MODULE *IscpBatchTimestamp*

asd

$$\forall i : a_i = b$$

Can we determine if proposal is invalid with regards to the timestamp? That is done after the ACS.

EXTENDS *Naturals*, *FiniteSets*, *TLAPS*
CONSTANT *Nodes* A set of node identifiers.
CONSTANT *Byzantine* A set of byzantine node identifiers.
CONSTANT *Time* A set of timestamps, represented as natural numbers to have $\leq$.
ASSUME $ConstantAssms \triangleq Byzantine \subseteq Nodes \land Time \subseteq Nat$
$Requests \triangleq Time$ Assume requests are identified by timestamps of their *TX* only.

VARIABLE *proposed* Was the proposal made?
VARIABLE *npRq* Node proposal: A set of requests.
VARIABLE *npTS* Node proposal: Timestamp.
$vars \triangleq \langle proposed, npRq, npTS \rangle$

$F \triangleq Cardinality(Byzantine)$
$N \triangleq Cardinality(Nodes)$
ASSUME $ByzantineAssm \triangleq N \geq 3 * F + 1$

$F1Quorums \triangleq \{q \in \text{SUBSET } Nodes : Cardinality(q) = F + 1\}$
$NFQuorums \triangleq \{q \in \text{SUBSET } Nodes : Cardinality(q) = N - F\}$

$BatchRq(rq) \triangleq \exists\, q \in F1Quorums : \forall\, n \in q : rq \in npRq[n]$
$BatchRqs \triangleq \{rq \in Requests : BatchRq(rq)\}$

$SubsetTS(s) \triangleq \{npTS[n] : n \in s\}$
$BatchTS(ts) \triangleq \exists\, q \in NFQuorums :$
$\quad \land\, ts \in SubsetTS(q)$
$\quad \land\, \forall\, x \in SubsetTS(q) : ts \geq x$
$\quad \land\, \forall\, x \in SubsetTS(Nodes \setminus q) : ts \leq x$

$ProposalValid(n) \triangleq \forall\, rq \in npRq[n] : rq \leq npTS[n]$
$Propose \triangleq \neg proposed \land proposed' = \text{TRUE}$
$\quad \land\, npRq' \in [Nodes \to (\text{SUBSET } Requests) \setminus \{\{\}\}]$ Some node non-empty proposals.
$\quad \land\, npTS' \in [Nodes \to Time]$ Some timestamps.
$\quad \land\, \forall\, n \in (Nodes \setminus Byzantine) : ProposalValid(n)'$ Fair node proposals are valid.

$Init \triangleq proposed = \text{FALSE} \land npRq = \{\} \land npTS = \{\}$ Dummy values, on init.
$Spec \triangleq Init \land \Box[Propose]_{vars}$ For model checking in *TLC*.
$Invariant \triangleq proposed \Rightarrow \forall\, ts \in Time, rq \in BatchRqs : BatchTS(ts) \Rightarrow rq \leq ts$

THEOREM $Spec \Rightarrow \Box Invariant$
PROOF OMITTED Checked with *TLC*.
THEOREM $Byzantine = \{\} \land Spec \Rightarrow \Box Invariant$

55 $\langle 1 \rangle 1.\ \mathit{Byzantine} = \{\} \wedge \mathit{Init} \Rightarrow \mathit{Invariant}$

56 $\langle 1 \rangle 2.\ \mathit{Invariant} \wedge \mathit{Byzantine} = \{\} \wedge \mathit{Propose} \Rightarrow \mathit{Invariant}'$

57 $\langle 1 \rangle \text{q}.$ QED BY $\langle 1 \rangle 1, \langle 1 \rangle 2, \mathit{PTL}, \mathit{ConstantAssms}$ DEF Spec

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Counter-example with $\mathit{Nodes} = 101 \,..\, 104$, $\mathit{Byzantine} = \{104\}$, $\mathit{Time} = 1 \,..\, 3$:
$\mathit{PropposedRq}$: (101 :> {1} @@ 102 :> {1} @@ 103 :> {2} @@ 104 :> {2}),
$\mathit{PropposedTS}$: (101 :> 1 @@ 102 :> 1 @@ 103 :> 2 @@ 104 :> 1),
$\mathit{BatchRq}$: {1, 2},
$\mathit{BatchTS}$: 1