

Smart Contract Program Analysis

whoami



□ Originally: security engineer at Trail of Bits

- Focuses: cryptanalysis, program analysis, mathy stuff
- Went to school for math, liked RE

Suddenly: blockchain bubble!

- Wrote a fuzzer/property based testing framework (more later!)
- Now audit smart contracts + dev tools
 - ▷ Still a good bit of the old stuff though

whatisthis



- Domain info
 - o EVM
 - Solidity
- Static Analysis
- Dynamic Analysis
 - Fuzzing
 - Symbolic Execution
- Questions

Domain Info

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Ethereum



- A state machine, in consensus
- Two types of entities
 - o People
 - "Contracts"
- □ Either make "transactions"
- Contracts run EVM code when transacted with
- Thus, "World Computer"



The Ethereum Virtual Machine



- Fairly ordinary Von Neumann machine
- Pay per operation
- Stack-based (ugh)
- Contract-local (but can see the whole chain)
- Storage and memory
- Has a notion of "reversion"
 - This transaction is messed up, don't let it change stuff
- Terribly designed
 - Example: higher memory addresses cost more to access

Interaction with the EVM



- Like hitting an API/ABI
- Defined set of function names and types
- Submit "ABI-encoded calldata" in a transaction
- This means no stdin
- Most patterns of calldata are meaningless
- Interaction looks like a web service, not a binary

Wanna know more?



- The Yellow paper is canonical
- The Jello paper might be better
- Best: download ethersplay and read some asm



The solidity language



- Solidity is the EVM lingua franca
 - Other languages exist (barely)
- It's terrible and beyond redemption
 - Too many flaws for 45 min, let alone this section
- Imagine javascript but less sensible
- Compiler has bizarre bugs somewhat frequently
 - Also just bad as a compiler
 - Computes 1 as 0x100 ** 0 sometimes
- Essentially every serious contract is written in this

Seriously, it's bad



- □ for (var i = 0; i < a.length; i ++) { a[i] = i; }</pre>
 - Infinite loop if a is larger than 255 elements
 - o i infers to byte, which is max 255
 - Overflows silently
 - Of course there's no warning here
 - i still takes up 256 bits, that's the minimum
- Stuff like this is everywhere in real code

Static Analysis

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Right now



- Currently, pretty easy
- Trail of Bits wrote a proprietary python script
- More or less a parser with grep bits
 - Iterate over functions, variables
 - Look for patterns
 - Do two things have the same name?
 - Are there uninitialized variables?
- No rich model, mostly heuristic
- Finds tons of IRL bugs
- Currently proprietary :(

The future



- Working on a lifter
 - Convert stack machine -> SSA
 - Finally get something you can really reverse
- Working on deeper binja integration
 - Use existing IL for richer analysis
 - Also you can reverse faster
- At some point, they might get a real compiler
 - Clang-style warnings supersede most of this
 - Clang-tidy could probably do it all
 - Not holding my breath though

Dynamic Analysis

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Dynamic Analysis: Surprisingly Easy



- □ Theory: solidity developers run code less
- You can't just execute, you must transact
- Ergo: running their code more finds bugs
- There exist debuggers
 - o I like hevm from dapphub
- There exist a plethora of symbolic execution tools
- There exists exactly one fuzzer



Symbolic execution: What?



- "Executing every possible path at once"
- Different notion of value
 - \circ Normal execution: values are concrete (x = 5)
 - Symbolic execution: values can be symbolic (x > 3, x < 6, x != 4)
 - Values accumulate constraints during execution
- Different notion of control flow
 - Forking on a symbolic value executes both paths in parallel
 - o Execution is a tree, not a line
- Typically solve for eventual values with a SAT solver

Applied to EVM



- Transactions are symbolic
 - Tons of symbolic data
 - Usually executors use heuristics
- Very fragile (must totally reimplement EVM)
- Often heuristic-heavy
 - Many (e.g. Oyente) assume all storage is symbolic
 - Leads to many false positives
- Great degree of assurance though
- Next steps: reliability and usability by mortals



Fuzzing



- EVM programs have no input!
 - Traditional fuzzers generate/mutate input
 - On EVM, random input is nonsensical
- Fuzzing needs an ABI and grammar
- EVM programs never segfault!
 - Traditional fuzzers know when they've found a bug
 - On EVM, more specification needed
- ${\scriptscriptstyle \square}$ ${}$ Thus, a fuzzer is just
 - ABI parser/tool to specify ABI
 - Tool to generate random type occupants
 - Random strings, random addresses, random arrays of ints, etc.
 - Some bug specification
 - VM implementation

Notes from the real world



□ Almost no programs need fancy tools

- Wrote deep VM introspection
- Wrote coverage guidance
- Wrote fancy state machine models
- None of it really mattered

You can share a fuzzer setup easily!

- o If two programs share an ABI, they fuzz the same
- Standardized ABI means massive scalability

Need more than just inputs for many real programs

- Temporal logic
- Complex multi-contract setup
- Many actors of different privilege levels
- Contracts can change mid-fuzz!
- All of these are next steps



Questions?

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