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# Operational notes

Document updated on **June 8, 2022**.

The following colors are **not** part of the final product, but serve as highlights in the editing/review process:

















































- text that needs attention from the Subject Matter Experts: Mirco, Anna,& Jan
- terms that have not yet been defined in the book
- things that need to be checked only at the very final typesetting stage (and it doesn't make sense to do them before)
- text that needs advice from the communications/marketing team: Aaron & Shane
- text that needs to be completed or otherwise edited (by Sylvia)

# 11 Todo list

12	 zero-knowledge proofs . . . . .	12
13	 finite field . . . . .	12
14	 elliptic curve . . . . .	12
15	 add references . . . . .	12
16	 methatical . . . . .	12
17	 We haven't really talked about what a ring is at this point . . . . .	14
18	 What's the significance of this distinction? . . . . .	15
19	 reverse . . . . .	15
20	 Turing machine . . . . .	15
21	 polynomial time . . . . .	15
22	 sub-exponentially, with $\mathcal{O}((1 + \varepsilon)^n)$ and some $\varepsilon > 0$ . . . . .	15
23	 Add text . . . . .	16
24	 $\mathbb{Q}$ of fractions . . . . .	16
25	 pseudocode . . . . .	18
26	 check reference . . . . .	18
27	 modular arithmetics . . . . .	18
28	 actual division . . . . .	18
29	 multiplicative inverses . . . . .	18
30	 factional numbers . . . . .	18
31	 exponentiation function . . . . .	20
32	 See XXX . . . . .	20
33	 This Sage snippet should be described in more detail. . . . .	21
34	 prime fields . . . . .	23
35	 residue class rings . . . . .	23
36	 check algorithm floating . . . . .	23
37	 $(-1)$ should be $(-a)$ ? . . . . .	26
38	 subtrahend . . . . .	33
39	 minuend . . . . .	33
40	 what does this mean? . . . . .	37
41	 Def Subgroup, Fundamental theorem of cyclic groups. . . . .	40
42	 add reference . . . . .	41
43	 Add real-life example of 0? . . . . .	41
44	 add reference . . . . .	41
45	 check reference . . . . .	42
46	 check references to previous examples . . . . .	43
47	 RSA crypto system . . . . .	43
48	 size 2048-bits . . . . .	43
49	 check reference . . . . .	43
50	 add reference: 27? . . . . .	43

51	check reference	44
52	polynomial time	44
53	exponential time	44
54	TODO: Fundamental theorem of finite cyclic groups	44
55	check reference	44
56	runtime complexity	45
57	add reference	45
58	S: what does “efficiently” mean here?	45
59	computational hardness assumptions	45
60	check reference	45
61	check reference	46
62	explain last sentence more	46
63	“equation”?	47
64	check reference	47
65	what’s the difference between $\mathbb{F}_p^*$ and $\mathbb{Z}_p^*$ ?	47
66	Legendre symbol	47
67	Euler’s formular	47
68	These are only explained later in the text, ‘4.27’	47
69	are these going to be relevant later? yes, they are used in various snark proof systems	48
70	TODO: theorem: every factor of order defines a subgroup...	48
71	Is there a term for this property?	49
72	a few examples?	51
73	check reference	51
74	TODO: DOUBLE CHECK THIS REASONING.	51
75	Mirco: We can do better than this	53
76	check reference	54
77	add reference	55
78	pseudorandom	55
79	oracle	55
80	check reference	55
81	add text on this	55
82	check reference	57
83	check reference	57
84	check reference	57
85	check reference	58
86	add more examples protocols of SNARK	58
87	check reference	58
88	add reference	58
89	Abelian groups	58
90	codomain	58
91	Check change of wording	59
92	add reference	60
93	Expand on this?	60
94	check reference	60
95	S: are we introducing elliptic curves in section 1 or 2?	61
96	check reference	62
97	check reference	62
98	add reference	62

99	check reference . . . . .	62
100	write paragraph on exponentiation . . . . .	63
101	add reference . . . . .	63
102	check reference . . . . .	63
103	add reference . . . . .	63
104	group pairings . . . . .	63
105	add reference . . . . .	64
106	check reference . . . . .	64
107	check reference . . . . .	67
108	add reference . . . . .	68
109	TODO: Elliptic Curve asymmetric cryptography examples. Private key, generator,	
110	public key. . . . .	70
111	add reference . . . . .	70
112	maybe remove this sentence? . . . . .	70
113	affine space . . . . .	70
114	cusps . . . . .	71
115	self-intersections . . . . .	71
116	check reference . . . . .	72
117	check reference . . . . .	73
118	jubjub . . . . .	73
119	check reference . . . . .	73
120	affine plane . . . . .	73
121	check reference . . . . .	74
122	check reference . . . . .	74
123	check reference . . . . .	75
124	sign . . . . .	75
125	more explanation of what the sign is . . . . .	75
126	check reference . . . . .	75
127	S: I don't follow this at all . . . . .	76
128	check reference . . . . .	76
129	add explanation of how this shows what we claim . . . . .	76
130	should this def. be moved even earlier? . . . . .	77
131	chord line . . . . .	77
132	tangential . . . . .	77
133	tangent line . . . . .	77
134	remove $Q$ ? . . . . .	77
135	where? . . . . .	78
136	check reference . . . . .	78
137	check reference . . . . .	78
138	check reference . . . . .	78
139	check reference . . . . .	79
140	check reference . . . . .	79
141	check reference . . . . .	80
142	check reference . . . . .	80
143	check reference . . . . .	81
144	add term . . . . .	81
145	add term . . . . .	81
146	add reference . . . . .	81

147	 cofactor clearing . . . . .	81
148	 add reference . . . . .	81
149	 check reference . . . . .	81
150	 check reference . . . . .	82
151	 add reference . . . . .	82
152	 add reference . . . . .	82
153	 check reference . . . . .	82
154	 check reference . . . . .	82
155	 check reference . . . . .	83
156	 check reference . . . . .	83
157	 check reference . . . . .	83
158	 Explain how . . . . .	83
159	 write example . . . . .	84
160	 check reference . . . . .	84
161	 add reference . . . . .	84
162	 check reference . . . . .	84
163	 add reference . . . . .	85
164	 check reference . . . . .	85
165	 add reference . . . . .	85
166	 check reference . . . . .	85
167	 add reference . . . . .	85
168	 check reference . . . . .	85
169	 add reference . . . . .	85
170	 add reference . . . . .	85
171	 add reference . . . . .	85
172	 check reference . . . . .	85
173	 check reference . . . . .	85
174	 Check if following Alg is floated too far . . . . .	86
175	 add reference . . . . .	86
176	 add reference . . . . .	86
177	 write up this part . . . . .	86
178	 is the label in L <sup>A</sup> T <sub>E</sub> X correct here? . . . . .	88
179	 check reference . . . . .	88
180	 check reference . . . . .	88
181	 check reference . . . . .	88
182	 check reference . . . . .	89
183	 check reference . . . . .	89
184	 check reference . . . . .	90
185	 check reference . . . . .	90
186	 check reference . . . . .	90
187	 check reference . . . . .	90
188	 check reference . . . . .	90
189	 add reference . . . . .	90
190	 check reference . . . . .	92
191	 check reference . . . . .	92
192	 check reference . . . . .	92
193	 check reference . . . . .	92
194	 check reference . . . . .	92

















































195	■ check reference . . . . .	94
196	■ check reference . . . . .	94
197	■ check reference . . . . .	95
198	■ either expand on this or delete it . . . . .	95
199	■ add reference . . . . .	95
200	■ check reference . . . . .	95
201	■ check reference . . . . .	95
202	■ check reference . . . . .	95
203	■ check reference . . . . .	95
204	■ check reference . . . . .	95
205	■ check reference . . . . .	96
206	■ check reference . . . . .	96
207	■ check reference . . . . .	97
208	■ check reference . . . . .	97
209	■ add reference . . . . .	97
210	■ add reference . . . . .	97
211	■ This needs to be written (in Algebra) . . . . .	97
212	■ add reference . . . . .	97
213	■ add reference . . . . .	97
214	■ check reference . . . . .	97
215	■ towers of curve extensions . . . . .	98
216	■ check reference . . . . .	98
217	■ check reference . . . . .	98
218	■ check reference . . . . .	98
219	■ check reference . . . . .	99
220	■ add reference . . . . .	99
221	■ check reference . . . . .	100
222	■ S: either add more explanation or move to a footnote . . . . .	100
223	■ type 3 pairing-based cryptography . . . . .	100
224	■ add references? . . . . .	100
225	■ check reference . . . . .	101
226	■ check reference . . . . .	101
227	■ check floating of algorithm . . . . .	102
228	■ add references . . . . .	102
229	■ check reference . . . . .	103
230	■ add reference . . . . .	103
231	■ check reference . . . . .	103
232	■ check reference . . . . .	103
233	■ add reference . . . . .	104
234	■ should all lines of all algorithms be numbered? . . . . .	104
235	■ check reference . . . . .	105
236	■ check reference . . . . .	105
237	■ check reference . . . . .	105
238	■ check if the algorithm is floated properly . . . . .	105
239	■ check reference . . . . .	105
240	■ again? . . . . .	107
241	■ check reference . . . . .	107
242	■ circuit . . . . .	107

243	signature schemes	107
244	add reference	107
245	check reference	108
246	check reference	108
247	add references	108
248	add reference	108
249	reference text to be written in Algebra	108
250	check reference	108
251	check reference	108
252	check reference	109
253	add reference	109
254	algebraic closures	109
255	check reference	109
256	check reference	110
257	check reference	110
258	check reference	110
259	check reference	111
260	disambiguate	111
261	add reference	111
262	unify terminology	111
263	check reference	112
264	actually make this a table?	112
265	exercise still to be written?	113
266	add reference	113
267	check reference	113
268	check reference	113
269	add reference	114
270	check reference	115
271	check reference	115
272	check reference	115
273	add reference	116
274	check reference	116
275	check reference	116
276	check reference	117
277	what does this mean? Maybe just delete it	118
278	write up this part	119
279	add reference	119
280	check reference	119
281	cyclotomic polynomial	119
282	Pholaard-rho attack	119
283	todo	119
284	why? Because in this book elliptic curves are only defined for fields of chracteristic $> 3$	120
285	check reference	120
286	check reference	120
287	what does this mean?	120
288	add reference	120
289	add reference	120
290	check reference	120

















































291	check reference	121
292	add reference	122
293	add exercise	122
294	check reference	123
295	add reference	123
296	add reference	123
297	add reference	123
298	check reference	124
299	check reference	124
300	add reference	124
301	add reference	124
302	add reference	125
303	check reference	125
304	add reference	125
305	add reference	125
306	finish writing this up	126
307	add reference	126
308	correct computations	126
309	fill in missing parts	126
310	add reference	127
311	check equation	127
312	Chapter 1?	128
313	"rigorous"?	128
314	"proving"?	128
315	Add example	129
316	M: 1:1 correspondence might actually be wrong	129
317	I'd delete this, too distracting	129
318	binary tuples	129
319	add reference	130
320	add reference	130
321	check reference	130
322	check reference	130
323	Are we using $w$ and $x$ interchangeably or is there a difference between them?	131
324	check reference	131
325	jubjub	131
326	check reference	131
327	check reference	131
328	check wording	131
329	check reference	131
330	check references	132
331	add reference	132
332	add reference	132
333	check reference	133
334	add reference	133
335	check reference	134
336	check reference	134
337	add reference	135
338	add reference	136



---

339	 Schur/Hadamard product . . . . .	136
340	 add reference . . . . .	136
341	 check reference . . . . .	136
342	 check reference . . . . .	137
343	 add reference . . . . .	138
344	 check reference . . . . .	139
345	 check reference . . . . .	139
346	 check reference . . . . .	139
347	 check reference . . . . .	139
348	 check reference . . . . .	140
349	 add reference . . . . .	140
350	 add reference . . . . .	141
351	 check reference . . . . .	141
352	 check reference . . . . .	142
353	 check reference . . . . .	142
354	 check reference . . . . .	142
355	 add reference . . . . .	143
356	 check reference . . . . .	145
357	 add reference . . . . .	145
358	 check reference . . . . .	146
359	 check reference . . . . .	146
360	 check reference . . . . .	146
361	 Should we refer to R1CS satisfiability (p. 139 here? . . . . .	147
362	 check reference . . . . .	148
363	 add reference . . . . .	148
364	 check reference . . . . .	148
365	 check reference . . . . .	149
366	 check reference . . . . .	149
367	 check reference . . . . .	150
368	 check reference . . . . .	152
369	 add reference . . . . .	153
370	 "by"? . . . . .	153
371	 check reference . . . . .	153
372	 check reference . . . . .	153
373	 add reference . . . . .	153
374	 add reference . . . . .	153
375	 check reference . . . . .	153
376	 add reference . . . . .	153
377	 clarify language . . . . .	155
378	 check reference . . . . .	156
379	 add reference . . . . .	156
380	 check reference . . . . .	156
381	 add reference . . . . .	156
382	 check references . . . . .	159
383	 add references to these languages? . . . . .	159
384	 check reference . . . . .	162
385	 check reference . . . . .	163
386	 check reference . . . . .	163












---

387		check reference . . . . .	164
388		check reference . . . . .	165
389		check reference . . . . .	165
390		check reference . . . . .	167
391		check reference . . . . .	167
392		check reference . . . . .	168
393		add reference . . . . .	168
394		check reference . . . . .	168
395		add reference . . . . .	168
396		add reference . . . . .	168
397		check reference . . . . .	169
398		check reference . . . . .	169
399		check reference . . . . .	169
400		check reference . . . . .	169
401		add reference . . . . .	169
402		check reference . . . . .	170
403		check reference . . . . .	171
404		"constraints" or "constrained"? . . . . .	171
405		check reference . . . . .	172
406		"constraints" or "constrained"? . . . . .	172
407		add reference . . . . .	172
408		"constraints" or "constrained"? . . . . .	172
409		add reference . . . . .	173
410		check references . . . . .	173
411		check reference . . . . .	173
412		add reference . . . . .	174
413		can we rotate this by 90°? . . . . .	174
414		check reference . . . . .	175
415		add reference . . . . .	175
416		add reference . . . . .	175
417		shift . . . . .	177
418		bishift . . . . .	178
419		add reference . . . . .	179
420		check reference . . . . .	180
421		Add example . . . . .	181
422		add reference . . . . .	182
423		add reference . . . . .	183
424		check reference . . . . .	184
425		add reference . . . . .	184
426		add reference . . . . .	184
427		check reference . . . . .	185
428		add reference . . . . .	185
429		add reference . . . . .	185
430		add reference . . . . .	187
431		check reference . . . . .	188
432		check reference . . . . .	189
433		common reference string . . . . .	189
434		simulation trapdoor . . . . .	189

---

435	check reference	189
436	check reference	189
437	add reference	190
438	check reference	190
439	check reference	190
440	check reference	190
441	"invariable"?	190
442	explain why	191
443	4 examples have the same title. Change it to be distinct	191
444	check reference	191
445	add reference	191
446	check reference	191
447	add reference	191
448	add reference	192
449	add reference	193
450	check reference	194
451	add reference	194
452	add reference	195
453	check reference	195
454	check reference	195
455	add reference	195
456	add reference	195
457	check reference	196
458	add reference	196
459	add reference	196
460	add reference	196
461	check reference	196
462	add reference	196
463	add reference	196
464	add reference	196
465	add reference	196
466	add reference	197
467	add reference	197
468	add reference	197
469	add reference	197
470	check reference	199
471	check reference	199
472	add reference	199
473	add reference	199
474	add reference	199
475	add reference	199
476	add reference	200
477	add reference	200
478	add reference	200
479	add reference	200
480	fix error	200
481	add reference	200
482	check reference	201

---

483	 add reference . . . . .	201
484	 add reference . . . . .	201
485	 add reference . . . . .	201
486	 add reference . . . . .	202
487	 add reference . . . . .	202
488	 add reference . . . . .	202
489	 add reference . . . . .	202
490	 add reference . . . . .	202
491	 add reference . . . . .	202
492	 add reference . . . . .	202
493	 add reference . . . . .	203

494

# MoonMath manual

495

TechnoBob and the Least Scruples crew

496

June 8, 2022

# Contents

498	<b>1 Introduction</b>	<b>5</b>
499	1.1 Target audience . . . . .	5
500	1.2 The Zoo of Zero-Knowledge Proofs . . . . .	6
501	To Do List . . . . .	8
502	Points to cover while writing . . . . .	8
503	<b>2 Preliminaries</b>	<b>9</b>
504	2.1 Preface and Acknowledgements . . . . .	9
505	2.2 Purpose of the book . . . . .	9
506	2.3 How to read this book . . . . .	10
507	2.4 Cryptological Systems . . . . .	10
508	2.5 SNARKS . . . . .	10
509	2.6 complexity theory . . . . .	10
510	2.6.1 Runtime complexity . . . . .	10
511	2.7 Software Used in This Book . . . . .	11
512	2.7.1 Sagemath . . . . .	11
513	<b>3 Arithmetics</b>	<b>12</b>
514	3.1 Introduction . . . . .	12
515	3.1.1 Aims and target audience . . . . .	12
516	3.1.2 The structure of this chapter . . . . .	13
517	3.2 Integer Arithmetics . . . . .	13
518	Euclidean Division . . . . .	16
519	The Extended Euclidean Algorithm . . . . .	18
520	3.3 Modular arithmetic . . . . .	19
521	Conguency . . . . .	20
522	Modular Arithmetics . . . . .	20
523	The Chinese Remainder Theorem . . . . .	23
524	Modular Inverses . . . . .	26
525	3.4 Polynomial Arithmetics . . . . .	29
526	Polynomial Arithmetics . . . . .	33
527	Euklidean Division . . . . .	34
528	Prime Factors . . . . .	36
529	Lange interpolation . . . . .	37
530	<b>4 Algebra</b>	<b>40</b>
531	4.1 Groups . . . . .	40
532	Commutative Groups . . . . .	41
533	Finite groups . . . . .	43

534		Generators . . . . .	43
535		The discrete Logarithm problem . . . . .	43
536	4.1.1	Cryptographic Groups . . . . .	44
537		The discrete logarithm assumption . . . . .	45
538		The decisional Diffie–Hellman assumption . . . . .	47
539		The computational Diffie–Hellman assumption . . . . .	47
540		Cofactor Clearing . . . . .	48
541	4.1.2	Hashing to Groups . . . . .	48
542		Hash functions . . . . .	48
543		Hashing to cyclic groups . . . . .	50
544		Hashing to modular arithmetics . . . . .	51
545		Pedersen Hashes . . . . .	55
546		MimC Hashes . . . . .	55
547		Pseudorandom Functions in DDH-A groups . . . . .	55
548	4.2	Commutative Rings . . . . .	55
549		Hashing to Commutative Rings . . . . .	58
550	4.3	Fields . . . . .	58
551		Prime fields . . . . .	60
552		Square Roots . . . . .	61
553		Exponentiation . . . . .	63
554		Hashing into prime fields . . . . .	63
555		Extension Fields . . . . .	63
556		Hashing into extension fields . . . . .	67
557	4.4	Projective Planes . . . . .	67
558	<b>5</b>	<b>Elliptic Curves</b>	<b>70</b>
559	5.1	Elliptic Curve Arithmetics . . . . .	70
560	5.1.1	Short Weierstraß Curves . . . . .	70
561		Affine short Weierstraß form . . . . .	71
562		Affine compressed representation . . . . .	75
563		Affine group law . . . . .	76
564		Scalar multiplication . . . . .	81
565		Projective short Weierstraß form . . . . .	84
566		Projective Group law . . . . .	85
567		Coordinate Transformations . . . . .	86
568	5.1.2	Montgomery Curves . . . . .	86
569		Affine Montgomery Form . . . . .	88
570		Affine Montgomery coordinate transformation . . . . .	89
571		Montgomery group law . . . . .	91
572	5.1.3	Twisted Edwards Curves . . . . .	92
573		Twisted Edwards Form . . . . .	92
574		Twisted Edwards group law . . . . .	94
575	5.2	Elliptic Curve Pairings . . . . .	95
576		Embedding Degrees . . . . .	95
577		Elliptic Curves over extension fields . . . . .	96
578		Full torsion groups . . . . .	97
579		Torsion subgroups . . . . .	100
580		The Weil pairing . . . . .	102

581	5.3	Hashing to Curves . . . . .	104
582		Try-and-increment hash functions . . . . .	105
583	5.4	Constructing elliptic curves . . . . .	107
584		The Trace of Frobenius . . . . .	108
585		The $j$ -invariant . . . . .	109
586		The Complex Multiplication Method . . . . .	110
587		The <i>BLS6_6</i> pen-and-paper curve . . . . .	119
588		Hashing to pairing groups . . . . .	126
589	<b>6</b>	<b>Statements</b>	<b>128</b>
590	6.1	Formal Languages . . . . .	128
591		Decision Functions . . . . .	129
592		Instance and Witness . . . . .	132
593		Modularity . . . . .	135
594	6.2	Statement Representations . . . . .	135
595	6.2.1	Rank-1 Quadratic Constraint Systems . . . . .	136
596		R1CS representation . . . . .	136
597		R1CS Satisfiability . . . . .	138
598		Modularity . . . . .	140
599	6.2.2	Algebraic Circuits . . . . .	140
600		Algebraic circuit representation . . . . .	140
601		Circuit Execution . . . . .	145
602		Circuit Satisfiability . . . . .	147
603		Associated Constraint Systems . . . . .	148
604	6.2.3	Quadratic Arithmetic Programs . . . . .	153
605		QAP representation . . . . .	153
606		QAP Satisfiability . . . . .	155
607	<b>7</b>	<b>Circuit Compilers</b>	<b>159</b>
608	7.1	A Pen-and-Paper Language . . . . .	159
609	7.1.1	The Grammar . . . . .	159
610	7.1.2	The Execution Phases . . . . .	161
611		The Setup Phase . . . . .	161
612		The Prover Phase . . . . .	163
613	7.2	Common Programing concepts . . . . .	163
614	7.2.1	Primitive Types . . . . .	163
615		The base-field type . . . . .	164
616		The Subtraction Constraint System . . . . .	167
617		The Inversion Constraint System . . . . .	168
618		The Division Constraint System . . . . .	169
619		The boolean Type . . . . .	170
620		The boolean Constraint System . . . . .	170
621		The AND operator constraint system . . . . .	171
622		The OR operator constraint system . . . . .	171
623		The NOT operator constraint system . . . . .	172
624		Modularity . . . . .	173
625		Arrays . . . . .	176
626		The Unsigned Integer Type . . . . .	176



627		The uN Constraint System . . . . .	177
628		The Unsigned Integer Operators . . . . .	178
629	7.2.2	Control Flow . . . . .	179
630		The Conditional Assignment . . . . .	179
631		Loops . . . . .	181
632	7.2.3	Binary Field Representations . . . . .	182
633	7.2.4	Cryptographic Primitives . . . . .	184
634		Twisted Edwards curves . . . . .	184
635		Twisted Edwards curve constraints . . . . .	184
636		Twisted Edwards curve addition . . . . .	185
637	<b>8</b>	<b>Zero Knowledge Protocols</b>	<b>187</b>
638	8.1	Proof Systems . . . . .	187
639	8.2	The “Groth16” Protocol . . . . .	188
640		The Setup Phase . . . . .	190
641		The Prover Phase . . . . .	195
642		The Verification Phase . . . . .	198
643		Proof Simulation . . . . .	200
644	<b>9</b>	<b>Exercises and Solutions</b>	<b>204</b>

# Chapter 1

## Introduction

This is dump from other papers as inspiration for the intro:

Zero knowledge proofs are a class of cryptographic protocols in which one can prove honest computation without revealing the inputs to that computation. A simple high-level example of a zero-knowledge proof is the ability to prove one is of legal voting age without revealing the respective age. In a typical zero knowledge proof system, there are two participants: a prover and a verifier. A prover will present a mathematical proof of computation to a verifier to prove honest computation. The verifier will then confirm whether the prover has performed honest computation based on predefined methods. Zero knowledge proofs are of particular interest to public blockchain activities as the verifier can be codified in smart contracts as opposed to trusted parties or third-party intermediaries.

Zero-knowledge proofs (ZKPs) are an important privacy-enhancing tool from cryptography. They allow proving the veracity of a statement, related to confidential data, without revealing any information beyond the validity of the statement. ZKPs were initially developed by the academic community in the 1980s, and have seen tremendous improvements since then. They are now of practical feasibility in multiple domains of interest to the industry, and to a large community of developers and researchers. ZKPs can have a positive impact in industries, agencies, and for personal use, by allowing privacy-preserving applications where designated private data can be made useful to third parties, despite not being disclosed to them.

ZKP systems involve at least two parties: a prover and a verifier. The goal of the prover is to convince the verifier that a statement is true, without revealing any additional information. For example, suppose the prover holds a birth certificate digitally signed by an authority. In order to access some service, the prover may have to prove being at least 18 years old, that is, that there exists a birth certificate, tied to the identity of the prover and digitally signed by a trusted certification authority, stating a birthdate consistent with the age claim. A ZKP allows this, without the prover having to reveal the birthdate.

### 1.1 Target audience

This book is accessible for both beginners and experienced developers alike. Concepts are gradually introduced in a logical and steady pace. Nonetheless, the chapters lend themselves rather well to being read in a different order. More experienced developers might get the most benefit by jumping to the chapters that interest them most. If you like to learn by example, then you should go straight to the chapter on Using Clarity.

It is assumed that you have a basic understanding of programming and the underlying logical concepts. The first chapter covers the general syntax of Clarity but it does not delve into what

programming itself is all about. If this is what you are looking for, then you might have a more difficult time working through this book unless you have an (undiscovered) natural affinity for such topics. Do not let that dissuade you though, find an introductory programming book and press on! The straightforward design of Clarity makes it a great first language to pick up.

## 1.2 The Zoo of Zero-Knowledge Proofs

First, a list of zero-knowledge proof systems:

1. Pinocchio (2013): Paper

– Notes: trusted setup

2. BCGTV (2013): Paper

– Notes: trusted setup, implementation

3. BCTV (2013): Paper

– Notes: trusted setup, implementation

4. Groth16 (2016): Paper

– Notes: trusted setup

– Other resources: Talk in 2019 by Georgios Konstantopoulos

5. GM17 (2017): Paper

– Notes: trusted setup

– Other resources: later Simulation extractability in ROM, 2018

6. Bulletproofs (2017): Paper

– Notes: no trusted setup

– Other resources: Polynomial Commitment Scheme on DL, 2016 and KZG10, Polynomial Commitment Scheme on Pairings, 2010

7. Ligero (2017): Paper

– Notes: no trusted setup

– Other resources:

8. Hyrax (2017): Paper

– Notes: no trusted setup

– Other resources:

9. STARKs (2018): Paper

– Notes: no trusted setup

– Other resources:

## 10. Aurora (2018): Paper

- Notes: transparent SNARK

- Other resources:

## 11. Sonic (2019): Paper

- Notes: SNORK - SNARK with universal and updateable trusted setup, PCS-based

- Other resources: Blog post by Mary Maller from 2019 and work on updateable and universal setup from 2018

## 12. Libra (2019): Paper

- Notes: trusted setup

- Other resources:

## 13. Spartan (2019): Paper

- Notes: transparent SNARK

- Other resources:

## 14. PLONK (2019): Paper

- Notes: SNORK, PCS-based

- Other resources: Discussion on Plonk systems and Awesome Plonk list

## 15. Halo (2019): Paper

- Notes: no trusted setup, PCS-based, recursive

- Other resources:

## 16. Marlin (2019): Paper

- Notes: SNORK, PCS-based

- Other resources: Rust Github

## 17. Fractal (2019): Paper

- Notes: Recursive, transparent SNARK

- Other resources:

## 18. SuperSonic (2019): Paper

- Notes: transparent SNARK, PCS-based

- Other resources: Attack on DARK compiler in 2021

## 19. Redshift (2019): Paper

- Notes: SNORK, PCS-based

- Other resources:

**Other resources on the zoo:** Awesome ZKP list on Github, ZKP community with the reference document

**To Do List**

- Make table for prover time, verifier time, and proof size
- Think of categories - *Achieved Goals*: Trusted setup or not, Post-quantum or not, ...
- Think of categories - *Mathematical background*: Polynomial commitment scheme, ...
- ... while we discuss the points above, we should also discuss a common notation/language for all these things. (E.g. transparent SNARK/no trusted setup/STARK)

**Points to cover while writing**

- Make a historical overview over the "discovery" of the different ZKP systems
- Make reader understand what paper is build on what result etc. - the tree of publications!
- Make reader understand the different terminology, e.g. SNARK/SNORK/STARK, PCS, R1CS, updateable, universal, ...
- Make reader understand the mathematical assumptions - and what this means for the zoo.
- Where will the development/evolution go? What are bottlenecks?

**Other topics I fell into while compiling this list**

- Vector commitments: <https://eprint.iacr.org/2020/527.pdf>
- Snark1: <http://ace.cs.ohio.edu/~gstewart/papers/snaark1.pdf>
- Virgo?: [https://people.eecs.berkeley.edu/~kubitron/courses/cs262a-F19/projects/reports/project5\\_report\\_ver2.pdf](https://people.eecs.berkeley.edu/~kubitron/courses/cs262a-F19/projects/reports/project5_report_ver2.pdf)

# Chapter 2

## Preliminaries

### 2.1 Preface and Acknowledgements

This book began as a set of lecture and notes accompanying the zk-Summit 0x and 0xx .... It arose from the desire to collect the scattered information of snarks [] and present them to an audience that does not have a strong background in cryptography []

### 2.2 Purpose of the book

The first version of this book is written by security auditors at Least Authority where we audited quite a few snark based systems. Its included "what we have learned" destilate of the time we spend on various audits.

We intend to let illustrative examples drive the discussion and present the key concepts of pairing computation with as little machinery as possible. For those that are fresh to pairing-based cryptography, it is our hope that this chapter might be particularly useful as a first read and prelude to more complete or advanced expositions (e.g. the related chapters in [Gal12]).

On the other hand, we also hope our beginner-friendly intentions do not leave any sophisticated readers dissatisfied by a lack of formality or generality, so in cases where our discussion does sacrifice completeness, we will at least endeavour to point to where a more thorough exposition can be found.

One advantage of writing a survey on pairing computation in 2012 is that, after more than a decade of intense and fast-paced research by mathematicians and cryptographers around the globe, the field is now racing towards full maturity. Therefore, an understanding of this text will equip the reader with most of what they need to know in order to tackle any of the vast literature in this remarkable field, at least for a while yet.

Since we are aiming the discussion at active readers, we have matched every example with a corresponding snippet of (hyperlinked) Magma [BCP97] code 1 , where we take inspiration from the helpful Magma pairing tutorial by Dominguez Perez et al. [DKS09].

Early in the book we will develop examples that we then later extend with most of the things we learn in each chapter. This way we incrementally build a few real world snarks but over full fledged cryptographic systems that are nevertheless simple enough to be computed by pen and paper to illustrate all steps in great detail.

## 2.3 How to read this book

Books and papers to read: XXXXXXXXXXXX

Software to try: XXXXXXXXXXXXXXXXXXXX

Correctly prescribing the best reading route for a beginner naturally requires individual diagnosis that depends on their prior knowledge and technical preparation.

## 2.4 Cryptological Systems

The science of information security is referred to as *cryptology*. In the broadest sense, it deals with encryption and decryption processes, with digital signatures, identification protocols, cryptographic hash functions, secrets sharing, electronic voting procedures and electronic money.  
EXPAND

## 2.5 SNARKS

## 2.6 complexity theory

Before we deal with the mathematics behind zero knowledge proof systems, we must first clarify what is meant by the runtime of an algorithm or the time complexity of an entire mathematical problem. This is particularly important for us when we analyze the various snark systems...

For the reader who is interested in complexity theory, we recommend, or example or , as well as the references contained therein.

### 2.6.1 Runtime complexity

The runtime complexity of an algorithm describes, roughly speaking, the amount of elementary computation steps that this algorithm requires in order to solve a problem, depending on the size of the input data.

Of course, the exact amount of arithmetic operations required depends on many factors such as the implementation, the operating system used, the CPU and many more. However, such accuracy is seldom required and is mostly meaningful to consider only the asymptotic computational effort.

In computer science, the runtime of an algorithm is therefore not specified in individual calculation steps, but instead looks for an upper limit which approximates the runtime as soon as the input quantity becomes very large. This can be done using the so-called *Landau notation* (also called big - $\mathcal{O}$ -notation) A precise definition would, however, go beyond the scope of this work and we therefore refer the reader to .

For us, only a rough understanding of transit times is important in order to be able to talk about the security of cryptographic systems. For example,  $\mathcal{O}(n)$  means that the running time of the algorithm to be considered is linearly dependent on the size of the input set  $n$ ,  $\mathcal{O}(n^k)$  means that the running time is polynomial and  $\mathcal{O}(2^n)$  stands for an exponential running time (chapter 2.4).

An algorithm which has a running time that is greater than a polynomial is often simply referred to as *slow*.

A generalization of the runtime complexity of an algorithm is the so-called *time complexity of a mathematical problem*, which is defined as the runtime of the fastest possible algorithm that can still solve this problem ( chapter 3.1).

Since the time complexity of a mathematical problem is concerned with the runtime analysis of all possible (and thus possibly still undiscovered) algorithms, this is often a very difficult and deep-seated question .

For us, the time complexity of the so-called discrete logarithm problem will be important. This is a problem for which we only know slow algorithms on classical computers at the moment, but for which at the same time we cannot rule out that faster algorithms also exist.

STUFF ON CRYPTOGRAPHIC HASH FUNCTIOND

## 2.7 Software Used in This Book

### 2.7.1 Sagemath

It order to provide an interactive learning experience, and to allow getting hands-on with the concepts described in this book, we give examples for how to program them in the Sage programming language. Sage is a dialect of the learning-friendly programming language Python, which was extended and optimized for computing with, in and over algebraic objects. Therefore, we recommend installing Sage before diving into the following chapters.

The installation steps for various system configurations are described on the sage websit <sup>1</sup>. Note however that we use Sage version 9, so if you are using Linux and your package manager only contains version 8, you may need to choose a different installation path, such as using prebuilt binaries.

We recommend the interested reader, who is not familiar with sagemath to read on the many tutorial before starting this book. For example

---

<sup>1</sup><https://doc.sagemath.org/html/en/installation/index.html>



# Chapter 3

## Arithmetics

### 3.1 Introduction

#### 3.1.1 Aims and target audience

The goal of this chapter is to enable a reader who is starting out with nothing more than basic high school algebra to be able to solve basic tasks in elliptic curve cryptography without the need of a computer.

How much mathematics do you need to understand **zero-knowledge proofs**? The answer, of course, depends on the level of understanding you aim for. It is possible to describe zero-knowledge proofs without using mathematics at all; however, to read a foundational paper like Groth [2016], some knowledge of mathematics is needed to be able to follow the discussion.

Without a solid grounding in mathematics, someone who is interested in learning the concepts of zero-knowledge proofs, but who has never seen or dealt with, say, a **finite field**, or an **elliptic curve**, may quickly become overwhelmed. This is not so much due to the complexity of the mathematics needed, rather because of the vast amount of technical jargon, unknown terms, and obscure symbols that quickly makes a text unreadable, even though the concepts themselves are not actually that hard. As a result, the reader might either lose interest, or pick up some incoherent bits and pieces of knowledge that, in the worst case scenario, result in immature code.

This is why we dedicated this chapter to explaining the mathematical foundations needed to understand the basic concepts underlying snark development. We encourage the reader who is not familiar with basic number theory and elliptic curves to take the time and read this and the following chapters, until they are able to solve at least a few exercises in each chapter.

If, on the other hand, you are already skilled in elliptic curve cryptography, feel free to skip this chapter and only come back to it for reference and comparison. Maybe the most interesting parts are XXX .

We start our explanations at a very basic level, and only assume pre-existing knowledge of fundamental concepts like integer arithmetics. At the same time, we'll attempt to teach you to "think mathematically", and to show you that there are numbers and **methatical** structures out there that appear to be very different from the things you learned about in high school, but on a deeper level, they are actually quite similar.

We want to stress, however that this introduction is informal, incomplete and optimized to enable the reader to understand zero-knowledge concepts as efficiently as possible. Our focus and design choices are to include as little theory as necessary, focusing on the wealth of numerical examples. We believe that such an informal, example-driven approach to learning

zero-knowledge proofs

finite field

elliptic curve

add references

methatical

mathematics may make it easier for beginners to digest the material in the initial stages.

For instance, as a beginner, you would probably find it more beneficial to first compute a simple toy **snark** with pen and paper all the way through, before actually developing real-world production-ready systems. In addition, it's useful to have a few simple examples in your head before getting started with reading actual academic papers.

However, in order to be able to derive these toy examples, some mathematical groundwork is needed. This chapter therefore will help you focus on what is important, accompanied by exercises that you are encouraged to recompute yourself. Every section usually ends with a list of additional exercises in increasing order of difficulty, to help the reader memorize and apply the concepts.

### 3.1.2 The structure of this chapter

We start with a brief recapitulation of basic integer arithmetics like long division, the greatest common divisor and Euclid's algorithm. After that, we introduce modular arithmetics as **the most important skill** to compute our pen-and-paper examples. We then introduce polynomials, compute their analogs to integer arithmetics and introduce the important concept of Lagrange interpolation.

After this practical warm up, we introduce some basic algebraic terms like groups and fields, because those terms are used very frequently in academic papers relating to zero-knowledge proofs. The beginner is advised to memorize those terms and think about them. We define these terms in the general abstract way of mathematics, hoping that the non mathematical trained reader will gradually learn to become comfortable with this style. We then give basic examples and do basic computations with these examples to get familiar with the concepts.

## 3.2 Integer Arithmetics

In a sense, integer arithmetics is at the heart of large parts of modern cryptography, because it provides the most basic tools for doing computations in those systems. Fortunately, most readers will probably remember integer arithmetics from school. It is, however, important that you can confidently apply those concepts to understand and execute computations in the many pen-and-paper examples that form an integral part of the MoonMath Manual. We will therefore recapitulate basic arithmetics concepts to refresh your memory and fill any knowledge gaps.

In what follows, we apply standard mathematical notations, and use the symbol  $\mathbb{Z}$  for the set of all **integers**: **S: I think it'd be useful to explain the difference between  $:=$  and  $=$  as well. We have a table on this in the ZKAPs whitepaper.**

**M: Yeah maybe we use the more suggestive leftarrows aks  $\leftarrow$  ? If a table of symbols is unavoidable then ok, I find I super ugly, though**

**S: The table below is similar to what we have in the ZKAPs whitepaper. I think it's easier to read than a wall of text explaining the same things. We can make it more visually different (e.g. typesetting it in "dark mode", and probably move it earlier in the chapter.**

## Notation used in this chapter

Symbol	Meaning of Symbol	Example	Explanation
=	equals	$a = r$	$a$ and $r$ have the same value
:=	defining a variable	$\mathbb{Z} := \{\dots, -2, -1, 0, 1, 2, \dots\}$	$\mathbb{Z}$ is the set of integers
$\in$	from the set	$a \in \mathbb{Z}$	$a$ is an integer

$$\mathbb{Z} := \{\dots, -3, -2, -1, 0, 1, 2, 3, \dots\} \quad (3.1)$$

Integers are also known as **whole numbers**, that is, numbers that can be written without fractional parts. Examples of numbers that are **not** integers are  $\frac{2}{3}$ , 1.2 and  $-1280.006$ .

If  $a \in \mathbb{Z}$  is an integer, then  $|a|$  stands for the **absolute value** of  $a$ , that is, the non-negative value of  $a$  without regard to its sign:

$$|4| = 4 \quad (3.2)$$

$$|-4| = 4 \quad (3.3)$$

In addition, we use the symbol  $\mathbb{N}$  for the set of all **counting numbers** (also called natural numbers). So whenever you see the symbol  $\mathbb{N}$ , think of the set of all non negative integers including the number 0:

$$\mathbb{N} := \{0, 1, 2, 3, \dots\} \quad (3.4)$$

Any number that is smaller than 0, that is, any number that has a minus sign, is not part of  $\mathbb{N}$ . All counting numbers are integers, but not the other way round. In other words, counting numbers are a subset of integers.

To make it easier to memorize new concepts and symbols, we might frequently link to definitions (See 3.1 for a definition of  $\mathbb{Z}$ ) in the beginning, but as to many links render a text unreadable, we will assume the reader will become familiar with definitions as the text proceeds at which point we will not link them anymore.

Both sets  $\mathbb{N}$  and  $\mathbb{Z}$  have a notion of addition and multiplication defined on them. Most of us are probably able to do many integer computations in our head, but this gets more and more difficult as these increase in complexity. We will frequently invoke the SageMath system (2.7.1) for more complicated computations. One way to invoke the integer type in Sage is: **We haven't really talked about what a ring is at this point**

```
sage: ZZ # A sage notation for the integer type
Integer Ring
sage: NN # A sage notation for the counting number type
Non negative integer semiring
sage: ZZ(5) # Get an element from the Ring of integers
5
sage: ZZ(5) + ZZ(3)
8
sage: ZZ(5) * NN(3)
15
sage: ZZ.random_element(10**50)
54428611290136105088662805064077040080301342920296
```

We haven't really talked about what a ring is at this point

```

956 sage: ZZ(27713).str(2) # Binary string representation      13
957 110110001000001                                           14
958 sage: NN(27713).str(2) # Binary string representation      15
959 110110001000001                                           16
960 sage: ZZ(27713).str(16) # Hexadecimal string representation 17
961 6c41                                                       18

```

One set of numbers that is of particular interest to us is **prime numbers**, which are counting numbers  $p \in \mathbb{N}$  with  $p \geq 2$ , which are only divisible by themselves and by 1. All prime numbers apart from the number 2 are called **odd** (since even numbers greater than 2 are all divisible by 2, they are not prime numbers). We write  $\mathbb{P}$  for the set of all prime numbers and  $\mathbb{P}_{\geq 3}$  for the set of all odd prime numbers.  $\mathbb{P}$  is infinite and can be ordered according to size, so that we can write them as follows:

$$2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37, 41, 43, 47, 53, 59, 61, 67, \dots \quad (3.5)$$

This is sequence A000040 in OEIS, the Online Encyclopedia of Integer Sequences. In particular, we can talk about small and large prime numbers.

As the **fundamental theorem of arithmetics** tells us, prime numbers are, in a certain sense, the basic building blocks from which all other natural numbers are composed. To see that, let  $n \in \mathbb{N}_{\geq 2}$  be any natural number. Then there are always prime numbers  $p_1, p_2, \dots, p_k \in \mathbb{P}$ , such that

$$n = p_1 \cdot p_2 \cdot \dots \cdot p_k. \quad (3.6)$$

This representation is unique for each natural number (except for the order of the factors) and is called the **prime factorization** of  $n$ .

*Example 1* (Prime Factorization). To see what we mean by prime factorization of a number, let's look at the number  $19214758032624000 \in \mathbb{N}$ . To get its prime factors, we can successively divide it by all prime numbers in ascending order starting with 2:

$$19214758032624000 = 2 \cdot 2 \cdot 2 \cdot 2 \cdot 2 \cdot 2 \cdot 2 \cdot 3 \cdot 3 \cdot 3 \cdot 5 \cdot 5 \cdot 5 \cdot 7 \cdot 11 \cdot 17 \cdot 17 \cdot 23 \cdot 43 \cdot 43 \cdot 47$$

We can double check our findings invoking Sage, which provides an algorithm to factor counting numbers:

```

978 sage: n = NN(19214758032624000)                             19
979 sage: factor(n)                                              20
980 2^7 * 3^3 * 5^3 * 7 * 11 * 17^2 * 23 * 43^2 * 47           21

```

This computation reveals an important observation: Computing the factorization of an integer is computationally expensive, while the reverse, that is, computing the product of given a set of prime numbers, is fast. **MIRCO: inverse process? Its not reversing something actually, but my english doesn't resolve to this level very well**

From this, an important question arises: How fast we can compute the prime factorization of a natural number? This is the famous **factorization problem** and, as far as we know, there is no method on a classical **Turing machine** that is able to compute this representation in **polynomial time**. The fastest algorithm known today run **sub-exponentially, with  $\mathcal{O}((1 + \epsilon)^n)$  and some  $\epsilon > 0$** .

It follows that number factorization  $\Leftrightarrow$  prime number multiplication is an example of a so-called **one-way function**: Something that is easy to compute in one direction, but hard to

What's the significance of this distinction?

reverse

Turing machine

polynomial time

sub-exponentially with  $\mathcal{O}((1 + \epsilon)^n)$  and

compute in the other direction. **The existence of one-way functions is a basic cryptographic assumptions that the security of many crypto systems is based on.**

It should be pointed out, however, that the American mathematician Peter Williston Shor developed an algorithm in 1994 which can calculate the prime factor representation of a natural number in polynomial time on a quantum computer. The consequence of this is that cryptosystems, which are based on the time complexity of the prime factor problem, are unsafe as soon as practically usable quantum computers become available. *Add text along the lines of “this is the best we got for now” Possibly something on when we can reasonably expect quantum computers to become accessible/usable enough*

Add text

*Exercise 1.* What is the absolute value of the integers  $-123$ ,  $27$  and  $0$ ?

*Exercise 2.* Compute the factorization of  $6469693230$  and double check your results using Sage.

*Exercise 3.* Consider the following equation  $4 \cdot x + 21 = 5$ . Compute the set of all solutions for  $x$  under the following alternative assumptions:

1. The equation is defined over the type of natural numbers.

2. The equation is defined over the type of integers.

*Exercise 4.* Consider the following equation  $2x^3 - x^2 - 2x = -1$ . Compute the set of all solutions  $x$  under the following assumptions:

1. The equation is defined over the type of natural numbers.

2. The equation is defined over the type of integers.

3. The equation is defined over the type  $\mathbb{Q}$  of fractions.

$\mathbb{Q}$  of fractions

**Euclidean Division** Division in the commonly understood sense is not defined for integers, as, for example,  $7$  divided by  $3$  will not be an integer again. However it is possible to divide any two integers with a remainder. So for example  $7$  divided by  $3$  is equal to  $2$  with a remainder of  $1$ , since  $7 = 2 \cdot 3 + 1$ .

Doing integer division like this is probably something many of us remember from school. It is usually called **Euclidean division**, or **division with a remainder**, and it is an essential technique to understand many concepts in this book. The precise definition is as follows:

Let  $a \in \mathbb{Z}$  and  $b \in \mathbb{Z}$  be two integers with  $b \neq 0$ . Then there is always another integer  $m \in \mathbb{Z}$  and a counting number  $r \in \mathbb{N}$ , with  $0 \leq r < |b|$  such that

$$a = m \cdot b + r \quad (3.7)$$

This decomposition of  $a$  given  $b$  is called **Euclidean division**, where  $a$  is called the **dividend**,  $b$  is called the **divisor**,  $m$  is called the **quotient** and  $r$  is called the **remainder**.

*Notation and Symbols* 1. Suppose that the numbers  $a, b, m$  and  $r$  satisfy equation (3.7). Then we often write

$$a \operatorname{div} b := m, \quad a \operatorname{mod} b := r \quad (3.8)$$

to describe the quotient and the remainder of the Euclidean division. We also say that an integer  $a$  is divisible by another integer  $b$  if  $a \operatorname{mod} b = 0$  holds. In this case we also write  $b|a$ .

So, in a nutshell Euclidean division is a process of dividing one integer by another in a way that produces a quotient and a non-negative remainder, the latter of which is smaller than the absolute value of the divisor. It can be shown that both the quotient and the remainder always exist and are unique, as long as the dividend is different from 0.

A special situation occurs whenever the remainder is zero, because in this case the dividend is divisible by the divisor. Our notation  $b|a$  reflects that.

*Example 2.* Applying Euclidean division and our previously defined notation 3.8 to the divisor  $-17$  and the dividend  $4$ , we get

$$-17 \operatorname{div} 4 = -5, \quad -17 \operatorname{mod} 4 = 3$$

because  $-17 = -5 \cdot 4 + 3$  is the Euclidean division of  $-17$  and  $4$  (the remainder is, by definition, a non-negative number). In this case  $4$  does not divide  $-17$ , as the remainder is not zero. The truth value of the expression  $4|-17$  therefore is FALSE. On the other hand, the truth value of  $4|12$  is TRUE, since  $4$  divides  $12$ , as  $12 \operatorname{mod} 4 = 0$ . We can invoke SageMath to do the computation for us. We get the following:

```
sage: ZZ(-17) // ZZ(4) # Integer quotient      22
-5                                              23
sage: ZZ(-17) % ZZ(4) # remainder             24
3                                              25
sage: ZZ(4).divides(ZZ(-17)) # self divides other 26
False                                         27
sage: ZZ(4).divides(ZZ(12))                  28
True                                         29
```

Methods to compute Euclidean division for integers are called **integer division algorithms**. Probably the best known algorithm is the so-called **long division**, which most of us might have learned in school. (It should be noted, however that there are faster methods like **Newton–Raphson division**.)

As long division is the standard method used for pen-and-paper division of multi-digit numbers expressed in decimal notation, the reader should become familiar with it as we use it throughout this book when we do simple pen-and-paper computations. However, instead of defining the algorithm formally, we rather give some examples that will hopefully make the process clear.

In a nutshell, the algorithm loops through the digits of the dividend from the left to right, subtracting the largest possible multiple of the divisor (at the digit level) at each stage; the multiples then become the digits of the quotient, and the remainder is the first digit of the dividend.

*Example 3 (Integer Long Division).* To give an example of integer long division algorithm, let's divide the integer  $a = 143785$  by the number  $b = 17$ . Our goal is therefore to find solutions to equation 3.7, that is, we need to find the quotient  $m \in \mathbb{Z}$  and the remainder  $r \in \mathbb{N}$  such that  $143785 = m \cdot 17 + r$ . Using a notation that is mostly used in Commonwealth countries, we



1063 compute as follows:

$$\begin{array}{r}
 8457 \\
 17 \overline{) 143785} \\
 \underline{136} \phantom{00} \\
 77 \phantom{00} \\
 \underline{68} \phantom{00} \\
 98 \phantom{00} \\
 \underline{85} \phantom{00} \\
 135 \phantom{00} \\
 \underline{119} \phantom{00} \\
 16
 \end{array}
 \tag{3.9}$$

1064 We therefore get  $m = 8457$  as well as  $r = 16$  and indeed we have  $143785 = 8457 \cdot 17 + 16$ ,  
 1065 which we can double check invoking Sage:

```

1066 sage: ZZ(143785).quo_rem(ZZ(17)) # Euclidean Division      30
1067 (8457, 16)                                                  31
1068 sage: ZZ(143785) == ZZ(8457)*ZZ(17) + ZZ(16) # check      32
1069 True                                                         33

```

1070 *Exercise 5* (Integer Long Division). Find an  $m \in \mathbb{Z}$  as well as an  $r \in \mathbb{N}$  such that  $a = m \cdot b +$   
 1071  $r$  holds for the following pairs  $(a, b) = (27, 5)$ ,  $(a, b) = (27, -5)$ ,  $(a, b) = (127, 0)$ ,  $(a, b) =$   
 1072  $(-1687, 11)$  and . In which cases are your solutions unique?

1073  $(a, b) = (0, 7)$

1074 *Exercise 6* (Long Division Algorithm). Write an algorithm in pseudocode that computes integer pseudocode  
 1075 long division, handling all edge cases properly.

1076 **The Extended Euclidean Algorithm** One of the most critical parts in this book is modular  
 1077 arithmetics XXX and its application in the computations in so-called **finite fields**, as we explain  
 1078 in section 4.3. In modular arithmetics, it is sometimes possible to define actual division and  
 1079 multiplicative inverses of numbers that is very different from inverses as we know them from  
 1080 other systems like factional numbers.

1081 However, to actually compute those inverses, we have to get familiar with the so-called  
 1082 **extended Euclidean algorithm**. A few more terms are necessary to explain the concept: The  
 1083 **greatest common divisor** (GCD) of two nonzero integers  $a$  and  $b$  is the greatest non-zero  
 1084 counting number  $d$  such that  $d$  divides both  $a$  and  $b$ , that is,  $d|a$  as well as  $d|b$ . We write  
 1085  $\gcd(a, b) := d$  for this number. In addition, two counting numbers are called **relative primes**  
 1086 or **coprimes**, if their greatest common divisor is 1.

1087 The extended Euclidean algorithm is a method to calculate the greatest common divisor of  
 1088 two counting numbers  $a$  and  $b \in \mathbb{N}$ , as well as two additional integers  $s, t \in \mathbb{Z}$ , such that the  
 1089 following equation holds:

$$\gcd(a, b) = s \cdot a + t \cdot b \tag{3.10}$$

1090 The following pseudocode shows in detail how to calculate these numbers with the extended  
 1091 Euclidean algorithm:

1092 The algorithm is simple enough to be done effectively in pen-and-paper examples, where  
 1093 it is common to write it as a table where the rows represent the while-loop and the columns  
 1094 represent the values of the the array  $r, s$  and  $t$  with index  $k$ . The following example provides a  
 1095 simple execution:

check  
reference

modular  
arith-  
metics

actual  
division

multiplicative  
inverses

factional  
numbers

**Algorithm 1** Extended Euclidean Algorithm**Require:**  $a, b \in \mathbb{N}$  with  $a \geq b$ **procedure** EXT-EUCLID( $a, b$ ) $r_0 \leftarrow a$  $r_1 \leftarrow b$  $s_0 \leftarrow 1$  $s_1 \leftarrow 0$  $k \leftarrow 1$ **while**  $r_k \neq 0$  **do** $q_k \leftarrow r_{k-1} \text{ div } r_k$  $r_{k+1} \leftarrow r_{k-1} - q_k \cdot r_k$  $s_{k+1} \leftarrow s_{k-1} - q_k \cdot s_k$  $k \leftarrow k + 1$ **end while****return**  $\gcd(a, b) \leftarrow r_{k-1}$ ,  $s \leftarrow s_{k-1}$  and  $t := (r_{k-1} - s_{k-1} \cdot a) \text{ div } b$ **end procedure****Ensure:**  $\gcd(a, b) = s \cdot a + t \cdot b$ 

1096 *Example 4.* To illustrate the algorithm, let's apply it to the numbers  $a = 12$  and  $b = 5$ . Since  
 1097  $12, 5 \in \mathbb{N}$  as well as  $12 \geq 5$  all requirements are met and we compute as follows:

k	$r_k$	$s_k$	$t_k = (r_k - s_k \cdot a) \text{ div } b$
0	12	1	0
1	5	0	1
2	2	1	-2
3	1	-2	5

1099 From this we can see that 12 and 5 are relatively prime (coprime), since their greatest common  
 1100 divisor is  $\gcd(12, 5) = 1$  and that the equation  $1 = (-2) \cdot 12 + 5 \cdot 5$  holds. We can also invoke  
 1101 sage to double check our findings:

1102 **sage:** `ZZ(12).xgcd(ZZ(5)) # (gcd(a,b), s, t)`  
 1103 `(1, -2, 5)`

34

35

1104 *Exercise 7* (Extended Euclidean Algorithm). Find integers  $s, t \in \mathbb{Z}$  such that  $\gcd(a, b) = s \cdot a +$   
 1105  $t \cdot b$  holds for the following pairs  $(a, b) = (45, 10)$ ,  $(a, b) = (13, 11)$ ,  $(a, b) = (13, 12)$ . What  
 1106 pairs  $(a, b)$  are coprime?

1107 *Exercise 8* (Towards Prime fields). Let  $n \in \mathbb{N}$  be a counting number and  $p$  a prime number, such  
 1108 that  $n < p$ . What is the greatest common divisor  $\gcd(p, n)$ ?

1109 *Exercise 9.* Find all numbers  $k \in \mathbb{N}$  with  $0 \leq k \leq 100$  such that  $\gcd(100, k) = 5$ .

1110 *Exercise 10.* Show that  $\gcd(n, m) = \gcd(n + m, m)$  for all  $n, m \in \mathbb{N}$ .

**3.3 Modular arithmetic**

1112 In mathematics, **modular arithmetic** is a system of arithmetic for integers, where numbers  
 1113 “wrap around” when reaching a certain value, much like calculations on a clock wrap around  
 1114 whenever the value exceeds the number 12. For example, if the clock shows that it is 11 o'clock,



then 20 hours later it will be 7 o'clock, not 31 o'clock. The number 31 has no meaning on a normal clock that shows hours.

The number at which the wrap occurs is called the **modulus**. Modular arithmetics generalizes the clock example to arbitrary moduli and studies equations and phenomena that arise in this new kind of arithmetics. It is of central importance for understanding most modern crypto systems, in large parts because the exponentiation function has an inverse with respect to certain moduli that is hard to compute. In addition, we will see that it provides the foundation of what is called finite fields ().

exponentiation  
function

See XXX

Although modular arithmetic appears very different from ordinary integer arithmetic that we are all familiar with, we encourage the interested reader to work through the example and to discover that, once they get used to the idea that this is a new kind of calculations, it will seem much less daunting.

**Congruency** In what follows, let  $n \in \mathbb{N}$  with  $n \geq 2$  be a fixed counting number that we will call the **modulus** of our modular arithmetics system. With such an  $n$  given, we can then group integers into classes, by saying that two integers are in the same class, whenever their Euclidean division 3.2 by  $n$  will give the same remainder. We then say that two numbers are **congruent** whenever they are in the same class.

*Example 5.* If we choose  $n = 12$  as in our clock example, then the integers  $-7, 5, 17$  and  $29$  are all congruent with respect to  $12$ , since all of them have the remainder  $5$  if we perform Euclidean division on them by  $12$ . In the picture of an analog 12-hour clock, starting at 5 o'clock, when we add 12 hours we are again at 5 o'clock, representing the number  $17$ . On the other hand, when we subtract 12 hours, we are at 5 o'clock again, representing the number  $-7$ .

We can formalize this intuition of what congruency should be into a proper definition utilizing Euclidean division (as explained previously in 3.2): Let  $a, b \in \mathbb{Z}$  be two integers and  $n \in \mathbb{N}$  a natural number. Then  $a$  and  $b$  are said to be **congruent with respect to the modulus  $n$** , if and only if the following equation holds

$$a \bmod n = b \bmod n \quad (3.11)$$

If, on the other hand, two numbers are not congruent with respect to a given modulus  $n$ , we call them **incongruent** w.r.t.  $n$ .

A **congruency** is then nothing but an equation "up to congruency", which means that the equation only needs to hold if we take the modulus on both sides. In which case we write

$$a \equiv b \pmod{n} \quad (3.12)$$

*Exercise 11.* Which of the following pairs of numbers are congruent with respect to the modulus 13:  $(5, 19), (13, 0), (-4, 9), (0, 0)$ .

*Exercise 12.* Find all integers  $x$ , such that the congruency  $x \equiv 4 \pmod{6}$  is satisfied.

**Modular Arithmetics** One particularly useful thing about congruencies is that we can do calculations (arithmetics), much like we can with integer equations. That is, we can add or multiply numbers on both sides. The main difference is probably that the congruency  $a \equiv b \pmod{n}$  is only equivalent to the congruency  $k \cdot a \equiv k \cdot b \pmod{n}$  for some non zero integer  $k \in \mathbb{Z}$ , whenever  $k$  and the modulus  $n$  are coprime. The following list gives a set of useful rules:

Suppose that the congruencies  $a_1 \equiv b_1 \pmod{n}$  as well as  $a_2 \equiv b_2 \pmod{n}$  are satisfied for integers  $a_1, a_2, b_1, b_2 \in \mathbb{Z}$  and that  $k \in \mathbb{Z}$  is another integer. Then:

- 1156 •  $a_1 + k \equiv b_1 + k \pmod{n}$  (compatibility with translation)
- 1157 •  $k \cdot a_1 \equiv k \cdot b_1 \pmod{n}$  (compatibility with scaling)
- 1158 •  $a_1 + a_2 \equiv b_1 + b_2 \pmod{n}$  (compatibility with addition)
- 1159 •  $a_1 \cdot a_2 \equiv b_1 \cdot b_2 \pmod{n}$  (compatibility with multiplication)

1160 Other rules, such as compatibility with subtraction and exponentiation, follow from the rules  
 1161 above. For example, compatibility with subtraction follows from compatibility with scaling by  
 1162  $k = -1$  and compatibility with addition.

1163 Note that the previous rules are implications, not equivalences, which means that you can  
 1164 not necessarily reverse those rules. The following rules makes this precise:

- 1165 • If  $a_1 + k \equiv b_1 + k \pmod{n}$ , then  $a_1 \equiv b_1 \pmod{n}$
- 1166 • If  $k \cdot a_1 \equiv k \cdot b_1 \pmod{n}$  and  $k$  is coprime with  $n$ , then  $a_1 \equiv b_1 \pmod{n}$
- 1167 • If  $k \cdot a_1 \equiv k \cdot b_1 \pmod{k \cdot n}$ , then  $a_1 \equiv b_1 \pmod{n}$

1168 Another property of congruencies, not known in the traditional arithmetics of integers is the  
 1169 **Fermat's Little Theorem**. In simple words, it states that, in modular arithmetics, every number  
 1170 raised to the power of a prime number modulus is congruent to the number itself. Or, to be more  
 1171 precise, if  $p \in \mathbb{P}$  is a prime number and  $k \in \mathbb{Z}$  is an integer, then:

$$k^p \equiv k \pmod{p}, \quad (3.13)$$

1172 If  $k$  is coprime to  $p$ , then we can divide both sides of this congruency by  $k$  and rewrite the  
 1173 expression into the equivalent form

$$k^{p-1} \equiv 1 \pmod{p} \quad (3.14)$$

1174 We can use Sage to compute examples for both  $k$  being coprime and not coprime to  $p$ :

```

1175 sage: ZZ(137).gcd(ZZ(64))
1176 1
1177 sage: ZZ(64)**ZZ(137) % ZZ(137) == ZZ(64) % ZZ(137)
1178 True
1179 sage: ZZ(64)**ZZ(137-1) % ZZ(137) == ZZ(1) % ZZ(137)
1180 True
1181 sage: ZZ(1918).gcd(ZZ(137))
1182 137
1183 sage: ZZ(1918)**ZZ(137) % ZZ(137) == ZZ(1918) % ZZ(137)
1184 True
1185 sage: ZZ(1918)**ZZ(137-1) % ZZ(137) == ZZ(1) % ZZ(137)
1186 False

```

This Sage  
 snippet  
 should be  
 described  
 in more  
 detail.

1187 Now, since for the sake of readers who have never encountered modular arithmetics before, let's  
 1188 compute an example that contains most of the concepts described in this section:

*Example 6.* Assume that we choose the modulus 6 and that our task is to solve the following congruency for  $x \in \mathbb{Z}$

$$7 \cdot (2x + 21) + 11 \equiv x - 102 \pmod{6}$$

As many rules for congruencies are more or less same as for integers, we can proceed in a similar way as we would if we had an equation to solve. Since both sides of a congruency contain ordinary integers, we can rewrite the left side as follows:  $7 \cdot (2x + 21) + 11 = 14x + 147 = 14x + 158$ . We can therefore rewrite the congruency into the equivalent form

$$14x + 158 \equiv x - 102 \pmod{6}$$

In the next step we want to shift all instances of  $x$  to left and every other term to the right. So we apply the “compatibility with translation” rules two times. In a first step we choose  $k = -x$  and in a second step we choose  $k = -158$ . Since “compatibility with translation” transforms a congruency into an equivalent form, the solution set will not change and we get

$$\begin{aligned} 14x + 158 &\equiv x - 102 \pmod{6} \Leftrightarrow \\ 14x - x + 158 - 158 &\equiv x - x - 102 - 158 \pmod{6} \Leftrightarrow \\ 13x &\equiv -260 \pmod{6} \end{aligned}$$

If our congruency would just be a normal integer equation, we would divide both sides by 13 to get  $x = -20$  as our solution. However, in case of a congruency, we need to make sure that the modulus and the number we want to divide by are coprime first – only then will we get an equivalent expression. So we need to find the greatest common divisor  $\gcd(13, 6)$ . Since 13 is prime and 6 is not a multiple of 13, we know that  $\gcd(13, 6) = 1$ , so these numbers are indeed coprime. We therefore compute

$$13x \equiv -260 \pmod{6} \Leftrightarrow x \equiv -20 \pmod{6}$$

Our task is now to find all integers  $x$ , such that  $x$  is congruent to  $-20$  with respect to the modulus 6. So we have to find all  $x$  such

$$x \bmod 6 = -20 \bmod 6$$

Since  $-4 \cdot 6 + 4 = -20$  we know  $-20 \bmod 6 = 4$  and hence we know that  $x = 4$  is a solution to this congruency. However, 22 is another solution since  $22 \bmod 6 = 4$  as well, and so is  $-20$ . In fact, there are infinitely many solutions given by the set

$$\{\dots, -8, -2, 4, 10, 16, \dots\} = \{4 + k \cdot 6 \mid k \in \mathbb{Z}\}$$

Putting all this together, we have shown that the every  $x$  from the set  $\{x = 4 + k \cdot 6 \mid k \in \mathbb{Z}\}$  is a solution to the congruency  $7 \cdot (2x + 21) + 11 \equiv x - 102 \pmod{6}$ . We double ckeck for, say,  $x = 4$  as well as  $x = 14 + 12 \cdot 6 = 86$  using sage:

```

1192 sage: (ZZ(7) * (ZZ(2) * ZZ(4) + ZZ(21)) + ZZ(11)) % ZZ(6) == (ZZ 48
1193       (4) - ZZ(102)) % ZZ(6)
1194 True
1195 sage: (ZZ(7) * (ZZ(2) * ZZ(76) + ZZ(21)) + ZZ(11)) % ZZ(6) == ( 50
1196       ZZ(76) - ZZ(102)) % ZZ(6)
1197 True

```

Readers who had not been familiar with modular arithmetics until now and who might be discouraged by how complicated modular arithmetics seems at this point, should keep two things in mind. First, computing congruencies in modular arithmetics is not really more complicated than computations in more familiar number systems (e.g. fractional numbers), it is just

1202 a matter of getting used to it. Second, the theory of **prime fields** (and more general **residue class**  
 1203 **rings**) takes a different view on modular arithmetics with the attempt to simplify matters. In  
 1204 other words, once we understand prime field arithmetics, things become conceptually cleaner  
 1205 and more easy to compute.

prime  
fieldsresidue  
class rings

1206 *Exercise 13.* Choose the modulus 13 and find all solutions  $x \in \mathbb{Z}$  to the following congruency  
 1207  $5x + 4 \equiv 28 + 2x \pmod{13}$

1208 *Exercise 14.* Choose the modulus 23 and find all solutions  $x \in \mathbb{Z}$  to the following congruency  
 1209  $69x \equiv 5 \pmod{23}$

1210 *Exercise 15.* Choose the modulus 23 and find all solutions  $x \in \mathbb{Z}$  to the following congruency  
 1211  $69x \equiv 46 \pmod{23}$

1212 **The Chinese Remainder Theorem** We have seen how to solve congruencies in modular  
 1213 arithmetic. However, one question that remains is how to solve systems of congruencies with  
 1214 different moduli? The answer is given by the **Chinese remainder theorem**, which states that  
 1215 for any  $k \in \mathbb{N}$  and coprime natural numbers  $n_1, \dots, n_k \in \mathbb{N}$  as well as integers  $a_1, \dots, a_k \in \mathbb{Z}$ , the  
 1216 so-called **simultaneous congruency**

$$\begin{aligned} x &\equiv a_1 \pmod{n_1} \\ x &\equiv a_2 \pmod{n_2} \\ &\dots \\ x &\equiv a_k \pmod{n_k} \end{aligned} \tag{3.15}$$

1217 has a solution, and all possible solutions of this congruence system are congruent modulo the  
 product  $N = n_1 \cdot \dots \cdot n_k$ .<sup>1</sup> In fact, the following algorithm computes the solution set:

check  
algorithm  
floating

---

#### Algorithm 2 Chinese Remainder Theorem

---

**Require:**  $n_0, \dots, n_{k-1} \in \mathbb{N}$  coprime

**procedure** CONGRUENCY-SYSTEMS-SOLVER( $k, a_0, \dots, a_{k-1}, n_0, \dots, n_{k-1}$ )

$N \leftarrow n_0 \cdot \dots \cdot n_{k-1}$

**while**  $j < k$  **do**

$N_j \leftarrow N / n_j$

$(\_, s_j, t_j) \leftarrow EXT - EUCLID(N_j, n_j)$

$$\triangleright 1 = s_j \cdot N_j + t_j \cdot n_j$$

**end while**

$x' \leftarrow \sum_{j=0}^{k-1} a_j \cdot s_j \cdot N_j$

$x \leftarrow x' \bmod N$

**return**  $\{x + m \cdot N \mid m \in \mathbb{Z}\}$

**end procedure**

**Ensure:**  $\{x + m \cdot N \mid m \in \mathbb{Z}\}$  is the complete solution set to 3.15.

---

1218

*Example 7.* To illustrate how to solve simultaneous congruences using the Chinese remainder theorem, let's look at the following system of congruencies:

$$\begin{aligned} x &\equiv 4 \pmod{7} \\ x &\equiv 1 \pmod{3} \\ x &\equiv 3 \pmod{5} \\ x &\equiv 0 \pmod{11} \end{aligned}$$

---

<sup>1</sup>This is the classical Chinese remainder theorem as it was already known in ancient China. Under certain circumstances, the theorem can be extended to non-coprime moduli  $n_1, \dots, n_k$  but this is beyond the scope of this book. Interested readers should consult XXX [add references](#)

Clearly all moduli are coprime and we have  $N = 7 \cdot 3 \cdot 5 \cdot 11 = 1155$ , as well as  $N_1 = 165$ ,  $N_2 = 385$ ,  $N_3 = 231$  and  $N_4 = 105$ . From this we calculate with the extended Euclidean algorithm

$$\begin{aligned} 1 &= 2 \cdot 165 + (-47) \cdot 7 \\ 1 &= 1 \cdot 385 + (-128) \cdot 3 \\ 1 &= 1 \cdot 231 + (-46) \cdot 5 \\ 1 &= 2 \cdot 105 + (-19) \cdot 11 \end{aligned}$$

so we have  $x = 4 \cdot 2 \cdot 165 + 1 \cdot 1 \cdot 385 + 3 \cdot 1 \cdot 231 + 0 \cdot 2 \cdot 105 = 2398$  as one solution. Because  $2398 \bmod 1155 = 88$  the set of all solutions is  $\{\dots, -2222, -1067, 88, 1243, 2398, \dots\}$ . In particular, there are infinitely many different solutions. We can invoke Sage's computation of the Chinese Remainder Theorem (CRT) to double check our findings:

```
1223 sage: CRT_list([4,1,3,0], [7,3,5,11]) 52
1224 88 53
```

As we have seen in various examples before, computing congruencies can be cumbersome and solution sets are large in general. It is therefore advantageous to find some kind of simplification for modular arithmetic.

Fortunately, this is possible and relatively straightforward once we consider all integers that have the same remainder with respect to a given modulus  $n$  in Euclidean division to be equivalent. Then we can go a step further, and identify each set of numbers with equal remainder with that remainder and call it a **remainder class** or **residue class** in modulo  $n$  arithmetics.

It then follows from the properties of Euclidean division that there are exactly  $n$  different remainder classes for every modulus  $n$  and that integer addition and multiplication can be projected to a new kind of addition and multiplication on those classes.

Roughly speaking, the new rules for addition and multiplication are then computed by taking any element of the first equivalence class and some element of the second, then add or multiply them in the usual way and see which equivalence class the result is contained in. The following example makes this abstract description more concrete:

*Example 8* (Arithmetics modulo 6). Choosing the modulus  $n = 6$ , we have six equivalence classes of integers which are congruent modulo 6 (they have the same remainder when divided by 6) and when we identify each of those remainder classes with the remainder, we get the following identification:

$$\begin{aligned} 0 &:= \{\dots, -6, 0, 6, 12, \dots\} \\ 1 &:= \{\dots, -5, 1, 7, 13, \dots\} \\ 2 &:= \{\dots, -4, 2, 8, 14, \dots\} \\ 3 &:= \{\dots, -3, 3, 9, 15, \dots\} \\ 4 &:= \{\dots, -2, 4, 10, 16, \dots\} \\ 5 &:= \{\dots, -1, 5, 11, 17, \dots\} \end{aligned}$$

Now to compute the addition of those equivalence classes, say  $2 + 5$ , one chooses arbitrary elements from both sets, say 14 and  $-1$ , adds those numbers in the usual way and then looks at the equivalence class of the result.

So we get  $14 + (-1) = 13$ , and 13 is in the equivalence class (of) 1. Hence we find that  $2 + 5 = 1$  in modular 6 arithmetics, which is a more readable way to write the congruency  $2 + 5 \equiv 1 \pmod{6}$ .

Applying the same reasoning to all equivalence classes, addition and multiplication can be transferred to equivalence classes. The results for modulus 6 arithmetics are summarized in the following addition and multiplication tables:

1248

+	0	1	2	3	4	5
0	0	1	2	3	4	5
1	1	2	3	4	5	0
2	2	3	4	5	0	1
3	3	4	5	0	1	2
4	4	5	0	1	2	3
5	5	0	1	2	3	4

·	0	1	2	3	4	5
0	0	0	0	0	0	0
1	0	1	2	3	4	5
2	0	2	4	0	2	4
3	0	3	0	3	0	3
4	0	4	2	0	4	2
5	0	5	2	3	2	1

1249 This way, we have defined a new arithmetic system that contains just 6 numbers and comes with  
 1250 its own definition of addition and multiplication. It is called **modular 6 arithmetics** and written  
 1251 as  $\mathbb{Z}_6$ .

1252 To see why such an identification of a congruency class with its remainder is useful and  
 1253 actually simplifies congruency computations a lot, lets go back to the congruency from example  
 1254 6 again:

$$7 \cdot (2x + 21) + 11 \equiv x - 102 \pmod{6} \quad (3.16)$$

1255 As shown in example 6, the arithmetics of congruencies can deviate from ordinary arith-  
 1256 metics: For example, division needs to check whether the modulus and the dividend are co-  
 1257 primes, and solutions are not unique in general.

We can rewrite this congruency as an **equation** over our new arithmetic type  $\mathbb{Z}_6$  by **project-  
 ing onto the remainder classes**. In particular, since  $7 \bmod 6 = 1$ ,  $21 \bmod 6 = 3$ ,  $11 \bmod 6 = 5$   
 and  $102 \bmod 6 = 0$  we have

$$7 \cdot (2x + 21) + 11 \equiv x - 102 \pmod{6} \text{ over } \mathbb{Z} \\ \Leftrightarrow 1 \cdot (2x + 3) + 5 = x \text{ over } \mathbb{Z}_6$$

1258 We can use the multiplication and addition table above to solves the equation on the right like  
 1259 we would solve normal integer equations:

$$\begin{aligned} 1 \cdot (2x + 3) + 5 &= x \\ 2x + 3 + 5 &= x && \# \text{ addition-table: } 3 + 5 = 2 \\ 2x + 2 &= x && \# \text{ add 4 and } -x \text{ on both sides} \\ 2x + 2 + 4 - x &= x + 4 - x && \# \text{ addition-table: } 2 + 4 = 0 \\ x &= 4 \end{aligned}$$

1260 As we can see, despite the somewhat unfamiliar rules of addition and multiplication, solving  
 1261 congruencies this way is very similar to solving normal equations. And, indeed, the solution  
 1262 set is identical to the solution set of the original congruency, since 4 is identified with the set  
 1263  $\{4 + 6 \cdot k \mid k \in \mathbb{Z}\}$ .

1264 We can invoke Sage to do computations in our modular 6 arithmetics type. This is particu-  
 1265 larly useful to double-check our computations:

```
1266 sage: Z6 = Integers(6)                                     54
1267 sage: Z6(2) + Z6(5)                                       55
1268 1                                                         56
1269 sage: Z6(7) * (Z6(2) * Z6(4) + Z6(21)) + Z6(11) == Z6(4) - Z6(102) 57
1270 True                                                       58
```



1271 *Jargon 1* ( $k$ -bit modulus). In cryptographic papers, we can sometimes read phrases like “[...]”  
 1272 using a 4096-bit modulus”. This means that the underlying modulus  $n$  of the modular arithmetic  
 1273 used in the system has a binary representation with a length of 4096 bits. In contrast, the  
 1274 number 6 has the binary representation 110 and hence our example 8 describes a 3-bit modulus  
 1275 arithmetics system.

1276 *Exercise 16.* Let  $a, b, k$  be integers, such that  $a \equiv b \pmod{n}$  holds. Show  $a^k \equiv b^k \pmod{n}$ .

1277 *Exercise 17.* Let  $a, n$  be integers, such that  $a$  and  $n$  are not coprime. For which  $b \in \mathbb{Z}$  does the  
 1278 congruency  $a \cdot x \equiv b \pmod{n}$  have a solution  $x$  and how does the solution set look in that  
 1279 case?

1280 **Modular Inverses** As we know, integers can be added, subtracted and multiplied so that the  
 1281 result is also an integer, but this is not true for the division of integers in general: for example,  
 1282  $3/2$  is not an integer anymore. To see why this is, from a more theoretical perspective, let us  
 1283 consider the definition of a multiplicative inverse first. When we have a set that has some kind  
 1284 of multiplication defined on it and we have a distinguished element of that set that behaves  
 1285 neutrally with respect to that multiplication (doesn’t change anything when multiplied with any  
 1286 other element), then we can define **multiplicative inverses** in the following way:

1287 Let  $S$  be our set that has some notion  $a \cdot b$  of multiplication and a **neutral element**  $1 \in S$ ,  
 1288 such that  $1 \cdot a = a$  for all elements  $a \in S$ . Then a **multiplicative inverse**  $a^{-1}$  of an element  $a \in S$   
 1289 is defined as follows:

$$a \cdot a^{-1} = 1 \quad (3.17)$$

1290 Informally speaking, the definition of a multiplicative inverse is means that it “cancels” the  
 1291 original element to give 1 when they are multiplied.

1292 Numbers that have multiplicative inverses are of particular interest, because they immedi-  
 1293 ately lead to the definition of division by those numbers. In fact, if  $a$  is number such that the  
 1294 multiplicative inverse  $a^{-1}$  exists, then we define **division** by  $a$  simply as multiplication by the  
 1295 inverse:

$$\frac{b}{a} := b \cdot a^{-1} \quad (3.18)$$

1296 *Example 9.* Consider the set of rational numbers, also known as fractions,  $\mathbb{Q}$ . For this set, the  
 1297 neutral element of multiplication is 1, since  $1 \cdot a = a$  for all rational numbers. For example,  
 1298  $1 \cdot 4 = 4$ ,  $1 \cdot \frac{1}{4} = \frac{1}{4}$ , or  $1 \cdot 0 = 0$  and so on.

1299 Every rational number  $a \neq 0$  has a multiplicative inverse, given by  $\frac{1}{a}$ . For example, the  
 1300 multiplicative inverse of 3 is  $\frac{1}{3}$ , since  $3 \cdot \frac{1}{3} = 1$ , the multiplicative inverse of  $\frac{5}{7}$  is  $\frac{7}{5}$ , since  $\frac{5}{7} \cdot \frac{7}{5} = 1$ ,  
 1301 and so on.

1302 *Example 10.* Looking at the set  $\mathbb{Z}$  of integers, we see that with respect to multiplication the  
 1303 neutral element is the number 1 and we notice that no integer  $a \neq 1$  has a multiplicative inverse,  
 1304 since the equation  $a \cdot x = 1$  has no integer solutions for  $a \neq 1$ .

1305 The definition of multiplicative inverse works verbatim for addition as well. In the case of  
 1306 integers, the neutral element with respect to addition is 0, since  $a + 0 = a$  for all integers  $a \in \mathbb{Z}$ .  
 1307 The additive inverse always exist and is given by the negative number  $-a$ , since  $a + (-a) = 0$ .

1308 *Example 11.* Looking at the set  $\mathbb{Z}_6$  of residual classes modulo 6 from example 8, we can use the  
 1309 multiplication table to find multiplicative inverses. To do so, we look at the row of the element  
 1310 and then find the entry equal to 1. If such an entry exists, the element of that column is the  
 1311 multiplicative inverse. If, on the other hand, the row has no entry equal to 1, we know that the  
 1312 element has no multiplicative inverse.

(-1)  
should  
be (-a)?

For example in  $\mathbb{Z}_6$  the multiplicative inverse of 5 is 5 itself, since  $5 \cdot 5 = 1$ . We can also see that 5 and 1 are the only elements that have multiplicative inverses in  $\mathbb{Z}_6$ .

Now, since 5 has a multiplicative inverse modulo 6, it makes sense to “divide” by 5 in  $\mathbb{Z}_6$ . For example

$$\frac{4}{5} = 4 \cdot 5^{-1} = 4 \cdot 5 = 2$$

From the last example, we can make the interesting observation that while 5 has no multiplicative inverse as an integer, it has a multiplicative inverse in modular 6 arithmetics.

The remaining question is to understand which elements have multiplicative inverses in modular arithmetics. The answer is that, in modular  $n$  arithmetics, a residue class  $r$  has a multiplicative inverse, if and only if  $n$  and  $r$  are coprime. Since  $\text{ggt}(n, r) = 1$  in that case, we know from the extended Euclidean algorithm that there are numbers  $s$  and  $t$ , such that

$$1 = s \cdot n + t \cdot r \quad (3.19)$$

If we take the modulus  $n$  on both sides, the term  $s \cdot n$  vanishes, which tells us that  $t \bmod n$  is the multiplicative inverse of  $r$  in modular  $n$  arithmetics.

*Example 12* (Multiplicative inverses in  $\mathbb{Z}_6$ ). In the previous example, we looked up multiplicative inverses in  $\mathbb{Z}_6$  from the lookup-table in Example 8. In real world examples, it is usually impossible to write down those lookup tables, as the modulus is way too large, and the sets occasionally contain more elements than there are atoms in the observable universe.

Now, trying to determine that  $2 \in \mathbb{Z}_6$  has no multiplicative inverse in  $\mathbb{Z}_6$  without using the lookup table, we immediately observe that 2 and 6 are not coprime, since their greatest common divisor is 2. It follows that equation 3.19 has no solutions  $s$  and  $t$ , which means that 2 has no multiplicative inverse in  $\mathbb{Z}_6$ .

The same reasoning works for 3 and 4, as neither of these are coprime with 6. The case of 5 is different, since  $\text{ggt}(6, 5) = 1$ . To compute the multiplicative inverse of 5, we use the extended Euclidean algorithm and compute the following:

k	$r_k$	$s_k$	$t_k = (r_k - s_k \cdot a) \bmod b$
0	6	1	0
1	5	0	1
2	1	1	-1
3	0	.	.

We get  $s = 1$  as well as  $t = -1$  and have  $1 = 1 \cdot 6 - 1 \cdot 5$ . From this, it follows that  $-1 \bmod 6 = 5$  is the multiplicative inverse of 5 in modular 6 arithmetics. We can double check using Sage:

```
sage: ZZ(6).xgcd(ZZ(5))
(1, 1, -1)
```

59  
60

At this point, the attentive reader might notice that the situation where the modulus is a prime number is of particular interest, because we know from exercise XXX that in these cases all remainder classes must have modular inverses, since  $\text{ggt}(r, n) = 1$  for prime  $n$  and  $r < n$ . In fact, Fermat’s little theorem provides a way to compute multiplicative inverses in this situation, since in case of a prime modulus  $p$  and  $r < p$ , we get the following:

$$\begin{aligned} r^p &\equiv r \pmod{p} \Leftrightarrow \\ r^{p-1} &\equiv 1 \pmod{p} \Leftrightarrow \\ r \cdot r^{p-2} &\equiv 1 \pmod{p} \end{aligned}$$



1339 This tells us that the multiplicative inverse of a residue class  $r$  in modular  $p$  arithmetic is pre-  
 1340 cisely  $r^{p-2}$ .

*Example 13* (Modular 5 arithmetics). To see the unique properties of modular arithmetics when-  
 ever the modulus is a prime number, we will replicate our findings from example 8, but this time  
 for the prime modulus 5. For  $n = 5$  we have five equivalence classes of integers which are con-  
 gruent modulo 5. We write this as follows:

$$\begin{aligned} 0 &:= \{\dots, -5, 0, 5, 10, \dots\} \\ 1 &:= \{\dots, -4, 1, 6, 11, \dots\} \\ 2 &:= \{\dots, -3, 2, 7, 12, \dots\} \\ 3 &:= \{\dots, -2, 3, 8, 13, \dots\} \\ 4 &:= \{\dots, -1, 4, 9, 14, \dots\} \end{aligned}$$

1341 Addition and multiplication can be transferred to the equivalence classes, in a way exactly  
 1342 parallel to Example 8. This results in the following addition and multiplication tables:

	+	0	1	2	3	4		·	0	1	2	3	4
	0	0	1	2	3	4		0	0	0	0	0	0
	1	1	2	3	4	0		1	0	1	2	3	4
1343	2	2	3	4	0	1		2	0	2	4	1	3
	3	3	4	0	1	2		3	0	3	1	4	2
	4	4	0	1	2	3		4	0	4	3	2	1

1344 Calling the set of remainder classes in modular 5 arithmetics with this addition and multiplica-  
 1345 tion  $\mathbb{Z}_5$ , we see some subtle but important differences to the situation in  $\mathbb{Z}_6$ . In particular, we  
 1346 see that in the multiplication table, every remainder  $r \neq 0$  has the entry 1 in its row and therefore  
 1347 has a multiplicative inverse. In addition, there are no non-zero elements such that their product  
 1348 is zero.

1349 To use Fermat's little theorem in  $\mathbb{Z}_5$  for computing multiplicative inverses (instead of using  
 1350 the multiplication table), let's consider  $3 \in \mathbb{Z}_3$ . We know that the multiplicative inverse is given  
 1351 by the remainder class that contains  $3^{5-2} = 3^3 = 3 \cdot 3 \cdot 3 = 4 \cdot 3 = 2$ . And indeed  $3^{-1} = 2$ , since  
 1352  $3 \cdot 2 = 1$  in  $\mathbb{Z}_5$ .

1353 We can invoke Sage to do computations in our modular 5 arithmetics type to double-check  
 1354 our computations:

```

1355 sage: Z5 = Integers(5)                                     61
1356 sage: Z5(3) ** (5-2)                                       62
1357 2                                                            63
1358 sage: Z5(3) ** (-1)                                        64
1359 2                                                            65
1360 sage: Z5(3) ** (5-2) == Z5(3) ** (-1)                     66
1361 True                                                         67

```

*Example 14.* To understand one of the principal differences between prime number modular  
 arithmetics and non-prime number modular arithmetics, consider the linear equation  $a \cdot x + b = 0$   
 defined over both types  $\mathbb{Z}_5$  and  $\mathbb{Z}_6$ . Since in  $\mathbb{Z}_5$  every non-zero element has a multiplicative  
 inverse, we can always solve these types of equations in  $\mathbb{Z}_5$ , which is not true in  $\mathbb{Z}_6$ . To see that,

consider the equation  $3x + 3 = 0$ . In  $\mathbb{Z}_5$  we have the following:

$$\begin{array}{ll}
 3x + 3 = 0 & \# \text{ add 2 and on both sides} \\
 3x + 3 + 2 = 2 & \# \text{ addition-table: } 2 + 3 = 0 \\
 3x = 2 & \# \text{ divide by 3} \\
 2 \cdot (3x) = 2 \cdot 2 & \# \text{ multiplication-table: } 2 + 2 = 4 \\
 x = 4 &
 \end{array}$$

So in the case of our prime number modular arithmetics, we get the unique solution  $x = 4$ . Now consider  $\mathbb{Z}_6$ :

$$\begin{array}{ll}
 3x + 3 = 0 & \# \text{ add 3 and on both sides} \\
 3x + 3 + 3 = 3 & \# \text{ addition-table: } 3 + 3 = 0 \\
 3x = 3 & \# \text{ no multiplicative inverse of 3 exists}
 \end{array}$$

So, in this case, we cannot solve the equation for  $x$  by dividing by 3. And, indeed, when we look at the multiplication table of  $\mathbb{Z}_6$  (Example 8), we find that there are three solutions  $x \in \{1, 3, 5\}$ , such that  $3x + 3 = 0$  holds true for all of them.

*Exercise 18.* Consider the modulus  $n = 24$ . Which of the integers 7, 1, 0, 805,  $-4255$  have multiplicative inverses in modular 24 arithmetics? Compute the inverses, in case they exist.

*Exercise 19.* Find the set of all solutions to the congruency  $17(2x + 5) - 4 \equiv 2x + 4 \pmod{5}$ . Then project the congruency into  $\mathbb{Z}_5$  and solve the resulting equation in  $\mathbb{Z}_5$ . Compare the results.

*Exercise 20.* Find the set of all solutions to the congruency  $17(2x + 5) - 4 \equiv 2x + 4 \pmod{6}$ . Then project the congruency into  $\mathbb{Z}_6$  and try to solve the resulting equation in  $\mathbb{Z}_6$ .

## 3.4 Polynomial Arithmetics

A polynomial is an expression consisting of variables (also called indeterminates) and coefficients that involves only the operations of addition, subtraction, multiplication, and non-negative integer exponentiation of variables. All coefficients of a polynomial must have the same type, e.g. being integers or fractions etc. To be more precise a *univariate polynomial* is an expression

$$P(x) := \sum_{j=0}^m a_j x^j = a_m x^m + a_{m-1} x^{m-1} + \cdots + a_1 x + a_0, \quad (3.20)$$

where  $x$  is called the **indeterminate**, each  $a_j$  is called a **coefficient**. If  $R$  is the type of the coefficients, then the set of all **univariate polynomials with coefficients in  $R$**  is written as  $R[x]$ . We often simply use **polynomial** instead of univariate polynomial, write  $P(x) \in R[x]$  for a polynomial and denote the constant term as  $P(0)$ .

A polynomial is called the **zero polynomial** if all coefficients are zero and a polynomial is called the **one polynomial** if the constant term is 1 and all other coefficients are zero.

If an univariate polynomial  $P(x) = \sum_{j=0}^m a_j x^j$  is given that is not the zero polynomial, we call  $\deg(P) := m$  the *degree* of  $P$  and define the degree of the zero polynomial to be  $-\infty$ , where  $-\infty$  (negative infinity) is a symbol with the property that  $-\infty + m = -\infty$  for all counting numbers  $m \in \mathbb{N}$ . In addition, we write

$$Lc(P) := a_m \quad (3.21)$$

1387 and call it the **leading coefficient** of the polynomial  $P$ . We can restrict the set  $R[x]$  of **all**  
 1388 polynomials with coefficients in  $R$ , to the set of all such polynomials that have a degree that  
 1389 does not exceed a certain value. If  $m$  is the maximum degree allowed, we write  $R_{\leq m}[x]$  for the  
 1390 set of all polynomials with a degree less than or equal to  $m$ .

*Example 15* (Integer Polynomials). The coefficients of a polynomial must all have the same type. The set of polynomials with integer coefficients is written as  $\mathbb{Z}[x]$ . Examples of such polynomials are:

$$\begin{array}{ll}
 P_1(x) = 2x^2 - 4x + 17 & \# \text{ with } \deg(P_1) = 2 \text{ and } Lc(P_1) = 2 \\
 P_2(x) = x^{23} & \# \text{ with } \deg(P_2) = 23 \text{ and } Lc(P_2) = 1 \\
 P_3(x) = x & \# \text{ with } \deg(P_3) = 1 \text{ and } Lc(P_3) = 1 \\
 P_4(x) = 174 & \# \text{ with } \deg(P_4) = 0 \text{ and } Lc(P_4) = 174 \\
 P_5(x) = 1 & \# \text{ with } \deg(P_5) = 0 \text{ and } Lc(P_5) = 1 \\
 P_6(x) = 0 & \# \text{ with } \deg(P_6) = -\infty \text{ and } Lc(P_6) = 0 \\
 P_7(x) = (x-2)(x+3)(x-5) & 
 \end{array}$$

In particular, every integer can be seen as an integer polynomial of degree zero.  $P_7$  is a polynomial, because we can expand its definition into  $P_7(x) = x^3 - 4x^2 - 11x + 30$ , which is polynomial of degree 3 and leading coefficient 1. The following expressions are not integer polynomials:

$$\begin{array}{l}
 Q_1(x) = 2x^2 + 4 + 3x^{-2} \\
 Q_2(x) = 0.5x^4 - 2x \\
 Q_3(x) = 1/x
 \end{array}$$

1391 We can invoke Sage to do computations with polynomials. To do so, we have to specify the  
 1392 symbol for the indeterminate and the type for the coefficients. Note, however that Sage defines  
 1393 the degree of the zero polynomial to be  $-1$ .

```

1394 sage: Zx = ZZ['x'] # integer polynomials with indeterminate x 68
1395 sage: Zt.<t> = ZZ[] # integer polynomials with indeterminate t 69
1396 sage: Zx 70
1397 Univariate Polynomial Ring in x over Integer Ring 71
1398 sage: Zt 72
1399 Univariate Polynomial Ring in t over Integer Ring 73
1400 sage: p1 = Zx([17,-4,2]) 74
1401 sage: p1 75
1402 2*x^2 - 4*x + 17 76
1403 sage: p1.degree() 77
1404 2 78
1405 sage: p1.leading_coefficient() 79
1406 2 80
1407 sage: p2 = Zt(t^23) 81
1408 sage: p2 82
1409 t^23 83
1410 sage: p6 = Zx([0]) 84
1411 sage: p6.degree() 85
1412 -1 86

```

*Example 16* (Polynomials over  $\mathbb{Z}_6$ ). Recall our definition of the residue classes  $\mathbb{Z}_6$  and their arithmetics as defined in Example 8. The set of all polynomials with indeterminate  $x$  and coefficients in  $\mathbb{Z}_6$  is symbolized as  $\mathbb{Z}_6[x]$ . Example of polynomials from  $\mathbb{Z}_6$  are:

$$\begin{array}{ll}
 P_1(x) = 2x^2 - 4x + 5 & \# \text{ with } \deg(P_1) = 2 \text{ and } Lc(P_1) = 2 \\
 P_2(x) = x^{23} & \# \text{ with } \deg(P_2) = 23 \text{ and } Lc(P_2) = 1 \\
 P_3(x) = x & \# \text{ with } \deg(P_3) = 1 \text{ and } Lc(P_3) = 1 \\
 P_4(x) = 3 & \# \text{ with } \deg(P_4) = 0 \text{ and } Lc(P_4) = 3 \\
 P_5(x) = 1 & \# \text{ with } \deg(P_5) = 0 \text{ and } Lc(P_5) = 1 \\
 P_6(x) = 0 & \# \text{ with } \deg(P_6) = -\infty \text{ and } Lc(P_6) = 0 \\
 P_7(x) = (x-2)(x+3)(x-5) & 
 \end{array}$$

Just like in the previous example,  $P_7$  is a polynomial. However, since we are working with coefficients from  $\mathbb{Z}_6$  now the expansion of  $P_7$  is computed differently, as we have to invoke addition and multiplication in  $\mathbb{Z}_6$  as defined in XXX. We get the following:

$$\begin{aligned}
 (x-2)(x+3)(x-5) &= (x+4)(x+3)(x+1) && \# \text{ additive inverses in } \mathbb{Z}_6 \\
 &= (x^2 + 4x + 3x + 3 \cdot 4)(x+1) && \# \text{ bracket expansion} \\
 &= (x^2 + 1x + 0)(x+1) && \# \text{ computation in } \mathbb{Z}_6 \\
 &= (x^3 + x^2 + x^2 + x) && \# \text{ bracket expansion} \\
 &= (x^3 + 2x^2 + x)
 \end{aligned}$$

1413 Again, we can use Sage to do computations with polynomials that have their coefficients in  $\mathbb{Z}_6$ .  
 1414 To do so, we have to specify the symbol for the indeterminate and the type for the coefficients:  
 1415

```

1416 sage: Z6 = Integers(6)                                     87
1417 sage: Z6x = Z6['x']                                         88
1418 sage: Z6x                                                    89
1419 Univariate Polynomial Ring in x over Ring of integers modulo 6 90
1420 sage: p1 = Z6x([5, -4, 2])                                   91
1421 sage: p1                                                     92
1422 2*x^2 + 2*x + 5                                             93
1423 sage: p1 = Z6x([17, -4, 2])                                  94
1424 sage: p1                                                     95
1425 2*x^2 + 2*x + 5                                             96
1426 sage: Z6x(x-2)*Z6x(x+3)*Z6x(x-5) == Z6x(x^3 + 2*x^2 + x) 97
1427 True                                                         98

```

1428 Given some element from the same type as the coefficients of a polynomial, the polynomial can be evaluated at that element, which means that we insert the given element for every  
 1429 occurrence of the indeterminate  $x$  in the polynomial expression.  
 1430

1431 To be more precise, let  $P \in R[x]$ , with  $P(x) = \sum_{j=0}^m a_j x^j$  be a polynomial with a coefficient  
 1432 of type  $R$  and let  $b \in R$  be an element of that type. Then the **evaluation** of  $P$  at  $b$  is given as  
 1433 follows:

$$P(b) = \sum_{j=0}^m a_j b^j \quad (3.22)$$

*Example 17.* Consider the integer polynomials from example XXX again. To evaluate them at given points, we have to insert the point for all occurrences of  $x$  in the polynomial expression. Inserting arbitrary values from  $\mathbb{Z}$ , we get:

$$\begin{aligned} P_1(2) &= 2 \cdot 2^2 - 4 \cdot 2 + 17 = 17 \\ P_2(3) &= 3^{23} = 94143178827 \\ P_3(-4) &= -4 = -4 \\ P_4(15) &= 174 \\ P_5(0) &= 1 \\ P_6(1274) &= 0 \\ P_7(-6) &= (-6-2)(-6+3)(-6+5) = -264 \end{aligned}$$

1434 Note, however, that it is not possible to evaluate any of those polynomial on values of different  
1435 type. For example, it is not strictly correct to write  $P_1(0.5)$ , since 0.5 is not an integer. We can  
1436 verify our computations using Sage:

```
1437 sage: Zx = ZZ['x'] 99
1438 sage: p1 = Zx([17, -4, 2]) 100
1439 sage: p7 = Zx(x-2)*Zx(x+3)*Zx(x-5) 101
1440 sage: p1(ZZ(2)) 102
1441 17 103
1442 sage: p7(ZZ(-6)) == ZZ(-264) 104
1443 True 105
```

*Example 18.* Consider the polynomials with coefficients in  $\mathbb{Z}_6$  from example XXX again. To evaluate them at given values from  $\mathbb{Z}_6$ , we have to insert the point for all occurrences of  $x$  in the polynomial expression. We get the following:

$$\begin{aligned} P_1(2) &= 2 \cdot 2^2 - 4 \cdot 2 + 5 = 2 - 2 + 5 = 5 \\ P_2(3) &= 3^{23} = 3 \\ P_3(-4) &= P_3(2) = 2 \\ P_5(0) &= 1 \\ P_6(4) &= 0 \end{aligned}$$

```
1444
1445 sage: Z6 = Integers(6) 106
1446 sage: Z6x = Z6['x'] 107
1447 sage: p1 = Z6x([5, -4, 2]) 108
1448 sage: p1(Z6(2)) == Z6(5) 109
1449 True 110
```

1450 *Exercise 21.* Compare both expansions of  $P_7$  from  $\mathbb{Z}[x]$  and from  $\mathbb{Z}_6[x]$  in example XXX and  
1451 example XXX, and consider the definition of  $\mathbb{Z}_6$  as given in example XXX. Can you see how  
1452 the definition of  $P_7$  over  $\mathbb{Z}$  projects to the definition over  $\mathbb{Z}_6$  if you consider the residue classes  
1453 of  $\mathbb{Z}_6$ ?

**Polynomial Arithmetics** Polynomials behave like integers in many ways. In particular, they can be added, subtracted and multiplied. In addition, they have their own notion of Euclidean division. Informally speaking, we can add two polynomials by simply adding the coefficients of the same index, and we can multiply them by applying the distributive property, that is, by multiplying every term of the left factor with every term of the right factor and adding the results together.

To be more precise let  $\sum_{n=0}^{m_1} a_n x^n$  and  $\sum_{n=0}^{m_2} b_n x^n$  be two polynomials from  $R[x]$ . Then the **sum** and the **product** of these polynomials is defined as follows:

$$\sum_{n=0}^{m_1} a_n x^n + \sum_{n=0}^{m_2} b_n x^n = \sum_{n=0}^{\max\{m_1, m_2\}} (a_n + b_n) x^n \quad (3.23)$$

$$\left( \sum_{n=0}^{m_1} a_n x^n \right) \cdot \left( \sum_{n=0}^{m_2} b_n x^n \right) = \sum_{n=0}^{m_1+m_2} \sum_{i=0}^n a_i b_{n-i} x^n \quad (3.24)$$

A rule for polynomial subtraction can be deduced from these two rules by first multiplying the **subtrahend** with (the polynomial)  $-1$  and then add the result to the **minuend**.

Regarding the definition of the degree of a polynomial, we see that the degree of the sum is always the maximum of the degrees of both summands, and the degree of the product is always the degree of the factors, since we defined  $-\infty \cdot m = \infty$  for every integer  $m \in \mathbb{Z}$ . Using Sage's definition of degree, this would not hold, as the zero polynomials degree is  $-1$  in Sage, which would violate this rule.

*Example 19.* To give an example of how polynomial arithmetics works, consider the following two integer polynomials  $P, Q \in \mathbb{Z}[x]$  with  $P(x) = 5x^2 - 4x + 2$  and  $Q(x) = x^3 - 2x^2 + 5$ . The sum of these two polynomials is computed by adding the coefficients of each term with equal exponent in  $x$ . This gives the following:

$$\begin{aligned} (P + Q)(x) &= (0 + 1)x^3 + (5 - 2)x^2 + (-4 + 0)x + (2 + 5) \\ &= x^3 + 3x^2 - 4x + 7 \end{aligned}$$

The product of these two polynomials is computed by multiplication of each term in the first factor with each term in the second factor. We get the following:

$$\begin{aligned} (P \cdot Q)(x) &= (5x^2 - 4x + 2) \cdot (x^3 - 2x^2 + 5) \\ &= (5x^5 - 10x^4 + 25x^2) + (-4x^4 + 8x^3 - 20x) + (2x^3 - 4x^2 + 10) \\ &= 5x^5 - 14x^4 + 10x^3 + 21x^2 - 20x + 10 \end{aligned}$$

```

1471 sage: Zx = ZZ['x']
1472 sage: P = Zx([2, -4, 5])
1473 sage: Q = Zx([5, 0, -2, 1])
1474 sage: P+Q == Zx(x^3 + 3*x^2 - 4*x + 7)
1475 True
1476 sage: P*Q == Zx(5*x^5 - 14*x^4 + 10*x^3 + 21*x^2 - 20*x + 10)
1477 True

```

subtrahend

minuend

*Example 20.* Let us consider the polynomials of the previous example but interpreted in modular 6 arithmetics. So we consider  $P, Q \in \mathbb{Z}_6[x]$  again with  $P(x) = 5x^2 - 4x + 2$  and  $Q(x) = x^3 - 2x^2 + 5$ . This time we get the following:

$$\begin{aligned}(P + Q)(x) &= (0 + 1)x^3 + (5 - 2)x^2 + (-4 + 0)x + (2 + 5) \\ &= (0 + 1)x^3 + (5 + 4)x^2 + (2 + 0)x + (2 + 5) \\ &= x^3 + 3x^2 + 2x + 1\end{aligned}$$

$$\begin{aligned}(P \cdot Q)(x) &= (5x^2 - 4x + 2) \cdot (x^3 - 2x^2 + 5) \\ &= (5x^2 + 2x + 2) \cdot (x^3 + 4x^2 + 5) \\ &= (5x^5 + 2x^4 + 1x^2) + (2x^4 + 2x^3 + 4x) + (2x^3 + 2x^2 + 4) \\ &= 5x^5 + 4x^4 + 4x^3 + 3x^2 + 4x + 4\end{aligned}$$

1478

```

1479 sage: Z6x = Integers(6) ['x']           118
1480 sage: P = Z6x([2, -4, 5])                119
1481 sage: Q = Z6x([5, 0, -2, 1])             120
1482 sage: P+Q == Z6x(x^3 +3*x^2 +2*x +1)    121
1483 True                                     122
1484 sage: P*Q == Z6x(5*x^5 +4*x^4 +4*x^3+3*x^2+4*x +4) 123
1485 True                                     124

```

1486 *Exercise 22.* Compare the sum  $P + Q$  and the product  $P \cdot Q$  from the previous two examples  
 1487 XXX and XXX and consider the definition of  $\mathbb{Z}_6$  as given in example XXX. How can we derive  
 1488 the computations in  $\mathbb{Z}_6[x]$  from the computations in  $\mathbb{Z}[x]$ ?

1489 **Euklidean Division** The ring of polynomials shares a lot of properties with integers. In par-  
 1490 ticular, the concept of Euclidean division and the algorithm of long division is also defined for  
 1491 polynomials. Recalling the Euclidean division of integers XXX, we know that, given two inte-  
 1492 gers  $a$  and  $b \neq 0$ , there is always another integer  $m$  and a counting number  $r$  with  $r < |b|$  such  
 1493 that  $a = m \cdot b + r$  holds.

1494 We can generalize this to polynomials whenever the leading coefficient of the dividend  
 1495 polynomial has a notion of multiplicative inverse. In fact, given two polynomials  $A$  and  $B \neq 0$   
 1496 from  $R[x]$  such that  $Lc(B)^{-1}$  exists in  $R$ , there exist two polynomials  $M$  (the quotient) and  $R$  (the  
 1497 remainder), such that the following equation holds:

$$A = M \cdot B + R \quad (3.25)$$

1498 and  $\deg(R) < \deg(B)$ . Similarly to integer Euclidean division, both  $M$  and  $R$  are uniquely  
 1499 defined by these relations.

1500 *Notation and Symbols 2.* Suppose that the polynomials  $A, B, M$  and  $R$  satisfy equation XX. We  
 1501 often use the following notation to describe the quotient and the remainder polynomials of the  
 1502 Euclidean division:

$$A \operatorname{div} B := M, \quad A \operatorname{mod} B := R \quad (3.26)$$

1503 We also say that a polynomial  $A$  is divisible by another polynomial  $B$  if  $A \operatorname{mod} B = 0$  holds. In  
 1504 this case, we also write  $B|A$  and call  $B$  a *factor* of  $A$ .

1505

1506

---

**Algorithm 3** Polynomial Euclidean Algorithm

**Require:**  $A, B \in R[x]$  with  $B \neq 0$ , such that  $Lc(B)^{-1}$  exists in  $R$

**procedure** POLY-LONG-DIVISION( $A, B$ )

$$M \leftarrow 0$$
$$R \leftarrow A$$
$$d \leftarrow \deg(B)$$
$$c \leftarrow Lc(B)$$

```

while  $\deg(R) \geq d$  do

```

$$S := L_C(R) \cdot c^{-1} \cdot x^{\deg(R)-d}$$
$$M \leftarrow M + S$$
$$R \leftarrow R - S \cdot B$$

end while

```

return (Q,R)

```

end procedure

**Ensure:**  $A = M \cdot B + R$

1507

1508

1510  
1511  
1512  
1513  
1514  
1515  
1516

[illegible]

1517  
1518  
1519

1520

1521

125

126



```

1522 sage: B = Zx([-1, 4, 1]) 127
1523 sage: M = Zx([-80, 19, -4, 1]) 128
1524 sage: R = Zx([-89, 339]) 129
1525 sage: A == M*B + R 130
1526 True 131

```

1527 *Example 22.* In the previous example, polynomial division gave a non-trivial (non-vanishing,  
 1528 i.e non-zero) remainder. Of special interest are divisions that don't give a remainder. Such  
 1529 divisors are called factors of the dividend.

1530 For example, consider the integer polynomial  $P_7$  from example XXX again. As we have  
 1531 shown, it can be written both as  $x^3 - 4x^2 - 11x + 30$  and as  $(x - 2)(x + 3)(x - 5)$ . From this,  
 1532 we can see that the polynomials  $F_1(x) = (x - 2)$ ,  $F_2(x) = (x + 3)$  and  $F_3(x) = (x - 5)$  are all  
 1533 factors of  $x^3 - 4x^2 - 11x + 30$ , since division of  $P_7$  by any of these factors will result in a zero  
 1534 remainder.

1535 *Exercise 23.* Consider the polynomial expressions  $P(x) := -3x^4 + 4x^3 + 2x^2 + 4$  and  $Q(x) =$   
 1536  $x^2 - 4x + 2$ . Compute the Euclidean division of  $P$  by  $Q$  in the following types:

- 1537 1.  $P, Q \in \mathbb{Z}[x]$
- 1538 2.  $P, Q \in \mathbb{Z}_6[x]$
- 1539 3.  $P, Q \in \mathbb{Z}_5[x]$

1540 Now consider the result in  $\mathbb{Z}[x]$  and in  $\mathbb{Z}_6[x]$ . How can we compute the result in  $\mathbb{Z}_6[x]$  from the  
 1541 result in  $\mathbb{Z}[x]$ ?

1542 *Exercise 24.* Show that the polynomial  $P(x) = 2x^4 - 3x + 4 \in \mathbb{Z}_5[x]$  is a factor of the polynomial  
 1543  $Q(x) = x^7 + 4x^6 + 4x^5 + x^3 + 2x^2 + 2x + 3 \in \mathbb{Z}_5[x]$  that is show  $P|Q$ . What is  $Q \div P$ ?

1544 **Prime Factors** Recall that the fundamental theorem of arithmetics XXX tells us that every  
 1545 number is the product of prime numbers. Something similar holds for polynomials, too.

1546 The polynomial analog to a prime number is a so called an **irreducible polynomial**, which  
 1547 is defined as a polynomial that cannot be factored into the product of two non-constant poly-  
 1548 nomials using Euclidean division. Irreducible polynomials are for polynomials what prime  
 1549 numbers are for integer: They are the basic building blocks from which all other polynomials  
 1550 can be constructed. To be more precise, let  $P \in R[x]$  be any polynomial. Then there are always  
 1551 irreducible polynomials  $F_1, F_2, \dots, F_k \in R[x]$ , such that the following holds:

$$P = F_1 \cdot F_2 \cdot \dots \cdot F_k. \quad (3.28)$$

1552 This representation is unique, except for permutations in the factors and is called the **prime**  
 1553 **factorization** of  $P$ .

1554 *Example 23.* Consider the polynomial expression  $P = x^2 - 3$ . When we interpret  $P$  as an integer  
 1555 polynomial  $P \in \mathbb{Z}[x]$ , we find that this polynomial is irreducible, since any factorization other  
 1556 than  $1 \cdot (x^2 - 3)$ , must look like  $(x - a)(x + a)$  for some integer  $a$ , but there is no integers  $a$  with  
 1557  $a^2 = 3$ .

```

1558 sage: Zx = ZZ['x'] 132
1559 sage: p = Zx(x^2-3) 133
1560 sage: p.roots() 134
1561 [] 135

```

```

1562 sage: p.factor() 136
1563 x^2 - 3 137

```

1564 On the other hand interpreting  $P$  as a polynomial  $P \in \mathbb{Z}_6[x]$  in modulo 6 arithmetics, we see that  
 1565  $P$  has two factors  $F_1 = (x-3)$  and  $F_2 = (x+3)$ , since  $(x-3)(x+3) = x^2 - 3x + 3 - 3 \cdot 3 = x^2 - 3$ .

1566 Finding prime factors of a polynomial is hard. As we have seen in example XXX, points  
 1567 where a polynomial evaluates to zero, i.e points  $x_0 \in R$  with  $P(x_0) = 0$  are of special interest,  
 1568 since it can be shown the polynomial  $F(x) = (x - x_0)$  is always a factor of  $P$ . The converse,  
 1569 however, is not necessarily true, because a polynomial can have irreducible prime factors.

1570 Points where a polynomial evaluates to zero are called the **roots** of the polynomial. To be  
 1571 more precise, let  $P \in R[x]$  be a polynomial. Then the set of all roots of  $P$  is defined as follows:

$$R_0(P) := \{x_0 \in R \mid P(x_0) = 0\} \quad (3.29)$$

1572 Finding the roots of a polynomial is sometimes called **solving the polynomial**. It is a hard  
 1573 problem and has been the subject of much research throughout history. In fact, it is well known  
 1574 that, for polynomials of degree 5 or higher, there is, in general, no closed expression, from  
 1575 which the roots can be deduced.

1576 It can be shown that if  $m$  is the degree of a polynomial  $P$ , then  $P$  can not have more than  $m$   
 1577 roots. However, in general, polynomials can have less than  $m$  roots.

1578 *Example 24.* Consider our integer polynomial  $P_7(x) = x^3 - 4x^2 - 11x + 30$  from example XXX  
 1579 again. We know that its set of roots is given by  $R_0(P_7) = \{-3, 2, 5\}$ .

1580 On the other hand, we know from example XXX that the integer polynomial  $x^2 - 3$  is irre-  
 1581 ducible. It follows that it has no roots, since every root defines a prime factor.

1582 *Example 25.* To give another example, consider the integer polynomial  $P = x^7 + 3x^6 + 3x^5 +$   
 1583  $x^4 - x^3 - 3x^2 - 3x - 1$ . We can invoke Sage to compute the roots and prime factors of  $P$ :

```

1584 sage: Zx = ZZ['x'] 138
1585 sage: p = Zx(x^7 + 3*x^6 + 3*x^5 + x^4 - x^3 - 3*x^2 - 3*x - 1) 139
1586 )
1587 sage: p.roots() 140
1588 [(1, 1), (-1, 4)] 141
1589 sage: p.factor() 142
1590 (x - 1) * (x + 1)^4 * (x^2 + 1) 143

```

We see that  $P$  has the root 1 and that the associated prime factor  $(x-1)$  occurs once in  $P$  and that it has the root  $-1$ , where the associated prime factor  $(x+1)$  occurs 4 times in  $P$ . This gives the prime following factorization:

$$P = (x-1)(x+1)^4(x^2+1)$$

1591 **Lange interpolation** One particularly useful property of polynomials is that a polynomial of  
 1592 degree  $m$  is completely determined on  $m+1$  evaluation points. Seeing this from a different  
 1593 angle, we can (sometimes) uniquely derive a polynomial of degree  $m$  from a set  $S$ :

$$S = \{(x_0, y_0), (x_1, y_1), \dots, (x_m, y_m) \mid x_i \neq x_j \text{ for all indices } i \text{ and } j\} \quad (3.30)$$

1594 This “few too many” property of polynomials is used in many places, like for example in erasure  
 1595 codes. It is also of importance in snarks and we therefore need to understand a method to  
 1596 actually compute a polynomial from a set of points.

what  
does this  
mean?

1597 If the coefficients of the polynomial we want to find have a notion of multiplicative inverse,  
 1598 it is always possible to find such a polynomial. One method for this is called **Lagrange in-**  
 1599 **terpolation**. It works as follows: Given a set like 3.30, a polynomial  $P$  of degree  $m + 1$  with  
 1600  $P(x_i) = y_i$  for all pairs  $(x_i, y_i)$  from  $S$  is given by the following algorithm:

---

**Algorithm 4** Lagrange Interpolation
 

---

**Require:**  $R$  must have multiplicative inverses

**Require:**  $S = \{(x_0, y_0), (x_1, y_1), \dots, (x_m, y_m) \mid x_i, y_i \in R, x_i \neq x_j \text{ for all indices } i \text{ and } j\}$

**procedure** LAGRANGE-INTERPOLATION( $S$ )

**for**  $j \in (0 \dots m)$  **do**

$$l_j(x) \leftarrow \prod_{i=0; i \neq j}^m \frac{x - x_i}{x_j - x_i} = \frac{(x - x_0)}{(x_j - x_0)} \cdots \frac{(x - x_{j-1})}{(x_j - x_{j-1})} \frac{(x - x_{j+1})}{(x_j - x_{j+1})} \cdots \frac{(x - x_m)}{(x_j - x_m)}$$

**end for**

$$P \leftarrow \sum_{j=0}^m y_j \cdot l_j$$

**return**  $P$

**end procedure**

**Ensure:**  $P \in R[x]$  with  $\deg(P) = m$

**Ensure:**  $P(x_j) = y_j$  for all pairs  $(x_j, y_j) \in S$

---

*Example 26.* Let us consider the set  $S = \{(0, 4), (-2, 1), (2, 3)\}$ . Our task is to compute a polynomial of degree 2 in  $\mathbb{Q}[x]$  with fractional number coefficients. Since  $\mathbb{Q}$  has multiplicative inverses, we can use the Lagrange interpolation algorithm from 4, to compute the polynomial.

$$\begin{aligned} l_0(x) &= \frac{x - x_1}{x_0 - x_1} \cdot \frac{x - x_2}{x_0 - x_2} = \frac{x + 2}{0 + 2} \cdot \frac{x - 2}{0 - 2} = -\frac{(x + 2)(x - 2)}{4} \\ &= -\frac{1}{4}(x^2 - 4) \\ l_1(x) &= \frac{x - x_0}{x_1 - x_0} \cdot \frac{x - x_2}{x_1 - x_2} = \frac{x - 0}{-2 - 0} \cdot \frac{x - 2}{-2 - 2} = \frac{x(x - 2)}{8} \\ &= \frac{1}{8}(x^2 - 2x) \\ l_2(x) &= \frac{x - x_0}{x_2 - x_0} \cdot \frac{x - x_1}{x_2 - x_1} = \frac{x - 0}{2 - 0} \cdot \frac{x + 2}{2 + 2} = \frac{x(x + 2)}{8} \\ &= \frac{1}{8}(x^2 + 2x) \\ P(x) &= 4 \cdot \left(-\frac{1}{4}(x^2 - 4)\right) + 1 \cdot \frac{1}{8}(x^2 - 2x) + 3 \cdot \frac{1}{8}(x^2 + 2x) \\ &= -x^2 + 4 + \frac{1}{8}x^2 - \frac{1}{4}x + \frac{3}{8}x^2 + \frac{3}{4}x \\ &= -\frac{1}{2}x^2 + \frac{1}{2}x + 4 \end{aligned}$$

1601 And, indeed, evaluation of  $P$  on the  $x$ -values of  $S$  gives the correct points, since  $P(0) = 4$ ,  
 1602  $P(-2) = 1$  and  $P(2) = 3$ .

*Example 27.* To give another example more relevant to the topics of this book, let us consider the same set  $S = \{(0, 4), (-2, 1), (2, 3)\}$  as in the previous example. This time, the task is to compute a polynomial  $P \in \mathbb{Z}_5[x]$  from this data. Since we know that multiplicative inverses exist in  $\mathbb{Z}_5$ , algorithm XXX applies and we can compute a unique polynomial of degree 2 in

$\mathbb{Z}_5[x]$  from  $S$ . We can use the lookup tables XXX for computation in  $\mathbb{Z}_5$  and get the following:

$$\begin{aligned} l_0(x) &= \frac{x-x_1}{x_0-x_1} \cdot \frac{x-x_2}{x_0-x_2} = \frac{x+2}{0+2} \cdot \frac{x-2}{0-2} = \frac{(x+2)(x-2)}{-4} = \frac{(x+2)(x+3)}{1} \\ &= x^2 + 1 \end{aligned}$$

$$\begin{aligned} l_1(x) &= \frac{x-x_0}{x_1-x_0} \cdot \frac{x-x_2}{x_1-x_2} = \frac{x-0}{-2-0} \cdot \frac{x-2}{-2-2} = \frac{x}{3} \cdot \frac{x+3}{1} = 2(x^2 + 3x) \\ &= 2x^2 + x \end{aligned}$$

$$\begin{aligned} l_2(x) &= \frac{x-x_0}{x_2-x_0} \cdot \frac{x-x_1}{x_2-x_1} = \frac{x-0}{2-0} \cdot \frac{x+2}{2+2} = \frac{x(x+2)}{3} = 2(x^2 + 2x) \\ &= 2x^2 + 4x \end{aligned}$$

$$\begin{aligned} P(x) &= 4 \cdot (x^2 + 1) + 1 \cdot (2x^2 + x) + 3 \cdot (2x^2 + 4x) \\ &= 4x^2 + 4 + 2x^2 + x + x^2 + 2x \\ &= 2x^2 + 3x + 4 \end{aligned}$$

1603 And, indeed, evaluation of  $P$  on the  $x$ -values of  $S$  gives the correct points, since  $P(0) = 4$ ,  
1604  $P(-2) = 1$  and  $P(2) = 3$ .

1605 *Exercise 25.* Consider example XXX and example XXX again. Why is it not possible to apply  
1606 algorithm XXX if we consider  $S$  as a set of integers, nor as a set in  $\mathbb{Z}_6$ ?

# Chapter 4

## Algebra

In the previous chapter, we gave an introduction to the basic computational skills needed for a pen-and-paper approach to SNARKs. This chapter provides a more abstract clarification of relevant mathematical terminology on **algebraic types** such as **groups**, **fields**, **rings** and similar.

In a nutshell, algebraic types define sets that are analogous to numbers in various aspects, in the sense that you can add, subtract, multiply or divide on those sets. We know many examples of sets that fall under those categories, such as natural numbers, integers, rational or the real numbers. In some sense, these are the most fundamental examples of such sets.

Papers on cryptography (and mathematical papers in general) frequently contain such terms, and it is necessary to get at least some understanding of these terms to be able to follow these papers. In this chapter, we therefore provide a short introduction to these concepts.

Def Sub-group, Fundamental theorem of cyclic groups.

### 4.1 Groups

Groups are abstractions that capture the essence of mathematical phenomena, like addition and subtraction, multiplication and division, permutations, or symmetries.

To understand groups, let us think back to when we learned about the addition and subtraction of integers (also called whole numbers) in school. We have learned that, whenever we add two integers, the result is guaranteed to be an integer as well. We have also learned that adding zero to any integer means that “nothing happens”, that is, the result of the addition is the same integer we started with. Furthermore, we have learned that the order in which we add two (or more) integers does not matter, that brackets have no influence on the result of addition, and that, for every integer, there is always another integer (the negative) such that we get zero when we add them together.

These conditions are the defining properties of a group, and mathematicians have recognized that the exact same set of rules can be found in very different mathematical structures. It therefore makes sense to give a formal definition of what a group should be, detached from any concrete examples. This lets us handle entities of very different mathematical origins in a flexible way, while retaining essential structural aspects of many objects in abstract algebra and beyond.

Distilling these rules to the smallest independent list of properties and making them abstract, we arrive at the definition of a group:

*Definition 4.1.0.1.* A **group**  $(\mathbb{G}, \cdot)$  is a set  $\mathbb{G}$ , together with a **map**  $\cdot$ . The map, also denoted as  $\mathbb{G} \times \mathbb{G} \rightarrow \mathbb{G}$  and called the **group law**, combines two elements of the set  $\mathbb{G}$  into a third one such that the following properties hold:

- **Existence of a neutral element:** There is a  $e \in \mathbb{G}$  for all  $g \in \mathbb{G}$ , such that  $e \cdot g = g$  as well as  $g \cdot e = g$ .
- **Existence of an inverse:** For every  $g \in \mathbb{G}$  there is a  $g^{-1} \in \mathbb{G}$ , such that  $g \cdot g^{-1} = e$  as well as  $g^{-1} \cdot g = e$ .
- **Associativity:** For every  $g_1, g_2, g_3 \in \mathbb{G}$  the equation  $g_1 \cdot (g_2 \cdot g_3) = (g_1 \cdot g_2) \cdot g_3$  holds.

Rephrasing the abstract definition in layman's terms, a group is something where we can do computations in a way that resembles the behavior of the addition of integers. Specifically, this means we can combine some element with another element into a new element in a way that is reversible and where the order of combining elements doesn't matter.

*Notation and Symbols 3.* Let  $(\mathbb{G}, \cdot)$  be a finite group. If there is no risk of ambiguity (about what the group law of that group is), we frequently drop the symbol  $\cdot$  and simply write  $\mathbb{G}$  as the notation for the group, keeping the group law implicit.

As we will see in XXX, groups are heavily used in cryptography and in SNARKs. But let us look at some more familiar examples first:

add reference

*Example 28* (Integer Addition and Subtraction). The set  $(\mathbb{Z}, +)$  of integers with integer addition is the archetypical example of a group, where the group law is traditionally written as  $+$  (instead of  $\cdot$ ). To compare integer addition against the abstract axioms of a group, we first see that the neutral element  $e$  is the number 0, since  $a + 0 = a$  for all integers  $a \in \mathbb{Z}$ . Furthermore, the inverse of a number is its negative counterpart, since  $a + (-a) = 0$ , for all  $a \in \mathbb{Z}$ . In addition, we know that  $(a + b) + c = a + (b + c)$ , so integers with addition are indeed a group in the abstract sense.

*Example 29* (The trivial group). The most basic example of a group is group with just one element  $\{\bullet\}$  and the group law  $\bullet \cdot \bullet = \bullet$ .

Add real-life example of 0?

**Commutative Groups** When we look at the general definition of a group, we see that it is somewhat different from what we know from integers. We know that the order in which we add two integers doesn't matter, as, for example,  $4 + 2$  is the same as  $2 + 4$ . However, we also know from example XXX that this is not the case for all groups.

add reference

This means that groups where the order in which the group law is executed doesn't matter are a special subcase of groups called **commutative groups**. To be more precise, a group is called commutative if  $g_1 \cdot g_2 = g_2 \cdot g_1$  holds for all  $g_1, g_2 \in \mathbb{G}$ .

*Notation and Symbols 4.* For commutative groups  $(\mathbb{G}, \cdot)$ , we frequently use the so-called **additive notation**  $(\mathbb{G}, +)$ , that is, we write  $+$  instead of  $\cdot$  for the group law, and  $-g := g^{-1}$  for the inverse of an element  $g \in \mathbb{G}$ .

*Example 30.* Consider the group of integers with integer addition again. Since  $a + b = b + a$  for all integers, this group is the archetypical example of a commutative group. Since there are infinitely many integers,  $(\mathbb{Z}, +)$  is not a finite group.

*Example 31.* Consider our definition of modulo 6 residue classes  $(\mathbb{Z}_6, +)$  as defined in the addition table from example 8. As we can see, the residue class 0 is the neutral element in modulo 6 arithmetics, and the inverse of a residue class  $r$  is given by  $6 - r$ , since  $r + (6 - r) = 6$ , which is congruent to 0, since  $6 \bmod 6 = 0$ . Moreover,  $(r_1 + r_2) + r_3 = r_1 + (r_2 + r_3)$  is inherited from integer arithmetic.

We therefore see that  $(\mathbb{Z}_6, +)$  is a group, and, since the addition table in example 8 is symmetrical, we see  $r_1 + r_2 = r_2 + r_1$ , which shows that  $(\mathbb{Z}_6, +)$  is commutative.

The previous example of a commutative group is a very important one for this book. Abstracting from this example and considering residue classes  $(\mathbb{Z}_n, +)$  for arbitrary moduli  $n$ , it can be shown that  $(\mathbb{Z}, +)$  is a commutative group with the neutral element 0 and the additive inverse  $n - r$  for any element  $r \in \mathbb{Z}_n$ . We call such a group the **remainder class group** of modulus  $n$ .

Of particular importance for pairing-based cryptography in general and SNARKs in particular are so-called **pairing maps** on commutative groups. To be more precise let  $\mathbb{G}_1$ ,  $\mathbb{G}_2$  and  $\mathbb{G}_3$  be three commutative groups. For historical reasons, we write the group law on  $\mathbb{G}_1$  and  $\mathbb{G}_2$  in additive notation and the group law on  $\mathbb{G}_3$  in multiplicative notation. Then a **pairing map** is a function

$$e(\cdot, \cdot) : \mathbb{G}_1 \times \mathbb{G}_2 \rightarrow \mathbb{G}_3 \quad (4.1)$$

This function takes pairs  $(g_1, g_2)$  (products) of elements from  $\mathbb{G}_1$  and  $\mathbb{G}_2$ , and maps them to elements from  $\mathbb{G}_3$ , such that the **bilinearity** property holds:

**Definition 4.1.0.2. Bilinearity**

For all  $g_1, g'_1 \in \mathbb{G}_1$  and  $g_2 \in \mathbb{G}_2$  we have  $e(g_1 + g'_1, g_2) = e(g_1, g_2) \cdot e(g'_1, g_2)$  and for all  $g_1 \in \mathbb{G}_1$  and  $g_2, g'_2 \in \mathbb{G}_2$  we have  $e(g_1, g_2 + g'_2) = e(g_1, g_2) \cdot e(g_1, g'_2)$ .

A pairing map is called **non-degenerate** if, whenever the result of the pairing is the neutral element in  $\mathbb{G}_3$ , one of the input values is the neutral element of  $\mathbb{G}_1$  or  $\mathbb{G}_2$ . To be more precise,  $e(g_1, g_2) = e_{\mathbb{G}_3}$  implies  $g_1 = e_{\mathbb{G}_1}$  or  $g_2 = e_{\mathbb{G}_2}$ .

Informally speaking, bilinearity means that it doesn't matter if we first execute the group law on one side and then apply the bilinear map, or if we first apply the bilinear map and then apply the group law. Moreover, non-degeneracy means that the result of the pairing is zero if and only if at least one of the input values is zero.

**Example 32.** One of the most basic examples of a non-degenerate pairing involves  $\mathbb{G}_1$ ,  $\mathbb{G}_2$  and  $\mathbb{G}_3$  all to be the groups of integers with addition  $(\mathbb{Z}, +)$ . Then the following map defines a non-degenerate pairing:

$$e(\cdot, \cdot) : \mathbb{Z} \times \mathbb{Z} \rightarrow \mathbb{Z} \quad (a, b) \mapsto a \cdot b$$

Note that bilinearity follows from the distributive law of integers, since for  $a, b, c \in \mathbb{Z}$ , we have  $e(a + b, c) = (a + b) \cdot c = a \cdot c + b \cdot c = e(a, c) + e(b, c)$  and the same reasoning is true for the second argument.

To see that  $e(\cdot, \cdot)$  is non-degenerate, assume that  $e(a, b) = 0$ . Then  $a \cdot b = 0$  implies that  $a$  or  $b$  must be zero.

**Exercise 26.** Consider example 13 again and let  $\mathbb{Z}_5^*$  be the set of all remainder classes from  $\mathbb{Z}_5$  without the class 0. Then  $\mathbb{Z}_5^* = \{1, 2, 3, 4\}$ . Show that  $(\mathbb{Z}_5^*, \cdot)$  is a commutative group.

check  
reference

**Exercise 27.** Generalizing the previous exercise, consider the general modulus  $n$ , and let  $\mathbb{Z}_n^*$  be the set of all remainder classes from  $\mathbb{Z}_n$  without the class 0. Then  $\mathbb{Z}_n^* = \{1, 2, \dots, n - 1\}$ . Provide a counter-example to show that  $(\mathbb{Z}_n^*, \cdot)$  is not a group in general.

Find a condition such that  $(\mathbb{Z}_n^*, \cdot)$  is a commutative group, compute the neutral element, give a closed form for the inverse of any element and prove the commutative group axioms.

**Exercise 28.** Consider the remainder class groups  $(\mathbb{Z}_n, +)$  for some modulus  $n$ . Show that the map

$$e(\cdot, \cdot) : \mathbb{Z}_n \times \mathbb{Z}_n \rightarrow \mathbb{Z}_n \quad (a, b) \mapsto a \cdot b$$

is bilinear. Why is it not a pairing in general and what condition must be imposed on  $n$ , such that the map will be a pairing?



**Finite groups** As we have seen in the previous examples, groups can either contain infinitely many elements (such as integers) or finitely many elements (as for example the remainder class groups  $(\mathbb{Z}_n, +)$ ). To capture this distinction, a group is called a **finite group** if the underlying set of elements is finite. In that case, the number of elements of that group is called its **order**.

*Notation and Symbols* 5. Let  $\mathbb{G}$  be a finite group. We write  $\text{ord}(\mathbb{G})$  or  $|\mathbb{G}|$  for the order of  $\mathbb{G}$ .

*Example 33.* Consider the remainder class groups  $(\mathbb{Z}_6, +)$  from example 8 and  $(\mathbb{Z}_5, +)$  from example 13, and the group  $(\mathbb{Z}_5^*, \cdot)$  from exercise 26. We can easily see that the order of  $(\mathbb{Z}_6, +)$  is 6, the order of  $(\mathbb{Z}_5, +)$  is 5 and the order of  $(\mathbb{Z}_5^*, \cdot)$  is 4.

To be more general, considering arbitrary moduli  $n$ , we know from Euclidean division that the order of the remainder class group  $(\mathbb{Z}_n, +)$  is  $n$ .

*Exercise 29.* The **RSA crypto system** is based on a modulus  $n$  that is typically the product of two prime numbers of **size 2048-bits**. What is the approximate order of the remainder class group  $(\mathbb{Z}_n, +)$  in this case?

check references to previous examples

RSA crypto system

size 2048-bits

**Generators** These are sets of elements that can be used to generate the entire group by applying the group law repeatedly to these elements or their inverses only. Generators are of particular interest when working with groups.

Of course, every group  $\mathbb{G}$  has a trivial set of generators, when we just consider every element of the group to be in the generator set. The more interesting question is to find the smallest possible set of generators for a given group. Of particular interest in this regard are groups that have a single generator, that is, there exists an element  $g \in \mathbb{G}$  such that every other element from  $\mathbb{G}$  can be computed by the repeated combination of  $g$  and its inverse  $g^{-1}$  only. Groups with a single generator are called **cyclic groups**.

*Example 34.* The most basic example of a cyclic group is the group of integers  $(\mathbb{Z}, +)$  with integer addition. 1 is a single generator of  $\mathbb{Z}$ , since every integer can be obtained by repeatedly adding either 1 or its inverse  $-1$  to itself. For example  $-4$  is generated by  $-1$ , since  $-4 = -1 + (-1) + (-1) + (-1)$ .

*Example 35.* Consider a modulus  $n$  and the remainder class groups  $(\mathbb{Z}_n, +)$  from example 33. These groups are cyclic, with a generator 1, since every other element of that group can be constructed by repeatedly adding the remainder class 1 to itself. Since  $\mathbb{Z}_n$  is also finite, we know that  $(\mathbb{Z}_n, +)$  is a finite cyclic group of order  $n$ .

*Example 36.* Let  $p \in \mathbb{P}$  be prime number and  $(\mathbb{F}_p^*, \cdot)$  the finite group from exercise XXX. Then  $(\mathbb{F}_p^*, \cdot)$  is cyclic and every element  $g \in \mathbb{F}_p^*$  is a generator.

check reference

add reference: 27?

**The discrete Logarithm problem** Observe that, when  $\mathbb{G}$  is a cyclic group of order  $n$  and  $g \in \mathbb{G}$  is a generator of  $\mathbb{G}$ , then there is a map with respect to the generator  $g$  with the following properties:

$$g^{(\cdot)} : \mathbb{Z}_n \rightarrow \mathbb{G} \quad x \mapsto g^x \quad (4.2)$$

In the map above,  $g^x$  means “multiply  $g$   $x$ -times by itself” and  $g^0 = e_{\mathbb{G}}$ . This map, called the **exponential map**, has the remarkable property that it maps the additive group law of the remainder class group  $(\mathbb{Z}_n, +)$  in a one-to-one correspondence to the group law of  $\mathbb{G}$ .

To see this, first observe that, since  $g^0 := e_{\mathbb{G}}$  by definition, the neutral element of  $\mathbb{Z}_n$  is mapped to the neutral element of  $\mathbb{G}$ , and, since  $g^{x+y} = g^x \cdot g^y$ , the map respects the group law.

Because the exponential map respects the group law, it doesn’t matter if we do our computation in  $\mathbb{Z}_n$  before we write the result into the exponent of  $g$  or afterwards: the result will be the



1762 same in both cases. This is usually referred to as doing computations “in the exponent”. In cryp-  
 1763 tography in general, and in SNARK development in particular, we often perform computations  
 1764 “in the exponent” of a generator.

*Example 37.* Consider the multiplicative group  $(\mathbb{F}_5^*, \cdot)$  from example 26. We know that  $\mathbb{F}_5^*$  is a cyclic group of order 4, and that every element is a generator. If we choose  $3 \in \mathbb{Z}_5^*$ , we then know that the following map respects the group law of addition in  $\mathbb{Z}_4$  and the group law of multiplication in  $\mathbb{Z}_5^*$ :

$$3^{(\cdot)} : \mathbb{Z}_4 \rightarrow \mathbb{Z}_5^* x \mapsto 3^x$$

Let us now perform a computation in the exponent:

$$\begin{aligned} 3^{2+3-2} &= 3^3 \\ &= 2 \end{aligned}$$

1765 This gives the same result as doing the same computation in  $\mathbb{F}_{*5}$ :

$$\begin{aligned} 3^{2+3-2} &= 3^2 \cdot 3^3 \cdot 3^{-2} \\ &= 4 \cdot 2 \cdot (-3)^2 \\ &= 3 \cdot 2^2 \\ &= 3 \cdot 4 \\ &= 2 \end{aligned}$$

1766 Since the exponential map is a one-to-one correspondence that respects the group law, it  
 1767 can be shown that this map has an inverse with respect to the base  $g$ , called the **base  $g$  discrete**  
 1768 **logarithm map**:

$$\log_g(\cdot) : \mathbb{G} \rightarrow \mathbb{Z}_n x \mapsto \log_g(x) \quad (4.3)$$

1769 Discrete logarithms are highly important in cryptography, because there are groups such that  
 1770 the exponential map and its inverse, the discrete logarithm, which are believed to be one-way  
 1771 functions, that is, while it is possible to compute the exponential map in **polynomial time**, com-  
 1772 puting the discrete log takes (sub)-**exponential time**. We have discussed this briefly following  
 1773 example 3.5 in the previous chapter, and will look at this and similar problems in more detail in  
 1774 the next section.

polynomial  
time

exponential  
time

### 1776 4.1.1 Cryptographic Groups

1777 In this section, we will look at families of groups that are believed to satisfy certain **compu-**  
 1778 **tational hardness assumptions**, namely that a particular problem cannot be solved efficiently  
 1779 (where efficiently typically means “in polynomial time of a given security parameter”) in the  
 1780 groups under consideration.

1781 *Example 38.* To highlight the concept of the computational hardness assumption, consider the  
 1782 group of integers  $\mathbb{Z}$  from example 3.5. One of the best known and most researched examples of  
 1783 computational hardness is the assumption that the factorization of integers into prime numbers  
 1784 cannot be solved by any algorithm in polynomial time with respect to the bit-length of the  
 1785 integer.

1786 To be more precise, the computational hardness assumption of integer factorization assumes  
 1787 that, given any integer  $z \in \mathbb{Z}$  with bit-length  $b$ , there is no integer  $k$  and no algorithm with the

TODO:  
Funda-  
mental  
theorem  
of finite  
cyclic  
groups

check  
reference

runtime complexity of  $\mathcal{O}(b^k)$  that is able to find the prime numbers  $p_1, p_2, \dots, p_j \in \mathbb{P}$ , such that  $z = p_1 \cdot p_2 \cdot \dots \cdot p_j$ .

This hardness assumption was proven to be false, since Shor's (1994) algorithm shows that integer factorization is at least efficiently possible on a quantum computer, since the runtime complexity of this algorithm is  $\mathcal{O}(b^3)$ . However, no such algorithm is known on a classical computer.

In the realm of classical computers, however, we still have to call the non-existence of such an algorithm an “assumption” because, to date, there is no proof that it is actually impossible to find one. The problem is that it is hard to reason about algorithms that we don't know.

So, despite the fact that there is currently no known algorithm that can factor integers efficiently on a classical computer, we cannot exclude that such an algorithm might exist in principle, and that someone eventually will discover it in the future.

However, what still makes the assumption plausible, despite the absence of any actual proof, is the fact that, after decades of extensive search, still no such algorithm has been found.

In what follows, we will describe a few computational hardness assumptions that arise in the context of groups in cryptography, because we will refer to them throughout the book.

**The discrete logarithm assumption** The so-called discrete logarithm problem is one of the most fundamental assumptions in cryptography. To define it, let  $\mathbb{G}$  be a finite cyclic group of order  $r$  and let  $g$  be a generator of  $\mathbb{G}$ . We know from 4.2 that there is an exponential map  $g^{(\cdot)} : \mathbb{Z}_r \rightarrow \mathbb{G} : x \mapsto g^x$  that maps the residue classes from modulo  $r$  arithmetic onto the group in a 1 : 1 correspondence. The **discrete logarithm problem** is the task of finding inverses to this map, that is, to find a solution  $x \in \mathbb{Z}_r$  to the following equation for some given  $h \in \mathbb{G}$ :

$$h = g^x \quad (4.4)$$

In other words, the **discrete logarithm assumption (DL-A)** is the assumption that there exists no algorithm with polynomial running time in the security parameter  $\log_2(r)$ , that is able to compute some  $x$  if only  $h$ ,  $g$  and  $g^x$  are given in  $\mathbb{G}$ . If this is the case for  $\mathbb{G}$ , we call  $\mathbb{G}$  a **DL-A group**.

Rephrasing the previous definition, DL-A groups are believed to have the property that it is infeasible to compute some number  $x$  that solves the equation  $h = g^x$  for a given  $h$  and  $g$ , assuming that the size of the group  $r$  is large enough.

*Example 39 (Public key cryptography).* One the most basic examples of an application for DL-A groups is in public key cryptography, where the parties publicly agree on some pair  $(\mathbb{G}, g)$  such that  $\mathbb{G}$  is a finite cyclic group of sufficiently large order  $r$ , where  $\mathbb{G}$  is believed to be a DL-A group, and  $g$  is a generator of  $\mathbb{G}$ .

In this setting, a secret key is some number  $sk \in \mathbb{Z}_r$  and the associated public key  $pk$  is the group element  $pk = g^{sk}$ . Since discrete logarithms are assumed to be hard, it is infeasible for an attacker to compute the secret key from the public key, since it is believed to be hard to find solutions  $x$  to the following equation:

$$pk = g^x \quad (4.5)$$

As the previous example shows, identifying DL-A groups is an important practical problem. Unfortunately, it is easy to see that it does not make sense to assume the hardness of the discrete logarithm problem in all finite cyclic groups: Counterexamples are common and easy to construct.

runtime complexity

S: what does “efficiently” mean here?

computational hardness assumptions

check reference

*Example 40* (Modular arithmetics for Fermat's primes). It is widely believed that the discrete logarithm problem is hard in multiplicative groups  $\mathbb{Z}_p^*$  of prime number modular arithmetics. However, this is not true in general. To see that, consider any so-called Fermat's prime, which is a prime number  $p \in \mathbb{P}$ , such that  $p = 2^n + 1$  for some number  $n$ .

We know from exercise 27 that in this case  $\mathbb{Z}_p^* = \{1, 2, \dots, p-1\}$  is a group with respect to integer multiplication in modular  $p$  arithmetics and since  $p = 2^n + 1$ , the order of  $\mathbb{Z}_p^*$  is  $2^n$ , which implies that the associated security parameter is given by  $\log_2(2^n) = n$ .

We show that, in this case,  $\mathbb{Z}_p^*$  is not a DL-A group, by constructing an algorithm, which is able compute some  $x \in \mathbb{Z}_{2^n}$  for any given generator  $g$  and arbitrary element  $h$  of  $\mathbb{F}_p^*$ , such that equation 4.6 holds, and the runtime complexity of the constructed algorithm is  $\mathcal{O}(n^2)$ , which is quadratic in the security parameter  $n = \log_2(2^n)$ .

$$h = g^x \quad (4.6)$$

To define such an algorithm, let us assume that the generator  $g$  is a public constant and that a group element  $h$  is given. Our task is to compute  $x$  efficiently.

The first thing to note is that, since  $x$  is a number in modular  $2^n$  arithmetic, we can write the binary representation of  $x$  as in 4.7, with binary coefficients  $c_j \in \{0, 1\}$ . In particular,  $x$  is an  $n$ -bit number if interpreted as an integer.

$$x = c_0 \cdot 2^0 + c_1 \cdot 2^1 + \dots + c_n \cdot 2^n \quad (4.7)$$

We then use this representation to construct an algorithm that computes the bits  $c_j$  one after another, starting at  $c_0$ . To see how this can be achieved, observe that we can determine  $c_0$  by raising the input  $h$  to the power of  $2^{n-1}$  in  $\mathbb{F}_p^*$ . We use the exponential laws and compute as follows:

$$\begin{aligned} h^{2^{n-1}} &= (g^x)^{2^{n-1}} \\ &= \left( g^{c_0 \cdot 2^0 + c_1 \cdot 2^1 + \dots + c_n \cdot 2^n} \right)^{2^{n-1}} \\ &= g^{c_0 \cdot 2^{n-1}} \cdot g^{c_1 \cdot 2^1 \cdot 2^{n-1}} \cdot g^{c_2 \cdot 2^2 \cdot 2^{n-1}} \dots g^{c_n \cdot 2^n \cdot 2^{n-1}} \\ &= g^{c_0 2^{n-1}} \cdot g^{c_1 2^0 \cdot 2^n} \cdot g^{c_2 2^1 \cdot 2^n} \dots g^{c_n 2^{n-1} \cdot 2^n} \end{aligned}$$

Now, since  $g$  is a generator and  $\mathbb{F}_p^*$  is cyclic of order  $2^n$ , we know  $g^{2^n} = 1$  and therefore  $g^{k \cdot 2^n} = 1^k = 1$ . From this, it follows that all but the first factor in the last expression are equal to 1 and we can simplify the expression into the following:

$$h^{2^{n-1}} = g^{c_0 2^{n-1}} \quad (4.8)$$

Now, in case  $c_0 = 0$ , we get  $h^{2^{n-1}} = g^0 = 1$ . In case  $c_0 = 1$ , we get  $h^{2^{n-1}} = g^{2^{n-1}} \neq 1$  (To see that  $g^{2^{n-1}} \neq 1$ , recall that  $g$  is a generator of  $\mathbb{F}_p^*$  and hence, is  $\mathbb{F}_p^*$  a cyclic group of order  $2^n$ , which implies  $g^y \neq 1$  for all  $y < 2^n$ ).

Raising  $h$  to the power of  $2^{n-1}$  determines  $c_0$ , and we can apply the same reasoning to the coefficient  $c_1$  by raising  $h \cdot g^{-c_0 \cdot 2^0}$  to the power of  $2^{n-2}$ . This approach can then be repeated until all the coefficients  $c_j$  of  $x$  are found.

Assuming that exponentiation in  $\mathbb{F}_p^*$  can be done in logarithmic runtime complexity  $\log(p)$ , it follows that our algorithm has a runtime complexity of  $\mathcal{O}(\log^2(p)) = \mathcal{O}(n^2)$ , since we have to execute  $n$  exponentiations to determine the  $n$  binary coefficients of  $x$ .

From this, it follows that whenever  $p$  is a Fermat's prime, the discrete logarithm assumption does not hold in  $\mathbb{F}_p^*$ .

check  
reference

explain  
last sen-  
tence  
more

**The decisional Diffie–Hellman assumption** Let  $\mathbb{G}$  be a finite cyclic group of order  $r$  and let  $g$  be a generator of  $\mathbb{G}$ . The decisional Diffie–Hellman assumption stipulates that there is no algorithm that has a polynomial runtime complexity in the security parameter  $s = \log(r)$  that is able to distinguish the so-called DDH- triple  $(g^a, g^b, g^{ab})$  from any triple  $(g^a, g^b, g^c)$  for randomly and independently chosen parameters  $a, b, c \in \mathbb{Z}_r$ . If this is the case for  $\mathbb{G}$ , we call  $\mathbb{G}$  a **DDH-A group**.

DDH-A is a stronger assumption than DL-A, in the sense that the discrete logarithm assumption is necessary for the decisional Diffie–Hellman assumption to hold, but not the other way around.

To see why this is the case, assume that the discrete logarithm assumption does not hold. In that case, given a generator  $g$  and a group element  $h$ , it is easy to compute some residue class  $x \in \mathbb{Z}_p$  with  $h = g^x$ . Then the decisional Diffie–Hellman assumption cannot hold, since given some triple  $(g^a, g^b, z)$ , one could efficiently decide whether  $z = g^{ab}$  is true by first computing the discrete logarithm  $b$  of  $g^b$ , then computing  $g^{ab} = (g^a)^b$  and deciding whether or not  $z = g^{ab}$ .

On the other hand, the following example shows that there are groups where the discrete logarithm assumption holds but the decisional Diffie–Hellman assumption does not.

*Example 41* (Efficiently computable pairings). Let  $\mathbb{G}$  be a finite, cyclic group of order  $r$  with generator  $g$ , such that the discrete logarithm assumption holds and there is a pairing map  $e(\cdot, \cdot) : \mathbb{G} \times \mathbb{G} \rightarrow \mathbb{G}_T$  for some target group  $\mathbb{G}_T$  that is computable in polynomial time of the parameter  $\log(r)$ .

In a setting like this, it is easy to show that DDH-A cannot hold, since given some triple  $(g^a, g^b, z)$ , it is possible to decide in polynomial time w.r.t  $\log(r)$  whether  $z = g^{ab}$  or not. To see that, check the following equation:

$$e(g^a, g^b) = e(g, z) \quad (4.9)$$

Since the bilinearity properties of  $e(\cdot, \cdot)$  imply  $e(g^a, g^b) = e(g, g)^{ab} = e(g, g^{ab})$ , and  $e(g, y) = e(g, y')$  implies  $y = y'$  due to the non-degenerate property, the equality means  $z = g^{ab}$ .

It follows that DDH-A is indeed weaker than DL-A, and groups with efficient pairings cannot be DDH-A groups. The following example shows another important class of groups where DDH-A does not hold: multiplicative groups of prime number residue classes.

*Example 42.* Let  $p$  be a prime number and  $\mathbb{Z}_p^* = \{1, 2, \dots, p-1\}$  the multiplicative group of modular  $p$  arithmetics as in exercise 27. As we have seen in XXX, this group is finite and cyclic of order  $p-1$  and every element  $g \neq 1$  is a generator.

To see that  $\mathbb{F}_p^*$  cannot be a DDH-A group, recall from XXX that the Legendre symbol  $\left(\frac{x}{p}\right)$  of any  $x \in \mathbb{F}_p^*$  is efficiently computable by Euler’s formular. But the Legendre symbol of  $g^a$  reveals whether  $a$  is even or odd. Given  $g^a, g^b$  and  $g^{ab}$ , one can thus efficiently compute and compare the least significant bit of  $a, b$  and  $ab$ , respectively, which provides a probabilistic method to distinguish  $g^{ab}$  from a random group element  $g^c$ .

**The computational Diffie–Hellman assumption** Let  $\mathbb{G}$  be a finite cyclic group of order  $r$  and let  $g$  be a generator of  $\mathbb{G}$ . The computational Diffie–Hellman assumption stipulates that, given randomly and independently chosen residue classes  $a, b \in \mathbb{Z}_r$ , it is not possible to compute  $g^{ab}$  if only  $g, g^a$  and  $g^b$  (but not  $a$  and  $b$ ) are known. If this is the case for  $\mathbb{G}$ , we call  $\mathbb{G}$  a CDH-A group.

In general, we don’t know if CDH-A is a stronger assumption than DL-A, or if both assumptions are equivalent. We know that DL-A is necessary for CDH-A, but the other direction

“equation”?

check  
reference

what’s the  
difference  
between  
 $\mathbb{F}_p^*$  and  
 $\mathbb{Z}_p^*$ ?

Legendre  
symbol

Euler’s  
formular

These  
are only  
explained  
later in  
the text,  
‘4.27’

is currently not well understood. In particular, there are no groups known where DL-A holds but CDH-A does not hold [Fifield, 2012].

To see why the discrete logarithm assumption is necessary, assume that it does not hold. So, given a generator  $g$  and a group element  $h$ , it is easy to compute some residue class  $x \in \mathbb{Z}_p$  with  $h = g^x$ . In that case, the computational Diffie–Hellman assumption cannot hold, since, given  $g$ ,  $g^a$  and  $g^b$ , one can efficiently compute  $b$  and hence is able to compute  $g^{ab} = (g^a)^b$  from this data.

The computational Diffie–Hellman assumption is a weaker assumption than the decisional Diffie–Hellman assumption, which means that there are groups where CDH-A holds and DDH-A does not hold, while there cannot be groups such that DDH-A holds but CDH-A does not hold. To see that, assume that it is efficiently possible to compute  $g^{ab}$  from  $g$ ,  $g^a$  and  $g^b$ . Then, given  $(g^a, g^b, z)$  it is easy to decide if  $z = g^{ab}$  holds or not.

Several variations and special cases of the CDH-A exist. For example, the **square computational Diffie–Hellman assumption** assumes that, given  $g$  and  $g^x$ , it is computationally hard to compute  $g^{x^2}$ . The **inverse computational Diffie–Hellman assumption** assumes that, given  $g$  and  $g^x$ , it is computationally hard to compute  $g^{x^{-1}}$ .

## Cofactor Clearing

### 4.1.2 Hashing to Groups

**Hash functions** Generally speaking, a hash function is any function that can be used to map data of arbitrary size to fixed-size values. Since binary strings of arbitrary length are a general way to represent arbitrary data, we can understand a general **hash function** as the following map where  $\{0, 1\}^*$  represents the set of all binary strings of arbitrary but finite length and  $\{0, 1\}^k$  represents the set of all binary strings that have a length of exactly  $k$  bits:

$$H : \{0, 1\}^* \rightarrow \{0, 1\}^k \quad (4.10)$$

In our definition, a hash function maps binary strings of arbitrary size onto binary strings of size exactly  $k$ . The **images** of  $H$ , that is, the values returned by the hash function  $H$ , are called **hash values**, **digests**, or simply **hashes**.

A hash function must be deterministic, that is, when we insert the same input  $x$  into  $H$ , the image  $H(x)$  must always be the same. In addition, a hash function should be as uniform as possible, which means that it should map input values as evenly as possible over its output range. In mathematical terms, every string of length  $k$  from  $\{0, 1\}^k$  should be generated with roughly the same probability.

*Example 43 ( $k$ -truncation hash).* One of the most basic hash functions  $H_k : \{0, 1\}^* \rightarrow \{0, 1\}^k$  is given by simply truncating every binary string  $s$  of size  $s.len() > k$  to a string of size  $k$  and by filling any string  $s'$  of size  $s'.len() < k$  with zeros. To make this hash function deterministic, we define that both truncation and filling should happen “on the left”.

For example, if  $k = 3$ ,  $x_1 = (0000101011101010011101010101)$  and  $x_2 = 1$ , then  $H(x_1) = (101)$  and  $H(x_2) = (001)$ . It is easy to see that this hash function is deterministic and uniform.

Of particular interest are so-called **cryptographic** hash functions, which are hash functions that are also **one-way functions**, which essentially means that, given a string  $y$  from  $\{0, 1\}^k$  it is practically infeasible to find a string  $x \in \{0, 1\}^*$  such that  $H(x) = y$  holds. This property is usually called **preimage-resistance**.

are these going to be relevant later? yes, they are used in various snark proof systems

TODO: theorem: every factor of order defines a subgroup...



In addition, it should be infeasible to find two strings  $x_1, x_2 \in \{0, 1\}^*$ , such that  $H(x_1) = H(x_2)$ , which is called **collision resistance**. It is important to note, though, that collisions always exist, since a function  $H : \{0, 1\}^* \rightarrow \{0, 1\}^k$  inevitably maps infinitely many values onto the same hash. In fact, for any hash function with digests of length  $k$ , finding a preimage to a given digest can always be done using a brute force search in  $2^k$  evaluation steps. It should just be practically impossible to compute those values, and statistically very unlikely to generate two of them by chance.

A third property of a cryptographic hash function is that small changes in the input string, like changing a single bit, should generate hash values that look completely different from each other.

Because cryptographically secure hash functions map tiny changes in input values onto large changes in the output, implementation errors that change the outcome are usually easy to spot by comparing them to expected output values. The definitions of cryptographically secure hash functions are therefore usually accompanied by some test vectors of common inputs and expected digests. Since the empty string  $''$  is the only string of length 0, a common test vector is the expected digest of the empty string.

*Example 44* ( $k$ -truncation hash). Consider the  $k$ -truncation hash from example 43. Since the empty string has length 0, it follows that the digest of the empty string is string of length  $k$  that only contains 0's:

$$H_k('') = (000 \dots 000) \quad (4.11)$$

It is pretty obvious from the definition of  $H_k$  that this simple hash function is not a cryptographic hash function. In particular, every digest is its own preimage, since  $H_k(y) = y$  for every string of size exactly  $k$ . Finding preimages is therefore easy, so the property of preimage resistance does not hold.

In addition, it is easy to construct collisions as all strings of size  $> k$  that share the same  $k$ -bits "on the right" are mapped to the same hash value, so this function is not collision resistant, either.

Finally, this hash function is not very chaotic, as changing bits that are not part of the  $k$  right-most bits don't change the digest at all.

Computing cryptographically secure hash functions in pen-and-paper style is possible but tedious. Fortunately, Sage can import the **hashlib** library, which is intended to provide a reliable and stable base for writing Python programs that require cryptographic functions. The following examples explain how to use hashlib in Sage.

*Example 45.* An example of a hash function that is generally believed to be a cryptographically secure hash function is the so-called **SHA256** hash, which, in our notation, is a function that maps binary strings of arbitrary length onto binary strings of length 256:

$$SHA256 : \{0, 1\}^* \rightarrow \{0, 1\}^{256} \quad (4.12)$$

To evaluate a proper implementation of the *SHA256* hash function, the digest of the empty string is supposed to be

$$SHA256('') = e3b0c44298fc1c149afbf4c8996fb92427ae41e4649b934ca495991b7852b855 \quad (4.13)$$

For better human readability, it is common practice to represent the digest of a string not in its binary form, but in a hexadecimal representation. We can use Sage to compute *SHA256* and freely transit between binary, hexadecimal and decimal representations. To do so, we import hashlib's implementation of *SHA256*:

Is there a term for this property?

```

1984 sage: import hashlib 144
1985 sage: test = 'e3b0c44298fc1c149afbf4c8996fb92427ae41e4649b934 145
1986         ca495991b7852b855'
1987 sage: hasher = hashlib.sha256(b'') 146
1988 sage: str = hasher.hexdigest() 147
1989 sage: type(str) 148
1990 <class 'str'> 149
1991 sage: d = ZZ('0x'+ str) # conversion to integer type 150
1992 sage: d.str(16) == str 151
1993 True 152
1994 sage: d.str(16) == test 153
1995 True 154
1996 sage: d.str(16) 155
1997 e3b0c44298fc1c149afbf4c8996fb92427ae41e4649b934ca495991b7852b8 156
1998 55
1999 sage: d.str(2) 157
2000 11100011101100001100010001000010100110001111110000011100000101 158
2001 0010011010111110111110100110010001001100101101111101110010
2002 01001000010011110101110010000011110010001100100100110111001
2003 00110100110010100100100101011001100100011011011110000101001
2004 01011100001010101
2005 sage: d.str(10) 159
2006 10298733624955409702953521232258132278979990064819803499337939 160
2007 7001115665086549

```

2008 **Hashing to cyclic groups** As we have seen in the previous paragraph, general hash functions  
2009 map binary strings of arbitrary length onto binary strings of length  $k$ . However, it is desirable  
2010 in various cryptographic primitives to not simply hash to binary strings of fixed length but to  
2011 hash into algebraic structures like groups, while keeping (some of) the properties like preimage  
2012 resistance or collision resistance.

2013 Hash functions like this can be defined for various algebraic structures, but, in a sense, the  
2014 most fundamental ones are hash functions that map into groups, because they can be easily  
2015 extended to map into other structures like rings or fields.

2016 To give a more precise definition, let  $\mathbb{G}$  be a group and  $\{0, 1\}^*$  the set of all finite, binary  
2017 strings, then a **hash-to-group** function is a deterministic map

$$H : \{0, 1\}^* \rightarrow \mathbb{G} \quad (4.14)$$

2018 Common properties of hash functions, like uniformity, are desirable but not always realized in  
2019 real-world instantiations of hash-to-group functions, so we skip those requirements for now and  
2020 keep the definition very general.

2021 As the following example shows, hashing to finite cyclic groups can be trivially achieved  
2022 for the price of some undesirable properties of the hash function:

2023 *Example 46* (Naive cyclic group hash). Let  $\mathbb{G}$  be a finite cyclic group. If the task is to implement  
2024 a hash-to-group function, one immediate approach can be based on the observation that binary  
2025 strings of size  $k$  can be interpreted as integers  $z \in \mathbb{Z}$  in the range  $0 \leq z < 2^k$ .

2026 To be more precise, choose an ordinary hash function  $H : \{0, 1\}^* \rightarrow \{0, 1\}^k$  for some pa-  
2027 rameter  $k$  and a generator  $g$  of  $\mathbb{G}$ . Then the expression below is a positive integer (where  $H(s)_j$

means the bit at the  $j$ -th position of  $H(s)$ ):

$$z_{H(s)} = H(s)_0 \cdot 2^0 + H(s)_1 \cdot 2^1 + \dots + H(s)_k \cdot 2^k \quad (4.15)$$

A hash-to-group function for the group  $\mathbb{G}$  can then be defined as a concatenation of the exponential map  $g^{(\cdot)}$  of  $g$  with the interpretation of  $H(s)$  as an integer:

$$H_g : \{0, 1\}^* \rightarrow \mathbb{G} : s \mapsto g^{z_{H(s)}} \quad (4.16)$$

Constructing a hash-to-group function like this is easy for cyclic groups, and it might be good enough in certain applications. It is, however, almost never adequate in cryptographic applications, as discrete log relations might be constructible between two given hash values  $H_g(s)$  and  $H_g(t)$ .

a few examples?

To see that, assume that  $\mathbb{G}$  is of order  $r$  and that  $z_{H(s)}$  has a multiplicative inverse in modular  $r$  arithmetics. In that case, we can compute  $x = z_{H(t)} \cdot z_{H(s)}^{-1}$  in  $\mathbb{Z}_r$  and find a discrete log relation between the group hash values, that is, find some  $x$  with  $H_g(t) = (H_g(s))^x$ :

$$\begin{aligned} H_g(t) &= (H_g(s))^x && \Leftrightarrow \\ g^{z_{H(t)}} &= g^{z_{H(s)} \cdot x} && \Leftrightarrow \\ g^{z_{H(t)}} &= g^{z_{H(t)}} \end{aligned}$$

Therefore applications where discrete log relations between hash values are undesirable need different approaches. Many of these approaches start with a way to hash into the set  $\mathbb{Z}_r$  of modular  $r$  arithmetics.

**Hashing to modular arithmetics** One of the most widely used applications of hash-into-group functions are hash functions that map into the set  $\mathbb{Z}_r$  of modular  $r$  arithmetics for some modulus  $r$ . Different approaches to construct such a function are known, but probably the most widely used ones are based on the insight that the images of arbitrary hash functions can be interpreted as binary representations of integers, as explained in example 46.

check reference

It follows from this interpretation that one simple method of hashing into  $\mathbb{Z}_r$  is constructed by observing that if  $r$  is a modulus with a bit-length of  $k = r.\text{nbits}()$ , then every binary string  $(b_0, b_1, \dots, b_{k-2})$  of length  $k - 1$  defines an integer  $z$  in the range  $0 \leq z < 2^{k-1} \leq r$ , by defining  $z$ :

$$z = b_0 \cdot 2^0 + b_1 \cdot 2^1 + \dots + b_{k-2} \cdot 2^{k-2} \quad (4.17)$$

Now, since  $z < r$ , we know that  $z$  is guaranteed to be in the set  $\{0, 1, \dots, r-1\}$ , and hence it can be interpreted as an element of  $\mathbb{Z}_r$ . From this it follows that if  $H : \{0, 1\}^* \rightarrow \{0, 1\}^{k-1}$  is a hash function, then a hash-to-group function can be constructed as follows (where  $H(s)_j$  means the  $j$ -th bit of the image binary string  $H(s)$  of the original binary hash function):

$$H_{r.\text{nbits}()-1} : \{0, 1\}^* \rightarrow \mathbb{Z}_r : s \mapsto H(s)_0 \cdot 2^0 + H(s)_1 \cdot 2^1 + \dots + H(s)_{k-2} \cdot 2^{k-2} \quad (4.18)$$

A drawback of this hash function is that the distribution of the hash values in  $\mathbb{Z}_r$  is not necessarily uniform. In fact, if  $r - 2^{k-1} \neq 0$ , then by design  $H_{r.\text{nbits}()-1}$  will never hash onto values  $z \geq 2^{k-1}$ . Good moduli  $r$  are therefore as close to  $2^{k-1}$  as possible, while less good moduli are closer to  $2^k$ . In the worst case, when  $r = 2^k - 1$ , it misses  $2^{k-1} - 1$ , that is, almost half of all elements, from  $\mathbb{Z}_r$ .

An advantage of this approach is that properties like preimage resistance or collision resistance of the original hash function  $H(\cdot)$  are preserved.

TODO:  
DOUBLE  
CHECK  
THIS  
REA-  
SONING.



*Example 47.* To give an implementation of the  $H_{r.nb\text{its}()-1}$  hash function, we use a 5-bit truncation of the *SHA256* hash from example 45 and define a hash into  $\mathbb{Z}_{16}$  as follows:

$$H_{16.nb\text{its}()-1} : \{0, 1\}^* \rightarrow \mathbb{Z}_{16} : s \mapsto \text{SHA256}(s)_0 \cdot 2^0 + \text{SHA256}(s)_1 \cdot 2^1 + \dots + \text{SHA256}(s)_4 \cdot 2^4$$

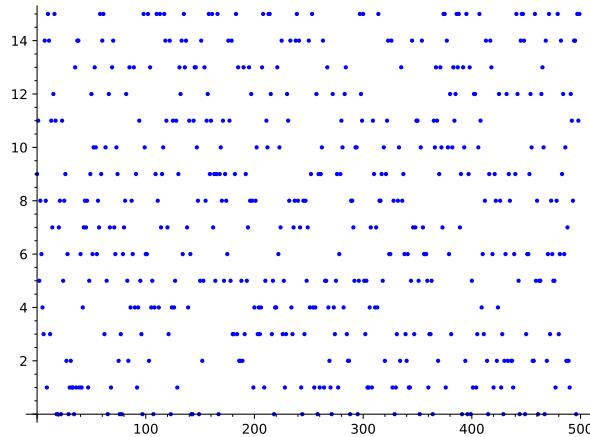
2058 Since  $k = 16.nb\text{its}() = 5$  and  $16 - 2^{k-1} = 0$ , this hash maps uniformly onto  $\mathbb{Z}_{16}$ . We can invoke  
 2059 Sage to implement it:

```

2060 sage: import hashlib                                161
2061 sage: def Hash5(x):                                162
2062     ....:     hasher = hashlib.sha256(x)            163
2063     ....:     digest = hasher.hexdigest()           164
2064     ....:     d = ZZ(digest, base=16)               165
2065     ....:     d = d.str(2)[-4:]                     166
2066     ....:     return ZZ(d, base=2)                  167
2067 sage: Hash5(b' ')                                  168
2068 5                                                    169

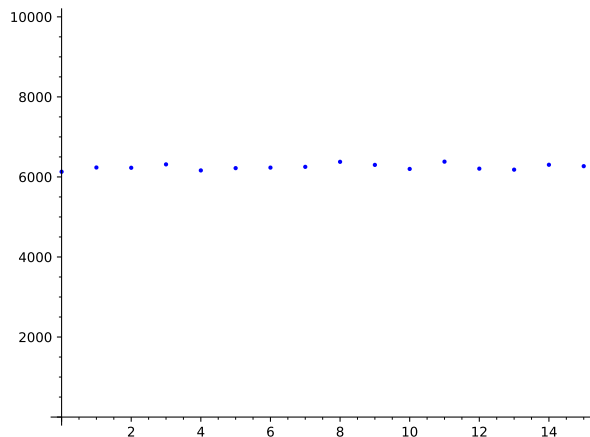
```

2069 We can then use Sage to apply this function to a large set of input values in order to plot a  
 2070 visualization of the distribution over the set  $\{0, \dots, 15\}$ . Executing over 500 input values gives  
 2071 the following plot:



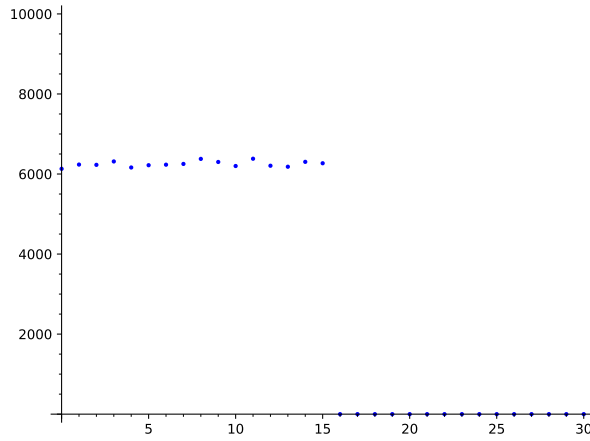
2072

2073 To get an intuition of uniformity, we can count the number of times the hash function  $H_{16.nb\text{its}()-1}$   
 2074 maps onto each number in the set  $\{0, 1, \dots, 15\}$  in a loop of 100000 hashes, and compare that  
 2075 to the ideal uniform distribution, which would map exactly 6250 samples to each element. This  
 2076 gives the following result:



2077

The uniformity of distribution problem becomes apparent if we want to construct a similar hash function for  $\mathbb{Z}_r$  for any  $r$  in the range  $17 \leq r \leq 31$ . In this case, the definition of the hash function is exactly the same as for  $\mathbb{Z}_{16}$ , and hence, the images will not exceed the value 16. So, for example, even in the case of hashing to  $\mathbb{Z}_{31}$ , the hash function never maps to any value larger than 16, leaving almost half of all numbers out of the image range.



The second widely used method of hashing into  $\mathbb{Z}_r$  is constructed by observing the following: If  $r$  is a modulus with a bit-length of  $r.bits() = k_1$  and  $H : \{0, 1\}^* \rightarrow \{0, 1\}^{k_2}$  is a hash function that produces digests of size  $k_2$ , with  $k_2 \geq k_1$ , then a hash-to-group function can be constructed by interpreting the image of  $H$  as a binary representation of an integer and then taking the modulus by  $r$  to map into  $\mathbb{Z}_r$ . This is formalized in the equation below, where  $H(s)_j$  means the  $j$ 'th bit of the image binary string  $H(s)$  of the original binary hash function.

$$H'_{mod_r} : \{0, 1\}^* \rightarrow \mathbb{Z}_r : s \mapsto \left( H(s)_0 \cdot 2^0 + H(s)_1 \cdot 2^1 + \dots + H(s)_{k_2} \cdot 2^{k_2} \right) \bmod r \quad (4.19)$$

A drawback of this hash function is that computing the modulus requires some computational effort. In addition, the distribution of the hash values in  $\mathbb{Z}_r$  might not be even, depending on the difference  $2^{k_2+1} - r$ . An advantage of it is that potential properties of the original hash function  $H(\cdot)$  (like preimage resistance or collision resistance) are preserved, and the distribution can be made almost uniform, with only negligible bias depending on what modulus  $r$  and images size  $k_2$  are chosen.

*Example 48.* To give an implementation of the  $H_{mod_r}$  hash function, we use  $k_2$ -bit truncation of the *SHA256* hash from example 45, and define a hash into  $\mathbb{Z}_{23}$  as follows:

$$H_{mod_{23}, k_2} : \{0, 1\}^* \rightarrow \mathbb{Z}_{23} : s \mapsto \left( SHA256(s)_0 \cdot 2^0 + SHA256(s)_1 \cdot 2^1 + \dots + SHA256(s)_{k_2} \cdot 2^{k_2} \right) \bmod 23$$

We want to use various instantiations of  $k_2$  to visualize the impact of truncation length on the distribution of the hashes in  $\mathbb{Z}_{23}$ . We can invoke Sage to implement it as follows:

```
sage: import hashlib 170
sage: Z23 = Integers(23) 171
sage: def Hash_mod23(x, k2): 172
.....:     hasher = hashlib.sha256(x.encode('utf-8')) 173
.....:     digest = hasher.hexdigest() 174
.....:     d = ZZ(digest, base=16) 175
```

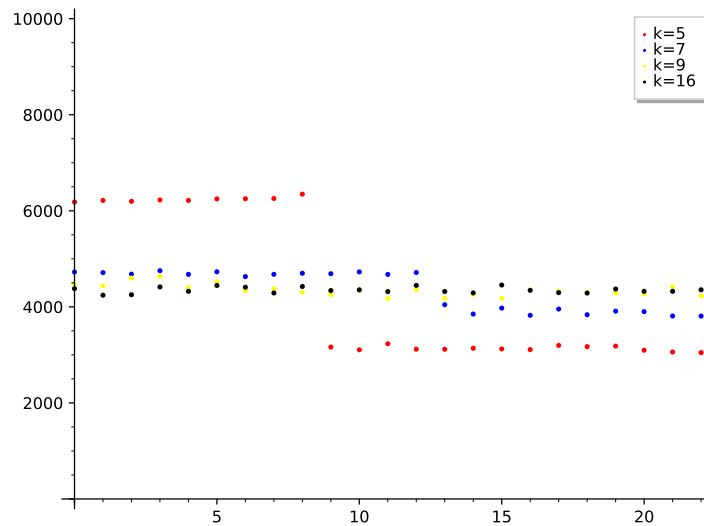
Mirco:  
We can  
do better  
than this

```

2105     ....:     d = d.str(2)[-k2:]
2106     ....:     d = ZZ(d, base=2)
2107     ....:     return ZZ3(d)

```

2108 We can then use Sage to apply this function to a large set of input values in order to plot  
 2109 visualizations of the distribution over the set  $\{0, \dots, 22\}$  for various values of  $k_2$ , by counting  
 2110 the number of times it maps onto each number in a loop of 100000 hashes. We get the following  
 2111 plot:



2112

2113 A third method that can sometimes be found in implementations is the so-called “**try-and-**  
 2114 **increment**” method. To understand this method, we define an integer  $z \in \mathbb{Z}$  from any hash  
 2115 value  $H(s)$  as we did in the previous methods, that is, we define  $z = H(s)_0 \cdot 2^0 + H(s)_1 \cdot 2^1 +$   
 2116  $\dots + H(s)_{k-1} \cdot 2^k$ .

2117 Hashing into  $\mathbb{Z}_r$  is then achievable by first computing  $z$ , and then trying to see if  $z \in \mathbb{Z}_r$ . If  
 2118 it is, then the hash is done; if not, the string  $s$  is modified in a deterministic way and the process  
 2119 is repeated until a suitable number  $z$  is found. A suitable, deterministic modification could be  
 2120 to concatenate the original string by some bit counter. A “try-and-increment” algorithm would  
 then work like in algorithm 5.

check  
reference

---

#### Algorithm 5 Hash-to- $\mathbb{Z}_n$

---

**Require:**  $r \in \mathbb{Z}$  with  $r.\text{nbits}() = k$  and  $s \in \{0, 1\}^*$

**procedure** TRY-AND-INCREMENT( $r, k, s$ )

$c \leftarrow 0$

**repeat**

$s' \leftarrow s || c.\text{bits}()$

$z \leftarrow H(s')_0 \cdot 2^0 + H(s')_1 \cdot 2^1 + \dots + H(s')_k \cdot 2^k$

$c \leftarrow c + 1$

**until**  $z < r$

**return**  $x$

**end procedure**

**Ensure:**  $z \in \mathbb{Z}_r$

---

2121

2122 Depending on the parameters, this method can be very efficient. In fact, if  $k$  is sufficiently  
 2123 large and  $r$  is close to  $2^{k+1}$ , the probability for  $z < r$  is very high and the repeat loop will almost

always be executed a single time only. A drawback is, however, that the probability of having to execute the loop multiple times is not zero.

Once some hash function into modular arithmetics is found, it can often be combined with additional techniques to hash into more general finite cyclic groups. The following paragraphs describe a few of those methods widely adopted in SNARK development.

**Pedersen Hashes** The so-called **Pedersen hash function** [Pedersen, 1992] provides a way to map binary inputs of fixed size  $k$  onto elements of finite cyclic groups that avoids discrete log relations between the images as they occur in the naive approach XXX. Combining it with a classical hash function provides a hash function that maps strings of arbitrary length onto group elements.

To be more precise, let  $j$  be an integer,  $\mathbb{G}$  a finite cyclic group of order  $r$  and  $\{g_1, \dots, g_j\} \subset \mathbb{G}$  a uniform randomly generated set of generators of  $\mathbb{G}$ . Then **Pedersen's hash function** is defined as follows:

$$H_{Ped} : (\mathbb{Z}_r)^j \rightarrow \mathbb{G} : (x_1, \dots, x_j) \mapsto \prod_{i=1}^j g_i^{x_i} \quad (4.20)$$

It can be shown that Pedersen's hash function is collision-resistant under the assumption that  $\mathbb{G}$  is a DL-A group. However, it is important to note that Pedersen hashes cannot be assumed to be **pseudorandom** and should therefore not be used where a hash function serves as an approximation of a random **oracle**. **will these be explained in the initial chapters?**

From an implementation perspective, it is important to derive the set of generators  $\{g_1, \dots, g_j\}$  in such a way that they are as uniform and random as possible. In particular, any known discrete log relation between two generators, that is, any known  $x \in \mathbb{Z}_r$  with  $g_h = (g_i)^x$  must be avoided.

To see how Pedersen hashes can be used to define an actual hash-to-group function according to our definition, we can use any of the hash-to- $\mathbb{Z}_r$  functions as we have derived them in equation 4.18.

**MimC Hashes** [Albrecht et al., 2016]

**Pseudorandom Functions in DDH-A groups** As noted in above, Pederson's hash function does not have the properties a random function and should therefore not be instantiated as such. To see an example of a random oracle function in groups where the decisional Diffie–Hellman construction is assumed to hold true, let  $\mathbb{G}$  be a DDH-A group of order  $r$  with generator  $g$  and  $\{a_0, a_1, \dots, a_k\} \subset \mathbb{Z}_r^*$  a uniform randomly generated set of numbers invertible in modular  $r$  arithmetics. Then a pseudo-random function is given by the as follows:

$$F_{rand} : \{0, 1\}^{k+1} \rightarrow \mathbb{G} : (b_0, \dots, b_k) \mapsto g^{b_0 \cdot \prod_{i=1}^k a_i^{b_i}} \quad (4.21)$$

Of course, if  $H : \{0, 1\}^* \rightarrow \{0, 1\}^{k+1}$  is a random oracle, then the concatenation of  $F_{rand}$  and  $H$  also defines a random oracle

$$H_{rand, \mathbb{G}} : \{0, 1\}^* \rightarrow \mathbb{G} : s \mapsto F_{rand}(H(s)) \quad (4.22)$$

## 4.2 Commutative Rings

Thinking back to operations on integers, we know that there are two of these: addition and multiplication. As we have seen, addition defines a group structure on the set of integers.

However, multiplication does not define a group structure, given that integers generally don't have multiplicative inverses.

Configurations like these constitute so-called **commutative rings with unit**. To be more precise, a commutative ring with unit  $(R, +, \cdot, 1)$  is a set  $R$  provided with two maps  $+: R \cdot R \rightarrow R$  and  $\cdot: R \cdot R \rightarrow R$ , called **addition** and **multiplication**, such that the following conditions hold:

**Definition 4.2.0.1. Commutative ring with unit**

- $(R, +)$  is a commutative group, where the neutral element is denoted with 0. **Commutativity of multiplication:**  $r_1 \cdot r_2 = r_2 \cdot r_1$  for all  $r_1, r_2 \in R$ .
- **Existence of a unit:** There is an element  $1 \in R$ , such that  $1 \cdot g$  holds for all  $g \in R$ ,
- **Associativity:** For every  $g_1, g_2, g_3 \in \mathbb{G}$  the equation  $g_1 \cdot (g_2 \cdot g_3) = (g_1 \cdot g_2) \cdot g_3$  holds.
- **Distributivity:** For all  $g_1, g_2, g_3 \in R$  the distributive laws  $g_1 \cdot (g_2 + g_3) = g_1 \cdot g_2 + g_1 \cdot g_3$  holds.

**Example 49** (The ring of integers). The set  $\mathbb{Z}$  of integers with the usual addition and multiplication is the archetypical example of a commutative ring with unit 1.

**Example 50** (Underlying commutative group of a ring). Every commutative ring with unit  $(R, +, \cdot, 1)$  gives rise to a group, if we disregard multiplication.

The following example is somewhat unusual, but we encourage you to think through it because it helps to detach the mind from familiar styles of computation and concentrate on the abstract algebraic explanation.

**Example 51.** Let  $S := \{\bullet, \star, \odot, \otimes\}$  be a set that contains four elements, and let addition and multiplication on  $S$  be defined as follows:

$\cup$	$\bullet$	$\star$	$\odot$	$\otimes$
$\bullet$	$\bullet$	$\star$	$\odot$	$\otimes$
$\star$	$\star$	$\odot$	$\otimes$	$\bullet$
$\odot$	$\odot$	$\otimes$	$\bullet$	$\star$
$\otimes$	$\otimes$	$\bullet$	$\star$	$\odot$

$\circ$	$\bullet$	$\star$	$\odot$	$\otimes$
$\bullet$	$\bullet$	$\bullet$	$\bullet$	$\bullet$
$\star$	$\bullet$	$\star$	$\odot$	$\otimes$
$\odot$	$\bullet$	$\odot$	$\bullet$	$\odot$
$\otimes$	$\bullet$	$\otimes$	$\odot$	$\star$

Then  $(S, \cup, \circ)$  is a ring with unit  $\star$  and zero  $\bullet$ . It therefore makes sense to ask for solutions to equations like this one: Find  $x \in S$  such that

$$\otimes \circ (x \cup \odot) = \star$$

To see how such a “moonmath equation” can be solved, we have to keep in mind that rings behaves mostly like normal numbers when it comes to bracketing and computation rules. The only differences are the symbols, and the actual way to add and multiply them. With this in

mind, we solve the equation for  $x$  in the “usual way”:

$$\begin{array}{ll}
 \otimes \circ (x \cup \odot) = \star & \# \text{ apply the distributive law} \\
 \otimes \circ x \cup \otimes \circ \odot = \star & \# \otimes \circ \odot = \odot \\
 \otimes \circ x \cup \odot = \star & \# \text{ concatenate the } \cup \text{ inverse of } \odot \text{ to both sides} \\
 \otimes \circ x \cup \odot \cup -\odot = \star \cup -\odot & \# \odot \cup -\odot = \bullet \\
 \otimes \circ x \cup \bullet = \star \cup -\odot & \# \bullet \text{ is the } \cup \text{ neutral element} \\
 \otimes \circ x = \star \cup -\odot & \# \text{ for } \cup \text{ we have } -\odot = \odot \\
 \otimes \circ x = \star \cup \odot & \# \star \cup \odot = \otimes \\
 \otimes \circ x = \otimes & \# \text{ concatenate the } \circ \text{ inverse of } \otimes \text{ to both sides} \\
 (\otimes)^{-1} \circ \otimes \circ x = (\otimes)^{-1} \circ \otimes & \# \text{ multiply with the multiplicative inverse} \\
 \star \circ x = \star & \\
 x = \star & 
 \end{array}$$

2181 So, even though this equation looked really alien at first glance, we could solve it basically  
 2182 exactly the way we solve “normal” equations containing numbers.

2183 Note, however, that whenever a multiplicative inverse would be needed to solve an equation  
 2184 in the usual way in a ring, things can be very different than most of us are used to. For example,  
 2185 the following equation cannot be solved for  $x$  in the usual way, since there is no multiplicative  
 2186 inverse for  $\odot$  in our ring.

$$\odot \circ x = \otimes \quad (4.23)$$

2187 We can confirm this by looking at the multiplication table to see that no such  $x$  exists.

2188 As another example, the following equation does not have a single solution but two:  $x \in$   
 2189  $\{\star, \otimes\}$ .

$$\odot \circ x = \odot \quad (4.24)$$

2190 Having no solution or two solutions is certainly not something to expect from types like  $\mathbb{Q}$   
 2191 (rational numbers). can we use another set as an example? we hardly talked about  $\mathbb{Q}$  so far

2192 *Example 52.* Considering polynomials again, we note from their definition that what we have  
 2193 called the type  $R$  of the coefficients must in fact be a commutative ring with a unit, since we  
 2194 need addition, multiplication, commutativity and the existence of a unit for  $R[x]$  to have the  
 2195 properties we expect.

2196 Considering  $R$  to be a ring with addition and multiplication of polynomials as defined in  
 2197 4.2.0.1 actually makes  $R[x]$  into a commutative ring with a unit, too, where the polynomial 1 is  
 2198 the multiplicative unit. check reference

2199 *Example 53.* Let  $n$  be a modulus and  $(\mathbb{Z}_n, +, \cdot)$  the set of all remainder classes of integers  
 2200 modulo  $n$ , with the projection of integer addition and multiplication as defined in 4.2.0.1. It can  
 2201 be shown that  $(\mathbb{Z}_n, +, \cdot)$  is a commutative ring with unit 1. check reference

Considering the exponential map from page 43 again, let  $\mathbb{G}$  be a finite cyclic group of order  
 $n$  with generator  $g \in \mathbb{G}$ . Then the ring structure of  $(\mathbb{Z}_n, +, \cdot)$  is mapped onto the group structure  
of  $\mathbb{G}$  in the following way: check reference

$$\begin{array}{ll}
 g^{x+y} = g^x \cdot g^y & \text{for all } x, y \in \mathbb{Z}_n \\
 g^{x \cdot y} = (g^x)^y & \text{for all } x, y \in \mathbb{Z}_n
 \end{array}$$

This is of particular interest in cryptography and SNARKs, as it allows for the evaluation of polynomials with coefficients in  $\mathbb{Z}_n$  to be evaluated “in the exponent”. To be more precise, let  $p \in \mathbb{Z}_n[x]$  be a polynomial with  $p(x) = a_m \cdot x^m + a_{m-1}x^{m-1} + \dots + a_1x + a_0$ . Then the previously defined exponential laws (example 37) imply the following:

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$$\begin{aligned} g^{p(x)} &= g^{a_m \cdot x^m + a_{m-1}x^{m-1} + \dots + a_1x + a_0} \\ &= \left(g^{x^m}\right)^{a_m} \cdot \left(g^{x^{m-1}}\right)^{a_{m-1}} \cdot \dots \cdot (g^x)^{a_1} \cdot g^{a_0} \end{aligned}$$

Hence, to evaluate  $p$  at some point  $s$  in the exponent, we can insert  $s$  into the right-hand side of the last equation and evaluate the product.

As we will see, this is a key insight to understanding many SNARK protocols like e.g. Groth16 [Groth, 2016] or XXX.

*Example 54.* To give an example of the evaluation of a polynomial in the exponent of a finite cyclic group, consider the exponential map from example 37:

$$3^{(\cdot)} : \mathbb{Z}_4 \rightarrow \mathbb{Z}_5^* \quad x \mapsto 3^x$$

Choosing the polynomial  $p(x) = 2x^2 + 3x + 1$  from  $\mathbb{Z}_4[x]$ , we can evaluate the polynomial at say  $x = 2$  in the exponent of 3 in two different ways. On the one hand, we can evaluate  $p$  at 2 and then write the result into the exponent as follows:

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$$\begin{aligned} 3^{p(2)} &= 3^{2 \cdot 2^2 + 3 \cdot 2 + 1} \\ &= 3^{2 \cdot 0 + 2 + 1} \\ &= 3^3 \\ &= 2 \end{aligned}$$

On the other hand, we can use the right-hand side of the equation to evaluate  $p$  at 2 in the exponent of 3 as follows:

$$\begin{aligned} 3^{p(2)} &= \left(3^{2^2}\right)^2 \cdot (3^2)^3 \cdot 3^1 \\ &= (3^0)^2 \cdot 3^3 \cdot 3 \\ &= 1^2 \cdot 2 \cdot 3 \\ &= 2 \cdot 3 \\ &= 2 \end{aligned}$$

**Hashing to Commutative Rings** As we have seen in XXX, various constructions for hashing-to-groups are known and used in applications. As commutative rings are **Abelian groups**, when we disregard the multiplicative structure, hash-to-group constructions can be applied for hashing into commutative rings, too. This is possible in general, as the **codomain** of a general hash function  $\{0, 1\}^*$  is just the set of binary strings of arbitrary but finite length, which has no algebraic structure that the hash function must respect.

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## 4.3 Fields

We started this chapter with the definition of a group (section 4.1), which we then expanded into the definition of a commutative ring with a unit (section 4.2). Such rings generalize the behavior

of integers. In this section, we will look at those special cases of commutative rings where every element other than the neutral element of addition has a multiplicative inverse. Those structures behave very much like the rational numbers  $\mathbb{Q}$ . Rational numbers are, in a sense, an extension of the ring of integers, that is, they are constructed by including newly defined multiplicative inverses (fractions) to the integers.

Now, considering the definition of a ring (4.2.0.1) again, we define a **field**  $(\mathbb{F}, +, \cdot)$  to be a set  $\mathbb{F}$ , together with two maps  $+: \mathbb{F} \cdot \mathbb{F} \rightarrow \mathbb{F}$  and  $\cdot: \mathbb{F} \cdot \mathbb{F} \rightarrow \mathbb{F}$ , called *addition* and *multiplication*, such that the following conditions hold:

**Definition 4.3.0.1. Field**

- $(\mathbb{F}, +)$  is a commutative group, where the neutral element is denoted by 0.
- $(\mathbb{F} \setminus \{0\}, \cdot)$  is a commutative group, where the neutral element is denoted by 1.
- (Distributivity) For all  $g_1, g_2, g_3 \in \mathbb{F}$  the distributive law  $g_1 \cdot (g_2 + g_3) = g_1 \cdot g_2 + g_1 \cdot g_3$  holds.

If a field is given and the definition of its addition and multiplication is not ambiguous, we will often simply write  $\mathbb{F}$  instead of  $(\mathbb{F}, +, \cdot)$  to denote the field. Moreover, we use  $\mathbb{F}^*$  to describe the multiplicative group of the field, that is, the set of elements with multiplication as the group law, excluding the neutral element of addition.

The **characteristic** of a field  $\mathbb{F}$ , represented as  $\text{char}(\mathbb{F})$ , is the smallest natural number  $n \geq 1$  for which the  $n$ -fold sum of 1 equals zero, i.e. for which  $\sum_{i=1}^n 1 = 0$ . If such an  $n > 0$  exists, the field is also said to have a **finite characteristic**. If, on the other hand, every finite sum of 1 is such that it is not equal to zero, then the field is defined to have characteristic 0. [S: Tried to disambiguate the scope of negation between 1. “It is true of every finite sum of 1 that it is not equal to 0” and 2. “It is not true of every finite sum of 1 that it is equal to 0” From the example below, it looks like 1. is the intended meaning here, correct?](#)

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*Example 55* (Field of rational numbers). Probably the best known example of a field is the set of rational numbers  $\mathbb{Q}$  together with the usual definition of addition, subtraction, multiplication and division. Since there is no natural number  $n \in \mathbb{N}$ , such that  $\sum_{j=0}^n 1 = 0$  in the set of rational numbers, the characteristic  $\text{char}(\mathbb{Q})$  of the field  $\mathbb{Q}$  is zero. In Sage, rational numbers are called as follows:

```
sage: QQ                                     179
Rational Field                               180
sage: QQ(1/5) # Get an element from the field of rational 181
         numbers
1/5                                           182
sage: QQ(1/5) / QQ(3) # Division           183
1/15                                         184
```

*Example 56* (Field with two elements). It can be shown that, in any field, the neutral element 0 of addition must be different from the neutral element 1 of multiplication, that is,  $0 \neq 1$  always holds in a field. From this, it follows that the smallest field must contain at least two elements. As the following addition and multiplication tables show, there is indeed a field with two elements, which is usually called  $\mathbb{F}_2$ :

Let  $\mathbb{F}_2 := \{0, 1\}$  be a set that contains two elements and let addition and multiplication on  $\mathbb{F}_2$  be defined as follows:



$+$	$0$	$1$	$\cdot$	$0$	$1$
$0$	$0$	$1$	$0$	$0$	$0$
$1$	$1$	$0$	$1$	$0$	$1$

Since  $1 + 1 = 0$  in the field  $\mathbb{F}_2$ , we know that the characteristic of  $\mathbb{F}_2$  is there, that is, we have  $\text{char}(\mathbb{F}_2) = 0$ .

For reasons we will understand better in XXX, Sage defines this field as a so-called Galois field with 2 elements. You can call it in Sage as follows:

```

sage: F2 = GF(2)
sage: F2(1) # Get an element from GF(2)
1
sage: F2(1) + F2(1) # Addition
0
sage: F2(1) / F2(1) # Division
1

```

*Example 57.* Both the real numbers  $\mathbb{R}$  as well as the complex numbers  $\mathbb{C}$  are well known examples of fields.

*Exercise 30.* Consider our remainder class ring  $(\mathbb{Z}_5, +, \cdot)$  and show that it is a field. What is the characteristic of  $\mathbb{Z}_5$ ?

**Prime fields** As we have seen in the various examples of the previous sections, modular arithmetics behaves similarly to the ordinary arithmetics of integers in many ways. This is due to the fact that remainder class sets  $\mathbb{Z}_n$  are commutative rings with units.

However, we have also seen in 36 that, whenever the modulus is a prime number, every remainder class other than the zero class has a modular multiplicative inverse. This is an important observation, since it immediately implies that, in case of a prime number, the remainder class set  $\mathbb{Z}_n$  is not just a ring but actually a **field**. Moreover, since  $\sum_{j=0}^n 1 = 0$  in  $\mathbb{Z}_n$ , we know that those fields have the finite characteristic  $n$ .

To distinguish this important case from arbitrary remainder class rings, we write  $(\mathbb{F}_p, +, \cdot)$  for the field of all remainder classes for a prime number modulus  $p \in \mathbb{P}$  and call it the **prime field** of characteristic  $p$ .

Prime fields are the foundation for many of the contemporary algebra-based cryptographic systems, as they have many desirable properties. One of them is that, since these sets are finite and a prime field of characteristic,  $p$  can be represented on a computer in roughly  $\log_2(p)$  amount of space without precision problems that are unavoidable for computer representations of infinite sets such as rational numbers or integers.

Since prime fields are special cases of remainder class rings, all computations remain the same. Addition and multiplication can be computed by first doing normal integer addition and multiplication, and then taking the remainder modulus  $p$ . Subtraction and division can be computed by adding or multiplying with the additive or the multiplicative inverse, respectively. The additive inverse  $-x$  of a field element  $x \in \mathbb{F}_p$  is given by  $p - x$ , and the multiplicative inverse of  $x \neq 0$  is given by  $x^{p-2}$ , or can be computed using the Extended Euclidean Algorithm.

Note that these computations might not be the fastest to implement on a computer. They are, however, useful in this book, as they are easy to compute for small prime numbers.

*Example 58.* The smallest field is the field  $\mathbb{F}_2$  of characteristic 2 as we have seen in example 56. It is the prime field of the prime number 2.

*Example 59.* To summarize the basic aspects of computation in prime fields, let us consider the prime field  $\mathbb{F}_5$  and simplify the following expression:

$$\left(\frac{2}{3} - 2\right) \cdot 2$$

A first thing to note is that since  $\mathbb{F}_5$  is a field, all rules are identical to the rules we learned in school when we were dealing with rational, real or complex numbers. This means we can use e.g. bracketing (distributivity) or addition as usual:

$$\begin{aligned} \left(\frac{2}{3} - 2\right) \cdot 2 &= \frac{2}{3} \cdot 2 - 2 \cdot 2 && \# \text{ distributive law} \\ &= \frac{2 \cdot 2}{3} - 2 \cdot 2 && 4 \bmod 5 = 4 \\ &= \frac{4}{3} - 4 && \# \text{ multiplicative inverse of 3 is } 3^{5-2} \bmod 5 = 2 \\ &= 4 \cdot 2 - 4 && \# \text{ additive inverse of 4 is } 5 - 4 = 1 \\ &= 4 \cdot 2 + 1 && 8 \bmod 5 = 3 \\ &= 3 + 1 && 4 \bmod 5 = 4 \\ &= 4 \end{aligned}$$

2300 In this computation, we computed the multiplicative inverse of 3 using the identity  $x^{-1} = x^{p-2}$   
 2301 in a prime field. This is impractical for large prime numbers. Recall that another way of  
 2302 computing the multiplicative inverse is the Extended Euclidean Algorithm (see 3.10 on page  
 2303 18). To refresh our memory, the task is to compute  $x^{-1} \cdot 3 + t \cdot 5 = 1$ , but  $t$  is actually irrelevant.  
 2304 We get

k	$r_k$	$x_k^{-1}$	$t_k = (r_k - s_k \cdot a) \operatorname{div} b$
0	3	1	.
1	5	0	.
2	3	1	.
3	2	-1	.
4	1	2	.

2306 So the multiplicative inverse of 3 in  $\mathbb{Z}_5$  is 2, and, indeed, if compute  $3 \cdot 2$ , we get 1 in  $\mathbb{F}_5$ .

2307 **Square Roots** In this part, we deal with square numbers, also called **quadratic residues** and  
 2308 **square roots** in prime fields. This is of particular importance in our studies on elliptic curves,  
 2309 as only square numbers can actually be points on an elliptic curve.

2310 To make the intuition of quadratic residues and roots precise, let  $p \in \mathbb{P}$  be a prime number  
 2311 and  $\mathbb{F}_p$  its associated prime field. Then a number  $x \in \mathbb{F}_p$  is called a **square root** of another  
 2312 number  $y \in \mathbb{F}_p$ , if  $x$  is a solution to the following equation:

$$x^2 = y \tag{4.25}$$

2313 In this case,  $y$  is called a **quadratic residue**. On the other hand, if  $y$  is given and the quadratic  
 2314 equation has no solution  $x$ , we call  $y$  a **quadratic non-residue**. For any  $y \in \mathbb{F}_p$ , we denote the  
 2315 set of all square roots of  $y$  in the prime field  $\mathbb{F}_p$  as follows:

$$\sqrt{y} := \{x \in \mathbb{F}_p \mid x^2 = y\} \tag{4.26}$$

S: are we introducing elliptic curves in section 1 or 2?

If  $y$  is a quadratic non-residue, then  $\sqrt{y} = \emptyset$  (an empty set), and if  $y = 0$ , then  $\sqrt{y} = \{0\}$ .

Informally speaking, quadratic residues are numbers such that we can take the square root of them, while quadratic non-residues are numbers that don't have square roots. The situation therefore parallels the familiar case of integers, where some integers like 4 or 9 have square roots and others like 2 or 3 don't (as integers).

It can be shown that, in any prime field, every non zero element has either no square root or two of them. We adopt the convention to call the smaller one (when interpreted as an integer) as the **positive** square root and the larger one as the **negative**. This makes sense, as the larger one can always be computed as the modulus minus the smaller one, which is the definition of the negative in prime fields.

*Example 60* (Quadratic (Non)-Residues and roots in  $\mathbb{F}_5$ ). Let us consider our example prime field  $\mathbb{F}_5$  again. All square numbers can be found on the main diagonal of the multiplication table in example 13 on page 28. As you can see, in  $\mathbb{F}_5$  only the numbers 0, 1 and 4 have square roots and we get  $\sqrt{0} = \{0\}$ ,  $\sqrt{1} = \{1, 4\}$ ,  $\sqrt{2} = \emptyset$ ,  $\sqrt{3} = \emptyset$  and  $\sqrt{4} = \{2, 3\}$ . The numbers 0, 1 and 4 are therefore quadratic residues, while the numbers 2 and 3 are quadratic non-residues.

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In order to describe whether an element of a prime field is a square number or not, the so-called **Legendre symbol** can sometimes be found in the literature, which is why we will summarize it here:

Let  $p \in \mathbb{P}$  be a prime number and  $y \in \mathbb{F}_p$  an element from the associated prime field. Then the *Legendre symbol* of  $y$  is defined as follows:

$$\left(\frac{y}{p}\right) := \begin{cases} 1 & \text{if } y \text{ has square roots} \\ -1 & \text{if } y \text{ has no square roots} \\ 0 & \text{if } y = 0 \end{cases} \quad (4.27)$$

*Example 61.* Looking at the quadratic residues and non residues in  $\mathbb{F}_5$  from example 13 again, we can deduce the following Legendre symbols, from example XXX.

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$$\left(\frac{0}{5}\right) = 0, \quad \left(\frac{1}{5}\right) = 1, \quad \left(\frac{2}{5}\right) = -1, \quad \left(\frac{3}{5}\right) = -1, \quad \left(\frac{4}{5}\right) = 1.$$

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The Legendre symbol provides a criterion to decide whether or not an element from a prime field has a quadratic root or not. This, however, is not just of theoretical use: The so-called **Euler criterion** provides a compact way to actually compute the Legendre symbol. To see that, let  $p \in \mathbb{P}_{\geq 3}$  be an odd prime number and  $y \in \mathbb{F}_p$ . Then the Legendre symbol can be computed as follows:

$$\left(\frac{y}{p}\right) = y^{\frac{p-1}{2}}. \quad (4.28)$$

*Example 62.* Looking at the quadratic residues and non residues in  $\mathbb{F}_5$  from example 13 again,

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we can compute the following Legendre symbols using the Euler criterion:

$$\begin{aligned}\left(\frac{0}{5}\right) &= 0^{\frac{5-1}{2}} = 0^2 = 0 \\ \left(\frac{1}{5}\right) &= 1^{\frac{5-1}{2}} = 1^2 = 1 \\ \left(\frac{2}{5}\right) &= 2^{\frac{5-1}{2}} = 2^2 = 4 = -1 \\ \left(\frac{3}{5}\right) &= 3^{\frac{5-1}{2}} = 3^2 = 4 = -1 \\ \left(\frac{4}{5}\right) &= 4^{\frac{5-1}{2}} = 4^2 = 1\end{aligned}$$

*Exercise 31.* Consider the prime field  $\mathbb{F}_{13}$ . Find the set of all pairs  $(x, y) \in \mathbb{F}_{13} \times \mathbb{F}_{13}$  that satisfy the equation

$$x^2 + y^2 = 1 + 7 \cdot x^2 \cdot y^2$$

2343 **Exponentiation** TO APPEAR...

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2344 **Hashing into prime fields** An important problem in SNARK development is the ability to  
2345 hash to (various subsets) of elliptic curves. As we will see in XXX, those curves are often  
2346 defined over prime fields, and hashing to a curve might start with hashing to the prime field. It  
2347 is therefore important to understand how to hash into prime fields.

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2348 On pages 51–55, we looked at a few methods of hashing into the residue class rings  $\mathbb{Z}_n$  for  
2349 arbitrary  $n > 1$ . As prime fields are just special instances of those rings, all methods for hashing  
2350 into  $\mathbb{Z}_n$  functions can be used for hashing into prime fields, too.

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2351 **Extension Fields** Prime fields, defined in the previous section, are the basic building blocks  
2352 for cryptography in general and SNARKs in particular.

2353 However, as we will see in XXX so-called **pairing-based** SNARK systems are crucially  
2354 dependent on **group pairings** XXX defined over the group of rational points of elliptic curves.  
2355 For those pairings to be non-trivial, the elliptic curve must not only be defined over a prime  
2356 field, but over a so-called **extension field** of a given prime field.

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2357 We therefore have to understand field extensions. First note that the field  $\mathbb{F}'$  is called an  
2358 **extension** of a field  $\mathbb{F}$  if  $\mathbb{F}$  is a subfield of  $\mathbb{F}'$ , that is,  $\mathbb{F}$  is a subset of  $\mathbb{F}'$  and restricting the  
2359 addition and multiplication laws of  $\mathbb{F}'$  to the subset  $\mathbb{F}$  recovers the appropriate laws of  $\mathbb{F}$ .

2360 Now it can be shown that whenever  $p \in \mathbb{P}$  is a prime and  $m \in \mathbb{N}$  a natural number, then there  
2361 is a field  $\mathbb{F}_{p^m}$  with characteristic  $p$  and  $p^m$  elements such that  $\mathbb{F}_{p^m}$  is an extension field of the  
2362 prime field  $\mathbb{F}_p$ .

2363 Similarly to the way prime fields  $\mathbb{F}_p$  are generated by starting with the ring of integers  
2364 and then dividing by a prime number  $p$  and keeping the remainder, prime field extensions  $\mathbb{F}_{p^m}$   
2365 are generated by starting with the ring  $\mathbb{F}_p[x]$  of polynomials and then dividing them by an  
2366 irreducible polynomial of degree  $m$  and keeping the remainder.

2367 To be more precise, let  $P \in \mathbb{F}_p[x]$  be an irreducible polynomial of degree  $m$  with coefficients  
2368 from the given prime field  $\mathbb{F}_p$ . Then the underlying set  $\mathbb{F}_{p^m}$  of the extension field is given by  
2369 the set of all polynomials with a degree less than  $m$ :

$$\mathbb{F}_{p^m} := \{a_{m-1}x^{m-1} + a_{m-2}x^{m-2} + \dots + a_1x + a_0 \mid a_i \in \mathbb{F}_p\} \quad (4.29)$$

This can be shown to be the set of all remainders when dividing any polynomial  $Q \in \mathbb{F}_p[x]$  by  $P$ , consequently, elements of the extension field are polynomials of degree less than  $m$ . This is analogous to how  $\mathbb{F}_p$  is the set of all remainders when dividing integers by  $p$ .

Addition is inherited from  $\mathbb{F}_p[x]$ , which means that addition on  $\mathbb{F}_{p^m}$  is defined like normal addition of polynomials:

$$+ : \mathbb{F}_{p^m} \times \mathbb{F}_{p^m} \rightarrow \mathbb{F}_{p^m}, \left( \sum_{j=0}^m a_j x^j, \sum_{j=0}^m b_j x^j \right) \mapsto \sum_{j=0}^m (a_j + b_j) x^j \quad (4.30)$$

We can see that the neutral element is (the polynomial) 0, and that the additive inverse is given by the polynomial with all negative coefficients.

Multiplication is inherited from  $\mathbb{F}_p[x]$ , too, but we have to divide the result by our modulus polynomial  $P$  whenever the degree of the resulting polynomial is equal or greater to  $m$ :

$$\cdot : \mathbb{F}_{p^m} \times \mathbb{F}_{p^m} \rightarrow \mathbb{F}_{p^m}, \left( \sum_{j=0}^m a_j x^j, \sum_{j=0}^m b_j x^j \right) \mapsto \left( \sum_{n=0}^{2m} \sum_{i=0}^n a_i b_{n-i} x^n \right) \bmod P \quad (4.31)$$

We can see that the neutral element is (the polynomial) 1. It is, however, not obvious from this definition how the multiplicative inverse looks.

We can easily see from the definition of  $\mathbb{F}_{p^m}$  that the field is of characteristic  $p$ , since the multiplicative neutral element 1 is equivalent to the multiplicative element 1 from the underlying prime field, and hence  $\sum_{j=0}^p 1 = 0$ . Moreover,  $\mathbb{F}_{p^m}$  is finite and contains  $p^m$  many elements, since elements are polynomials of degree  $< m$ , and every coefficient  $a_j$  can have a  $p$  number of different values. In addition, we see that the prime field  $\mathbb{F}_p$  is a subfield of  $\mathbb{F}_{p^m}$  that occurs when we restrict the elements of  $\mathbb{F}_p$  to polynomials of degree zero.

One key point is that the construction of  $\mathbb{F}_{p^m}$  depends on the choice of an irreducible polynomial, and, in fact, different choices will give different multiplication tables, since the remainders from dividing a product by  $P$  will be different.

It can, however, be shown that the fields for different choices of  $P$  are **isomorphic**, which means that there is a one-to-one correspondence between all of them. Consequently, from an abstract point of view, they are the same thing. From an implementations point of view, however, some choices are preferable to others because they allow for faster computations.

To summarize, we have seen that when a prime field  $\mathbb{F}_p$  is given, any field  $\mathbb{F}_{p^m}$  constructed in the above manner is a field extension of  $\mathbb{F}_p$ . To be more general, a field  $\mathbb{F}_{p^{m_2}}$  is a field extension of a field  $\mathbb{F}_{p^{m_1}}$ , if and only if  $m_1$  divides  $m_2$ . From this, we can deduce that, for any given fixed prime number, there are nested sequences of fields whenever the power  $m_j$  divides the power  $m_{j+1}$ , such that  $\mathbb{F}_{p^{m_j}}$  is a subfield of  $\mathbb{F}_{p^{m_{j+1}}}$ :

$$\mathbb{F}_p \subset \mathbb{F}_{p^{m_1}} \subset \cdots \subset \mathbb{F}_{p^{m_k}} \quad (4.32)$$

To get a more intuitive picture of this, we construct an extension field of the prime field  $\mathbb{F}_3$  in the following example, and we can see how  $\mathbb{F}_3$  sits inside that extension field.

*Example 63* (The Extension field  $\mathbb{F}_{32}$ ). In (XXX) we have constructed the prime field  $\mathbb{F}_3$ . In this example, we apply the definition of a field extension (page 63) to construct  $\mathbb{F}_{32}$ . We start by choosing an irreducible polynomial of degree 2 with coefficients in  $\mathbb{F}_3$ . We try  $P(t) = t^2 + 1$ . Possibly the fastest way to show that  $P$  is indeed irreducible is to just insert all elements from

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$\mathbb{F}_3$  to see if the result is ever zero. We compute as follows:

$$P(0) = 0^2 + 1 = 1$$

$$P(1) = 1^2 + 1 = 2$$

$$P(2) = 2^2 + 1 = 1 + 1 = 2$$

This implies that  $P$  is irreducible. The set  $\mathbb{F}_{3^2}$  contains all polynomials of degrees lower than two, with coefficients in  $\mathbb{F}_3$ , which are precisely as listed below:

$$\mathbb{F}_{3^2} = \{0, 1, 2, t, t+1, t+2, 2t, 2t+1, 2t+2\}$$

As expected, our extension field contains 9 elements. Addition is defined as addition of polynomials; for example  $(t+2) + (2t+2) = (1+2)t + (2+2) = 1$ . Doing this computation for all elements gives the following addition table

+	0	1	2	t	t+1	t+2	2t	2t+1	2t+2
0	0	1	2	t	t+1	t+2	2t	2t+1	2t+2
1	1	2	0	t+1	t+2	t	2t+1	2t+2	2t
2	2	0	1	t+2	t	t+1	2t+2	2t	2t+1
t	t	t+1	t+2	2t	2t+1	2t+2	0	1	2
t+1	t+1	t+2	t	2t+1	2t+2	2t	1	2	0
t+2	t+2	t	t+1	2t+2	2t	2t+1	2	0	1
2t	2t	2t+1	2t+2	0	1	2	t	t+1	t+2
2t+1	2t+1	2t+2	2t	1	2	0	t+1	t+2	t
2t+2	2t+2	2t	2t+1	2	0	1	t+2	t	t+1

As we can see, the group  $(\mathbb{F}_3, +)$  is a subgroup of the group  $(\mathbb{F}_{3^2}, +)$ , obtained by only considering the first three rows and columns of this table.

As it was the case in previous examples, we can use the table to deduce the negative of any element from  $\mathbb{F}_{3^2}$ . For example, in  $\mathbb{F}_{3^2}$  we have  $-(2t+1) = t+2$ , since  $(2t+1) + (t+2) = 0$

Multiplication needs a bit more computation, as we first have to multiply the polynomials, and whenever the result has a degree  $\geq 2$ , we have to divide it by  $P$  and keep the remainder. To see how this works, let us compute the product of  $t+2$  and  $2t+2$  in  $\mathbb{F}_{3^2}$ :

$$\begin{aligned}
 (t+2) \cdot (2t+2) &= (2t^2 + 2t + t + 1) \bmod (t^2 + 1) \\
 &= (2t^2 + 1) \bmod (t^2 + 1) & \# 2t^2 + 1 : t^2 + 1 &= 2 + \frac{2}{t^2 + 1} \\
 &= 2
 \end{aligned}$$

This means that the product of  $t+2$  and  $2t+2$  in  $\mathbb{F}_{3^2}$  is 2. Performing this computation for all elements gives the following multiplication table:

·	0	1	2	t	t+1	t+2	2t	2t+1	2t+2
0	0	0	0	0	0	0	0	0	0
1	0	1	2	t	t+1	t+2	2t	2t+1	2t+2
2	0	2	1	2t	2t+2	2t+1	t	t+2	t+1
t	0	t	2t	2	t+2	2t+2	1	t+1	2t+1
t+1	0	t+1	2t+2	t+2	2t	1	2t+1	2	t
t+2	0	t+2	2t+1	2t+2	1	t	t+1	2t	2
2t	0	2t	t	1	2t+1	t+1	2	2t+2	t+2
2t+1	0	2t+1	t+2	t+1	2	2t	2t+2	t	1
2t+2	0	2t+2	t+1	2t+1	t	2	t+2	1	2t

As it was the case in previous examples, we can use the table to deduce the multiplicative inverse of any non-zero element from  $\mathbb{F}_{3^2}$ . For example, in  $\mathbb{F}_{3^2}$  we have  $(2t+1)^{-1} = 2t+2$ , since  $(2t+1) \cdot (2t+2) = 1$ .

From the multiplication table, we can also see that the only quadratic residues in  $\mathbb{F}_{3^2}$  are from the set  $\{0, 1, 2, t, 2t\}$ , with  $\sqrt{0} = \{0\}$ ,  $\sqrt{1} = \{1, 2\}$ ,  $\sqrt{2} = \{t, 2t\}$ ,  $\sqrt{t} = \{t+2, 2t+1\}$  and  $\sqrt{2t} = \{t+1, 2t+2\}$ .

Since  $\mathbb{F}_{3^2}$  is a field, we can solve equations as we would for other fields, (such as rational numbers). To see that, let us find all  $x \in \mathbb{F}_{3^2}$  that solve the quadratic equation  $(t+1)(x^2 + (2t+2)) = 2$ . We compute as follows:

$$\begin{aligned}
 (t+1)(x^2 + (2t+2)) &= 2 && \# 2 \text{ distributive law} \\
 (t+1)x^2 + (t+1)(2t+2) &= 2 \\
 (t+1)x^2 + (t) &= 2 && \# 2 \text{ add the additive inverse of } t \\
 (t+1)x^2 + (t) + (2t) &= (2) + (2t) \\
 (t+1)x^2 &= 2t+2 && \# \text{ multiply with the multiplicative invers of } t+1 \\
 (t+2)(t+1)x^2 &= (t+2)(2t+2) && \# \text{ multiply with the multiplicative invers of } t+1 \\
 x^2 &= 2 && \# 2 \text{ is quadratic residue. Take the roots.} \\
 x &\in \{t, 2t\}
 \end{aligned}$$

Computations in extension fields are arguably on the edge of what can reasonably be done with pen and paper. Fortunately, Sage provides us with a simple way to do the computations.

```

sage: Z3 = GF(3) # prime field
sage: Z3t.<t> = Z3[] # polynomials over Z3
sage: P = Z3t(t^2+1)
sage: P.is_irreducible()
True
sage: F3_2.<t> = GF(3^2, name='t', modulus=P)
sage: F3_2
Finite Field in t of size 3^2
sage: F3_2(t+2)*F3_2(2*t+2) == F3_2(2)
True
sage: F3_2(2*t+2)^(-1) # multiplicative inverse
2*t + 1
sage: # verify our solution to (t+1)(x^2 + (2t+2)) = 2
sage: F3_2(t+1)*(F3_2(t)**2 + F3_2(2*t+2)) == F3_2(2)
True
sage: F3_2(t+1)*(F3_2(2*t)**2 + F3_2(2*t+2)) == F3_2(2)
True

```

*Exercise 32.* Consider the extension field  $\mathbb{F}_{3^2}$  from the previous example and find all pairs of elements  $(x, y) \in \mathbb{F}_{3^2}$ , for which the following equation holds:

$$y^2 = x^3 + 4 \quad (4.33)$$

*Exercise 33.* Show that the polynomial  $P = x^3 + x + 1$  from  $\mathbb{F}_5[x]$  is irreducible. Then consider the extension field  $\mathbb{F}_{5^3}$  defined relative to  $P$ . Compute the multiplicative inverse of  $(2t^2 + 4) \in$



2441  $\mathbb{F}_{5^3}$  using the extended Euclidean algorithm. Then find all  $x \in \mathbb{F}_{5^3}$  that solve the following  
 2442 equation:

$$(2t^2 + 4)(x - (t^2 + 4t + 2)) = (2t + 3) \quad (4.34)$$

2443 **Hashing into extension fields** On page 63, we have seen how to hash into prime fields. As  
 2444 elements of extension fields can be seen as polynomials over prime fields, hashing into extension  
 2445 fields is therefore possible if every coefficient of the polynomial is hashed independently.

check  
reference

## 2446 4.4 Projective Planes

2447 Projective planes are particular geometric constructs defined over a given field. In a sense,  
 2448 projective planes extend the concept of the ordinary Euclidean plane by including “points at  
 2449 infinity.”

2450 Such an inclusion of infinity points makes projective planes particularly useful in the de-  
 2451 scription of elliptic curves, as the description of such a curve in an ordinary plane needs an  
 2452 additional symbol “the point at infinity” to give the set of points on the curve the structure of  
 2453 a group. Translating the curve into projective geometry includes this “point at infinity” more  
 2454 naturally into the set of all points on a projective plane.

2455 To understand the idea of constructing of projective planes, note that in an ordinary Eu-  
 2456 clidean plane, two lines either intersect in a single point or are parallel. In the latter case, both  
 2457 lines are either the same, that is, they intersect in all points, or do not intersect at all. A projec-  
 2458 tive plane can then be thought of as an ordinary plane, but equipped with additional “point at  
 2459 infinity” such that two different lines always intersect in a single point. Parallel lines intersect  
 2460 “at infinity”.

2461 To be more precise, let  $\mathbb{F}$  be a field,  $\mathbb{F}^3 := \mathbb{F} \times \mathbb{F} \times \mathbb{F}$  the set of all three tuples over  $\mathbb{F}$  and  
 2462  $x \in \mathbb{F}^3$  with  $x = (X, Y, Z)$ . Then there is exactly one *line* in  $\mathbb{F}^3$  that intersects both  $(0, 0, 0)$  and  
 2463  $x$ . This line is given as follows:

$$[X : Y : Z] := \{(k \cdot X, k \cdot Y, k \cdot Z) \mid k \in \mathbb{F}\} \quad (4.35)$$

2464 A **point** in the **projective plane** over  $\mathbb{F}$  is defined as such a **line**, and the projective plane is the  
 2465 set of all such points:

$$\mathbb{FP}^2 := \{[X : Y : Z] \mid (X, Y, Z) \in \mathbb{F}^3 \text{ with } (X, Y, Z) \neq (0, 0, 0)\} \quad (4.36)$$

2466 It can be shown that a projective plane over a finite field  $\mathbb{F}_{p^m}$  contains  $p^{2m} + p^m + 1$  number of  
 2467 elements.

2468 To understand why  $[X : Y : Z]$  is called a line, consider the situation where the underlying  
 2469 field  $\mathbb{F}$  is the set of real numbers  $\mathbb{R}$ . In this case,  $\mathbb{R}^3$  can be seen as the three-dimensional space,  
 2470 and  $[X : Y : Z]$  is an ordinary line in this 3-dimensional space that intersects zero and the point  
 2471 with coordinates  $X, Y$  and  $Z$ .

2472 The key observation here is that points in the projective plane are lines in the 3-dimensional  
 2473 space  $\mathbb{F}^3$ . Additionally, for finite fields, the terms **space** and **line** share very little visual simi-  
 2474 larity with their counterparts over the set of real numbers.

2475 It follows from this that points  $[X : Y : Z] \in \mathbb{FP}^2$  are not simply described by fixed co-  
 2476 ordinates  $(X, Y, Z)$ , but by **sets of coordinates**, where two different coordinates  $(X_1, Y_1, Z_1)$   
 2477 and  $(X_2, Y_2, Z_2)$  describe the same point if and only if there is some field element  $k$  such that  
 2478  $(X_1, Y_1, Z_1) = (k \cdot X_2, k \cdot Y_2, k \cdot Z_2)$ . Points  $[X : Y : Z]$  are called **projective coordinates**.



2479 *Notation and Symbols 6* (Projective coordinates). Projective coordinates of the form  $[X : Y : 1]$   
 2480 are descriptions of so-called **affine points**. Projective coordinates of the form  $[X : Y : 0]$  are  
 2481 descriptions of so-called **points at infinity**. In particular, the projective coordinate  $[1 : 0 : 0]$   
 2482 describes the so-called **line at infinity**.

2483 *Example 64.* Consider the field  $\mathbb{F}_3$  from example XXX. As this field only contains three ele-  
 2484 ments, it does not take too much effort to construct its associated projective plane  $\mathbb{F}_3\mathbb{P}^2$ , as we  
 2485 know that it only contains 13 elements.

 add refer-  
ence

To find  $\mathbb{F}_3\mathbb{P}^2$ , we have to compute the set of all lines in  $\mathbb{F}_3 \times \mathbb{F}_3 \times \mathbb{F}_3$  that intersect  $(0, 0, 0)$ . Since those lines are parameterized by tuples  $(x_1, x_2, x_3)$ , we compute as follows:

$$\begin{aligned}
 [0 : 0 : 1] &= \{(k \cdot x_1, k \cdot x_2, k \cdot x_3) \mid k \in \mathbb{F}_3\} = \{(0, 0, 1), (0, 0, 2)\} \\
 [0 : 0 : 2] &= \{(k \cdot x_1, k \cdot x_2, k \cdot x_3) \mid k \in \mathbb{F}_3\} = \{(0, 0, 2), (0, 0, 1)\} = [0 : 0 : 1] \\
 [0 : 1 : 0] &= \{(k \cdot x_1, k \cdot x_2, k \cdot x_3) \mid k \in \mathbb{F}_3\} = \{(0, 1, 0), (0, 2, 0)\} \\
 [0 : 1 : 1] &= \{(k \cdot x_1, k \cdot x_2, k \cdot x_3) \mid k \in \mathbb{F}_3\} = \{(0, 1, 1), (0, 2, 2)\} \\
 [0 : 1 : 2] &= \{(k \cdot x_1, k \cdot x_2, k \cdot x_3) \mid k \in \mathbb{F}_3\} = \{(0, 1, 2), (0, 2, 1)\} \\
 [0 : 2 : 0] &= \{(k \cdot x_1, k \cdot x_2, k \cdot x_3) \mid k \in \mathbb{F}_3\} = \{(0, 2, 0), (0, 1, 0)\} = [0 : 1 : 0] \\
 [0 : 2 : 1] &= \{(k \cdot x_1, k \cdot x_2, k \cdot x_3) \mid k \in \mathbb{F}_3\} = \{(0, 2, 1), (0, 1, 2)\} = [0 : 1 : 2] \\
 [0 : 2 : 2] &= \{(k \cdot x_1, k \cdot x_2, k \cdot x_3) \mid k \in \mathbb{F}_3\} = \{(0, 2, 2), (0, 1, 1)\} = [0 : 1 : 1] \\
 [1 : 0 : 0] &= \{(k \cdot x_1, k \cdot x_2, k \cdot x_3) \mid k \in \mathbb{F}_3\} = \{(1, 0, 0), (2, 0, 0)\} \\
 [1 : 0 : 1] &= \{(k \cdot x_1, k \cdot x_2, k \cdot x_3) \mid k \in \mathbb{F}_3\} = \{(1, 0, 1), (2, 0, 2)\} \\
 [1 : 0 : 2] &= \{(k \cdot x_1, k \cdot x_2, k \cdot x_3) \mid k \in \mathbb{F}_3\} = \{(1, 0, 2), (2, 0, 1)\} \\
 [1 : 1 : 0] &= \{(k \cdot x_1, k \cdot x_2, k \cdot x_3) \mid k \in \mathbb{F}_3\} = \{(1, 1, 0), (2, 2, 0)\} \\
 [1 : 1 : 1] &= \{(k \cdot x_1, k \cdot x_2, k \cdot x_3) \mid k \in \mathbb{F}_3\} = \{(1, 1, 1), (2, 2, 2)\} \\
 [1 : 1 : 2] &= \{(k \cdot x_1, k \cdot x_2, k \cdot x_3) \mid k \in \mathbb{F}_3\} = \{(1, 1, 2), (2, 2, 1)\} \\
 [1 : 2 : 0] &= \{(k \cdot x_1, k \cdot x_2, k \cdot x_3) \mid k \in \mathbb{F}_3\} = \{(1, 2, 0), (2, 1, 0)\} \\
 [1 : 2 : 1] &= \{(k \cdot x_1, k \cdot x_2, k \cdot x_3) \mid k \in \mathbb{F}_3\} = \{(1, 2, 1), (2, 1, 2)\} \\
 [1 : 2 : 2] &= \{(k \cdot x_1, k \cdot x_2, k \cdot x_3) \mid k \in \mathbb{F}_3\} = \{(1, 2, 2), (2, 1, 1)\} \\
 [2 : 0 : 0] &= \{(k \cdot x_1, k \cdot x_2, k \cdot x_3) \mid k \in \mathbb{F}_3\} = \{(2, 0, 0), (1, 0, 0)\} = [1 : 0 : 0] \\
 [2 : 0 : 1] &= \{(k \cdot x_1, k \cdot x_2, k \cdot x_3) \mid k \in \mathbb{F}_3\} = \{(2, 0, 1), (1, 0, 2)\} = [1 : 0 : 2] \\
 [2 : 0 : 2] &= \{(k \cdot x_1, k \cdot x_2, k \cdot x_3) \mid k \in \mathbb{F}_3\} = \{(2, 0, 2), (1, 0, 1)\} = [1 : 0 : 1] \\
 [2 : 1 : 0] &= \{(k \cdot x_1, k \cdot x_2, k \cdot x_3) \mid k \in \mathbb{F}_3\} = \{(2, 1, 0), (1, 2, 0)\} = [1 : 2 : 0] \\
 [2 : 1 : 1] &= \{(k \cdot x_1, k \cdot x_2, k \cdot x_3) \mid k \in \mathbb{F}_3\} = \{(2, 1, 1), (1, 2, 2)\} = [1 : 2 : 2] \\
 [2 : 1 : 2] &= \{(k \cdot x_1, k \cdot x_2, k \cdot x_3) \mid k \in \mathbb{F}_3\} = \{(2, 1, 2), (1, 2, 1)\} = [1 : 2 : 1] \\
 [2 : 2 : 0] &= \{(k \cdot x_1, k \cdot x_2, k \cdot x_3) \mid k \in \mathbb{F}_3\} = \{(2, 2, 0), (1, 1, 0)\} = [1 : 1 : 0] \\
 [2 : 2 : 1] &= \{(k \cdot x_1, k \cdot x_2, k \cdot x_3) \mid k \in \mathbb{F}_3\} = \{(2, 2, 1), (1, 1, 2)\} = [1 : 1 : 2] \\
 [2 : 2 : 2] &= \{(k \cdot x_1, k \cdot x_2, k \cdot x_3) \mid k \in \mathbb{F}_3\} = \{(2, 2, 2), (1, 1, 1)\} = [1 : 1 : 1]
 \end{aligned}$$

These lines define the 13 points in the projective plane  $\mathbb{F}_3\mathbb{P}$ :

$$\begin{aligned}
 \mathbb{F}_3\mathbb{P} = \{ & [0 : 0 : 1], [0 : 1 : 0], [0 : 1 : 1], [0 : 1 : 2], [1 : 0 : 0], [1 : 0 : 1], \\
 & [1 : 0 : 2], [1 : 1 : 0], [1 : 1 : 1], [1 : 1 : 2], [1 : 2 : 0], [1 : 2 : 1], [1 : 2 : 2] \}
 \end{aligned}$$

2486 This projective plane contains 9 affine points, three points at infinity and one line at infinity.

2487 To understand the ambiguity in projective coordinates a bit better, let us consider the point  
2488  $[1 : 2 : 2]$ . As this point in the projective plane is a line in  $\mathbb{F}_3^3$ , it has the projective coordinates  
2489  $(1, 2, 2)$  as well as  $(2, 1, 1)$ , since the former coordinate gives the latter when multiplied in  $\mathbb{F}_3$   
2490 by the factor 2. In addition, note that, for the same reasons, the points  $[1 : 2 : 2]$  and  $[2 : 1 : 1]$   
2491 are the same, since their underlying sets are equal.

2492 *Exercise 34.* Construct the so-called **Fano plane**, that is, the projective plane over the finite  
2493 field  $\mathbb{F}_2$ .

# Chapter 5

## Elliptic Curves

Generally speaking, elliptic curves are “curves” defined in geometric planes like the Euclidean or the projective plane over some given field. One of the key features of elliptic curves over finite fields from the point of view of cryptography is that their set of points has a group law such that the resulting group is finite and cyclic, and it is believed that the discrete logarithm problem on these groups is hard.

A special class of elliptic curves are so-called **pairing-friendly curves**, which have a notation of a group pairing as defined in XXX. This pairing has cryptographically advantageous properties. Those curve are useful in the development of SNARKs, since they allow to compute so-called RICS-satisfiability “in the exponent” **MIRCO: (THIS HAS TO BE REWRITTEN WITH WAY MORE DETAIL)**

In this chapter, we introduce epileptic curves as they are used in pairing-based approaches to the construction of SNARKs. The elliptic curves we consider are all defined over prime fields or prime field extensions and the reader should be familiar with the contend of the previous section on those fields.

In its most generality elliptic curves are defined as a smooth projective curve of genus 1 defined over some field  $\mathbb{F}$  with a distinguished  $\mathbb{F}$ -rational point, but this definition is not very useful for the introductory character of this book. We will therefore look at 3 more practical definitions in the following sections, by introducing Weierstraß, Montgomery and Edwards curves. All of them are widely used in cryptography, and understanding them is crucial to being able to follow the rest of this book.

### 5.1 Elliptic Curve Arithmetics

#### 5.1.1 Short Weierstraß Curves

In this section, we introduce **short Weierstraß** curves, which are the most general types of curves over finite fields of characteristic greater than 3.

We start with their representation in **affine space**. This representation has the advantage that affine points correspond to pairs of numbers, which makes it more accessible for beginners. However, it has the disadvantage that a special “point at infinity”, that is not a point on the curve, is necessary to describe the group structure. We introduce the elliptic curve group law and describe elliptic curve scalar multiplication, which is an instantiation of the exponential map from general cyclic groups.

Then we look at the projective representation of short Weierstraß curves. This has the advantage that no special symbol is necessary to represent the point at infinity but comes with

TODO:  
Elliptic  
Curve  
asymmet-  
ric cryp-  
tography  
examples.  
Private  
key, gen-  
erator,  
public  
key.

add refer-  
ence

maybe re-  
move this  
sentence?

affine  
space

the drawback that projective points are classes of numbers, which might be a bit unusual for a beginner.

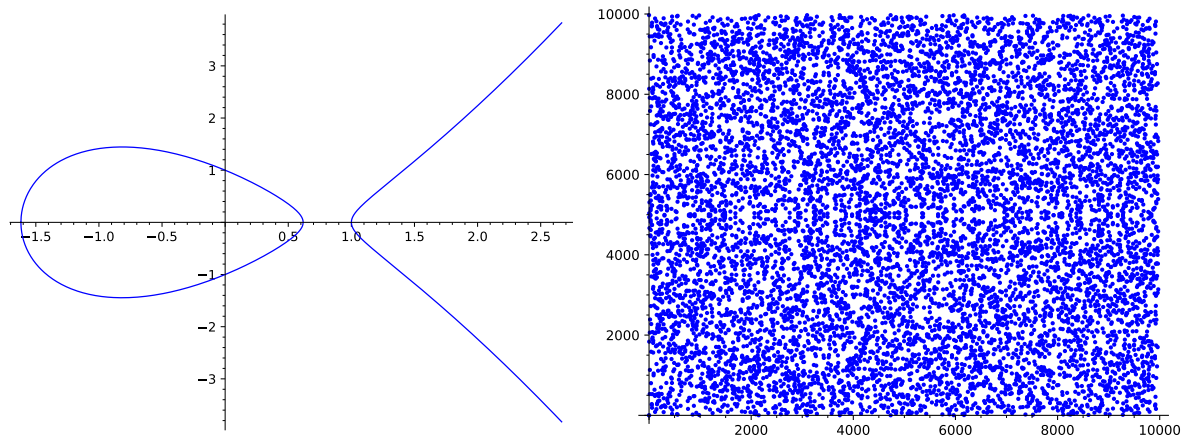
We finish this section with an explicit equivalence that transforms affine representations into projective ones and vice versa.

**Affine short Weierstraß form** Probably the least abstract and most straight-forward way to introduce elliptic curves for non-mathematicians and beginners is the so-called affine representation of a short Weierstraß curve. To see what this is, let  $\mathbb{F}$  be a finite field of order  $q$  and  $a, b \in \mathbb{F}$  two field elements such that  $4a^3 + 27b^2 \bmod q \neq 0$ . Then a **short Weierstraß elliptic curve**  $E(\mathbb{F})$  over  $\mathbb{F}$  in its affine representation is the set of all pairs of field elements  $(x, y) \in \mathbb{F} \times \mathbb{F}$  that satisfy the short Weierstraß cubic equation  $y^2 = x^3 + a \cdot x + b$ , together with a distinguished symbol  $\mathcal{O}$ , called the **point at infinity**:

$$E(\mathbb{F}) = \{(x, y) \in \mathbb{F} \times \mathbb{F} \mid y^2 = x^3 + a \cdot x + b\} \cup \{\mathcal{O}\} \quad (5.1)$$

*Notation and Symbols 7.* In the literature, the set  $E(\mathbb{F})$ , which includes the symbol  $\mathcal{O}$ , is often called the set of **rational points** of the elliptic curve, in which case the curve itself is usually written as  $E/\mathbb{F}$ . However, in what follows, we will frequently identify an elliptic curve with its set of rational points and therefore use the notation  $E(\mathbb{F})$  instead. This is possible in our case, since we only the group structure of the curve in consideration is relevant for us.

The term “curve” is used here because, in the ordinary 2 dimensional plane  $\mathbb{R}^2$ , the set of all points  $(x, y)$  that satisfy  $y^2 = x^3 + a \cdot x + b$  looks like a curve. We should note however that visualizing elliptic curves over finite fields as “curves” has its limitations, and we will therefore not stress the geometric picture too much, but focus on the computational properties instead. To understand the visual difference, consider the following two elliptic curves:



Both elliptic curves are defined by the same short Weierstraß equation  $y^2 = x^3 - 2x + 1$ , but the first curve is defined in the real affine plane  $\mathbb{R}^2$ , that is, the pair  $(x, y)$  contains real numbers, while the second one is defined in the affine plane  $\mathbb{F}_{9973}^2$ , which means that both  $x$  and  $y$  are from the prime field  $\mathbb{F}_{9973}$ . Every blue dot represents a pair  $(x, y)$ , that is a solution to  $y^2 = x^3 - 2x + 1$ . As we can see, the second curve hardly looks like a geometric structure one would naturally call a curve. This shows that our geometric intuitions from  $\mathbb{R}^2$  are obfuscated in curves over finite fields.

The identity  $6 \cdot (4a^3 + 27b^2) \bmod q \neq 0$  ensures that the curve is non-singular, which basically means that the curve has no **cusps** or **self-intersections**.

cusps

self-intersections

Throughout this book, the reader is advised to do as many computations in a pen-and-paper fashion as possible, as this helps getting a deeper understanding of the details. However, when dealing with elliptic curves, computations can quickly become cumbersome and tedious, and one might get lost in the details. Fortunately, Sage is very helpful in dealing with elliptic curves. This book introduces the reader to the great elliptic curve capabilities of Sage. The following snippet shows a way to define elliptic curves and work with them in Sage:

```

sage: F5 = GF(5) # define the base field
sage: a = F5(2) # parameter a
sage: b = F5(4) # parameter b
sage: # check non-singularity
sage: F5(6)*(F5(4)*a^3+F5(27)*b^2) != F5(0)
True
sage: # short Weierstrass curve
sage: E = EllipticCurve(F5,[a,b]) # y^2 == x^3 + ax +b
sage: P = E(0,2) # 2^2 == 0^3 + 2*0 + 4
sage: P.xy() # affine coordinates
(0, 2)
sage: INF = E(0) # point at infinity
sage: try: # point at infinity has no affine coordinates
.....:     INF.xy()
.....: except ZeroDivisionError:
.....:     pass
sage: P = E.plot() # create a plotted version

```

The following three examples give a more practical understanding of what an elliptic curve is and how we can compute it. The reader is advised to read them carefully, and ideally, to also carry out the computation themselves. We will repeatedly build on these examples in this chapter, and use the second example throughout this book.

*Example 65.* To provide the reader with an example of a small elliptic curve where all computation can be done with pen and paper, consider the prime field  $\mathbb{F}_5$  from example 59 (page 61), which should be quite familiar to readers who had worked through the examples and exercises in the previous chapter.

To define an elliptic curve over  $\mathbb{F}_5$ , we have to choose two numbers  $a$  and  $b$  from that field. Assuming we choose  $a = 1$  and  $b = 1$  then  $4a^3 + 27b^2 \equiv 1 \pmod{5}$  from which follows that the corresponding elliptic curve  $E_1(\mathbb{F}_5)$  is given by the set of all pairs  $(x, y)$  from  $\mathbb{F}_5$  that satisfy the equation  $y^2 = x^3 + x + 1$ , together with the special symbol  $\mathcal{O}$ , which represents the “point at infinity”.

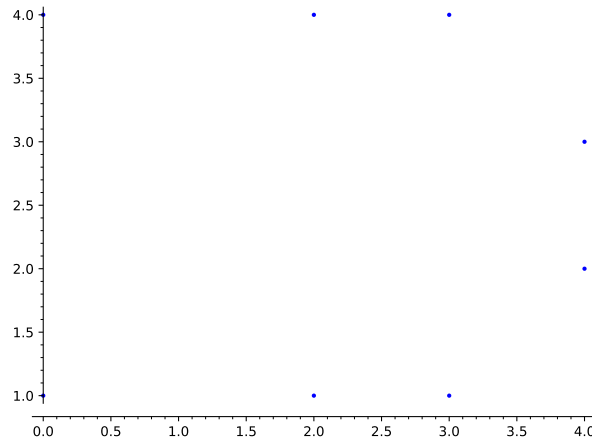
To get a better understanding of that curve, observe that if we choose arbitrarily the pair  $(x, y) = (1, 1)$ , we see that  $1^2 \neq 1^3 + 1 + 1$  and hence  $(1, 1)$  is not an element of the curve  $E_1(\mathbb{F}_5)$ . On the other hand choosing for example  $(x, y) = (2, 1)$  gives  $1^2 = 2^3 + 2 + 1$  and hence the pair  $(2, 1)$  is an element of  $E_1(\mathbb{F}_5)$  (Remember that all computations are done in modulo 5 arithmetics).

Now since the set  $\mathbb{F}_5 \times \mathbb{F}_5$  of all pairs  $(x, y)$  from  $\mathbb{F}_5$  contains only  $5 \cdot 5 = 25$  pairs, we can compute the curve, by just inserting every possible pair  $(x, y)$  into the short Weierstraß equation  $y^2 = x^3 + x + 1$ . If the equation holds, the pair is a curve point, if not that means that the point is not on the curve. Combining the result of this computation with the point at infinity gives the curve as follows:

$$E_1(\mathbb{F}_5) = \{\mathcal{O}, (0, 1), (2, 1), (3, 1), (4, 2), (4, 3), (0, 4), (2, 4), (3, 4)\}$$

check  
reference

2600 This means that our elliptic curve is a set of 9 elements, 8 of which are pairs of numbers and  
 2601 one special symbol  $\mathcal{O}$ . Visualizing  $E1$  gives the following plot:



2602

2603 In the development of SNARKs, it is sometimes necessary to do elliptic curve cryptography  
 2604 “in a circuit”, which basically means that the elliptic curves need to be implemented in a certain  
 2605 SNARK-friendly way. We will look at what this means in chapter 7. To be able to do this  
 2606 efficiently, it is desirable to have curves with special properties. The following example is a  
 2607 pen-and-paper version of such a curve, called **Baby-jubjub**, which resembles cryptographically  
 2608 secure curves extensively used in real-world SNARKs. The interested reader is advised to study  
 2609 this example carefully, as we will use it and build on it in various places throughout the book. S:  
 2610 I feel like a lot of people won’t get the Lewis Carroll reference unless we make it more explicit.  
 2611 M: The term Baby-JubJub is actually the name of a curve used in zCash and Ethereum a lot.  
 2612 IDK why they choosed that name.

check  
reference

jubjub

2613 *Example 66 (Tiny-Jubjub).* Consider the prime field  $\mathbb{F}_{13}$  from exercise 4.3 (page 63. If we  
 2614 choose  $a = 8$  and  $b = 8$ , then  $4a^3 + 27b^2 \equiv 6 \pmod{13}$  and the corresponding elliptic curve  
 2615 is given by all pairs  $(x, y)$  from  $\mathbb{F}_{13}$  such that  $y^2 = x^3 + 8x + 8$  holds. We call this curve the  
 2616 **Tiny-jubjub** curve, or  $TJJ\_13$  for short.

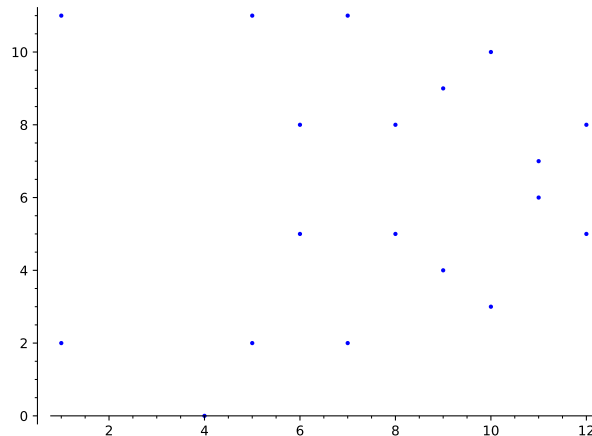
check  
reference

2617 Now, since the set  $\mathbb{F}_{13} \times \mathbb{F}_{13}$  of all pairs  $(x, y)$  from  $\mathbb{F}_{13}$  contains only  $13 \cdot 13 = 169$  pairs,  
 2618 we can compute the curve by just inserting every possible pair  $(x, y)$  into the short Weierstraß  
 2619 equation  $y^2 = x^3 + 8x + 8$ . We get the following result:

$$TJJ\_13 = \{\mathcal{O}, (1, 2), (1, 11), (4, 0), (5, 2), (5, 11), (6, 5), (6, 8), (7, 2), (7, 11), (8, 5), (8, 8), \\ (9, 4), (9, 9), (10, 3), (10, 10), (11, 6), (11, 7), (12, 5), (12, 8)\} \quad (5.2)$$

2620 As we can see, the curve consists of 20 points; 19 points from the affine plane and the point at  
 2621 infinity. To get a visual impression of the  $TJJ\_13$  curve, we might plot all of its points (except  
 2622 the point at infinity) in the  $\mathbb{F}_{13} \times \mathbb{F}_{13}$  affine plane. We get the following plot:

affine  
plane



As we will see in what follows, this curve is rather special, as it is possible to represent it in two alternative forms called the **Montgomery** and the **twisted Edwards form** (See sections 5.1.2 and 5.1.3, respectively).

check  
reference

Now that we have seen two pen-and-paper friendly elliptic curves, let us look at a curve, that is used in actual cryptography. Cryptographically secure elliptic curves are not **qualitatively** different from the curves we looked at so far, but the prime number modulus of their prime field is much larger. Typical examples use prime numbers that have binary representations in the magnitude of more than double the size of the desired security level. If, for example, a security of 128 bits is desired, a prime modulus of binary size  $\geq 256$  is chosen. The following example provides such a curve.

check  
reference

*Example 67* (Bitcoin's Secp256k1 curve). To give an example of a real-world, cryptographically secure curve, let us look at curve Secp256k1, which is famous for being used in the public key cryptography of Bitcoin. The prime field  $\mathbb{F}_p$  of Secp256k1 is defined by the following prime number:

$$p = 115792089237316195423570985008687907853269984665640564039457584007908834671663$$

The binary representation of this number needs 256 bits, which implies that the prime field  $\mathbb{F}_p$  contains approximately  $2^{256}$  many elements, which is considered quite large. To get a better impression of how large the base field is, consider that the number  $2^{256}$  is approximately in the same order of magnitude as the estimated number of atoms in the observable universe.

The curve Secp256k1 is defined by the parameters  $a, b \in \mathbb{F}_p$  with  $a = 0$  and  $b = 7$ . Since  $4 \cdot 0^3 + 27 \cdot 7^2 \bmod p = 1323$ , those parameters indeed define an elliptic curve given as follows:

$$\text{Secp256k1} = \{(x, y) \in \mathbb{F}_p \times \mathbb{F}_p \mid y^2 = x^3 + 7\}$$

Clearly, the Secp256k1 curve is too large to do computations by hand, since it can be shown that the number of its elements is a prime number  $r$  that also has a binary representation of 256 bits:

$$r = 115792089237316195423570985008687907852837564279074904382605163141518161494337$$

Cryptographically secure elliptic curves are therefore not useful in pen-and-paper computations. Fortunately, Sage handles large curves efficiently:

```
sage: p = 1157920892373161954235709850086879078532699846656405 226
      64039457584007908834671663
```



```

2642 sage: # Hexadecimal representation
2643 sage: p.str(16)
2644 ffffffffffffffffffffffffffffffffffffffffffffffffffefffffc
2645 2f
2646 sage: p.is_prime()
2647 True
2648 sage: p.nbits()
2649 256
2650 sage: Fp = GF(p)
2651 sage: Secp256k1 = EllipticCurve(Fp, [0, 7])
2652 sage: r = Secp256k1.order() # number of elements
2653 sage: r.str(16)
2654 ffffffffffffffffffffffffffffffebaaedce6af48a03bbfd25e8cd03641
2655 41
2656 sage: r.is_prime()
2657 True
2658 sage: r.nbits()
2659 256

```

2660 *Exercise 35.* Look up the definition of curve BLS12-381, implement it in Sage and compute its  
 2661 order.

2662 **Affine compressed representation** As we have seen in example 67, cryptographically secure  
 2663 elliptic curves are defined over large prime fields, where elements of those fields typically need  
 2664 more than 255 bits of storage on a computer. Since elliptic curve points consist of pairs of those  
 2665 field elements, they need double that amount of storage.

2666 However, we can reduce the amount of space needed to represent a curve point by using  
 2667 a technique called **point compression**. Note that, up to a sign, the y coordinate of a curve  
 2668 point can be computed from the  $x$  coordinate by simply inserting  $x$  into the Weierstraß equation  
 2669 and then computing the roots of the result. This gives two results, and it means that we can  
 2670 represent a curve point in **compressed form** by simply storing the  $x$  coordinate together with  
 2671 a single sign bit only, the latter of which deterministically decides which of the two roots to  
 2672 choose. One convention could be to always choose the root closer to 0 when the sign bit is 0,  
 2673 and the root closer to the order of  $\mathbb{F}$  when the sign bit is 1. In case the  $y$  coordinate is zero, both  
 2674 sign bits give the same result.

*Example 68 (Tiny-jubjub).* To understand the concept of compressed curve points a bit better,  
 consider the *TJJ\_13* curve from example 66 again. Since this curve is defined over the prime  
 field  $\mathbb{F}_{13}$ , and numbers between 0 and 13 need approximately 4 bits to be represented, each  
*TJJ\_13* point on this curve needs 8 bits of storage in uncompressed form. The following set  
 represents the uncompressed form of the points on this curve:

$$TJJ\_13 = \{\mathcal{O}, (1, 2), (1, 11), (4, 0), (5, 2), (5, 11), (6, 5), (6, 8), (7, 2), (7, 11), \\ (8, 5), (8, 8), (9, 4), (9, 9), (10, 3), (10, 10), (11, 6), (11, 7), (12, 5), (12, 8)\}$$

Using the technique of point compression, we can reduce the bits needed to represent the points  
 on this curve to 5 per point. To achieve this, we can replace the  $y$  coordinate in each  $(x, y)$  pair  
 by a sign bit indicating whether or not  $y$  is closer to 0 or to 13. As a result  $y$  values in the range  
 $[0, \dots, 6]$  will have the sign bit 0, while  $y$ -values in the range  $[7, \dots, 12]$  will have the sign bit 1.

check  
reference

sign

more ex-  
planation  
of what  
the sign is

check  
reference



Applying this to the points in *TJJ\_13* gives the compressed representation as follows:

$$TJJ_{13} = \{\mathcal{O}, (1,0), (1,1), (4,0), (5,0), (5,1), (6,0), (6,1), (7,0), (7,1), \\ (8,0), (8,1), (9,0), (9,1), (10,0), (10,1), (11,0), (11,1), (12,0), (12,1)\}$$

2675 Note that the numbers  $7, \dots, 12$  are the negatives (additive inverses) of the numbers  $1, \dots, 6$  in  
 2676 modular 13 arithmetics and that  $-0 = 0$ . Calling the compression bit a “sign bit” therefore  
 2677 makes sense.

2678 To recover the uncompressed counterpart of, say, the compressed point  $(5, 1)$ , we insert the  
 2679  $x$  coordinate 5 into the Weierstraß equation and get  $y^2 = 5^3 + 8 \cdot 5 + 8 = 4$ . As expected, 4 is a  
 2680 quadratic residue in  $\mathbb{F}_{13}$  with roots  $\sqrt{4} = \{2, 11\}$ . Since the sign bit of the point is 1, we have  
 2681 to choose the root closer to the modulus 13, which is 11. The uncompressed point is therefore  
 2682  $(5, 11)$ .

2683 Looking at the previous examples, the compression rate does not look very impressive.  
 2684 However, looking at the real-life example of the Secp256k1 curve shows that compression is  
 2685 has significant practical advantages.

2686 *Example 69.* Consider the Secp256k1 curve from example 67 again. The following code in-  
 2687 vokes Sage to generate a random affine curve point, then applies our compression method to it:

```

2688
2689 sage: P = Secp256k1.random_point().xy()
2690 sage: P
2691 (8336069568093324890567563480087255358308762674236378487270893
2692 5158363574021275, 45281020659746011620809576891522057974164
2693 397548442100074006135181816198851495)
2694 sage: # uncompressed affine point size
2695 sage: ZZ(P[0]).nbits()+ZZ(P[1]).nbits()
2696 511
2697 sage: # compute the compression
2698 sage: if P[1] > Fp(-1)/Fp(2):
2699 .....:     PARITY = 1
2700 .....: else:
2701 .....:     PARITY = 0
2702 sage: PCOMPRESSED = [P[0], PARITY]
2703 sage: PCOMPRESSED
2704 [8336069568093324890567563480087255358308762674236378487270893
2705 5158363574021275, 0]
2706 sage: # compressed affine point size
2707 sage: ZZ(PCOMPRESSED[0]).nbits()+ZZ(PCOMPRESSED[1]).nbits()
2708 256
2709

```

2710 **Affine group law** One of the key properties of an elliptic curve is that it is possible to define  
 2711 a group law on the set of its rational points such that the point at infinity serves as the neutral  
 2712 element and inverses are reflections on the  $x$ -axis.

2713 The origin of this law can be understood in a geometric picture and is known as the **chord-**  
 2714 **and-tangent rule**. In the affine representation of a short Weierstraß curve, the rule can be  
 2715 described in the following way:

S: I don't  
follow  
this at all

check  
reference

add ex-  
planation  
of how  
this shows  
what we  
claim

### Definition 5.1.1.1. Chord-and-tangent rule

- (Point at infinity) We define the point at infinity  $\mathcal{O}$  as the neutral element of addition, that is, we define  $P + \mathcal{O} = P$  for all points  $P \in E(\mathbb{F})$ .
- (Point addition) Let  $P, Q \in E(\mathbb{F}) \setminus \{\mathcal{O}\}$  with  $P \neq Q$  be two distinct points on an elliptic curve, neither of them the point at infinity. The sum of  $P$  and  $Q$  is defined as follows:  
Consider the line  $l$  which intersects the curve in  $P$  and  $Q$ . If  $l$  intersects the elliptic curve at a third point  $R'$ , define the sum  $R = P \oplus Q$  of  $P$  and  $Q$  as the reflection of  $R'$  at the  $x$ -axis. If the line  $l$  does not intersect the curve at a third point, define the sum to be the point at infinity  $\mathcal{O}$ . It can be shown that no such **chord line** will intersect the curve in more than three points, so addition is not ambiguous.
- (Point doubling) Let  $P \in E(\mathbb{F}) \setminus \{\mathcal{O}\}$  be a point on an elliptic curve, that is not the point at infinity. The sum of  $P$  with itself (the doubling of  $P$ ) is defined as follows:  
Consider the line which is **tangential** to the elliptic curve at  $P$ . If this line intersects the elliptic curve at a second point  $R'$ , the sum  $2P = P + P$  is the reflection of  $R'$  at the  $x$ -axis. If it does not intersect the curve at a third, point define the sum to be the point at infinity  $\mathcal{O}$ . It can be shown that no such **tangent line** will intersect the curve in more than two points, so addition is not ambiguous.

It can be shown that the points of an elliptic curve form a commutative group with respect to the tangent-and-chord rule such that  $\mathcal{O}$  acts the neutral element, and the inverse of any element  $P \in E(\mathbb{F})$  is the reflection of  $P$  on the  $x$ -axis.

To translate the geometric description into algebraic equations, first observe that, for any two given curve points  $(x_1, y_1), (x_2, y_2) \in E(\mathbb{F})$ , it can be shown that the identity  $x_1 = x_2$  implies  $y_2 = \pm y_1$ , which shows that the following rules are a complete description of the affine addition law.

### Definition 5.1.1.2. Chord-and-tangent rule: algebraic equations

- (Neutral element) The point at infinity  $\mathcal{O}$  is the neutral element.
- (Additive inverse ) The additive inverse of  $\mathcal{O}$  is  $\mathcal{O}$ . For any other curve point  $(x, y) \in E(\mathbb{F}) \setminus \{\mathcal{O}\}$ , the additive inverse is given by  $(x, -y)$ .
- (Addition rule) For any two curve points  $P, Q \in E(\mathbb{F})$ , addition is defined by one of the following three cases:
  1. (Adding the neutral element) If  $Q = \mathcal{O}$ , then the sum is defined as  $P \oplus Q = P$ .
  2. (Adding inverse elements) If  $P = (x, y)$  and  $Q = (x, -y)$ , the sum is defined as  $P \oplus Q = \mathcal{O}$ .
  3. (Adding non-self-inverse equal points) If  $P = (x, y)$  and  $Q = (x, y)$  with  $y \neq 0$ , the sum  $2P = (x', y')$  is defined as follows: **We only referred to  $P$  in the definition of point doubling above so  $Q$  seems a bit confusing here even though it's defined as equal to  $P$**
  4. (Adding non-inverse different points) If  $P = (x_1, y_1)$  and  $Q = (x_2, y_2)$  such that  $x_1 \neq x_2$ , the sum  $R = P + Q$  with  $R = (x_3, y_3)$  is defined as follows:

$$x' = \left( \frac{3x^2 + a}{2y} \right)^2 - 2x \quad , \quad y' = \left( \frac{3x^2 + a}{2y} \right)^2 (x - x') - y$$

$$x_3 = \left( \frac{y_2 - y_1}{x_2 - x_1} \right)^2 - x_1 - x_2 \quad , \quad y_3 = \left( \frac{y_2 - y_1}{x_2 - x_1} \right) (x_1 - x_3) - y_1$$

2749 Note that short Weierstraß curve points  $P$  with  $P = (x, 0)$  are inverses of themselves, which  
 2750 implies  $2P = \mathcal{O}$  in this case.

2751 *Notation and Symbols 8.* Let  $\mathbb{F}$  be a field and  $E(\mathbb{F})$  be an elliptic curve over  $\mathbb{F}$ . We write  $\oplus$  for  
 2752 the group law on  $E(\mathbb{F})$  and  $(E(\mathbb{F}), \oplus)$  for the group of rational points.

2753 As we can see, it is very efficient to compute inverses on elliptic curves. However, com-  
 2754 puting the addition of elliptic curve points in the affine representation needs to consider many  
 2755 cases and involves extensive finite field divisions. As we will see in the next paragraph, this can  
 2756 be simplified in projective coordinates. where?

2757 To get some practical impression of how the group law on an elliptic curve is computed,  
 2758 let's look at some actual cases:

*Example 70.* Consider the elliptic curve  $E_1(\mathbb{F}_5)$  from example 65 again. As we have seen, the  
 set of rational points contains 9 elements: check  
reference

$$E_1(\mathbb{F}_5) = \{\mathcal{O}, (0, 1), (2, 1), (3, 1), (4, 2), (4, 3), (0, 4), (2, 4), (3, 4)\}$$

2759 We know that this set defines a group, so we can add any two elements from  $E_1(\mathbb{F}_5)$  to get a  
 2760 third element.

To give an example, consider the elements  $(0, 1)$  and  $(4, 2)$ . Neither of these elements is  
 the neutral element  $\mathcal{O}$ , and since, the  $x$  coordinate of  $(0, 1)$  is different from the  $x$  coordinate of  
 $(4, 2)$ , we know that we have to use the chord rule, that is, rule number 4 from definition 5.1.1.2,  
 to compute the sum  $(0, 1) \oplus (4, 2)$ : check  
reference

$$\begin{aligned} x_3 &= \left( \frac{y_2 - y_1}{x_2 - x_1} \right)^2 - x_1 - x_2 && \# \text{ insert points} \\ &= \left( \frac{2 - 1}{4 - 0} \right)^2 - 0 - 4 && \# \text{ simplify in } \mathbb{F}_5 \\ &= \left( \frac{1}{4} \right)^2 + 1 = 4^2 + 1 = 1 + 1 = 2 \end{aligned}$$

$$\begin{aligned} y_3 &= \left( \frac{y_2 - y_1}{x_2 - x_1} \right) (x_1 - x_3) - y_1 && \# \text{ insert points} \\ &= \left( \frac{2 - 1}{4 - 0} \right) (0 - 2) - 1 && \# \text{ simplify in } \mathbb{F}_5 \\ &= \left( \frac{1}{4} \right) \cdot 3 + 4 = 4 \cdot 3 + 4 = 2 + 4 = 1 \end{aligned}$$

So, in our elliptic curve  $E_1(\mathbb{F}_5)$  we get  $(0, 1) \oplus (4, 2) = (2, 1)$ , and, indeed, the pair  $(2, 1)$  is an  
 element of  $E_1(\mathbb{F}_5)$  as expected. On the other hand,  $(0, 1) \oplus (0, 4) = \mathcal{O}$ , since both points have  
 equal  $x$  coordinates and inverse  $y$  coordinates, rendering them inverses of each other. Adding  
 the point  $(4, 2)$  to itself, we have to use the tangent rule, that is, rule 3 from definition 5.1.1.2: check  
reference

$$\begin{aligned}
 x' &= \left( \frac{3x^2 + a}{2y} \right)^2 - 2x && \# \text{ insert points} \\
 &= \left( \frac{3 \cdot 4^2 + 1}{2 \cdot 2} \right)^2 - 2 \cdot 4 && \# \text{ simplify in } \mathbb{F}_5 \\
 &= \left( \frac{3 \cdot 1 + 1}{4} \right)^2 + 3 \cdot 4 = \left( \frac{4}{4} \right)^2 + 2 = 1 + 2 = 3
 \end{aligned}$$

$$\begin{aligned}
 y' &= \left( \frac{3x^2 + a}{2y} \right)^2 (x - x') - y && \# \text{ insert points} \\
 &= \left( \frac{3 \cdot 4^2 + 1}{2 \cdot 2} \right)^2 (4 - 3) - 2 && \# \text{ simplify in } \mathbb{F}_5 \\
 &= 1 \cdot 1 + 3 = 4
 \end{aligned}$$

2761 So, in our elliptic curve  $E_1(\mathbb{F}_5)$ , we get the doubling of  $(4, 2)$ , that is,  $(4, 2) \oplus (4, 2) = (3, 4)$ ,  
 2762 and, indeed the pair  $(3, 4)$  is an element of  $E_1(\mathbb{F}_5)$  as expected. The group  $E_1(\mathbb{F}_5)$  has no self-  
 2763 inverse points other than the neutral element  $\mathcal{O}$ , since no point has 0 as its y coordinate. We can  
 2764 invoke Sage to double-check the computations.

```

2765 sage: F5 = GF(5)                                260
2766 sage: E1 = EllipticCurve(F5, [1, 1])             261
2767 sage: INF = E1(0) # point at infinity             262
2768 sage: P1 = E1(0, 1)                              263
2769 sage: P2 = E1(4, 2)                              264
2770 sage: P3 = E1(0, 4)                              265
2771 sage: R1 = E1(2, 1)                              266
2772 sage: R2 = E1(3, 4)                              267
2773 sage: R1 == P1+P2                                268
2774 True                                              269
2775 sage: INF == P1+P3                               270
2776 True                                              271
2777 sage: R2 == P2+P2                                272
2778 True                                              273
2779 sage: R2 == 2*P2                                  274
2780 True                                              275
2781 sage: P3 == P3 + INF                             276
2782 True                                              277

```

*Example 71* (Tiny-jubjub). Consider the *TJJ\_13*-curve from example 66 again and recall that its group of rational points is given as follows:

check  
reference

$$\begin{aligned}
 TJJ_{13} = \{ &\mathcal{O}, (1, 2), (1, 11), (4, 0), (5, 2), (5, 11), (6, 5), (6, 8), (7, 2), (7, 11), \\
 &(8, 5), (8, 8), (9, 4), (9, 9), (10, 3), (10, 10), (11, 6), (11, 7), (12, 5), (12, 8) \}
 \end{aligned}$$

2783 In contrast to the group from the previous example, this group contains a self-inverse point,  
 2784 which is different from the neutral element, defined by  $(4, 0)$ . To see what this means, observe  
 2785 that we cannot add  $(4, 0)$  to itself using the tangent rule 3 from definition 5.1.1.2, as the y  
 2786 coordinate is zero. Instead, we have to use rule 2, since  $0 = -0$ . We therefore get  $(4, 0) \oplus$

check  
reference

2787  $(4,0) = \mathcal{O}$  in *TJJ\_13*. The point  $(4,0)$  is therefore the inverse of itself, as adding it to itself  
 2788 results in the neutral element.

```

2789 sage: F13 = GF(13)                                278
2790 sage: MJJ = EllipticCurve(F13, [8, 8])              279
2791 sage: P = MJJ(4, 0)                                280
2792 sage: INF = MJJ(0) # Point at infinity              281
2793 sage: INF == P+P                                     282
2794 True                                                283
2795 sage: INF == 2*P                                     284
2796 True                                                285

```

2797 *Example 72.* Consider the Secp256k1 curve from example 67 again. The following code in-  
 2798 vokes Sage to generate a random affine curve point, then applies our compression method:

check  
reference

```

2799 sage: P = Secp256k1.random_point()                  286
2800 sage: Q = Secp256k1.random_point()                  287
2801 sage: INF = Secp256k1(0)                             288
2802 sage: R1 = -P                                         289
2803 sage: R2 = P + Q                                     290
2804 sage: R3 = Secp256k1.order()*P                       291
2805 sage: P.xy()                                          292
2806 (1130868227573853511381389042467690570841619515611544178521669 293
2807     69177999631234327, 2646197175708723149603620823295226311661
2808     1980265674790621032498974977304647728)
2809 sage: Q.xy()                                          294
2810 (2414716439929438296244703333226322200280520019278069504600440 295
2811     1608813793760166, 65427070121420533449306869796077019361909
2812     78848410309765214049291272511088410)
2813 sage: (ZZ(R1[0]).str(16), ZZ(R1[1]).str(16))         296
2814 ('fa04df99557bc86e6c90b89ea2059dd18748a3a1cbaa03ec88bb588c007 297
2815     cb517', 'c57f0e03fe321bc8d68839578ebd8e40db1b0dc6321997f116
2816     fd08d60152a3ff')
2817 sage: R2.xy()                                          298
2818 (1109262839214235453283782379764994478529835823600618133536533 299
2819     17110811012954221, 8154853520569337014391887865411569741540
2820     4403832529631694427435540217815448611)
2821 sage: R3 == INF                                       300
2822 True                                                  301
2823 sage: P[1]+R1[1] == Fp(0) # -(x,y) = (x,-y)         302
2824 True                                                  303

```

2825 *Exercise 36.* Consider the *TJJ\_13*-curve from example 66.

check  
reference

- 2826 1. Compute the inverse of  $(10,10)$ ,  $\mathcal{O}$ ,  $(4,0)$  and  $(1,2)$ .
- 2827 2. Compute the expression  $3 \cdot (1,11) - (9,9)$ .
- 2828 3. Solve the equation  $x + 2(9,4) = (5,2)$  for some  $x \in TJJ_{13}$
- 2829 4. Solve the equation  $x \cdot (7,11) = (8,5)$  for  $x \in \mathbb{Z}$

**Scalar multiplication** As we have seen in the previous section, elliptic curves  $E(\mathbb{F})$  have the structure of a commutative group associated to them. Moreover, It can moreover be shown that this group is finite and cyclic whenever the field is finite.

To understand elliptic curve scalar multiplication, recall from page 43 that every finite cyclic group of order  $q$  has a generator  $g$  and an associated exponential map  $g^{(\cdot)} : \mathbb{Z}_q \rightarrow \mathbb{G}$ , where  $g^n$  is the  $n$ -fold product of  $g$  with itself.

check  
reference

Elliptic curve scalar multiplication is the exponential map written in additive notation. To be more precise, let  $\mathbb{F}$  be a finite field,  $E(\mathbb{F})$  an elliptic curve of order  $r$ , and  $P$  a generator of  $E(\mathbb{F})$ . Then the **elliptic curve scalar multiplication** with base  $P$  is defined as follows (where  $[0]P = \mathcal{O}$  and  $[m]P = P + P + \dots + P$  is the  $m$ -fold sum of  $P$  with itself):

$$[\cdot]P : \mathbb{Z}_r \rightarrow E(\mathbb{F}); m \mapsto [m]P$$

therefore, elliptic curve scalar multiplication is an instantiation of the general exponential map using additive instead of multiplicative notation. This map is a homomorph of groups, which means that  $[n + m]P = [n]P \oplus [m]P$ .

As with all finite, cyclic groups, the inverse of the exponential map exists and is usually called the **elliptic curve discrete logarithm map**. However, elliptic curves are believed to be XXX-groups, which means that we don't know of any efficient way to actually compute this map.

add term

Scalar multiplication and its inverse, the elliptic curve discrete logarithm, define the **elliptic curve discrete logarithm problem**, which consists of finding solutions  $m \in \mathbb{Z}_r$  such that the following equation holds:

$$P = [m]Q \tag{5.3}$$

Any solution  $m$  is usually called a **discrete logarithm relation** between  $P$  and  $Q$ . If  $Q$  is a generator of the curve, then there is a discrete logarithm relation between  $Q$  and any other point, since  $Q$  generates the group by repeatedly adding  $Q$  to itself. Therefore, we know that some discrete logarithm relation exists for generator  $Q$  and point  $P$ . However, since elliptic curves are believed to be XXX-groups, finding actual relations  $m$  is computationally hard, with runtimes being approximately the size of the order of the group. In practice, we often need the assumption that a discrete logarithm relation exists, while the relation itself is not known.

add term

One useful property of the exponential map in regard to the examples in this book is that it can be used to greatly simplify pen-and-paper computations. As we have seen in example XXX, computing the elliptic curve addition law takes quite a bit of effort when done without a computer. However, when  $g$  is a generator of a small pen-and-paper elliptic curve group of order  $r$ , we can use the exponential map to write the group using cofactor clearing, which implies that  $[r]g = \mathcal{O}$ :

add refer-  
encecofactor  
clearing

$$\mathbb{G} = \{[1]g \rightarrow [2]g \rightarrow [3]g \rightarrow \dots \rightarrow [r-1]g \rightarrow \mathcal{O}\} \tag{5.4}$$

“Logarithmic ordering” like this greatly simplifies complicated elliptic curve addition to the much simpler case of modular  $r$  addition. In order to add two curve points  $P$  and  $Q$ , we only have to look up their discrete log relations with the generator, say  $P = [n]g$  and  $Q = [m]g$ , and compute the sum as  $P \oplus Q = [n + m]g$ . This is, of course, only possible for small groups where we can keep a clear overview, such as XXX.

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ence

In the following example, we will look at some implications of the fact that elliptic curves are finite cyclic groups. We will apply the fundamental theorem of finite cyclic groups and look how it reflects on the curves in consideration.

*Example 73.* Consider the elliptic curve group  $E_1(\mathbb{F}_5)$  from example 65. Since it is a finite cyclic group of order 9, and the prime factorization of 9 is  $3 \cdot 3$ , we can use the fundamental

check  
reference



theorem of finite cyclic groups to reason about all its subgroups. In fact, since the only prime factor of 9 is 3, we know that  $E_1(\mathbb{F}_5)$  has the following subgroups:

- $\mathbb{G}_1 = E_1(\mathbb{F}_5)$  is a subgroup of order 9. By definition, any group is a subgroup of itself.
- $\mathbb{G}_2 = \{(2, 1), (2, 4), \mathcal{O}\}$  is a subgroup of order 3. This is the subgroup associated to the prime factor 3.
- $\mathbb{G}_3 = \{\mathcal{O}\}$  is a subgroup of order 1. This is the trivial subgroup.

Moreover, since  $E_1(\mathbb{F}_5)$  and all its subgroups are cyclic, we know from page 43 that they must have generators. For example, the curve point  $(2, 1)$  is a generator of the order 3 subgroup  $\mathbb{G}_2$ , since every element of  $\mathbb{G}_2$  can be generated by repeatedly adding  $(2, 1)$  to itself:

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reference

$$\begin{aligned}[1](2, 1) &= (2, 1) \\ [2](2, 1) &= (2, 4) \\ [3](2, 1) &= \mathcal{O}\end{aligned}$$

Since  $(2, 1)$  is a generator, we know from XXX that it gives rise to an exponential map from the finite field  $\mathbb{F}_3$  onto  $\mathbb{G}_2$  defined by scalar multiplication:

add refer-  
ence

$$[\cdot](2, 1) : \mathbb{F}_3 \rightarrow \mathbb{G}_2 : x \mapsto [x](2, 1)$$

To give an example of a generator that generates the entire group  $E_1(\mathbb{F}_5)$ , consider the point  $(0, 1)$ . Applying the tangent rule repeatedly, we compute as follows:

$$\begin{array}{ll} [0](0, 1) = \mathcal{O} & [1](0, 1) = (0, 1) \\ [2](0, 1) = (4, 2) & [3](0, 1) = (2, 1) \\ [4](0, 1) = (3, 4) & [5](0, 1) = (3, 1) \\ [6](0, 1) = (2, 4) & [7](0, 1) = (4, 3) \\ [8](0, 1) = (0, 4) & [9](0, 1) = \mathcal{O} \end{array}$$

Again, since  $(2, 1)$  is a generator, we know from XXX that it gives rise to an exponential map. However, since the group order is not a prime number, the exponential map does not map from any field, but from the residue class ring  $\mathbb{Z}_9$  only:

add refer-  
ence

$$[\cdot](0, 1) : \mathbb{Z}_9 \rightarrow \mathbb{G}_1 : x \mapsto [x](0, 1)$$

Using the generator  $(0, 1)$  and its associated exponential map, we can write  $E(\mathbb{F}_1)$  i logarithmic order with respect to  $(0, 1)$  as explained in equation 5.4. We get the following:

check  
reference

$$E_1(\mathbb{F}_5) = \{(0, 1) \rightarrow (4, 2) \rightarrow (2, 1) \rightarrow (3, 4) \rightarrow (3, 1) \rightarrow (2, 4) \rightarrow (4, 3) \rightarrow (0, 4) \rightarrow \mathcal{O}\}$$

This indicates that the first element is a generator, and the  $n$ -th element is the scalar product of  $n$  and the generator. To see how logarithmic orders like this simplify the computations in small elliptic curve groups, consider example 70 again. In that example, we use the chord-and-tangent rule to compute  $(0, 1) \oplus (4, 2)$ . Now, in the logarithmic order of  $E_1(\mathbb{F})$ , we can compute that sum much easier, since we can directly see that  $(0, 1) = [1](0, 1)$  and  $(4, 2) = [2](0, 1)$ . We can then deduce  $(0, 1) \oplus (4, 2) = (2, 1)$  immediately, since  $[1](0, 1) \oplus [2](0, 1) = [3](0, 1) = (2, 1)$ .

check  
reference

To give another example, we can immediately see that  $(3, 4) \oplus (4, 3) = (4, 2)$ , without doing any expensive elliptic curve addition, since we know  $(3, 4) = [4](0, 1)$  as well as  $(4, 3) =$

2883  $[7](0, 1)$  from the logarithmic representation of  $E_1(\mathbb{F}_5)$ . Since  $4 + 7 = 2$  in  $\mathbb{Z}_9$ , the result must  
 2884 be  $[2](0, 1) = (4, 2)$ .

2885 Finally we can use  $E_1(\mathbb{F}_5)$  as an example to understand the concept of cofactor clearing from  
 2886 5.4. Since the order of  $E_1(\mathbb{F}_5)$  is 9, we only have a single factor, which happen to be the cofactor  
 2887 as well. Cofactor clearing then implies that we can map any element from  $E_1(\mathbb{F}_5)$  onto its prime  
 2888 factor group  $\mathbb{G}_2$  by scalar multiplication with 3. For example, taking the element  $(3, 4)$ , which  
 2889 is not in  $\mathbb{G}_2$ , and multiplying it with 3, we get  $[3](3, 4) = (2, 1)$ , which is an element of  $\mathbb{G}_2$  as  
 2890 expected.

check  
reference

2891 In the following example, we will look at the subgroups of our tiny-jubjub curve, define  
 2892 generators, and compute the logarithmic order for pen-and-paper computations. Then we take  
 2893 another look at the principle of cofactor clearing.

2894 *Example 74.* Consider the tiny-jubjub curve  $TJJ\_13$  from example 66 again. Since the order of  
 2895  $TJJ\_13$  is 20, and the prime factorization of 20 is  $2^2 \cdot 5$ , we know that the  $TJJ\_13$  contains a  
 2896 “large” prime-order subgroup of size 5 and a small prime order subgroup of size 2.

check  
reference

2897 To compute those groups, we can apply the technique of cofactor clearing in a try-and-repeat  
 2898 loop. We start the loop by arbitrarily choosing an element  $P \in TJJ\_13$ , then multiplying that  
 2899 element with the cofactor of the group that we want to compute. If the result is  $\mathcal{O}$ , we try a  
 2900 different element and repeat the process until the result is different from the point at infinity  $\mathcal{O}$ .

2901 To compute a generator for the small prime-order subgroup  $(TJJ\_13)_2$ , first observe that the  
 2902 cofactor is 10, since  $20 = 2 \cdot 10$ . We then arbitrarily choose the curve point  $(5, 11) \in TJJ\_13$   
 2903 and compute  $[10](5, 11) = \mathcal{O}$ . Since the result is the point at infinity, we have to try another  
 2904 curve point, say  $(9, 4)$ . We get  $[10](9, 4) = (4, 0)$  and we can deduce that  $(4, 0)$  is a generator  
 2905 of  $(TJJ\_13)_2$ . Logarithmic order then gives  $(TJJ\_13)_2 = \{(4, 0) \rightarrow \mathcal{O}\}$  as expected, since we  
 2906 know from example 71 that  $(4, 0)$  is self-inverse, with  $(4, 0) \oplus (4, 0) = \mathcal{O}$ . We double check the  
 2907 computations using Sage:

check  
reference

2908	<b>sage:</b> <code>F13 = GF(13)</code>	304
2909	<b>sage:</b> <code>TJJ = EllipticCurve(F13, [8, 8])</code>	305
2910	<b>sage:</b> <code>P = TJJ(5, 11)</code>	306
2911	<b>sage:</b> <code>INF = TJJ(0)</code>	307
2912	<b>sage:</b> <code>10*P == INF</code>	308
2913	<b>True</b>	309
2914	<b>sage:</b> <code>Q = TJJ(9, 4)</code>	310
2915	<b>sage:</b> <code>R = TJJ(4, 0)</code>	311
2916	<b>sage:</b> <code>10*Q == R</code>	312
2917	<b>True</b>	313

We can apply the same reasoning to the “large” prime-order subgroup  $(TJJ\_13)_5$ , which contains 5 elements. To compute a generator for this group, first observe that the associated cofactor is 4, since  $20 = 5 \cdot 4$ . We choose the curve point  $(9, 4) \in TJJ\_13$  again, and compute  $[4](9, 4) = (7, 11)$ . We can deduce that  $(7, 11)$  is a generator of  $(TJJ\_13)_5$ . Using the generator  $(7, 11)$ , we compute the exponential map  $[\cdot](7, 11) : \mathbb{F}_5 \rightarrow TJJ\_13$  and get the following:

Explain  
how

$$\begin{aligned}
 [0](7, 11) &= \mathcal{O} \\
 [1](7, 11) &= (7, 11) \\
 [2](7, 11) &= (8, 5) \\
 [3](7, 11) &= (8, 8) \\
 [4](7, 11) &= (7, 2)
 \end{aligned}$$



We can use this computation to write the large-order prime group  $(TJJ\_13)_5$  of the tiny-jubjub curve in logarithmic order, which we will use quite frequently in what follows. We get the following:

$$(TJJ\_13)_5 = \{(7, 11) \rightarrow (8, 5) \rightarrow (8, 8) \rightarrow (7, 2) \rightarrow \mathcal{O}\} \quad (5.5)$$

From this, we can immediately see, for example that  $(8, 8) \oplus (7, 2) = (8, 5)$ , since  $3 + 4 = 2$  in  $\mathbb{F}_5$ .

From the previous two examples, the reader might get the impression that elliptic curve computation can be largely replaced by modular arithmetics. This however, is not true in general, but only an artifact of small groups, where it is possible to write the entire group in a logarithmic order. The following example gives some understanding of why this is not possible in cryptographically secure groups.

*Example 75.* SEKTP BICOIN. DISCRETE LOG HARDNESS PROHIBITS ADDITION IN THE FIELD...

write example

**Projective short Weierstraß form** As we have seen in the previous section, describing elliptic curves as pairs of points that satisfy a certain equation is relatively straight-forward. However, in order to define a group structure on the set of points, we had to add a special point at infinity to act as the neutral element.

Recalling from the definition of projective planes (section 4.4), we know that points at infinity are handled as ordinary points in projective geometry. Therefore, it makes sense to look at the definition of a short Weierstraß curve in projective geometry.

check reference

To see what a short Weierstraß curve in projective coordinates is, let  $\mathbb{F}$  be a finite field of order  $q$  and characteristic  $> 3$ , let  $a, b \in \mathbb{F}$  be two field elements such that  $4a^3 + 27b^2 \bmod q \neq 0$  and let  $\mathbb{FP}^2$  be the projective plane over  $\mathbb{F}$ . Then a **short Weierstraß elliptic curve** over  $\mathbb{F}$  in its projective representation is the set of all points  $[X : Y : Z] \in \mathbb{FP}^2$  from the projective plane that satisfy the **homogenous** cubic equation  $Y^2 \cdot Z = X^3 + a \cdot X \cdot Z^2 + b \cdot Z^3$ :

$$E(\mathbb{FP}^2) = \{[X : Y : Z] \in \mathbb{FP}^2 \mid Y^2 \cdot Z = X^3 + a \cdot X \cdot Z^2 + b \cdot Z^3\} \quad (5.6)$$

To understand how the point at infinity is unified in this definition, recall from XXX that, in projective geometry, points at infinity are given by homogeneous coordinates  $[X : Y : 0]$ . Inserting representatives  $(x_1, y_1, 0) \in [X : Y : 0]$  from those classes into the defining homogenous cubic equations gives the following:

add reference

$$\begin{aligned} y_1^2 \cdot 0 &= x_1^3 + a \cdot x_1 \cdot 0^2 + b \cdot 0^3 \\ 0 &= x_1^3 \end{aligned} \quad \Leftrightarrow$$

This shows that the only point at infinity, that is also a point on a projective short Weierstraß curve is the class  $[0, 1, 0] = \{(0, y, 0) \mid y \in \mathbb{F}\}$ .

This point is the projective representation of  $\mathcal{O}$ . The projective representation of a short Weierstraß curve, therefore, has the advantage that it does not need a special symbol to represent the point at infinity  $\mathcal{O}$  from the affine definition.

*Example 76.* To get an intuition of how an elliptic curve in projective geometry looks, consider curve  $E_1(\mathbb{F}_5)$  from example (65). We know that, in its affine representation, the set of rational points is given as follows:

check reference

$$E_1(\mathbb{F}_5) = \{\mathcal{O}, (0, 1), (2, 1), (3, 1), (4, 2), (4, 3), (0, 4), (2, 4), (3, 4)\} \quad (5.7)$$

2950 This is defined as the set of all pairs  $(x, y) \in \mathbb{F}_5 \times \mathbb{F}_5$  such that the affine short Weierstraß  
 2951 equation  $y^2 = x^3 + ax + b$  with  $a = 1$  and  $b = 1$  is satisfied.

2952 To find the projective representation of a short Weierstraß curve with the same parameters  
 2953  $a = 1$  and  $b = 1$ , we have to compute the set of projective points  $[X : Y : Z]$  from the projec-  
 2954 tive plane  $\mathbb{F}_5\mathbb{P}^2$  that satisfy the following homogenous cubic equation for any representative  
 2955  $(x_1, y_1, z_1) \in [X : Y : Z]$ :

$$y_1^2 z_1 = x_1^3 + 1 \cdot x_1 z_1^2 + 1 \cdot z_1^3 \quad (5.8)$$

2956 We know from XXX that the projective plane  $\mathbb{F}_5\mathbb{P}^2$  contains  $5^2 + 5 + 1 = 31$  elements, so we  
 2957 can take the effort and insert all elements into equation 5.8 and see if both sides match.

For example, consider the projective point  $[0 : 4 : 1]$ . We know from XXX that this point in  
 the projective plane represents the following line in the three-dimensional space  $\mathbb{F}^3$ :

$$[0 : 4 : 1] = \{(0, 0, 0), (0, 4, 1), (0, 3, 2), (0, 2, 3), (0, 1, 4)\}$$

To check whether or not  $[0 : 4 : 1]$  satisfies 5.8, we can insert any representative, in other words,  
 any element from XXX. Each element satisfies the equation if and only if all other elements  
 satisfy the equation. So, we insert  $(0, 4, 1)$  and get the following result:

$$1^2 \cdot 1 = 0^3 + 1 \cdot 0 \cdot 1^2 + 1 \cdot 1^3$$

This tells us that the affine point  $[0 : 4 : 1]$  is indeed a solution to the equation 5.8, but we  
 could just as well have inserted any other representative. For example, inserting  $(0, 3, 2)$  also  
 satisfies 5.8:

$$3^2 \cdot 2 = 0^3 + 1 \cdot 0 \cdot 2^2 + 1 \cdot 2^3$$

2958 To find the projective representation of  $E_1$ , we first observe that the projective line at infinity  
 2959  $[1 : 0 : 0]$  is not a curve point on any projective short Weierstraß curve, since it cannot satisfy  
 2960 XXX for any parameter  $a$  and  $b$ . Therefore, we can exclude it from our consideration.

2961 Moreover, a point at infinity  $[X : Y : 0]$  can only satisfy equation XXX for any  $a$  and  $b$ , if  
 2962  $X = 0$ , which implies that the only point at infinity relevant for short Weierstraß elliptic curves  
 2963 is  $[0 : 1 : 0]$ , since  $[0 : k : 0] = [0 : 1 : 0]$  for all  $k$  from the finite field. Therefore, we can exclude  
 2964 all points at infinity except the point  $[0 : 1 : 0]$ .

2965 All points that remain are the affine points  $[X : Y : 1]$ . Inserting all of them into XXX, we  
 2966 get the set of all projective curve points as follows:

$$E_1(\mathbb{F}_5\mathbb{P}^2) = \{[0 : 1 : 0], [0 : 1 : 1], [2 : 1 : 1], [3 : 1 : 1], \\ [4 : 2 : 1], [4 : 3 : 1], [0 : 4 : 1], [2 : 4 : 1], [3 : 4 : 1]\}$$

2967 If we compare this with the affine representation, we see that there is a 1:1 correspondence  
 2968 between the points in the affine representation in 5.7 and the affine points in projective geometry,  
 2969 and that the point  $[0 : 1 : 0]$  represents the additional point  $\mathcal{O}$  in the projective representation.

2970 *Exercise 37.* Compute the projective representation of the tiny-jubjub curve and the logarithmic  
 2971 order of its large prime-order subgroup with respect to the generator  $(7, 11)$ .

2972 **Projective Group law** As we have seen on page 70, one of the key properties of an elliptic  
 2973 curve is that it comes with a definition of a group law on the set of its rational points, described  
 2974 geometrically by the chord-and-tangent rule (definition 5.1.1.1). This rule was kind of intuitive,

with the exception of the distinguished point at infinity, which appeared whenever the chord or the tangent did not have a third intersection point with the curve.

One of the key features of projective coordinates is that, in projective space, it is guaranteed that any chord will always intersect the curve in three points, and any tangent will intersect it in two points including the tangent point. So, the geometric picture simplifies, as we don't need to consider external symbols and associated cases.

Again, it can be shown that the points of an elliptic curve in projective space form a commutative group with respect to the tangent-and-chord rule such that the projective point  $[0 : 1 : 0]$  is the neutral element, and the additive inverse of a point  $[X : Y : Z]$  is given by  $[X : -Y : Z]$ . The addition law is usually described by the following algorithm, minimizing the number of necessary additions and multiplications in the base field.

*Exercise 38.* Compare the affine addition law for short Weierstraß curves with the projective addition rule. Which branch in the projective rule corresponds to which case in the affine law?

Check if following Alg is floated too far

**Coordinate Transformations** As we have seen in example XXX, there was a close relation between the affine and the projective representation of a short Weierstraß curve. This was not a coincidence. In fact, from a mathematical point of view, projective and affine short Weierstraß curves describe the same thing, as there is a one-to-one correspondence (an isomorphism) between both representations for any arbitrary parameters  $a$  and  $b$ .

add reference

To specify the isomorphism, let  $E(\mathbb{F})$  and  $E(\mathbb{F}\mathbb{P}^2)$  be an affine and a projective short Weierstraß curve defined for the same parameters  $a$  and  $b$ . Then the map in 5.9 maps points from the affine representation to points from the projective representation of a short Weierstraß curve. In other words, if the pair of points  $(x, y)$  satisfies the affine equation  $y^2 = x^3 + ax + b$ , then all homogeneous coordinates  $(x_1, y_1, z_1) \in [x : y : 1]$  satisfy the projective equation  $y_1^2 \cdot z_1 = x_1^3 + ay_1 \cdot z_1^2 + b \cdot z_1^3$ .

$$\Phi : E(\mathbb{F}) \rightarrow E(\mathbb{F}\mathbb{P}^2) : \begin{array}{ll} (x, y) & \mapsto [x : y : 1] \\ \mathcal{O} & \mapsto [0 : 1 : 0] \end{array} \quad (5.9)$$

The inverse is given by the following map:

$$\Phi^{-1} : E(\mathbb{F}\mathbb{P}^2) \rightarrow E(\mathbb{F}) : [X : Y : Z] \mapsto \begin{cases} (\frac{X}{Z}, \frac{Y}{Z}) & \text{if } Z \neq 0 \\ \mathcal{O} & \text{if } Z = 0 \end{cases} \quad (5.10)$$

Note that the only projective point  $[X : Y : Z]$  with  $Z \neq 0$  that satisfies XXX is given by the class  $[0 : 1 : 0]$ .

add reference

One key feature of  $\Phi$  and its inverse is that it respects the group structure, which means that  $\Phi((x_1, y_1) \oplus (x_2, y_2))$  is equal to  $\Phi(x_1, y_1) \oplus \Phi(x_2, y_2)$ . The same holds true for the inverse map  $\Phi^{-1}$ .

Maps with these properties are called **group isomorphisms**, and, from a mathematical point of view, the existence of  $\Phi$  implies that these two definitions are equivalent, and implementations can choose freely between these representations.

## 5.1.2 Montgomery Curves

History and use of them (optimized scalar multiplication)

write up this part

**Algorithm 6** Projective Weierstraß Addition Law

---

**Require:**  $[X_1 : Y_1 : Z_1], [X_2 : Y_2 : Z_2] \in E(\mathbb{F}_{\mathbb{P}}^2)$

**procedure** ADD-RULE( $[X_1 : Y_1 : Z_1], [X_2 : Y_2 : Z_2]$ )

**if**  $[X_1 : Y_1 : Z_1] == [0 : 1 : 0]$  **then**

$[X_3 : Y_3 : Z_3] \leftarrow [X_2 : Y_2 : Z_2]$

**else if**  $[X_2 : Y_2 : Z_2] == [0 : 1 : 0]$  **then**

$[X_3 : Y_3 : Z_3] \leftarrow [X_1 : Y_1 : Z_1]$

**else**

$U_1 \leftarrow Y_2 \cdot Z_1$

$U_2 \leftarrow Y_1 \cdot Z_2$

$V_1 \leftarrow X_2 \cdot Z_1$

$V_2 \leftarrow X_1 \cdot Z_2$

**if**  $V_1 == V_2$  **then**

**if**  $U_1 \neq U_2$  **then**  $[X_3 : Y_3 : Z_3] \leftarrow [0 : 1 : 0]$

**else**

**if**  $Y_1 == 0$  **then**  $[X_3 : Y_3 : Z_3] \leftarrow [0 : 1 : 0]$

**else**

$W \leftarrow a \cdot Z_1^2 + 3 \cdot X_1^2$

$S \leftarrow Y_1 \cdot Z_1$

$B \leftarrow X_1 \cdot Y_1 \cdot S$

$H \leftarrow W^2 - 8 \cdot B$

$X' \leftarrow 2 \cdot H \cdot S$

$Y' \leftarrow W \cdot (4 \cdot B - H) - 8 \cdot Y_1^2 \cdot S^2$

$Z' \leftarrow 8 \cdot S^3$

$[X_3 : Y_3 : Z_3] \leftarrow [X' : Y' : Z']$

**end if**

**end if**

**else**

$U = U_1 - U_2$

$V = V_1 - V_2$

$W = Z_1 \cdot Z_2$

$A = U^2 \cdot W - V^3 - 2 \cdot V^2 \cdot V_2$

$X' = V \cdot A$

$Y' = U \cdot (V^2 \cdot V_2 - A) - V^3 \cdot U_2$

$Z' = V^3 \cdot W$

$[X_3 : Y_3 : Z_3] \leftarrow [X' : Y' : Z']$

**end if**

**end if**

**return**  $[X_3 : Y_3 : Z_3]$

**end procedure**

**Ensure:**  $[X_3 : Y_3 : Z_3] == [X_1 : Y_1 : Z_1] \oplus [X_2 : Y_2 : Z_2]$

---

**Affine Montgomery Form** To see what a Montgomery curve in affine coordinates is, let  $\mathbb{F}$  be a finite field of characteristic  $> 2$ , and let  $A, B \in \mathbb{F}$  be two field elements such that  $B \neq 0$  and  $A^2 \neq 4$ . A **Montgomery elliptic curve**  $M(\mathbb{F})$  over  $\mathbb{F}$  in its affine representation is the set of all pairs of field elements  $(x, y) \in \mathbb{F} \times \mathbb{F}$  that satisfy the Montgomery cubic equation  $B \cdot y^2 = x^3 + A \cdot x^2 + x$ , together with a distinguished symbol  $\mathcal{O}$ , called the **point at infinity**.

$$M(\mathbb{F}) = \{(x, y) \in \mathbb{F} \times \mathbb{F} \mid B \cdot y^2 = x^3 + A \cdot x^2 + x\} \cup \{\mathcal{O}\} \quad (5.11)$$

Despite the fact that Montgomery curves look different from short Weierstraß curves, they are just a special way to describe certain short Weierstraß curves. In fact, every curve in affine Montgomery form can be transformed into an elliptic curve in Weierstraß form. To see that, assume that a curve is given in Montgomery form  $By^2 = x^3 + Ax^2 + x$ . The associated Weierstraß form is then as follows:

$$y^2 = x^3 + \frac{3 - A^2}{3B^2} \cdot x + \frac{2A^3 - 9A}{27B^3} \quad (5.12)$$

On the other hand, an elliptic curve  $E(\mathbb{F})$  over base field  $\mathbb{F}$  in Weierstraß form  $y^2 = x^3 + ax + b$  can be converted to Montgomery form if and only if the following conditions hold:

**Definition 5.1.2.1. Requirements for Montgomery curves**

- The number of points on  $E(F)$  is divisible by 4
- The polynomial  $z^3 + az + b \in \mathbb{F}[z]$  has at least one root  $z_0 \in \mathbb{F}$
- $3z_0^2 + a$  is a quadratic residue in  $\mathbb{F}$ .

When these conditions are satisfied, then for  $s = (\sqrt{3z_0^2 + a})^{-1}$ , the equivalent Montgomery curve is defined by the following equation:

$$sy^2 = x^3 + (3z_0s)x^2 + x \quad (5.13)$$

In the following example we will look at our tiny-jubjub curve again, and show that it is actually a Montgomery curve.

*Example 77.* Consider the prime field  $\mathbb{F}_{13}$  and the tiny-jubjub curve *TJJ\_13* from example 66. To see that it is a Montgomery curve, we have to check the requirements from 5.1.2.1:

Since the order of *TJJ\_13* is 20, which is divisible by 4, the first requirement is met.

Next, since  $a = 8$  and  $b = 8$ , we have to check if the polynomial  $P(z) = z^3 + 8z + 8$  has a root in  $\mathbb{F}_{13}$ . We simply evaluate  $P$  at all numbers  $z \in \mathbb{F}_{13}$ , and find that  $P(4) = 0$ , so a root is given by  $z_0 = 4$ .

In the last step, we have to check that  $3 \cdot z_0^2 + a$  has a root in  $\mathbb{F}_{13}$ . We compute as follows:

$$\begin{aligned} 3z_0^2 + a &= 3 \cdot 4^2 + 8 \\ &= 3 \cdot 3 + 8 \\ &= 9 + 8 \\ &= 4 \end{aligned}$$

To see if 4 is a quadratic residue, we can use Euler's criterion (4.28) to compute the Legendre symbol of 4. We get the following:

is the label in L<sup>A</sup>T<sub>E</sub>X correct here?

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$$\left(\frac{4}{13}\right) = 4^{\frac{13-1}{2}} = 4^6 = 1$$

3039 This means that 4 does have a root in  $\mathbb{F}_{13}$ . In fact, computing a root of 4 in  $\mathbb{F}_{13}$  is easy, since  
 3040 the integer root 2 of 4 is also one of its roots in  $\mathbb{F}_{13}$ . The other root is given by  $13 - 4 = 9$ .

Since all requirements are met, we have now shown that *TJJ\_13* is indeed a Montgomery curve, and we can use 5.13 to compute its associated Montgomery form. We compute as follows:

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$$\begin{aligned} s &= \left( \sqrt{3 \cdot z_0^2 + 8} \right)^{-1} \\ &= 2^{-1} && \# \text{ Fermat's little theorem} \\ &= 2^{13-2} && \# 2048 \bmod 13 = 7 \\ &= 7 \end{aligned}$$

The defining equation for the Montgomery form of our tiny-jubjub curve is then given by the following equation:

$$\begin{aligned} sy^2 &= x^3 + (3z_0s)x^2 + x && \Rightarrow \\ 7 \cdot y^2 &= x^3 + (3 \cdot 4 \cdot 7)x^2 + x && \Leftrightarrow \\ 7 \cdot y^2 &= x^3 + 6x^2 + x \end{aligned}$$

3041 So, we get the defining parameters as  $B = 7$  and  $A = 6$ , and we can write the tiny-jubjub curve  
 3042 in its affine Montgomery representation as follows:

$$TJJ\_13 = \{(x, y) \in \mathbb{F}_{13} \times \mathbb{F}_{13} \mid 7 \cdot y^2 = x^3 + 6x^2 + x\} \cup \{\mathcal{O}\} \quad (5.14)$$

Now that we have the abstract definition of our tiny-jubjub curve in Montgomery form, we can compute the set of points by inserting all pairs  $(x, y) \in \mathbb{F}_{13} \times \mathbb{F}_{13}$  similarly to how we computed the curve points in its Weierstraß representation. We get the following:

$$TJJ\_13 = \{\mathcal{O}, (0, 0), (1, 4), (1, 9), (2, 4), (2, 9), (3, 5), (3, 8), (4, 4), (4, 9), (5, 1), (5, 12), (7, 1), (7, 12), (8, 1), (8, 12), (9, 2), (9, 11), (10, 3), (10, 10)\}$$

3043

```

3044 sage: F13 = GF(13)                                     314
3045 sage: L_MTJJ = []                                       315
3046 ....: for x in F13:                                     316
3047 ....:     for y in F13:                                  317
3048 ....:         if F13(7)*y^2 == x^3 + F13(6)*x^2 + x:    318
3049 ....:             L_MTJJ.append((x, y))                 319
3050 sage: MTJJ = Set(L_MTJJ)                                320
3051 sage: # does not compute the point at infinity          321

```

3052 **Affine Montgomery coordinate transformation** Comparing the Montgomery representa-  
 3053 tion of the previous example (equation 5.14) with the Weierstraß representation of the same  
 3054 curve (equation 5.2), we see that there is a 1:1 correspondence between the curve points in both

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examples. This is no accident. In fact, if  $M_{A,B}$  is a Montgomery curve, and  $E_{a,b}$  a Weierstraß curve with  $a = \frac{3-A^2}{3B^2}$  and  $b = \frac{2A^2-9A}{27B^3}$  then the following function maps all points in Montgomery representation onto the points in Weierstraß representation:

$$\Phi : M_{A,B} \rightarrow E_{a,b} : (x, y) \mapsto \left( \frac{3x+A}{3B}, \frac{y}{B} \right) \quad (5.15)$$

This map is a 1:1 correspondence (an isomorphism), and its inverse map is given by the following equation (where  $z_0$  is a root of the polynomial  $z^3 + az + b \in \mathbb{F}[z]$  and  $s = (\sqrt{3z_0^2 + a})^{-1}$ ).

$$\Phi^{-1} : E_{a,b} \rightarrow M_{A,B} : (x, y) \mapsto (s \cdot (x - z_0), s \cdot y) \quad (5.16)$$

Using this map, it is therefore possible for implementations of Montgomery curves to freely transit between the Weierstraß and the Montgomery representation. However, as we saw in definition 5.1.2.1, not every Weierstraß curve is a Montgomery curve, as all criteria in 5.1.2.1 have to be satisfied. This means that the map  $\Phi^{-1}$  does not always exist.

*Example 78.* Consider our tiny-jubjub curve again. In equation 5.2 we derived its Weierstraß representation and in example 5.14, we derived its Montgomery representation.

To see how coordinate transformation  $\Phi$  works in this example, let's map points from the Montgomery representation onto points from the Weierstraß representation. Inserting, for example, the point  $(0, 0)$  from the Montgomery representation 5.14 into  $\Phi$  gives the following:

$$\begin{aligned} \Phi(0, 0) &= \left( \frac{3 \cdot 0 + A}{3B}, \frac{0}{B} \right) \\ &= \left( \frac{3 \cdot 0 + 6}{3 \cdot 7}, \frac{0}{7} \right) \\ &= \left( \frac{6}{8}, 0 \right) \\ &= (4, 0) \end{aligned}$$

As we can see, the Montgomery point  $(0, 0)$  maps to the self-inverse point  $(4, 0)$  of the Weierstraß representation. On the other hand, we can use our computations of  $s = 7$  and  $z_0 = 4$  from XXX to compute the inverse map  $\Phi^{-1}$ , which maps points on the Weierstraß representation to points on the Montgomery form. Inserting, for example,  $(4, 0)$  we get the following:

$$\begin{aligned} \Phi^{-1}(4, 0) &= (s \cdot (4 - z_0), s \cdot 0) \\ &= (7 \cdot (4 - 4), 0) \\ &= (0, 0) \end{aligned}$$

As expected, the inverse map maps the Weierstraß point back to where it originated in the Montgomery form. We can invoke Sage to check that our computation of  $\Phi$  is correct:

```

3069 sage: # Compute PHI of Montgomery form: 322
3070 sage: L_PHI_MTJJ = [] 323
3071 sage: for (x,y) in L_MTJJ: # LMJJ as defined previously 324
3072     . . . . . v = (F13(3)*x + F13(6)) / (F13(3)*F13(7)) 325
3073     . . . . . w = y/F13(7) 326
3074     . . . . . L_PHI_MTJJ.append((v,w)) 327

```

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```

3075 sage: PHI_MTJJ = Set(L_PHI_MTJJ) 328
3076 sage: # Computation Weierstrass form 329
3077 sage: C_WTJJ = EllipticCurve(F13, [8, 8]) 330
3078 sage: L_WTJJ = [P.xy() for P in C_WTJJ.points() if P.order() > 331
3079 1]
3080 sage: WTJJ = Set(L_WTJJ) 332
3081 sage: # check PHI(Montgomery) == Weierstrass 333
3082 sage: WTJJ == PHI_MTJJ 334
3083 True 335
3084 sage: # check the inverse map PHI^(-1) 336
3085 sage: L_PHIINV_WTJJ = [] 337
3086 sage: for (v,w) in L_WTJJ: 338
3087 .....:     x = F13(7)*(v-F13(4)) 339
3088 .....:     y = F13(7)*w 340
3089 .....:     L_PHIINV_WTJJ.append((x,y)) 341
3090 sage: PHIINV_WTJJ = Set(L_PHIINV_WTJJ) 342
3091 sage: MTJJ == PHIINV_WTJJ 343
3092 True 344

```

3093 **Montgomery group law** We have seen that Montgomery curves special cases of short Weier-  
3094 straß curves. As such, they have a group structure defined on the set of their points, which can  
3095 also be derived from the chord-and-tangent rule. In accordance with short Weierstraß curves, it  
3096 can be shown that the identity  $x_1 = x_2$  implies  $y_2 = \pm y_1$ , meaning that the following rules are a  
3097 complete description of the affine addition law.

3098 *Definition 5.1.2.2. Montgomery group law*

- 3099 • (Neutral element) Point at infinity  $\mathcal{O}$  is the neutral element.
- 3100 • (Additive inverse ) The additive inverse of  $\mathcal{O}$  is  $\mathcal{O}$ . For any other curve point  $(x,y) \in$   
3101  $M(\mathbb{F}_q) \setminus \{\mathcal{O}\}$ , the additive inverse is given by  $(x, -y)$ .
- 3102 • (Addition rule) For any two curve points  $P, Q \in M(\mathbb{F}_q)$ , addition is defined by one of the  
3103 following cases:
  - 3104 1. (Adding the neutral element) If  $Q = \mathcal{O}$ , then the sum is defined as  $P + Q = P$ .
  - 3105 2. (Adding inverse elements) If  $P = (x,y)$  and  $Q = (x, -y)$ , the sum is defined as  $P +$   
3106  $Q = \mathcal{O}$ .
  3. (Adding non-self-inverse equal points) If  $P = (x,y)$  and  $Q = (x,y)$  with  $y \neq 0$ , the  
sum  $2P = (x', y')$  is defined as follows:

$$x' = \left( \frac{3x_1^2 + 2Ax_1 + 1}{2By_1} \right)^2 \cdot B - (x_1 + x_2) - A, \quad y' = \frac{3x_1^2 + 2Ax_1 + 1}{2By_1} (x_1 - x') - y_1$$

4. (Adding non-inverse different points) If  $P = (x_1, y_1)$  and  $Q = (x_2, y_2)$  such that  $x_1 \neq$   
 $x_2$ , the sum  $R = P + Q$  with  $R = (x_3, y_3)$  is defined as follows:

$$x' = \left( \frac{y_2 - y_1}{x_2 - x_1} \right)^2 B - (x_1 + x_2) - A, \quad y' = \frac{y_2 - y_1}{x_2 - x_1} (x_1 - x') - y_1$$



### 5.1.3 Twisted Edwards Curves

As we have seen in 5.1.2.2 both Weierstraß and Montgomery curves have somewhat complicated addition and doubling laws, as many cases have to be distinguished. Those various cases translate to branches in computer programs.

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In the context of SNARK development, two computational models for bounded computations are used, called **circuits** and **rank-1 constraint systems**. Program branches are undesirably costly when implemented in those models. It is therefore advantageous to look for curves with an addition/doubling rule that requires no branches and as few field operations as possible.

**Twisted Edwards curves** are particularly useful here, as a subclass of these curves has a compact and easily implementable addition law that works for all points including the point at infinity. Implementing this law needs no branching.

**Twisted Edwards Form** To see what an affine **twisted Edwards curve** looks like, let  $\mathbb{F}$  be a finite field of characteristic  $> 2$ , and let  $a, d \in \mathbb{F} \setminus \{0\}$  be two non-zero field elements with  $a \neq d$ . A **twisted Edwards elliptic curve** in its affine representation is the set of all pairs  $(x, y)$  from  $\mathbb{F} \times \mathbb{F}$  that satisfy the twisted Edwards equation  $a \cdot x^2 + y^2 = 1 + d \cdot x^2 y^2$ , given below:

$$E(\mathbb{F}) = \{(x, y) \in \mathbb{F} \times \mathbb{F} \mid a \cdot x^2 + y^2 = 1 + d \cdot x^2 y^2\} \quad (5.17)$$

A twisted Edwards curve is called an **Edwards curve (non-twisted)**, if the parameter  $a$  is equal to 1, and it is called a **SNARK-friendly twisted Edwards curve** if the parameter  $a$  is a quadratic residue and the parameter  $d$  is a quadratic non-residue.

As we can see from the definition, affine twisted Edwards curves look somewhat different from Weierstraß curves, as their affine representation does not need a special symbol to represent the point at infinity. In fact, we will see that the pair  $(0, 1)$  is always a point on any twisted Edwards curve, and that it takes the role of the point at infinity.

Despite their different appearances however, twisted Edwards curves are equivalent to Montgomery curves in the sense that, for every twisted Edwards curve, there is a Montgomery curve, and a way to map the points of one curve in a 1:1 correspondence onto the other and vice versa. To see that, assume that a curve in twisted Edwards form  $a \cdot x^2 + y^2 = 1 + d \cdot x^2 y^2$  is given. The associated Montgomery curve is then defined by the Montgomery equation:

$$\frac{4}{a-d} y^2 = x^3 + \frac{2(a+d)}{a-d} \cdot x^2 + x \quad (5.18)$$

On the other hand, a Montgomery curve  $By^2 = x^3 + Ax^2 + x$  with  $B \neq 0$  and  $A^2 \neq 4$  can give rise to a twisted Edwards curve defined by the following equation:

$$\left(\frac{A+2}{B}\right)x^2 + y^2 = 1 + \left(\frac{A-2}{B}\right)x^2 y^2 \quad (5.19)$$

As we have seen in equation 5.12 and the following discussion, Montgomery curves are just a special class of Weierstraß curves. Furthermore we now know that twisted Edwards curves are special Weierstraß curves too. This means that the more general way to describe elliptic curves is as Weierstraß curves.

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*Example 79.* Consider the tiny-jubjub curve from example 66 again. We know from example 77 that it is a Montgomery curve, and, since Montgomery curves are equivalent to twisted Edwards curves, we want to write this curve in twisted Edwards form. We use equation 5.19,

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and compute the parameters  $a$  and  $d$  as follows:

$$\begin{aligned}
 a &= \frac{A+2}{B} && \# \text{ insert } A=6 \text{ and } B=7 \\
 &= \frac{8}{7} = 3 && \# 7^{-1} = 2 \\
 \\ 
 d &= \frac{A-2}{B} \\
 &= \frac{4}{7} = 8
 \end{aligned}$$

Thus, we get the defining parameters as  $a = 3$  and  $d = 8$ . Since our goal is to use this curve later on in implementations of pen-and-paper SNARKs, let us show that tiny-jubjub is also a **SNARK-friendly** twisted Edwards curve. To see that, we have to show that  $a$  is a quadratic residue and  $d$  is a quadratic non-residue. We therefore compute the Legendre symbols of  $a$  and  $d$  using Euler's criterion. We get the following:

$$\begin{aligned}
 \left( \frac{3}{13} \right) &= 3^{\frac{13-1}{2}} \\
 &= 3^6 = 1
 \end{aligned}$$

$$\begin{aligned}
 \left( \frac{8}{13} \right) &= 8^{\frac{13-1}{2}} \\
 &= 8^6 = 12 = -1
 \end{aligned}$$

3140 This proves that tiny-jubjub is SNARK-friendly. We can write the tiny-jubjub curve in its  
3141 affine twisted Edwards representation as follows:

$$TJJ_{13} = \{(x, y) \in \mathbb{F}_{13} \times \mathbb{F}_{13} \mid 3 \cdot x^2 + y^2 = 1 + 8 \cdot x^2 \cdot y^2\} \quad (5.20)$$

3142 Now that we have the abstract definition of our tiny-jubjub curve in twisted Edwards form,  
3143 we can compute the set of points by inserting all pairs  $(x, y) \in \mathbb{F}_{13} \times \mathbb{F}_{13}$ , similarly to how we  
3144 computed the curve points in its Weierstraß or Edwards representation. We get the following:

$$\begin{aligned}
 TJJ_{13} = \{ &(0, 1), (0, 12), (1, 2), (1, 11), (2, 6), (2, 7), (3, 0), (5, 5), (5, 8), (6, 4), \\
 &(6, 9), (7, 4), (7, 9), (8, 5), (8, 8), (10, 0), (11, 6), (11, 7), (12, 2), (12, 11) \} \quad (5.21)
 \end{aligned}$$

3145

```

3146 sage: F13 = GF(13)                                     345
3147 sage: L_ETJJ = []                                       346
3148 .....: for x in F13:                                     347
3149 .....:     for y in F13:                                 348
3150 .....:         if F13(3)*x^2 + y^2 == 1+ F13(8)*x^2*y^2: 349
3151 .....:             L_ETJJ.append((x, y))                 350
3152 sage: ETJJ = Set(L_ETJJ)                                351

```

**Twisted Edwards group law** As we have seen, twisted Edwards curves are equivalent to Montgomery curves, and, as such, also have a group law. However, in contrast to Montgomery and Weierstraß curves, the group law of SNARK-friendly twisted Edwards curves can be described by a single computation that works in all cases, no matter if we add the neutral element, the inverse, or if we have to double a point. To see what the group law looks like, first observe that the point  $(0, 1)$  is a solution to  $a \cdot x^2 + y^2 = 1 + d \cdot x^2 \cdot y^2$  for any curve. The sum of any two points  $(x_1, y_1), (x_2, y_2)$  on an Edwards curve  $E(\mathbb{F})$  is then given by the following equation:

$$(x_1, y_1) \oplus (x_2, y_2) = \left( \frac{x_1 y_2 + y_1 x_2}{1 + d x_1 x_2 y_1 y_2}, \frac{y_1 y_2 - a x_1 x_2}{1 - d x_1 x_2 y_1 y_2} \right) \quad (5.22)$$

and it can be shown that the point  $(0, 1)$  serves as the neutral element and the inverse of a point  $(x_1, y_1)$  is given by  $(-x_1, y_1)$ .

*Example 80.* Lets look at the tiny-jubjub curve in Edwards form from example 5.20 again. As we have seen, this curve is given by

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$$TJJ\_13 = \{(0, 1), (0, 12), (1, 2), (1, 11), (2, 6), (2, 7), (3, 0), (5, 5), (5, 8), (6, 4), (6, 9), (7, 4), (7, 9), (8, 5), (8, 8), (10, 0), (11, 6), (11, 7), (12, 2), (12, 11)\}$$

To get an understanding of the twisted Edwards addition law, let's first add the neutral element  $(0, 1)$  to itself. We apply the group law 5.22 and get the following:

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$$\begin{aligned} (0, 1) \oplus (0, 1) &= \left( \frac{0 \cdot 1 + 1 \cdot 0}{1 + 8 \cdot 0 \cdot 0 \cdot 1 \cdot 1}, \frac{1 \cdot 1 - 3 \cdot 0 \cdot 0}{1 - 8 \cdot 0 \cdot 0 \cdot 1 \cdot 1} \right) \\ &= (0, 1) \end{aligned}$$

So, as expected, the neutral element added to itself gives the neutral element again. Now let's add the neutral element to some other curve point. We get the following:

$$\begin{aligned} (0, 1) \oplus (8, 5) &= \left( \frac{0 \cdot 5 + 1 \cdot 8}{1 + 8 \cdot 0 \cdot 8 \cdot 1 \cdot 5}, \frac{1 \cdot 5 - 3 \cdot 0 \cdot 8}{1 - 8 \cdot 0 \cdot 8 \cdot 1 \cdot 5} \right) \\ &= (8, 5) \end{aligned}$$

Again, as expected, adding the neutral element to any element will result in that element again. Given any curve point  $(x, y)$ , we know that its inverse is given by  $(-x, y)$ . To see how the addition of a point to its inverse works, we compute as follows:

$$\begin{aligned} (5, 5) \oplus (8, 5) &= \left( \frac{5 \cdot 5 + 5 \cdot 8}{1 + 8 \cdot 5 \cdot 8 \cdot 5 \cdot 5}, \frac{5 \cdot 5 - 3 \cdot 5 \cdot 8}{1 - 8 \cdot 5 \cdot 8 \cdot 5 \cdot 5} \right) \\ &= \left( \frac{12 + 1}{1 + 5}, \frac{12 - 3}{1 - 5} \right) \\ &= \left( \frac{0}{6}, \frac{12 + 10}{1 + 8} \right) \\ &= \left( 0, \frac{9}{9} \right) \\ &= (0, 1) \end{aligned}$$

Adding a curve point to its inverse gives the neutral element, as expected. As we have seen from these examples, the twisted Edwards addition law handles edge cases particularly well and in a unified way.

## 5.2 Elliptic Curve Pairings

As we have seen in equation 4.1, some groups come with the notation of a so-called pairing map, which is a non-degenerate bilinear map from two groups into another group.

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In this section, we discuss **pairings on elliptic curves**, which form the basis of several zk-SNARKs and other zero-knowledge proof schemes. The SNARKs derived from pairings have the advantage of constant proof sizes, which is crucial to blockchains.

We start out by defining elliptic curve pairings and discussing a simple application which bears some resemblance to more advanced SNARKs. We then introduce the pairings arising from elliptic curves and describe Miller's algorithm, which makes these pairings practical rather than just theoretically interesting.

Elliptic curves have a few structures, like the Weil or the Tate map that qualifies as pairing.

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**Embedding Degrees** As we will see in what follows, every elliptic curve gives rise to a pairing map. However, we will also see in example XXX that not every such pairing can be efficiently computed. In order to distinguish curves with efficiently computable pairings from the rest, we need to start with an introduction to the so-called **embedding degree** of a curve.

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### Definition 5.2.0.1. Embedding degree

Let  $\mathbb{F}$  be a finite field, let  $E(\mathbb{F})$  be an elliptic curve over  $\mathbb{F}$ , and let  $n$  be a prime number that divides the order of  $E(\mathbb{F})$ . The embedding degree of  $E(\mathbb{F})$  with respect to  $n$  is then the smallest integer  $k$  such that  $n$  divides  $q^k - 1$ .

Fermat's little theorem (page 21 ff.) implies that every curve has at least **some** embedding degree  $k$ , since at least  $k = n - 1$  is always a solution to the congruency  $q^k \equiv 1 \pmod{n}$ . This implies that the remainder of the integer division of  $q^k - 1$  by  $n$  is 0.

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*Example 81.* To get a better intuition of the embedding degree, let's consider the elliptic curve  $E_1(\mathbb{F}_5)$  from example 65. We know from 65 that the order of  $E_1(\mathbb{F}_5)$  is 9, and, since the only prime factor of 9 is 3, we compute the embedding degree of  $E_1(\mathbb{F}_5)$  with respect to 3.

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To find the embedding degree, we have to find the smallest integer  $k$  such that 3 divides  $q^k - 1 = 5^k - 1$ . We try and increment until we find a proper  $k$ .

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$$k = 1: 5^1 - 1 = 4$$

not divisible by 3

$$k = 2: 5^2 - 1 = 24$$

divisible by 3

Now we know that the embedding degree of  $E_1(\mathbb{F}_5)$  is 2 relative to the prime factor 3.

*Example 82.* Let us consider the tiny-jubjub curve  $TJJ\_13$  from example 66. We know from 66 that the order of  $TJJ\_13$  is 20, and that the order therefore has two prime factors. A "large" prime factor 5 and a small prime factor 2.

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We start by computing the embedding degree of  $TJJ\_13$  with respect to the large prime factor 5. To find that embedding degree, we have to find the smallest integer  $k$  such that 5 divides  $q^k - 1 = 13^k - 1$ . We try and increment until we find a proper  $k$ .

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$$k = 1: 13^1 - 1 = 12$$

not divisible by 5

$$k = 2: 13^2 - 1 = 168$$

not divisible by 5

$$k = 3: 13^3 - 1 = 2196$$

not divisible by 5

$$k = 4: 13^4 - 1 = 28560$$

divisible by 5

Now we know that the embedding degree of  $TJJ\_13$  is 4 relative to the prime factor 5.

In real-world applications, like on pairing-friendly elliptic curves such as BLS\_12-381, usually only the embedding degree of the large prime factor is relevant, which in the case of our tiny-jubjub curve is represented by 5. It should be noted, however that every prime factor of a curve's order has its own notation of embedding degree despite the fact that this is mostly irrelevant in applications.

To find the embedding degree of the small prime factor 2, we have to find the smallest integer  $k$  such that 2 divides  $q^k - 1 = 13^k - 1$ . We try and increment until we find a proper  $k$ .

$$k = 1: 13^1 - 1 = 12 \quad \text{divisible by 2}$$

Now we know that the embedding degree of  $TJJ\_13$  is 1 relative to the prime factor 2. As we have seen, different prime factors can have different embedding degrees in general.

```

sage: p = 13
sage: # large prime factor
sage: n = 5
sage: for k in range(1,5): # Fermat's little theorem
.....:     if (p^k-1)%n == 0:
.....:         break
sage: k
4
sage: # small prime factor
sage: n = 2
sage: for k in range(1,2): # Fermat's little theorem
.....:     if (p^k-1)%n == 0:
.....:         break
sage: k
1

```

*Example 83.* To give an example of a cryptographically secure real-world elliptic curve that does not have a small embedding degree, let's look at curve Secp256k1 again. We know from 67 that the order of this curve is a prime number, so we only have a single embedding degree.

To test potential embedding degrees  $k$ , say, in the range  $1 \dots 1000$ , we can invoke Sage and compute as follows:

```

sage: p = 1157920892373161954235709850086879078532699846656405
      64039457584007908834671663
sage: n = 1157920892373161954235709850086879078528375642790749
      04382605163141518161494337
sage: for k in range(1,1000):
.....:     if (p^k-1)%n == 0:
.....:         break
sage: k
999

```

We see that Secp256k1 has at least no embedding degree  $k < 1000$ , which renders Secp256k1 a curve that has no small embedding degree. This property will be of importance later on.

**Elliptic Curves over extension fields** Suppose that  $p$  is a prime number, and  $\mathbb{F}_p$  its associated prime field. We know from equation 4.29 that the fields  $\mathbb{F}_{p^m}$  are extensions of  $\mathbb{F}_p$  in the sense

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that  $\mathbb{F}_p$  is a subfield of  $\mathbb{F}_{p^m}$ . This implies that we can extend the affine plane that an elliptic curve is defined on by changing the base field to any extension field. To be more precise, let  $E(\mathbb{F}) = \{(x, y) \in \mathbb{F} \times \mathbb{F} \mid y^2 = x^3 + a \cdot x + b\}$  be an affine short Weierstraß curve, with parameters  $a$  and  $b$  taken from  $\mathbb{F}$ . If  $\mathbb{F}'$  is an extension field of  $\mathbb{F}$ , then we extend the domain of the curve by defining  $E(\mathbb{F}')$  as follows:

$$E(\mathbb{F}') = \{(x, y) \in \mathbb{F}' \times \mathbb{F}' \mid y^2 = x^3 + a \cdot x + b\} \quad (5.23)$$

While we did not change the defining parameters, we consider curve points from the affine plane over the extension field now. Since  $\mathbb{F} \subset \mathbb{F}'$ , it can be shown that the original elliptic curve  $E(\mathbb{F})$  is a sub-curve of the extension curve  $E(\mathbb{F}')$ .

*Example 84.* Consider the prime field  $\mathbb{F}_5$  from example 59 and the elliptic curve  $E_1(\mathbb{F}_5)$  from example 65. Since we know from XXX that  $\mathbb{F}_{5^2}$  is an extension field of  $\mathbb{F}_5$ , we can extend the definition of  $E_1(\mathbb{F}_5)$  to define a curve over  $\mathbb{F}_{5^2}$ :

$$E_1(\mathbb{F}_{5^2}) = \{(x, y) \in \mathbb{F} \times \mathbb{F} \mid y^2 = x^3 + x + 1\}$$

Since  $\mathbb{F}_{5^2}$  contains 25 points, in order to compute the set  $E_1(\mathbb{F}_{5^2})$ , we have to try  $25 \cdot 25 = 625$  pairs, which is probably a bit too much for the average motivated reader. Instead, we invoke Sage to compute the curve for us. To do, we so choose the representation of  $\mathbb{F}_{5^2}$  from XXX. We get:

3249	<code>sage: F5= GF(5)</code>	374
3250	<code>sage: F5t.&lt;t&gt; = F5[]</code>	375
3251	<code>sage: P = F5t(t^2+2)</code>	376
3252	<code>sage: P.is_irreducible()</code>	377
3253	<code>True</code>	378
3254	<code>sage: F5_2.&lt;t&gt; = GF(5^2, name='t', modulus=P)</code>	379
3255	<code>sage: E1F5_2 = EllipticCurve(F5_2, [1,1])</code>	380
3256	<code>sage: E1F5_2.order()</code>	381
3257	<code>27</code>	382

The curve  $E_1(\mathbb{F}_{5^2})$  consist of 27 points, in contrast to curve  $E_1(\mathbb{F}_5)$ , which consists of 9 points. Printing the points gives the following:

$$\begin{aligned} E_1(\mathbb{F}_{5^2}) = \{ & \mathcal{O}, (0, 4), (0, 1), (3, 4), (3, 1), (4, 3), (4, 2), (2, 4), (2, 1), \\ & (4t + 3, 3t + 4), (4t + 3, 2t + 1), (3t + 2, t), (3t + 2, 4t), \\ & (2t + 2, t), (2t + 2, 4t), (2t + 1, 4t + 4), (2t + 1, t + 1), \\ & (2t + 3, 3), (2t + 3, 2), (t + 3, 2t + 4), (t + 3, 3t + 1), \\ & (3t + 1, t + 4), (3t + 1, 4t + 1), (3t + 3, 3), (3t + 3, 2), (1, 4t) \} \end{aligned}$$

As we can see, curve  $E_1(\mathbb{F}_5)$  sits inside curve  $E(\mathbb{F}_{5^2})$ , which is implied by  $\mathbb{F}_5$  being a subfield of  $\mathbb{F}_{5^2}$ .

**Full torsion groups** The fundamental theorem of finite cyclic groups XXX implies that every prime factor  $n$  of a cyclic group's order defines a subgroup of the size of the prime factor. Such a subgroup is called an  $n$ -torsion group. We have seen many of those subgroups in the examples XXX and XXX.

When we consider elliptic curve extensions as defined in 5.23, we could ask what happens to the  $n$ -torsion groups in the extension. One might intuitively think that their extension just



parallels the extension of the curve. For example, when  $E(\mathbb{F}_p)$  is a curve over prime field  $\mathbb{F}_p$ , with some  $n$ -torsion group  $\mathbb{G}$  and when we extend the curve to  $E(\mathbb{F}_{p^m})$ , then there is a bigger  $n$ -torsion group such that  $\mathbb{G}$  is a subgroup. This might make intuitive sense, as  $E(\mathbb{F}_p)$  is a sub-curve of  $E(\mathbb{F}_{p^m})$ .

However, the actual situation is a bit more surprising than that. To see that, let  $\mathbb{F}_p$  be a prime field and let  $E(\mathbb{F}_p)$  be an elliptic curve of order  $r$ , with embedding degree  $k$  and  $n$ -torsion group  $E(\mathbb{F}_p)[n]$  for the same prime factor  $n$  of  $r$ . Then it can be shown that the  $n$ -torsion group  $E(\mathbb{F}_{p^m})[n]$  of a curve extension is equal to  $E(\mathbb{F}_p)[n]$ , as long as the power  $m$  is less than the embedding degree  $k$  of  $E(\mathbb{F}_p)$ .

However, for the prime power  $p^m$ , for any  $m \geq k$ ,  $E(\mathbb{F}_{p^m})[n]$  is strictly larger than  $E(\mathbb{F}_p)[n]$  and contains  $E(\mathbb{F}_p)[n]$  as a subgroup. We call the  $n$ -torsion group  $E(\mathbb{F}_{p^k})[n]$  of the extension of  $E$  over  $\mathbb{F}_{p^k}$  the **full  $n$ -torsion group** of that elliptic curve. It can be shown that it contains  $n^2$  many elements and consists of  $n + 1$  subgroups, one of which is  $E(\mathbb{F}_p)[n]$ .

So, roughly speaking, when we consider **towers of curve extensions**  $E(\mathbb{F}_{p^m})$  ordered by the prime power  $m$ , then the  $n$ -torsion group stays constant for every level  $m$ , that is smaller than the embedding degree, while it suddenly blossoms into a larger group on level  $k$  with  $n + 1$  subgroups, and then stays like that for any level  $m$  larger than  $k$ . In other words, once the extension field is big enough to find one more point of order  $n$  (that is not defined over the base field), then we actually find all of the points in the full torsion group.

*Example 85.* Consider curve  $E_1(\mathbb{F}_5)$  again. We know that it contains a 3-torsion group and that the embedding degree of 3 is 2. From this we can deduce that we can find the full 3-torsion group  $E_1[3]$  in the curve extension  $E_1(\mathbb{F}_{5^2})$ , the latter of which we computed in example 84.

Since that curve is small, in order to find the full 3-torsion, we can loop through all elements of  $E_1(\mathbb{F}_{5^2})$  and check the defining equation  $[3]P = \mathcal{O}$ . Invoking Sage, we compute as follows:

```
sage: INF = E1F5_2(0) # Point at infinity      383
sage: L_E1_3 = []      384
sage: for p in E1F5_2:      385
....:     if 3*p == INF:      386
....:         L_E1_3.append(p)      387
sage: E1_3 = Set(L_E1_3) # Full 3-torsion set      388
```

We get the following result:

$$E_1[3] = \{\mathcal{O}, (1, t), (1, 4t), (2, 1), (2, 4), (2t + 1, t + 1), (2t + 1, 4t + 4), (3t + 1, t + 4), (3t + 1, 4t + 1)\}$$

*Example 86.* Consider the tiny-jubjub curve from example 66. We know from example 82 that it contains a 5-torsion group and that the embedding degree of 5 is 4. This implies that we can find the full 5-torsion group  $TJJ\_13[5]$  in the curve extension  $TJJ\_13(\mathbb{F}_{13^4})$ .

To compute the full torsion, first observe that, since  $\mathbb{F}_{13^4}$  contains 28561 elements, computing  $TJJ\_13(\mathbb{F}_{13^4})$  means checking  $28561^2 = 815730721$  elements. From each of these curve points  $P$ , we then have to check the equation  $[5]P = \mathcal{O}$ . Doing this for 815730721 is a bit too slow even on a computer.

Fortunately, Sage has a way to loop through points of a given order efficiently. The following Sage code provides a way to compute the full torsion group:

```
sage: # define the extension field      389
sage: F13= GF(13) # prime field      390
sage: F13t.<t> = F13[] # polynomials over t      391
```

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```

3309 sage: P = F13t(t^4+2) # irreducible polynomial of degree 4      392
3310 sage: P.is_irreducible()                                         393
3311 True                                                             394
3312 sage: F13_4.<t> = GF(13^4, name='t', modulus=P) # F_{13^4}      395
3313 sage: TJJF13_4 = EllipticCurve(F13_4,[8,8]) # tiny-jubjub       396
3314     extension
3315 sage: # compute the full 5-torsion                                397
3316 sage: L_TJJF13_4_5 = []                                          398
3317 sage: INF = TJJF13_4(0)                                          399
3318 sage: for P in INF.division_points(5): # [5]P == INF            400
3319     ....:     L_TJJF13_4_5.append(P)                             401
3320 sage: len(L_TJJF13_4_5)                                          402
3321 25                                                                403
3322 sage: TJJF13_4_5 = Set(L_TJJF13_4_5)                            404

```

3323 As expected, we get a group that contains  $5^2 = 25$  elements. As it's rather tedious to write  
 3324 this group down, and as we don't need it in what follows, we forgo doing this. To see that the  
 3325 embedding degree 4 is actually the smallest prime power to find the full 5-torsion group, let's  
 3326 compute the 5-torsion group over of the tiny-jubjub curve of the extension field  $\mathbb{F}_{13^3}$ . We get  
 3327 the following:

```

3328 sage: # define the extension field                                405
3329 sage: P = F13t(t^3+2) # irreducible polynomial of degree 3      406
3330 sage: P.is_irreducible()                                         407
3331 True                                                             408
3332 sage: F13_3.<t> = GF(13^3, name='t', modulus=P) # F_{13^3}      409
3333 sage: TJJF13_3 = EllipticCurve(F13_3,[8,8]) # tiny-jubjub       410
3334     extension
3335 sage: # compute the 5-torsion                                    411
3336 sage: L_TJJF13_3_5 = []                                          412
3337 sage: INF = TJJF13_3(0)                                          413
3338 sage: for P in INF.division_points(5): # [5]P == INF            414
3339     ....:     L_TJJF13_3_5.append(P)                             415
3340 sage: len(L_TJJF13_3_5)                                          416
3341 5                                                                417
3342 sage: TJJF13_3_5 = Set(L_TJJF13_3_5) # full 5-torsion           418

```

3343 As we can see, the 5-torsion group of tiny-jubjub over  $\mathbb{F}_{13^3}$  is equal to the 5-torsion group  
 3344 of tiny-jubjub over  $\mathbb{F}_{13}$  itself.

3345 *Example 87.* Let's look at the curve Secp256k1. We know from example 67 that the curve is of  
 3346 some prime order  $r$ . Because of this, the only  $n$ -torsion group to consider is the curve itself, so  
 3347 the curve group is the  $r$ -torsion.

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However, in order to find the full  $r$ -torsion of Secp256k1, we need to compute the embed-  
 ding degree  $k$ . And as we have seen in XXX it is at least not small. However, we know from  
 Fermat's little theorem (page 21 ff.) that a finite embedding degree must exist. It can be shown  
 that it is given by the following 256-bit number:

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$$k = 192986815395526992372618308347813175472927379845817397100860523586360249056$$

3348 This means that the embedding degree is **very large**, which implies that the field extension  $\mathbb{F}_{p^k}$  is  
 3349 very large too. To understand how big  $\mathbb{F}_{p^k}$  is, recall that an element of  $\mathbb{F}_{p^m}$  can be represented as



3350 a string  $[x_0, \dots, x_m]$  of  $m$  elements, each containing a number from the prime field  $\mathbb{F}_p$ . Now, in  
 3351 the case of Secp256k1, such a representation has  $k$ -many entries, each of them 256 bits in size.  
 3352 So, without any optimizations, representing such an element would need  $k \cdot 256$  bits, which is  
 3353 too much to be represented in the observable universe.

3354 **Torsion subgroups** As we have stated above, any full  $n$ -torsion group contains  $n + 1$  cyclic  
 3355 subgroups, two of which are of particular interest in pairing-based elliptic curve cryptography.  
 3356 To characterize these groups, we need to consider the so-called **Frobenius endomorphism** of  
 3357 an elliptic curve  $E(\mathbb{F})$  over some finite field  $\mathbb{F}$  of characteristic  $p$ :

$$\pi : E(\mathbb{F}) \rightarrow E(\mathbb{F}) : \begin{array}{ccc} (x, y) & \mapsto & (x^p, y^p) \\ \mathcal{O} & \mapsto & \mathcal{O} \end{array} \quad (5.24)$$

3358 It can be shown that  $\pi$  maps curve points to curve points. The first thing to note is that, in case  
 3359  $\mathbb{F}$  is a prime field, the Frobenius endomorphism acts trivially, since  $(x^p, y^p) = (x, y)$  on prime  
 3360 fields due to Fermat's little theorem (page 21 ff.). This means that the Frobenius map is more  
 3361 interesting over prime field extensions.

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3362 With the Frobenius map at hand, we can characterize two important subgroups of the full  
 3363  $n$ -torsion. The first subgroup is the  $n$ -torsion group that already exists in the curve over the  
 3364 base field. In pairing-based cryptography, this group is usually written as  $\mathbb{G}_1$ , assuming that the  
 3365 prime factor  $n$  in the definition is implicitly given. Since we know that the Frobenius map acts  
 3366 trivially on curves over the prime field, we can define  $\mathbb{G}_1$  as follows:

$$\mathbb{G}_1[n] := \{(x, y) \in E[n] \mid \pi(x, y) = (x, y)\} \quad (5.25)$$

3367 In more mathematical terms, this definition means that  $\mathbb{G}_1$  is the **Eigenspace** of the Frobenius  
 3368 map with respect to the **Eigenvalue 1**.

3369 It can be shown that there is another subgroup of the full  $n$ -torsion group that can be char-  
 3370 acterized by the Frobenius map. In the context of so-called **type 3 pairing-based cryptography**,  
 3371 this subgroup is usually called  $\mathbb{G}_2$  and it is defined as follows:

$$\mathbb{G}_2[n] := \{(x, y) \in E[n] \mid \pi(x, y) = [p](x, y)\} \quad (5.26)$$

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3372 In mathematical terms,  $\mathbb{G}_2$  is the **Eigenspace** of the Frobenius map with respect to the  
 3373 **Eigenvalue  $p$** .

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3374 *Notation and Symbols 9.* If the prime factor  $n$  of a curve's order is clear from the context, we  
 3375 sometimes simply write  $\mathbb{G}_1$  and  $\mathbb{G}_2$  to mean  $\mathbb{G}_1[n]$  and  $\mathbb{G}_2[n]$ , respectively.

3376 It should be noted, however that other definitions of  $\mathbb{G}_2$  also exists in the literature. However,  
 3377 in the context of pairing-based cryptography, this is the most common one. It is particularly  
 3378 useful because we can define hash functions that map into  $\mathbb{G}_2$ , which is not possible for all  
 3379 subgroups of the full  $n$ -torsion.

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3380 *Example 88.* Consider the curve  $E_1(\mathbb{F}_5)$  from example 65 again. As we have seen, this curve  
 3381 has the embedding degree  $k = 2$ , and a full 3-torsion group is given as follows:

$$\begin{aligned} E_1[3] = \{ & \mathcal{O}, (2, 1), (2, 4), (1, t), (1, 4t), (2t + 1, t + 1), \\ & (2t + 1, 4t + 4), (3t + 1, t + 4), (3t + 1, 4t + 1) \} \end{aligned} \quad (5.27)$$

3382 According to the general theory,  $E_1[3]$  contains 4 subgroups, and we can characterize the  
 3383 subgroups  $\mathbb{G}_1$  and  $\mathbb{G}_2$  using the Frobenius endomorphism. Unfortunately, at the time of writing,

Sage does not have a predefined Frobenius endomorphism for elliptic curves, so we have to use the Frobenius endomorphism of the underlying field as a temporary workaround. We compute as follows:

```

3387 sage: L_G1 = []
3388 sage: for P in E1_3:
3389     ....:     PiP = E1F5_2([a.frobenius() for a in P]) # pi(P)
3390     ....:     if P == PiP:
3391     ....:         L_G1.append(P)
3392 sage: G1 = Set(L_G1)

```

As expected, the group  $\mathbb{G}_1 = \{\mathcal{O}, (2, 4), (2, 1)\}$  is identical to the 3-torsion group of the (unextended) curve over the prime field  $E_1(\mathbb{F}_5)$ . We can use almost the same algorithm to compute the group  $\mathbb{G}_2$  and get the following:

```

3396 sage: L_G2 = []
3397 sage: for P in E1_3:
3398     ....:     PiP = E1F5_2([a.frobenius() for a in P]) # pi(P)
3399     ....:     pP = 5*P # [5]P
3400     ....:     if pP == PiP:
3401     ....:         L_G2.append(P)
3402 sage: G2 = Set(L_G2)

```

Thus, we have computed the the second subgroup of the full 3-torsion group of curve  $E_1$  as the set  $\mathbb{G}_2 = \{\mathcal{O}, (1, t), (1, 4t)\}$ .

*Example 89.* Consider the tiny-jubjub curve *TJJ\_13* from example 66. In example 86, we computed its full 5 torsion, which is a group that has 6 subgroups. We compute  $G_1$  using Sage as follows:

```

3408 sage: L_TJJ_G1 = []
3409 sage: for P in TJJF13_4_5:
3410     ....:     PiP = TJJF13_4([a.frobenius() for a in P]) # pi(P)
3411     ....:     if P == PiP:
3412     ....:         L_TJJ_G1.append(P)
3413 sage: TJJ_G1 = Set(L_TJJ_G1)

```

We get  $\mathbb{G}_1 = \{\mathcal{O}, (7, 2), (8, 8), (8, 5), (7, 11)\}$

```

3415 sage: L_TJJ_G1 = []
3416 sage: for P in TJJF13_4_5:
3417     ....:     PiP = TJJF13_4([a.frobenius() for a in P]) # pi(P)
3418     ....:     pP = 13*P # [5]P
3419     ....:     if pP == PiP:
3420     ....:         L_TJJ_G1.append(P)
3421 sage: TJJ_G1 = Set(L_TJJ_G1)

```

$\mathbb{G}_2 = \{\mathcal{O}, (9t^2 + 7, t^3 + 11t), (9t^2 + 7, 12t^3 + 2t), (4t^2 + 7, 5t^3 + 10t), (4t^2 + 7, 8t^3 + 3t)\}$

*Example 90.* Consider Bitcoin's curve *Secp256k1* again. Since the group  $\mathbb{G}_1$  is identical to the torsion group of the unextended curve, and since *Secp256k1* has prime order, we know that, in this case,  $\mathbb{G}_1$  is identical to *Secp256k1*. It is however, infeasible not to compute not only  $\mathbb{G}_2$  itself, but to even compute an average element of  $\mathbb{G}_2$ , as elements need too much storage to be

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3427 representable in this universe.

3428 **The Weil pairing** In this part, we consider a pairing function defined on the subgroups  $\mathbb{G}_1[r]$   
 3429 and  $\mathbb{G}_2[r]$  of the full  $r$ -torsion  $E[r]$  of a short Weierstraß elliptic curve. To be more precise, let  
 3430  $E(\mathbb{F}_p)$  be an elliptic curve of embedding degree  $k$  such that  $r$  is a prime factor of its order. Then  
 3431 the **Weil pairing** is a bilinear, non-degenerate map:

$$e(\cdot, \cdot) : \mathbb{G}_1[r] \times \mathbb{G}_2[r] \rightarrow \mathbb{F}_{p^k} ; (P, Q) \mapsto (-1)^r \cdot \frac{f_{r,P}(Q)}{f_{r,Q}(P)} \quad (5.28)$$

The extension field elements  $f_{r,P}(Q), f_{r,Q}(P) \in \mathbb{F}_{p^k}$  are computed by **Miller's algorithm**:

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algorithm

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**Algorithm 7** Miller's algorithm for short Weierstraß curves  $y^2 = x^3 + ax + b$

---

**Require:**  $r > 3, P \in E[r], Q \in E[r]$  and

$b_0, \dots, b_t \in \{0, 1\}$  with  $r = b_0 \cdot 2^0 + b_1 \cdot 2^1 + \dots + b_t \cdot 2^t$  and  $b_t = 1$

**procedure** MILLER'S ALGORITHM( $P, Q$ )

**if**  $P = \mathcal{O}$  or  $Q = \mathcal{O}$  or  $P = Q$  **then**

**return**  $f_{r,P}(Q) \leftarrow (-1)^r$

**end if**

$(x_T, y_T) \leftarrow (x_P, y_P)$

$f_1 \leftarrow 1$

$f_2 \leftarrow 1$

**for**  $j \leftarrow t - 1, \dots, 0$  **do**

$m \leftarrow \frac{3 \cdot x_T^2 + a}{2 \cdot y_T}$

$f_1 \leftarrow f_1^2 \cdot (y_Q - y_T - m \cdot (x_Q - x_T))$

$f_2 \leftarrow f_2^2 \cdot (x_Q + 2x_T - m^2)$

$x_{2T} \leftarrow m^2 - 2x_T$

$y_{2T} \leftarrow -y_T - m \cdot (x_{2T} - x_T)$

$(x_T, y_T) \leftarrow (x_{2T}, y_{2T})$

**if**  $b_j = 1$  **then**

$m \leftarrow \frac{y_T - y_P}{x_T - x_P}$

$f_1 \leftarrow f_1 \cdot (y_Q - y_T - m \cdot (x_Q - x_T))$

$f_2 \leftarrow f_2 \cdot (x_Q + (x_P + x_T) - m^2)$

$x_{T+P} \leftarrow m^2 - x_T - x_P$

$y_{T+P} \leftarrow -y_T - m \cdot (x_{T+P} - x_T)$

$(x_T, y_T) \leftarrow (x_{T+P}, y_{T+P})$

**end if**

**end for**

$f_1 \leftarrow f_1 \cdot (x_Q - x_T)$

**return**  $f_{r,P}(Q) \leftarrow \frac{f_1}{f_2}$

**end procedure**

---

3432 Understanding how the algorithm works in detail requires the concept of **divisors**, which is  
 3433 outside of the scope this book. The interested reader might look at XXX.

3434 In real-world applications of pairing-friendly elliptic curves, the embedding degree is usu-  
 3435 ally a small number like 2, 4, 6 or 12, and the number  $r$  is the largest prime factor of the curve's  
 3436 order.  
 3437

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3438 *Example 91.* Consider curve  $E_1(\mathbb{F}_5)$  from example 65. Since the only prime factor of the  
 3439 group's order is 3, we cannot compute the Weil pairing on this group using our definition of  
 3440 Miller's algorithm. In fact, since  $\mathbb{G}_1$  is of order 3, executing the if statement on line XXX will  
 3441 lead to a "division by zero" error in the computation of the slope  $m$ .

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ence

*Example 92.* Consider the tiny-jubjub curve  $TJJ\_13(\mathbb{F}_{13})$  from example 66 again. We want to  
 instantiate the general definition of the Weil pairing for this example. To do so, recall that, as we  
 have see in example 82, its embedding degree is 4, and that we have the following type-3 pairing  
 groups (where  $\mathbb{G}_1$  and  $\mathbb{G}_2$  are subgroups of the full 5-torsion found in the curve  $TJJ\_13(\mathbb{F}_{13^4})$ ):

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$$\mathbb{G}_1 = \{\mathcal{O}, (7, 2), (8, 8), (8, 5), (7, 11)\}$$

$$\mathbb{G}_2 = \{\mathcal{O}, (9t^2 + 7, t^3 + 11t), (9t^2 + 7, 12t^3 + 2t), (4t^2 + 7, 5t^3 + 10t), (4t^2 + 7, 8t^3 + 3t)\}$$

3442 The type-3 Weil pairing is a map  $e(\cdot, \cdot) : \mathbb{G}_1 \times \mathbb{G}_2 \rightarrow \mathbb{F}_{13^4}$ . From the first if-statement in  
 3443 Miller's algorithm, we can deduce that  $e(\mathcal{O}, Q) = 1$  as well as  $e(P, \mathcal{O}) = 1$  for all arguments  
 3444  $P \in \mathbb{G}_1$  and  $Q \in \mathbb{G}_2$ . In order to compute a non-trivial Weil pairing, we choose the arguments  
 3445  $P = (7, 2)$  and  $Q = (9t^2 + 7, 12t^3 + 2t)$ .

3446 To compute the pairing  $e((7, 2), (9t^2 + 7, 12t^3 + 2t))$ , we have to compute the extension field  
 3447 elements  $f_{5,P}(Q)$  and  $f_{5,Q}(P)$  by applying Miller's algorithm. Do do so, observe that we have  
 3448  $5 = 1 \cdot 2^0 + 0 \cdot 2^1 + 1 \cdot 2^2$ , so we get  $t = 2$  as well as  $b_0 = 1, b_1 = 0$  and  $b_2 = 1$ . The loop therefore  
 3449 needs to be executed two times.

Computing  $f_{5,P}(Q)$ , we initiate  $(x_T, y_T) = (7, 2)$  as well as  $f_1 = 1$  and  $f_2 = 1$ . Then we  
 proceed as follows:

$j$	$b_j$	$m$	$f_1$	$f_2$	$x_{2T}$	$y_{2T}$	$x_{T+P}$	$y_{T+P}$
1	.							

$$\begin{aligned}
m &= \frac{3 \cdot x_T^2 + a}{2 \cdot y_T} \\
&= \frac{3 \cdot 2^2 + 1}{2 \cdot 4} = \frac{3}{3} \\
&= 1
\end{aligned}$$

$$\begin{aligned}
f_1 &= f_1^2 \cdot (y_Q - y_T - m \cdot (x_Q - x_T)) \\
&= 1^2 \cdot (t - 4 - 1 \cdot (1 - 2)) = t - 4 + 1 \\
&= t + 2
\end{aligned}$$

$$\begin{aligned}
f_2 &= f_2^2 \cdot (x_Q + 2x_T - m^2) \\
&= 1^2 \cdot (1 + 2 \cdot 2 - 1^2) = (1 + 4 - 1) \\
&= 4
\end{aligned}$$

$$\begin{aligned}
x_{2T} &= m^2 - 2x_T \\
&= 1^2 - 2 \cdot 2 = -3 \\
&= 2
\end{aligned}$$

$$\begin{aligned}
y_{2T} &= -y_T - m \cdot (x_{2T} - x_T) \\
&= -4 - 1 \cdot (2 - 2) = -4 \\
&= 1
\end{aligned}$$

We update  $(x_T, y_T) = (2, 1)$  and, since  $b_0 = 1$ , we have to execute the if statement on line XXX in the **for** loop. However, since we only loop a single time, we don't need to compute  $y_{T+P}$ , since we only need the updated  $x_T$  in the final step. We get:

$$\begin{aligned}
m &= \frac{y_T - y_P}{x_T - x_P} \\
&= \frac{1 - 4}{2 - x_P}
\end{aligned}$$

$$f_1 = f_1 \cdot (y_Q - y_T - m \cdot (x_Q - x_T))$$

$$f_2 = f_2 \cdot (x_Q + (x_P + x_T) - m^2)$$

$$x_{T+P} = m^2 - x_T - x_P$$

add reference

should all lines of all algorithms be numbered?

## 3450 5.3 Hashing to Curves

3451 Elliptic curve cryptography frequently requires the ability to hash data onto elliptic curves. If  
3452 the order of the curve is not a prime number, hashing to prime number subgroups is also of

importance. In the context of pairing-friendly curves, it is also sometimes necessary to hash specifically onto the group  $\mathbb{G}_1$  or  $\mathbb{G}_2$ .

As we have seen in section 4.1.2, many general methods are known for hashing into groups in general, and finite cyclic groups in particular. As elliptic groups are cyclic, those methods can be utilized in this case, too. However, in what follows we want to describe some methods specific to elliptic curves that are frequently used in real-world applications.

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**Try-and-increment hash functions** One of the most straight-forward ways of hashing a bit-string onto an elliptic curve point in a secure way is to use a cryptographic hash function together with one of the methods we described in section 4.1.2 to hash to the modular arithmetic base field of the curve. Ideally, the hash function generates an image that is at least one bit longer than the bit representation of the base field modulus.

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The image in the base field can then be interpreted as the  $x$  coordinate of the curve point, and the two possible  $y$  coordinates are derived from the curve equation, while one of the bits that exceeded the modulus determines which of the two  $y$  coordinates to choose.

Such an approach would be deterministic and easy to implement, and it would conserve the cryptographic properties of the original hash function. However, not all  $x$  coordinates generated in such a way will result in quadratic residues when inserted into the defining equation. It follows that not all field elements give rise to actual curve points. In fact, on a prime field, only half of the field elements are quadratic residues. Hence, assuming an even distribution of the hash values in the field, this method would fail to generate a curve point in about half of the attempts.

One way to account for this problem is the so-called **try-and-increment** method. Its basic assumption is that, when hashing different values, the result will eventually lead to a valid curve point.

Therefore, instead of simply hashing a string  $s$  to the field, we hash the concatenation of  $s$  with additional bytes to the field instead. In other words, we use a try-and-increment hash as described in 5. If the first try of hashing to the field does not result in a valid curve point, the counter is incremented, and the hashing is repeated again. This is done until a valid curve point is found.

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This method has a number of advantages: It is relatively easy to implement in code, and it maintains the cryptographic properties of the original hash function. However, it is not guaranteed to find a valid curve point, as there is a chance that all possible values in the chosen size of the counter will fail to generate a quadratic residue. Fortunately, it is possible to make the probability for this arbitrarily small by choosing large enough counters and relying on the (approximate) uniformity of the hash-to-field function.

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rithm is  
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If the curve is not of prime order, the result will be a general curve point that might not be in the “large” prime-order subgroup. In this case, a **cofactor clearing** step is then necessary to project the curve point onto the subgroup. This is done by scalar multiplication with the cofactor of prime order with respect to the curves order.

*Example 93.* Consider the tiny-jubjub curve from example 66. We want to construct a try-and-increment hash function that hashes a binary string  $s$  of arbitrary length onto the large prime-order subgroup of size 5.

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Since the curve, as well as our targeted subgroup, is defined over the field  $\mathbb{F}_{13}$ , and the binary representation of 13 is  $13.bits() = 1101$ , we apply SHA256 from Sage’s hashlib library on the concatenation  $s||c$  for some binary counter string, and use the first 4 bits of the image to try to hash into  $\mathbb{F}_{13}$ . In case we are able to hash to a value  $z$  such that  $z^3 + 8 \cdot z + 8$  is a quadratic residue in  $\mathbb{F}_{13}$ , we use the 5-th bit to decide which of the two possible roots of  $z^3 + 8 \cdot z + 8$  we

**Algorithm 8** Hash-to- $E(\mathbb{F}_r)$ **Require:**  $r \in \mathbb{Z}$  with  $r.\text{nbits}() = k$  and  $s \in \{0, 1\}^*$ **Require:** Curve equation  $y^2 = x^3 + ax + b$  over  $\mathbb{F}_r$ **procedure** TRY-AND-INCREMENT( $r, k, s$ ) $c \leftarrow 0$ **repeat** $s' \leftarrow s || c.\text{bits}()$  $z \leftarrow H(s')_0 \cdot 2^0 + H(s')_1 \cdot 2^1 + \dots + H(s')_k \cdot 2^k$  $x \leftarrow z^3 + a \cdot z + b$  $c \leftarrow c + 1$ **until**  $z < r$  and  $x^{\frac{r-1}{2}} \bmod r = 1$ **if**  $H(s')_{k+1} == 0$  **then** $y \leftarrow \sqrt{x} \#(\text{root in } \mathbb{F}_r)$ **else** $y \leftarrow r - \sqrt{x} \#(\text{root in } \mathbb{F}_r)$ **end if****return**  $(x, y)$ **end procedure****Ensure:**  $(x, y) \in E(\mathbb{F}_r)$ 

will choose as the  $y$  coordinate. The result is a curve point different from the point at infinity. To project it to a point of  $\mathbb{G}_1$ , we multiply it with the cofactor 4. If the result is still not the point at infinity, it is the result of the hash.

To make this concrete, let  $s = '10011001111010110100000111'$  be our binary string that we want to hash onto  $\mathbb{G}_1$ . We use a 4-bit binary counter starting at zero, that is, we choose  $c = 0000$ . Invoking Sage, we define the try-hash function as follows:

```

sage: import hashlib
sage: def try_hash(s, c):
.....:     s_1 = s+c
.....:     hasher = hashlib.sha256(s_1.encode('utf-8'))
.....:     digest = hasher.hexdigest()
.....:     d = Integer(digest, base=16)
.....:     sign = d.str(2)[-5:-4]
.....:     d = d.str(2)[-4:]
.....:     z = Integer(d, base=2)
.....:     return (z, sign)
sage: try_hash('10011001111010110100000111', '0000')
(15, '1')
```

As we can see, our first attempt to hash into  $\mathbb{F}_{13}$  was not successful, as 15 is not a number in  $\mathbb{F}_{13}$ , so we increment the binary counter by 1 and try again:

```

sage: try_hash('10011001111010110100000111', '0001')
(3, '0')
```

With this try, we found a hash into  $\mathbb{F}_{13}$ . However, this point is not guaranteed to define a curve point. To see that, we insert  $z = 3$  into the right side of the Weierstraß equation of the tiny-jubjub curve, and compute  $3^3 + 8 * 3 + 8 = 7$ . However, 7 is not a quadratic residue in



3525  $\mathbb{F}_{13}$ , since  $7^{\frac{13-1}{2}} = 7^6 = 12 = -1$ . This means that 3 is not a suitable point, and we have to  
 3526 increment the counter two more times:

```

3527 sage: try_hash('10011001111010110100000111', '0010')      459
3528 (3, '0')                                                    460
3529 sage: try_hash('10011001111010110100000111', '0011')      461
3530 (6, '1')                                                    462

```

Since  $6^3 + 8 \cdot 6 + 8 = 12$ , and we have  $\sqrt{12} \in \{5, 8\}$ , we finally found the valid  $x$  coordinate  $x = 6$  for the curve point hash. Now, since the sign bit of this hash is 1, we choose the larger root  $y = 8$  as the  $y$  coordinate and get the following hash which is a valid curve point point on the tiny-jubjub curve:

$$H('10011001111010110100000111') = (6, 8)$$

In order to project this onto the “large” prime-order subgroup, we have to do cofactor clearing, that is, we have to multiply the point with the cofactor 4. We get the following:

$$[4](6, 8) = \mathcal{O}$$

3531 This means that the hash value is still not right. We therefore have to increment the counter  
 3532 two more times again, until we finally find a correct hash to  $\mathbb{G}_1$ :

```

3533 sage: try_hash('10011001111010110100000111', '0100')      463
3534 (0, '1')                                                    464
3535 sage: try_hash('10011001111010110100000111', '0101')      465
3536 (12, '0')                                                  466

```

Since  $12^3 + 8 \cdot 12 + 8 = 12$ , and we have  $\sqrt{12} \in \{5, 8\}$ , we found another valid  $x$  coordinate  $x = 12$  for the curve point hash. Since the sign bit of this hash is 0, we choose the smaller root  $y = 5$  as the  $y$  coordinate, and get the following hash, which is a valid curve point point on the tiny-jubjub curve:

$$H('10011001111010110100000111') = (12, 5)$$

In order to project this onto the “large” prime-order subgroup we have to do cofactor clearing, again? that is, we have to multiply the point with the cofactor 4. We get the following:

$$[4](12, 5) = (8, 5)$$

3537 So, hashing the binary string '10011001111010110100000111' onto  $\mathbb{G}_1$  gives the hash value  
 3538 (8, 5) as a result.

## 3539 5.4 Constructing elliptic curves

3540 Cryptographically secure elliptic curves like Secp256k1 from example 67 have been known for  
 3541 quite some time. Given the latest advancements of cryptography, however, it is often necessary  
 3542 to design and instantiate elliptic curves from scratch that satisfy certain very specific properties.

3543 For example, in the context of SNARK development, it was necessary to design a curve that  
 3544 can be efficiently implemented inside of a so-called **circuit** in order to enable primitives like  
 3545 elliptic curve **signature schemes** in a zero-knowledge proof. Such a curve is given by the Baby-  
 3546 jubjub curve in XXX, and we have paralleled its definition by introducing the tiny-jubjub curve

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from example 66. As we have seen, those curves are instances of so-called twisted Edwards curves, and as such have easy to implement addition laws that work without branching. However, we introduced the tiny-jubjub curve out of thin air, as we just gave the curve parameters without explaining how we came up with them.

Another requirement in the context of many so-called **pairing-based zero-knowledge proofing systems** is the existence of a suitable, pairing-friendly curve with a specified security level and a low embedding degree as defined in 5.2.0.1. Famous examples are the BLS\_12 and the NMT curves.

The major goal of this section is to explain the most important method of designing elliptic curves with predefined properties from scratch, called the **complex multiplication method**. We will apply this method in section XXX to synthesize a particular BLS\_6 curve, which is one of the most insecure curves, but it will serve as the main curve to build our pen-and-paper SNARKs on. As we will see, this curve has a “large” prime factor subgroup of order 13, which implies that we can use our tiny-jubjub curve to implement certain elliptic curve cryptographic primitives in circuits over that BLS\_6 curve.

Before we introduce the complex multiplication method, we have to explain a few properties of elliptic curves that are of key importance in understanding the complex multiplication method.

**The Trace of Frobenius** To understand the complex multiplication method of elliptic curves, we have to define the so-called **trace** of an elliptic curve first.

We know from XXX that elliptic curves over finite fields are products of cyclic groups of finite order. Therefore, an interesting question is whether it is possible to estimate the number of elements that this curve contains. Since an affine short Weierstraß curve consists of pairs  $(x, y)$  of elements from a finite field  $\mathbb{F}_q$  plus the point at infinity, and the field  $\mathbb{F}_q$  contains  $q$  elements, the number of curve points cannot be arbitrarily large, since it can contain at most  $q^2 + 1$  many elements.

There is however, a more precise estimation, usually called the **Hasse bound**. To understand it, let  $E(\mathbb{F}_q)$  be an affine short Weierstraß curve over a finite field  $\mathbb{F}_q$  of order  $q$ , and let  $|E(\mathbb{F}_q)|$  be the order of the curve. Then there is an integer  $t \in \mathbb{Z}$ , called the **trace of Frobenius** of the curve, such that  $|t| \leq 2\sqrt{q}$  and the following equation holds:

$$|E(\mathbb{F}_q)| = q + 1 - t \quad (5.29)$$

A positive trace, therefore, implies that the curve contains less points than the underlying field, whereas a negative trace means that the curve contains more points. However, the estimation  $|t| \leq 2\sqrt{q}$  implies that the difference is not very large in either direction, and the number of elements in an elliptic curve is always approximately in the same order of magnitude as the size of the curve’s base field.

*Example 94.* Consider the elliptic curve  $E_1(\mathbb{F}_5)$  from example 65. We know that it contains 9 curve points. Since the order of  $\mathbb{F}_5$  is 5, we compute the trace of  $E_1(\mathbb{F})$  to be  $t = -3$ , since the Hasse bound is given by the following equation:

$$9 = 5 + 1 - (-3)$$

Indeed, we have  $|t| \leq 2\sqrt{q}$ , since  $\sqrt{5} > 2.23$  and  $|-3| = 3 \leq 4.46 = 2 \cdot 2.23 < 2 \cdot \sqrt{5}$ .

*Example 95.* To compute the trace of the tiny-jubjub curve, recall from example 74 that the order of  $TJJ_{13}$  is 20. Since the order of  $\mathbb{F}_{13}$  is 13, we can therefore use the Hasse bound and compute the trace as  $t = -6$ :

$$20 = 13 + 1 - (-6) \quad (5.30)$$

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3586 Again, we have  $|t| \leq 2\sqrt{q}$ , since  $\sqrt{13} > 3.60$  and  $|-6| = 6 \leq 7.20 = 2 \cdot 3.60 < 2 \cdot \sqrt{13}$ .

*Example 96.* To compute the trace of Secp256k1, recall from example 67 [that this curve is](#) defined over a prime field with  $p$  elements, and that the order of that group is given by  $r$ :

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$p = 115792089237316195423570985008687907853269984665640564039457584007908834671663$   
 $r = 115792089237316195423570985008687907852837564279074904382605163141518161494337$

Using the Hesse bound  $r = p + 1 - t$ , we therefore compute  $t = p + 1 - r$ , which gives the trace of curve Secp256k1 as follows:

$t = 432420386565659656852420866390673177327$

3587 As we can see, Secp256k1 contains less elements than its underlying field. However, the  
 3588 difference is tiny, since the order of Secp256k1 is in the same order of magnitude as the order  
 3589 of the underlying field. Compared to  $p$  and  $r$ ,  $t$  is tiny.

```
3590 sage: p = 1157920892373161954235709850086879078532699846656405 467
3591         64039457584007908834671663
3592 sage: r = 1157920892373161954235709850086879078528375642790749 468
3593         04382605163141518161494337
3594 sage: t = p + 1 - r 469
3595 sage: t.nbits() 470
3596 129 471
3597 sage: abs(RR(t)) <= 2*sqrt(RR(p)) 472
3598 True 473
```

3599 **The  $j$ -invariant** As we have seen in XXX, [two elliptic curves  \$E\_1\(\mathbb{F}\)\$  defined by  \$y^2 = x^3 + ax + b\$  and  \$E\_2\(\mathbb{F}\)\$  defined by  \$y^2 + a'x + b'\$  are strictly isomorphic if and only if there is a quadratic residue  \$d \in \mathbb{F}\$  such that  \$a' = ad^2\$  and  \$b' = bd^3\$ .](#)

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ence

3602 There is, however, a more general way to classify elliptic curves over finite fields  $\mathbb{F}_q$ , based  
 3603 on the so-called  **$j$ -invariant** of an elliptic curve with  $j(E(\mathbb{F}_q)) \in \mathbb{F}_q$ , as defined below:

$$j(E(\mathbb{F}_q)) = (1728 \bmod q) \frac{4 \cdot a^3}{4 \cdot a^3 + (27 \bmod q) \cdot b^2} \quad (5.31)$$

3604 A detailed description of the  $j$ -invariant is beyond the scope of this book. For our present  
 3605 purposes, it is sufficient to note that two elliptic curves  $E_1(\mathbb{F})$  and  $E_2(\mathbb{F}')$  are isomorphic over  
 3606 the **algebraic closures** of  $\mathbb{F}$  and  $\mathbb{F}'$ , if and only if  $\overline{\mathbb{F}} = \overline{\mathbb{F}'}$  and  $j(E_1) = j(E_2)$ .

algebraic  
closures

3607 So, the  $j$ -invariant is an important tool to classify elliptic curves and it is needed in the com-  
 3608 plex multiplication method to decide on an actual curve instantiation that implements abstractly  
 3609 chosen properties.

*Example 97.* Consider the elliptic curve  $E_1(\mathbb{F}_5)$  from example 65. [We compute its  \$j\$ -invariant](#) as follows:

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$$\begin{aligned} j(E_1(\mathbb{F}_5)) &= (1728 \bmod 5) \frac{4 \cdot 1^3}{4 \cdot 1^3 + (27 \bmod 5) \cdot 1^2} \\ &= 3 \frac{4}{4 + 2} \\ &= 3 \cdot 4 = 2 \end{aligned}$$

*Example 98.* Consider the elliptic curve *TJJ\_13* from example 66. We compute its *j*-invariant as follows:

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$$\begin{aligned}
 j(E_1(\mathbb{F}_5)) &= (1728 \bmod 13) \frac{4 \cdot 8^3}{4 \cdot 8^3 + (27 \bmod 13) \cdot 8^2} \\
 &= 12 \cdot \frac{4 \cdot 5}{4 \cdot 5 + 1 \cdot 12} \\
 &= 12 \cdot \frac{7}{7 + 12} \\
 &= 12 \cdot 7 \cdot 6^{-1} \\
 &= 12 \cdot 7 \cdot 11 \\
 &01
 \end{aligned}$$

*Example 99.* Consider *Sepc256k1* from example *Sepc256k1*. We compute its *j*-invariant using Sage:

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```

3612 sage: p = 1157920892373161954235709850086879078532699846656405 474
3613       64039457584007908834671663
3614 sage: F = GF(p) 475
3615 sage: j = F(1728) * ( (F(4) * F(0) ^3) / (F(4) * F(0) ^3 + F(27) * F(7) ^2) ) 476
3616 sage: j == F(0) 477
3617 True 478

```

**The Complex Multiplication Method** As we have seen in the previous sections, elliptic curves have various defining properties, like their order, their prime factors, the embedding degree, or the cardinality (number of elements) of the base field. The **complex multiplication** (CM) method provides a practical way of constructing elliptic curves with pre-defined restrictions on the order and the base field.

The method usually starts by choosing a base field  $\mathbb{F}_q$  of the curve  $E(\mathbb{F}_q)$  we want to construct such that  $q = p^m$  for some prime number  $p$ , and “ $m \in \mathbb{N}$  with  $m \geq 1$ . We assume  $p > 3$  to simplify things in what follows.

Next, the trace of Frobenius  $t \in \mathbb{Z}$  of the curve is chosen such that  $p$  and  $t$  are coprime, that is,  $\gcd(p, t) = 1$  holds true. The choice of  $t$  also defines the curve’s order  $r$ , since  $r = p + 1 - t$  by the Hasse bound (equation 5.29), so choosing  $t$  will define the large order subgroup as well as all small cofactors.  $r$  has to be defined in such a way that the elliptic curve meets the security requirements of the application it is designed for.

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Note that the choice of  $p$  and  $t$  also determines the embedding degree  $k$  of any prime-order subgroup of the curve, since  $k$  is defined as the smallest number such that the prime order  $n$  divides the number  $q^k - 1$ .

$$\begin{aligned}
 D &< 0 \\
 D \bmod 4 &= 0 \text{ or } D \bmod 4 = 1 \\
 4q &= t^2 + |D|v^2
 \end{aligned} \tag{5.32}$$

In order for the complex multiplication method to work, neither  $q$  nor  $t$  can be arbitrary, but must be chosen in such a way that two additional integers  $D \in \mathbb{Z}$  and  $v \in \mathbb{Z}$  exist and the following conditions hold:

If such numbers exist, we call  $D$  the **CM-discriminant**, and we know that we can construct a curve  $E(\mathbb{F}_q)$  over a finite field  $\mathbb{F}_q$  such that the order of the curve is  $|E(\mathbb{F}_q)| = q + 1 - t$ .

It is the content of the complex multiplication method to actually construct such a curve, that is finding the parameters  $a$  and  $b$  from  $\mathbb{F}_q$  in the defining Weierstraß equation such that the curve has the desired order  $r$ .

Finding solutions to equation 5.29,= can be achieved in different ways, but we will forego the fine detail here. In general, it can be said that there are well-known constraints for elliptic curve families (e.g. the BLS (ECT) families) that provides families of solutions. In what follows, we will look at one type curve in the BLS-family, which gives an entire range of solutions. Are we looking at a subtype of BLS or is BLS the specific type we're referring to?

Assuming that the proper parameters  $q$ ,  $t$ ,  $D$  and  $v$  are found, we have to compute the so-called **Hilbert class polynomial**  $H_D \in \mathbb{Z}[x]$  of the CM-discriminant  $D$ , which is a polynomial with integer coefficients. To do so, we first have to compute the following set:

$$ICG(D) = \{(A, B, C) \mid A, B, C \in \mathbb{Z}, D = B^2 - 4AC, \gcd(A, B, C) = 1,$$

$$|B| \leq A \leq \sqrt{\frac{|D|}{3}}, A \leq C, \text{ if } B < 0 \text{ then } |B| < A < C\}$$

One way to compute this set is to first compute the integer  $A_{\max} = \text{Floor}(\sqrt{\frac{|D|}{3}})$ , then loop through all the integers  $A$  in the range  $[0, \dots, A_{\max}]$ , as well as through all the integers  $B$  in the range  $[-A_{\max}, \dots, A_{\max}]$ , then see if there is an integer  $C$  that satisfies  $D = B^2 - 4AC$  and the rest of the requirements in XXX.

To compute the Hilbert class polynomial, the so-called  **$j$ -function** (or  $j$ -invariant) is needed, which is a complex function defined on the upper half  $\mathbb{H}$  of the complex plane  $\mathbb{C}$ , usually written as follows: is this the same as equation eq:j-invariant1 No they just sound very similar. But they are very different?

$$j : \mathbb{H} \rightarrow \mathbb{C} \quad (5.33)$$

Roughly speaking, what this means is that the  $j$ -functions takes complex numbers  $(x + i \cdot y)$  with a positive imaginary part  $y > 0$  as inputs and returns a complex number  $j(x + i \cdot y)$  as a result.

For the purposes of this book, it is not important to understand the  $j$ -function in detail, and we can use Sage to compute it in a similar way that we would use Sage to compute any other well-known function. It should be noted, however, that the computation of the  $j$ -function in Sage is sometimes prone to precision errors. For example, the  $j$ -function has a root in  $\frac{-1+i\sqrt{3}}{2}$ , which Sage only approximates. Therefore, when using Sage to compute the  $j$ -function, we need to take precision loss into account and possibly round to the nearest integer.

```
sage: z = ComplexField(100) (0, 1) 479
sage: z # (0+1i) 480
1.00000000000000000000000000000000*I 481
sage: elliptic_j(z) 482
1728.000000000000000000000000000000 483
sage: # j-function only defined for positive imaginary 484
arguments
sage: z = ComplexField(100) (1, -1) 485
sage: try: 486
.....:     elliptic_j(z) 487
.....: except PariError: 488
```

```

3675     ....:      pass                                     489
3676     sage: # root at (-1+i sqrt(3))/2                  490
3677     sage: z = ComplexField(100)(-1, sqrt(3))/2        491
3678     sage: elliptic_j(z)                                492
3679     -2.6445453750358706361219364880e-88              493
3680     sage: elliptic_j(z).imag().round()                 494
3681     0                                                    495
3682     sage: elliptic_j(z).real().round()                  496
3683     0                                                    497

```

3684 With a way to compute the  $j$ -function and the precomputed set  $ICG(D)$  at hand, we can now  
 3685 compute the Hilbert class polynomial as follows:

$$H_D(x) = \prod_{(A,B,C) \in ICG(D)} \left( x - j \left( \frac{-B + \sqrt{D}}{2A} \right) \right) \quad (5.34)$$

3686 In other words, we loop over all elements  $(A, B, C)$  from the set  $ICG(D)$  and compute the  
 3687  $j$ -function at the point  $\frac{-B + \sqrt{D}}{2A}$ , where  $D$  is the CM-discriminant that we chose in a previous  
 3688 step. The result defines a factor of the Hilbert class polynomial and all factors are multiplied  
 3689 together.

3690 It can be shown that the Hilbert class polynomial is an integer polynomial, but actual com-  
 3691 putations need high-precision arithmetics to avoid approximation errors that usually occur in  
 3692 computer approximations of the  $j$ -function (as shown above). So, in case the calculated Hilbert  
 3693 class polynomial does not have integer coefficients, we need to round the result to the nearest  
 3694 integer. Given that the precision we used was high enough, the result will be correct.

In the next step, we use the Hilbert class polynomial  $H_D \in \mathbb{Z}[x]$ , and project it to a poly-  
 nomial  $H_{D,q} \in \mathbb{F}_q[x]$  with coefficients in the base field  $\mathbb{F}_q$  as chosen in the first step. We do  
 this by simply computing the new coefficients as the old coefficients modulus  $p$ , that is, if  
 $H_D(x) = a_m x^m + a_{m-1} x^{m-1} + \dots + a_1 x + a_0$ , we compute the  $q$ -modulus of each coefficient  
 $\tilde{a}_j = a_j \bmod p$ , which defines the **projected Hilbert class polynomial** as follows:

$$H_{D,p}(x) = \tilde{a}_m x^m + \tilde{a}_{m-1} x^{m-1} + \dots + \tilde{a}_1 x + \tilde{a}_0$$

3695 We then search for roots of  $H_{D,p}$ , since every root  $j_0$  of  $H_{D,p}$  defines a family of elliptic curves  
 3696 over  $\mathbb{F}_q$ , which all have a  $j$ -invariant 5.31 or 5.33 equal to  $j_0$ . We can pick any root, since all of  
 3697 them will lead to proper curves eventually.

check  
reference

3698 However, some of the curves with the correct  $j$ -invariant might have an order different from  
 3699 the one we initially decided on. Therefore, we need a way to decide on a curve with the correct  
 3700 order.

3701 To compute such a curve, we have to distinguish a few different cases based on our choice  
 3702 of the root  $j_0$  and of the CM-discriminant  $D$ . If  $j_0 \neq 0$  or  $j_0 \neq 1728 \bmod q$ , we compute  
 3703  $c_1 = \frac{j_0}{(1728 \bmod q) - j_0}$ , then we chose some arbitrary quadratic non-residue  $c_2 \in \mathbb{F}_q$ , and some  
 3704 arbitrary cubic non-residue  $c_3 \in \mathbb{F}_q$ .

3705 The following table is guaranteed to define a curve with the correct order  $r = q + 1 - t$  for  
 3706 the trace of Frobenius  $t$  we initially decided on:

actually  
make this  
a table?

3707 **Definition 5.4.0.1.** • Case  $j_0 \neq 0$  and  $j_0 \neq 1728 \bmod q$ . A curve with the correct order is  
 3708 defined by one of the following equations:

$$y^2 = x^3 + 3c_1 x + 2c_1 \quad \text{or} \quad y^2 = x^3 + 3c_1 c_2^2 x + 2c_1 c_2^3 \quad (5.35)$$

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3710

- Case  $j_0 = 0$  and  $D \neq -3$ . A curve with the correct order is defined by one of the following equations:

$$y^2 = x^3 + 1 \quad \text{or} \quad y^2 = x^3 + c_2^3 \quad (5.36)$$

- Case  $j_0 = 0$  and  $D = -3$ . A curve with the correct order is defined by one of the following equations:

$$\begin{aligned} y^2 &= x^3 + 1 \quad \text{or} \quad y^2 = x^3 + c_2^3 \quad \text{or} \\ y^2 &= x^3 + c_3^2 \quad \text{or} \quad y^2 = c_3^2 c_2^3 \quad \text{or} \\ y^2 &= x^3 + c_3^{-2} \quad \text{or} \quad y^2 = x^3 + c_3^{-2} c_2^3 \end{aligned}$$

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- Case  $j_0 = 1728 \bmod q$  and  $D \neq -4$ . A curve with the correct order is defined by one of the following equations:

$$y^2 = x^3 + x \quad \text{or} \quad y^2 = x^3 + c_2^2 x \quad (5.37)$$

- Case  $j_0 = 1728 \bmod q$  and  $D = -4$ . A curve with the correct order is defined by one of the following equations:

$$\begin{aligned} y^2 &= x^3 + x \quad \text{or} \quad y^2 = x^3 + c_2 x \quad \text{or} \\ y^2 &= x^3 + c_2^2 x \quad \text{or} \quad y^2 = x^3 + c_2^3 x \end{aligned}$$

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To decide the proper defining Weierstraß equation, we therefore have to compute the order of any of the potential curves above, and then choose the one that fits our initial requirements. Since it can be shown that the Hilbert class polynomials for the CM-discriminants  $D = -3$  and  $D = -4$  are given by  $H_{-3,q}(x) = x$  and  $H_{-4,q}(x) = x - (1728 \bmod q)$  (EXERCISE), the previous cases are exhaustive.

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To summarize, using the complex multiplication method, it is possible to synthesize elliptic curves with predefined order over predefined base fields from scratch. However, the curves that are constructed this way are just some representatives of a larger class of curves, all of which have the same order. Therefore, in real-world applications, it is sometimes more advantageous to choose a different representative from that class. To do so recall from XXX that any curve defined by the Weierstraß equation  $y^2 = x^3 + ax + b$  is isomorphic to a curve of the form  $y^2 = x^3 + ad^2x + bd^3$  for some quadratic residue  $d \in \mathbb{F}_q$ .

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In order to find a suitable representative (e.g. with small parameters  $a$  and  $b$ ) in the last step, the curve designer might choose a quadratic residue  $d$  such that the transformed curve has the properties they wanted.

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*Example 100.* Consider curve  $E_1(\mathbb{F}_5)$  from example 65. We want to use the complex multiplication method to derive that curve from scratch. Since  $E_1(\mathbb{F}_5)$  is a curve of order  $r = 9$  over the prime field of order  $q = 5$ , we know from example 94 that its trace of Frobenius is  $t = -3$ , which also implies that  $q$  and  $|t|$  are coprime.

We then have to find parameters  $D, v \in \mathbb{Z}$  such that the criteria in 5.32 hold. We get the following:

$$\begin{aligned} 4q &= t^2 + |D|v^2 && \Rightarrow \\ 20 &= (-3)^2 + |D|v^2 && \Leftrightarrow \\ 11 &= |D|v^2 \end{aligned}$$

exercise  
still to be  
written?

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ence

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reference

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Now, since 11 is a prime number, the only solution is  $|D| = 11$  and  $v = 1$  here. With  $D = -11$  and the Euclidean division of  $-11$  by 4 being  $-11 = -3 \cdot 4 + 1$ , we have  $-11 \bmod 4 = 1$ , which shows that  $D = -11$  is a proper choice.

In the next step, we have to compute the Hilbert class polynomial  $H_{-11}$ . To do so, we first have to find the set  $ICG(D)$ . To compute that set, observe that, since  $\sqrt{\frac{|D|}{3}} \approx 1.915 < 2$ , we know from  $A \leq \sqrt{\frac{|D|}{3}}$  and  $A \in \mathbb{Z}$  that  $A$  must be either 0 or 1.

For  $A = 0$ , we know  $B = 0$  from the constraint  $|B| \leq A$ . However, in this case, there could be no  $C$  satisfying  $-11 = B^2 - 4AC$ . So we try  $A = 1$  and deduce  $B \in \{-1, 0, 1\}$  from the constraint  $|B| \leq A$ . The case  $B = -1$  can be excluded, since then  $B < 0$  has to imply  $|B| < A$ . The case  $B = 0$  can also be excluded, as there cannot be an integer  $C$  with  $-11 = -4C$ , since 11 is a prime number.

This leaves the case  $B = 1$ , and we compute  $C = 3$  from the equation  $-11 = 1^2 - 4C$ , which gives the solution  $(A, B, C) = (1, 1, 3)$ :

$$ICG(D) = \{(1, 1, 3)\}$$

With the set  $ICG(D)$  at hand, we can compute the Hilbert class polynomial of  $D = -11$ . To do so, we have to insert the term  $\frac{-1+\sqrt{-11}}{2}$  into the  $j$ -function. To do so, first observe that  $\sqrt{-11} = i\sqrt{11}$ , where  $i$  is the imaginary unit, defined by  $i^2 = -1$ . Using this, we can invoke Sage to compute the  $j$ -invariant and get the following:

$$H_{-11}(x) = x - j\left(\frac{-1+i\sqrt{11}}{2}\right) = x + 32768$$

As we can see, in this particular case, the Hilbert class polynomial is a linear function with a single integer coefficient. In the next step, we have to project it onto a polynomial from  $\mathbb{F}_5[x]$  by computing the modular 5 remainder of the coefficients 1 and 32768. We get  $32768 \bmod 5 = 3$ , from which it follows that the projected Hilbert class polynomial is considered a polynomial from  $\mathbb{F}_5[x]$ :

$$H_{-11,5}(x) = x + 3$$

As we can see, the only root of this polynomial is  $j = 2$ , since  $H_{-11,5}(2) = 2 + 3 = 0$ . We therefore have a situation with  $j \neq 0$  and  $j \neq 1728$ , which tells us that we have to compute the parameter  $c_1$  in modular 5 arithmetics:

$$c_1 = \frac{2}{1728 - 2}$$

Since  $1728 \bmod 5 = 3$ , we get  $c_1 = 2$ .

Next, we have to check if the curve  $E(\mathbb{F}_5)$  defined by the Weierstraß equation  $y^2 = x^3 + 3 \cdot 2x + 2 \cdot 2$  has the correct order. We invoke Sage, and find that the order is indeed 9, so it is a curve with the required parameters. Thus, we have successfully constructed the curve with the desired properties.

Note, however, that in real-world applications, it might be useful to choose parameters  $a$  and  $b$  that have certain properties, e.g. to be as small as possible. As we know from XXX, choosing any quadratic residue  $d \in \mathbb{F}_5$  gives a curve of the same order defined by  $y^2 = x^2 + ak^2x + bk^3$ . Since 4 is a quadratic residue in  $\mathbb{F}_4$ , we can transform the curve defined by  $y^2 = x^3 + x + 4$  into the curve  $y^2 = x^3 + 4^2 + 4 \cdot 4^3$  which gives the following:

$$y^2 = x^3 + x + 1$$

add reference

3748 This is the curve  $E_1(\mathbb{F}_5)$  that we used extensively throughout this book. Thus, using the  
 3749 complex multiplication method, we were able to derive a curve with specific properties from  
 3750 scratch.

3751 *Example 101.* Consider the tiny-jubjub curve  $TJJ\_13$  from example 66. We want to use the  
 3752 complex multiplication method to derive that curve from scratch. Since  $TJJ\_13$  is a curve of  
 3753 order  $r = 20$  over the prime field of order  $q = 13$ , we know from example 95 that its trace of  
 3754 Frobenius is  $t = -6$ , which also implies that  $q$  and  $|t|$  are coprime.

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referencecheck  
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We then have to find parameters  $D, v \in \mathbb{Z}$  such that 5.32 holds. We get the following:

$$\begin{aligned} 4q &= t^2 + |D|v^2 && \Rightarrow \\ 4 \cdot 13 &= (-6)^2 + |D|v^2 && \Rightarrow \\ 52 &= 36 + |D|v^2 && \Leftrightarrow \\ 16 &= |D|v^2 \end{aligned}$$

3755 This equation has two solutions for  $(D, v)$ , namely  $(-4, \pm 2)$  and  $(-16, \pm 1)$ . Looking at the  
 3756 first solution, we know that  $D = -4$  implies  $j = 1728$ , and the constructed curve is defined by  
 3757 a Weierstraß equation 5.1 that has a vanishing parameter  $b = 0$ . We can therefore conclude that  
 3758 choosing  $D = -4$  will not help us reconstructing  $TJJ\_13$ . It will produce curves with order 20,  
 3759 just not the one we are looking for.

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reference

So we choose the second solution  $D = -16$ . In the next step, we have to compute the Hilbert class polynomial  $H_{-16}$ . To do so, we first have to find the set  $ICG(D)$ . To compute that set, observe that since  $\sqrt{\frac{|-16|}{3}} \approx 2.31 < 3$ , we know from  $A \leq \sqrt{\frac{|-16|}{3}}$  and  $A \in \mathbb{Z}$  that  $A$  must be in the range  $0..2$ . So we loop through all possible values of  $A$  and through all possible values of  $B$  under the constraints  $|B| \leq A$ , and if  $B < 0$  then  $|B| < A$ . Then we compute potential  $C$ 's from  $-16 = B^2 - 4AC$ . We get the following two solutions for  $ICG(D)$ : we get

$$ICG(D) = \{(1, 0, 4), (2, 0, 2)\}$$

With the set  $ICG(D)$  at hand, we can compute the Hilbert class polynomial of  $D = -16$ . We can invoke Sage to compute the  $j$ -invariant and get the following:

$$\begin{aligned} H_{-16}(x) &= \left( x - j \left( \frac{i\sqrt{16}}{2} \right) \right) \left( x - j \left( \frac{i\sqrt{16}}{4} \right) \right) \\ &= (x - 287496)(x - 1728) \end{aligned}$$

As we can see, in this particular case, the Hilbert class polynomial is a quadratic function with two integer coefficients. In the next step, we have to project it onto a polynomial from  $\mathbb{F}_5[x]$  by computing the modular 5 remainder of the coefficients 1, 287496 and 1728. We get  $287496 \bmod 13 = 1$  and  $1728 \bmod 13 = 2$ , which means that the projected Hilbert class polynomial is as follows:

$$H_{-11,5}(x) = (x - 1)(x - 12) = (x + 12)(x + 1)$$

3760 This is considered a polynomial from  $\mathbb{F}_5[x]$ . Thus, we have two roots, namely  $j = 1$  and  $j =$   
 3761 12. We already know that  $j = 12$  is the wrong root to construct the tiny-jubjub curve, since  
 3762  $1728 \bmod 13 = 2$ , and that case is not compatible with a curve with  $b \neq 0$ . So we choose  $j = 1$ .



Another way to decide the proper root is to compute the  $j$ -invariant of the tiny-jubjub curve. We get the following:

$$\begin{aligned}
 j(TJJ\_13) &= 12 \frac{4 \cdot 8^3}{4 \cdot 8^3 + 1 \cdot 8^2} \\
 &= 12 \frac{4 \cdot 5}{4 \cdot 5 + 12} \\
 &= 12 \frac{7}{7 + 12} \\
 &= 12 \frac{7}{7 + 12} \\
 &= 1
 \end{aligned}$$

This is equal to the root  $j = 1$  of the Hilbert class polynomial  $H_{-16,13}$  as expected. We therefore have a situation with  $j \neq 0$  and  $j \neq 1728$ , which tells us that we have to compute the parameter  $c_1$  in modular 5 arithmetics:

$$c_1 = \frac{1}{12 - 1} = 6$$

Since  $1728 \bmod 13 = 12$ , we get  $c_1 = 6$ . Then we have to check if the curve  $E(\mathbb{F}_5)$  defined by the Weierstraß equation  $y^2 = x^3 + 3 \cdot 6x + 2 \cdot 6$ , which is equivalent to  $y^2 = x^3 + 5x + 12$ , has the correct order. We invoke Sage and find that the order is 8, which implies that the trace of this curve is 6, not  $-6$  as required. So we have to consider the second possibility, and choose some quadratic non-residue  $c_2 \in \mathbb{F}_{13}$ . We choose  $c_2 = 5$  and compute the Weierstraß equation  $y^2 = x^3 + 5c_2^2 + 12c_2^3$  as follows:

$$y^2 = x^3 + 8x + 5$$

We invoke Sage and find that the order is 20, which is indeed the correct one. As we know from XXX, choosing any quadratic residue  $d \in \mathbb{F}_5$  gives a curve of the same order defined by  $y^2 = x^2 + ad^2x + bd^3$ . Since 12 is a quadratic residue in  $\mathbb{F}_{13}$ , we can transform the curve defined by  $y^2 = x^3 + 8x + 5$  into the curve  $y^2 = x^3 + 12^2 \cdot 8 + 5 \cdot 12^3$  which gives the following:

$$y^2 = x^3 + 8x + 8$$

add reference

3763 This is the tiny-jubjub curve that we used extensively throughout this book. So using the  
3764 complex multiplication method, we were able to derive a curve with specific properties from  
3765 scratch.

*Example 102.* To consider a real-world example, we want to use the complex multiplication method in combination with Sage to compute Secp256k1 from scratch. So based on example 67, we decided to compute an elliptic curve over a prime field  $\mathbb{F}_p$  of order  $r$  for the following security parameters:

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$$\begin{aligned}
 p &= 115792089237316195423570985008687907853269984665640564039457584007908834671663 \\
 r &= 115792089237316195423570985008687907852837564279074904382605163141518161494337
 \end{aligned}$$

According to example 96, this gives the following trace of Frobenius:

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$$t = 43242038656559659652420866390673177327$$

3766 We also decided that we want a curve of the form  $y^2 = x^3 + b$ , that is, we want the parameter  
3767  $a$  to be zero. This implies that the  $j$ -invariant of our curve must be zero.

In a first step, we have to find a CM-discriminant  $D$  and some integer  $v$  such that the equation  $4p = t^2 + |D|v^2$  is satisfied. Since we aim for a vanishing  $j$ -invariant, the first thing to try is  $D = -3$ . In this case, we can compute  $v^2 = (4p - t^2)$ , and if  $v^2$  happens to be an integer that has a square root  $v$ , we are done. Invoking Sage we compute as follows:

```

3772 sage: D = -3
3773 sage: p = 1157920892373161954235709850086879078532699846656405
3774         64039457584007908834671663
3775 sage: r = 1157920892373161954235709850086879078528375642790749
3776         04382605163141518161494337
3777 sage: t = p+1-r
3778 sage: v_sqr = (4*p - t^2)/abs(D)
3779 sage: v_sqr.is_integer()
3780 True
3781 sage: v = sqrt(v_sqr)
3782 sage: v.is_integer()
3783 True
3784 sage: 4*p == t^2 + abs(D)*v^2
3785 True
3786 sage: v
3787 303414439467246543595250775667605759171

```

The pair  $(D, v) = (-3, 303414439467246543595250775667605759171)$  does indeed solve the equation, which tells us that there is a curve of order  $r$  over a prime field of order  $p$ , defined by a Weierstraß equation  $y^2 = x^3 + b$  for some  $b \in \mathbb{F}_p$ . Now we need to compute  $b$ .

For  $D = -3$ , we already know that the associated Hilbert class polynomial is given by  $H_{-3}(x) = x$ , which gives the projected Hilbert class polynomial as  $H_{-3,p} = x$  and the  $j$ -invariant of our curve is guaranteed to be  $j = 0$ . Now, looking at 5.4.0.1, we see that there are 6 possible cases to construct a curve with the correct order  $r$ . In order to construct the curves in question, we have to choose some arbitrary quadratic and cubic non-residue. So we loop through  $\mathbb{F}_p$  to find them, invoking Sage:

```

3797 sage: F = GF(p)
3798 sage: for c2 in F:
3799     ....:     try: # quadratic residue
3800     ....:         _ = c2.nth_root(2)
3801     ....:     except ValueError: # quadratic non-residue
3802     ....:         break
3803 sage: c2
3804 3
3805 sage: for c3 in F:
3806     ....:     try:
3807     ....:         _ = c3.nth_root(3)
3808     ....:     except ValueError:
3809     ....:         break
3810 sage: c3
3811 2

```

We found the quadratic non-residue  $c_2 = 3$  and the cubic non-residue  $c_3 = 2$ . Using those numbers, we check the six cases against the the expected order  $r$  of the curve we want to

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reference

3814 synthesize:

```

3815 sage: C1 = EllipticCurve(F, [0, 1])           527
3816 sage: C1.order() == r                         528
3817 False                                         529
3818 sage: C2 = EllipticCurve(F, [0, c2^3])         530
3819 sage: C2.order() == r                         531
3820 False                                         532
3821 sage: C3 = EllipticCurve(F, [0, c3^2])         533
3822 sage: C3.order() == r                         534
3823 False                                         535
3824 sage: C4 = EllipticCurve(F, [0, c3^2*c2^3])   536
3825 sage: C4.order() == r                         537
3826 False                                         538
3827 sage: C5 = EllipticCurve(F, [0, c3^(-2)])     539
3828 sage: C5.order() == r                         540
3829 False                                         541
3830 sage: C6 = EllipticCurve(F, [0, c3^(-2)*c2^3]) 542
3831 sage: C6.order() == r                         543
3832 True                                          544

```

As expected, we found an elliptic curve of the correct order  $r$  over a prime field of size  $p$ . In principle, we are done, as we have found a curve with the same basic properties as Secp256k1. However, the curve is defined by the following equation, which uses a very large parameter  $b_1$ , and so it might perform too slowly in certain algorithms.

$$y^2 = x^3 + 86844066927987146567678238756515930889952488499230423029593188005931626003754$$

It is also not very elegant to be written down by hand. It might therefore be advantageous to find an isomorphic curve with the smallest possible parameter  $b_2$ . In order to find such a  $b_2$ , we have to choose a quadratic residue  $d$  such that  $b_2 = b_1 \cdot d^3$  is as small as possible. To do so, we rewrite the last equation into the following form:

$$d = \sqrt[3]{\frac{b_2}{b_1}}$$

what does this mean? Maybe just delete it

3833 Then we invoke Sage to loop through values  $b_2 \in \mathbb{F}_p$  until it finds some number such that  
 3834 the quotient  $\frac{b_2}{b_1}$  has a cube root  $d$  and this cube root itself is a quadratic residue.

```

3835 sage: b1=86844066927987146567678238756515930889952488499230423  545
3836       029593188005931626003754
3837 sage: for b2 in F:                                           546
3838     ....:     try:                                           547
3839     ....:         d = (b2/b1).nth_root(3)                   548
3840     ....:         try:                                       549
3841     ....:             __ = d.nth_root(2)                     550
3842     ....:             if d != 0:                             551
3843     ....:                 break                               552
3844     ....:         except ValueError:                          553
3845     ....:             pass                                    554
3846     ....:     except ValueError:                             555
3847     ....:         pass                                       556

```

3848 **sage: b2**  
 3849 **7**

557

558

3850 Indeed, the smallest possible value is  $b_2 = 7$  and the defining Weierstraß equation of a curve  
 3851 over  $\mathbb{F}_p$  with prime order  $r$  is  $y^2 = x^3 + 7$ , which we might call Secp256k1. As we have just  
 3852 seen, the complex multiplication method is powerful enough to derive cryptographically secure  
 3853 curves like Secp256k1 from scratch.

3854 **The BLS6\_6 pen-and-paper curve** In this paragraph, we summarize our understanding of  
 3855 elliptic curves to derive our main pen-and-paper example for the rest of the book. To do so, we  
 3856 want to use the complex multiplication method to derive a pairing-friendly elliptic curve that  
 3857 has similar properties to curves that are used in actual cryptographic protocols. However, we  
 3858 design the curve specifically to be useful in pen-and-paper examples, which mostly means that  
 3859 the curve should contain only a few points so that we are able to derive exhaustive addition and  
 3860 pairing tables.

3861 A well-understood family of pairing-friendly curves is the the group of BLS curves (STUFF  
 3862 ABOUT THE HISTORY AND THE NAMING CONVENTION), which are derived in [XXX].  
 3863 BLS curves are particularly useful in our case if the embedding degree  $k$  satisfies  $k \equiv 6 \pmod{0}$ .  
 3864 Of course, the smallest embedding degree  $k$  that satisfies this congruency is  $k = 6$  and we there-  
 3865 fore aim for a BLS6 curve as our main pen-and-paper example.

write up  
this part

add refer-  
ence

3866 To apply the complex multiplication method from page 110 ff., recall that this method starts  
 3867 with a definition of the base field  $\mathbb{F}_{p^m}$ , as well as the trace of Frobenius  $t$  and the order of the  
 3868 curve. If the order  $p^m + 1 - t$  is not a prime number, then the order  $r$  of the largest prime factor  
 3869 group needs to be controlled.

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reference

3870 In the case of BLS\_6 curves, the parameter  $m$  is chosen to be 1, which means that the  
 3871 curves are defined over prime fields. All relevant parameters  $p$ ,  $t$  and  $r$  are then themselves  
 3872 parameterized by the following three polynomials:

$$\begin{aligned} r(x) &= \Phi_6(x) \\ t(x) &= x + 1 \\ q(x) &= \frac{1}{3}(x-1)^2(x^2 - x + 1) + x \end{aligned} \tag{5.38}$$

3873 In the equations above,  $\Phi_6$  is the 6-th cyclotomic polynomial and  $x \in \mathbb{N}$  is a parameter  
 3874 that the designer has to choose in such a way that the evaluation of  $p$ ,  $t$  and  $r$  at the point  $x$   
 3875 gives integers that have the proper size to meet the security requirements of the curve that they  
 3876 want to design. It is then guaranteed that the complex multiplication method can be used in  
 3877 combination with those parameters to define an elliptic curve with CM-discriminant  $D = -3$ ,  
 3878 embedding degree  $k = 6$ , and curve equation  $y^2 = x^3 + b$  for some  $b \in \mathbb{F}_p$ .

cyclotomic  
polyno-  
mial

3879 For example, if the curve should target the 128-bit security level, due to the Pholaard-rho  
 3880 attack (TODO) the parameter  $r$  should be prime number of at least 256 bits.

Pholaard-  
rho attack

3881 In order to design the smallest BLS\_6 curve, we therefore have to find a parameter  $x$  such  
 3882 that  $r(x)$ ,  $t(x)$  and  $q(x)$  are the smallest natural numbers that satisfy  $q(x) > 3$  and  $r(x) > 3$ .<sup>1</sup>

todo

We therefore initiate the design process of our BLS6 curve by looking up the 6-th cyclo-  
 tomic polynomial, which is  $\Phi_6 = x^2 - x + 1$ , and then insert small values for  $x$  into the defining

<sup>1</sup>The smallest BLS curve will also be the most insecure BLS curve. However, since our goal with this curve is ease of pen-and-paper computation rather than security, it fits the purposes of this book.

polynomials  $r, t, q$ . We get the following results:

$$\begin{array}{lll} x = 1 & (r(x), t(x), q(x)) & (1, 2, 1) \\ x = 2 & (r(x), t(x), q(x)) & (3, 3, 3) \\ x = 3 & (r(x), t(x), q(x)) & (7, 4, \frac{37}{3}) \\ x = 4 & (r(x), t(x), q(x)) & (13, 5, 43) \end{array}$$

Since  $q(1) = 1$  is not a prime number, the first  $x$  that gives a proper curve is  $x = 2$ . However, such a curve would be defined over a base field of characteristic 3, and we would rather like to avoid that. We therefore find  $x = 4$ , which defines a curve over the prime field of characteristic 43 that has a trace of Frobenius  $t = 5$  and a larger order prime group of size  $r = 13$ .

Since the prime field  $\mathbb{F}_{43}$  has 43 elements and 43's binary representation is  $43_2 = 101011$ , which consists of 6 digits, the name of our pen-and-paper curve should be *BLS6\_6*, since it is common to name a BLS curve by its embedding degree and the bit-length of the modulus in the base field. We call *BLS6\_6* the **moon-math-curve**.

Based on 5.29, we know that the Hasse bound implies that *BLS6\_6* will contain exactly 39 elements. Since the prime factorization of 39 is  $39 = 3 \cdot 13$ , we have a “large” prime factor group of size 13, as expected, and a small cofactor group of size 3. Fortunately, a subgroup of order 13 is well suited for our purposes, as 13 elements can be easily handled in the associated addition, scalar multiplication and pairing tables in a pen-and-paper style.

We can check that the embedding degree is indeed 6 as expected, since  $k = 6$  is the smallest number  $k$  such that  $r = 13$  divides  $43^k - 1$ .

```
sage: for k in range(1, 42): # Fermat's little theorem
.....:     if (43^k - 1) % 13 == 0:
.....:         break
sage: k
6
```

In order to compute the defining equation  $y^2 = x^3 + ax + b$  of *BLS6\_6*, we use the complex multiplication method as described in 5.4. The goal is to find  $a, b \in \mathbb{F}_{43}$  representations that are particularly nice to work with. The authors of XXX showed that the CM-discriminant of every BLS curve is  $D = -3$  and, indeed, the following equation has the four solutions  $(D, v) \in \{(-3, -7), (-3, 7), (-49, -1), (-49, 1)\}$  if  $D$  is required to be negative, as expected:

$$\begin{array}{ll} 4p = t^2 + |D|v^2 & \Rightarrow \\ 4 \cdot 43 = 5^2 + |D|v^2 & \Rightarrow \\ 172 = 25 + |D|v^2 & \Leftrightarrow \\ 49 = |D|v^2 & \end{array}$$

This means that  $D = -3$  is indeed a proper CM-discriminant, and we can deduce that the parameter  $a$  has to be 0, and that the Hilbert class polynomial is given by  $H_{-3, 43}(x) = x$ .

This implies that the  $j$ -invariant of *BLS6\_6* is given by  $j(\text{BLS6\_6}) = 0$ . We therefore have to look at case XXX in table 5.4.0.1 to derive a parameter  $b$ . To decide the proper case for  $j_0 = 0$  and  $D = -3$ , we therefore have to choose some arbitrary quadratic non-residue  $c_2$  and cubic non-residue  $c_3$  in  $\mathbb{F}_{43}$ . We choose  $c_2 = 5$  and  $c_3 = 36$ . We check these with Sage:

```
sage: F43 = GF(43)
```

why? Because in this book elliptic curves are only defined for fields of characteristic  $> 3$

check reference

559  
560  
561  
562  
563

check reference

what does this mean?

add reference

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check reference

```

3915 sage: c2 = F43(5) 565
3916 .....: try: # quadratic residue 566
3917 .....:     c2.nth_root(2) 567
3918 .....: except ValueError: # quadratic non-residue 568
3919 .....:     c2 569
3920 sage: c3 = F43(36) 570
3921 .....: try: 571
3922 .....:     c3.nth_root(3) 572
3923 .....: except ValueError: 573
3924 .....:     c3 574

```

3925 Using those numbers we check the six possible cases from 5.4.0.1 against the the expected  
3926 order 39 of the curve we want to synthesize:

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```

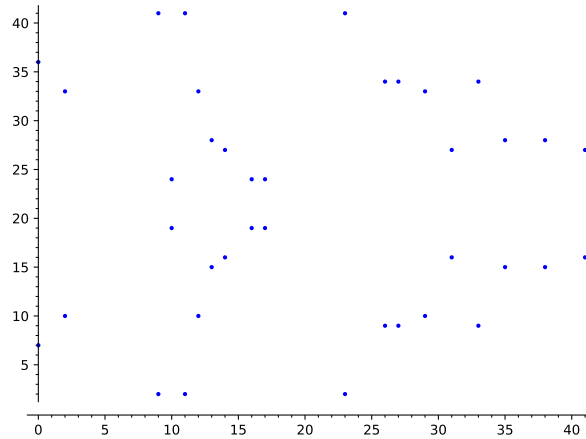
3927 sage: BLS61 = EllipticCurve(F43, [0, 1]) 575
3928 sage: BLS61.order() == 39 576
3929 False 577
3930 sage: BLS62 = EllipticCurve(F43, [0, c2^3]) 578
3931 sage: BLS62.order() == 39 579
3932 False 580
3933 sage: BLS63 = EllipticCurve(F43, [0, c3^2]) 581
3934 sage: BLS63.order() == 39 582
3935 True 583
3936 sage: BLS64 = EllipticCurve(F43, [0, c3^2*c2^3]) 584
3937 sage: BLS64.order() == 39 585
3938 False 586
3939 sage: BLS65 = EllipticCurve(F43, [0, c3^(-2)]) 587
3940 sage: BLS65.order() == 39 588
3941 False 589
3942 sage: BLS66 = EllipticCurve(F43, [0, c3^(-2)*c2^3]) 590
3943 sage: BLS66.order() == 39 591
3944 False 592
3945 sage: BLS6 = BLS63 # our BLS6 curve in the book 593

```

3946 As expected, we found an elliptic curve of the correct order 39 over a prime field of size 43,  
3947 defined by the following equation:

$$BLS6\_6 := \{(x, y) \mid y^2 = x^3 + 6 \text{ for all } x, y \in \mathbb{F}_{43}\} \quad (5.39)$$

3948 There are other choices for  $b$ , such as  $b = 10$  or  $b = 23$ , but all these curves are isomorphic,  
3949 and hence represent the same curve in a different way. Since BLS6-6 only contains 39 points, it  
3950 is possible to give a visual impression of the curve:



3951

3952 As we can see, our curve has some desirable properties: it does not contain self-inverse  
 3953 points, that is, points with  $y = 0$ . It follows that the addition law can be optimized, since the  
 3954 branch for those cases can be eliminated.

3955 Summarizing the previous procedure, we have used the method of Barreto, Lynn and Scott  
 3956 to construct a pairing-friendly elliptic curve of embedding degree 6. However, in order to do  
 3957 elliptic curve cryptography on this curve, note that, since the order of  $BLS6\_6$  is 39, its group  
 3958 of rational points is not a finite cyclic group of prime order. We therefore have to find a suitable  
 3959 subgroup as our main target. Since  $39 = 13 \cdot 3$ , we know that the curve must contain a “large”  
 3960 prime-order group of size 13 and a small cofactor group of order 3.

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3961 The following step is to construct this group. One way to do so is to find a generator. We  
 3962 can achieve this by choosing an arbitrary element of the group that is not the point at infinity,  
 3963 and then multiply that point with the cofactor of the group’s order. If the result is not the point  
 3964 at infinity, the result will be a generator. If it is the point at infinity we have to choose a different  
 3965 element.

In order to find a generator for the large order subgroup of size 13, we first notice that the cofactor of 13 is 3, since  $39 = 3 \cdot 13$ . We then need to construct an arbitrary element from  $BLS6\_6$ . To do so in a pen-and-paper style, we can choose some *arbitrary*  $x \in \mathbb{F}_{43}$  and see if there is some solution  $y \in \mathbb{F}_{43}$  that satisfies the defining Weierstraß equation  $y^2 = x^3 + 6$ . We choose  $x = 9$ , and check that  $y = 2$  is a proper solution:

$$\begin{aligned} y^2 &= x^3 + 6 && \Rightarrow \\ 2^2 &= 9^3 + 6 && \Leftrightarrow \\ 4 &= 4 \end{aligned}$$

3966 This implies that  $P = (9, 2)$  is therefore a point on  $BLS6\_6$ . To see if we can project this  
 3967 point onto a generator of the large order prime group  $BLS6\_6[13]$ , we have to multiply  $P$  with  
 3968 the cofactor, that is, we have to compute  $[3](9, 2)$ . After some computation (EXERCISE) we  
 3969 get  $[3](9, 2) = (13, 15)$ . Since this is not the point at infinity, we know that  $(13, 15)$  must be a  
 3970 generator of  $BLS6\_6[13]$ . The generator  $g_{BLS6\_6[13]}$ , which we will use in pairing computations  
 3971 in the remainder of this book, is given as follows:

add exer-  
cise

$$g_{BLS6\_6[13]} = (13, 15) \tag{5.40}$$

3972 Since  $g_{BLS6\_6[13]}$  is a generator, we can use it to construct the subgroup  $BLS6\_6[13]$  by re-  
 3973 peatedly adding the generator to itself. Using Sage, we get the following:

3974 **sage:** `P = BLS6(9, 2)`

594



```

3975 sage: Q = 3*P
3976 sage: Q.xy()
3977 (13, 15)
3978 sage: BLS6_13 = []
3979 sage: for x in range(0,13): # cyclic of order 13
3980 .....:     P = x*Q
3981 .....:     BLS6_13.append(P)

```

Repeatedly adding a generator to itself, as we just did, will generate small groups in logarithmic order with respect to the generator as, explained on page 43 ff. We therefore get the following description of the large prime-order subgroup of  $BLS6_6$ :

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reference

$$\begin{aligned}
 BLS6_6[13] = \\
 \{ (13, 15) \rightarrow (33, 34) \rightarrow (38, 15) \rightarrow (35, 28) \rightarrow (26, 34) \rightarrow (27, 34) \rightarrow \\
 (27, 9) \rightarrow (26, 9) \rightarrow (35, 15) \rightarrow (38, 28) \rightarrow (33, 9) \rightarrow (13, 28) \rightarrow \mathcal{O} \} \quad (5.41)
 \end{aligned}$$

Having a logarithmic description of this group is tremendously helpful in pen-and-paper computations. To see that, observe that we know from XXX that there is an exponential map from the scalar field  $\mathbb{F}_{13}$  to  $BLS6_6[13]$  with respect to our generator, which generates the group in logarithmic order:

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$$[\cdot]_{(13,15)} : \mathbb{F}_{13} \rightarrow BLS6_6[13] ; x \mapsto [x](13, 15)$$

So, for example, we have  $[1]_{(13,15)} = (13, 15)$ ,  $[7]_{(13,15)} = (27, 9)$  and  $[0]_{(13,15)} = \mathcal{O}$  and so on. The relevant point here is that we can use this representation to do computations in  $BLS6_6[13]$  efficiently in our head using XXX, as in the following example:

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$$\begin{aligned}
 (27, 34) \oplus (33, 9) &= [6](13, 15) \oplus [11](13, 15) \\
 &= [6 + 11](13, 15) \\
 &= [4](13, 15) \\
 &= (35, 28)
 \end{aligned}$$

So XXX is really all we need to do computations in  $BLS6_6[13]$  in this book efficiently. However, out of convenience, the following picture lists the entire addition table of that group, as it might be useful in pen-and-paper computations:

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$\oplus$	$\mathcal{O}$	(13, 15)	(33, 34)	(38, 15)	(35, 28)	(26, 34)	(27, 34)	(27, 9)	(26, 9)	(35, 15)	(38, 28)	(33, 9)	(13, 28)
$\mathcal{O}$	$\mathcal{O}$	(13, 15)	(33, 34)	(38, 15)	(35, 28)	(26, 34)	(27, 34)	(27, 9)	(26, 9)	(35, 15)	(38, 28)	(33, 9)	(13, 28)
(13, 15)	(13, 15)	(33, 34)	(38, 15)	(35, 28)	(26, 34)	(27, 34)	(27, 9)	(26, 9)	(35, 15)	(38, 28)	(33, 9)	(13, 28)	$\mathcal{O}$
(33, 34)	(33, 34)	(38, 15)	(35, 28)	(26, 34)	(27, 34)	(27, 9)	(26, 9)	(35, 15)	(38, 28)	(33, 9)	(13, 28)	$\mathcal{O}$	(13, 15)
(38, 15)	(38, 15)	(35, 28)	(26, 34)	(27, 34)	(27, 9)	(26, 9)	(35, 15)	(38, 28)	(33, 9)	(13, 28)	$\mathcal{O}$	(13, 15)	(33, 34)
(35, 28)	(35, 28)	(26, 34)	(27, 34)	(27, 9)	(26, 9)	(35, 15)	(38, 28)	(33, 9)	(13, 28)	$\mathcal{O}$	(13, 15)	(33, 34)	(38, 15)
(26, 34)	(26, 34)	(27, 34)	(27, 9)	(26, 9)	(35, 15)	(38, 28)	(33, 9)	(13, 28)	$\mathcal{O}$	(13, 15)	(33, 34)	(38, 15)	(35, 28)
(27, 34)	(27, 34)	(27, 9)	(26, 9)	(35, 15)	(38, 28)	(33, 9)	(13, 28)	$\mathcal{O}$	(13, 15)	(33, 34)	(38, 15)	(35, 28)	(26, 34)
(27, 9)	(27, 9)	(26, 9)	(35, 15)	(38, 28)	(33, 9)	(13, 28)	$\mathcal{O}$	(13, 15)	(33, 34)	(38, 15)	(35, 28)	(26, 34)	(27, 34)
(26, 9)	(26, 9)	(35, 15)	(38, 28)	(33, 9)	(13, 28)	$\mathcal{O}$	(13, 15)	(33, 34)	(38, 15)	(35, 28)	(26, 34)	(27, 34)	(27, 9)
(35, 15)	(35, 15)	(38, 28)	(33, 9)	(13, 28)	$\mathcal{O}$	(13, 15)	(33, 34)	(38, 15)	(35, 28)	(26, 34)	(27, 34)	(27, 9)	(26, 9)
(38, 28)	(38, 28)	(33, 9)	(13, 28)	$\mathcal{O}$	(13, 15)	(33, 34)	(38, 15)	(35, 28)	(26, 34)	(27, 34)	(27, 9)	(26, 9)	(35, 15)
(33, 9)	(33, 9)	(13, 28)	$\mathcal{O}$	(13, 15)	(33, 34)	(38, 15)	(35, 28)	(26, 34)	(27, 34)	(27, 9)	(26, 9)	(35, 15)	(38, 28)
(13, 28)	(13, 28)	$\mathcal{O}$	(13, 15)	(33, 34)	(38, 15)	(35, 28)	(26, 34)	(27, 34)	(27, 9)	(26, 9)	(35, 15)	(38, 28)	(33, 9)



Now that we have constructed a “large” cyclic prime-order subgroup of  $BLS6\_6$  suitable for many pen-and-paper computations in elliptic curve cryptography, we have to look at how to do pairings in this context. We know that  $BLS6\_6$  is a pairing-friendly curve by design, since it has a small embedding degree  $k = 6$ . It is therefore possible to compute Weil pairings efficiently. However, in order to do so, we have to decide the groups  $\mathbb{G}_1$  and  $\mathbb{G}_2$  as explained in exercise 73.

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Since  $BLS6\_6$  has two non-trivial subgroups, it would be possible to use any of them as the  $n$ -torsion group. However, in cryptography, the only secure choice is to use the large prime-order subgroup, which in our case is  $BLS6\_6[13]$ . We therefore decide to consider the 13-torsion and define  $G_1[13]$  as the first argument for the Weil pairing function:

$$\mathbb{G}_1[13] = \{(13, 15) \rightarrow (33, 34) \rightarrow (38, 15) \rightarrow (35, 28) \rightarrow (26, 34) \rightarrow (27, 34) \rightarrow (27, 9) \rightarrow (26, 9) \rightarrow (35, 15) \rightarrow (38, 28) \rightarrow (33, 9) \rightarrow (13, 28) \rightarrow \mathcal{O}\}$$

In order to construct the domain for the second argument, we need to construct  $\mathbb{G}_2[13]$ , which, according to the general theory, should be defined by those elements  $P$  of the full 13-torsion group  $BLS6\_6[13]$  that are mapped to  $43 \cdot P$  under the Frobenius endomorphism (equation 5.24).

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To compute  $\mathbb{G}_2[13]$ , we therefore have to find the full 13-torsion group first. To do so, we use the technique from XXX, which tells us that the full 13-torsion can be found in the curve extension over the extension field  $\mathbb{F}_{43^6}$ , since the embedding degree of  $BLS6\_6$  is 6:

$$BLS6\_6 := \{(x, y) \mid y^2 = x^3 + 6 \text{ for all } x, y \in \mathbb{F}_{43^6}\} \quad (5.42)$$

Thus, we have to construct  $\mathbb{F}_{43^6}$ , a field that contains 6321363049 elements. In order to do so, we use the procedure of XXX and start by choosing a non-reducible polynomial of degree 6 from the ring of polynomials  $\mathbb{F}_{43}[t]$ . We choose  $p(t) = t^6 + 6$ . Using Sage, we get the following:

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```
sage: F43 = GF(43) 602
sage: F43t.<t> = F43[] 603
sage: p = F43t(t^6+6) 604
sage: p.is_irreducible() 605
True 606
sage: F43_6.<v> = GF(43^6, name='v', modulus=p) 607
```

Recall from XXX that elements  $x \in \mathbb{F}_{43^6}$  can be seen as polynomials  $a_0 + a_1v + a_2v^2 + \dots + a_5v^5$  with the usual addition of polynomials and multiplication modulo  $t^6 + 6$ .

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In order to compute  $\mathbb{G}_2[13]$ , we first have to extend  $BLS6\_6$  to  $\mathbb{F}_{43^6}$ , that is, we keep the defining equation, but expand the domain from  $\mathbb{F}_{43}$  to  $\mathbb{F}_{43^6}$ . After that, we have to find at least one element  $P$  from that curve that is not the point at infinity, is in the full 13-torsion and satisfies the identity  $\pi(P) = [43]P$ . We can then use this element as our generator of  $\mathbb{G}_2[13]$  and construct all other elements by repeatedly adding the generator to itself.

Since  $BLS6(\mathbb{F}_{43^6})$  contains 6321251664 elements, it's not a good strategy to simply loop through all elements. Fortunately, Sage has a way to loop through elements from the torsion group directly:

```
sage: BLS6 = EllipticCurve(F43_6, [0, 6]) # curve extension 608
sage: INF = BLS6(0) # point at infinity 609
```

```

4024 sage: for P in INF.division_points(13): # full 13-torsion           610
4025     ....: # PI(P) == [q]P                                           611
4026     ....:     if P.order() == 13: # exclude point at infinity      612
4027     ....:         PiP = BLS6([a.frobenius() for a in P])           613
4028     ....:         qP = 43*P                                           614
4029     ....:         if PiP == qP:                                       615
4030     ....:             break                                           616
4031 sage: P.xy()                                                         617
4032 (7*v^2, 16*v^3)                                                     618

```

4033 We found an element from the full 13-torsion that is in the Eigenspace of the Eigenvalue 43,  
 4034 which implies that it is an element of  $\mathbb{G}_2[13]$ . As  $\mathbb{G}_2[13]$  is cyclic of prime order, this element  
 4035 must be a generator:

$$g_{\mathbb{G}_2[13]} = (7v^2, 16v^3) \quad (5.43)$$

4036 We can use this generator to compute  $\mathbb{G}_2$  in logarithmic order with respect to  $g_{\mathbb{G}_2[13]}$ . Using  
 4037 Sage we get the following:

```

4038 sage: Q = BLS6(7*v^2, 16*v^3)                                       619
4039 sage: BLS6_13_2 = []                                                620
4040 sage: for x in range(0, 13):                                         621
4041     ....:     P = x*Q                                                 622
4042     ....:     BLS6_13_2.append(P)                                     623

```

$$\begin{aligned} \mathbb{G}_2 = \{ & (7v^2, 16v^3) \rightarrow (10v^2, 28v^3) \rightarrow (42v^2, 16v^3) \rightarrow (37v^2, 27v^3) \rightarrow \\ & (16v^2, 28v^3) \rightarrow (17v^2, 28v^3) \rightarrow (17v^2, 15v^3) \rightarrow (16v^2, 15v^3) \rightarrow \\ & (37v^2, 16v^3) \rightarrow (42v^2, 27v^3) \rightarrow (10v^2, 15v^3) \rightarrow (7v^2, 27v^3) \rightarrow \emptyset \} \end{aligned}$$

Again, having a logarithmic description of  $\mathbb{G}_2[13]$  is tremendously helpful in pen-and-paper computations, as it reduces complicated computation in the extended curves to modular 13 arithmetics, as in the following example:

$$\begin{aligned} (17v^2, 28v^3) \oplus (10v^2, 15v^3) &= [6](7v^2, 16v^3) \oplus [11](7v^2, 16v^3) \\ &= [6 + 11](7v^2, 16v^3) \\ &= [4](7v^2, 16v^3) \\ &= (37v^2, 27v^3) \end{aligned}$$

4043 So XXX is really all we need to do computations in  $\mathbb{G}_2[13]$  in this book efficiently.

4044 To summarize the previous steps, we have found two subgroups,  $\mathbb{G}_1[13]$  and  $\mathbb{G}_2[13]$  suit-  
 4045 able to do Weil pairings on *BLS6\_6* as explained in 5.28. Using the logarithmic order XXX  
 4046 of  $\mathbb{G}_1[13]$ , the logarithmic order XXX of  $\mathbb{G}_2[13]$  and the bilinearity in 5.44, we can do Weil  
 4047 pairings on *BLS6\_6* in a pen-and-paper style:

$$e([k_1]g_{BLS6\_6[13]}, [k_2]g_{\mathbb{G}_2[13]}) = e(g_{BLS6\_6[13]}, g_{\mathbb{G}_2[13]})^{k_1 \cdot k_2} \quad (5.44)$$

4048 Observe that the Weil pairing between our two generators is given by the following identity:

$$e(g_{BLS6\_6[13]}, g_{\mathbb{G}_2[13]}) = 5v^5 + 16v^4 + 16v^3 + 15v^2 + 3v + 41 \quad (5.45)$$

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ence

```

4050 sage: g1 = BLS6([13,15]) 624
4051 sage: g2 = BLS6([7*v^2, 16*v^3]) 625
4052 sage: g1.weil_pairing(g2,13) 626
4053 5*v^5 + 16*v^4 + 16*v^3 + 15*v^2 + 3*v + 41 627

```

4054 **Hashing to pairing groups** We give various constructions to hash into  $\mathbb{G}_1$  and  $\mathbb{G}_2$ .

4055 We start with hashing to the scalar field... **TO APPEAR**

4056 None of these techniques work for hashing into  $\mathbb{G}_2$ . We therefore implement Pederson's

4057 Hash for BLS6.

We start with  $\mathbb{G}_1$ . Our goal is to define an 12-bit bounded hash function:

$$H_1 : \{0,1\}^{12} \rightarrow \mathbb{G}_1$$

Since  $12 = 3 \cdot 4$  we “randomly” select 4 uniformly distributed generators  $\{(38,15), (35,28), (27,34), (38,28)\}$  from  $\mathbb{G}_1$  and use the pseudo-random function from XXX. Therefore, we have to choose a set of 4 randomly generated invertible elements from  $\mathbb{F}_{13}$  for every generator. We choose the following:

$$\begin{aligned}
 (38,15) &: \{2,7,5,9\} \\
 (35,28) &: \{11,4,7,7\} \\
 (27,34) &: \{5,3,7,12\} \\
 (38,28) &: \{6,5,1,8\}
 \end{aligned}$$

4058 Our hash function is then computed as follows:

$$\begin{aligned}
 H_1(x_{11}, x_1, \dots, x_0) = & [2 \cdot 7^{x_{11}} \cdot 5^{x_{10}} \cdot 9^{x_9}](38,15) + [11 \cdot 4^{x_8} \cdot 7^{x_7} \cdot 7^{x_6}](35,28) + \\
 & [5 \cdot 3^{x_5} \cdot 7^{x_4} \cdot 12^{x_3}](27,34) + [6 \cdot 5^{x_2} \cdot 1^{x_1} \cdot 8^{x_0}](38,28)
 \end{aligned}$$

4059 Note that  $a^x = 1$  when  $x = 0$ . Hence, those terms can be omitted in the computation. In  
 4060 particular, the hash of the 12-bit zero string is given as follows:

**WRONG – ORDERING – REDO**

$$\begin{aligned}
 H_1(0) = & [2](38,15) + [11](35,28) + [5](27,34) + [6](38,28) = \\
 & (27,34) + (26,34) + (35,28) + (26,9) = (33,9) + (13,28) = (38,28)
 \end{aligned}$$

The hash of 011010101100 is given as follows:

$$\begin{aligned}
 H_1(011010101100) = & \text{WRONG – ORDERING – REDO} \\
 & [2 \cdot 7^0 \cdot 5^1 \cdot 9^1](38,15) + [11 \cdot 4^0 \cdot 7^1 \cdot 7^0](35,28) + [5 \cdot 3^1 \cdot 7^0 \cdot 12^1](27,34) + [6 \cdot 5^1 \cdot 1^0 \cdot 8^0](38,28) = \\
 & [2 \cdot 5 \cdot 9](38,15) + [11 \cdot 7](35,28) + [5 \cdot 3 \cdot 12](27,34) + [6 \cdot 5](38,28) = \\
 & [12](38,15) + [12](35,28) + [11](27,34) + [4](38,28) =
 \end{aligned}$$

**TO APPEAR**

We can use the same technique to define a 12-bit bounded hash function in  $\mathbb{G}_2$ :

$$H_2 : \{0,1\}^{12} \rightarrow \mathbb{G}_2$$

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fill in  
missing  
parts

Again, we “randomly” select 4 uniformly distributed generators  $\{(7v^2, 16v^3), (42v^2, 16v^3), (17v^2, 15v^3), (10v^2, 15v^3)\}$  from  $\mathbb{G}_2$ , and use the pseudo-random function from XXX. Therefore, we have to choose a set of 4 randomly generated invertible elements from  $\mathbb{F}_{13}$  for every generator:

add reference

$$\begin{aligned} (7v^2, 16v^3) &: \{8, 4, 5, 7\} \\ (42v^2, 16v^3) &: \{12, 1, 3, 8\} \\ (17v^2, 15v^3) &: \{2, 3, 9, 11\} \\ (10v^2, 15v^3) &: \{3, 6, 9, 10\} \end{aligned}$$

Our hash function is then computed like this:

$$H_1(x_{11}, x_{10}, \dots, x_0) = [8 \cdot 4^{x_{11}} \cdot 5^{x_{10}} \cdot 7^{x_9}](7v^2, 16v^3) + [12 \cdot 1^{x_8} \cdot 3^{x_7} \cdot 8^{x_6}](42v^2, 16v^3) + [2 \cdot 3^{x_5} \cdot 9^{x_4} \cdot 11^{x_3}](17v^2, 15v^3) + [3 \cdot 6^{x_2} \cdot 9^{x_1} \cdot 10^{x_0}](10v^2, 15v^3)$$

We extend this to a hash function that maps unbounded bitstrings to  $\mathbb{G}_2$  by precomposing with an actual hash function like MD5, and feed the first 12 bits of its outcome into our previously defined hash function, with  $TinyMD5_{\mathbb{G}_2}(s) = H_2(MD5(s)_{11}, \dots, MD5(s)_0)$ :

$$TinyMD5_{\mathbb{G}_2} : \{0, 1\}^* \rightarrow \mathbb{G}_2$$

For example, since  $MD5(“”) =$

$0xd41d8cd98f00b204e9800998ecf8427e$ , and the binary representation of the hexadecimal number  $0x27e$  is  $001001111110$ , we compute  $TinyMD5_{\mathbb{G}_2}$  of the empty string as follows:

$$TinyMD5_{\mathbb{G}_2}(“”) = H_2(MD5(s)_{11}, \dots, MD5(s)_0) = H_2(001001111110) =$$

4061

check equation

# Chapter 6

## Statements

As we have seen in the informal introduction XXX, a SNARK is a short non-interactive argument of knowledge, where the knowledge-proof attests to the correctness of statements like “The prover knows the prime factorization of a given number” or “The prover knows the preimage to a given SHA2 digest value” and similar things. However, human-readable statements like these are imprecise and not very useful from a formal perspective.

Chapter 1?

In this chapter we therefore look more closely at ways to formalize statements in mathematically rigorous ways, useful for SNARK development. We start by introducing formal languages as a way to define statements properly (section 6.1). We will then look at algebraic circuits and rank-1 constraint systems as two particularly useful ways to define statements in certain formal languages (section 6.2). After that, we will have a look at fundamental building blocks of compilers that compile high-level languages to circuits and associated rank-1 constraint systems.

Proper statement design should be of high priority in the development of SNARKs, since unintended true statements can lead to potentially severe and almost undetectable security vulnerabilities in the applications of SNARKs.

### 6.1 Formal Languages

Formal languages provide the theoretical background in which statements can be formulated in a logically rigorous way and where proving the correctness of any given statement can be realized by computing words in that language.

"rigorous"?

One might argue that the understanding of formal languages is not very important in SNARK development and associated statement design, but terms from that field of research are standard jargon in many papers on zero-knowledge proofs. We therefore believe that at least some introduction to formal languages and how they fit into the picture of SNARK development is beneficial, mostly to give developers a better intuition about where all this is located in the bigger picture of the logic landscape. In addition, formal languages give a better understanding of what a formal proof for a statement actually is.

"proving"?

Roughly speaking, a formal language (or just language for short) is nothing but a set of words, *th*. Words, in turn, are strings of letters taken from some alphabet and formed according to some defining rules of the language.

To be more precise, let  $\Sigma$  be any set and  $\Sigma^*$  the set of all finite **tuples** (ordered lists)  $(x_1, \dots, x_n)$  of elements  $x_j$  from  $\Sigma$  including the empty tuple  $() \in \Sigma^*$ . Then, a **language**  $L$ , in its most general definition, is nothing but a subset of  $\Sigma^*$ . In this context, the set  $\Sigma$  is called the **alphabet** of the language  $L$ , elements from  $\Sigma$  are called letters and elements from  $L$  are called **words**. The rules that specify which tuples from  $\Sigma^*$  belong to the language and which don't,

are called the **grammar** of the language. S: I suggest adding an example based on English, e.g. “tea” and “eat” are words of English, but “aet” and “tae” are not

Add ex-  
ample

If  $L_1$  and  $L_2$  are two formal languages over the same alphabet, we call  $L_1$  and  $L_2$  **equivalent** if they generate the same set of words.

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might ac-  
tually be  
wrong

**Decision Functions** Our previous definition of formal languages is very general and many subclasses of languages like regular languages or context-free languages are known in the literature. However, in the context of SNARK development, languages are commonly defined as **decision problems** where a so-called **deciding relation**  $R \subset \Sigma^*$  decides whether a given tuple  $x \in \Sigma^*$  is a word in the language or not. If  $x \in R$  then  $x$  is a word in the associated language  $L_R$  and if  $x \notin R$  then not. The relation  $R$  therefore summarizes the grammar of language  $L_R$ .

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this, too  
distracting

Unfortunately, in some literature on proof systems,  $x \in R$  is often written as  $R(x)$ , which is misleading since in general  $R$  is not a function but a relation in  $\Sigma^*$ . For the sake of this book, we therefore adopt a different point of view and work with what we might call a **decision function** instead:

$$R : \Sigma^* \rightarrow \{true, false\} \quad (6.1)$$

Decision functions decide if a tuple  $x \in \Sigma^*$  is an element of a language or not. In case a decision function is given, the associated language itself can be written as the set of all tuples that are decided by  $R$ , i.e as the set:

$$L_R := \{x \in \Sigma^* \mid R(x) = true\} \quad (6.2)$$

In the context of formal languages and decision problems, a **statement**  $S$  is the claim that language  $L$  contains a word  $x$ , i.e a statement claims that there exist some  $x \in L$ . A constructive **proof** for statement  $S$  is given by some string  $P \in \Sigma^*$  and such a proof is **verified** by checking  $R(P) = true$ . In this case,  $P$  is called an **instance** of the statement  $S$ .

While the term **language** might suggest a deeper relation to the well known **natural languages** like English, formal languages and natural languages differ in many ways. The following examples will provide some intuition about formal languages, highlighting the concepts of statements, proofs and instances:

*Example 103* (Alternating Binary strings). To consider a very basic formal language with an almost trivial grammar, consider the set  $\{0, 1\}$  of the two letters 0 and 1 as our alphabet  $\Sigma$  and imply the rule that a proper word must consist of alternating binary letters of arbitrary length.

Then, the associated language  $L_{alt}$  is the set of all finite binary tuples, where a 1 must follow a 0 and vice versa. So, for example,  $(1, 0, 1, 0, 1, 0, 1, 0, 1) \in L_{alt}$  is a proper word in this languages as is  $(0) \in L_{alt}$  or the empty word  $( ) \in L_{alt}$ . However, the binary tuple  $(1, 0, 1, 0, 1, 0, 1, 1, 1) \in \{0, 1\}^*$  is not a proper word, as it violates the grammar of  $L_{alt}$ : the last 3 letters are all 1. Furthermore, the tuple  $(0, A, 0, A, 0, A, 0)$  is not a proper word, as not all its letters are not from the proper alphabet: we defined the alphabet  $\Sigma$  as the set  $\{0, 1\}$ , and  $A$  is not part of that set.

Attempting to write the grammar of this language in a more formal way, we can define the following decision function:

$$R : \{0, 1\}^* \rightarrow \{true, false\} ; (x_0, x_1, \dots, x_n) \mapsto \begin{cases} true & x_{j-1} \neq x_j \text{ for all } 1 \leq j \leq n \\ false & \text{else} \end{cases}$$

We can use this function to decide if arbitrary binary tuples are words in  $L_{alt}$  or not. Some examples are given below:

binary  
tuples

- 4134 •  $R(1, 0, 1) = \text{true}$ ,
- 4135 •  $R(0) = \text{true}$ ,
- 4136 •  $R() = \text{true}$ ,
- 4137 • but  $R(1, 1) = \text{false}$ .

4138 Inside our language  $L_{alt}$ , it makes sense to claim the following statement: “There exists an  
 4139 alternating string.” One way to prove this statement constructively is by providing an actual  
 4140 instance, that is, finding actual alternating string like  $x = (1, 0, 1)$ . Constructing string  $(1, 0, 1)$   
 4141 therefore proves the statement “There exists an alternating string.”, because it is easy to verify  
 4142 that  $R(1, 0, 1) = \text{true}$ .

4143 *Example 104 (Programming Language).* Programming languages are a very important class of  
 4144 formal languages. For these languages, the alphabet is usually (a subset) of the ASCII table,  
 4145 and the grammar is defined by the rules of the programming language’s compiler. Words, then,  
 4146 are nothing but properly written computer programs that the compiler accepts. The compiler  
 4147 can therefore be interpreted as the decision function.

4148 To give an unusual example strange enough to highlight the point, consider the program-  
 4149 ming language Malbolge as defined in XXX. This language was specifically designed to be  
 4150 almost impossible to use and writing programs in this language is a difficult task. An inter-  
 4151 esting claim is therefore the statement: “There exists a computer program in Malbolge”. As it  
 4152 turned out, proving this statement constructively, that is, by providing an actual instance of such  
 4153 a program, was not an easy task, as it took two years after the introduction of Malbolge to write  
 4154 a program that its compiler accepts. So, for two years, no one was able to prove the statement  
 4155 constructively.

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ence

To look at this high-level description more formally, we write  $L_{Malbolge}$  for the language that  
 uses the ASCII table as its alphabet and its words are tuples of ASCII letters that the Malbolge  
 compiler accepts. Proving the statement “There exists a computer program in Malbolge” is then  
 equivalent to the task of finding some word  $x \in L_{Malbolge}$ . The string

$(=<\#9] 6ZY327Uv4-QsqpMn\&+Ij\prime'E\%e\{Ab w=_:]Kw\%o44Uqp0/Q?xNvL:'H\%c\#DD2\wedge WV>gY;dts76qKJImZk j$

4156 is an example of such a proof, as it is excepted by the Malbolge compiler and is compiled to  
 4157 an executable binary that displays “Hello, World.” (See XXX). In this example, the Malbolge  
 4158 compiler therefore serves as the verification process.

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*Example 105 (The Empty Language).* To see that not every language has even one word, con-  
 sider the alphabet  $\Sigma = \mathbb{Z}_6$ , where  $\mathbb{Z}_6$  is the ring of modular 6 arithmetics as derived in example  
 8 in chapter 3, together with the following decision function

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reference

$$R_\emptyset : \mathbb{Z}_6^* \rightarrow \{\text{true}, \text{false}\} ; (x_1, \dots, x_n) \mapsto \begin{cases} \text{true} & n = 4 \text{ and } x_1 \cdot x_1 = 2 \\ \text{true} & \text{else} \end{cases}$$

4159 We write  $L_\emptyset$  for the associated language. As we can see from the multiplication table of  $\mathbb{Z}_6$   
 4160 in example 8 in chapter 3, the ring  $\mathbb{Z}_6$  does not contain any element  $x$  such that  $x^2 = 2$ , which  
 4161 implies  $R_\emptyset(x_1, \dots, x_n) = \text{false}$  for all tuples  $(x_1, \dots, x_n) \in \Sigma^*$ . The language therefore does  
 4162 not contain any words. Proving the statement “There exists a word in  $L_\emptyset$ ” constructively by  
 4163 providing an instance is therefore impossible. The verification will never check any tuple.

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*Example 106 (3-Factorization).* We will use the following simple example repeatedly throughout this book. The task is to develop a SNARK that proves knowledge of three factors of an element from the finite field  $\mathbb{F}_{13}$ . There is nothing particularly useful about this example from an application point of view, however, in a sense, it is the most simple example that gives rise to a non trivial SNARK in some of the most common zero-knowledge proof systems.

Formalizing the high-level description, we use  $\Sigma := \mathbb{F}_{13}$  as the underlying alphabet of this problem and define the language  $L_{3.fac}$  to consists of those tuples of field elements from  $\mathbb{F}_{13}$  that contain exactly 4 letters  $w_1, w_2, w_3, w_4$  which satisfy the equation  $w_1 \cdot w_2 \cdot w_3 = w_4$ .

So, for example, the tuple  $(2, 12, 4, 5)$  is a word in  $L_{3.fac}$ , while neither  $(2, 12, 11)$ , nor  $(2, 12, 4, 7)$  nor  $(2, 12, 7, 168)$  are words in  $L_{3.fac}$  as they don't satisfy the grammar or are not define over the proper alphabet.

We can describe the language  $L_{3.fac}$  more formally by introducing a decision function (as described in equation 6.1):

$$R_{3.fac} : \mathbb{F}_{13}^* \rightarrow \{true, false\} ; (x_1, \dots, x_n) \mapsto \begin{cases} true & n = 4 \text{ and } x_1 \cdot x_2 \cdot x_3 = x_4 \\ false & else \end{cases}$$

Having defined the language  $L_{3.fac}$ , it then makes sense to claim the statement "There is a word in  $L_{3.fac}$ ". The way  $L_{3.fac}$  is designed, this statement is equivalent to the statement "There are four elements  $w_1, w_2, w_3, w_4$  from the finite field  $\mathbb{F}_{13}$  such that the equation  $w_1 \cdot w_2 \cdot w_3 = w_4$  holds."

Proving the correctness of this statement constructively means to actually find some concrete field elements like  $x_1 = 2, x_2 = 12, x_3 = 4$  and  $x_4 = 5$  that satisfy the relation  $R_{3.fac}$ . The tuple  $(2, 12, 4, 5)$  is therefore a constructive proof for the statement and the computation  $R_{3.fac}(2, 12, 4, 5) = true$  is a verification of that proof. In contrast, the tuple  $(2, 12, 4, 7)$  is not a proof of the statement, since the check  $R_{3.fac}(2, 12, 4, 7) = false$  does not verify the proof.

*Example 107 (Tiny-jubjub Membership).* In our main example, we derive a SNARK that proves a pair  $(x, y)$  of field elements from  $\mathbb{F}_{13}$  to be a point on the tiny-jubjub curve in its **Edwards form** (see section 5.1.3).

In the first step, we define a language such that points on the tiny-jubjub curve are in 1:1 correspondence with words in that language.

Since the tiny-jubjub curve is an elliptic curve over the field  $\mathbb{F}_{13}$ , we choose the alphabet  $\Sigma = \mathbb{F}_{13}$ . In this case, the set  $\mathbb{F}_{13}^*$  consists of all finite strings of field elements from  $\mathbb{F}_{13}$ . To define the grammar, recall from 66 that a point on the tiny-jubjub curve is a pair  $(x, y)$  of field elements such that  $3 \cdot x^2 + y^2 = 1 + 8 \cdot x^2 \cdot y^2$ . We can use this equation to derive the following decision function:

$$R_{tiny.jj} : \mathbb{F}_{13}^* \rightarrow \{true, false\} ; (x_1, \dots, x_n) \mapsto \begin{cases} true & n = 2 \text{ and } 3 \cdot x_1^2 + x_2^2 = 1 + 8 \cdot x_1^2 \cdot x_2^2 \\ false & else \end{cases}$$

The associated language  $L_{tiny.jj}$  is then given as the set of all strings from  $\mathbb{F}_{13}^*$  that are mapped onto *true* by  $R_{tiny.jj}$ . We get

$$L_{tiny.jj} = \{(x_1, \dots, x_n) \in \mathbb{F}_{13}^* \mid R_{tiny.jj}(x_1, \dots, x_n) = true\}$$

We can claim the statement "There is a word in  $L_{tiny.jj}$ " and because  $L_{tiny.jj}$  is defined by  $R_{tiny.jj}$ , this statement is equivalent to the claim "The tiny-jubjub curve in its Edwards form has curve a point."

A constructive proof for this statement is a pair  $(x, y)$  of field elements that satisfies the Edwards equation. Example 66 therefore implies that the tuple  $(11, 6)$  is a constructive proof

Are we using  $w$  and  $x$  interchangeably or is there a difference between them?

check reference

jubjub

check reference

check reference

check wording

check reference



and the computation  $R_{tiny.jj}(11,6) = true$  is a proof verification. In contrast, the tuple  $(1,1)$  is not a proof of the statement, since the check  $R_{tiny.jj}(1,1) = false$  does not verify the proof.

*Exercise 39.* Consider exercise XXX again. Define a decision function such that the associated language  $L_{Exercise_{XXX}}$  consist precisely of all solutions to the equation  $5x + 4 = 28 + 2x$  over  $\mathbb{F}_{13}$ . Provide a constructive proof for the claim: “There exist a word in  $L_{Exercise_{XXX}}$  and verify the proof.

*Exercise 40.* Consider the modular 6 arithmetics  $\mathbb{Z}_6$  from example 8 in chapter 3, the alphabet  $\Sigma = \mathbb{Z}_6$  and the decision function

$$R_{example\_8} : \Sigma^* \rightarrow \{true, false\} ; x \mapsto \begin{cases} true & x.len() = 1 \text{ and } 3 \cdot x + 3 = 0 \\ false & \text{else} \end{cases}$$

Compute all words in the associated language  $L_{example\_8}$ , provide a constructive proof for the statement “There exist a word in  $L_{example\_8}$ ” and verify the proof.

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**Instance and Witness** As we have seen in the previous paragraph, statements provide membership claims in formal languages, and instances serve as constructive proofs for those claims. However, in the context of **zero-knowledge** proof systems, our naive notion of constructive proofs is refined in such a way that its possible to hide parts of the proof instance and still be able to prove the statement. In this context, it is therefore necessary to split a proof into a **public part** called the **instance** and a private part called a **witness**.

To account for this separation of a proof instance into a public and a private part, our previous definition of formal languages needs a refinement in the context of zero-knowledge proof systems. Instead of a single alphabet, the refined definition considers two alphabets  $\Sigma_I$  and  $\Sigma_W$ , and a decision function defined as follows:

$$R : \Sigma_I^* \times \Sigma_W^* \rightarrow \{true, false\} ; (i; w) \mapsto R(i; w) \quad (6.3)$$

Words are therefore tuples  $(i; w) \in \Sigma_I^* \times \Sigma_W^*$  with  $R(i; w) = true$ . The refined definition differentiates between public inputs  $i \in \Sigma_I$  and private inputs  $w \in \Sigma_W$ . The public input  $i$  is called an **instance** and the private input  $w$  is called a **witness** of  $R$ .

If a decision function is given, the associated language is defined as the set of all tuples from the underlying alphabet that are verified by the decision function:

$$L_R := \{(i; w) \in \Sigma_I^* \times \Sigma_W^* \mid R(i; w) = true\} \quad (6.4)$$

In this refined context, a **statement**  $S$  is a claim that, given an instance  $i \in \Sigma_I^*$ , there is a witness  $w \in \Sigma_W^*$  such that language  $L$  contains a word  $(i; w)$ . A constructive **proof** for statement  $S$  is given by some string  $P = (i; w) \in \Sigma_I^* \times \Sigma_W^*$  and a proof is **verified** by checking  $R(P) = true$ .

It is worth understanding the difference between statements as defined in XXX and the refined notion of statements from this paragraph. While statements in the sense of the previous paragraph can be seen as membership claims, statements in the refined definition can be seen as knowledge-proofs, where a prover claims knowledge of a witness for a given instance. For a more detailed discussion on this topic see [XXX sec 1.4]

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*Example 108* (SHA256 – Knowledge of Preimage). One of the most common examples in the context of zero-knowledge proof systems is the **knowledge-of-a-preimage proof** for some cryptographic hash function like *SHA256*, where a publicly known *SHA256* digest value is

given, and the task is to prove knowledge of a preimage for that digest under the  $SHA256$  function, without revealing that preimage.

To understand this problem in detail, we have to introduce a language able to describe the knowledge-of-preimage problem in such a way that the claim “Given digest  $i$ , there is a preimage  $w$  such that  $SHA256(w) = i$ ” becomes a statement in that language. Since  $SHA256$  is a function

$$SHA256 : \{0, 1\}^* \rightarrow \{0, 1\}^{256}$$

that maps binary strings of arbitrary length onto binary strings of length 256 and we want to prove knowledge of preimages, we have to consider binary strings of size 256 as instances and binary strings of arbitrary length as witnesses.

An appropriate alphabet  $\Sigma_I$  for the set of all instances and an appropriate alphabet  $\Sigma_W$  for the set of all witnesses is therefore given by the set  $\{0, 1\}$  of the two binary letters and a proper decision function is given by:

$$R_{SHA256} : \{0, 1\}^* \times \{0, 1\}^* \rightarrow \{true, false\};$$

$$(i; w) \mapsto \begin{cases} true & i.len() = 256, i = SHA256(w) \\ false & else \end{cases}$$

We write  $L_{SHA256}$  for the associated language and note that it consists of words, which are tuples  $(i; w)$  such that the instance  $i$  is the  $SHA256$  image of the witness  $w$ .

Given some instance  $i \in \{0, 1\}^{256}$ , a statement in  $L_{SHA256}$  is the claim “Given digest  $i$ , there is a preimage  $w$  such that  $SHA256(w) = i$ ”, which is exactly what the knowledge-of-preimage problem is about. A constructive proof for this statement is therefore given by a preimage  $w$  to the digest  $i$  and proof verification is achieved by checking that  $SHA256(w) = i$ .

*Example 109 (3-factorization).* To give an intuition about the implication of refined languages, consider  $L_{3, fac}$  from example 106 again. As we have seen, a constructive proof in  $L_{3, fac}$  is given by 4 field elements  $x_1, x_2, x_3$  and  $x_4$  from  $\mathbb{F}_{13}$  such that the product in modular 13 arithmetics of the first three elements is equal to the 4'th element.

Splitting words from  $L_{3, fac}$  into private and public parts, we can reformulate the problem and introduce different levels of privacy into the problem. For example, we could reformulate the membership statement of  $L_{3, fac}$  into a statement where all factors  $x_1, x_2, x_3$  of  $x_4$  are private and only the product  $x_4$  is public. A statement for this reformulation is then expressed by the claim: “Given a publicly known field element  $x_4$ , there are three private factors of  $x_4$ ”. Assuming some instance  $x_4$ , a constructive proof for the associated knowledge claim is then provided by any tuple  $(x_1, x_2, x_3)$  such that  $x_1 \cdot x_2 \cdot x_3 = x_4$ .

At this point, it is important to note that, while constructive proofs in the refinement don't look very different from constructive proofs in the original language, we will see in XXX that there are proof systems able to prove the statement (at least with high probability) without revealing anything about the factors  $x_1, x_2$ , or  $x_3$ . This is why the importance of the refinement only becomes clear once more elaborate proofing methods beyond naive constructive proofs are provided.

We can formalize this new language, which we might call  $L_{3, fac\_zk}$ , by defining the following decision function:

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reference

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Definition 6.1.0.1.

$$R_{3.fac\_zk} : \mathbb{F}_{13}^* \times \mathbb{F}_{13}^* \rightarrow \{true, false\};$$

$$((i_1, \dots, i_n); (w_1, \dots, w_m)) \mapsto \begin{cases} true & n = 1, m = 3, i_1 = w_1 \cdot w_2 \cdot w_3 \\ false & else \end{cases}$$

The associated language  $L_{3.fac\_zk}$  is defined by all tuples from  $\mathbb{F}_{13}^* \times \mathbb{F}_{13}^*$  that are mapped onto *true* under the decision function  $R_{3.fac\_zk}$ .

Considering the distinction we made between the instance and the witness part in  $L_{3.fac\_zk}$ , one might ask why we chose the factors  $x_1, x_2$  and  $x_3$  to be the witness and the product  $x_4$  to be the instance and why we didn't choose another combination? This was an arbitrary choice in the example. Every other combination of private and public factors would be equally valid. For example, it would be possible to declare all variables as private or to declare all variables as public. Actual choices are determined by the application only.

*Example 110 (The Tiny-Jubjub Curve).* Consider the language  $L_{tiny.jj}$  from example 107. As we have seen, a constructive proof in  $L_{tiny.jj}$  is given by a pair  $(x_1, x_2)$  of field elements from  $\mathbb{F}_{13}$  such that the pair is a point of the tiny-jubjub curve in its Edwards representation.

We look at a reasonable splitting of words from  $L_{tiny.jj}$  into private and public parts. The two obvious choices are to either choose both coordinates  $x_1$  as  $x_2$  as public inputs, or to choose both coordinates  $x_1$  as  $x_2$  as private inputs.

In case both coordinates are public, we define the grammar of the associated language by introducing the following decision function:

$$R_{tiny.jj.1} : \mathbb{F}_{13}^* \times \mathbb{F}_{13}^* \rightarrow \{true, false\};$$

$$(I_1, \dots, I_n; W_1, \dots, W_m) \mapsto \begin{cases} true & n = 2, m = 0 \text{ and } 3 \cdot I_1^2 + I_2^2 = 1 + 8 \cdot I_1^2 \cdot I_2^2 \\ false & else \end{cases}$$

The language  $L_{tiny.jj.1}$  is defined as the set of all strings from  $\mathbb{F}_{13}^* \times \mathbb{F}_{13}^*$  that are mapped onto *true* by  $R_{tiny.jj.1}$ .

In case both coordinates are private, we define the grammar of the associated refined language by introducing the following decision function:

$$R_{tiny.jj\_zk} : \mathbb{F}_{13}^* \times \mathbb{F}_{13}^* \rightarrow \{true, false\};$$

$$(I_1, \dots, I_n; W_1, \dots, W_m) \mapsto \begin{cases} true & n = 0, m = m \text{ and } 3 \cdot W_1^2 + W_2^2 = 1 + 8 \cdot W_1^2 \cdot W_2^2 \\ false & else \end{cases}$$

The language  $L_{tiny.jj\_zk}$  is defined as the set of all strings from  $\mathbb{F}_{13}^* \times \mathbb{F}_{13}^*$  that are mapped onto *true* by  $R_{tiny.jj\_zk}$ .

*Exercise 41.* Consider the modular 6 arithmetics  $\mathbb{Z}_6$  from example 8 in chapter 3 as alphabets  $\Sigma_I$  and  $\Sigma_W$  and the following decision function

$$R_{linear} : \Sigma^* \times \Sigma^* \rightarrow \{true, false\};$$

$$(i; w) \mapsto \begin{cases} true & i.len() = 3 \text{ and } w.len() = 1 \text{ and } i_1 \cdot w_1 + i_2 = i_3 \\ false & else \end{cases}$$

Which of the following instances  $(i_1, i_2, i_3)$  has a proof of knowledge in  $L_{linear}$ ?

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• (3, 3, 0)

• (2, 1, 0)

• (4, 4, 2)

*Exercise 42* (Edwards Addition on Tiny-Jubjub). Consider the tiny-jubjub curve together with its Edwards addition law from example XXX. Define an instance alphabet  $\Sigma_I$ , a witness alphabet  $\Sigma_W$  and a decision function  $R_{add}$  with associated language  $L_{add}$  such that a string  $(i; w) \in \Sigma_I^* \times \Sigma_W^*$  is a word in  $L_{add}$  if and only if  $i$  is a pair of curve points on the tiny-jubjub curve in Edwards form and  $w$  is the Edwards sum of those curve points.

Choose some instance  $i \in \Sigma_I^*$ , provide a constructive proof for the statement “There is a witness  $w \in \Sigma_W^*$  such that  $(i; w)$  is a word in  $L_{add}$ ” and verify that proof. Then find some instance  $i \in \Sigma_I^*$  such that  $i$  has no knowledge proof in  $L_{add}$ .

add reference

**Modularity** From a developers perspective, it is often useful to construct complex statements and their representing languages from simple ones. In the context of zero-knowledge proof systems, those simple building blocks are often called **gadgets**, and gadget libraries usually contain representations of atomic types like booleans, integers, various hash functions, elliptic curve cryptography and many more. In order to synthesize statements, developers then combine predefined gadgets into complex logic. We call the ability to combine statements into more complex statements **modularity**.

To understand the concept of modularity on the level of formal languages defined by decision functions, we need to look at the **intersection** of two languages, which exists whenever both languages are defined over the same alphabet. In this case, the intersection is a language that consists of strings which are words in both languages.

To be more precise, let  $L_1$  and  $L_2$  be two languages defined over the same instance and witness alphabets  $\Sigma_I$  and  $\Sigma_W$ . Then the intersection  $L_1 \cap L_2$  of  $L_1$  and  $L_2$  is defined as

$$L_1 \cap L_2 := \{x \mid x \in L_1 \text{ and } x \in L_2\} \quad (6.5)$$

If both languages are defined by decision functions  $R_1$  and  $R_2$ , the following function is a decision function for the intersection language  $L_1 \cap L_2$ :

$$R_{L_1 \cap L_2} : \Sigma_I^* \times \Sigma_W^* \rightarrow \{true, false\}; (i, w) \mapsto R_1(i, w) \text{ and } R_2(i, w) \quad (6.6)$$

Thus, the intersection of two decision-function-based languages is a also decision-function-based language. This is important from an implementations point of view: It allows us to construct complex decision functions, their languages and associated statements from simple building blocks. Given a publicly known instance  $i \in \Sigma_I^*$  a statement in an intersection language then claims knowledge of a witness that satisfies all relations simultaneously.

## 6.2 Statement Representations

As we have seen in the previous section, formal languages and their definitions by decision functions are a powerful tool to describe statements in a formally rigorous manner.

However, from the perspective of existing zero-knowledge proof systems, not all ways to actually represent decision functions are equally useful. Depending on the proof system, some are more suitable than others. In this section, will describe two of the most common ways to represent decision functions and their statements.

### 6.2.1 Rank-1 Quadratic Constraint Systems

Although decision functions are expressible in various ways, many contemporary proofing systems require the deciding relation to be expressed in terms of a system of quadratic equations over a finite field. This is true in particular for pairing-based proofing systems like XXX, roughly because it is possible to check solutions to those equations “in the exponent” of pairing-friendly cryptographic groups.

add reference

In this section, we will therefore have a closer look at a particular type of quadratic equation called **rank-1 quadratic constraints systems**, which are a common standard in zero-knowledge proof systems. We will start with a general introduction to those systems and then look at their relation to formal languages. We will look into a common way to compute solutions to those systems, and then show how a simple compiler might derive rank-1 constraint systems from more high-level programming code.

**R1CS representation** To understand what **rank-1 (quadratic) constraint systems** are in detail, let  $\mathbb{F}$  be a field,  $n, m$  and  $k \in \mathbb{N}$  three numbers and  $a_j^i, b_j^i$  and  $c_j^i \in \mathbb{F}$  constants from  $\mathbb{F}$  for every index  $0 \leq j \leq n+m$  and  $1 \leq i \leq k$ . Then a rank-1 constraint system (R1CS) is defined as follows:

*Definition 6.2.1.1. R1CS representation*

$$\begin{aligned} (a_0^1 + \sum_{j=1}^n a_j^1 \cdot I_j + \sum_{j=1}^m a_{n+j}^1 \cdot W_j) \cdot (b_0^1 + \sum_{j=1}^n b_j^1 \cdot I_j + \sum_{j=1}^m b_{n+j}^1 \cdot W_j) &= c_0^1 + \sum_{j=1}^n c_j^1 \cdot I_j + \sum_{j=1}^m c_{n+j}^1 \cdot W_j \\ &\vdots \\ (a_0^k + \sum_{j=1}^n a_j^k \cdot I_j + \sum_{j=1}^m a_{n+j}^k \cdot W_j) \cdot (b_0^k + \sum_{j=1}^n b_j^k \cdot I_j + \sum_{j=1}^m b_{n+j}^k \cdot W_j) &= c_0^k + \sum_{j=1}^n c_j^k \cdot I_j + \sum_{j=1}^m c_{n+j}^k \cdot W_j \end{aligned}$$

If a rank-1 constraint system is given, the parameter  $k$  is called the **number of constraints**. If a tuple  $(I_1, \dots, I_n; W_1, \dots, W_m)$  of field elements satisfies these equations,  $(I_1, \dots, I_n)$  is called an **instance** and  $(W_1, \dots, W_m)$  is called an associated **witness** of the system.

*Remark 1 (Matrix notation).* The presentation of rank-1 constraint systems can be simplified using the notation of vectors and matrices, which abstracts over the indices. In fact if  $x = (1, I, W) \in \mathbb{F}^{1+n+m}$  is a  $(n+m+1)$ -dimensional vector,  $A, B, C$  are  $(n+m+1) \times k$ -dimensional matrices and  $\odot$  is the **Schur/Hadamard product**, then a R1CS can be written as

Schur/Hadamard product

$$Ax \odot Bx = Cx$$

However, since we did not introduce matrix calculus in the book, we use XXX as the defining equation for rank-1 constraints systems. We only highlighted the matrix notation, because it is sometimes used in the literature.

add reference

Generally speaking, the idea of a rank-1 constraint system is to keep track of all the values that any variable can assume during a computation and to bind the relationships among all those variables that are implied by the computation itself. Enforcing relations between all the steps of a computer program, the execution is then constrained to be computed in exactly the expected way without any opportunity for deviations. In this sense, solutions to rank-1 constraint systems are proofs of proper program execution.

*Example 111 (3-Factorization).* To provide a better intuition of rank-1 constraint systems, consider the language  $L_{3\text{-}fac\_zk}$  from example 106 again. As we have seen,  $L_{3\text{-}fac\_zk}$  consists of words  $(I_1; W_1, W_2, W_3)$  over the alphabet  $\mathbb{F}_{13}$  such that  $I_1 = W_1 \cdot W_2 \cdot W_3$ . We show how to rewrite the decision function as a rank-1 constraint system.

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Since R1CS are systems of quadratic equations, expressions like  $W_1 \cdot W_2 \cdot W_3$  which contain products of more than two factors (which are therefore not quadratic) have to be rewritten in a process often called **flattening**. To flatten the defining equation  $I_1 = W_1 \cdot W_2 \cdot W_3$  of  $L_{3, \text{fac\_zk}}$  we introduce a new variable  $W_4$ , which captures two of the three multiplications in  $W_1 \cdot W_2 \cdot W_3$ . We get the following two constraints

$$W_1 \cdot W_2 = W_4 \quad \text{constraint 1}$$

$$W_4 \cdot W_3 = I_1 \quad \text{constraint 2}$$

Given some instance  $I_1$ , any solution  $(W_1, W_2, W_3, W_4)$  to this system of equations provides a solution to the original equation  $I_1 = W_1 \cdot W_2 \cdot W_3$  and vice versa. Both equations are therefore equivalent in the sense that solutions are in a 1:1 correspondence.

Looking at both equations, we see how each constraint enforces a step in the computation. In fact, the first constraint forces any computation to multiply the witness  $W_1$  and  $W_2$  first. Otherwise it would not be possible to compute the witness  $W_4$ , which is needed to solve the second constraint. Witness  $W_4$  therefore expresses the constraining of an intermediate computational state.

At this point, one might ask why equation 1 constrains the system to compute  $W_1 \cdot W_2$  first, since computing  $W_2 \cdot W_3$ , or  $W_1 \cdot W_3$  in the beginning and then multiplying with the remaining factor gives the exact same result. The reason is that the way we designed the R1CS prohibits any of these alternative computations, which shows that R1CS are in general **not unique** descriptions of a language: many different R1CS are able to describe the same problem.

To see that the two quadratic equations qualify as a rank-1 constraint system, choose the parameter  $n = 1$ ,  $m = 4$  and  $k = 2$  as well as

$$\begin{array}{cccccc} a_0^1 = 0 & a_1^1 = 0 & a_2^1 = 1 & a_3^1 = 0 & a_4^1 = 0 & a_5^1 = 0 \\ a_0^2 = 0 & a_1^2 = 0 & a_2^2 = 0 & a_3^2 = 0 & a_4^2 = 0 & a_5^2 = 1 \end{array}$$

$$\begin{array}{cccccc} b_0^1 = 0 & b_1^1 = 0 & b_2^1 = 0 & b_3^1 = 1 & b_4^1 = 0 & b_5^1 = 0 \\ b_0^2 = 0 & b_1^2 = 0 & b_2^2 = 0 & b_3^2 = 0 & b_4^2 = 1 & b_5^2 = 0 \end{array}$$

$$\begin{array}{cccccc} c_0^1 = 0 & c_1^1 = 0 & c_2^1 = 0 & c_3^1 = 0 & c_4^1 = 0 & c_5^1 = 1 \\ c_0^2 = 0 & c_1^2 = 1 & c_2^2 = 0 & c_3^2 = 0 & c_4^2 = 0 & c_5^2 = 0 \end{array}$$

With this choice, the rank-1 constraint system of our 3-factorization problem can be written in its most general form as follows:

$$\begin{aligned} (a_0^1 + a_1^1 I_1 + a_2^1 W_1 + a_3^1 W_2 + a_4^1 W_3 + a_5^1 W_4) \cdot (b_0^1 + b_1^1 I_1 + b_2^1 W_1 + b_3^1 W_2 + b_4^1 W_3 + b_5^1 W_4) &= (c_0^1 + c_1^1 I_1 + c_2^1 W_1 + c_3^1 W_2 + c_4^1 W_3 + c_5^1 W_4) \\ (a_0^2 + a_1^2 I_1 + a_2^2 W_2 + a_3^2 W_2 + a_4^2 W_3 + a_5^2 W_4) \cdot (b_0^2 + b_1^2 I_1 + b_2^2 W_2 + b_3^2 W_2 + b_4^2 W_3 + b_5^2 W_4) &= (c_0^2 + c_1^2 I_1 + c_2^2 W_2 + c_3^2 W_2 + c_4^2 W_3 + c_5^2 W_4) \end{aligned}$$

*Example 112* (The Tiny-Jubjub curve). Consider the languages  $L_{\text{tiny}, jj, 1}$  from example 107, which consist of words  $(I_1, I_2)$  over the alphabet  $\mathbb{F}_{13}$  such that  $3 \cdot I_1^2 + I_2^2 = 1 + 8 \cdot I_1^2 \cdot I_2^2$ .

We derive a rank-1 constraint system such that its associated language is equivalent to  $L_{\text{tiny}, jj, 1}$ . To achieve this, we first rewrite the defining equation:

$$\begin{aligned} 3 \cdot I_1^2 + I_2^2 &= 1 + 8 \cdot I_1^2 \cdot I_2^2 && \Leftrightarrow \\ 0 &= 1 + 8 \cdot I_1^2 \cdot I_2^2 - 3 \cdot I_1^2 - I_2^2 && \Leftrightarrow \\ 0 &= 1 + 8 \cdot I_1^2 \cdot I_2^2 + 10 \cdot I_1^2 + 12 \cdot I_2^2 \end{aligned}$$

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Since R1CSs are systems of quadratic equations, we have to reformulate this expression into a system of quadratic equations. To do so, we have to introduce new variables that constrain intermediate steps in the computation and we have to decide if those variables should be public or private. We decide to declare all new variables as private and get the following constraints

$$\begin{array}{ll}
 I_1 \cdot I_1 = W_1 & \text{constraint 1} \\
 I_2 \cdot I_2 = W_2 & \text{constraint 2} \\
 (8 \cdot W_1) \cdot W_2 = W_3 & \text{constraint 3} \\
 (12 \cdot W_2 + W_3 + 10 \cdot W_1 + 1) \cdot 1 = 0 & \text{constraint 4}
 \end{array}$$

To see that these four quadratic equations qualify as a rank-1 constraint system according to definition XXX, choose the parameter  $n = 2$ ,  $m = 3$  and  $k = 4$ :

add reference

$$\begin{array}{llllll}
 a_0^1 = 0 & a_1^1 = 1 & a_2^1 = 0 & a_3^1 = 0 & a_4^1 = 0 & a_5^1 = 0 \\
 a_0^2 = 0 & a_1^2 = 0 & a_2^2 = 1 & a_3^2 = 0 & a_4^2 = 0 & a_5^2 = 0 \\
 a_0^3 = 0 & a_1^3 = 0 & a_2^3 = 0 & a_3^3 = 8 & a_4^3 = 0 & a_5^3 = 0 \\
 a_0^4 = 1 & a_1^4 = 0 & a_2^4 = 0 & a_3^4 = 10 & a_4^4 = 12 & a_5^4 = 1 \\
 \\ 
 b_0^1 = 0 & b_1^1 = 1 & b_2^1 = 0 & b_3^1 = 0 & b_4^1 = 0 & b_5^1 = 0 \\
 b_0^2 = 0 & b_1^2 = 0 & b_2^2 = 1 & b_3^2 = 0 & b_4^2 = 0 & b_5^2 = 0 \\
 b_0^3 = 0 & b_1^3 = 0 & b_2^3 = 0 & b_3^3 = 0 & b_4^3 = 1 & b_5^3 = 0 \\
 b_0^4 = 1 & b_1^4 = 0 & b_2^4 = 0 & b_3^4 = 0 & b_4^4 = 0 & b_5^4 = 0 \\
 \\ 
 c_0^1 = 0 & c_1^1 = 0 & c_2^1 = 0 & c_3^1 = 1 & c_4^1 = 0 & c_5^1 = 0 \\
 c_0^2 = 0 & c_1^2 = 0 & c_2^2 = 0 & c_3^2 = 0 & c_4^2 = 1 & c_5^2 = 0 \\
 c_0^3 = 0 & c_1^3 = 0 & c_2^3 = 0 & c_3^3 = 0 & c_4^3 = 0 & c_5^3 = 1 \\
 c_0^4 = 0 & c_1^4 = 0 & c_2^4 = 0 & c_3^4 = 0 & c_4^4 = 0 & c_5^4 = 0
 \end{array}$$

With this choice, the rank-1 constraint system of our tiny-jubjub curve point problem can be written in its most general form as follows:

$$\begin{aligned}
 (a_0^1 + a_1^1 I_1 + a_2^1 I_2 + a_3^1 W_1 + a_4^1 W_2 + a_5^1 W_3) \cdot (b_0^1 + b_1^1 I_1 + b_2^1 I_2 + b_3^1 W_1 + b_4^1 W_2 + b_5^1 W_3) &= (c_0^1 + c_1^1 I_1 + c_2^1 I_2 + c_3^1 W_1 + c_4^1 W_2 + c_5^1 W_3) \\
 (a_0^2 + a_1^2 I_1 + a_2^2 I_2 + a_3^2 W_1 + a_4^2 W_2 + a_5^2 W_3) \cdot (b_0^2 + b_1^2 I_1 + b_2^2 I_2 + b_3^2 W_1 + b_4^2 W_2 + b_5^2 W_3) &= (c_0^2 + c_1^2 I_1 + c_2^2 I_2 + c_3^2 W_1 + c_4^2 W_2 + c_5^2 W_3) \\
 (a_0^3 + a_1^3 I_1 + a_2^3 I_2 + a_3^3 W_1 + a_4^3 W_2 + a_5^3 W_3) \cdot (b_0^3 + b_1^3 I_1 + b_2^3 I_2 + b_3^3 W_1 + b_4^3 W_2 + b_5^3 W_3) &= (c_0^3 + c_1^3 I_1 + c_2^3 I_2 + c_3^3 W_1 + c_4^3 W_2 + c_5^3 W_3) \\
 (a_0^4 + a_1^4 I_1 + a_2^4 I_2 + a_3^4 W_1 + a_4^4 W_2 + a_5^4 W_3) \cdot (b_0^4 + b_1^4 I_1 + b_2^4 I_2 + b_3^4 W_1 + b_4^4 W_2 + b_5^4 W_3) &= (c_0^4 + c_1^4 I_1 + c_2^4 I_2 + c_3^4 W_1 + c_4^4 W_2 + c_5^4 W_3)
 \end{aligned}$$

4363 In what follows, we write  $L_{jubjub}$  for the associated language that consists of solutions to the  
 4364 R1CS.

4365 To see that  $L_{jubjub}$  is equivalent to  $L_{tiny.jj.1}$ , let  $(I_1, I_2; W_1, W_2, W_3)$  be a word in  $L_{jubjub}$ , then  
 4366  $(I_1, I_2)$  is a word in  $L_{tiny.jj.1}$ , since the defining R1CS of  $L_{jubjub}$  implies that  $I_1$  and  $I_2$  satisfy the  
 4367 Edwards equation of the tiny-jubjub curve. On the other hand, let  $(I_1, I_2)$  be a word in  $L_{tiny.jj.1}$ .  
 4368 Then  $(I_1, I_2; I_1^2, I_2^2, 8 \cdot I_1^2 \cdot I_2^2)$  is a word in  $L_{jubjub}$  and both maps are inverses of each other.

4369 *Exercise 43.* Consider the language  $L_{tiny.jj\_zk}$  and define a rank-1 constraint relation with a  
 4370 decision function such that the associated language is equivalent to  $L_{tiny.jj\_zk}$ .

4371 **R1CS Satisfiability** To understand how rank-1 constraint systems define formal languages,  
 4372 observe that every R1CS over a field  $\mathbb{F}$  defines a decision function over the alphabet  $\Sigma_I \times \Sigma_W =$

4373  $\mathbb{F} \times \mathbb{F}$  in the following way:

$$R_{\text{R1CS}} : \mathbb{F}^* \times \mathbb{F}^* \rightarrow \{true, false\} ; (I; W) \mapsto \begin{cases} true & (I; W) \text{ satisfies R1CS} \\ false & \text{else} \end{cases} \quad (6.7)$$

4374 Every R1CS therefore defines a formal language. The grammar of this language is encoded  
 4375 in the constraints, words are solutions to the equations and a **statement** is a knowledge claim  
 4376 “Given instance  $I$ , there is a witness  $W$  such that  $(I; W)$  is a solution to the rank-1 constraint  
 4377 system”. A constructive proof to this claim is therefore an assignment of a field element to every  
 4378 witness variable, which is verified whenever the set of all instance and witness variables solves  
 4379 the R1CS.

*Remark 2* (R1CS satisfiability). It should be noted that in our definition, every R1CS defines its own language. However, in more theoretical approaches, another language usually called **R1CS satisfiability** is often considered, which is useful when it comes to more abstract problems like expressiveness or the computational complexity of the class of **all** R1CS. From our perspective, the R1CS satisfiability language is obtained by the union of all R1CS languages that are in our definition. To be more precise, let the alphabet  $\Sigma = \mathbb{F}$  be a field. Then

$$L_{\text{R1CS\_SAT}}(\mathbb{F}) = \{(i; w) \in \Sigma^* \times \Sigma^* \mid \text{there is a R1CS } R \text{ such that } R(i; w) = true\}$$

4380 *Example 113* (3-Factorization). Consider the language  $L_{3, \text{fac\_zk}}$  from example 106 and the  
 4381 R1CS defined in example ex:3-factorization-r1cs. As we have seen in ex:3-factorization-r1cs,  
 4382 solutions to the R1CS are in 1:1 correspondence with solutions to the decision function of  
 4383  $L_{3, \text{fac\_zk}}$ . Both languages are therefore equivalent in the sense that there is a 1:1 correspon-  
 4384 dence between words in both languages.

To give an intuition of what constructive proofs in  $L_{3, \text{fac\_zk}}$  look like, consider the instance  $I_1 = 11$ . To prove the statement “There exists a witness  $W$  such that  $(I_1; W)$  is a word in  $L_{3, \text{fac\_zk}}$ ” constructively, a proof has to provide assignments to all witness variables  $W_1, W_2, W_3$  and  $W_4$ . Since the alphabet is  $\mathbb{F}_{13}$ , an example assignment is given by  $W = (2, 3, 4, 6)$  since  $(I_1; W)$  satisfies the R1CS

$$\begin{array}{ll} W_1 \cdot W_2 = W_4 & \# 2 \cdot 3 = 6 \\ W_4 \cdot W_3 = I_1 & \# 6 \cdot 4 = 11 \end{array}$$

4385 A proper constructive proof is therefore given by  $P = (2, 3, 4, 6)$ . Of course,  $P$  is not the only  
 4386 possible proof for this statement. Since factorization is not unique in a field in general, another  
 4387 constructive proof is given by  $P' = (3, 5, 12, 2)$ .

*Example 114* (The tiny-jubjub curve). Consider the language  $L_{\text{jubjub}}$  from example 107 and  
 its associated R1CS. To see how constructive proofs in  $L_{\text{jubjub}}$  look like, consider the instance  
 $(I_1, I_2) = (11, 6)$ . To prove the statement “There exists a witness  $W$  such that  $(I_1, I_2; W)$  is a  
 word in  $L_{\text{jubjub}}$ ” constructively, a proof has to provide assignments to all witness variables  $W_1,$   
 $W_2$  and  $W_3$ . Since the alphabet is  $\mathbb{F}_{13}$ , an example assignment is given by  $W = (4, 10, 8)$  since  
 $(I_1, I_2; W)$  satisfies the R1CS

$$\begin{array}{ll} I_1 \cdot I_1 = W_1 & 11 \cdot 11 = 4 \\ I_2 \cdot I_2 = W_2 & 6 \cdot 6 = 10 \\ (8 \cdot W_1) \cdot W_2 = W_3 & (8 \cdot 4) \cdot 10 = 8 \\ (12 \cdot W_2 + W_3 + 10 \cdot W_1 + 1) \cdot 1 = 0 & 12 \cdot 10 + 8 + 10 \cdot 4 + 1 = 0 \end{array}$$

4388 A proper constructive proof is therefore given by  $P = (4, 10, 8)$ , which shows that the instance  
 4389  $(11, 6)$  is a point on the tiny-jubjub curve.

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**Modularity** As we discussed on page 135 XXX, it is often useful to construct complex statements and their representing languages from simple ones. Rank-1 constraint systems are particularly useful for this, as the intersection of two R1CS over the same alphabet results in a new R1CS over that same alphabet.

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To be more precise, let  $S_1$  and  $S_2$  be two R1CS over  $\mathbb{F}$ , then a new R1CS  $S_3$  is obtained by the intersection  $S_3 = S_1 \cap S_2$  of  $S_1$  and  $S_2$ . In this context, intersection means that both the equations of  $S_1$  **and** the equations of  $S_2$  have to be satisfied in order to provide a solution for the system  $S_3$ .

As a consequence, developers are able to construct complex R1CS from simple ones and this modularity provides the theoretical foundation for many R1CS compilers, as we will see in XXX.

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ence

### 6.2.2 Algebraic Circuits

As we have seen in the previous paragraphs, rank-1 constraint systems are quadratic equations such that solutions are knowledge proofs for the existence of words in associated languages. From the perspective of a proofer, it is therefore important to solve those equations efficiently.

However, in contrast to systems of linear equation, no general methods are known that solve systems of quadratic equations efficiently. Rank-1 constraint systems are therefore impractical from a proofers perspective and auxiliary information is needed that helps to compute solutions efficiently.

Methods which compute R1CS solutions are sometimes called **witness generator functions**. To provide a common example, we introduce another class of decision functions called **algebraic circuits**. As we will see, every algebraic circuit defines an associated R1CS and also provides an efficient way to compute solutions for that R1CS.

It can be shown that every space- and time-bounded computation is expressible as an algebraic circuit. Transforming high-level computer programs into those circuits is a process often called **flattening**.

To understand this in more detail, we will introduce our model for algebraic circuits and look at the concept of circuit execution and valid assignments. After that, we will show how to derive rank-1 constraint systems from circuits and how circuits are useful to compute solutions to their R1CS efficiently.

**Algebraic circuit representation** To see what algebraic circuits are, let  $\mathbb{F}$  be a field. An algebraic circuit is then a directed acyclic (multi)graph that computes a polynomial function over  $\mathbb{F}$ . Nodes with only outgoing edges (source nodes) represent the variables and constants of the function and nodes with only incoming edges (sink nodes) represent the outcome of the function. All other nodes have exactly two incoming edges and represent the defining field operations **addition** as well as **multiplication**. Graph edges represent the flow of the computation along the nodes.

To be more precise, we call a directed acyclic multi-graph  $C(\mathbb{F})$  an **algebraic circuit** over  $\mathbb{F}$  in this book if the following conditions hold:

*Definition 6.2.2.1. Algebraic circuit*

- The set of edges has a total order.
- Every source node has a label that represents either a variable or a constant from the field  $\mathbb{F}$ .

- Every sink node has exactly one incoming edge and a label that represents either a variable or a constant from the field  $\mathbb{F}$ .
- Every node that is neither a source nor a sink has exactly two incoming edges and a label from the set  $\{+, *\}$  that represents either addition or multiplication in  $\mathbb{F}$ .
- All outgoing edges from a node have the same label.
- Outgoing edges from a node with a label that represents a variable have a label.
- Outgoing edges from a node with a label that represents multiplication have a label, if there is at least one labeled edge in both input path.
- All incoming edges to sink nodes have a label.
- If an edge has two labels  $S_i$  and  $S_j$  it gets a new label  $S_i = S_j$ .
- No other edge has a label.
- Incoming edges to sink nodes that are labeled with a constant  $c \in \mathbb{F}$  are labeled with the same constant. Every other edge label is taken from the set  $\{W, I\}$  and indexed compatible with the order of the edge set.

It should be noted that the details in the definitions of algebraic circuits vary between different sources. We use this definition as it is conceptually straightforward and well-suited for pen-and-paper computations.

To get a better intuition of our definition, let  $C(\mathbb{F})$  be an algebraic circuit. Source nodes are the inputs to the circuit and either represent variables or constants. In a similar way, sink nodes represent termination points of the circuit and are either output variables or constants. Constant sink nodes enforce computational outputs to take on certain values.

Nodes that are neither source nodes nor sink nodes are called **arithmetic gates**. Arithmetic gates that are decorated with the “+”-label are called **addition-gates** and arithmetic gates that are decorated with the “·”-label are called **multiplication-gates**. Every arithmetic gate has exactly two inputs, represented by the two incoming edges.

Since the set of edges is ordered, we can write it as  $\{E_1, E_2, \dots, E_n\}$  for some  $n \in \mathbb{N}$  and we use those indices to index the edge labels, too. Edge labels are therefore either constants or symbols like  $I_j$ ,  $W_j$  or  $S_j$ , where  $j$  is an index compatible with the edge order. Labels  $I_j$  represent instance variables, labels  $W_j$  witness variables. Labels on the outgoing edges of input variables constrain the associated variable to that edge. Every other edge defines a constraining equation in the associated R1CS. We will explain this in more detail in XXX.

*Notation and Symbols 10.* In synthesizing algebraic circuits, assigning instance  $I_j$  or witness  $W_j$  labels to appropriate edges is often the final step. It is therefore convenient to not distinguish these two types of edges in previous steps. To account for that, we often simply write  $S_j$  for an edge label, indicating that the private/public property of the label is unspecified and it might represent an instance or a witness label.

*Example 115 (Generalized factorization SNARK).* To give a simple example of an algebraic circuit, consider our 3-factorization problem from example 106 again. To express the problem in the algebraic circuit model, consider the following function

$$f_{3.fac} : \mathbb{F}_{13} \times \mathbb{F}_{13} \times \mathbb{F}_{13} \rightarrow \mathbb{F}_{13}; (x_1, x_2, x_3) \mapsto x_1 \cdot x_2 \cdot x_3$$

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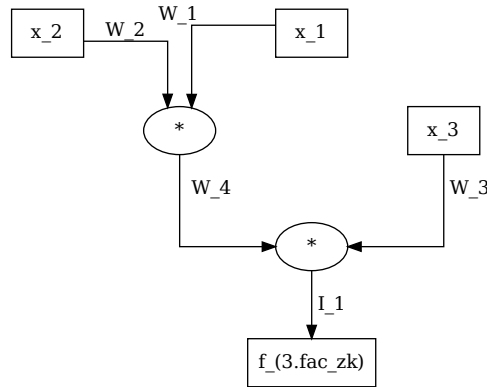
4472 Using this function, we can describe the zero-knowledge 3-factorization problem from 106,  
 4473 in the following way: Given instance  $I_1 \in \mathbb{F}_{13}$ , a valid witness is a preimage of  $f_{3, \text{fac}}$  at  
 4474 the point  $I_1$ , i.e., a valid witness consists of three values  $W_1, W_2$  and  $W_3$  from  $\mathbb{F}_{13}$  such that  
 4475  $f_{3, \text{fac}}(W_1, W_2, W_3) = I_1$ .

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reference

To see how this function can be transformed into an algebraic circuit over  $\mathbb{F}_{13}$ , it is a common first step to introduce brackets into the function's definition and then write the operations as binary operators, in order to highlight how exactly every field operation acts on its two inputs. Due to the associativity laws in a field, we have several choices. We choose

$$\begin{aligned} f_{3, \text{fac}}(x_1, x_2, x_3) &= x_1 \cdot x_2 \cdot x_3 && \# \text{ bracket choice} \\ &= (x_1 \cdot x_2) \cdot x_3 && \# \text{ operator notation} \\ &= \text{MUL}(\text{MUL}(x_1, x_2), x_3) \end{aligned}$$

4476 Using this expression, we can write an associated algebraic circuit by first constraining the  
 4477 variables to edge labels  $W_1 = x_1, W_2 = x_2$  and  $W_3 = x_3$  as well as  $I_1 = f_{3, \text{fac}}(x_1, x_2, x_3)$ , taking  
 4478 the distinction between private and public inputs into account. We then rewrite the operator  
 4479 representation of  $f_{3, \text{fac}}$  into circuit nodes and get the following:



4480

4481 In this case, the directed acyclic multi-graph is a binary tree with three leaves (the source  
 4482 nodes) labeled by  $x_1, x_2$  and  $x_3$ , one root (the single sink node) labeled by  $f(x_1, x_2, x_3)$  and two  
 4483 internal nodes, which are labeled as multiplication gates.

4484 The order we use to label the edges is chosen to make the edge labeling consistent with  
 4485 the choice of  $W_4$  as defined in definition 6.2.2.1. This order can be obtained by a depth-first  
 4486 right-to-left-first traversal algorithm.

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*Example 116.* To give a more realistic example of an algebraic circuit, look at the defining equation of the tiny-jubjub curve (66) again. A pair of field elements  $(x, y) \in \mathbb{F}_{13}^2$  is a curve point, precisely if the following equation holds:

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$$3 \cdot x^2 + y^2 = 1 + 8 \cdot x^2 \cdot y^2$$

To understand how one might transform this identity into an algebraic circuit, we first rewrite this equation by shifting all terms to the right. We get the following:

$$\begin{aligned} 3 \cdot x^2 + y^2 &= 1 + 8 \cdot x^2 \cdot y^2 && \Leftrightarrow \\ 0 &= 1 + 8 \cdot x^2 \cdot y^2 - 3 \cdot x^2 - y^2 && \Leftrightarrow \\ 0 &= 1 + 8 \cdot x^2 \cdot y^2 + 10 \cdot x^2 + 12 \cdot y^2 \end{aligned}$$

Then we use this expression to define a function such that all points of the tiny-jubjub curve are characterized as the function preimages at 0.

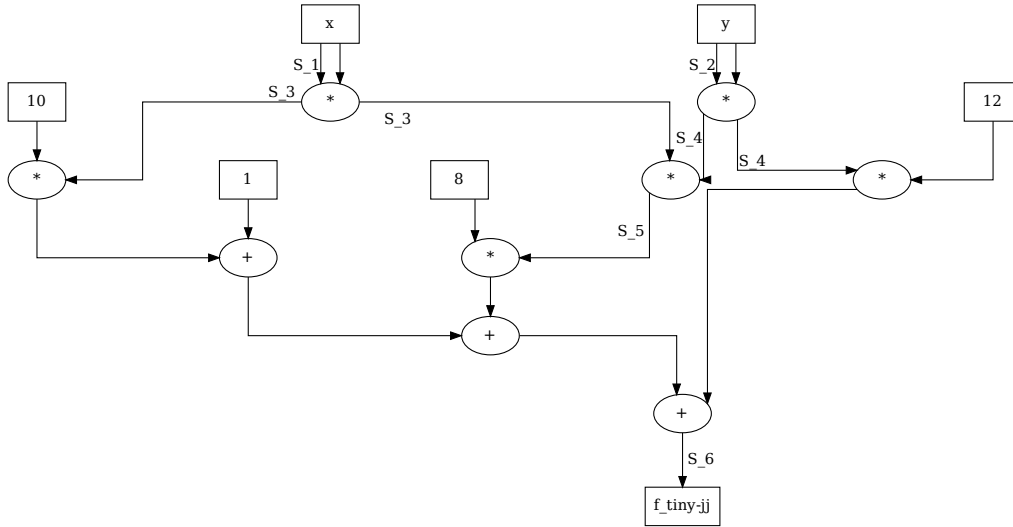
$$f_{\text{tiny-jj}} : \mathbb{F}_{13} \times \mathbb{F}_{13} \rightarrow \mathbb{F}_{13} ; (x, y) \mapsto 1 + 8 \cdot x^2 \cdot y^2 + 10 \cdot x^2 + 12 \cdot y^2$$

Every pair of points  $(x, y) \in \mathbb{F}_{13}^2$  with  $f_{\text{tiny-jj}}(x, y) = 0$  is a point on the tiny-jubjub curve, and there are no other curve points. The preimage  $f_{\text{tiny-jj}}^{-1}(0)$  is therefore a complete description of the tiny-jubjub curve.

We can transform this function into an algebraic circuit over  $\mathbb{F}_{13}$ . We first introduce brackets into potentially ambiguous expressions and then rewrite the function in terms of binary operators. We get the following:

$$\begin{aligned} f_{\text{tiny-jj}}(x, y) &= 1 + 8 \cdot x^2 \cdot y^2 + 10 \cdot x^2 + 12y^2 && \Leftrightarrow \\ &= ((8 \cdot ((x \cdot x) \cdot (y \cdot y))) + (1 + 10 \cdot (x \cdot x))) + (12 \cdot (y \cdot y)) && \Leftrightarrow \\ &= \text{ADD}(\text{ADD}(\text{MUL}(8, \text{MUL}(\text{MUL}(x, x), \text{MUL}(y, y))), \text{ADD}(1, \text{MUL}(10, \text{MUL}(x, x)))), \text{MUL}(12, \text{MUL}(y, y))) \end{aligned}$$

Since we haven't decided which part of the computation should be public and which part should be private, we use the unspecified symbol  $S$  to represent edge labels. Constraining all variables to edge labels  $S_1 = x$ ,  $S_2 = y$  and  $S_6 = f_{\text{tiny-jj}}$ , we get the following circuit, representing the function  $f_{\text{tiny-jj}}$ , by inductively replacing binary operators with their associated arithmetic gates:



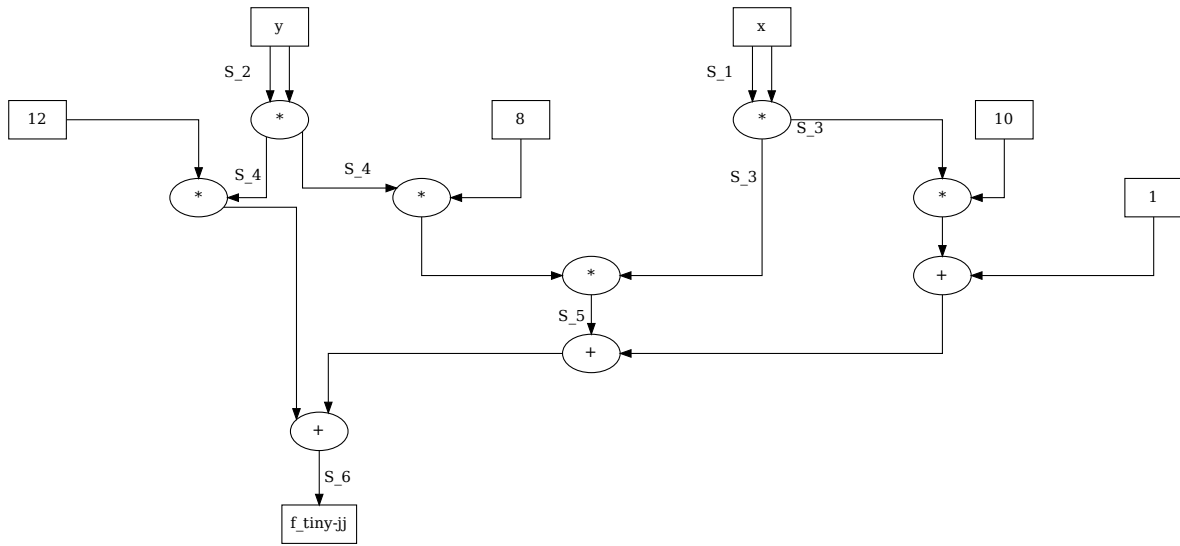
4495

This circuit is not a graph, but a multigraph, since there is more than one edge between some of the nodes.

In the process of designing of circuits from functions, it should be noted that circuit representations are not unique in general. In case of the function  $f_{\text{tiny-jj}}$ , the circuit shape is dependent on our choice of bracketing in XXX. An alternative design is, for example, given by the following circuit, which occurs when the bracketed expression  $8 \cdot ((x \cdot x) \cdot (y \cdot y))$  is replaced by the expression  $(x \cdot x) \cdot (8 \cdot (y \cdot y))$ .

4503

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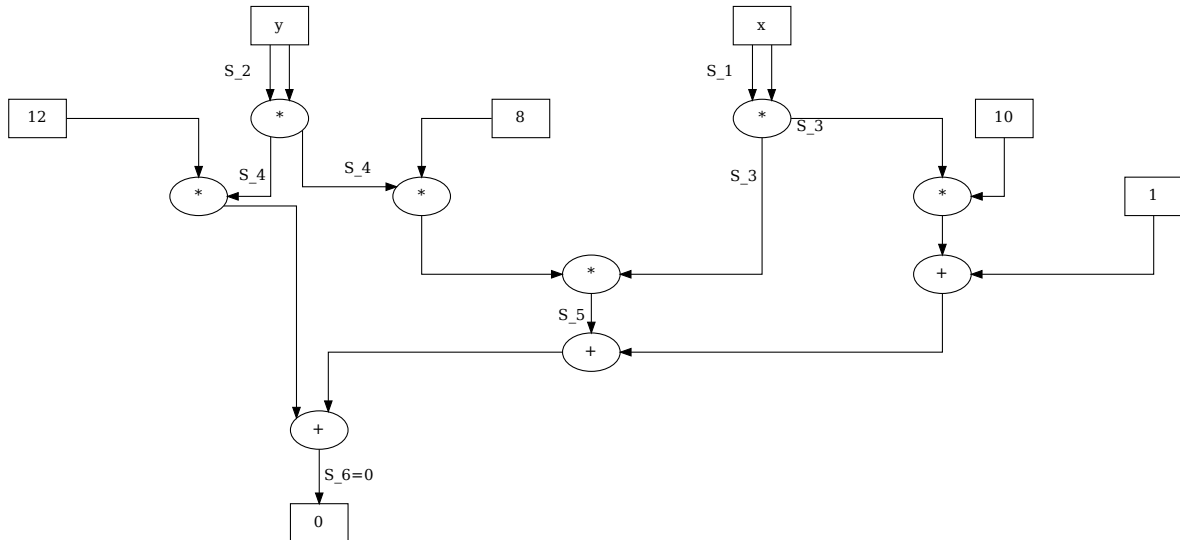
4504

4505

4506 Of course, both circuits represent the same function, due to the associativity and commutativity  
 4507 laws that hold true in any field.

4508 With a circuit that represents the function  $f_{\text{tiny-jj}}$ , we can now proceed to derive a circuit  
 4509 that constrains arbitrary pairs  $(x, y)$  of field elements to be points on the tiny-jubjub curve. To do  
 4510 so, we have to constrain the output to be zero, that is, we have to constrain  $S_6 = 0$ . To indicate  
 4511 this in the circuit, we replace the output variable by the constant 0 and constrain the related edge  
 4512 label accordingly. We get the following:

4513

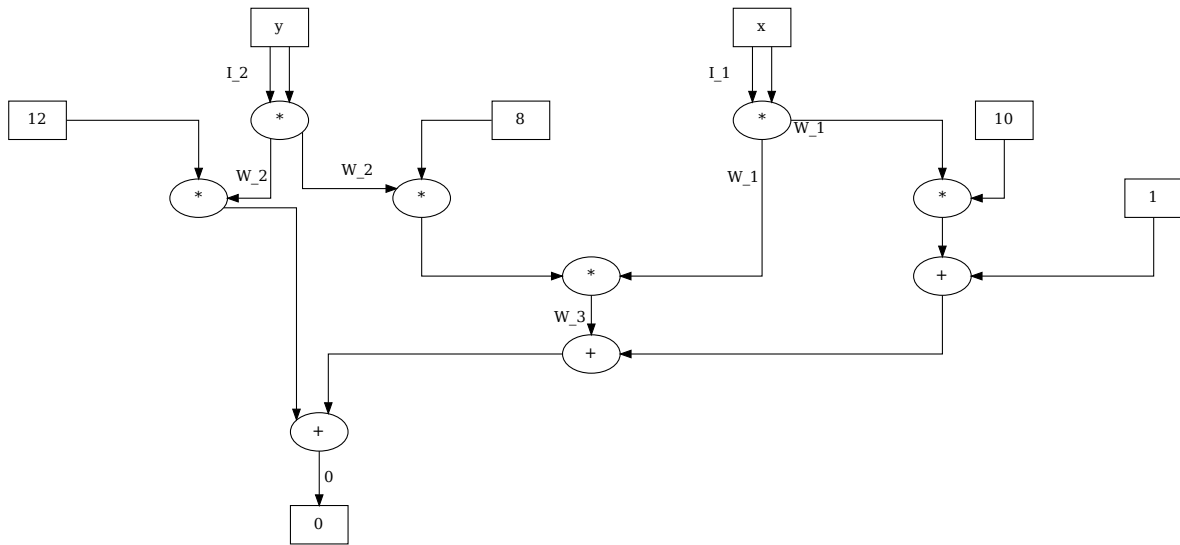


4514

4515

4516 The previous circuit enforces input values assigned to the labels  $S_1$  and  $S_2$  to be points on the  
 4517 tiny-jubjub curve. However, it does not specify which labels are considered public and which  
 4518 are considered private. The following circuit defines the inputs to be public, while all other  
 4519 labels are private:

4520



4521

4522

4523 It can be shown that every space- and time-bounded computation can be transformed into  
 4524 an algebraic circuit. We call any process that transforms a bounded computation into a circuit  
 4525 **flattening**.

4526 **Circuit Execution** Algebraic circuits are directed, acyclic multi-graphs, where nodes repre-  
 4527 sent variables, constants, or addition and multiplication gates. In particular, every algebraic  
 4528 circuit with  $n$  input nodes decorated with variable symbols and  $m$  output nodes decorated with  
 4529 variables can be seen a function that transforms an input tuple  $(x_1, \dots, x_n)$  from  $\mathbb{F}^n$  into an out-  
 4530 put tuple  $(f_1, \dots, f_m)$  from  $\mathbb{F}^m$ . The transformation is done by sending values associated to  
 4531 nodes along their outgoing edges to other nodes. If those nodes are gates, then the values are  
 4532 transformed according to the gate label and the process is repeated along all edges until a sink  
 4533 node is reached. We call this computation **circuit execution**.

4534 When executing a circuit, it is possible to not only compute the output values of the circuit  
 4535 but to derive field elements for all edges, and, in particular, for all edge labels in the circuit. The  
 4536 result is a tuple  $(S_1, S_2, \dots, S_n)$  of field elements associated to all labeled edges, which we call a  
 4537 **valid assignment** to the circuit. In contrast, any assignment  $(S'_1, S'_2, \dots, S'_n)$  of field elements to  
 4538 edge labels that can not arise from circuit execution is called an **invalid assignment**.

4539 Valid assignments can be interpreted as **proofs for proper circuit execution** because they  
 4540 keep a record of the computational result as well as intermediate computational steps.

4541 *Example 117 (3-factorization).* Consider the 3-factorization problem from example 106 and its  
 4542 representation as an algebraic circuit from XXX. We know that the set of edge labels is given  
 4543 by  $S := \{I_1, W_1, W_2, W_3, W_4\}$ .

4544 To understand how this circuit is executed, consider the variables  $x_1 = 2$ ,  $x_2 = 3$  as well as  
 4545  $x_3 = 4$ . Following all edges in the graph, we get the assignments  $W_1 = 2$ ,  $W_2 = 3$  and  $W_3 = 4$ .  
 4546 Then the assignments of  $W_1$  and  $W_2$  enter a multiplication gate and the output of the gate is  
 4547  $2 \cdot 3 = 6$ , which we assign to  $W_4$ , i.e.  $W_4 = 6$ . The values  $W_4$  and  $W_3$  then enter the second  
 4548 multiplication gate and the output of the gate is  $6 \cdot 4 = 11$ , which we assign to  $I_1$ , i.e.  $I_1 = 11$ .

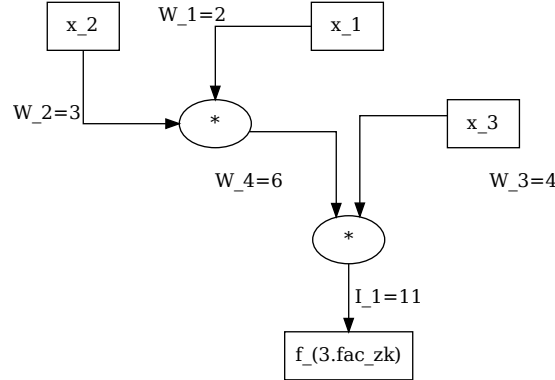
4549 A valid assignment to the 3-factorization circuit  $C_{3, fac}(\mathbb{F}_{13})$  is therefore given by the fol-

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4550 lowing set:

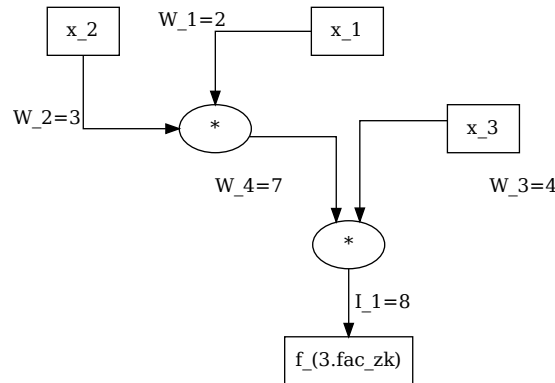
$$S_{\text{valid}} := \{11; 2, 3, 4, 6\} \quad (6.8)$$

4551 We can visualise this assignment in the circuit as follows:



4552

4553 To see what an invalid assignment looks like, consider the assignment  $S_{\text{err}} := \{8; 2, 3, 4, 7\}$ . In  
 4554 this assignment, the input values are the same as in the previous case. The associated circuit is:



4555

4556 This assignment is invalid, as the assignments of  $I_1$  and  $W_4$  cannot be obtained by executing the  
 4557 circuit.

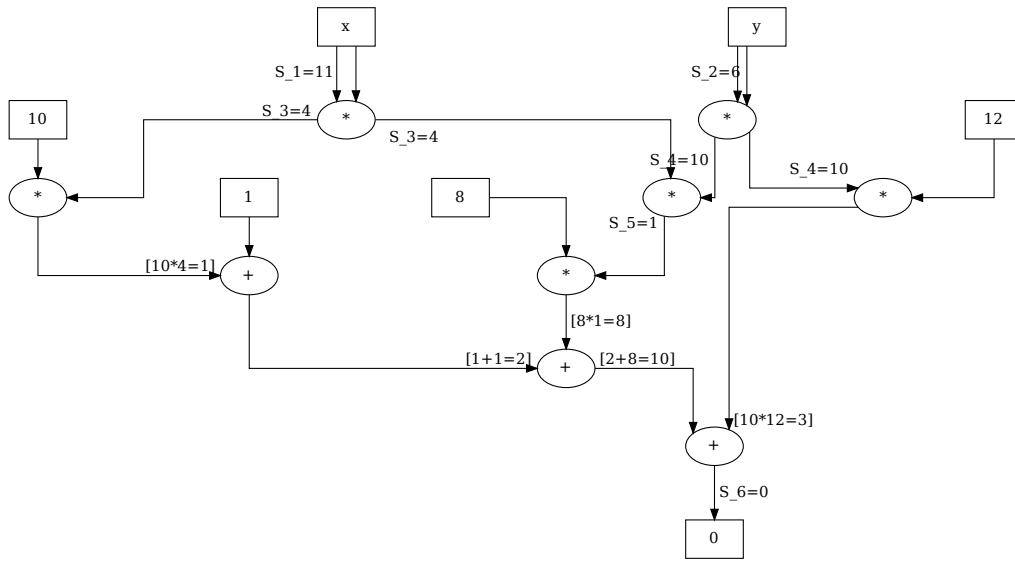
4558 *Example 118.* To compute a more realistic algebraic circuit execution, consider the defining  
 4559 circuit  $C_{\text{tiny-jj}}(\mathbb{F}_{13})$  from example 114 again. We already know from the way this circuit is  
 4560 constructed that any valid assignment with  $S_1 = x$ ,  $S_2 = y$  and  $S_6 = 0$  will ensure that the pair  
 4561  $(x, y)$  is a point on the tiny-jubjub curve in its Edwards representation (equation 5.20).

4562 From example 114, we know that the pair  $(11, 6)$  is a proper point on the tiny-jubjub curve  
 4563 and we use this point as input to a circuit execution. We get the following:

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4564

Executing the circuit, we indeed compute  $S_6 = 0$  as expected, which proves that  $(11, 6)$  is a point on the tiny-jubjub curve in its Edwards representation. A valid assignment of  $C_{\text{tiny-jj}}(\mathbb{F}_{13})$  is therefore given by the following equation:

$$S_{\text{tiny-jj}} = \{S_1, S_2, S_3, S_4, S_5, S_6\} = \{11, 6, 4, 10, 1, 0\}$$

4565 **Circuit Satisfiability** To understand how algebraic circuits give rise to formal languages, ob-  
 4566 serve that every algebraic circuit  $C(\mathbb{F})$  over a fields  $\mathbb{F}$  defines a decision function over the al-  
 4567 phabet  $\Sigma_I \times \Sigma_W = \mathbb{F} \times \mathbb{F}$  in the following way:

$$R_{C(\mathbb{F})} : \mathbb{F}^* \times \mathbb{F}^* \rightarrow \{true, false\} ; (I; W) \mapsto \begin{cases} true & (I; W) \text{ is valid assignment to } C(\mathbb{F}) \\ false & \text{else} \end{cases} \quad (6.9)$$

4568 Every algebraic circuit therefore defines a formal language. The grammar of this language is  
 4569 encoded in the shape of the circuit, words are assignments to edge labels that are derived from  
 4570 circuit execution, and **statements** are knowledge claims “Given instance  $I$ , there is a witness  
 4571  $W$  such that  $(I; W)$  is a valid assignment to the circuit”. A constructive proof to this claim  
 4572 is therefore an assignment of a field element to every witness variable, which is verified by  
 4573 executing the circuit to see if the assignment of the execution meets the assignment of the  
 4574 proof.

4575 In the context of zero-knowledge proof systems, executing circuits is also often called **wit-**  
 4576 **ness generation**, since in applications the instance part is usually public, while its the task of a  
 4577 prover to compute the witness part.

*Remark 3 (Circuit satisfiability).* It should be noted that, in our definition, every circuit defines its own language. However, in more theoretical approaches another language usually called **circuit satisfiability** is often considered, which is useful when it comes to more abstract problems like expressiveness, or computational complexity of the class of **all** algebraic circuits over a given field. From our perspective the circuit satisfiability language is obtained by union of all circuit languages that are in our definition. To be more precise, let the alphabet  $\Sigma = \mathbb{F}$  be a field. Then

Should we refer to RICS satisfiability (p. 139 here?)

$$L_{\text{CIRCUIT\_SAT}(\mathbb{F})} = \{(i; w) \in \Sigma^* \times \Sigma^* \mid \text{there is a circuit } C(\mathbb{F}) \text{ such that } (i; w) \text{ is valid assignment}\}$$



4578 *Example 119 (3-Factorization).* Consider the circuit  $C_{3, fac}$  from equation 6.8 again. We call the  
 4579 associated language  $L_{3, fac\_circ}$ .

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4580 To understand how a constructive proof of a statement in  $L_{3, fac\_circ}$  looks like, consider the  
 4581 instance  $I_1 = 11$ . To provide a proof for the statement “There exist a witness  $W$  such that  $(I_1; W)$   
 4582 is a word in  $L_{3, fac\_circ}$ ” a proof therefore has to consists of proper values for the variables  $W_1$ ,  
 4583  $W_2$ ,  $W_3$  and  $W_4$ . Any proofer therefore has to find input values for  $W_1$ ,  $W_2$  and  $W_3$  and then  
 4584 execute the circuit to compute  $W_4$  under the assumption  $I_1 = 11$ .

4585 Example XXX implies that  $(2, 3, 4, 6)$  is a proper constructive proof and in order to verify  
 4586 the proof a verifier needs to execute the circuit with instance  $I_1 = 11$  and inputs  $W_1 = 2$ ,  $W_2 = 3$   
 4587 and  $W_3 = 4$  to decide whether the proof is a valid assignment or not.

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4588 **Associated Constraint Systems** As we have seen in 6.2.1.1, rank-1 constraint systems define  
 4589 a way to represent statements in terms of a system of quadratic equations over finite fields,  
 4590 suitable for pairing-based zero-knowledge proof systems. However, those equations provide no  
 4591 practical way for a proofer to actually compute a solution. On the other hand, algebraic circuits  
 4592 can be executed in order to derive valid assignments efficiently.

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4593 In this paragraph, we show how to transform any algebraic circuit into a rank-1 constraint  
 4594 system such that valid circuit assignments are in 1:1 correspondence with solutions to the asso-  
 4595 ciated R1CS.

4596 To see this, let  $C(\mathbb{F})$  be an algebraic circuit over a finite field  $\mathbb{F}$ , with a set of edge labels  
 4597  $\{S_1, S_2, \dots, S_n\}$ . Then one of the following steps is executed for every edge label  $S_j$  from that  
 4598 set:

4599 *Definition 6.2.2.2.* • If the edge label  $S_j$  is an outgoing edge of a multiplication gate, the  
 4600 R1CS gets a new quadratic constraint

$$(\text{left input}) \cdot (\text{right input}) = S_j \quad (6.10)$$

4601 where (left input) respectively (right input) is the output from the symbolic execution of  
 4602 the subgraph that consists of the left respectively right input edge of this gate, and all  
 4603 edges and nodes that have this edge in their path, starting with constant inputs or labeled  
 4604 outgoing edges of other nodes.

4605 • If the edge label  $S_j$  is an outgoing edge of an addition gate, the R1CS gets a new quadratic  
 4606 constraint

$$(\text{left input} + \text{right input}) \cdot 1 = S_j \quad (6.11)$$

4607 where (left input) respectively (right input) is the output from the symbolic execution of  
 4608 the subgraph that consists of the left respectively right input edge of this gate and all  
 4609 edges and nodes that have this edge in their path, starting with constant inputs or labeled  
 4610 outgoing edges of other nodes.

4611 • No other edge label adds a constraint to the system.

4612 The result of this method is a rank-1 constraint system, and, in this sense, every algebraic  
 4613 circuit  $C(\mathbb{F})$  generates a R1CS  $R$ , which we call the **associated R1CS** of the circuit. It can be  
 4614 shown that a tuple of field elements  $(S_1, S_2, \dots, S_n)$  is a valid assignment to a circuit if and only  
 4615 if the same tuple is a solution to the associated R1CS. Circuit executions therefore compute  
 4616 solutions to rank-1 constraints systems efficiently.

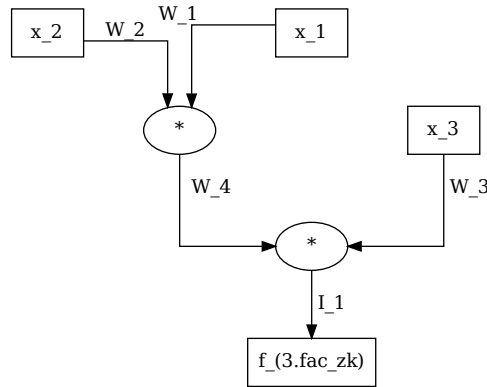
4617 To understand the contribution of algebraic gates to the number of constraints, note that, by  
 4618 definition, multiplication gates have labels on their outgoing edges if and only if there is at least

one labeled edge in both input paths, or if the outgoing edge is an input to a sink node. This implies that multiplication with a constant is essentially free in the sense that it doesn't add a new constraint to the system, as long as that multiplication gate is not an input to an output node.

Moreover, addition gates have labels on their outgoing edges if and only if they are inputs to sink nodes. This implies that addition is essentially free in the sense that it doesn't add a new constraint to the system, as long as that addition gate is not an input to an output node.

*Example 120 (3-factorization).* Consider our 3-factorization problem from equation 6.8 and the associated circuit  $C_{3.fac}(\mathbb{F}_{13})$ . Our task is to transform this circuit into an equivalent rank-1 constraint system.

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We start with an empty R1CS, and, in order to generate all constraints, we have to iterate over the set of edge labels  $\{I_1; W_1, W_2, W_3, W_4\}$ .

Starting with the edge label  $I_1$ , we see that it is an outgoing edge of a multiplication gate, and, since both input edges are labeled, we have to add the following constraint to the system:

$$\begin{aligned} (\text{left input}) \cdot (\text{right input}) &= I_1 \\ W_4 \cdot W_3 &= I_1 \end{aligned} \quad \Leftrightarrow$$

Next, we consider the edge label  $W_1$  and, since, it's not an outgoing edge of a multiplication or addition label, we don't add a constraint to the system. The same holds true for the labels  $W_2$  and  $W_3$ .

For edge label  $W_4$ , we see that it is an outgoing edge of a multiplication gate, and, since both input edges are labeled, we have to add the following constraint to the system:

$$\begin{aligned} (\text{left input}) \cdot (\text{right input}) &= W_4 \\ W_2 \cdot W_1 &= W_4 \end{aligned} \quad \Leftrightarrow$$

Since there are no more labeled edges, all constraints are generated, and we have to combine them to get the associated R1CS of  $C_{3.fac}(\mathbb{F}_{13})$ :

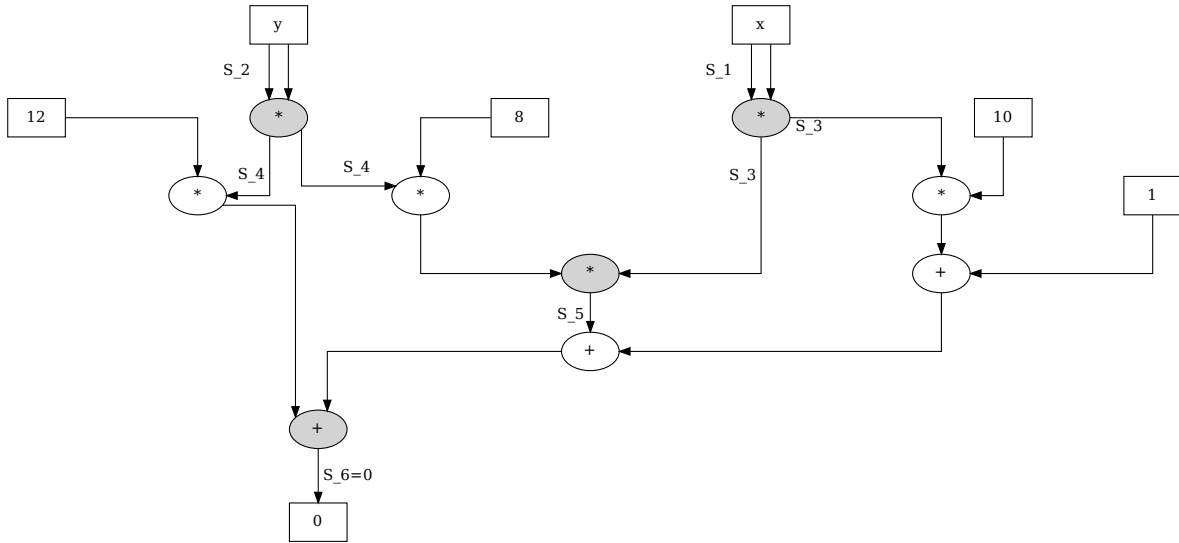
$$\begin{aligned} W_4 \cdot W_3 &= I_1 \\ W_2 \cdot W_1 &= W_4 \end{aligned}$$

This system is equivalent to the R1CS we derived in example 111. The languages  $L_{3.fac\_zk}$  and  $L_{3.fac\_circ}$  are therefore equivalent and both the circuit as well as the R1CS are just two different ways of expressing the same language.

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*Example 121.* To consider a more general transformation, we consider the tiny-jubjub circuit from example 114 again. A proper circuit is given by

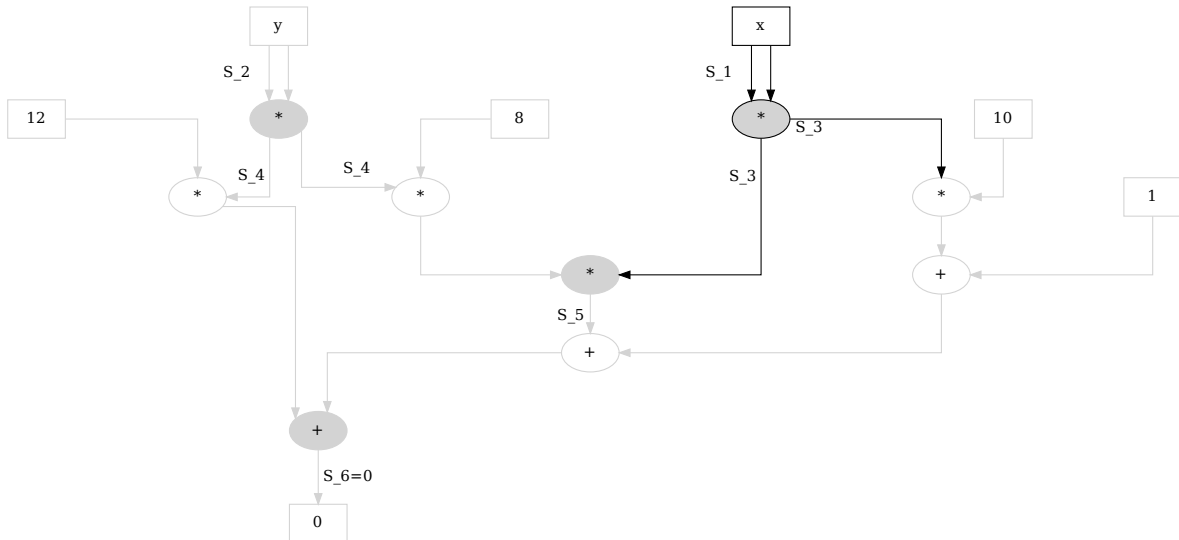
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To compute the number of constraints, observe that we have 3 multiplication gates that have labels on their outgoing edges and 1 addition gate that has a label on its outgoing edge. We therefore have to compute 4 quadratic constraints.

In order to derive the associated R1CS, we have start with an empty R1CS and then iterate over the set  $\{S_1, S_2, S_3, S_4, S_5, S_6 = 0\}$  of all edge labels, in order to generate the constraints.

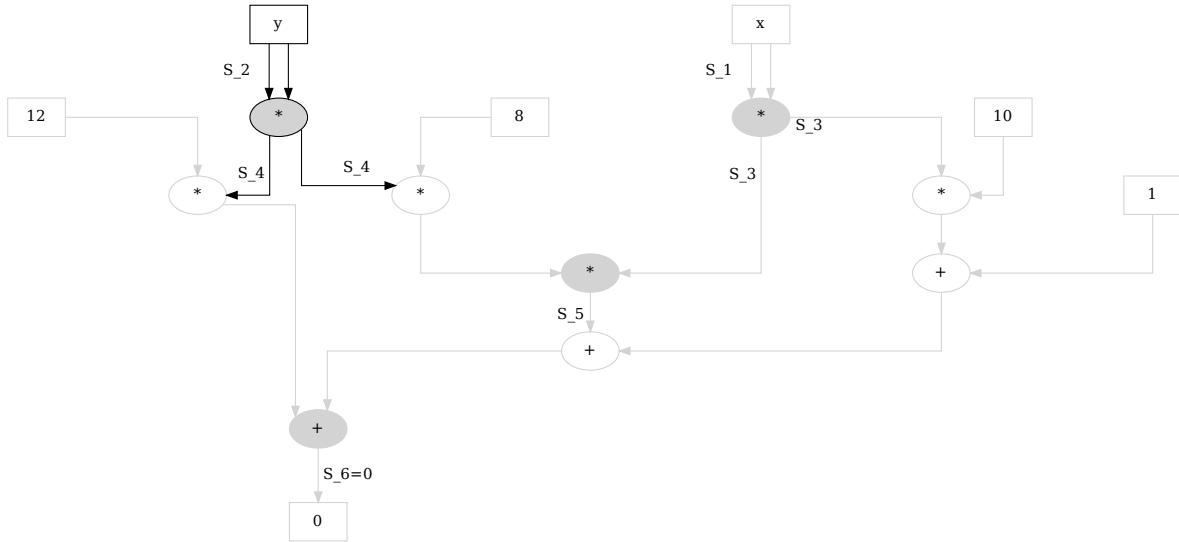
Considering edge label  $S_1$ , we see that the associated edges are not outgoing edges of any algebraic gate, and we therefore have to add no new constraint to the system. The same holds true for edge label  $S_2$ . Looking at edge label  $S_3$ , we see that the associated edges are outgoing edges of a multiplication gate and that the associated subgraph is given by:



Both the left and the right input to this multiplication gate are labeled by  $S_1$ . We therefore have to add the following constraint to the system:

$$S_1 \cdot S_1 = S_3$$

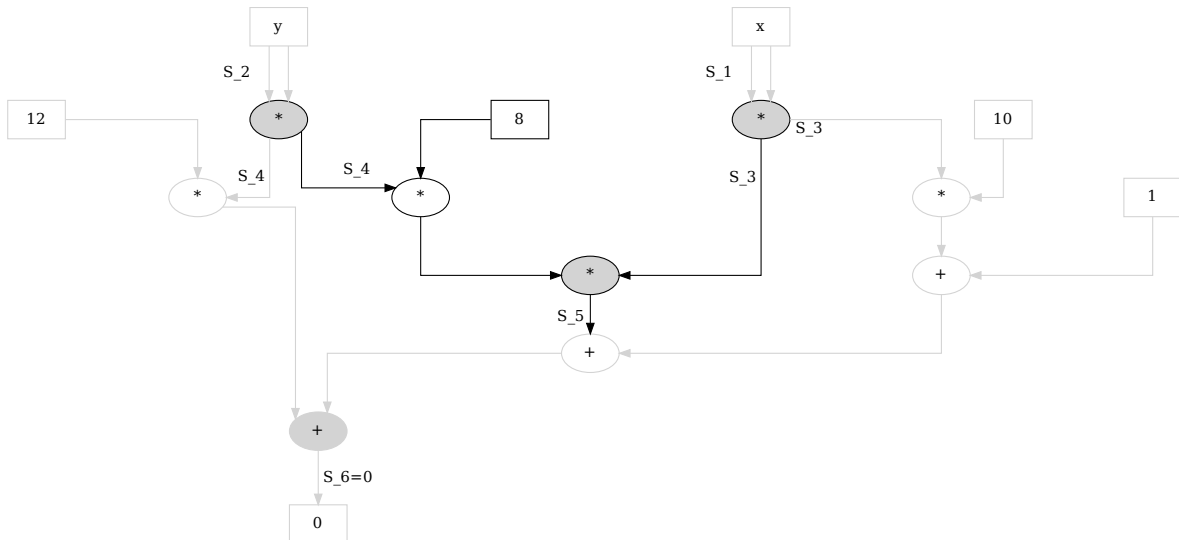
Looking at edge label  $S_4$ , we see that the associated edges are outgoing edges of a multiplication gate and that the associated subgraph is given by:



Both the left and the right input to this multiplication gate are labeled by  $S_2$  and we therefore have to add the following constraint to the system:

$$S_2 \cdot S_2 = S_4$$

Edge label  $S_5$  is more interesting. To see if it implies a constraint, we have to construct the associated subgraph first, which consists of all edges and all nodes in all path starting either at a constant input or a labeled edge. We get



4665

The right input to the associated multiplication gate is given by the labeled edge  $S_3$ . However, the left input is not a labeled edge, but has a labeled edge in one of its path. This implies that we have to add a constraint to the system. To compute the left factor of that constraint, we have to compute the output of subgraph associated to the left edge, which is  $8 \cdot W_2$ . This gives the constraint

$$(S_4 \cdot 8) \cdot S_3 = S_5$$

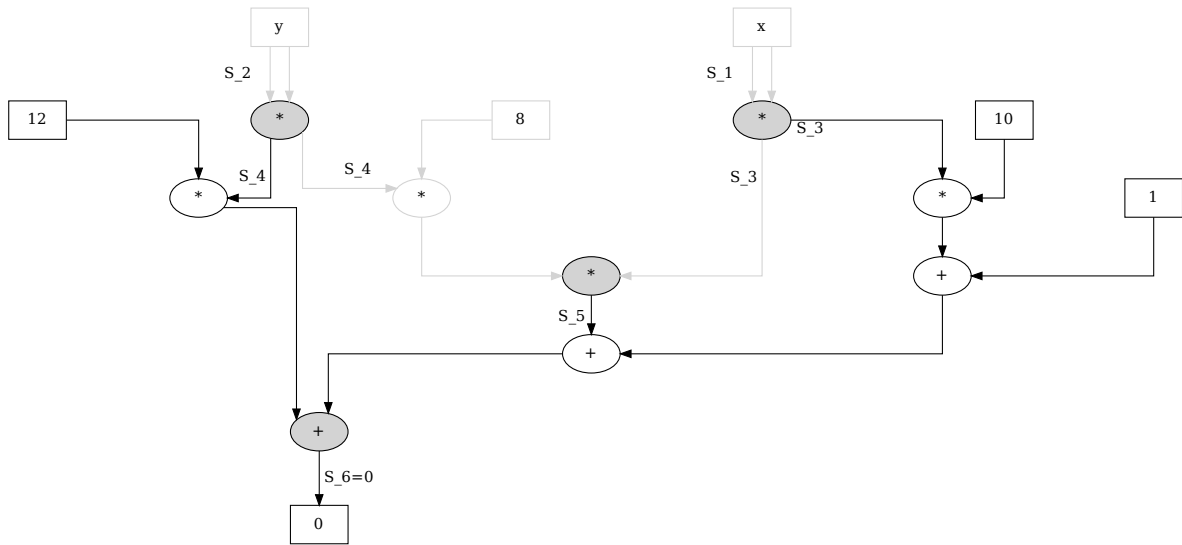
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The last edge label is the constant  $S_6 = 0$ . To see if it implies a constraint, we have to construct the associated subgraph, which consists of all edges and all nodes in all path starting either at a constant input or a labeled edge. We get

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Both the left and the right input are unlabeled, but have a labeled edges in their path. This implies that we have to add a constraint to the system. Since the gate is an addition gate, the right factor in the quadratic constraint is always 1 and the left factor is computed by symbolically executing all inputs to all gates in sub-circuit. We get

$$(12 \cdot S_4 + S_5 + 10 \cdot S_3 + 1) \cdot 1 = 0$$

Since there are no more labeled outgoing edges, we are done deriving the constraints. Combining all constraints together, we get the following R1CS:

$$S_1 \cdot S_1 = S_3$$

$$S_2 \cdot S_2 = S_4$$

$$(S_4 \cdot 8) \cdot S_3 = S_5$$

$$(12 \cdot S_4 + S_5 + 10 \cdot S_3 + 1) \cdot 1 = 0$$

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which is equivalent to the R1CS we derived in example 114. The languages  $L_{3, fac\_zk}$  and  $L_{3, fac\_circ}$  are therefore equivalent and both the circuit as well as the R1CS are just two different ways to express the same language.

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### 6.2.3 Quadratic Arithmetic Programs

We have introduced algebraic circuits and their associated rank-1 constraints systems as two particular models able to represent space- and time-bounded computation. Both models define formal languages, and associated membership as well as knowledge claims can be constructively proved by executing the circuit in order to compute solutions to its associated R1CS.

One reason why those systems are useful in the context of succinct zero-knowledge proof systems is because any R1CS can be transformed into another computational model called **quadratic arithmetic programs** (QAP), which serve as the basis for some of the most efficient succinct non-interactive zero-knowledge proof generators that currently exist.

As we will see, proving statements for languages that have checking relations defined by quadratic arithmetic programs can be achieved by providing certain polynomials, and those proofs can be verified by checking a particular divisibility property.

**QAP representation** To understand what quadratic arithmetic programs are in detail, let  $\mathbb{F}$  be a field and  $R$  a rank-1 constraints system over  $\mathbb{F}$  such that the number of non-zero elements in  $\mathbb{F}$  is strictly larger than the number  $k$  of constraints in  $R$ . Moreover, let  $a_j^i, b_j^i$  and  $c_j^i \in \mathbb{F}$  for every index  $0 \leq j \leq n + m$  and  $1 \leq i \leq k$ , be the defining constants of the R1CS and  $m_1, \dots, m_k$  be arbitrary, invertible and distinct elements from  $\mathbb{F}$ .

Then a **quadratic arithmetic program** [QAP] of the R1CS is the following set of polynomials over  $\mathbb{F}$ :

$$QAP(R) = \left\{ T \in \mathbb{F}[x], \{A_j, B_j, C_j \in \mathbb{F}[x]\}_{h=0}^{n+m} \right\} \quad (6.12)$$

In the equation above,  $T(x) := \prod_{l=1}^k (x - m_l)$  is a polynomial of degree  $k$ , called the **target polynomial** of the QAP and  $A_j, B_j$  as well as  $C_j$  are the unique degree  $k - 1$  polynomials defined by the following equation:

$$A_j(m_i) = a_j^i \quad B_j(m_i) = b_j^i \quad C_j(m_i) = c_j^i \quad j = 1, \dots, n + m + 1, i = 1, \dots, k \quad (6.13)$$

Given some rank-1 constraint system, an associated quadratic arithmetic program is therefore nothing but a set of polynomials, computed from the constants in the R1CS. To see that the polynomials  $A_j, B_j$  and  $C_j$  are uniquely defined by the equations in XXX, recall that a polynomial of degree  $k - 1$  is completely determined on  $k$  evaluation points and the equation 4 precisely determines those  $k$  evaluation points.

Since we only consider polynomials over fields, Lagrange's interpolation method from 3.30 in chapter 3 can be used to derive the polynomials  $A_j, B_j$  and  $C_j$  from their defining equations XXX. A practical method to compute a QAP from a given R1CS therefore consists of two steps. If the R1CS consists of  $k$  constraints, first choose  $k$  invertible and mutually different points from the underlying field. Every choice defines a different QAP for the same R1CS. Then use Lagrange's method and equation XXX to compute the polynomials  $A_j, B_j$  and  $C_j$  for every  $1 \leq j \leq k$ .

*Example 122* (Generalized factorization SNARK). To provide a better intuition of quadratic arithmetic programs and how they are computed from their associated rank-1 constraint systems, consider the language  $L_{3, fac\_zk}$  from example 106 and its associated R1CS:

$$\begin{array}{ll} W_1 \cdot W_2 = W_4 & \text{constraint 1} \\ W_4 \cdot W_3 = I_1 & \text{constraint 2} \end{array}$$

In this example we want to transform this R1CS into an associated QAP. In a first step, we have to compute the defining constants  $a_j^i, b_j^i$  and  $c_j^i$  of the R1CS. According to XXX, we have

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$$\begin{array}{cccccc} a_0^1 = 0 & a_1^1 = 0 & a_2^1 = 1 & a_3^1 = 0 & a_4^1 = 0 & a_5^1 = 0 \\ a_0^2 = 0 & a_1^2 = 0 & a_2^2 = 0 & a_3^2 = 0 & a_4^2 = 0 & a_5^2 = 1 \end{array}$$

$$\begin{array}{cccccc} b_0^1 = 0 & b_1^1 = 0 & b_2^1 = 0 & b_3^1 = 1 & b_4^1 = 0 & b_5^1 = 0 \\ b_0^2 = 0 & b_1^2 = 0 & b_2^2 = 0 & b_3^2 = 0 & b_4^2 = 1 & b_5^2 = 0 \end{array}$$

$$\begin{array}{cccccc} c_0^1 = 0 & c_1^1 = 0 & c_2^1 = 0 & c_3^1 = 0 & c_4^1 = 0 & c_5^1 = 1 \\ c_0^2 = 0 & c_1^2 = 1 & c_2^2 = 0 & c_3^2 = 0 & c_4^2 = 0 & c_5^2 = 0 \end{array}$$

Since the R1CS is defined over the field  $\mathbb{F}_{13}$  and has two constraining equations, we need to choose two arbitrary but distinct elements  $m_1$  and  $m_2$  from  $\mathbb{F}_{13}$ . We choose  $m_1 = 5$ , and  $m_2 = 7$  and with this choice we get the target polynomial

$$\begin{aligned} T(x) &= (x - m_1)(x - m_2) && \# \text{ Definition of } T \\ &= (x - 5)(x - 7) && \# \text{ Insert our choice} \\ &= (x + 8)(x + 6) && \# \text{ Negatives in } \mathbb{F}_{13} \\ &= x^2 + x + 9 && \# \text{ expand} \end{aligned}$$

4709 Then we have to compute the polynomials  $A_j$ ,  $B_j$  and  $C_j$  by their defining equation from the  
4710 R1CS coefficients. Since the R1CS has two constraining equations, those polynomials are of  
4711 degree 1 and they are defined by their evaluation at the point  $m_1 = 5$  and the point  $m_2 = 7$ .

At point  $m_1$ , each polynomial  $A_j$  is defined to be  $a_j^1$  and at point  $m_2$ , each polynomial  $A_j$  is defined to be  $a_j^2$ . The same holds true for the polynomials  $B_j$  as well as  $C_j$ . Writing all these equations now, we get:

$$\begin{array}{cccccc} A_0(5) = 0, & A_1(5) = 0, & A_2(5) = 1, & A_3(5) = 0, & A_4(5) = 0, & A_5(5) = 0 \\ A_0(7) = 0, & A_1(7) = 0, & A_2(7) = 0, & A_3(7) = 0, & A_4(7) = 0, & A_5(7) = 1 \end{array}$$

$$\begin{array}{cccccc} B_0(5) = 0, & B_1(5) = 0, & B_2(5) = 0, & B_3(5) = 1, & B_4(5) = 0, & B_5(5) = 0 \\ B_0(7) = 0, & B_1(7) = 0, & B_2(7) = 0, & B_3(7) = 0, & B_4(7) = 1, & B_5(7) = 0 \end{array}$$

$$\begin{array}{cccccc} C_0(5) = 0, & C_1(5) = 0, & C_2(5) = 0, & C_3(5) = 0, & C_4(5) = 0, & C_5(5) = 1 \\ C_0(7) = 0, & C_1(7) = 1, & C_2(7) = 0, & C_3(7) = 0, & C_4(7) = 0, & C_5(7) = 0 \end{array}$$

4712 Lagrange's interpolation implies that a polynomial of degree  $k$ , that is, that zero on  $k + 1$  points  
4713 has to be the zero polynomial. Since our polynomials are of degree 1 and determined on 2  
4714 points, we therefore know that the only non-zero polynomials in our QAP are  $A_2$ ,  $A_5$ ,  $B_3$ ,  $B_4$ ,  
4715  $C_1$  and  $C_5$ , and that we can use Lagrange's interpolation to compute them.

To compute  $A_2$  we note that the set  $S$  in our version of Lagrange's method is given by  $S = \{(x_0, y_0), (x_1, y_1)\} = \{(5, 1), (7, 0)\}$ . Using this set we get:

$$\begin{aligned} A_2(x) &= y_0 \cdot l_0 + y_1 \cdot l_1 \\ &= y_0 \cdot \left( \frac{x - x_1}{x_0 - x_1} \right) + y_1 \cdot \left( \frac{x - x_0}{x_1 - x_0} \right) = 1 \cdot \left( \frac{x - 7}{5 - 7} \right) + 0 \cdot \left( \frac{x - 5}{7 - 5} \right) \\ &= \frac{x - 7}{-2} = \frac{x - 7}{11} && \# 11^{-1} = 6 \\ &= 6(x - 7) = 6x + 10 && \# -7 = 6 \text{ and } 6 \cdot 6 = 10 \end{aligned}$$

To compute  $A_5$ , we note that the set  $S$  in our version of Lagrange's method is given by  $S = \{(x_0, y_0), (x_1, y_1)\} = \{(5, 0), (7, 1)\}$ . Using this set we get:

$$\begin{aligned}
 A_5(x) &= y_0 \cdot l_0 + y_1 \cdot l_1 \\
 &= y_0 \cdot \left(\frac{x - x_1}{x_0 - x_1}\right) + y_1 \cdot \left(\frac{x - x_0}{x_1 - x_0}\right) = 0 \cdot \left(\frac{x - 7}{5 - 7}\right) + 1 \cdot \left(\frac{x - 5}{7 - 5}\right) \\
 &= \frac{x - 5}{2} \quad \# 2^{-1} = 7 \\
 &= 7(x - 5) = 7x + 4 \quad \# -5 = 8 \text{ and } 7 \cdot 8 = 4
 \end{aligned}$$

4716 Using Lagrange's interpolation, we can deduce that  $A_2 = B_3 = C_5$  as well as  $A_5 = B_4 = C_1$ ,  
 4717 since they are polynomials of degree 1 that evaluate to same values on 2 points. Using this, we  
 4718 get the following set of polynomials

$A_0(x) = 0$	$B_0(x) = 0$	$C_0(x) = 0$
$A_1(x) = 0$	$B_1(x) = 0$	$C_1(x) = 7x + 4$
$A_2(x) = 6x + 10$	$B_2(x) = 0$	$C_2(x) = 0$
$A_3(x) = 0$	$B_3(x) = 6x + 10$	$C_3(x) = 0$
$A_4(x) = 0$	$B_4(x) = 7x + 4$	$C_4(x) = 0$
$A_5(x) = 7x + 4$	$B_5(x) = 0$	$C_5(x) = 6x + 10$

4720 We can invoke Sage to verify our computation. In sage every polynomial ring has a function  
 4721 `lagrange_polynomial` that takes the defining points as inputs and the associated Lagrange  
 4722 polynomial as output.

```

4723 sage: F13 = GF(13)                                     628
4724 sage: F13t.<t> = F13[]                                   629
4725 sage: T = F13t((t-5)*(t-7))                             630
4726 sage: A2 = F13t.lagrange_polynomial([(5, 1), (7, 0)])    631
4727 sage: A5 = F13t.lagrange_polynomial([(5, 0), (7, 1)])    632
4728 sage: T == F13t(t^2 + t + 9)                             633
4729 True                                                       634
4730 sage: A2 == F13t(6*t + 10)                                635
4731 True                                                       636
4732 sage: A5 == F13t(7*t + 4)                                 637
4733 True                                                       638

```

4734 Combining this computation with the target polynomial we derived earlier, a quadratic arith-  
 4735 metic program associated to the rank-1 constraint system  $R_{3.fac\_zk}$  is given as follows:

$$\begin{aligned}
 QAP(R_{3.fac\_zk}) &= \{x^2 + x + 9, \\
 \{0, 0, 6x + 10, 0, 0, 7x + 4\}, \{0, 0, 0, 6x + 10, 7x + 4, 0\}, \{0, 7x + 4, 0, 0, 0, 6x + 10\}\}
 \end{aligned} \quad (6.14)$$

4736 **QAP Satisfiability** One of the major points of quadratic arithmetic programs in proofing sys-  
 4737 tems is that solutions of their associated rank-1 constraints systems are in 1:1 correspondence  
 4738 with certain polynomials  $P$  such that  $P$  is divisible by the target polynomial  $T$  of the QAP if and  
 4739 only if the solution is a solution. Verifying solutions to the R1CS and hence, checking proper  
 4740 circuit execution is then achievable by polynomial division of  $P$  by  $T$ .

4741 To be more specific, let  $R$  be some rank-1 constraints system with associated assignment  
 4742 variables  $(I_1, \dots, I_n; W_1, \dots, W_m)$  and let  $QAP(R)$  be a quadratic arithmetic program of  $R$ . Then

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the tuple  $(I_1, \dots, I_n; W_1, \dots, W_m)$  is a solution to the R1CS if and only if the following polynomial is divisible by the target polynomial  $T$ :

$$P_{(I;W)} = (A_0 + \sum_j^n I_j \cdot A_j + \sum_j^m W_j \cdot A_{n+j}) \cdot (B_0 + \sum_j^n I_j \cdot B_j + \sum_j^m W_j \cdot B_{n+j}) - (C_0 + \sum_j^n I_j \cdot C_j + \sum_j^m W_j \cdot C_{n+j}) \quad (6.15)$$

Every tuple  $(I; W)$  defines a polynomial  $P_{(I;W)}$ , and, since each polynomial  $A_j$ ,  $B_j$  and  $C_j$  is of degree  $k-1$ ,  $P_{(I;W)}$  is of degree  $(k-1) \cdot (k-1) = k^2 - 2k + 1$ .

To understand how quadratic arithmetic programs define formal languages, observe that every QAP over a field  $\mathbb{F}$  defines a decision function over the alphabet  $\Sigma_I \times \Sigma_W = \mathbb{F} \times \mathbb{F}$  in the following way:

$$R_{QAP} : \mathbb{F}^* \times \mathbb{F}^* \rightarrow \{true, false\} ; (I; W) \mapsto \begin{cases} true & P_{(I;W)} \text{ is divisible by } T \\ false & \text{else} \end{cases} \quad (6.16)$$

This means that every QAP defines a formal language, and, if the QAP is associated to an R1CS, it can be shown that the two languages are equivalent. A **statement** is a membership claim “There is a word  $(I; W)$  in  $L_{QAP}$ ”. A proof to this claim is therefore a polynomial  $P_{(I;W)}$ , which is verified by dividing  $P_{(I;W)}$  by  $T$ .

Note the structural similarity to the definition of an R1CS in 6.2.1.1 and the different ways of computing proofs in both systems. For circuits and their associated rank-1 constraints systems, a constructive proof consists of a valid assignment of field elements to the edges of the circuit, or the variables in the R1CS. However, in the case of QAPs, a valid proof consists of a polynomial  $P_{(I;W)}$ .

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To compute a proof for a statement in  $L_{QAP}$  given some instance  $I$ , a proofer first needs to compute a constructive proof  $W$ , e.g. by executing the circuit. With  $(I; W)$  at hand, the proofer can then compute the polynomial  $P_{(I;W)}$  and publish it as proof.

Verifying a constructive proof in the case of a circuit is achieved by executing the circuit, comparing the result to the given proof, and verifying the same proof in the R1CS picture means checking if the elements of the proof satisfy all equation.

In contrast, verifying a proof in the case of a QAP is done by polynomial division of the proof  $P$  by the target polynomial  $T$  of the QAP. The proof checks out if and only if  $P$  is divisible by  $T$ .

*Example 123.* Consider the quadratic arithmetic program  $QAP(R_{3, fac\_zk})$  from example XXX, and its associated R1CS from equation 6.1.0.1. To give an intuition of how proofs in the language  $L_{QAP(R_{3, fac\_zk})}$  lets consider the instance  $I_1 = 11$ . As we know from example XXX,  $(W_1, W_2, W_3, W_5) = (2, 3, 4, 6)$  is a proper witness, since  $(I_1; W_1, W_2, W_3, W_5) = (11; 2, 3, 4, 6)$  is a valid circuit assignment and hence, a solution to  $R_{3, fac\_zk}$  and a constructive proof for language  $L_{R_{3, fac\_zk}}$ .

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In order to transform this constructive proof into a membership proof in language  $L_{QAP(R_{3, fac\_zk})}$  a proofer has to use the elements of the constructive proof, to compute the polynomial  $P_{(I;W)}$ .

In the case of  $(I_1; W_1, W_2, W_3, W_5) = (11; 2, 3, 4, 6)$ , the associated proof is computed as fol-

lows:

$$\begin{aligned}
P_{(I;W)} &= (A_0 + \sum_j^n I_j \cdot A_j + \sum_j^m W_j \cdot A_{n+j}) \cdot (B_0 + \sum_j^n I_j \cdot B_j + \sum_j^m W_j \cdot B_{n+j}) - (C_0 + \sum_j^n I_j \cdot C_j + \sum_j^m W_j \cdot C_{n+j}) \\
&= (2(6x + 10) + 6(7x + 4)) \cdot (3(6x + 10) + 4(7x + 4)) - (11(7x + 4) + 6(6x + 10)) \\
&= ((12x + 7) + (3x + 11)) \cdot ((5x + 4) + (2x + 3)) - ((12x + 5) + (10x + 8)) \\
&= (2x + 5) \cdot (7x + 7) - (9x) \\
&= (x^2 + 2 \cdot 7x + 5 \cdot 7x + 5 \cdot 7) - (9x) \\
&= (x^2 + x + 9x + 9) - (9x) \\
&= x^2 + x + 9
\end{aligned}$$

4776 Given instance  $I_1 = 11$  a proofer therefore provides the polynomial  $x^2 + x + 9$  as proof. To verify  
4777 this proof, any verifier can then look up the target polynomial  $T$  from the QAP and divide  $P_{(I;W)}$   
4778 by  $T$ . In this particular example,  $P_{(I;W)}$  is equal to the target polynomial  $T$ , and hence, it is  
4779 divisible by  $T$  with  $P/T = 1$ . The verification therefore checks the proof.

```

4780 sage: F13 = GF(13) 639
4781 sage: F13t.<t> = F13[] 640
4782 sage: T = F13t(t^2 + t + 9) 641
4783 sage: P = F13t((2*(6*t+10)+6*(7*t+4))*(3*(6*t+10)+4*(7*t+4)) 642
4784             -(11*(7*t+4)+6*(6*t+10)))
4785 sage: P == T 643
4786 True 644
4787 sage: P % T # remainder 645
4788 0 646

```

To give an example of a false proof, consider the tuple  $(I_1; W_1, W_2, W_3, W_4) = (11, 2, 3, 4, 8)$ . Executing the circuit, we can see that this is not a valid assignment and not a solution to the R1CS, and hence, not a constructive knowledge proof in  $L_{3.fac\_zk}$ . However, a proofer might use these values to construct a false proof  $P_{(I;W)}$ :

$$\begin{aligned}
P'_{(I;W)} &= (A_0 + \sum_j^n I_j \cdot A_j + \sum_j^m W_j \cdot A_{n+j}) \cdot (B_0 + \sum_j^n I_j \cdot B_j + \sum_j^m W_j \cdot B_{n+j}) - (C_0 + \sum_j^n I_j \cdot C_j + \sum_j^m W_j \cdot C_{n+j}) \\
&= (2(6x + 10) + 8(7x + 4)) \cdot (3(6x + 10) + 4(7x + 4)) - (8(6x + 10) + 11(7x + 4)) \\
&= 8x^2 + 6
\end{aligned}$$

Given instance  $I_1 = 11$ , a proofer therefore provides the polynomial  $8x^2 + 6$  as proof. To verify this proof, any verifier can look up the target polynomial  $T$  from the QAP and divide  $P_{(I;W)}$  by  $T$ . However, polynomial division has the following remainder:

$$(8x^2 + 6)/(x^2 + x + 9) = 8 + \frac{5x + 12}{x^2 + x + 9}$$

4789 This implies that  $P_{(I;W)}$  is not divisible by  $T$ , and hence, the verification does not check the  
4790 proof. Any verifier can therefore show that the proof is false.

```

4791 sage: F13 = GF(13) 647
4792 sage: F13t.<t> = F13[] 648
4793 sage: T = F13t(t^2 + t + 9) 649
4794 sage: P = F13t((2*(6*t+10)+8*(7*t+4))*(3*(6*t+10)+4*(7*t+4))-( 650
4795             8*(6*t+10)+11*(7*t+4)))

```

4796	<b>sage:</b> $P == F13t(8*t^2 + 6)$	651
4797	<b>True</b>	652
4798	<b>sage:</b> $P \% T$ <b># remainder</b>	653
4799	$5*t + 12$	654

# Chapter 7

## Circuit Compilers

As we have seen in the previous chapter, statements can be formalized as membership or knowledge claims in formal language, and algebraic circuits as well as rank-1 constraint systems are two practically important ways to define those languages.

However, both algebraic circuits and rank-1 constraint systems are not ideal from a developers point of view, because they deviate substantially from common programming paradigms. Writing real-world applications as circuits and the associated verification in terms of rank-1 constraint systems is at least as troublesome as writing any other low-level language like assembler code. To allow for complex statement design, it is therefore necessary to have some kind of compiler framework, capable of transforming high-level languages into arithmetic circuits and associated rank-1 constraint systems.

As we have seen in chapter 6 and in 6.2.1.1., both arithmetic circuits and rank-1 constraint systems have a modularity property by which it is possible to synthesize complex circuits from simple ones. A basic approach taken by many circuit/R1CS compilers is therefore to provide a library of atomic and simple circuits and then define a way to combine those basic building blocks into arbitrary complex systems.

In this chapter, we provide an introduction to basic concepts of so-called **circuit compilers** and derive a toy language which we can “compile” in a pen-and-paper approach into algebraic circuits and their associated rank-1 constraint systems.

We start with a general introduction to our language, and then introduce atomic types like booleans and unsigned integers. Then we define the fundamental control flow primitives like the if-then-else conditional and the bounded loop. We will look at basic functionality primitives like elliptic curve cryptography. Primitives like these are often called **gadgets** in the literature.

### 7.1 A Pen-and-Paper Language

To explain basic concepts of circuit compilers and their associated high-level languages, we derive an informal toy language and associated “brain-compiler” which we name PAPER (**Pen-And-Paper Execution Rules**). PAPER allows programmers to define statements in Rust-like pseudo-code. The language is inspired by ZOKRATES and `circom`.

#### 7.1.1 The Grammar

In PAPER, any statement is defined as an ordered list of functions, where any function has to be declared in the list before it is called in another function of that list. The last entry in a statement has to be a special function, called `main`. Functions take a list of typed parameters as inputs

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and compute a tuple of typed variables as output, where types are special functions that define how to transform that type into another type, ultimately transforming any type into elements of the base field where the circuit is defined over.

Any statement is parameterized over the field that the circuit will be defined on, and has additional optional parameters of unsigned type, needed to define the size of array or the counter of bounded loops. The following definition makes the grammar of a statement precise using a command line language like description:

```
statement <Name> {F:<Field> [ , <N_1: unsigned>, ... ] } {
  [fn <Name>([[pub]<Arg>:<Type>, ...]) -> (<Type>, ...)] {
    [let [pub] <Var>:<Type> ; ... ]
    [let const <Const>:<Type>=<Value> ; ... ]
    Var<==>(fn ([<Arg>|<Const>|<Var>, ...]) | (<Arg>|<Const>|<Var>)) ;
    return (<Var>, ...) ;
  } ; ...]
  fn main([[pub]<Arg>:<Type>, ...]) -> (<Type>, ...) {
    [let [pub] <Var>:<Type> ; ... ]
    [let const <Const>:<Type>=<Value> ; ... ]
    Var<==>(fn ([<Arg>|<Const>|<Var>, ...]) | (<Arg>|<Const>|<Var>)) ;
    return (<Var>, ...) ;
  } ;
}
```

Function arguments and variables are private by default, but can be declared as public by the `pub` specifier. Declaring arguments and variables as public always overwrites any previous or conflicting private declarations. Every argument, constant or variable has a type, and every type is defined as a function that transforms that type into another type:

```
type <TYPE>( t1 : <TYPE_1>) -> TYPE_2{
  let t2: TYPE_2 <== fn(TYPE_1)
  return t2
}
```

Many real-world circuit languages are based on a similar, but of course more sophisticated approach than PAPER. The purpose of PAPER is to show basic principles of circuit compilers and their associated high-level languages.

*Example 124.* To get a better understanding of the grammar of PAPER, the following constitutes proper high-level code that follows the grammar of the PAPER language, assuming that all types in that code have been defined elsewhere.

```
statement MOCK_CODE {F: F_43, N_1 = 1024, N_2 = 8} {
  fn foo(in_1 : F, pub in_2 : TYPE_2) -> F {
    let const c_1 : F = 0 ;
    let const c_2 : TYPE_2 = SOME_VALUE ;
    let pub out_1 : F ;
    out_1<== c_1 ;
    return out_1 ;
  } ;

  fn bar(pub in_1 : F) -> F {
    let out_1 : F ;
    out_1<==foo(in_1);
    return out_1 ;
  } ;
}
```

```

4881   } ;
4882
4883   fn main(in_1 : TYPE_1) -> (F, TYPE_2) {
4884       let const c_1 : TYPE_1 = SOME_VALUE ;
4885       let const c_2 : F = 2;
4886       let const c_3 : TYPE_2 = SOME_VALUE ;
4887       let pub out_1 : F ;
4888       let out_2 : TYPE_2 ;
4889       c_1 <== in_1 ;
4890       out_1 <== foo(c_2) ;
4891       out_2 <== TYPE_2 ;
4892       return (out_1, out_2) ;
4893   } ;
4894 }

```

### 7.1.2 The Execution Phases

In contrast to normal executable programs, programs for circuit compilers have two modes of execution. The first mode, usually called the **setup phase**, is executed in order to generate the circuit and its associated rank-1 constraint system, the latter of which is then usually used as input to some zero-knowledge proof system.

The second mode of execution is usually called the **prover phase**. In this phase, a prover usually computes a valid assignment to the circuit. Depending on the use case, this valid assignment is then either directly used as constructive proof for proper circuit execution or is transferred as input to the proof generation algorithm of some zero-knowledge proof system, where the full-sized, non hiding constructive proof is processed into a succinct proof with various levels of zero knowledge.

Modern circuit languages and their associated compilers abstract over those two phases and provide a unified **interphase** to the developer, who then writes a single program that can be used in both phases.

To give the reader a clear, conceptual distinction between the two phases, PAPER keeps them separated. Code can be “brain-compiled” during the **setup-phase** in a pen-and-paper approach into visual circuits. Once a circuit is derived, it can be executed in a **prover phase** to generate a valid assignment. The valid assignment is then interpreted as a constructive proof for a knowledge claim in the associated language.

**The Setup Phase** In PAPER, the task of the setup phase is to compile code in the PAPER language into a visual representation of an algebraic circuit. Deriving the circuit from the code in a pen-and-paper style is what we call **brain-compiling**.

Given some statement description that adheres to the correct grammar, we start circuit development with an empty circuit, compile the main function first and then inductively compile all other functions as they are called during the process.

For every function we compile, we draw a box-node for every argument, every variable and every constant of that function. If the node represents a variable, we label it with that variable’s name, and if it represents a constant, we label it with that constant’s value. We group arguments into a subgraph labeled “inputs” and return values into a subgraph labeled “outputs”. We then group everything into a subgraph and label that subgraph with the function’s name.

After this is done, we have to do a consistency and type check for every occurrence of the

assignment operator `<==`. We have to ensure that the expression on the right side of the operator is well defined and that the types of both side match.

Then we compile the right side of every occurrence of the assignment operator `<==`. If the right side is a constant or variable defined in this function, we draw a dotted line from the box-node that represents the left side of `<==` to the box node that represents the right side of the same operator. If the right side represents an argument of that function we draw a line from the box-node that represents the left side of `<==` to the box node that represents the right side of the same operator.

If the right side of the `<==` operator is a function, we look into our database, find its associated circuit and draw it. If no circuit is associated to that function yet, we repeat the compilation process for that function, drawing edges from the function's argument to its input nodes and from the functions output nodes to the nodes on the right side of `<==`.

During that process, edge labels are drawn according to the rules from 6.2.2.1. If the associated variable represents a private value, we use the *W* label to indicate a witness, and if it represents a public value, we use the *I* label to indicate an instance.

Once this is done, we compile all occurring types in a function, by compiling the function of each type. We do this inductively until we reach the type of the base field. Circuits have no notion of types, only of field elements; hence, every type needs to be compiled to the field type in a sequence of compilation steps.

The compilation stops once we have inductively replaced all functions by their circuits. The result is a circuit that contains many unnecessary box nodes. In a final optimization step, all box nodes that are directly linked to each other are collapsed into a single node, and all box nodes that represent the same constants are collapsed into a single node.

Of course, PAPER's brain-compiler is not properly defined in any formal manner. Its purpose is to highlight important steps that real-world compilers undergo in their setup phases.

*Example 125 (A trivial Circuit).* To give an intuition of how to write and compile circuits in the PAPER language, consider the following statement description:

```
statement trivial_circuit {F:F_13} {
  fn main{F}(in1 : F, pub in2 : F) -> (F,F) {
    let const outc1 : F = 0 ;
    let const incl : F = 7 ;
    let out1 : F ;
    let out2 : F ;
    out1 <== incl;
    out2 <== in1;
    outc1 <== in2;
    return (out1, out2) ;
  }
}
```

To brain-compile this statement into an algebraic circuit with PAPER, we start with an empty circuit and evaluate function `main`, which is the only function in this statement.

We draw box-nodes for every argument, every constant and every variable of the function and label them with their names or values, respectively. Then we do a consistency and type check for every `<==` operator in the function. Since the circuit only wires inputs to outputs and all elements have the same type, the check is valid.

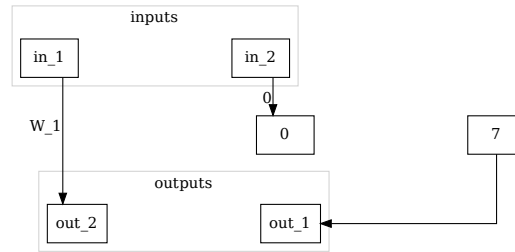
Then we evaluate the right side of the assignment operators. Since, in our case, the right side of each operator is not a function, we draw edges from the box-nodes on the right side to the associated box node on the left side. To label those edges, we use the general rules of

check  
reference



algebraic circuits as defined in 6.2.2.1. According to those rules, every incoming edge of a sink node has a label and every outgoing edge of a source node has a label, if the node is labeled with a variable. Since nodes that represent constants are implicitly assumed to be private, and since the public specifier determines if an edge is labeled with  $W$  or  $I$ , we get the following circuit:

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**The Prover Phase** In PAPER, a so-called **prover phase** can be executed once the setup phase has generated a circuit image from its associated high-level code. This is done by executing the circuit while assigning proper values to all input nodes of the circuit. However, in contrast to most real-world compilers, PAPER does not tell the prover how to find proper input values to a given circuit. Real-world programming languages usually provide this data by computations that are done outside of the circuit.

*Example 126.* Consider the circuit from example 125. Valid assignments to this circuit are constructive proofs that the pair of inputs  $(S_1, S_2)$  is a point on the tiny-jubjub curve. However, the circuit does not provide a way to actually compute proper values for  $S_1$  and  $S_2$ . Any real-world system therefore needs an auxiliary computation that provides those values.

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## 7.2 Common Programing concepts

In this section, we cover concepts that appear in almost every programming language, and see how they can be implemented in circuit compilers.

### 7.2.1 Primitive Types

Primitive data types like booleans, (unsigned) integers, or strings are the most basic building blocks one can expect to find in every general high-level programming language. In order to write statements as computer programs that compile into circuits, it is therefore necessary to implement primitive types as constraint systems, and define their associated operations as circuits.

In this section, we look at some common ways to achieve this. After a recapitulation of the atomic type of prime field elements, we start with an implementation of the boolean type and its associated boolean algebra as circuits. After that, we define unsigned integers based on the boolean type, and leave the implementation of signed integers as an exercise to the reader.

It should be noted, however, that while primitive data types in common programming languages (like C, Go, or Rust) have a one-to-one correspondence with objects in the computer's memory, this is not the case for most languages that compile into algebraic circuits. As we will see in the following paragraphs, common primitives like booleans or unsigned integers require many constraints and memory. Primitives different from the underlying field elements can be expensive.

### The base-field type

Since both algebraic circuits and their associated rank-1 constraint systems are defined over a finite field, elements from that field are the atomic informational units in those models. In this sense, field elements  $x \in \mathbb{F}$  are for algebraic circuits what bits are for computers.

In PAPER, we write  $F$  for this type and specify the actual field instance for every statement in curly brackets after the name of that statement. Two functions are associated to this type, which are induced by the **addition** and **multiplication** law in the field  $F$ . We write

$$\text{MUL} : F \times F \rightarrow F ; (x, y) \mapsto \text{MUL}(x, y) \quad (7.1)$$

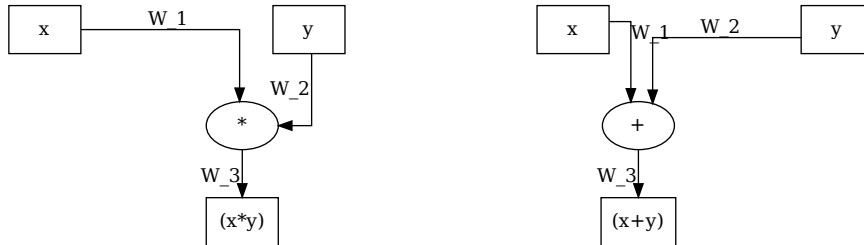
$$\text{ADD} : F \times F \rightarrow F ; (x, y) \mapsto \text{ADD}(x, y) \quad (7.2)$$

Circuit compilers have to compile these functions into algebraic gates, as explained in 6.2.2.2. Every other function has to be expressed in terms of them and proper wiring.

To represent addition and multiplication in the PAPER language, we define the following two functions:

```
fn MUL(x : F, y : F) -> (MUL(x, y) : F) {}
fn ADD(x : F, y : F) -> (ADD(x, y) : F) {}
```

The compiler then compiles every occurrence of the MUL or the ADD function into the following circuits:



*Example 127* (Basic gates). To give an intuition of how a real-world compiler might transform addition and multiplication in algebraic expressions into a circuit, consider the following PAPER statement:

```
statement basic_ops {F:F_13} {
  fn main(in_1 : F, pub in_2 : F) -> (out_1:F, out_2:F) {
    out_1 <== MUL(in_1, in_2) ;
    out_2 <== ADD(in_1, in_2) ;
  }
}
```

To compile this into an algebraic circuit, we start with an empty circuit and evaluate the function `main`, which is the only function in this statement.

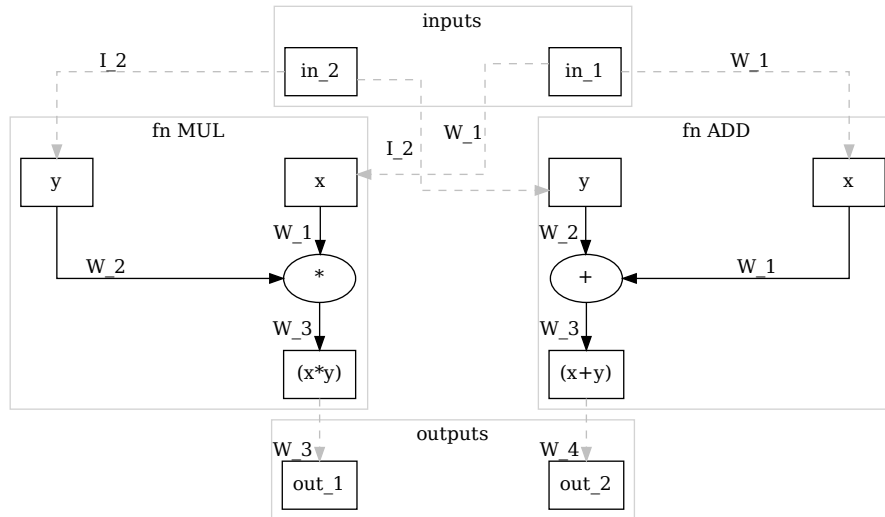
We draw an inputs subgraph containing box-nodes for every argument of the function, and an outputs subgraph containing box-nodes for every factor in the return value. Since all of these nodes represent variables of the `field` type, we don't have to add any type constraints to the circuit.

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We check the validity of every expression on the right side of every `<==` operator including a type check. In our case, every variable is of the `field` type and hence the types match the types of the `MUL` as well as the `ADD` function and the type of the left sides of `<==` operators.

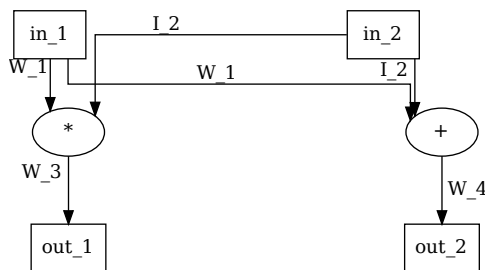
We evaluate the expressions on the right side of every `<==` operator inductively, replacing every occurrence of a function with a subgraph that represents its associated circuit.

According to `PAPER`, every occurrence of the `public` specifier overwrites the associate `private` default value. Using the appropriate edge labels we get:



Any real-world compiler might process its associated high-level language in a similar way, replacing functions, or gadgets by predefined associated circuits. This process is often followed by various optimization steps that try to reduce the number of constraints as much as possible.

In `PAPER`, we optimize this circuit by collapsing all box nodes that are directly connected to other box nodes, adhering to the rule that a variable's `public` specifier overwrites any `private` specifier. Reindexing edge labels, we get the following circuit as our pen and pencil compiler output:



*Example 128 (3-factorization).* Consider our 3-factorization problem from example 106 and the associated circuit  $C_{3.fac\_zk}(\mathbb{F}_{13})$  we provided in equation 6.8. To understand the process of replacing high-level functions by their associated circuits inductively, we want define a `PAPER` statement that we brain-compile into an algebraic circuit equivalent to  $C_{3.fac\_zk}(\mathbb{F}_{13})$ :

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```

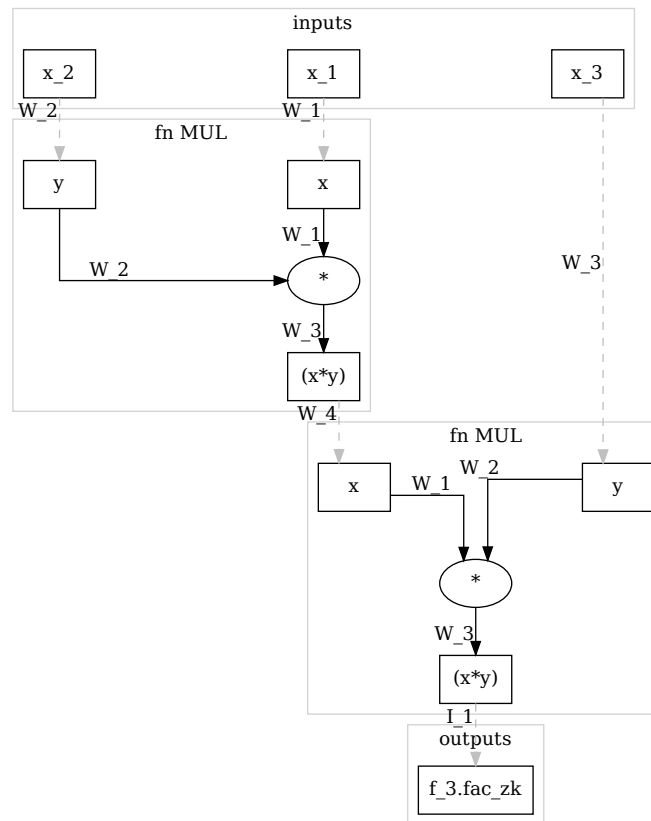
5059 statement 3_fac_zk {F:F_13} {
5060   fn main(x_1 : F, x_2 : F, x_3 : F) -> (pub 3_fac_zk : F){
5061     f_3.fac_zk <== MUL( MUL( x_1 , x_2 ) , x_3 ) ;
5062   }
5063 }

```

Using PAPER, we start with an empty circuit and then add 3 input nodes to the input subgraph as well as 1 output node to the output subgraph. All these nodes are decorated with the associated variable names. Since all of these nodes represent variables of the `field` type, we don't have to add any type constraints to the circuit.

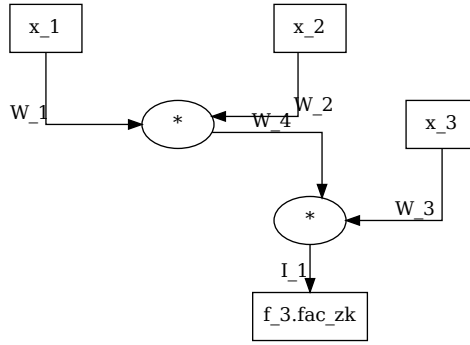
We check the validity of every expression on the right side of the single `<==` operator including a type check.

We evaluate the expressions on the right side of every `<==` operator inductively. We have two nested multiplication functions and we replace them by the associated multiplication circuits, starting with the most outer function. We get:



5073

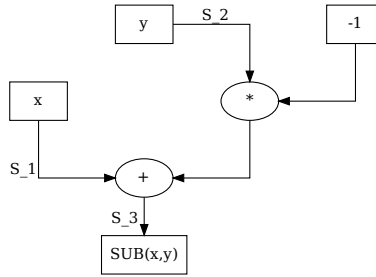
In a final optimization step, we collapse all box nodes directly connected to other box nodes, adhering to the rule that a variables `public` specifier overwrites any `private` specifier. Reindexing edge labels we get the following circuit:



5077

**The Subtraction Constraint System** By definition, algebraic circuits only contain addition and multiplication gates, and it follows that there is no single gate for field subtraction, despite the fact that subtraction is a native operation in every field.

High-level languages and their associated circuit compilers, therefore, need another way to deal with subtraction. To see how this can be achieved, recall that subtraction is defined by addition with the additive inverse, and that the inverse can be computed efficiently by multiplication with  $-1$ . A circuit for field subtraction is therefore given by



5085

Using the general method from 6.2.1.1, the circuits associated rank-1 constraint system is given by:

$$(S_1 + (-1) \cdot S_2) \cdot 1 = S_3 \quad (7.3)$$

Any valid assignment  $\{S_1, S_2, S_3\}$  to this circuit therefore enforces the value  $S_3$  to be the difference  $S_1 - S_2$ .

Real-world compilers usually provide a gadget or a function to abstract over this circuit such that programers can use subtraction as if it were native to circuits. In PAPER, we define the following subtraction function that compiles to the previous circuit:

```
fn SUB(x : F, y : F) -> (SUB(x,y) : F) {
  constant c : F = -1 ;
  SUB <== ADD(x , MUL( y , c ) );
}
```

In the setup phase of a statement, we compile every occurrence of the SUB function into an instance of its associated subtraction circuit, and edge labels are generated according to the rules from 6.2.2.1.

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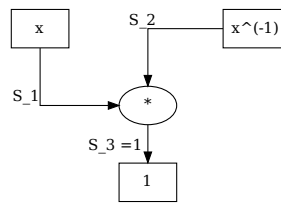
**The Inversion Constraint System** By definition, algebraic circuits only contain addition and multiplication gates, and it follows that there is no single gate for field inversion, despite the fact that inversion is a native operation in every field.

If the underlying field is a prime field, one approach would be to use Fermat’s little theorem 3.3 to compute the multiplicative inverse inside the circuit. To see how this works, let  $\mathbb{F}_p$  be the prime field. The multiplicative inverse  $x^{-1}$  of a field element  $x \in \mathbb{F}$  with  $x \neq 0$  is then given by  $x^{-1} = x^{p-2}$ , and computing  $x^{p-2}$  in the circuit therefore computes the multiplicative inverse.

Unfortunately, real-world primes  $p$  are large and computing  $x^{p-2}$  by repeated multiplication of  $x$  with itself is infeasible. A “double and multiply” approach (as described in XXX) is faster, as it computes the power in roughly  $\log_2(p)$  steps, but still adds a lot of constraints to the circuit.

Computing inverses in the circuit makes no use of the fact that inversion is a native operation in any field. A more constraints friendly approach is therefore to compute the multiplicative inverse outside of the circuit and then only enforce correctness of the computation in the circuit.

To understand how this can be achieved, observe that a field element  $y \in \mathbb{F}$  is the multiplicative inverse of a field element  $x \in \mathbb{F}$  if and only if  $x \cdot y = 1$  in  $\mathbb{F}$ . We can use this, and define a circuit that has two inputs,  $x$  and  $y$ , and enforces  $x \cdot y = 1$ . It is then guaranteed that  $y$  is the multiplicative inverse of  $x$ . The price we pay is that we can not compute  $y$  by circuit execution, but auxiliary data is needed to tell any prover which value of  $y$  is needed for a valid circuit assignment. The following circuit defines the constraint



Using the general method from 6.2.1.1, the circuit is transformed into the following rank-1 constraint system:

$$S_1 \cdot S_2 = 1 \quad (7.4)$$

Any valid assignment  $\{S_1, S_2\}$  to this circuit enforces that  $S_2$  is the multiplicative inverse of  $S_1$ , and, since there is no field element  $S_2$  such that  $0 \cdot S_2 = 1$ , it also handles the fact that the multiplicative inverse of 0 is not defined in any field.

Real-world compilers usually provide a gadget or a function to abstract over this circuit, and those functions compute the inverse  $x^{-1}$  as part of their witness generation process. Programmers then don’t have to care about providing the inverse as auxiliary data to the circuit. In PAPER, we define the following inversion function that compiles to the previous circuit:

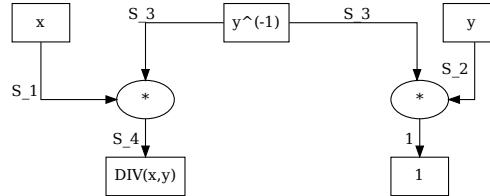
```
fn INV(x : F, y : F) -> (x_inv : F) {
  constant c : F = 1 ;
  c <== MUL( x , y ) ;
  x_inv <== y ;
}
```

As we see, this functions takes two inputs, the field value and its inverse. It therefore does not handle the computation of the inverse by itself. This is to keep PAPER as simple as possible.

In the setup phase, we compile every occurrence of the INV function into an instance of the inversion circuit XXX, and edge labels are generated according to the rules from XXX.

**The Division Constraint System** By definition, algebraic circuits only contain addition and multiplication gates, and it follows that there is no single gate for field division, despite the fact that division is a native operation in every field.

Implementing division as a circuit, we use the fact that division is multiplication with the multiplicative inverse. We therefore define division as a circuit using the inversion circuit and constraint system from the previous paragraph. Expensive inversion is computed outside of the circuit and then provided as circuit input. We get



Using the method from 6.2.1.1, we transform this circuit into the following rank-1 constraint system:

$$\begin{aligned} S_2 \cdot S_3 &= 1 \\ S_1 \cdot S_3 &= S_4 \end{aligned}$$

Any valid assignment  $\{S_1, S_2, S_3, S_4\}$  to this circuit enforces  $S_4$  to be the field division of  $S_1$  by  $S_2$ . It handles the fact that division by 0 is not defined, since there is no valid assignment in case  $S_2 = 0$ .

In PAPER, we define the following division function that compiles to the previous circuit:

```
fn DIV(x : F, y : F, y_inv : F) -> (DIV : F) {
  DIV <== MUL( x , INV( y, y_inv ) ) ;
}
```

In the setup phase, we compile every occurrence of the binary INV operator into an instance of the inversion circuit.

**Exercise 44.** Let  $F$  be the field  $\mathbb{F}_5$  of modular 5 arithmetics from example 13. Brain-compile the following PAPER statement into an algebraic circuit:

```
statement STUPID_CIRC {F: F_5} {
  fn foo(in_1 : F, in_2 : F)->(out_1 : F, out_2 : F){
    constant c_1 : F = 3 ;
    out_1<== ADD( MUL( c_1 , in_1 ) , in_1 ) ;
    out_2<== INV( c_1 , in_2 ) ;
  } ;

  fn main(in_1 : F, in_2 : F)->(out_1 : F, out_2 : TYPE_2){
    constant (c_1,c_2) : (F,F) = (3,2) ;
    (out_1,out_2) <== foo(in_1, in_2) ;
  } ;
}
```

**Exercise 45.** Consider the tiny-jubjub curve from example 66 and its associated circuit 125. Write a statement in PAPER that brain-compiles the statement into a circuit equivalent to the one derived in XXX, assuming that curve points are instances and every other assignment is a witness.



5173 *Exercise 46.* Let  $F = \mathbb{F}_{13}$  be the modular 13 prime field and  $x \in F$  some field element. Define a  
 5174 statement in PAPER such that given instance  $x$  a field element  $y \in F$  is a witness for the statement  
 5175 if and only if  $y$  is the square root of  $x$ .

5176 Brain-compile the statement into a circuit and derive its associated rank-1 constraint system.  
 5177 Consider the instance  $x = 9$  and compute a constructive proof for the statement.

## 5178 The boolean Type

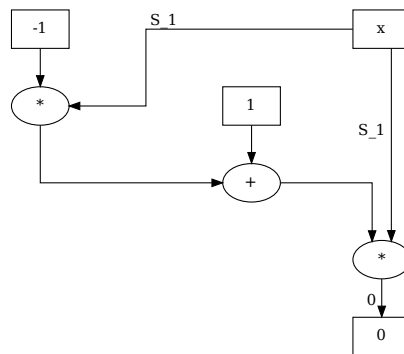
5179 Booleans are a classical primitive type, implemented by virtually every higher programing lan-  
 5180 guage. It is therefore important to implement booleans in circuits. One of the most common  
 5181 ways to do this is by interpreting the additive and multiplicative neutral element  $\{0, 1\} \subset \mathbb{F}$  as  
 5182 the two boolean values such that 0 represents *false* and 1 represents *true*. boolean operators  
 5183 like *and*, *or*, or *xor* are then expressible as algebraic computations inside  $\mathbb{F}$ .

5184 Representing booleans this way is convenient, because the elements 0 and 1 are defined in  
 5185 any field. The representation is therefore independent of the actual field in consideration.

5186 To fix boolean algebra notation, we write 0 to represent *false* and 1 to represent *true*, and  
 5187 we write  $\wedge$  to represent the boolean AND as well as  $\vee$  to represent the boolean OR operator.  
 5188 The boolean NOT operator is written as  $\neg$ .

5189 **The boolean Constraint System** To represent booleans by the additive and multiplicative  
 5190 neutral elements of a field, a constraint is required to actually enforce variables of boolean type  
 5191 to be either 1 or 0. In fact, many of the following circuits that represent boolean functions are  
 5192 only correct under the assumption that their input variables are constrained to be either 0 or 1.  
 5193 Not constraining boolean variables is a common problem in circuit design.

5194 In order to constrain an arbitrary field element  $x \in \mathbb{F}$  to be 1 or 0, the key observation is  
 5195 that the equation  $x \cdot (1 - x) = 0$  has only two solutions 0 and 1 in any field. Implementing this  
 5196 equation as a circuit therefore generates the correct constraint:



5197

Using the method from 6.2.1.1, we transform this circuit into the following rank-1 constraint system:

$$S_1 \cdot (1 - S_1) = 0$$

check  
reference

5198 Any valid assignment  $\{S_1\}$  to this circuit enforces  $S_1$  to be either 0 or 1.

5199 Some real-world circuit compilers (like ZOKRATES or BELLMAN) are typed, while others  
 5200 (like circom) are not. However, all of them have their way of dealing with the binary con-  
 5201 straint. In PAPER, we define the following boolean type that compiles to the previous circuit:

```

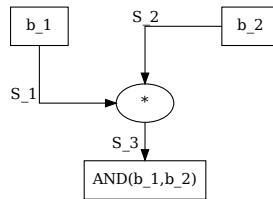
5202 type BOOL(b : BOOL) -> (x : F) {
5203     constant c1 : F = 0 ;
5204     constant c2 : F = 1 ;
5205     constant c3 : F = -1 ;
5206     c1 <== MUL( x , ADD( c2 , MUL( x , c3 ) ) ) ;
5207     x <== b ;
5208 }

```

5209 In the setup phase of a statement, we compile every occurrence of a variable of boolean type  
 5210 into an instance of its associated boolean circuit.

5211 **The AND operator constraint system** Given two field elements  $b_1$  and  $b_2$  from  $\mathbb{F}$  that are  
 5212 constrained to represent boolean variables, we want to find a circuit that computes the logical  
 5213 **and** operator  $AND(b_1, b_2)$  as well as its associated R1CS that enforces  $b_1, b_2, AND(b_1, b_2)$  to  
 5214 satisfy the constraint system if and only if  $b_1 \wedge b_2 = AND(b_1, b_2)$  holds true.

5215 The key insight here is that, given three boolean constraint variables  $b_1, b_2$  and  $b_3$ , the  
 5216 equation  $b_1 \cdot b_2 = b_3$  is satisfied in  $\mathbb{F}$  if and only if the equation  $b_1 \wedge b_2 = b_3$  is satisfied in  
 5217 boolean algebra. The logical operator  $\wedge$  is therefore implementable in  $\mathbb{F}$  by field multiplication  
 5218 of its arguments and the following circuit computes the  $\wedge$  operator in  $\mathbb{F}$ , assuming all inputs are  
 5219 restricted to be 0 or 1:



5220

5221 The associated rank-1 constraint system can be deduced from the general process 6.2.1.1 and  
 5222 consists of the following constraint:

$$S_1 \cdot S_2 = S_3 \quad (7.5)$$

 check  
reference

5223 Common circuit languages typically provide a gadget or a function to abstract over this circuit  
 5224 such that programers can use the  $\wedge$  operator without caring about the associated circuit. In  
 5225 PAPER, we define the following function that compiles to the  $\wedge$ -operator's circuit:

```

5226 fn AND(b_1 : BOOL, b_2 : BOOL) -> AND(b_1, b_2) : BOOL{
5227     AND(b_1, b_2) <== MUL( b_1 , b_2 ) ;
5228 }

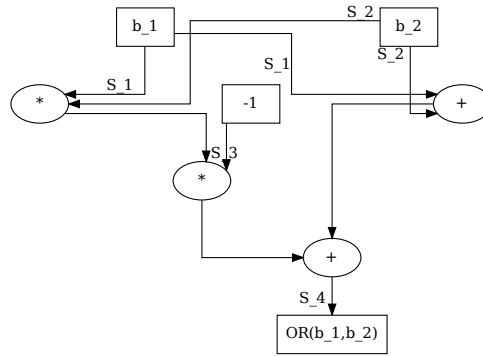
```

5229 In the setup phase of a statement, we compile every occurrence of the AND function into an  
 5230 instance of its associated  $\wedge$ -operator's circuit.

5231 **The OR operator constraint system** Given two field elements  $b_1$  and  $b_2$  from  $\mathbb{F}$  that are  
 5232 constrained to represent boolean variables, we want to find a circuit that computes the logical  
 5233 **or** operator  $OR(b_1, b_2)$  as well as its associated R1CS that enforces  $b_1, b_2, OR(b_1, b_2)$  to satisfy  
 5234 the constraint system if and only if  $b_1 \vee b_2 = OR(b_1, b_2)$  holds true.

5235 Assuming that three variables  $b_1, b_2$  and  $b_3$  are boolean constraint, the equation  $b_1 + b_2 - b_1 \cdot$   
 5236  $b_2 = b_3$  is satisfied in  $\mathbb{F}$  if and only if the equation  $b_1 \vee b_2 = b_3$  is satisfied in boolean algebra.  
 5237 The logical operator  $\vee$  is therefore implementable in  $\mathbb{F}$  by the following circuit, assuming all  
 5238 inputs are restricted to be 0 or 1:

 "constraints"  
or "con-  
strained"?



5239

The associated rank-1 constraint system can be deduced from the general process 6.2.1.1 and consists of the following constraints:

check  
reference

$$\begin{aligned} S_1 \cdot S_2 &= S_3 \\ (S_1 + S_2 - S_3) \cdot 1 &= S_4 \end{aligned}$$

Common circuit languages typically provide a gadget or a function to abstract over this circuit such that programers can use the  $\vee$  operator without caring about the associated circuit. In PAPER, we define the following function that compiles to the  $\vee$ -operator's circuit:

```
5240 fn OR(b_1 : BOOL, b_2 : BOOL) -> OR(b_1,b_2) : BOOL{
5241   constant c1 : F = -1 ;
5242   OR(b_1,b_2) <== ADD(ADD(b_1,b_2),MUL(c1,MUL(b_1,b_2))) ;
5243 }
5244
```

In the setup phase of a statement, we compile every occurrence of the OR function into an instance of its associated  $\vee$ -operator's circuit.

*Exercise 47.* Let  $\mathbb{F}$  be a finite field and let  $b_1$  as well as  $b_2$  two boolean constraint variables from  $\mathbb{F}$ . Show that the equation  $OR(b_1, b_2) = 1 - (1 - b_1) \cdot (1 - b_2)$  holds true.

Use this equation to derive an algebraic circuit with ingoing variables  $b_1$  and  $b_2$  and outgoing variable  $OR(b_1, b_2)$  such that  $b_1$  and  $b_2$  are boolean **constraint** and the circuit has a valid assignment, if and only if  $OR(b_1, b_2) = b_1 \vee b_2$ .

"constraints"  
or "con-  
strained"?

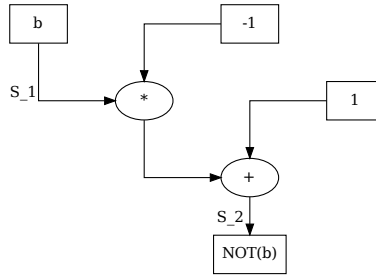
Use the technique from XXX to transform this circuit into a rank-1 constraint system and find its full solution set. Define a PAPER function that brain-compiles into the circuit.

add refer-  
ence

**The NOT operator constraint system** Given a field element  $b$  from  $\mathbb{F}$  that is constrained to represent a boolean variable, we want to find a circuit that computes the logical **NOT** operator  $NOT(b)$  as well as its associated RICS that enforces  $b, NOT(b)$  to satisfy the constraint system if and only if  $\neg b = NOT(b)$  holds true.

Assuming that two variables  $b_1$  and  $b_2$  are boolean **constraint**, the equation  $(1 - b_1) = b_2$  is satisfied in  $\mathbb{F}$  if and only if the equation  $\neg b_1 = b_2$  is satisfied in boolean algebra. The logical operator  $\neg$  is therefore implementable in  $\mathbb{F}$  by the following circuit, assuming all inputs are restricted to be 0 or 1:

"constraints"  
or "con-  
strained"?



5264

The associated rank-1 constraint system can be deduced from the general process XXX and consists of the following constraints

add reference

$$(1 - S_1) \cdot 1 = S_2$$

Common circuit languages typically provide a gadget or a function to abstract over this circuit such that programers can use the  $\neg$  operator without caring about the associated circuit. In PAPER, we define the following function that compiles to the  $\neg$ -operator's circuit:

```

5268 fn NOT(b : BOOL -> NOT(b) : BOOL{
5269   constant c1 = 1 ;
5270   constant c2 = -1 ;
5271   NOT(b_1) <== ADD( c1 , MUL( c2 , b ) ) ;
5272 }

```

In the setup phase of a statement, we compile every occurrence of the NOT function into an instance of its associated  $\neg$ -operator's circuit.

*Exercise 48.* Let  $\mathbb{F}$  be a finite field. Derive the algebraic circuit and associated rank-1 constraint system for the following operators: NOR, XOR, NAND, EQU.

**Modularity** As we have seen in chapter 6,, both algebraic circuits and R1CS have a modularity property, and as we have seen in this section, all basic boolean functions are expressible in circuits. Combining those two properties, show that it is possible to express arbitrary boolean functions as algebraic circuits.

check references

This shows that the expressiveness of algebraic circuits and therefore rank-1 constraint systems is as general as the expressiveness of boolean circuits. An important implication is that the languages  $L_{R1CS-SAT}$  and  $L_{Circuit-SAT}$  as defined in 2, are as general as the famous language  $L_{3-SAT}$ , which is known to be  $\mathcal{NP}$ -complete.

check reference

*Example 129.* To give an example of how a compiler might construct complex boolean expressions in algebraic circuits from simple ones and how we derive their associated rank-1 constraint systems, let's look at the following PAPER statement:

```

5288 statement BOOLEAN_STAT {F: F_p} {
5289   fn main(b_1:BOOL,b_2:BOOL,b_3:BOOL,b_4:BOOL )-> pub b_5:BOOL {
5290     b_5 <== AND( OR( b_1 , b_2 ) , AND( b_3 , NOT( b_4 ) ) ) ;
5291   } ;
5292 }

```

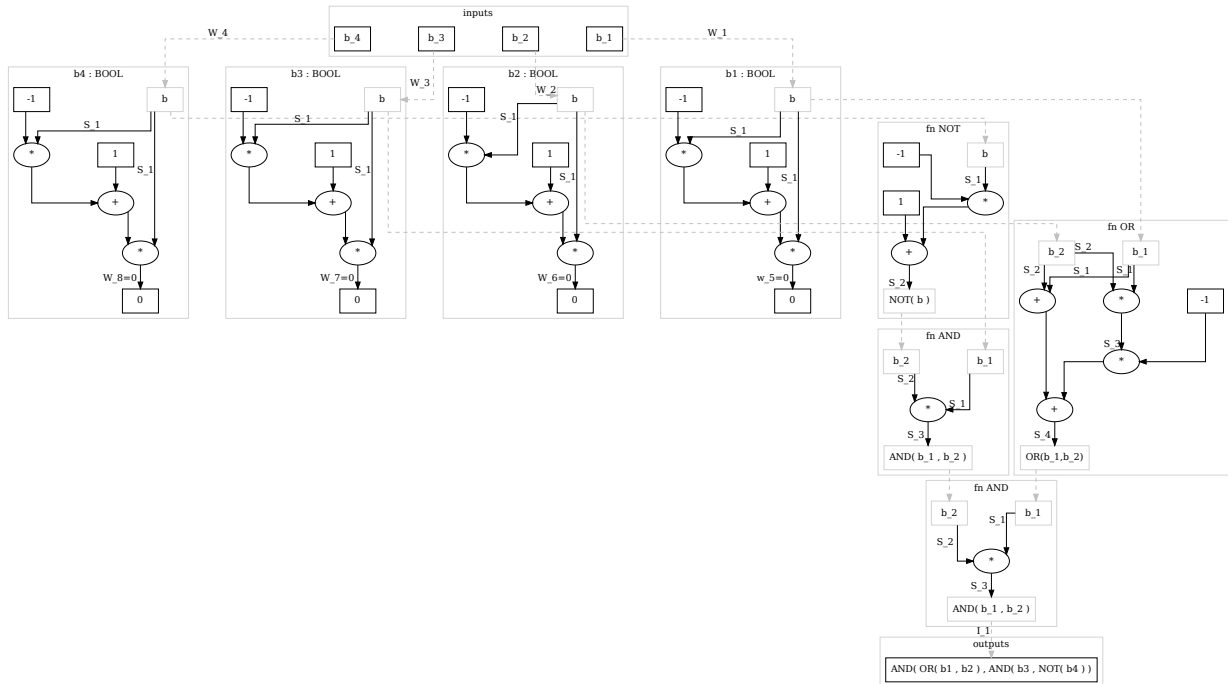
The code describes a circuit that takes four private inputs  $b_1, b_2, b_3$  and  $b_4$  of boolean type and computes a public output  $b_5$  such that the following boolean expression holds true:

$$(b_1 \vee b_2) \wedge (b_3 \wedge \neg b_4) = b_5$$

During a setup-phase, a circuit compiler transforms this high-level language statement into a circuit and associated rank-1 constraint systems and hence defines a language  $L_{BOOLEAN\_STAT}$ .

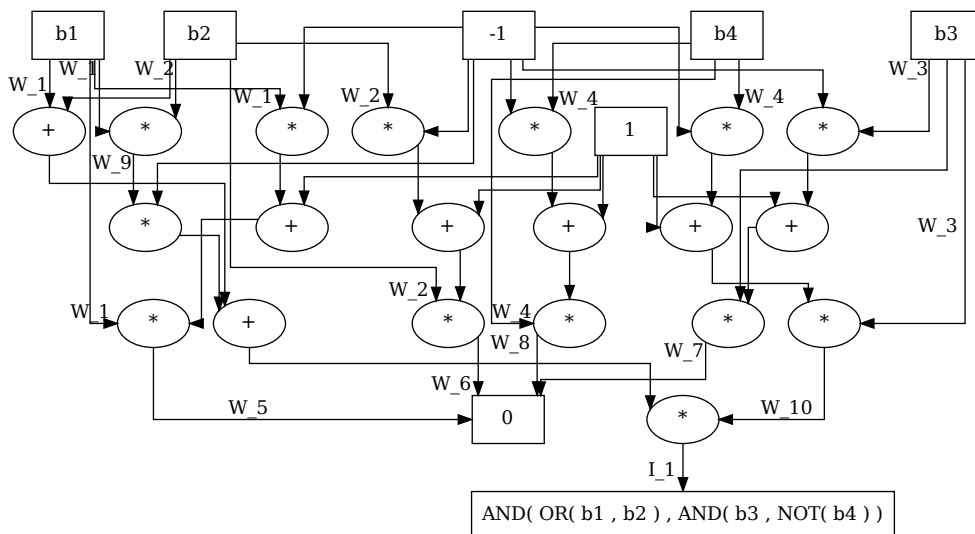
To see how this might be achieved, we use PAPER as an example to execute the setup-phase and compile `BOOLEAN_STAT` into a circuit. Taking the definition of the boolean constraint `XXX` as well as the definitions of the appropriate boolean operators into account, we get the following circuit:

add reference



Simple optimization then collapses all box-nodes that are directly linked and all box nodes that represent the same constants. After relabeling the edges, the following circuit represents the circuit associated to the `BOOLEAN_STAT` statement:

can we rotate this by 90°?



5307 Given some public input  $I_1$  from  $\mathbb{F}_{13}$ , a valid assignment to this circuit consists of private inputs  
 5308  $W_1, W_2, W_3, W_4$  from  $\mathbb{F}_{13}$  such that the equation  $I_1 = (W_1 \vee W_2) \wedge (W_3 \wedge \neg W_4)$  holds true. In  
 5309 addition, a valid assignment also has to contain private inputs  $W_5, W_6, W_7, W_8, W_9$  and  $W_{10}$ ,  
 5310 which can be derived from circuit execution. The inputs  $W_5, \dots, W_8$  ensure that the first four  
 5311 private inputs are either 0 or 1 but not any other field element, and the others enforce the boolean  
 5312 operations in the expression.

To compute the associated R1CS, we can use the general method from 6.2.1.1 and look at every labeled outgoing edge not coming from a source node. We declare the edges coming from input nodes as well as the edge going to the single output node as public, and every other edge as private input. In this case we get:

check  
reference

$$\begin{array}{ll}
 W_5 : W_1 \cdot (1 - W_1) = 0 & \text{boolean constraints} \\
 W_6 : W_2 \cdot (1 - W_2) = 0 & \\
 W_7 : W_3 \cdot (1 - W_3) = 0 & \\
 W_8 : W_4 \cdot (1 - W_4) = 0 & \\
 W_9 : W_1 \cdot W_2 = W_9 & \text{first OR-operator constraint} \\
 W_{10} : W_3 \cdot (1 - W_4) = W_{10} & \text{AND(.,NOT(.))-operator constraints} \\
 I_1 : (W_1 + W_2 - W_9) \cdot W_{10} = I_1 & \text{AND-operator constraints}
 \end{array}$$

5313 The reason why this R1CS only contains a single constraint for the multiplication gate in the  
 5314 OR-circuit, while the general definition XXX requires two constraints, is that the second con-  
 5315 straint in XXX only appears because the final addition gate is connected to an output node. In  
 5316 this case, however, the final addition gate from the OR-circuit is enforced in the left factor of  
 5317 the  $I_1$  constraint. Something similar holds true for the negation circuit.

add refer-  
ence

add refer-  
ence

During a prover-phase, some public instance  $I_5$  must be given. To compute a constructive proof for the statement of the associated languages with respect to instance  $I_5$ , a prover has to find four boolean values  $W_1, W_2, W_3$  and  $W_4$  such that

$$(W_1 \vee W_2) \wedge (W_3 \wedge \neg W_4) = I_5$$

5318 holds true. In our case neither the circuit nor the PAPER statement specifies how to find those  
 5319 values, and it is a problem that any prover has to solve outside of the circuit. This might or  
 5320 might not be true for other problems, too. In any case, once the prover found those values, they  
 5321 can execute the circuit to find a valid assignment.

To give a concrete example, let  $I_1 = 1$  and assume  $W_1 = 1, W_2 = 0, W_3 = 1$  and  $W_4 = 0$ . Since  $(1 \vee 0) \wedge (1 \wedge \neg 0) = 1$ , those values satisfy the problem and we can use them to execute the circuit. We get

$$\begin{array}{l}
 W_5 = W_1 \cdot (1 - W_1) = 0 \\
 W_6 = W_2 \cdot (1 - W_2) = 0 \\
 W_7 = W_3 \cdot (1 - W_3) = 0 \\
 W_8 = W_4 \cdot (1 - W_4) = 0 \\
 W_9 = W_1 \cdot W_2 = 0 \\
 W_{10} = W_3 \cdot (1 - W_4) = 1 \\
 I_1 = (W_1 + W_2 - W_9) \cdot W_{10} = 1
 \end{array}$$

5322 A constructive proof of knowledge of a witness, for instance,  $I_1 = 1$ , is therefore given by the  
 5323 tuple  $P = (W_5, W_6, W_7, W_8, W_9, W_{10}) = (0, 0, 0, 0, 0, 1)$ .

## Arrays

The `array` type represents a fixed-size collection of elements of equal type, each selectable by one or more indices that can be computed at run time during program execution.

Arrays are a classical type, implemented by many higher programming languages that compile to circuits or rank-1 constraint systems. However, most high-level circuit languages support **static** arrays, i.e., arrays whose length is known at compile time only.

The most common way to compile arrays to circuits is to transform any array of a given type  $\tau$  and size  $N$  into  $N$  circuit variables of type  $\tau$ . Arrays are therefore **syntactic sugar**, that is, parts of the formal language that makes the code easier for humans to read, which the compiler transforms into input nodes, much like any other variable. In PAPER, we define the following array type:

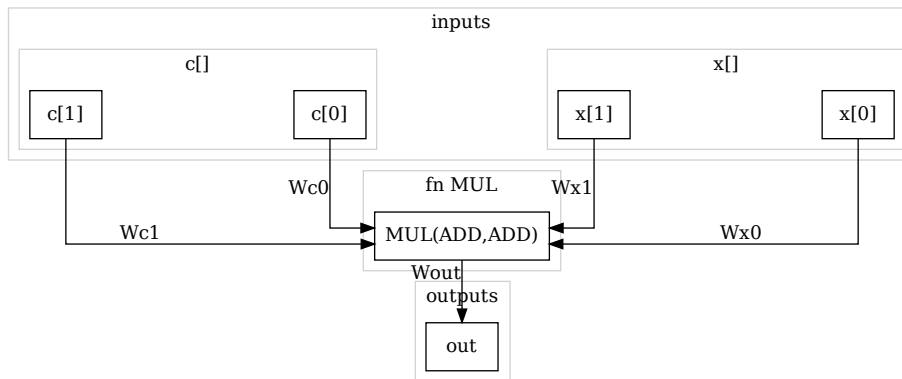
```
type <Name>: <Type>[N : unsigned] -> (Type,...) {
  return (<Name>[0],...)
}
```

In the setup phase of a statement, we compile every occurrence of an array of size  $N$  that contains elements of type `Type` into  $N$  variables of type `Type`.

*Example 130.* To give an intuition of how a real-world compiler might transform arrays into circuit variables, consider the following PAPER statement:

```
statement ARRAY_TYPE {F: F_5} {
  fn main(x: F[2])-> F {
    let constant c: F[2] = [2,4] ;
    let out:F <== MUL(ADD(x[1],c[0]),ADD(x[0],c[1])) ;
    return out ;
  } ;
}
```

During a setup phase, a circuit compiler might then replace any occurrence of the array type by a tuple of variables of the underlying type, and then use those variables in the circuit synthesis process instead. To see how this can be achieved, we use PAPER as an example. Abstracting over the sub-circuit of the computation, we get the following circuit:



## The Unsigned Integer Type

Unsigned integers of size  $N$ , where  $N$  is usually a power of two, represent non-negative integers in the range  $0 \dots 2^N - 1$ . They have a notion of addition, subtraction and multiplication, defined



5357 by modular  $2^N$  arithmetics. If some  $N$  is given, we write  $uN$  for the associated type.

5358 **The  $uN$  Constraint System** Many high-level circuit languages define the the various  $uN$  types  
 5359 as arrays of size  $N$ , where each element is of boolean type. This is similar to their representa-  
 5360 tion on common computer hardware and allows for efficient and straightforward definition of  
 5361 common operators, like the various **shift**, or logical operators.

shift

5362 If some unsigned integer  $N$  is known at compile time in PAPER, we define the following  $uN$   
 5363 type:

```
5364 type uN -> BOOL[N] {
5365   let base2 : BOOL[N] <== BASE_2 (uN) ;
5366   return base2 ;
5367 }
```

To enforce an  $N$ -tuple of field elements  $(b_0, \dots, b_{N-1})$  to represent an element of type  $uN$  we therefore need  $N$  boolean constraints

$$\begin{aligned} S_0 \cdot (1 - S_0) &= 0 \\ S_1 \cdot (1 - S_1) &= 0 \\ &\dots \\ S_{N-1} \cdot (1 - S_{N-1}) &= 0 \end{aligned}$$

5368 In the setup phase of a statement, we compile every occurrence of the  $uN$  type by a size  $N$  array  
 5369 of boolean type. During a=the prover phase, actual elements of the  $uN$  type are first transformed  
 5370 into binary representation and then this binary representation is assigned to the boolean array  
 5371 that represents the  $uN$  type.

5372 *Remark 4.* Representing the  $uN$  type as boolean arrays is conceptually clean and works over  
 5373 generic base fields. However, representing unsigned integers in this way requires a lot of space  
 5374 as every bit is represented as a field element and if the base field is large, those field elements  
 5375 require considerable space in hardware.

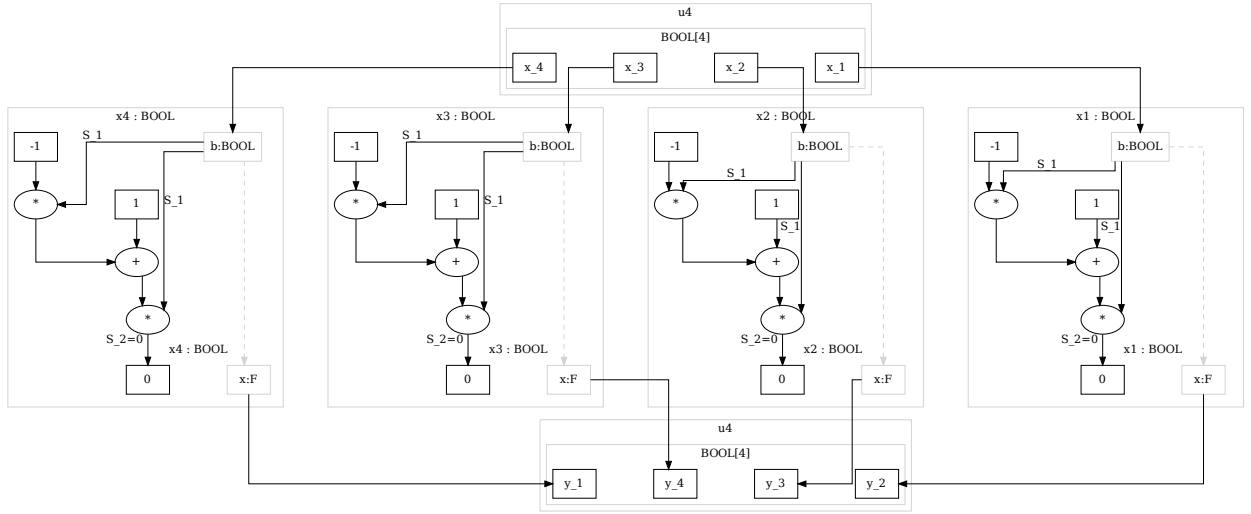
5376 It should be noted that, in some cases, there is another, more space- and constraint-efficient  
 5377 approach for representing unsigned integers that can be used whenever the underlying base field  
 5378 is sufficiently large. To understand this, recall that addition and multiplication in a prime field  
 5379  $\mathbb{F}_p$  is equal to addition and multiplication of integers, as long as the sum or the product does  
 5380 not exceed the modulus  $p$ . It is therefore possible to represent the  $uN$  type inside the base-field  
 5381 type whenever  $N$  is small enough. In this case, however, care has to be taken to never overflow  
 5382 the modulus. It is also important to make sure that, in the case of subtraction, the subtrahend is  
 5383 never larger than the minuend.

5384 *Example 131.* To give an intuition of how a real-world compiler might transform unsigned  
 5385 integers into circuit variables, consider the following PAPER statement:

```
5386 statement RING_SHIFT{F: F_p, N=4} {
5387   fn main(x: uN) -> uN {
5388     let y:uN <== [x[1], x[2], x[3], x[0]] ;
5389     return y ;
5390   } ;
5391 }
```

5392 During the setup-phase, a circuit compiler might then replace any occurrence of the  $uN$  type  
 5393 by  $N$  variables of `boolean` type. Using the definition of booleans, each of these variables is

then transformed into the `field` type and a boolean constraint system. To see how this can be achieved, we use PAPER as an example and get the following circuit:



During the prover phase, the function `main` is called with an actual input of `u4` type, say  $x=14$ . The high-level language then has to transform the decimal value 14 into its 4-bit binary representation  $14_2 = (0, 1, 1, 1)$  outside of the circuit. Then the array of field values  $x[4] = [0, 1, 1, 1]$  is used as an input to the circuit. Since all 4 field elements are either 0 or 1, the four boolean constraints are satisfiable and the output is an array of the four field elements  $[1, 1, 1, 0]$ , which represents the `u4` element 7.

**The Unigned Integer Operators** Since elements of `uN` type are represented as boolean arrays, shift operators are implemented in circuits simply by rewiring the boolean input variables to the output variables accordingly.

Logical operators, like AND, OR, or NOT are defined on the `uN` type by invoking the appropriate boolean operators bitwise to every bit in the boolean array that represents the `uN` element.

Addition and multiplication can be represented similarly to how machines represent those operations. Addition can be implemented by first defining the **full adder** circuit and then combining  $N$  of these circuits into a circuit that adds two elements from the `uN` type.

**Exercise 49.** Let  $F = \mathbb{F}_{13}$  and  $N=4$  be fixed. Define circuits and associated R1CS for the left and right **bishift** operators  $x \ll 2$  as well as  $x \gg 2$  that operate on the `uN` type. Execute the associated circuit for  $x : u4 = 11$ .

**Exercise 50.** Let  $F = \mathbb{F}_{13}$  and  $N=2$  be fixed. Define a circuit and associated R1CS for the addition operator  $\text{ADD} : F \times F \rightarrow F$ . Execute the associated circuit to compute  $\text{ADD}(2, 7)$ .

**Exercise 51.** Brain-compile the following PAPER code into a circuit and derive the associated R1CS.

```
statement MASK_MERGE {F:F_5, N=4} {
  fn main(pub a : uN, pub b : uN) -> F {
    let constant mask : uN = 10 ;
    let r : uN <== XOR(a, AND(XOR(a, b), mask)) ;
    return r ;
  }
}
```

```
5425     }
5426 }
```

5427 Let  $L_{mask\_merge}$  be the language defined by the circuit. Provide a constructive knowledge proof  
 5428 in  $L_{mask\_merge}$  for the instance  $I = (I_a, I_b) = (14, 7)$ .

## 5429 7.2.2 Control Flow

5430 Most programming languages of the imperative or functional style have some notion of basic  
 5431 control structures to direct the order in which instructions are evaluated. Contemporary circuit  
 5432 compilers usually provide a single thread of execution and provide basic flow constructs that  
 5433 implement control flow in circuits.

### 5434 The Conditional Assignment

5435 Writing high-level code that compiles to circuits, it is often necessary to have a way for condi-  
 5436 tional assignment of values or computational output to variables.

5437 One way to realize this in many programming languages is in terms of the conditional  
 5438 ternary assignment operator  $?:$  that branches the control flow of a program according to some  
 5439 condition and then assigns the output of the computed branch to some variable.

```
5440 variable = condition ? value_if_true : value_if_false
```

5441 In this description, `condition` is a boolean expression and `value_if_true` as well as  
 5442 `value_if_false` are expressions that evaluate to the same type as `variable`.

5443 In programming languages like Rust, another way to write the conditional assignment oper-  
 5444 ator that is more familiar to many programmers is given by

```
5445 variable = if condition then {
5446     value_if_true
5447 } else {
5448     value_if_false
5449 }
```

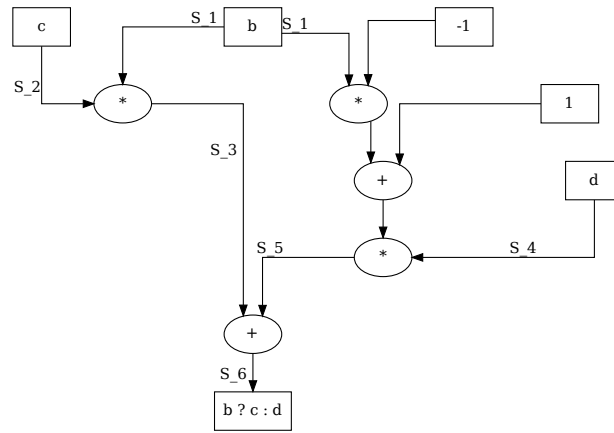
5450 In most programming languages, it is a key property of the ternary assignment operator that  
 5451 the expression `value_if_true` is only evaluated if `condition` evaluates to true and the  
 5452 expression `value_if_false` is only evaluated if `condition` evaluates to false. In fact,  
 5453 computer programs would turn out to be very inefficient if the ternary operator would evaluate  
 5454 both expressions regardless of the value of `condition`.

5455 A simple way to implement conditional assignment operator as a circuit can be achieved  
 5456 if the requirement that only one branch of the conditional operator is executed is dropped. To  
 5457 see that, let  $b$ ,  $c$  and  $d$  be field elements such that  $b$  is a boolean constraint. In this case, the  
 5458 following equation enforces a field element  $x$  to be the result of the conditional assignment  
 5459 operator:

$$x = b \cdot c + (1 - b) \cdot d \quad (7.6)$$

5460 Expressing this equation in terms of the addition and multiplication operators from XXX, we  
 5461 can flatten it into the following algebraic circuit:

add refer-  
ence



5462

5463 Note that, in order to compute a valid assignment to this circuit, both  $S_2$  as well as  $S_4$  are  
 5464 necessary. If the inputs to the nodes  $c$  and  $d$  are circuits themselves, both circuits need valid  
 5465 assignments and therefore have to be executed. As a consequence, this implementation of  
 5466 the conditional assignment operator has to execute all branches of all circuits, which is very  
 5467 different from the execution of common computer programs and contributes to the increased  
 5468 computational effort any prover has to invest, in contrast to the execution in other programming  
 5469 models.

We can use the general technique from 6.2.1.1 to derive the associated rank-1 constraint system of the conditional assignment operator. We get the following:

check  
reference

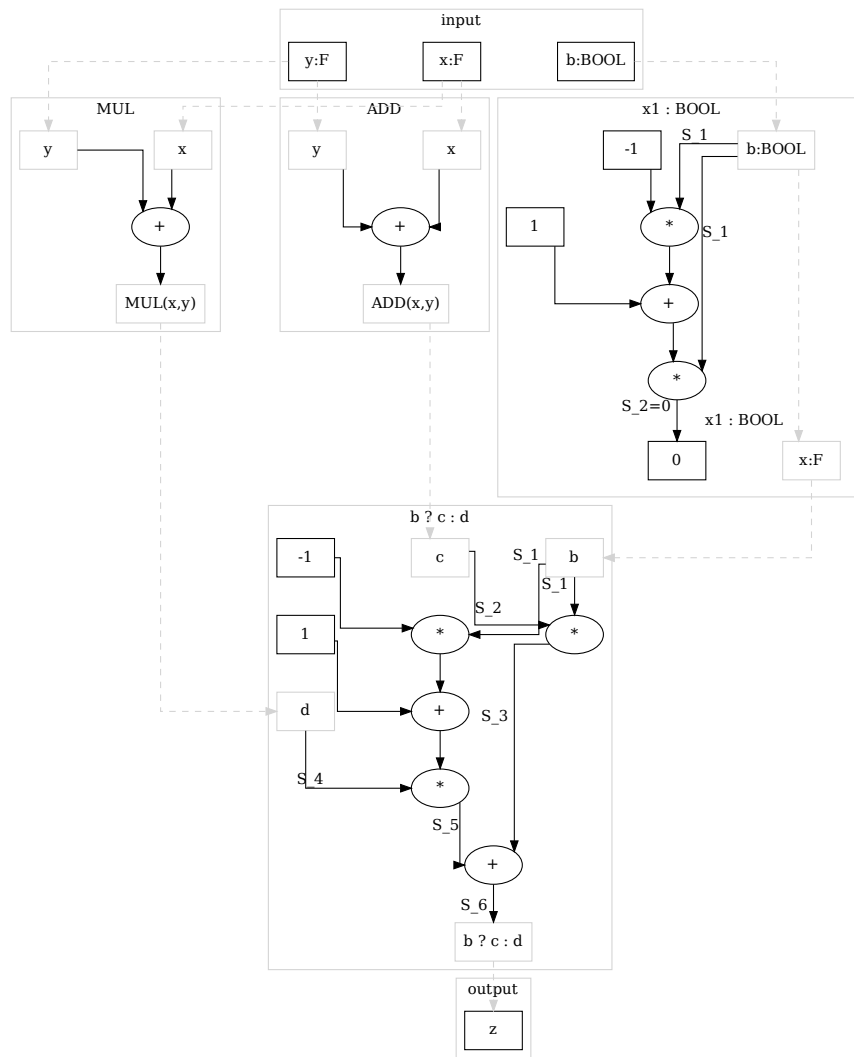
$$\begin{aligned} S_1 \cdot S_2 &= S_3 \\ (1 - S_1) \cdot S_4 &= S_5 \\ (S_3 + S_5) \cdot 1 &= S_6 \end{aligned}$$

5470 *Example 132.* To give an intuition of how a real-world circuit compiler might transform any  
 5471 high-level description of the conditional assignment operator into a circuit, consider the follow-  
 5472 ing PAPER code:

```

5473 statement CONDITIONAL_OP {F:F_p} {
5474   fn main(x : F, y : F, b : BOOL) -> F {
5475     let z : F <== if b then {
5476       ADD(x,y)
5477     } else {
5478       MUL(x,y)
5479     } ;
5480     return z ;
5481   }
5482 }
```

5483 Brain-compiling this code into a circuit, we first draw box nodes for all input and output vari-  
 5484 ables, and then transform the boolean type into the field type together with its associated con-  
 5485 straint. Then we evaluate the assignments to the output variables. Since the conditional assign-  
 5486 ment operator is the top level function, we draw its circuit and then draw the circuits for both  
 5487 conditional expressions. We get the following:



5488

5489 **Loops**

5490 In many programming languages, various loop control structures are defined that allow devel-  
 5491 opers to execute expressions with a specified number of repetitions or arguments. In particular,  
 5492 it is often possible to implement unbounded loops like the loop structure give below, where the  
 5493 number of executions depends on execution inputs and is therefore unknown at compile time:

5494 **while** true do {  
 5495

5496 **M:** Add another example where the loop actually depends on the input.

Add ex-  
ample

5497 In contrast to this, it should be noted that algebraic circuits and rank-1 constraint systems  
 5498 are not general enough to express arbitrary computation, but bounded computation only. As a  
 5499 consequence, it is not possible to implement unbounded loops, or loops with bounds that are  
 5500 unknown at compile time in those models. This can be easily seen since circuits are acyclic  
 5501 by definition, and implementing an unbounded loop as an acyclic graph requires a circuits of  
 5502 unbounded size.

5503 However, circuits are general enough to express bounded loops, where the upper bound on  
 5504 its execution is known at compile time. Those loop can be implemented in circuits by enrolling  
 5505 the loop.

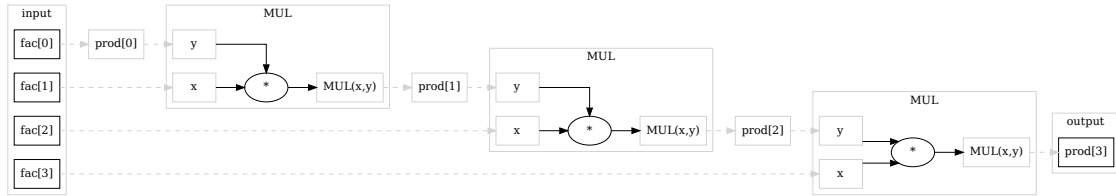
As a consequence, any programming language that compiles to algebraic circuits can only provide loop structures where the bound is a constant known at compile time. This implies that loops cannot depend on execution inputs, but on compile time parameters only.

*Example 133.* To give an intuition of how a real-world circuit compiler might transform any high-level description of a bounded `for` loop into a circuit, consider the following PAPER code:

```
statement FOR_LOOP {F:F_p, N: unsigned = 4} {
  fn main(fac : F[N]) -> F {
    let prod[N] : F ;
    prod[0] <== fac[0] ;
    for unsigned i in 1..N do [{
      prod[i] <== MUL(fac[i], prod[i-1]) ;
    }]
    return prod[N] ;
  }
}
```

Note that, in a program like this, the loop counter `i` has no expression in the derived circuit. It is pure syntactic sugar, telling the compiler how to unroll the loop.

Brain-compiling this code into a circuit, we first draw box nodes for all input and output variables, noting that the loop counter is not represented in the circuit. Since all variables are of `field` type, we don't have to compile any type constraints. Then we evaluate the assignments to the output variables by unrolling the loop into 3 individual assignment operators. We get:



### 7.2.3 Binary Field Representations

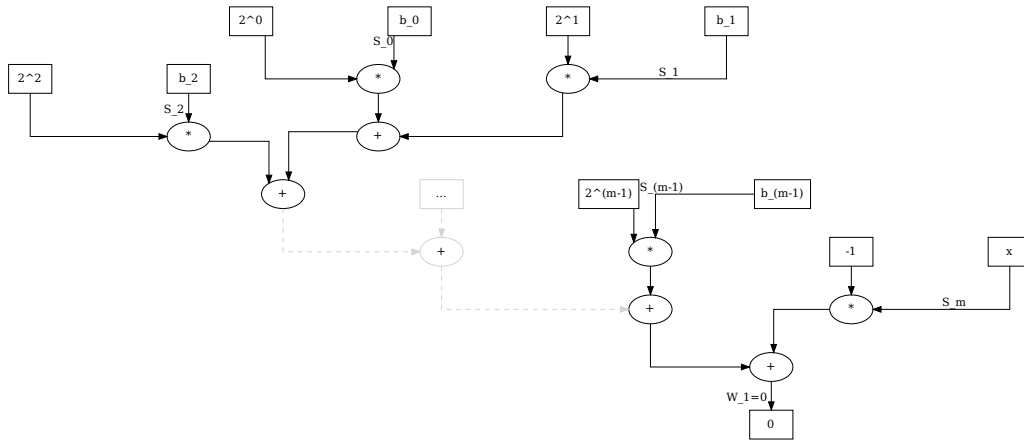
In applications, it is often necessary to enforce a binary representation of elements from the `field` type. To derive an appropriate circuit over a prime field  $\mathbb{F}_p$ , let  $m = \lceil \log_2 p \rceil$  be the smallest number of bits necessary to represent the prime modulus  $p$ . Then a bitstring  $(b_0, \dots, b_{m-1}) \in \{0, 1\}^m$  is a binary representation of a field element  $x \in \mathbb{F}_p$ , if and only if

$$x = b_0 \cdot 2^0 + b_1 \cdot 2^1 + \dots + b_{m-1} \cdot 2^{m-1}$$

In this expression, addition and exponentiation is considered to be executed in  $\mathbb{F}_p$ , which is well defined since all terms  $2^j$  for  $0 \leq j < m$  are elements of  $\mathbb{F}_p$ . Note, however, that in contrast to the binary representation of unsigned integers  $n \in \mathbb{N}$ , this representation is not unique in general, since the modular  $p$  equivalence class might contain more than one binary representative.

Considering that the underlying prime field is fixed and the most significant bit of the prime modulus is  $m$ , the following circuit flattens equation XXX, assuming all inputs  $b_1, \dots, b_m$  are of boolean type.

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Applying the general transformation rule to compute the associated rank-1 constraint systems, we see that we actually only need a single constraint to enforce some binary representation of any field element. We get

$$(S_0 \cdot 2^0 + S_1 \cdot 2^1 + \dots + S_{m-1} \cdot 2^{m-1} - S_m) \cdot 1 = 0$$

5538 Given an array `BOOL[N]` of  $N$  boolean constraint field elements and another field element  $x$ ,  
 5539 the circuit enforces `BOOL[N]` to be one of the binary representations of  $x$ . If `BOOL[N]` is not  
 5540 a binary representation of  $x$ , no valid assignment and hence no solution to the associated R1CS  
 5541 can exist.

5542 *Example 134.* Consider the prime field  $\mathbb{F}_{13}$ . To compute binary representations of elements  
 5543 from that field, we start with the binary representation of the prime modulus 13, which is  $|13|_2 =$   
 5544  $(1, 0, 1, 1)$  since  $13 = 1 \cdot 2^0 + 0 \cdot 2^1 + 1 \cdot 2^2 + 1 \cdot 2^3$ . So  $m = 4$  and we need up to 4 bits to represent  
 5545 any element  $x \in \mathbb{F}_{13}$ .

To see that binary representations are not unique in general, consider the element  $2 \in \mathbb{F}_{13}$ . It has the binary representations  $|2|_2 = (0, 1, 0, 0)$  and  $|2|_2 = (1, 1, 1, 1)$ , since in  $\mathbb{F}_{13}$  we have

$$2 = \begin{cases} 0 \cdot 2^0 + 1 \cdot 2^1 + 0 \cdot 2^2 + 0 \cdot 2^3 \\ 1 \cdot 2^0 + 1 \cdot 2^1 + 1 \cdot 2^2 + 1 \cdot 2^3 \end{cases}$$

5546 This is because the unsigned integers 2 and 15 are both in the modular 13 remainder class of 2  
 5547 and hence are both representatives of 2 in  $\mathbb{F}_{13}$ .

To see how circuit XXX works, we want to enforce the binary representation of  $7 \in \mathbb{F}_{13}$ . Since  $m = 4$  we have to enforce a 4-bit representation for 7, which is  $(1, 1, 1, 0)$ , since  $7 = 1 \cdot 2^0 + 1 \cdot 2^1 + 1 \cdot 2^2 + 0 \cdot 2^3$ . A valid circuit assignment is therefore given by  $(S_0, S_1, S_2, S_3, S_4) = (1, 1, 1, 0, 7)$  and, indeed, the assignment satisfies the required 5 constraints including the 4 boolean constraints for  $S_0, \dots, S_3$ :

$$\begin{aligned} 1 \cdot (1 - 1) &= 0 & // \text{boolean constraints} \\ 1 \cdot (1 - 1) &= 0 \\ 1 \cdot (1 - 1) &= 0 \\ 0 \cdot (1 - 0) &= 0 \\ (1 + 2 + 4 + 0 - 7) \cdot 1 &= 0 & // \text{binary rep. constraint} \end{aligned}$$

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$$\begin{aligned} S_1 \cdot S_1 &= S_3 \\ S_2 \cdot S_2 &= S_4 \\ (S_4 \cdot 8) \cdot S_3 &= S_5 \\ (0 \cdot S_3 + 1) \cdot 1 &= 0 \end{aligned}$$

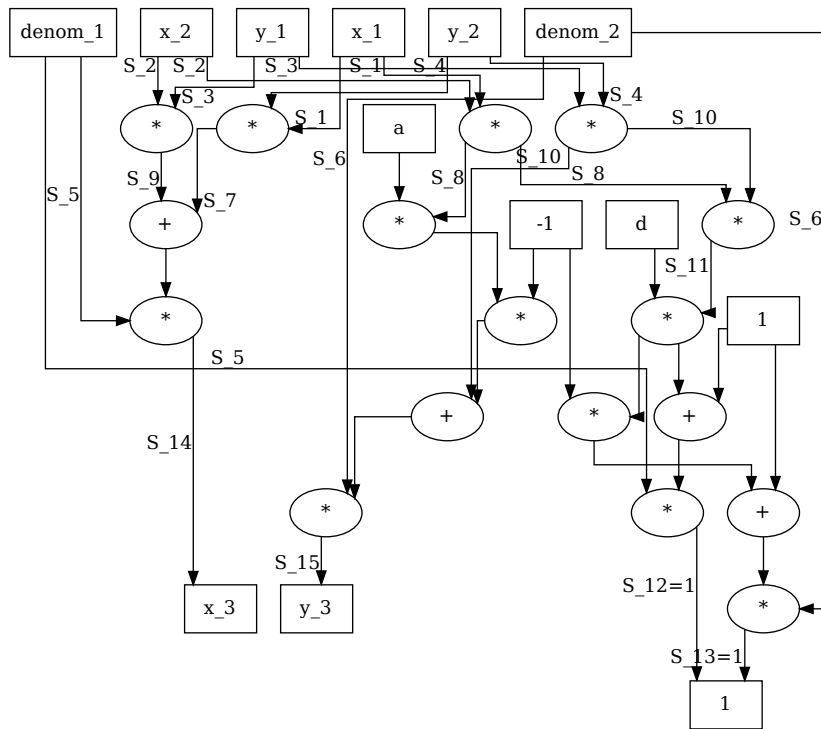
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**Twisted Edwards curve addition** As we have seen in 5.1.3, a major advantage of twisted Edwards curves is the existence of an addition law that contains no branching and is valid for all curve points. Moreover, the neutral element is not “at infinity” but the actual curve point  $(0, 1)$ . In fact, given two points  $(x_1, y_1)$  and  $(x_2, y_2)$  on a twisted Edwards curve, their sum is defined as

$$(x_3, y_3) = \left( \frac{x_1 y_2 + y_1 x_2}{1 + d \cdot x_1 x_2 y_1 y_2}, \frac{y_1 y_2 - a \cdot x_1 x_2}{1 - d \cdot x_1 x_2 y_1 y_2} \right)$$

5570 We can use the division circuit from XXX to flatten this equation into an algebraic circuit.  
 5571 Inputs to the circuit are then not only the two curve points  $(x_1, y_1)$  and  $(x_2, y_2)$ , but also the two  
 5572 denominators  $denum_1 = 1 + d \cdot x_1 x_2 y_1 y_2$  as well as  $denum_2 = 1 - d \cdot x_1 x_2 y_1 y_2$ , which any prover  
 5573 needs to compute outside of the circuit. We get



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Using the general technique from XXX to derive the associated rank-1 constraint system, we get the following result:

$$\begin{aligned} S_1 \cdot S_4 &= S_7 \\ S_1 \cdot S_2 &= S_8 \\ S_2 \cdot S_3 &= S_9 \\ S_3 \cdot S_4 &= S_{10} \\ S_8 \cdot S_{10} &= S_{11} \\ S_5 \cdot (1 + d \cdot S_{11}) &= 1 \\ S_6 \cdot (1 - d \cdot S_{11}) &= 1 \\ S_5 \cdot (S_9 + S_7) &= S_{14} \\ S_6 \cdot (S_{10} - a \cdot S_8) &= S_{15} \end{aligned}$$

5575 *Exercise 53.* Let  $\mathbb{F}$  be a field. Define a circuit that enforces field inversion for a point of a  
5576 twisted Edwards curve over  $\mathbb{F}$ .

## Chapter 8

# Zero Knowledge Protocols

A so-called **zero-knowledge protocol** is a set of mathematical rules by which one party (usually called **the prover**) can convince another party (usually called **the verifier**) that a given statement is true, while not revealing any additional information apart from the truth of the statement.

As we have seen in chapter 6, given some language  $L$  and instance  $I$ , the knowledge claim “there is a witness  $W$  such that  $(I; W)$  is a word in  $L$ ” is constructively provable by providing  $W$  to the verifier. However, the challenge for a zero-knowledge protocol is to prove knowledge of a witness without revealing any information beyond its bare existence.

In this chapter, we look at various systems that exist to solve this task. We start with an introduction to the basic concepts and terminology in zero-knowledge proving systems and then introduce the so-called Groth\_16 protocol as one of the most efficient systems. We plan to update this chapter with new inventions in future versions of this book.

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## 8.1 Proof Systems

From an abstract point of view, a proof system is a set of rules which models the generation and exchange of messages between two parties: a prover and a verifier. Its task is to ascertain whether a given string belongs to a formal language or not.

Proof systems are often classified by certain trust assumptions and the computational capabilities of both parties. In its most general form, the prover usually possesses unlimited computational resources but cannot be trusted, while the verifier has bounded computation power but is assumed to be honest.

Proving the membership statement for some string is executed by the generation of certain messages that are sent between prover and verifier, until the verifier is convinced that the string is an element of the language in consideration.

To be more specific, let  $\Sigma$  be an alphabet, and let  $L$  be a formal language defined over  $\Sigma$ . Then a **proof system** for language  $L$  is a pair of probabilistic interactive algorithms  $(P, V)$ , where  $P$  is called the **prover** and  $V$  is called the **verifier**.

Both algorithms are able to send messages to one another, and each algorithm has its own state, some shared initial state and access to the messages. The verifier is bounded to a number of steps which is polynomial in the size of the shared initial state, after which it stops and outputs either `accept` or `reject` indicating that it accepts or rejects a given string to be in  $L$ . In contrast, there are bounds on the computational power of the prover.

When the execution of the verifier algorithm stops the following conditions are required to hold:

- (Completeness) If the tuple  $x \in \Sigma^*$  is a word in language  $L$  and both prover and verifier follow the protocol, the verifier outputs `accept`.
- (Soundness) If the tuple  $x \in \Sigma^*$  is not a word in language  $L$  and the verifier follows the protocol, the verifier outputs `reject`, except with some small probability.

In addition, a proof system is called **zero-knowledge** if the verifier learns nothing about  $x$  other than  $x \in L$ .

The previous definition of proof systems is very general, and many subclasses of proving systems are known in the field. The type of languages that a proof system can support crucially depends on the abilities of the verifier (for example, whether it can make random choices) or on the nature and number of the messages that can be exchanged. If the system only requires to send a single message from the prover to the verifier, the proof system is called **non-interactive**, because no interaction other than sending the actual proof is required. In contrast, any other proof system is called **interactive**.

A proof system is usually called **succinct** if the size of the proof is shorter than the witness necessary to generate the proof. Moreover, a proof system is called **computationally sound** if soundness only holds under the assumption that the computational capabilities of the prover are polynomially bound. To distinguish general proofs from computationally sound proofs, the latter are often called **arguments**. Zero-knowledge, succinct, non-interactive arguments of knowledge claims are often abbreviated **zk-SNARKs**.

*Example 135 (Constructive Proofs for Algebraic Circuits).* To formalize our previous notion of constructive proofs for algebraic circuits, let  $\mathbb{F}$  be a finite field, and let  $C(\mathbb{F})$  be an algebraic circuit over  $\mathbb{F}$  with associated language  $L_{C(\mathbb{F})}$ . A non-interactive proof system for  $L_{C(\mathbb{F})}$  is given by the following two algorithms:

Given some instance  $I$ , the prover algorithm  $P$  uses its unlimited computational power to compute a witness  $W$  such that the pair  $(I; W)$  is a valid assignment to  $C(\mathbb{F})$  whenever the circuit is satisfiable for  $I$ . The prover then sends the constructive proof  $(I; W)$  to the verifier.

On receiving a message  $(I; W)$ , the verifier algorithm  $V$  assigns the constructive proof  $(I; W)$  to circuit  $C(\mathbb{F})$ , and decides whether the assignment is valid by executing all gates in the circuit. The runtime is polynomial in the number of gates. If the assignment is valid, the verifier returns `accepts`, if not, it returns `reject`.

To see that this proof system has the completeness and soundness properties, let  $C(\mathbb{F})$  be a circuit of the field  $\mathbb{F}$ , and let  $I$  be an instance. The circuit may or may not have a witness  $W$  such that  $(I; W)$  is a valid assignment to  $C(\mathbb{F})$ .

If no  $W$  exists,  $I$  is not part of any word in  $L_{C(\mathbb{F})}$ , and there is no way for  $P$  to generate a valid assignment. It follows that the verifier will not accept any claimed proof sent by  $P$ , which implies that the system has **soundness**.

If, on the other hand,  $W$  exists and  $P$  is honest,  $P$  can use its unlimited computational power to compute  $W$  and send  $(I; W)$  to  $V$ , which  $V$  will accept in polynomial time. This implies that the system has **completeness**.

The system is non-interactive because the prover only sends a single message to the verifier, which contains the proof itself. Because in this simple system the witness itself is the proof, the proof system is **not** succinct.

## 8.2 The “Groth16” Protocol

In chapter 6, we have introduced algebraic circuits, their associated rank-1 constraint systems and their induced quadratic arithmetic programs. These models define formal languages, and

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associated memberships and knowledge claims can be constructively proofed by executing the circuit to compute a solution to its associated R1CS. The solution can then be transformed into a polynomial such that the polynomial is divisible by another polynomial if and only if the solution is correct.

In Groth [2016], Jens Groth provides a method that can transform those proofs into zero-knowledge succinct non-interactive arguments of knowledge. Assuming that the pairing groups  $(\mathbb{G}_1, \mathbb{G}_2, \mathbb{G}_T, b)$  are given, the arguments are of constant size and consist of 2 elements from  $\mathbb{G}_1$  and a single element from  $\mathbb{G}_2$ , regardless of the size of the witness. They are zero-knowledge in the sense that the verifier learns nothing about the witness beside the fact that the instance-witness pair is a proper word in the language of the problem.

Verification is non-interactive and needs to compute a number of exponentiations proportional to the size of the instance, together with 3 group pairings in order to check a single equation.

The generated argument is perfectly zero-knowledge has perfect completeness and soundness in the generic bilinear group model, assuming the existence of a trusted third party that executes a preprocessing phase to generate a **common reference string** and a **simulation trapdoor**. This party must be trusted to delete the simulation trapdoor, since everyone in possession of it can simulate proofs.

To be more precise, let  $R$  be a rank-1 constraint system defined over some finite field  $\mathbb{F}_r$ . Then **Groth\_16 parameters** for  $R$  are given by the following set:

$$\text{Groth}_{16} - \text{Param}(R) = (r, \mathbb{G}_1, \mathbb{G}_2, e(\cdot, \cdot), g_1, g_2) \quad (8.1)$$

In the equation above,  $\mathbb{G}_1$  and  $\mathbb{G}_2$  are finite cyclic groups of order  $r$ ,  $g_1$  is a generator of  $\mathbb{G}_1$ ,  $g_2$  is a generator of  $\mathbb{G}_2$  and  $e : \mathbb{G}_1 \times \mathbb{G}_2 \rightarrow \mathbb{G}_T$  is a non-degenerate, bilinear pairing for some target group  $\mathbb{G}_T$ . In real-world applications, the parameter set is usually agreed on in advance.

Given some Groth\_16 parameters, a **Groth\_16 protocol** is then a quadruple of probabilistic polynomial algorithms (SETUP, PROVE, VFY, SIM) such that the following conditions hold:

- (Setup-Phase):  $(CRS, \tau) \leftarrow \text{Setup}(R)$ : Algorithm Setup takes the R1CS  $R$  as input and computes a common reference string  $CRS$  and a simulation trapdoor  $\tau$ .
- (Prover-Phase):  $\pi \leftarrow \text{Prove}(R, CRS, I, W)$ : Given a constructive proof  $(I, W)$  for  $R$ , algorithm Prove takes the R1CS  $R$ , the common reference string  $CRS$  and the constructive proof  $(I, W)$  as input and computes an zk-SNARK  $\pi$ .
- Verify:  $\{\text{accept}, \text{reject}\} \leftarrow \text{Vfy}(R, CRS, I, \pi)$ : Algorithm Vfy takes the R1CS  $R$ , the common reference string  $CRS$ , the instance  $I$  and the zk-SNARK  $\pi$  as input and returns `reject` or `accept`.
- $\pi \leftarrow \text{Sim}(R, \tau, CRS, I)$ : Algorithm Sim takes the R1CS  $R$ , the common reference string  $CRS$ , the simulation trapdoor  $\tau$  and the instance  $I$  as input and returns a zk-SNARK  $\pi$ .

We will explain these algorithms together with detailed examples in the remainder of this section.

Assuming a trusted third party for the setup, the protocol is able to compute a zk-SNARK from a constructive proof for  $R$ , provided that  $r$  is sufficiently large, and, in particular, larger than the number of constraints in the associated R1CS.

**Example 136** (The 3-Factorization Problem). Consider the 3-factorization problem from 106 and its associated algebraic circuit and rank-1 constraint system from 6.8. In this example,

we want to agree on a parameter set  $(R, r, \mathbb{G}_1, \mathbb{G}_2, e(\cdot, \cdot), g_1, g_2)$  in order to use the Groth\_16 protocol for our 3-factorization problem.

To find proper parameters, first observe that the circuit XXX, as well as its associated R1CS  $R_{3, \text{fac\_zk}}$  ex:3-factorization-r1cs and the derived QAP 6.14, are defined over the field  $\mathbb{F}_{13}$ . We therefore have  $r = 13$  and need pairing groups  $\mathbb{G}_1$  and  $\mathbb{G}_2$  of order 13.

We know from 5.4 that the moon-math curve BLS6\_6 has two subgroups  $\mathbb{G}_1[13]$  and  $\mathbb{G}_2[13]$ , which are both of order 13. The associated Weil pairing  $b$  (5.45) is a proper bilinear map. We therefore choose those groups and the Weil pairing together with the generators  $g_1 = (13, 15)$  and  $g_2 = (7v^2, 16v^3)$  of  $\mathbb{G}_1[13]$  and  $\mathbb{G}_2[13]$ , as a parameter:

$$\text{Groth\_16} - \text{Param}(R_{3, \text{fac\_zk}}) = (r, \mathbb{G}_1[13], \mathbb{G}_2[13], e(\cdot, \cdot), (13, 15), (7v^2, 16v^3))$$

It should be noted that our choice is not unique. Every pair of finite cyclic groups of order 13 that has a proper bilinear pairing qualifies as a Groth\_16 parameter set. The situation is similar to real-world applications, where SNARKs with equivalent behavior are defined over different curves, used in different applications.

**The Setup Phase** To generate zk-SNARKs from constructive knowledge proofs in the Groth16 protocol, a preprocessing phase is required. This has to be executed a single time for every rank-1 constraint system and any associated quadratic arithmetic program. The outcome of this phase is a common reference string that prover and verifier need in order to generate and verify the zk-SNARK. In addition, a simulation trapdoor is produced that can be used to simulate proofs.

To be more precise, let  $L$  be a language defined by some rank-1 constraint system  $R$  such that a constructive proof of knowledge for an instance  $(I_1, \dots, I_n)$  in  $L$  consists of a witness  $(W_1, \dots, W_m)$ . Let  $QAP(R) = \{T \in \mathbb{F}[x], \{A_j, B_j, C_j \in \mathbb{F}[x]\}_{j=0}^{n+m}\}$  be a quadratic arithmetic program associated to  $R$ , and let  $\{\mathbb{G}_1, \mathbb{G}_2, e(\cdot, \cdot), g_1, g_2, \mathbb{F}_r\}$  be the set of Groth\_16 parameters.

The setup phase then samples 5 random, **inverible** elements  $\alpha, \beta, \gamma, \delta$  and  $s$  from the scalar field  $\mathbb{F}_r$  of the protocol and outputs the **simulation trapdoor**  $\tau$ :

$$\tau = (\alpha, \beta, \gamma, \delta, s) \quad (8.2)$$

In addition, the setup phase uses those 5 random elements together with the two generators  $g_1$  and  $g_2$  and the quadratic arithmetic program to generate a **common reference string**  $CRS_{QAP} = (CRS_{\mathbb{G}_1}, CRS_{\mathbb{G}_2})$  of language  $L$ :

$$CRS_{\mathbb{G}_1} = \left\{ g_1^\alpha, g_1^\beta, g_1^\delta, \left( g_1^{s^j}, \dots \right)_{j=0}^{\deg(T)-1}, \left( g_1^{\frac{\beta \cdot A_j(s) + \alpha \cdot B_j(s) + C_j(s)}{\gamma}}, \dots \right)_{j=0}^n, \left( g_1^{\frac{\beta \cdot A_{j+n}(s) + \alpha \cdot B_{j+n}(s) + C_{j+n}(s)}{\delta}}, \dots \right)_{j=1}^m, \left( g_1^{\frac{s^j \cdot T(s)}{\delta}}, \dots \right)_{j=0}^{\deg(T)-2} \right\}$$

$$CRS_{\mathbb{G}_2} = \left\{ g_2^\beta, g_2^\gamma, g_2^\delta, \left( g_2^{s^j}, \dots \right)_{j=0}^{\deg(T)-1} \right\}$$

Common reference strings depend on the simulation trapdoor, and are therefore not unique to the problem. Any language can have more than one common reference string. The size of a common reference string is linear in the size of the instance and the size of the witness.

If a simulation trapdoor  $\tau = (\alpha, \beta, \gamma, \delta, s)$  is given, we call the element  $s$  a **secret evaluation point** of the protocol, because if  $\mathbb{F}_r$  is the scalar field of the finite cyclic groups  $\mathbb{G}_1$  and  $\mathbb{G}_2$ , then



5723 a key feature of any common reference string is that it provides data to compute the evaluation  
 5724 of any polynomial  $P \in \mathbb{F}_r[x]$  of degree  $\deg(P) < \deg(T)$  at the point  $s$  in the exponent of the  
 5725 generator  $g_1$  or  $g_2$ , without knowing  $s$ .

To be more precise, let  $s$  be the secret evaluation point and let  $P(x) = a_0 \cdot x^0 + a_1 \cdot x^1 + \dots a_k \cdot x^k$  be a polynomial of degree  $k < \deg(T)$  with coefficients in  $\mathbb{F}_r$ . Then we can compute  $g_1^{P(s)}$  without knowing what the actual value of  $s$  is:

$$\begin{aligned} g_1^{P(s)} &= g_1^{a_0 \cdot s^0 + a_1 \cdot s^1 + \dots a_k \cdot s^k} \\ &= g_1^{a_0 \cdot s^0} \cdot g_1^{a_1 \cdot s^1} \cdot \dots \cdot g_1^{a_k \cdot s^k} \\ &= \left(g_1^{s^0}\right)^{a_0} \cdot \left(g_1^{s^1}\right)^{a_1} \cdot \dots \cdot \left(g_1^{s^k}\right)^{a_k} \end{aligned}$$

5726 In this expression, all group points  $g_1^{s^j}$  are part of the common reference string, hence, they can  
 5727 be used to compute the result. The same holds true for the evaluation of  $g_2^{P(s)}$ , since the  $\mathbb{G}_2$  part  
 5728 of the common reference string contains the points  $g_2^{s^j}$ .

5729 In real-world applications, the simulation trapdoor is often called the **toxic waste** of the  
 5730 setup-phase, while a common reference string is also-called the pair of **prover and verifier**  
 5731 **key**.

explain why

5732 In order to make the protocol secure, the setup needs to be executed in a way that guarantees  
 5733 that the simulation trapdoor is deleted. Anyone in possession of it can generate arguments  
 5734 without knowledge of a constructive proof. The most simple approach to achieving deletion of  
 5735 the toxic waste is by a so-called **trusted third party**, where the trust assumption is that that  
 5736 the party generates the common reference string precisely as defined and deletes the simulation  
 5737 backdoor afterwards.

5738 However, as trusted third parties are not easy to find, more sophisticated protocols exist  
 5739 in real-world applications. They execute the setup phase as a multi party computation, where  
 5740 the proper execution can be publicly verified and the simulation trapdoor is deleted if at least  
 5741 one participant deletes their individual contribution to the randomness. Each participant only  
 5742 possesses a fraction of the simulation trapdoor, so it can only be recovered if all participants  
 5743 collude and share their fraction.

*Example 137* (The 3-factorization Problem). To see how the setup phase of a Groth\_16 zk-SNARK can be computed, consider the 3-factorization problem from 106 and the parameters from page 190. As we have seen in 6.14, an associated quadratic arithmetic program is given as follows:

$$\begin{aligned} QAP(R_{3, fac\_zk}) &= \{x^2 + x + 9, \\ &\quad \{0, 0, 6x + 10, 0, 0, 7x + 4\}, \{0, 0, 0, 6x + 10, 7x + 4, 0\}, \{0, 7x + 4, 0, 0, 0, 6x + 10\}\} \end{aligned}$$

4 ex-amples have the same title. Change it to be distinct

To transform this QAP into a common reference string, we choose the field elements  $\alpha = 6$ ,  $\beta = 5$ ,  $\gamma = 4$ ,  $\delta = 3$ ,  $s = 2$  from  $\mathbb{F}_{13}$ . In real-world applications, it is important to sample those values randomly from the scalar field, but in our approach, we choose those non-random values to make them more memorizable, which helps in pen-and-paper computations. Our simulation trapdoor is then given as follows:

$$\tau = (6, 5, 4, 3, 2)$$

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5744 We keep this secret in order to simulate proofs later on, but we are careful though to hide  $\tau$   
 5745 from anyone who hasn't read this book. Then we instantiate the common reference string XXX

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5746 from those values. Since our groups are subgroups of the  $\text{BLS}_{6\_6}$  elliptic curve, we use scalar  
5747 product notation instead of exponentiation.

To compute the  $\mathbb{G}_1$  part of the common reference string, we use the logarithmic order of the group  $\mathbb{G}_1$  XXX, the generator  $g_1 = (13, 15)$ , as well as the values from the simulation backdoor. Since  $\deg(T) = 2$ , we get the following:

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$$\begin{aligned} [\alpha]_{g_1} &= [6](13, 15) = (27, 34) \\ [\beta]_{g_1} &= [5](13, 15) = (26, 34) \\ [\delta]_{g_1} &= [3](13, 15) = (38, 15) \end{aligned}$$

To compute the rest of the  $\mathbb{G}_1$  part of the common reference string, we expand the indexed tuples and insert the secret random elements from the simulation backdoor. We get the following:

$$\begin{aligned} \left( [s^j]_{g_1}, \dots \right)_{j=0}^1 &= \left( [2^0](13, 15), [2^1](13, 15) \right) \\ &= \left( (13, 15), (33, 34) \right) \\ \left( \left[ \frac{\beta A_j(s) + \alpha B_j(s) + C_j(s)}{\gamma} \right]_{g_1}, \dots \right)_{j=0}^1 &= \left( \left[ \frac{5A_0(2) + 6B_0(2) + C_0(2)}{4} \right](13, 15), \right. \\ &\quad \left. \left[ \frac{5A_1(2) + 6B_1(2) + C_1(2)}{4} \right](13, 15) \right) \\ \left( \left[ \frac{\beta A_{j+n}(s) + \alpha B_{j+n}(s) + C_{j+n}(s)}{\delta} \right]_{g_1}, \dots \right)_{j=1}^4 &= \left( \left[ \frac{5A_2(2) + 6B_2(2) + C_2(2)}{3} \right](13, 15), \right. \\ &\quad \left[ \frac{5A_3(2) + 6B_3(2) + C_3(2)}{3} \right](13, 15), \\ &\quad \left[ \frac{5A_4(2) + 6B_4(2) + C_4(2)}{3} \right](13, 15), \\ &\quad \left. \left[ \frac{5A_5(2) + 6B_5(2) + C_6(2)}{3} \right](13, 15) \right) \\ \left( \left[ \frac{s^j \cdot T(s)}{\delta} \right]_{g_1} \right)_{j=0}^0 &= \left( \left[ \frac{2^0 \cdot T(2)}{3} \right](13, 15) \right) \end{aligned}$$

To compute the curve points on the right side of these expressions, we need the polynomials from the associated quadratic arithmetic program and evaluate them on the secret point  $s = 2$ .

Since  $4^{-1} = 10$  and  $3^{-1} = 9$  in  $\mathbb{F}_{13}$ , we get the following:

$$\left[ \frac{5A_0(2) + 6B_0(2) + C_0(2)}{4} \right](13, 15) = [(5 \cdot 0 + 6 \cdot 0 + 0) \cdot 10](13, 15) = [0](13, 15) =$$

$\emptyset$

$$\left[ \frac{5A_1(2) + 6B_1(2) + C_1(2)}{4} \right](13, 15) = [(5 \cdot 0 + 6 \cdot 0 + (7 \cdot 2 + 4)) \cdot 10](13, 15) = [11](13, 15) =$$

$(33, 9)$

$$\left[ \frac{5A_2(2) + 6B_2(2) + C_2(2)}{3} \right](13, 15) = [(5 \cdot (6 \cdot 2 + 10) + 6 \cdot 0 + 0) \cdot 9](13, 15) = [2](13, 15) =$$

$(33, 34)$

$$\left[ \frac{5A_3(2) + 6B_3(2) + C_3(2)}{3} \right](13, 15) = [(5 \cdot 0 + 6 \cdot (6 \cdot 2 + 10) + 0) \cdot 9](13, 15) = [5](13, 15) =$$

$(26, 34)$

$$\left[ \frac{5A_4(2) + 6B_4(2) + C_4(2)}{3} \right](13, 15) = [(5 \cdot 0 + 6 \cdot (7 \cdot 2 + 4) + 0) \cdot 9](13, 15) = [10](13, 15) =$$

$(38, 28)$

$$\left[ \frac{5A_5(2) + 6B_5(2) + C_5(2)}{3} \right](13, 15) = [(5 \cdot (7 \cdot 2 + 4) + 6 \cdot 0 + 0) \cdot 9](13, 15) = [4](13, 15) =$$

$(35, 28)$

$$\left[ \frac{2^0 \cdot T(2)}{3} \right](13, 15) = [1 \cdot (2^2 + 2 + 9) \cdot 9](13, 15) = [5](13, 15) =$$

$(26, 34)$

Putting all those values together, we see that the  $\mathbb{G}_1$  part of the common reference string is given by the following set of 12 points from the  $\text{BLS6\_6}$  13-torsion group  $\mathbb{G}_1$ :

$$CRS_{\mathbb{G}_1} = \left\{ \begin{array}{l} (27, 34), (26, 34), (38, 15), \left( (13, 15), (33, 34) \right), \left( \emptyset, (33, 9) \right) \\ \left( (33, 34), (26, 34), (38, 28), (35, 28) \right), \left( (26, 34) \right) \end{array} \right\}$$

To compute the  $\mathbb{G}_2$  part of the common reference string, we use the logarithmic order of the group  $\mathbb{G}_2$  XXX, the generator  $g_2 = (7v^2, 16v^3)$ , as well as the values from the simulation backdoor. Since  $\deg(T) = 2$ , we get the following:

$$[\beta]g_2 = [5](7v^2, 16v^3) = (16v^2, 28v^3)$$

$$[\gamma]g_2 = [4](7v^2, 16v^3) = (37v^2, 27v^3)$$

$$[\delta]g_2 = [3](7v^2, 16v^3) = (42v^2, 16v^3)$$

To compute the rest of the  $\mathbb{G}_2$  part of the common reference string, we expand the indexed tuple and insert the secret random elements from the simulation backdoor. We get the following:

$$\begin{aligned} \left( [s^j]g_2, \dots \right)_{j=0}^1 &= ([2^0](7v^2, 16v^3), [2^1](7v^2, 16v^3)) \\ &= ((7v^2, 16v^3), (10v^2, 28v^3)) \end{aligned}$$

Putting all these values together, we see that the  $\mathbb{G}_2$  part of the common reference string is given by the following set of 5 points from the  $\text{BLS6\_6}$  13-torsion group  $\mathbb{G}_2$ :

$$CRS_{\mathbb{G}_2} = \left\{ (16v^2, 28v^3), (37v^2, 27v^3), (42v^2, 16v^3), (7v^2, 16v^3), (10v^2, 28v^3) \right\}$$

add reference

Given the simulation trapdoor  $\tau$  and the quadratic arithmetic program 6.14, the associated common reference string of the 3-factorization problem is as follows:

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$$\begin{aligned} CRS_{\mathbb{G}_1} &= \left\{ (27, 34), (26, 34), (38, 15), \left( (13, 15), (33, 34) \right), \left( \mathcal{O}, (33, 9) \right) \right\} \\ &\quad \left\{ \left( (33, 34), (26, 34), (38, 28), (35, 28) \right), \left( (26, 34) \right) \right\} \\ CRS_{\mathbb{G}_2} &= \left\{ (16v^2, 28v^3), (37v^2, 27v^3), (42v^2, 16v^3), \left( 7v^2, 16v^3 \right), (10v^2, 28v^3) \right\} \end{aligned}$$

5748 We then publish this data to everyone who wants to participate in the generation of a zk-SNARK  
5749 or its verification in the 3-factorization problem.

5750 To understand how this common reference string can be used to evaluate polynomials at  
5751 the secret evaluation point in the exponent of a generator, let's assume that we have deleted the  
5752 simulation trapdoor. In that case, we have no way to know the secret evaluation point anymore,  
5753 hence, we cannot evaluate polynomials at that point. However, we can evaluate polynomials of  
5754 smaller degree than the degree of the target polynomial in the exponent of both generators at  
5755 that point.

To see that, consider e.g. the polynomials  $A_2(x) = 6x + 10$  and  $A_5(x) = 7x + 4$  from the QAP of this problem. To evaluate these polynomials in the exponent of  $g_1$  and  $g_2$  at the secret point  $s$  without knowing the value of  $s$  (which is 2), we can use the common reference string and equation XXX. Using the scalar product notation instead of exponentiation, we get the following:

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$$\begin{aligned} [A_2(s)]g_1 &= [6 \cdot s^1 + 10 \cdot s^0]g_1 \\ &= [6](33, 34) + [10](13, 15) & \# [s^0]g_1 = (13, 15), [s^1]g_1 = (33, 34) \\ &= [6 \cdot 2](13, 15) + [10](13, 15) = [9](13, 15) & \# \text{logarithmic order on } \mathbb{G}_1 \\ &= (35, 15) \\ [A_5(s)]g_1 &= [7 \cdot s^1 + 4 \cdot s^0]g_1 \\ &= [7](33, 34) + [4](13, 15) \\ &= [7 \cdot 2](13, 15) + [4](13, 15) = [5](13, 15) \\ &= (26, 34) \end{aligned}$$

Indeed, we are able to evaluate the polynomials in the exponent at a secret evaluation point, because that point is encrypted in the curve point  $(33, 34)$  and its secrecy is protected by the discrete logarithm assumption. Of course, in our computation, we recovered the secret point  $s = 2$ , but that was only possible because we have a group of logarithmic order in order to simplify our pen-and-paper computations. Such an order is infeasible for computing in cryptographically secure curves. We can do the same computation on  $\mathbb{G}_2$  and get the following:

$$\begin{aligned} [A_2(s)]g_2 &= [6 \cdot s^1 + 10 \cdot s^0]g_2 \\ &= [6](10v^2, 28v^3) + [10](7v^2, 16v^3) \\ &= [6 \cdot 2](7v^2, 16v^3) + [10](7v^2, 16v^3) = [9](7v^2, 16v^3) \\ &= (37v^2, 16v^3) \\ [A_5(s)]g_2 &= [7 \cdot s^1 + 4 \cdot s^0]g_2 \\ &= [7](10v^2, 28v^3) + [4](7v^2, 16v^3) \\ &= [7 \cdot 2](7v^2, 16v^3) + [4](7v^2, 16v^3) = [5](7v^2, 16v^3) \\ &= (16v^2, 28v^3) \end{aligned}$$

Apart from the target polynomial  $T$ , all other polynomials of the quadratic arithmetic program can be evaluated in the exponent this way.

**The Prover Phase** Given some rank-1 constraint system  $R$  and instance  $I = (I_1, \dots, I_n)$ , the task of the prover phase is to convince any verifier that a prover knows a witness  $W$  to instance  $I$  such that  $(I; W)$  is a word in the language  $L_R$  of the system, without revealing anything about  $W$ .

To achieve this in the Groth\_16 protocol, we assume that any prover has access to the rank-1 constraint system of the problem, in addition to some algorithm that tells the prover how to compute constructive proofs for the R1CS. In addition, the prover has access to a common reference string and its associated quadratic arithmetic program.

In order to generate a zk-SNARK for this instance, the prover first computes a valid constructive proof as explained in XXX, that is, the prover generates a proper witness  $W = (W_1, \dots, W_m)$  such that  $(I_1, \dots, I_n; W_1, \dots, W_m)$  is a solution to the rank-1 constraint system  $R$ .

The prover then uses the quadratic arithmetic program and computes the polynomial  $P_{(I;W)}$ , as explained in 6.15. They then divide  $P_{(I;W)}$  by the target polynomial  $T$  of the quadratic arithmetic program. Since  $P_{(I;W)}$  is constructed from a valid solution to the R1CS, we know from 6.15 that it is divisible by  $T$ . This implies that polynomial division of  $P$  by  $T$  generates another polynomial  $H := P/T$ , with  $\deg(H) < \deg(T)$ .

The prover then evaluates the polynomial  $(H \cdot T)\delta^{-1}$  in the exponent of the generator  $g_1$  at the secret point  $s$ , as explained in XXX. To see how this can be achieved, let  $H(x)$  be the quotient polynomial  $P/T$ :

$$H(x) = H_0 \cdot x^0 + H_1 \cdot x^1 + \dots + H_k \cdot x^k \quad (8.3)$$

To evaluate  $H \cdot T$  at  $s$  in the exponent of  $g_1$ , the prover uses the common reference string and computes as follows:

$$g_1^{\frac{H(s) \cdot T(s)}{\delta}} = \left(g_1^{\frac{s^0 \cdot T(s)}{\delta}}\right)^{H_0} \cdot \left(g_1^{\frac{s^1 \cdot T(s)}{\delta}}\right)^{H_1} \dots \left(g_1^{\frac{s^k \cdot T(s)}{\delta}}\right)^{H_k}$$

After this has been done, the prover samples two random field elements  $r, t \in \mathbb{F}_r$ , and uses the common reference string, the instance variables  $I_1, \dots, I_n$  and the witness variables  $W_1, \dots, W_m$  to compute the following curve points:

$$\begin{aligned} g_1^W &= \left(g_1^{\frac{\beta \cdot A_{1+n}(s) + \alpha \cdot B_{1+n}(s) + C_{1+n}(s)}{\delta}}\right)^{W_1} \dots \left(g_1^{\frac{\beta \cdot A_{m+n}(s) + \alpha \cdot B_{m+n}(s) + C_{m+n}(s)}{\delta}}\right)^{W_m} \\ g_1^A &= g_1^\alpha \cdot g_1^{A_0(s)} \cdot \left(g_1^{A_1(s)}\right)^{I_1} \dots \left(g_1^{A_n(s)}\right)^{I_n} \cdot \left(g_1^{A_{n+1}(s)}\right)^{W_1} \dots \left(g_1^{A_{n+m}(s)}\right)^{W_m} \cdot \left(g_1^\delta\right)^r \\ g_1^B &= g_1^\beta \cdot g_1^{B_0(s)} \cdot \left(g_1^{B_1(s)}\right)^{I_1} \dots \left(g_1^{B_n(s)}\right)^{I_n} \cdot \left(g_1^{B_{n+1}(s)}\right)^{W_1} \dots \left(g_1^{B_{n+m}(s)}\right)^{W_m} \cdot \left(g_1^\delta\right)^t \\ g_2^B &= g_2^\beta \cdot g_2^{B_0(s)} \cdot \left(g_2^{B_1(s)}\right)^{I_1} \dots \left(g_2^{B_n(s)}\right)^{I_n} \cdot \left(g_2^{B_{n+1}(s)}\right)^{W_1} \dots \left(g_2^{B_{n+m}(s)}\right)^{W_m} \cdot \left(g_2^\delta\right)^t \\ g_1^C &= g_1^W \cdot g_1^{\frac{H(s) \cdot T(s)}{\delta}} \cdot \left(g_1^A\right)^t \cdot \left(g_1^B\right)^r \cdot \left(g_1^\delta\right)^{-r \cdot t} \end{aligned}$$

In this computation, the group elements  $g_1^{A_j(s)}$ ,  $g_1^{B_j(s)}$  and  $g_2^{B_j(s)}$  can be derived from the common reference string and the quadratic arithmetic program of the problem, as we have seen in XXX. In fact, those points only have to be computed once, and can be published and reused

for multiple proof generations because they are the same for all instances and witnesses. All other group elements are part of the common reference string.

After all these computations have been done, a valid zero-knowledge succinct non-interactive argument of knowledge  $\pi$  in the Groth\_16 protocol is given by the following three curve points:

$$\pi = (g_1^A, g_1^C, g_2^B) \quad (8.4)$$

As we can see, a Groth\_16 zk-SNARK consists of 3 curve points, two points from  $\mathbb{G}_1$  and 1 point from  $\mathbb{G}_2$ . The argument is specifically designed this way because, in typical applications,  $\mathbb{G}_1$  is a torsion group of an elliptic curve over some prime field, while  $\mathbb{G}_2$  is a subgroup of a torsion group over an extension field. Elements from  $\mathbb{G}_1$  therefore need less space to be stored, and computations in  $\mathbb{G}_1$  are typically faster than in  $\mathbb{G}_2$ .

Since the witness is encoded in the exponent of a generator of a cryptographically secure elliptic curve, it is hidden from anyone but the prover. Moreover, since any proof is randomized by the occurrence of the random field elements  $r$  and  $t$ , proofs are not unique to any given witness. This is an important feature because, if all proofs for the same witness would be the same, knowledge of a witness would destroy the zero-knowledge property of those proofs.

*Example 138* (The 3-factorization Problem). To see how a prover might compute a zk-SNARK, consider the 3-factorization problem from 106, our protocol parameters from XXX as well as the common reference string from XXX.

Our task is to compute a zk-SNARK for the instance  $I_1 = 11$  and its constructive proof  $(W_1, W_2, W_3, W_4) = (2, 3, 4, 6)$  as computed in XXX. As we know from 6.15, the associated polynomial  $P_{(I;W)}$  of the quadratic arithmetic program from XXX is given by the following equation:

$$P_{(I;W)} = x^2 + x + 9$$

Since  $P_{(I;W)}$  is identical to the target polynomial  $T(x) = x^2 + x + 9$  in this example, we know from XXX that the quotient polynomial  $H = P/T$  is the constant degree 0 polynomial:

$$H(x) = H_0 \cdot x^0 = 1 \cdot x^0$$

We therefore use  $[\frac{s^0 \cdot T(s)}{\delta}]_{g_1} = (26, 34)$  from our common reference string XXX of the 3-factorization problem and compute as follows:

$$\begin{aligned} [\frac{H(s) \cdot T(s)}{\delta}]_{g_1} &= [H_0](26, 34) = [1](26, 34) \\ &= (26, 34) \end{aligned}$$

In the next step, we have to compute all group elements required for a proper Groth16 zk-SNARK. We start with  $g_1^W$ . Using scalar products instead of the exponential notation, and  $\oplus$  for the group law on the BLS6\_6 curve, we have to compute the point  $[W]g_1$ :

$$\begin{aligned} [W]g_1 &= [W_1]g_1 \frac{\beta \cdot A_2(s) + \alpha \cdot B_2(s) + C_2(s)}{\delta} \oplus [W_2]g_1 \frac{\beta \cdot A_3(s) + \alpha \cdot B_3(s) + C_3(s)}{\delta} \oplus [W_3]g_1 \frac{\beta \cdot A_4(s) + \alpha \cdot B_4(s) + C_4(s)}{\delta} \\ &\quad \oplus [W_4]g_1 \frac{\beta \cdot A_5(s) + \alpha \cdot B_5(s) + C_5(s)}{\delta} \end{aligned}$$

To compute this point, we have to remember that a prover should not be in possession of the simulation trapdoor, hence, they do not know what  $\alpha$ ,  $\beta$ ,  $\delta$  and  $s$  are. In order to compute this group element, the prover therefore needs the common reference string. Using the logarithmic order from XXX and the witness, we get the following:

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$$\begin{aligned}
[W]_{g_1} &= [2](33, 34) \oplus [3](26, 34) \oplus [4](38, 28) \oplus [6](35, 28) \\
&= [2 \cdot 2](13, 15) \oplus [3 \cdot 5](13, 15) \oplus [4 \cdot 10](13, 15) \oplus [6 \cdot 4](13, 15) \\
&= [2 \cdot 2 + 3 \cdot 5 + 4 \cdot 10 + 6 \cdot 4](13, 15) = [5](13, 15) \\
&= (26, 34)
\end{aligned}$$

5801 In a next step, we compute  $g_1^A$ . We sample the random point  $r = 11$  from  $\mathbb{F}_{13}$ , using scalar  
 5802 products instead of the exponential notation, and  $\oplus$  for the group law on the BLS6\_6 curve.  
 5803 We then have to compute the following expression:

$$\begin{aligned}
[A]_{g_1} &= [\alpha]_{g_1} \oplus [A_0(s)]_{g_1} \oplus [I_1][A_1(s)]_{g_1} \oplus [W_1][A_2(s)]_{g_1} \oplus [W_2][A_3(s)]_{g_1} \\
&\quad \oplus [W_3][A_4(s)]_{g_1} \oplus [W_4][A_5(s)]_{g_1} \oplus [r][\delta]_{g_1}
\end{aligned}$$

Since we don't know what  $\alpha$ ,  $\delta$  and  $s$  are, we look up  $[\alpha]_{g_1}$  and  $[\delta]_{g_1}$  from the common reference string. Recall from XXX that we can evaluate  $[A_j(s)]_{g_1}$  without knowing the secret evaluation point  $s$ . According to XXX, we have  $[A_2(s)]_{g_1} = (35, 15)$ ,  $[A_5(s)]_{g_1} = (26, 34)$  and  $[A_j(s)]_{g_1} = \mathcal{O}$  for all other indices  $0 \leq j \leq 5$ . Since  $\mathcal{O}$  is the neutral element on  $\mathbb{G}_1$ , we get the following:

$$\begin{aligned}
[A]_{g_1} &= (27, 34) \oplus \mathcal{O} \oplus [11]\mathcal{O} \oplus [2](35, 15) \oplus [3]\mathcal{O} \oplus [4]\mathcal{O} \oplus [6](26, 34) \oplus [11](38, 15) \\
&= (27, 34) \oplus [2](35, 15) \oplus [6](26, 34) \oplus [11](38, 15) \\
&= [6](13, 15) \oplus [2 \cdot 9](13, 15) \oplus [6 \cdot 5](13, 15) \oplus [11 \cdot 3](13, 15) \\
&= [6 + 2 \cdot 9 + 6 \cdot 5 + 11 \cdot 3](13, 15) = [9](13, 15) \\
&= (35, 15)
\end{aligned}$$

5804 In order to compute the two curve points  $[B]_{g_1}$  and  $[B]_{g_2}$ , we sample another random element  
 5805  $t = 4$  from  $\mathbb{F}_{13}$ . Using the scalar product instead of the exponential notation, and  $\oplus$  for the group  
 5806 law on the BLS6\_6 curve, we have to compute the following expressions:

$$\begin{aligned}
[B]_{g_1} &= [\beta]_{g_1} \oplus [B_0(s)]_{g_1} \oplus [I_1][B_1(s)]_{g_1} \oplus [W_1][B_2(s)]_{g_1} \oplus [W_2][B_3(s)]_{g_1} \\
&\quad \oplus [W_3][B_4(s)]_{g_1} \oplus [W_4][B_5(s)]_{g_1} \oplus [t][\delta]_{g_1} \\
[B]_{g_2} &= [\beta]_{g_2} \oplus [B_0(s)]_{g_2} \oplus [I_1][B_1(s)]_{g_2} \oplus [W_1][B_2(s)]_{g_2} \oplus [W_2][B_3(s)]_{g_2} \\
&\quad \oplus [W_3][B_4(s)]_{g_2} \oplus [W_4][B_5(s)]_{g_2} \oplus [t][\delta]_{g_2}
\end{aligned}$$

Since we don't know what  $\beta$ ,  $\delta$  and  $s$  are, we look up the associated group elements from the common reference string. Recall from XXX that we can evaluate  $[B_j(s)]_{g_1}$  without knowing the secret evaluation point  $s$ . Since  $B_3 = A_2$  and  $B_4 = A_5$ , we have  $[B_3(s)]_{g_1} = (35, 15)$ ,  $[B_4(s)]_{g_1} = (26, 34)$  according to XXX, and  $[B_j(s)]_{g_1} = \mathcal{O}$  for all other indices  $0 < j < 5$ . Since  $\mathcal{O}$  is the neutral element on  $\mathbb{G}_1$ , we get the following:

$$\begin{aligned}
[B]_{g_1} &= (26, 34) \oplus \mathcal{O} \oplus [11]\mathcal{O} \oplus [2]\mathcal{O} \oplus [3](35, 15) \oplus [4](26, 34) \oplus [6]\mathcal{O} \oplus [4](38, 15) \\
&= (26, 34) \oplus [3](35, 15) \oplus [4](26, 34) \oplus [4](38, 15) \\
&= [5](13, 15) \oplus [3 \cdot 9](13, 15) \oplus [4 \cdot 5](13, 15) \oplus [4 \cdot 3](13, 15) \\
&= [5 + 3 \cdot 9 + 4 \cdot 5 + 4 \cdot 3](13, 15) = [12](13, 15) \\
&= (13, 28)
\end{aligned}$$



$$\begin{aligned}
 [B]g_2 &= (16v^2, 28v^3) \oplus \mathcal{O} \oplus [11]\mathcal{O} \oplus [2]\mathcal{O} \oplus [3](37v^2, 16v^3) \oplus [4](16v^2, 28v^3) \oplus [6]\mathcal{O} \oplus [4](42v^2, 16v^3) \\
 &= (16v^2, 28v^3) \oplus [3](37v^2, 16v^3) \oplus [4](16v^2, 28v^3) \oplus [4](42v^2, 16v^3) \\
 &= [5](7v^2, 16v^3) \oplus [3 \cdot 9](7v^2, 16v^3) \oplus [4 \cdot 5](7v^2, 16v^3) \oplus [4 \cdot 3](7v^2, 16v^3) \\
 &= [5 + 3 \cdot 9 + 4 \cdot 5 + 4 \cdot 3](7v^2, 16v^3) = [12](7v^2 + 16v^3) \\
 &= (7v^2, 27v^3)
 \end{aligned}$$

In a last step, we combine the previous computations to compute the point  $[C]g_1$  in the group  $\mathbb{G}_1$  as follows:

$$\begin{aligned}
 [C]g_1 &= [W]g_1 \oplus \left[ \frac{H(s) \cdot T(s)}{\delta} \right]g_1 \oplus [t][A]g_1 \oplus [r][B]g_1 \oplus [-r \cdot t][\delta]g_1 \\
 &= (26, 34) \oplus (26, 34) \oplus [4](35, 15) \oplus [11](13, 28) \oplus [-11 \cdot 4](38, 15) \\
 &= [5](13, 15) \oplus [5](13, 15) \oplus [4 \cdot 9](13, 15) \oplus [11 \cdot 12](13, 15) \oplus [-11 \cdot 4 \cdot 3](13, 15) \\
 &= [5 + 5 + 4 \cdot 9 + 11 \cdot 12 - 11 \cdot 4 \cdot 3](13, 15) = [7](13, 15) \\
 &= (27, 9)
 \end{aligned}$$

Given the instance  $I_1 = 11$ , we can now combine these computations and see that the following 3 curve points are a zk-SNARK for the witness  $(W_1, W_2, W_3, W_4) = (2, 3, 4, 6)$ :

$$\pi = ((35, 15), (27, 9), (7v^2, 27v^3))$$

5807 We can now publish this zk-SNARK, or send it to a designated verifier. Note that, if we had  
 5808 sampled different values for  $r$  and  $t$ , we would have computed a different SNARK for the same  
 5809 witness. The SNARK, therefore, hides the witness perfectly, which means that it is impossible  
 5810 to reconstruct the witness from the SNARK.

5811 **The Verification Phase** Given some rank-1 constraint system  $R$ , instance  $I = (I_1, \dots, I_n)$  and  
 5812 zk-SNARK  $\pi$ , the task of the verification phase is to check that  $\pi$  is indeed an argument for  
 5813 a constructive proof. Assuming that the simulation trapdoor does not exist anymore and the  
 5814 verification checks the proof, the verifier can be convinced that someone knows a witness  $W =$   
 5815  $(W_1, \dots, W_m)$  such that  $(I; W)$  is a word in the language of  $R$ .

To achieve this in the Groth16 protocol, we assume that any verifier is able to compute the pairing map  $e(\cdot, \cdot)$  efficiently, and has access to the common reference string used to produce the SNARK  $\pi$ . In order to verify the SNARK with respect to the instance  $(I_1, \dots, I_n)$ , the verifier computes the following curve point:

$$g_1^I = \left( g_1^{\frac{\beta \cdot A_0(s) + \alpha \cdot B_0(s) + C_0(s)}{\gamma}} \right) \cdot \left( g_1^{\frac{\beta \cdot A_1(s) + \alpha \cdot B_1(s) + C_1(s)}{\gamma}} \right)^{I_1} \cdots \left( g_1^{\frac{\beta \cdot A_n(s) + \alpha \cdot B_n(s) + C_n(s)}{\gamma}} \right)^{I_n}$$

5816 With this group element, the verifier is able to verify the SNARK  $\pi = (g_1^A, g_1^C, g_2^B)$  by checking  
 5817 the following equation using the pairing map:

$$e(g_1^A, g_2^B) = e(g_1^\alpha, g_2^\beta) \cdot e(g_1^I, g_2^\gamma) \cdot e(g_1^C, g_2^\delta) \quad (8.5)$$

5818 If the equation holds true, the SNARK is accepted. If the equation does not hold, the  
 5819 SNARK is rejected.



*Remark 5.* We know from chapter 5 that computing pairings in cryptographically secure pairing groups is computationally expensive. As we can see, in the Groth16 protocol, 3 pairings are required to verify the SNARK, because the pairing  $e(g_1^\alpha, g_2^\beta)$  is independent of the proof, meaning that can be computed once and then stored as an amendment to the verifier key.

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In Groth [2016], the author showed that 2 is the minimal amount of pairings that any protocol with similar properties has to use. This protocol is therefore close to the theoretical minimum. In the same paper, the author outlined an adaptation that only uses 2 pairings. However, that reduction comes with the price of much more overhead computation. Having 3 pairings is therefore a compromise that gives the overall best performance. To date, the Groth16 protocol is the most efficient in its class.

*Example 139* (The 3-factorization Problem). To see how a verifier might check a zk-SNARK for some given instance  $I$ , consider the 3-factorization problem from 106, our protocol parameters from XXX, the common reference string from XXX as well as the zk-SNARK  $\pi = ((35, 15), (27, 9), (7v^2, 27v^3))$ , which claims to be an argument of knowledge for a witness for the instance  $I_1 = 11$ .

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In order to verify the zk-SNARK for that instance, we first compute the curve point  $g_1^I$ . Using scalar products instead of the exponential notation, and  $\oplus$  for the group law on the BLS6\_6 curve, we have to compute the point  $[I]g_1$  as follows:

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$$[I]g_1 = \left[ \frac{\beta \cdot A_0(s) + \alpha \cdot B_0(s) + C_0(s)}{\gamma} \right]_{g_1} \oplus [I_1] \left[ \frac{\beta \cdot A_1(s) + \alpha \cdot B_1(s) + C_1(s)}{\gamma} \right]_{g_1}$$

To compute this point, we have to remember that a verifier should not be in possession of the simulation trapdoor, which means that they do not know what  $\alpha$ ,  $\beta$ ,  $\gamma$  and  $s$  are. In order to compute this group element, the verifier therefore needs the common reference string. Using the logarithmic order from XXX and instance  $I_1$ , we get the following:

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$$\begin{aligned} [I]g_1 &= \left[ \frac{\beta \cdot A_0(s) + \alpha \cdot B_0(s) + C_0(s)}{\gamma} \right]_{g_1} \oplus [I_1] \left[ \frac{\beta \cdot A_1(s) + \alpha \cdot B_1(s) + C_1(s)}{\gamma} \right]_{g_1} \\ &= \mathcal{O} \oplus [11](33, 9) \\ &= [11 \cdot 11](13, 15) = [4](13, 15) \\ &= (35, 28) \end{aligned}$$

In the next step, we have to compute all the pairings involved in equation XXX. Using the logarithmic order on  $\mathbb{G}_1$  and  $\mathbb{G}_2$  as well as the bilinearity property of the pairing map we get the following:

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$$\begin{aligned}
e([A]_{g_1}, [B]_{g_2}) &= e((35, 15), (7v^2, 27v^3)) = e([9](13, 15), [12](7v^2, 16v^3)) \\
&= e((13, 15), (7v^2, 16v^3))^{9 \cdot 12} \\
&= e((13, 15), (7v^2, 16v^3))^{108} \\
e([\alpha]_{g_1}, [\beta]_{g_2}) &= e((27, 34), (16v^2, 28v^3)) = e([6](13, 15), [5](7v^2, 16v^3)) \\
&= e((13, 15), (7v^2, 16v^3))^{6 \cdot 5} \\
&= e((13, 15), (7v^2, 16v^3))^{30} \\
e([I]_{g_1}, [\gamma]_{g_2}) &= e((35, 28), (37v^2, 27v^3)) = e([4](13, 15), [4](7v^2, 16v^3)) \\
&= e((13, 15), (7v^2, 16v^3))^{4 \cdot 4} \\
&= e((13, 15), (7v^2, 16v^3))^{16} \\
e([C]_{g_1}, [\delta]_{g_2}) &= e((27, 9), (42v^2, 16v^3)) = e([7](13, 15), [3](7v^2, 16v^3)) \\
&= e((13, 15), (7v^2, 16v^3))^{7 \cdot 3} \\
&= e((13, 15), (7v^2, 16v^3))^{21}
\end{aligned}$$

5845 In order to check equation XXX, observe that the target group  $\mathbb{G}_T$  of the Weil pairing is  
 5846 a finite cyclic group of order 13. Exponentiation is therefore done in modular 13 arithmetics.  
 5847 Accordingly, since  $108 \bmod 13 = 4$ , we evaluate the left side of equation XXX as follows:

$$e([A]_{g_1}, [B]_{g_2}) = e((13, 15), (7v^2, 16v^3))^{108} = e((13, 15), (7v^2, 16v^3))^4$$

5848 Similarly, we evaluate the right side of equation XXX using modular 13 arithmetics and the  
 5849 exponential law  $a^x \cdot a^y = a^{x+y}$ :

$$\begin{aligned}
&e([\alpha]_{g_1}, [\beta]_{g_2}) \cdot e([I]_{g_1}, [\gamma]_{g_2}) \cdot e([C]_{g_1}, [\delta]_{g_2}) = \\
&e((13, 15), (7v^2, 16v^3))^{30} \cdot e((13, 15), (7v^2, 16v^3))^{16} \cdot e((13, 15), (7v^2, 16v^3))^{21} = \\
&e((13, 15), (7v^2, 16v^3))^4 \cdot e((13, 15), (7v^2, 16v^3))^3 \cdot e((13, 15), (7v^2, 16v^3))^8 = \\
&e((13, 15), (7v^2, 16v^3))^{4+3+8} = \\
&e((13, 15), (7v^2, 16v^3))^2
\end{aligned}$$

5850 As we can see, both the left and the right side of equation XXX are identical, which implies  
 5851 that the verification process accepts the simulated proof.

5852 **NOTE: UNFORTUNATELY NOT! :-( (HENCE THERE IS AN ERROR SOMEWHERE ...**  
 5853 **NEED TO FIX IT AFTER VACATION**

5854 **Proof Simulation** During the execution of a setup phase, a common reference string is gen-  
 5855 erated, along with a simulation trapdoor, the latter of which must be deleted at the end of the  
 5856 setup-phase. As an alternative, a more complicated multi-party protocol like [XXX] can be  
 5857 used to split the knowledge of the simulation trapdoor among many different parties.

5858 In this paragraph, we will show why knowledge of the simulation trapdoor is problematic,  
 5859 and how it can be used to generate zk-SNARKs for a given instance without any knowledge or  
 5860 the existence of associated witness.

5861 To be more precise, let  $I$  be an instance for some R1CS language  $L_R$ . We call a zk-SNARK  
 5862 for  $L_R$  **forged** or **simulated** if it passes a verification but its generation does not require the  
 5863 existence of a witness  $W$  such that  $(I; W)$  is a word in  $L_R$ .

To see how simulated zk-SNARKs can be computed, assume that a forger has knowledge of proper Groth\_16 parameters, a quadratic arithmetic program of the problem, a common reference string and its associated simulation trapdoor  $\tau$ :

$$\tau = (\alpha, \beta, \gamma, \delta, s) \quad (8.6)$$

Given some instance  $I$ , the forger’s task is to generate a zk-SNARK for this instance that passes the verification process, without having access to any other zk-SNARKs for this instance and without knowledge of a valid witness  $W$ .

To achieve this in the Groth\_16 protocol, the forger can use the simulation trapdoor in combination with the QAP and two arbitrary field elements  $A$  and  $B$  from the scalar field  $\mathbb{F}_r$  of the pairing groups to  $g_1^C$  compute for the instance  $(I_1, \dots, I_n)$  as follows:

$$g_1^C = g_1^{\frac{A \cdot B}{\delta}} \cdot g_1^{-\frac{\alpha \cdot \beta}{\delta}} \cdot g_1^{-\frac{\beta A_0(s) + \alpha B_0(s) + C_0(s)}{\delta}} \cdot \left( g_1^{-\frac{\beta A_1(s) + \alpha B_1(s) + C_1(s)}{\delta}} \right)^{I_1} \cdots \left( g_1^{-\frac{\beta A_n(s) + \alpha B_n(s) + C_n(s)}{\delta}} \right)^{I_n}$$

The forger then publishes the zk-SNARK  $\pi_{\text{forged}} = (g_1^A, g_1^C, g_2^B)$ , which will pass the verification process and is computable without the existence of a witness  $(W_1, \dots, W_m)$ .

To see that the simulation trapdoor is necessary and sufficient to compute the simulated proof  $\pi_{\text{forged}}$ , first observe that both generators  $g_1$  and  $g_2$  are known to the forger, as they are part of the common reference string, encoded as  $g_1^{s^0}$  and  $g_2^{s^0}$ . The forger is therefore able to compute  $g_1^{A \cdot B}$ . Moreover, since the forger knows  $\alpha, \beta, \delta$  and  $s$  from the trapdoor, they are able to compute all factors in the computation of  $g_1^C$ .

If, on the other hand, the simulation trapdoor is unknown, it is not possible to compute  $g_1^C$ , since, for example, the computational Diffie-Hellman assumption makes the derivation of  $g_1^{\alpha \cdot \beta}$  from  $g_1^\alpha$  and  $g_1^\beta$  infeasible.

*Example 140* (The 3-factorization Problem). To see how a forger might simulate a zk-SNARK for some given instance  $I$ , consider the 3-factorization problem from 106, our protocol parameters from XXX, the common reference string from XXX and the simulation trapdoor  $\tau = (6, 5, 4, 3, 2)$  of that CRS.

In order to forge a zk-SNARK for instance  $I_1 = 11$ , we don’t need a constructive proof for the associated rank-1 constraint system, which implies that we don’t have to execute the circuit  $C_{3.\text{fac}}(\mathbb{F}_{13})$ . Instead, we have to choose 2 arbitrary elements  $A$  and  $B$  from  $\mathbb{F}_{13}$ , and compute  $g_1^A, g_2^B$  and  $g_1^C$  as defined in XXX. We choose  $A = 9$  and  $B = 3$ , and, since  $\delta^{-1} = 3$ , we compute as follows:

$$\begin{aligned} [A]g_1 &= [9](13, 15) = (35, 15) \\ [B]g_2 &= [3](7v^2, 16v^3) = (42v^2, 16v^3) \\ [C]g_1 &= \left[ \frac{A \cdot B}{\delta} \right]g_1 \oplus \left[ -\frac{\alpha \cdot \beta}{\delta} \right]g_1 \oplus \left[ -\frac{\beta A_0(s) + \alpha B_0(s) + C_0(s)}{\delta} \right]g_1 \oplus \\ &\quad [I_1] \left[ -\frac{\beta A_1(s) + \alpha B_1(s) + C_1(s)}{\delta} \right]g_1 \\ &= [(9 \cdot 3) \cdot 9](13, 15) \oplus [-(6 \cdot 5) \cdot 9](13, 15) \oplus [0](13, 15) \oplus [11][-(7 \cdot 2 + 4) \cdot 9](13, 15) \\ &= [9](13, 15) \oplus [3](13, 15) \oplus [12](13, 15) = [11](13, 15) \\ &= (33, 9) \end{aligned}$$

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This is all we need to generate our forged proof for the 3-factorization problem. We publish the simulated zk-SNARK:

$$\pi_{fake} = ((35, 15), (33, 9), (42v^2, 16v^3))$$

Despite the fact that this zk-SNARK was generated without knowledge of a proper witness, it is indistinguishable from a zk-SNARK that proves knowledge of a proper witness.

To see that, we show that our forged SNARK passes the verification process. In order to verify  $\pi_{fake}$ , we proceed as in XXX, and compute the curve point  $g_1^I$  for the instance  $I_1 = 11$ . Since the instance is the same as in example XXX, we can parallel the computation from XXX:

$$\begin{aligned} [I]g_1 &= \left[ \frac{\beta \cdot A_0(s) + \alpha \cdot B_0(s) + C_0(s)}{\gamma} \right]_{g_1} \oplus [I_1] \left[ \frac{\beta \cdot A_1(s) + \alpha \cdot B_1(s) + C_1(s)}{\gamma} \right]_{g_1} \\ &= (35, 28) \end{aligned}$$

In a next step we have to compute all the pairings involved in equation XXX. Using the logarithmic order on  $\mathbb{G}_1$  and  $\mathbb{G}_2$  as well as the bilinearity property of the pairing map we get

$$\begin{aligned} e([A]g_1, [B]g_2) &= e((35, 15), (42v^2, 16v^3)) = e([9](13, 15), [3](7v^2, 16v^3)) \\ &= e((13, 15), (7v^2, 16v^3))^{9 \cdot 3} \\ &= e((13, 15), (7v^2, 16v^3))^{27} \\ e([\alpha]g_1, [\beta]g_2) &= e((27, 34), (16v^2, 28v^3)) = e([6](13, 15), [5](7v^2, 16v^3)) \\ &= e((13, 15), (7v^2, 16v^3))^{6 \cdot 5} \\ &= e((13, 15), (7v^2, 16v^3))^{30} \\ e([I]g_1, [\gamma]g_2) &= e((35, 28), (37v^2, 27v^3)) = e([4](13, 15), [4](7v^2, 16v^3)) \\ &= e((13, 15), (7v^2, 16v^3))^{4 \cdot 4} \\ &= e((13, 15), (7v^2, 16v^3))^{16} \\ e([C]g_1, [\delta]g_2) &= e((33, 9), (42v^2, 16v^3)) = e([11](13, 15), [3](7v^2, 16v^3)) \\ &= e((13, 15), (7v^2, 16v^3))^{11 \cdot 3} \\ &= e((13, 15), (7v^2, 16v^3))^{33} \end{aligned}$$

In order to check equation XXX, observe that the target group  $\mathbb{G}_T$  of the Weil pairing is a finite cyclic group of order 13. Exponentiation is therefore done in modular 13 arithmetics. Using this, we evaluate the left side of equation XXX as follows:

$$e([A]g_1, [B]g_2) = e((13, 15), (7v^2, 16v^3))^{27} = e((13, 15), (7v^2, 16v^3))^1$$

since  $27 \bmod 13 = 1$ . Similarly, we evaluate the right side of equation XXX using modular 13 arithmetics and the exponential law  $a^x \cdot a^y = a^{x+y}$ . We get

$$\begin{aligned} e([\alpha]g_1, [\beta]g_2) \cdot e([I]g_1, [\gamma]g_2) \cdot e([C]g_1, [\delta]g_2) &= \\ e((13, 15), (7v^2, 16v^3))^{30} \cdot e((13, 15), (7v^2, 16v^3))^{16} \cdot e((13, 15), (7v^2, 16v^3))^{33} &= \\ e((13, 15), (7v^2, 16v^3))^4 \cdot e((13, 15), (7v^2, 16v^3))^3 \cdot e((13, 15), (7v^2, 16v^3))^7 &= \\ e((13, 15), (7v^2, 16v^3))^{4+3+7} &= \\ e((13, 15), (7v^2, 16v^3))^1 & \end{aligned}$$

5897 As we can see, both the left and the right side of equation XXX are identical, which implies  
5898 that the verification process accepts the simulated proof.  $\pi_{fake}$  therefore convinces the veri-  
5899 fier that a witness to the 3-factorization problem exists. However, no such witness was really  
5900 necessary to generate the proof.

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## 5901 **Chapter 9**

## 5902 **Exercises and Solutions**

5903 TODO: All exercises we provided should have a solution, which we give here in all detail.

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