asm.js

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November 8, 2012

1 Abstract syntax

$$
\begin{array}{rcl}
b, e, f, g, x, y, z & \in & \textit{Identifier} \\
\texttt{arguments}, \texttt{eval} & \notin & \textit{Identifier}
\end{array}
$$

$$
\begin{array}{rcl}
P & ::= & \texttt{function } [g]\texttt{(}[e[, b]]\texttt{) \{ "use asm"; } \overline{imp_x}\ \overline{fn_f}\ exp \texttt{ \}} \\
imp_x & ::= & \texttt{var } x \texttt{ = } e\texttt{.}y\texttt{;} \\
 & | & \texttt{var } x \texttt{ = new } e\texttt{.}y\texttt{(}b\texttt{);} \\
exp & ::= & \texttt{return } f\texttt{;} \\
 & | & \texttt{return \{ } \overline{x\texttt{:}f} \texttt{ \};} \\
fn_f & ::= & \texttt{function } f\texttt{(}\overline{x}\texttt{) \{ } \overline{x \texttt{ = } \kappa_x\texttt{;}}\ \texttt{var } \overline{y \texttt{ = } v}\texttt{; } ss \texttt{ \}}
\end{array}
$$

$$
\begin{array}{rcl}
s & ::= & \texttt{\{ } ss \texttt{ \}} \\
 & | & e\texttt{;} \\
 & | & \texttt{;} \\
 & | & \texttt{if (}e\texttt{) } s \\
 & | & \texttt{if (}e\texttt{) } s \texttt{ else } s \\
 & | & \texttt{return } e\texttt{;} \\
 & | & \texttt{while (}e\texttt{) } s \\
 & | & \texttt{do } s \texttt{ while (}e\texttt{);} \\
 & | & \texttt{for (}[e]\texttt{; }[e]\texttt{; }[e]\texttt{) } s \\
 & | & \texttt{switch (}e\texttt{) \{ } \overline{c} \texttt{ \}} \\
 & | & \texttt{switch (}e\texttt{) \{ } \overline{c}\ d \texttt{ \}} \\
 & | & \texttt{break;} \\
 & | & \texttt{break } lab\texttt{;} \\
 & | & \texttt{continue;} \\
 & | & \texttt{continue } lab\texttt{;} \\
 & | & lab\texttt{: } s
\end{array}
$$

$$
ss ::= \overline{s}
$$

$$
\begin{array}{rcl}
c & ::= & \texttt{case } e\texttt{: } ss \\
d & ::= & \texttt{default: } ss \\
cd & ::= & c \mid d
\end{array}
$$

$$\kappa_x ::= \texttt{~~}x \mid \texttt{+}x \mid x\texttt{|0} \mid x\texttt{>>>0}$$

$$\begin{array}{rcl}
v & ::= & r \mid n \\
\\
e & ::= & v \\
& \mid & \mathit{lval} \\
& \mid & \mathit{lval}\ \texttt{=}\ e \\
& \mid & f\texttt{(}\overline{e}\texttt{)} \\
& \mid & \mathit{unop}\ e \\
& \mid & e\ \mathit{binop}\ e \\
& \mid & e\ \texttt{?}\ e\ \texttt{:}\ e \\
& \mid & \texttt{(}\overline{e}\texttt{)} \\
\\
\mathit{unop} & ::= & \texttt{+} \mid \texttt{~} \mid \texttt{!} \\
\\
\mathit{binop} & ::= & \texttt{+} \mid \texttt{-} \mid \texttt{*} \mid \texttt{/} \mid \texttt{\%} \\
& \mid & \texttt{|} \mid \texttt{\&} \mid \texttt{\^{}} \mid \texttt{<<} \mid \texttt{>>} \mid \texttt{>>>} \\
& \mid & \texttt{<} \mid \texttt{<=} \mid \texttt{>} \mid \texttt{>=} \mid \texttt{!=} \mid \texttt{==} \\
\mathit{lval} & ::= & x \mid x\texttt{[(}e\ \texttt{\&}\ m\texttt{)}\ \texttt{>>}\ n\texttt{]}
\end{array}$$

2 Type rules

$$\sigma, \tau ::= \texttt{bit} \mid \texttt{double} \mid \texttt{int} \mid \texttt{signed} \mid \texttt{unsigned} \mid \texttt{boolish} \mid \texttt{intish} \mid \texttt{void} \mid \texttt{unknown}$$

$$\rho ::= \tau \mid \texttt{array}^n_\tau \mid \texttt{imul} \mid ((\overline{\sigma}) \to \tau) \wedge \ldots \wedge ((\overline{\sigma'}) \to \tau') \mid \texttt{function}$$

$$\begin{array}{rcl}
\ell & ::= & \mathit{lab} \mid \epsilon \\
L & ::= & \{\overline{\ell}\} \\
\varepsilon & ::= & L \mid \mathsf{return}
\end{array}$$

$$\begin{array}{rcl}
L\ ;\ L' & = & L \cup L' \\
\emptyset\ ;\ \mathsf{return} & = & \mathsf{return} \\
\{\ell, \overline{\ell'}\}\ ;\ \mathsf{return} & = & \{\ell, \overline{\ell'}\} \\
\mathsf{return}\ ;\ L & = & \mathsf{return}
\end{array}$$

$$\begin{array}{rcl}
L \cup \mathsf{return} & = & L \\
\mathsf{return} \cup L & = & L \\
\mathsf{return} \cup \mathsf{return} & = & \mathsf{return}
\end{array}$$

$$
\begin{aligned}
\mathit{type}(\texttt{\~{}\~{}}X) &= \texttt{int} \\
\mathit{type}(\texttt{+}X) &= \texttt{double} \\
\mathit{type}(n) &= \texttt{int} \\
\mathit{type}(r) &= \texttt{double} \\
\mathit{type}(X\,\texttt{|}\,\texttt{0}) &= \texttt{signed} \\
\mathit{type}(X\texttt{>>>0}) &= \texttt{unsigned}
\end{aligned}
$$

$$
\begin{aligned}
\texttt{constant} &<: \texttt{signed}, \texttt{unsigned} \\
\texttt{signed}, \texttt{unsigned} &<: \texttt{int}, \texttt{extern} \\
\texttt{bit}, \texttt{int} &<: \texttt{boolish} \\
\texttt{double} &<: \texttt{extern} \\
\texttt{unknown}, \texttt{int} &<: \texttt{intish}
\end{aligned}
$$

$$
\begin{aligned}
M(\texttt{imul}) &: \texttt{imul} \\
M(\texttt{ceil}), M(\texttt{sin}), M(\texttt{cos}) &: (\texttt{double}) \to \texttt{double}
\end{aligned}
$$

$$
\begin{aligned}
A(\texttt{Uint8Array}), A(\texttt{Int8Array}) &= \texttt{array}^{8}_{\texttt{int}} \\
A(\texttt{Uint16Array}), A(\texttt{Int16Array}) &= \texttt{array}^{16}_{\texttt{int}} \\
A(\texttt{Uint32Array}), A(\texttt{Int32Array}) &= \texttt{array}^{32}_{\texttt{int}} \\
A(\texttt{Float32Array}) &= \texttt{array}^{32}_{\texttt{double}} \\
A(\texttt{Float64Array}) &= \texttt{array}^{64}_{\texttt{double}}
\end{aligned}
$$

$$
\begin{aligned}
\texttt{+}, \texttt{-} &: \quad (\texttt{double}, \texttt{double}) \to \texttt{double} \\
&\quad \wedge (\texttt{int}, \texttt{int}) \to \texttt{intish} \\
\texttt{*} &: \quad (\texttt{double}, \texttt{double}) \to \texttt{double} \\
\texttt{/}, \texttt{\%} &: \quad (\texttt{double}, \texttt{double}) \to \texttt{double} \\
&\quad \wedge (\texttt{signed}, \texttt{signed}) \to \texttt{intish} \\
&\quad \wedge (\texttt{unsigned}, \texttt{unsigned}) \to \texttt{intish} \\
\\
\texttt{|}, \texttt{\&}, \texttt{\^{}}, \texttt{<<}, \texttt{>>} &: \quad (\texttt{intish}, \texttt{intish}) \to \texttt{signed} \\
\texttt{>>>} &: \quad (\texttt{intish}, \texttt{intish}) \to \texttt{unsigned} \\
\\
\texttt{<}, \texttt{<=}, \texttt{>}, \texttt{>=}, \texttt{==}, \texttt{!=} &: \quad (\texttt{signed}, \texttt{signed}) \to \texttt{bit} \\
&\quad \wedge (\texttt{unsigned}, \texttt{unsigned}) \to \texttt{bit} \\
&\quad \wedge (\texttt{double}, \texttt{double}) \to \texttt{bit} \\
\\
\texttt{+} &: \quad (\texttt{intish}) \to \texttt{double} \\
\texttt{\~{}} &: \quad (\texttt{intish}) \to \texttt{signed} \\
\texttt{!} &: \quad (\texttt{boolish}) \to \texttt{bit}
\end{aligned}
$$

$$
\Gamma ::= \{\overline{x : \rho}\} \mid \Gamma, \{\overline{x : \rho}\}
$$

Program checking $\boxed{\vdash P\ \mathbf{ok}}$

[T-PROGRAM]

$$\frac{\begin{array}{c}\{\overline{x}\} \cap \{\overline{f}\} = \emptyset \qquad \{\overline{x}\} \cap \{[g],[e],[b]\} = \emptyset \qquad \{\overline{f}\} \cap \{[g],[e],[b]\} = \emptyset \\ \forall i.[e];[b];\Gamma_0 \vdash \mathit{imp}_x\ \mathbf{ok} \qquad \forall i.\Gamma_0,\Gamma_1 \vdash \mathit{fn}_f\ \mathbf{ok} \qquad \forall i.\Gamma_0,\Gamma_1 \vdash \mathit{exp}\ \mathbf{ok}\end{array}}{\vdash \texttt{function}\ [g]\texttt{(}[e[,b]]\texttt{)}\ \texttt{\{}\ \overline{\mathit{imp}_x}\ \overline{\mathit{fn}_f}\ \mathit{exp}\ \texttt{\}}\ \mathbf{ok}}$$

Import checking $\boxed{[e];[b];\Gamma \vdash \mathit{imp}\ \mathbf{ok}}$

[T-IMPORTSTD]

$$\frac{\Gamma(x) = M(y)}{e;[b];\Gamma \vdash \texttt{var}\ x\ \texttt{=}\ e\texttt{.}y\texttt{;}\ \mathbf{ok}}$$

[T-IMPORTFFI]

$$\frac{y \notin \mathit{dom}(M), \mathit{dom}(A) \qquad \Gamma(x) = \texttt{function}}{e;[b];\Gamma \vdash \texttt{var}\ x\ \texttt{=}\ e\texttt{.}y\texttt{;}\ \mathbf{ok}}$$

[T-NEWVIEW]

$$\frac{\Gamma(x) = \texttt{array}^{n}_{A(y)}}{e;b;\Gamma \vdash \texttt{var}\ x\ \texttt{=}\ \texttt{new}\ e\texttt{.}y(b)\texttt{;}\ \mathbf{ok}}$$

Function checking $\boxed{\Gamma \vdash \mathit{fn}\ \mathbf{ok}}$

[T-FUNCTION]

$$\frac{\begin{array}{c}\{\overline{x}\} \cap \{\overline{y}\} = \emptyset \qquad \Gamma(f) = (\overline{\sigma}) \to \tau \qquad \overline{\sigma} = \overline{\mathit{type}(\kappa_x)} \qquad \tau \neq \texttt{void} \\ \Gamma, \{\overline{x : \sigma}, \overline{y : \mathit{type}(v)}\}; \emptyset \vdash \mathit{ss} : \tau/\mathsf{return}\end{array}}{\Gamma \vdash \texttt{function}\ f\texttt{(}\overline{x}\texttt{)}\ \texttt{\{}\ \overline{x\ \texttt{=}\ \kappa_x\texttt{;}}\ \texttt{var}\ \overline{y\ \texttt{=}\ v}\texttt{;}\ \mathit{ss}\ \texttt{\}}\ \mathbf{ok}}$$

[T-VOIDFUNCTION]

$$\frac{\begin{array}{c}\{\overline{x}\} \cap \{\overline{y}\} = \emptyset \qquad \Gamma(f) = (\overline{\sigma}) \to \texttt{void} \qquad \overline{\sigma} = \overline{\mathit{type}(\kappa_x)} \\ \Gamma, \{\overline{x : \sigma}, \overline{y : \mathit{type}(v)}\}; \emptyset \vdash \mathit{ss} : \texttt{void}/\varepsilon\end{array}}{\Gamma \vdash \texttt{function}\ f\texttt{(}\overline{x}\texttt{)}\ \texttt{\{}\ \overline{x\ \texttt{=}\ \kappa_x\texttt{;}}\ \texttt{var}\ \overline{y\ \texttt{=}\ v}\texttt{;}\ \mathit{ss}\ \texttt{\}}\ \mathbf{ok}}$$

Export checking $\boxed{\Gamma \vdash \mathit{exp}\ \mathbf{ok}}$

[T-SINGLETON]

$$\frac{\Gamma(f) = (\overline{\sigma}) \to \tau \qquad \tau <: \texttt{extern}}{\Gamma \vdash \texttt{return}\ f\texttt{;}\ \mathbf{ok}}$$

[T-MODULE]

$$\frac{\forall f.(\Gamma(f) = (\overline{\sigma}) \to \tau \land \tau <: \texttt{extern})}{\Gamma \vdash \texttt{return}\ \texttt{\{}\ \overline{x\texttt{:}f}\ \texttt{\};}\ \mathbf{ok}}$$

Statement list checking $\boxed{\Gamma; L \vdash ss : \tau/\varepsilon}$

$$\text{[T-NoStatements]}\ \frac{}{\Gamma; L \vdash \epsilon : \tau/\emptyset} \qquad \text{[T-Statements]}\ \frac{\forall i. \Gamma; L \vdash s_i : \tau/\varepsilon_i \qquad n > 0 \qquad \varepsilon = \varepsilon_1 \,;\, \ldots \,;\, \varepsilon_n}{\Gamma; L \vdash \overline{s} : \tau/\varepsilon}$$

Statement checking $\boxed{\Gamma; L \vdash s : \tau/\varepsilon}$

$$\text{[T-Block]}\ \frac{\Gamma; \emptyset \vdash ss : \tau/\varepsilon}{\Gamma; L \vdash \texttt{\{}\ ss\ \texttt{\}} : \tau/\varepsilon} \qquad \text{[T-ExprStmt]}\ \frac{\Gamma \vdash e : \sigma}{\Gamma; L \vdash e\texttt{;} : \tau/\emptyset}$$

$$\text{[T-If]}\ \frac{\Gamma \vdash e : \texttt{boolish} \qquad \Gamma; \emptyset \vdash s : \tau/\varepsilon \qquad \varepsilon' = \varepsilon \cup \emptyset}{\Gamma; L \vdash \texttt{if (}e\texttt{)}\ s : \tau/\varepsilon'} \qquad \text{[T-IfElse]}\ \frac{\Gamma \vdash e : \texttt{boolish} \qquad \Gamma; \emptyset \vdash s_1 : \tau/\varepsilon_1 \qquad \Gamma; \emptyset \vdash s_2 : \tau/\varepsilon_2 \qquad \varepsilon = \varepsilon_1 \cup \varepsilon_2}{\Gamma; L \vdash \texttt{if (}e\texttt{)}\ s_1\ \texttt{else}\ s_2 : \tau/\varepsilon}$$

$$\text{[T-ReturnExpr]}\ \frac{\mathit{type}(e) <: \tau \qquad \Gamma \vdash e : \tau}{\Gamma; L \vdash \texttt{return}\ e\texttt{;} : \tau/\mathsf{return}} \qquad \text{[T-ReturnVoid]}\ \frac{}{\Gamma; L \vdash \texttt{return;} : \texttt{void}/\mathsf{return}}$$

$$\text{[T-While]}\ \frac{\Gamma \vdash e : \texttt{boolish} \qquad \Gamma; L \cup \{\epsilon\} \vdash s : \tau/\varepsilon \qquad \varepsilon' = \emptyset \cup \varepsilon - (L \cup \{\epsilon\})}{\Gamma; L \vdash \texttt{while (}e\texttt{)}\ s : \tau/\varepsilon'} \qquad \text{[T-DoWhile]}\ \frac{\Gamma; L \cup \{\epsilon\} \vdash s : \tau/\varepsilon \qquad \Gamma \vdash e : \texttt{boolish} \qquad \varepsilon' = \varepsilon - (L \cup \{\epsilon\})}{\Gamma; L \vdash \texttt{do}\ s\ \texttt{while (}e\texttt{);} : \tau/\varepsilon'}$$

$$\text{[T-For]}\ \frac{[\Gamma \vdash e_1 : \sigma_1] \qquad [\Gamma \vdash e_2 : \texttt{boolish}] \qquad [\Gamma \vdash e_3 : \sigma_3] \qquad \Gamma; L \cup \{\epsilon\} \vdash s : \tau/\varepsilon \qquad \varepsilon' = \emptyset \cup \varepsilon - (L \cup \{\epsilon\})}{\Gamma; L \vdash \texttt{for (}[e_1]\texttt{; }[e_2]\texttt{; }[e_3]\texttt{)}\ s : \tau/\varepsilon'}$$

Statement checking (cont'd) $\boxed{\Gamma; L \vdash s : \tau/\varepsilon}$

$$\frac{}{\Gamma; L \vdash \texttt{break;} : \tau/\{\epsilon\}}\ [\text{T-Break}] \qquad \frac{}{\Gamma; L \vdash \texttt{break}\ \ lab\texttt{;} : \tau/\{lab\}}\ [\text{T-BreakLabel}]$$

$$\frac{}{\Gamma; L \vdash \texttt{continue;} : \tau/\emptyset}\ [\text{T-Continue}] \qquad \frac{}{\Gamma; L \vdash \texttt{continue}\ \ lab\texttt{;} : \tau/\emptyset}\ [\text{T-ContinueLabel}]$$

$$\frac{\begin{array}{c}\Gamma; L \cup \{lab\} \vdash s : \tau/\varepsilon \\ \varepsilon' = \varepsilon - (L \cup \{lab\})\end{array}}{\Gamma; L \vdash lab\texttt{:}\ s : \tau/\varepsilon'}\ [\text{T-Label}]$$

$$\frac{\begin{array}{c}\Gamma \vdash e : \sigma \\ \forall i. \Gamma; L \cup \{\epsilon\} \vdash c_i : \sigma, \tau/\varepsilon_i \\ \Gamma; L \cup \{\epsilon\} \vdash cd : \sigma, \tau/\varepsilon \\ \varepsilon \neq \mathsf{return} \vee \exists i. \varepsilon_i \cup \emptyset \neq \emptyset \\ \varepsilon' = (\varepsilon \cup \bigcup_i \varepsilon_i) - (L \cup \{\epsilon\})\end{array}}{\Gamma; L \vdash \texttt{switch}\ \ \texttt{(}e\texttt{)}\ \ \texttt{\{}\ \ \overline{c}\ \ cd\ \ \texttt{\}} : \tau/\varepsilon'}\ [\text{T-Switch}]$$

$$\frac{\begin{array}{c}\Gamma \vdash e : \sigma \\ \forall i. \Gamma; L \cup \{\epsilon\} \vdash c_i : \sigma, \tau/\varepsilon_i \\ \forall i. \varepsilon_i \cup \emptyset = \emptyset \\ \Gamma; L \cup \{\epsilon\} \vdash cd : \sigma, \tau/\mathsf{return}\end{array}}{\Gamma; L \vdash \texttt{switch}\ \ \texttt{(}e\texttt{)}\ \ \texttt{\{}\ \ \overline{c}\ \ cd\ \ \texttt{\}} : \tau/\mathsf{return}}\ [\text{T-SwitchReturn}]$$

$$\frac{\Gamma \vdash e : \sigma}{\Gamma; L \vdash \texttt{switch}\ \ \texttt{(}e\texttt{)}\ \ \texttt{\{}\quad \texttt{\}} : \tau/\emptyset}\ [\text{T-EmptySwitch}] \qquad \frac{}{\Gamma; L \vdash \texttt{;} : \tau/\emptyset}\ [\text{T-EmptyStatement}]$$

Case checking $\boxed{\Gamma; L \vdash cd : \sigma, \tau/\varepsilon}$

$$\frac{\begin{array}{c}\Gamma \vdash e : \sigma \\ \Gamma; L \vdash ss : \tau/\varepsilon\end{array}}{\Gamma; L \vdash \texttt{case}\ \ e\texttt{:}\ ss : \sigma, \tau/\varepsilon}\ [\text{T-Case}] \qquad \frac{\Gamma; L \vdash ss : \tau/\varepsilon}{\Gamma; L \vdash \texttt{default:}\ ss : \sigma, \tau/\varepsilon}\ [\text{T-Default}]$$

Expression checking $\boxed{\Gamma \vdash e : \tau}$

$$\text{[T-Constant]}\quad \frac{-2^{31} \leq n < 2^{32}}{\Gamma \vdash n : \texttt{constant}} \qquad \text{[T-Double]}\quad \frac{}{\Gamma \vdash r : \texttt{double}}$$

$$\text{[T-VarRef]}\quad \frac{\Gamma(x) = \tau}{\Gamma \vdash x : \tau} \qquad \text{[T-Assign]}\quad \frac{\Gamma \vdash e : \Gamma(x)}{\Gamma \vdash x \texttt{ = } e : \tau}$$

$$\text{[T-Load]}\quad \frac{m = 2^k - 1 \qquad \Gamma(x) = \texttt{array}^n_\tau \qquad \Gamma \vdash e : \texttt{int}}{\Gamma \vdash x\texttt{[(}e \texttt{ \& } m\texttt{) >> } n/8\texttt{]} : \tau}$$

$$\text{[T-Store]}\quad \frac{m = 2^k - 1 \qquad \Gamma(x) = \texttt{array}^n_\tau \qquad \Gamma \vdash e_1 : \texttt{int} \qquad \Gamma \vdash e_2 : \tau}{\Gamma \vdash x\texttt{[(}e_1 \texttt{ \& } m\texttt{) >> } n/8\texttt{] = } e_2 : \tau}$$

$$\text{[T-IMul]}\quad \frac{\Gamma(f) = \texttt{imul} \qquad \forall i.\Gamma \vdash e_i : \texttt{intish}}{\Gamma \vdash f\texttt{(}e_1, e_2\texttt{)} : \texttt{signed}} \qquad \text{[T-FunCall]}\quad \frac{\Gamma(f) = \_ \wedge (\overline{\sigma}) \to \tau \wedge \_ \qquad \forall i.\Gamma \vdash e_i : \sigma_i}{\Gamma \vdash f\texttt{(}\overline{e}\texttt{)} : \tau} \qquad \text{[T-FFI]}\quad \frac{\Gamma(f) = \texttt{function} \qquad \forall i.\Gamma \vdash e_i : \texttt{extern}}{\Gamma \vdash f\texttt{(}\overline{e}\texttt{)} : \texttt{unknown}}$$

$$\text{[T-Conditional]}\quad \frac{\Gamma \vdash e_1 : \texttt{boolish} \qquad \Gamma \vdash e_2 : \tau \qquad \Gamma \vdash e_3 : \tau}{\Gamma \vdash e_1 \texttt{ ? } e_2 \texttt{ : } e_3 : \tau} \qquad \text{[T-Paren]}\quad \frac{\forall i \leq n.\Gamma \vdash e_i : \tau_i}{\Gamma \vdash \texttt{(}\overline{e}\texttt{)} : \tau_n}$$

$$\text{[T-Unop]}\quad \frac{\mathit{unop} : \_ \wedge (\sigma) \to \tau \wedge \_ \qquad \Gamma \vdash e : \sigma}{\Gamma \vdash \mathit{unop}\ e : \tau} \qquad \text{[T-Binop]}\quad \frac{\mathit{binop} : \_ \wedge (\sigma_1, \sigma_2) \to \tau \wedge \_ \qquad \Gamma \vdash e_1 : \sigma_1 \qquad \Gamma \vdash e_2 : \sigma_2}{\Gamma \vdash e_1\ \mathit{binop}\ e_2 : \tau}$$

$$\text{[T-Sub]}\quad \frac{\Gamma \vdash e : \sigma \qquad \sigma <: \tau}{\Gamma \vdash e : \tau} \qquad \text{[T-Cast]}\quad \frac{\Gamma \vdash e : \texttt{double}}{\Gamma \vdash \texttt{\~{}\~{}}e : \texttt{signed}}$$