# Chapter 6: Architecture

# **Recursive Functions**

- Function that calls itself
- When converting to assembly code:
- In the first pass, treat recursive calls as if it's calling a different function and ignore overwritten registers.
- Then save/restore registers on stack as needed.

#### Factorial function:

```
- factorial(n) = n!
= n*(n-1)*(n-2)*(n-3)...*1
```

```
- Example: factorial(6) = 6!
= 6*5*4*3*2*1
= 720
```

# High-Level Code int factorial(int n) { if (n <= 1) return 1; else return (n\*factorial(n-1)); Thus, }</pre> factorial(2): returns 3\*factorial(2) factorial(2): returns 2\*factorial(1) factorial(1): returns 1 factorial(2): returns 2\*1 = 2 factorial(3): returns 3\*2 = 6

```
High-Level Code
                                    RISC-V Assembly
                                    factorial:
int factorial(int n) {
  if (n <= 1)
    return 1;
  else
    return (n*factorial(n-1));
Pass 1. Treat as if calling another
function. Ignore stack.
Pass 2. Save overwritten registers
(needed after function call) on the
stack before call.
```

#### High-Level Code

```
int factorial(int n) {
  if (n <= 1)
    return 1;
 else
```

Pass 1. Treat as if calling another function. Ignore stack.

Pass 2. Save overwritten registers (needed after function call) on the stack before call.

#### RISC-V Assembly

```
factorial:
                          addi sp, sp, -8 # save regs
                          sw a0, 4(sp)
                          sw ra, 0(sp)
                          addi t0, zero, 1 # temporary = 1
                          bgt a0, t0, else # if n>1, go to else
                          addi a0, zero, 1 # otherwise, return 1
                          addi sp, sp, 8 # restore sp
                          jr
                                        # return
                               ra
                         else:
return (n*factorial(n-1)); addi a0, a0, -1 # n = n - 1
                          jal factorial # recursive call
                          lw t1, 4(sp) # restore n into t1
                          lw ra, 0(sp) # restore ra
                          addi sp, sp, 8 # restore sp
                          mul a0, t1, a0 # a0=n*factorial(n-1)
                          jr
                                           # return
                               ra
```

Note: n is restored from stack into t.1 so it doesn't overwrite return value in a0.

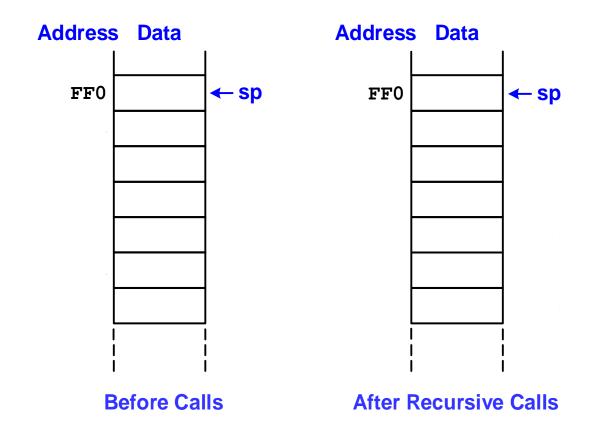
#### Recursive Functions

```
0x8500 factorial: addi sp, sp, -8 # save registers
0x8504
                  sw a0, 4(sp)
0x8508
                  sw ra, 0(sp)
0x850C
                  addi t0, zero, 1 # temporary = 1
0 \times 8510
                  bqt a0, t0, else \# if n > 1, qo to else
0 \times 8514
                  addi a0, zero, 1 # otherwise, return 1
0 \times 8518
                  addi sp, sp, 8 # restore sp
                                 # return
0x851C
                  jr
                       ra
0x8520 else: addi a0, a0, -1 # n = n - 1
                  jal factorial # recursive call
0 \times 8524
0 \times 8528
                  lw t1, 4(sp) # restore n into t1
                  lw ra, 0(sp) # restore ra
0 \times 852C
0 \times 8530
                  addi sp, sp, 8 # restore sp
0 \times 8534
                  mul a0, t1, a0 \# a0 = n*factorial(n-1)
0 \times 8538
                  jr
                                  # return
                       ra
```

PC+4 = 0x8528 when factorial is called recursively.

## Stack During Recursive Function

When **factorial(3)** is called:



## Chapter 6: Architecture

# More on Jumps & Pseudoinstructions

#### Jumps

- RISC-V has two types of unconditional jumps
  - Jump and link (jal rd, imm<sub>20:0</sub>)
    - rd = PC+4; PC = PC + imm
  - jump and link register (jalr rd, rs,  $imm_{11:0}$ )
    - rd = PC+4; PC = [rs] + SignExt(imm)

#### Pseudoinstructions

- Pseudoinstructions are not actual RISC-V instructions but they are often more convenient for the programmer.
- Assembler converts them to real RISC-V instructions.

#### Jump Pseudoinstructions

RISC-V has four jump psuedoinstructions

```
-j imm jal x0, imm
-jal imm jal ra, imm
-jr rs jalr x0, rs, 0
-ret jr ra (i.e., jalr x0, ra, 0)
```

#### Labels

- Label indicates where to jump
- Represented in jump as immediate offset
  - imm = # bytes past jump instruction
  - In example, below, **imm** = (51C-300) = 0x21C

```
-jal simple = jal ra, 0x21C
```

#### RISC-V assembly code

#### Long Jumps

- The immediate is limited in size
  - 20 bits for jal, 12 bits for jalr
  - Limits how far a program can jump
- Special instruction to help jumping further
  - auipc rd, imm: add upper immediate to PC
    - rd = PC +  $\{imm_{31:12}, 12'b0\}$
- Pseudoinstruction: call imm<sub>31:0</sub>
  - Behaves like jal imm, but allows 32-bit immediate offset

```
auipc ra, imm_{31:12} jalr ra, ra, imm_{11:0}
```

#### More RISC-V Pseudoinstructions

Pseudoinstruction	RISC-V Instructions
j label	jal zero, label
jr ra	jalr zero, ra, 0
mv t5, s3	addi t5, s3, 0
not s7, t2	xori s7, t2, -1
nop	addi zero, zero, 0
li s8, 0x56789DEF	lui s8, 0x5678A
	addi s8, s8, 0xDEF
bgt s1, t3, L3	blt t3, s1, L3
bgez t2, L7	bge t2, zero, L7
call L1	auipc ra, imm <sub>31:12</sub>
	jalr ra, ra, imm $_{11:0}$
ret	jalr zero, ra, 0

See Appendix B for more pseudoinstructions.