

ECCS-3631

Networks and Data Communications

Module 2-1

Network Layer, Routing Fundamentals, Link State Routing Algorithm

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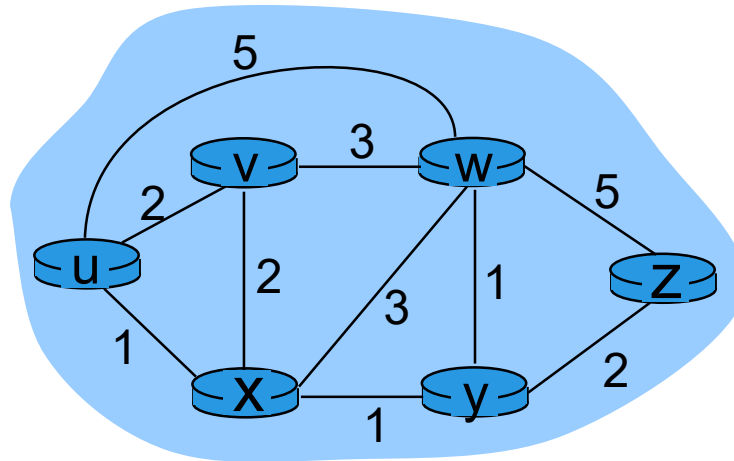
Routing Goal

Routing goal: determine “good” paths (equivalently, routes), from sending hosts to receiving host, through network of routers

path: sequence of routers packets will traverse in going from given initial source host to given final destination host

“*good*”: least “cost”, “fastest”, “least congested”

Graph abstraction of the network

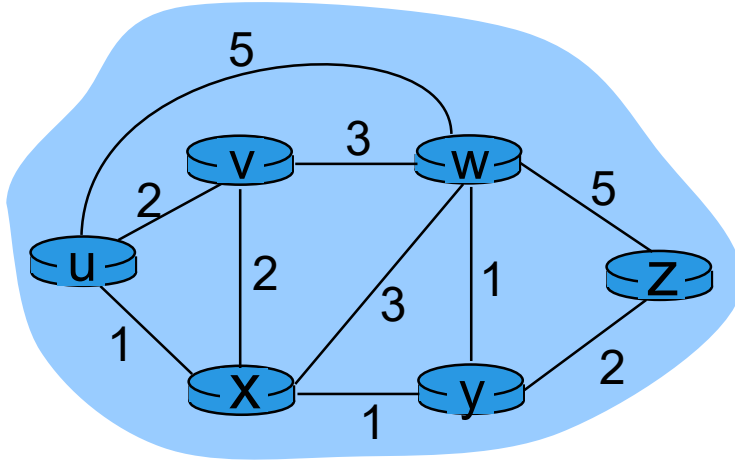


graph: $G = (N, E)$

$N = \text{set of routers} = \{ u, v, w, x, y, z \}$

$E = \text{set of links} = \{ (u,v), (u,x), (v,x), (v,w), (x,w), (x,y), (w,y), (w,z), (y,z) \}$

Graph abstraction: costs



$c(x, x') = \text{cost of link } (x, x')$

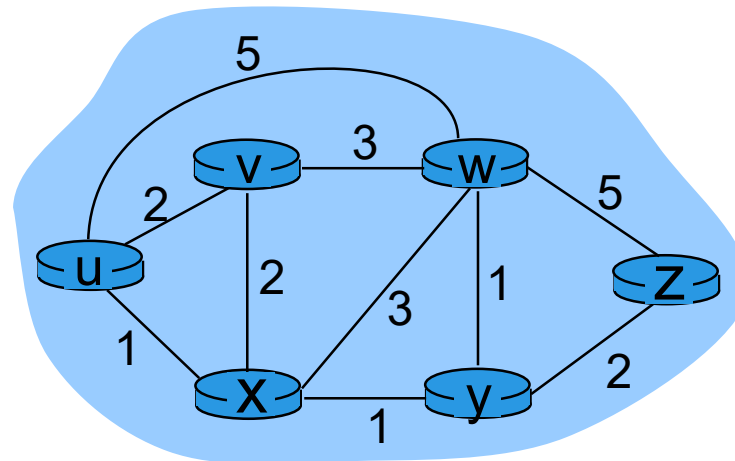
e.g., $c(w, z) = 5$, $c(u, v) = 2$

cost could always be 1, or
inversely related to bandwidth, or
inversely related to congestion

cost of path $(x_1, x_2, x_3, \dots, x_p) = c(x_1, x_2) + c(x_2, x_3) + \dots + c(x_{p-1}, x_p)$

Review Question

- What is the least-cost path between u and z?
- How to find that least-cost path?



The least-cost path between u and z is 4.

Routing Algorithms are used to find that least-cost path.

Network Routing: Algorithms & Protocols

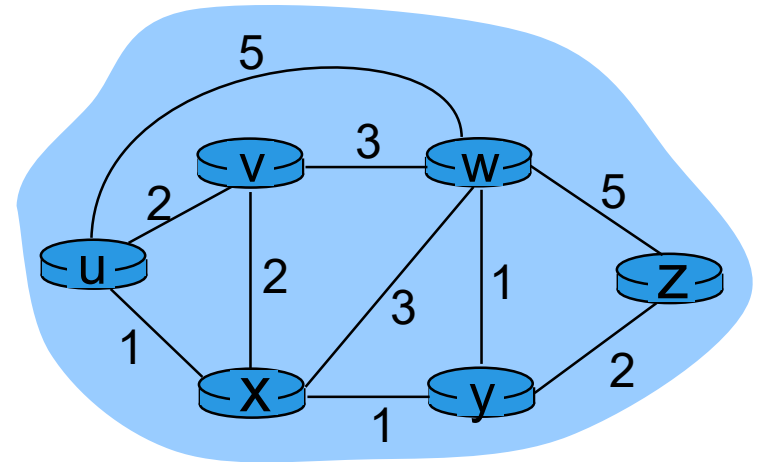
Routing Algorithms

Link-State (Dijkstra)

- each router knows complete topology & link cost information
- Run routing algorithm to calculate shortest path to each destination

Distance-Vector (Bellman-Ford)

- Each router knows direct neighbors & link costs to neighbors
- Calculate the shortest path to each destination through an *iterative process based on the neighbors distances* to each destination



Routing Protocols

- ❑ define the format of routing information exchanges
- ❑ define the computation upon receiving routing updates
- ❑ network topology changes over time, routing protocol must continuously update the routers with latest changes

Link-State Routing Algorithm

Dijkstra's Algorithm

net topology, link costs known to all nodes

- accomplished via “link state broadcast”
- all nodes have same info

computes least cost paths from one node (‘Source’) to all other nodes

- Creates *forwarding table* for that node

iterative: after k iterations, source knows least cost path to k destinations

notation:

$c(x,y)$: link cost from node x to y ; (∞ if not direct neighbors)

$D(v)$: current value of cost of path from source to destination v

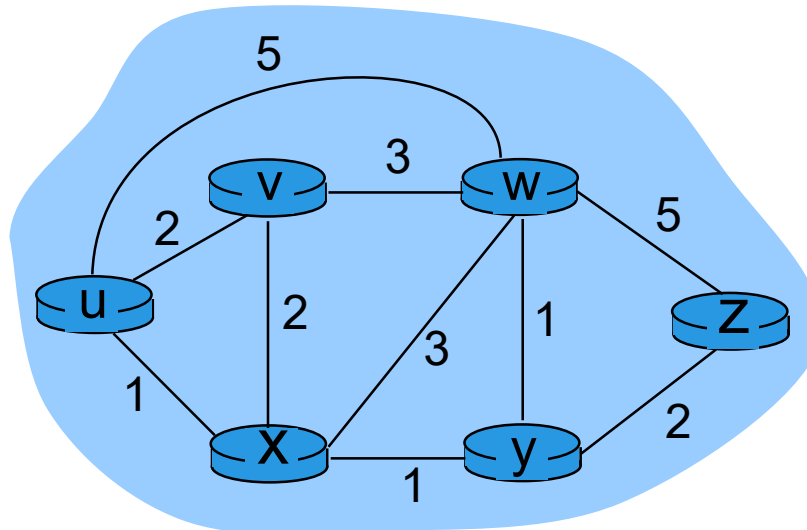
$p(v)$: predecessor node along path from source to v (neighbor of v)

N' : set of nodes whose least cost path already known

Dijkstra's Algorithm: Example

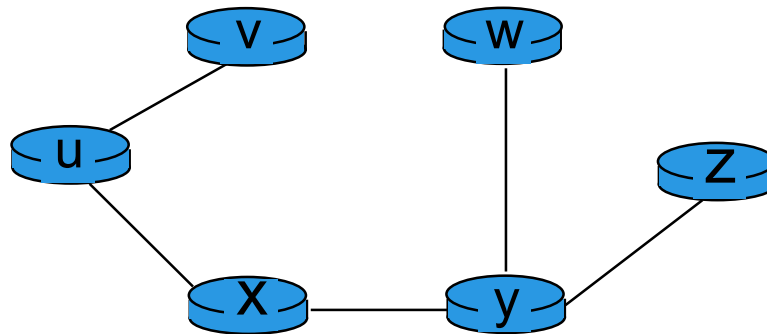
Computer the short-path from u to all network nodes using Dijkstra's Algorithm

Step	N'	D(v),p(v)	D(w),p(w)	D(x),p(x)	D(y),p(y)	D(z),p(z)
0	u	2,u	5,u	1,u	∞	∞
1	ux	2,u	4,x		2,x	∞
2	uxy	2,u	3,y			4,y
3	uxyv		3,y			4,y
4	uxyvw					4,y
5	uxyvwz					



Dijkstra's Algorithm: Example

resulting shortest-path tree from u:



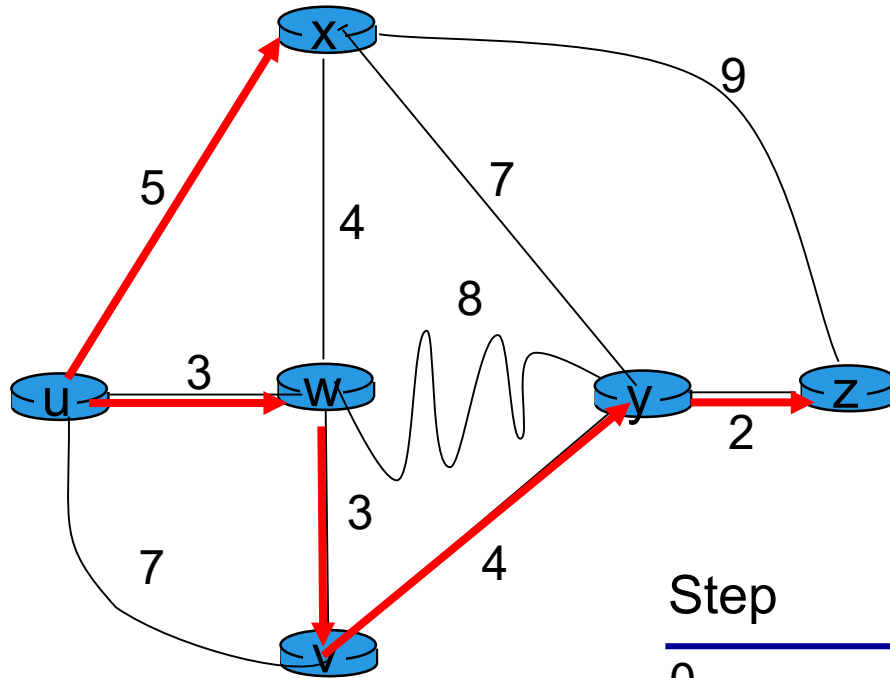
resulting forwarding table at u:

destination	link
v	(u,v)
x	(u,x)
y	(u,x)
w	(u,x)
z	(u,x)

Dijkstra's Algorithm, Discussion

- **Algorithm complexity:** n nodes
- each iteration: need to check all nodes, w , not in N
- $n(n+1)/2$ comparisons:
- more efficient implementations possible, Distance-Vector Algorithm

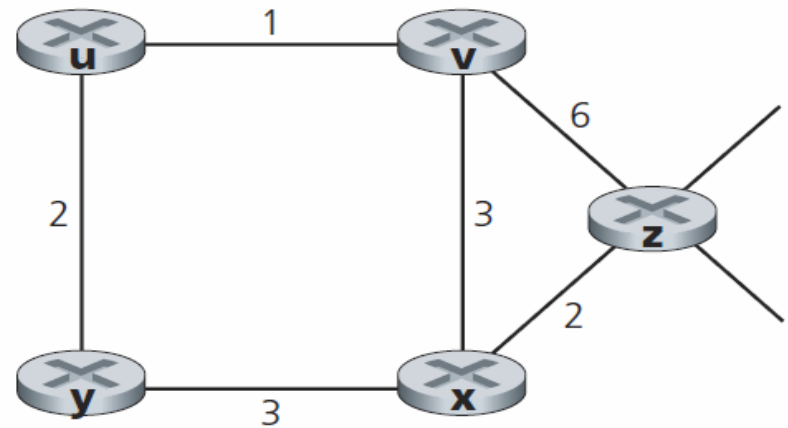
Dijkstra's Algorithm: another example



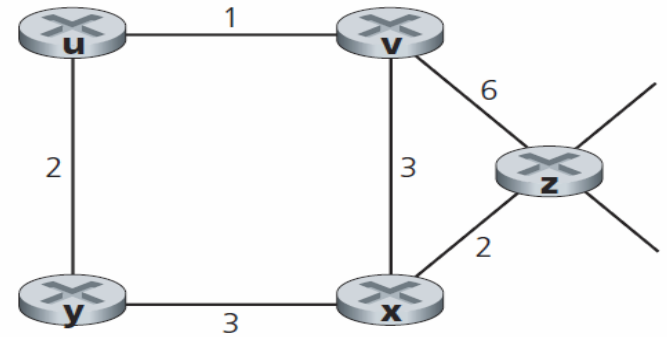
Step	N'	D(v) p(v)	D(w) p(w)	D(x) p(x)	D(y) p(y)	D(z) p(z)
0	u	7,u	3,u	5,u	∞	∞
1	uw	6,w		5,u	11,w	∞
2	uwx	6,w			11,w	14,x
3	uwxv				10,v	14,x
4	uwxvy					12,y
5	uwxvyz					

Dijkstra's Algorithm, Practice Problem 1

Compute the shortest path from **z** to all network nodes, using Dijkstra Algorithm.

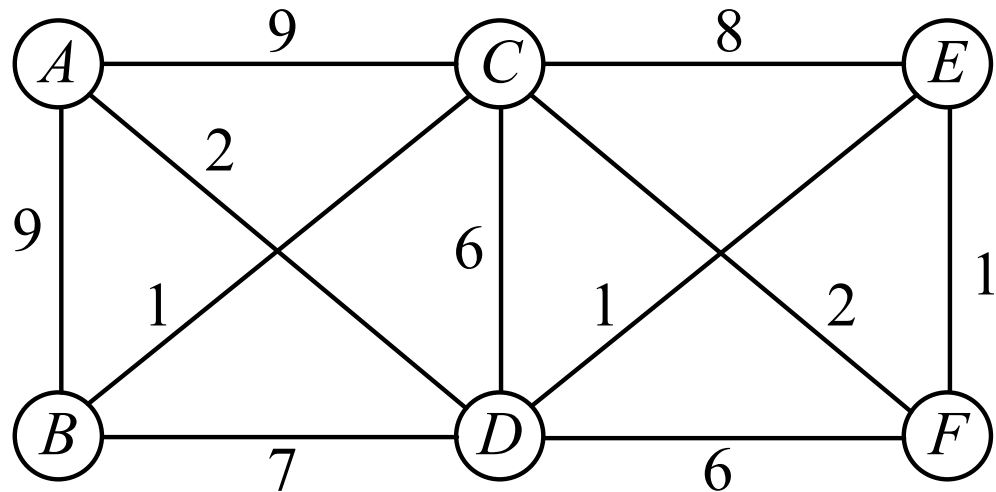


Dijkstra's Algorithm, Practice Problem 1



Dijkstra's Algorithm, Practice Problem 2

Consider the following network. With the indicated link costs, compute the shortest path from A to all network nodes using Dijkstra Algorithm.



Dijkstra's Algorithm, Practice Problem 2

