Framing Signals a return to portable shellcode



memory corruption,

the problem that just won't go away

25+ years after the morris worm and still going strong

return addr



return addr

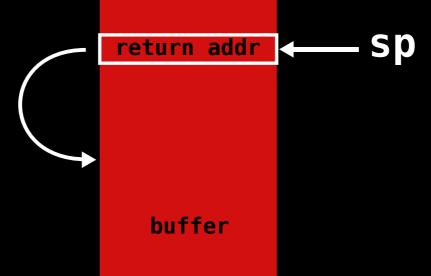


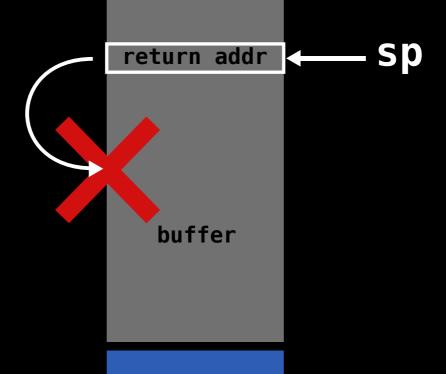
return addr

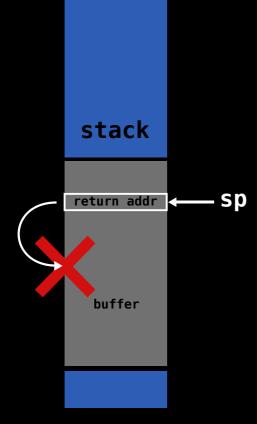


return addr

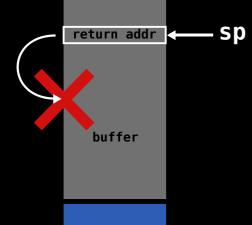
—— sp







stack



code



return addr

buffer

code

return oriented programming / ret2libc stack return addr buffer gadget code

return oriented programming / ret2libc return addr return addr return addr buffer gadget gadget code gadget

Return Oriented Programming

- dependent on available gadgets
- non-trivial to program
- chains may differ greatly between different binaries
- Turing complete

Sigreturn Oriented Programming

- minimal number of gadgets
- constructing shellcode by chaining system calls
- easy to change functionality of shellcode
- shellcode portable (gadgets are always present)
- Turing complete

unix signals stack —— sp



ucontext



ucontext

siginfo

ucontext

siginfo

sigreturn ← S

good:

kernel agnostic about signal handlers

ucontext

siginfo

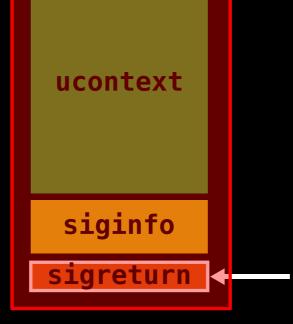
sigreturn

sp

bad:

kernel agnostic about signal handlers

(we can fake 'em)



two gadgets

- call to sigreturn
- syscall & return

forged signal frame

sigreturn

program counter

forged signal frame

sigreturn

program counter stack pointer

forged signal frame

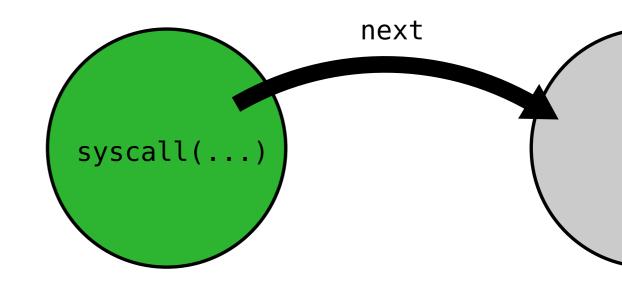
sigreturn

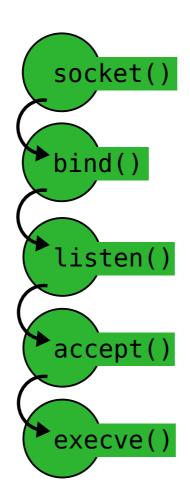
program counter stack pointer RAX **RDI** RSI RDX **R10** R8 R9 sigreturn

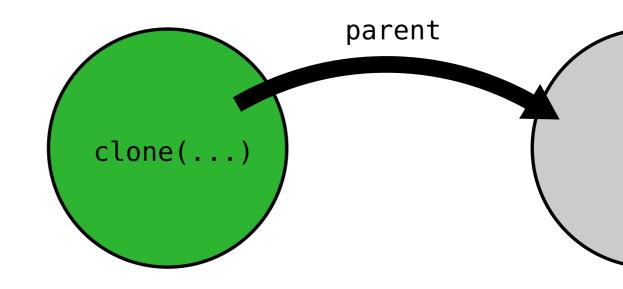
program counter stack pointer syscall number arg1 arg2 arg3 arg4 arg5 arg6 sigreturn

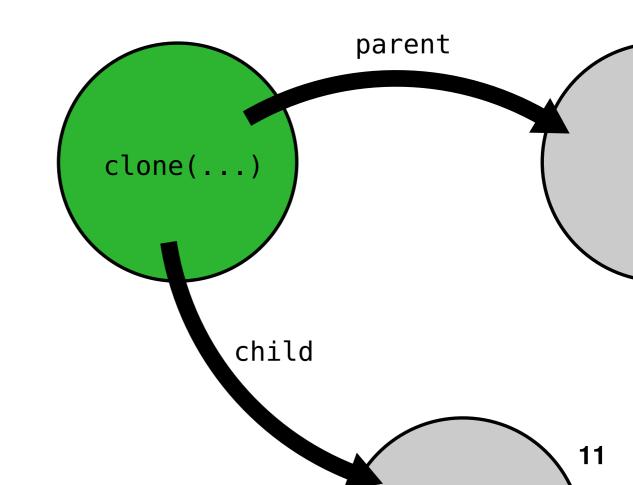
syscall & return stack pointer syscall number arg1 arg2 arg3 arg4 arg5 arg6 sigreturn

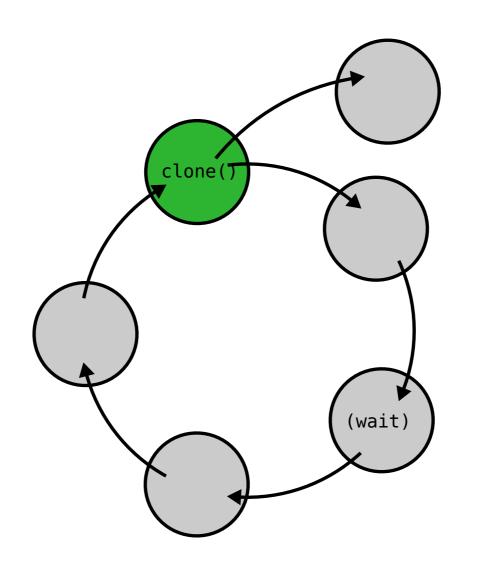
syscall & return next sigframe syscall number arg1 arg2 arg3 arg4 arg5 arg6 sigreturn











usage scenarios

- stealthy backdoor
- code signing circumvention
- generic shellcode for exploitation

usage scenarios

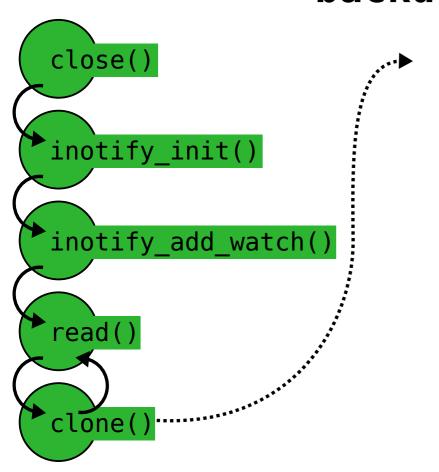
- stealthy backdoor
- code signing circumvention
- generic shellcode for exploitation

stealthy backdoor

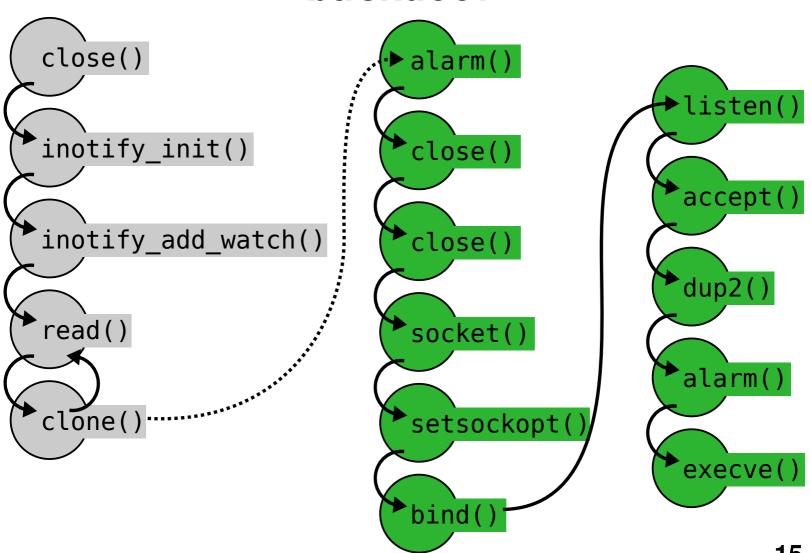
basic idea:

- use the inotify API to wait for a file to be read
- when this file is read: open a listen socket to spawn a shell
- terminate the listening socket quickly if nobody connects

backdoor



backdoor



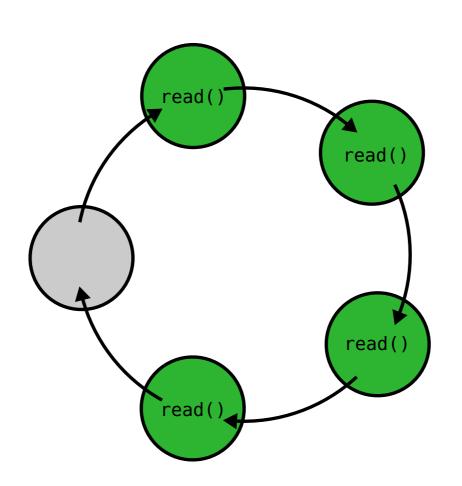
usage scenarios

- stealthy backdoor
- code signing circumvention
- generic shellcode for exploitation

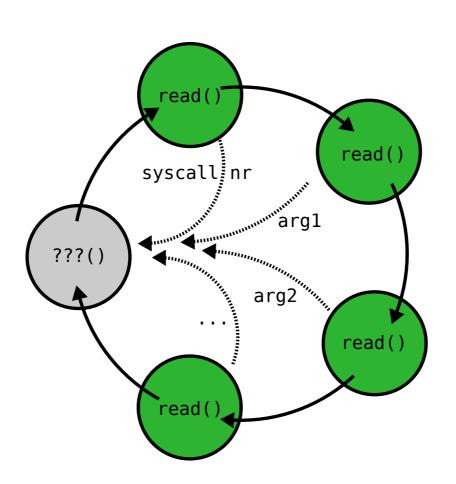
code signing circumvention

- serialize system calls over a socket
- write into our own signal frames
 useful to bypass code-signing restrictions

system call proxy



system call proxy



and... It's turing complete

usage scenarios

- stealthy backdoor
- code signing circumvention
- generic shellcode for exploitation

SROP exploit on x86-64

we have:

- a stack buffer overflow
- not a single gadget from the binary

assumption:

- we can guess/leak the location
 of a writable address (any address!)
- we have some control over RAX (function's return value)

two gadgets

- call to sigreturn
- syscall & return

two gadgets

- call to sigreturn: RAX = 15 + syscall
- syscall & return

one gadget

- RAX = 15
- syscall & return

kernel memory

kernel memory

user memory

kernel memory

user memory

linux memory layout kernel memory vsyscall user memory

[vsyscall]

[vsyscall]

[vsyscall]

0f05 syscall c3 return

syscall(arg1, arg2, arg3, ...) = result

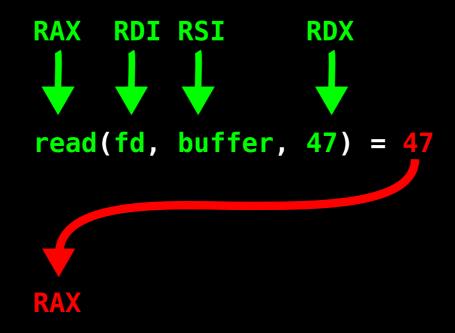
```
execve("/bin/sh",
["/bin/sh", "--", "...", NULL],
NU'Z)
```

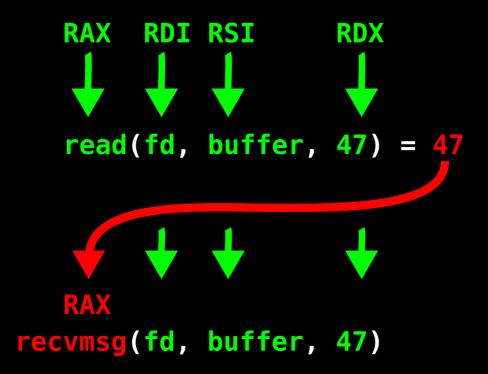
syscall(arg1, arg2, arg3, ...) = result

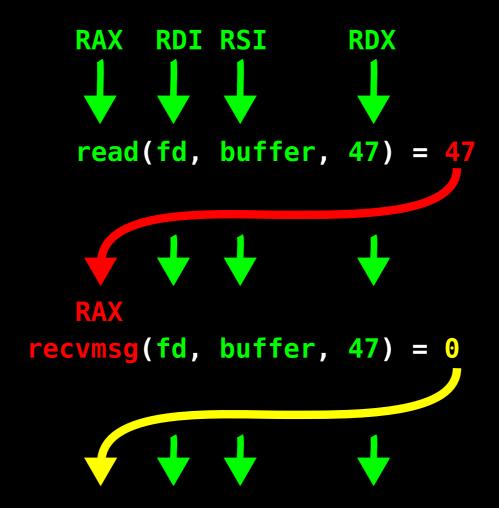
read(fd, buffer, 47)











syscall(arg1, arg2, arg3, ...) = result

read(fd, addr, ...) = result

read(fd, stack_addr, ...) = result

 $read(fd, stack_addr, 306) = 306$

```
read(fd, stack_addr, 306) = 306

RAX == 306 == __NR_syncfs
top of stack points to syscall & return
```

```
read(fd, stack_addr, 306) = 306

RAX == 306 == __NR_syncfs
  top of stack points to syscall & return
syncfs(fd) = ...
```

```
read(fd, stack_addr, 306) = 306

RAX == 306 == __NR_syncfs
  top of stack points to syscall & return
syncfs(fd) = 0
```

```
read(fd, stack_addr, 306) = 306

RAX == 306 == __NR_syncfs
  top of stack points to syscall & return
syncfs(fd) = 0

RAX == 0 == __NR_read
  top of stack points to syscall & return
```

```
read(fd, stack_addr, 306) = 306

RAX == 306 == __NR_syncfs
  top of stack points to syscall & return
syncfs(fd) = 0

RAX == 0 == __NR_read
  top of stack points to syscall & return
read(fd, stack addr, 306) = ...
```

```
read(fd, stack_addr, 306) = 306

RAX == 306 == __NR_syncfs
  top of stack points to syscall & return
syncfs(fd) = 0

RAX == 0 == __NR_read
  top of stack points to syscall & return
read(fd, stack addr, 306) = 15
```

```
read(fd, stack addr, 306) = 306
    RAX == 306 == NR syncfs
    top of stack points to syscall & return
syncfs(fd) = 0
    RAX == 0 == NR read
    top of stack points to syscall & return
read(fd, stack addr, 306) = 15
    RAX == 15 == NR rt sigreturn
    top of stack points to syscall & return
```

```
read(fd, stack addr, 306) = 306
    RAX == 306 == NR syncfs
    top of stack points to syscall & return
syncfs(fd) = 0
    RAX == 0 == NR read
    top of stack points to syscall & return
read(fd, stack addr, 306) = 15
    RAX == 15 == NR rt sigreturn
    top of stack points to syscall & return
mprotect(stack addr, 0x1000,
         PROT READ | PROT WRITE | PROT EXEC)
```

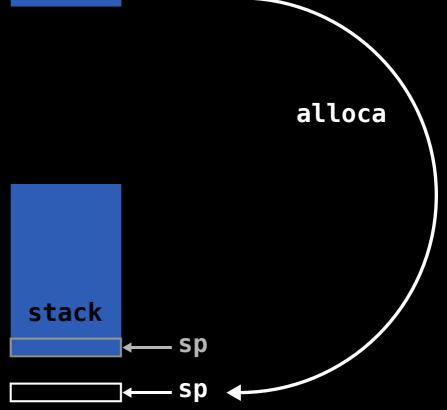
```
read(fd, stack addr, 306) = 306
    RAX == 306 == NR syncfs
    top of stack points to syscall & return
syncfs(fd) = 0
    RAX == 0 == NR read
    top of stack points to syscall & return
read(fd, stack addr, 306) = 15
    RAX == 15 == NR rt sigreturn
    top of stack points to syscall & return
mprotect(stack addr, 0x1000,
         PROT READ | PROT WRITE | PROT EXEC)
    top of stack points to our code
```



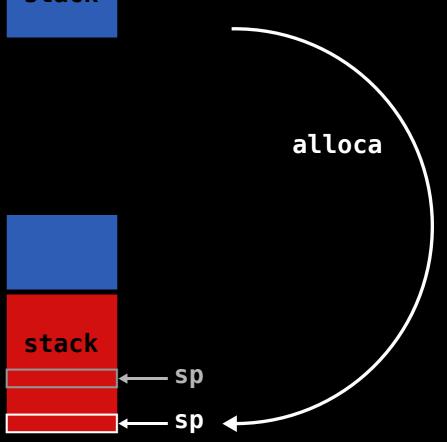
stack ←— sp

← sp

6 (asterisk)
stack



stack



On some systems SROP gadgets are randomised, on others, they are not

Operating system	Gadget	Memory map
Linux i386	sigreturn	[vdso]
Linux < 3.11 ARM	sigreturn	[vectors] 0xffff0000
Linux < 3.3 x86-64	syscall & return	[vsyscall] 0xffffffff600000
Linux ≥ 3.3 x86-64	syscall & return	Libc
Linux x86-64	sigreturn	Libc
FreeBSD 9.2 x86-64	sigreturn	0x7ffffff000
Mac OSX x86-64	sigreturn	Libc
iOS ARM	sigreturn	Libsystem
iOS ARM	syscall & return	Libsystem

On some systems SROP gadgets are randomised, on others, they are not android

n	on - A	SLR :-(alidi otd	
		Operating system	Gadget	Memory map
		Linux i386	sigreturn	[vdso]
	\rightarrow	Linux < 3.11 ARM	sigreturn	[vectors] 0xffff0000
		Linux < 3.3 x86-64	syscall & return	[vsyscall] 0xffffffff600000
		Linux ≥ 3.3 x86-64	syscall & return	Libc
		Linux x86-64	sigreturn	Libc
		FreeBSD 9.2 x86-64	sigreturn	0x7fffffff000
		Mac OSX x86-64	sigreturn	Libc
		iOS ARM	sigreturn	Libsystem
		iOS ARM	syscall & return	Libsystem

mitigation:

```
It may be useful to disable
vsyscall

vsyscall=emulate
(default from Linux 3.3 onward)
or

vsyscall=none
```

mitigation:

- Counting signals in progress

mitigation:

- Counting signals in progress

- Signal frame canaries

stack canary stack

return addr



buffer



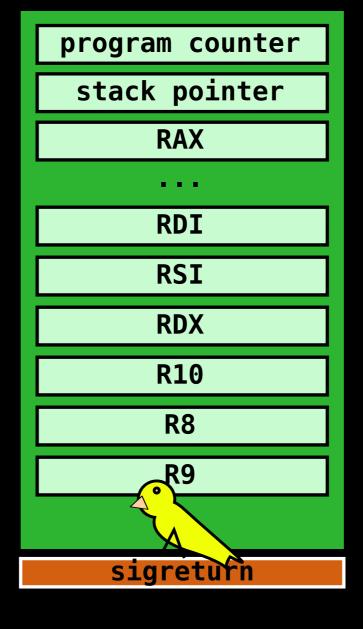
stack canary stack

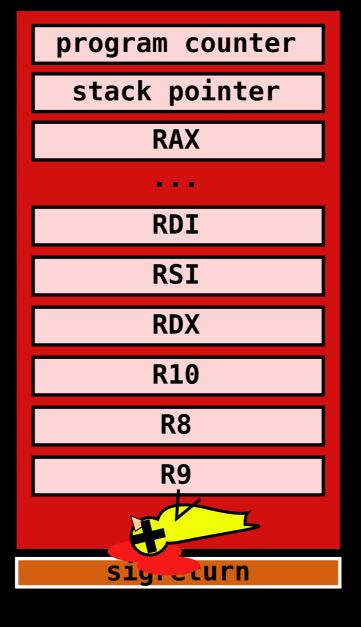
return addr



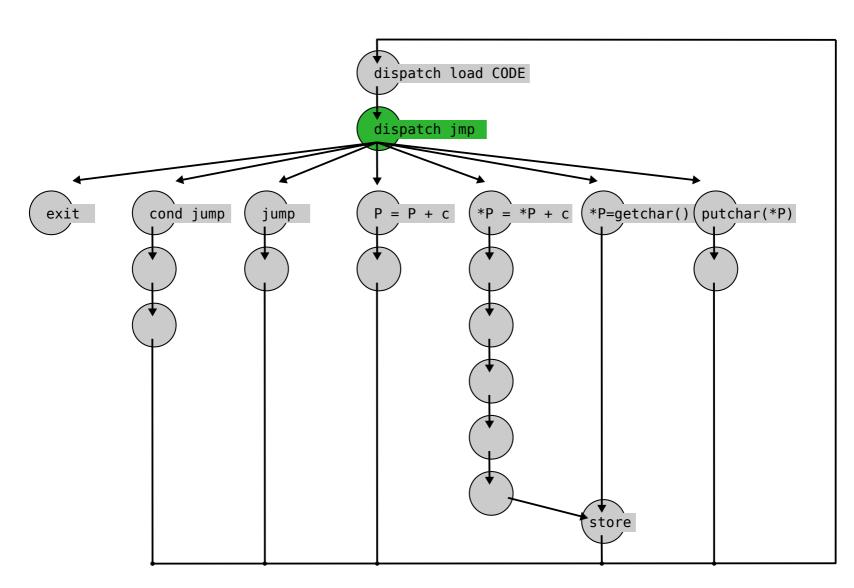
buffer







questions?



```
code = open("/proc/self/mem",0_RDWR);
p = open("/proc/self/mem",0_RDWR);
a = open("/proc/self/mem",0_RDWR);
```

```
code = open("/proc/self/mem", 0_RDWR);
p = open("/proc/self/mem", 0_RDWR);
a = open("/proc/self/mem", 0_RDWR);
instruction dispatch:
```

read(code, &ucontext.sp, sizeof(long));

```
code = open("/proc/self/mem", 0_RDWR);
p = open("/proc/self/mem", 0_RDWR);
a = open("/proc/self/mem", 0_RDWR);
instruction dispatch:
   read(code, &ucontext.sp, sizeof(long));

pointer ops:
   p++ -> lseek(p, 1, SEEK_CUR);
```

```
code = open("/proc/self/mem",0 RDWR);
p = open("/proc/self/mem", 0 RDWR);
a = open("/proc/self/mem", O RDWR);
instruction dispatch:
  read(code, &ucontext.sp, sizeof(long));
pointer ops:
  p++ -> lseek(p, 1, SEEK CUR);
addition:
  lseek(a, &identity table x2, SEEK SET);
  lseek(a, val1, SEEK SET);
  lseek(a, val2, SEEK SET);
  read(a, dest, 1);
```