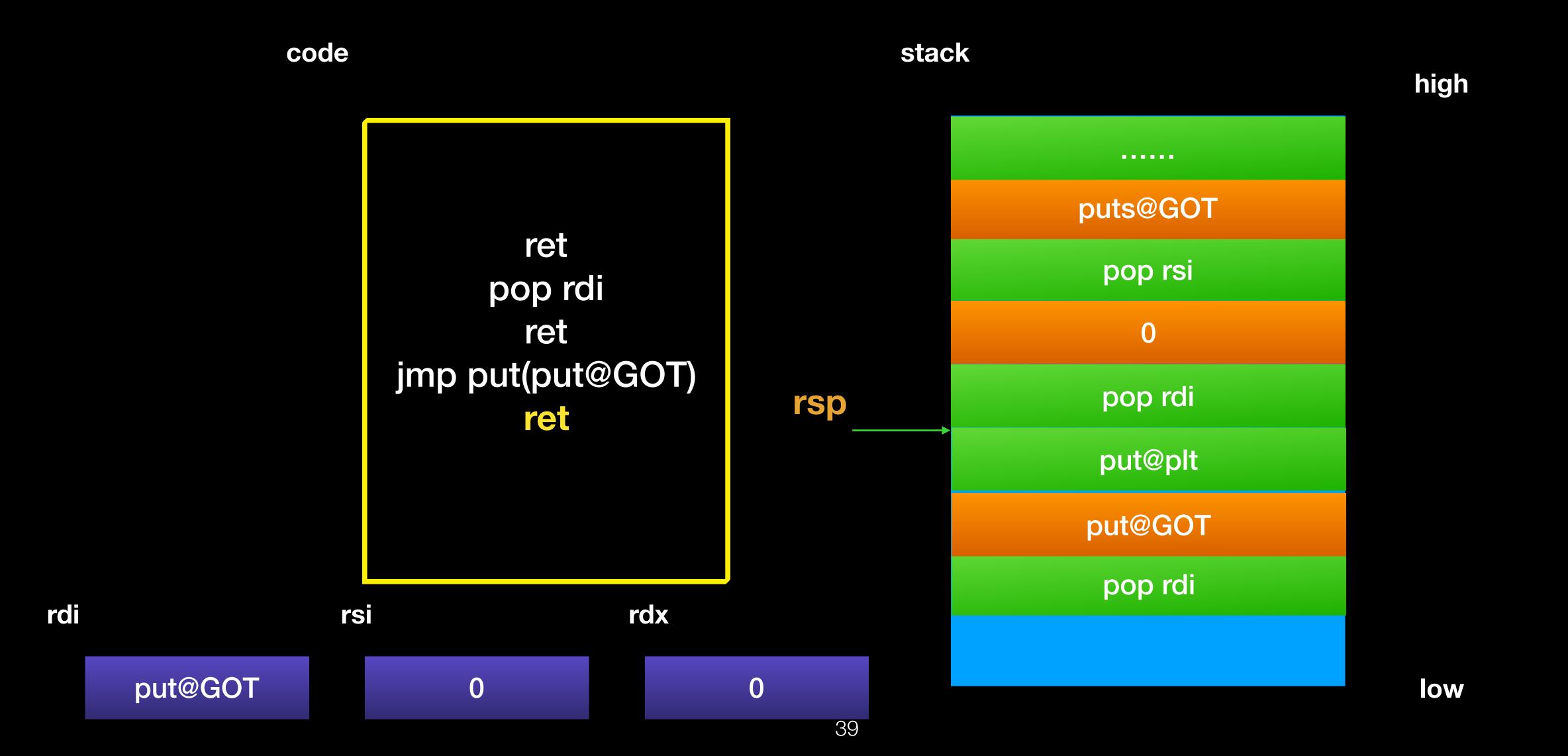


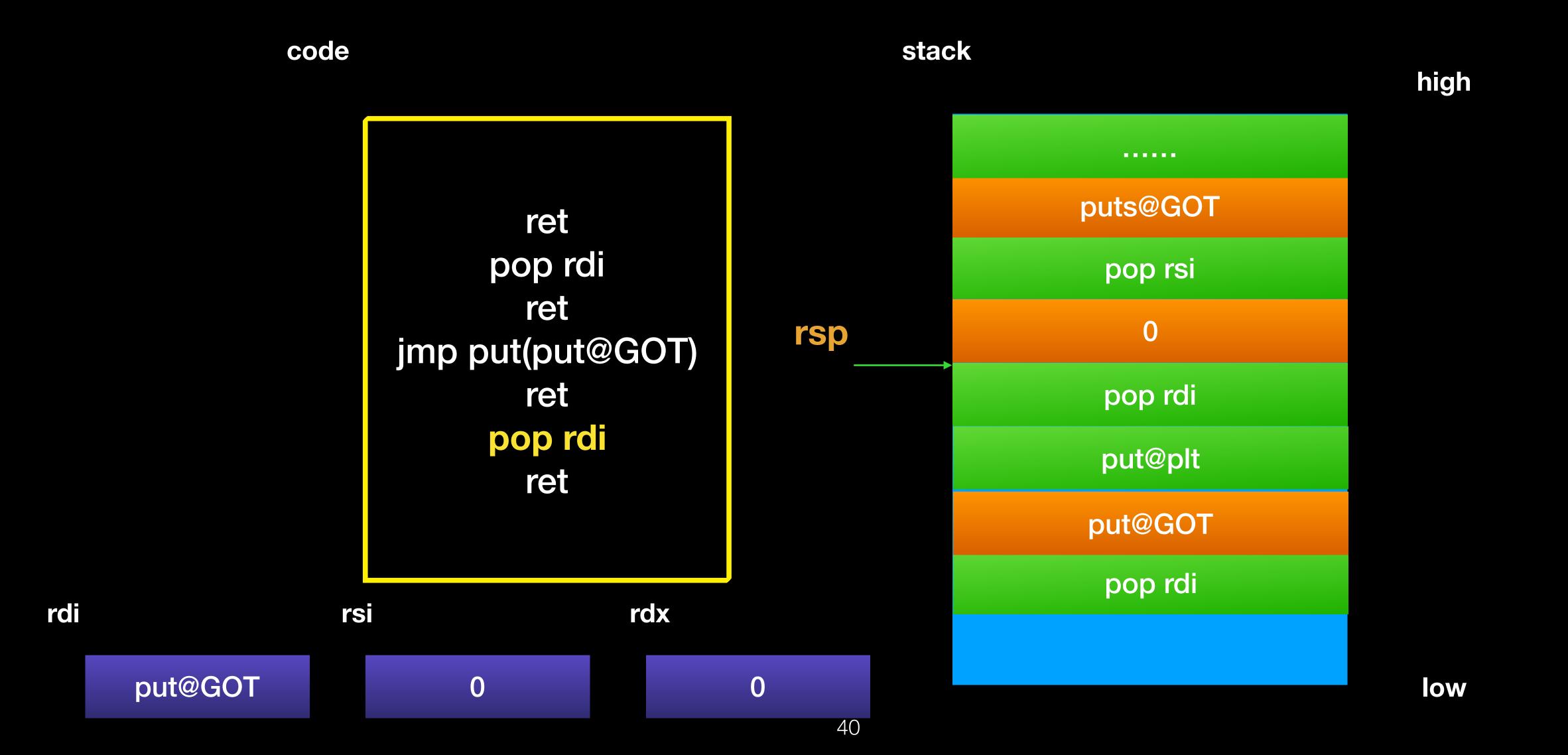


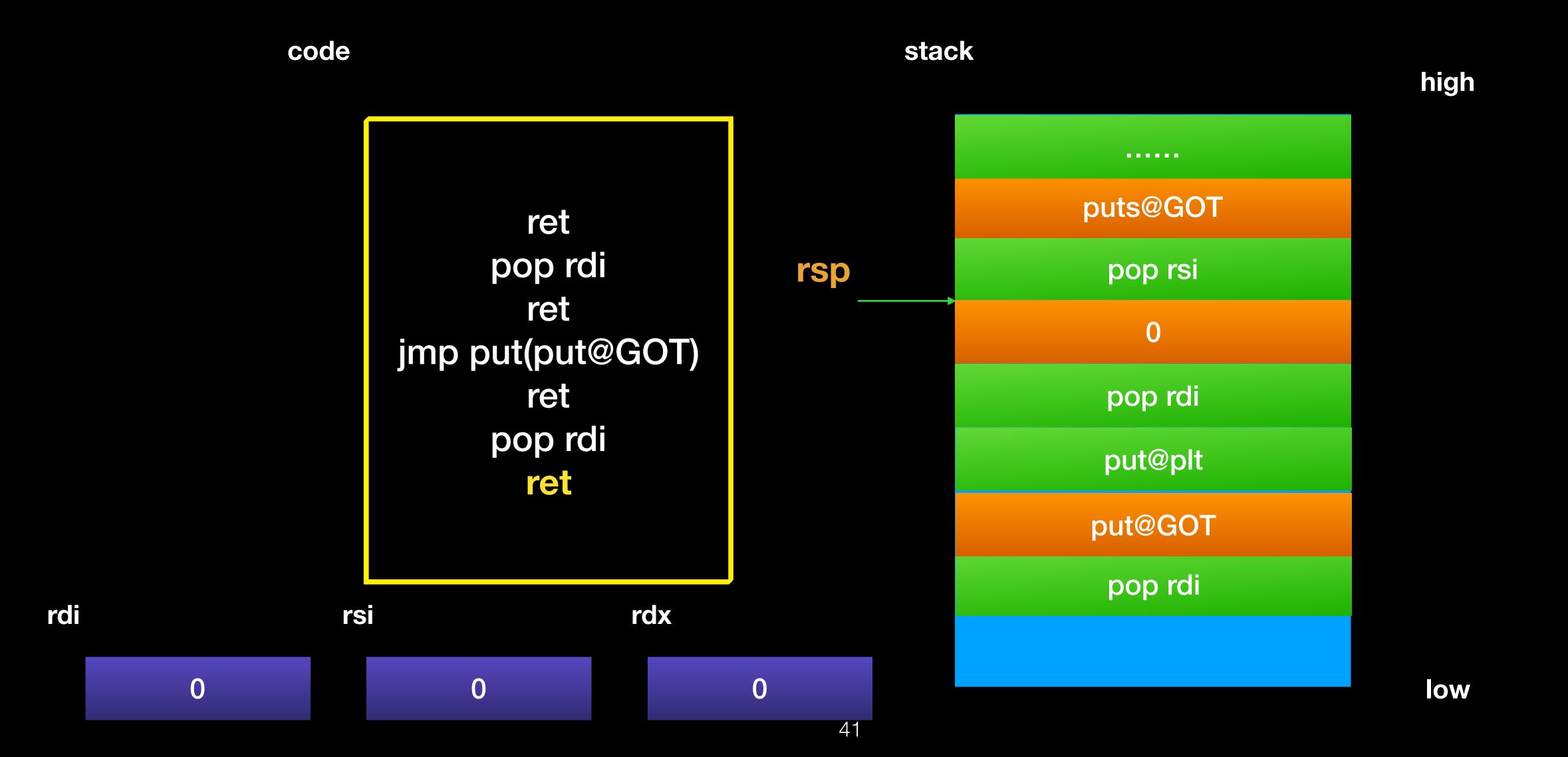


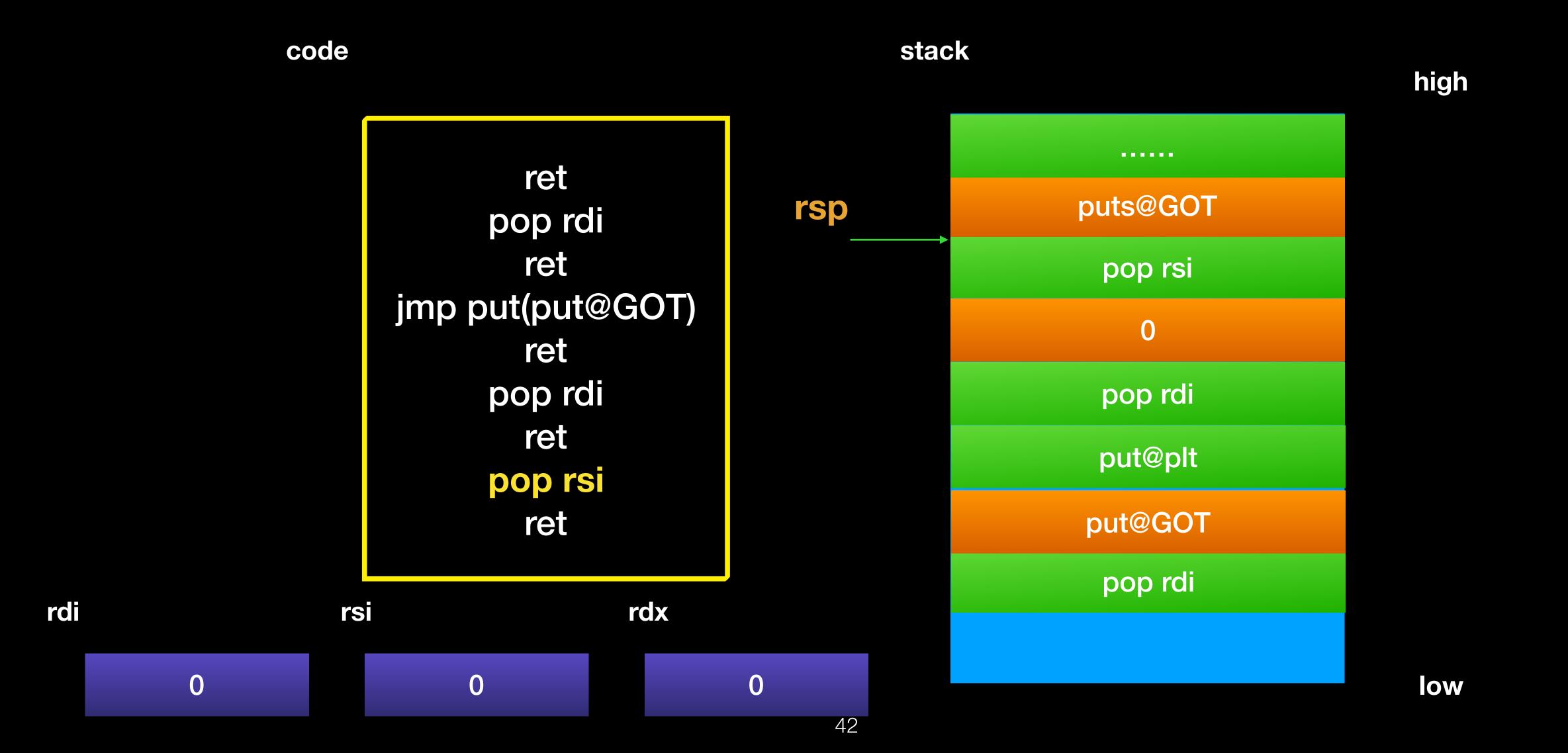


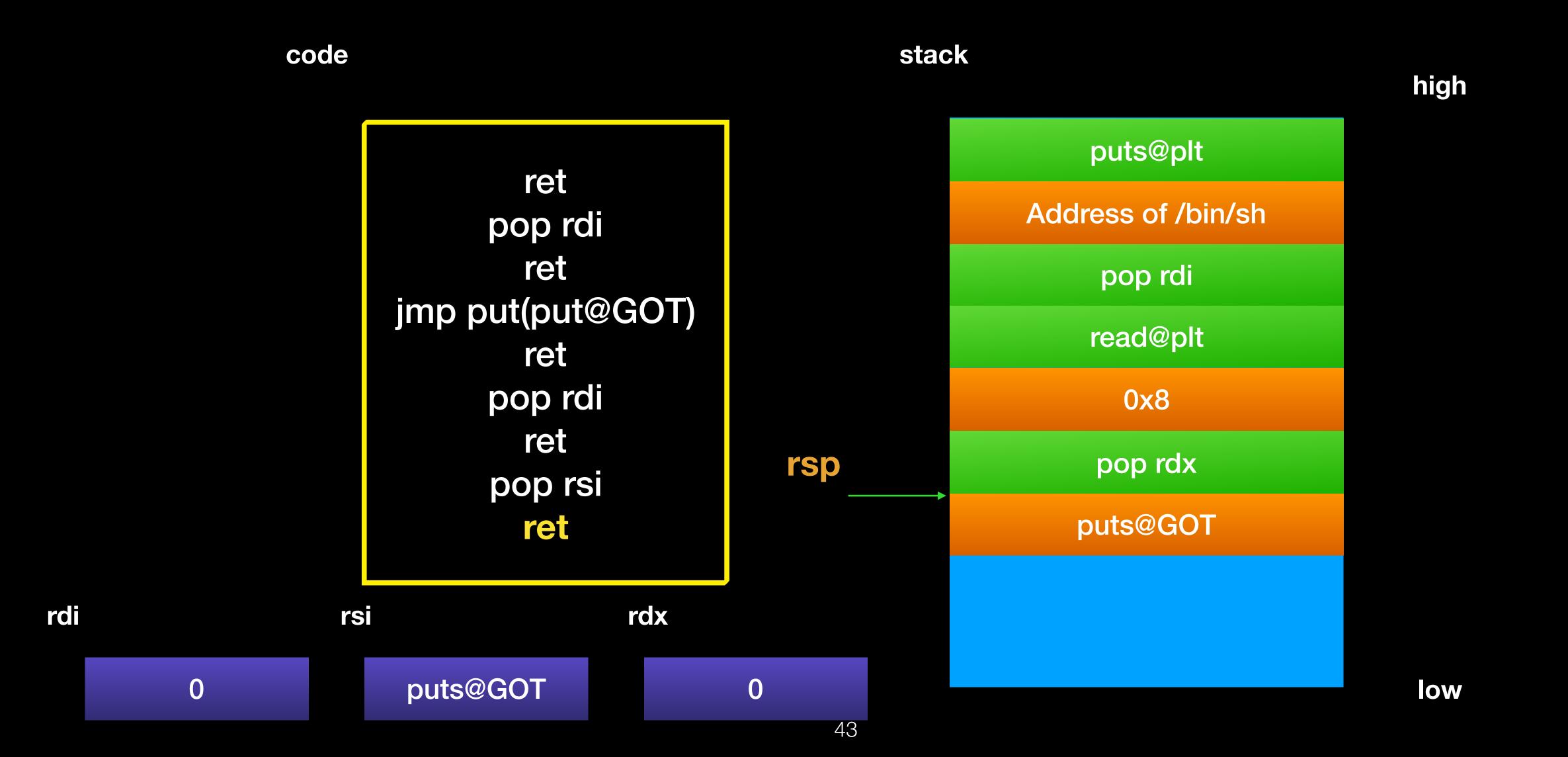


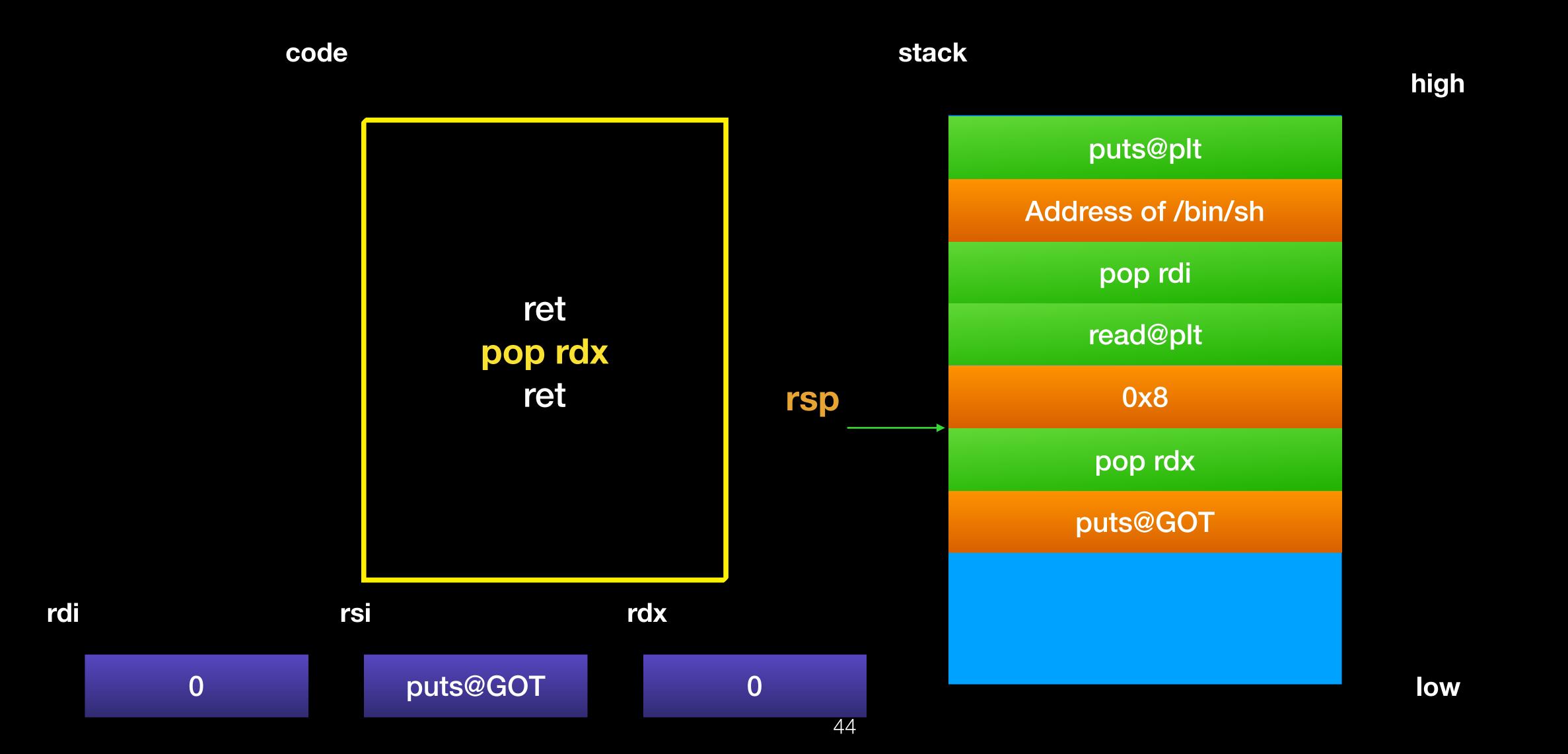


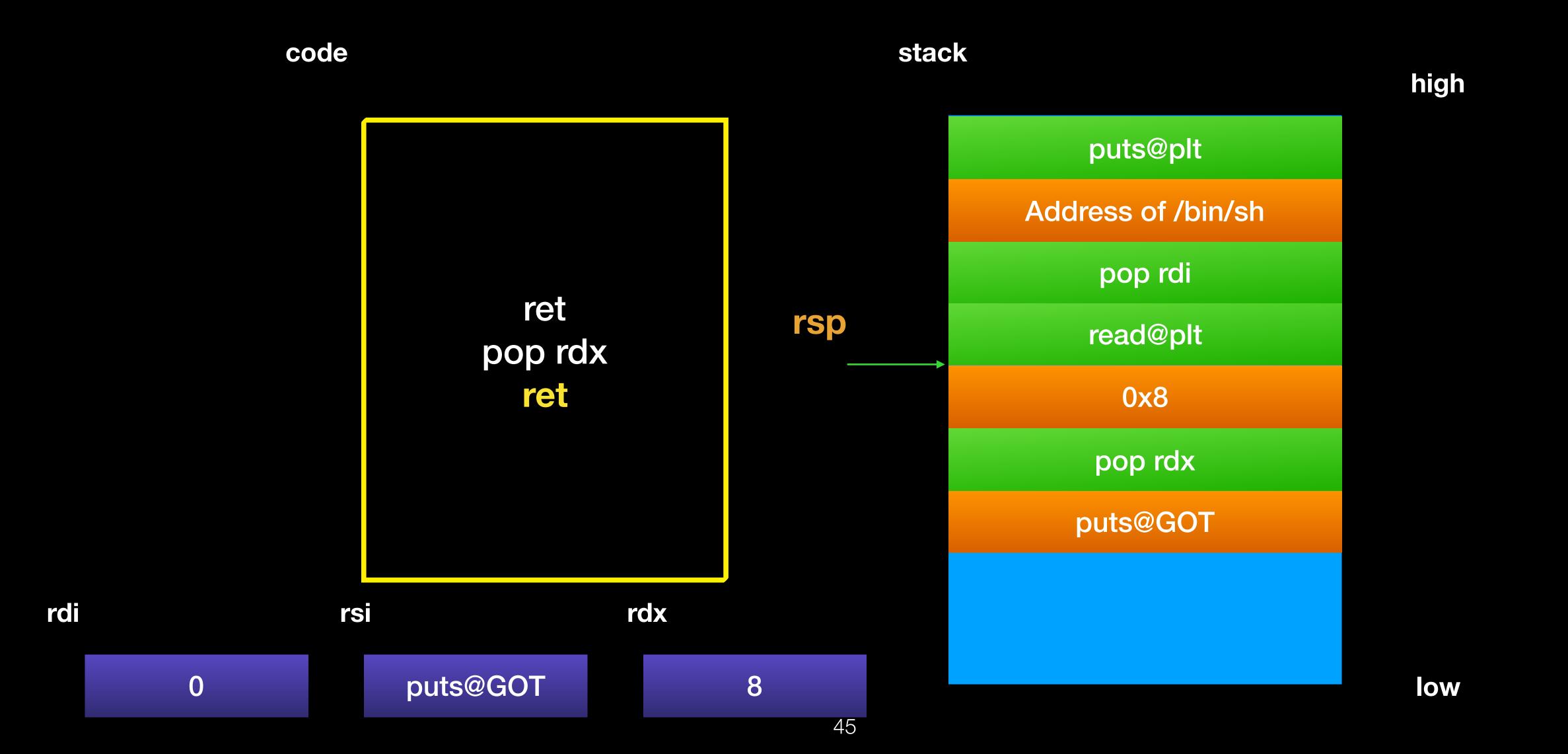


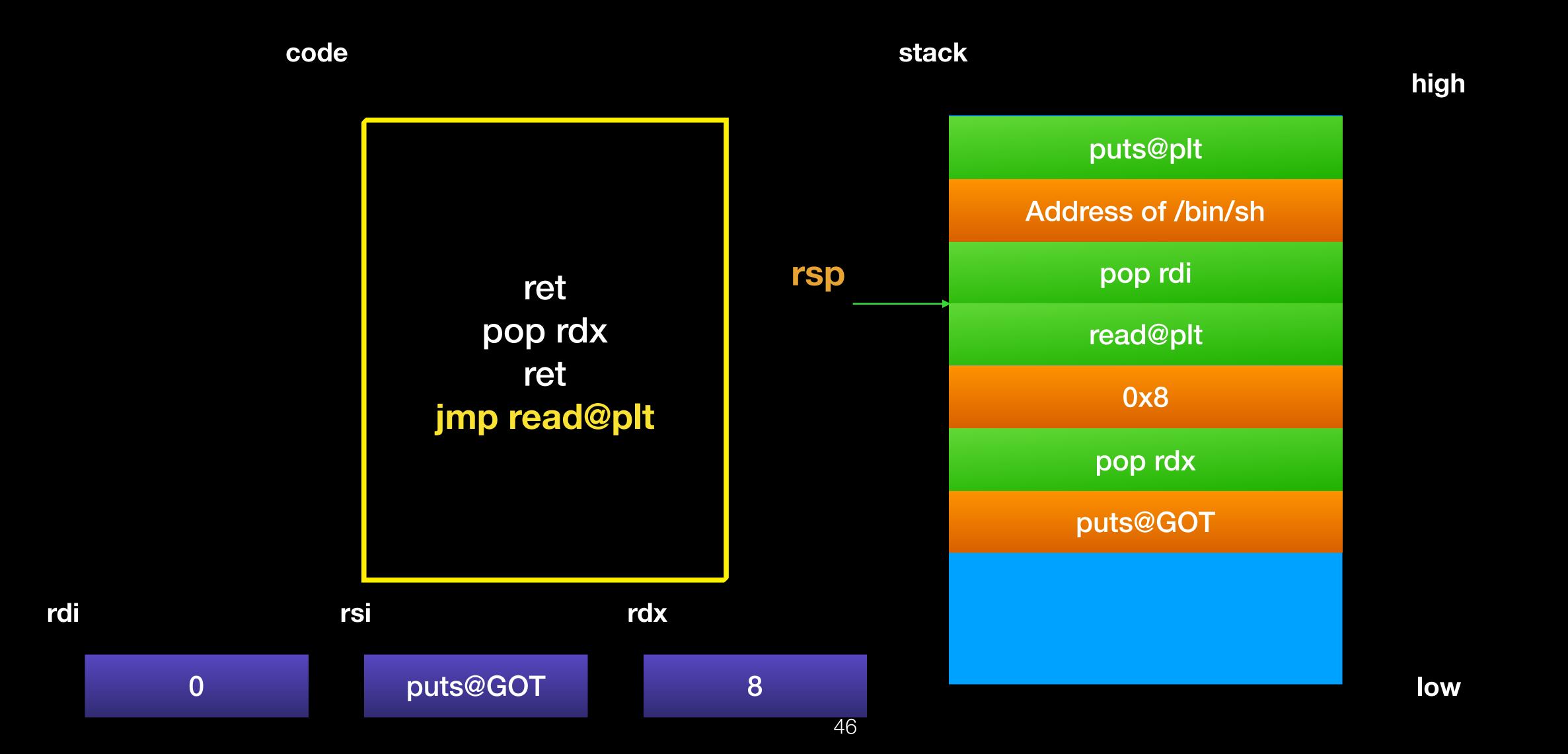


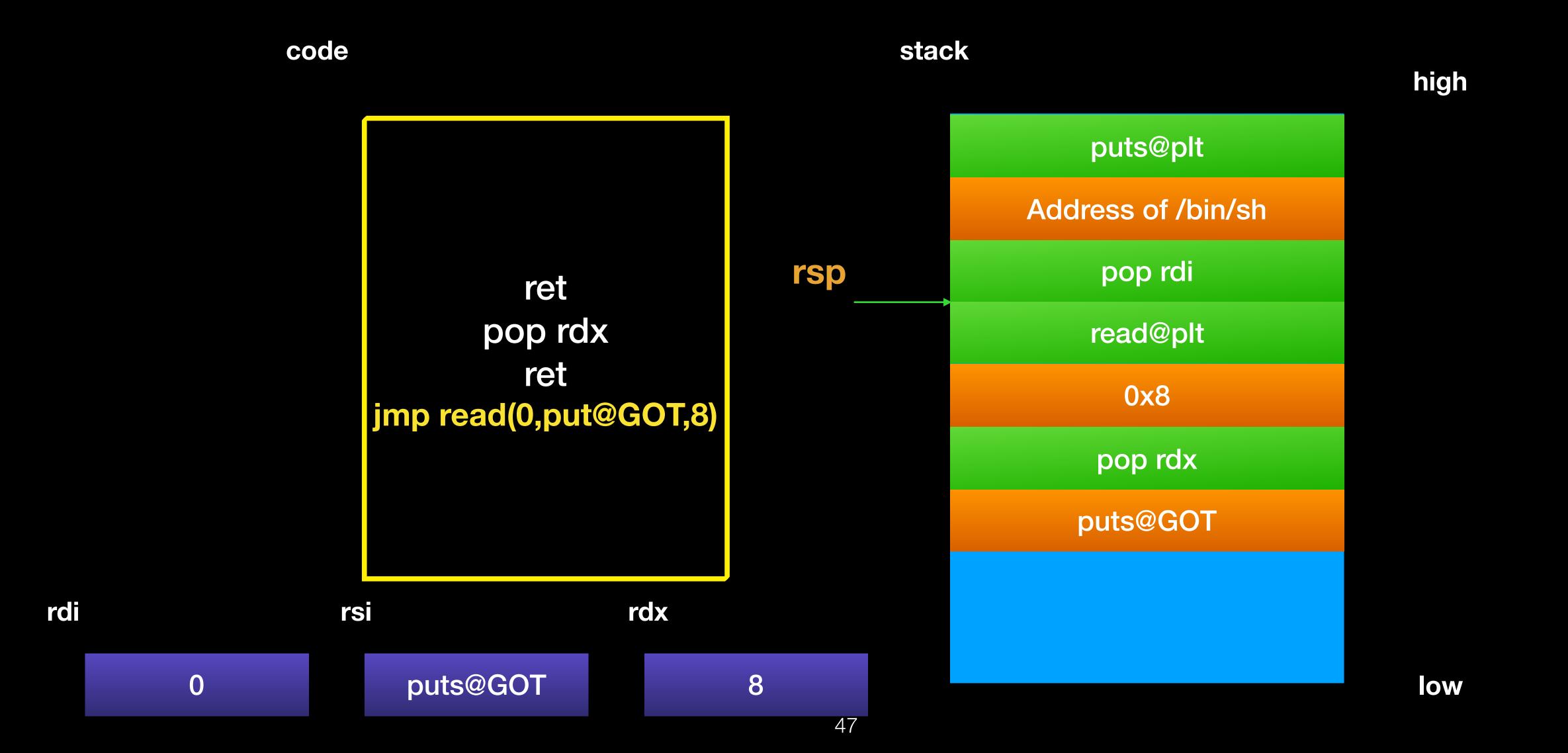


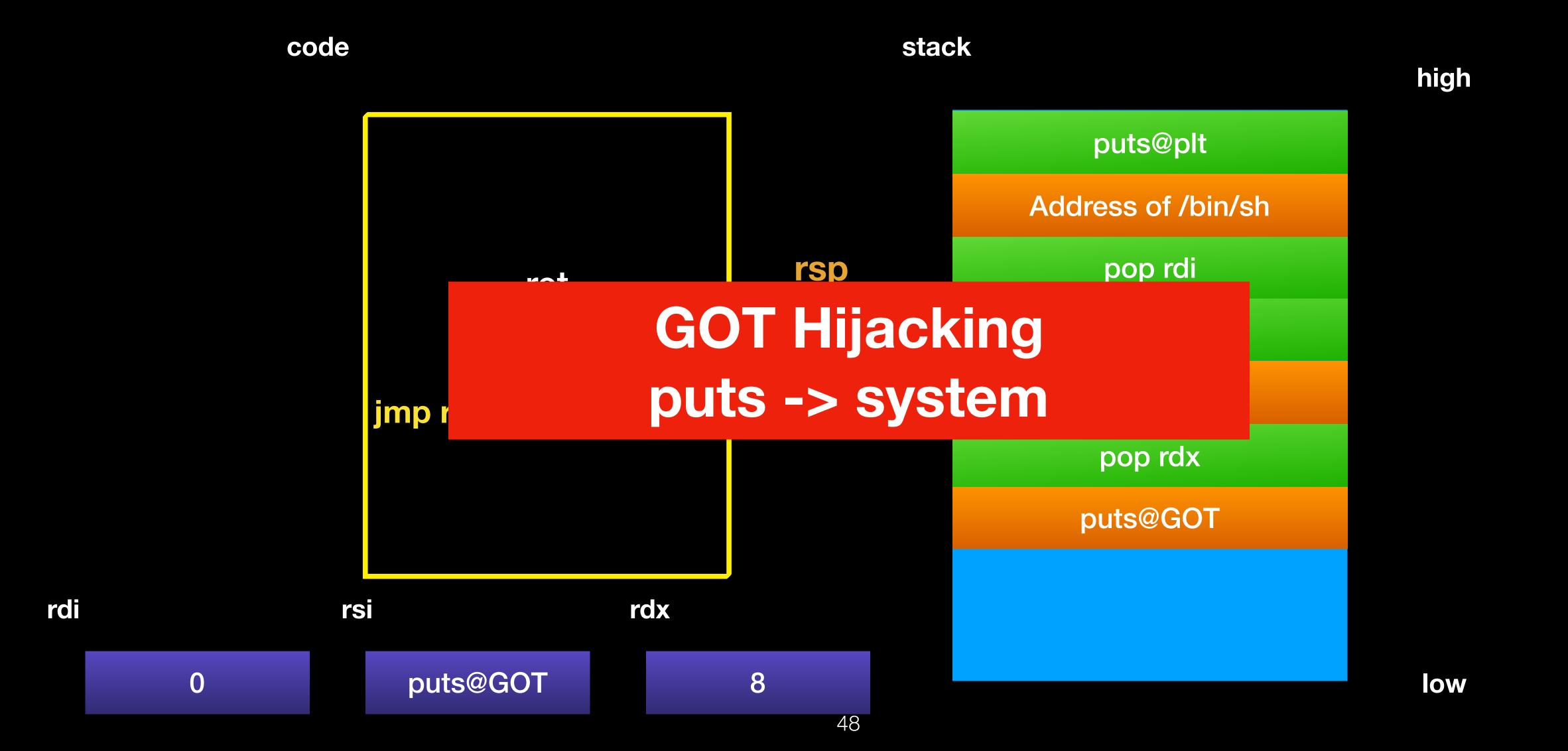


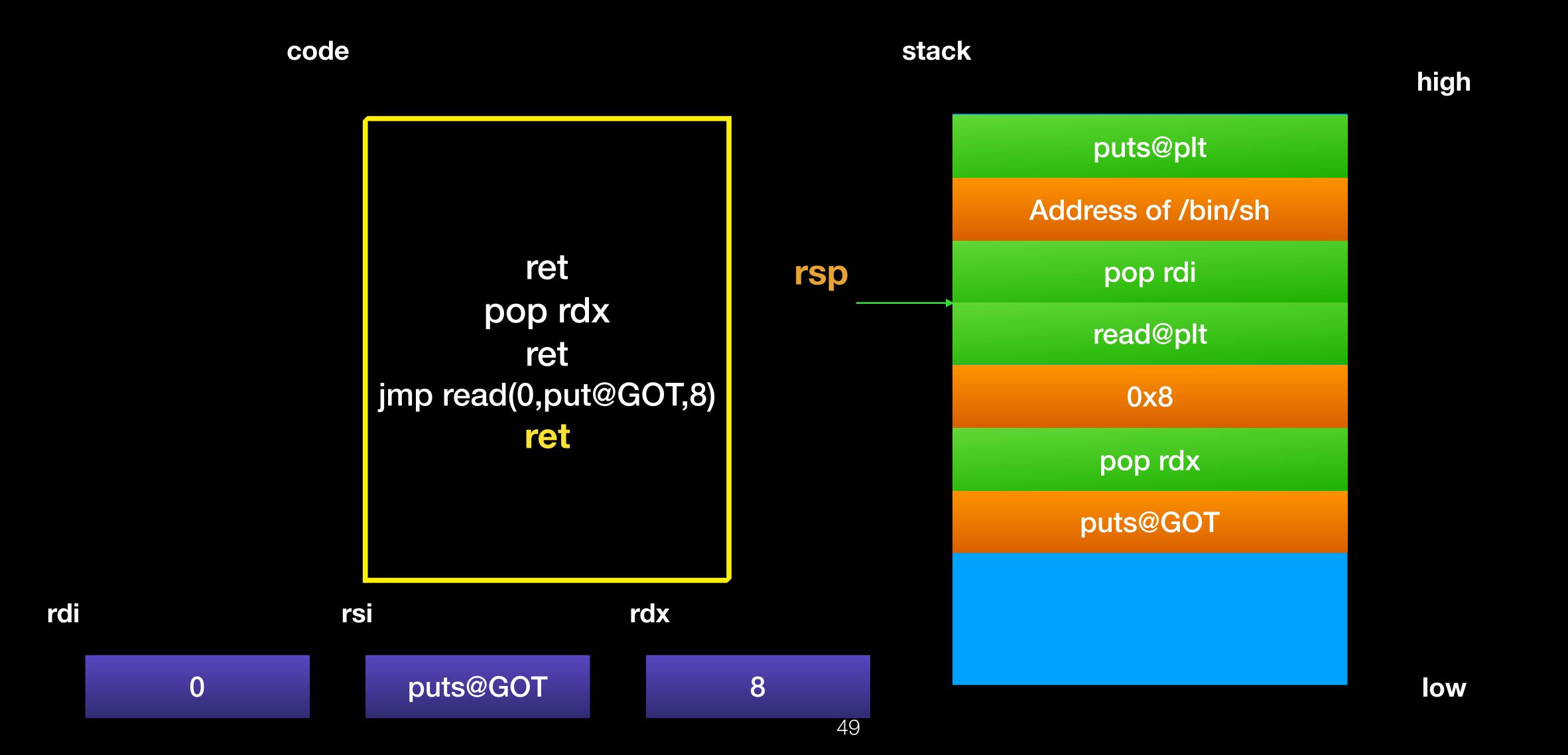


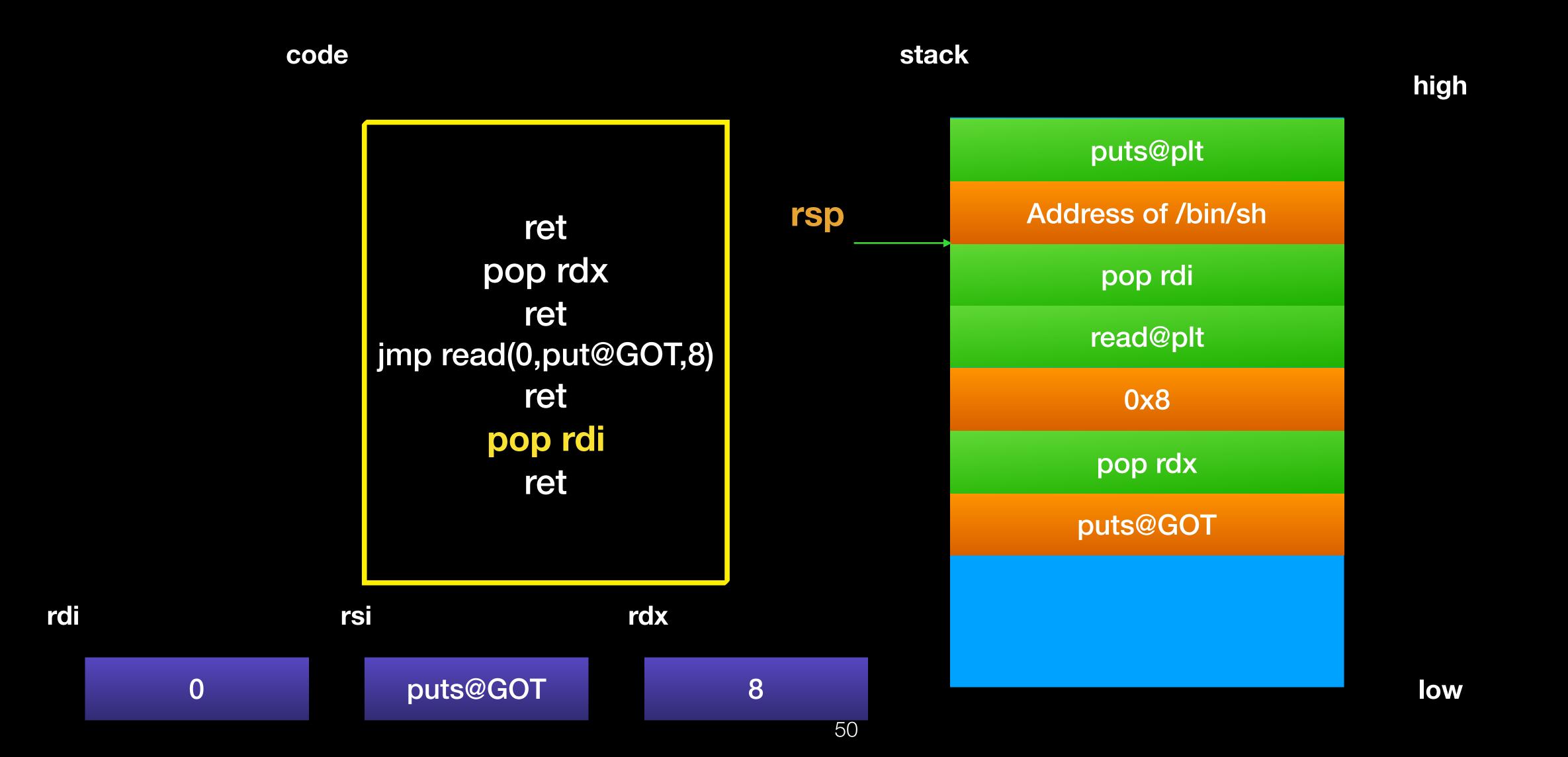


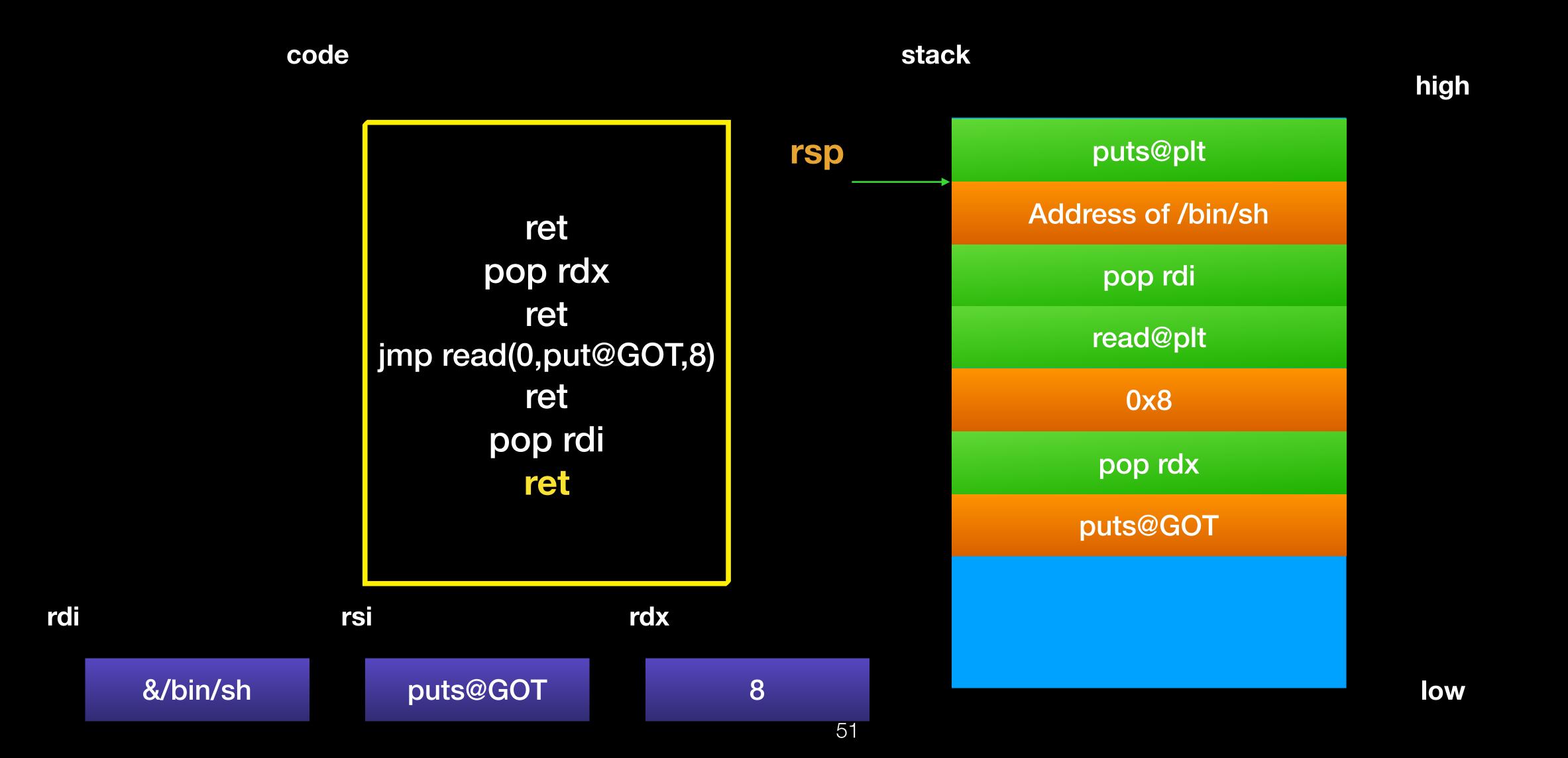


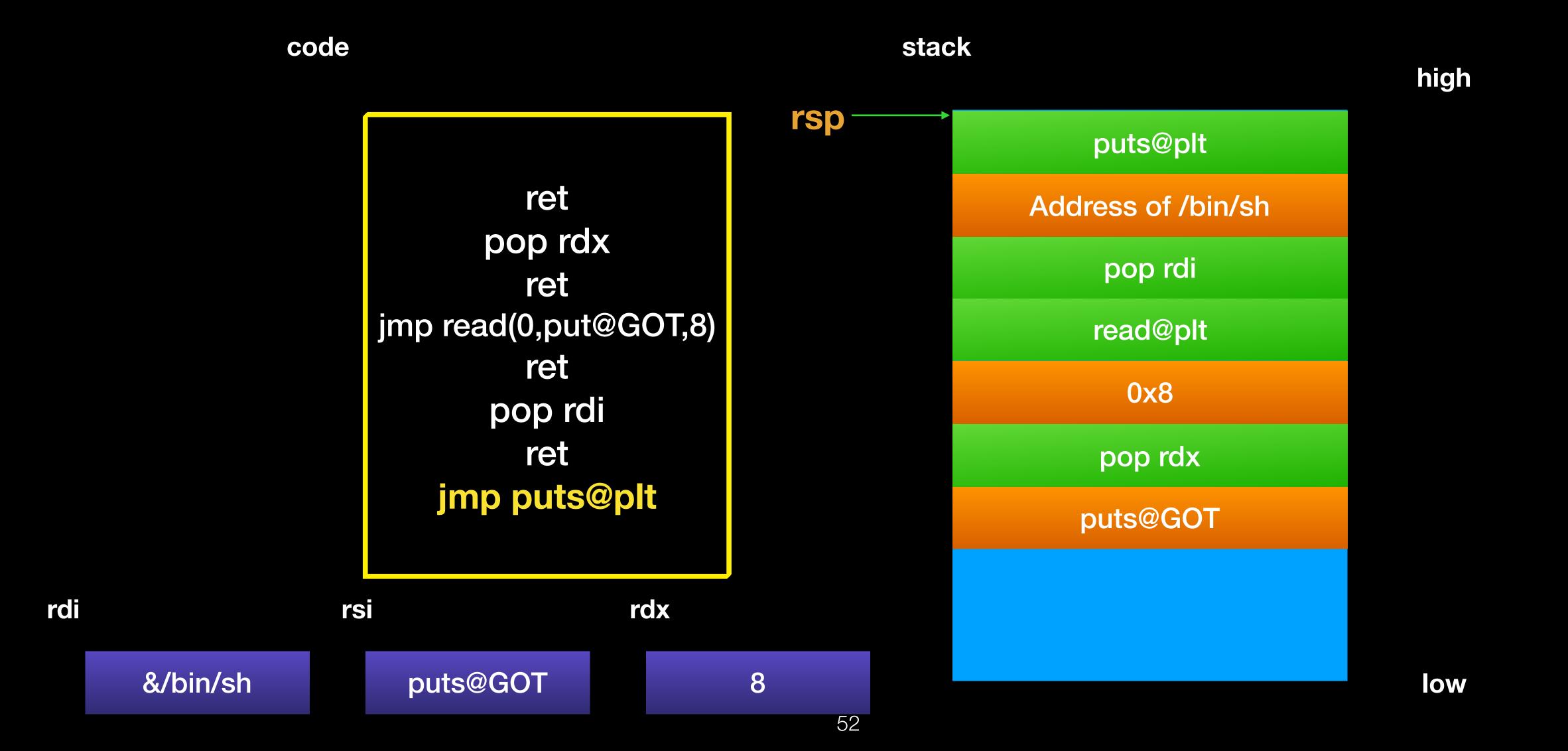


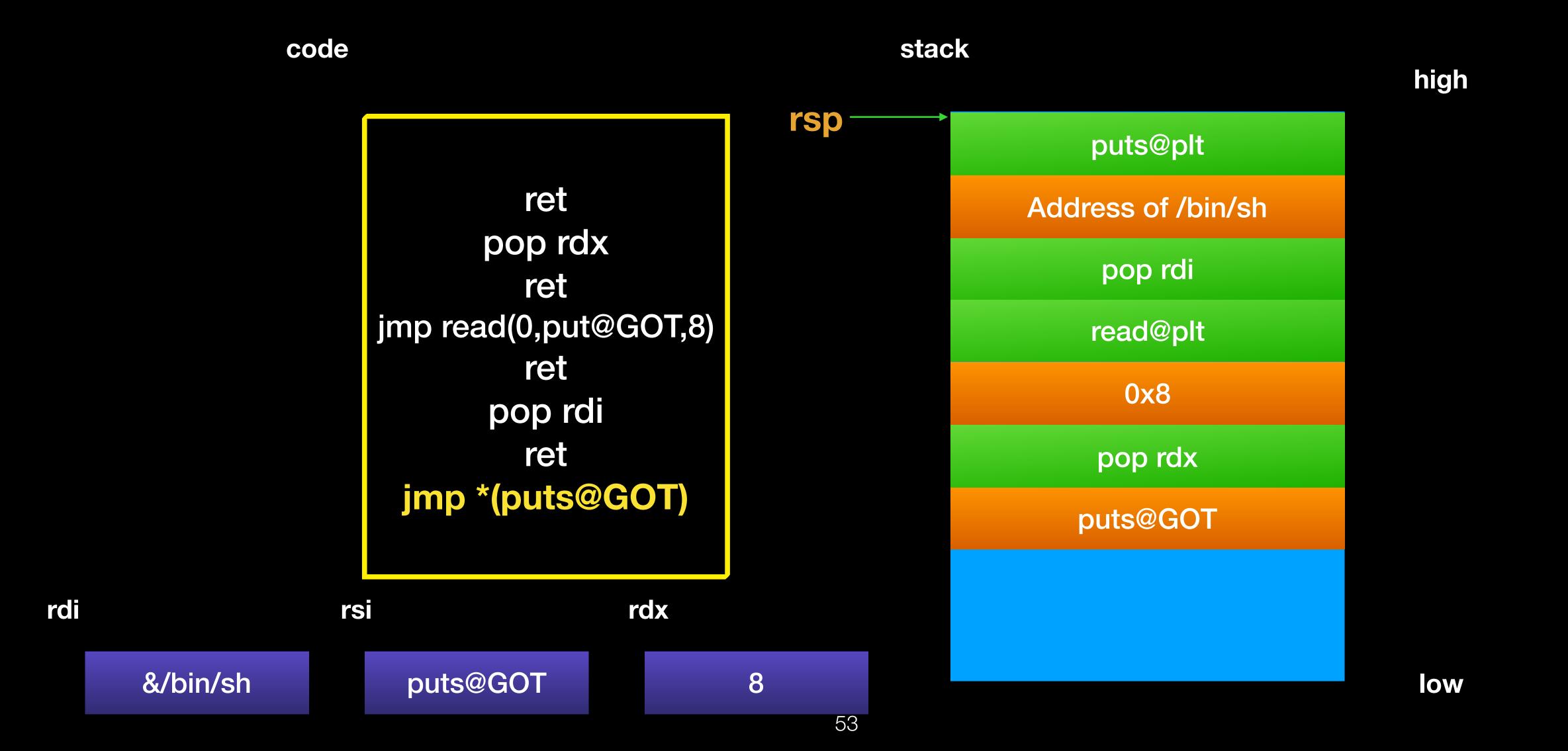


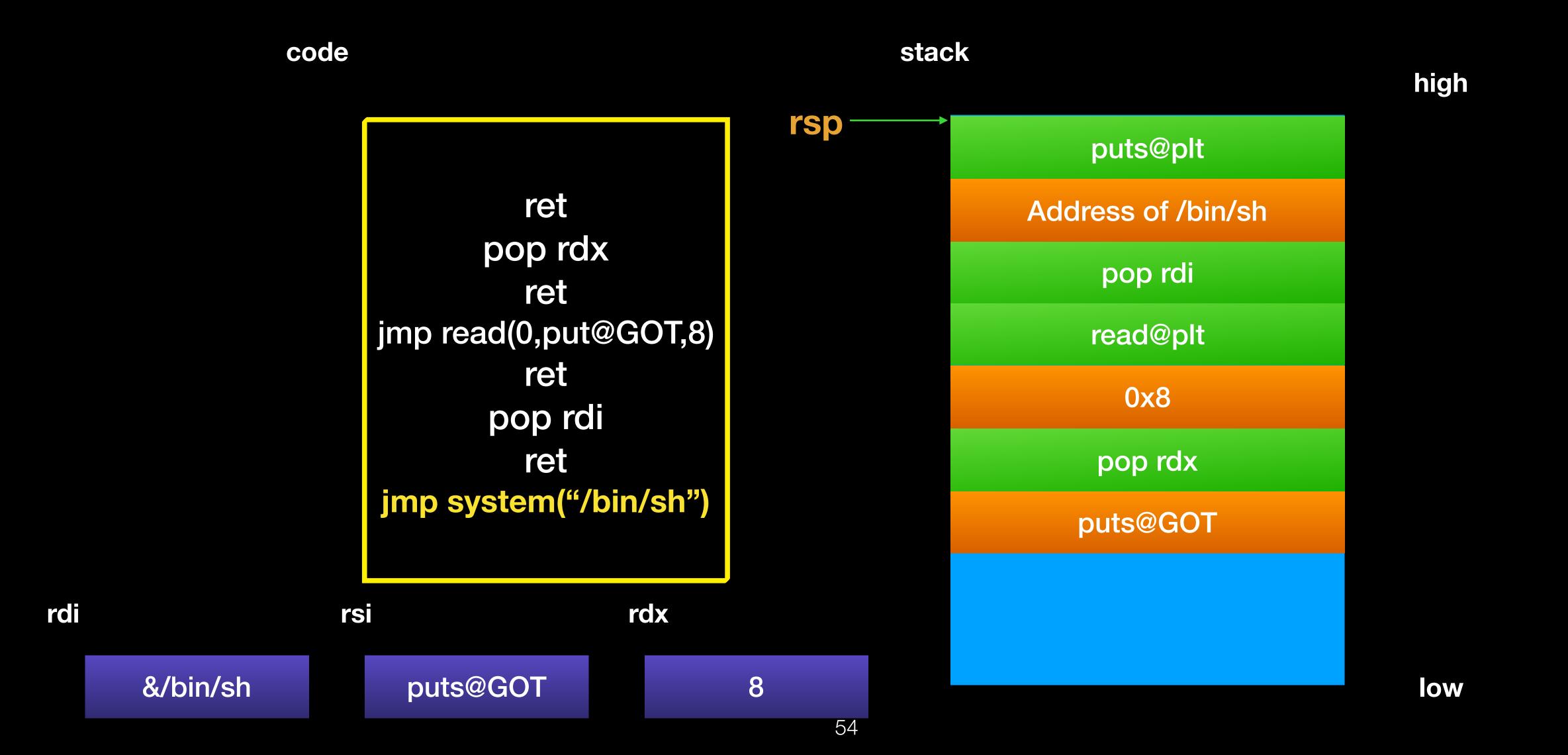


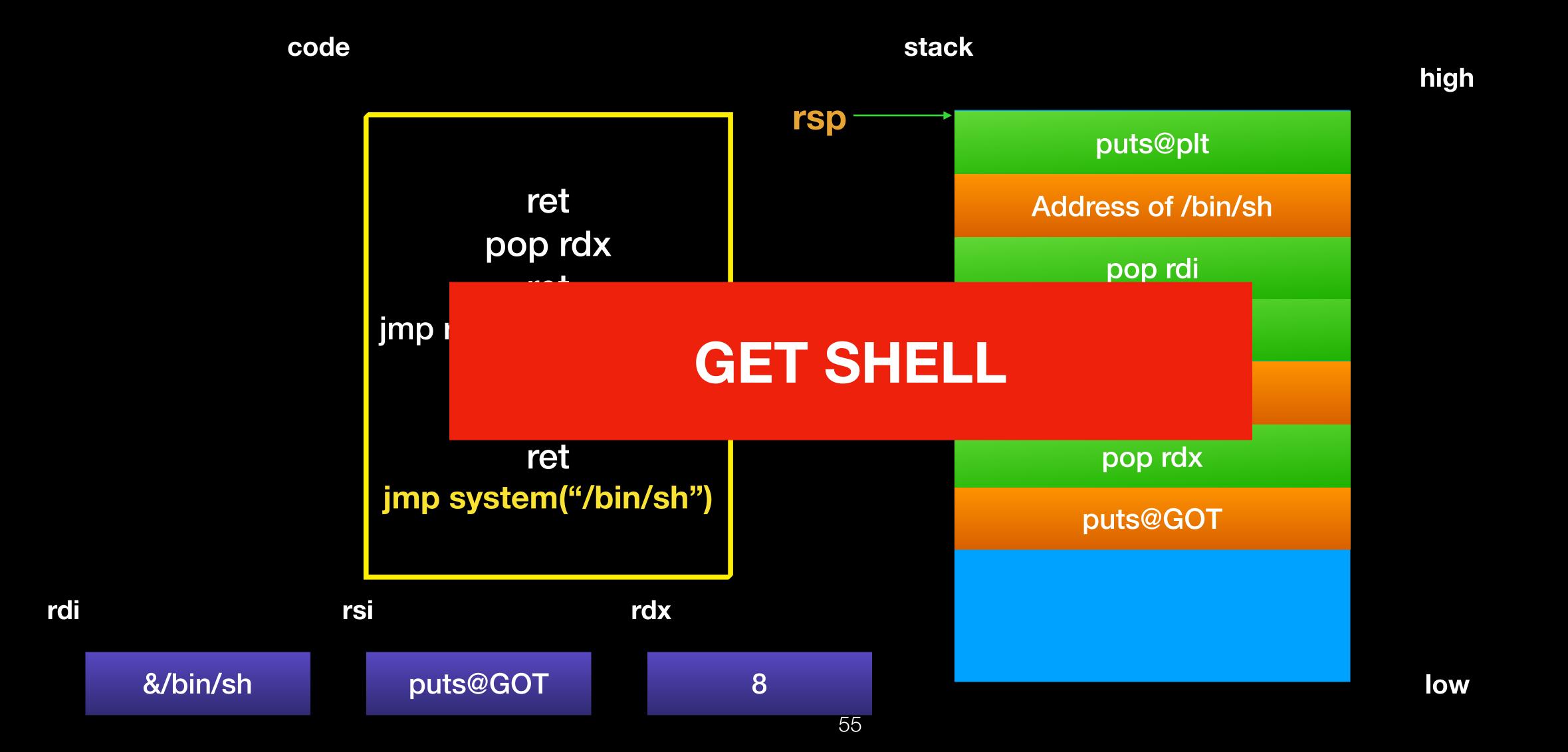












- Bypass PIE
 - 必須比平常多 leak 一個 code 段的位置,藉由這個值算出 code base 進而推出所有 GOT 等資訊
 - 有了 code base 之後其他就跟沒有 PIE 的情況下一樣

- Bypass StackGuard
 - canary 只有在 function return 時做檢查
 - 只檢查 canary 值時否一樣
 - 所以可以先想辦法 leak 出 canary 的值,塞一模一樣的內容就可 bypass,或是想辦法不要改到 canary 也可以

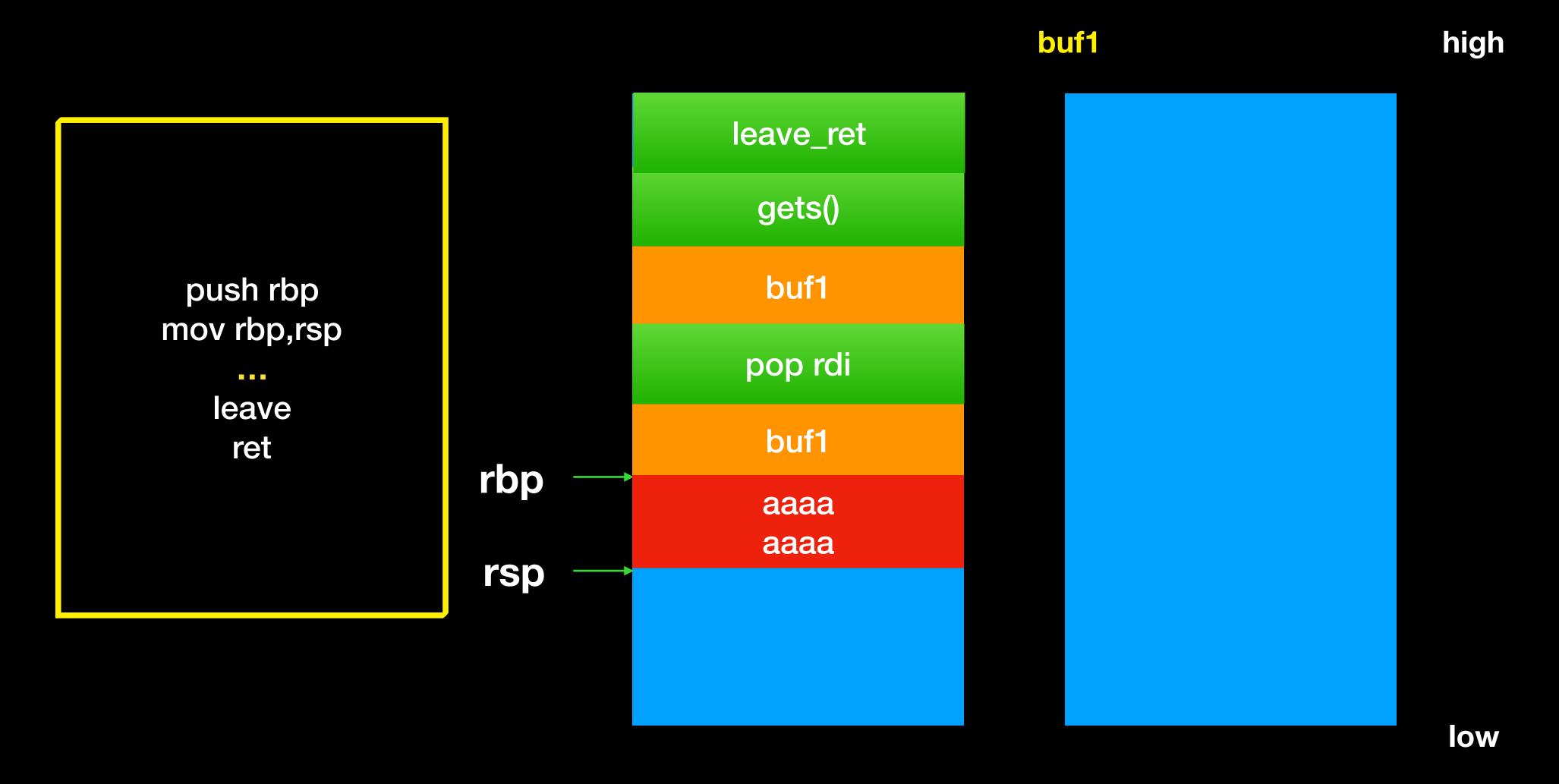
- Weakness in fork
 - canary and memory mappings are same as parent.

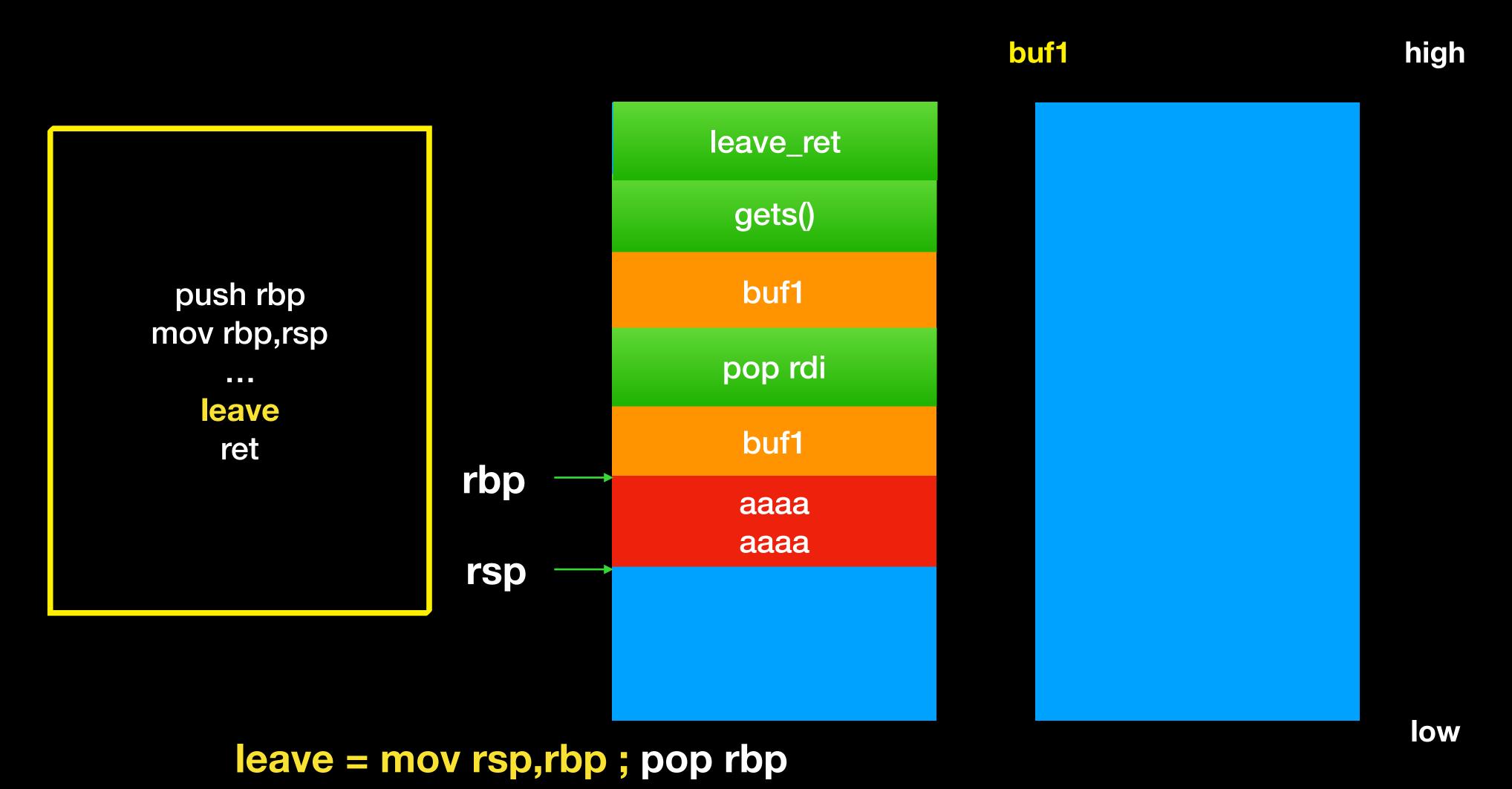
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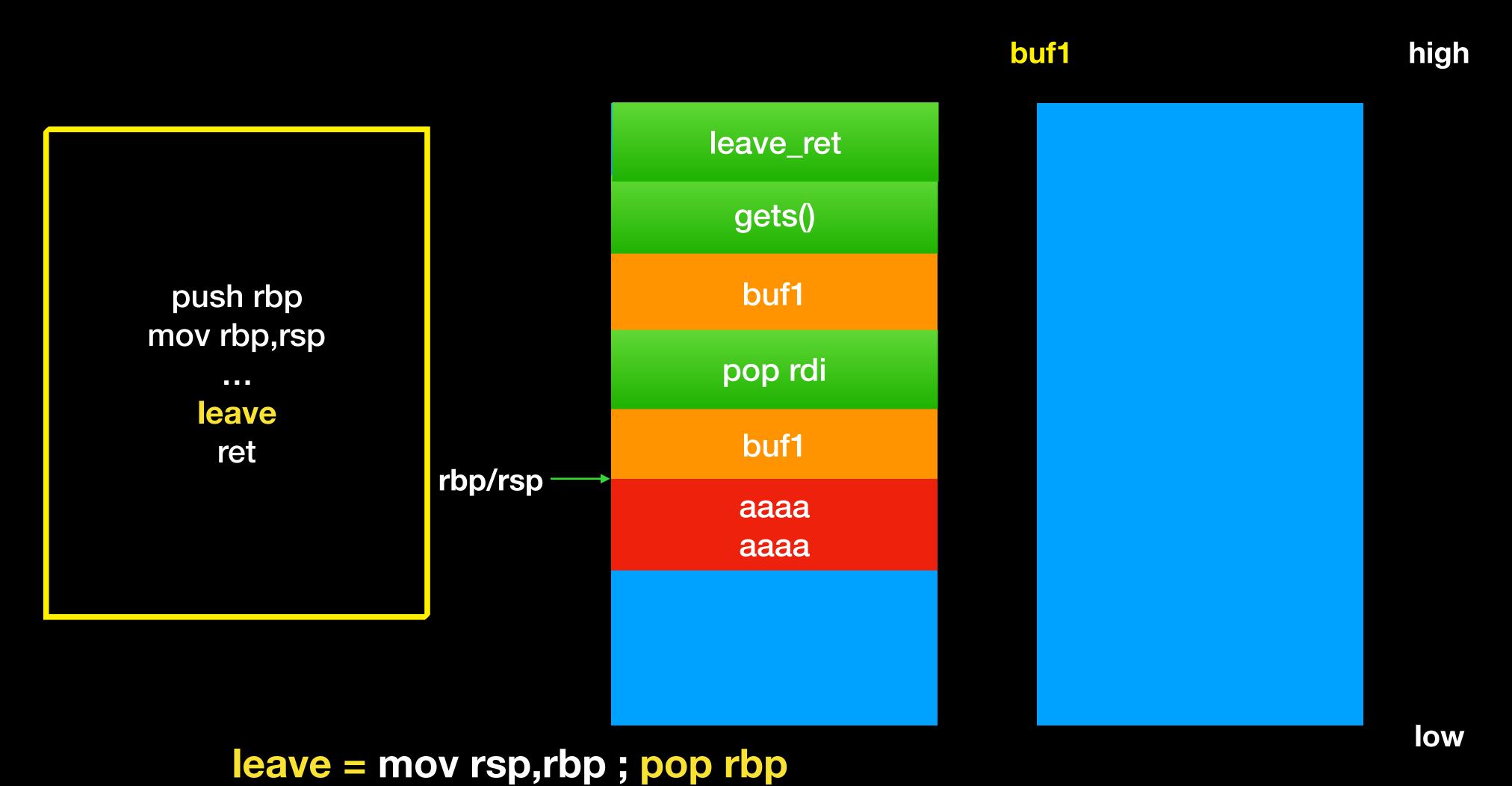
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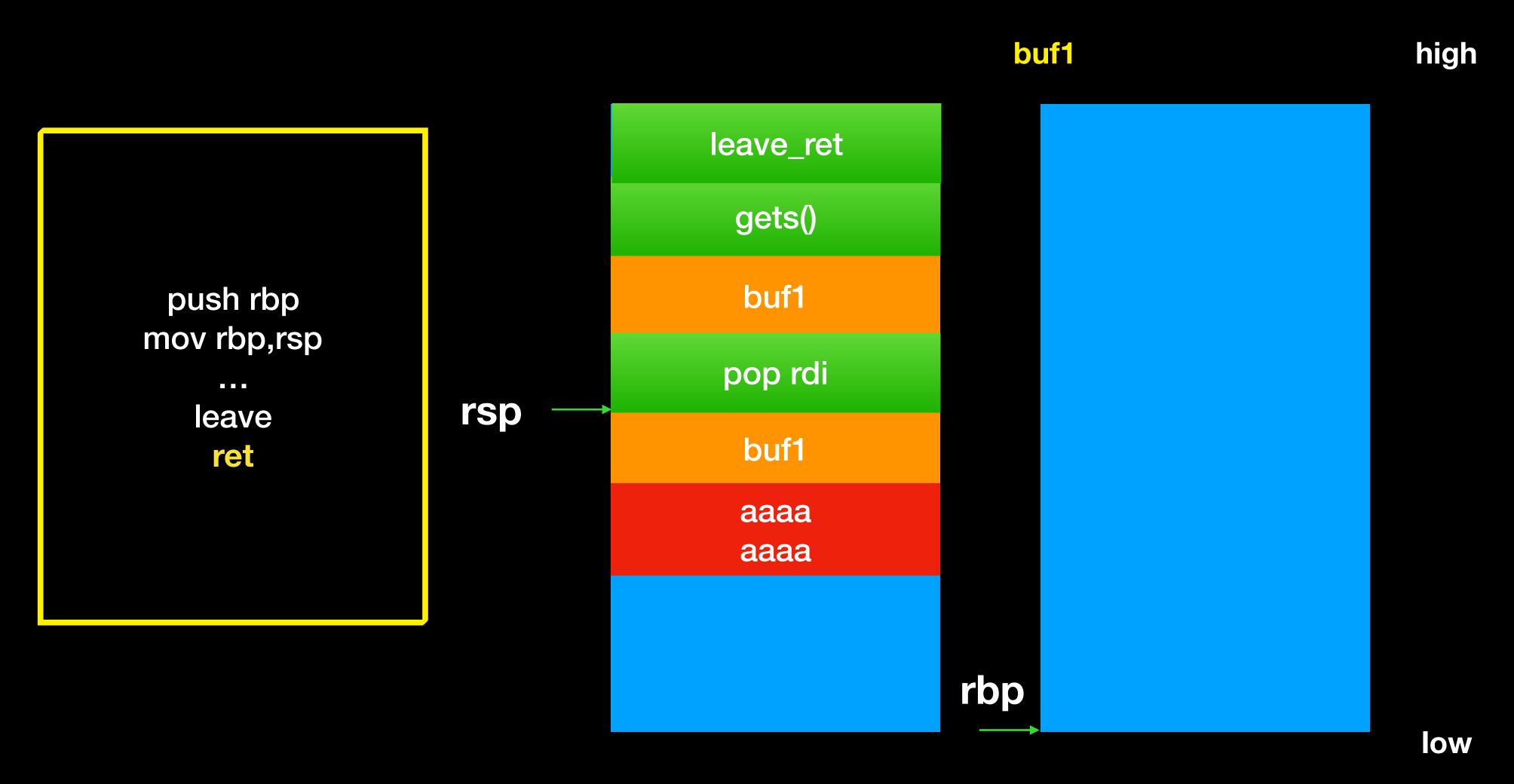
- 將 ROP Chain 寫在已知固定位置上
- 再利用 leave 搬移 Stack 位置到已知位置
- 可無限接 ROP Chain
- 必須注意到 Migration 之後 stack 要留大一點,有些 function 可能會需要很大的 stack frame,太小可能會存取到唯獨區域,導致 Segmentation Fault

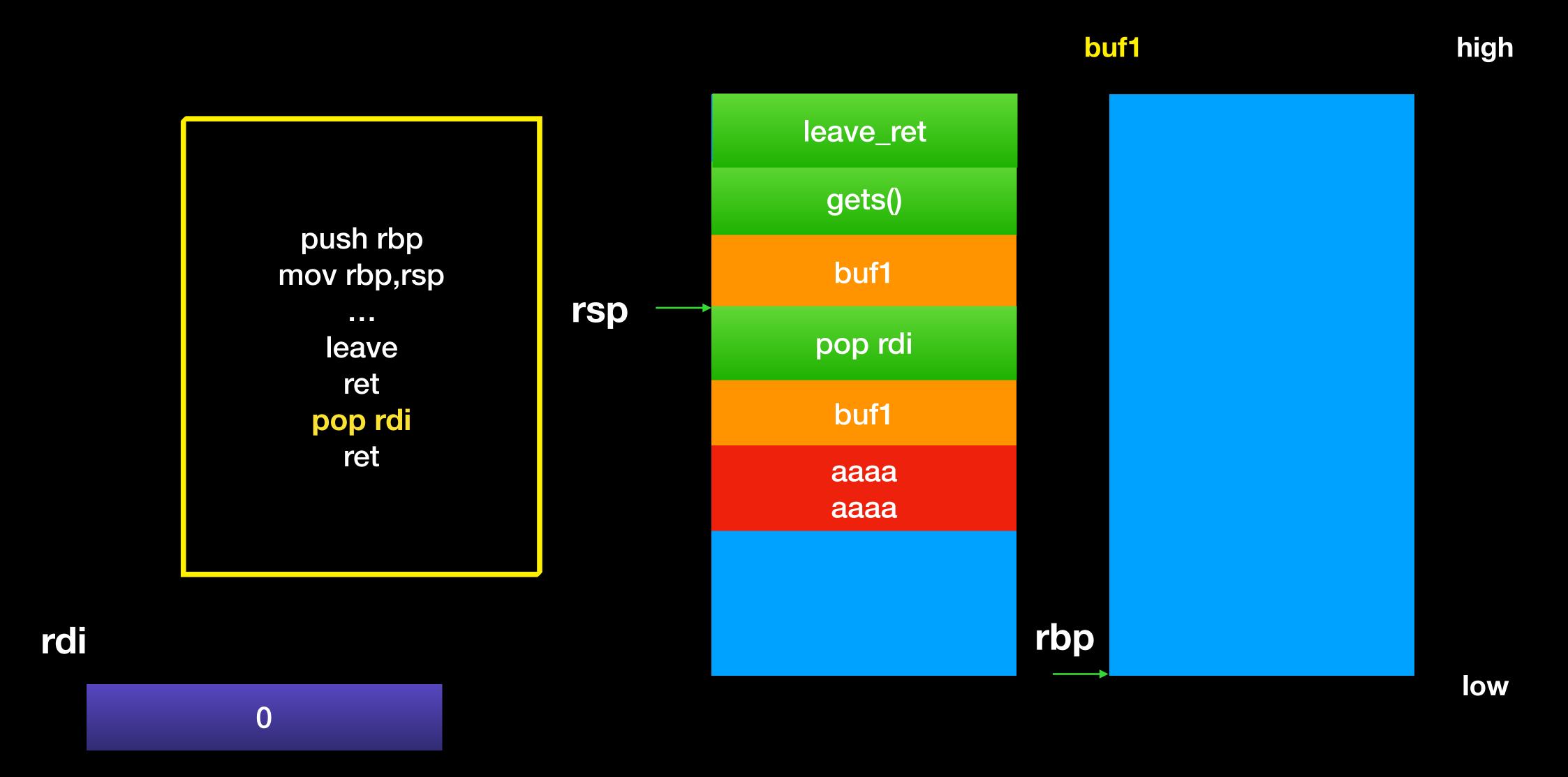
high return address rbp rbp/rsp push rbp mov rbp,rsp leave ret low

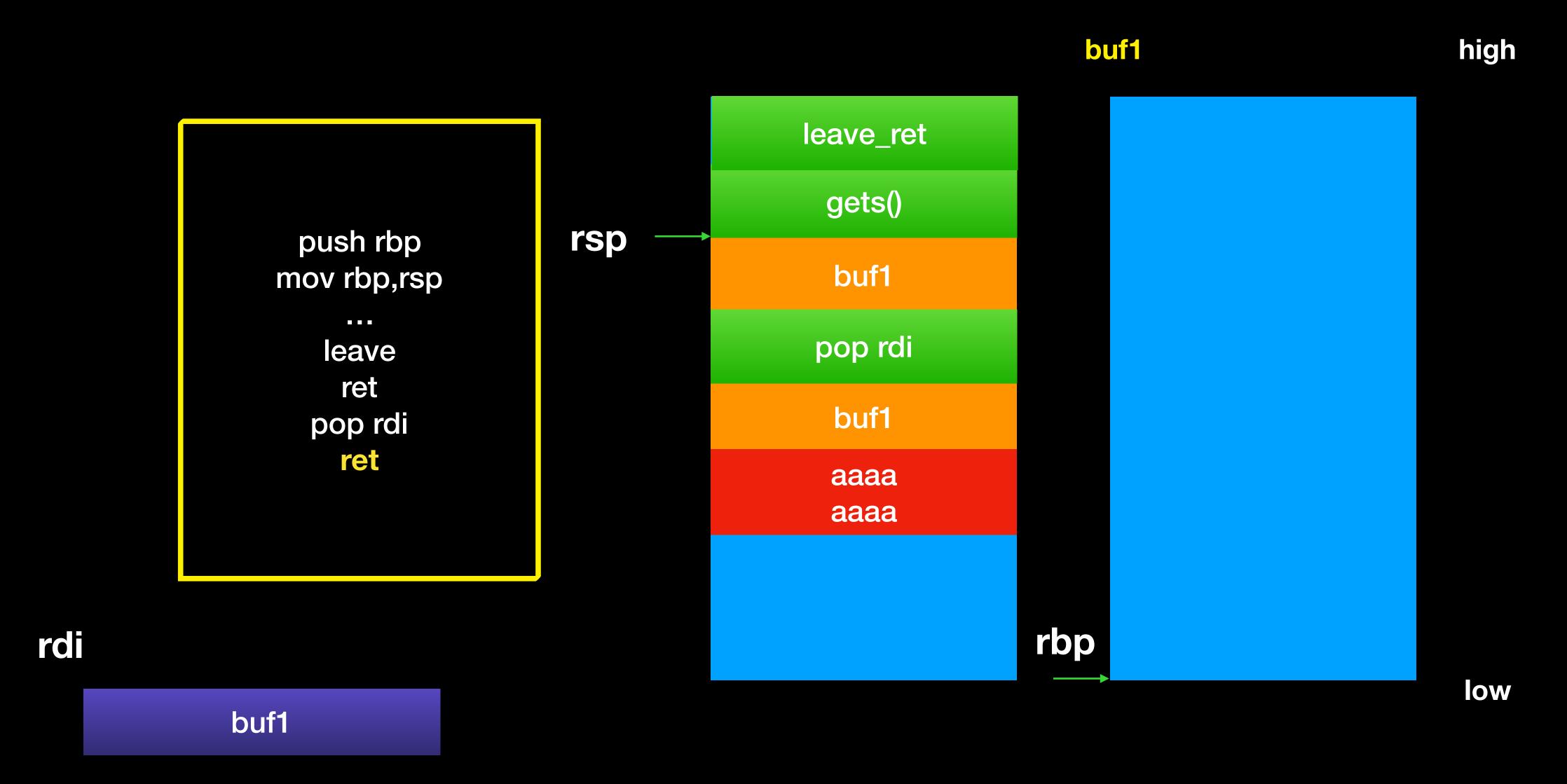


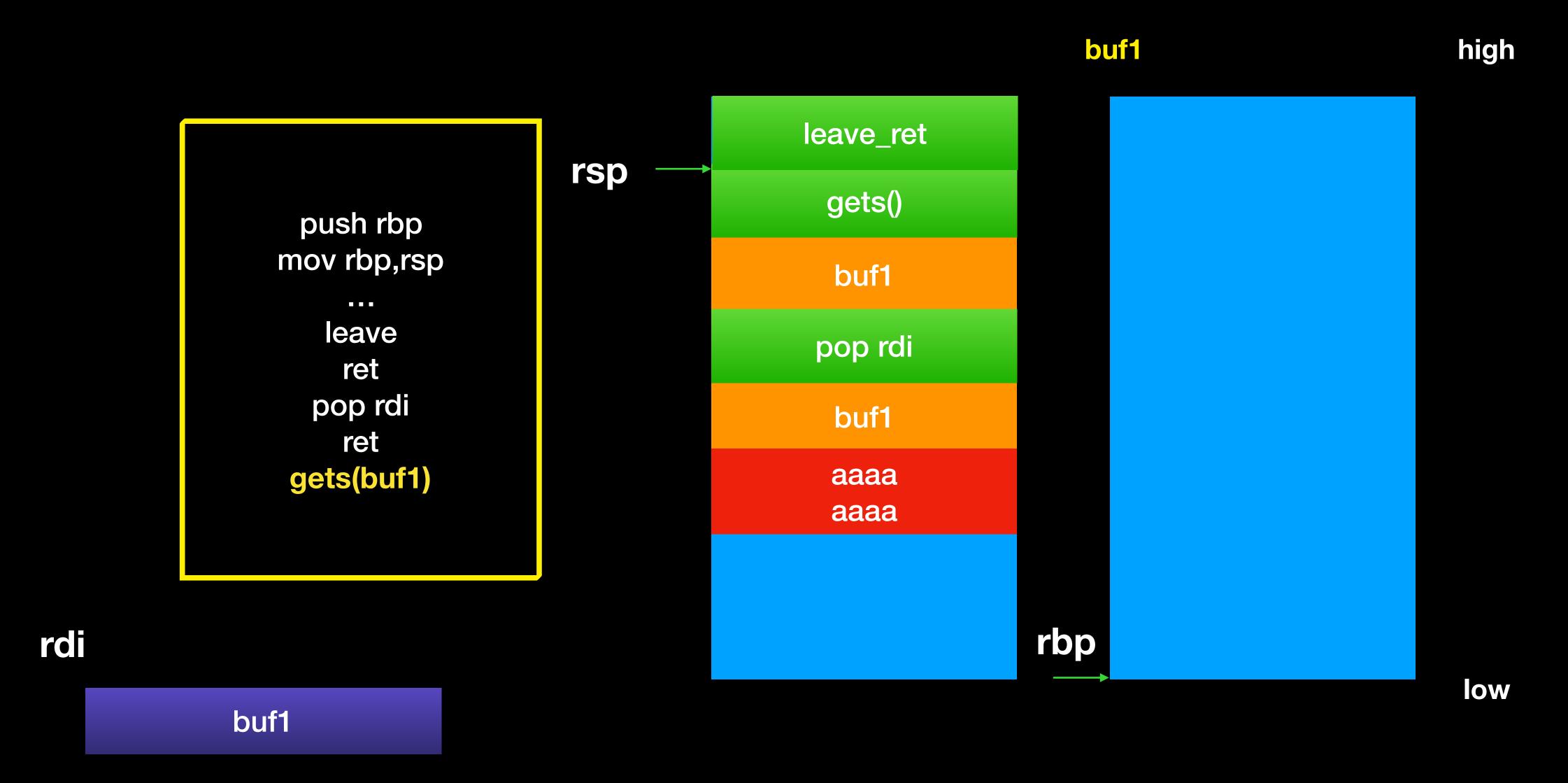


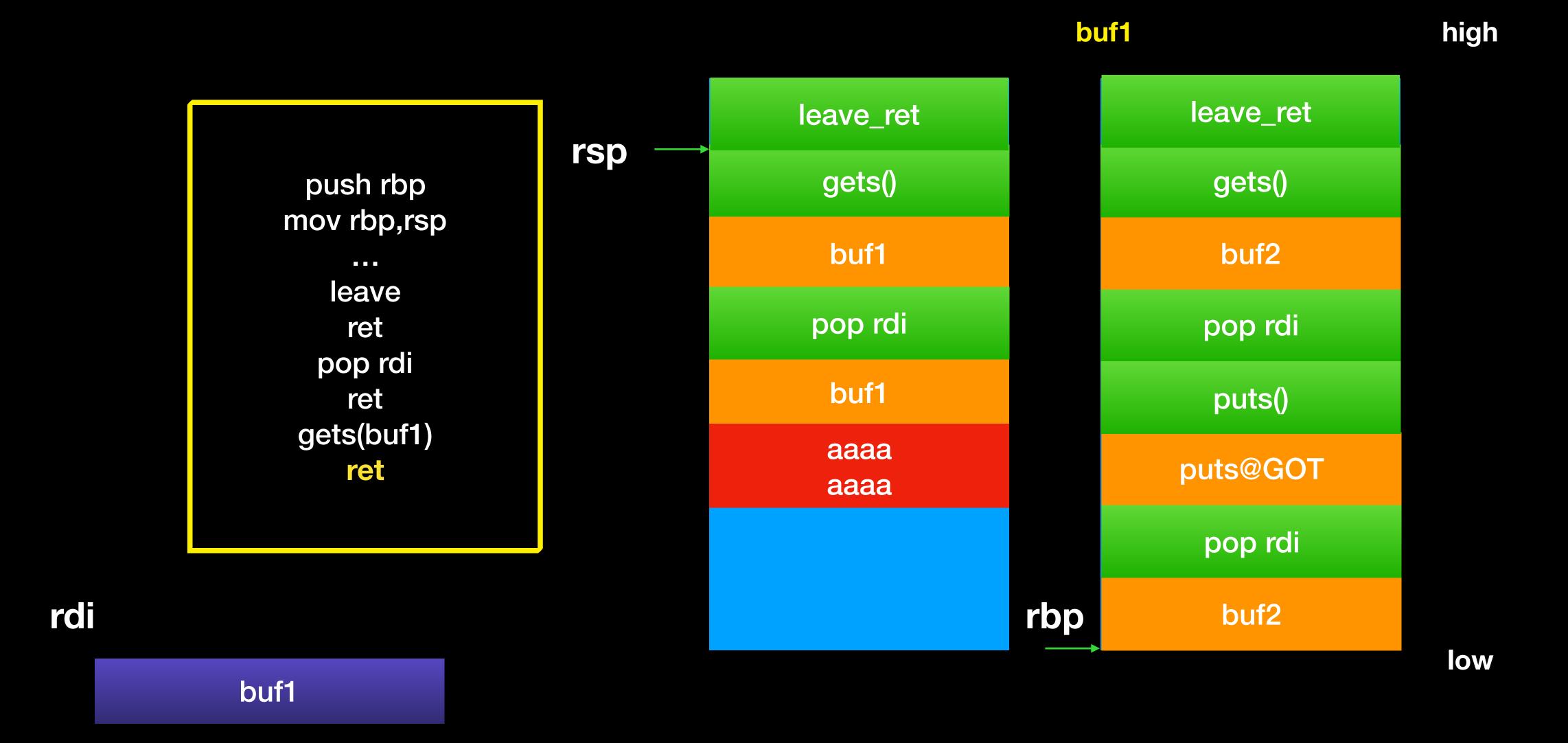


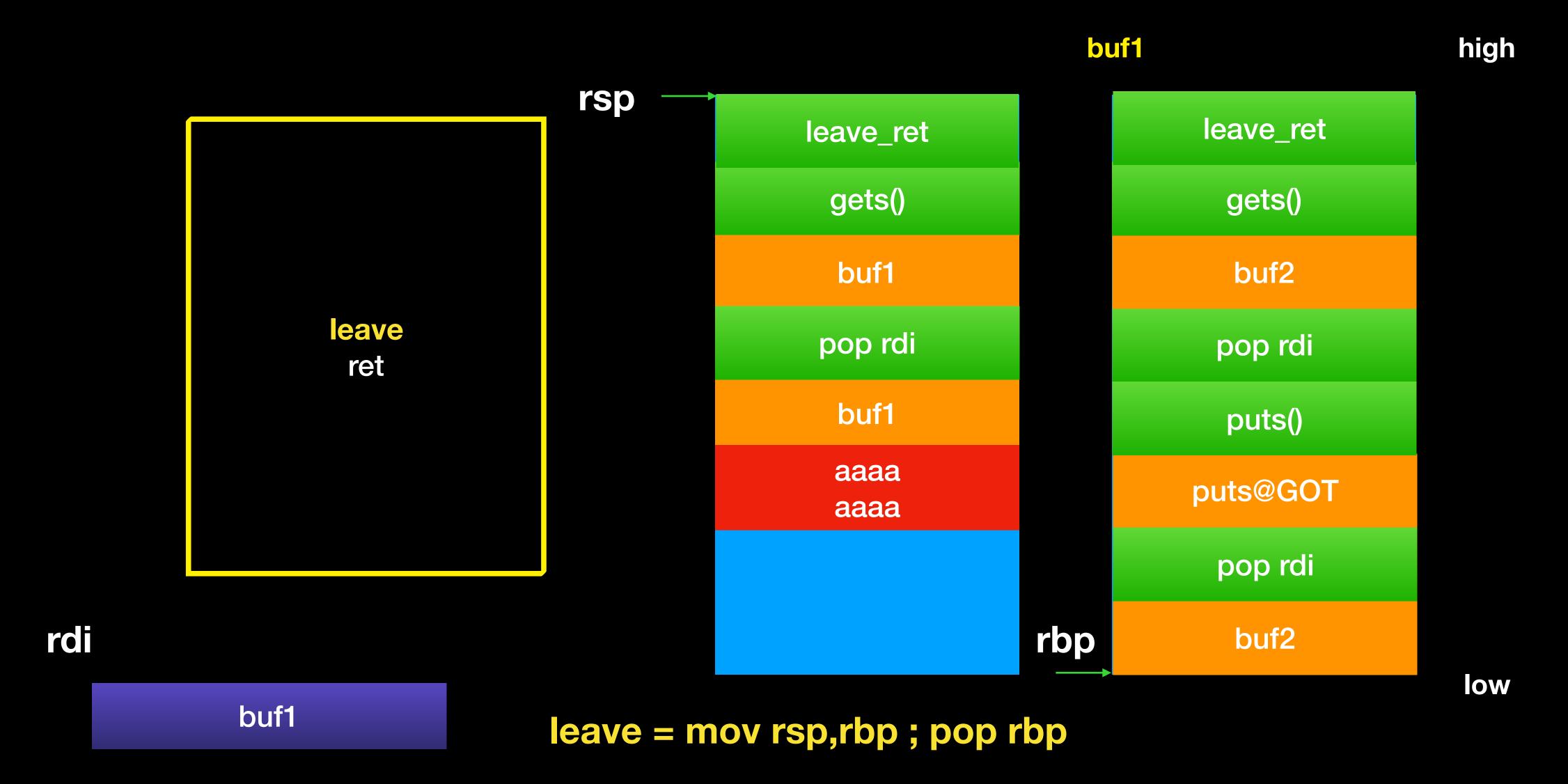


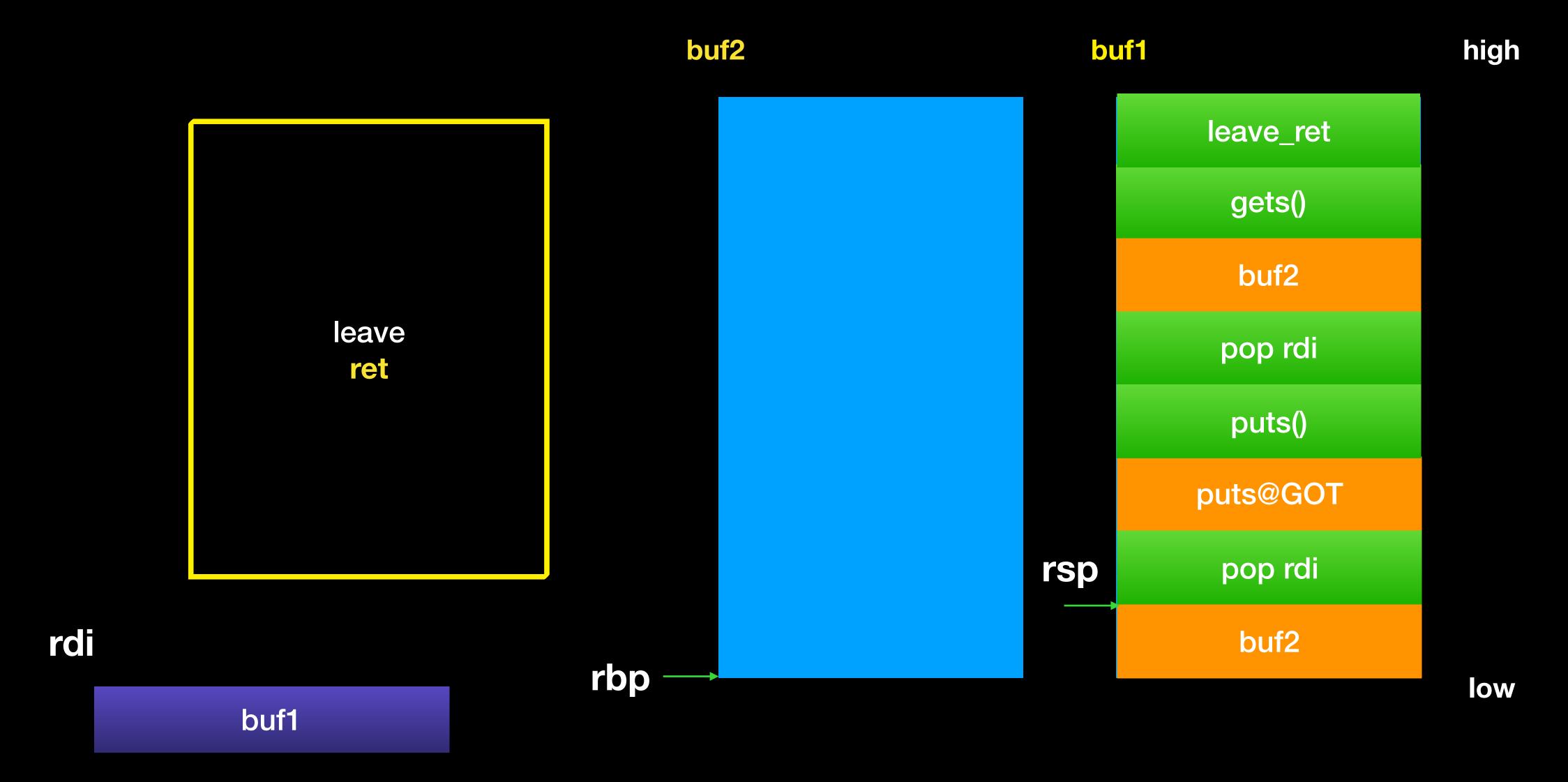


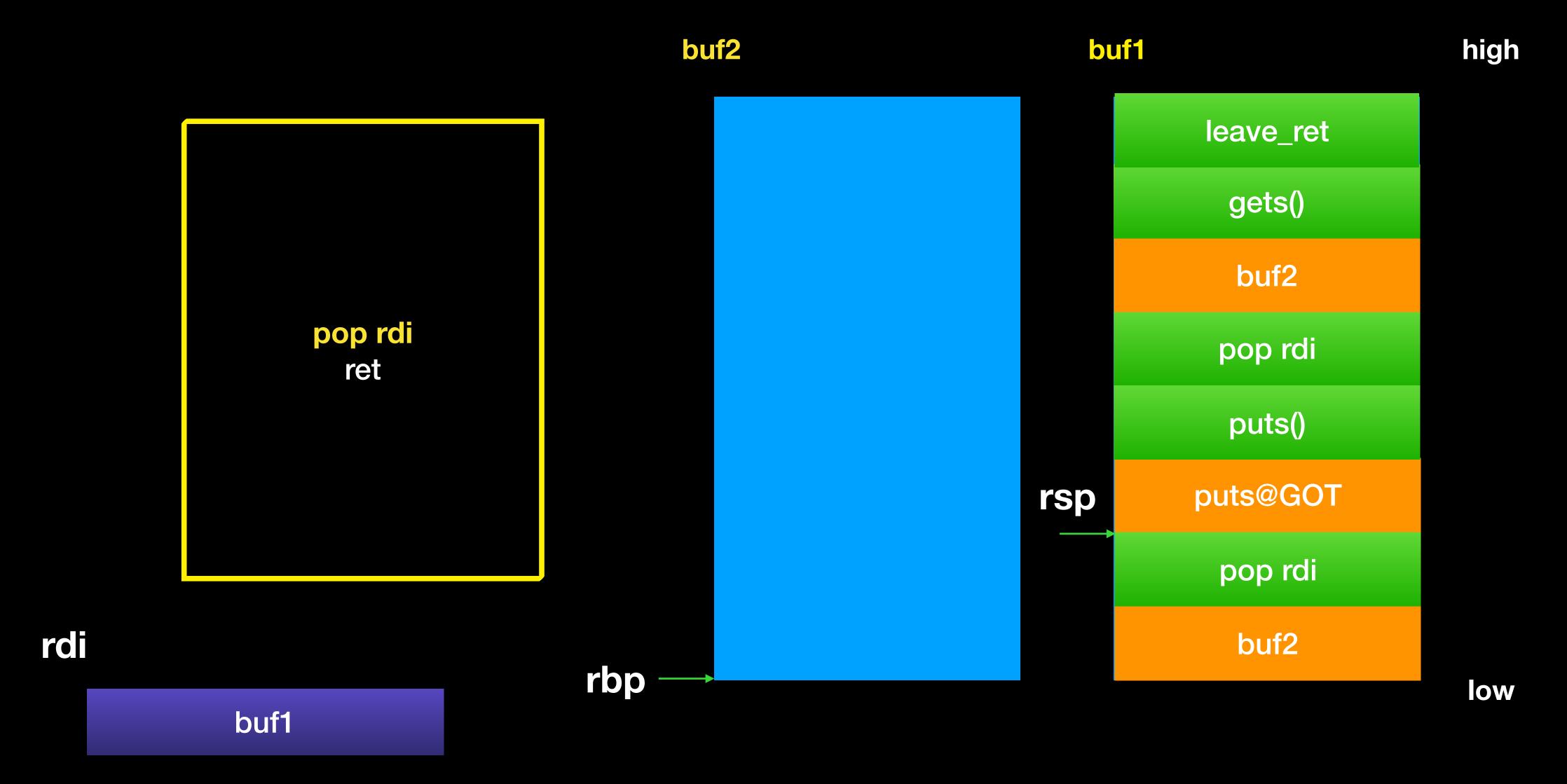


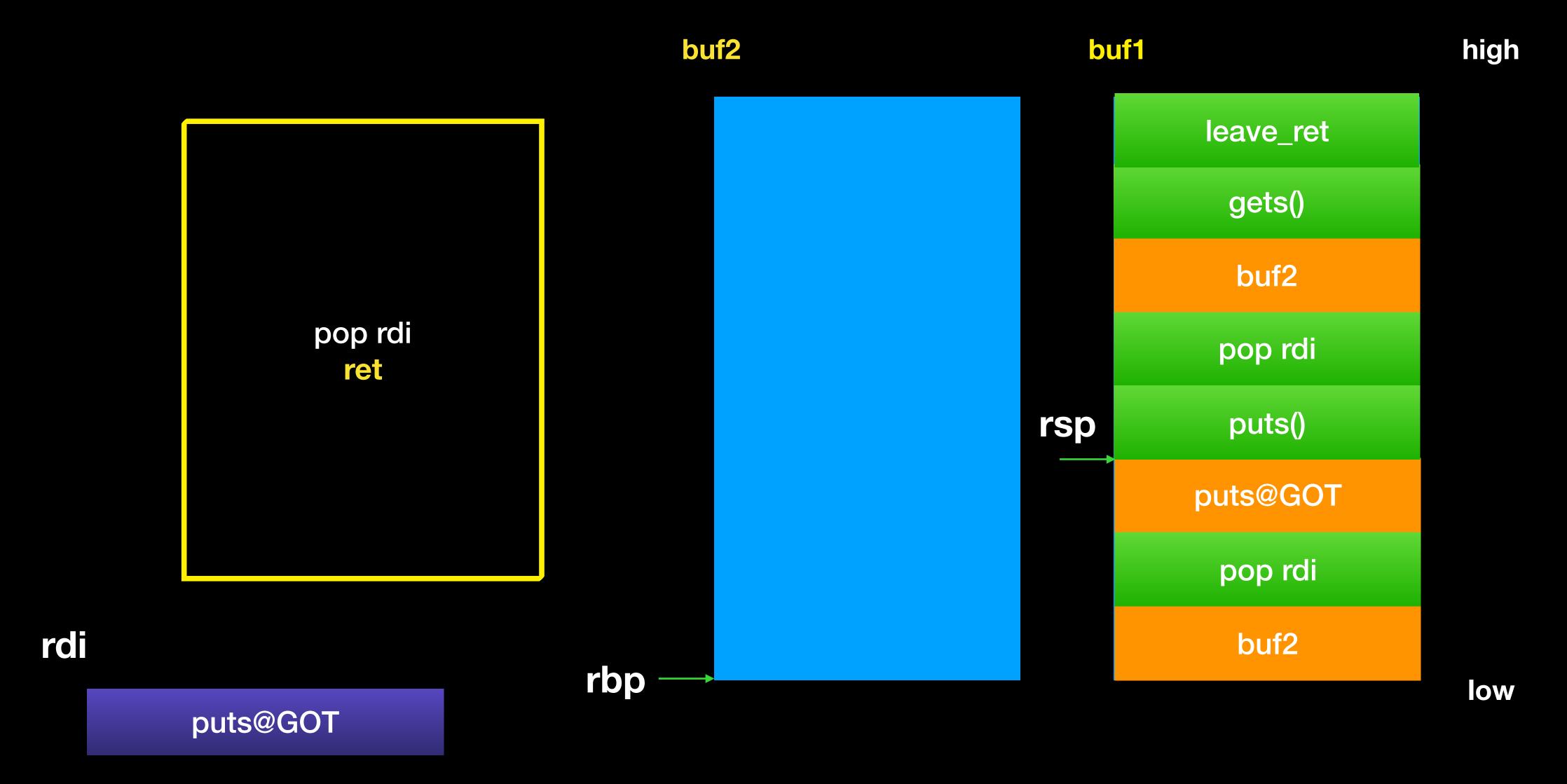


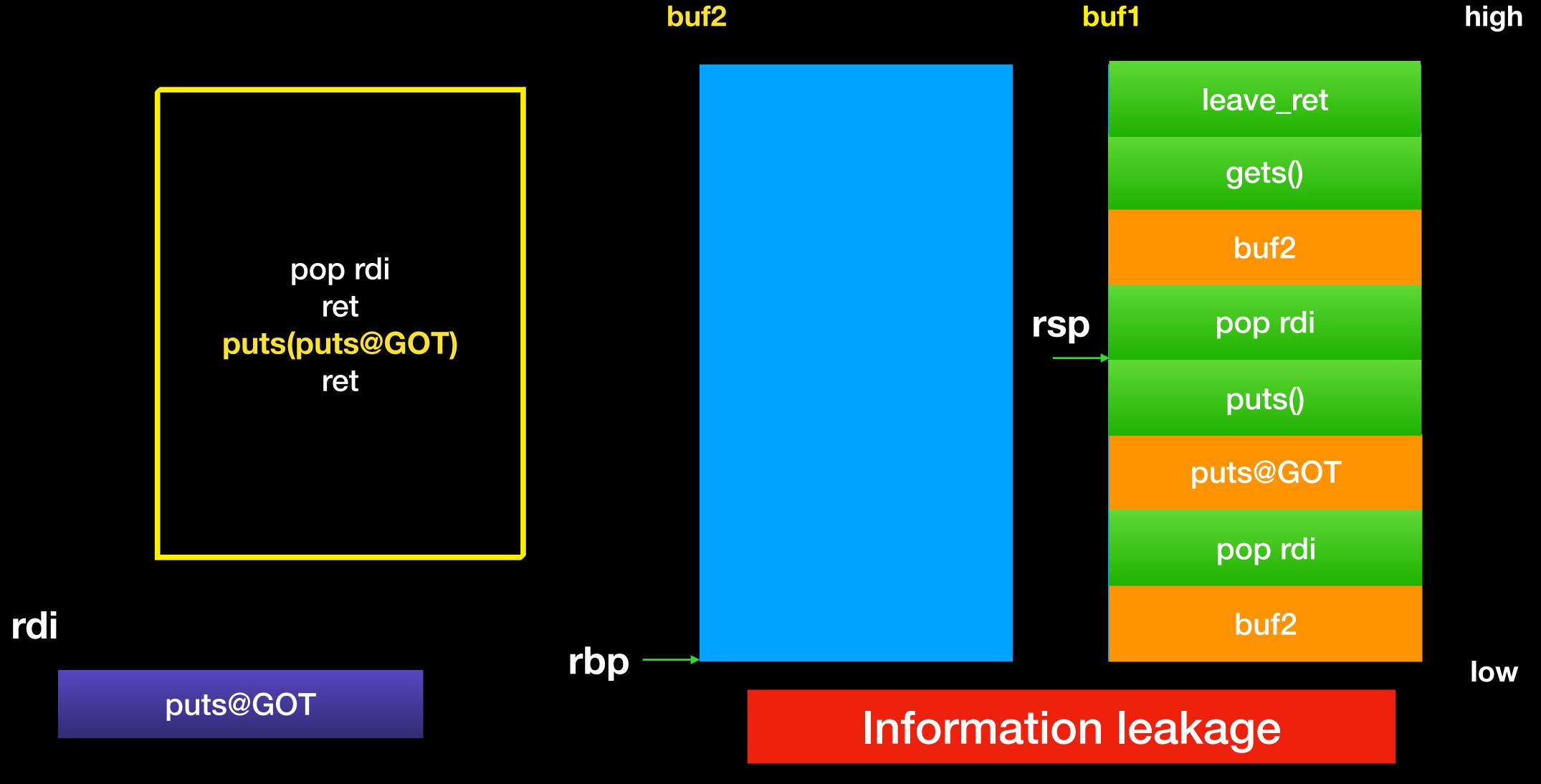


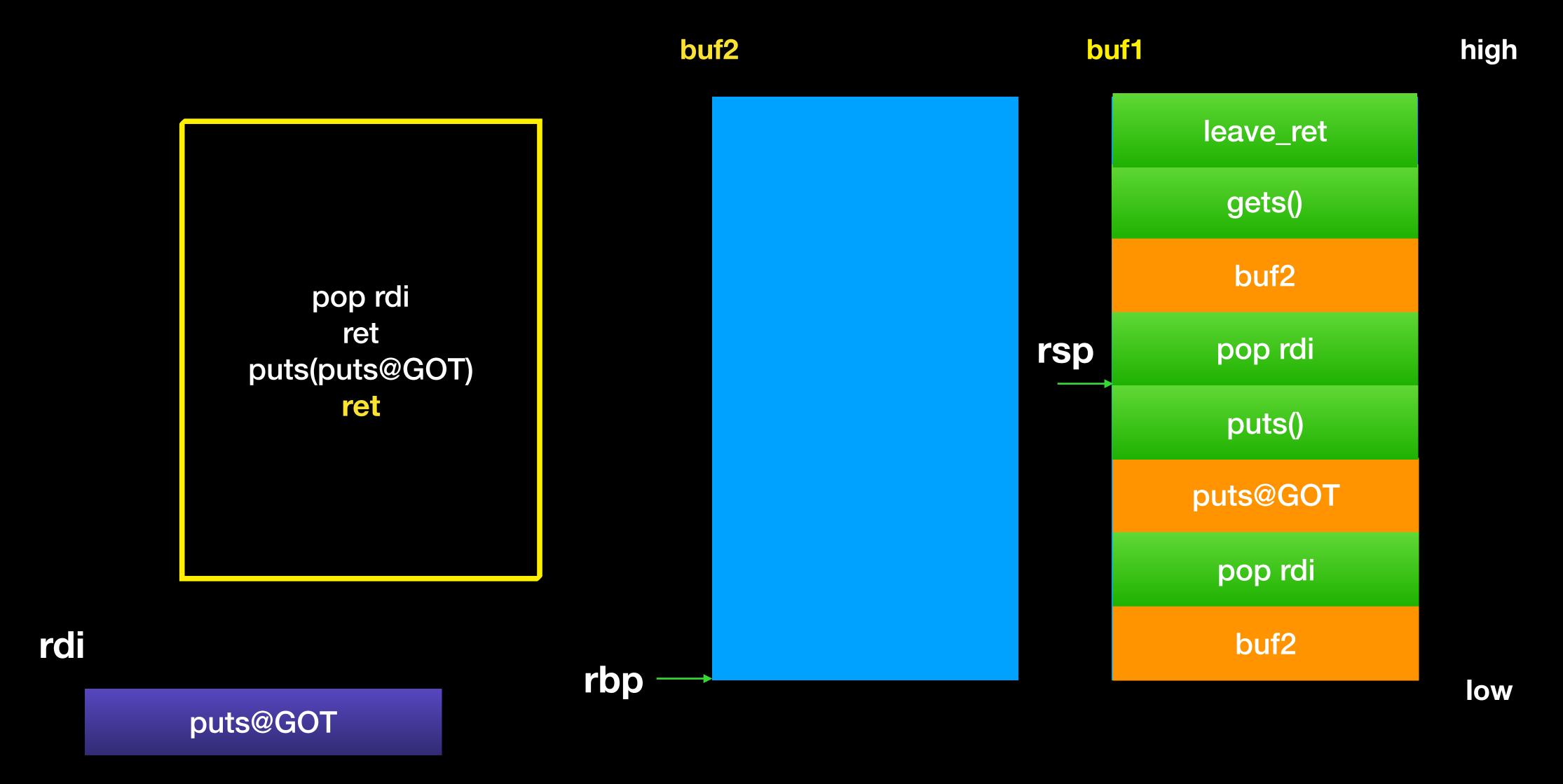


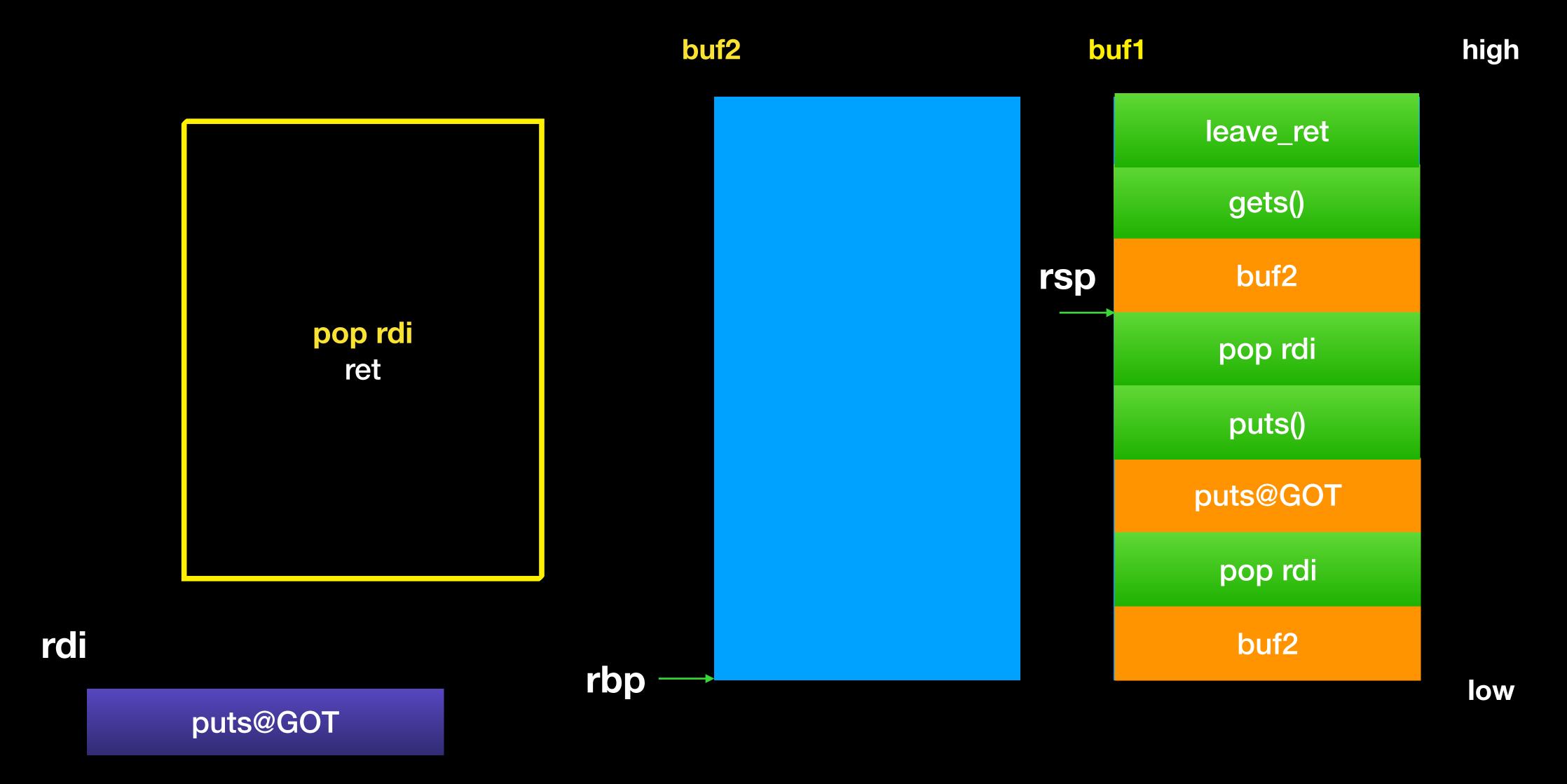


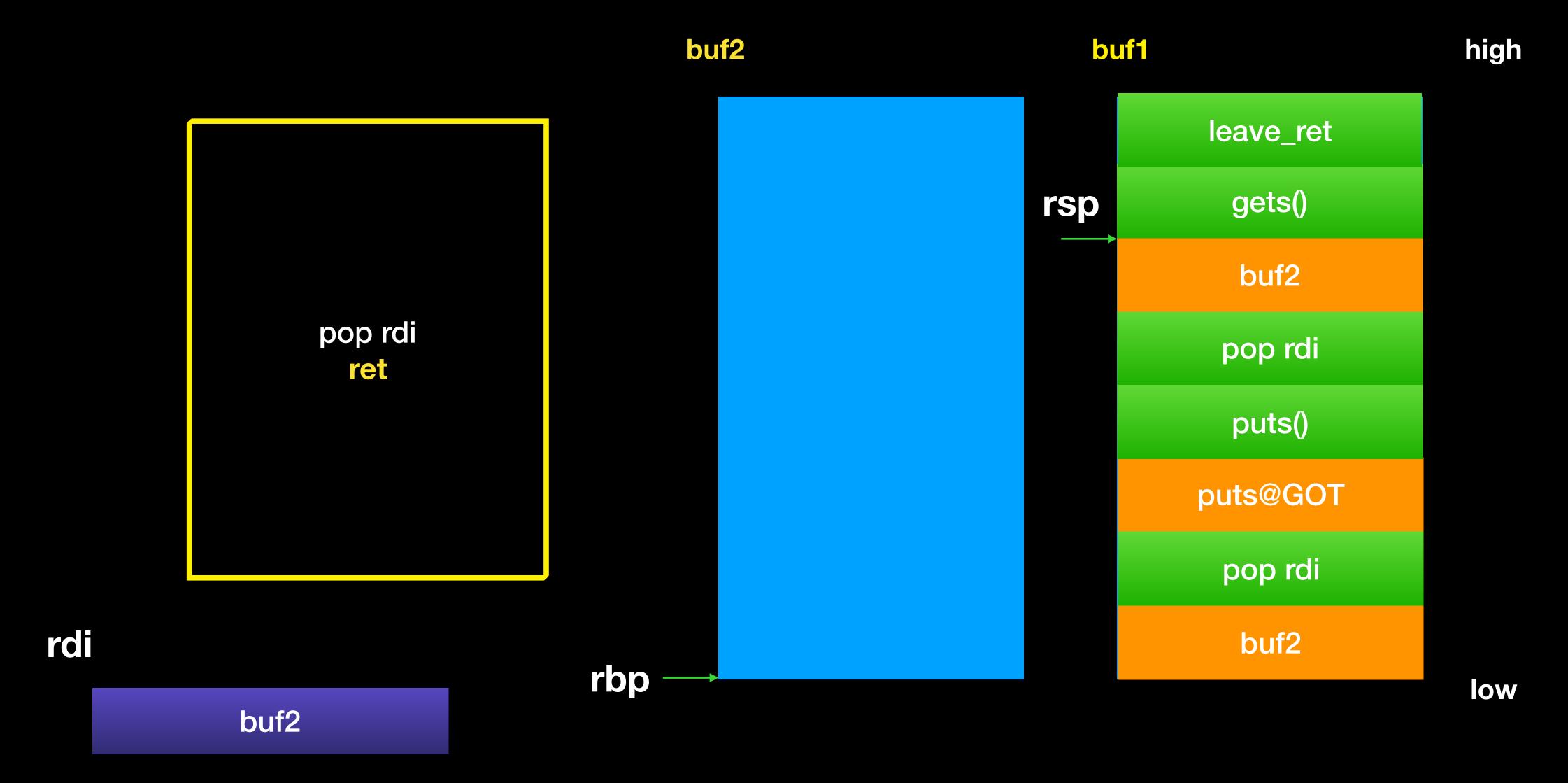


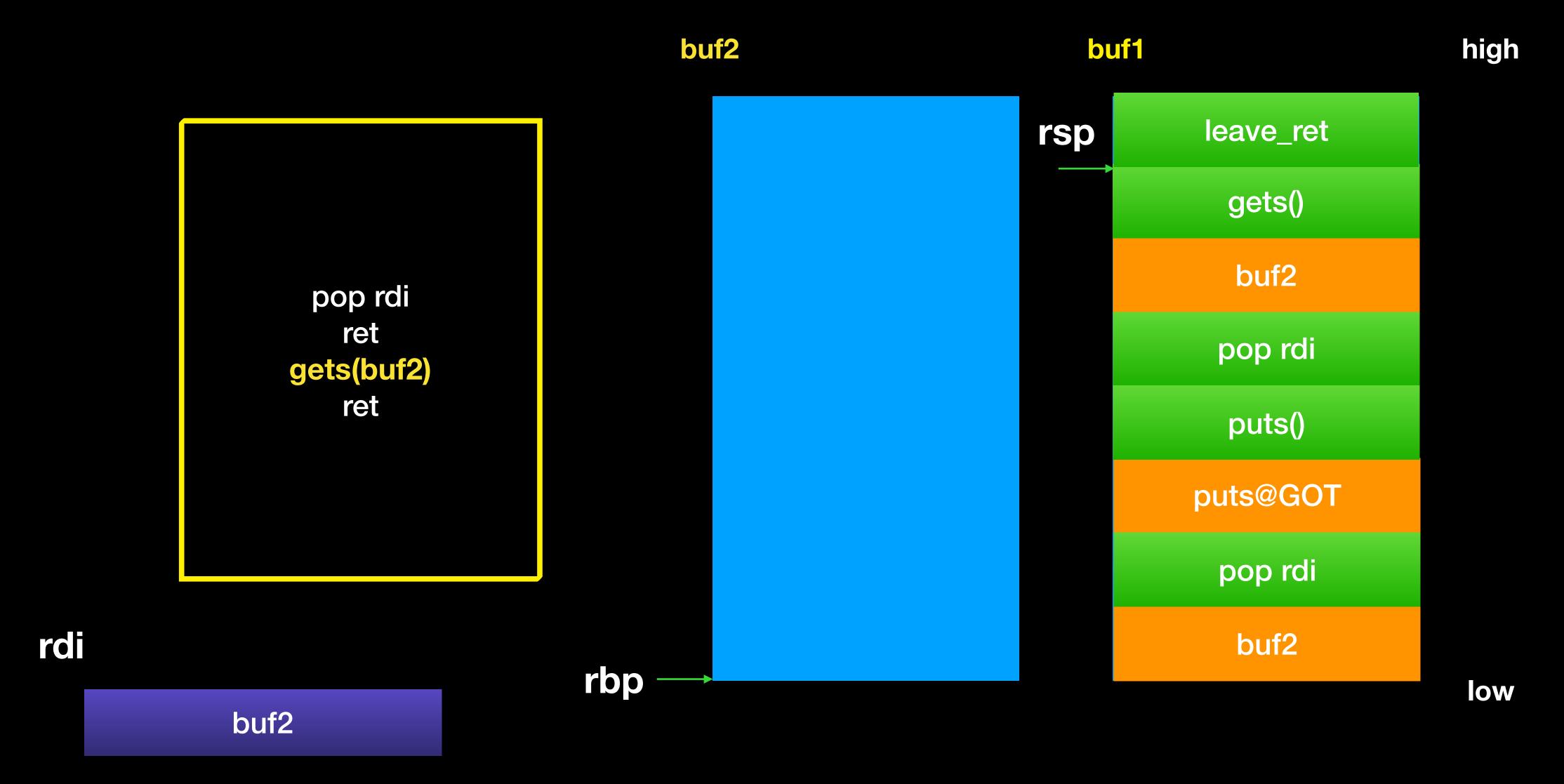


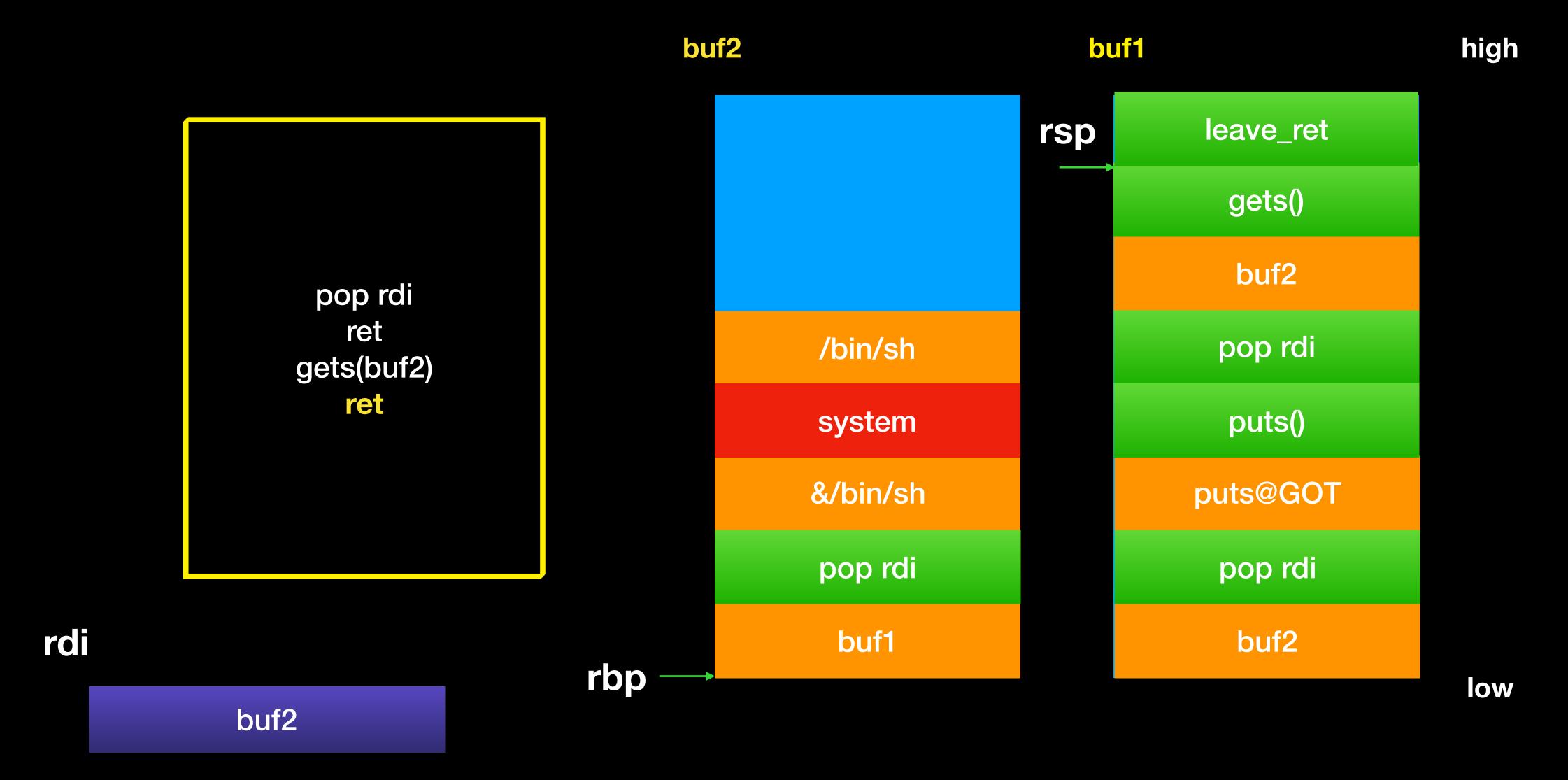


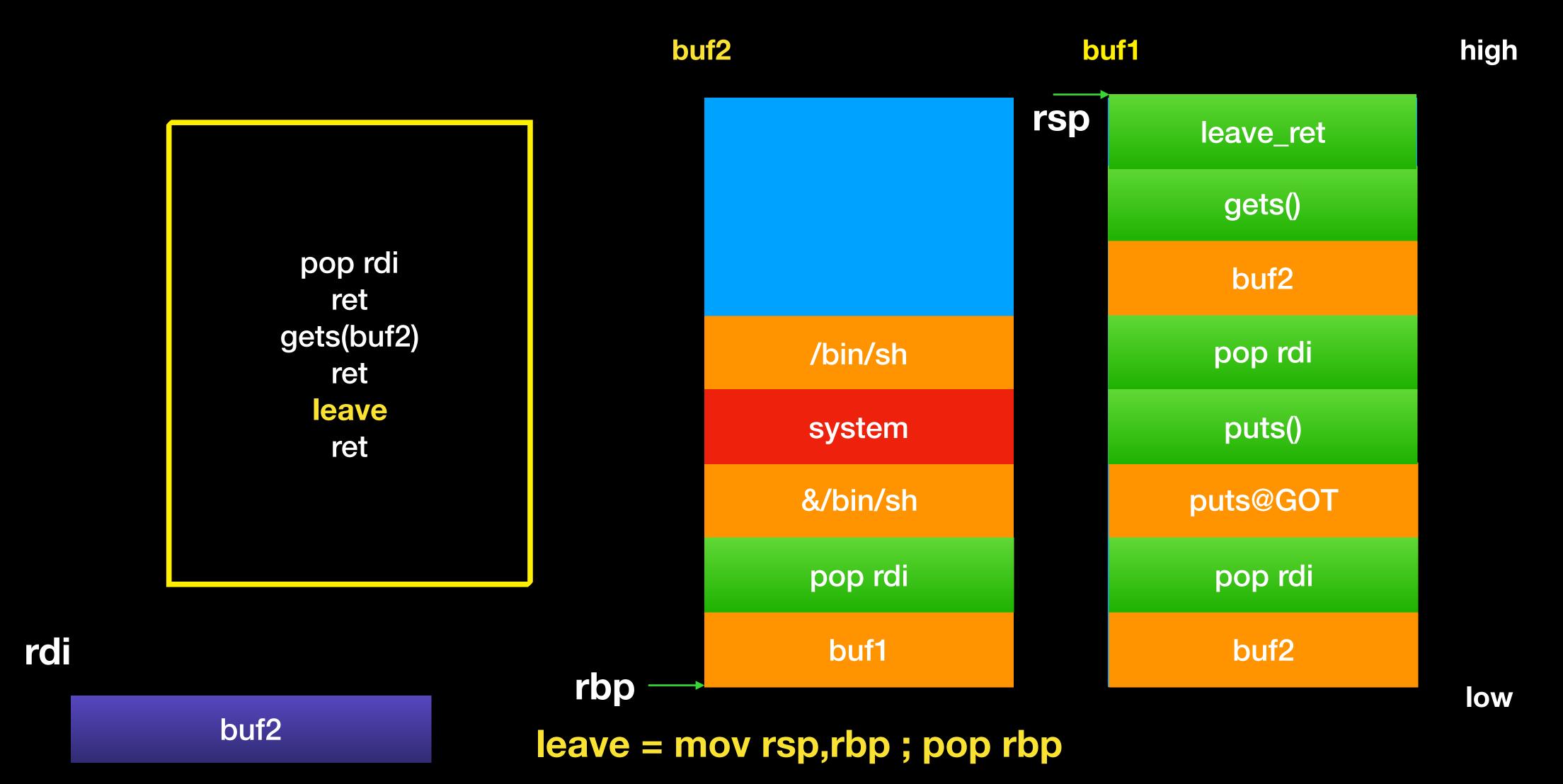


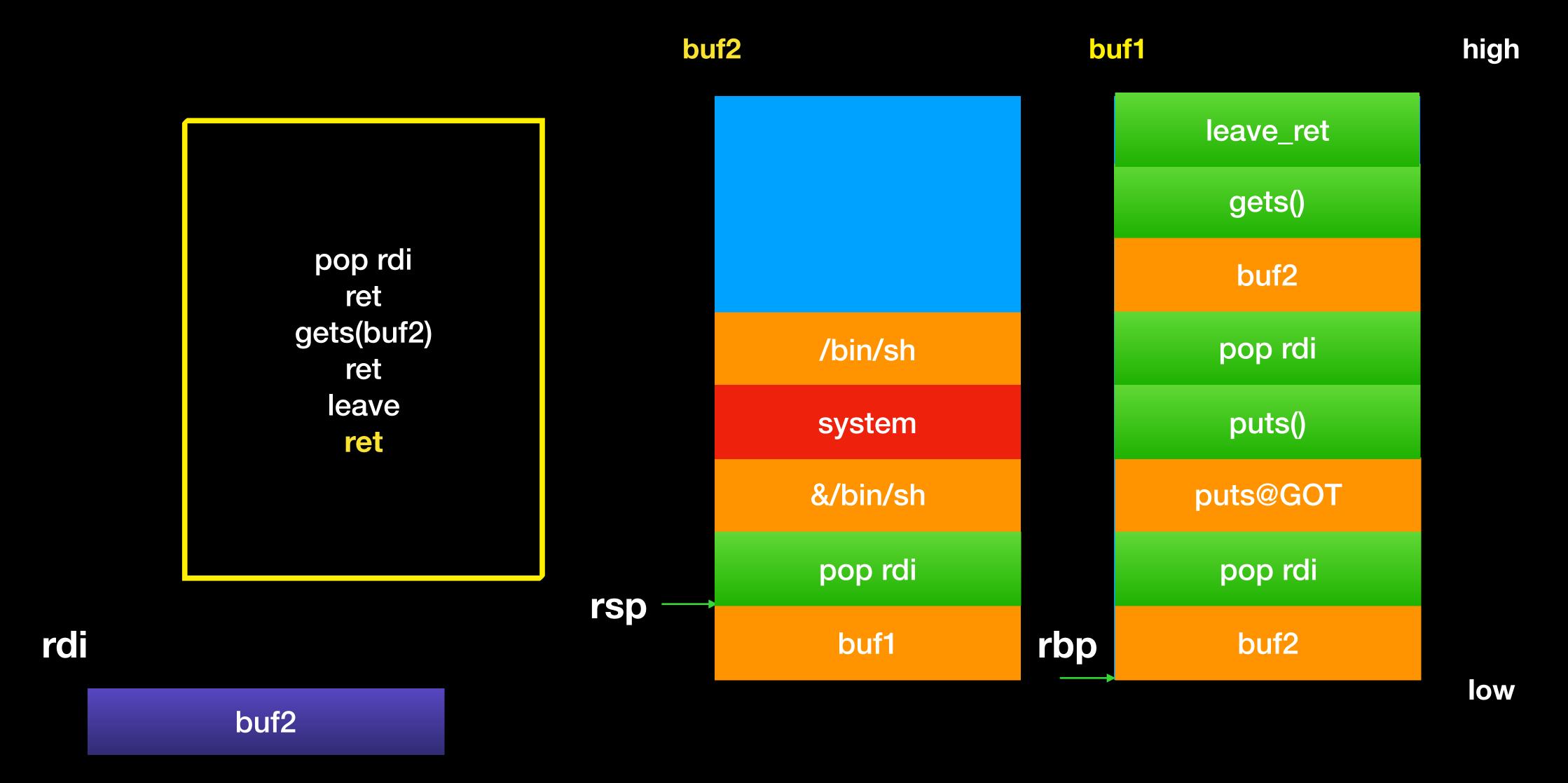


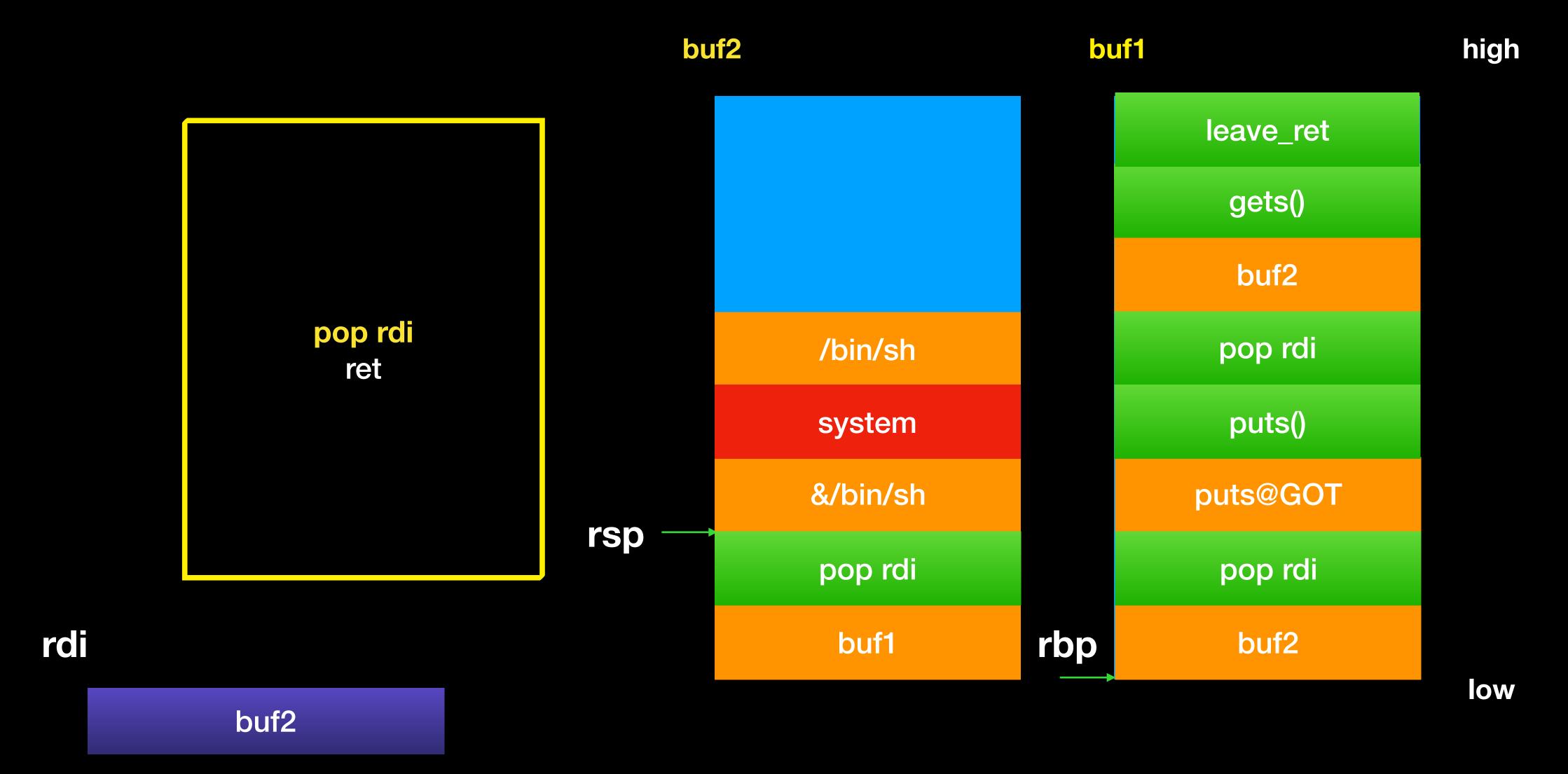


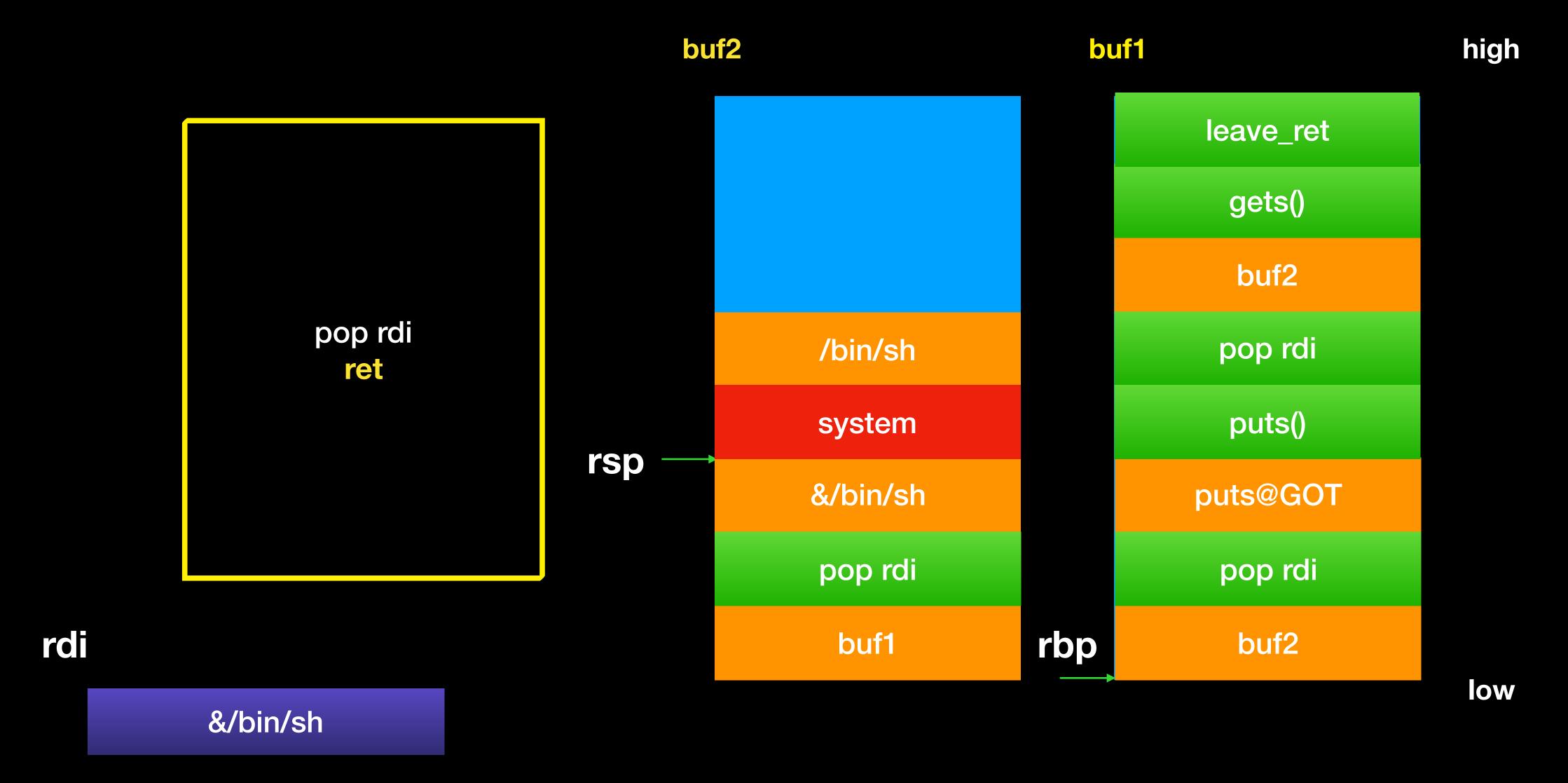


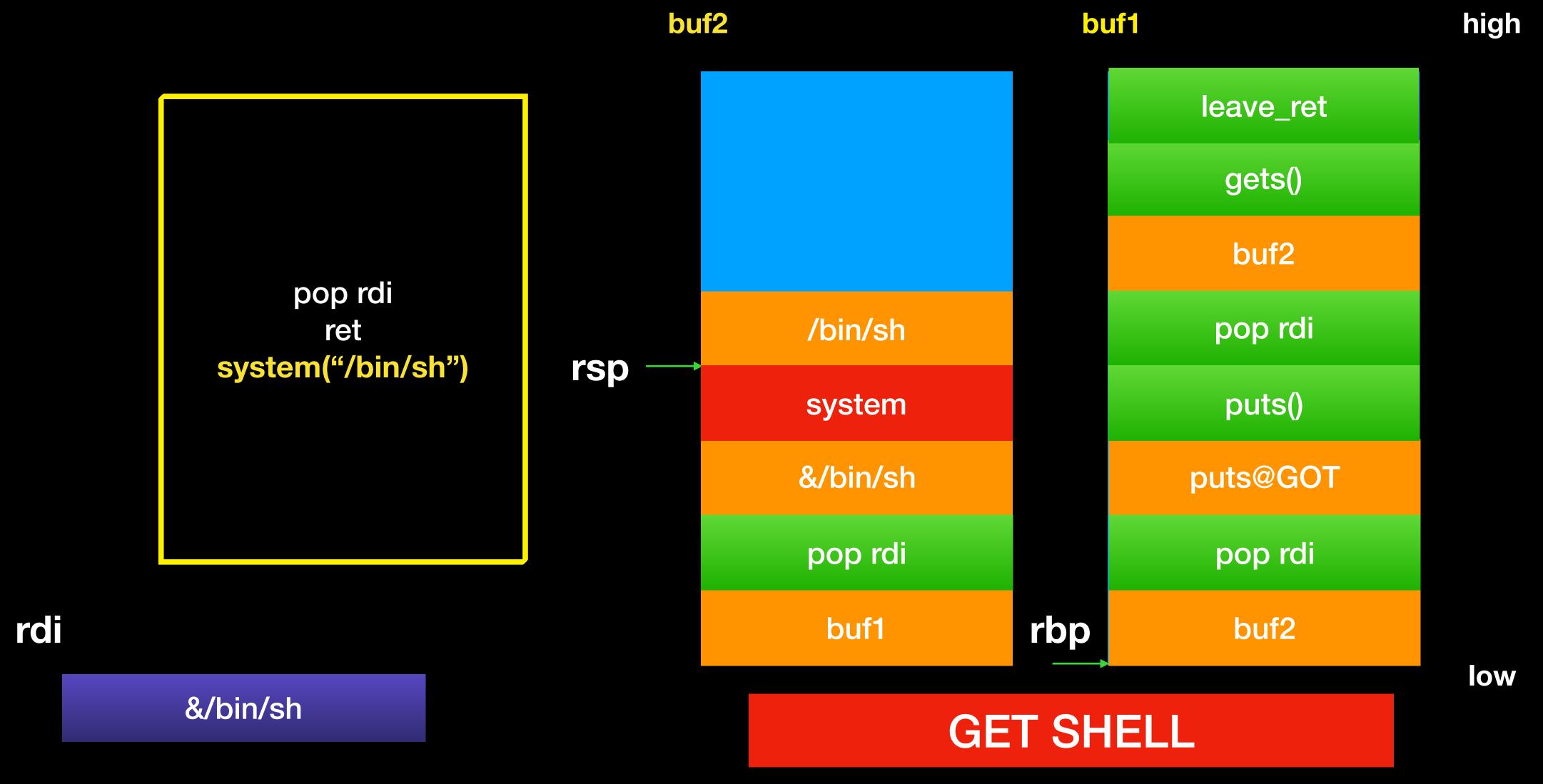






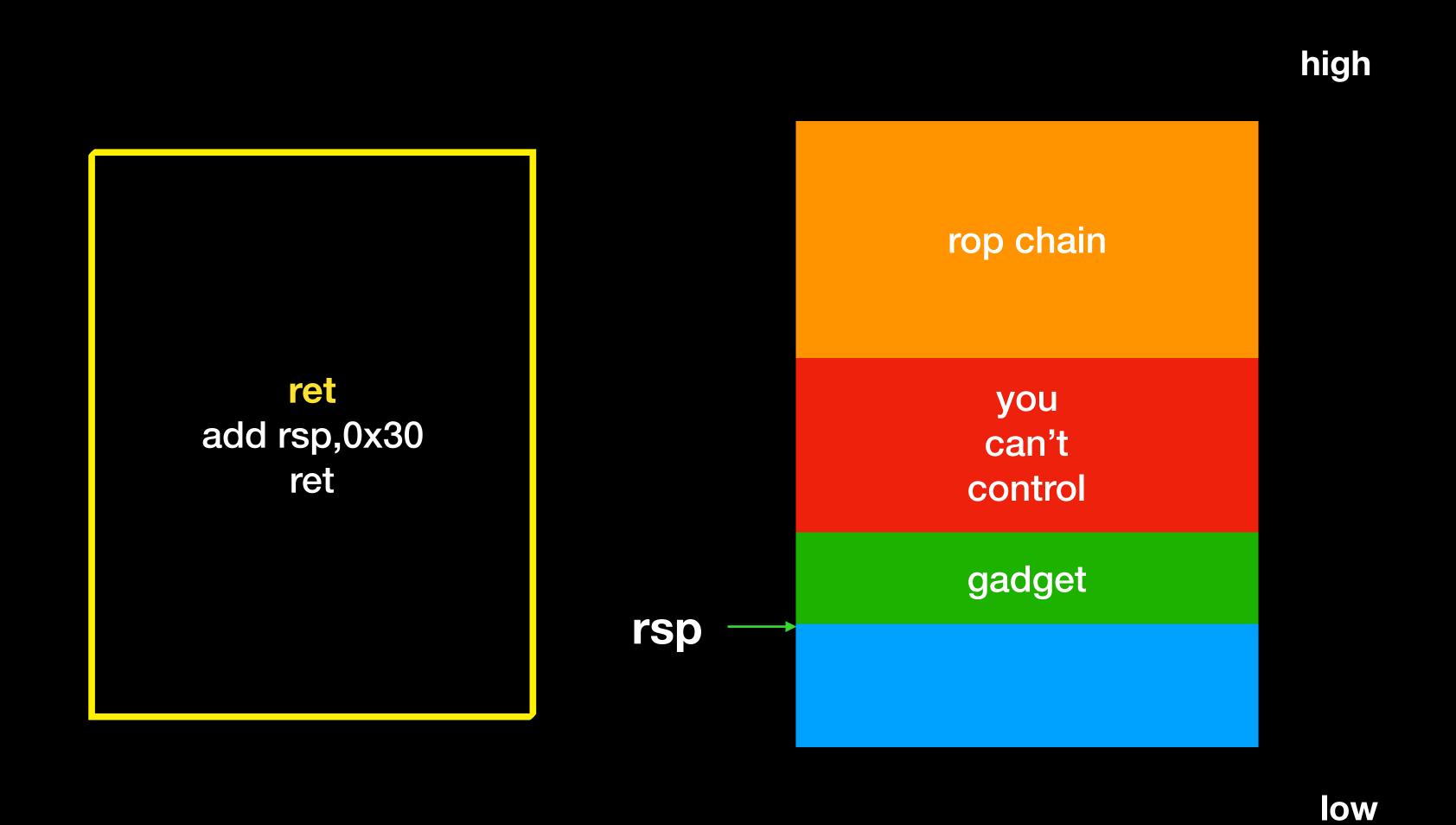


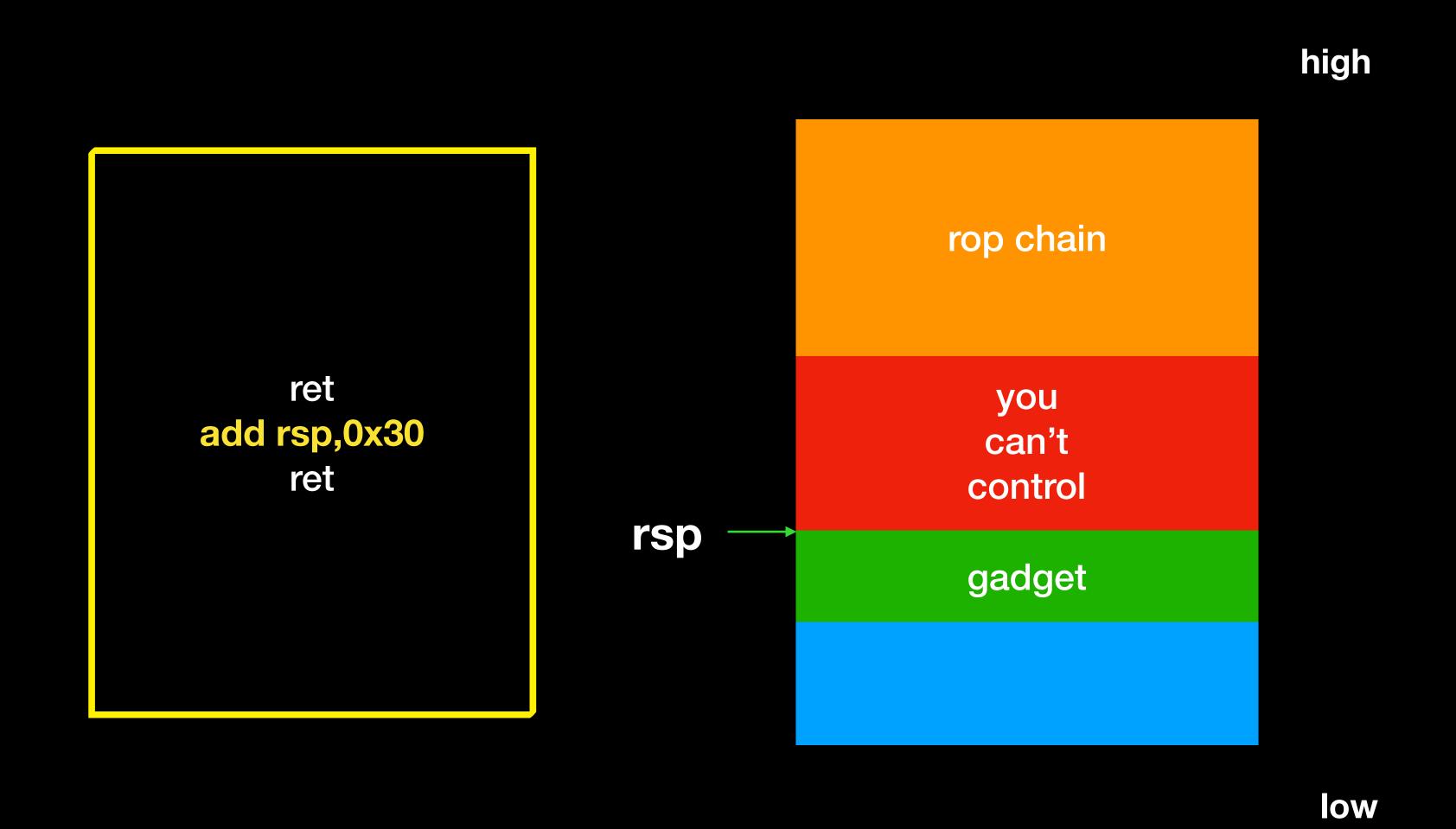


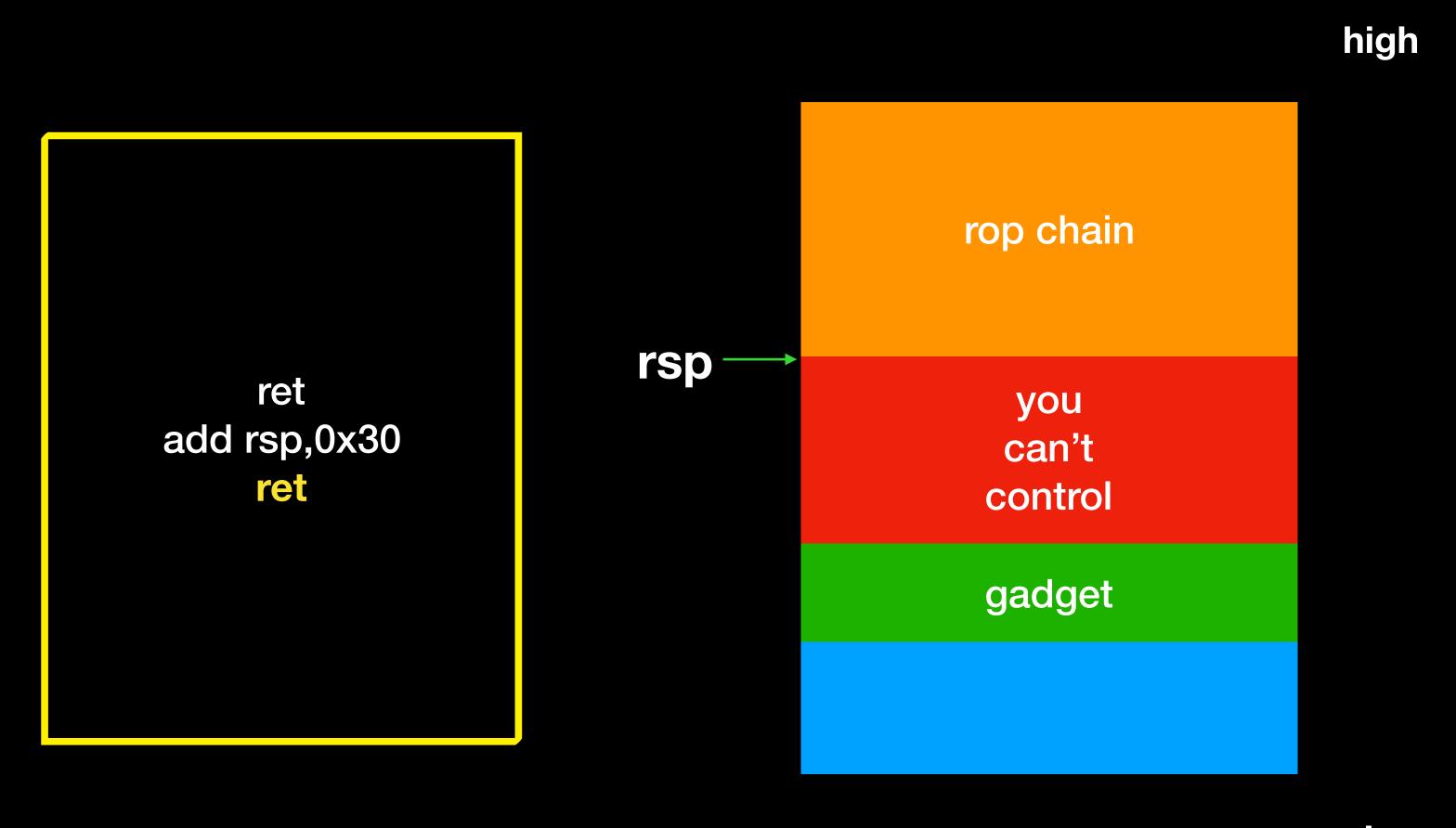


- 若能妥善利用,在沒有 libc 的情況下,有機會將整個 libc 給 dump 出來, 更有機會直接找出 system 的位置
- 無限 ROP ,幾乎可以做出所有事情,但唯一要注意的是 buf 大小要控制好,盡量選 bss 後半段位置,否則可能因為 stack 不夠大而 segfault

- Other migration gadget
 - add rsp,0xNN; ret
 - sub rsp,0xNN; ret
 - ret 0xNN
 - xchg rsp,exx; ret
 - partial overwrite rbp







Reference

• http://www.slideshare.net/hackstuff/rop-40525248

Q & A