

-- Day3 --

***Everything in a Linux is a FILE!

***There are 3 roots:

1: root user, uid = 0

2: / directory, root file system

3: /root/ directory, root home directory

cd command can work with the

1: absolute/full path

2: relative path:

a) cd . : current working directory

b) cd .. : parent working directory

*** cd ../../ to get the root, will work as cd ../../../../ ../../ .. etc

because the root directory doesn't have a parent, it is the parent of itself

-- Day 4 --

command [options] [arguments]

an option changes the behavior of the command.

an argument gives additional info to the command

*** Don't use rm -rf , use it very carefully!!

*** Every file in Linux started with . , is a hidden file

*** exit = logout = ctrl + D

*** reset = ctrl + L

*** using bg to continue the paused commands and processes, make th bg command not able to be paused again

*** shutdown -r now = systemctl reboot = reboot

*** Users & Groups ***

When a user created, automatically a "primary group" with its name is created and it will be its private group!

A user can be member in more than a group, but only secondary groups!

A user can be in one primary group only!

U can't use the user till you set him a passwd!

A normal user can't set a weak password for himself, but only the root can!!

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A normal user has no access to the system, except his home directory

*** When adding a user to another secondary group (while he is in another secondary group):

we use the command : `usermod -a -G [groupname] [username]`

without adding -a option to the command, the user will be deleted from any secondary group and added to the new selected

There is a big difference between the userid and groupid

Primary group is so important for Ownership!

Because when you create a file, the file's ownership goes to the user and its primary group

It doesn't make any sense to make the user's private group has more than one user!!

When changing the primary group of an existed user, it is deleted from the ex-primary group, and also deleted from it as a secondary group

`user:x:uid:gid::/home/user:/bin/bash`

x here means the password, but it is hidden and hashed in another file, `/etc/shadow`

The space between the home directory of the user and its primary group is for the user description,

It also called "GECOS"

The `/bin/bash` in the end, is the user shell!

Password are hashed in MD5 due to RHEL 5 and CentOS 5, now we use SHA

There are many shells/terminals in Linux, but the default one is /bin/bash

The exclamation sign ! means there is no password set yet

groupname:!:members ---> from: # cat /etc/gshadow

The space between the group members and the hashed password, means the group admin!

When you delete a user, his files still existed in his home directory!

After that you can delete his home directory, `rm -rf /home/userdel/`

Or you can only use: `userdel -r [username]`

Note: If the user's primary group is also a primary group of another user, and you delete the user,

The primary group won't be deleted, as it is a primary group of another user!

This would happen unless it is a secondary group!

When you are asked to add a user to group1, group2 and group3:

If you're not told to add it to a primary group, then all groups are secondary groups!

If you add ex-user and his primary group was still existed, the add process will fail!

Because you must choose his primary group first using: `usermod -g [primarygroup] [username]`

A new user created can't login in the terminal until you set it a password!

The user will be active on the system

Permissions !!

/*There are three things you can do to files in Linux:

You can r=read, w=write, x=execute */

There are 3 persons who can deal with files on a system:

1) Owner/User (r,w,x)

2) User's group (r,w,x)

3) Other/Default Category (r,w,x)

Any one of U,G,Other can (read, write, or execute)

This is not the *Actual Permissions, but the *Probability of permissions due to the Policies

Every permission can be represented by 1 bit, so we get 3 bits to all of them (User, Group, Others)

- rwx rwx rwx

The 10th bit is represented with the file type, and is the first bit of the 10 bits!

- = normal file

d = directory

b = block device

c = character device

l = link file

There are two types of device files:

1) Block device: the device which can be read/write to, block by block or chunk by chunk! like the hard-disk!

2) Character device: the device which can be read/write to, character by character! like the keyboard or the tty!

tty is represented on linux as a character device file on the /dev/

Link file means shortcut on windows!

If the permission is not given in any bit, we represent it by - sign!

The . sign in the end of 10 bits, means that SELinux is protecting the file!

The number after the permissions and . sign, means the number of links/shortcuts of the file (link count)

After the number came the user, group, file size, last accessed time and file name

Color of the files in terminal:

1) White: normal file

2) Blue: directory

3) Lemon: executable file

Executable file means that you can put code, binary or shell in the file and execute it!

It looks like the .bat or .exe in windows!

You can make a combination with chmod command!

If you don't select the permission to whom (u,g,o), it will execute it to all of them

When you apply a permission on a directory, It won't be applied to the files inside it!!!!

Or You can use the -R option to the chmod command!

In copy command "cp", we use -r, but in "chmod" command we use -R, both for recursively!!

Meaning of Permissions:

File - :

r --> to read or view its content

w --> write/edit/delete/overwrite

x --> execute or run the file

Directory d :

r --> to view the content inside the directory, to make (ls)

w --> to add/delete files, or delete the directory, like using (rm, cp, etc..)

x --> to make (cd) or (ls-l) to the directory

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Permissions:

****We can append two permissions in one command like:**

ex: `chmod u+r, g-w, o-wx file1`

****The Method of writing permissions is named: /Symbolic Method/ = Using symbols!**

*******We can use = to set permission, It is called Setting Permissions**

Which means, forget everything and set these permissions as new permissions

ex: `chmod u=rwx, g=rx, o=rx file2`

*****There is another Method called /Numeric Method/ = Using numbers only!**

Here we represent every permission with a number:

$r = 4$

$w = 2$

$x = 1$

The addition of any combination of these numbers will not ever conflict with other!!

Remember: `chmod [ugo] [rwx] [filename]`

In the /Numeric Method/: we get the addition of the permissions we want to give to the UGO

and then write it in its bit, like: $\text{user}=\text{rwx}=4+2+1=7$, $\text{group}=\text{rw}=4+2=6$, $\text{o}=\text{rw}=4+1=5$

So we here get user permissions = 7, group = 6, other = 5

Note: There is no any confliction between any permission and the other!!

ex: `chmod [ugo] [filename] ---> chmod 765 textfile`

Also Remember: We can use -R to recursive the command to the whole directory with its content!!

End of Permissions :) !

****\\ Redirection /****

****In Redirection we use three things: Input, Output and Error**

Your default Input is Keyboard, and the default Output is the Screen!

each one is represented by a number:

Input = 0, Output = 1, Error = 2;

Each number is called a descriptor!

We use also 2 symbols: to represent whether it is an input, output or error!

(<) --> to input, (>) --> to output/error

****Note: Error is an /Exceptional case/, so it must be passed**

Only Input or Output which could be passed by default!!

*****Sometimes one command can result both Errors and Outputs!!**

"cat" command takes its input from a file by default, and doesn't take it from the Keyboard!!!

cat file.txt do the same as cat <file.txt

To send input to the 'cat' command, use: cat <<EndOfText>> filetest.txt , to append the input in the file

'EndOfText' here means, add the inputs to the file and when you reach this word, terminate the input operation!!

We can take input from a file to another file

ex: cat <file1> file2.txt

The name of the file doesn't have any value in Linux!

What's matter only the content inside the file!!

ex: \$ cp file2 music.mp3

\$ file music.mp3

\$ music.mp3: ASCII text

Every file has a signature, it looks like a header to the file which tells what its type!!

Linux only looks for the file' content not the file name or its extension!

We can put the error and the output in the same file, There are 2 ways:

1) ls file1 file1001 >>outerror.txt 2>>outerror.txt

2) use '&>' to append both output and error

The difference between 'less' and 'more',

is that more terminates after reach the end of the file, but less needs to 'q' quit

We can redirect the output as an input to another command using '|' pipe command

ls -lR /usr/ | less

'tee' always overwrite the file of the results,

but we can use the '-a' option to append the result of the command to the file

Remember: TTY is the terminal when you have physical access to the machine

PTS = Psudo terminal screen = Terminal Emulator

as you emulate that you have a physical access or console access on the machine, but you connect remotely!!

connect over telnet, ssh, or the terminal opened in the graphical user interface

gdm = Gnome Display Manager : The Graphical user interface

gdm is the only service user is logged in!!

:0 refers to the first graphical monitor!

:1 second, and :2 the third ... etc

-- Day 7 --

*** Kernel is dealing with everything in Linux with numbers ***

*** Raw Space or Raw Disk:

Is the new disk you recently bought

MBR : Master Boot Record

MBR is necessarily to be existed, to manage and control the Disk Space

It is used to know and map where the partition Begin and End.

MBR, is not the only way to manage and control the space on the disk, there is also GPT.

MBR: is consisted on 3 things:

1) Partition Table: to know the Begin and the End of partitions!

2)

3)

* If the Partition table corrupted or MBR, the Disk will get back 'Raw' and the data will be lost!!

FileSystem: is the way of partitioning every partition.

FileSystem organize and manage the Store and Retrieve of data on the partition.

You can choose the way of partitioning due to your needs!

Every filesystem has both advantages and disadvantages

FAT32:

The file size can't be bigger than 4 GBs

It doesn't support Encryption or Compression

Supported on any Operating System

Other examples:

NTFS, ext2, ext3, ext4, XFS, ZFS, BTRFS

Every Partition is a block of blocks

Every Partition has an 'inode table'

inode table: is a table that tell all information about the file, where the block it stored and its privileges etc.

You can say the 'Metadata' about the files

Or 'Data about Data'

Metadata is stored in the inode table

Note: Partition table store the partition info not the Metadata itself!

Partition table is for the whole Disk, but the inode table is only for the filesystem/partition .

inode table also existed in Windows, but with different structure.

If inode table is corrupted, it affects only its Partition.

If Partition table is corrupted, the whole Disk will be affected.

Format: is the process of creating the Filesystem for every partition

Directory is a 'special' file, pointing to other files.

These are all the metadata stored about a file in the inode table:

- inode number (the number of the block where the file stored)
- permission
- owner
- access time
- creation time
- modification time
- link count

If it is Directory:

you'll see the same info mentioned above,

but with a metadata (pointers) about the files with their numbers which the directory is pointing to.

Every inode number is pointing to a block.

By default, every inode number is pointing to a block with 128-byte size.

* We can change the size of the blocks.

We change the size of blocks while Formatting the Disk

demanding on the Filesystem and what will be stored on it.

Every file is pointed to it by only 1 inode number.

If the inode number is pointing to 128-byte file size, and you're storing files with only 50 bytes:

the rest of the space will be wasted!!!

If the file size is larger than the block size, The file size will be separated into

a suitable number of blocks, until it is enough to the file size

and the inode number will point to the first block of the blocks

and the rest of blocks will be flagged as it is 'used'.

Defragmentation: is the process of assembling the Miscellaneous Units/Blocks to the same place
to reduce the seek time to make the reading and writing of the file faster

when you get the disk free space (df command), you'll find real and virtual filesystems.

It is very important to know about the Inode numbers,

Sometimes when you try to copy or write files on a free space on disk, and you surprised it will not

That's because the Free space already existed, but there aren't enough Inode numbers to point to it!!

This also mean that you made the block size on the filesystem big, or every Inode pointing to a very large space!!

and you stored a small-size files, so now you have a very big wasted space..

In this case, you must be able to know the Inodes, and see how many used and available

To ensure the issue isn't from the Inode or No!

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Continue **Inodes**:

It is impossible that 2 different blocks have the same Inode number,
because this represents a physical place on the disk

When you copy a file on the same Filesystem (Partition), It reserves an Inode number and then takes a copy of the data blocks
and write it on another data blocks, and makes a Pointer to point to it in the directory

**

When you cut/move a file on the same Filesystem (Partition), already written data blocks still as it, and the Inode number stills also

and what changes is the Reference or the Pointer which is what was pointing to the data blocks in the directory!!

Briefly, the data blocks still a it and also the Inode number, but just changing the Pointer from a directory to another.

That's why the Move is so fast!

When you copy a file on a different Filesystem (Partition), It reserves an Inode number and then takes a copy of the data blocks

and write it on another data blocks on the other Partition, and makes a Pointer to point to it in the directory

**

When you cut/move a file on a different Filesystem (Partition), It reserves a new Inode number and then takes a copy of the data blocks

and write it on another data blocks on the other Partition, and makes a new Pointer on the new Partition!!

and then gets back and deletes the Pinter/s which was pointing to the data and deletes its Reference!!

It doesn't delete the Inode, It just makes it free for future use!!

That's why the Move isn't fast!

When you remove a file, It just make the Inode number free!!

Minute 16:00 -

Never overwrite a data you recently deleted!

There could be 2 files(data blocks) in different filesystem with the same Inode number!

**

*

Soft Links:-

Reference to Reference (A shortcut to a file)!

If you make a soft link to 'file1' with the name 'file2', 'file2' will point to the filename

and not its Inode number!!

When you set 'file1' Inode' free, 'file2' will not be able to access the datablocks of 'file1'

Note: after creating 'file2', It took an "1 Inode number" to be stored on it,

But the 'file2' real data, is a reference to the 'file1' !!

You could make a Soft Link between two different filesystems!

*

**

Hard Links:-

Another reference to the same data-blocks!

Hard link take the same Inode number with the one that datablocks took!

When you delete a reference, the other reference(s) still can work and access these datablocks!!

Hard Link must be existed in the same Filesystem, because datablocks existed in one filesystem!!

Hard Link doesn't exist in Windows!!

The color of SoftLink is baby-blue

But after deleting the parent file, the color becomes 'red'.

It is a must to use the 'absolute' path of the source and the shortcut when you make a SoftLink

If you don't do this, It may work! but still wrong!!!

In Hard Link, you can not use the 'absolute' path, because it points to an Inode number!

Hard Link is not used in the Real-World scenarios nowadays!

Hard Link is not allowed for directory, because directory is a special file pointing to a references!!!

When you edit the Soft Link, it will edit the source file!!

But when you delete the Soft Link, the source file stills existed!!

Disks & Partitions:-

MBR takes only 512 Bytes from the space, and Partition table takes only 64 byte from the MBR !!

Every Partition can store its information in 16 bytes !!

Maximum number of Partitions on the Disk must be only 4 Primary Partitions !

We can solve the problem by making the 4th Primary Partition an Extended Partition,
and 'inside' it we can make Logical Partitions!!

MBR will have information about the 4 Partition only,

but the Extended Partition (4th Partition), will have a special Partition Table has information about the Logical Partitions inside it!!

Logical Partition start counting from number 5, no matter how many Primary Partitions existed!!

Revise some things, but later!!

These are all devices to read CDs (cdrom - cdrw - dvdrom - dvdrw - sr0)

sr0 : the reader of BlueRay Disks (the disks with huge spaces)

The BlueRay can read any CD with any space!!

If you have more than one cd reader or writer, it will have a differet number (sr0, sr1, sr2 etc)

In Linux, all those devices are accessible with '/dev/sr0'

Note: Linux is pointing to sr0 with the cdrom/dvd you have on the comp

-- Day 9 --

Creating and Formatting FileSystems:-

**

To show the recent connected disk, you must restart the machine

Or we can use the Package S3_Utils

**

Note: MBR is called 'dos' when it comes to 'Disk label type', in some Linux systems dos = 'disk operating system'

**

Every Partition has a different System ID, even between the normal partitions !!

There are a 'normal partitions' like = FAT, NTFS, extf3 ... etc.

There're also a non-normal partitions to be assembled from more then one disk and make a big 'Virtual Storage Space'

There are 2 famous types of Virtual Disks: RAID and LVM

These partitions are 'special', because you can assemble them form more than one disk !!

**

Warning: when you work with 'fdisk' command, you're dealing with a 'Disk' not a 'Partition'

Be careful before you write the changes after finish using fdisk command!!

Always verify from your work, Especially when you're deleting a partition !!

**

Linux by default using 'MBR schema' for making new partitions on a RAW disk !!

**

Adding new Partition means that you've added a Partition Table or you've edited it!!

**

Kernel is reading the MBR / Partition Table, after attaching it or booting the device. (after attaching, that's bcause there are auto-swappable disks)

After applying anything with 'fdisk', you must make the Kernel re-read MBr or PT !!

**

ext2 filesystem is one of the old version being used by linux!!

Mostly nowadays, people use 'ext4' or 'xfs'

Warning: when working with the 'mkfs', we're dealing with the partition not the disk !!!

When trying to make a new filesystem using 'mkfs' command without selecting the type, it will be ext2 by default !!

inode table of every partition will be called 'Super blocks'

You can not access any file on the partition until you know what block it is pointing to !!

If the Super Block is deleted, how would you access files on the disk, or how'd you know their blocks !!??

That's why it is so important !!

The system creates the 'Super Block' during the formatting process

While the importance of the inode table, the system makes copies of the super blocks and stores it in Miscellaneous parts

This happens that when any version of Super blocks is corrupted, you still can access it !!

The number of Super blocks is created depending on the space of the partition !!

Every inode number is representing 1 block of the data-blocks on the disk

No matter what the block size is, 1Kb, 1Mb or even 1GB !!!

Partition table still exists after deleting the partitions

To delete the Partition table or the whole MBR we need to write a random data on it !!

The difference between '/dev/random' and '/dev/urandom' :

that random may corrupt the input and output of the terminal screen (you should use 'reset')

And we use 'em to destroy the MBR, or to overwrite it!

We can use the command to overwrite the disk:

```
cat /dev/random >> /dev/sdb
```

This way is uncontrolled and informal way! Because you don't know how much data is overwritten

and you also don't know if the MBR is deleted only, or the data too !!

**

Use the command 'dd' to write random data, but with a selected range!!

You mustn't pass the size of blocks with b or B appended, it deals with the size of block by bytes

But you should pass if it was by Megabyte

```
dd if=/dev/random of=/dev/sdb bs=[1] count=[nofblocks]
```

You can make the number of blocks 512, to delete the MBR !!!!

**

MBR : Master Boot Record, takes 512 bytes

1--> Partition Table takes 64 byte

2--> Boot Loader takes 446 bytes

3--> Magic number takes 2 bytes

Magic number is like check sum for the rest of MBR, and used in data recovery !!

Format = Zero Fill !!

When trying to create a file using random data like using `/dev/random` or `urandom`

the file will be created but with a smaller space than we thought !!!

If we tried to make a file with 100Mbs and every block is 1Mb, the file will be at almost 8Kbs !!

This happens because, the 'random/urandom' write a random data that its output will be 100Mbs

but not its size when storing it !!!!!

'random/urandom' writes random data with every character used in any human language!!

that's why it is difficult to store it with Normal Encoding extension !!

Like the difference between ANSI Encoding and the UTF-8 Encoding !!

UTF-8 needs more space to store the data inside it, but ANSI not !!

Because UTF-8 can read the ASCII code of any character of any human language !!

That's why the process of creating the file isn't failed! but it doesn't work like we hope!!

To solve this problem, we use another file `'/dev/zero'` which is filling these blocks with zero !!

Remember: There are Super blocks still existed in the Partition, even after deleting it !!

So, when we re-create a new filesystem, it reads that there was a really-filesystem on it!!

So, we use the option `-f` to force the operation

**

We must make the partition accessible by mounting it!

After we unmount the partition, the data still existed but we need to re-mount the disk to access it !!

**

Data on partition is gone after you format it!!

**

We can get the partition info using 'dumpe2fs', but it works only with ext2, ext3, ext4

dump extended to filesystem

xfs filesystem is highly recommended than ext4, because xfs is so fast rather than ext4 may take a bunch of minutes

even if the hard wasn't SSD type !!

**

If the Electricity off, you might have partial data corruption

And then to solve this problem you need to check disk....

You must unmount the partition before you check the disk !!!!

If you didn't unmount, you'll have data corruption

**

Note: You must take a backup to the partition before you -force it !!

It might happen when you check, there may be a data corruption,
so you need force to check the partition, and it gives error in a block

it will offer you to fix the problem and you tell it yes !!

**

Filesystem check, checks only the inode tables !!

It checks to see if the inode table is consistent, if not it will only change the entry point
of the inode table

**

To take a backup of the whole partition even when you can not mount it!!

'dd if=/dev/sdb1 of=/sdb1-backup'

**

Revision Notes:-

- 1) If you mounted a partition without setting it a filesystem, you will get an Error !!
- 2) While working with a mounted-partition, we work with its mounted point if we were working with the files related to it
But when we work with partition itself, we must use its real path with its disk like /dev/sdb1

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Mounting:-

Root is the parent of itself,

Mounting in Linux means, linking the partition you mount with the related drive point on disk
and make the partition accessible across that point

Mount Point, is the point that across it you can access the partition

Linux is unmounting the root filesystem when you shutdown, and while booting it searches
the root filesystem and mount it automatically

In Linux, you mount any partition under the / root filesystem !!

When you mount any partition, physically, the mount point takes the space of that partition !!

If you mounted 2 partitions in the same mount point, the mount point will have the space of the latest partition mounted, and note:

- 1) If you unmounted the second partition, the space will be as the first partition!
- 2) If you unmounted the first partition, the mount point will get back to the / root filesystem !!

Never format the Extended Partition !!!

You just format the the Logical Partitions inside it !!!

When you unmount any partition, the command 'umount', is unmounting the mounted partitions in reversing order !!!!

mount -t [type] [source path] [mount point destination]

The public command of 'mount' command

*That was happening in the past, you should had to send the filesystem

but now this happens automatically !!

CentOS and RHEL 4 and 5 doesn't have that auto-detect, but 6 and 7 has !!

When unmounting any device, you can pass the device or its mount point !!

ex: 'umount /dev/sdb1' or 'umount /media/', while as '/media/' is the mount point of the device '/dev/sdb1'

The Problem with the mounting, is changing the place of devices !!

When you connect your device to a machine you're using a specific slot with a specific symbol !!

but when you re-connect it in another slot, you're now having a new symbol on the disk !!

if the disk first was 'sda', it may be after 'sdc'

When booting you'll get an Error-Message!! because you're mounting 'sda1' in '/'

and 'sda2' in '/media/' etc, and now there are a big problem because they're not found as 'sda' !!!!!

The Boot will fail and break, and Error-Message says that you must enter the root password to maintain the operation !!!!!

To overcome this problem we have 2 methods:

- 1) We use 'blkid' command:

When you format a partition, it gets a random UUID automatically

UUID length is too long, and its generation contains letters and characters!

That's why You will not see 2 partitions have the same UUID :)

So we mounting/unmounting the device using its UUID, ex: 'mount UUID="" /media/'

UUID is the best and preferred method !!

2) Using the Human readable name:

We use the command 'e2label'

The command is given the partition an attribute called LABEL with name we set to the partition !!

Now we can mount/unmount with the LABEL, ex: 'umount LABEL=P4rtition'

This method isn't preferred!

*

Typically in Data-Centers, we make an identical setup !!

Say that you have about 100 VMs or servers, The best way is that all of these 100 servers is to have

a standard-setup (/ takes 10GB, /var/ takes 5GB, /Oracle/ takes 100GB and is the Database)

rather than every machine have a special-setup

*

If you have a disk on the server1 and the Database '/dev/sda3' is named 'Oracle'

when you connect that disk to server2 which have a Database too '/dev/sdb3' is named 'Oracle' too

When you try to mount which one will you mount, both have the same LABEL !!??

That's why the UUID method is preferred :)

When you reboot the system, you'll lose the mounting you have had before !!

To overcome this problem,

Everything you do on a Linux system without setting it on a Configuration file,

doesn't have any value until you save it!!

To make the system auto mount the Disks you connect after booting:

You must save these changes in a configuration file in the '/etc/fstab'

The Syntax of the FSTAB is:

[device] [mount point] [type] [mount options] [dump order] [FS check order]

Mount Options:

rw = read write, ro = read only, exec = execute any binary or script on the mount point

no exec = do not execute any binary or script on the mount point

(By Default it executes the binaries or scripts on the mount point) and also read write!!

FS Check order:

When force shutdown or improperly shutdown happens, we use the FS check order to tell which partition will be checked first !!

its valid value is (0 : 9)

0 means don't check anything !

1 - 9 means the order of check you choose to the partitions !

If there are 2 partitions that have the same order, it will check the first one in the FSTAB file !!

Dump Order:

"Dump" is used to achieve the meaning of "Backup", still used but no more like in the past

You can configure the system to take the dump or the backup of the partitions while booting

0 means don't dump anything

1 - 9 means the order of dump you choose to the partitions !

Be Careful!!! When you editing or configuring the '/etc/fstab' file !!!

You might forget anything or make mistakes in the FSTAB file:

It will warn you that there is a problem while booting the system!!

In RHEL 6 or before, it will not give you a 1 min 30 sec timeout before it needs the root password !!

Only in RHEL or CentOS 7, it give you the timeout and then asks for the root password to maintain the problem !!

It will then take the root password and then you can edit the FSTAB file !!

To avoid this problem: We have 2 options:

1) We use the command 'mount -a', it will mount everything in the FSTAB

and will warn you if there is something wrong!

2) We can 'cat /etc/mtab' and take the line of command related to the partition we need to mount

and paste it in the '/etc/fstab', that line is the same as the syntax of the FSTAB !!

Before you delete any partition, be sure it doesn't exist in the '/etc/fstab'

and delete it from the file and then 'mount -a' and delete that partition !!

Mistake I made and Mr. Mustafa:

Never delete a Mounted Partition!! Unmount it first!!

To write a mkfs (To make a filesystem) you must unmount the partition first !!

Never Fucking use the (fdisk / cfdisk) without selecting the disk __!_

Archiving and Compressing:-

There are 2 common packages to compress files:

gzip ---> fast

bzip2 --> higher compression ratio

You cannot read a compressed file on any Operating System.

What happens in Windows with WinRAR: is that it decompress it in the temporary files !!

Video and Audio files aren't compressed in these ways like the text files. It's completely different !!

Difference between Compression and Archiving:-

Compression is used to reduce the size of the files.

Archiving is used to link or assemble many files in one file only, without reducing the size of these files !!

Red color means that it's an archived or compressed file.

To compress a directory and archive it at the same time we use the option (z) appended to cf/xf

ex: 'tar cvfz' or 'tar xvfz' and 'tar cvfj' or 'tar xvfj'

while 'z' is gzip package, and 'j' is bzip2 package !!

and the extension is appended as 'filename.tar.gz' and 'filename.tar.bz2' !!

'tar' is one of the few commands that its options don't have - sign before it !!

It may work also with the - sign, but sometimes it may give another meaning !!

Freelancing last 3 minutes !!

-- Day 11 --

Process Management:-

The Difference between the Process and the Program/Application:

*The Program is the binary existed on the system without running...

*The Procss is the binary when it starts to consume the CPU, to take a space on the memory, to allocate data on the disk !!

Kernel controls the Processes through the 'Process ID' or the 'PID'

and every process has a 'Parent Process ID' or 'PPID' !!

PPID is the number of the Process ID of the Parent Process

Why we need PPID? for 2 reasons:

1) When an Application is open on Linux, like 'Shell', and you opened 'nano'

Linux here is using "Copy on Write" method,

The best thing here is not ('allocating nano a piece by piece'),

The best thing is to Copy on Write the Shell and customize it due to the needs of nano

and what still needed from nano is allocated then !!

2)

Every running-process on Linux has Memory pieces called 'Pages' or 'Memory Pages'

Most Apps on Linux is using 'Shared Libraries'

If Shell and nano were written in C Language, Shell is written with 20 Libs and nano 12

there could be some shared libraries or memory pages between them !!

So the best practice here is to make a 'Common/Shared Memory Space' and make both of them pointing to it !!

rather than making a different Memory Space to everyone of them

For example, If nano wants to write on the Shared Memory, Linux will give a new Space to nano

to write what he needs in !!

So, instead of making a special memory space for every program, Linux solves it with the shared library method

and if one of the programs wanted to write or edit or delete, it will get its special space !!

So the Second 2nd reason to know the PPID is:

If nano is not responding, you can terminate it, but if it doesn't responding too???

You can go then and terminate its Parent Process !! as nano was built on the memory of Shell !!

The first 1st Process running on the System is 'systemd' and takes PID '1' !!

This only in RHEL7 or CentOS7

The first 1st Process running on the System is 'init' and takes PID '1' !!

This in RHEL5/6 and CentOS5/6

The job of 'systemd' is calling the Processes and Services on the System !!

if you 'ps aux' and the tty is '?', this means the process is working in background

'kill' command is sending a 'signal' to the process saying 'stop'

use 'kill -l' to list all the signals

the default signal used with kill is 'SIGTERM', signal terminate number -15

SIGTERM is so polite, but if the process doesn't respond we can use SIGKILL number -9

SIGHUP makes the process re-read or re-load the configuration file and number -1

Using 'kill -9' directly will cause Disastrous results !!

Kill, which is kill -15 by default, is asking the process to terminate, but kill -9 force it to

so if the process was writing data on the disk and you force it to stop, you'll have data loss !!!

Kernel itself is pointing to 'systemd' and 'kthreadd', that's why it's PPID is = 0;

Kernel is the Parent of 'systemd' and 'kthreadd' process, as it starts them !!

'systemd' or 'init' is the Parent of all processes

Try to avoid killing the process with its name, just use the PID of it

Using the 'pkill' command can make Disastrous results, if you are running instances of the process

the 'pkill' will terminate the all, because you killed them by their uni-name.

-- Day 12 --

Process Management Cont:-

Process working in foreground is a normal process

Process working in the background is called 'daemon' or 'service' like in Windows !!

The Parent of daemon process is the process started it,

if you started it from the GUI, the Parent will be GNOME Display Manager (Graphical Interface)

if you started it in booting by the call of systemd, its Parent will be systemd

If the daemon process working in the **background**, and you terminated its Parent,

Its Parent automatically moves to be 'systemd'

Example:

If you started the 'sshd' service from the terminal in the GUI, and you terminated the terminal

the 'sshd' won't terminated, because the process control goes to 'systemd'

Service: is the process waiting the request from the user to reply with his needs or to give him a specific order or to serve a specific job to him !! this is in common.

Services typically listen to a port !!

Services could also work locally on the machine

Any application can work in foreground or background as a daemon !!

'nice value', is the value controlling the priority of the running-processes

its range is from {-20 : 19}, and its default value is = 0, means normal priority

the highest priority is = -20 , and the lowest priority = 19;

**

Any normal user can give his process a lower priority, but can't give it heigher priority !!

and he can'e even get it back on high priority after he give it lower priority !!!!!

If you try to give the process a higher priority, access will be denied !!

Root can do anything :D !!

Zombie process is a running process consuming resources but not working in the real !!

You can start the service/process with a specific priority,
and you can also change this priority to real-time running process !!

Typically, when you take back-up for the Database you need to pause it temporarily till you finish

If you want to take a back-up and need it in fast, you must give the process the highest priority

as it can consume the machine resources as needed !!

ctrl + C terminates the process with signal 15 (SIGTERM)

ctrl + Z interrupt the process with signal 2 (SIGINT)

Q quit the opened session with signal 3 (SIGQUIT)

Searching and Locating Files and Directories:-

We can search commands with 'whatis' and 'whereis'

To search files we have 'locate' and 'find' !!

The problem with 'locate' is: that it depends on a Database when it searches for files ..!

Linux System while working is trying to track what files/data have been added or deleted to/from the Filesystems

This update happens everyday, about midnight ...

We can force the Database to synchronizing with 'updatedb'

As it depends on a Database, the locate is so fast !!

'find' is a utility like 'locate'

actually it is an updated-locate, that can search in the real-time not in the database ..

Note: find searches the name only, but we can search and get results and variation about the name

just by appending asteric * to the last of the filename !!!

'find' by default searches in the pwd, Present Working Directory !

To make a file with a name of two word or more, we write the name between double quotes !!

or we can use a '\' = escape sequence after between the words of the name !!

We can take the result from the 'find' command and put it in a special file:

ex 'find /etc/ -iname network -exec cp {} ./work/Day12 \''

where -exec means execute another command and \ means don't change anything and take it as it was like !!

Note: '-exec' command needs ';' after it, as it link between two commands,

that's why we've added '\' or escape sequence between it to ignore it !!

You can use any command with '-exec' command, like move or delete .. etc.

You can use pipe | more than one time in the same command line

-- Day 13 --

VIM Editor:-

Vim is an Advanced Text Editor

Vim is an advanced 'Vi' = 'Visual Editor'

Vim is an enhanced Vi

Vi was like a nano

Vim has 3 main modes:-

1) Command Mode:- Since you open a file you are in the Command Mode!

- Allows you to only read the file

- You must click 'Esc' to go to the Command Mode

2) Insert Mode:- To be able to edit the file

- You must click 'insert' or 'i'

3) Execution Mode:- To exit or to save the file

- Click ':' to go to the Execution Mode

- write 'w' to save and write

- write 'q' to quit

- write 'wq' to save and quit

- write 'q!' to quit without saving

If you wanna search a word in the Vim editor, just click / and then enter the word !

press 'n', for the next result !

press 'N', for the next result !

press 'l', to move the cursor to the right !

press 'h', to move the cursor to the left !

press 'j', to move the cursor to the down !

press 'k', to move the cursor to the up !

press 'o', to make new line under the cursor and go Insert Mode !

press 'O', to make new line above the cursor and go Insert Mode !

press 'yy', to copy the line, y = yank = copy :D !

press 'nyy', to copy more than one line, n = number of lines!

press 'dd', to cut the line !

press 'nndd', to cut more than one line, n = number of line !

press 'p', to paste the line/lines !

press 'u', to Undo !

press 'dw', to delete the word where the cursor is, if in the middle, delete from the cursor to the end !!

press 'dl;', to delete the letter where the cursor is !!

press '.', to redo the last command you did

press 'gg', to go to the begin of the file

press 'G', to go to the end of the file

press 'd\$', to delete from the cursor to the end of the file !

press 'd0', to delete from the cursor to the begin of the file !

Note: Last search result still being highlighted !!

To see the number of lines, go to Execution Mode and write 'set number' !

go to Execution Mode and write 'set nonumber' to cancel !

To move to selected line, go to Execution Mode and write the number of the line you wanna move to !!

To delete more than one line, go to Execution Mode and write '7,10d'

from line 7 to line 10, delete !

To delete from the cursor to the end of the file, go to Execution and write '.,\$d'

'.' where the cursor, '\$' the end of the file

To search a word and replace it, go to Execution Mode and write '%s/sync/replicate'

'%s' to substitute /word/replacedword

Note: this change happens to the first result existed on the same line only !!

If the results are splitted into many lines, without repetetion, all lines will be changed !!

To replace all the substituted words, go to Execution Mode and add '/g'

ex '%s /sync/replicate/g', while g means global

To read and copy the content of the file to the cursor place, go to Execution Mode and write 'r file1'

'r' means read the file

To save the file to a new file , go to Execution Mode and write 'w filename'

To append the content to another file, go to Execution Mode and write 'w >> filename'

while '>>' to append the data without override the file !!

To highlight/unhighlight the text, go to Execution Mode and write 'synatx on/off'

To go to shell or execute a command in hurry, go to Execution Mode and write '![command]'

To add the result of a command to the cursor place or to the file, go to Execution Mode and write

'!.date', while '.' means where the cursor is !!

'r ![command]' :) !! By me :)

End of the first 1st Track () ;

-- Day 13 --

Boot Loader:-

The process of booting the Linux Machine consists of 2 stages:

- 1) Before the Operating system booting
- 2) After the Operating system booting

BIOS : Basic Input/Output System

After clicking the Power button, the BIOS checks the hardware components

- 1- and the first thing it makes is **POST** power on-self test, to check the minimum hardware like processor, RAM, HD and the machine capabilities
- 2- and then the BIOS detects the Bootable devices and compare it with its settings, like the 1st and 2nd boot device and then boots from these devices.
- 3- Load the **IPL**, Initial Program Loader, its job is to execute the **BootLoader**.
- 4- **IPL** then load/start the **BootLoader**.
- 5- **Boot Loader** loads the Kernel.
- 6- And then the Kernel starts the Process (**systemd** / **init**)
- 7- **systemd** then calls/starts the services/daemons.

Our Boot Loader in Linux is called **GRUB**, Grand Unified Boot Loader

Boot Loader is about 446 bytes stored in the MBR

Boot Loader consists of 2 stages:

- 1) 1st stage: the real Boot Loader, stored in MBR
- 2) 2nd stage: the files related to it, stored on the disk, in **/boot**

The 1st stage job is to call the 2nd stage 😊

This isn't a shortage of Boot Loader, but the MBR !!

GRUB is so smart, and can detect other Operating Systems like Windows, Solaris, Mac and can load them together.

Windows has **NTLDR**, NT Loader, can't detect any other operating system.

If you have a Machine with RHEL and Windows installed

If you got a virus on Windows, and you re-installed it, the NTLDR will over-write the 1st stage of the Boot Loader in the MBR !!

NTLDR doesn't care about the other operating systems

But Note that the 2nd stage of the Boot Loader still existed with its files on the RHEL, don't try to re-install the full RHEL

Just re-install the 1st stage of the Boot Loader, this will overwrite the NTLDR and install the 1st stage

And the 1st stage will notice that you have Windows and Linux and both of them will work together again !!

Note: when you installing Windows never try to format the MBR!

The configuration file of the GRUB is in the **/etc** as a soft link to itself in the **/boot/grub2**

That's why we do as the rule of 'Every configuration file must be in the /etc'

KERNEL :-

Kernel in Linux consists of 2 parts:

- 1- Kernel itself
- 2- initramfs (initrd in RHEL 6 and above!) initial ram drive = initrd

Kernel is splitted into 2 parts because the Kernel developers make the **Core** in the kernel file, but the modules and drivers in the initramfs.

If you want to access the Wi-Fi, network, tape device, the disk or any other thing.

Splitting the Kernel allows you to :-

- 1- the **Memory Footprint** of the Kernel will be too small, to reduce the memory used by the Kernel
- 2- In the booting process, loading Kernel with reduced memory size like 5MBs will be faster, rather than if it was 40 MBs !
- 3- If you installing a driver and the installation was corrupted, you'll need to re-generate the **initramfs** or to re-build your drivers rather than re-generate the whole Kernel, because the Kernel will be untouched!

You will not see that when you install a driver and the installation corrupted, that the entire system is corrupted and you need to re-install your machine, this doesn't happen in Linux!

But may happen in Windows!

Because the Kernel of Windows is just 1 file! One file containing kernel, drivers and modules and many other things!

If **initramfs** was corrupted, you can re-generate it with only 1 command!

Kernel while loading :-

Kernel mounts the root file system / in **ro** (read only) mode and try to ensure that there are no problems in the system!

If you had a problem, it will fail the load.

If Everything is OK, the kernel re-mount the root file system / in **rw** (read - write) mode.

Because if the mounting process was in the **rw** mode, and you have a problem, you'll end up have more than one problem!

Summary:

- 1- Kernel Loading
- 2- Kernel loads some of the drivers from the initramfs
- 3- Kernel mounts the root file system /
- 4- Kernel calls the systemd / init

These are the cases you need to restart your running-server:

- 1- If you re-enabled the SELinux if it was disabled
 - 2- If you upgrade the Kernel
- Needing to restart the server as you need to upgrade the Kernel will be unnecessary soon.

Because there is K-Patch or Kernel Patch, which allows you to upgrade the Kernel during running the machine!

-- Day 14 --

GRUB supports:

- 1- More than one Operating System
- 2- Making Security settings to it, like making password to it to prevent anyone from editing

systemd :-

init and systemd similarities:

Always takes number 1
Calling the services to run

init and systemd differences:

- 1- When init calls the services to run, it calls the 1st service and wait for it to respond (respond if failed or worked), and then calls the 2nd service and also wait for it to respond, and then calls the 3rd service .. etc
- 2- If every process needs about 3 seconds to run, the 3 services needs 9 seconds to work which is a waste of time !
- 3- On the other side, systemd, run the services **in parallel**, calls the services in the same time and wait for them to work !
- 4- If every process of the 3 processes needs 3 seconds to run, the 3 processes together will need to 3 seconds to work !
- 5-
- 6- init to call the services with a "init script", every service has its own "init script", written in Shell Script!
- 7- systemd has "Unit files", normal text files like scripts but without codes, these "Unit files" contain only texts like: var = value; what makes the debugging and troubleshooting much easier than "init scripts" !
- 8-
- 9- systemd is compatible with init, because systemd allows init to work with it in the same time
- 10- That's why if you have a service and it was written in init shell script, you don't need to run these service on an earlier version of RHEL and also you don't have to write a Unit file for that service to work with systemd!
- 11-
- 12- init has a different command for every task you need to do, but systemd has 1 command gives you a full control on every service you have ! '**systemctl**' system control
- 13-
- 14- We had a service called syslog and was developed to rsyslog was to recording logs, so if the system is booting in RHEL 6 and there was a service running before syslog, you will not be able to see its logs and if there was a problem you also will never be able to know what the problem exactly is with the service!
- 15- To avoid that problem with init in systemd, they wrote **journald**, also to recording the logs.
- 16- When you want to search the logs in syslog or rsyslog, you were searching any problem in the logs, you were searching the whole file without filtering!
- 17- **journald** has a command '**journalctl**', you can search with this command about a selected-service with filters!

The development between init and systemd didn't happen overnight!

There were many implementations, like **upstart** developed by Ubuntu team pushing people to use **upstart** and it also uses **Parallel Booting System**!

After that Fedora developers wrote systemd and both Fedora and Ubuntu developers agreed to use systemd !

systemd also compatible with upstart.

Apache service is called httpd

If the service status is active, this means that the service is currently working.
enabled or disabled status is telling the machine if the service will work once the machine rebooting or not!
You will need to make it run manually after rebooting!

Every service you install comes with a default configuration, may be it is enabled to work after booting by default or not!
You can change it.

Real World Scenario

If you have more than one implementation for a service like **Apache** and **nginx**

These two services can be installed on the same time on the machine, but they can't work on the same time as they listen on the same port **80** !

Apache has many features than **nginx**, but also slower than **nginx**!

Mr.Mustafa stopped Apache and installed nginx as the web server was too heavy and also disabled the Apache service.

The Genius working with him reboots the system and found that the web service wasn't running, so he turned on Apache!

This is a خازوق !

To avoid this problem:

You can **mask** the service, preventing the service from running even if you start it!

After that you can test the new service you installed and if you want to get back to the service, you can **unmask** it again!

systemd calls the mode of GUI.

So the Basic is CLI not the GUI!

You can't work only in GUI without the CLI

As systemd can work in more than one mode, these modes are called **targets**.

init has modes also but called **run levels**.

Windows' has modes also like: **Safe Mode**, **Safe Mode with Networking**, **Recovery Mode** ..etc

Windows has the same concept of ren levels or modes.

init run levels:

0: poweroff / halt

1: "Single User Mode"

2: "Multi User Mode", but without GUI and NFS (Network File System)

3: CLI, only ttys, was called "Multi User Mode".

4: unused, left for future and not used

5: "Multi User Mode", with GUI too.

6: reboot

NFS: is a service used to share files between Linux and Linux or Linux and Unix Machines .. etc.

Usually, Servers were working in the 3rd run level of init and the Desktop Machines in the 5th run level.

GUI is the most thing consuming resources.

The 1st run level, Single User Mode, was for troubleshooting, it allows you to get **root** privileges without asking for the password.

systemd targets:

1. multi-user.target: the 2nd and 3rd run level of init.

2. graphical.target: the 5th run level of init.

3. rescue.target: the 1st run level of init, equals to Single Mode.

4. Reboot: run level **0** of init.

5. Poweroff: run level **6** of init.

The **default target** is the target started when you boot the machine.

The result of the 'runlevel' command is like "N 3", which N is the ex-runlevel and 3 is the current run level.

But N here means "Not available", as it doesn't know what the run level before.

You can change the current target while running the machine.

The difference between 'isolate' and 'set-default' is that set-default changes after rebooting, but isolate changes the running session.

init command is still existed on Linux, just for compatibility.

The main source to most of the files of the systemd is in `/usr/lib/systemd/`

To enable/disable any service, the system just make a soft link to it in the `/etc` if you want to enable it, or delete the soft link if you want to disable the service.

init start in Sequential Boot, but systemd in Parallel Boot.

One of the most important features in systemd is: Dependency-Based Boot
That allows to **only** run a service after a specific service.

And tells that a service doesn't run **only if** a specific service run before it.

Example: To don't run Apache while the Network is down, to prevent consuming the machine resources

If the service state was "static", this means that this isn't running in the background, but the system used it and it doesn't in need in the real time.

The difference between the 'restart' and 'reload' to the service when using **systemctl** command:

Is that restart, turn off the service and run it again.

But reload, makes the service re-read the configuration file and apply it.

-- Day 15 --

GRUB Boot Loader:-

Menuentry in the GRUB cfg file, is a selectable line to allows you to choose which Kernel version you choose to boot from, or to choose another Operating System if you have.

You can install more than one Kernel version in Linux system.

But when booting you can choose only one version to boot from it.

This also allows you to test the new version of Kernel you recently installed.

Every Kernel version has a version of initramfs.

Remember:-

Kernel naming the disks with its physical type, like giving **sda**, **sdb** for the SATA, SSD and SCSI disks

and **hda**, **hdb** for ATA and Parallel ATA and **vd** for virtual disks like on the virtual machines. All these names are human readable.

Kernel is which giving the names of partitions like "sda1, sda2, sda3, sdb1, sdb2 ..etc", but GRUB is so low level than Kernel!

That's why the GRUB is reading all Disks with the name "**hd**" but without lettering it, it numbers it with numbers like "**hd0**, **hd1**, **hd2** .etc", this happens because the GRUB can't read the name of Disks because it loads before the Kernel.

And also, the GRUB identifies the Partitions with MBR, with the name "**msdos1**, **msdos2**, **msdos3**"

So the complete name for **sda1** → **hd0,msdos1** and **sda2** → **hd0,msdos2** and **sdb3** → **hd1,msdos3** . (hd0,gpt1)

When GRUB is installed, it detects the disks on the machine and detects their type, rather it was **dos=MBR** or **GPT**

And then the GRUB auto-detect the Operating Systems installed on the disks and which partition where it's installed.

The Difference between the MBR and GPT:-

- 1- MBR is written at the first of the disk, but GPT is written at the first and at the end of the disk.
- 2- MBR can have only 4 primary partitions, but GPT can have **128** partitions!
- 3- MBR has a limit space reach to **2TB**, but the GPT can reach to **9 ZB**!!

Note: If you make the disk MBR and you wanted to turn it to GPT you will have Data Loss! And you will need to make the disk raw first and then install the GPT. And vice versa is right, if it was GPT and you need it MBR !!

Some Programs take a backup copy of the MBR and put it in the Memory and then delete the MBR and install GPT, and after that restore the begin and end of every partition and puts it in the GPT ! It works with some programs but not all!!

You may not see any program that converts the GPT to MBR, because MBR is 32 and GPT is 64 bit and you can not put the 64 bit in the 32 bit 😊 but vice versa is right!

MBR = msdos = LBA (Logical Block Addressing)

GPT = GUID Partitioning = EFI = uEFI (Universal Extensible Firmware Interface)

UUID of the partition changes in one case, if you format that Partition because it is generated randomly!

In the Menuentry, the system searches the root file system with the name of it “**hd0,msdos1**” and its UUID, if it isn’t found 😞 It then searches with UUID only, as it is a unique key to the partition!

In the GRUB cfg file, you can delete “rhgb quiet”, to not see the graphical boot and to not be quiet means to be able to see what happens in the background while booting and what the services starting.

Being able to see the boot log can help you in many situations, like if the Apache fail to run, you can see in the boot log what is the problem with the Apache service!

****You can re-generate the GRUB cfg file with the command ‘grub2-mkconfig’**

And Note: that everything you’ve edited before will get back to its main source!!

Everything you’ve edited will be overwritten immediately after re-generating the GRUB.

Sometimes you need to make these edits still, even after re-generating the GRUB cfg file.

To make this, you need to edit the cfg file of the RUB in the “**/etc/sysconfig/grub**”

After editing the file, you **must** re-generate the GRUB config!

GRUB start counting the Menu Entries from number **0**.

You make re-generate to the GRUB in the case of:-

- 1- You installed a new version of Kernel.
- 2- You installed/updated the package of the version of GRUB2.
- 3- You edited this file manually to edit its settings.

SELinux:-

SELinux classifies the files on the system and gives every file a LABEL.

LABEL’s job is to give every file a specific access.

If you forget the root password and you get into the Recovery Mode to re-set it, you’ll edit the “**/etc/shadow**” you’ll change the LABEL of the file or maybe you’ll delete it!!

If the LABEL was edited, and Kernel is booting and look to LABEL, SELinux will prevent the Kernel that this LABEL isn’t recognized.

You may have changed the root password, but root can’t log in as the Kernel can’t read the file **/etc/shadow** because of SELinux.

SELinux gives every file a Label, and if the Label doesn’t match the **Policy**, the Kernel will not be able to do anything to this.

And the **Policy** says that **/etc/shadow** must be with label **shadow_t**

NOTE: Recovery Mode in Linux doesn’t support the **SELinux**!

➔ After changing the root password, you must make the system re-set the Label to its origin!!

Re-setting root password:

We can edit the Menu entry from the machine while booting by clicking “e”.

We add “**rd.break**” in the last of the line before loading the **initramfs**, which will make the kernel fail the boot after it loads the **initramfs** file and after this you’ll get root access on the machine!

Kernel has a Mini-Version of a File System Looks Like **/**, but not the main one, this version is the version Mounted in the memory when you enter the Recover Mode! And your real partition of the disk **sda1** will be mounted in **/sysroot**

Remember: Kernel is loading the root FS as **read-only** mode! To switch to the real File System, you need first to mount it in the **read-write** mode!!

To change the mount mode we use the command “**mount -o remount,rw /sysroot/**”

And then we change the **root** filesystem to the **/sysroot/** with the command ‘**chroot /sysroot/**’ !

After that we change the password with ‘**passwd**’ command and type the new password!

If we take a look to the Label of ‘**/etc/shadow**’ we will see it as ‘?’ so we must change it.

To change we need to make it manually and this is the hard way.

Or we can do it automatically by making a file ‘**touch /.autorelabel**’ this will search all files and try to find what the Policies of these files must be and apply or re-label these files due to the Policy!

If we created that file under the **/**, it ensures that SELinux checks the files and re-label it before the Kernel starts!

Note: this file tells the SELinux that there is a file has some changes please re-label it again due to the Policies.

You must make a Password for the GRUB to avoid the change of the physical change of the root password!!

If you forget the GRUB password too, you can attach your disk with another machine and do the same steps!

Or you can start booting from a CD and entering the Recovery Mode!

You can encrypt the Hard Disk to Prevent this process because it has many risks, and someone can steal your data.

-- Day 16 --

Disable Ctrl + Alt + Delete:-

While working in the machine, If you accidentally clicked **Alt + Ctrl + Del**, the Machine will restart!!

This causes a disastrous problem in the Production Environment!

Your role as a **Sys Admin** is to prevent these **Alt + Ctrl + Del** from working:

What causes the **Ctrl + Alt + Del** to reboot is the soft link found in the ‘**/usr/lib/systemd/system/ctrl-alt-del.target**’ and is a soft link to the file ‘**../systemd/system/reboot.target**’.

The file was controlling **init** to make the **Ctrl Alt Del** to reboot was the file ‘**/etc/inittab**’.

If you want to prevent this in RHEL 5 or 6 you just comment # the file of the **inittab**.

But now as this file is a soft link to a real **systemd target**, you just can **mask** it with

the command ‘**systemctl mask ctrl-alt-del.target**’.

Now if you clicked **Ctrl + Alt + Delete**, the system will not re-boot!!

We **masked** a target not a service! So, we must do ‘**systemctl daemon-reload**’ after that to reload all the configuration files that we’ve edited recently!

After these steps, if you ‘**ls -l**’ to the **ctrl-alt-del.target** you will see that it stills pointing to **reboot.target**!!

But if you see the file in the ‘**/etc/systemd/system/**’ you’ll see the file pointing to ‘**/dev/null**’ as the configurations must be edited in the **/etc** not in the **/usr/lib/systemd/system/**, so the **mask** didn’t edit it, it just take a soft link from it in the ‘**/etc/systemd/system/**’ because what existed in the real **/usr/lib** can’t be edited!

Note: We must ‘**systemctl daemon-reload**’ after editing the configuration files especially editing the targets.

What the **masking** really do, is that it makes the file as a soft link to the Special File `‘/dev/null’`.

If you want to delete a file in Secure Way, just make it pointing to the `‘/dev/null’!!!!`

NOTE: **systemd** while working it takes its configurations and settings from `/etc/systemd/` and it has the priority! If it not found there, it takes it from the `/usr/lib/systemd/`.

➔ If you want to disable it in the RHEL6 you can go to the file `‘/etc/init/control-alt-delete.conf’` and inside the file, you can comment **#** the execution command!

Network Manager:-

We have 2 implementations to Network on Linux:

- Network (Old)
- Network Manager (New)

Network was so easy to configure, but if you want to configure VLAN you’ll need other commands with other modules to configure it. If you want to configure Bonding, you need an additional package installed first!

Network Manager allows you to:-

- Configure Basic TCP/IP
- Configure VLANs
- Configure Bonding / Nig Teaming
- Support advanced features like connecting to VPN connections.
- It consists of **Building Blocks**, If you want to change the VLAN implementation you don’t have to change the whole code.

The Network Manager has a feature makes it easier to use like the systemd, which is the one command used `‘nmcli’`.

Every PC has an implementation to the TCP/IP Stack will have a **LoopBack Interface** as a Virtual/Software Interface.

The **LoopBack**’s job is to test the TCP/IP implementation, whether it works correctly or not.

The IP 127.0.0.1 is locally and existed on every PC.

It allows you to run local services on your machine like: MySQL, Apache or NginX.

The second Part of this Day was about Networking and Ips and Subnet mask...

-- Day 17 --

The Default Gateway is the gateway that connect you to the Internet by default!

But the normal* Gateway is the way to a specific network.

DNS stands for **Domain Name Server**.

Default-gateway typically used with Layer2 Switches.

Default-route typically used with Layer3 Switches, Routers and Firewalls.

This because the L2 Switches can’t do Layer3 Packet Forwarding or (Routing), That’s why it is called Default-gateway! And L3 Switches, Routers and Firewalls can do Packet Forwarding or (Routing), so they called Default-route!

Network Manager makes you work with **Profiles**.

You can make a **Profile** to every connection of your connections and then active the **Profile** you want due to your needs!

You make a Profile only one time and active it when need it.

You can not active 2 Profiles at the same time.!

When adding a connection with the Network Manager, if you didn't select the type, it will choose **Ethernet** automatically. **autoconnect** option to choose yes or no to bring this connection up automatically, and when you re-boot to connect automatically or no, If you don't make it yes you'll need to make it manually later!!

When adding a connection and you didn't the way of Linking, it will automatically take an **IP** from the **DHCP** by default!

DHCP stands for **Dynamic Host Configuration Protocol**.

It's job to give any device connected on the Network an IP address from the range of the IPs.

If you modified a connection, the settings of it will last as it.

You need to re-activate your connection to apply the new modifications you have added. Make down and up to the connection!

****You can make more than one profile on the same Network Device!**

-- Day 18 --

