# Using Binary Instrumentation for Vulnerability Discovery (or even mitigation!)

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#### Binary Instrumentation and Vulnerability Discovery

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- Writing a bug finding tool
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#### What is instrumentation?

 According to Wikipedia, it's the ability to monitor or measure the level of a product's performance, to diagnose errors and to write trace information.

#### There are 2 types of instrumentation approaches:

- Source instrumentation. Not possible to use for auditing closed sourced products.
- Binary instrumentation. An approach available for any software as long as we can execute it.

Binary Instrumentation is often called Dynamic Binary Instrumentation (DBI) as it instruments a program while executing it at binary level.

- Basically, it inserts instrumentation code before and after a piece of code (a basic block) is executed.
- Instrumentation can be inserted at a user-defined granularity level like:
  - o Instruction level. For every single assembly instruction.
  - Basic block level. For each basic block that is discovered by the DBI tool.
  - Function level. Only at the beginning, usually, of each function the DBI tool discovered.

As DBI lets us go as deep as we want to instrument a program, we can monitor the full execution of a program at whatever granularity level we want. This is great, but it also has some drawbacks:

- A program instrumented with a DBI tool will always run slower than running the program alone.
- Depending on the granularity we choose and the amount of instrumentation inserted, the program can go so slow that DBI turns out to be impossible to use. Or at least turn out to be "not so useful".
  - Some examples I have tried: Oracle database and modern PC games.

#### Introduction: Tools

There are various DBI toolkits out there but I'll just talk about the most commonly used ones:

- DynamoRIO. Open Source and very fast. However, tends to fail for not an easy reason to understand and there are softwares that you simply cannot instrument with it. Besides, the API is neither easy nor pretty, in my opinion.
- Intel PIN. Commercial software with a ridiculous license and very slow.
   However, the only DBI tool that is reliable and it has a super easy API.

Other DBI tools that I haven't used yet:

 Frida. Open Source and pretty fast. However, it seems not to support statically compiled binaries.

## DynamoRIO

#### Introduction: DynamoRIO

Example DynamoRIO instrumentation tool (part I):

```
uint num dyn instrs;
static void event init(void);
static void event exit(void);
static dr emit flags t event basic block(void *drcontext, void *tag, instrlist t *ilist,
                                            bool for trace, bool translating);
DR EXPORT void dr init(client id t id) {
  /* register events */
  dr_register_bb_event (event_basic_block);
  dr_register_exit_event(event_exit);
  /* process initialization event */
  event_init();
```

#### Introduction: DynamoRIO

Example DynamoRIO instrumentation tool (part II):

```
static void event_init(void) {
  num dyn instrs = 0;
static void event_exit(void) {
  dr_printf("Total number of instruction executed: %u\n", num_dyn_instrs);
static dr_emit_flags_t_event_basic_block(void *drcontext, void *tag, instrlist_t *ilist,
                                            bool for_trace, bool translating) {
  int num_instrs;
  num_instrs = ilist_num_instrs(ilist);
  insert count code(drcontext, ilist, num_instrs);
  return DR EMIT DEFAULT;
```

#### Introduction: DynamoRIO

Example DynamoRIO instrumentation tool (part III):

```
static int ilist num instrs(instrlist t *ilist) {
  instr t *instr;
          num instrs = 0;
  int
  /* iterate over instruction list to count number of instructions */
  for (instr = instrlist_first(ilist); instr != NULL; instr = instr_get_next(instr))
     num instrs++;
  return num instrs;
static void do ins count(int num_instrs) { num_dyn_instrs += num_instrs; }
static void insert count code(void * drcontext, instrlist t * ilist, int num instrs) {
  dr_insert_clean_call(drcontext, ilist, instrlist_first(ilist),
                        do_ins_count, false, 1,
                         OPND_CREATE_INT32(num_instrs));
```

## Intel PIN

#### Introduction: Intel PIN

Example Intel PIN instrumentation tool (part I):

```
/* Print Help Message
   -----*/
INT32 Usage()
   cerr << "This tool counts the number of dynamic instructions executed" << endl;
   cerr << endl << KNOB BASE::StringKnobSummary() << endl;</pre>
   return -1:
    argc, argv are the entire command line: pin -t <toolname> -- ...
/* -----*/
int main(int argc, char * argv[])
   // Initialize pin
   if (PIN Init(argc, argv)) return Usage();
   OutFile.open(KnobOutputFile.Value().c str());
   // Register Instruction to be called to instrument instructions
   INS AddInstrumentFunction(Instruction, 0);
   // Register Fini to be called when the application exits
   PIN AddFiniFunction(Fini, 0);
   // Start the program, never returns
   PIN StartProgram();
   return Θ;
```

#### Introduction: Intel PIN

#### Example Intel PIN instrumentation tool (part II):

```
ofstream OutFile:
// The running count of instructions is kept here
// make it static to help the compiler optimize docount
static UINT64 icount = 0;
// This function is called before every instruction is executed
VOID docount() { icount++; }
// Pin calls this function every time a new instruction is encountered
VOID Instruction(INS ins, VOID *v)
    // Insert a call to docount before every instruction, no arguments are passed
    INS InsertCall(ins, IPOINT BEFORE, (AFUNPTR)docount, IARG END);
KNOB<string> KnobOutputFile(KNOB MODE WRITEONCE, "pintool",
    "o", "inscount.out", "specify output file name");
// This function is called when the application exits
VOID Fini(INT32 code, VOID *v)
   // Write to a file since cout and cerr maybe closed by the application
   OutFile.setf(ios::showbase):
    OutFile << "Count " << icount << endl;
    OutFile.close();
```

## Introduction: Intel PIN and DynamoRIO

As we can see, it's way easier to write an Intel PIN tool than a DynamoRIO tool.

However, the Intel PIN license is a "no go" for many use cases:

• Example: do you want to distribute a binary build? You can't.

The choice of a DBI toolkit will largely depend on what do you want it for.

- If you want to make a commercial tool or distribute binary versions, use DynamoRIO. If it works in your case.
- If you want to use a FOSS tool, use DynamoRIO. Again, if it works for you.
- For anything else, use Intel Pin.

#### Introduction: Intel PIN

Because of its stability and reliability, most reverse engineers use Intel PIN.

After all, most of the time you just want something for your specific use case and most likely aren't going to distribute anything because is of no value for others.

And in the case of distributing the tool, most reverse engineers are fine with distributing source code. Even companies doing commercial products:

For example: IDA's Intel Pin debugger module is distributed in source form.

For the example tool I've written for this talk, I will use Intel PIN.

DBI tools are very handy for writing bug finding tools. We can instrument at whatever granularity level we want and we can pretty much instrument/monitor anything the application does.

In my case, I have written an Intel PIN tool with the aim of detecting (and even mitigating!) the following common memory management bugs:

- Double free-s.
- Invalid free-s.
- Use-after-free-s.
- Writes to freed memory.

What we have to do in order to find such bugs?

- Find the dynamic memory allocation functions and hook them.
- Monitor how the buffers are manipulated.
- Report when a vulnerability is found.
- ...
- Profit!

However, is not that simple and little problems appear while coding such tools.

During this talk I will talk about some of them.

As I said, we first need to hook (instrument) the allocation functions. Which functions are they?

- Malloc
- Calloc
- Realloc
- Free

How can we do hook these functions using Intel Pin?

- First, we will need to check which module exports these functions.
- And then, just instrument these functions.

The memory allocator functions will be "stored" in some specific module.

• Like the libc, in most cases.

We will need to find the appropriate module and hook these functions. How?

We will need to install a callback to instrument "image loading":

```
int main(int argc, char *argv[])
 // Initialize symbols
 PIN InitSymbols();
 // Initialize PIN library. Print help message if -h(elp) is specified
 // in the command line or the command line is invalid
 if ( PIN Init(argc,argv) )
   return usage();
  q mitigate = (knob mitigate.Value() != 0);
  if ( knob memory tracker.Value() != 0 )
   q hooks = new CHooksInstaller();
   g checker = new CMemoryChecker();
   q checker->break always = (knob break always.Value() != 0);
  // Register Image to be called to instrument functions.
  IMG AddInstrumentFunction(image cbk, 0);
 IMG AddUnloadFunction(image unload cbk, 0);
  // Register function to be called to instrument instructions
 INS AddInstrumentFunction(ins cbk, 0);
  // Register function to be called when the application exits
 PIN AddFiniFunction(fini cbk, 0);
 // Start the program, never returns
 PIN StartProgram();
  return 0;
```

Then, we will need to install our hooks in the appropriate module (part I):

```
//-
static VOID image_cbk(IMG img, VOID *v)

{
   if ( g_hooks != NULL )
      g_hooks->install_hooks(img);

   if ( g_checker != NULL )
      g_checker->add_image(img);
}
```

Then, we will need to install our hooks in the appropriate module (part II):

```
void CHooksInstaller::install hooks(IMG img)
 str vec t malloc funcs:
 malloc funcs.push back("malloc");
 malloc funcs.push back(" malloc");
 str vec t calloc funcs;
 calloc funcs.push back("calloc");
 calloc funcs.push back(" calloc");
 str vec t free funcs;
 free funcs.push back("free");
 free funcs.push back(" free");
 str vec t realloc funcs:
 realloc funcs.push back("realloc");
 realloc funcs.push back(" realloc");
 hook functions(img, malloc funcs, FTT MALLOC);
 hook functions(img, calloc funcs, FTT CALLOC);
 hook functions(img, free funcs, FTT FREE);
 hook functions(img, realloc funcs, FTT REALLOC);
 hook memory access():
```

Then, we will need to install our hooks in the appropriate module (part III):

```
//-
void CHooksInstaller::hook_functions(
   IMG img,
   const str_vec_t &funcs,
   FUNC_TYPE_T type)
{
   str_vec_t::const_iterator it;
   str_vec_t::const_iterator end = funcs.end();
   for ( it = funcs.begin(); it != end; ++it )
   {
      if ( hook_one_function(img, (*it).c_str(), type) )
          break;
   }
}
```

Then, we will need to install our hooks in the appropriate module (part IV):

```
bool CHooksInstaller::hook one function(
 IMG img,
 const char *func,
 FUNC TYPE T type)
 bool ret = false;
 RTN target rtn = RTN FindByName(img, func);
 if (RTN Valid(target rtn))
   RTN Open(target rtn);
    RTN InsertCall(target rtn, IPOINT BEFORE, (AFUNPTR)mem before cbk,
                   IARG FAST ANALYSIS CALL.
                  IARG FUNCARG ENTRYPOINT VALUE, 0,
                   IARG FUNCARG ENTRYPOINT VALUE, 1,
                   IARG ADDRINT, type,
                   IARG FUNCARG ENTRYPOINT REFERENCE, 0,
                   IARG END);
    RTN InsertCall(target rtn, IPOINT AFTER, (AFUNPTR)mem after cbk,
                   IARG FAST ANALYSIS CALL,
                   IARG FUNCRET EXITPOINT VALUE,
                   IARG ADDRINT, type,
                   IARG END);
    RTN Close(target rtn);
    ret = true;
  return ret;
```

In the previous slide we simply search for the specific functions to be hooked (malloc, \_malloc, free, \_free, etc...) using RTN\_FindByName.

If any of these functions is found, we 'hook' it by RTN\_InsertCall.

In my example tool, I'm doing it \*before\* and \*after\* the functions are called.

- The arguments to the functions will only be available before the functions are called.
- The result of the functions will only be available after the functions are called.

Then, in the callbacks we just installed we will handle the casuistic for the 2 functions we're going to support:

- malloc
- free

Well, we're actually hooking 4 functions (malloc, calloc, realloc and free) but we can consider the other 2 functions just specialized versions of malloc.

Handling the (m|c)alloc/free functions from the instrumentation callback (part I):

```
static void PIN FAST ANALYSIS CALL mem before cbk(
 ADDRINT al,
 ADDRINT a2,
 FUNC TYPE T type,
  ADDRINT *ref al)
 if ( type == FTT MALLOC )
   if ( g checker != NULL )
      q checker->check before malloc(a1, type);
 else if ( type == FTT CALLOC )
   if ( g checker != NULL )
      g checker->check before malloc(a1 * a2, type);
 else if ( type == FTT FREE )
   if ( g checker != NULL )
      q checker->check before free(a1, a2, type, ref a1);
```

Handling the (m|c)alloc/free functions from the instrumentation callback (part II):

```
void CMemoryChecker::check before malloc(size t size, FUNC TYPE T type)
 PIN LockClient();
 if ( (signed) size < 0 )
   printf("WARNING! Negative size given to a malloc call!\n");
   add breakpoint("Negative size given to a malloc call");
 else if ( size == 0 )
   printf("WARNING! Zero allocation detected!\n");
    add breakpoint("WARNING! Zero allocation detected!\n");
 mem area t area;
 area.size = size;
 area.available = true;
 area.ignore = false;
 area.ignore write = false;
 area.tid = PIN GetTid();
 areas.push back(area);
 PIN UnlockClient();
```

In the class method CMemoryChecker::check\_before\_malloc() we're finally adding our first rules to find vulnerabilities.

By just checking the argument "size" (of type size\_t) given to malloc/calloc we can find 2 different but related vulnerability types:

- Negative mallocs. For example, in 32 bits platforms malloc(-1) will try to allocate 4 GB of memory, which is almost always wrong, will most likely fail and is a clear sign of a vulnerability.
- **Zero allocations**. A zero allocation usually returns a valid pointer with as much memory as the size of a pointer in the specific architecture. If more than the reserved bytes are written to the buffer, a heap overflow happens.

One interesting thing in the previously shown code is the calls to add\_breakpoint.

This function sets a flag to call the API PIN\_ApplicationBreakpoint from an installed instruction level instrumentation callback.

We can connect from GDB, IDA or Visual Studio to the Intel PIN remote debugger and when PIN\_ApplicationBreakpoint is called, a breakpoint is triggered in it.

Executions stops exactly at the point where the violation is made and one can check the trace back, arguments, etc...

By just hooking malloc and calloc functions we have added support to dynamically discover 2 bug classes.

For discovering use-after-frees, double-frees, etc... we need to do some more things:

- Basically, keep a vector of all the allocations, the memory address returned and their sizes.
- When free is called, check if the pointer is valid, was previously freed, etc...

Let's see the implementation in my example tool...

#### Finding double frees (part I):

```
void CMemoryChecker::check before malloc(size t size, FUNC TYPE T type)
  PIN LockClient();
  if ( (signed) size < 0 )
    printf("WARNING! Negative size given to a malloc call!\n");
    add breakpoint("Negative size given to a malloc call");
  else if ( size == 0 )
    printf("WARNING! Zero allocation detected!\n");
    add breakpoint("WARNING! Zero allocation detected!\n");
  mem area t area;
  area.size = size:
  area.available = true:
  area.ignore = false;
  area.ignore write = false;
  area.tid = PIN GetTid();
  areas.push back(area);
  PIN UnlockClient();
```

#### Finding double frees (part II):

```
void CMemoryChecker::check after malloc(ADDRINT ret, FUNC TYPE T type)
 PIN LockClient();
 mem area vec t::reverse iterator it;
 for ( it = areas.rbegin(); it != areas.rend(); ++it )
   if ( it->tid == PIN GetTid() )
     ADDRINT ea = it->ea:
     if ( ea == ret || (ret >= ea && ret <= (ea + it->size)) )
       areas.erase((it+1).base()):
 for ( it = areas.rbegin(); it != areas.rend(); ++it )
   if ( it->tid == PIN GetTid() )
     if ( ret == 0 )
       areas.erase((it+1).base());
     else
       it->ea = ret;
     break:
 PIN UnlockClient();
```

#### Finding double frees (part III):

```
void CMemoryChecker::check before free(ADDRINT al,
 ADDRINT a2.
 FUNC TYPE T type.
 ADDRINT *ref al)
 PIN LockClient();
 bool found = false:
 mem area vec t::reverse iterator it;
 for ( it = areas.rbegin(); it != areas.rend() && al != 0; ++it )
   if ( it->tid == PIN GetTid() && it->ea == al )
      found = true:
     if (!it->available)
       printf("WARNING! Freeing already available memory (0x" EA FMT " - 0x" EA FMT ")!\n", it->ea, it->ea + it->size);
       dump areas();
#endif
       if ( break always )
         add breakpoint("Freeing already available memory!");
       if ( g mitigate )
          printf("NOTICE: Mitigation is enabled, returning a null pointer...\n");
          *ref a1 = 0:
        it->ignore = true;
     it->available = false:
     break:
```

In the last code slide we can see how we can detect double frees and stop at the debugger when we detect one by triggering a breakpoint.

But we can even mitigate such a bug!

- Well, kind of. We just send a null pointer to free.
- It's documented that free will \*NOT\* fail when given a null pointer and the execution should continue "happily".
- However, where there is a bug, there can be others, and the application can crash at some other, probably half-random, point.

We can also detect other memory related bugs involving calls to free:

• Freeing invalid pointers.

From an Intel PIN tool we can check if an address is valid so, in addition to detect double frees, we can detect invalid frees.

Let us see how...

#### Finding invalid frees:

```
if ( !found && al != 0 )
   // Trick to avoid false positives: try to read 4 bytes from that
   // memory page we don't have information about.
   char buf[4]:
   size t bytes = PIN SafeCopy(buf, (ADDRINT*)al, sizeof(buf));
   if (bytes < 4)
     printf("WARNING! Freeing an invalid memory page at 0x" EA FMT "!\n", a1);
    add_breakpoint("Freeing an invalid memory page");
     if ( q mitigate )
       printf("NOTICE: Mitigation is enabled, returning a null pointer #2...\n");
       *ref al = 0;
   else
     //printf("WARNING! An unknown memory page has been freed!\n");
     mem area t area;
     area.ea = al:
     area.size = sizeof(void*);
     area.available = false:
     area.ignore = false;
     area.ignore write = false;
     area.tid = PIN GetTid();
     areas.push back(area);
// The libc actually needs to write to freed pages so, while inside
 // the "free" implementation, ignore any error.
inside alloc[PIN ThreadId()] = true;
PIN UnlockClient();
```

As we can see, we can detect invalid frees and, as before, we can even kind of mitigate them.

• Same problems as with double frees apply, though.

Now, time to detect even more bug types: write-after-frees.

Basically, we will detect when a freed page is being written to.

Let's see how...

Detecting writes to freed pages (part I):

```
static void ins cbk(INS ins, void *v)
 INS InsertCall(ins, IPOINT BEFORE, (AFUNPTR)ins instruction cbk,
                 IARG FAST ANALYSIS CALL, IARG CONST CONTEXT,
                 IARG END);
 // Iterate over each memory operand of the instruction.
 UINT32 mem ops = INS MemoryOperandCount(ins);
  for ( UINT32 mem op = \theta; mem op < mem ops; mem op++)
    if ( INS MemoryOperandIsWritten(ins, mem op) )
      INS InsertPredicatedCall(
          ins, IPOINT_BEFORE, (AFUNPTR) memory access callback,
          IARG INST PTR.
          IARG MEMORYOP EA, mem op,
          IARG END);
```

Detecting writes to freed pages (part II):

```
//----
static VOID memory_access_callback(VOID *ip, VOID *addr)
{
   if ( g_checker != NULL )
       g_checker->check_write((ADDRINT)addr, 0);
}
```

#### Detecting writes to freed pages (part III):

```
size t CMemoryChecker::check write(ADDRINT ea, size t size)
 PIN LockClient();
 bool in alloc = is in allocator();
  size t ret = size:
  bool found = false:
  mem area vec t::iterator it;
  for ( it = areas.end()-1; it != areas.begin() && ea != 0; --it )
   if ( ea >= it->ea && ea <= (it->ea + it->size) )
     found = true:
     if (!it->available && !it->ignore write)
       if (!in alloc)
          // Only report once per area
         it->ignore write = true;
          add breakpoint("Writing into a freed page");
          printf("WARNING: Writing into a freed page, write to 0x" EA FMT
#if 0
          dump areas();
#endif
     break:
```

Once the malloc/realloc/free tracker is working, it's possible to add other rules to detect other types of bugs like, for example, calls to realloc passing an already freed pointer.

However, getting the tracker to work is... non trivial. And most likely, specific to each heap allocator:(

Let's talk about the problems of this approach to find bugs...

The most obvious problem of using this approach for finding bugs is the following:

A dynamic analyzer will only analyse whatever is executed.

If there is a bug in function do\_buggy\_things() but with the inputs given to the program we don't reach that portion of code, we will never see that bug.

Writing a malloc/realloc/free tracker is hard. It might sound easy, but is not. Some examples:

- Inside free(), the heap manager might decide to do "things" with more than a single page.
- Inside malloc/realloc(), the heap manager might decide to merge previously freed pages into one single page.
- Detecting writes to freed memory means having to determine when is it done by the heap manager and when is it done by the application.

Such a tool requires to have a perfectly working tracker or most of the bugs will be false positives otherwise.

Another problem of this approach is that the tracker is most likely specific to each heap manager that we want to support.

Each heap manager will behave differently.

There are many heap managers out there:

- Windows' own heap, specific for each Windows version.
- The GNU allocator.
- Jemalloc.
- ...

Some of them aren't generic but program specific, actually.

Another problem inherent to such tools is analysing the outputs of it.

Analyzing bugs found with such a tool can be summarized, often, to this message from a friend of mine:



Some bugs are pretty easy to analyse:

Zero allocations. You can break in the debugger at malloc/realloc and check the size.
 As easy as it sounds.

Others aren't that easy to analyse:

- An example: a call to realloc passing a freed pointer that was freed in a totally different part of the code.
- Another example, a write to a previously freed memory page: Is it the heap allocator itself? Did the allocator reused that page and is available now? Where was the freed memory page first reserved, where it was freed and what is the relationship with the final write we observed?

# And some examples of discovered bugs

Time for a DEMO!

## Example bugs found

A zero allocation found in radare2 (fixed in <2 minutes because... pancake):

```
WARNING! Zero allocation detected!

Program received signal SIGTRAP, Trace/breakpoint trap.
__GI___libc_malloc (bytes=0) at malloc.c:2902
2902 malloc.c: No such file or directory.
(gdb) back
#0 __GI___libc_malloc (bytes=0) at malloc.c:2902
#1 0x00007f0c3adacd3f in cons_stack_dump (recreate=true) at cons.c:71
#2 0x00007f0c3adadf16 in r_cons_push () at cons.c:572
#3 0x00007f0c3b52b87f in r_core_cmd_str (core=0x55c31a823580, cmd=0x55c31a6220e8 "ieq") at cmd.c:3648
```

### Example bugs found

WARNING! Zero allocation detected!

A zero allocation found in NVIDIA 390 user-land libraries:

## Conclusions

### Conclusions

DBI toolkits are extremely useful tools for analyzing programs and discovering vulnerabilities.

However, while it might look easy at first to write such a tool, it's actually pretty hard.

Analysing the resulting bug candidates is very time consuming and automating it is not trivial.

Writing a DBI tool by itself is not enough: a fuzzer, for example, is required to discover bugs hidden deep in the code.

Using a DBI tool with a coverage guided fuzzer looks like a very good idea.

#### Conclusions

From an exploitability point of view, most of the bugs found are not interesting:

- Bugs at start-up.
- Bugs, like zero allocations, that cannot be exploited.

Also, most of the software we use is so plagued with bugs that the amount of bugs that such a tool can discover in a normal application is too big to analyze.

All of that said, however, perhaps such a tool can be mixed with other program analysis techniques at least to lift the results and determine what looks like interesting and what doesn't. I haven't explored this option yet.

### Questions?

And that's it! The source code of MemBugTool (call me original...) is available here:

http://github.com/joxeankoret/membugtool