machine.js

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1 Abstract syntax

```
\rho ::= runtime lexical environment
                          \overset{\cdot}{\Delta} \quad ::= \quad \{\overline{f:(\overline{\sigma}) \to \tau}\}
                          \Gamma ::= \{\overline{x:\tau}\}
\begin{array}{rcl} capsule & ::= & \text{function } f(\overline{x}) \ \left\{ \begin{array}{c} \overline{fn} \text{ return } f; \\ \hline \\ & \mid & \text{function } f(\overline{x}) \end{array} \right\} \ \left\{ \begin{array}{c} \overline{fn} \text{ return } f; \\ \hline \\ fn & ::= & \text{function } f(\overline{x}) \end{array} \right\} \ \left\{ \begin{array}{c} \overline{x} = \kappa_x; \text{ var } \overline{y} = \overline{v}; \text{ ss } \end{array} \right\}
                       s ::= { ss }
                                       e;
                                         if (e) s
                                         if (e) s else s
                                           return v;
                                           while (e) s
                                           do s while (e);
                                           for (e; e; e) s
                                           switch (e) \{ \bar{c} \}
                                           switch (e) { \bar{c} default: ss }
                                           break;
                                           break lab;
                                           continue;
                                           continue lab;
                                            lab:s
                     ss ::= \overline{s}
                       c \ ::= \ \operatorname{case} \ e \colon ss
```

2 Type rules

```
\begin{array}{lll} \sigma,\tau & ::= & \texttt{boolean} \\ & & \texttt{uint8} \mid \texttt{uint16} \mid \texttt{uint32} \\ & & \texttt{int8} \mid \texttt{int16} \mid \texttt{int32} \\ & & \texttt{float32} \mid \texttt{float64} \\ & & \texttt{array}_\tau \mid \texttt{Function} \\ & & \texttt{any} \mid \texttt{undefined} \mid \texttt{null} \mid \texttt{string} \mid \texttt{number} \mid \texttt{object} \end{array}
```

```
\ell ::= lab \mid \epsilon
         L ::= \{\overline{\ell}\}
         \varepsilon ::= L \mid \mathsf{return}
              L; L' = L \cup L'
          \emptyset; return = return
    \{\ell, \overline{\ell'}\}; return = \{\ell, \overline{\ell'}\}
         \mathsf{return} \; ; L \;\; = \;\; \mathsf{return}
         L \cup \mathsf{return} \ = \ L
         \mathsf{return} \cup L \quad = \quad L
    \mathsf{return} \cup \mathsf{return} \quad = \quad \mathsf{return}
           type(bool) = boolean
           type(str) = string
          type(null) = null
  type(X \& Oxff) =
                              int8
type(X \& Oxffff) =
                              int16
       type(X \mid 0) = int32
    type(X >>> 0) = uint32
           type(+X) = float64
     type(X + ",") = string
     type(X + "") = string
          type(!!X) = boolean
uint8, uint16, uint32 <: int
```

int8, int16, int32 <: int

Function checking

 $\rho; \Delta \vdash fn \ \mathbf{ok}$

[T-FUNCTION]

$$\frac{\forall i.type(\kappa_{x_i}) = \sigma_i}{\rho; \Delta; \{\overline{x} : \sigma, \overline{y} : type(v)\}; \emptyset \vdash ss : \Delta(f)/\text{return}}$$
$$\frac{\rho; \Delta \vdash \text{function } f(\overline{x}) \{ \overline{x} = \kappa_x; \text{var } \overline{y} = \overline{v}; ss \} \text{ ok}}{\sigma; \Delta \vdash \text{function } f(\overline{x}) \{ \overline{x} = \kappa_x; \text{var } \overline{y} = \overline{v}; ss \} \text{ ok}}$$

Statement list checking

$$\rho; \Delta; \Gamma; L \vdash ss : \tau/\varepsilon$$

$$\frac{\forall i.\rho; \Delta; \Gamma; L \vdash s_i : \tau/\varepsilon_i \qquad \varepsilon = \varepsilon_1 ; \dots ; \varepsilon_n}{\rho; \Delta; \Gamma; L \vdash \overline{s} : \tau/\varepsilon}$$

Statement checking

$$\rho; \Delta; \Gamma; L \vdash s : \tau/\varepsilon$$

$$\begin{array}{ll} \text{[T-Block]} & & \text{[T-ExprStmt]} \\ \rho; \Delta; \Gamma; \emptyset \vdash ss : \tau/\varepsilon & & \rho; \Delta; \Gamma \vdash e : \sigma \\ \rho; \Delta; \Gamma; L \vdash \{\ ss\ \} : \tau/\varepsilon & & \rho; \Delta; \Gamma; L \vdash e; : \tau/\emptyset \end{array}$$

 $[T-IF] \hspace{1cm} [T-IFELSE]$

$$\begin{array}{ll} \rho; \Delta; \Gamma \vdash e : \mathtt{boolean} \\ \rho; \Delta; \Gamma; \emptyset \vdash s : \tau/\varepsilon \\ \varepsilon' = \varepsilon \cup \emptyset \\ \hline \rho; \Delta; \Gamma; L \vdash \mathtt{if} \ \ (e) \ \ s : \tau/\varepsilon' \end{array} \qquad \begin{array}{l} \rho; \Delta; \Gamma \vdash e : \mathtt{boolean} \\ \rho; \Delta; \Gamma; \emptyset \vdash s : \tau/\varepsilon_1 \\ \hline \varepsilon = \varepsilon_1 \cup \varepsilon_2 \\ \hline \rho; \Delta; \Gamma; L \vdash \mathtt{if} \ \ (e) \ \ s : \tau/\varepsilon \end{array}$$

[T-RETURN]

$$\rho; \Delta; \Gamma; L \vdash \mathtt{return} \ v; : type(v)/\mathsf{return}$$

$$\begin{array}{ll} \text{[T-DoWhile]} \\ \rho; \Delta; \Gamma \vdash e : \text{boolean} \\ \rho; \Delta; \Gamma; L \cup \{\epsilon\} \vdash s : \tau/\varepsilon \\ \varepsilon' = \emptyset \cup \varepsilon - (L \cup \{\epsilon\}) \\ \hline \rho; \Delta; \Gamma; L \vdash \text{while } (e) \ s : \tau/\varepsilon' \end{array} \qquad \begin{array}{ll} \text{[T-DoWhile]} \\ \rho; \Delta; \Gamma; L \cup \{\epsilon\} \vdash s : \tau/\varepsilon \\ \rho; \Delta; \Gamma \vdash e : \text{boolean} \\ \varepsilon' = \varepsilon - (L \cup \{\epsilon\}) \\ \hline \rho; \Delta; \Gamma; L \vdash \text{do } s \text{ while } (e) \text{; } : \tau/\varepsilon' \end{array}$$

$$\begin{split} & \forall i \in \{1,2,3\}.\rho; \Delta; \Gamma \vdash e_i : \sigma_i \\ & \rho; \Delta; \Gamma; L \cup \{\epsilon\} \vdash s : \tau/\varepsilon \\ & \varepsilon' = \emptyset \cup \varepsilon - (L \cup \{\epsilon\}) \\ \hline & \rho; \Delta; \Gamma; L \vdash \text{for } (e_1; \ e_2; \ e_3) \ s : \tau/\varepsilon' \end{split}$$

$$\rho; \Delta; \Gamma; L \vdash s : \tau/\varepsilon$$

$$\frac{\varepsilon = \{\epsilon\}}{\rho; \Delta; \Gamma; L \vdash \mathtt{break}; : \tau/\varepsilon}$$

[T-BreakLabel]

$$\frac{\varepsilon = \{lab\}}{\rho; \Delta; \Gamma; L \vdash \mathtt{break}\ lab; : \tau/\varepsilon}$$

[T-CONTINUE]

[T-CONTINUELABEL]

$$\rho; \Delta; \Gamma; L \vdash \text{continue}; : \tau/\emptyset$$

$$\overline{
ho;\Delta;\Gamma;L\vdash \mathtt{continue}\ lab;: au/\emptyset}$$

[T-LABEL]
$$\rho: \Delta: \Gamma: L$$

$$\frac{\rho; \Delta; \Gamma; L \cup \{lab\} \vdash s : \tau/\varepsilon}{\varepsilon' = \varepsilon - (L \cup \{lab\})}$$
$$\frac{\rho; \Delta; \Gamma; L \vdash lab \colon s : \tau/\varepsilon'}{\rho; \Delta; \Gamma; L \vdash lab \colon s : \tau/\varepsilon'}$$

[T-SWITCH]

$$\begin{split} & \rho; \Delta; \Gamma \vdash e : \sigma \\ \forall i.\rho; \Delta; \Gamma; L \cup \{\epsilon\} \vdash c_i : \sigma, \tau/\varepsilon_i \\ & \varepsilon = (\bigcup_i \varepsilon_i) - (L \cup \{\epsilon\}) \\ \hline & \rho; \Delta; \Gamma; L \vdash \text{switch } (e) ~ \{ \ \overline{c} \ \} : \tau/\varepsilon \end{split}$$

[T-SWITCHDEFAULT]

$$\begin{split} \rho; \Delta; \Gamma \vdash e : \sigma \\ \forall i. \rho; \Delta; \Gamma; L \cup \{\epsilon\} \vdash c_i : \sigma, \tau/\varepsilon_i \\ \rho; \Delta; \Gamma; L \cup \{\epsilon\} \vdash ss : \tau/\varepsilon \\ \varepsilon \neq \mathsf{return} \lor \exists i. \varepsilon_i \cup \emptyset \neq \emptyset \\ \varepsilon' = (\varepsilon \cup \bigcup_i \varepsilon_i) - (L \cup \{\epsilon\}) \end{split}$$

 $\overline{\rho; \Delta; \Gamma; L \vdash \text{switch } (e) \ \{ \ \overline{c} \ \text{default: } ss \ \} : \tau/\varepsilon'}$

[T-SWITCHRETURN]

$$\begin{array}{c} \rho; \Delta; \Gamma \vdash e : \sigma \\ \forall i. \rho; \Delta; \Gamma; L \cup \{\epsilon\} \vdash c_i : \sigma, \tau/\varepsilon_i \\ \forall i. \varepsilon_i \cup \emptyset = \emptyset \\ \rho; \Delta; \Gamma; L \cup \{\epsilon\} \vdash ss : \tau/\text{return} \\ \hline \rho; \Delta; \Gamma; L \vdash \text{switch } (e) ~ \{ \ \overline{c} \ \text{default: } ss \ \} : \tau/\text{return} \end{array}$$

Case checking

$$\rho; \Delta; \Gamma; L \vdash c : \sigma, \tau/\varepsilon$$

[T-Case]

$$\frac{\rho; \Delta; \Gamma \vdash e : \sigma}{\rho; \Delta; \Gamma; L \vdash ss : \tau/\varepsilon} \\ \frac{\rho; \Delta; \Gamma; L \vdash case \ e : ss : \sigma, \tau/\varepsilon}{\rho; \Delta; \Gamma; L \vdash case \ e : ss : \sigma, \tau/\varepsilon}$$

$$\rho; \Delta; \Gamma \vdash e : \tau$$

$$\overline{\rho; \Delta; \Gamma \vdash v : type(v)}$$

$$\frac{\Gamma\text{-VarRef}]}{\Gamma(x) = \tau} \qquad \frac{\Gamma(x) = \tau}{\rho; \Delta; \Gamma \vdash x : \tau} \qquad \frac{\Gamma(x) = \tau}{\rho; \Delta; \Gamma \vdash x = e : \tau}$$

[T-Store]

$$\begin{aligned} &type(\rho(x)) = \text{array}_{\tau} \\ &\rho; \Delta; \Gamma \vdash e : \sigma \quad \sigma <: \text{int} \end{aligned}$$

$$egin{aligned} & ype(
ho(x)) = \mathtt{array}_{ au} \ & ; \Gamma dash e : \sigma \qquad \sigma <: \mathtt{int} \ &
ho; \Delta; \Gamma dash x \llbracket e
rlaph : au \end{aligned}$$

$$\begin{array}{c} type(\rho(x)) = \texttt{array}_{\tau} \\ \rho; \Delta; \Gamma \vdash e_1 : \sigma \quad \sigma <: \texttt{int} \\ \rho; \Delta; \Gamma \vdash e_2 : \tau \\ \hline \rho; \Delta; \Gamma \vdash x \llbracket e_1 \rrbracket \ = \ e_2 : \tau \end{array}$$

[T-FFI]

$$\begin{array}{l} \text{[T-FunCall]} \\ \Delta(f) = (\overline{\sigma}) \rightarrow \tau \\ \forall i.\rho; \Delta; \Gamma \vdash e_i : \sigma_i \\ \hline \rho; \Delta; \Gamma \vdash f(\overline{e}) : \tau \end{array}$$

$$\begin{split} f \not\in dom(\Delta) & f \not\in dom(\Gamma) \\ type(\rho(f)) &= \texttt{Function} \\ \frac{\forall i.\rho; \Delta; \Gamma \vdash e_i : \sigma_i}{\rho; \Delta; \Gamma \vdash f(\overline{e}) : \texttt{any}} \end{split}$$

$$\frac{type(unop) = \sigma \to \tau \qquad \rho; \Delta; \Gamma \vdash e : \sigma}{\rho; \Delta; \Gamma \vdash unop \ e : \tau}$$

[T-BinaryOp]

$$\frac{\forall i \in \{1, 2\}.\rho; \Delta; \Gamma \vdash e_i : \sigma_i \quad type(binop) = \sigma_1 \times \sigma_2 \to \tau}{\rho; \Delta; \Gamma \vdash e_1 \ binop \ e_2 : \tau}$$

[T-CONDITIONAL]

$$\begin{array}{l} \rho; \Delta; \Gamma \vdash e_1 : \mathtt{boolean} \\ \forall i \in \{2,3\}. \rho; \Delta; \Gamma \vdash e_i : \tau \\ \hline \rho; \Delta; \Gamma \vdash e_1 ? e_2 : e_3 : \tau \end{array}$$

[T-Paren]

$$\frac{\forall i \leq n.\rho; \Delta; \Gamma \vdash e_i : \tau_i}{\rho; \Delta; \Gamma \vdash (\overline{e}) : \tau_n}$$